

Chapter 6: Process Synchronization

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Chapter 6: Process Synchronization

Background

The Critical-Section Problem

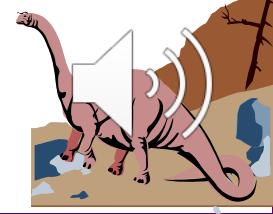
Synchronization Hardware

Semaphores

Classical Problems of Synchronization

Conditional Variables and Monitors

Synchronization Examples





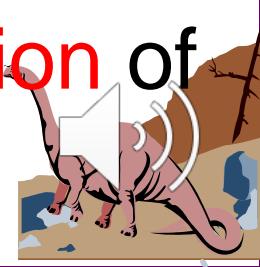
Background

Concurrent access to shared data may result in data inconsistency.

Let us recall the concept of **race condition**

- ◆ Several processes (threads) access and manipulate the same data concurrently and the outcome of the execution depends on the particular order in which the access takes place.

Maintaining **data consistency** requires mechanisms to ensure the **orderly execution** of cooperating processes.





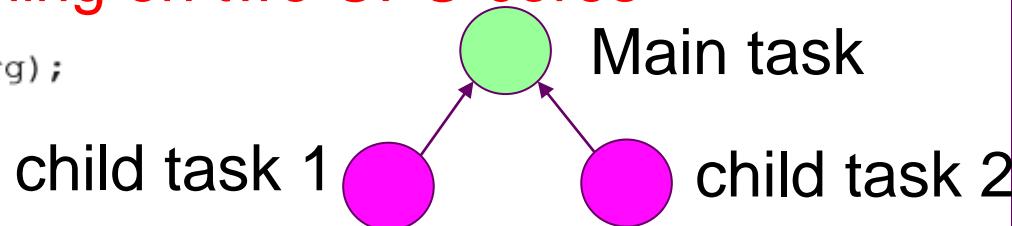
A Previously Used Example

// volatile keyword forces the compiler to avoid caching the variable in CPU register. It always firstly read the data from memory.

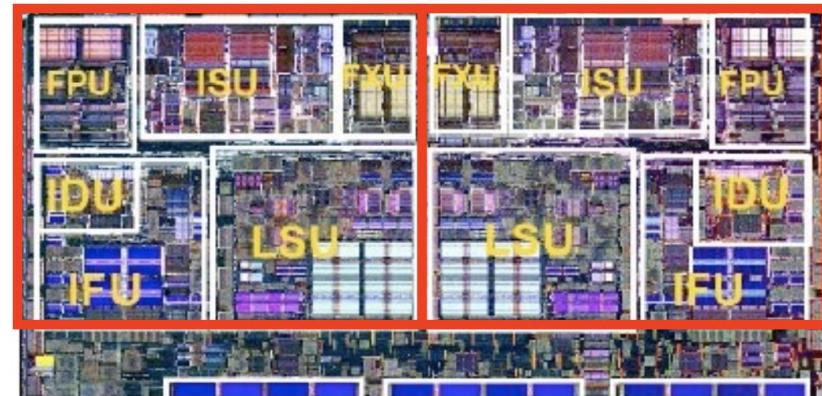
```
volatile int counter = 0;
```

```
14 void *  
15 mythread(void *arg)  
16 {  
17     printf("%s: begin\n", (char *) arg);  
18     int i;  
19     for (i = 0; i < 1e7; i++) {  
20         counter = counter + 1;  
21     }  
22     printf("%s: done\n", (char *) arg)  
23     return NULL;  
24 }
```

We split the for-loop across two threads running on two CPU cores

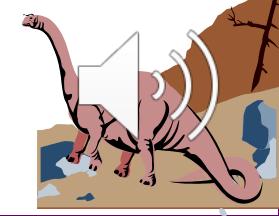


#1



#2

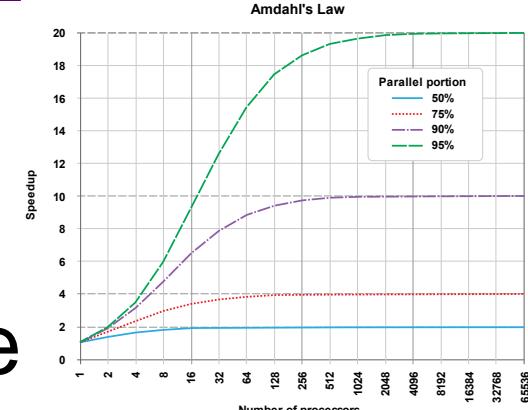
```
32 int  
33 main(int argc, char *argv[])  
34 {  
35     pthread_t p1, p2;  
36     printf("main: begin (counter = %d)\n", counter);  
37     Pthread_create(&p1, NULL, mythread, "A");  
38     Pthread_create(&p2, NULL, mythread, "B");  
39  
40     // join waits for the threads to finish  
41     Pthread_join(p1, NULL);  
42     Pthread_join(p2, NULL);  
43     printf("main: done with both (counter = %d)\n", counter);  
44     return 0;  
45 }
```





How much faster?

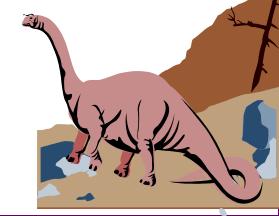
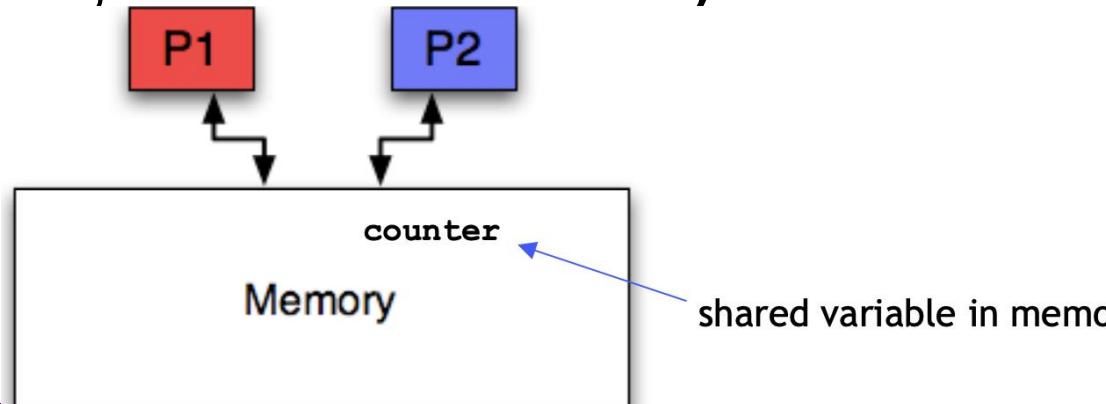
We're expecting a speedup of 2
OK, perhaps a little less because
of Amdahl's Law (to predict the theoretical
speedup when using multiple processors)



- ◆ overhead for forking and joining multiple threads

But it is actually slower!! Why??

