

# Chapter 10: File System Interface

肖卿俊

办公室：九龙湖校区计算机楼212室

电邮：[csqjxiao@seu.edu.cn](mailto:csqjxiao@seu.edu.cn)

主页：<https://csqjxiao.github.io/PersonalPage>

电话：025-52091022



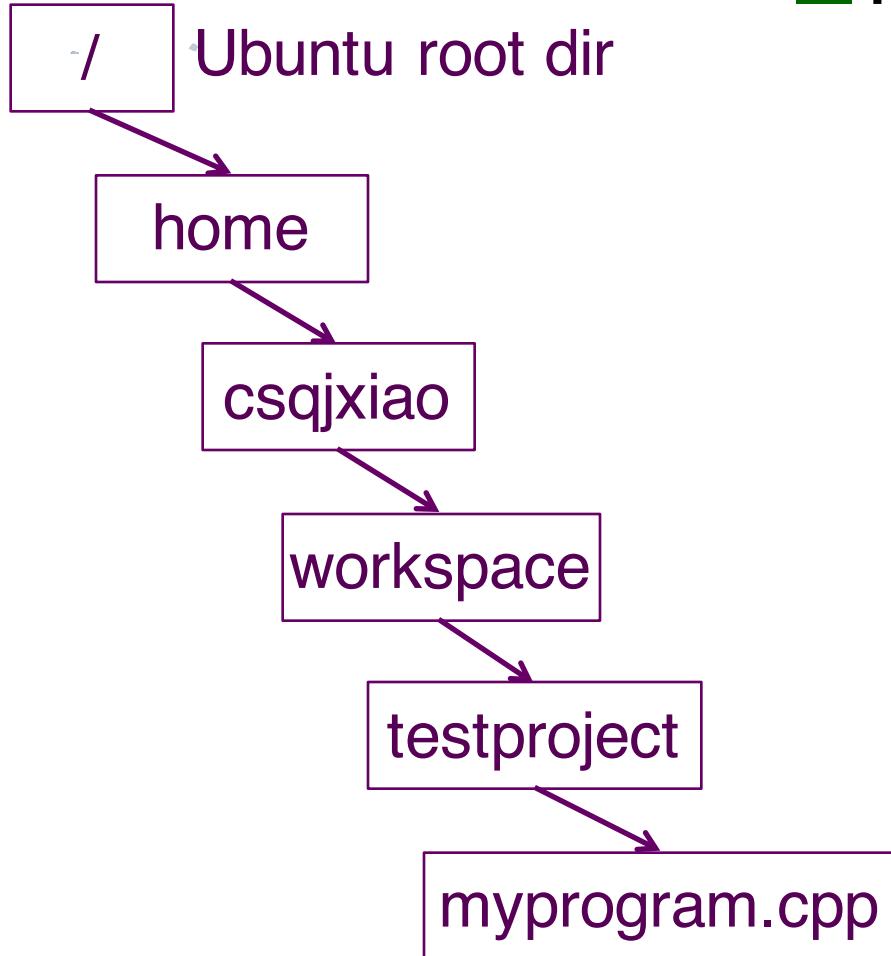
# Chapter 10: File-System Interface

- File Concept
- Access Methods
- Directory Structure
- File System Mounting
- File Sharing
- Protection



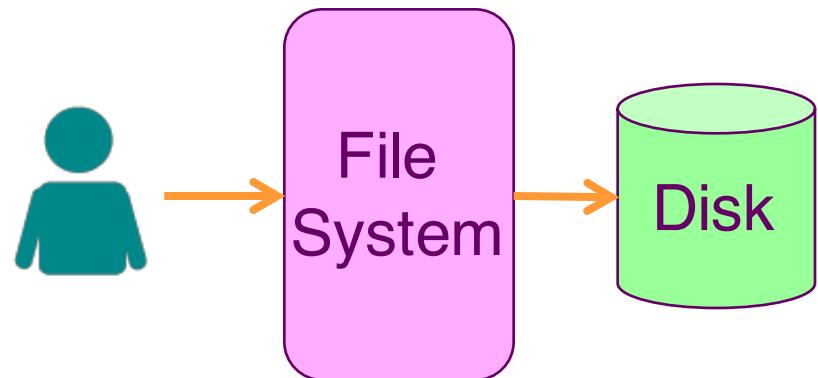


# File System Concept

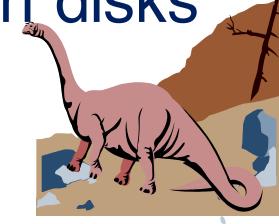


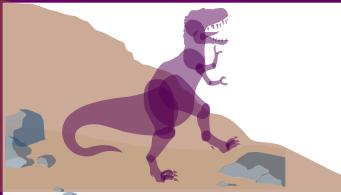
## ■ Key Abstraction

- ◆ File
- ◆ Filename
- ◆ Directory tree (folders)

