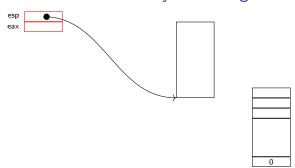
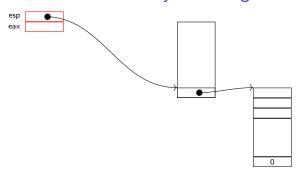


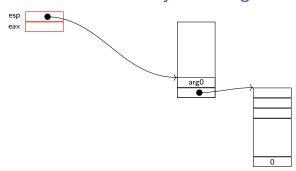
Carmi Merimovich

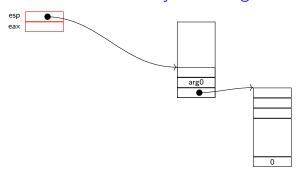
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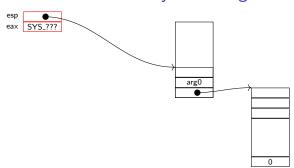
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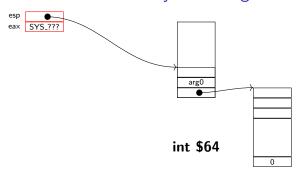




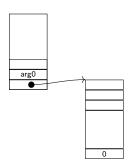


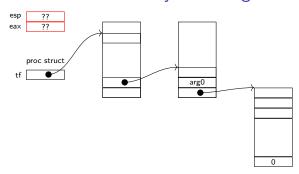


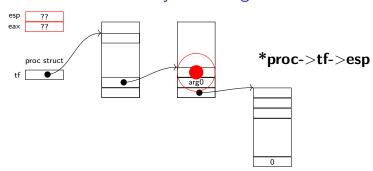


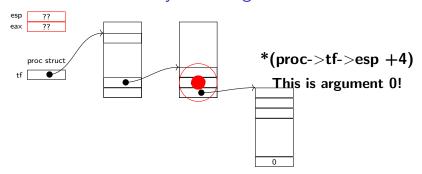


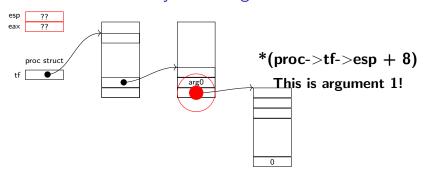












sys_exec

```
int sys_exec(void) {
 char *path , *argv[MAXARG];
 int i:
 uint uargv, uarg;
 if (argstr(0,&path)<0 || argint(1,(int*)&uargv)<0)</pre>
 return -1:
 memset(argv, 0, sizeof(argv));
 for (i=0; i++) {
  if (i >= NELEM(argv)) return -1;
  if (fetchint(uargv+4*i,(int*)&uarg)<0) return -1;
  if (uarg == 0) {
   argv[i] = 0:
  break:
 if (fetchstr(uarg, &argv[i]) < 0) return -1;
 return exec(path, argv);
```

exec

There are four steps in **exec**'s implementation:

- 1. ELF file loading.
- 2. Allocating user stack and guard page.
- 3. Passing argv[] from the old address space to the new one.
- 4. Switching address spaces and fixing trapframe.

exec: step 1.

- We hold the knowldge to implement steps 2–4 above.
- What knowledge is crucially missing in order to implement step 1?
 - How to read from files.
 - The ELF.

Reading from file in kernel mode

Kernel file system interface for reading file

For now we use this kpi, just as user mode programmers use cluelessly the **open/read/close** api

```
void begin_op(void);
struct inode *namei(char *name);
void
               ilock(struct inode *ip);
int
               readi(struct inode *ip,
                         char *buf, uint len,
                         uint pos);
biov
               iunlockput(struct inode *ip);
void end_op(void);
```

Open/Read loop/Close

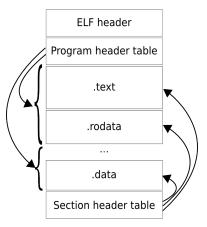
```
begin_op();
ip = namei(filename);
ilock(ip);
for (...) {
 error = readi(ip, buf, bufsiz, pos);
 if (error \ll 0) ...
iunlockput(ip);
end_op();
```

