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The **fork** system call

Carmi Merimovich

Tel-Aviv Academic College

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sys_fork

```
3760  int sys_fork(void) {
        return fork();
      }
```

`fork()`

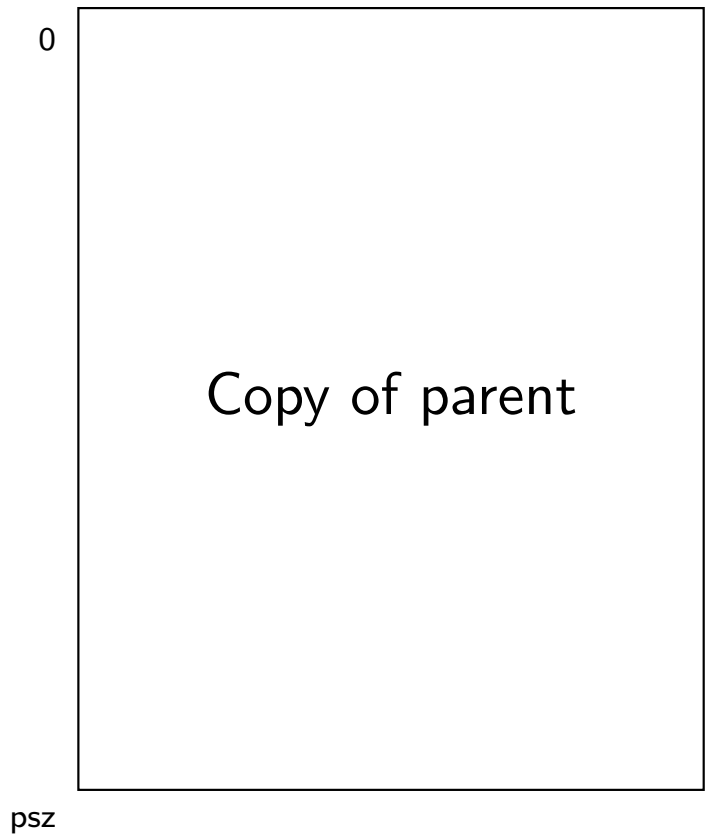
Recall:

- A child (of the invoker) process is created.
- The `pid` of the child process is returned to the invoker.
- The child process is (almost) identical to the parent process.
  - To the child, the return value of the system call is zero.
- So, how do we begin?

Child process state needed

eax	0
ebx	pebx
ecx	pecx
edx	pedx
ebp	pebp
esi	pesi
edi	pedi
esp	pesp
eip	peip

cs	pcs
ds	pds
ss	pss
es	pes
fs	pfs
gs	pgs



proc struct

How do we fill the fields of the new process?

```
uint sz; // @proc->sz@
pde_t* pgdir; // @Serious replication needed@
char *kstack; // @probably allocproc()@
enum procstate state; // @RUNNABLE@
volatile int pid; // @allocproc()@
struct proc *parent; // @proc@
struct trapframe *tf; // @allocproc()@
struct context *context; // @allocproc()@
void *chan; // @0@
int killed; // @0@
struct file *ofile[NOFILE]; // @filedup()@ (when st
struct inode *cwd; // @idup()@ (when studying fs)
char name[16]; // @proc->name@
};
```

Needed work

- proc struct and friends.
- Filling `pgdir` requires considerable code replication.
- Replicatin ofile and cwd requires help from the relevant modules.

proc struct and friends

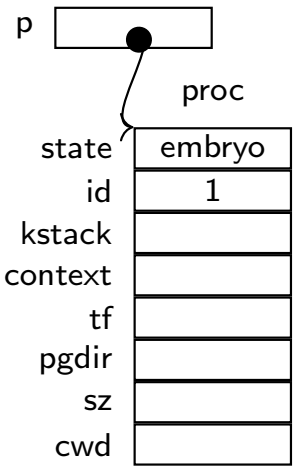
allocproc(): (1) Finding unused `proc` structure

```
2473 static struct proc *allocproc(void) {
      struct proc *p;
      char *sp;

      acquire(&ptable.lock);
      for(p = ptable.proc; p < &ptable.proc[NPROC]; p++)
       if (p->state == UNUSED)
        goto found;
      release(&ptable.lock);
      return 0;

     found:
      p->state = EMBRYO;
      p->pid = nextpid++;
      release(&ptable.lock);
```

allocproc(): (1) Operation




```
void forkret(void) {
static int first = 1;
release(&ptable.lock);
if (first) {
   first = 0;
   initlog();
   }
}
```

```
trapret:
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $8,%esp
iret
```

allocproc: (2) Initialize process kernel stack

```
2494  if((p->kstack = kalloc()) == 0){
        p->state = UNUSED;
        return 0;
      }
      sp = p->kstack + KSTACKSIZE;
      sp -= sizeof *p->tf;
      p->tf = (struct trapframe*)sp;
      sp -= 4;
      *(uint*)sp = (uint)trapret;
      sp -= sizeof *p->context;
      p->context = (struct context*)sp;
      memset(p->context, 0, sizeof *p->context);
      p->context->eip = (uint)forkret;
      return p;
    }
```

allocproc: (2) Operation

proc

state	embryo
id	1
kstack	
context	
tf	
pgdir	
sz	
cwd	

```
void forkret(void) {
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if (first) {
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}
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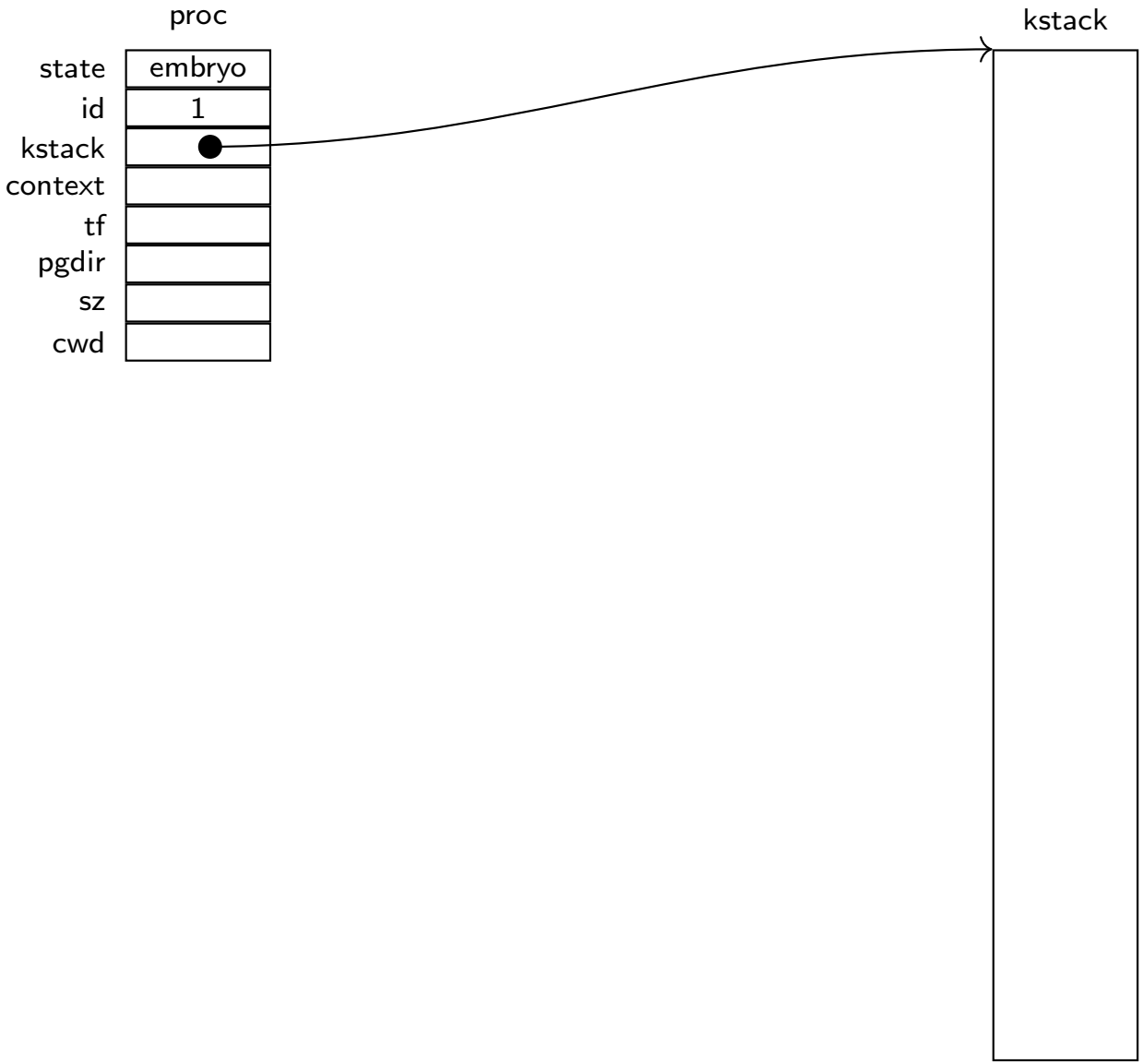
allocproc: (2) Operation

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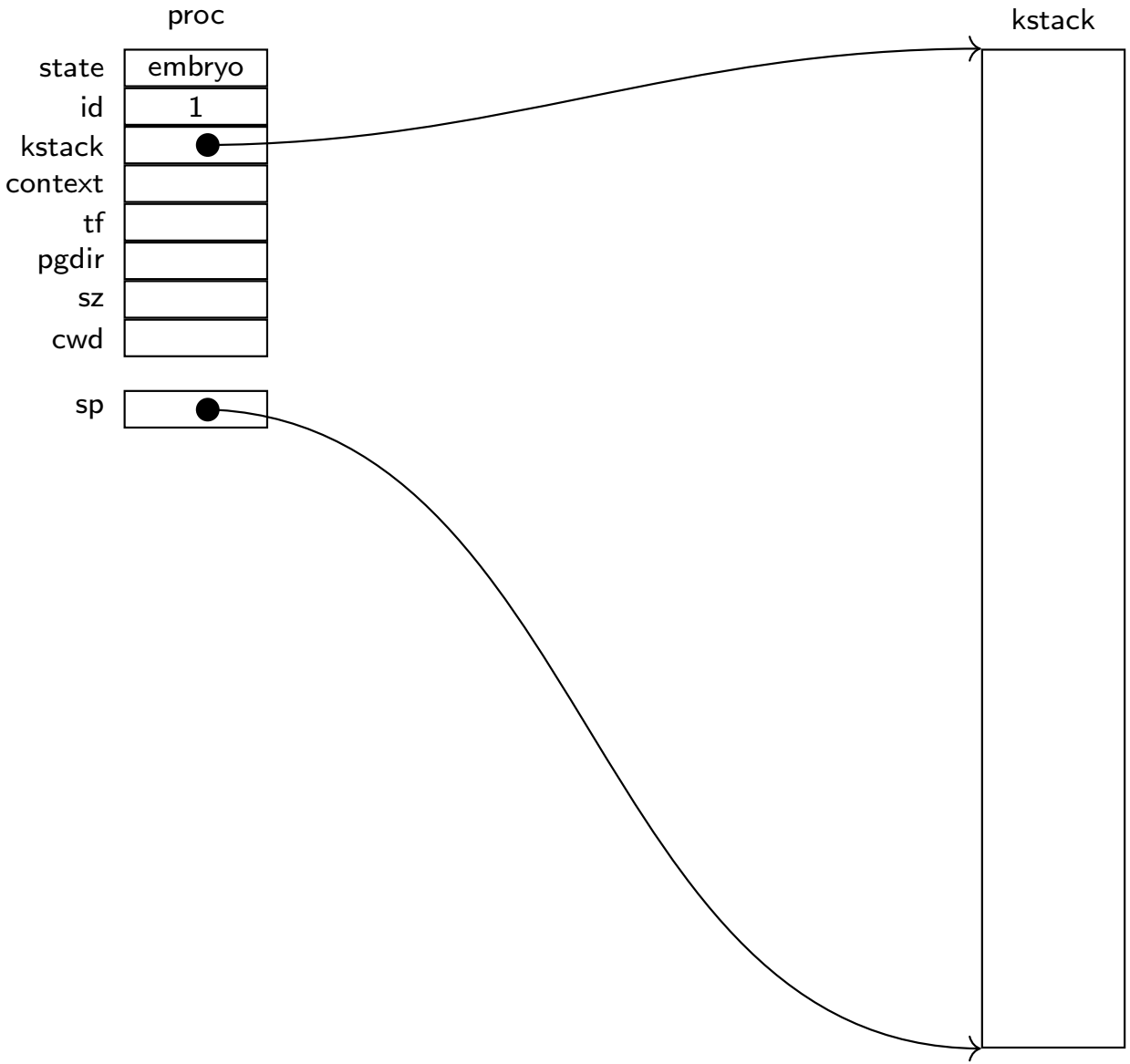
allocproc: (2) Operation

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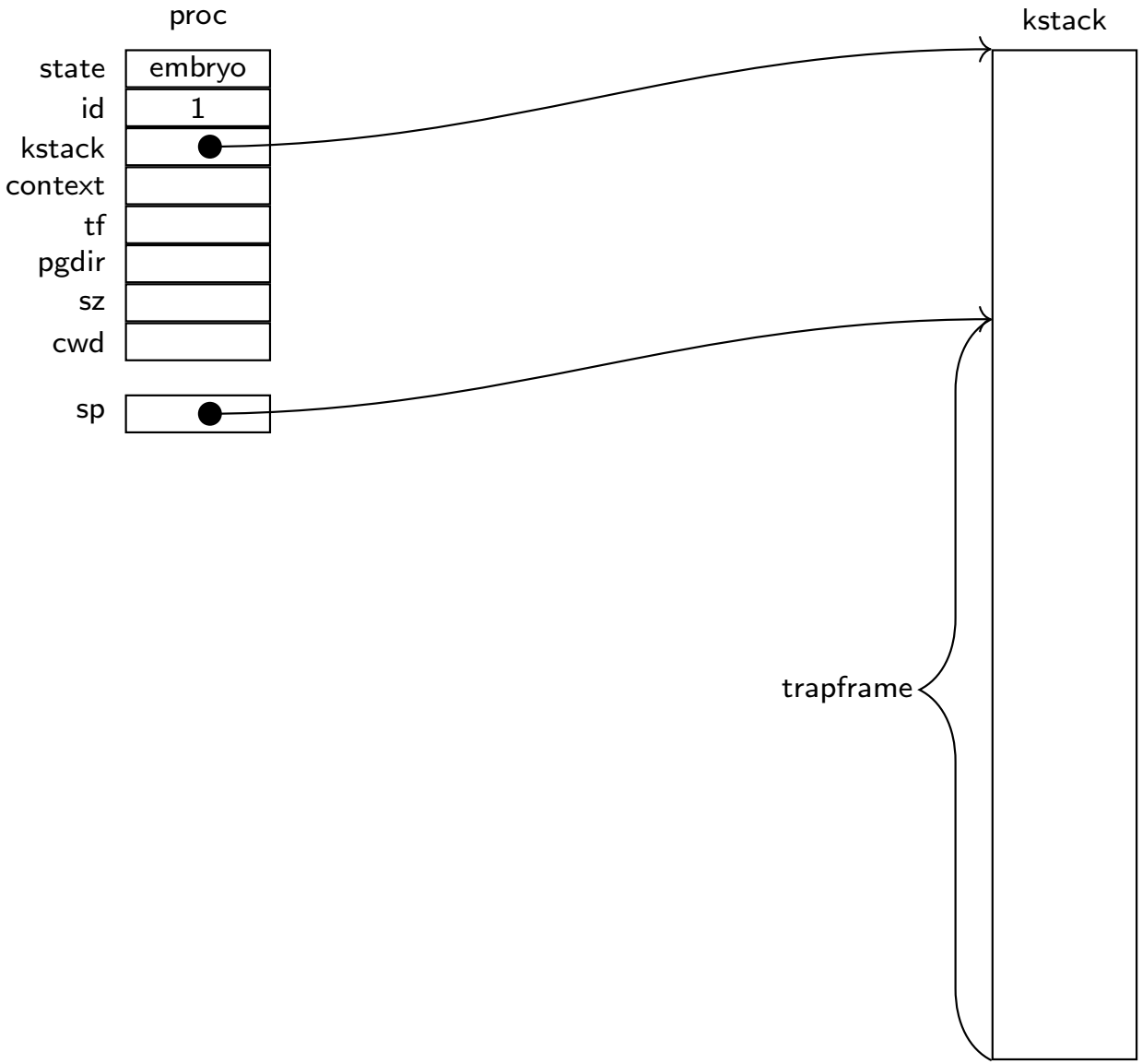
allocproc: (2) Operation

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}
```

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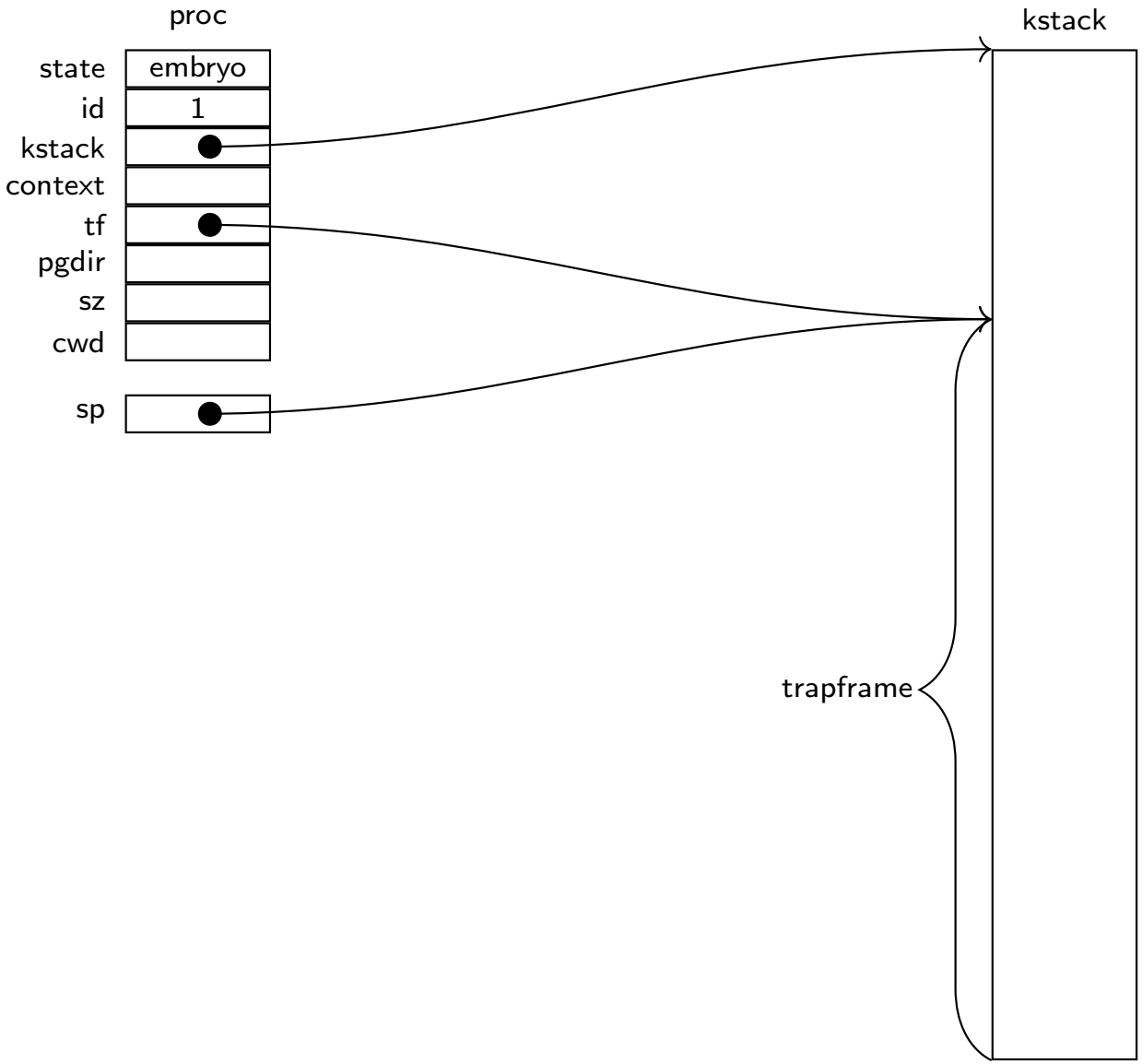
allocproc: (2) Operation

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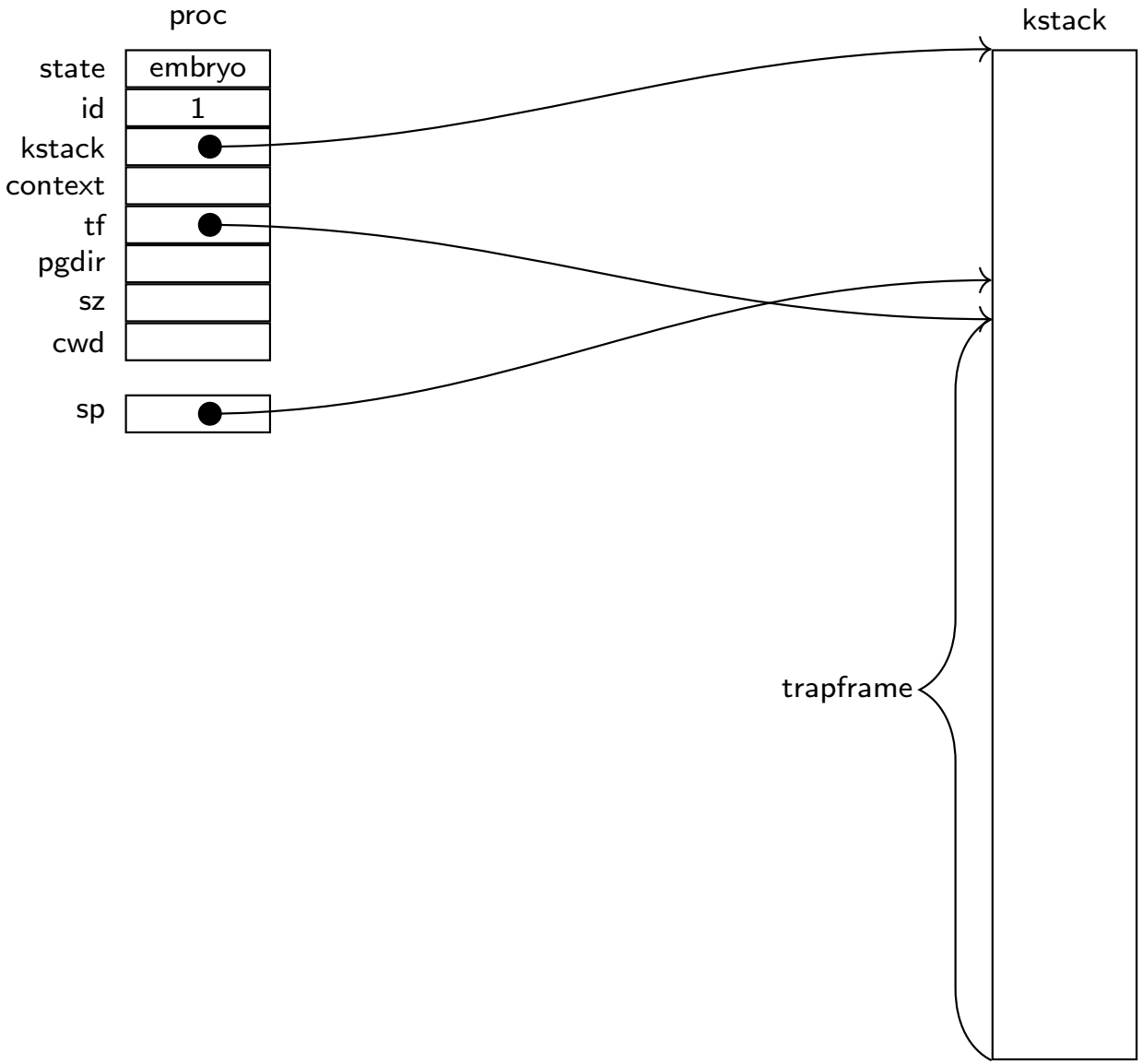


