

Section 6.4 - Reflections

Mittwoch, 21. Februar 2024 15:00

• black text: this stems from the book

• blue text: deviates from book bc of types etc

Def: \underline{B} τ -algebra, A set, $h: B \rightarrow A$ and $g: A \rightarrow B$. The **reflection** of \underline{B} w.r.t. g and h is \underline{A} where for all $f \in \tau$ of arity n and all $x_1, \dots, x_n \in A$,

$$f^{\underline{A}}(x_1, \dots, x_n) := h(f^{\underline{B}}(g(x_1), \dots, g(x_n)))$$

For a class \mathcal{C} of τ -algebras, the class of their reflections is denoted by $\text{Refl}(\mathcal{C})$.

Analogous to the HSP-Lemma, the following holds:

Lemma \mathcal{C} class of τ -algebras.

- The smallest class of τ -algebras containing \mathcal{C} and closed under Refl , H , S and P is $\text{Refl } P(\mathcal{C})$.
- The smallest class of τ -algebras containing \mathcal{C} and closed under Refl , H , S and P^{fin} is $\text{Refl } P^{\text{fin}}(\mathcal{C})$.

Proof: We show first statement. Suffices to show closeness of $\text{Refl } P(\mathcal{C})$ under Refl, H, S, P .

Refl: By the definition of Reflection, the following diagram commutes.

$$\begin{array}{ccccc} \underline{B} & \xrightarrow{h_1} & \underline{A} & \xrightarrow{h_2} & \underline{C} \\ \text{Refl} \downarrow & & \text{Refl} \downarrow & & \text{Refl} \downarrow \\ g_1 \uparrow & & g_2 \uparrow & & \\ \text{Refl} \quad \text{Refl} \quad \text{Refl} & & & & \\ \text{Refl} \quad \text{Refl} \quad \text{Refl} & & & & \\ g_1 \circ g_2 \quad h_2 \circ h_1 & & & & \\ \end{array} \sim h_2(h_1(f^{\underline{A}}(g_1(g_2(x))))) \sim (h_2 \circ h_1)(f^{\underline{A}}((g_1 \circ g_2)(x)))$$

Thus, $\text{Refl}(\text{Refl}(\mathcal{K})) = \text{Refl}(\mathcal{K})$ for any class \mathcal{K} .

H: Show that $H(\mathcal{K}) \subseteq \text{Refl}(\mathcal{K})$ for any class \mathcal{K} .

Given $\underline{A} \in H(\mathcal{K})$, there is a surjective homomorphism h from some $\underline{B} \in \mathcal{K}$

to \underline{A} . Let $g: A \rightarrow B$ s.t. $h \circ g = \text{id}_A$. It holds for $f \in \tau$ of arity n and $x_1, \dots, x_n \in A$ that

$$h(f^{\underline{B}}(g(x_1), \dots, g(x_n))) = f^{\underline{A}}(h(g(x_1)), \dots, h(g(x_n))) = f^{\underline{A}}(x_1, \dots, x_n).$$

Thus, \underline{A} is the reflection of \underline{B} w.r.t. h and g , so $\underline{A} \in \text{Refl}(\mathcal{K})$.

S: Show that $H(\mathcal{K}) \subseteq \text{Refl}(\mathcal{K})$ for any class \mathcal{K} .

Given $\underline{A} \in H(\mathcal{K})$, set $g: A \rightarrow B$ to be the identity and $h: B \rightarrow A$ any extension of g to B . Then \underline{A} is the reflection of \underline{B} w.r.t. h and g , analogous to H.

P: Show that $P(\text{Refl } P(\mathcal{K})) \subseteq \text{Refl}(P(\mathcal{K}))$.

Given a set I , algebras $(\underline{B}_i)_{i \in I} \in (P(\mathcal{K}))^I$ and τ -algebras $(A_i)_{i \in I}$ s.t.

wonderland: for every operation f of \underline{B} then apply Refl to operation clones.
Need this at one point in proof of the theorem.

Given a set I , algebras $(\underline{B}_i)_{i \in I} \in (P(K))^I$ and τ -algebras $(A_i)_{i \in I}$ s.t. for all i , A_i is the reflection of \underline{B}_i w.r.t. h_i and g_i .

Then $h: (\underline{b}_i)_{i \in I} \mapsto (h_i(\underline{b}_i))_{i \in I}$ and $g: (a_i)_{i \in I} \mapsto (g_i(a_i))_{i \in I}$ witness that $\prod_{i \in I} A_i$ is a reflection of $\prod_{i \in I} \underline{B}_i$, which proves

$$P(\text{Ref}(P(K))) \subseteq \text{Ref}(P(K)). \quad \square$$

Notation for a class \mathcal{C} of relational structures, write $H(\mathcal{C})$ for the class of all homomorphically equivalent structures.

