xv6(c)-rev10 (Copyright Frans Kaashoek, Robert Morris, and Russ Cox.) First Process Creation, II

Carmi Merimovich

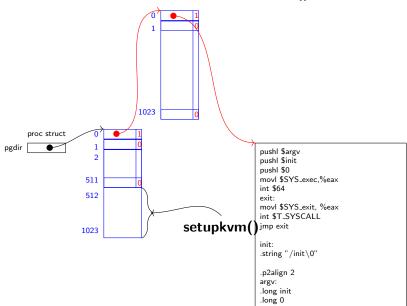
Tel-Aviv Academic College

November 20, 2017

Context

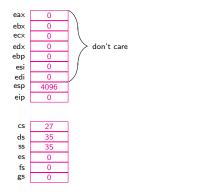
```
kinit1 (end, P2V(4*1024*1024)); // phys page alloca
kvmalloc(); // kernel page table
seginit(); // set up segments
pinit();
           // process table
kinit2(P2V(4*1024*1024), P2V(PHYSTOP)); // must co
userinit(); // first user process
mpmain();
```

Context in **userinit()**



The Road to User Mode

Initial user mode state for first process



```
4095
```

The naive approach for the scheduler

```
movw $35,%ax
movw %ax . ds
movw %ax,ss # Interrupts use the stack. We just ruined
movl $4096,%esp
movl $0,%eax # Much better to move eax to other
movl $0.%ecx
movl $0.%edx
movl $0.%ebx
movl $0,%ebp
movl $0.% esi
movl $0.%edi
movw %ax, %es
movw %ax,%fs
movw %ax, %gs
```

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The problem

It is not possible to find a reasonable order to load:

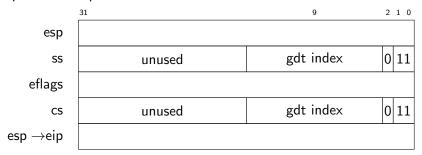
- cs, eip.
- ss, esp.
- eflags (Containing the Interrupt flag).

What the hardware should provide us with?

• The **iret** instruction fits the bill.

iret instruction in kernel mode

Implied stack operands:



The operands are loaded into the appropriate registers and POPped of the stack.

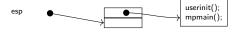
Working Approach for the scheduler!!

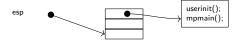
```
pushl
       $4096
pushl
      $35
pushl
      $512
      $27
pushl
pushl
       $0
```

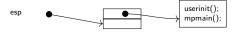
iret

We are not the scheduler!







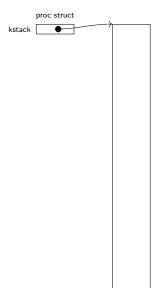


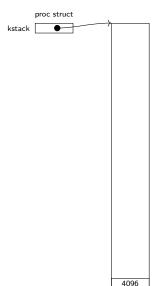


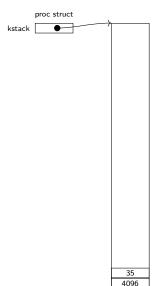


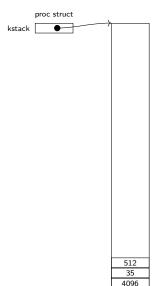
What do we do??

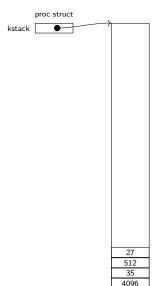
proc struct kstack

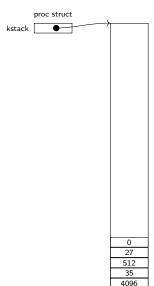


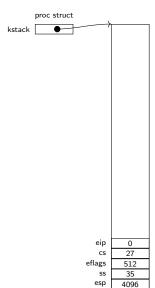


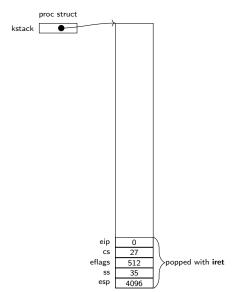








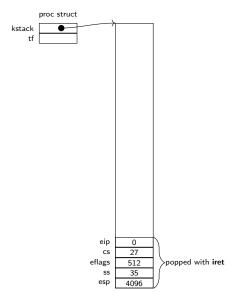


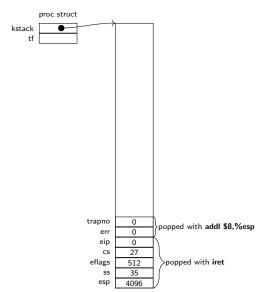


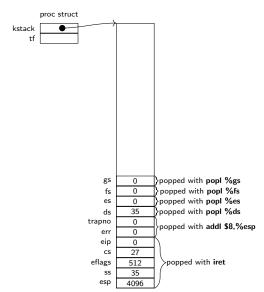
- We expect the scheduler to execute iret with its esp pointing to the eip field.
- We began the discussion with more registers.
- What about those other registers?
- The scheduler should load them with the prescribed values.
 - Ugly. Registers are being setup in two different places.