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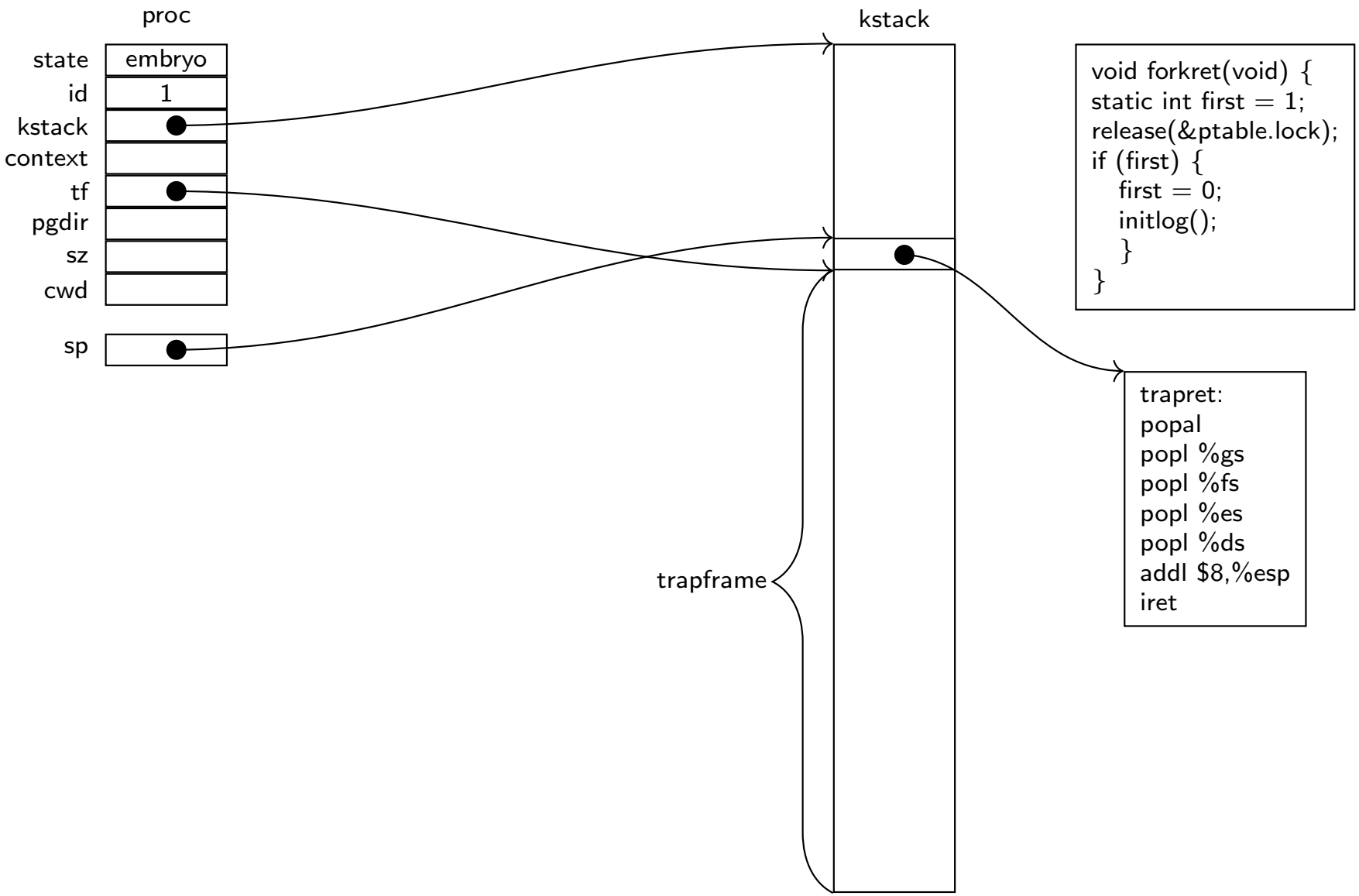
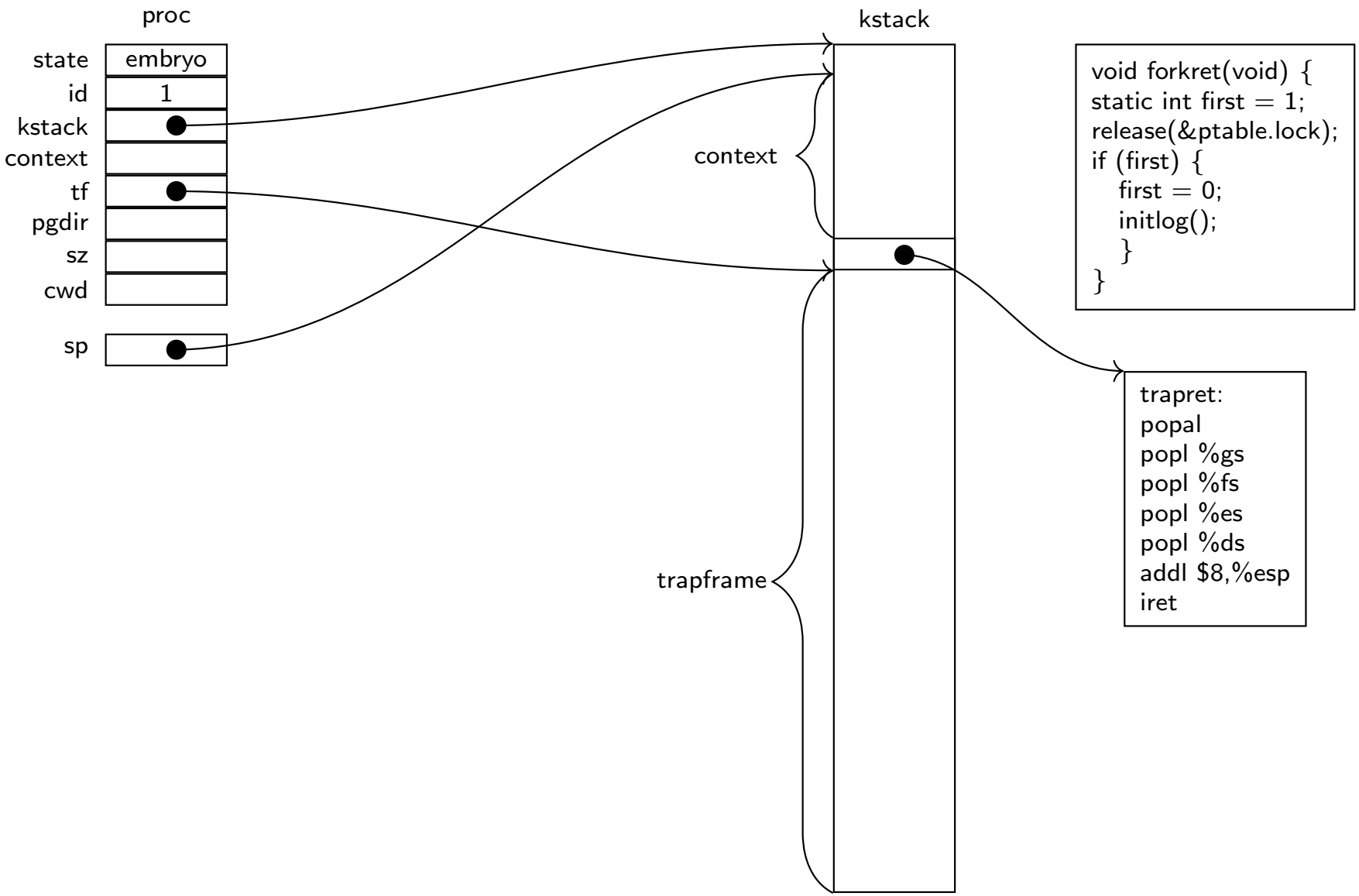
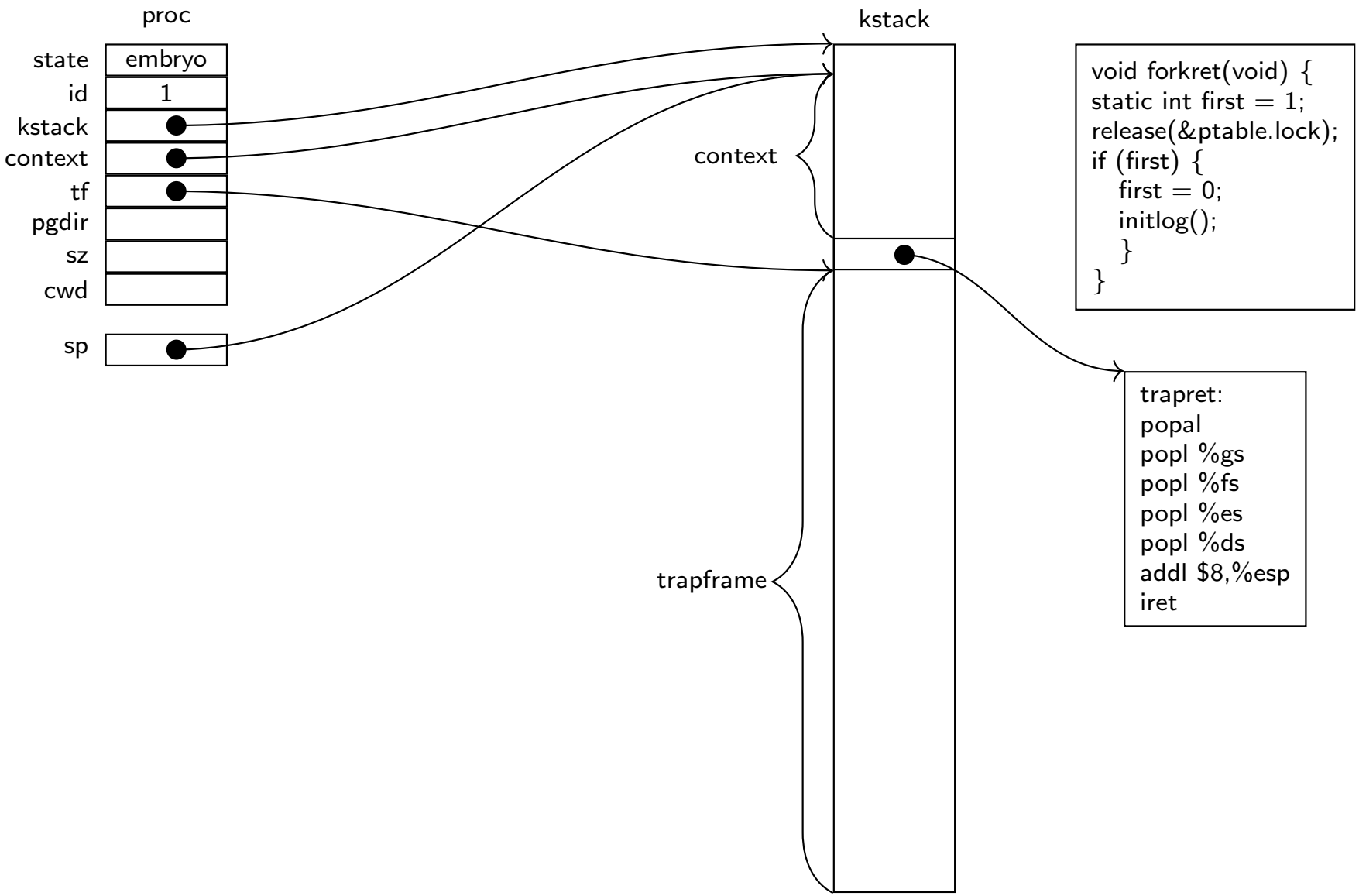
allocproc: (2) Operation



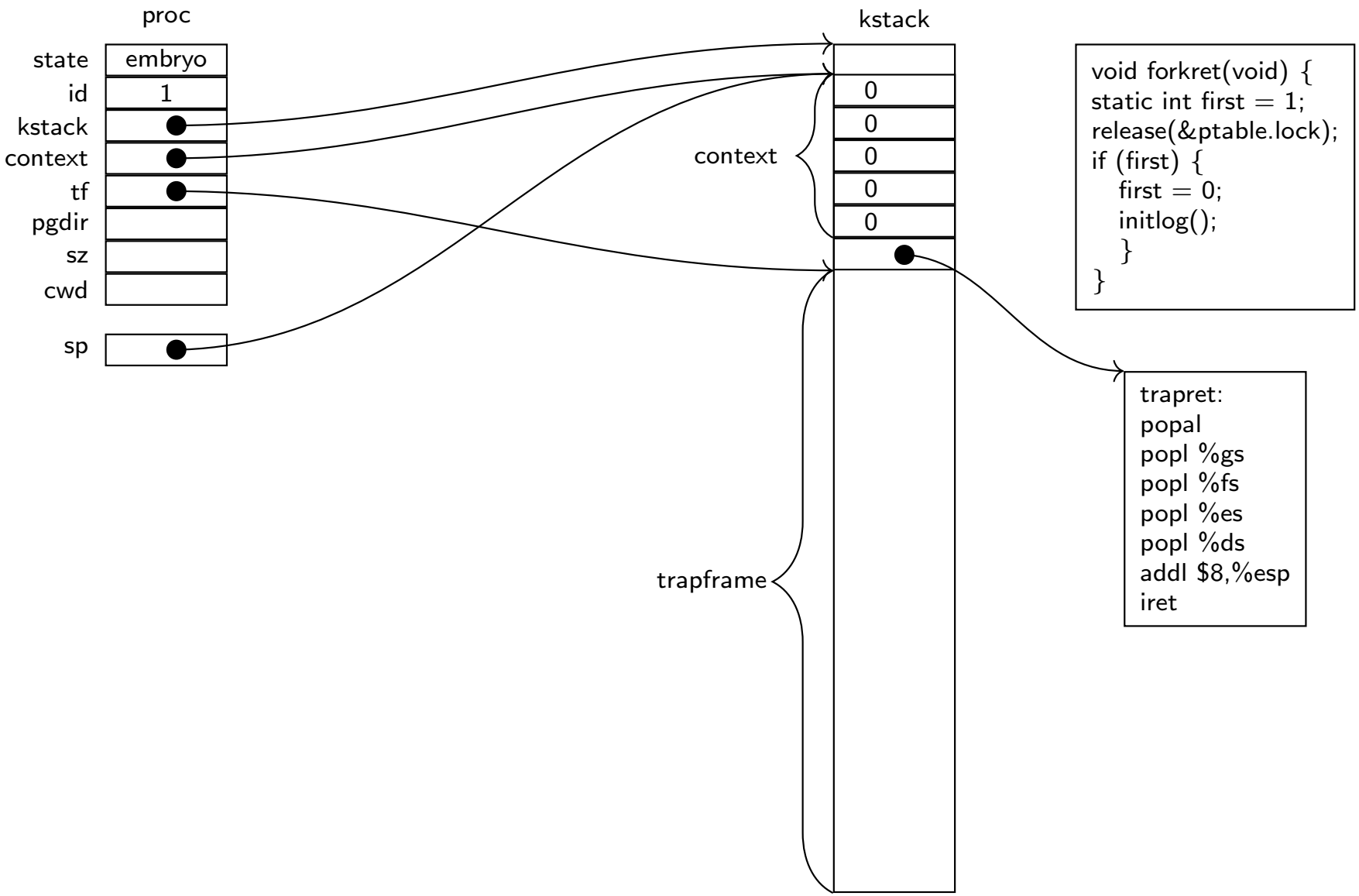
allocproc: (2) Operation



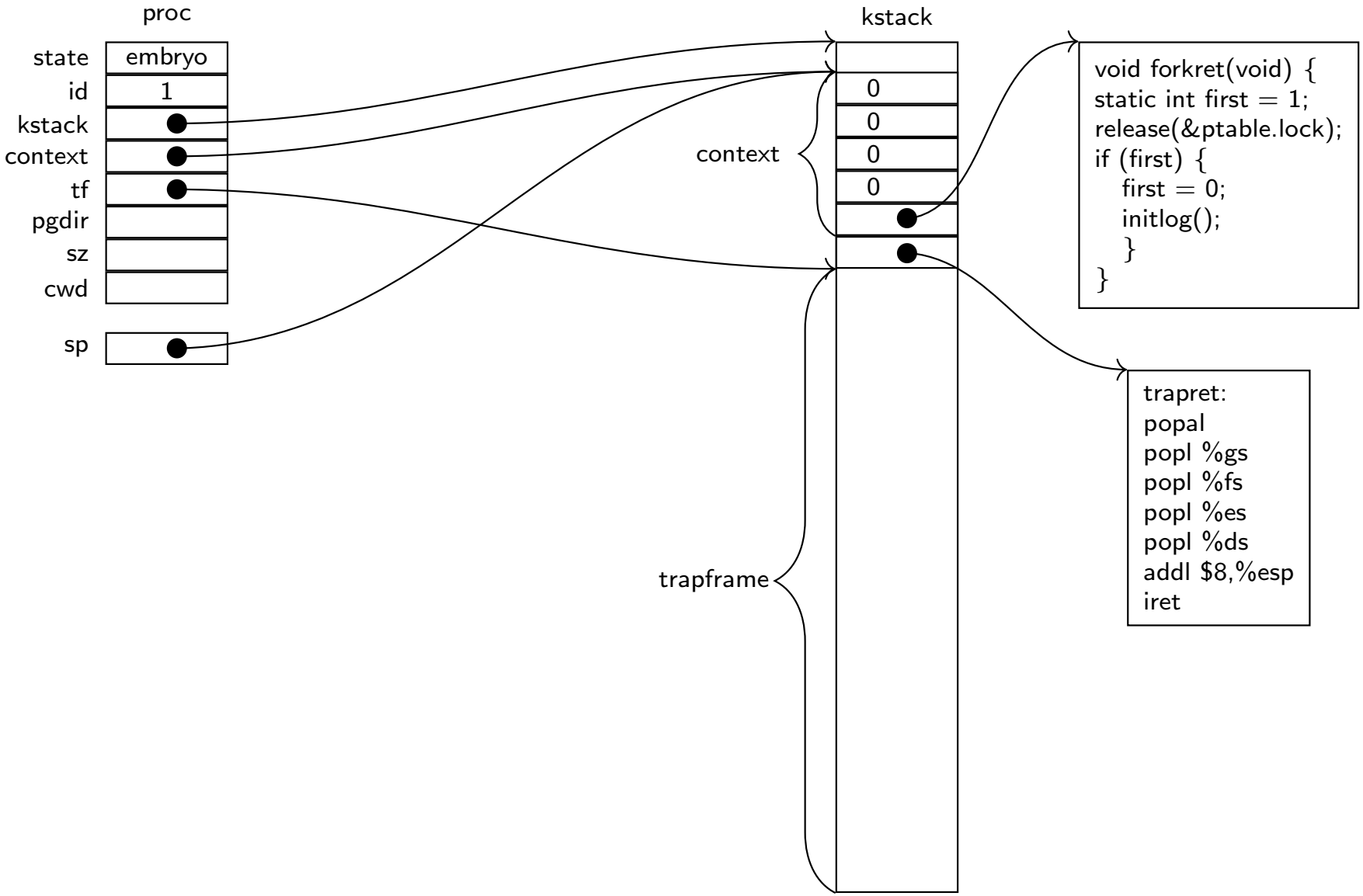
allocproc: (2) Operation



allocproc: (2) Operation



allocproc: (2) Operation



User space replication

(two methods)

Replcation page 0 of current process only

Replicating current process first page

- Allocating new block of memory:

```
dst = kalloc();
```

- Copying. One of the following is possible:

1.
```
memmove(dst,0,4096);
```

```
p = walkpgdir(myproc()->pgdir,0,0);
memmove(dst,p2v(PTE_ADDR(*p)),4096);
```

2. This is of course useless as it is.

Create new address space

Replcation page 0 of current process only

Add mapping rules in the new address space to the replication

New address space and mapping

- Creating a new address space:

```
pgdir = setupkvm();
```

- Allocate and copy:

```
dst = kalloc();
p = walkpgdir(myproc()->pgdir,0,0);
memmove(dst,p2v(PTE_ADDR(*p)),4096);
```

- Adding translation rule:

```
mappages(pgdir, 0,4096,v2p(dst), (*p) & 4095);
```

Replcating means doing the copy and translation for each page.

Replicate and Map ALL user space pages

Replicating currnet process pages

NO ERROR CHECKING IN HERE!

```
pgdir = setupkvm();

for (va=0; va<myproc()->sz; va += PGSIZE) {
 kva = kalloc();
 memmov(kva, va, PGSIZE);
 pte = walkpgdir(myproc()->pgdir,va,0);
 mappages(pgdir,va,PGSIZE,v2p(kva),(*pte) & 4095);
}
```

- If there is allocation error, all previous allocations must be freed!
- We show freeing on the next slide.

Freeing address space

```
for (i = 0; i < 512; i++) {
  if ((pgdir[i] & PTE_P) == 0)
    continue;
  pgtbl = p2v(pgdir[i] & -4096);
  for (j=0; j < 1024; j++) {
    if (pgtbl[j] & PTE_P) {
      kfree(p2v(pgtbl[j] & ~4095));
    }
  }
  kfree(pgtbl);
}
for (i = 512; i < 1023; i++) {
  if ((pgdir[i] & PTE_P) == 0)
    continue;
  pgtbl = p2v(pgdir[i] & -4096);
  kfree(pgtbl);
}
kfree(pgdir);
```

xv6 code for user space replication

(More general than needed)

xv6 replicating and freeing address space

The following xv6 code replicates arbitrary address space.

- The code checks for allocation errors.
- It deallocates all previous allocation in case of failure.

copyuvm()

```
2035  pde_t* copyuvm(pde_t *pgdir, uint sz) {
       pde_t *d; pte_t *pte;
       uint pa, i;
       char *mem;
       if ((d = setupkvm()) == 0)  return 0;
       for (i = 0; i < sz; i += PGSIZE) {
        if ((pte = walkpgdir(pgdir, (void *) i, 0)) == 0) panic
        if (!(*pte & PTE_P)) panic("copyuvm: page not present")
        pa = PTE_ADDR(*pte);
        if ((mem = kalloc()) == 0)  goto bad;
        memmove(mem, (char*)p2v(pa), PGSIZE);
        if (mappages(d, (void*)i, PGSIZE, v2p(mem),
              PTE_FLAGS(*pte)) < 0) goto bad;
       }
       return d;
      bad:
       freevm(d);
       return 0;
      }
```

freevm()

```
2003 void freevm(pde_t *pgdir) {
  uint i;

  if (pgdir == 0)
    panic("freevm: no pgdir");
  deallocuvm(pgdir, KERNBASE, 0);
  for (i = 0; i < NPDENTRIES; i++) {
  if (pgdir[i] & PTE_P) {
    char * v = p2v(PTE_ADDR(pgdir[i]));
    kfree(v);
    }
  }
  kfree((char*)pgdir);
}
```

deallocuvm

```
1961 deallocuvm(pde_t *pgdir, uint oldsz, uint newsz) {
  pte_t *pte;
  uint a, pa;
  if (newsz >= oldsz)  return oldsz;
  a = PGROUNDUP(newsz);
  for (; a < oldsz; a += PGSIZE) {
   pte = walkpgdir(pgdir, (char*)a, 0);
   if (!pte) a += (NPTENTRIES - 1) * PGSIZE;
   else if ((*pte & PTE_P) != 0) {
    pa = PTE_ADDR(*pte);
    if (pa == 0) panic("kfree");
    char *v = p2v(pa);
    kfree(v);
    *pte = 0;
   }
  }
  return newsz;
```

fork

- `allocproc`:
  - An `EMBRYO` `proc` struct is constructed.
  - A kernel stack is allocated.
  - An uninitialized `trapframe` is allocated.
  - An artificial `context` is constructed.
- The user space memory of the caller is replicated.
- A new matching page table is constructed.
- The `trapframe` of the caller is copied to the uninitialized `trapframe`.
- The `eax` field of the new `trapframe` is cleared.
- File pointers are replicated.
- Rest of the caller **proc** struct fields are copied to the new **proc** struct.

fork (1)

```
2580  int fork(void) {
       int i, pid;
       struct proc *np;

       if ((np = allocproc()) == 0)
        return -1;

       if ((np->pgdir=copyuvm(myproc()->pgdir,myproc()->sz
                        == 0) {
        kfree(np->kstack);
        np->kstack = 0;
        np->state = UNUSED;
        return -1;
       }
```

fork (2)

```
  np->sz = myproc()->sz;
  np->parent = myproc();
  *np->tf = *myproc()->tf;

  np->tf->eax = 0;

  for(i = 0; i < NOFILE; i++)
    if(myproc()->ofile[i])
      np->ofile[i] = filedup(myproc()->ofile[i]);
  np->cwd = idup(myproc()->cwd);

  safestrcpy(np->name, myproc()->name, sizeof(myproc()
  pid = np->pid;
  np->state = RUNNABLE;
  return pid;
}
```

How and when the child runs?!

