

xv6©-rev10
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xv6 Interrupt Servicing

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Interrupt Service routines

- In xv6 the C-routine **trap()** is the basis for interrupt servicing.
- Interrupt service routines are necessarily hardware dependent, thus
 - ISRs cannot be written **totally** in C.
 - (as usual) We aim to write whatever is possible in C.
- In the following we look at what the hardware forces us to do.

Infrastructure of interrupt service routines in C

Using **trap0()** issues

- Assume **trap0** is a C routine to handle interrupt request 0.
- Issues we need to deal with:
 - Can IDT[0] point directly to **trap0**? Yes, but returning will crash.
 - What about the interrupted code registers? Probably will be destroyed.
 - What about our segment registers? Unexpected if was at user mode.

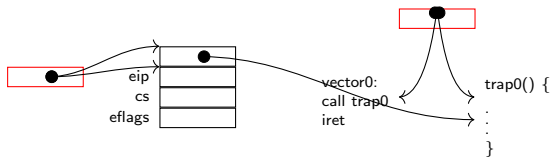
We solve these issues one by one.

IDT[0] pointing to a good place

- **trap0()** ends with **ret**, no way around it.
- So we wrap **trap0** inside a routine ending with **iret**:

```
vector0 :  
    call trap0  
    iret
```

- Then we let **IDT[0]** point to **vector0**.



Saving interrupted code GPRs

- Registers destroyed by the interrupt service routine should be saved.
- **eip**, **eflags**, and possibly **esp** where already saved.
- The registers destroyed by **trap0()** are not known, thus save them all:

```
vector0 :  
    pushal  
    call trap0  
    popal  
    iret
```

Our segment registers?

- **cs** was saved and then loaded from the IDT.
- **ss**:
 - User mode interrupted: **ss** was saved and then loaded from the TSS.
 - Kernel mode interrupted: **ss** already contains the correct value (SEG_KDATA).
- **ds, es, fs, gs**: Code needed:

```
movw $SEG_KDATA<<3,%ax
movw %ax,%ds
movw %ax,%es
movw $0,%ax
movw %ax,%fs
movw %ax,%gs
```


Infrastructure (1) for trap0()

```
vector0:  
    pushal  
    movw $SEG_KDATA<<3,%ax  
    movw %ax,%ds  
    movw %ax,%es  
    movw $0,%ax  
    movw %ax,%fs  
    movw %ax,%gs
```

```
    call trap0  
    popal  
    iret
```

- We might have just destroyed user mode registers!

Infrastructure (2) for trap0()

vector0:

pushl %ds

pushl %es

pushl %fs

pushl %gs

pushal

movw \$SEG_KDATA<<3,%ax

movw %ax,%ds

movw %ax,%es

movw \$0,%ax

movw %ax,%fs

movw %ax,%gs

call trap0

popal

popl %gs

popl %fs

popl %es

popl %ds

iret

Interrupted code registers content

- ISRs might be interested in accessing the interrupted code registers.
- (This is reasonable only for internal interrupt requests!)
- The register values were pushed onto the stack.
- A slight addition allows easy access to them from C.
- (elective) The values are accessible also without the change.

Infrastructure (3) for trap0()

```
vector0 :  
    pushl %ds  
    pushl %es  
    pushl %fs  
    pushl %gs  
    pushal  
    movw $SEG_KDATA<<3,%ax  
    movw %ax,%ds  
    movw %ax,%es  
    movw $0,%ax  
    movw %ax,%fs  
    movw %ax,%gs
```

```
    pushl %esp  
    call trap0  
    addl $4,%esp  
    popal  
    popl %gs  
    popl %fs  
    popl %es  
    popl %ds  
    iret
```

frame struct

```
struct frame {  
    long edi;  
    long esi;  
    long ebp;  
    long esp_irrelevant;  
    long ebx;  
    long edx;  
    long ecx;  
    long eax;  
  
    short gs;  
    short padding1;  
    short fs;  
    short padding2;  
    short es;  
    short padding3;
```

```
    short ds;  
    short padding4;  
  
    long eip;  
    short cs;  
    short padding5;  
    long eflags;  
  
    long esp;  
    short ss;  
    short padding6;  
}
```

Accessing interrupted code registers

- Possible declaration of `trap0()`:

```
void trap0(struct frame *f) {  
    f->eax = someValue(f);  
    :  
}
```

- Note: The **esp** and **ss** fields exist only if user mode was interrupted.

Infrastructure for **trap8()** (serving IRQ 8)

- IRQ8 leaves error code below **eip**.
- Hence:
 - The **frame** struct does not give the correct picture.
 - There is a superfluous word before the **iret** instruction.
- So:
 - A new structure replacing **frame** struct is needed.
 - Removing a word from the stack is needed.

