# Knowledge representation and reasoning Lecture 20: Planning 3

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#### Plan of the lecture

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- SET and SEM
- Partial order planning (based on Russell and Norvig, Russell's slides)
- Overview of planning algorithms and systems

#### Outline

- ♦ Totally vs partially ordered plans
- ♦ Partial-order planning
- $\Diamond$  Examples

#### Totally vs partially ordered plans

So far we produced a linear sequence of actions (totally ordered plan)

Often it does not matter in which order some of the actions are executed

For problems with independent subproblems often easier to find a **partially** ordered plan: a plan which is a set of actions and a set of constraints  $Before(a_i, a_i)$ 

Partially ordered plans are created by a search through a space of plans (rather than the state space)

#### Partially ordered plans

Partially ordered collection of steps with

Start step has the initial state description as its effect Finish step has the goal description as its precondition causal links from outcome of one step to precondition of another temporal ordering between pairs of steps

Open condition = precondition of a step not yet causally linked

A plan is complete iff every precondition is achieved

A precondition is achieved iff it is the effect of an earlier step and no possibly intervening step undoes it

# Example

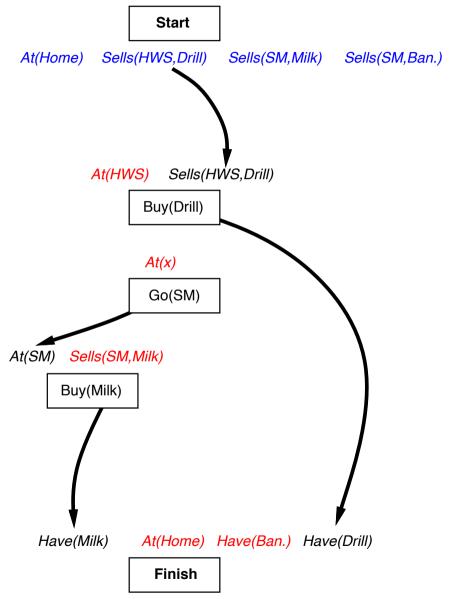
Start

At(Home) Sells(HWS,Drill) Sells(SM,Milk) Sells(SM,Ban.)

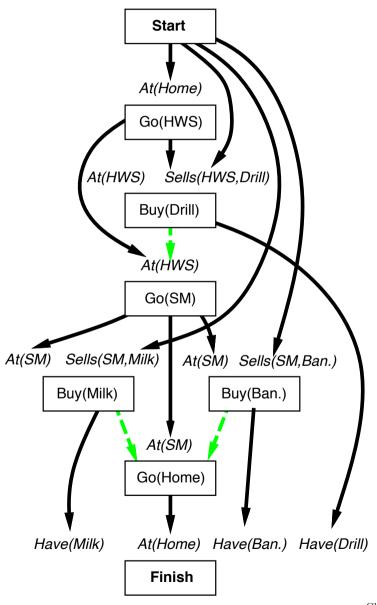
Have(Milk) At(Home) Have(Ban.) Have(Drill)

**Finish** 

### Example



### Example



#### Planning process

Operators on partial plans:

add a link from an existing action to an open condition add a step to fulfill an open condition order one step wrt another to remove possible conflicts

Gradually move from incomplete/vague plans to complete, correct plans

Backtrack if an open condition is unachievable or if a conflict is unresolvable

#### POP algorithm sketch

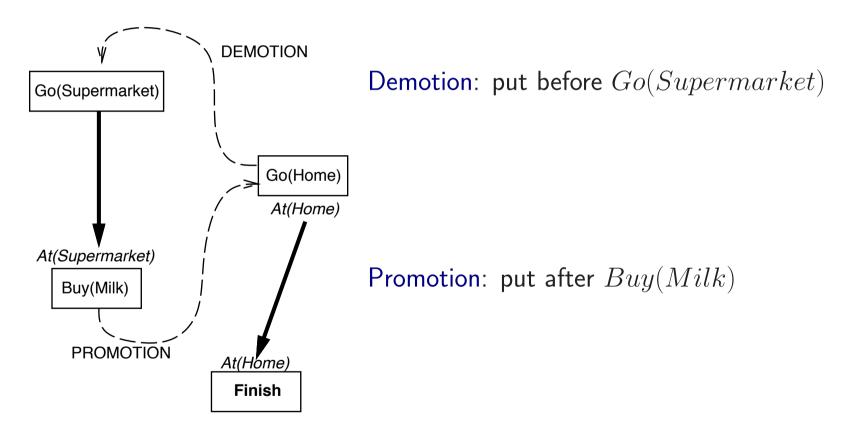
```
function POP(initial, goal, actions) returns plan
   plan \leftarrow Make-Minimal-Plan(initial, goal)
   loop do
       if SOLUTION? (plan) then return plan
       S_{need}, c \leftarrow \text{Select-Subgoal}(plan)
       Choose-Action (plan, actions, S_{need}, c)
       RESOLVE-THREATS (plan)
   end
function Select-Subgoal (plan) returns S_{need}, c
   pick a plan step S_{need} from STEPS( plan)
       with a precondition c that has not been achieved
   return S_{need}, c
```