Recall the scheduler

scheduler

```
2758  void scheduler(void) {
       struct proc *p;
       struct cpu *c = mycpu();
       c->proc = 0;
       for(;;) { sti();
        acquire(&ptable.lock);
        for(p = ptable.proc; p < &ptable.proc[NPROC]; p++) {
         if (p->state != RUNNABLE) continue;

         c->proc = p;
         switchuvm(p);
         p->state = RUNNING;
         swtch(&c->scheduler, p->context);
         switchkvm();
         c->proc = 0;
        }
        release(&ptable.lock);
       }
      }
```

switchuvm

If switching to user mode is expected:

- The `tr` register should contain the index of a TSS descriptor.
- The `TSS` descriptor should point to a `taskstate` structure.
- The `ss0` and `esp0` fields should point to a valid kernel stack top.
- The above is ESSENTIAL for proper interrupt service in user mode.

switchuvm

```
1860  void switchuvm(struct proc *p) {
        pushcli();
        mycpu()->gdt[SEG_TSS] = SEG16(STS_T32A,
                                      &mycpu()->ts,
                                      sizeof(mycpu()->ts)-1, 0)
        mycpu()->gdt[SEG_TSS].s = 0;
        mycpu()->ts.ss0 = SEG_KDATA<<3;
        mycpu()->ts.esp0 = (uint)p->kstack + KSTACKSIZE;
        mycpu()->ts.iomb = (ushort) 0xFFFF;
        ltr(SEG_TSS << 3);
        if (p->pgdir == 0)
          panic("switchuvm: no pgdir");
        lcr3(v2p(p->pgdir)); // switch to new address space
        popcli();
      }
```

taskstate (hardware structure)

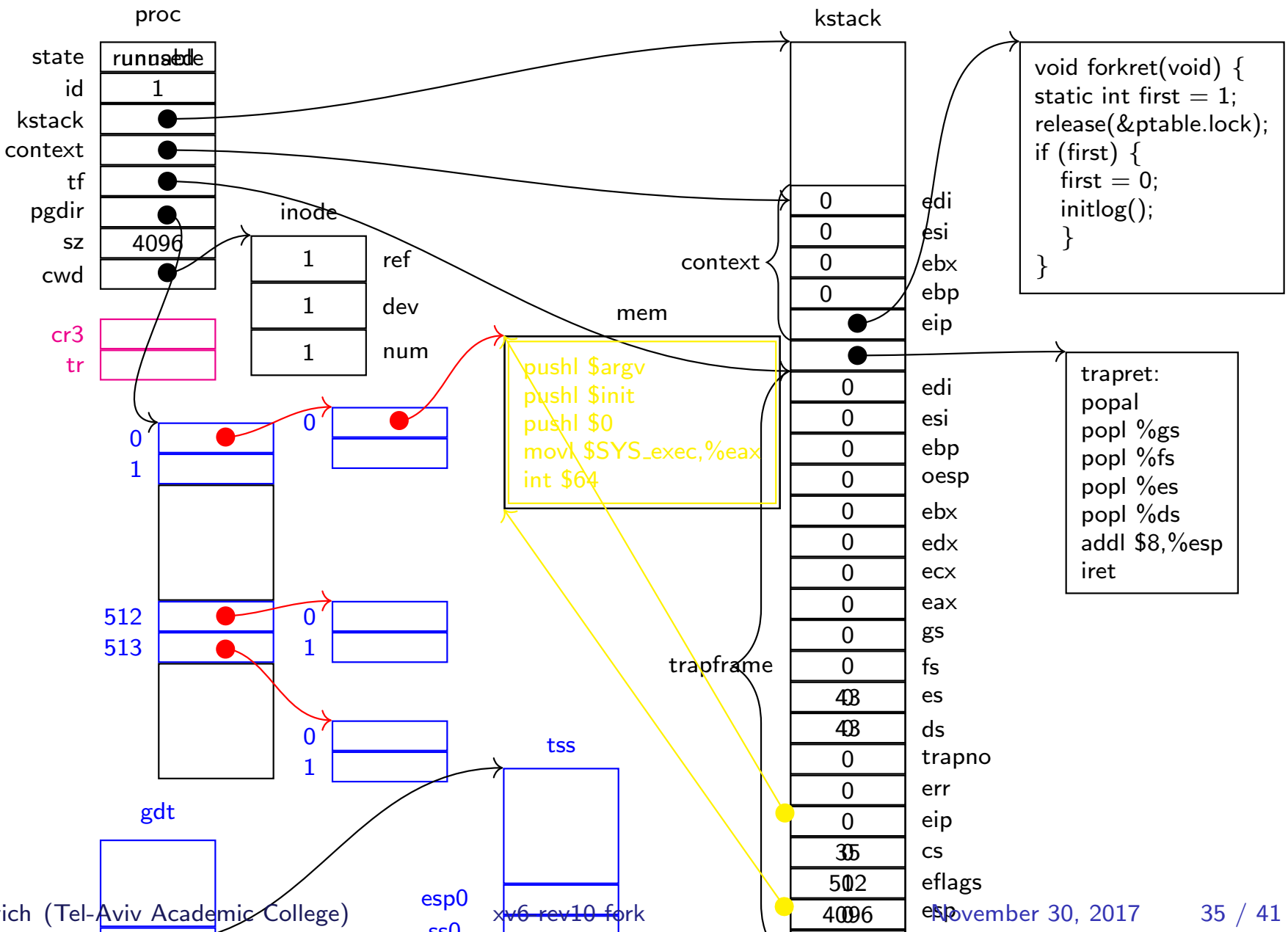
link	
esp0	
	ss0
esp1	
	ss1
esp2	
	ss2
cr3	
eip	
eflags	
eax	
ecx	
edx	
ebx	
esp	
ebp	
esi	
edi	

	es
	cs
	ss
	ds
	fs
	gs
	ldt
	t
iomb	

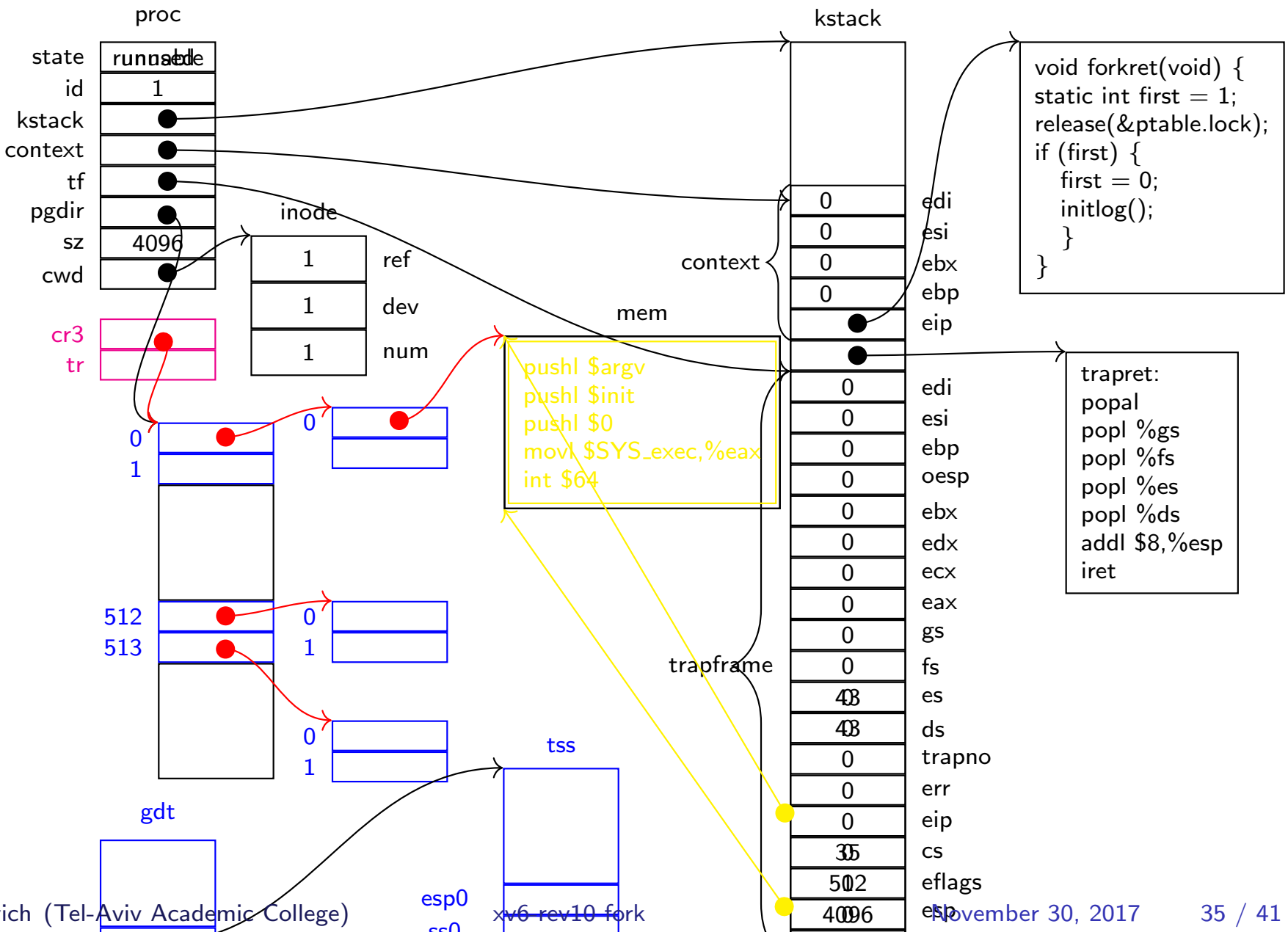
taskstate in C

```
851 struct taskstate
    {
     uint link;
     uint esp0;
     ushort ss0;
     ushort padding1;
     uint *esp1;
     ushort ss1;
     ushort padding2;
     uint *esp2;
     ushort ss2;
     ushort padding3;
     void *cr3;
     uint *eip;
     uint eflags;
     uint eax;
     uint ecx;
     uint edx;
     uint ebx;
     uint *esp;
     uint *ebp;
     uint esi;
     uint edi;
     ushort es;
     ushort padding4;
     ushort cs;
     ushort padding5;
     ushort ss;
     ushort padding6;
     ushort ds;
     ushort padding7;
     ushort fs;
     ushort padding8;
     ushort gs;
     ushort padding9;
     ushort ldt;
     ushort padding10;
     ushort t;
     ushort iomb;
    };
```

switchuvm


proc
state
id
kstack
context
tf
pgdir
sz
cwd
runnable
1
4096
inode
1
1
1
ref
dev
num
cr3
tr
kstack
context
edi
esi
ebx
ebp
eip
edi
esi
ebp
oesp
ebx
edx
ecx
eax
gs
fs
es
ds
trapno
err
eip
cs
eflags
esp
trapframe
mem
pushl $argv
pushl $init
pushl $0
movl $SYS_exec,%eax
int $64
void forkret(void) {
static int first = 1;
release(&ptable.lock);
if (first) {
first = 0;
initlog();
}
}
trapret:
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $8,%esp
iret
0
1
512
513
gdt
tss
esp0
ss0

switchuvm


proc
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inode
1
1
1
ref
dev
num
cr3
tr
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pushl $argv
pushl $init
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movl $SYS_exec,%eax
int $64
kstack
context
edi
esi
ebx
ebp
eip
trapframe
edi
esi
ebp
oesp
ebx
edx
ecx
eax
gs
fs
es
ds
trapno
err
eip
cs
eflags
esp
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1
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0
1
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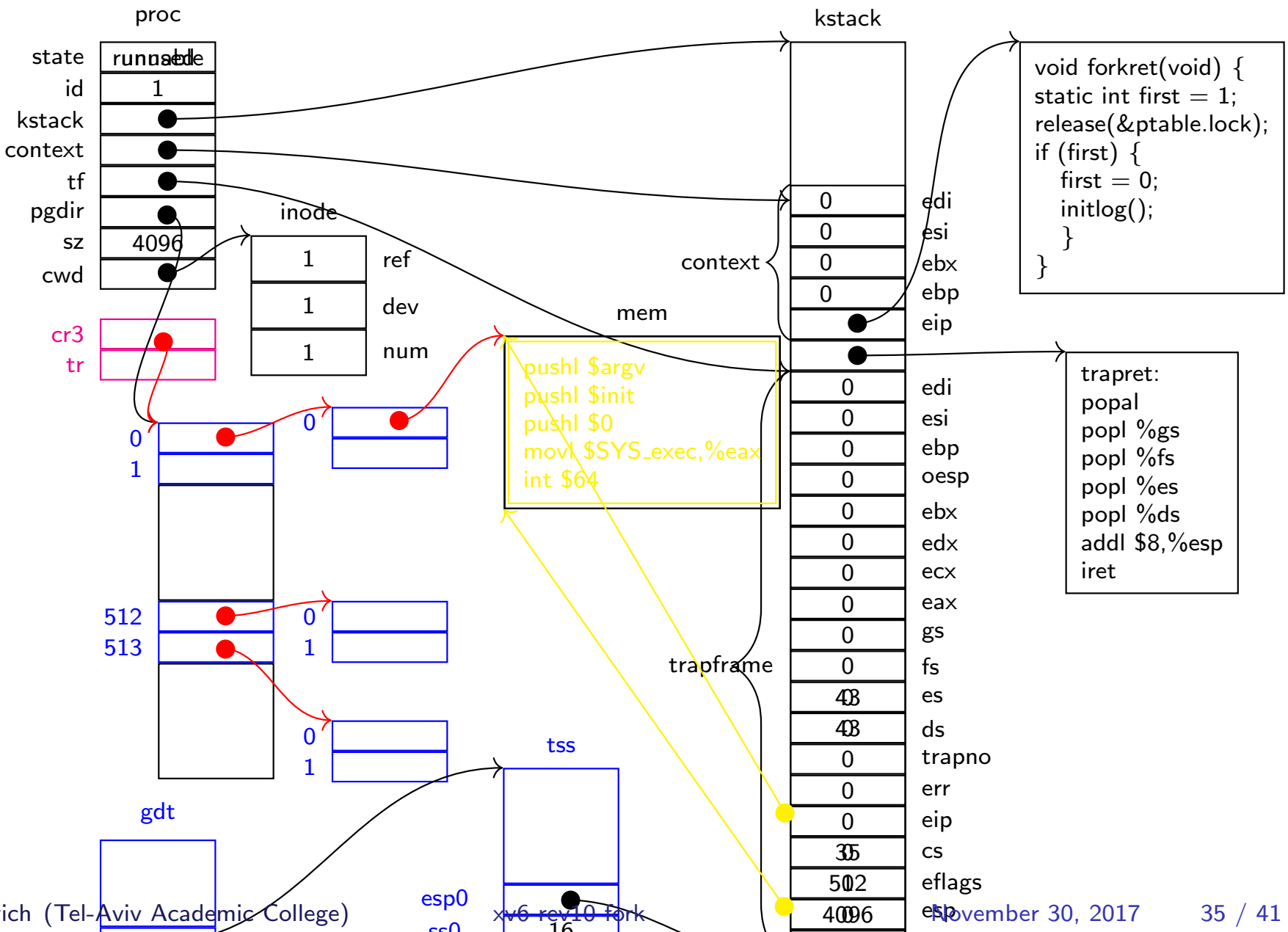
iret

switchuvm


proc
state
id
kstack
context
tf
pgdir
sz
cwd
runnable
1
4096
inode
1
1
1
ref
dev
num
cr3
tr
mem
pushl $argv
pushl $init
pushl $0
movl $SYS_exec,%eax
int $64
kstack
context
edi
esi
ebx
ebp
eip
trapframe
edi
esi
ebp
oesp
ebx
edx
ecx
eax
gs
fs
es
ds
trapno
err
eip
cs
eflags
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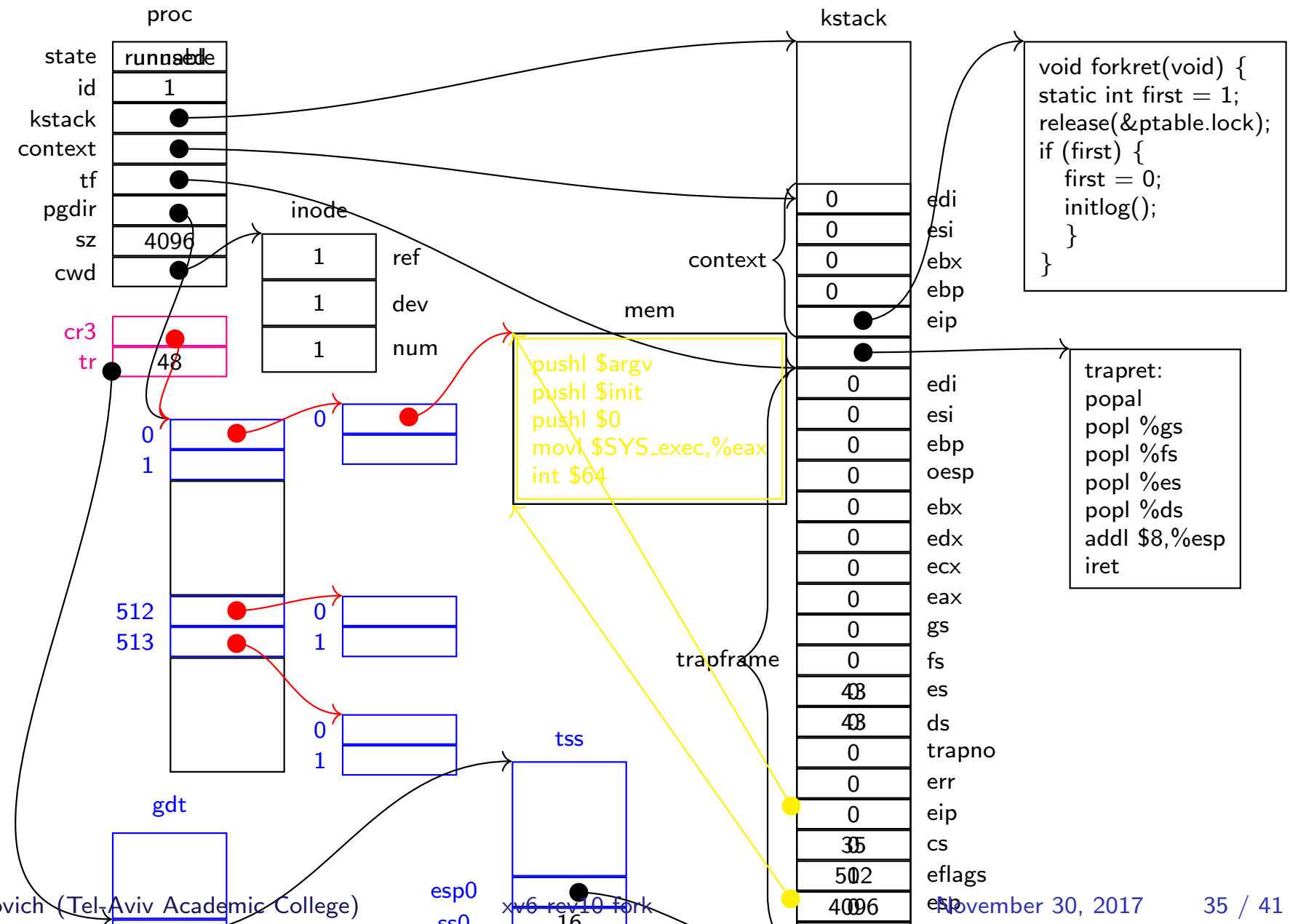
popl %es
popl %ds
addl $8,%esp
iret
0
1
512
513
tss
gdt
esp0
ss0

switchuvm


proc
state
id
kstack
context
tf
pgdir
sz
cwd
runnable
1
4096
inode
1
1
1
ref
dev
num
cr3
tr
48
mem
pushl $argv
pushl $init
pushl $0
movl $SYS_exec,%eax
int $64
kstack
context
edi
esi
ebx
ebp
eip
trapframe
edi
esi
ebp
oesp
ebx
edx
ecx
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err
eip
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if (first) {
first = 0;
initlog();
}
}
trapret:
popal
popl %gs
popl %fs
popl %es
popl %ds
addl $8,%esp
iret
0
1
512
513
tss
gdt
esp0
ss0
16

swtch

co-routines

- The `scheduler` switches to a process by using:

```
2478      swtch(&c->scheduler, p->context);
```

- A process leaves the cpu by returning to the `scheduler` using:

```
2516      swtch(&p->context, mycpu()->scheduler);
```

- We have here co-routines.

