Real-world Functional Programming

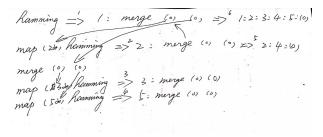
Coursework Part I Report

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1 Task I.1

```
1
3 hamming :: [Int]
4 hamming :: [Int]
5 merge :: [Int] -> [Int] -> [Int]
7 merge xsse(xxs) ysse(y;ys)
8 | x == y = x : merge xs ys
9 | x < y = x : merge xs ys
10 | x > y = x : merge xs ys
11
```

(a) Hamming Function Definition



(b) Cyclic Graph

Figure 1: TaskI.1

Hence, the hamming function can be defined as a infinite list in a recursive manner. As Figure 1 (a) shows, the type of hamming function is a list of Int. This implementation using map function to calculate 2x, 3x and 5x and then merge then together as Figure 1 (b) demonstrated how it was evaluated.

2 Task I.2

Figure 2: Extended Evaluator

Figure 2 demostrates how evalCell function was extended to support Sum and Avg expression. The key idea of this implementation is to given to CellRef c1 and c2, find every cell in between. Next, lookup corresponding values in the given sheet s. Then put all the values found in sheet s in list l. Finally using foldr to calculate the sum of all values.

The similiar idea was used to implementate the evaluation of Avg expression but a further step was taken to yield the average value.

The most obvious weakness of this evaluator is that it will stuck in infinite loop when there are cyclic references in the given sheet. As the type signature indicates, the evalCell function will return a Double for every given sheet s and an Exp, which means when the Exp is Ref, it will keep evaluating until a Double can be returned. If there is a cyclic reference, this evaluator will keep evaluating without return.

One simple solution to this problem is to maintain a seperate list to keep track of the evaluation of Ref Exp. This list works like stack but disallowing duplicated elements. A cell will of pushed to the stack if it contains Ref Exp as well as the reference cell. Then continue to evaluate that reference cell, if it returns, which means there is no cyclic reference. Hence those cell can be popped off the stack. If a cell that contains Ref Exp already exists in the stack that there must be a cyclic reference. Therefore, the evaluator can just return fail without continue evaluating.

3 Task I.3

In skew binary random access list, a number is represented by a sequence of digits starting from least significant bit to the most significant bit. Each digit is associated with weight predefined by the positional system.

```
66 drop :: Int -> RList a -> RList a
67 drop _ [] = []
68 drop 0 t = t
69 drop i (t:ts)
      | i == fst t = ts
| i > fst t = drop (i - fst t) ts
| otherwise = drop' i (fst t) (snd t) ++ ts
74 drop' :: Int -> Int -> Tree a -> RList a
75 drop' 0 _ (Leaf a) = [(1, (Leaf a))]
76 drop¹
              (Leaf _) = []
77 drop' 0 w n@(Node _ _ _) = [(w, n)]
78 drop' i w (Node l _ r)
      | i <= half = drop' (i - 1) half l ++ [(half, r)]
| otherwise = drop' (i - half -1) half r
      where
        half = w `div` 2
82
84 testDrop :: Int -> Bool
85 testDrop i = l1 == l2
87
         l1 = last $ take i $ iterate (cons 1) empty
        l2 = drop i $ last $ take (i*2) $ iterate (cons 1) empty
89
    - 2.2k SBRAL.hs Haskell @YSPEeK Git-master
```

Figure 3: Drop Function

Figure 3 shows the implementation of drop function which drops the first n elements from the list. The drop function takes an integer and a RList as input and return a RList without the first n elements as output. If the tree is empty which means nothing can be dropped, hence empty list is returned. If the input integer is zero which means nothing needs to be dropped, then the same list taken as input will be returned. If the size of the first element in the list is equals to the number of elements need to be dropped, then the tail of the input list will be returned. If the number of elements need to be dropped is greater that the size of first element in the list, then the first element will be dropped and the drop function will be called recursively to continue processing the remaining list. The operations mentioned above would run id constant time.

The next case of the drop function determines the overall time complexity. If the number of elements need to be dropped is less than the size of first element in the list, then a helper function drop' will be called. The drop' function takes two integers and a tree as input and returns a RList. Similar to the first two cases of drop function, the first three cases of drop' function are practically doing nothing. Since the input tree is complete binary leaf tree, the size of tree will be halved for every element dropped from the tree. Therefore, the overall time complexity is O(logn). If the number of elements need to be dropped is less than or equals to the half size of the input tree, then only the left children of the input tree will be processed. Otherwise, the left children will be dropped and the right children will be processed accordingly. The drop' function will be called recursively untill the drop is finished or there nothing to drop and return the final result.