Red(\mathcal{C}) for the class of all pp-reducts A of structures B in \mathcal{C} , i.e. same domain and all relations pp-definable.
I(\mathcal{C}) for the class of structures with a pp-interpretation in a structure from \mathcal{C} .

Remark Chapter 3.6: Given a class \mathcal{C} of structures, $HI(\mathcal{C})$ is the class of structures that can be pp-constructed from \mathcal{C} .

Theorem: Let B, C be at most countable ω -categorical structures and let \subseteq be a polymorphism algebra of C . Then:

- i) $B \in HR\text{ed}(C)$ iff there exists $\underline{B} \in \text{Exp Refl}(C)$ s.t. $\text{Clo}(\underline{B}) = \text{Pol}(B)$
- ii) $B \in HI(C)$ iff there exists $\underline{B} \in \text{Exp Refl}^{\text{fin}}(\subseteq)$ s.t. $\text{Clo}(\underline{B}) = \text{Pol}(B)$

Proof: i) "⇒" Let $C' \in \text{Red}(C)$ s.t. $B \in H(C')$ witnessed by $h: C' \rightarrow B$ as well as $g: B \rightarrow C'$. Set \subseteq' to be the expansion of \subseteq obtained by adding $\text{Pol}(C') \setminus \text{Clo}(\subseteq)$ to the signature, making \subseteq' a polymorphism algebra of C' .

Let \underline{B}' be the reflection of \subseteq' w.r.t. h and g .

All operations on \underline{B}' are obtained as compositions of operations of the form $f^{\underline{B}'}(x_1, \dots, x_n) := h(f^{\subseteq'}(g(x_1), \dots, g(x_n)))$. As h, g are

All operations on \underline{B} are obtained as composition of operations of the form $f^{\underline{B}'}(x_1, \dots, x_n) := h(f^{\underline{C}'}(g(x_1), \dots, g(x_n)))$. As h, g are homomorphisms and $f^{\underline{C}'} \in \text{Pol}(\underline{C})$, they preserve all relations of \underline{B} , which gives $\text{Clo}(\underline{B}') \subseteq \text{Pol}(\underline{B})$.

Extending \underline{B}' by adding $\text{Pol}(\underline{B}) \setminus \text{Clo}(\underline{B}')$ to its signature gives us $\underline{B} \in \text{Exp}(\underline{B}') \subseteq \text{Exp} \text{Refl}(\underline{C}')$

$$\begin{aligned} &\subseteq \text{Exp} \text{Refl} \text{Exp}(\underline{C}) \\ &= \text{Exp} \text{Refl}(\underline{C}) \end{aligned}$$

" $\underline{B} \in H \text{Ref}(\underline{C}) \Leftrightarrow$ there exists $\underline{B} \in \text{Exp} \text{Refl}(\underline{C})$ s.t. $\text{Clo}(\underline{B}) = \text{Pol}(\underline{B})$ "

Suppose that the reflection \underline{B} of \underline{C} w.r.t. $h: C \rightarrow B$ and $g: B \rightarrow C$ is such that $\text{Clo}(\underline{B}) \subseteq \text{Pol}(\underline{B})$.

Set \underline{C}' to be the structure with domain C , the same signature as \underline{B} and for all k -ary rel.symb. R of \underline{B} the relation

$$R^{\underline{C}'} := \left\{ (f(g(b'_1), \dots, g(b'_k)), \dots, f(g(b'_1), \dots, g(b'_k))) \mid k \in \mathbb{N}, f \in \text{Pol}(\underline{C})^{(e)}, \overline{b}_1, \dots, \overline{b}_k \in R^{\underline{B}} \right\}$$

These relations are preserved by $\text{Pol}(\underline{C})$ and therefore pp-def. in $\underline{C} \Rightarrow \underline{C}' \in \text{Ref}(\underline{C})$.

As $\text{Pol}(\underline{C})$ contains the projections, g is a hom from \underline{B} to \underline{C}' .

h is a homomorphism from \underline{C}' to \underline{B} :

Let $f(g(b^1), \dots, g(b^e)) \in R^{\underline{C}'}$ (non-ary case works analogously)

then $g: (x^1, \dots, x^e) \mapsto h(f(g(x^1), \dots, g(x^e)))$ is an operation

on $\underline{B} \in \text{Ref}(\underline{C})$. Since $\text{Clo}(\underline{B}) \subseteq \text{Pol}(\underline{B})$, g is a polymorphism and thus $h(f(g(b^1), \dots, g(b^e))) \in R^{\underline{B}}$.