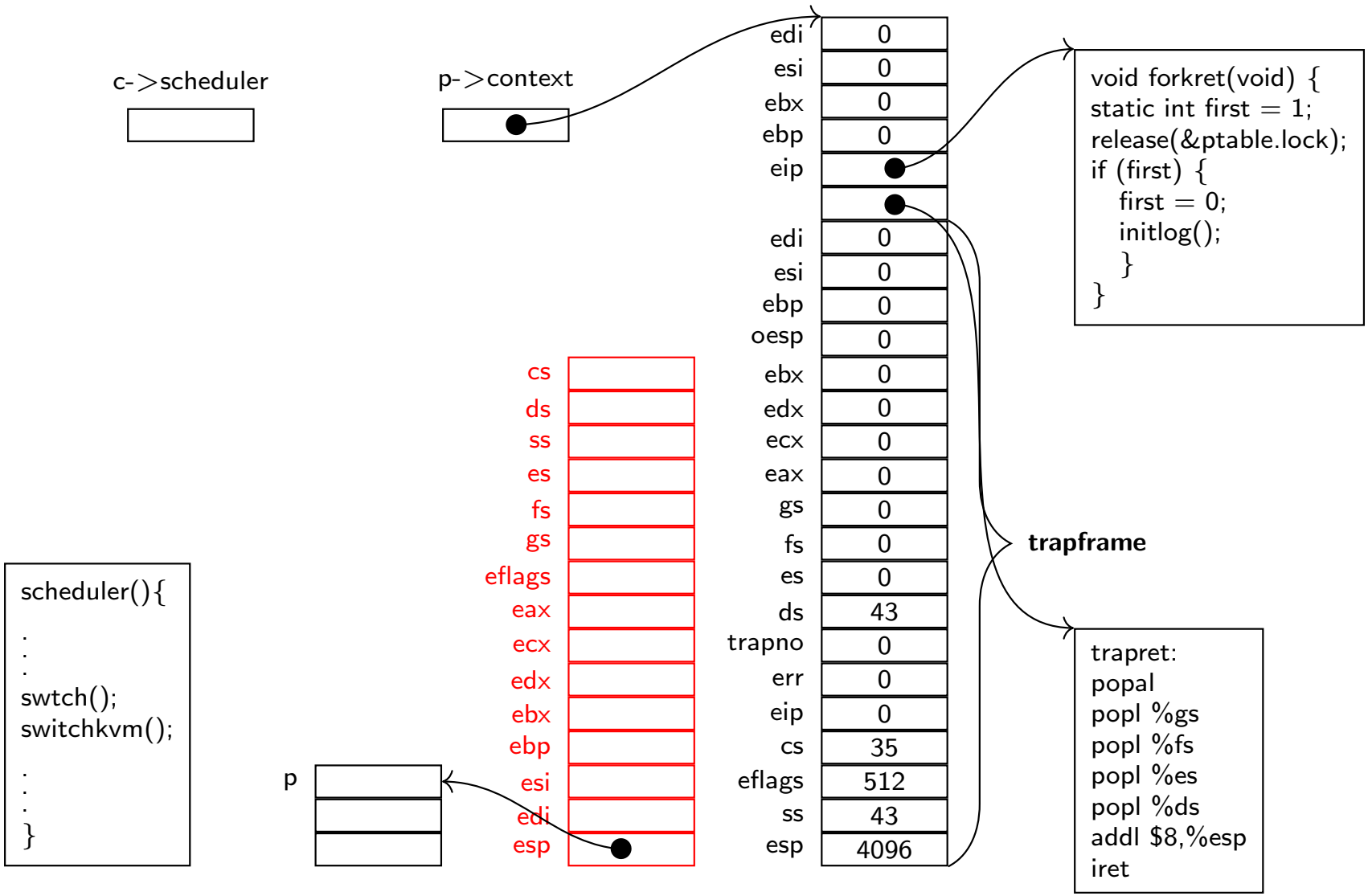
swtch()

```
3058 .globl swtch
swtch:
  movl 4(%esp), %eax
  movl 8(%esp), %edx

  pushl %ebp
  pushl %ebx
  pushl %esi
  pushl %edi

  movl %esp, (%eax)
  movl %edx, %esp
```

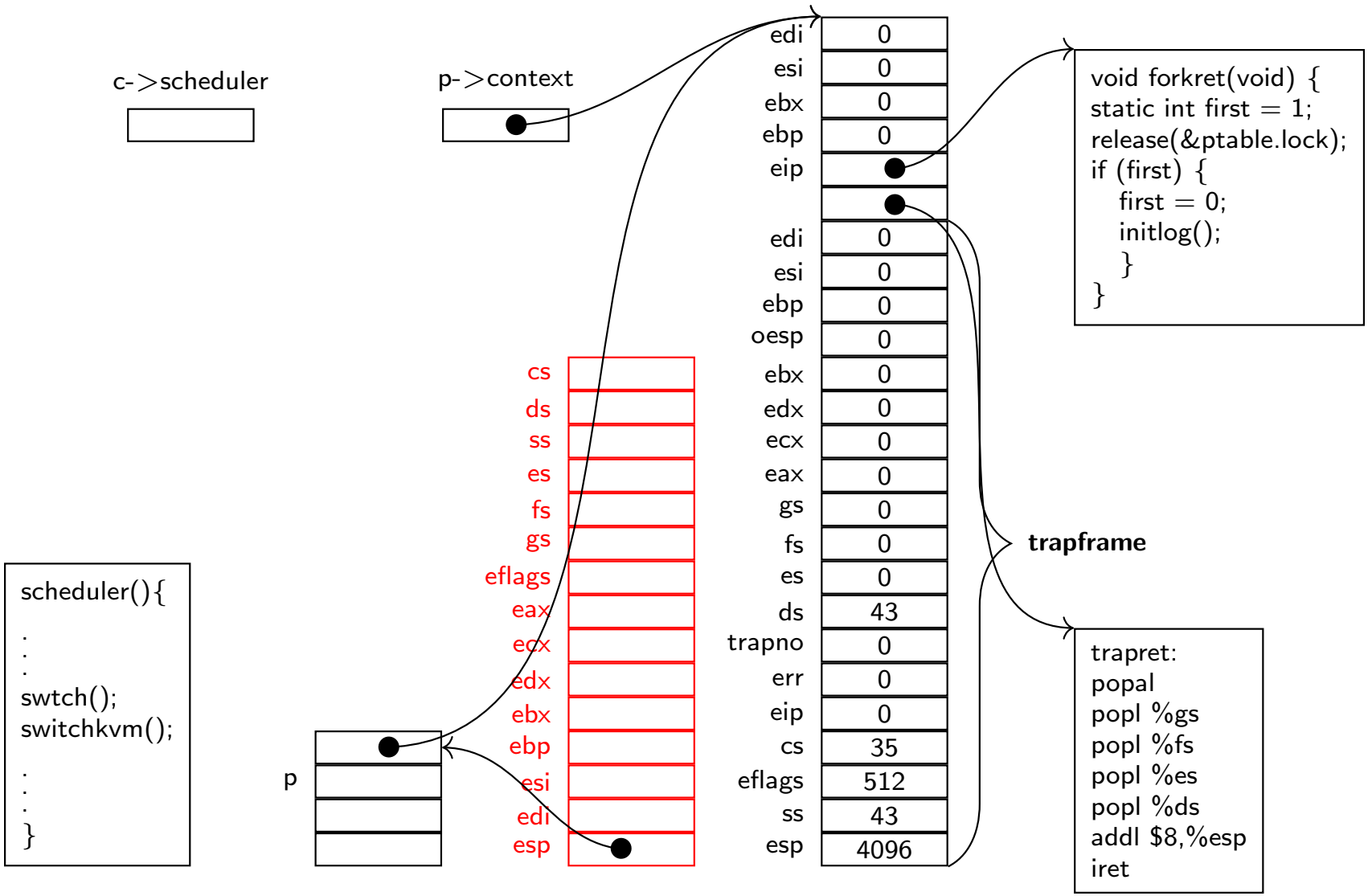
```
  popl %edi
  popl %esi
  popl %ebx
  popl %ebp
  ret
```

The `eip` field of `context` is generated by the instruction calling `swtch`.

swtch() operation



swtch() operation



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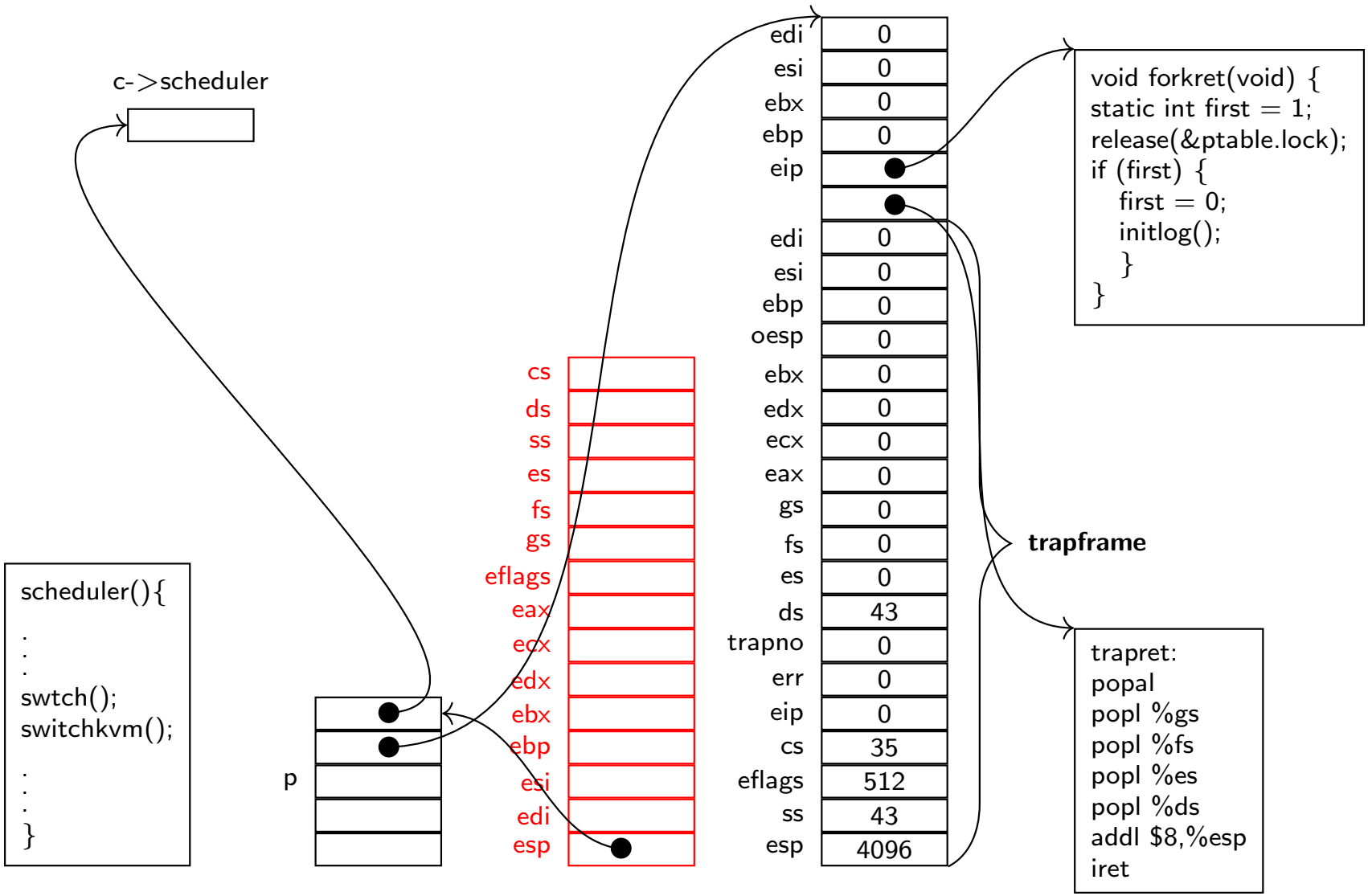

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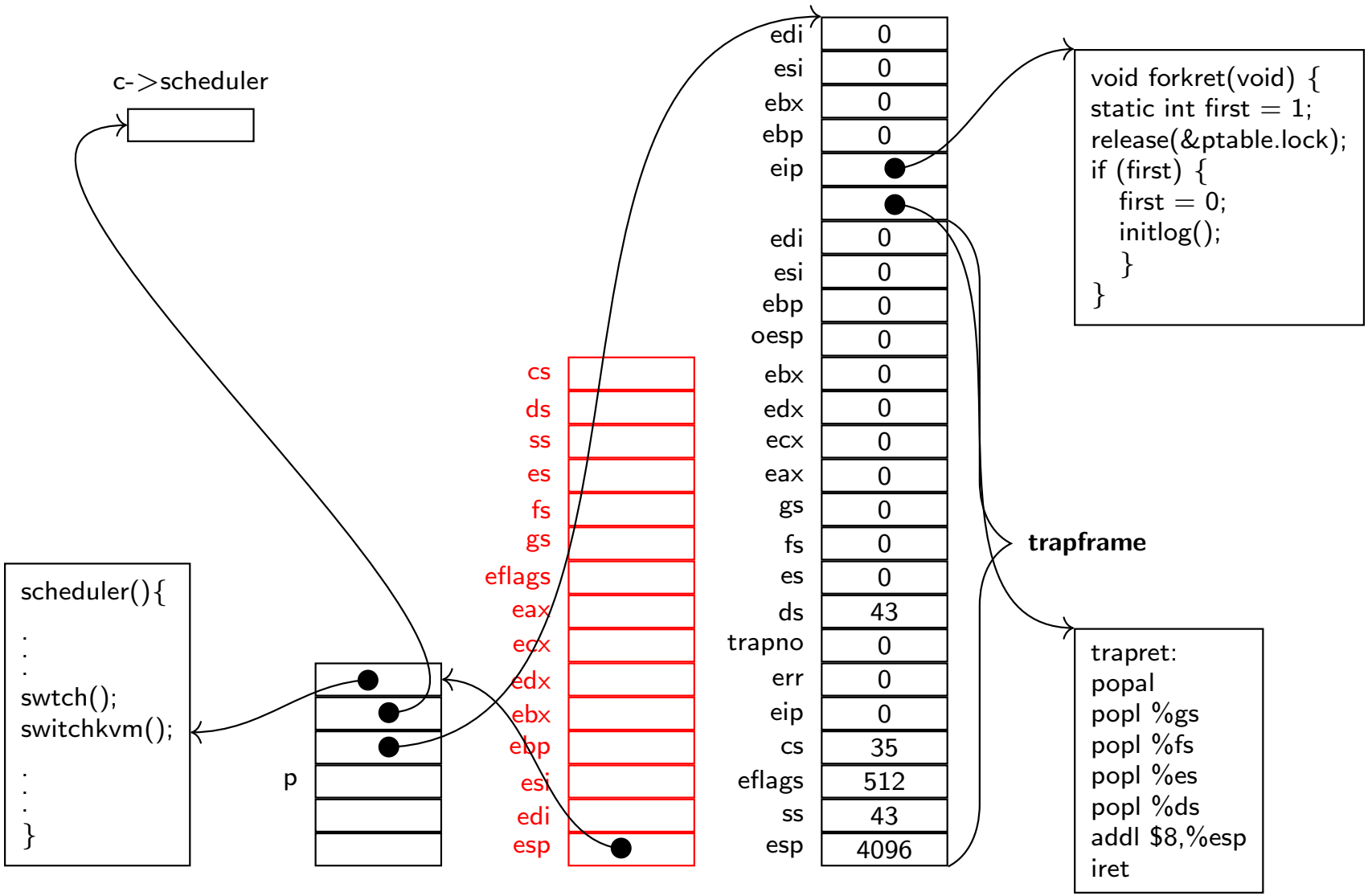

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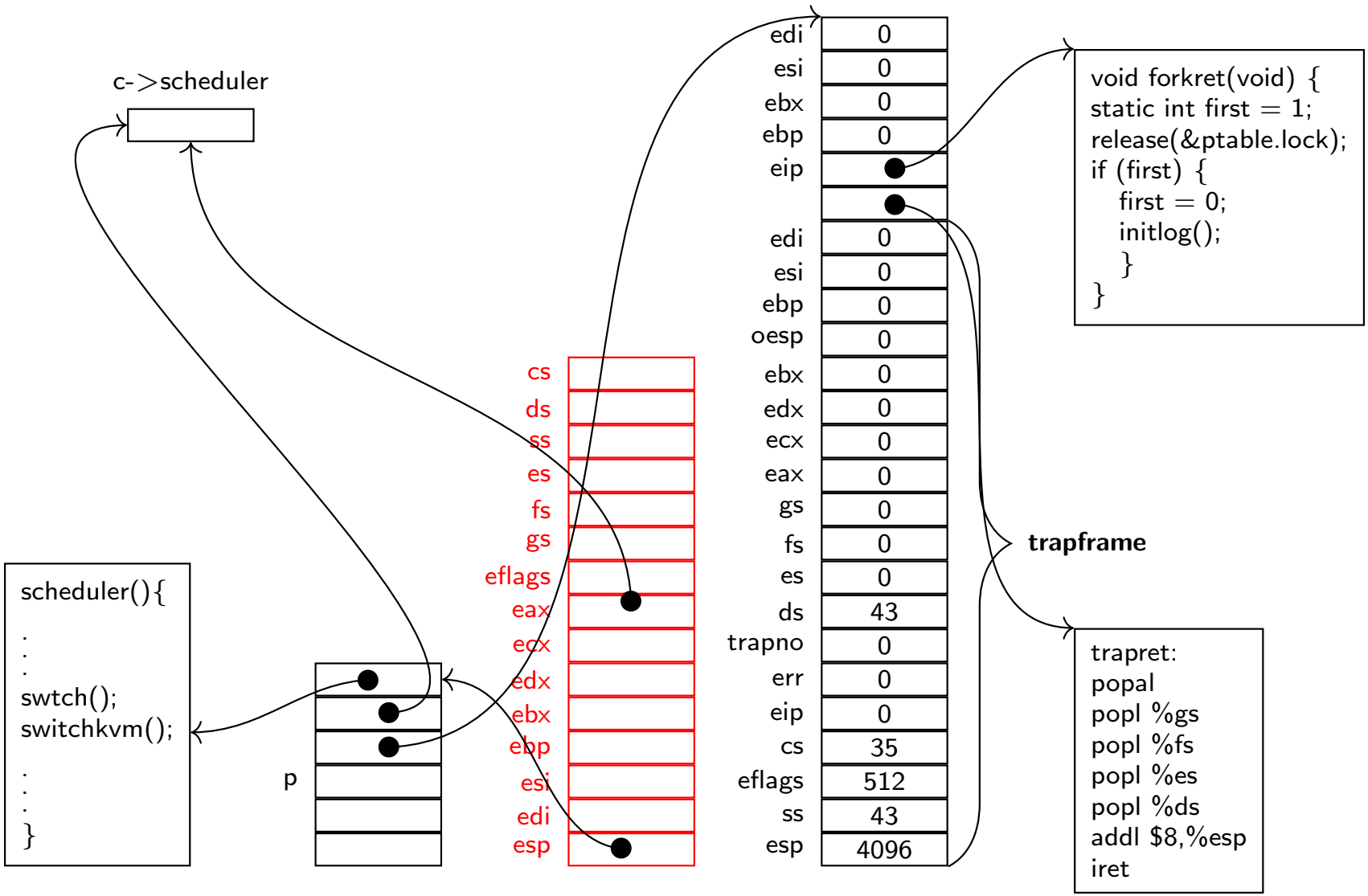

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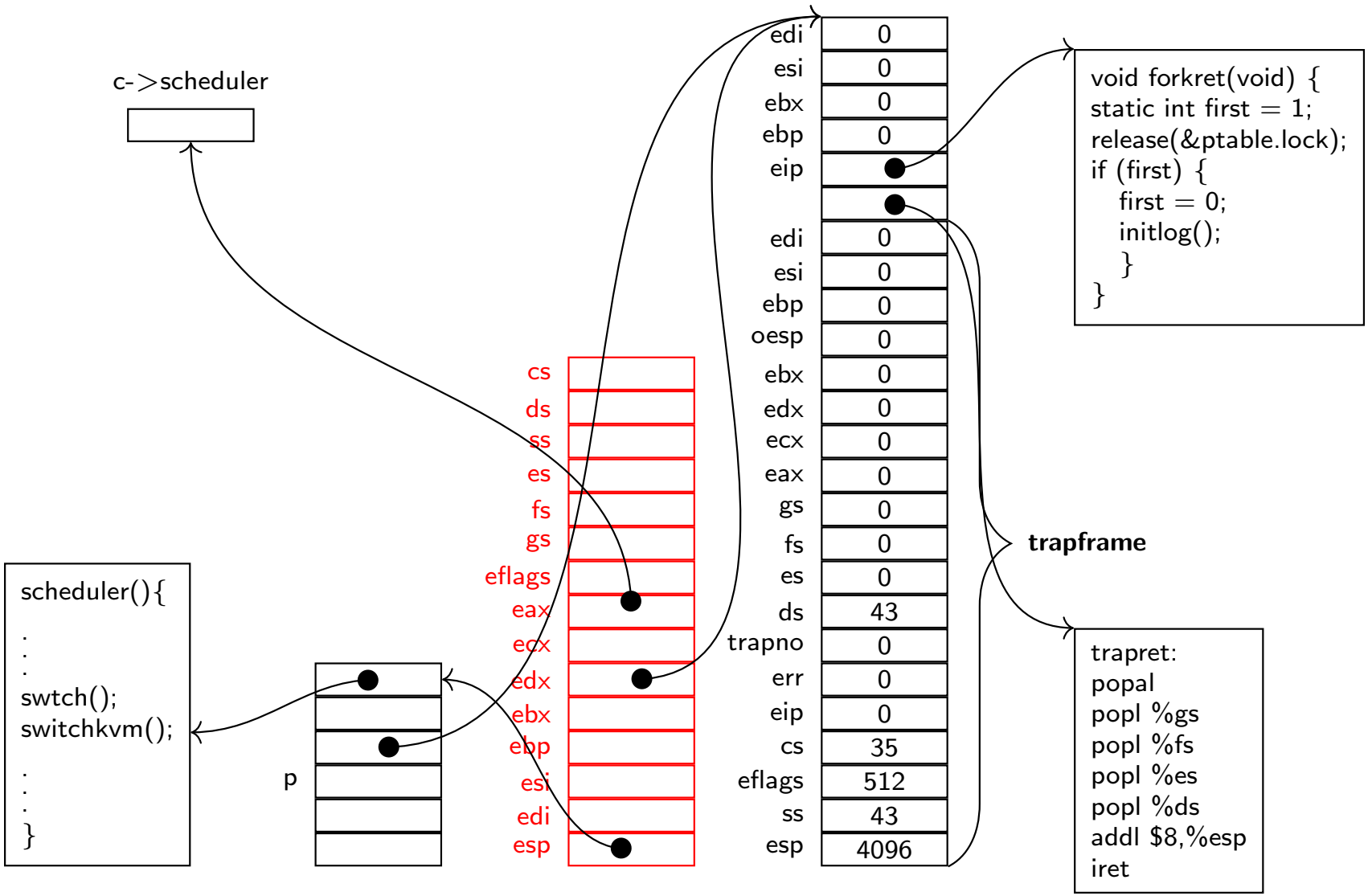

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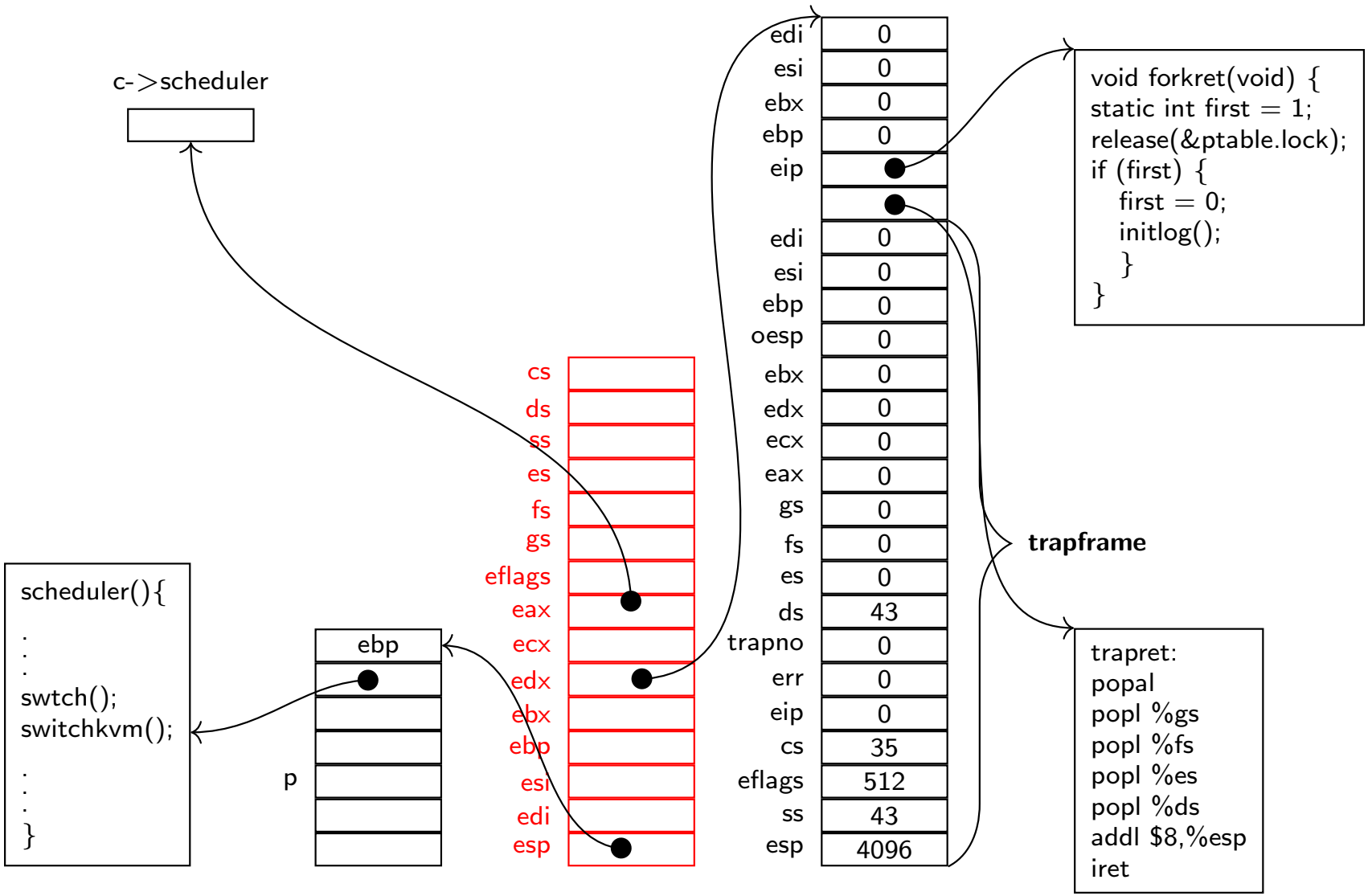