Infrastructure (1) for `trap8()`

`vector8:`

```
pushl %ds
pushl %es
pushl %fs
pushl %gs
pushal
movw $SEG_KDATA<<3,%ax
movw %ax,%ds
movw %ax,%es
movw $0,%ax
movw %ax,%fs
movw %ax,%gs
```

```
pushl %esp
call trap8
addl $4,%esp
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $4,%esp
iret
```


eframe struct

```
struct eframe {  
    long edi;  
    long esi;  
    long ebp;  
    long esp_irrelevant;  
    long ebx;  
    long edx;  
    long ecx;  
    long eax;  
  
    short gs;  
    short padding1;  
    short fs;  
    short padding2;  
    short es;  
    short padding3;
```

```
    short ds;  
    short padding4;  
  
    long err;  
    long eip;  
    short cs;  
    short padding5;  
    long eflags;  
  
    long esp;  
    short ss;  
    short padding6;  
}
```

Accessing error code

- Possible declaration of trap8():

```
void trap8(struct eframe *f) {  
    if (f->err)  
        f->eax = someValue(f);  
    else  
        f->eax = someOtherValue(f);  
    :  
}
```

Uniform behaviour would be useful

- The above solution is feasible.
- However, if one uniform routine `trap()` is asked for, we tinker.
- How?

Infrastructure (4) for `trap0()`

`vector0:`

`pushl $0`

`pushl %ds`

`pushl %es`

`pushl %fs`

`pushl %gs`

`pushal`

`movw $SEG_KDATA<<3,%ax`

`movw %ax,%ds`

`movw %ax,%es`

`movw $0,%ax`

`movw %ax,%fs`

`movw %ax,%gs`

`pushl %esp`

`call trap0`

`addl $4,%esp`

`popal`

`popl %gs`

`popl %fs`

`popl %es`

`popl %ds`

`addl $4,%esp`

`iret`

The **trapframe** struct (1)

```
struct trapframe {  
    long edi;  
    long esi;  
    long ebp;  
    long esp_irrelevant;  
    long ebx;  
    long edx;  
    long ecx;  
    long eax;  
  
    short gs;  
    short padding1;  
    short fs;  
    short padding2;  
    short es;  
    short padding3;
```

```
    short ds;  
    short padding4;  
  
    long err;  
    long eip;  
    short cs;  
    short padding5;  
    long eflags;  
  
    long esp;  
    short ss;  
    short padding6;  
}
```

All isr's can use the **trapframe** struct

```
void trap0(struct trapframe *f) {  
    :  
}
```

```
void trap8(struct trapframe *f) {  
    :  
}
```

- Let us reorganize the entry code.

Infrastructure (5) for **trap0()**

vector0:

pushl \$0

jmp alltraps

alltraps:

pushl %ds

pushl %es

pushl %fs

pushl %gs

pushal

movw \$SEG_KDATA<<3,%ax

movw %ax,%ds

movw %ax,%es

movw \$0,%ax

movw %ax,%fs

movw %ax,%gs

pushl %esp

call trap0

addl \$4,%esp

popal

popl %gs

popl %fs

popl %es

popl %ds

addl \$4,%esp

iret

Infrastructure (5) for **trap8()**

```
vector8:  
    jmp alltraps
```

```
alltraps:  
    pushl %ds  
    pushl %es  
    pushl %fs  
    pushl %gs  
    pushal  
    movw $SEG_KDATA<<3,%ax  
    movw %ax,%ds  
    movw %ax,%es  
    movw $0,%ax  
    movw %ax,%fs  
    movw %ax,%gs
```

```
    pushl %esp  
    call trap8  
    addl $4,%esp  
    popal  
    popl %gs  
    popl %fs  
    popl %es  
    popl %ds  
    addl $4,%esp  
    iret
```


Infrastructure (1) for **trap()**, IRQ0

```
vector0:
    pushl    $0
    jmp     alltraps
vector8:
    jmp     alltraps
alltraps:
    pushl    %ds
    pushl    %es
    pushl    %fs
    pushl    %gs
    pushal
    movw     $SEG_KDATA<<3,%ax
```

What is lost??

```
movw    %ax,%ds
movw    %ax,%es
movw    $0,%ax
movw    %ax,%fs
movw    %ax,%gs
pushl    %esp
call    trap
addl    $4,%esp
popal
popl    %gs
popl    %fs
popl    %es
popl    %ds
addl    $4,%esp
iret
```

Infrastructure (2) for `trap()`

```
vector0:
    pushl    $0
    pushl    $0
    jmp     alltraps
vector8:
    pushl    $8
    jmp     alltraps
alltraps:
    pushl    %ds
    pushl    %es
    pushl    %fs
    pushl    %gs
    pushal
    movw     $SEG_KDATA<<3,%ax
```

```
    movw     %ax,%ds
    movw     %ax,%es
    movw     $0,%ax
    movw     %ax,%fs
    movw     %ax,%gs
    pushl    %esp
    call     trap
    addl     $4,%esp
    popal
    popl     %gs
    popl     %fs
    popl     %es
    popl     %ds
    addl     $8,%esp
    iret
```