Here's the mental picture that we have – two
processors, shared memory

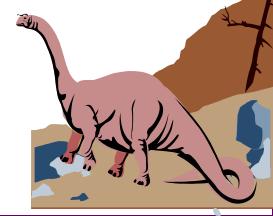
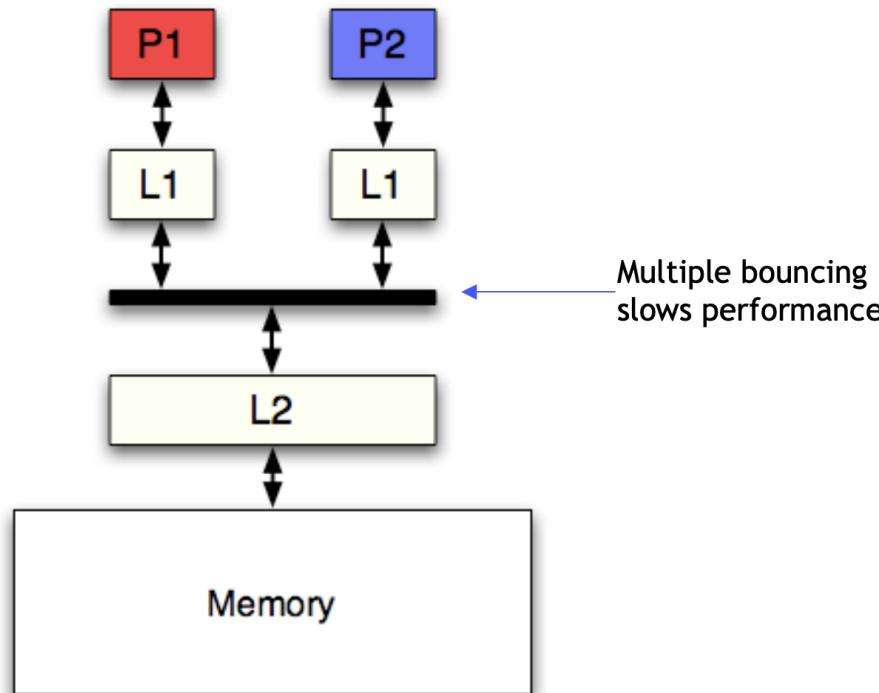




This mental picture is wrong!

We've forgotten about caches!

- ◆ The memory may be shared, but each processor has its own L1 cache
- ◆ As each processor updates *counter*, it bounces between L1 caches





The code is not only slow, it's WRONG!

Since the variable *counter* is shared, we can get a data race
Increment operation: *counter++*

MIPS equivalent:

```
lw $t0, counter  
addi $t0, $t0, 1  
sw $t0, counter
```

A data race occurs when data is accessed and manipulated by multiple processors, and the outcome depends on the sequence or timing of these events.

Sequence 1

Processor 1

```
lw $t0, counter  
addi $t0, $t0, 1  
sw $t0, counter
```

Processor 2

```
lw $t0, counter  
addi $t0, $t0, 1  
sw $t0, counter
```

Sequence 2

Processor 1

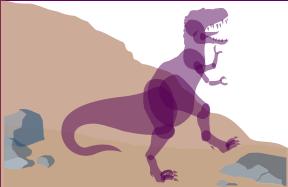
```
lw $t0, counter  
addi $t0, $t0, 1  
sw $t0, counter
```

Processor 2

```
lw $t0, counter  
addi $t0, $t0, 1  
sw $t0, counter
```

counter increases by 2

counter increases by 1 !!



Another Example: Revisit the Producer Consumer Problem

Shared-memory solution to bounded-buffer problem (Chapter 3)

- ◆ The code can only use N-1 items in the buffer

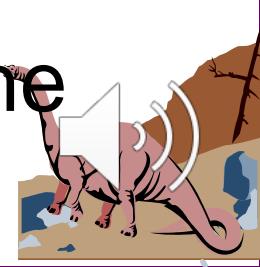
Producer:

```
while (1) {  
    while (((in+1) % BUF_SIZE) == out) ;  
    .....  
    in = (in+1) % BUF_SIZE;  
}
```

Consumer:

```
while (1) {  
    while (in == out) ;  
    .....  
    out = (out+1) % BUF_SIZE;  
}
```

We modify the above code by adding a variable *counter*, such that all items in the buffer can be used

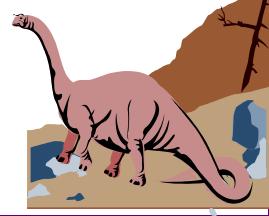




Bounded-Buffer Solution

Shared data

```
#define BUF_SIZE 10
class Item {
    ...
    Item & operator=(const Item & ) { ... }
}
Item buffer[BUF_SIZE];
int in = 0;
int out = 0;
int counter = 0;
```





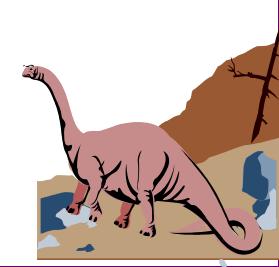
Bounded-Buffer Solution

Producer process

```
Item nextProduced;  
  
while (1) {  
    while (counter == BUF_SIZE)  
        ; /* do nothing */  
    buffer[in] = nextProduced;  
    in = (in + 1) % BUF_SIZE;  
    counter++;  
}
```

Consumer process

```
Item nextConsumed;  
  
while (1) {  
    while (counter == 0)  
        ; /* do nothing */  
    nextConsumed = buffer[out];  
    out = (out + 1) % BUF_SIZE;  
    counter--;  
}
```





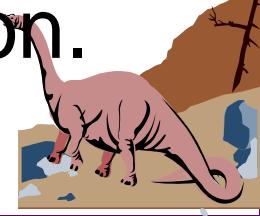
Critical Shared Data

Counter is a piece of critical shared data
The statements

counter++;
counter--;

must be performed *atomically*.

Atomic operation means an operation that completes in its entirety without interruption.





Difficult to Implement the Atomic Guarantee

However, the statement “**count++**” may be implemented in machine language as:

register1 = counter

register1 = register1 + 1

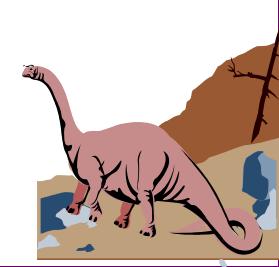
counter = register1

The statement “**count--**” may be implemented in machine language as:

register2 = counter

register2 = register2 - 1

counter = register2

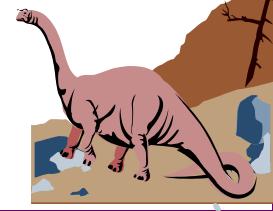




Potential Data Inconsistency

If both the producer and consumer attempt to update the buffer concurrently, the assembly language statements may get interleaved.

Interleaving depends upon how the producer and consumer processes are scheduled.