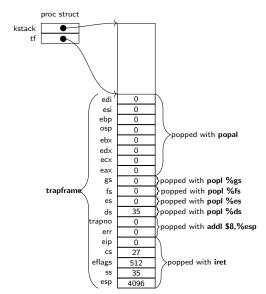
OR

We will put reigsters values on stack for all the registers!









trapframe structure

```
struct trapframe
 uint edi:
 uint esi;
 uint ebp;
 uint oesp;
 uint ebx;
 uint edx:
 uint ecx;
 uint eax;
 ushort gs;
 ushort padding1;
 ushort fs:
 ushort padding2;
```

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```
ushort es:
ushort padding3;
ushort ds:
ushort padding4;
uint trapno;
uint err;
uint eip;
ushort cs:
ushort padding5;
uint eflags;
uint esp;
ushort ss:
ushort padding6;
```

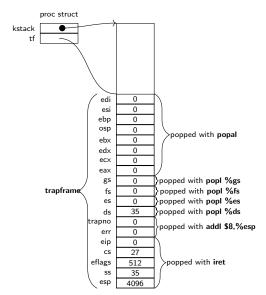
Building trapframe example

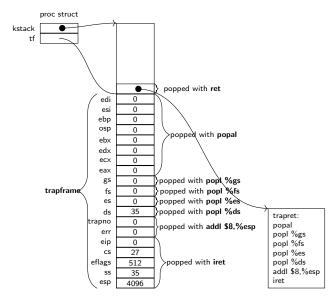
```
sp = p \rightarrow kstack + KSTACKSIZE;
sp = -sizeof *p -> tf;
p->tf = (struct trapframe *)sp;
memset(p\rightarrow tf, 0, sizeof *p\rightarrow tf);
p\rightarrow tf\rightarrow cs = (SEG\_UCODE << 3) \mid DPL\_USER:
p\rightarrow tf\rightarrow ds = (SEG\_UDATA << 3) \mid DPL\_USER;
p\rightarrow tf\rightarrow es = p\rightarrow tf\rightarrow ds:
p\rightarrow tf\rightarrow ss = p\rightarrow tf\rightarrow ds:
p\rightarrow tf\rightarrow eflags = FL_IF;
p\rightarrow tf\rightarrow esp = PGSIZE:
p\rightarrow tf\rightarrow eip=0; // beginning of initcode. S
```

• So, the code the scheduler is supposed to execute is:

```
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $8,%esp
iret
```

 For the sake of generality, we leave the address of this code on the stack.





trapret pointer setting

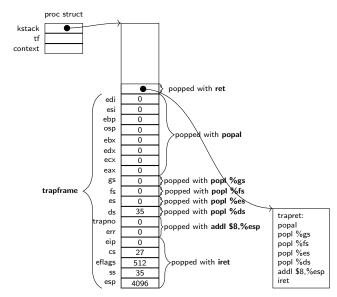
```
sp = 4;
*sp = trapret;
```

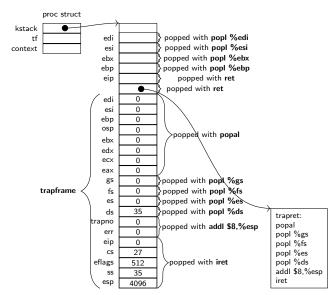
A word of caution

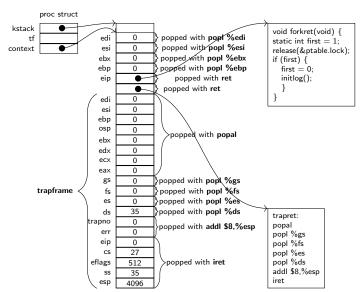
- The real reason behind the trapframe construction is the interrupt servicing code.
- Wait for the real reasoning until there.

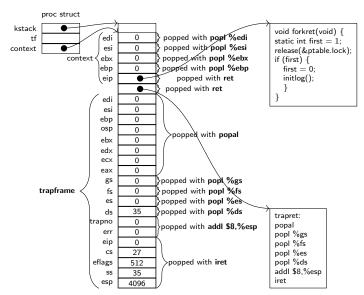
The **context** struct

- The **xv6** machinery is more general than the above.
- The scheduler has a context of its own.
- Using just **trapframe** will cause the scheduler to loose context.
- So, a context struct is pushed on the stack in order to preserve context.









context structure

