xv6(c)-rev10 (Copyright Frans Kaashoek, Robert Morris, and Russ Cox.) Initialization, Dynamic Allocation Subsystem

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Our context

```
extern char end[];
kinit1(end, P2V(4*1024*1024)); // phys page alloc
kvmalloc(); // kernel page table
seginit();
          // set up segments
pinit();
               // process table
kinit2(P2V(4*1024*1024), P2V(PHYSTOP)); // must c
userinit(); // first user process
mpmain();
```

Why dynamic memory subsystem first?

- Many initialization routines uses utility kernel routines.
- The utility routines use the dynamic allocation subsystem.
- Ergo...

Prerequisites

- freerange
- kfree.

It will be clearer to explain first the dynamic allocation system, and then explain the initialization.

Services supplied by the dynamic allocation

- void *kalloc(void): Returns a pointer to a non-used page.
- kfree(void *p): Marks the **page** pointed to by p as unused.

Implementation

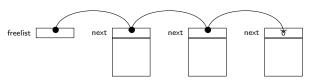
The unused pages are in a linked list headed by **kmem**:

```
struct run {
    struct run *next;
};

struct {
    struct spinlock lock;
    int use_lock;
    struct run *freelist;
} kmem;
```

The first longword in each page points to the next free page.

The freelist



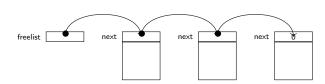
- kalloc removes from the head of the list.
- kfree inserts at the begining of the list.

kalloc

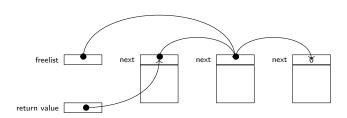
```
char *kalloc(void) {
    struct run *r;

acquire(&kmem.lock);
    r = kmem.freelist;
    if (r)
        kmem.freelist = r->next;
    release(&kmem.lock);
    return (char*)r;
}
```

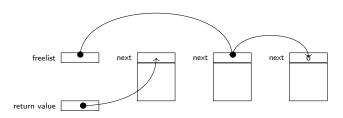
kalloc in action



kalloc in action



kalloc in action

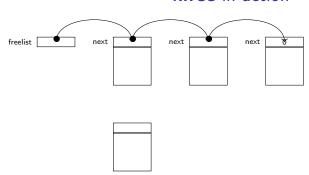


Logically we mean this, of course!

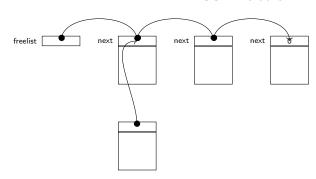
kfree

```
3163 void kfree(char *v) {
    struct run *r:
    if ((uint)v \% PGSIZE || v < end || v2p(v)>=PHYSTOP)
      panic("kfree");
    memset(v, 1, PGSIZE);
    acquire(&kmem.lock);
    r = (struct run*)v;
    r\rightarrow next = kmem. freelist;
    kmem. freelist = r;
     release(&kmem.lock);
```

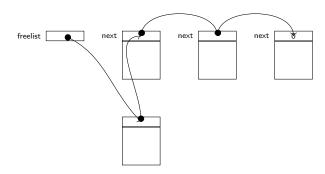
kfree in action



kfree in action



kfree in action



Loading freelist

• What is in the **freelist** when entering main?

Loading freelist

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 - The empty list.
- What addresses can be considered free?

Loading freelist

- What is in the freelist when entering main?
 - · The empty list.
- What addresses can be considered free?
 - Not valid addresses are, well, not valid!
 - Thus addresses in [0x80000000,0x80400000) without kernel addresses.
- xv6 is linked to address 0x80100000.
- The linker defines _end to where the kernel ends.
- Thus the following two ranges are free:
 - [0×80000000,0×800A0000).
 - [_end,0x80400000).
- xv6 'frees' only the range [_end,0x80400000).
- Since we work with pages, _end should be rounded up.
- (elective) There is info we need at the first 64KB.

freerange

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freerange

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 - No. Why?
- Can we, in the xv6 kernel, **kfree** a page not **kalloc**'ed?

freerange

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 No. Why?
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```
• Yes. Why?
```

```
#define PGROUNDUP(sz) (((sz)+PGSIZE-1) &
                                     ~(PGSIZE-1))
   void freerange(void *vstart, void *vend) {
3151
    char *p;
    p = (char*)PGROUNDUP((uint)vstart);
    for (; p + PGSIZE \le (char*)vend; p += PGSIZE)
     kfree(p):
   void kinit1(void *vstart, void *vend) {
3131
    initlock(&kmem.lock, "kmem");
    freerange(vstart, vend);
```

Dynamic allocation initialization

• Step 1:

```
#define KERNBASE 0x80000000
#define PHYSTOP 0x0E000000
#define P2V(a) (((void *) (a)) + KERNBASE)
extern char end[];
 kinit1(end, P2V(4*1024*1024));
```

• Step 2:

```
kinit2(P2V(4*1024*1024), P2V(PHYSTOP));
```