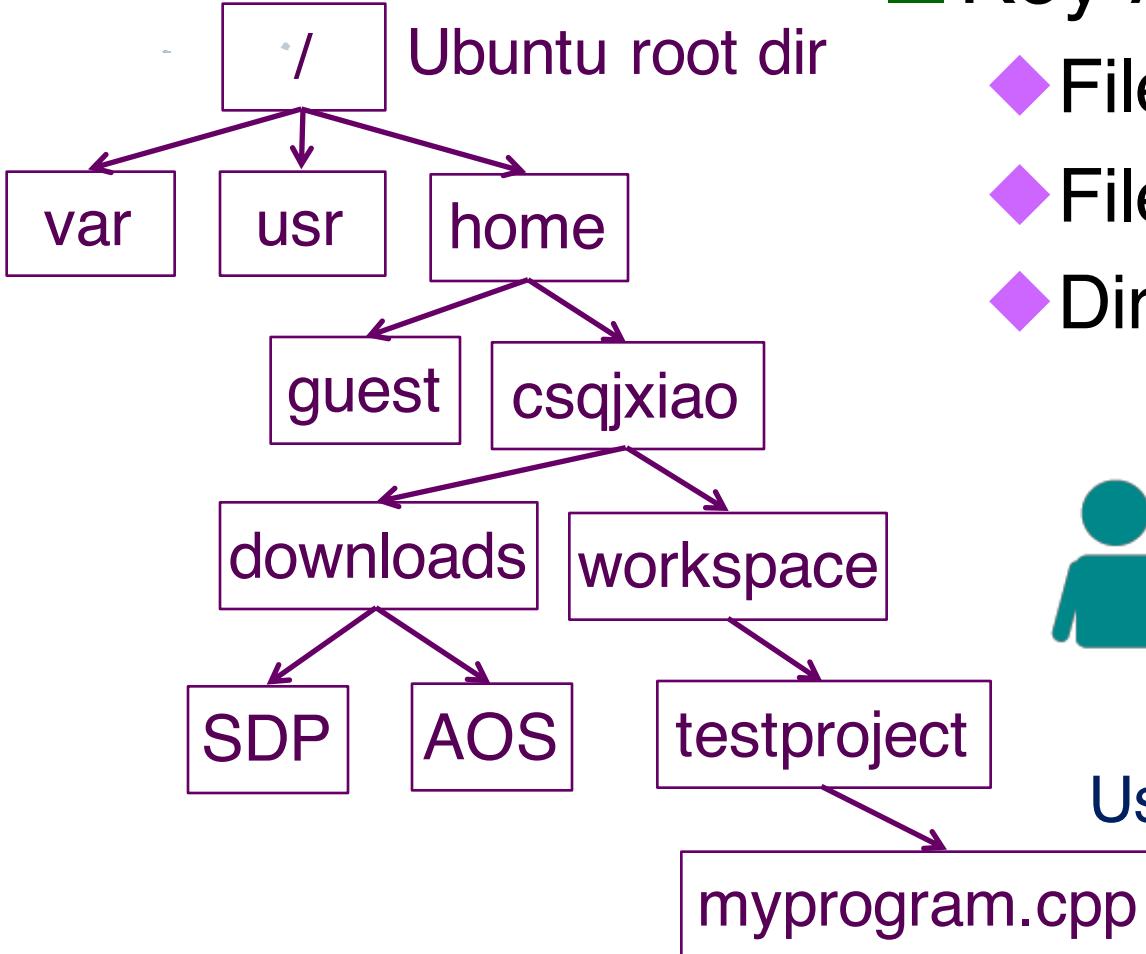


Users do not access directly  
the file blocks on disks



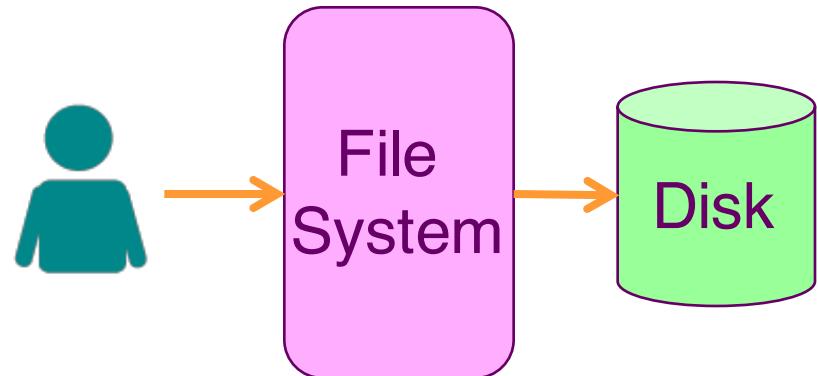


# File Path and Directory Tree

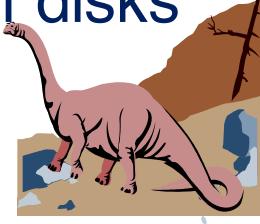


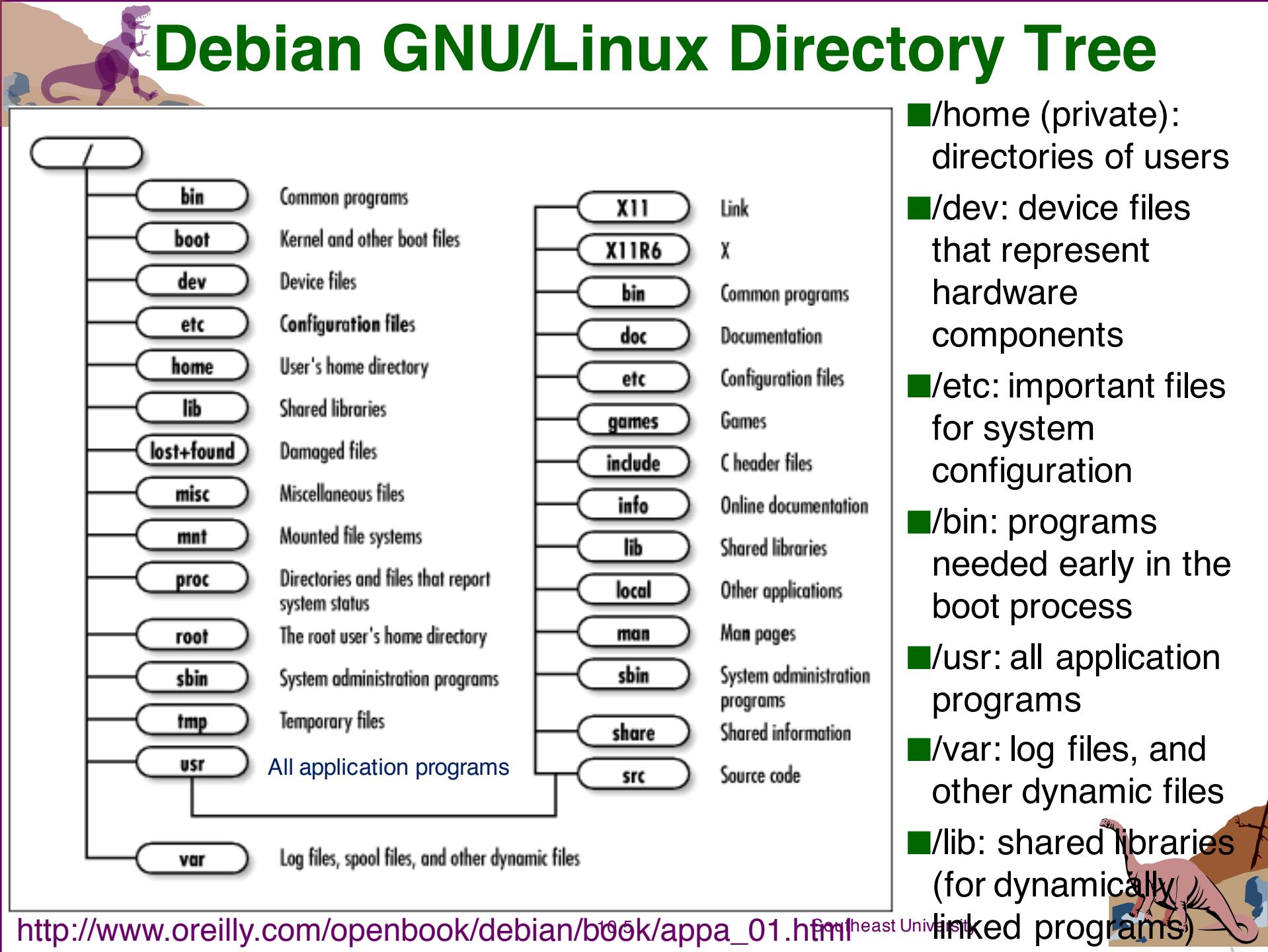
## Key Abstraction

- ◆ File
- ◆ Filename
- ◆ Directory tree (folders)



Users do not access directly  
the file blocks on disks







# File Concept

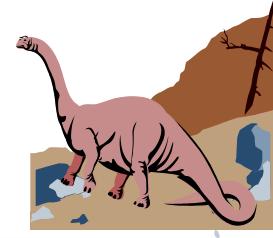
- Contiguous logical address space

- Types:

- ◆ Data

- ✓ numeric
    - ✓ character
    - ✓ binary

- ◆ Program





# File Structure

- None - sequence of words, bytes
- Simple record structure
  - ◆ Lines
  - ◆ Fixed length
  - ◆ Variable length
- Complex Structures
  - ◆ Formatted document
  - ◆ Relocatable load file
- Can simulate last two with the first method by inserting appropriate control characters.
- Who decides the file structure?



# File Attributes

- **Name** – only information kept in human-readable form.
- **Type** – needed for systems that support different types.
- **Location** – pointer to file location on device.
- **Size** – current file size.

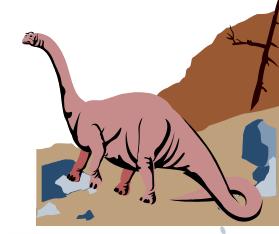
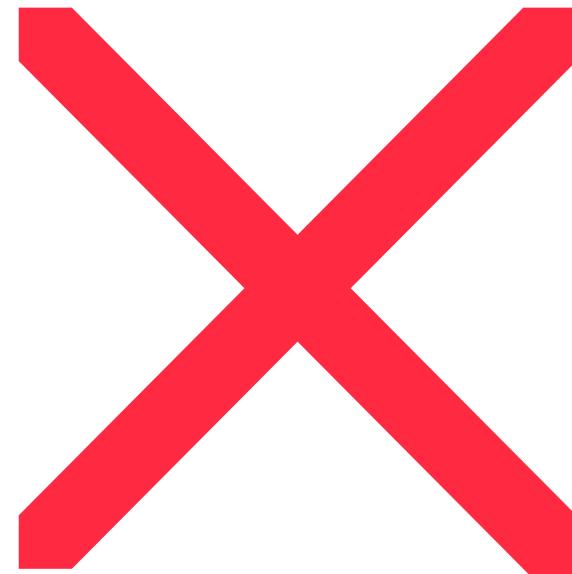


# File Attributes (Cont.)

- **Protection** – controls who can do reading, writing, executing.
- **Time, date, and user identification** – data for protection, security, and usage monitoring.
- All these information about files are kept in the directory structure, which is maintained on the disk.



