

## Section 6.4 - Reflections

Mittwoch, 21. Februar 2024 15:00

black text: this stems from the book

blue text: deviates from book bc of types etc

**Def:**  $\underline{B}$   $\tau$ -algebra,  $A$  set,  $h: B \rightarrow A$  and  $g: A \rightarrow B$ . The **reflection** of  $\underline{B}$  w.r.t.  $g$  and  $h$  is  $\underline{A}$  where for all  $f \in \tau$  of arity  $n$  and all  $x_1, \dots, x_n \in A$ ,

$$f^{\underline{A}}(x_1, \dots, x_n) := h(f^{\underline{B}}(g(x_1), \dots, g(x_n)))$$

For a class  $\mathcal{C}$  of  $\tau$ -algebras, the class of their reflections is denoted by **Ref $\mathcal{I}$ ( $\mathcal{C}$ )**.

Analogous to the HSP-Lemma, the following holds:

**Lemma**  $\mathcal{C}$  class of  $\tau$ -algebras.

- The smallest class of  $\tau$ -algebras containing  $\mathcal{C}$  and closed under Ref $\mathcal{I}$ , H, S and P is Ref $\mathcal{I}$ P( $\mathcal{C}$ ).
- The smallest class of  $\tau$ -algebras containing  $\mathcal{C}$  and closed under Ref $\mathcal{I}$ , H, S and P<sub>fin</sub> is Ref $\mathcal{I}$ P<sub>fin</sub>( $\mathcal{C}$ ).

**Proof:** We show first statement. Suffices to show closeness of Ref $\mathcal{I}$ P( $\mathcal{C}$ ) under Ref $\mathcal{I}$ , H, S, P.

**Ref $\mathcal{I}$ :** By the definition of Reflection, the following diagram commutes.

$$\begin{array}{ccccc} \textcircled{B} & \xrightarrow{h_1} & \textcircled{A} & \xrightarrow{h_2} & \textcircled{C} \\ \text{Ref}\mathcal{I} \downarrow & & \text{Ref}\mathcal{I} \downarrow & & \downarrow \\ g_1 \uparrow & & g_2 \uparrow & & \\ \textcircled{B} & \xrightarrow{h_2 \circ h_1} & \textcircled{A} & \xrightarrow{g_1 \circ g_2} & \textcircled{C} \end{array}$$

$$\sim h_2(h_1(f^{\underline{A}}(g_1(g_2(x)))))$$

$$\sim (h_2 \circ h_1)(f^{\underline{A}}((g_1 \circ g_2)(x)))$$

Thus, Ref $\mathcal{I}$ (Ref $\mathcal{I}$ ( $\mathcal{K}$ )) = Ref $\mathcal{I}$ ( $\mathcal{K}$ ) for any class  $\mathcal{K}$ .

**H:** Show that H( $\mathcal{K}$ )  $\subseteq$  Ref $\mathcal{I}$ ( $\mathcal{K}$ ) for any class  $\mathcal{K}$ .

Given  $\underline{A} \in H(\mathcal{K})$ , there is a surjective homomorphism  $h$  from some  $\underline{B} \in \mathcal{K}$  to  $\underline{A}$ . Let  $g: A \rightarrow B$  s.t.  $h \circ g = \text{id}_A$ . It holds for  $f \in \tau$  of arity  $n$  and  $x_1, \dots, x_n \in A$  that

$$h(f^{\underline{B}}(g(x_1), \dots, g(x_n))) = f^{\underline{A}}(h(g(x_1)), \dots, h(g(x_n))) = f^{\underline{A}}(x_1, \dots, x_n).$$

Thus,  $\underline{A}$  is the reflection of  $\underline{B}$  w.r.t.  $h$  and  $g$ , so  $\underline{A} \in \text{Ref}\mathcal{I}(\mathcal{K})$ .

**S:** Show that H( $\mathcal{K}$ )  $\subseteq$  Ref $\mathcal{I}$ ( $\mathcal{K}$ ) for any class  $\mathcal{K}$ .

Given  $A \leq B \in \mathcal{C}$ , set  $g: A \rightarrow B$  to be the identity and  $h: B \rightarrow A$  any extension of  $g$  to  $B$ . Then  $\underline{A}$  is the reflection of  $\underline{B}$  w.r.t.  $h$  and  $g$ , analogous to H.