Potential Data Inconsistency

Assume **counter** is initially 5. One interleaving of statements is:

producer: **register1 = counter** (*register1 = 5*)

producer: **register1 = register1 + 1** (*register1 = 6*)

consumer: **register2 = counter** (*register2 = 5*)

consumer: **register2 = register2 - 1** (*register2 = 4*)

producer: **counter = register1** (*counter = 6*)

consumer: **counter = register2** (*counter = 4*)

The value of **count** may be either 4 or 6, where the correct result should be 5.





Producer

register1 = counter

register1 = register1 + 1

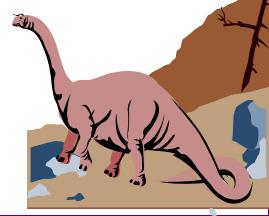
counter = register1

Consumer

register2 = counter

register2 = register2 – 1

counter = register2



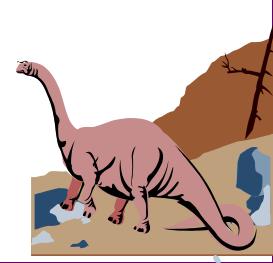


Concept of Race Condition

Race condition occurs, if:

- ◆ **two or more** processes/threads access and manipulate the **same** data **concurrently**, and
- ◆ the outcome of the execution **depends on the particular order** in which the access takes place.

To prevent race conditions, concurrent processes must be **synchronized**.





Chapter 6: Process Synchronization

Background

The Critical-Section Problem

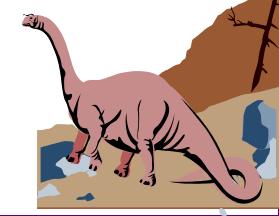
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Synchronization Examples





Three Typical Mechanisms of Process Synchronization

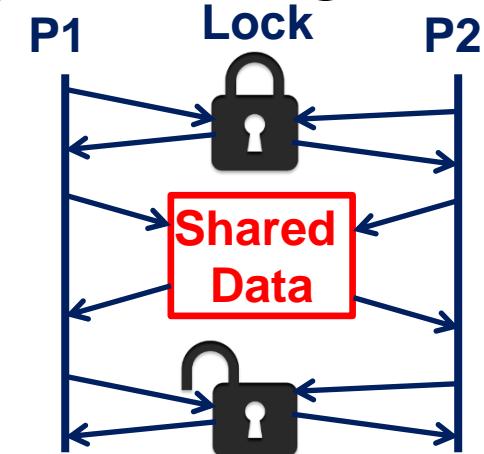
Locks for shared memory programming

- ◆ Exclusive Lock

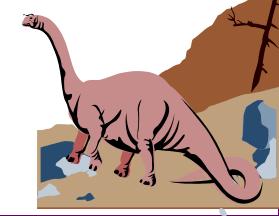
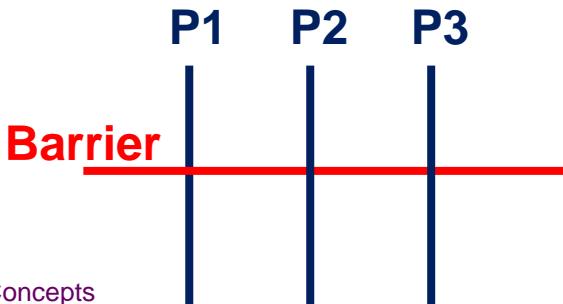


- ◆ Shared Lock:

- ✓ Readers can share a lock, but readers and writers must access the data mutually exclusively



There are other synchronization primitives for shared memory programming, e.g., Barrier





OS Support to Implement an Exclusive Lock for Threads

Using Mutex: Sleeping of threads Instead of spinning that wastes CPU cycles

APIs of POSIX Thread for exclusive lock

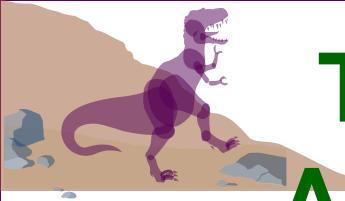
```
int pthread_mutex_lock(pthread_mutex_t* mutex)  
int pthread_mutex_unlock(pthread_mutex_t* mutex)
```

An Example

```
pthread_mutex_t lock = PTHREAD_MUTEX_INITIALIZER;  
pthread_mutex_lock(&lock);  
counter = counter+1; // or whatever your critical section is  
pthread_mutex_unlock(&lock);
```

Give a demonstration



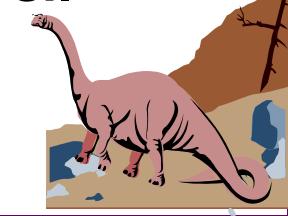


The Critical-Section Problem: An Use Case of Exclusive Lock

Multiple processes all competing to use some shared data

Each process has a code segment, called *critical section*, in which the shared data is accessed.

Problem – ensure that when one process is executing in its critical section, no other process is allowed to execute in its critical section.

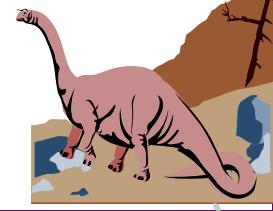


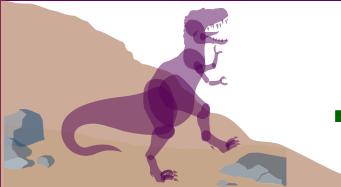


Critical Section and Mutual Exclusion

Thus, the execution of critical sections must be *mutually exclusive* (e.g., at most one process can be in its critical section at any time).

The *critical-section problem* is to design a protocol that processes can use to cooperate.





The Critical Section Protocol

```
do {
```

```
    :::
```

```
        entry section
```

```
        critical section
```

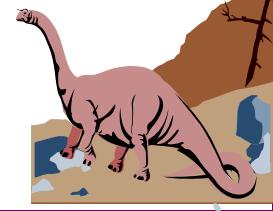
```
        exit section
```

```
    :::
```

```
} while (1);
```

A critical section protocol consists of two parts: an *entry section (or lock)* and an *exit section (or unlock)*.

Between them is the critical section that must run in a **mutually exclusive** way.



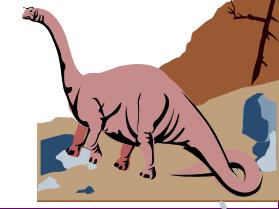


Solution to Critical-Section Problem

Any solution to the critical section problem must satisfy the following three conditions:

- ◆ **Mutual Exclusion**
- ◆ **Progress**
- ◆ **Bounded Waiting**

Moreover, the correctness of the solution cannot depend on **relative speed** of processes and **scheduling policy**.



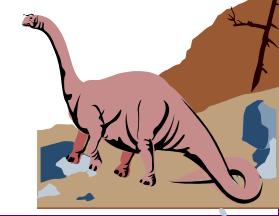


Mutual Exclusion

If a process P is executing in its critical section, then *no* other processes can be executing in their critical sections.

The **entry protocol** should be capable of blocking processes that wish to enter but cannot.