```
2623 struct context {
    uint edi;
    uint esi;
    uint ebx;
    uint ebp;
    uint eip;
```

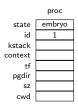
Bulding the **context** structure

```
sp -= sizeof *p->context;
p->context = sp;
memset(p->context, 0, sizeof(*p->context));
p->context->eip = forkret;
```



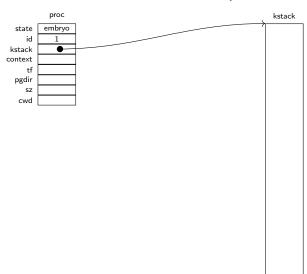
```
void forkret(void) {
static int first = 1:
release(&ptable.lock);
if (first) {
  first = 0:
  initlog();
```

```
trapret:
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $8,%esp
iret
```



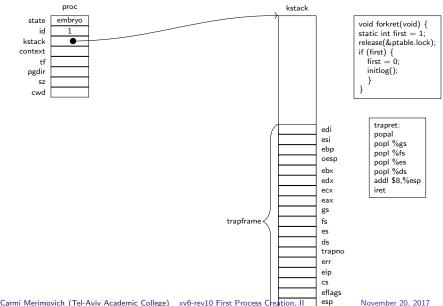
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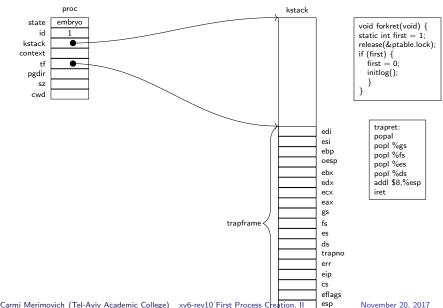
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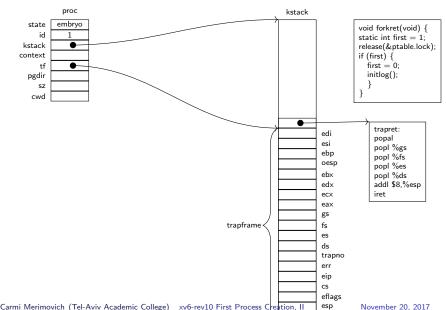


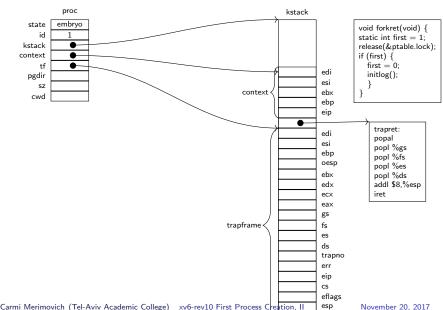
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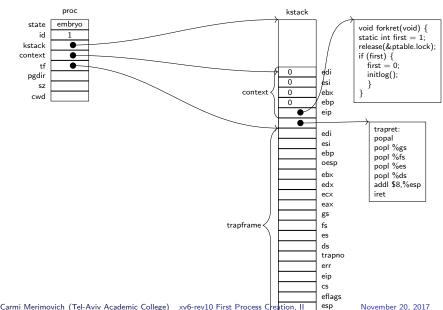
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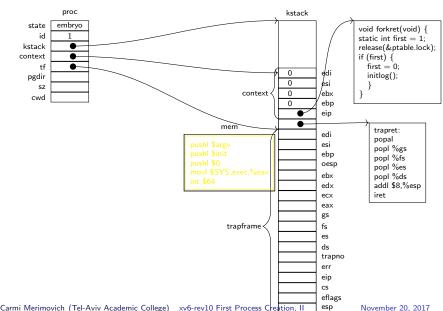