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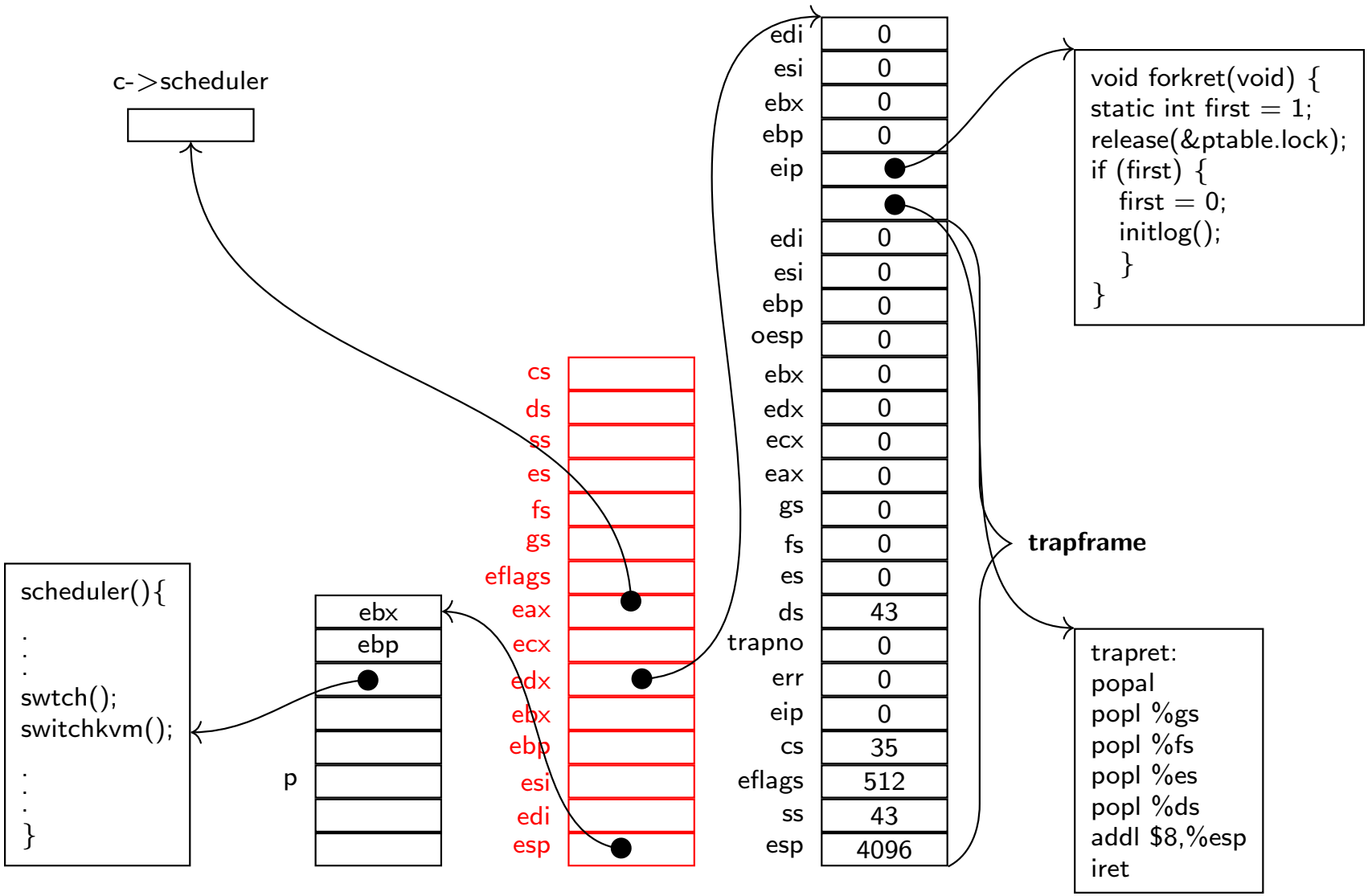

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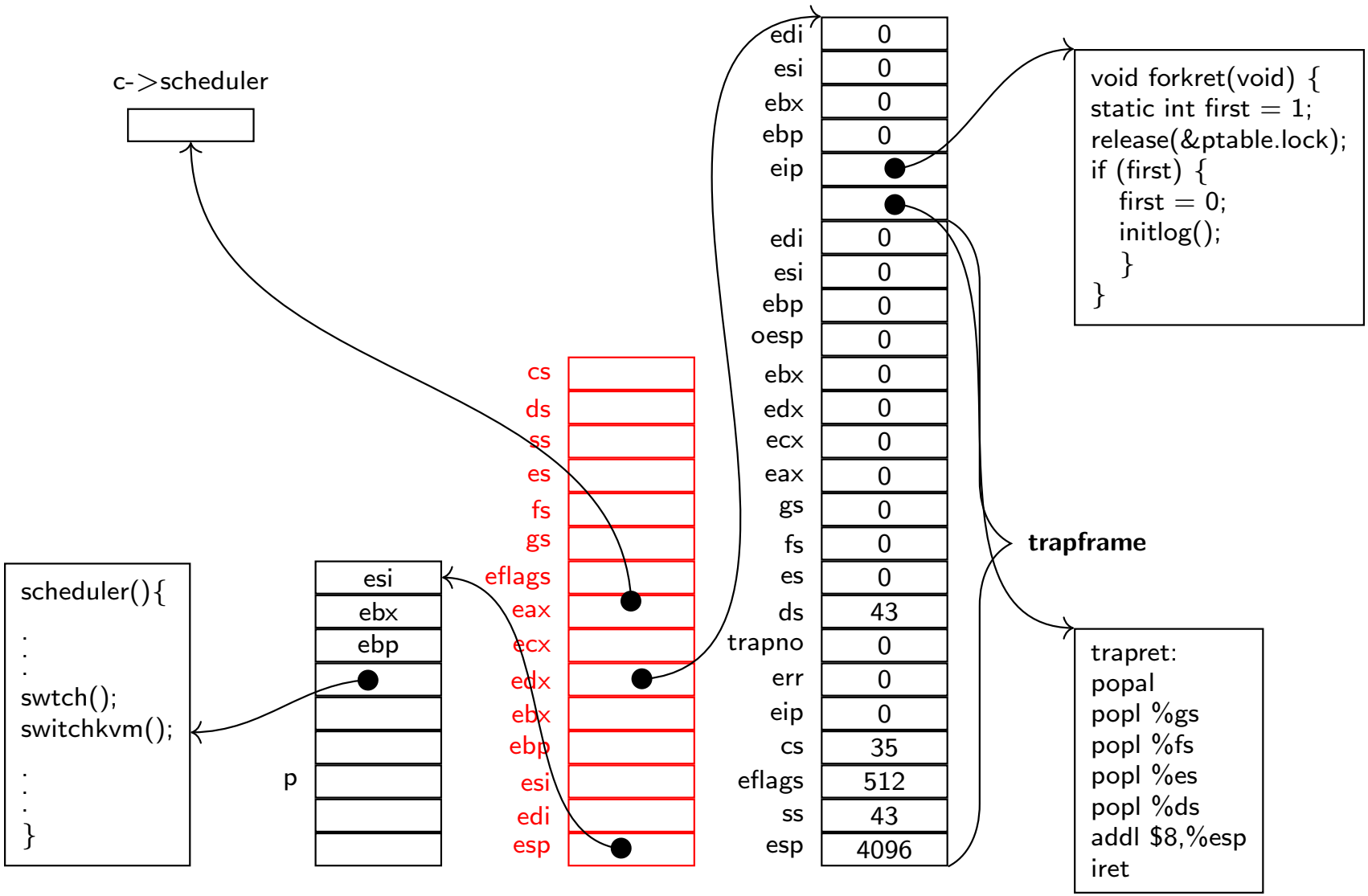

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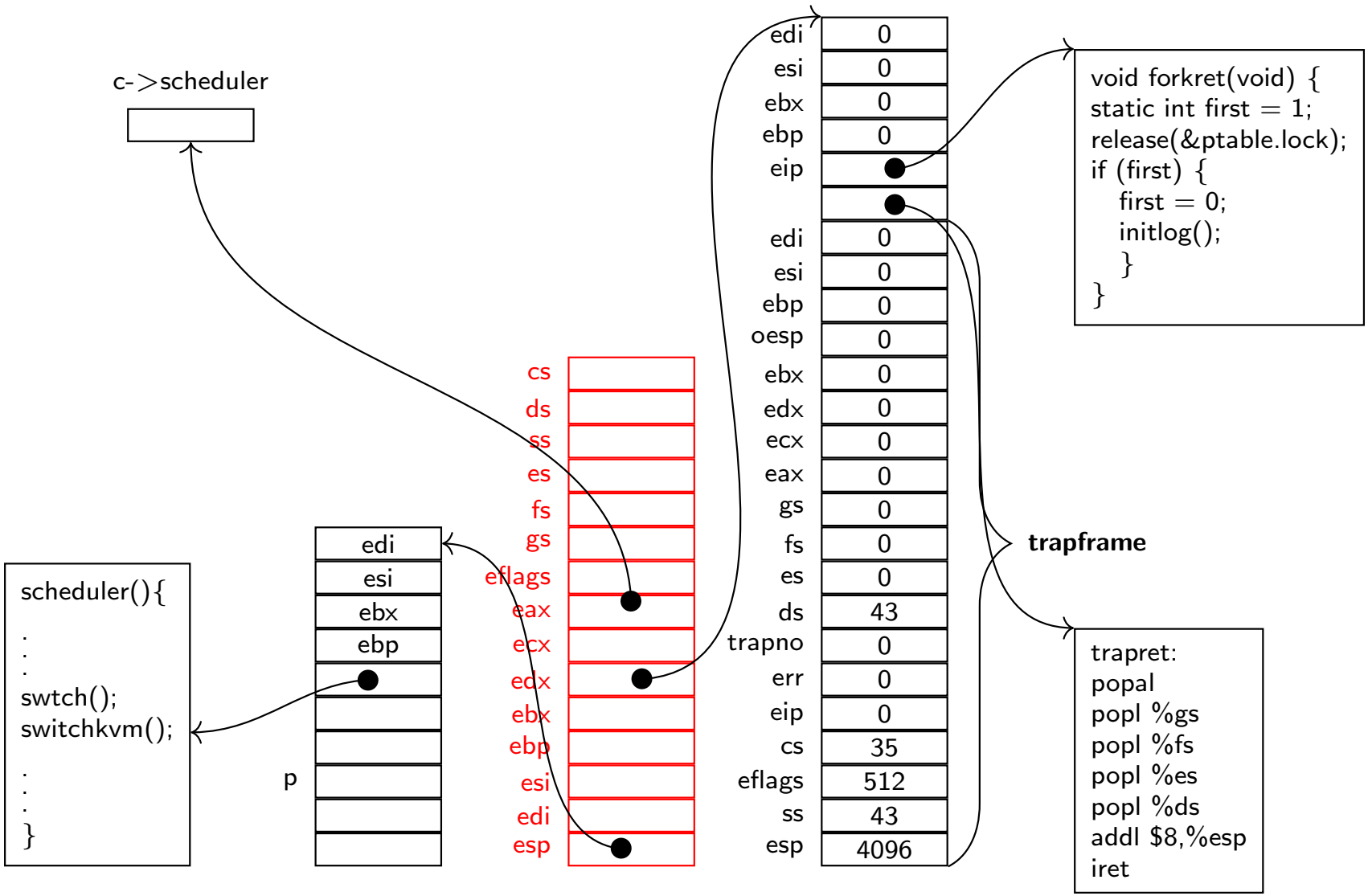

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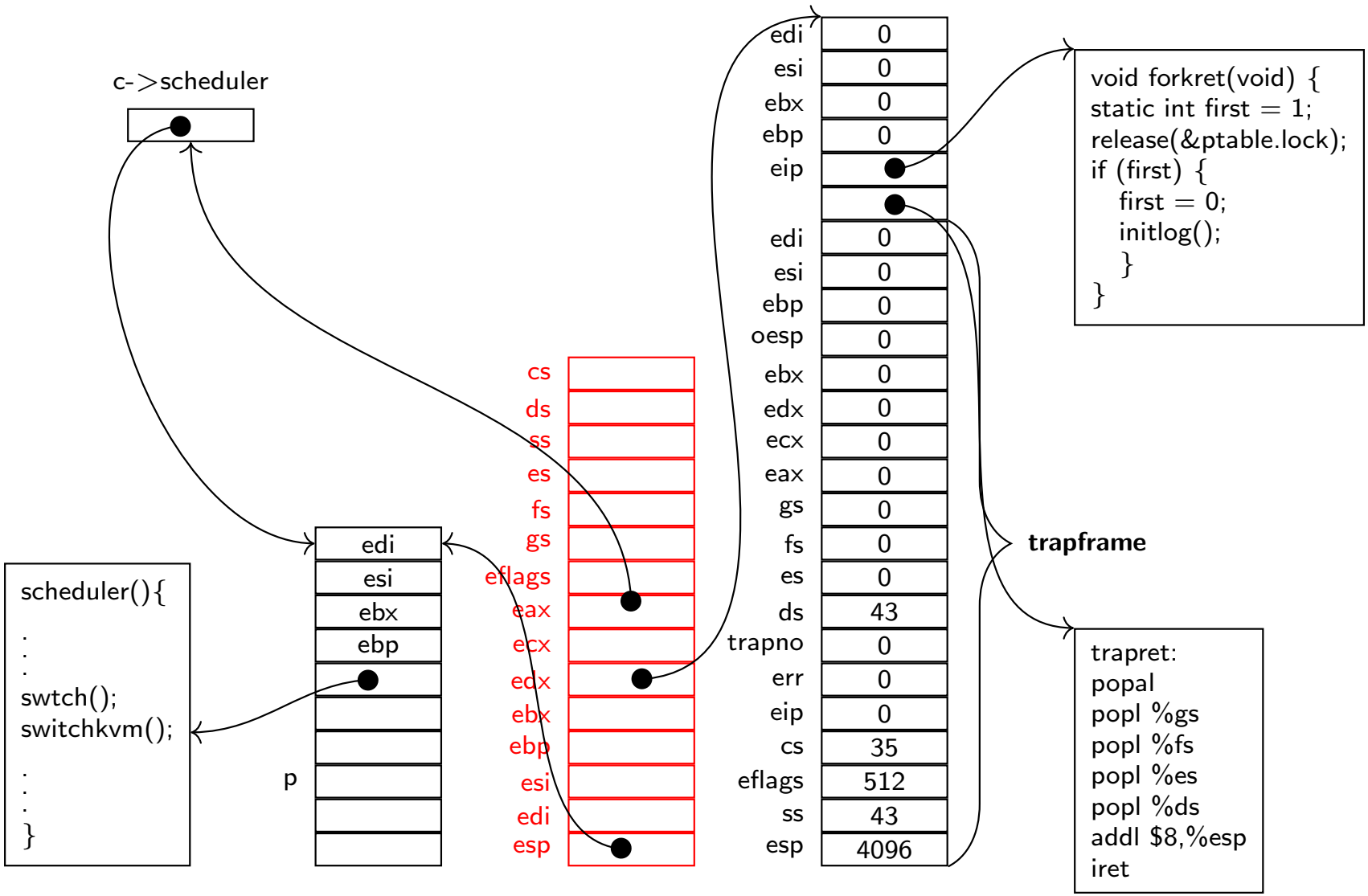

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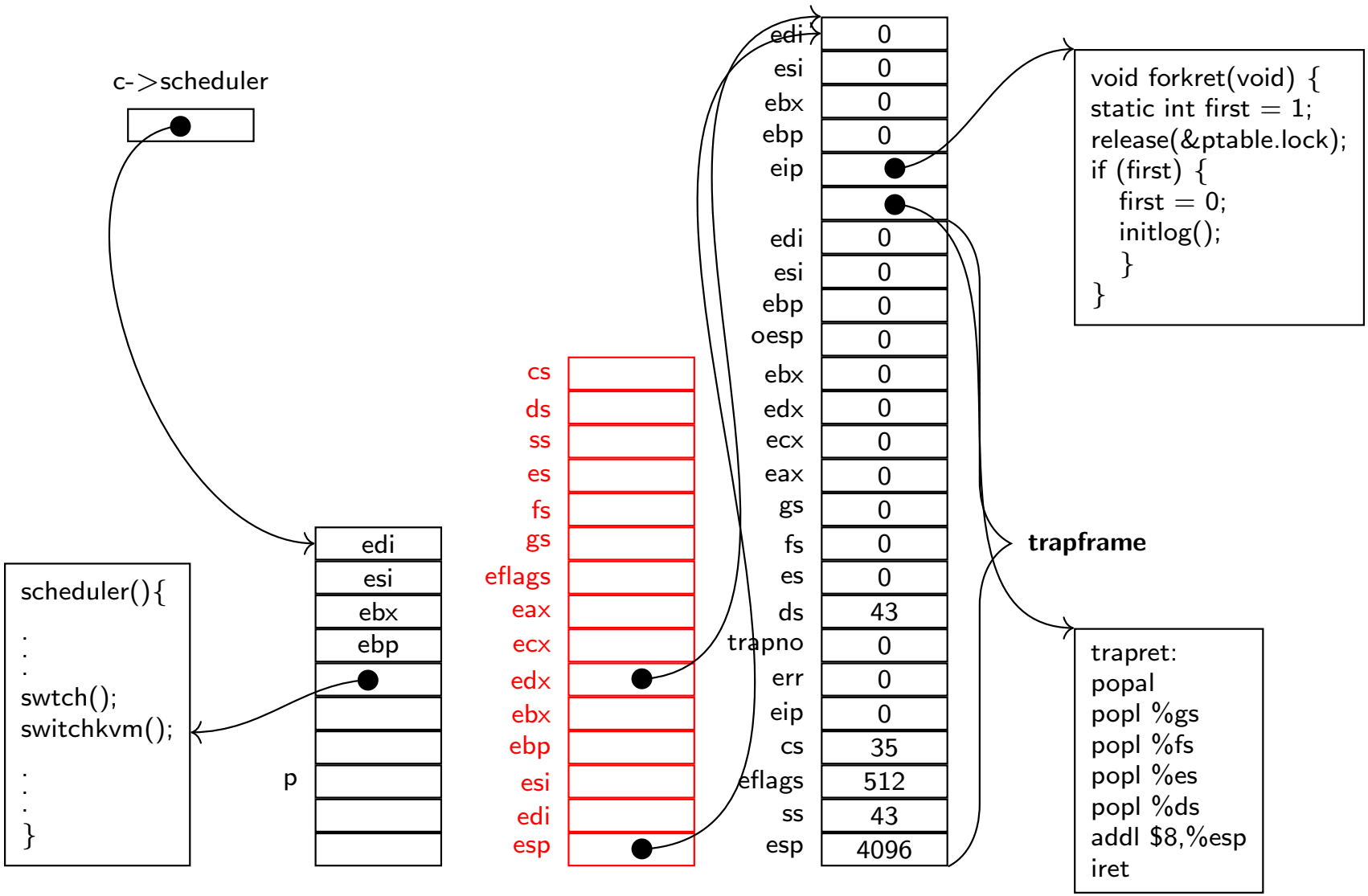

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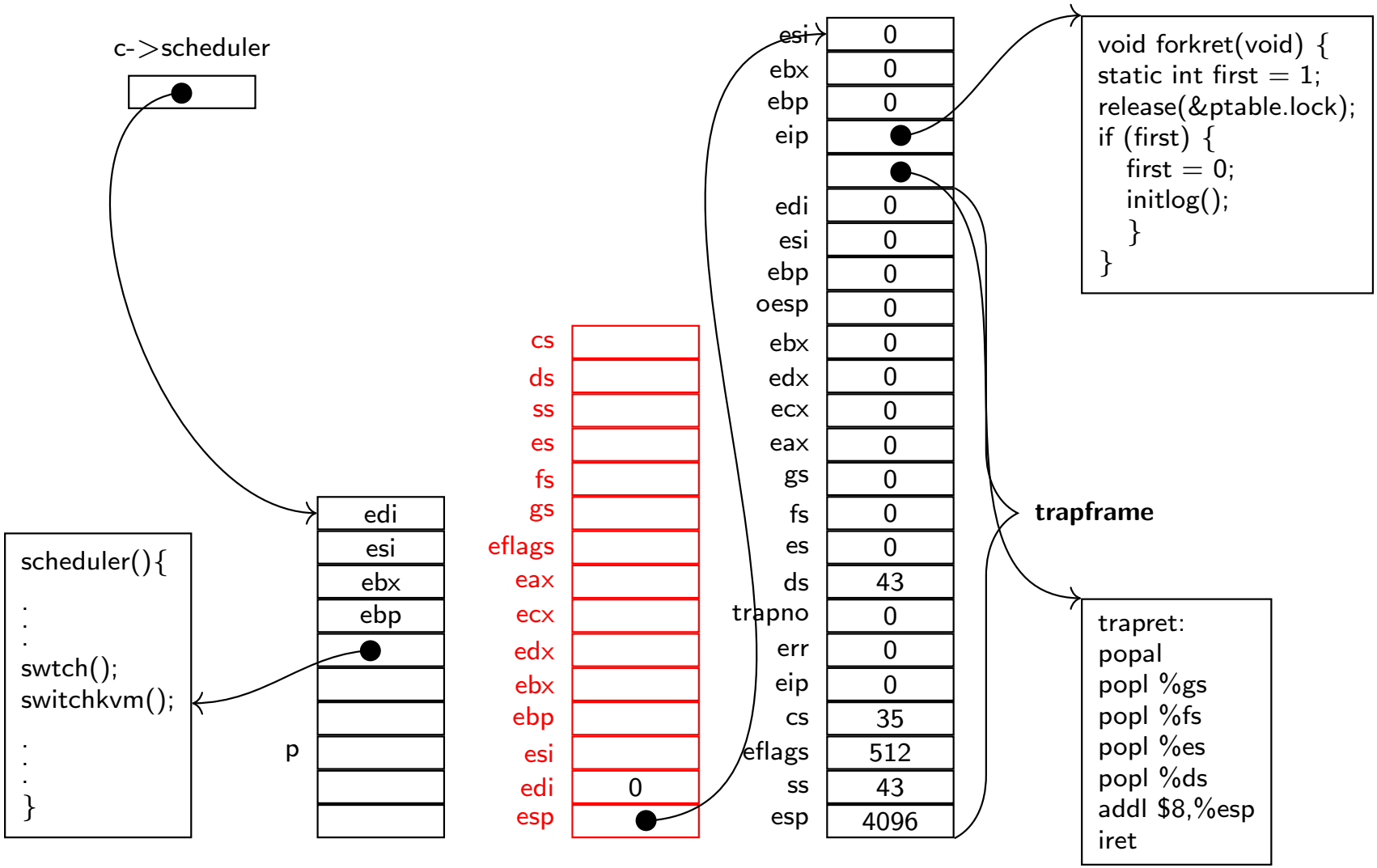