Moreover, when the process that is executing in its critical section exits, the **entry protocol** must be able to know this fact and allows a waiting process to enter.

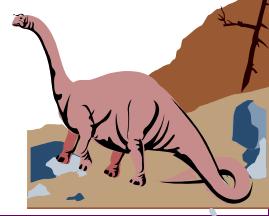




Progress

If no process is executing in its critical section and some processes wish to enter their critical sections, then

- ◆ Only those processes that are waiting to enter can participate in the competition (to enter their critical sections).
- ◆ No other process can influence this decision.
- ◆ This decision cannot be postponed indefinitely.





Bounded Waiting

After a process made a request to enter its critical section and before it is granted the permission to enter, there exists a *bound* on the number of times that other processes are allowed to enter.

Hence, even though a process may be blocked by other waiting processes, it will not be waiting forever.

- Assume that each process executes at a nonzero speed
- No assumption concerning relative speed of the n processes





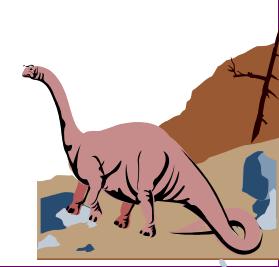
Solve the Problem without any OS Support

Consider a simple case of only 2 processes,
 P_0 and P_1

General structure of process P_i (and P_j)

```
do {   entry section
        critical section
        exit section
        remainder section
    } while (1);
```

Processes may share some common
variables to synchronize their actions.





Consider Three Test Cases

Case 1: One process repeatedly attempts to enter the critical section (CS)

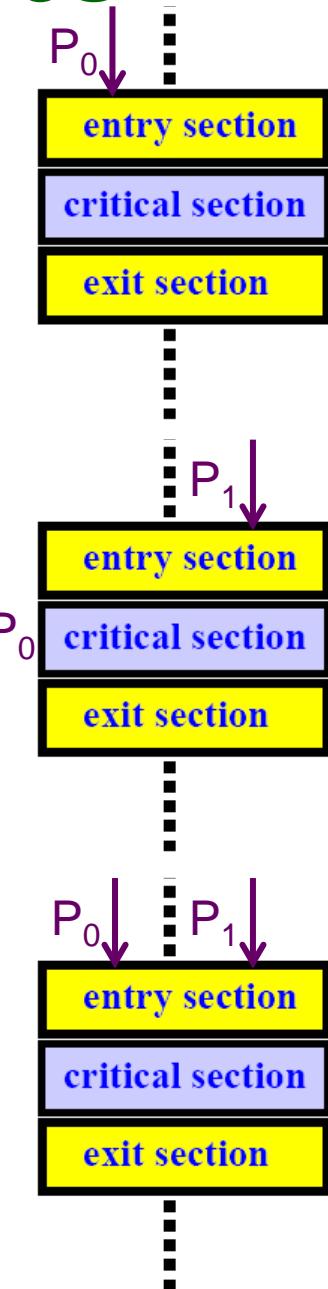
- ◆ Progress Test: Whether P_0 's repeated entering of CS is independent of P_1 's attempt

Case 2: One process is already in the critical section, and meanwhile the other process attempts to enter

- ◆ Mutual Exclusive Test: P_0 safely blocks P_1 out

Case 3: Two processes try to enter the critical section simultaneously

- ◆ Progress Test: Whether it is possible for the two processes to block each other's entry
- ◆ Mutual Exclusive Test: Whether it is possible for them to both enter the section





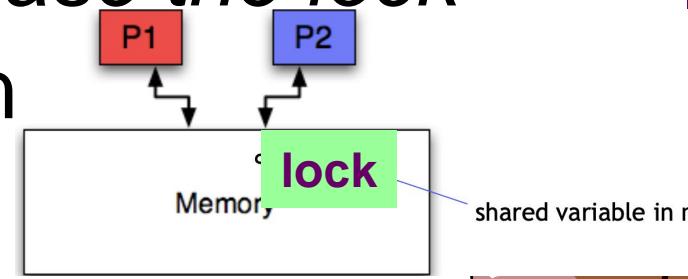
Our First Attempt: Algorithm 1

Shared variables:

- ◆ **boolean lock;** // initially **lock = false**
- ◆ **lock = true** ⇒ the critical section has been locked

Process P_i :

```
do {   while (lock) ; // if locked then wait  
      lock = true; // acquire the lock  
      critical section  
      lock = false; // release the lock  
      remainder section  
} while (1);
```



Does not satisfy **mutual exclusion**. Why?





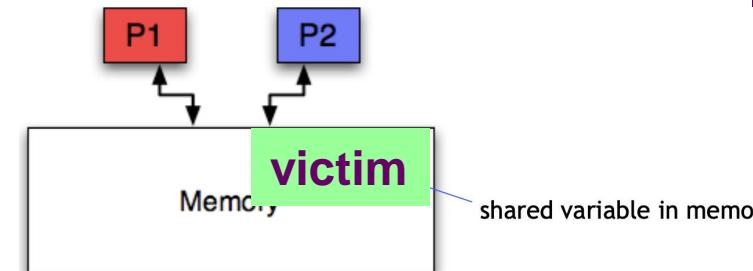
Our Second Attempt: Algorithm 2

Shared variables:

- ◆ **int victim;** initially **victim = 0 (or victim = 1)**

Process P_i :