The **trapframe** struct (2)

```
struct trapframe {  
    long edi;  
    long esi;  
    long ebp;  
    long esp_irrelevant;  
    long ebx;  
    long edx;  
    long ecx;  
    long eax;  
  
    short gs;  
    short padding1;  
    short fs;  
    short padding2;  
    short es;  
    short padding3;
```

```
    short ds;  
    short padding4;  
  
    long trapno;  
    long err;  
  
    long eip;  
    short cs;  
    short padding5;  
    long eflags;  
  
    long esp;  
    short ss;  
    short padding6;  
}
```

trap()

```
void trap(struct trapframe *tf) {  
    switch (tf->trapno) {  
        case 0:  
            trap0(tf);  
            break;  
  
        :  
  
        case 8:  
            trap8(tf);  
            break;  
  
        :  
    }  
}
```

Retro groking on process creation

- Note the **trapframe** structure and recall **userinit()**.
- Now we can understand why things are as they are:
 - There are always interrupts.
 - User mode code is always going to be interrupted.
 - So there must be a framework for:
 - switching from user mode to kernel mode.
 - Return safely from kernel mode to user mode.
 - Running new user code builds on this framework.
 - New user code has no history (It is new!).
 - So a fake history is build for it.
 - And then the kenel 'returns' to the code entry point.

- The xv6 code.

Kernel entry points

We combine the previous techniques to get:

- Kernel entry points **vector0** upto **vector255** follows on the next slide.
- Each entry will contains one or two **push** instructions followed by
`jmp alltraps`
- Use some scripting language to generate the .S file with the above code.
- (elective) (or by a smart.s and some fancy assembly directives).

all vectors

```
vector0:
    pushl $0
    pushl $0
    jmp alltraps
:
vector7:
    pushl $0
    pushl $7
    jmp alltraps
vector8:
    pushl $8
    jmp alltraps
vector9:
    pushl $0
    pushl $9
    jmp alltraps
```

```
vector10:
    pushl $10
    jmp alltraps
:
vector14:
    pushl $14
    jmp alltraps
vector15:
    pushl $0
    pushl $15
    jmp alltraps
vector16:
    pushl $0
    pushl $16
    jmp alltraps
```

```
vector17:
    pushl $17
    jmp alltraps
vector18:
    pushl $0
    pushl $18
    jmp alltraps
:
vector255:
    pushl $0
    pushl $255
    jmp alltraps
```

Tedious!

alltraps

3304

```
alltraps::  
    pushl %ds  
    pushl %es  
    pushl %fs  
    pushl %gs  
    pushal  
    movw $SEG_KDATA<<3,%ax  
    movw %ax,%ds  
    movw %ax,%es  
    movw $0,%ax  
    movw %ax,%fs  
    movw %ax,%gs
```

```
    pushl %esp  
    call trap  
    addl $4,%esp  
trapret:  
    popal  
    popl %gs  
    popl %fs  
    popl %es  
    popl %ds  
    addl $8,%esp  
    iret
```

vectors[]

- How do we loop on **vector0** through **vector255** to initialize IDT?
 - We cannot.
- So, we build a vector of addresses:

```
vectors ::  
    .long vector0  
    .long vector1  
    :  
    .long vector255
```

- Writing this is also TEDIOUS.

vectors.pl

2960

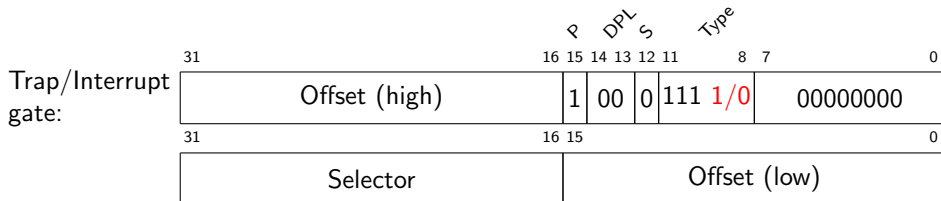
```
for (my $i = 0; $i < 256; $i++) {  
    print "vector$i:\n";  
    if (!( $i == 8 ||  
          ( $i >= 10 && $i <= 14) || $i == 17)) {  
        print "_pushl_\n";  
    }  
    print "_pushl_\n";  
    print "_jmp_alltraps\n";  
}  
print ".data\n";  
print "vectors::\n";  
for (my $i = 0; $i < 256; $i++) {  
    print "_long_vector$i\n";  
}
```