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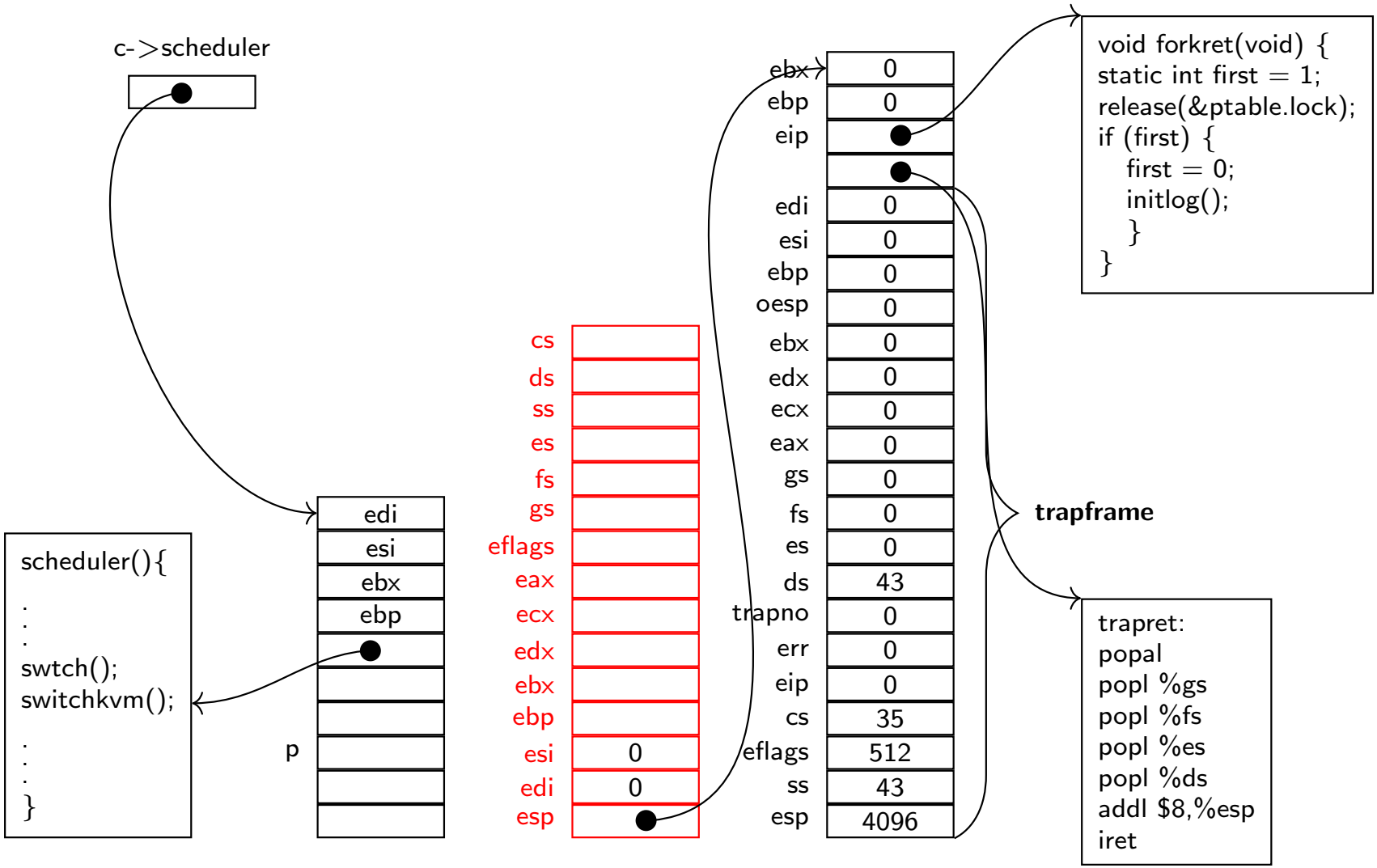

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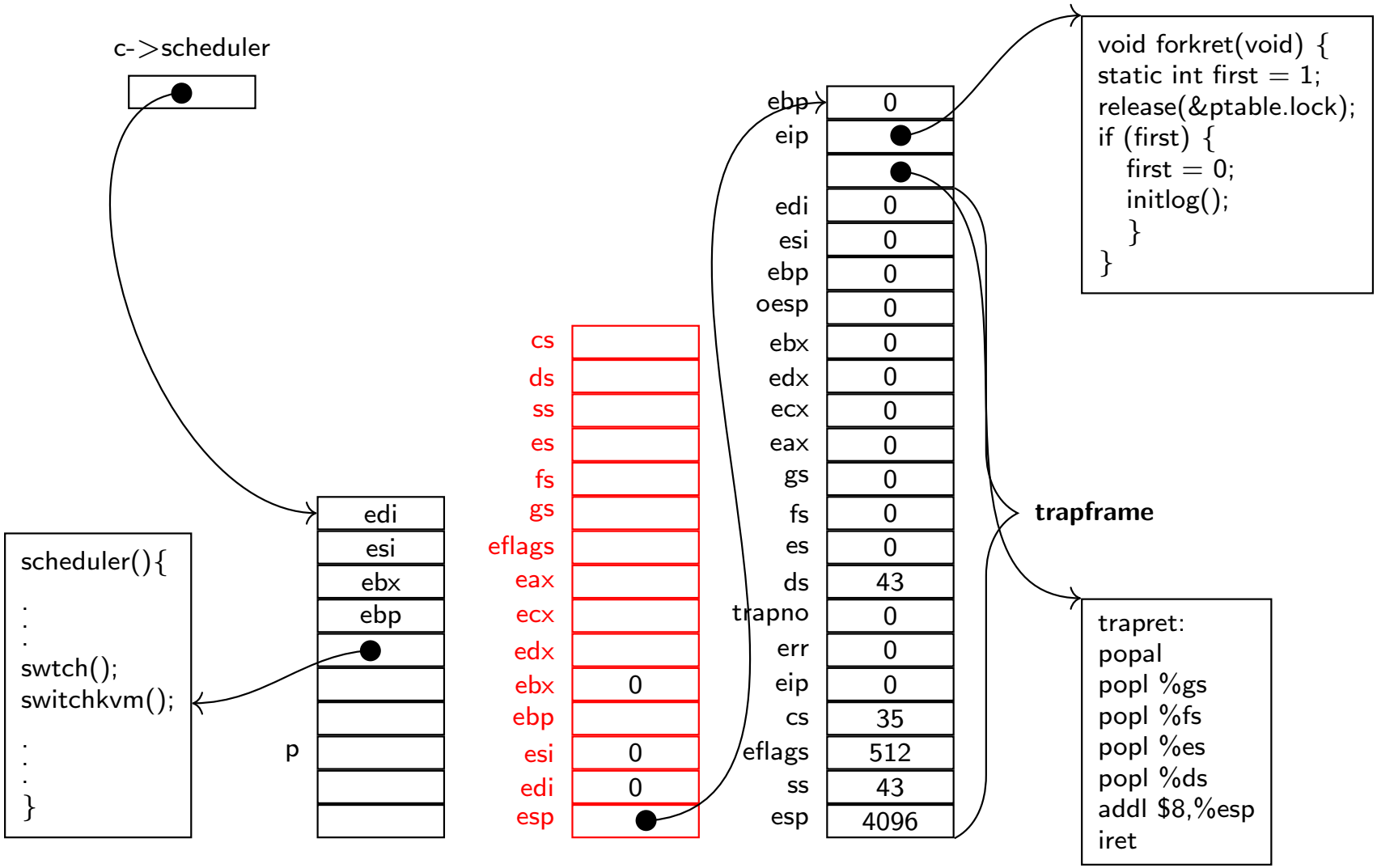

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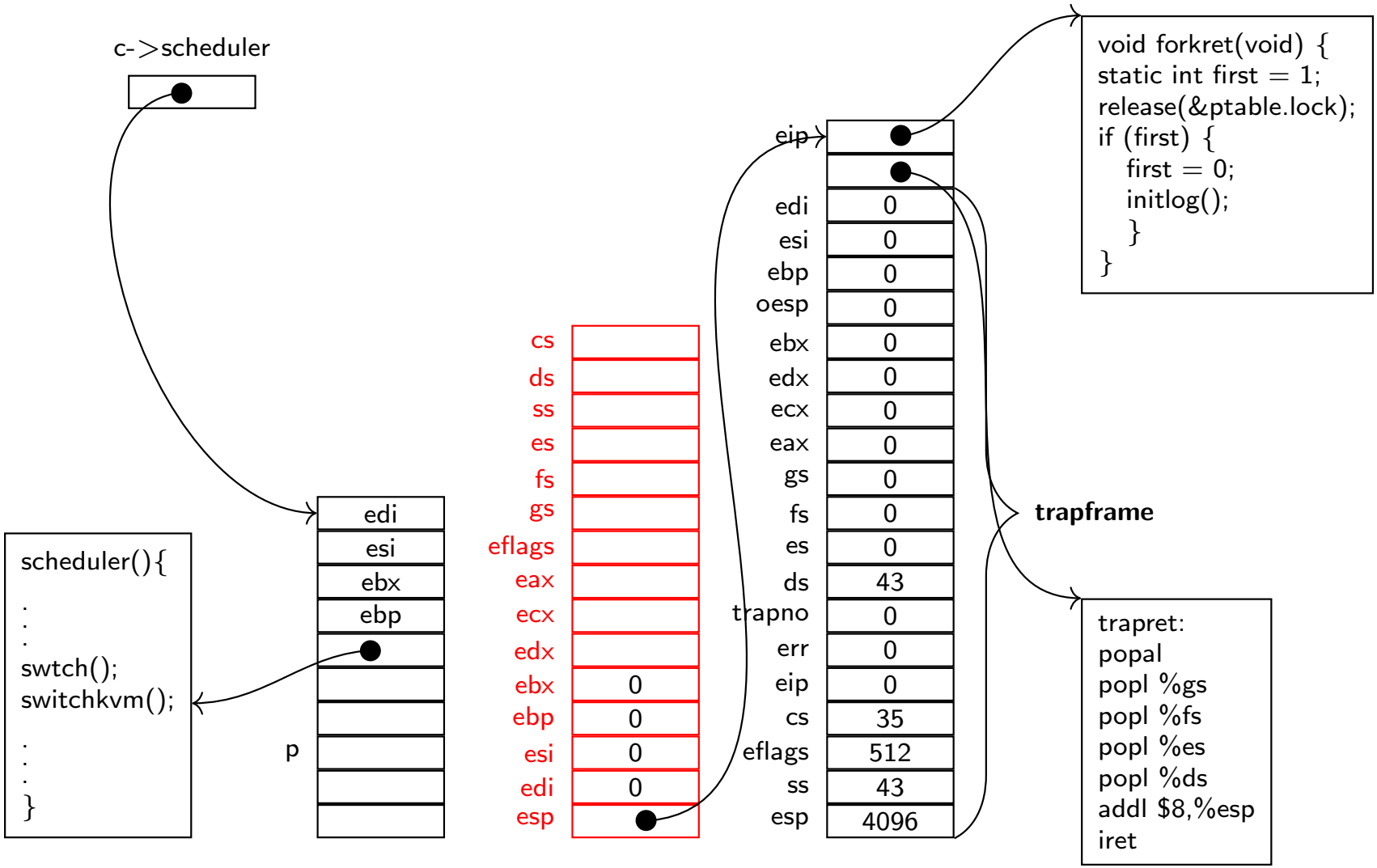

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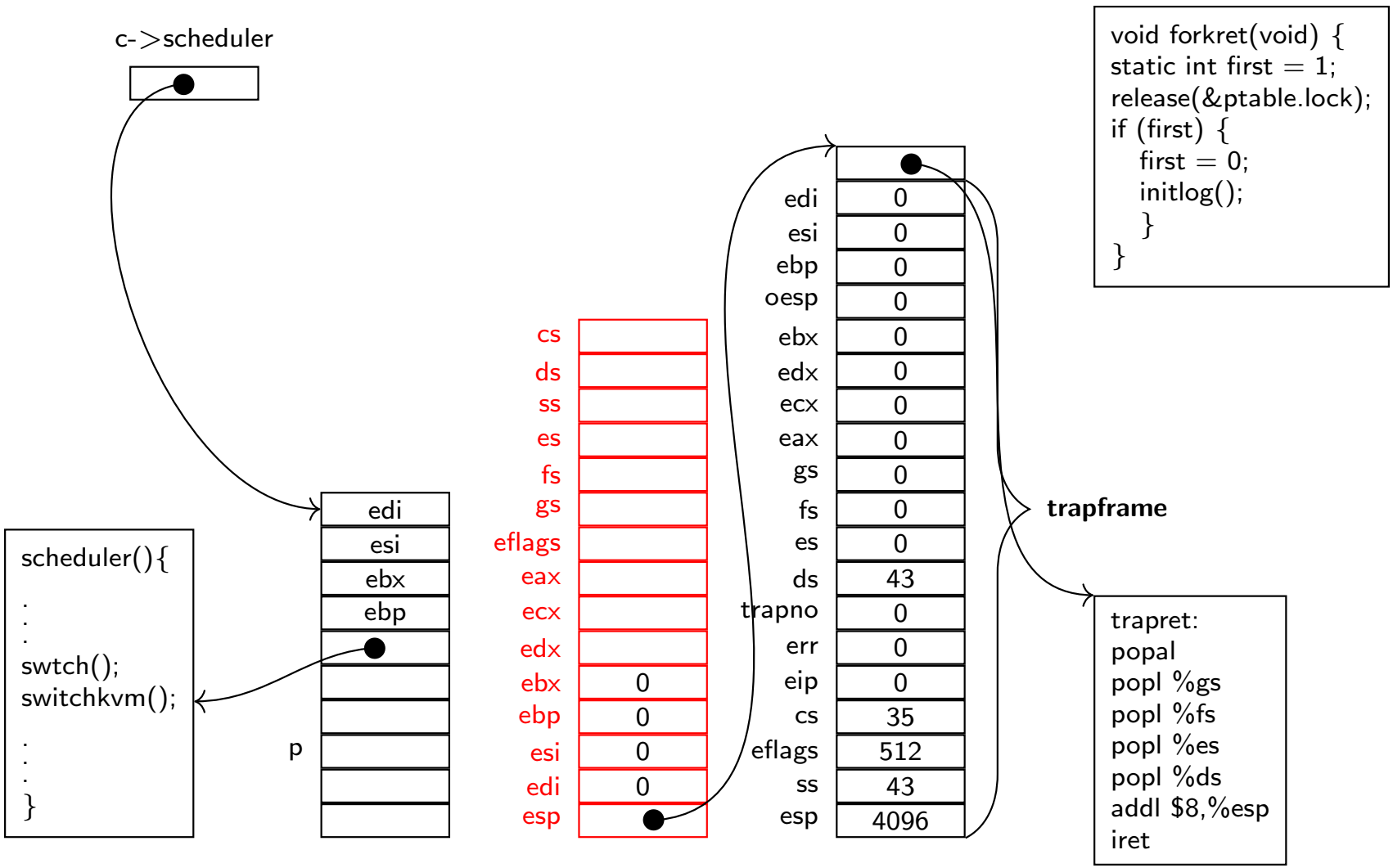

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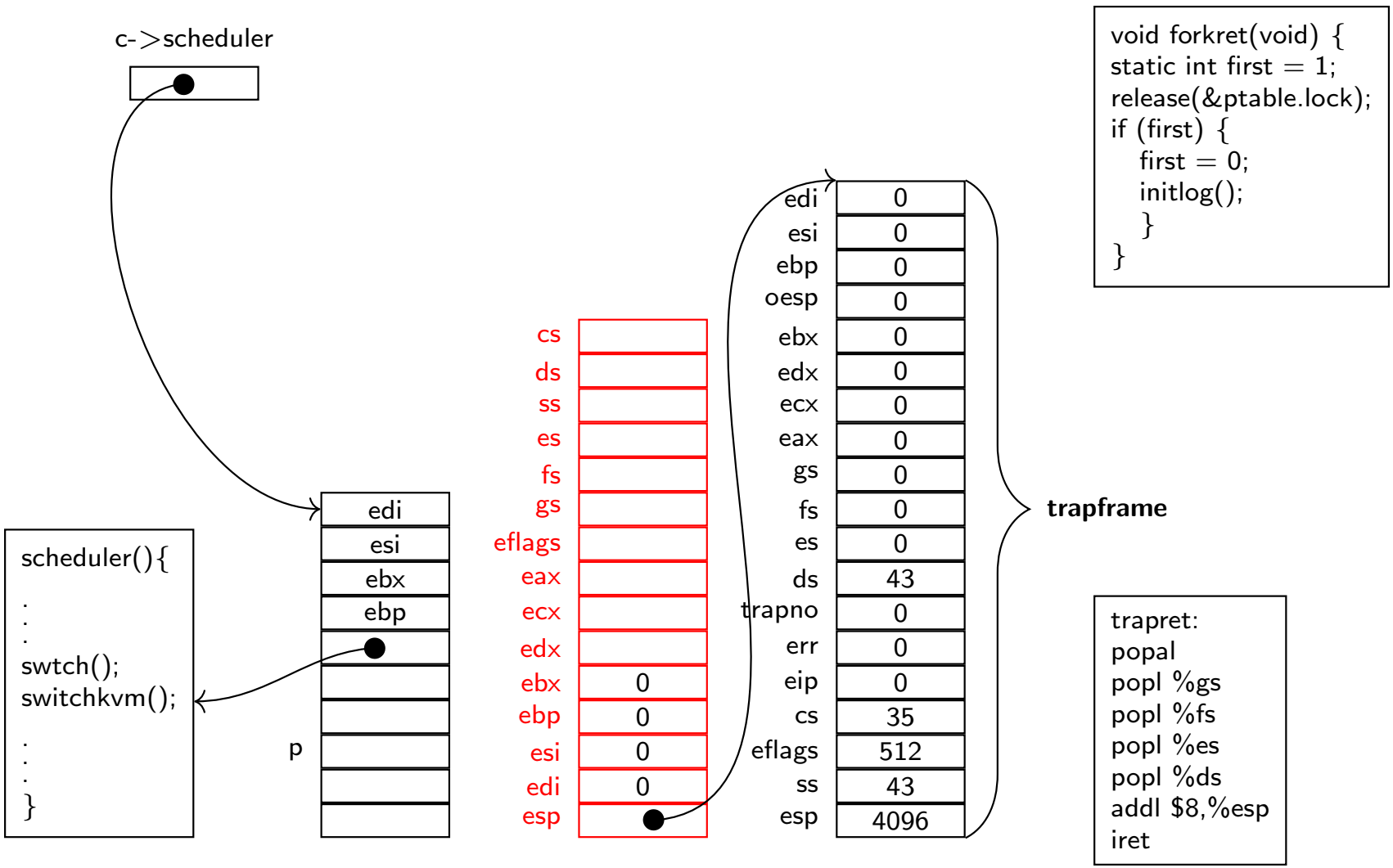

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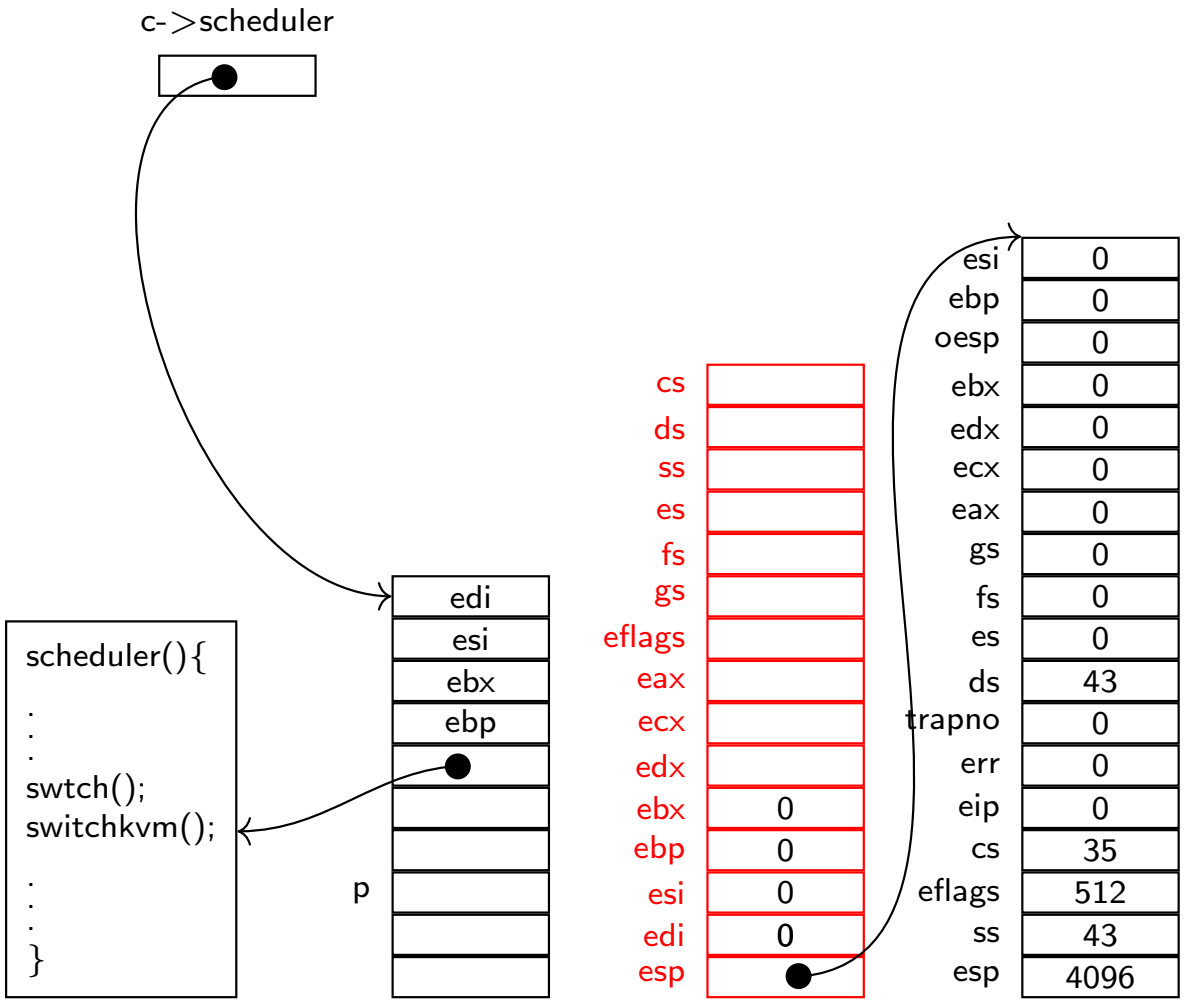