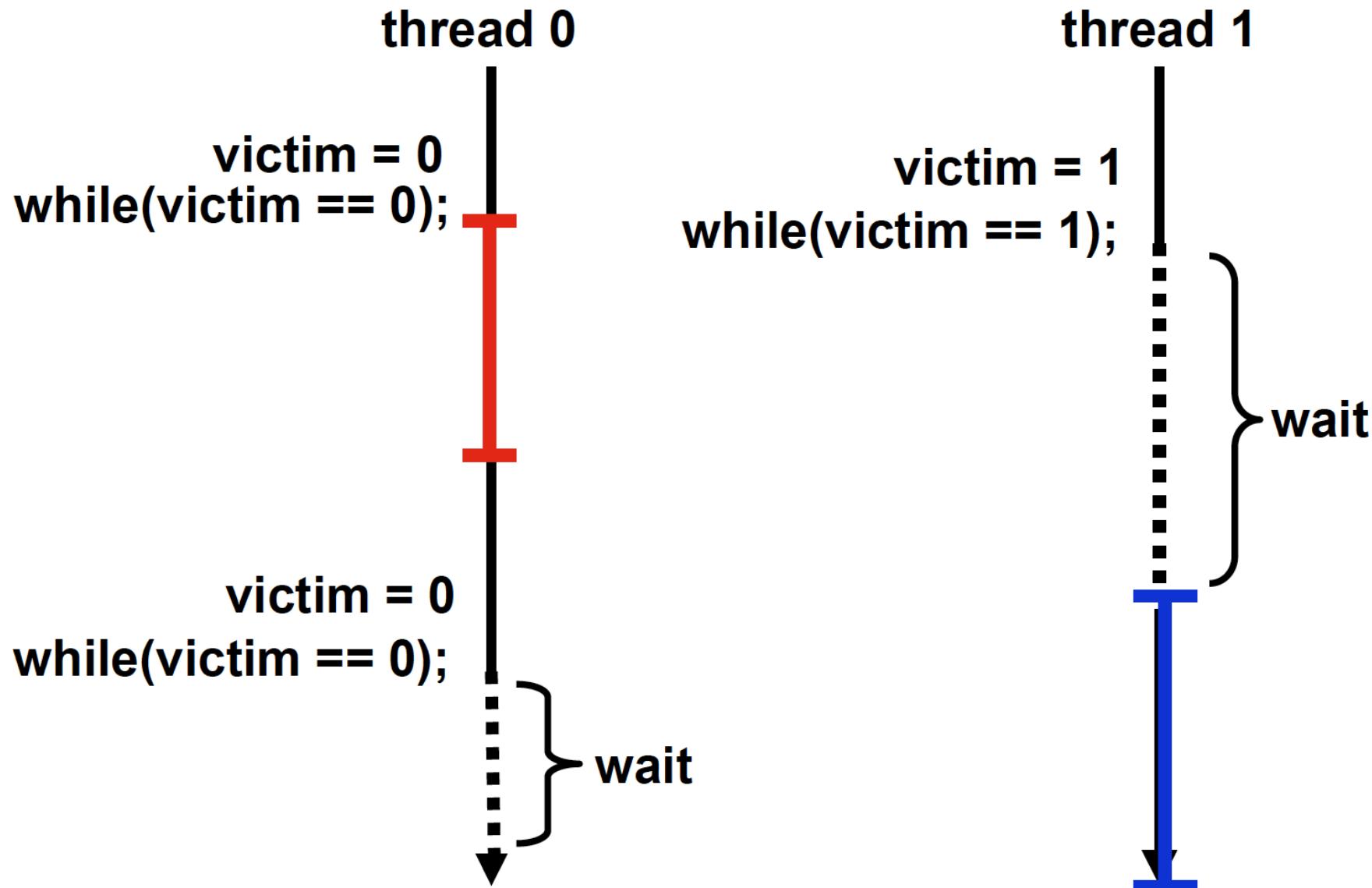
```
do {victim = i;          // determine who is the victim
    while (victim == i); // if I am victim, then wait
    critical section // assume empty
    // do nothing for CS exit
    remainder section
} while (1);
```



Processes are forced to run in an alternating way
Satisfies mutual exclusion, but not progress

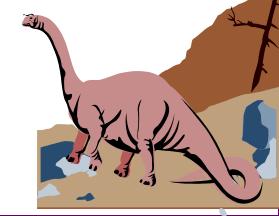
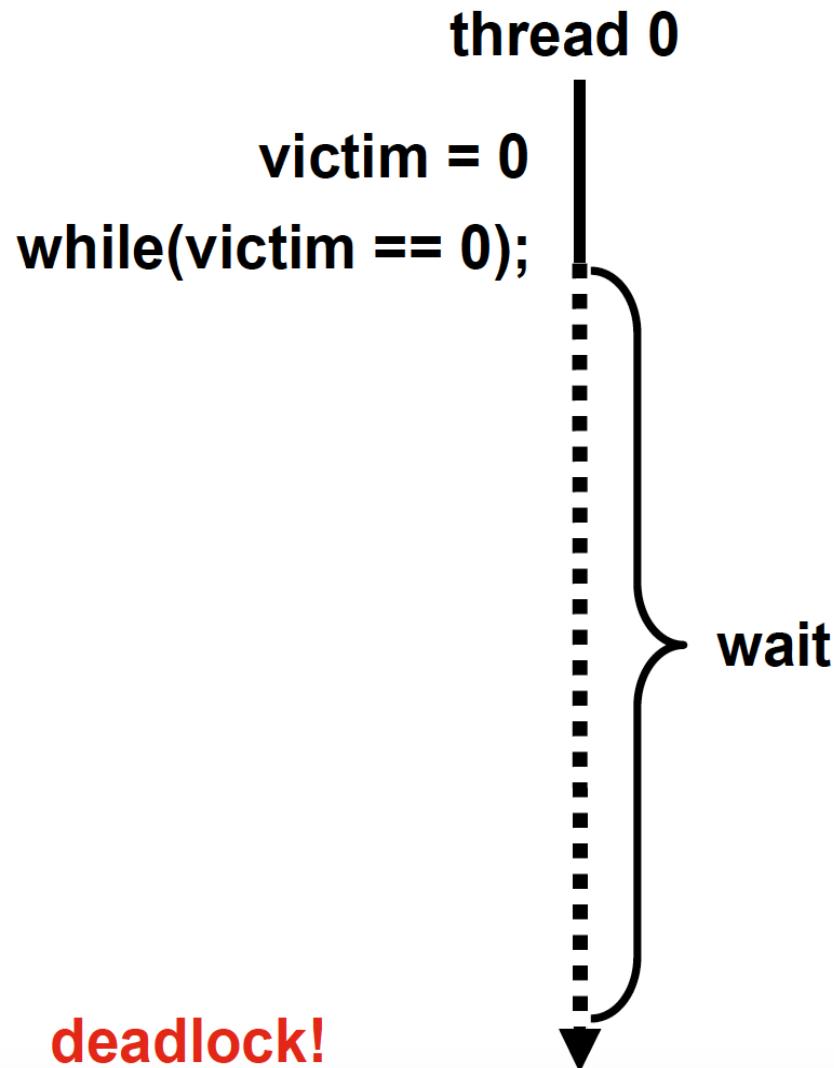


Alternating and Atomic Execution of Algorithm 2





Deadlock of Algorithm 2





Another Failed Attempt: Algorithm 3

Shared variables:

- ◆ **boolean flag[2]; // initially flag[0] = flag[1] = false**
- ◆ **flag[i] = true** $\Rightarrow P_i$ wants to enter its critical section