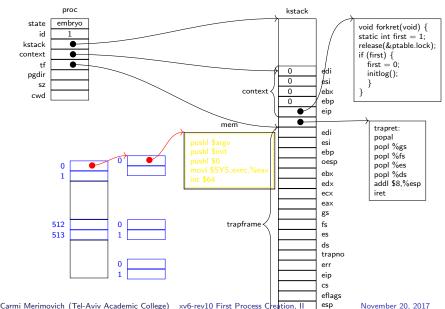


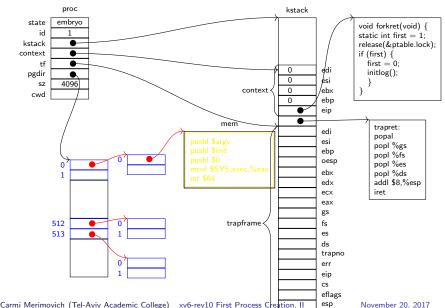


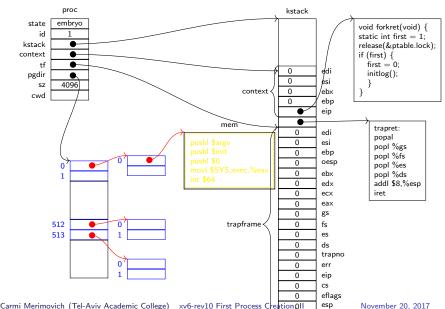


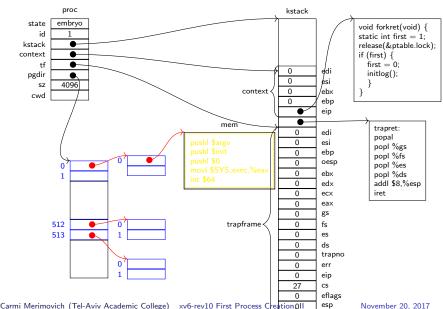


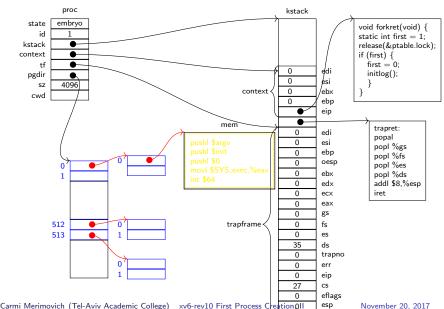


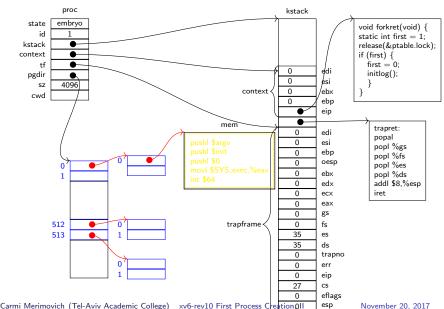


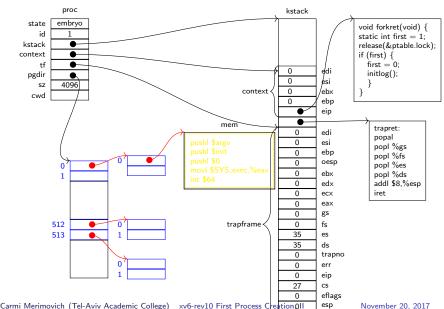


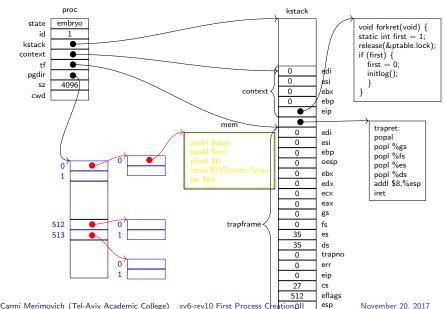


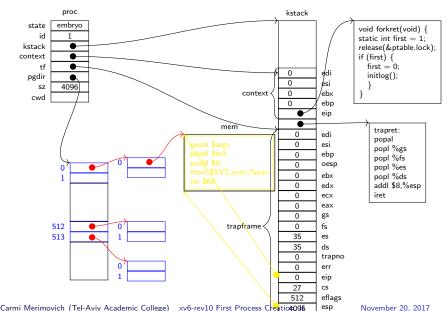


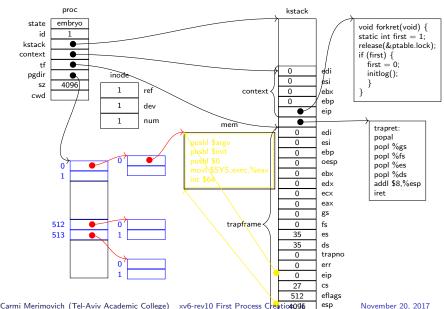


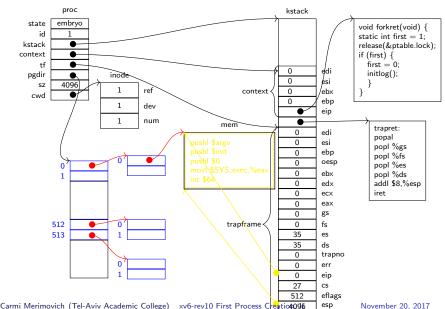


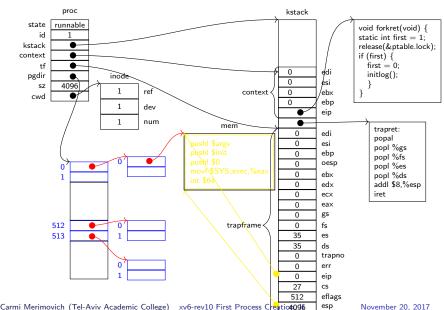












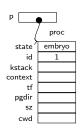
userinit()

- allocproc(): Allocate proc structure.
 - Allocate process kernel stack.
 - Leave space for a **trapframe** structure.
 - Build context structure on the stack.
- setupkvm: Create page table.
- inituvm: User space:
 - Allocate one page.
 - Copy user code to the allocated page.
 - Modify page table to use the allocated page.
- Set the trapframe so the user process will be able to run.

allocproc(): (1) Finding unused proc structure