#### POP algorithm contd.

```
procedure Choose-Action(plan, actions, S_{need}, c)
   choose a step S_{add} from actions or STEPS(plan) that has c as an effect
   if there is no such step then fail
   add the causal link S_{add} \xrightarrow{c} S_{need} to Links (plan)
   add the ordering constraint S_{add} \prec S_{need} to Orderings (plan)
   if S_{add} is a newly added step from actions then
        add S_{add} to STEPS( plan)
        add Start \prec S_{add} \prec Finish to Orderings (plan)
procedure Resolve-Threats(plan)
   for each S_{threat} that threatens a link S_i \xrightarrow{c} S_j in LINKS (plan) do
        choose either
              Demotion: Add S_{threat} \prec S_i to Orderings (plan)
              Promotion: Add S_i \prec S_{threat} to Orderings (plan)
        if not Consistent (plan) then fail
   end
```

#### Clobbering and promotion/demotion

A clobberer is a potentially intervening step that destroys the condition achieved by a causal link. E.g., Go(Home) clobbers At(Supermarket):



#### Properties of POP

Nondeterministic algorithm: backtracks at choice points on failure:

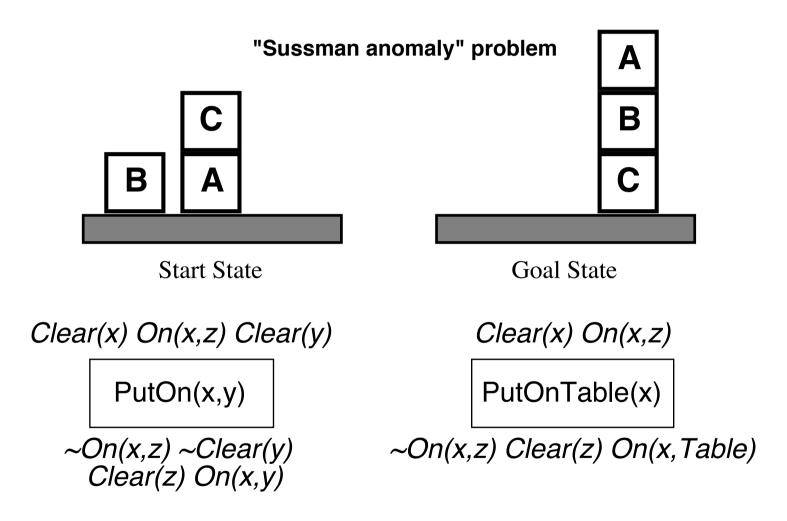
- choice of  $S_{add}$  to achieve  $S_{need}$
- choice of demotion or promotion for clobberer
- selection of  $S_{need}$  is irrevocable

POP is sound, complete, and systematic (no repetition)

Can be made efficient with good heuristics derived from problem description

Particularly good for problems with many loosely related subgoals

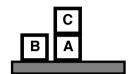
#### Example: Blocks world



+ several inequality constraints

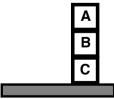
START

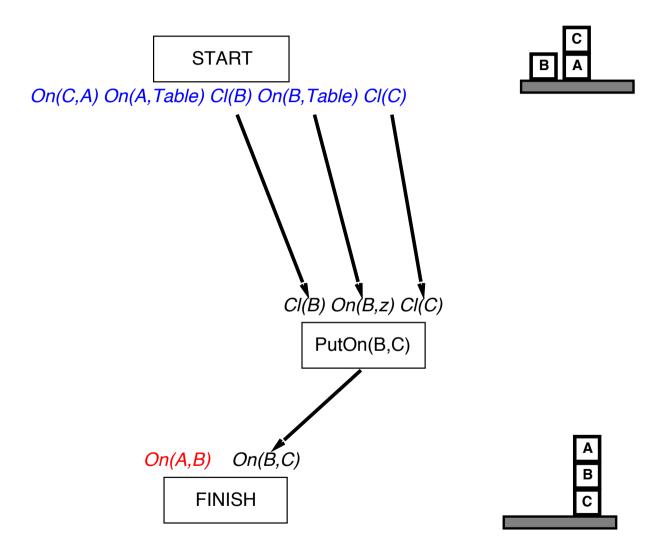
On(C,A) On(A, Table) Cl(B) On(B, Table) Cl(C)