A test function called testDrop id created to test the drop function. The key idea is the RList which represents n should equal to the RList that represents 2n without the first n elements.

4 Task I.4

Figure 4: Intervals

Figure 4 shows how Ivl was made an instance of Num and Fractional class. As an instance of Num class, (+), (-) and abs functions were implementated according the rule of interval arithmetic.

As for (*) function, there are a few cases that NaN may be yield which means the arithmetic operation is illegal. Hence, the (*) function will return error every time NaN occurs. The signum function takes an Ivl as input and

return an Ivl with the signum value of lower bound and signum value of upper bound. the from Integral function takes an int as input and returns an Ivl with the same value converted from input.

the from Rational function in Fractional class works in a similar way as the from Integral function. The (/) function also employed the same manner to deal with NaN. Furthermore, in this implementation, divided by zero is not allowed. Hence, if divided by zero ever occurs, an error will be returned.

As for the (+/-) function, the key idea is to calculate the symmetric interval around a specific number. The input numbers may be negative, to deal with this situation, the lower bound of the resulting Ivl is the smaller value of the result of plus and minus of the two input value while the upper bound is the larger value of the result of plus and minus of the two input values. A few test cases is provided in the testIvl function to test the result of the (+/-) function.

5 Task I.5

```
38 statistics :: Drawing Object → Statistics
39 statistics dobject =
40 Statistics {
41    avgArea = asSumArea accum / (fromIntegral $ asCount accum),
42    avgCircumference = asSumCircumference accum / (fromIntegral $ asCount accum),
43    maxArea = asMaxArea accum,
44    maxCircumference = asMaxCircumference accum }
45    where
46    accum = accumStats (AccumStats 0 0 0 0 0) dobject
47
48
49    accumStats :: AccumStats → Drawing Object → AccumStats
50    accumStats s: AccumStats → Drawing Object → AccumStats
62    asCount = asCount accum + 1,
63    asSumArea = asSumArea accum + area object,
64    asSumCircumference = asSumCircumference accum + circum object,
65    asMaxArea = max (asMaxArea object),
65    asMaxArea = max (asMaxArea object),
66    asMaxCircumference = max (asMaxCircumference accum) (circum object) }
75    accumStats accum (Group []) = accum
66    accumStats accum (Group (x:xs)) = accumStats accum x) (Group xs)
67    accumStats accum (Group (x:xs)) = accumStats accum x) (Group xs)
68    accumStats accum (Group (x:xs)) = accumStats accum x) (Group xs)
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62    accumStats (accumStats accum x) (Group xs)
63    accum
```

(a) Recursive Statistics

(b) Statistics using foldMap

Figure 5: TaskI.5 1 2

Figure 5 shows how the statistics functions were implementated using two different approach. Figure 5 (a) illustrates how the statistics function was defined using recursion. The statistics function takes a Drawing Object as input and return Statistics as result by calling the accumStats function to do the computation. The accumStats was implementated using recursion by taking the result of last computation and the remaining Drawing Object as input and eventually return the final result. The area and circum functions were defined to aid the computation. They calculated the area and circumference of a single object separately.

Figure 5 (b) shows a different approach of defining the statistics function by making Drawing an instance of Foldable and making AccumStats an instance of Semigroup. By making Drawing an instance of Foldable, a function can be applied to every object in drawing and the results can be combined. By making AccumStats an instance of Semigroup, it provided a way of merging two AccumStats to form a new AccumStats.

As for the statistics using foldMap, the overall structure is similar to the recursive version, but the accumStats was

defined differently. Instead processing all objects in Drawing, the accumStats now only need to convert one object to AccumStats and the results for all object in Drawing can be combined using foldMap.

```
9 newtype Length = Length { getLength :: Double } deriving (Eq, Ord, Num, Show)
10 newtype Area = Area { getArea :: Double } deriving (Eq, Ord, Num, Show)
11 12 instance Bounded Length where
13 minBound = 0
14 maxBound = Length $ read "Infinity"
15 16 instance Bounded Area where
17 minBound = 0
18 maxBound = Area $ read "Infinity"
19
```

Figure 6: newtype Bounded

Figure 6 shows the implementation details of Bounded class for Length and Area. They both bounded below by 0 and above by +Infinity.

(a) Recursive Statistics for newtype

(b) Statistics using foldMap for newtype

Figure 7: TaskI.5 3

Figure 7 demostrates the two version of statistics function implementated according newtype Length and Area. They work exactly the same as the version showed in figure 5 but there are some work arounds according to the type definitions. For example, some values may wrapped in Area then they need to be extracted out and do some computation and put back to Area. The same processes works for values in Length, Sum and Max.