```
<sub>2473</sub> static struct proc *allocproc(void) {
    struct proc *p;
    char *sp;
    acquire(&ptable.lock);
    for(p = ptable.proc; p < &ptable.proc[NPROC]; p++)
      if (p->state == UNUSED)
      goto found;
     release(&ptable.lock);
    return 0:
   found:
    p—>state = EMBRYO:
    p->pid = nextpid++:
```

release(&ptable.lock);

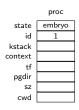


```
void forkret(void) {
static int first = 1:
release(&ptable.lock);
if (first) {
  first = 0;
  initlog();
```

```
trapret:
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $8,%esp
iret
```

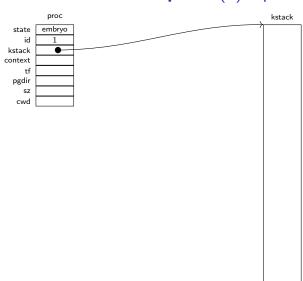
allocproc: (2) Initialize process kernel stack

```
_{2494} if ((p\rightarrow)kstack = kalloc()) == 0) {
     p\rightarrow state = UNUSED:
     return 0;
    sp = p->kstack + KSTACKSIZE:
    sp = -sizeof *p > tf;
    p->tf = (struct trapframe *)sp;
    sp = 4:
    *(uint*)sp = (uint)trapret;
    sp = -sizeof *p -> context:
    p->context = (struct context*)sp;
    memset(p->context, 0, sizeof *p->context);
    p->context->eip = (uint) forkret;
    return p:
```



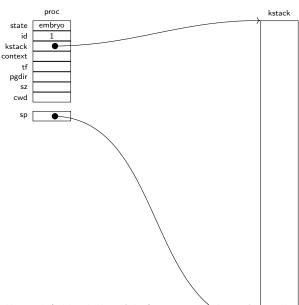
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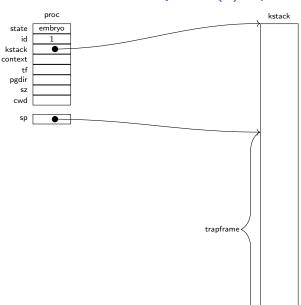
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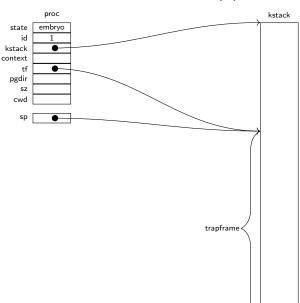
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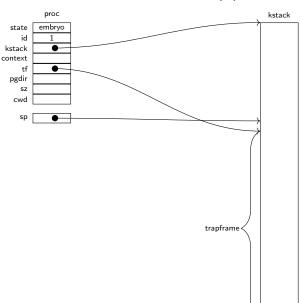
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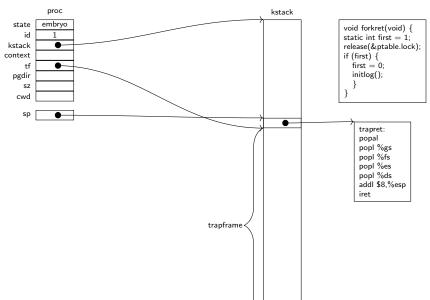
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```

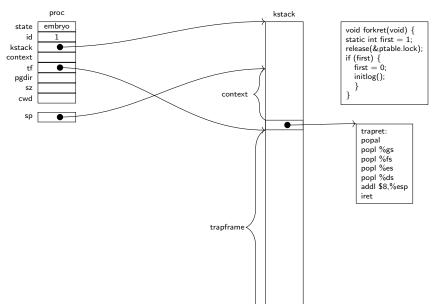
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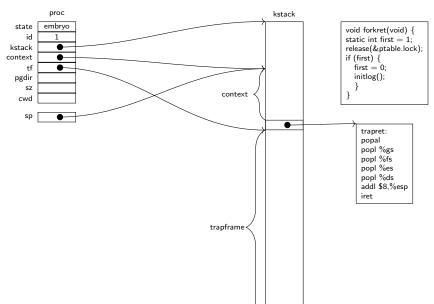


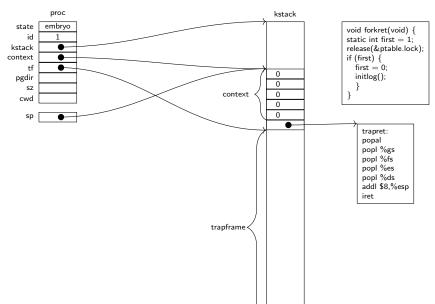
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if (first) {
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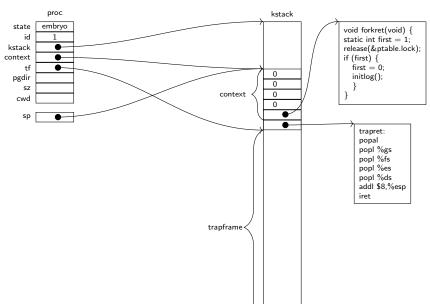
```
trapret:
popal
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addl $8,%esp
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```