**P:** Show that P(Ref $\mathcal{I}$ P( $\mathcal{K}$ ))  $\subseteq$  Ref $\mathcal{I}$ (P( $\mathcal{K}$ )).

Given a set  $I$ , algebras  $(\underline{B}_i)_{i \in I} \in (P(\mathcal{K}))^I$  and  $\tau$ -algebras  $(A_i)_{i \in I}$  s.t. for all  $i$ ,  $A_i$  is the reflection of  $\underline{B}_i$  w.r.t.  $h_i$  and  $g_i$ .

Then  $h: (b_i)_{i \in I} \mapsto (h_i(b_i))_{i \in I}$  and  $g: (\alpha_i)_{i \in I} \mapsto (g_i(\alpha_i))$  witness

wonderland: for every operation  $f$  of  $\underline{B}$  then apply Ref $\mathcal{I}$  to operation clones.  
Need this at one point in proof of the theorem.

For all  $i$ ,  $A_i$  is the reflection of  $B_i$  w.r.t.  $h_i$  and  $g_i$ .

Then  $h: (b_i)_{i \in I} \mapsto (h_i(b_i))_{i \in I}$  and  $g: (a_i)_{i \in I} \mapsto (g_i(a_i))_{i \in I}$  witness that  $\prod_{i \in I} A_i$  is a reflection of  $\prod_{i \in I} B_i$ , which proves  $P(\text{Refl } P(K)) = \text{Refl } P(K)$ .  $\square$

Notation for a class  $\mathcal{C}$  of relational structures, write  $H(\mathcal{C})$  for the class of all homomorphically equivalent structures.

$\text{Red}(\mathcal{C})$  for the class of all pp-reducts  $A$  of structures  $B$  in  $\mathcal{C}$ , i.e. same domain and all relations pp-definable.

$I(\mathcal{C})$  for the class of structures with a pp-interpretation in a structure from  $\mathcal{C}$ .

**Remark Chapter 3.6:** Given a class  $\mathcal{C}$  of structures,  $HI(\mathcal{C})$  is the class of structures that can be pp-constructed from  $\mathcal{C}$ .

**Theorem:** Let  $B, C$  be at most countable  $\omega$ -categorical structures and let  $\subseteq$  be a polymorphism algebra of  $C$ . Then:

- $B \in H\text{Red}(C)$  iff there exists  $\underline{B} \in \text{Exp Refl}(\subseteq)$  s.t.  $\text{Clo}(\underline{B}) = \text{Pol}(IB)$
- $B \in HI(C)$  iff there exists  $\underline{B} \in \text{Exp Refl } p^{\text{fin}}(\subseteq)$  s.t.  $\text{Clo}(\underline{B}) = \text{Pol}(IB)$

**Proof: i) "⇒"** Let  $C' \in \text{Red}(C)$  s.t.  $IB \in H(C')$  witnessed by  $h: C' \rightarrow B$  as well as  $g: B \rightarrow C'$ . Set  $\subseteq'$  to be the expansion of

  $\subseteq$  obtained by adding  $\text{Pol}(C') \setminus \text{Clo}(\subseteq)$  to the signature, making  $\subseteq'$  a polymorphism algebra of  $C'$ .

Let  $\underline{B}'$  be the reflection of  $\subseteq'$  w.r.t.  $h$  and  $g$ .

All operations on  $\underline{B}'$  are obtained as compositions of operations of the form  $f^{\underline{B}'}(x_1, \dots, x_n) := h(f^{\subseteq'}(g(x_1), \dots, g(x_n)))$ . As  $h, g$  are homomorphisms and  $f^{\subseteq'} \in \text{Pol}(C')$ , they preserve all relations of  $IB$ , which gives  $\text{Clo}(\underline{B}') \subseteq \text{Pol}(IB)$ .

which gives  $\text{Clo}(\underline{B}') \subseteq \text{Pol}(\text{IB})$ .

Extending  $\underline{B}'$  by adding  $\text{Pol}(\text{IB}) \setminus \text{Clo}(\underline{B}')$  to its signature gives us

$$\underline{B} \in \text{Exp}(\underline{B}') \subseteq \text{Exp Refl}(\underline{C}')$$

$$\subseteq \text{Exp Refl}(\text{Exp}(\underline{C}))$$

$$= \text{Exp Refl}(\subseteq)$$

" $\text{B} \in H\text{Red}(\mathcal{C}) \Leftrightarrow$  there exists  $\underline{B} \in \text{Exp Refl}(\underline{C})$  s.t.  $\text{Clo}(\underline{B}) = \text{Pol}(\text{IB})$ "

Suppose that the reflection  $\underline{B}$  of  $\underline{C}$  w.r.t.  $h: C \rightarrow B$  and  $g: B \rightarrow C$  is such that  $\text{Clo}(\underline{B}) \subseteq \text{Pol}(\text{IB})$ .

