

## 1 Load and Store

- 1.1 For each line of RISC-V code, answer what will be the value saved into the registers involved or the effect on memory. Assume that x8 contains a valid address in memory, such that  $\text{Mem}[\text{R}[\text{x8}]] = 0\text{x}00000180$ .

(a) `lw x9, 0(x8)` What value does x9 now store?

`0x00000180`

(b) `lb x10, 1(x8)` What value does x10 store?

`0x00000001`

(c) `lb x11, 0(x8)` What value does x11 store?

`0xFFFFFFFF80`

(d) `sb x11, 3(x8)` What's the effect on  $\text{Mem}[\text{R}[\text{x8}]]$ ?

`Mem[R[x8]] = 0x80000180`

(e) `lbu x12, 0(x8)` What value does x12 store?

`0x00000080`

## 2 More RISC-V

- 2.1 You wish to speed up one of your programs by implementing it directly in assembly. Your partner started translating the function `is_substr()` from C to RISC-V, but didn't finish. Please complete the translation by filling in the lines below with RISC-V assembly. The prologue and epilogue have been written correctly but are not shown.

Note: `strlen()`, both as a C function and RISC-V procedure, takes in one string as an argument and returns the length of the string (not including the null terminator).

```
/* Returns 1 if s2 is a substring of s1, and 0 otherwise. */
int is_substr(char* s1, char* s2) {
    int len1 = strlen(s1);
    int len2 = strlen(s2);
    int offset = len1 - len2;
    while (offset >= 0) {
        int i = 0;
        while (s1[i + offset] == s2[i]) {
            i += 1;
            if (s2[i] == '\0')
                return 1;
        }
        offset -= 1;
    }
    return 0;
}
```

2.2 Fill in the following RISC-V code based on the given C code:

```

1. is_substr:
2.     mv s1, a0
3.     mv s2, a1
4.     jal ra, strlen
5.     mv s3, a0
6.     mv a0, s2
7.     jal ra, strlen
8.     sub s3, s3, a0
9. Outer_Loop:
10.    _____, _____, _____, False
11.    add t0, x0, x0
12. Inner_Loop:
13.    add t1, t0, s3
14.    add t1, s1, t1
15.    lbu t1, 0(t1)
16.    _____
17.    _____
18.    _____, t1, _____, Update_Offset
19.    addi t0, t0, 1
20.    add t2, t0, s2
21.    _____
22.    beq t2, _____, _____,
23.    jal x0 Inner_Loop
24. Update_Offset:
25.    addi s3, s3, -1
26.    _____
27. False:
28.    xor a0, a0, _____
29.    jal x0, End
30. True:
31.    addi a0, x0, 1
32. End: _____

```

```

1. is_substr:
2.     mv s1, a0
3.     mv s2, a1
4.     jal ra, strlen
5.     mv s3, a0
6.     mv a0, s2
7.     jal ra, strlen

```

```
8.      sub s3, s3, a0
9. Outer_Loop:
10.     blt s3, x0, False
11.     add t0, x0, x0
12. Inner_Loop:
13.     add t1, t0, s3
14.     add t1, s1, t1
15.     lbu t1, 0(t1)
16.     add t2 s2 t0
17.     lbu t2 0(t2)
18.     bne t1, t2, Update_Offset
19.     addi t0, t0, 1
20.     add t2, t0, s2
21.     lbu t2 0(t2)
22.     beq t2, x0, True
23.     jal x0 Inner_Loop
24. Update_Offset:
25.     addi s3, s3, -1
26.     jal x0 Outer_Loop
27. False:
28.     xor a0, a0, a0
29.     jal x0, End
30. True:
31.     addi a0, x0, 1
32. End: ret
```

### 3 A Fib-ulous Question

- 3.1 Fill in the following RISC-V function that recursively calculates the n-th fibonacci number.

```

1. fib:
2.     addi t0 x0 2
3.     blt __ t0 exit
4.     addi __ __ -8 #prologue
5.     sw __ 4(sp)
6.     sw s0 0(sp)
7.     addi __ __ -1
8.     mv s0 a0
9.     jal ra fib
10.    mv t0 a0
11.    mv a0 __
12.    mv s0 t0
13.    addi a0 a0 __
14.    -----
15.    add a0 a0 s0
16.    lw __ 4(sp)
17.    lw s0 0(sp)
18.    addi __ __ 8
19. exit:
20.    jr ra

```

```

1. fib:
2.     addi t0 x0 2
3.     blt a0 t0 exit
4.     addi sp sp -8
5.     sw ra 4(sp)
6.     sw s0 0(sp)
7.     addi a0 a0 -1
8.     mv s0 a0
9.     jal ra fib
10.    mv t0 a0
11.    mv a0 s0
12.    mv s0 t0
13.    addi a0 a0 -1
14.    jal ra fib
15.    add a0 a0 s0
16.    lw ra 4(sp)
17.    lw s0 0(sp)

```

## 6 *RISC-V*

```
18.    addi sp sp 8
19. exit:
20.    jr ra
```

## 4 Linked List Reversals in RISC-V

4.1 Assume we have the following linked list node struct:

```
struct node{
    int val;
    struct node * next;
};
```

Also, recall the function to reverse a linked list iteratively, given a pointer to the head of the linked list.

```
void reverse(struct node * head){
    struct node * prev = NULL;
    struct node * next;
    struct node * curr = head;
    while(curr != NULL){
        next = curr->next;
        curr->next = prev;
        prev = curr;
        curr = next;
    }
}
```

- 4.2 Now assume `a0` contains the address of the head of a linked list. Fill in the function below to reverse a linked list. Assume ‘reverse’ follows calling conventions. ‘reverse’ doesn’t return anything. You may not need all lines.

```

1. reverse: _____
2. _____
3. _____
4. _____
5. add s0 a0 x0 #s0 corresponds to curr
6. xor s2 s2 s2 #s2 corresponds to the pointer ‘prev’
7. loop: ___ s0 x0 exit
8. _____
9. _____
10. add s2 s0 x0
11. add s0 s1 x0
12. j loop
13. exit: _____
14. _____
15. _____
16. addi sp sp 12
17. jr ra

```

```

1. reverse: addi sp sp -12
2. sw s0 0(sp)
3. sw s1 4(sp)
4. sw s2 8(sp)

```

```

7. beq s0 x0 exit
8. lw s1 4(s0)
9. sw s2 4(s0)

```

```

13. exit:lw s0 0(sp)
14. lw s1 4(sp)
15. lw s2 8(sp)

```

Lines 1-4 and 13-15 is just following calling conventions for RISC-V, since `s0-s11` by convention, must be preserved when someone calls ‘reverse’. We notice that `s0`, `s1`, and `s2` are all being modified.