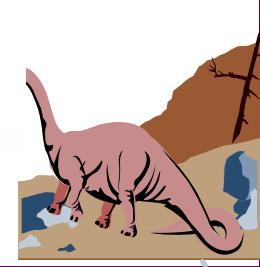
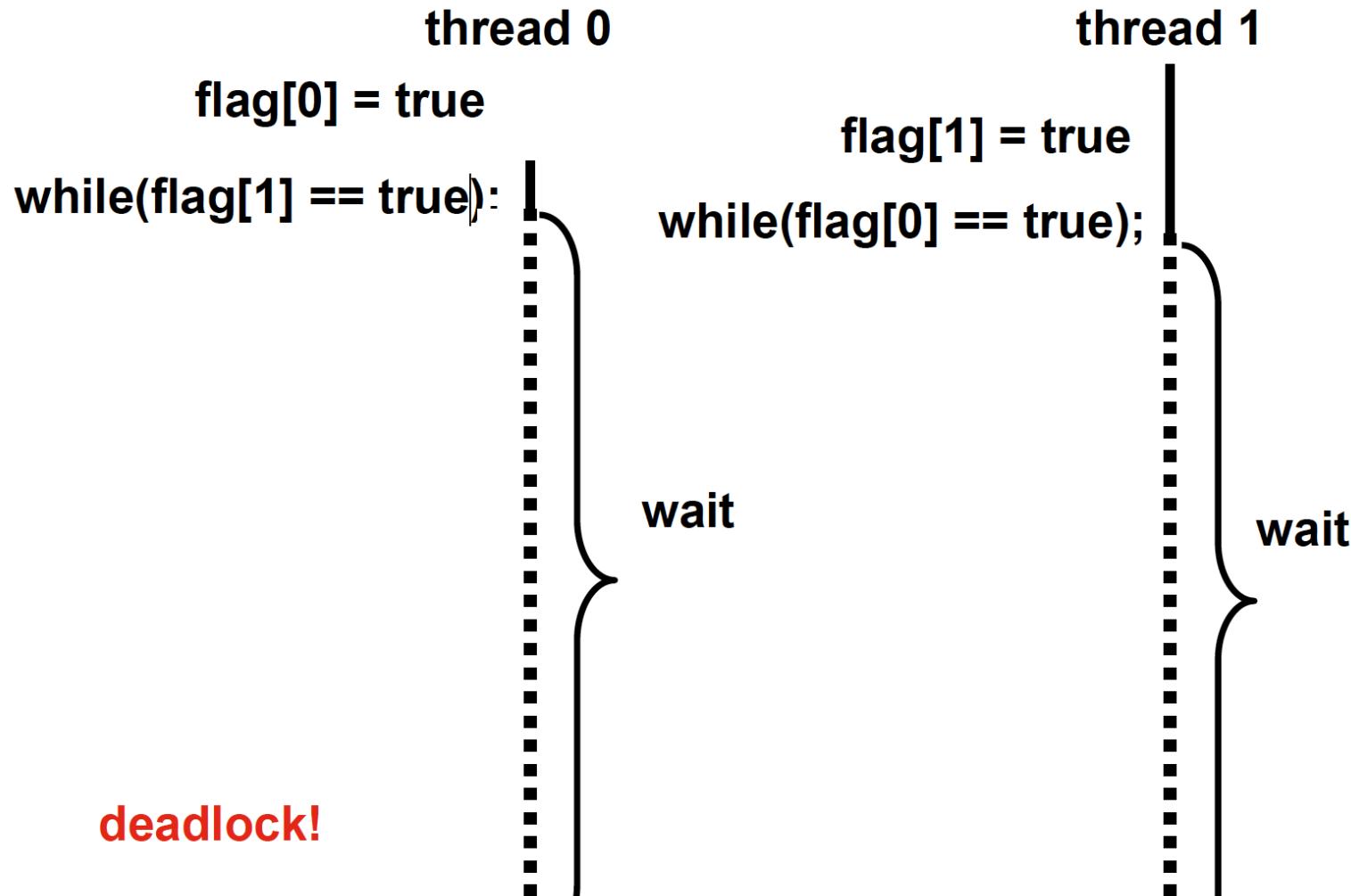
Process P_i

```
do {flag[i] = true;    // I want to enter
     while (flag[1-i]); // If you also want, then I wait
                         critical section
     flag[i] = false;   // I leave
                         remainder section
} while (1);
```

Can satisfy mutual exclusion, but not progress requirement. Why?



Deadlock Problem of Algorithm 3





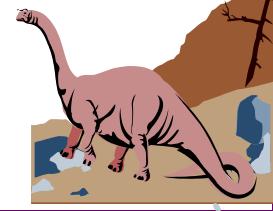
Is the Following Algorithm Correct?

What if we change the location of the statement: **flag[i] = true**?

Process P_i :

```
do {    while (flag[1-i]) ;
        flag[i] = true;
        critical section
        flag[i] = false;
        remainder section
    } while (1);
```

Does not satisfy mutual exclusion





Comparison of Algorithms 1, 2, 3

Critical Section Algorithms	Test Case 1: P_0 serialized enter	Test Case 2: P_0, P_1 serialized enter	Test Case 3: P_0, P_1 concurrent enter
Algorithm 1 with a shared <i>lock</i> variable	✓	✓	✗ (ME)
Algorithm 2 with a shared <i>victim</i> variable	✗ (Progress)	✓	✓
Algorithm 3 with two shared <i>flag[2]</i> variables	✓	✓	✗ (Progress)
Peterson's Algorithm, with a shared <i>victim</i> variable and two shared <i>flag[2]</i> variables	✓	✓	✓

Combine the advantages of
Algorithms 2 and 3





Peterson's Algorithm Give a demonstration

Combined shared variables of algorithms 2, 3.

Process P_i

```
do {   flag[i] = true; // I'm interested  
      victim = i;      // you go first  
      while (flag[1-i] and victim == i) ;  
            critical section
```

```
      flag[i] = false; // I'm not interested  
                        // any more
```