Thus, we have $\underline{B} \in H(\underline{C}') \subseteq H(\text{Ref}(\underline{C}))$.

" $\underline{B} \in H I(\underline{C}) \Rightarrow$ there exists $\underline{B} \in \text{Exp} \text{Refl}^{\text{fin}}(\underline{C})$ s.t. $\text{Clo}(\underline{B}) = \text{Pol}(\underline{B})$ "

Let $\underline{B} \in H I(\underline{C})$.

$\Rightarrow \exists \underline{D} \in I(\underline{C})$ s.t. $\underline{B} \in H(\underline{D})$

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$$\Rightarrow \exists \underline{D} \in \text{Exp HSP}^{\text{fin}}(\subseteq) \text{ s.t. } \text{Clo}(\underline{D}) = \text{Pol}(\underline{D})$$

$$\stackrel{i)}{\Rightarrow} \exists \underline{B} \in \text{Exp Refl}(\underline{D}) \text{ s.t. } \text{Clo}(\underline{B}) = \text{Pol}(\underline{B})$$

$$\underline{B} \in \text{Exp Refl}(\underline{D}) \subseteq \text{Exp Refl}(\text{Exp HSP}^{\text{fin}}(\subseteq)) \\ = \text{Exp Refl P}^{\text{fin}}(\subseteq).$$

Below is a proof of part ii) of the theorem (both implications), as the " \Leftarrow " seems to be missing in the current version of the book. A different proof, using techniques from [28] and results of the third chapter is given at the end of the document.

Recall theorem 3.6.2, which states that, given a class \mathcal{C} of structures, $\text{HI}(\mathcal{C}) = \text{HRed P}_{\text{full}}^{\text{fin}}(\mathcal{C})$.

Lemma: Let $\underline{A}, \underline{B}$ be at most countable ω -cat. structures, and let \underline{B} be a polymorphism algebra of \underline{B} . Then, the following holds:

$$i) \underline{A} \in \text{P}_{\text{full}}^{\text{fin}}(\underline{B}) \Rightarrow \exists \underline{A} \in \text{Exp P}^{\text{fin}}(\underline{B}): \text{Clo}(\underline{A}) = \text{Pol}(\underline{A})$$

$$ii) \underline{A} \in \text{P}_{\text{full}}^{\text{fin}}(\underline{B}) \Leftarrow \exists \underline{A} \in \text{P}^{\text{fin}}(\underline{B}): \text{Clo}(\underline{A}) = \text{Pol}(\underline{A})$$

Proof (Lemma):

\Leftarrow Suppose that $\underline{A} = \underline{B}^n$. Then, the domain A of \underline{A} is B^n .

- Let R be any n -ary, pp-def. relation on A .

Then, R can be viewed as an nk -ary relation on B via

$$R^B = \{(b_1^1, \dots, b_n^1, \dots, b_1^k, \dots, b_n^k) \mid (b_1^i, \dots, b_n^i) \in R^A\}$$

This relation is pp-def. from \underline{B} :

Let $f^B \in \text{Pol}(\underline{B})^{(e)}$ and $(b_1^j, \dots, b_n^j) \in B^k$, $i = 1, \dots, n$, $j = 1, \dots, k$
s.t. $R^B(b_1^1, \dots, b_n^1), \dots, R^B(b_1^k, \dots, b_n^k)$

By our assumption, R^A is pp-def, and thus f^A , which acts componentwise like f^B , preserves it:

$$R^A(\underbrace{(b_1^1, \dots, b_n^1), \dots, b_1^k, \dots, b_n^k}_{=: b^1}, \dots, R^A(b_1^k, \dots, b_n^k))$$

$$\Rightarrow R^A(f^A(\underbrace{b_1^1, \dots, b_n^1}_{=: b^1}, \dots, b_1^k, \dots, b_n^k), \dots, f^A(b_1^k, \dots, b_n^k)) \\ = (f^B(b_1^1, \dots, b_n^1), \dots, f^B(b_1^k, \dots, b_n^k))$$

$$\begin{aligned}
&= (f^{\text{IB}}(b_1^1, \dots, b_n^1), \dots, f^{\text{IB}}(b_1^k, \dots, b_n^k)) \\
\Leftrightarrow &\text{R}^{\text{B}}(f^{\text{IB}}(b_1^1, \dots, b_n^1), \dots, f^{\text{IB}}(b_1^k, \dots, b_n^k)) \\
\Rightarrow &\text{R}^{\text{B}} \text{ is pp-def from B.}
\end{aligned}$$

- With a similar proof we also see that a relation is pp-def from A if it (as un-ary rel.) is from A.