# File Types – Name, Extension





# File Operations from Developer's Perspective

- Create
- Write
- Read
- Reposition within file – file seek
- Delete
- Truncate





# File Operations from Developer's Perspective (cont.)

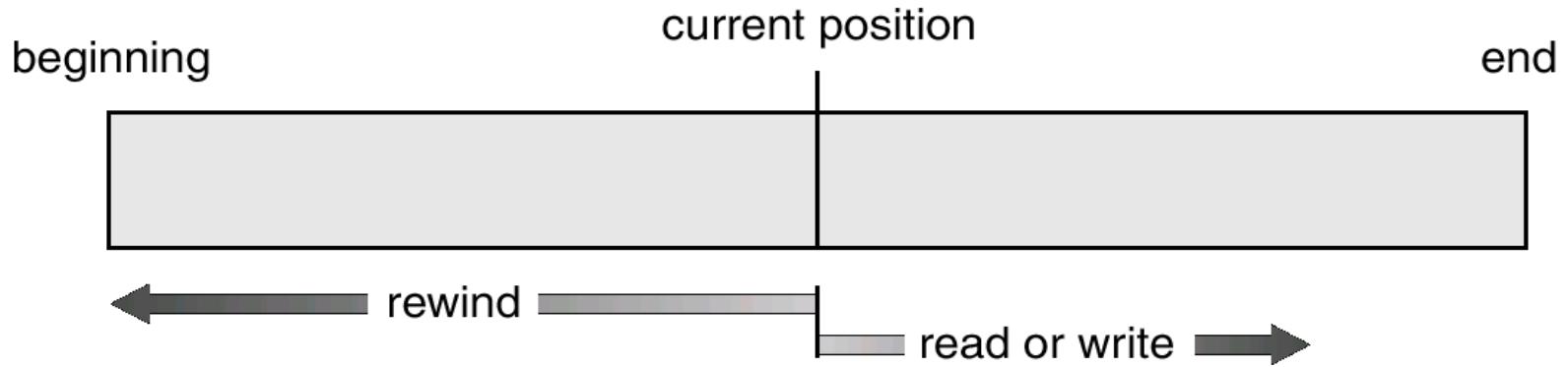
- $\text{open}(F_i)$  – search the directory structure on disk for entry  $F_i$ , and move the content of the entry from disk to memory.
- $\text{close}(F_i)$  – persist the content of entry  $F_i$  in memory to directory structure on disk.
- $\text{read}(F_i)$  – read the file content
- $\text{write}(F_i)$  – write to the file
- $\text{fseek}(F_i)$  – reposition the file cursor



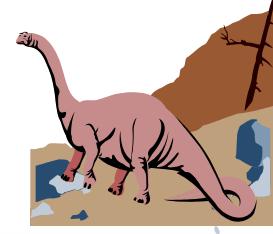


# File Content Access Methods

## ■ Sequential Access



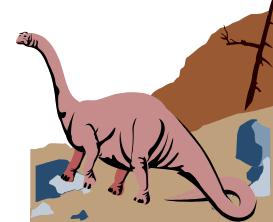
## ■ Direct Access





# Simulation of Sequential Access on a Direct-Access File

sequential access	implementation for direct access
<i>reset</i>	$cp = 0;$
<i>read next</i>	$read cp;$ $cp = cp+1;$
<i>write next</i>	$write cp;$ $cp = cp+1;$

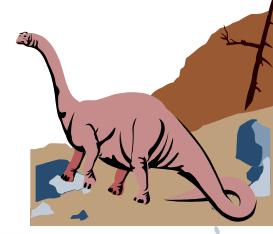




# Code Modifying a Key-Value Pair

KEY	VALUE
integer	integer

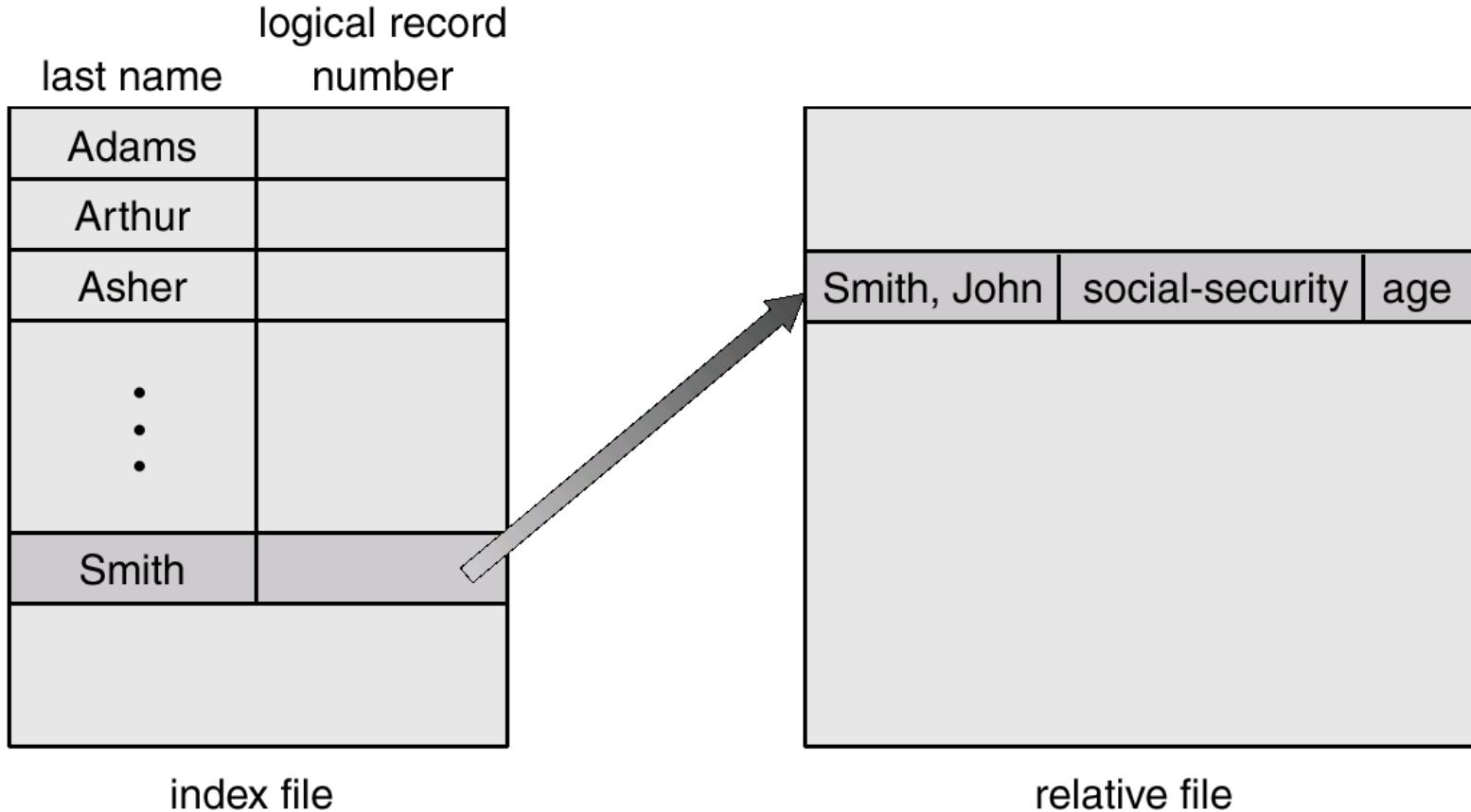
```
ssize_t len;
char * filename;
int key, srch_key, new_value;
filename = argv[1];
srch_key = strtol(argv[2], NULL, 10);
new_value = strtol(argv[3], NULL, 10);
int fd = open(filename, O_RDWR);
while(sizeof(int) == read(fd, key, sizeof(int))) {
    if(key != srch_key)
        lseek(fd, sizeof(int), SEEK_CUR);
    else {
        write(fd, &new_value, sizeof(int));
        close(fd);
        return EXIT_SUCCESS;
    }
}
fprintf(stderr, "key not found!");
return EXIT_FAILURE;
```





