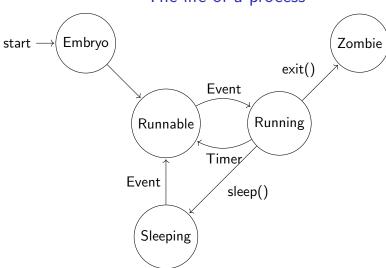
# **Unix Process Programming**

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# The life of a process



All transition are due to INTERRUPTS.

## Processes organization

- The processes are organized into a tree structure.
- The first process is created by the kernel as part of the initialization.
- Except for the first process, processes are created only by processes.
- The fork() system call creates a new process.
- The process invoking the **fork()** is called the **parent** process.
- The created process is called the **child** process.

Process creation/exiting system calls.

# exit()

- Process resources are freed. (mostly)
- Process enters the ZOMBIE state.
- Children of the process are adopted by the first process.
- Process really dies when its parent wait()s on it.

# pid=wait()

- If there are no child process returns error.
- If there are ZOMBIE children:
  - one of them (really) dies.
  - The **id** of the dead process is returned.
- If there are no ZOMBIE children:
  - wait for one of the children to become ZOMBIE.

# exec(filename, argv)

- The code/data/stack of the current process is freed.
- The executable at filename is loaded and begins runing at main.
- The **argv** parameter of **exec** is supplied to the new executable:
  - main(argc, argv)
- Nothing else changes:
  - Files open.
  - Current working directory.
  - etc.
- NOTE: This is NOT a call.
- The new code/data **replaces** the previous code/data.

# pid=fork()

- A new child process is created.
- The parent process proceeds with the return value being the process id of the child process.
- The child process begins as a replication of its parent with one difference:
  - The return value is zero.

# pid=fork() finegrained

#### Hypothetical forking code:

movl \$1,%eax int \$64 movl %eax, pid

#### Run time result:

movl \$1,%eax int \$64	
Parent continues	Child created
movl %eax, pid	movl %eax, pid
(pid=child process id)	( pid==0)

# fork()/exec()/exit()

```
pid = fork();
if (pid < 0)  // fork failed
    exit();
if (pid == 0) { // Child code
    char *argv[] = {"Is", 0};
    exec("Is", argv);
    exit();  // exec failed
} else { // Parent process executes here
}</pre>
```

First process and the /init program

## The first process

- The kernel sets the initial state to:
  - cwd is "/".
  - No file is open.
- Sets standard input, standard output, and standard error, to the console device.
- Creates a process to run the shell (sh).
- Enters an infinite loop of wait()'s.
- It uses the system calls:
  - open.
  - mknod.
  - dup.
  - fork.
  - exec.
  - exit.

# /init (1): Sets standard input and output

```
int main(void) {
  int pid, wpid;

if (open("console", O_RDWR) < 0) {
    mknod("console", 1, 1);
    open("console", O_RDWR);
  }
  dup(0); // stdout
  dup(0); // stderr</pre>
```

# /init (2): Forks to shell, wait loop

```
for (;;) {
  printf(1, "init:_starting_sh\n");
  pid = fork();
  if(pid < 0) {
    printf(1, "init:_fork_failed\n");
    exit():
  if (pid == 0) \{
    char *argv[] = { "sh", 0 };
    exec("sh", argv);
    printf(1, "init:_exec_sh_failed\n");
    exit();
  while ((wpid=wait()) >= 0 \&\& wpid!= pid)
    printf(1, "zombie! \ n");
```

# sh main functionality

#### sh main loop

```
while (read(0, cmd, ...) > 0) {
  if (cmd is internal command)
    executeInternalCmd(cmd);
  else
    forkExternalCmd(cmd);
}
exit();
```

- Internal cmd "cd" causes execution of the chdir system call.
- External commands are assumed to be executble files.

```
Typing:
  ١s
will use the following code,
   pid = fork();
   if (pid = 0) {
    char *argv[] = {" | s", 0};
    exec("ls", argv);
    exit();
   wait();
```

```
Typing:
```

١s

will use the following code, where the parent sh executes:

```
pid = fork();
if (pid == 0) {
  char *argv[] = {"ls", 0};
  exec("ls", argv);
  exit();
}
wait();
```

```
Typing:
```

١s

will use the following code, and the child sh executes:

```
pid = fork();
if (pid == 0) {
  char *argv[] = {"ls", 0};
  exec("ls", argv);
  exit();
}
wait();
```

```
Typing:
  |s-|
will use the code.
   pid = fork();
   if (pid = 0) {
    char *argv[] = {"Is", "-I", 0};
    exec("ls", argv);
    exit();
   wait();
```

```
Typing:
```

```
|s-|
```

will use the code, where the parent sh executes:

```
pid = fork();
if (pid == 0) {
  char *argv[] = {"ls", "-l",0};
  exec("ls", argv);
  exit();
}
wait();
```

```
Typing:
```

```
ls - l
```