(elective) a direct vectors.S

```
.globl alltraps , vectors
.data
vectors:
    v = 0
.rept 256
    .text
0:
    .if !(v == 8 || (10 <= v && v <= 14) || v == 17)
        pushl $0
    .endif
    pushl $v
    jmp alltraps
    .data
    .long 0b
    v = v + 1
.endr
end
```

IDT initializing

Gate descriptor



```

901 struct gatedesc {
    uint off_15_0 : 16; // low 16 bits of offset in segment
    uint cs : 16; // code segment selector
    uint args : 5; // # args, 0 for interrupt/trap gates
    uint rsv1 : 3; // reserved (should be zero I guess)
    uint type : 4; // type (STS_{TG,IG32,TG32})
    uint s : 1; // must be 0 (system)
    uint dpl : 2; // descriptor (meaning new) privilege level
    uint p : 1; // Present
    uint off_31_16 : 16; // high bits of offset in segment
};
    
```

SETGATE macro

```
921 #define SETGATE(gate, istrap, sel, off, d) \ { \
    (gate).off_15_0 = (uint)(off) & 0xffff; \
    (gate).cs = (sel); \
    (gate).args = 0; \
    (gate).rsv1 = 0; \
    ((gate).type = (istrap) ? STS_TG32 : STS_IG32; \
    (gate).s = 0; \
    (gate).dpl = (d); \
    (gate).p = 1; \
    (gate).off_31_16 = (uint)(off) >> 16; \
}
```

IDT construction

```
#define T_SYSCALL 64
```

```
3361 struct gatedesc idt[256];  
extern uint vectors[]; // in vectors.S: array of 256  
  
3367 tvinit(void) {  
    for (int i = 0; i < 256; i++)  
        SETGATE(idt[i], 0, SEG_KCODE<<3, vectors[i], 0);  
  
    SETGATE(idt[T_SYSCALL], 1, SEG_KCODE<<3,  
            vectors[T_SYSCALL], DPL_USER);  
}  
  
3378 void idtinit(void) {  
    lidt(idt, sizeof(idt));  
}
```


Finally, we can investigate **trap()**.

trap()

The code of **trap()** is composed of three parts:

1. Handle **killed** and dispatch to **syscall()** if invoked by **int \$64**.
2. Dispatch to code handling other IRQ's.
3. Code handling **killed** and **yield()**.

Note about **killed**:

- **killed** set means the process should terminate.
- The **killed** field should be checked at the following locations:
 - Entering the kernel.
 - Leaving the kernel.
 - Whenever the process might enter **SLEEPING** for a long time.

trap() part 1

```
3401 void trap(struct trapframe *tf) {  
    if (tf->trapno == T_SYSCALL) {  
        if (myproc()->killed)  
            exit();  
        myproc()->tf = tf;  
        syscall();  
        if (myproc()->killed)  
            exit();  
        return;  
    }  
}
```

- **syscall()** might be long (i.e., invokes **swtch**), hence the double check on **myproc()->killed**.

trap() part 2, controller interrupts

3413

```
switch (tf->trapno) {
case T_IRQ0+IRQ_TIMER:
    if (cpuid() == 0) {
        acquire(&tickslock);
        ticks++;
        wakeup(&ticks);
        release(&tickslock);
    }
    lapiceoi();
    break;
case T_IRQ0+IRQ_IDE:
    ideintr();
    lapiceoi();
    break;
case T_IRQ0+IRQ_IDE+1:
    break;
```

```
case T_IRQ0+IRQ_KBD:
    kbdintr();
    lapiceoi();
    break;
case T_IRQ0+IRQ_COM1:
    uartintr();
    lapiceoi();
    break;
case T_IRQ0+7:
case T_IRQ0+IRQ_SPURIOUS:
    cprintf("cpu%d: _spurious_\n",
            cpuid(), tf->cs, tf->eip);
    lapiceoi();
    break;
```

trap() part 2, unexpected interrupt

3450

default :

```
if (myproc() == 0 || (tf->cs&3) == 0) {
    cprintf("unexpected trap %d from cpu %d \
            eip %x (cr2=0x%x)\n",
            tf->trapno, mycpu()->id, tf->eip, rcr2());
    panic("trap");
}
cprintf("pid %d %s: trap %d err %d on cpu %d \
        eip 0x%x addr 0x%x      kill proc\n",
        myproc()->pid, myproc()->name, tf->trapno, tf->err,
        cpuid(), tf->eip,
        rcr2());
myproc()->killed = 1;
}
```

trap() part 3

3468

```
if (myproc() && myproc()->killed &&
    (tf->cs&3) == DPL_USER)
    exit();

if (myproc() && myproc()->state == RUNNING &&
    tf->trapno == T_IRQ0+IRQ_TIMER)
    yield();

if (myproc() && myproc()->killed &&
    (tf->cs&3) == DPL_USER)
    exit();
}
```