# Example of Index and Relative Files

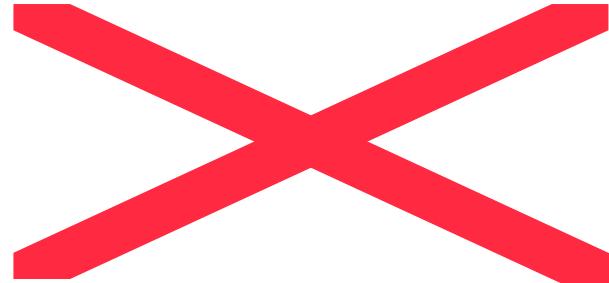
- Store keys in the index file
- Store values (or records) in the relative file
- How to quickly locate the record of John Smith





# B+ Tree for Index File

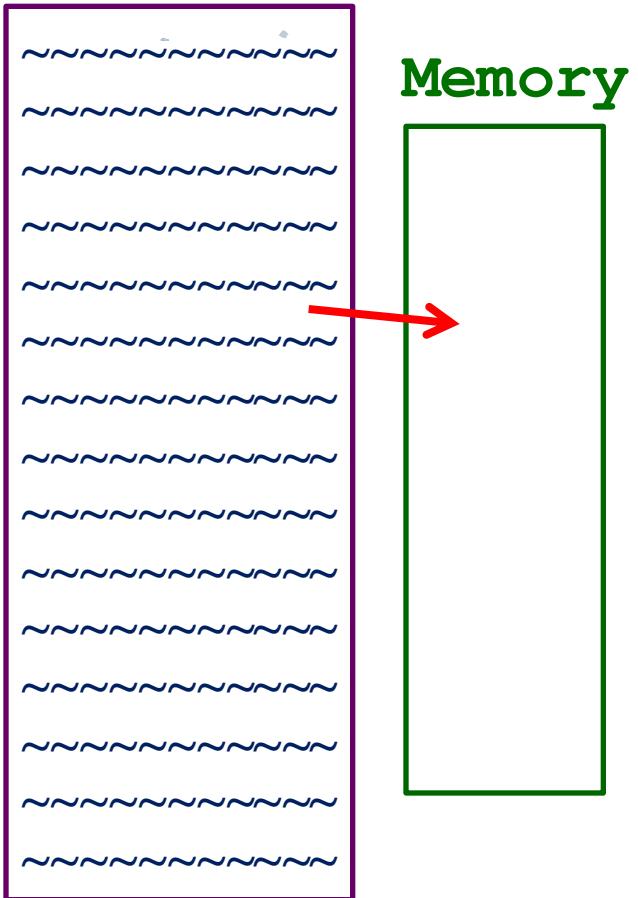
- The primary value of a B+ tree is in storing data for efficient retrieval in a block-oriented storage context – in particular, filesystems. Unlike binary search trees, B+ trees have very high fanout (number of pointers to child nodes in a node, typically on the order of 100 or more), which reduces the number of I/O operations required to find an element in the tree.





# File Content Direct Access by Memory Mapped File

File.txt



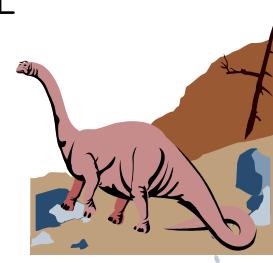
- **mmap()** creates a new mapping in the virtual address space of the calling process
- **munmap()** system call deletes the mappings for the specified address range, and causes further references to addresses within the range to generate invalid memory references

```
fd = open("file.txt", ...);  
buffer = mmap(..., fd, ...);  
  
// manipulate the buffer
```

```
munmap(buf, ...);
```

```
close(fd);
```

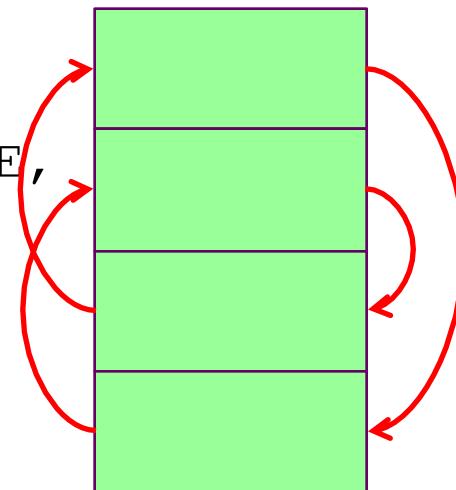
Southeast University





# An Example of Memory Mapped File: Shuffle Blocks within a File

```
filename = argv[1];
card_size = strtol(argv[2], NULL, 10);
fd = open(filename, O_RDWR);
len = lseek(fd, 0, SEEK_END);
lseek(fd, 0, SEEK_SET);
buf = mmap(NULL, len, PROT_READ | PROT_WRITE,
           MAP_FILE | MAP_SHARED, fd, 0);
if( buf == (void*) -1) {
    fprintf(stderr, "mmap failed.\n");
    exit(EXIT_FAILURE);
}
memshuffle(buf, len, card_size);
munmap(buf, len);
close(fd);
return EXIT_SUCCESS;
```



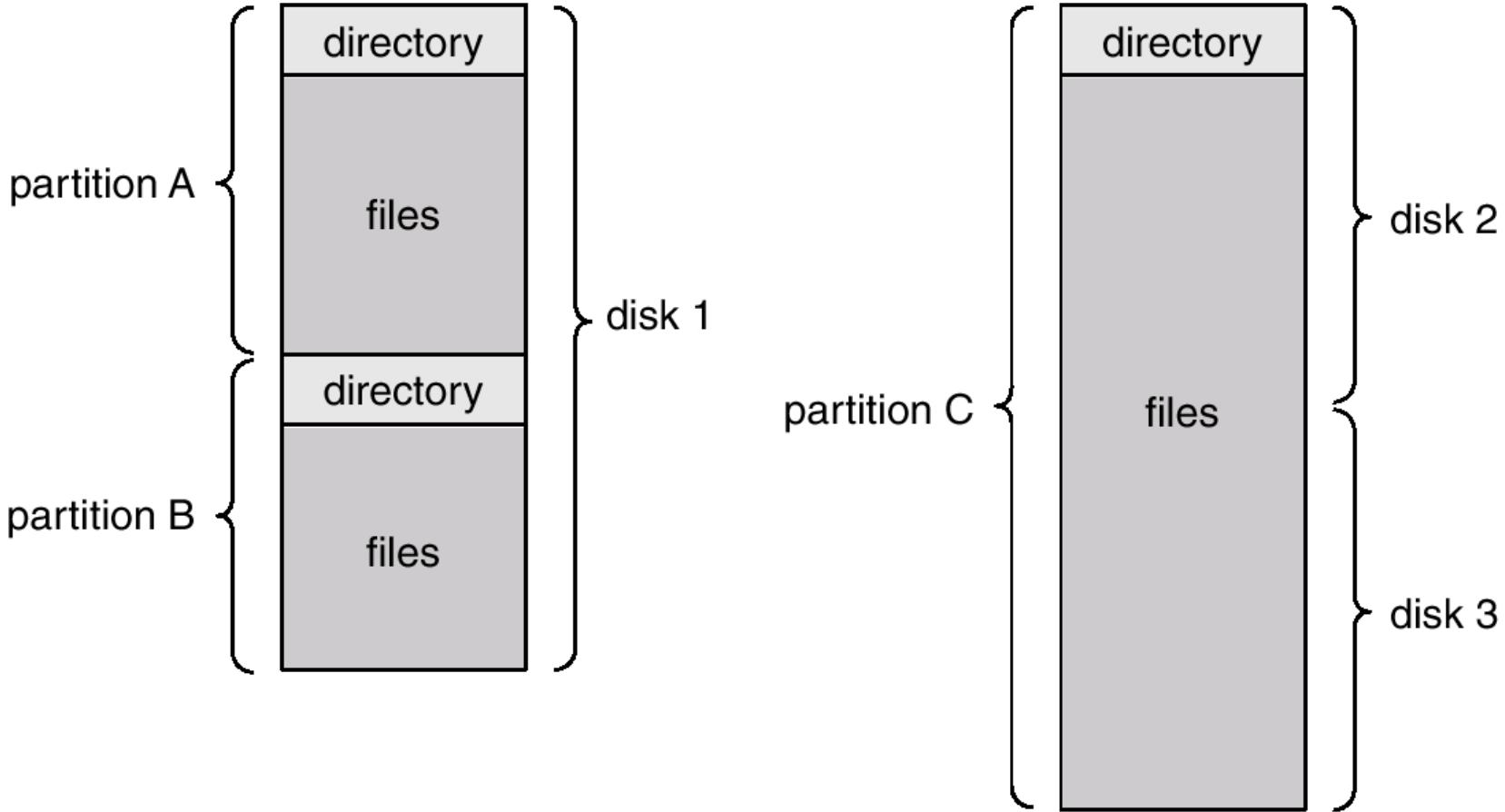


# Directory Structure

- disks are split into one or more partitions.
- each partition contains information about files within it
- The information is kept in entries in a device directory or volume table of contents