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```

```
trapret:
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popl %gs
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popl %ds
addl $8,%esp
iret
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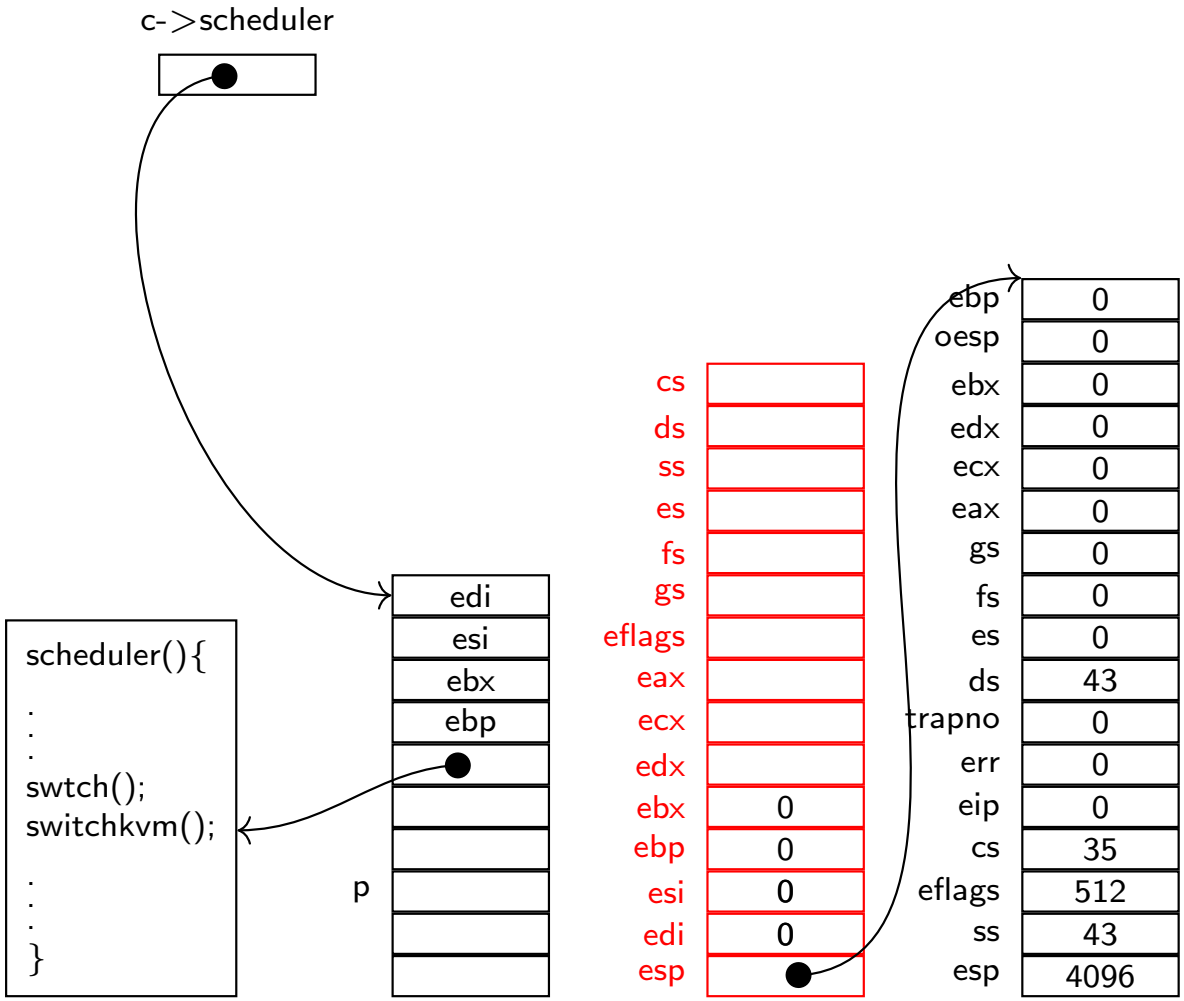


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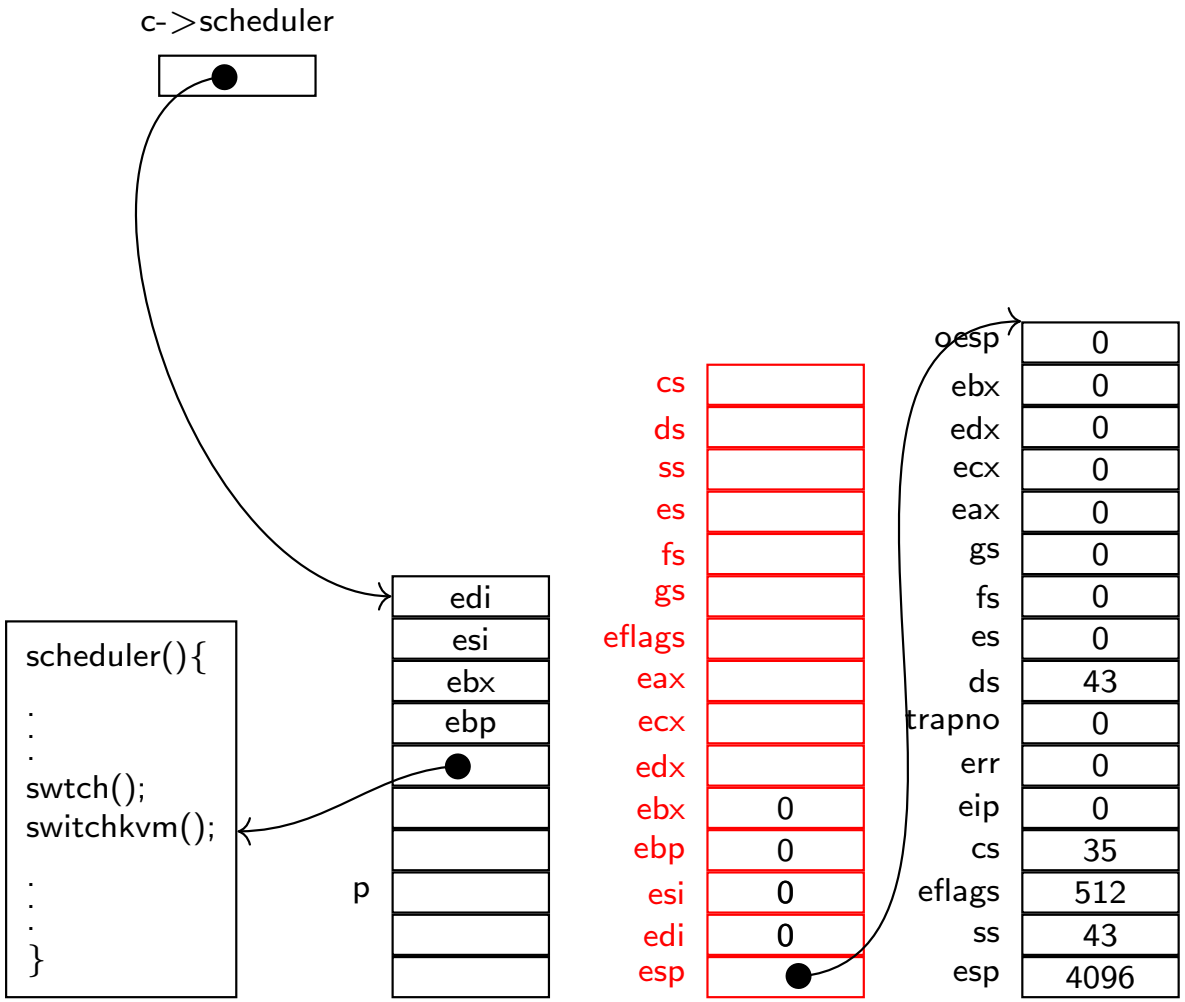


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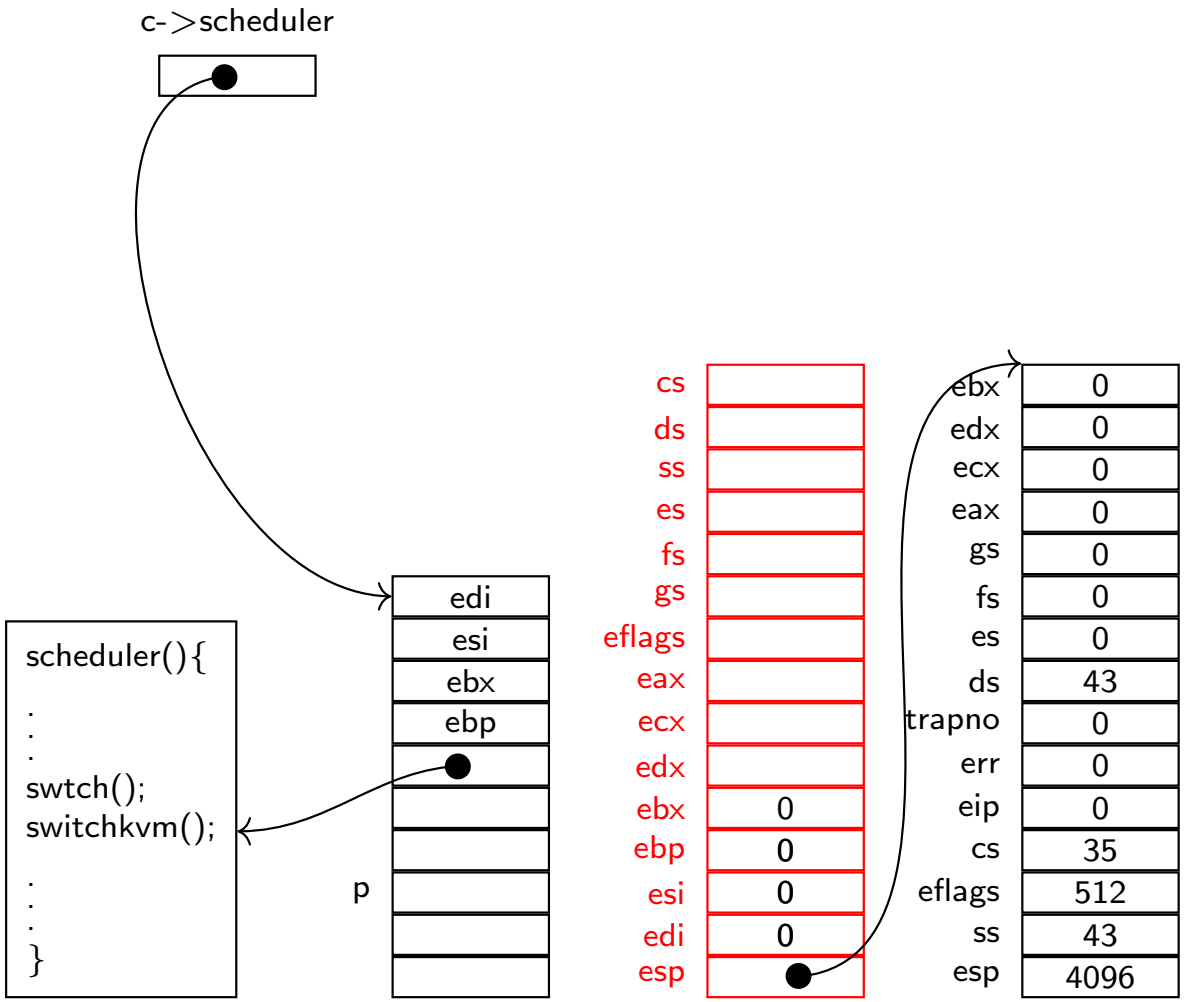


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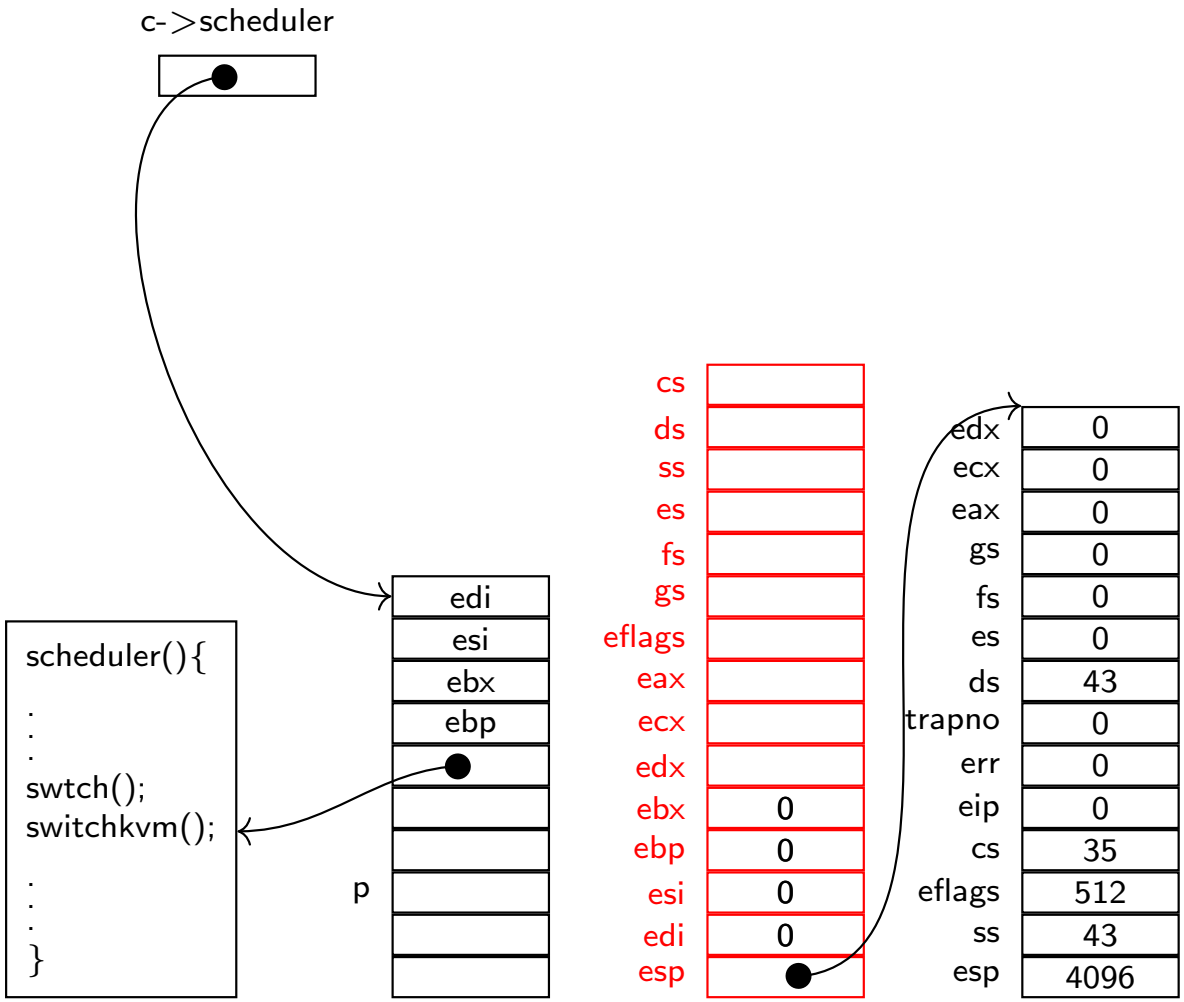


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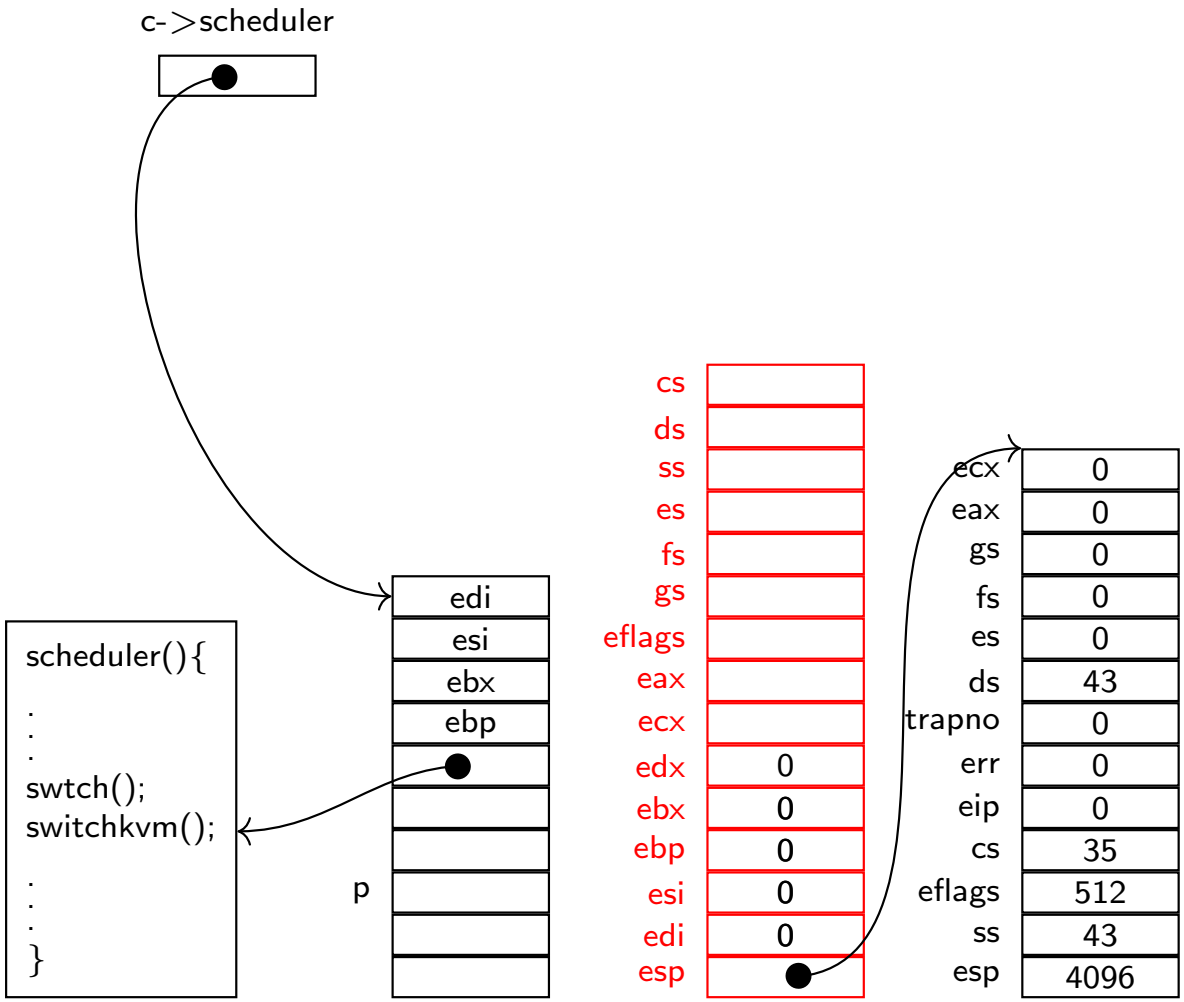


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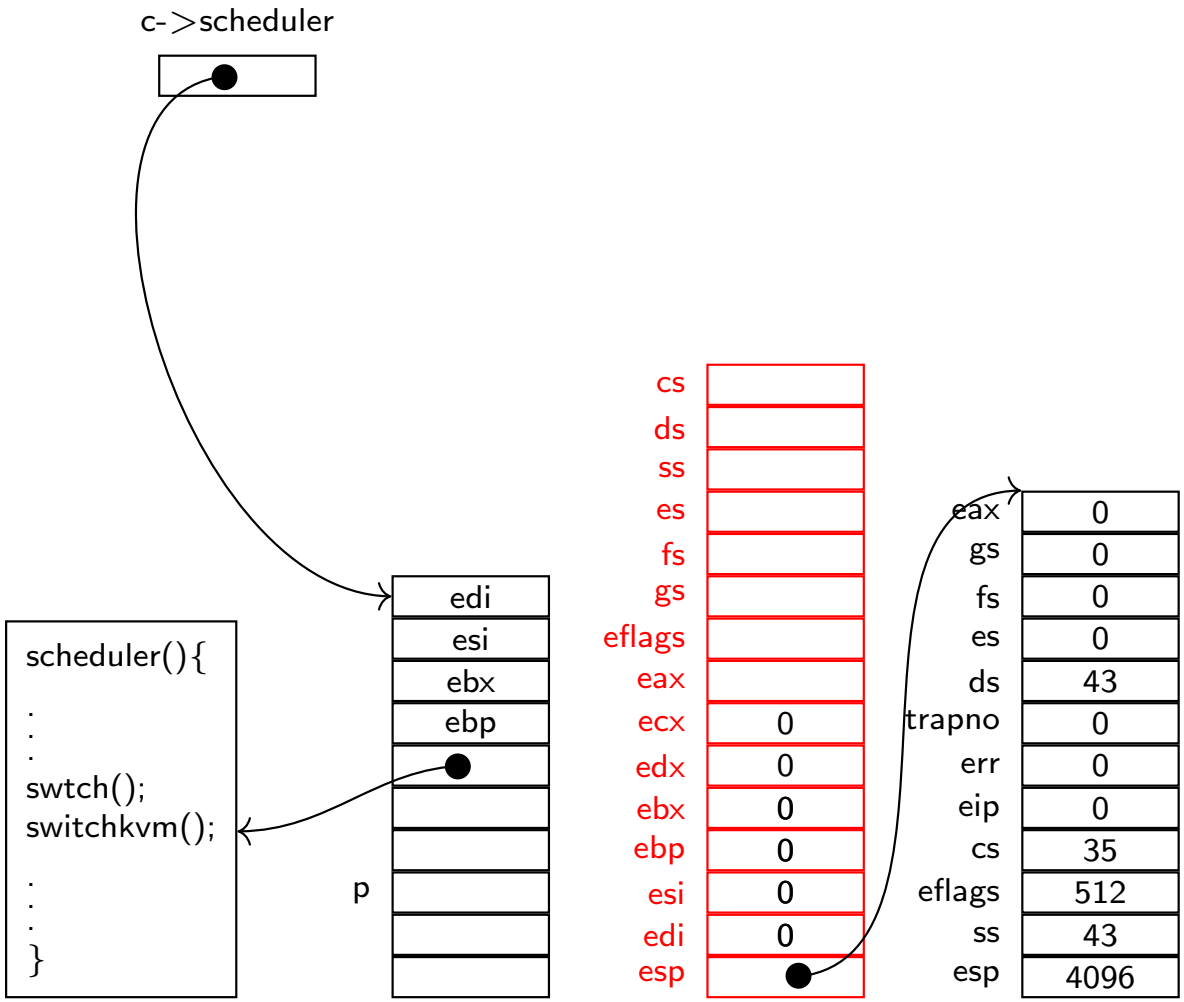


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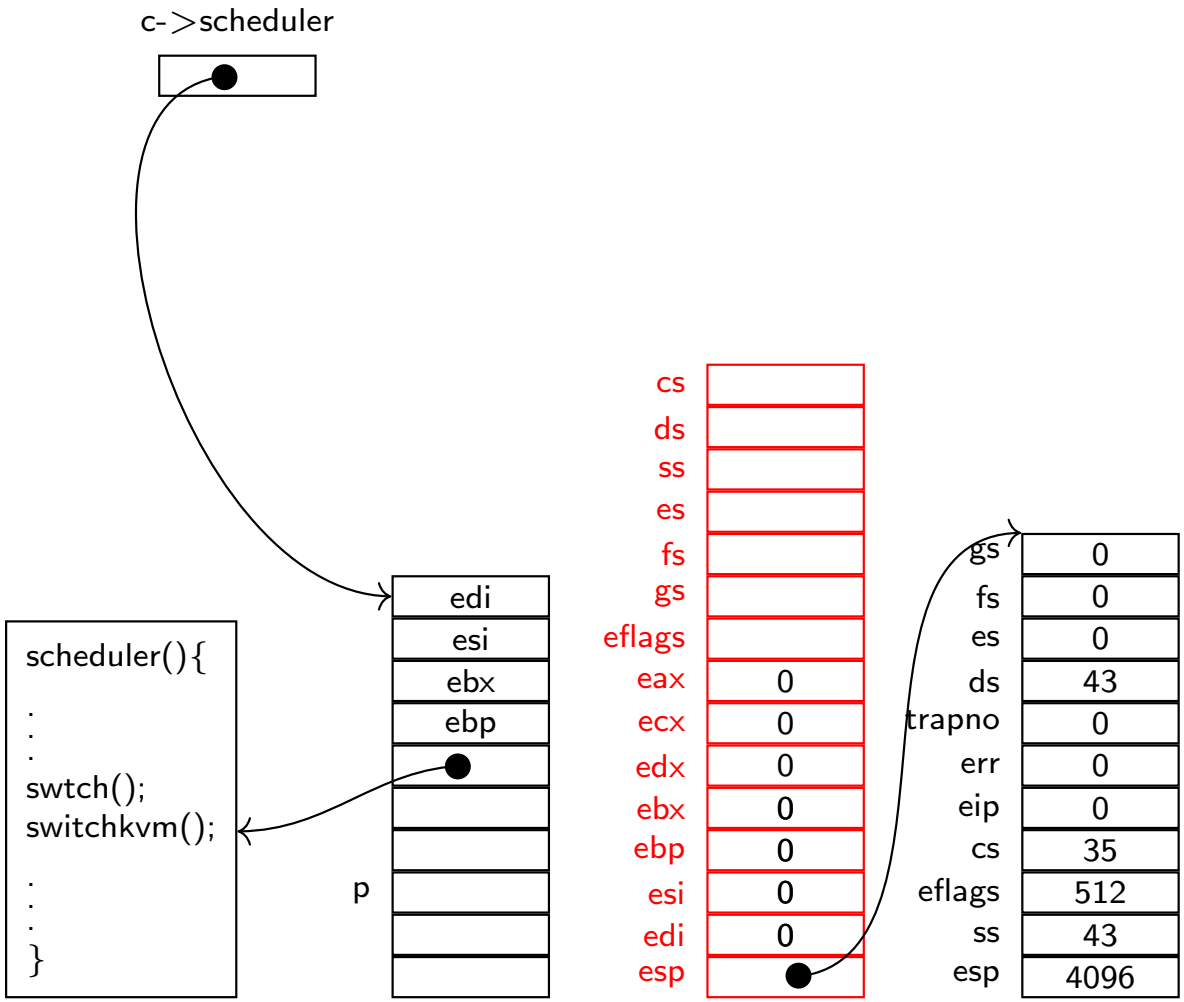


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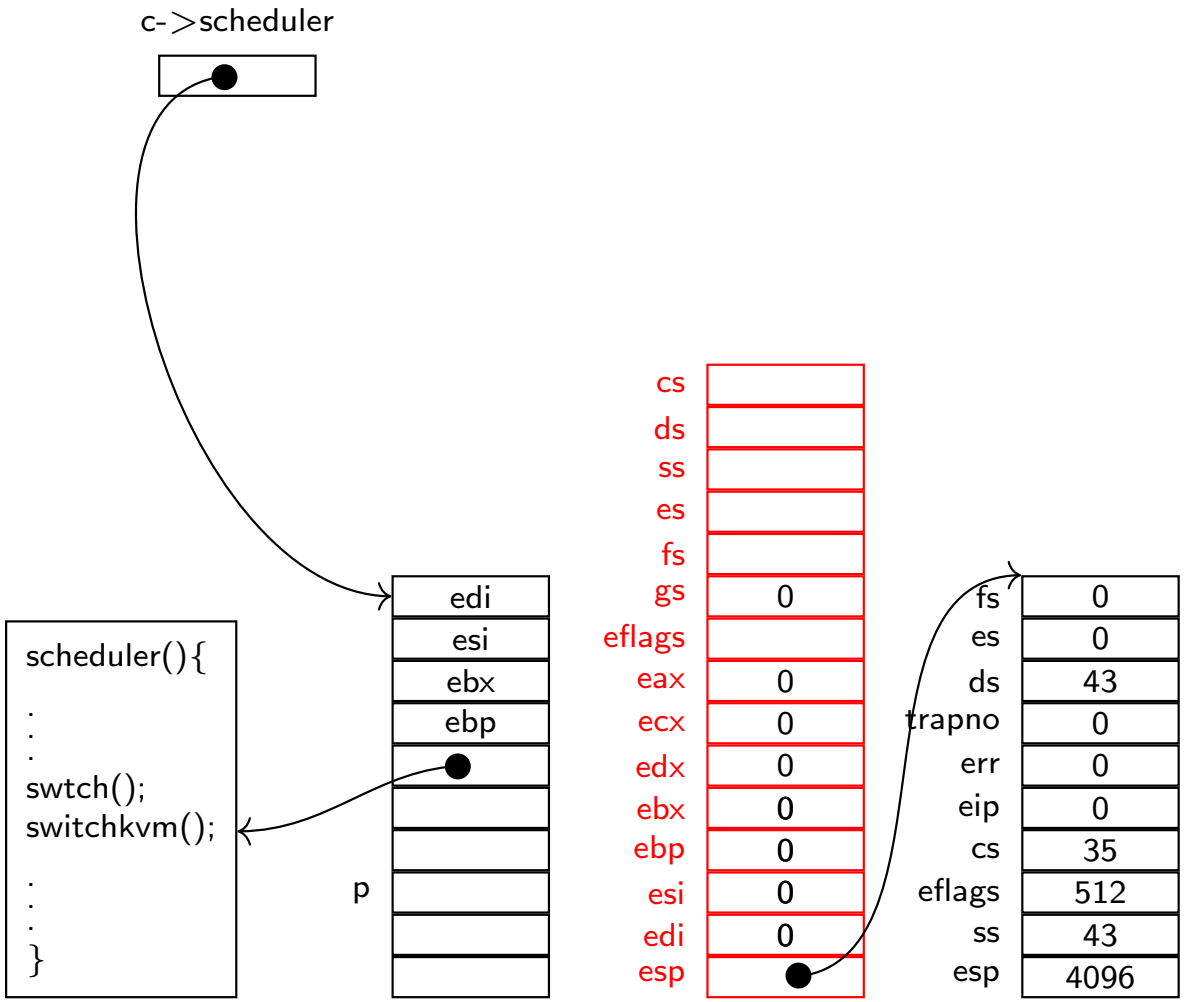


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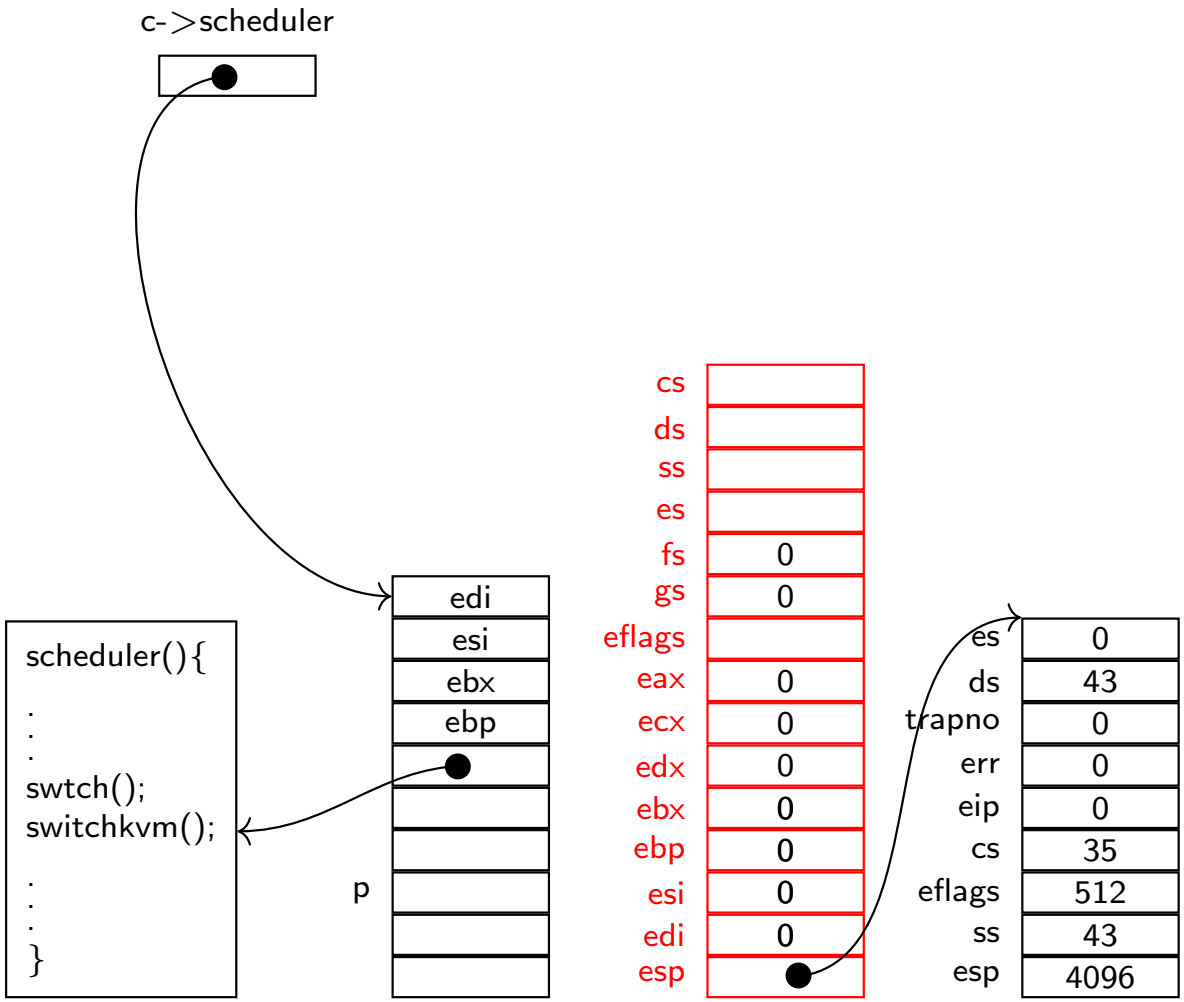


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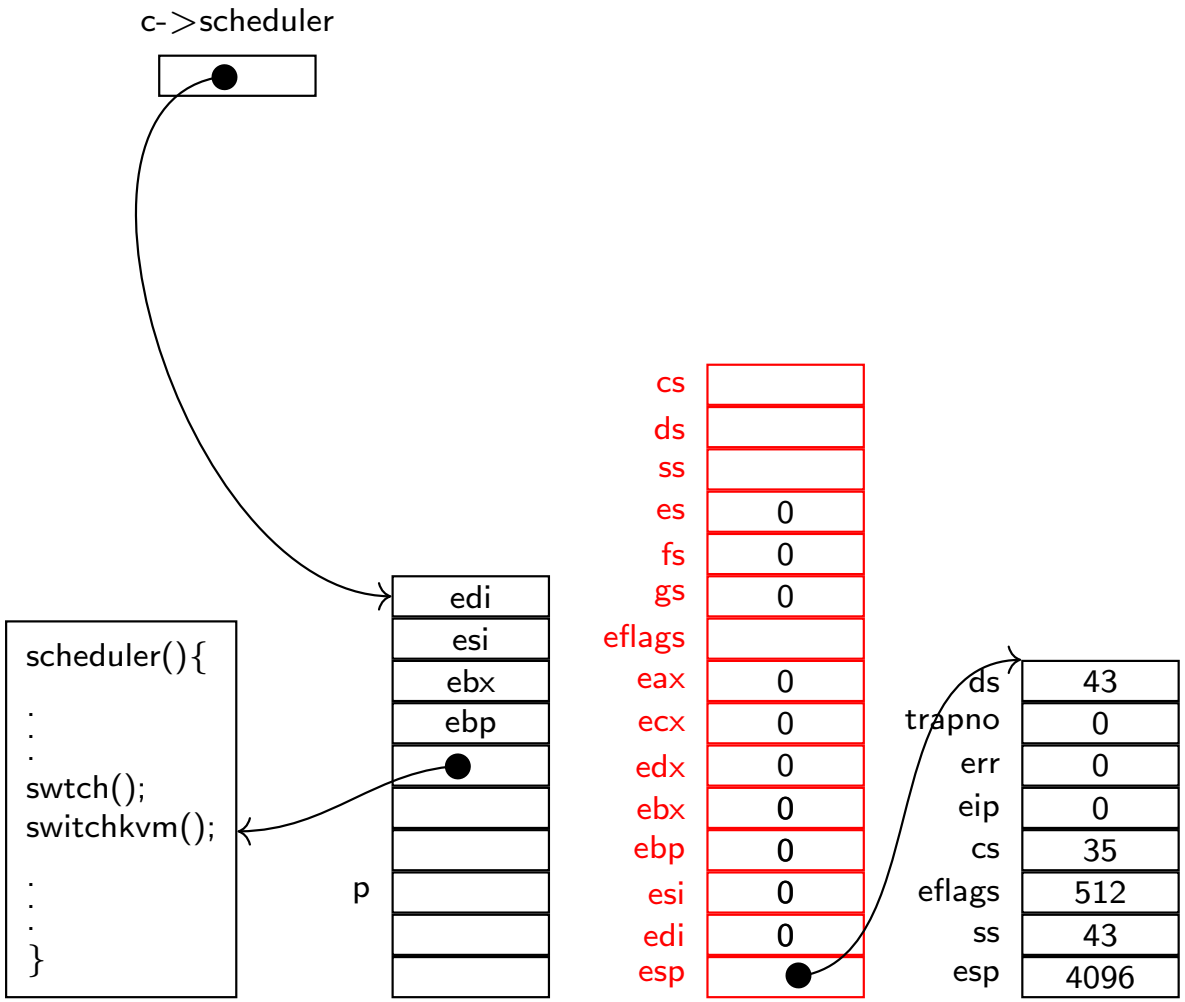