(static) Executlabe and Linkable Format (ELF)

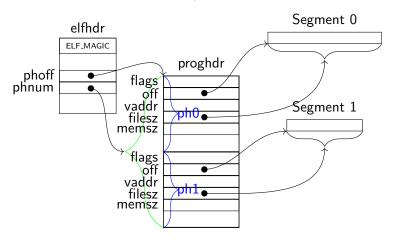
ELF components

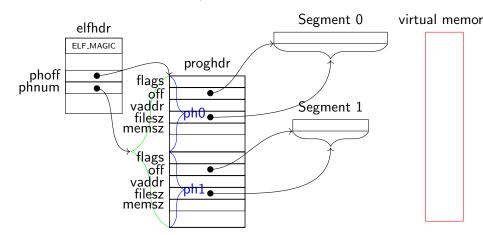
- 1. ELF header. Must begin at byte zero of the file.
- 2. PROGHDR vector.
- 3. Program segments.

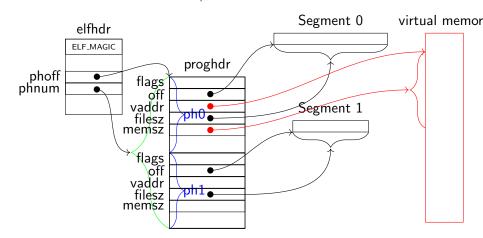
ELF file, very abstract

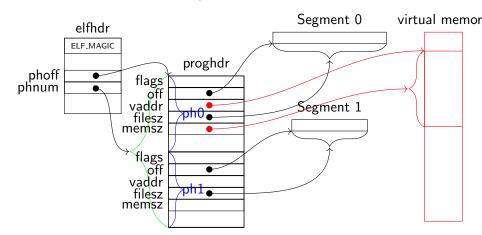


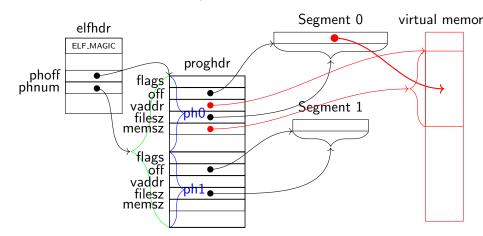
By Suruea - Own work, CC BY-SA 3.0, https://commons.wikimedia.org/w/index.php?curid=2922605

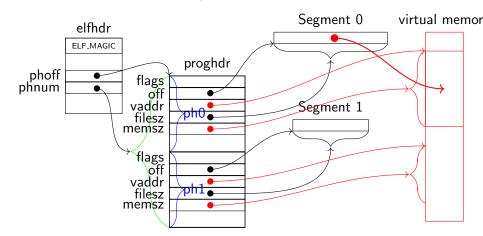


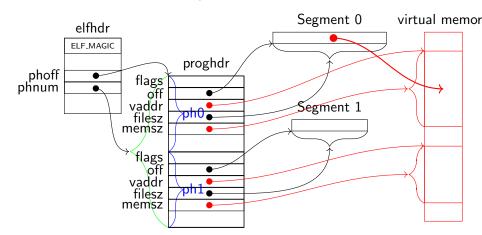


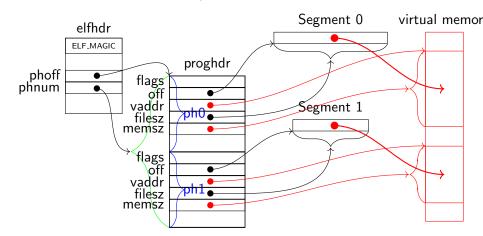












ELFHDR

```
955 struct elfhdr {
    uint magic; // must equal ELF_MAGIC
    uchar elf[12];
    ushort type;
    ushort machine;
    uint version;
    uint entry; // Entry point
    uint phoff; // (File) Location of PROGHDR vectors
    uint shoff;
    uint flags; // flag
    ushort ehsize:
    ushort phentsize;
    ushort phnum; // Length of PROGHDR vector
    ushort shentsize;
    ushort shnum;
    ushort shstrndx:
```

ELFHDR fields we are interested in

- magic: Should be ELF_MAGIC (0x464C457F).
- entry: Virtual address the program is starting at.
- **phoff**: File offset the Program segments Headers vector begins at.
- **phnum**: Number of elements in the Program segments Headers vector.