# A Typical File-system Organization





# Operations Performed on Directory

- Search for a file
- Create a file
- Delete a file
- List a directory
- Rename a file
- Traverse the file system



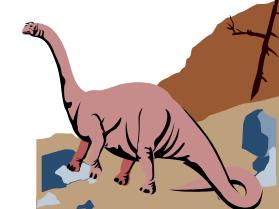
# Organize the Directory (Logically) to Obtain

■ **Efficiency** – locating a file quickly.

■ **Naming** – convenient to users.

- ◆ Two users can have the same name for different files.
- ◆ The same file can have several different names.

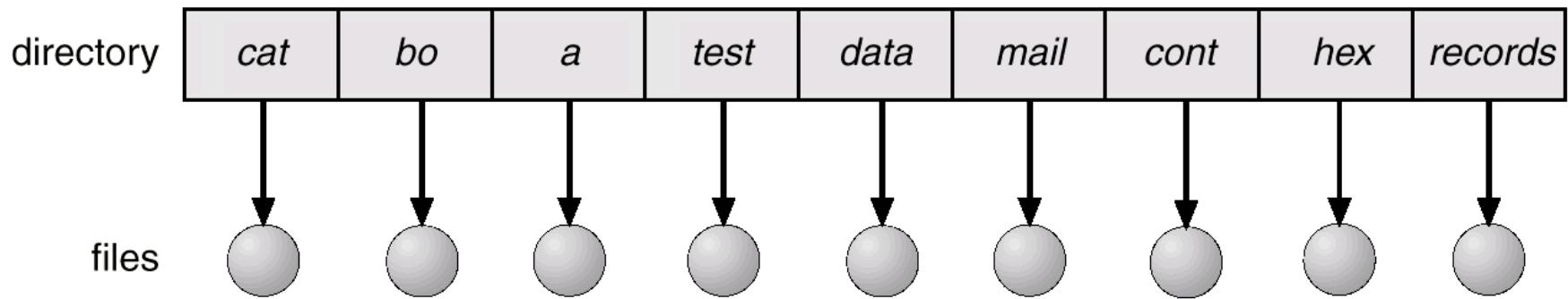
■ **Grouping** – logical grouping of files by properties, (e.g., all Java programs, all games, ...)





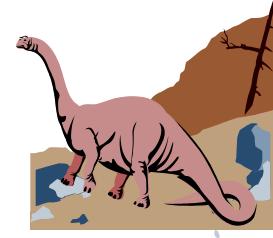
# Single-Level Directory

- A single directory for all users.



Naming problem

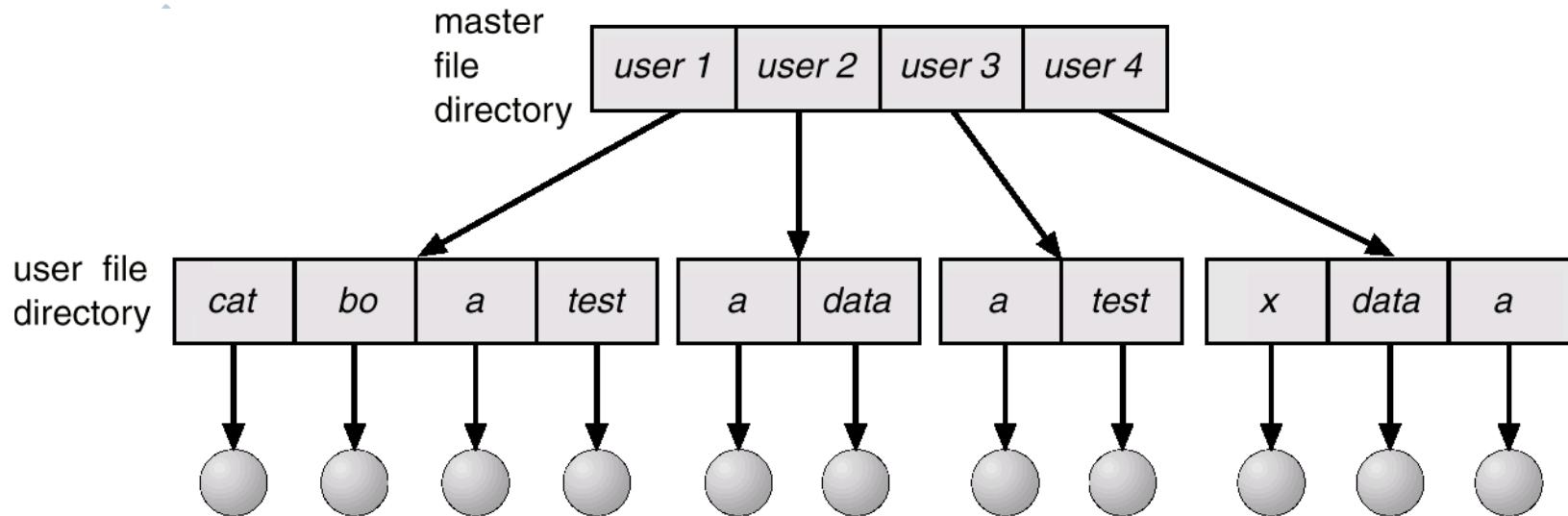
Grouping problem



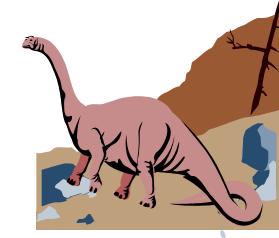


# Two-Level Directory

- Separate directory for each user.

