# COMPUTER SCIENCE MENTORS 61A

November 4 – November 8, 2024

# Scheme

Scheme is a functional language, as opposed to Python, which is an imperative language. A Python program is comprised of statements or "instructions," each of which directs the computer to take some action. (An example of a statement would be something like x = 3 in Python. This does not evaluate to any value but just instructs Python to create a variable x with the value 3.) In contrast, a Scheme program is composed solely of (often heavily nested) expressions, each of which simply evaluates to a value.

The three basic types of expressions in Scheme are atomics/primitive expressions, call expressions, and special forms.

#### 1.1 Atomics

Atomics are the simplest expressions. Some atomics, such as numbers and booleans, are called "selfevaluating" because they evaluate to themselves:

- 123 → 123
- $-3.14 \rightarrow -3.14$
- $\#t \rightarrow \#t$  ;booleans in scheme are #t and #f

Symbols (variables) are also atomic expressions; they evaluate to the values to which they are bound. For example, the symbol + evaluates to the addition **procedure** (function) # [+].

Let's practice identifying atomic expressions in Scheme.

```
True or False: 3.14 is an atomic expression.
True. Floats are atomic expressions.
    True or False: pi is an atomic expression.
False. Pi in most cases isn't naturally defined within the programming
language -- rather, it's a variable name. As we haven't defined pi in this
case, Scheme doesn't know what to interpret pi as, so we get an Error if we
try to evaluate this on its own.
    True or False: - is an atomic expression.
True. - is an operator for subtraction.
    True or False: // is an atomic expression.
False. Scheme can only interpret the basic symbols of +, -, *, /. We'll have
to define our own floor division function!
    True or False: 'b' is an atomic expression.
True. b is a character, which is an atomic data type!
    True or False: "is this atomic?" is an atomic expression.
True. Even though a string is a list of characters, Scheme still treats them
as an atomic. It's important to note that in Scheme's case, strings can only
be denoted with ".
```

#### 1.2 Call Expressions

A call expression is denoted with parentheses and is formed like so:

```
(<operator> <operand0> <operand1> ... <operand>)
```

Each "element" of a call expression is an expression itself and is separated from its neighbors by whitespace. All call expressions are evaluated in the same way:

- 1. Evaluate the operator, which will return a procedure.
- 2. Evaluate the operands.
- 3. Apply operator on operands.

For example, to evaluate the call expression (+ (+ 1 2) 2), we first evaluate the operator +, which returns the procedure #[+]. Then we evaluate the first operand (+ 1 2), which returns 3. Then, we evaluate the last operand 2, receiving 2. Finally, we apply the #[+] procedure to 3 and 2, which returns 5. So this call expression evaluates to 5.

Note that in order to add two numbers, we had to call a function. In Python, + is a binary operator that can add two numbers without calling a function. Scheme has no such constructs, so even the most basic arithmetic requires you to call a function. The other math operators, including - (both subtraction and negation), \*, /, expt (exponentiation), =, <, >, <=, and >= function in the same way.

Let's do some practice with call expressions. What would Scheme do for each expression?

```
scm> (+ 1 (* 3 4))

13
    scm> (* 1 (+ 3 4))

7
    scm> (/ 2)

0.5 || When Scheme sees only one operand in the case of multiplication and division, it fills in <operand0> to be 1.
    scm> (- 2)

-2 || When Scheme sees only one operand in the case of addition and subtraction, it fills in <operand0> to be 0.
    scm> (* 4 (+ 3 (- 2 (/ 1))))
16
```

#### 1.3 Special Forms

Special forms *look* just like call expressions but are distinct in two ways:

- 1. One of the following keywords appears in the operator slot: define, if, cond, and, or, let, begin, lambda, quote, quasiquote, unquote, mu, define-macro, expect, unquote-splicing, delay, cons-stream, set!
- 2. They do not follow the evaluation rules for call expressions.

Below, we will go through a few commonly seen special forms.

#### 1.3.1 if expression

```
(if cate> <true-expr> <false-expr>)
```

An **if** expression is similar to a Python **if** statement. First, evaluate predicate>. The general structure follows (if predicate> <false-expr>).