Set  $\mathcal{C}'$  to be the structure with domain  $C$ , the same signature as  $\text{IB}$  and for all  $n$ -ary rel.symb.  $R$  of  $\text{IB}$  the relation

$$R^{\mathcal{C}'} := \left\{ (f(g(b_1'), \dots, g(b_n')), \dots, f(g(b_1'), \dots, g(b_n'))) \mid \right.$$

$$\left. l \in \mathbb{N}, f \in \text{Pol}(\mathcal{C})^{(l)}, \overline{b}_1', \dots, \overline{b}_n' \in R^{\underline{B}} \right\}$$

These relations are preserved by  $\text{Pol}(\mathcal{C})$  and therefore pp-def. in  $\mathcal{C} \Rightarrow \mathcal{C}' \in \text{Red}(\mathcal{C})$ .

As  $\text{Pol}(\mathcal{C})$  contains the projections,  $g$  is a hom from  $\text{IB}$  to  $\mathcal{C}'$ .

$h$  is a homomorphism from  $\mathcal{C}'$  to  $\text{IB}$ :

Let  $f(g(b_1'), \dots, g(b_n')) \in R^{\mathcal{C}'}$  (non-ary case works analogously)

then  $g: (x_1', \dots, x_n') \mapsto h(f(g(x_1'), \dots, g(x_n')))$  is an operation

on  $\underline{B} \in \text{Ref}(C)$ . Since  $\text{Clo}(\underline{B}) \subseteq \text{Pol}(\text{IB})$ ,  $g$  is

a polymorphism and thus  $h(f(g(b_1'), \dots, g(b_n'))) \in R^{\underline{B}}$ .

Thus, we have  $\text{B} \in H(\mathcal{C}') \subseteq H(\text{Red}(\mathcal{C}))$ .

ii) " $\text{B} \in HI(\mathcal{C}) \Rightarrow$  there exists  $\underline{B} \in \text{Exp Refl Pfin}(\underline{C})$  s.t.  $\text{Clo}(\underline{B}) = \text{Pol}(\text{IB})$ "

Let  $\text{B} \in HI(\mathcal{C})$ .

$\Rightarrow \exists \text{ D} \in I(\mathcal{C})$  s.t.  $\text{IB} \in H(\text{D})$

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$\Rightarrow \exists \underline{D} \in \text{Exp HCPfin}(\underline{C})$  s.t.  $\text{Clo}(\underline{D}) = \text{Pol}(\text{D})$

$$\stackrel{i)}{\Rightarrow} \exists \underline{B} \in \text{ExpRef}(D) \text{ s.t. } \text{Clo}(\underline{B}) = \text{Pol}(\underline{B})$$

$$\underline{B} \in \text{ExpRef}(D) \subseteq \text{ExpRef}(\text{ExpHSP}^{\text{fin}}(\subseteq)) \\ = \text{ExpRef}(\text{P}^{\text{fin}}(\subseteq)).$$

ii) " $\underline{B} \in \text{HI}(C) \Leftrightarrow \text{there exists } \underline{B} \in \text{ExpRef}(\text{P}^{\text{fin}}(\subseteq)) \text{ s.t. } \text{Clo}(\underline{B}) = \text{Pol}(\underline{B})$ "

This part can currently only be found in "Graph Homomorphisms and Universal Algebra".

Suppose that there exists  $\underline{B} \in \text{ExpRef}(\text{P}^{\text{fin}}(\subseteq))$  s.t.  $\text{Clo}(\underline{B}) = \text{Pol}(\underline{B})$ .

$$\Rightarrow \exists \underline{D} \in \text{P}^{\text{fin}}(\subseteq) \text{ s.t. } \underline{B} \in \text{ExpRef}(\underline{D}). \text{ Let } n \in \mathbb{N} \text{ s.t. } \underline{D} = \underline{C}^n$$

Set  $\text{ID}$  to be the structure with domain  $\underline{C}^n$  with relations  $\text{Inv}(\text{Clo}(\underline{D}))$ . It follows that  $\text{ID}$  has an  $n$ -dimensional pp-interpretation in  $C$  and by Prop. 6.1.5 and \*,  $\text{Pol}(\text{ID}) = \text{Clo}(\underline{D})$ .

$$\text{By i), } \underline{B} \in \text{HRed}(\text{ID}) = \text{HRed}(\text{P}^{\text{fin}}(C)) \subseteq \text{HI}(C).$$

\* Note that as  $\text{Clo}(C)$  is locally closed (by Cor. 6.1.6), so is  $\text{Clo}(\text{ID})$ : □

$$f \in \overline{\text{Clo}(\text{ID})}^{(n)} \Rightarrow \forall F \subseteq \text{ID}^n \text{ finite } \exists g \in \text{Clo}(\text{ID})^{(n)}: f|_F = g|_F.$$

$g$  acts componentwise like some  $\tilde{g} \in \text{Clo}(C)^{(n)}$

$\Rightarrow f$  acts componentwise and the same for all components.