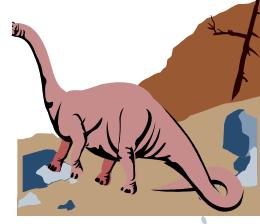
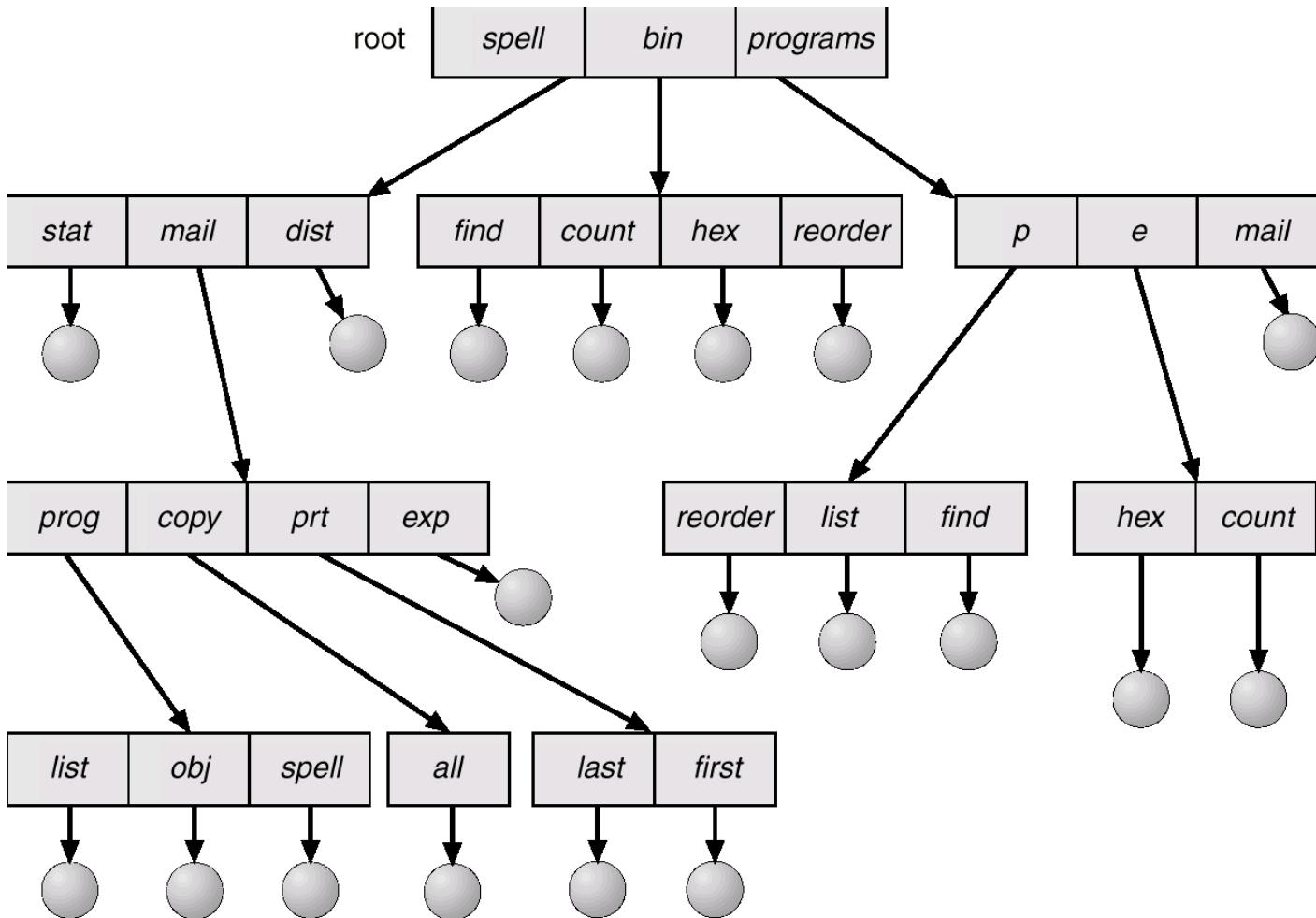


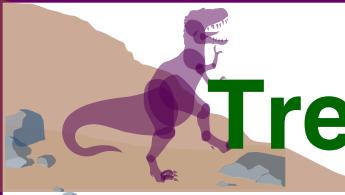
- Efficient searching
- Support path name, so can have the same file name for different users
- No grouping capability





# Tree-Structured Directories





# Tree-Structured Directories (cont.)

## ■ Efficient searching

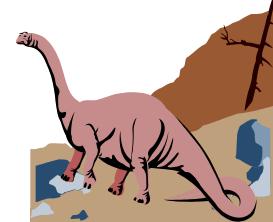
## ■ Convenient naming

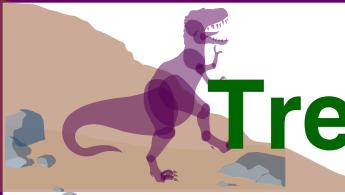
- ◆ Two users can have the same name for different files.

## ■ Grouping capability

## ■ Current directory (working directory)

- ◆ `cd /spell/mail/prog`
- ◆ `type list`





# Tree-Structured Directories (cont.)

- Absolute or relative path name
- Creating a new file can be done in current directory.
- Delete a file

**rm <file-name>**





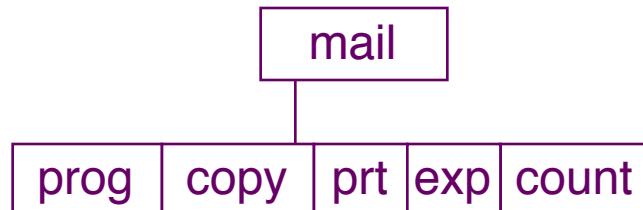
# Tree-Structured Directories (cont.)

- Creating a new subdirectory is done in current directory.

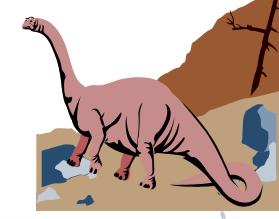
**mkdir <dir-name>**

Example: if in current directory **/mail**

**mkdir count**



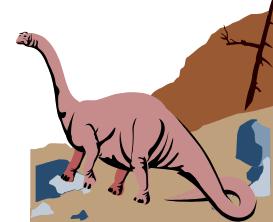
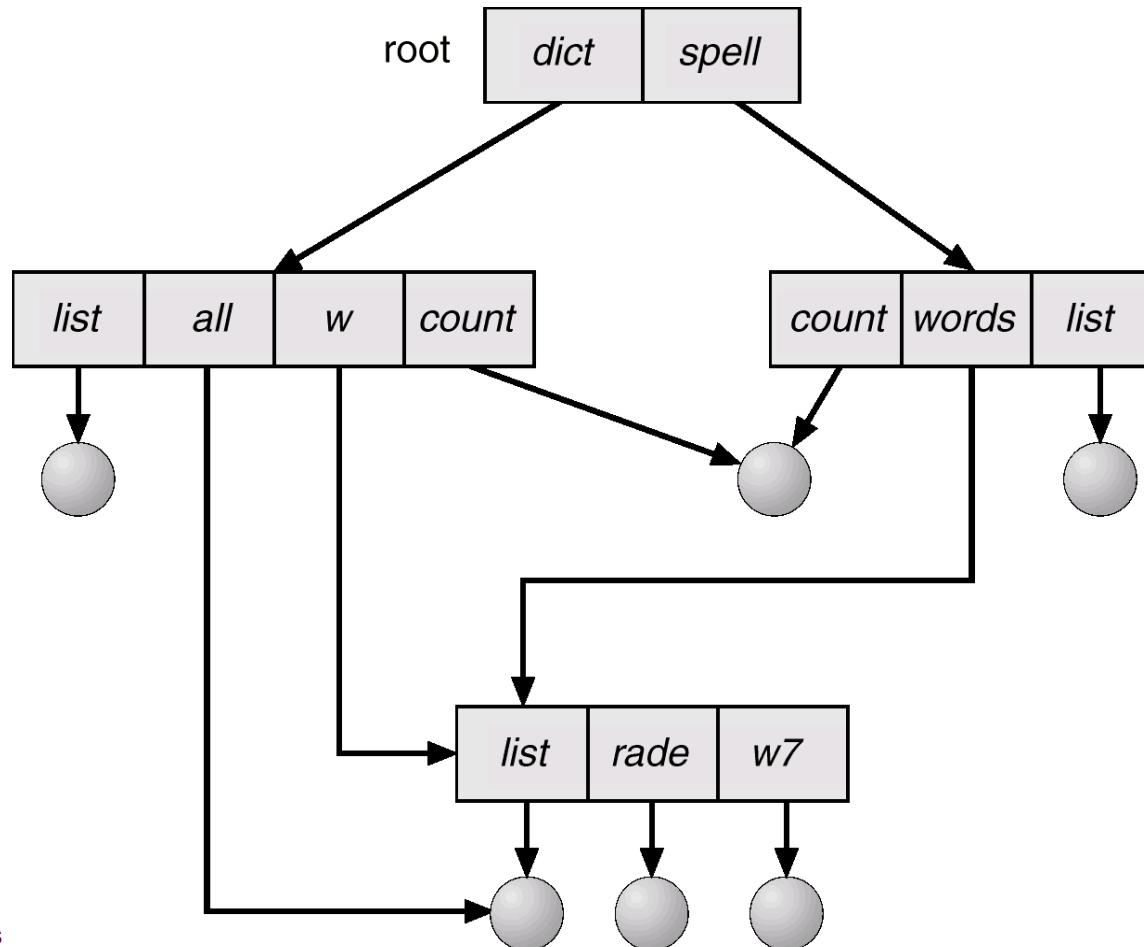
Deleting “mail” ⇒ deleting the entire subtree rooted by “mail”.





# Acyclic-Graph Directories

- Have shared subdirectories and files.  
The same file can have several different paths.