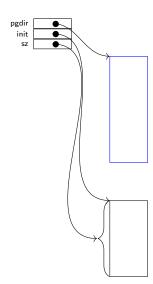


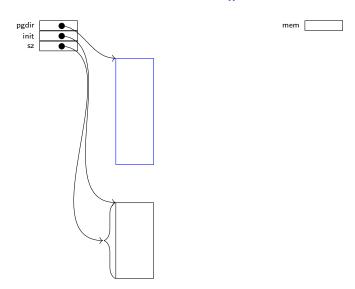
inituvm

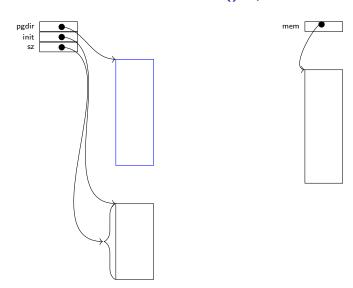
```
inituvm(pde_t *pgdir, char *init, uint sz)
 char *mem;
 if (sz >= PGSIZE)
  panic("inituvm: _more_than_a_page");
 mem = kalloc();
 memset (mem, 0, PGSIZE);
 mappages(pgdir, 0, PGSIZE, v2p(mem), PTE_W|PTE_U);
 memmove(mem, init, sz);
```

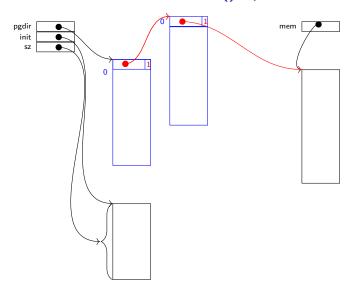
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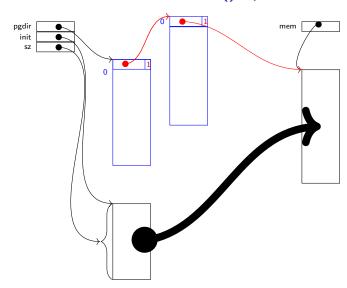