- If redicate> is true, evaluate and return <true-expr>.
- If redicate> is false, evaluate and return <false-expr>.

Note that everything in Scheme is truthy (including 0) except for #f.

Also note that in Python, **if** is a statement, whereas in Scheme, **if** is an expression that evaluates to a value like any other expression would. In Scheme, you can then write something like this:

```
scm> (+ 1 (if #t 9 99))
10
```

Other special forms are also expressions that evaluate to values. Therefore, when we say "returns x," we mean "the special form evaluates to x."

#### 1.3.2 cond expression

```
(cond
    (<predicate1> <expr1>)
    ...
    (<predicateN> <exprN>)
    (else <else-expr>))
```

A **cond** expression is similar to a Python **if**-elif-**else** statement. It is an alternative to using many nested **if** expressions.

- Evaluate cate1>. If it is true, evaluate and return <expr1>.
- Otherwise, continue down the list by evaluating cate2>. If it is true, evaluate and return <expr2>.
- Continue in this fashion down the list until you hit a true predicate.
- If every predicate is false, return <else-expr>.

Let's practice using ifs and conds to evaluate Scheme problems!

### 1.3.3 and expression

```
(and <expr1> ... <exprN>)
```

and in Scheme works similarly to and in Python. Evaluate the expressions in order and return the value of the first false expression. If all of the values are true, return the last value. If no operands are provided, return #t.

#### 1.3.4 or expression

```
(or <expr1> ... <exprN>)
```

**or** in Scheme works similarly to **or** in Python. Evaluate the expressions in order and return the value of the first true expression. If all of the values are false, return the last value. If no operands are provided, return #f.

Let's practice and/or statements!

```
scm> (and #t (= 3 3) (> (- 61 42) (+ 61 42)))
#f
scm> (or #f (< 3 3) (< (- 61 42) (+ 61 42)))
#t</pre>
```

## 1.3.5 define expression

**define** does two things. It can define variables, similar to the Python = assignment operator:

```
(define <symbol> <expr>)
```

This will evaluate <expr> and bind the resulting value to <symbol> in the current frame.

**define** is also used to define procedures.

This code will create a new procedure that takes in the formal parameters  $<op1> \ldots <opN>$  and bind it to <symbol> in the current frame. When that procedure is called, the <body>, which may have multiple expressions, will be executed with the provided arguments bound to  $<op1> \ldots <opN>$ . The value of the final expression of <body> will be returned.

With either version of **define**, <symbol> is returned.

Dealing with the different types of **define** can be tricky. Scheme differentiates between the two by whether the first operand is a symbol or a list:

```
(define (x) 1); like x = lambda: 1 (define x 1); like x = 1
```

# 1.3.6 lambda expressions

Returns a new procedure that takes in the formal parameters <op1> ... <opN>. When that procedure is called, the <body>, which may have multiple expressions, will be executed with the provided arguments bound to <pp>< ... <pre><opN>. The value of the final expression of <body> will be returned.

#### 1.3.7 begin special form

Evaluates <expr1>, <expr2>, ..., <exprN> in order in the current environment. Returns the value of <exprN>.

### 1.3.8 let special form

Evaluates <exprl>, ..., <exprN> in the current environment. Then, creates a new frame as a child of the current frame and binds the values of <exprl>, ..., <exprN> to <symboll>, ..., <symbolN>, respectively, in that new frame. Finally, Scheme evaluates the <body>, which may have multiple expressions, in the new frame. The value of the final expression of <body> is be returned.

### 1.3.9 quote special form

```
(quote <expr>)
'<expr> ; shorthand syntax
```

Returns an expression that evaluates to <expr> in its unevaluated form. In other words, if you put '<expr> into the Scheme interpreter, you should get <expr> out exactly.