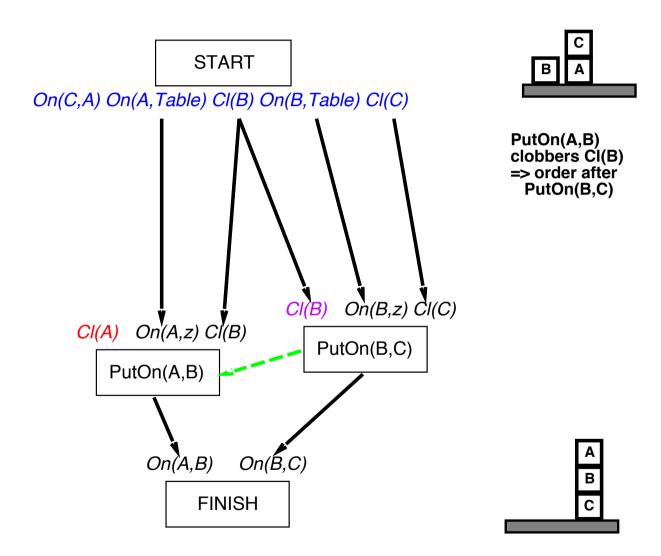


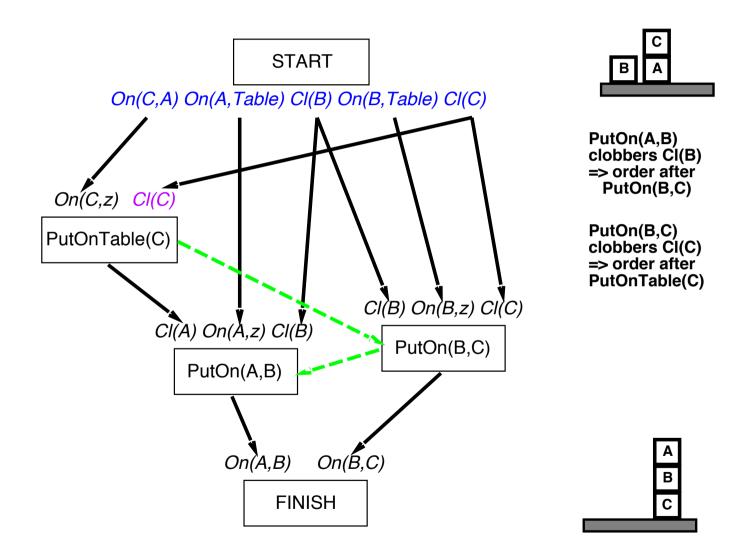
On(A,B) On(B,C)

**FINISH** 









#### Plan of the lecture

#### Other planning approaches

- SatPlan: reduction to propositional reasoning, and checking for satisfiability.
- Can check whether a plan of fixed length exists
- Hierarchical planning (planning with abstraction): first find a plan with high-level abstract actions, then refine it
- Planning with time and resources, and scheduling plans

#### The complexity of classical planning

PlanSAT is the question whether there exists any plan that solves a given planning problem

Bounded PlanSAT is the question whether there exists a plan of length k or less

PlanSAT is about **satisficing** (want any solution, not necessarily the cheapest or the shortest)

Bounded PlanSAT can be used to ask for the optimal solution

If in the PDDL language we do not allow functional symbols, both problems are decidable

Complexity of both problems is PSPACE (can be solved by a Turing machine which uses polynomial amount of space)

 $NP \subseteq PSPACE$  (PSPACE is even harder than NP)

### Some of the top-performing systems

International planning competition winners (from Russell and Norvig, 3rd edition):

Year	Track	Winning systems (approaches)
2008	Optimal	Gamer (symbolic bi-directional search)
2008	Satisficing	LAMA (fast forward search with FF heuristic)
2006	Optimal	SATPlan, MAXPlan (boolean satisfiability)
2006	Satisficing	SGPlan (forward search, partition into independent subproblems
2004	Optimal	SATPlan (boolean satisfiability)
2004	Satisficing	Fast Diagonally Forward (forward search with causal graph)
2002	Automated	LPG (local search, constraint satisfaction)
2002	Hand-coded	TLPLan (temporal action logic with control rules for forward searc
2000	Automated	FF (forward search)
2000	Hand-coded	TalPlanner (temporal action logic with control rules for forward sea
1998	Automated	IPP (planning graphs); HSP (forward search)

#### What follows for the algorithms

If a planning system based on a particular planning algorithm is very fast it does not mean necessarily that the algorithm is 'better'

Sat solvers are very fast, but one may argue that it is not very practical to propositionalise planning problems

Forward planning with good heuristics can be very fast but again may be not always practically possible

Partial order planners are considered to be very flexible and generally useful, although they don't feature in the winners table...

#### Plan of the lecture

#### **Next lecture**

- Reasoning with uncertainty
- Brachman and Levesque, chapter Chapter 12, slides 202-216.
- Revision advice
- I can also answer revision questions during the lab on Monday
- No lecture next Thursday (13th of December)