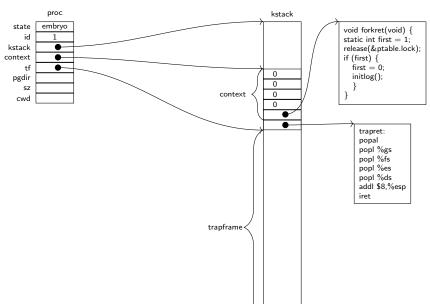


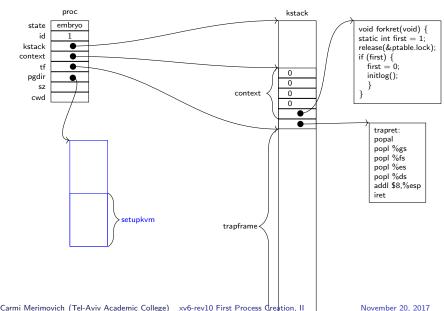


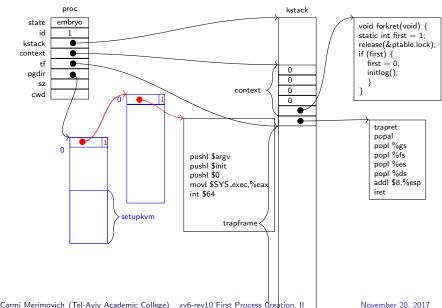
userinit(): (1) Invoking allocproc & inituvm

```
userinit (void) {
 struct proc *p;
 extern char _binary_initcode_start[],
              _binary_initcode_size[];
 p = allocproc();
 initproc = p;
 if ((p->pgdir = setupkvm()) == 0)
  panic("userinit: _out_of_memory?");
 inituvm(p—>pgdir, _binary_initcode_start,
                (int)_binary_initcode_size);
 p->sz = PGSIZE:
```

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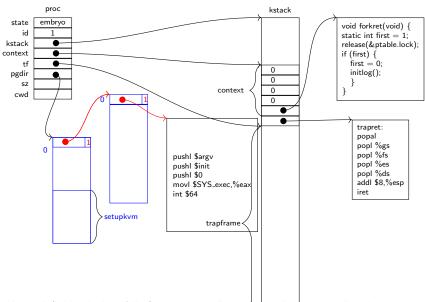


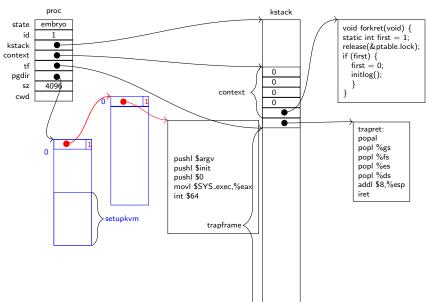


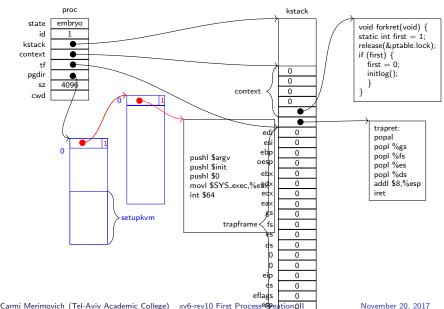
userinit(): (2) trapframe initialization

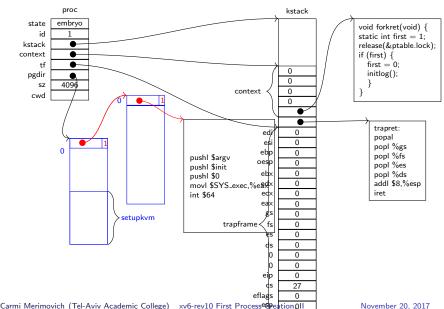
```
memset(p\rightarrow tf, 0, sizeof(*p\rightarrow tf));
p\rightarrow tf\rightarrow cs = (SEG\_UCODE << 3) \mid DPL\_USER;
p\rightarrow tf\rightarrow ds = (SEG\_UDATA << 3) \mid DPL\_USER;
p\rightarrow tf\rightarrow es = p\rightarrow tf\rightarrow ds:
p\rightarrow tf\rightarrow ss = p\rightarrow tf\rightarrow ds:
p\rightarrow tf\rightarrow eflags = FL_IF:
p\rightarrow tf\rightarrow esp = PGSIZE:
p\rightarrow tf\rightarrow eip=0; // beginning of initcode. S
p\rightarrow cwd = namei("/");
p->state = RUNNABLE:
```