## 1.3.10 Summary of special forms

We have presented the main details of the most important special forms here, but this account is not comprehensive. Please see https://cs6la.org/articles/scheme-spec/ for a fuller explanation of the Scheme language.

behavior	syntax
if/else	(if <predicate> <true-expr> <false-expr>)</false-expr></true-expr></predicate>
if/elif/else	(cond ( <predicate1> <expr1>)</expr1></predicate1>
	 ( <predicaten> <exprn>)</exprn></predicaten>
	( <b>else</b> <else-expr>))</else-expr>
and	(and <expr1> <exprn>)</exprn></expr1>
or	(or <expr1> <exprn>)</exprn></expr1>
variable assignment	(define <symbol> <expr>)</expr></symbol>
function definition	<pre>(define (<symbol> <op1> <opn>)</opn></op1></symbol></pre>
lambdas	( <b>lambda</b> ( <op1> <opn>)</opn></op1>
evaluate many lines	(begin
	<expr1></expr1>
	•••
	<exprn>)</exprn>
temporary environment	(let (( <symbol1> <expr1>)</expr1></symbol1>
	<pre> (<symboln> <exprn>)) <body>)</body></exprn></symboln></pre>
quote	(quote <expr>) or '<expr></expr></expr>

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### 1. What will Scheme output?.

```
(a) (if 0 (/ 1 0) 1)
```

```
Error: Zero Division
```

Recall that 0 is a Truth-y value in Scheme. Thus (/ 1 0) evaluates to a Zero Division Error

```
(b) (and 1 #f (/ 1 0))
```

#f

Short-circuiting rules apply. This means that and returns the first False-y value or the last Truth-y value. In this case, the first False-y value is #f.

```
(c) (and 1 2 3)
```

3

Short-circuiting rules apply. This means that and returns the first False-y value or the last Truth-y value. In this case, the last Truth-y value is 3.

```
(d) (or #f #f 0 #f (/ 1 0))
```

0

Short-circuiting rules apply. This means that or returns the first Truth-y value or the last False-y value. In this case, the first Truth-y value is 0.

```
(e) (or #f #f (/ 1 0) 3 4)
```

```
Error: Zero Division
```

Short-circuiting rules apply. This means that or returns the first Truth-y value or the last False-y value. In this case, (/10) evaluates in a Zero Division Error.

```
(f) (and (and) (or))
```

#f

The special form or without any arguments evaluates to #f. The special form and without any arguments evaluates to #t. Also, short-circuiting rules apply. This means that and returns the first False-y value or the last Truth-y value. In this case, the first False-y value is #f.

Let's practice using **if**s and **cond**s to evaluate Scheme problems!

```
scm> (if 2 3 4)
  3
      scm> (if 0 3 4)
  3
      scm> (- 5 (if #f 3 4))
  1
      scm> (cond ((< -5 -7) 3)
            (else 4))
2. What will Scheme output?
  scm> (define c 2)
  scm> (eval 'c)
  scm> '(cons 1 nil)
  (cons 1 nil)
  scm> (eval '(cons 1 nil))
  (1)
  scm> (eval (list 'if '(even? c) 1 2))
  1
```

3. Define a procedure called hailstone, which takes in two numbers seed and n and returns the nth number in the hailstone sequence starting at seed. Assume the hailstone sequence starting at seed has a length of at least n. As a reminder, to get the next number in the sequence, divide by 2 if the current number is even. Otherwise, multiply by 3 and add 1.

## Useful procedures

• quotient: floor divides, much like // in python (quotient 103 10) outputs 10

(remainder 103 10) outputs 3

• remainder: takes two numbers and computes the remainder of dividing the first number by the second

```
; The hailstone sequence starting at seed = 10 would be
; 10 => 5 => 16 => 8 => 4 => 2 => 1

; Doctests
> (hailstone 10 0)
10
> (hailstone 10 1)
5
> (hailstone 10 2)
16
> (hailstone 5 1)
16

(define (hailstone seed n)
```

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)

```
(define (hailstone seed n)
    (if (= n 0)
        seed
        (if (= 0 (remainder seed 2))
            (hailstone
            (quotient seed 2)
            (- n 1))
          (hailstone
          (+ 1 (* seed 3))
          (- n 1)))))
; Alternative solution with cond
(define (hailstone seed n)
    (cond
        ((= n 0) seed)
        ((= 0 (remainder seed 2))
          (hailstone
          (quotient seed 2)
         (- n 1)))
        (else
          (hailstone
          (+ 1 (* seed 3))
          (- n 1)))))
```