remainder section

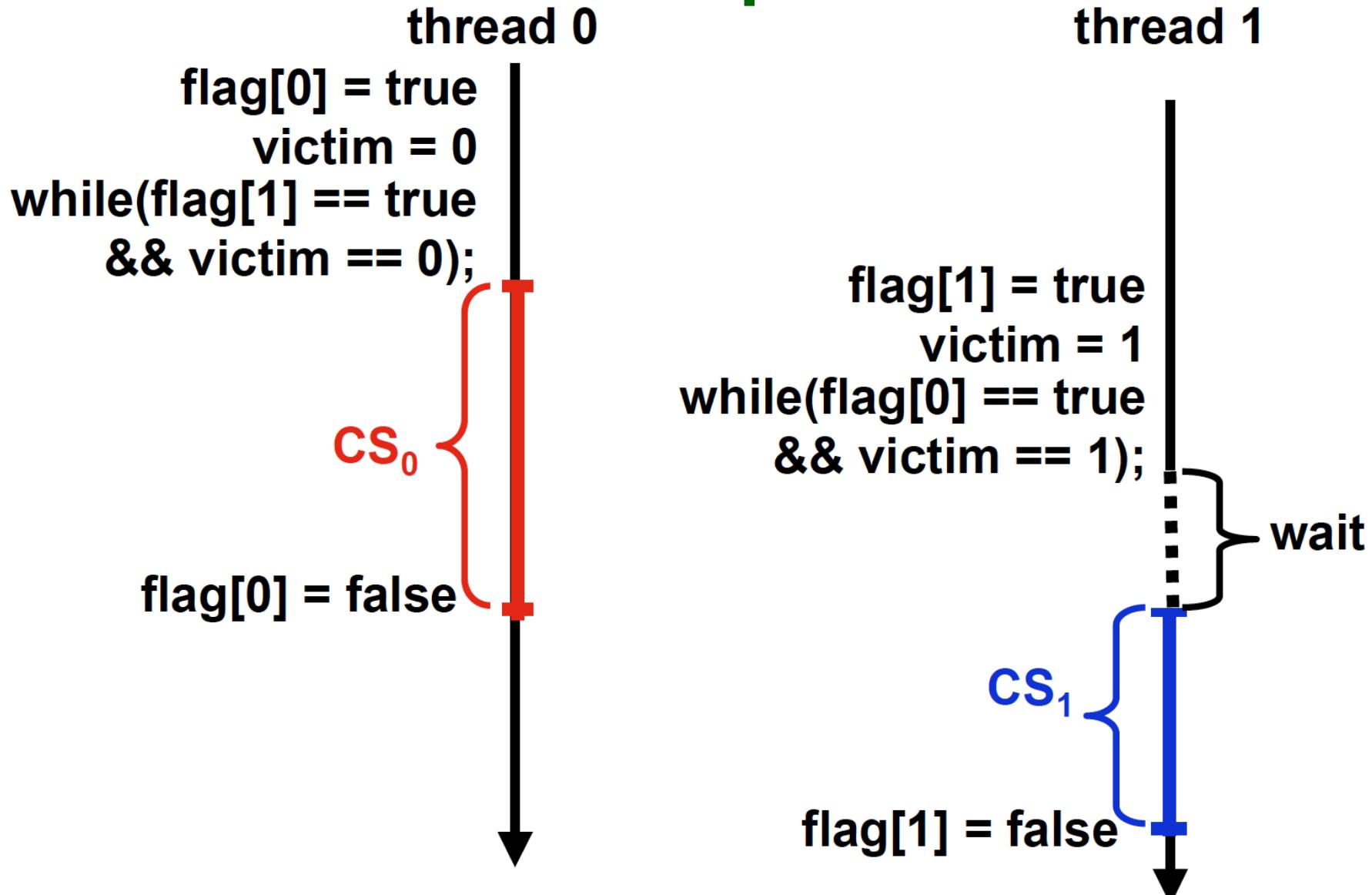
```
} while (1);
```

Gary Peterson. Myths about the Mutual Exclusion Problem.
Information Processing Letters, 12(3):115-116, 1981.

Meet all the three requirements; Can solve the critical-section problem for two processes.

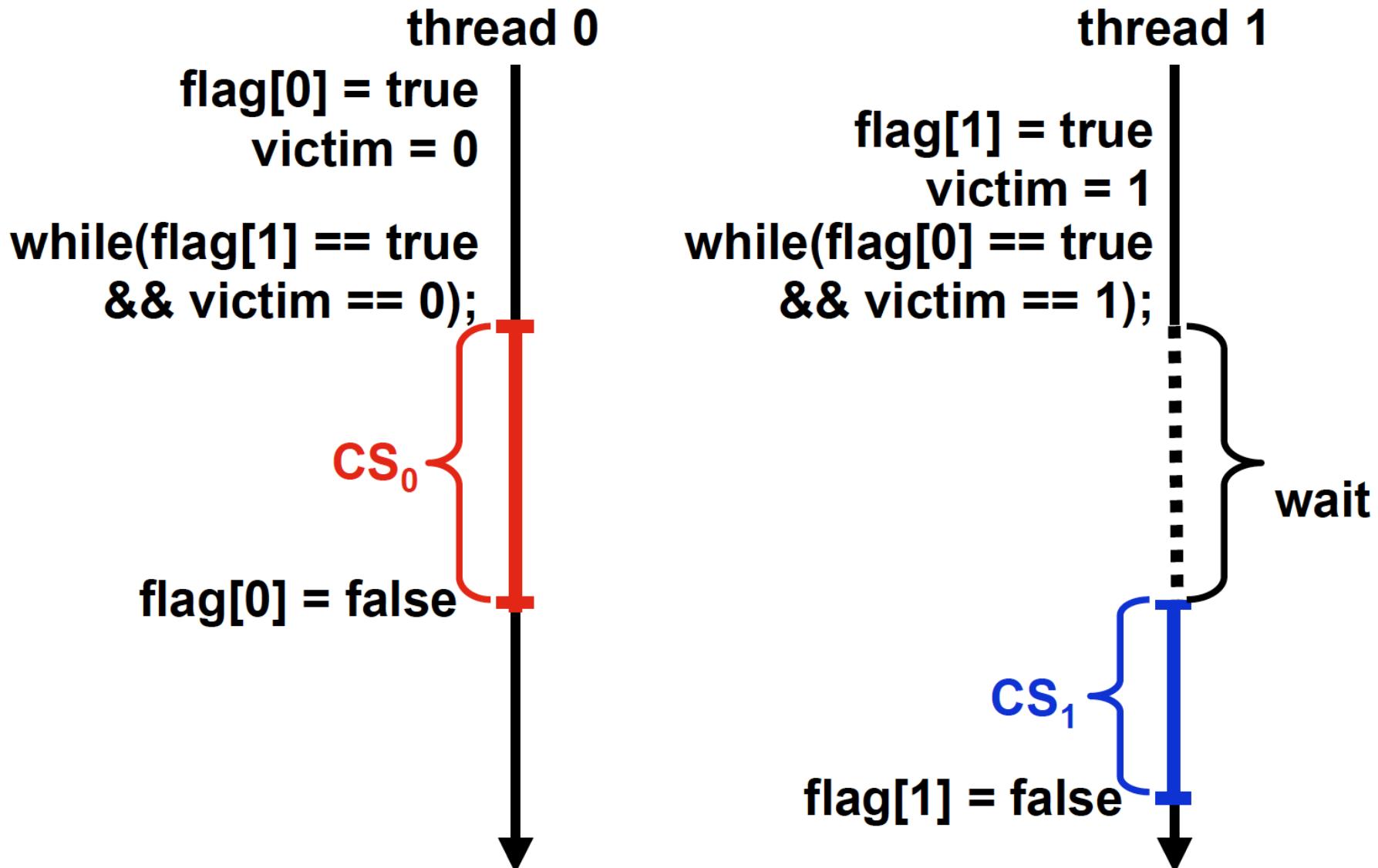


Peterson's Lock: Serialized Acquires





Peterson's Lock: Concurrent Acquires



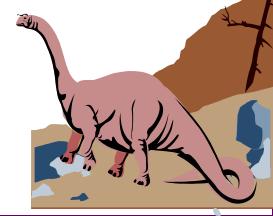
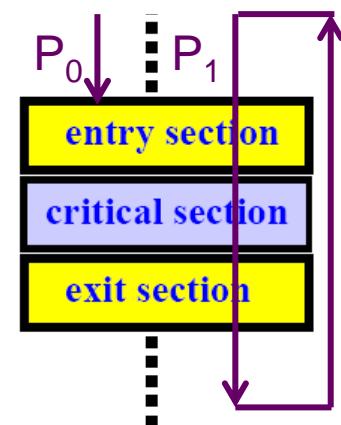


Test the Bounded Waiting Property

Recall: After a process made a request to enter its critical section and before it is granted the permission to enter, there exists a *bound* on the number of times that other processes are allowed to enter.