# Acyclic-Graph Directories (cont.)

- Two different names (aliasing)
- If *dict* deletes *count* ⇒ dangling pointer.

Solutions:

- ◆ Backpointers, so we can delete all pointers.
- ◆ Entry-hold-count solution.

- These links we talked about are hard links in UNIX/Linus





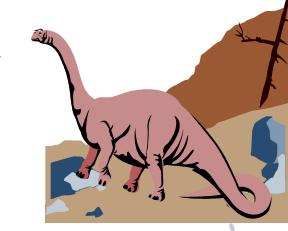
# In Linux/Unix Shortcuts are known as Link

## ■ Soft Links (symbolic links )

- ◆ You can make a link for either a file or a folder
- ◆ You can create link (shortcut) on different partition
- ◆ You got a different inode number from original.
- ◆ If real copy is deleted the link will not work.

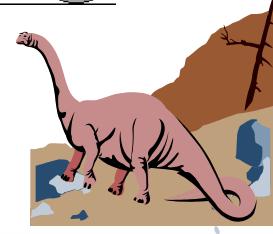
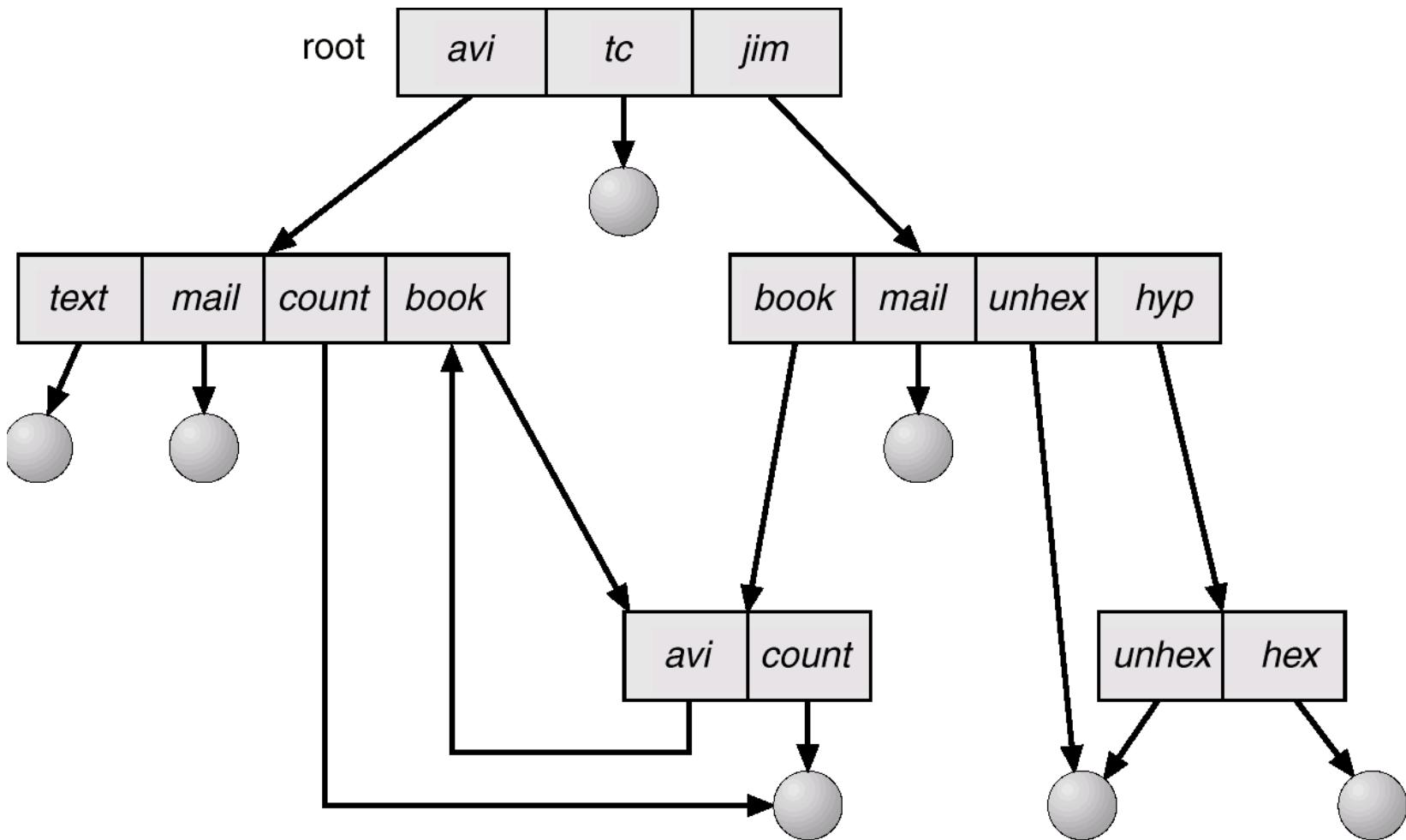
## ■ Hard Links

- ◆ For files only, and you cannot create a hard link on different partition (it should be on same partition)
- ◆ You got the same inode number as original
- ◆ If the real copy is deleted the link will work (because it act as original file )





# General Graph Directory





# General Graph Directory (cont.)

## ■ How do we guarantee no cycles?

- ◆ Allow only links to file not subdirectories.
- ◆ Garbage collection.
- ◆ Every time a new link is added, use a cycle detection algorithm to determine whether it is OK.





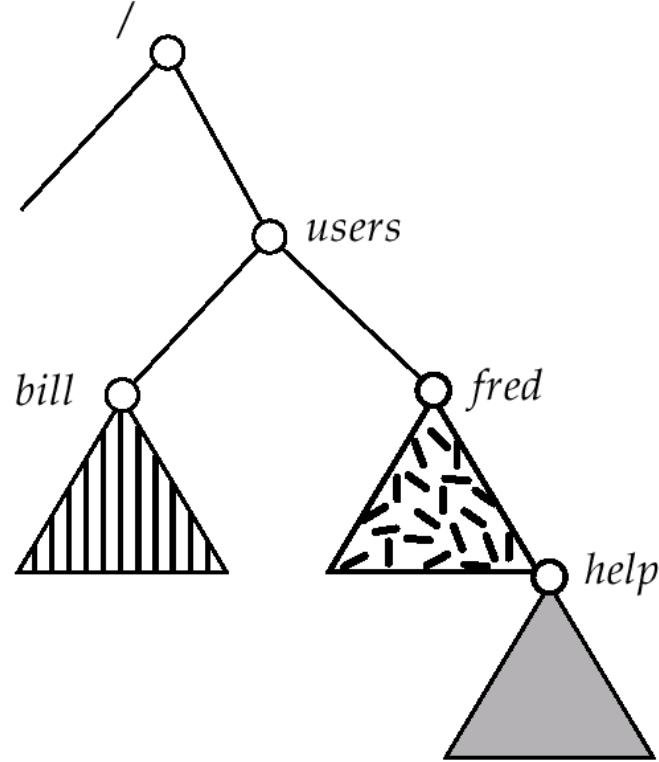
# File System Mounting

- A file system must be **mounted** before it can be accessed.
- An unmounted file system (I.e. Fig. 11-11(b)) is mounted at a **mount point**.

