

### G53KRR 2018 Answer to Informal Exercise 3

1. Rewrite the following propositional formulas to clausal form (conjunction of disjunctions of literals):

(a)  $p \vee q \supset r$

(b)  $r \supset s$

(c)  $p$

*Answer.*

The answer uses the following definitions:

**Definition of  $\supset$**   $\alpha \supset \beta$  is equivalent to  $\neg\alpha \vee \beta$

**De Morgan law**  $\neg(\alpha \vee \beta)$  is equivalent to  $\neg\alpha \wedge \neg\beta$

(the second law:  $\neg(\alpha \wedge \beta) \equiv \neg\alpha \vee \neg\beta$ )

**Distributivity**  $\alpha \vee (\beta \wedge \gamma) \equiv (\alpha \vee \beta) \wedge (\alpha \vee \gamma)$

(another instance of Distributivity is  $\alpha \wedge (\beta \vee \gamma) \equiv (\alpha \wedge \beta) \vee (\alpha \wedge \gamma)$ )

$p \vee q \supset r$  is by definition of  $\supset$  equivalent to  $\neg(p \vee q) \vee r$ .  $\neg(p \vee q) \vee r$  is by de Morgan's law equivalent to  $(\neg p \wedge \neg q) \vee r$ . By distributivity,  $(\neg p \wedge \neg q) \vee r$  is equivalent to  $(\neg p \vee r) \wedge (\neg q \vee r)$ . The last formula is in clausal form.

$r \supset s$  is equivalent to  $\neg r \vee s$ .

$p$  is already in clausal form.

So we get:  $\neg p \vee r$  (a1),  $\neg q \vee r$  (a2),  $\neg r \vee s$  (b),  $p$  (c).

2. Prove by resolution that the sentences above entail  $s$ .

*Answer.* We derive a contradiction from the sentences above and *the negation of  $s$* ,  $\neg s$  (d):

**1**  $\neg p \vee r$  (a1)

**2**  $p$  (c)

**3**  $r$  resolution 1,2

**4**  $\neg r \vee s$  (b)

**5**  $s$  resolution 3,4

**6**  $\neg s$  (d)

**7**  $\emptyset$

3. Show by resolution that the following set of clauses is unsatisfiable:

(a)  $p \vee q \vee r$ ,

(b)  $p \vee q \vee \neg r$ ,

(c)  $p \vee \neg q \vee r$ ,

(d)  $p \vee \neg q \vee \neg r$ ,

- (e)  $\neg p \vee q \vee r$ ,
- (f)  $\neg p \vee q \vee \neg r$ ,
- (g)  $\neg p \vee \neg q \vee r$ ,
- (h)  $\neg p \vee \neg q \vee \neg r$

*Answer.*

- 1**  $p \vee q$  by resolution from (a) and (b) (we also use  $\alpha \vee \alpha \equiv \alpha$ ; usually resolution is applied to sets rather than disjunctions)
- 2**  $p \vee \neg q$  by resolution from (c) and (d)
- 3**  $p$  by resolution from 1 and 2
- 4**  $\neg p \vee q$  by resolution from (e) and (f)
- 5**  $\neg p \vee \neg q$  by resolution from (g) and (h)
- 6**  $\neg p$  by resolution from 4 and 5
- 7**  $\emptyset$  by resolution from 3 and 6