For each segment there is PROGHDR

```
struct proghdr {
 uint type; // Only PROG_LOAD matters to us
 uint off; // Section location in file
 uint vaddr; // Virtual address of section
 uint paddr; // Physical address of section
 uint filesz; // Section size in file
 uint memsz; // Section size in memory
 uint flags;
 uint align;
```

ELF file loading

- 1. Create a new address space.
- 2. For each segment with a PROG_LOAD type do the following:
 - 2.1 If necessary, allocate memory and add mapping rules so that address vaddr up to vaddr + memsz - 1 will be legal in the new address space.
 - 2.2 Read the file from position off to off + filesz 1 into memory addresses vaddr up to vaddr + memsz 1.
 - Observe: The new address space is not active. Hence we have to read page after page.

allocuvm(): Mission 2.1 above

```
allocuvm(pde_t *pgdir, uint oldsz, uint newsz) {
1927
    char *mem:
    uint a:
    if (newsz >= KERNBASE) return 0;
    if (newsz < oldsz) return oldsz;</pre>
    a = PGROUNDUP(oldsz);
    for (; a < newsz; a += PGSIZE) {
     mem = kalloc();
     if (mem = 0) {
      deallocuvm(pgdir, newsz, oldsz);
      return 0:
     memset(mem, 0, PGSIZE);
     mappages (pgdir, a, PGSIZE, v2p (mem), PTE_W | PTE_U);
    return newsz:
```

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loaduvm: Mission 2.2 above

```
loaduvm(pde_t *pgdir, char *addr,
1903
            struct inode *ip, uint offset, uint sz) {
    uint i, pa, n;
    pte_t *pte:
    if ((uint) addr % PGSIZE != 0) panic("loaduvm: _addr
    for (i = 0; i < sz; i += PGSIZE) {
     if ((pte = walkpgdir(pgdir, addr+i, 0)) == 0) pani
     pa = PTE_ADDR(*pte);
     if (sz - i < PGSIZE)
      n = sz - i;
     else
      n = PGSIZE:
     if (readi(ip, p2v(pa), offset+i, n) != n) return -
    return 0;
```

loaduvm vs. readi

readif struct inode *ip, chai *bui, unit ien, unit pos

- loaduvm is not tied at all to ELF loading!
- It is a generalized readi.
- It allows reading into non-active virtual address spaces!

exec: ELFHDR loading

```
int exec(char *path, char **argv) {
struct elfhdr elf;
struct inode *ip;
begin_op()
if ((ip = namei(path)) == 0) \{end_op();
return -1:
ilock(ip);
pde_t * pgdir pgdir = 0;
if (readi(ip, (char*)&elf, 0, sizeof(elf)) !=
                               sizeof(elf))
goto bad;
if (elf.magic != ELF_MAGIC) goto bad;
if ((pgdir = setupkvm()) == 0) goto bad;
sz=0:
```

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exec: Program segments loading

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```
struct proghdr ph;
for (i=0, off=elf.phoff; i < elf.phnum; i++,
                                off += sizeof(ph)) {
 if (readi(ip, (char*)&ph, off, sizeof(ph)) !=
                             sizeof(ph)) goto bad;
 if (ph.type != ELF_PROG_LOAD) continue;
 if (ph.memsz < ph.filesz) goto bad;</pre>
 if (ph.vaddr + ph.memsz < ph.vaddr) goto bad;</pre>
 if ((sz = allocuvm(pgdir, sz,
             ph.vaddr + ph.memsz)) == 0) goto bad;
 if (ph.vaddr % PGSIZE != 0) goto bad;
 if (loaduvm(pgdir, (char*)ph.vaddr, ip,
                ph.off, ph.filesz) < 0) goto bad;
iunlockput(ip);
end_{-}op(); ip = 0;
```

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