" \Rightarrow " Let $A \in P_{\text{full}}^{\text{fin}}(\text{IB})$ be an n-th full power. Set $\underline{A} := \underline{B}^n$. Clearly, $\text{Clo}(\underline{A}) \cong \text{Pcl}(A)$, as elements of $\text{Clo}(\underline{A})$ act componentwise \square

We get the following corollary from the proof of ii):

Corollary: Let IB be an at most countable, ω -categorical structure,

and let \underline{B} be a polymorphism algebra of IB . Let $A \in P_{\text{full}}^{\text{fin}}(\underline{B})$. Then A is a polymorphism algebra of a structure $A \in P_{\text{full}}^{\text{fin}}(\text{IB})$.

Proof theorem ii)

$$\begin{aligned}
"\Rightarrow" \text{ BI}(\text{C}) &\Leftrightarrow \text{BI} \in \text{HRel } P_{\text{full}}^{\text{fin}}(\text{C}) \\
&\stackrel{\text{3.6.2}}{\Rightarrow} \exists \underline{B} \in \text{Exp Refl} \mid \text{Exp } P_{\text{full}}^{\text{fin}}(\text{C}) : \text{Clo}(\underline{B}) = \text{Pcl}(\text{BI}). \\
&\quad \stackrel{\text{Lemma}}{\Rightarrow} \exists \underline{B} \in \text{Exp Refl} \mid \text{Exp } P_{\text{full}}^{\text{fin}}(\text{C}) : \text{Clo}(\underline{B}) = \text{Pcl}(\text{BI}).
\end{aligned}$$

$$\text{"}\Leftarrow\text{" } \exists \underline{B} \in \text{Exp Refl} \mid \text{Exp } P_{\text{full}}^{\text{fin}}(\text{C}) : \text{Clo}(\underline{B}) = \text{Pcl}(\text{BI}).$$

$\Rightarrow \exists A \in P_{\text{full}}^{\text{fin}}(\text{C}) : \underline{B} \in \text{Exp Refl}(A)$. By the previous corollary,
 A is a polymorphism algebra of $A \in P_{\text{full}}^{\text{fin}}(\text{C})$.

Thus: $\text{BI} \in \text{HRel } P_{\text{full}}^{\text{fin}}(\text{C}) = \text{BI}(\text{C})$. \square

Corollary: Let B be an at most countable, ω -categorical structure and let \underline{B} be a polymorphism algebra of B . TFAE:

- $\text{HI}(\text{B})$ contains IK_3 ;
- $\text{HT}(\text{B})$ contains all finite structures;

- i) $\text{HI}(\underline{B})$ contains IK_3 ;
- ii) $\text{HI}(\underline{B})$ contains all finite structures;
- iii) $\text{HI}(\underline{B})$ contains $(\{0, 1\}; 1 \text{IN} 3)$;
- iv) $\text{Refl } \text{P}^{\text{fin}}(\underline{B})$ contains an algebra of size ≥ 2 all of whose operations are projections;
- v) $\text{Refl } \text{P}^{\text{fin}}(\underline{B})$ contains for every finite set A an algebra on A all of whose operations are projections.

If these conditions apply, \underline{B} has a finite-signature reduct with an NP-hard CSP.

Proof: i) \Rightarrow ii) $\text{I}(\text{IK}_3)$ contains all finite structures (Cor. 3.2.1)
and $\text{I} \text{H} \text{I}(\underline{B}) \subseteq \text{HI}(\underline{B})$ (Th 3.6.2)

ii) \Rightarrow iii) trivial

iii) \Rightarrow iv) The polymorphisms of $1 \text{IN} 3$ only contain projections.

Theorem $\Rightarrow \exists \underline{A} \in \text{Exp Refl } \text{P}^{\text{fin}}(\underline{B}) : \text{Clo}(\underline{A}) \subseteq \text{Pol}(1 \text{IN} 3)$

iv) \Rightarrow v) Let $\underline{A} \in \text{Ref} \text{P}^{\text{fin}}(\underline{B})$ s.t. $|A| \geq 2$ and $\text{Clo}(\underline{A}) \subseteq \text{Proj}_A$
Th. 6.3.10 $\Rightarrow \text{HSP}^{\text{fin}}(\underline{A})$ contains for every finite set S an
alg. \underline{S} on S s.t. $\text{Clo}(\underline{S}) \subseteq \text{Proj}_S$.

