CSM 61C Control Spring 2019

Single-Cycle Datapath and

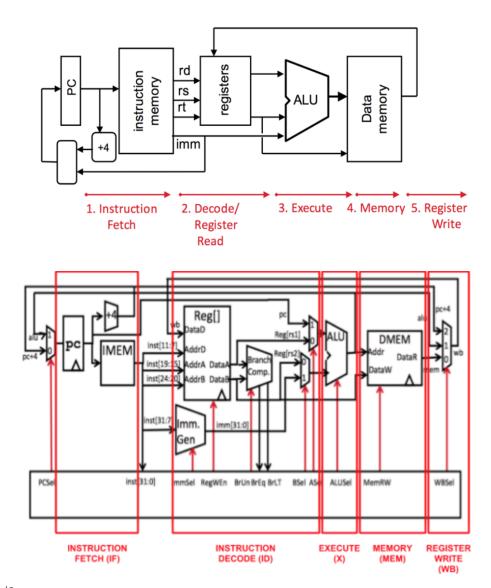
Mentoring 7: March 3, 2019

1 Single-Cycle Datapath and Control

1.1 5 Stages of a Single Cycle CPU:

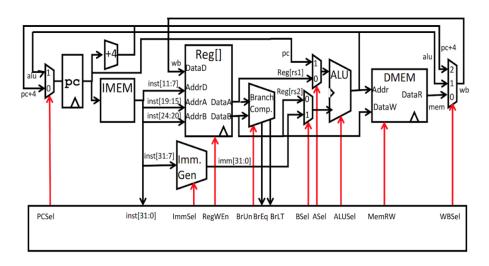
- Instruction Fetch (IF) Fetch from memory (IMEM)
- Instruction Decode (ID) Decode instruction
- Execute (EX) Execute operation (arithmetic, shifting, etc) using ALU
- Memory Access (MEM) Load and store instructions access memory
- Write Back to Register (WB) Write instruction back to RegFile

Datapaths: A Visual Approach



1.2 Control Logic

A controller send signals to our circuit, telling which pieces to perform what operations. Not all control signals matter for every instruction: for example, R-type instructions ignores the output from the immediate generator. Control signals are used to pick between mux inputs in order to perform the correct operation. They are embedded within the actual machine code for an instruction.



Control Inputs

Signal:	inst[31:0]	BrEq	BrLT		
Purpose:	Sends the current instruction to control	(DataA == DataB) ? 1 : 0	(DataA; DataB)?1:0		

Control Outputs

Signal:	PCSel	ALUSel	RegWEn	ImmSel
Purpose:	Next instruction location.	What operation to perform.	Do we change a registers value?	Format the
Signal:	MemRW	WBSel	BrUn	ASel/BSe
Purpose:	Read or write to mem.	What value to write back.	Branch signed or unsigned	Pick between

2 Game of Signals

2.1 Fill out the control signals for the following instructions (put an X if the signal does not matter). For ImmSel, write the corresponding instruction type.

Instr	BrEq	PCSel	ImmSel	BrUn	ASel	BSel	ALUSel	MemRW	RegWEn	WBSel
add	X	0	X	X	0	0	Add	0	1	1
lw										
bge										
sw										
auipc										
jal										

- 2.2 We want to expand our instruction set from the base RISC-V ISA (RV32I) to support some new instructions. You can find the canonical single-cycle datapath above. For the proposed instruction below, choose ONE of the options below.
 - 1. Can be implemented without changing datapath wiring, only changes in control signals are needed. (i.e. change existing control signals to recognize the new instruction)
 - 2. Can be implemented, but needs changes in datapath wiring, only additional wiring, logical gates and muxes are needed.
 - Can be implemented, but needs change in datapath wiring, and additional arithmetic units are needed (e.g. comparators, adders, shifters etc.). item Cannot be implemented.

(Note that the options from 1 to 3 gradually add complexity; thus, selecting 2 implies that 1 is not sufficient. You should select the option that changes the datapath the least (e.g. do not select 3 if 2 is sufficient). You can assume that necessary changes in the control signals will be made if the datapath wiring is changed.)

- (a) Allowing software to deal with 2s complement is very prone to error. Instead, we want to implement the negate instruction, neg rd, rs1, which puts -R[rs1] in R[rd].
- (b) Sometimes, it is necessary to allow a program to self-destruct. Implement segfault rs1, offset(rs2). This instruction compares the value in R[rs1] and the value in MEM[R[rs2]+offset]. If the two values are equal, write 0 into the PC; otherwise treat this instruction as a NOP.

3 Single Cycle CPU Timing Practice

The delays of a circuit elements are given as follows:

Clk-to-Q	RegFile Read	Mux	ALU	MEM Read
5ns	20ns	5ns	100ns	150ns

MEM Write Branch Comp		Imm Gen	RegFile Setup	
200ns	$50 \mathrm{ns}$	25ns	4ns	

- 3.1 Ignoring the length of a clock cycle, how long does it take to execute the instruction:
 - (a) lui t0, 0x1234
 - (b) jal ra, 0b1100
 - (c) beq x0, x5, 0b1100
- 3.2 What is the length of the critical path in the CPU? Which instruction exercises the critical path? Highlight its path on the diagram above starting from the PC.
- 3.3 What is the fastest possible clock for this datapath without any violations?

movz and movnz

Consider adding the following instruction to RISC-V (disregard any existing/similar definition on the green sheet):

Instruction	Operation
movz rd, rs1, rs2	if (R[rs1] == 0) R[rd]<-R[rs2]
movnz rd, rs1, rs2	if (R[rs1] != 0) R[rd]<-R[rs2]

Translate the following C code using movz and movnz. Do not use branches.

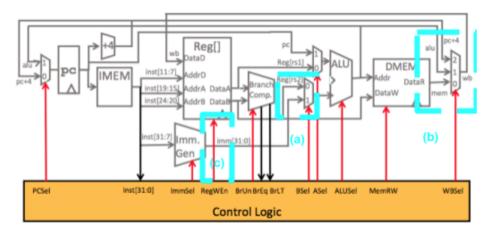
C Code

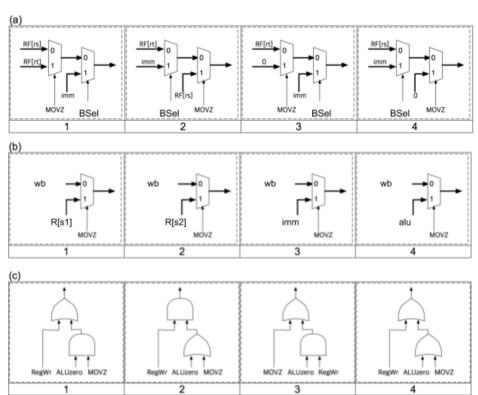
// a ->
$$s0$$
, b-> $s1$, c-> $s2$ int a = $b < c ? b : c;$

RISC-V Code

slt t0, s1, s2

4.2 Implement **movz** (but not movnz) in the datapath. Choose the correct implementation for (a), (b), and (c). Note that you do not need to use all the signals provided to each box, and the control signal MOVZ is 1 if and only if the instruction is **movz**. In the following diagrams, RF[rs] refers to R[rs1] and RF[rt] refers to R[rs2]. ALUZero = 1 if the A input of the ALU is zero.





4.3 Generate the control signals for movz. The values should be 0, 1, or X (dont care) terms. You must use dont care terms where possible.

MOVZ	PCSel	ImmSel	RegWEn	BrUn	BSel	ASel	ALUSel	MemRW	WBSel
1									