$\Rightarrow \forall i \in \{1, \dots, n\}: \pi_i \circ f \circ \left[ (x_1, \dots, x_n) \mapsto ((x_1, \dots, x_n)_1, \dots, (x_1, \dots, x_n)_n) \right]$   
is equal to some  $\tilde{h} \in \text{Clo}(C)^{(n)}$  as  $\text{Clo}(C)$  is closed.

$\Rightarrow f$  acts like  $\tilde{h}$  on all components.

**Corollary:** Let  $\underline{B}$  be an at most countable,  $\omega$ -categorical structure and let  $\underline{B}$  be a polymorphism algebra of  $(\mathbb{Q}, \text{TFAE})$ :

- i)  $\text{HI}(\underline{B})$  contains  $\text{K}_3$ ;
- ii)  $\text{HI}(\underline{B})$  contains all finite structures;
- iii)  $\text{HI}(\underline{B})$  contains  $(\{0, 1\}; 1 \mid N 3)$ ;

- iii)  $\text{HI}(\underline{B})$  contains  $(\{0,1\}; 1 \text{IN} 3)$ ;
- iv)  $\text{RefI } \text{P}^{\text{fin}}(\underline{B})$  contains an algebra of size  $\geq 2$  all of whose operations are projections;
- v)  $\text{RefI } \text{P}^{\text{fin}}(\underline{B})$  contains for every finite set  $A$  an algebra on  $A$  all of whose operations are projections.

If these conditions apply,  $\underline{B}$  has a finite-signature reduct with an NP-hard CSP.

**Proof:** i)  $\Rightarrow$  ii)  $\text{I}(\underline{K}_3)$  contains all finite structures (Cor. 3.2.1) and  $\text{I} \cap \text{I}(\underline{B}) \subseteq \text{HI}(\underline{B})$  (Th 3.6.2)

ii)  $\Rightarrow$  iii) trivial

iii)  $\Rightarrow$  iv) The polymorphisms of  $1 \text{IN} 3$  only contain projections.

Theorem  $\Rightarrow \exists \underline{A} \in \text{Exp RefI } \text{P}^{\text{fin}}(\underline{B}) : \text{Clo}(\underline{A}) = \text{Pol}(1 \text{IN} 3)$

iv)  $\Rightarrow$  v) Let  $\underline{A} \in \text{RefI } \text{P}^{\text{fin}}(\underline{B})$  s.t.  $|A| \geq 2$  and  $\text{Clo}(\underline{A}) \subseteq \text{Proj}_{\underline{A}}$ .  
Th.6.3.10  $\Rightarrow \text{HSP}^{\text{fin}}(\underline{A})$  contains for every finite set  $S$  an alg.  $\underline{S}$  on  $S$  s.t.  $\text{Clo}(\underline{S}) \subseteq \text{Proj}_S$ .

Now  $\underline{S} \in \text{HSP}^{\text{fin}}(\underline{A}) \subseteq \text{HSP}^{\text{fin}}(\text{RefI } \text{P}^{\text{fin}}(\underline{B})) \stackrel{\text{Lemma}}{\subseteq} \text{RefI } \text{P}^{\text{fin}}(\underline{B})$

v)  $\Rightarrow$  i) Let  $\underline{A} \in \text{RefI } \text{P}^{\text{fin}}(\underline{B})$  s.t.  $A = \{0, 1, 2\}$  and  $\text{Clo}(\underline{A}) \subseteq \text{Proj}_{\underline{A}}$ . Then  $\text{Clo}(\underline{A}) \subseteq \text{Pol}(\underline{K}_3)$   
 $\Rightarrow \exists \underline{C} \in \text{Exp RefI } \text{P}^{\text{fin}}(\underline{B})$  s.t.  $\text{Clo}(\underline{C}) = \text{Pol}(\underline{K}_3)$   
 $\stackrel{\text{Theorem}}{\Rightarrow} \underline{K}_3 \in \text{HI}(\underline{B})$

The hardness follows from Cor. 3.7.1 ( $\underline{K}_3 \in \text{HI}(\underline{B}) \Rightarrow$  finite sign. reduct whose CSP is NP-hard). □