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2a. Code (incomplete, using c\_hanoiH as recursive helper):

```
def c_hanoi(n):
    long(n, 0, 2)

def c_hanoiHr(n, a, b):
    if n == 0:
        return
    c = 3-a-b
    print(n)
    print('a: {}'.format(a))
    print('b: {}'.format(b))
    print('c: {}'.format(c))

    c_hanoiHr(n-1, a, c)
    print('Move disk from {} to {}'.format(a, b))
    c_hanoiHr(n-1, c, b)

def c_hanoiH(n, a, b):
    if n == 0:
        return
    c = 3-a-b

    if (b == 2):
        print('Move disk from {} to {}'.format(a, c))
        print('Move disk from {} to {}'.format(c, b))
        c_hanoiH(n-1, a, b)
    if (b == 1):
        print('Move disk from {} to {}'.format(b, a))
        c_hanoiH(n-1, b, c)

def c_hanoiH2(n, a, b):
    if n == 0:
        return
    c = 3-a-b

    c_hanoiH2(n-1, a, b)
    print('Move disk from {} to {}'.format(a, b))
    c_hanoiH2(n-1, a, c)
    print('Move disk from {} to {}'.format(c, b))
    c_hanoiH2(n-1, a, c)
```

Result (not correct):

```
>>> c_hanoi(4)
Move disk from 0 to 1
Move disk from 1 to 2
Move disk from 0 to 1
Move disk from 1 to 2
Move disk from 0 to 1
Move disk from 1 to 2
Move disk from 0 to 1
Move disk from 1 to 2
\\ \ I
```

2b.

3.

a.

- i.  $t(n) = t(n/2) + t(n/2) = 2t(n/2) + n = 2(2t(n/4) + 2n/4) + n = 4t(n/4) + 2n \dots 2^d$   
 $+(n/2^d) + dn$  where  $d = \log_2 n$ ,  $0 + dn = n \log n$ , same as mergesort
- ii. Wolfram alpha output:

$$t(n) = t(n/2) + t(n/2) + n, t(1) = 0$$

 Extended Keyboard  Upload

Input:

$$t(n) = 2t\left(\frac{n}{2}\right) + n \mid t(1) = 0$$

Result:

$$\left\{ t(n) = n + 2t\left(\frac{n}{2}\right), t(1) = 0 \right\}$$

Recurrence equation solution:

$$t(n) = \frac{n \log(n)}{\log(2)}$$

iii.  $O(n \log n)$

- b. I would use merge sort to start splitting lists up, when I have the list of elements smaller than the right side that are added to the out array, I can get the number of inversions for that list by subtracting the length of the left side minus the iterator for left side
- c. In my code it went to 1000000 before it took long
- d. Code:

```

def invert(lst):
    return inversions(lst)[1]
def inversions(lst):
    #merge and mergesort code from week2.py
    if len(lst) <= 1:
        return lst, 0
    mid = len(lst)//2
    L, l = inversions(lst[:mid])
    R, r = inversions(lst[mid:])

    i, j = 0, 0
    out = []
    count = 0 + l + r;

    while i < len(L) and j < len(R):
        if L[i] < R[j]:
            out.append(L[i])
            i += 1
        else:
            out.append(R[j])
            j += 1
            count+= (len(L)-i)
    if i == len(L):
        out += R[j:]
    else:
        out += L[i:]
    return out, count

```

Results:

```

>>> invert([4,3,5,2,1])
8
>>> invert([3,1,2])
2
>>> invert([1,2])
0
>>> invert([2,1])
1

```

- e. Make length of the longlist  $n$   
The inversion number will be around  $n \log n$