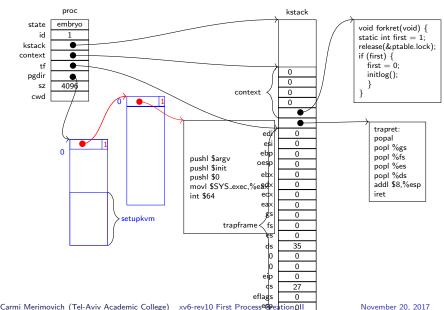
2532

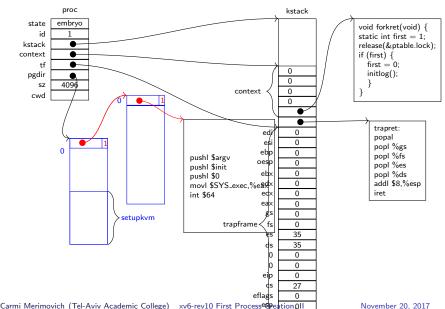


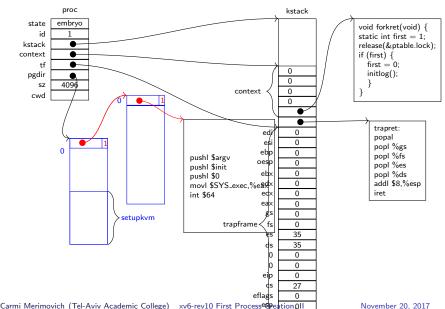


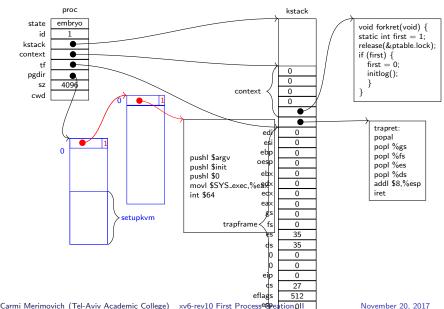


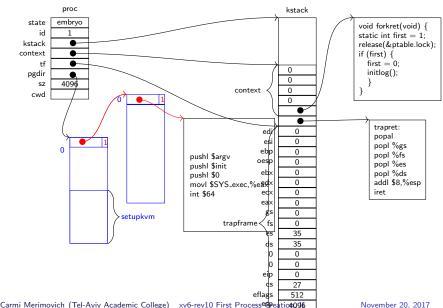


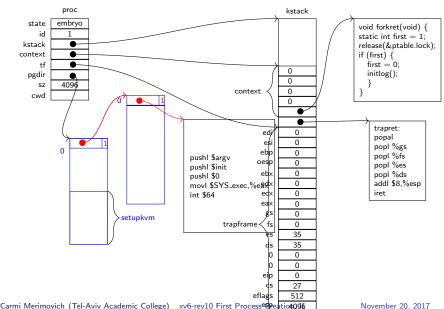


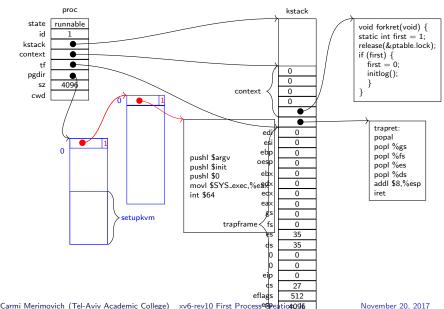












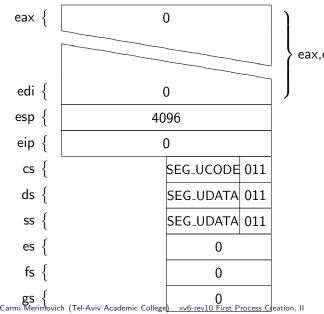
trapret

```
trapret:
 popal
 popl %gs
 popl %fs
 popl %es
 popl %ds
 addl $0x8, %esp \# trapno and errcode
 iret
```

What is the registers state after the iret instruction?

3324

Registers after iret



eax,ecx,edx,ebx,ebp,esi,edi

Informing the CPU about safe stack

The PUSHAL/POPAL instructions

pushl %eax pushl %ecx pushl %edx pushl %ebx pushl %esp !!! pushl %ebp pushl %esi pushl %edi

```
popl %edi
popl %esi
popl %ebp
popl %esp
popl %ebx
popl %edx
popl %ecx
popl %eax
```

What now?

- How the kernel will return to run?
 - Only by interrupt service routine.
- Where is the interrupt service routine located?
 - Later topic,
- What happens to the stack when interrupt is delivered??
 - GOOD POINT.
 - Before going to user mode we set stack address for interrupts.

```
state unused id kstack context tf pgdir sz cwd cr3 tr
```

```
void forkret(void) {
static int first = 1;
release(&ptable.lock)
if (first) {
    first = 0;
    initlog();
    }
}
```

trapret:

```
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $8,%esp
iret
```

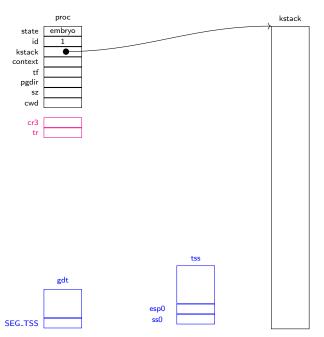


```
state embryo id 1  
kstack context tf pgdir sz cwd cr3 tr
```

```
void forkret(void) {
  static int first = 1;
  release(&ptable.lock)
  if (first) {
    first = 0;
    initlog();
    }
}
```







void forkret(void) {
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trapret: popal popl %gs popl %fs popl %es popl %ds addl \$8,%esp iret