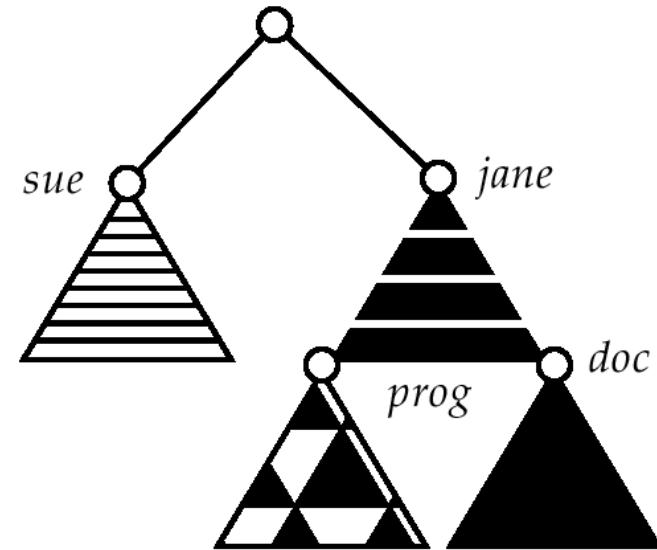




# (a) Existing (b) Unmounted Partition

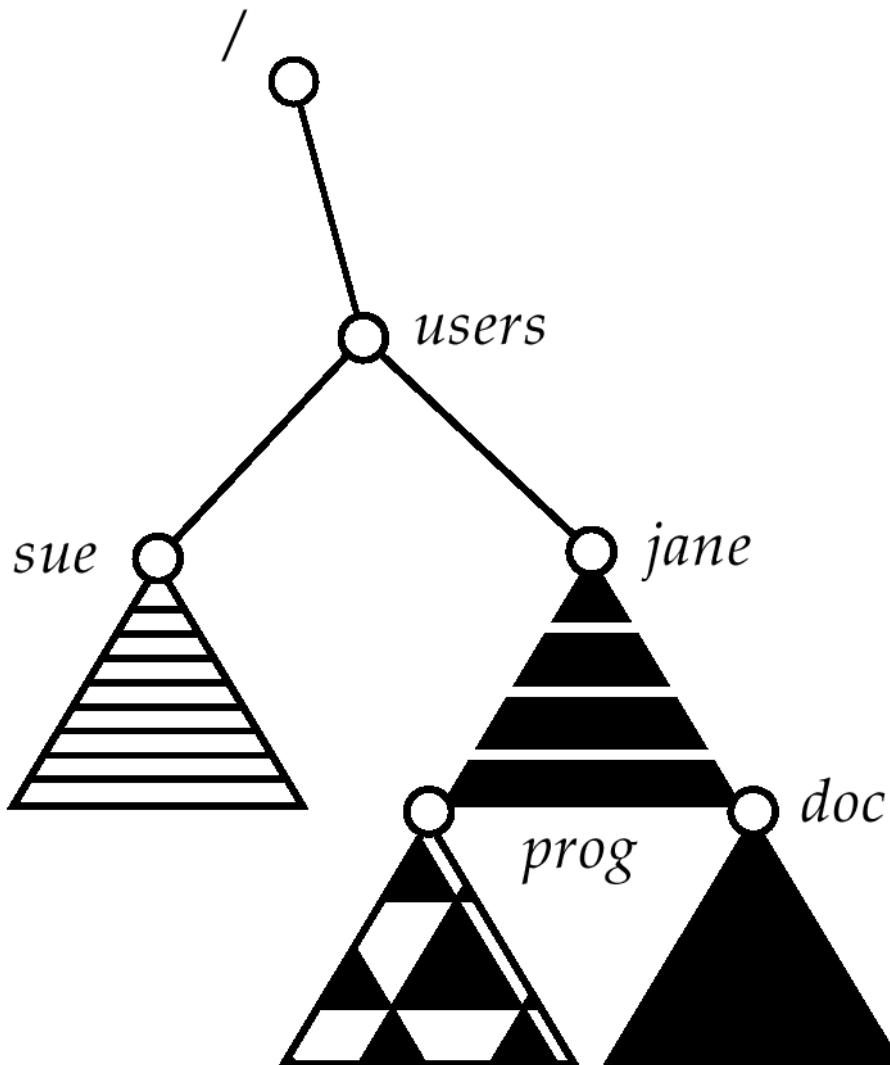


(a)



(b)

# Mount Point





# File Sharing

- Sharing of files on multi-user systems is desirable.
- Sharing may be done through a *protection* scheme.
- On distributed systems, files may be shared across a network.
- Network File System (NFS) is a common distributed file-sharing method.



# Protection

■ File owner/creator should be able to control:

- ◆ what can be done
- ◆ by whom

■ Types of access

- ◆ Read
- ◆ Write
- ◆ Execute
- ◆ Append
- ◆ Delete
- ◆ List



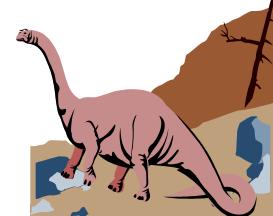
# Access Lists and Groups

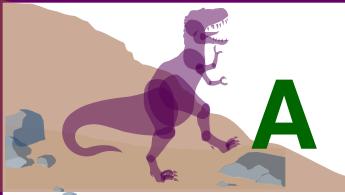
- Mode of access: read, write, execute
- Three classes of users RWX
  - a) **owner access** 7  $\Rightarrow$  1 1 1
  - b) **group access** 6  $\Rightarrow$  1 1 0
  - c) **public access** 1  $\Rightarrow$  0 0 1
- Ask manager to create a group (unique name), say G, and add some users to the group.
- For a particular file (say *game*) or subdirectory, define an appropriate access.

Attach a group to a file

owner    group    public  
              |  
              chgrp  
              |  
              761  
              |  
              game

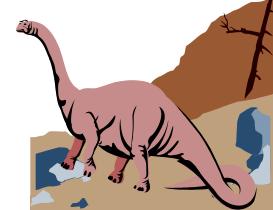
chgrp    G    game





# A Sample UNIX Directory Listing

-rw-rw-r--	1	pbg	staff	31200	Sep 3 08:30	intro.ps
drwx-----	5	pbg	staff	512	Jul 8 09.33	private/
drwxrwxr-x	2	pbg	staff	512	Jul 8 09:35	doc/
drwxrwx---	2	pbg	student	512	Aug 3 14:13	student-proj/
-rw-r--r--	1	pbg	staff	9423	Feb 24 2003	program.c
-rwxr-xr-x	1	pbg	staff	20471	Feb 24 2003	program
drwx--x--x	4	pbg	faculty	512	Jul 31 10:31	lib/
drwx-----	3	pbg	staff	1024	Aug 29 06:52	mail/
drwxrwxrwx	3	pbg	staff	512	Jul 8 09:35	test/





# Windows XP Access-Control List Management

10.tex Properties

General Security Summary

Group or user names:

- Administrators (PBG-LAPTOP\Administrators)
- Guest (PBG-LAPTOP\Guest)**
- pbg (CTI\pbg)
- SYSTEM
- Users (PBG-LAPTOP\Users)

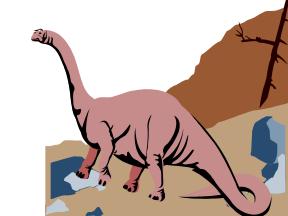
Add... Remove

Permissions for Guest	Allow	Deny
Full Control	<input type="checkbox"/>	<input checked="" type="checkbox"/>
Modify	<input type="checkbox"/>	<input checked="" type="checkbox"/>
Read & Execute	<input type="checkbox"/>	<input checked="" type="checkbox"/>
Read	<input type="checkbox"/>	<input checked="" type="checkbox"/>
Write	<input type="checkbox"/>	<input checked="" type="checkbox"/>
Special Permissions	<input type="checkbox"/>	<input type="checkbox"/>

For special permissions or for advanced settings, click Advanced.

Advanced

OK Cancel Apply





# Question about File Access-Control

■ Which of the following will generate a permission error?

- cat foo.txt
- cat dir/bar.txt
- touch dir/new.txt

```
$ ls -l ./
```

<b>Permission</b>	<b>user</b>	<b>group</b>	....	<b>Filename</b>
drw-r--r--	me	me		dir
-rw-r--r--	other	other		foo.txt

```
$ sudo ls -l dir
```

<b>Permission</b>	<b>user</b>	<b>group</b>	....	<b>Filename</b>
-rw-r--r--	me	me		bar.txt





# Another Question

■ Which of the following will generate a permission error?

- cat foo.txt
- cat dir/bar.txt
- touch dir/new.txt

```
$ ls -l ./
```

<b>Permission</b>	<b>user</b>	<b>group</b>	....	<b>Filename</b>
d--xr--r--	me	me		dir
-rw-r--r--	other	other		foo.txt

```
$ sudo ls -l dir
```

<b>Permission</b>	<b>user</b>	<b>group</b>	....	<b>Filename</b>
-rw-r--r--	me	me		bar.txt