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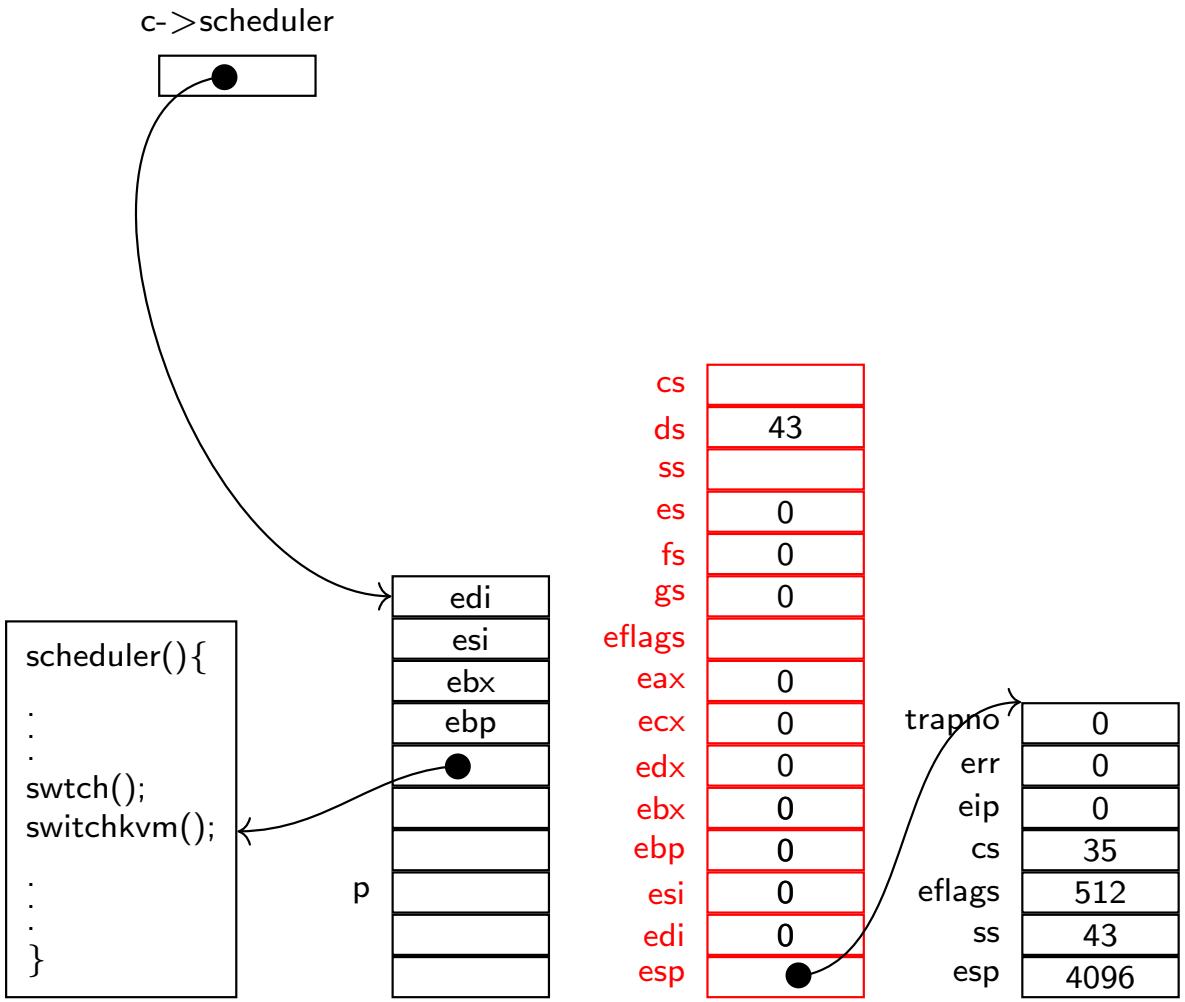


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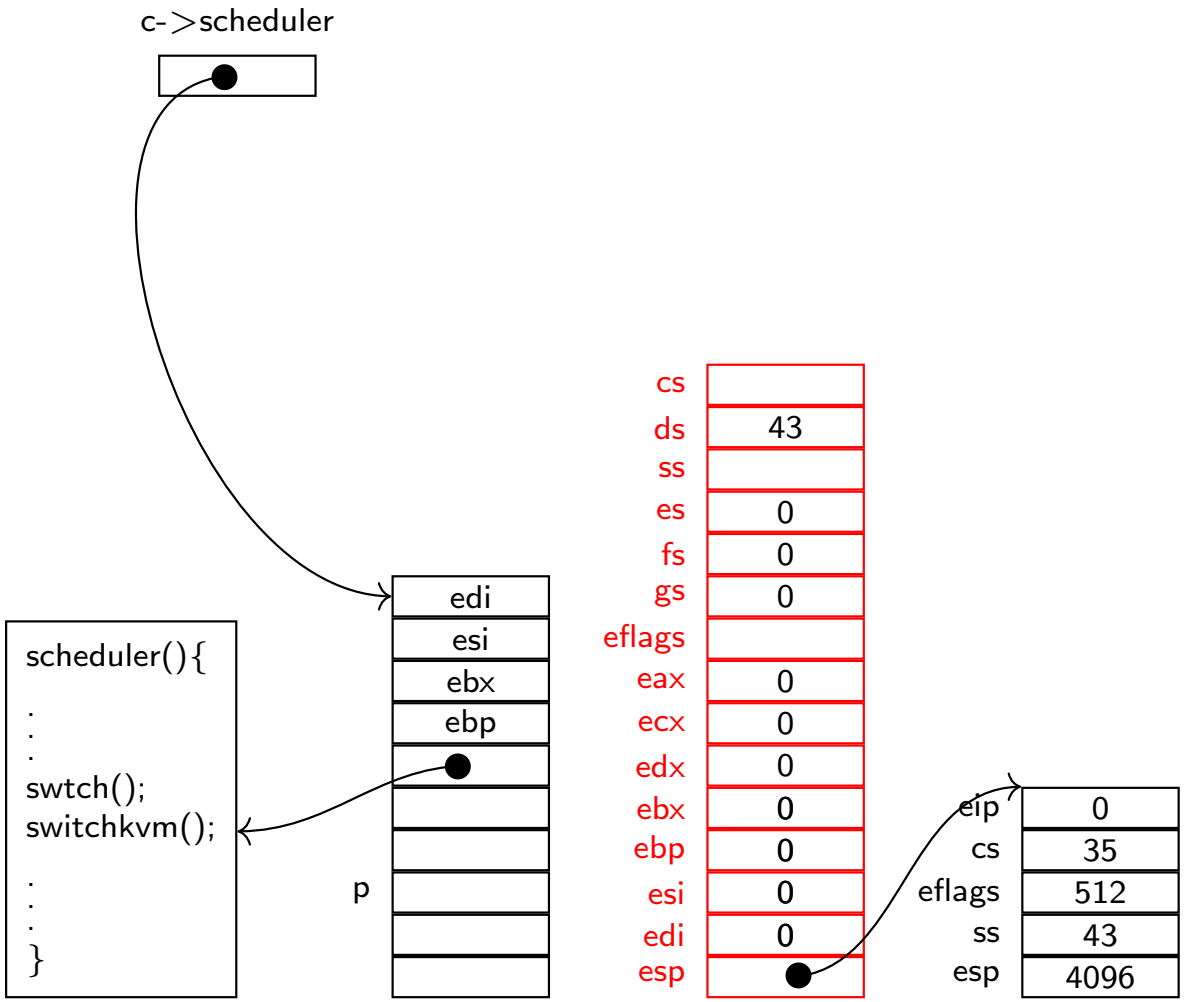


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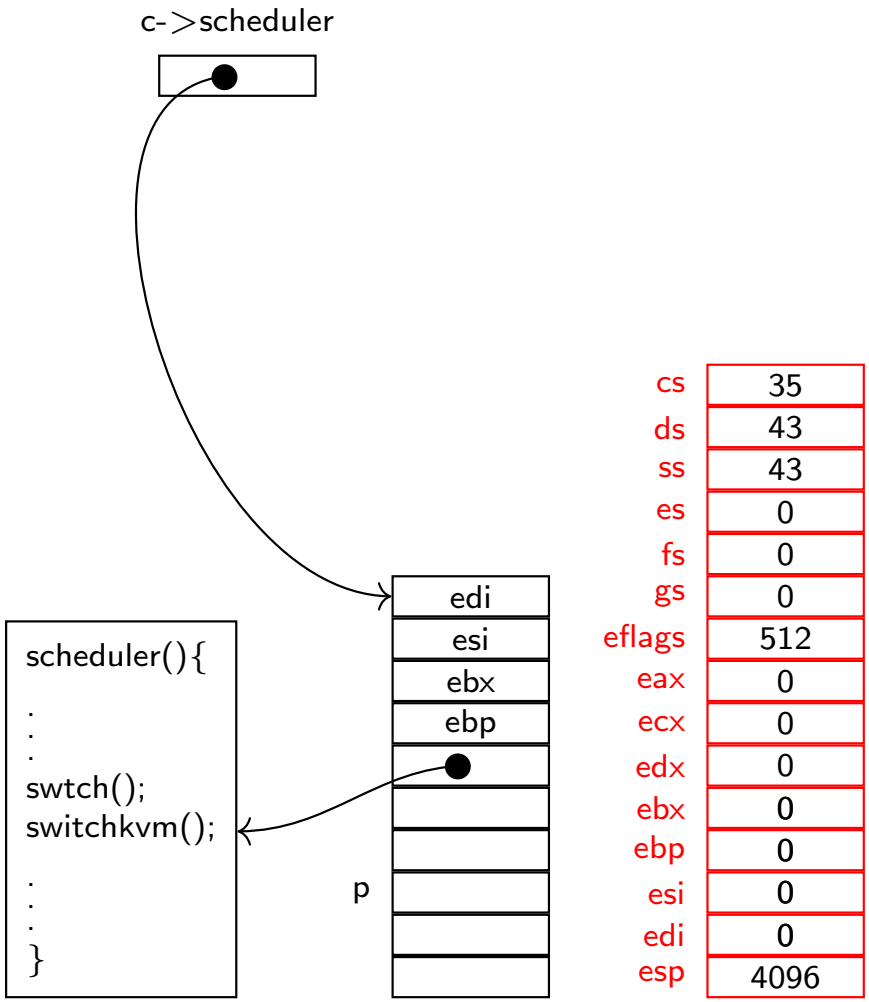


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Context switch

- Calling **swtch()**:
  - Creates a **context** structure on the current stack.
  - Stores the **cotext** structure address created in the first argument.
  - Load the **context** structure pointed to by the second argument.
- We are switching KERNEL contexts.
- User mode context of a process is loaded by the kernel side of the process.

Kernel context seems too small

- Where are **eax**, **ecx**, **edx**????

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  - The base and limit fields are identical across all kernel sides.
  - So, they need to be loaded only on aech kernel entering.
- Where is **gdtr**????
  - The address and size fields are different across processors.
  - The base and limit MUST NOT change between kernel sides on the same CPU.
  - Since **gdtr** is privileged, it needs to be loaded ONLY on kernel initialization.