Backwards reasoning is from (potential) conclusions to facts instead of from facts to conclusions.

Prolog uses it because the space of possible conclusions grows too quickly in the number of premises.

You can use the consult and reconsult predicates, or use bracket notation.

Bracket notation allows re-consultation of several files at once:

```
['file1.pl', file2.pl']
```

A Prolog database knows everything it needs to know.

A name followed by zero or more arguments. Parens are omitted if there are no arguments.

A structure terminated by a period. It represents a simple fact.

A structure followed by a turnstile and a list of structures separated by commas. It represents a rule.

# A **collection** of clauses with the same *functor* (name) and arity.

# A collection of predicates, in any order.

Use single quotes, which does **not** make a string.

# They indicate a list of ASCII values.

Use double quotes or an escaped single quote to use a single quote.

Other escapes use backslash as in other languages.

call --> | | --> exit fail <-- | | <-- redo -----

exit ports connect to call ports. fail ports connect to redo ports.

Facts can be added at any time using the assert predicate.

Facts can be removed at any time using the retract predicate.

Such rules are dynamic.

To force a rule to be interpreted as a single argument, since rules contain commas.

```
assert((loves(chuck, X) :- female(X), rich(X))).
```

- Output can't be undone.
- assert and retract can't be undone either.

write predicate outputs its single argument to stdout.

listing(predicate)	
listing	

It's an anonymous variable. It can unify with anything.

fail doesn't force other the entire predicate to fail. Other clauses will be tried.

Using a cut creates a commit point, preventing backtracking past the commit point and preventing attempts on other clauses.

The predicate as a whole fails.

To tell Prolog you are only going to use it once, but don't wish to use the anonymous variable.

You can still use it again if you really want.

For onymous variables, the single clause in which it appears.

The same order in which they were defined.

- Predicates with multiple clauses that have "tests" in them.
- Recursion.

### Use a fail loop.

```
my_func(X) :-
    Some(),
    Imperative(),
    Calls(),
    fail.
my func()
```

It's generally bad style to do this.

A parameter (conventionally the final one) of a parameter list can be used for output.

```
?- assert((first([Head | Tail], X) :- X = Head)). true.
```

?- first([1,2,3], X).

X = 1.

## You can store state in the database as facts.

```
bump_count :-
   retract(count(X)),
   Y is X + 1,
   assert(count(Y)).
```

It's generally bad style to do this.

The "univ" operator. It converts between structures and lists.

loves (chuck, X) = .. [loves, chuck, X]

```
?-member(1, [1, 2]).
true .
?- append([1], [2], [1, 2]).
true.
```

Make the final case a "dummy" with no body that uses anonymous variables for all the parameters.

wrapper : potentially\_failing\_call.
wrapper.

Now any call to wrapper will succeed.

They are like assert, but assert a quarantees the added clause will come before any clauses with the same functor.

Likewise assertz guarantees the added clause will be the last case in its predicate.

It removes *all* clauses of the predicate with the given functor and arity.

abolish (somePred, arity).

+, -, \*, /, and mod have their normal meanings, but *only* when evaluated.

They may not be evaluated when you want, leading to strange outcomes like:

?-2+2=4.

Static clauses are the default and cannot be later modified using assert/retract.

Marking clauses <code>dynamic</code> (before they are defined) allows you to change the definition during program execution.

:- dynamic somePredArityTwo/2, somePredArityOne/1.

To actually force arithmetic to be performed you can use is or comparison operators like =:=,=/=,>,>=,<,<=.