will use the code, and the child sh executes:

```
pid = fork();
if (pid == 0) {
  char *argv[] = {"ls", "-l",0};
  exec("ls", argv);
  exit();
}
wait();
```

```
Typing:
```

```
|s>a.txt
will use the code.
   pid = fork();
   if (pid = 0) {
    close (1);
    open("a.txt", O_CREAT);
    char *argv[] = {" | s", 0};
    exec("ls", argv);
    exit();
   wait();
```

```
Typing:
```

```
|s>a.txt
will use the code, where the parent sh executes:
   pid = fork();
   if (pid == 0) \{
    close (1);
    open("a.txt", O_CREAT);
    char *argv[] = {" | s", 0};
    exec("ls", argv);
    exit();
   wait();
```

```
Typing:
```

```
|s>a.txt
will use the code, and the child sh executes:
   pid = fork();
   if (pid == 0) {
    close (1);
    open("a.txt", O_CREAT);
    char *argv[] = {" | s", 0};
    exec("Is", argv);
    exit();
   wait();
```

```
Typing:
```

```
|s - l| > b.txt
will use the code.
   pid = fork();
   if (pid = 0) {
    close (1);
    open("b.txt", O_CREAT);
    char *argv[] = {" | s", "-|", 0};
    exec("ls", argv);
    exit();
   wait();
```

```
Typing:
```

```
Is -I > b \cdot txt will use the code, where the parent sh executes:
```

```
pid = fork();
if (pid == 0) {
  close (1);
  open("b.txt", O_CREAT);
  char *argv[] = {"Is", "-I", 0};
  exec("Is", argv);
  exit();
}
wait();
```

```
Typing:
```

```
|s-1>b.txt
will use the code. and the child sh executes:
   pid = fork();
   if (pid == 0) {
    close (1);
    open("b.txt", O_CREAT);
    char *argv[] = {" | s", "-|", 0};
    exec("Is", argv);
    exit();
   wait();
```

```
Typing:
  sh < b.txt
will use the code.
   pid = fork();
   if (pid = 0) {
    close (0);
    open("b.txt", O_RONLY);
    char *argv[] = {"sh", 0};
    exec("sh", argv);
    exit();
   wait();
```

```
Typing:
  sh < b.txt
will use the code, where the parent sh executes:
   pid = fork();
   if (pid == 0) {
    close (0);
    open("b.txt", O_RONLY);
    char *argv[] = {"sh", 0};
    exec("sh", argv);
    exit();
   wait();
```

```
Typing:
  sh < b.txt
will use the code, and the child sh executes:
   pid = fork();
   if (pid == 0) {
    close (0);
    open("b.txt", O_RONLY);
    char *argv[] = {"sh", 0};
    exec("sh", argv);
    exit();
```

wait();

Unix Process programming

```
Typing:

cat a.bat | sh

will use the code:
```

```
int p[2];
pipe(p);
pid = fork();
if (pid = 0) {
close (1):
dup(p[1]);
 close(p[0]):
 close(p[1]);
 char *argv[] = {"cat", 0};
 exec("cat", argv);
 exit();
```

```
pid = fork()
if (pid = 0) {
close (0):
dup(p[0]);
 close(p[0]);
 close(p[1]);
 char *argv[] = {"sh", 0}
 exec("sh", argv);
 exit();
close(p[0]);
close(p[1]);
wait();
```

```
Typing:
```

```
cat a.bat | sh
will use the code: where the parent sh executes:
```

```
int p[2];
pipe(p);
pid = fork();
if (pid == 0) {
close (1):
dup(p[1]);
 close(p[0]);
 close(p[1]);
 char *argv[] = \{"cat", 0\};
 exec("cat", argv);
 exit();
```

```
pid = fork()
if (pid = 0) {
close (0):
dup(p[0]);
close(p[0]);
 close(p[1]);
 char *argv[] = \{"sh", 0\}
 exec("sh", argv);
 exit();
close(p[0]);
close(p[1]);
wait();
```

```
Typing:
```

```
cat a.bat | sh
```

will use the code: the first child **sh** executes:

```
int p[2];
                             pid = fork()
pipe(p);
                             if (pid == 0) {
pid = fork():
                              close (0);
if (pid == 0) {
                              dup(p[0]);
close (1):
                              close(p[0]);
dup(p[1]);
                              close(p[1]);
                              char *argv[] = {"sh", 0}
 close(p[0]);
 close(p[1]);
                              exec("sh", argv);
 char *argv[] = {"cat", |0}; exit();
 exec("cat", argv);
 exit();
                             close(p[0]);
                             close(p[1]);
                             wait();
```

Unix Process programming ()

```
Typing:
```

```
cat a.bat | sh
```

will use the code: the second child **sh** executes:

```
int p[2];
pipe(p);
pid = fork();
if (pid == 0) {
close (1):
 dup(p[1]);
 close(p[0]);
 close(p[1]);
 char *argv[] = {"cat", 0}; exit();
 exec("cat", argv);
 exit();
```

```
pid = fork()
if (pid == 0) {
close (0):
dup(p[0]);
close(p[0]):
 close(p[1]);
 char *argv[] = {"sh", 0}
 exec("sh", argv);
close(p[0]);
close(p[1]);
```