Now $\underline{S} \in \text{HSP}^{\text{fin}}(\underline{A}) \subseteq \text{HSP}^{\text{fin}}(\text{Ref} \text{P}^{\text{fin}}(\underline{B})) \stackrel{\text{Lemma}}{\subseteq} \text{Ref} \text{C P}^{\text{fin}}(\underline{B})$

v) \Rightarrow i) Let $\underline{A} \in \text{Ref} \text{P}^{\text{fin}}(\underline{B})$ s.t. $A = \{0, 1, 2\}$ and

$\text{Clo}(\underline{A}) \subseteq \text{Proj}_A$. Then $\text{Clo}(\underline{A}) \subseteq \text{Pol}(\text{IK}_3)$

$\Rightarrow \exists \underline{C} \in \text{Exp Refl } \text{P}^{\text{fin}}(\underline{B})$ s.t. $\text{Clo}(\underline{C}) = \text{Pol}(\text{IK}_3)$

$\stackrel{\text{Theorem}}{\Rightarrow} \text{IK}_3 \in \text{HI}(\underline{B})$

The hardness follows from Cor. 3.7.1 ($\text{IK}_3 \in \text{HI}(\underline{B}) \Rightarrow$ finite sign.
reduct whose CSP is NP-hard).

□

This concludes the section. A different proof of Thm. ii) is given below.

The following definition stems from "The wonderland of reflections" [28].
The lemma thereafter is a shortened version of Lemma 3.8,

The following definition stems from "The wonderland of reflections" [28].

The lemma thereafter is a shortened version of Lemma 3.8,

Lemma 3.9 and Corollary 3.10 in that same paper;

using our knowledge from section 3.6

Def: Let A, B be relational structures. We say that B is a pp-power of A if B is isomorphic to a structure C with domain A^n for some $n \geq 1$ all of whose relations are pp-definable from A (viewing unary relations on A^n as n -ary relations on A .)

Lemma 1 Let \mathcal{C} be a class of relational structures. Then $H\mathcal{H}(\mathcal{C}) = H\mathcal{P}_{pp}(\mathcal{C})$

Proof: The inclusion " \subseteq " is trivial.

To proof the other inclusion first note that $H\mathcal{H}(C) = H(C)$ for any rel str. C . Therefore, it suffices to show that

$I(\mathcal{C}) \subseteq H\mathcal{P}_{pp}(\mathcal{C})$. Let $B \in I(\mathcal{C})$ as well as $A \in \mathcal{C}$ s.t. $B \in I(A)$. Let $J: A^n \rightarrow B$ be the coordinate map of the pp-interpretation. Now set C to be the rel. str. with domain A^n , the same signature as B and relations $R^C := J^{-1}(R^B)$ in the sense of Def. 2.4.1., which are pp-definable from A .

Set h to be any extension of J to the entirety of A^n and g to be any right-inverse of J . Now $C \in \mathcal{P}_{pp}(A)$ and $h: C \rightarrow B$ as well as $g: B \rightarrow C$ are homomorphisms, which implies $B \in H\mathcal{P}_{pp}(A)$. □

The following fact is part of Prop. 5.6 in [28], and an easy observation when keeping in mind that, given at most countable ω -cat structures A and B as well as a polymorphism algebra \underline{B} of B ,
 $A \in \text{Red}(B) \Leftrightarrow \exists \underline{A} \in \text{Exp}(\underline{B}): \text{Clo}(\underline{A}) = \text{Pol}(A)$.

Lemma 2 Let A, B be at most countable ω -cat. structures, and let \underline{B} be a polymorphism algebra of B . Then, the following holds:

$A \in \mathcal{P}_{pp}(B) \Leftrightarrow \exists \underline{A} \in \text{Exp}(\mathcal{P}^{fin}(\underline{B})): \text{Clo}(\underline{A}) = \text{Pol}(A)$.

" $\underline{B} \in \text{HI}(C) \Leftrightarrow$ there exists $\underline{B} \in \text{Exp Refl } P^{\text{fin}}(\subseteq)$ s.t. $\text{Clo}(\underline{B}) = \text{Pol}(\underline{B})$ "

$\underline{B} \in \text{HI}(C) \Leftrightarrow$ $\underline{B} \in \text{HP}_{\text{pp}}(C)$ Lemma ①

[Keeping in mind that a reduct of a pp-power is a reduct, Part 1)
of the theorem + Lemma ②) show us that

$\Leftrightarrow \exists \underline{B} \in \text{Exp Refl } \text{Exp } P^{\text{fin}}(\subseteq) : \text{Clo}(\underline{B}) \circ \text{Pol}(\underline{B})$
 $\text{Exp Refl } P^{\text{fin}}(\subseteq)$

□