Test Case: Two processes attempt to enter the critical section simultaneously

- ◆ Assume P0 is fast, while P1 is slow
- ◆ Can P0 repeatedly grab the exclusive lock, causing P1 to starve?
 - ✓ If yes/no, the solution of critical section cannot/can satisfy bounded waiting property





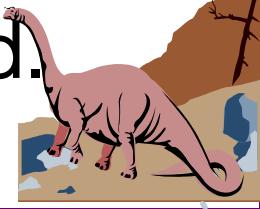
Proof of Peterson's Algorithm

The mutual exclusion requirement is assured.

The progress requirement is assured. The victim variable is only considered when both processes are using, or trying to use, the resource.

Deadlock is not possible. If both processes are testing the while condition, one of them must be the victim. The other process will proceed.

Finally, bounded waiting is assured. When a process that has exited the CS reenters, it will mark itself as the victim. If the other process is already waiting, it will be the next to proceed.



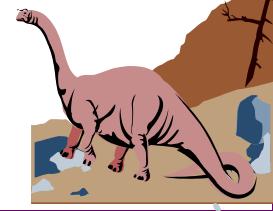


Quiz: Is the following code correct?

What if we change **victim** = **i** to **victim** = **1-i**?

```
do {      flag[i] = true; // I'm interested  
          victim = 1-i;    // I go first  
          while (flag[1-i] and victim == i) ;  
                  critical section  
          flag[i] = false; // I'm not interested  
                  remainder section  
} while (1);
```

- ◆ Can the code satisfy mutual exclusion?
- ◆ Can the code satisfy progress?
- ◆ Can the code satisfy bounded waiting?



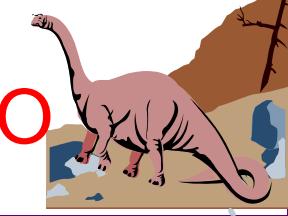


Quiz: Is the following code correct?

What if we change **victim = i** to **victim = 1-i**?

```
do {      flag[i] = true; // I'm interested  
         victim = 1-i;    // I go first  
         while (flag[1-i] and victim == i) ;  
                  critical section  
         flag[i] = false; // I'm not interested  
                  remainder section  
 } while (1);
```

- ◆ Can the code satisfy mutual exclusion? NO
- ◆ Can the code satisfy progress? YES
- ◆ Can the code satisfy bounded waiting? NO





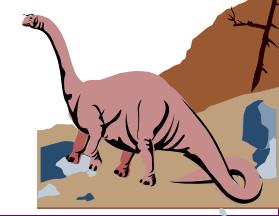
Memory Fence

Give a C-code demo of Peterson's algorithm

A memory barrier, also known as a membar, memory fence or fence instruction, is a type of barrier instruction that causes a central processing unit (CPU) or compiler to enforce an ordering constraint on memory operations issued before and after the barrier instruction. https://en.wikipedia.org/wiki/Memory_barrier

Operations issued prior to the barrier are guaranteed to be performed before operations issued after the barrier.

<https://gcc.gnu.org/onlinedocs/gcc-4.6.2/gcc/Atomic-Builtins.html>

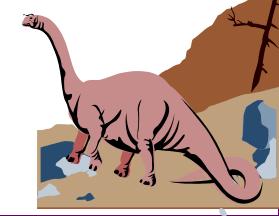




Guarantee Memory Access Ordering

Insert full memory barrier at multiple points

```
do {      flag[i] = true; // I'm interested  
         __sync_synchronize(); // full memory barrier  
         victim = i;          // You go first  
         while (flag[j] and victim == i) ;  
         __sync_synchronize(); // full memory barrier  
         critical section  
         __sync_synchronize(); // full memory barrier  
         flag[i] = false; // I'm not interested  
         remainder section  
 } while (1);
```





Lamport's Bakery Algorithm

Solve the critical section problem for an arbitrary number of processes

Before entering its critical section, process receives a number. Holder of the smallest number enters the critical section.

If processes P_i and P_j receive the same number, if $i < j$, then P_i is served first; else P_j is served first.

The numbering scheme always generates numbers in increasing order of enumeration; i.e., 1,2,3,3,3,3,4,5...





Bakery Algorithm

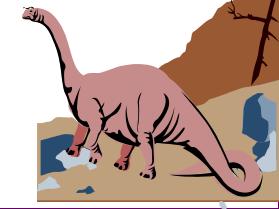
Notation

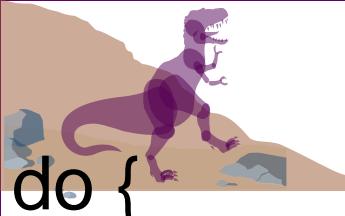
- ◆ $(a,b) < (c,d)$ if $a < c$ or if $a = c$ and $b < d$
- ◆ $\max(a_0, \dots, a_{n-1})$ is a number, k , such that $k \geq a_i$ for $i = 0, \dots, n - 1$

Shared data

```
boolean choosing[n];
int number[n];
```

Data structures are initialized to **false** and **0** respectively





do {

Bakery Algorithm

```
choosing[i] = true;           Get a ticket first
number[i] = max(number[0], number[1], ..., number[n - 1]) + 1;
choosing[i] = false;
for (j = 0; j < n; j++) {
    while (choosing[j]) ;
    while ( (number[j] != 0) &&
            ((number[j], j) < (number[i], i)) ) ;
}
critical section
number[i] = 0;                 What is the use of choosing[]?
                                Which parts are the entry and exit sections?
                                remainder section
} while (1);
```

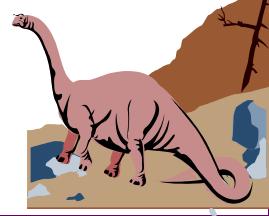




What if we remove the choosing[]?

```
do {  
    number[i] = max(number[0], number[1], ...,  
                    number[n - 1]) + 1;  
    for (j = 0; j < n; j++) {  
        while ( (number[j] != 0) &&  
                ((number[j], j) < (number[i], i)) ) ;  
    }  
    critical section  
    number[i] = 0;  
    remainder section  
} while (1);
```

Can the above code satisfy the mutual exclusion and why?
Hint: two processes computed the same number[i].





Chapter 6: Process Synchronization

Background

The Critical-Section Problem

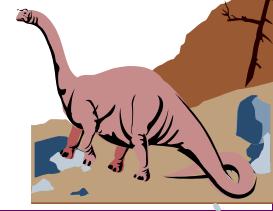
Synchronization Hardware to Build a Lock

Semaphores

Classical Problems of Synchronization

Conditional Variables and Monitors

Synchronization Examples

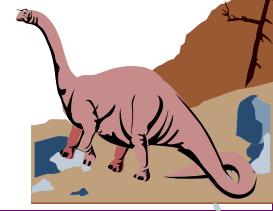




Hardware Support

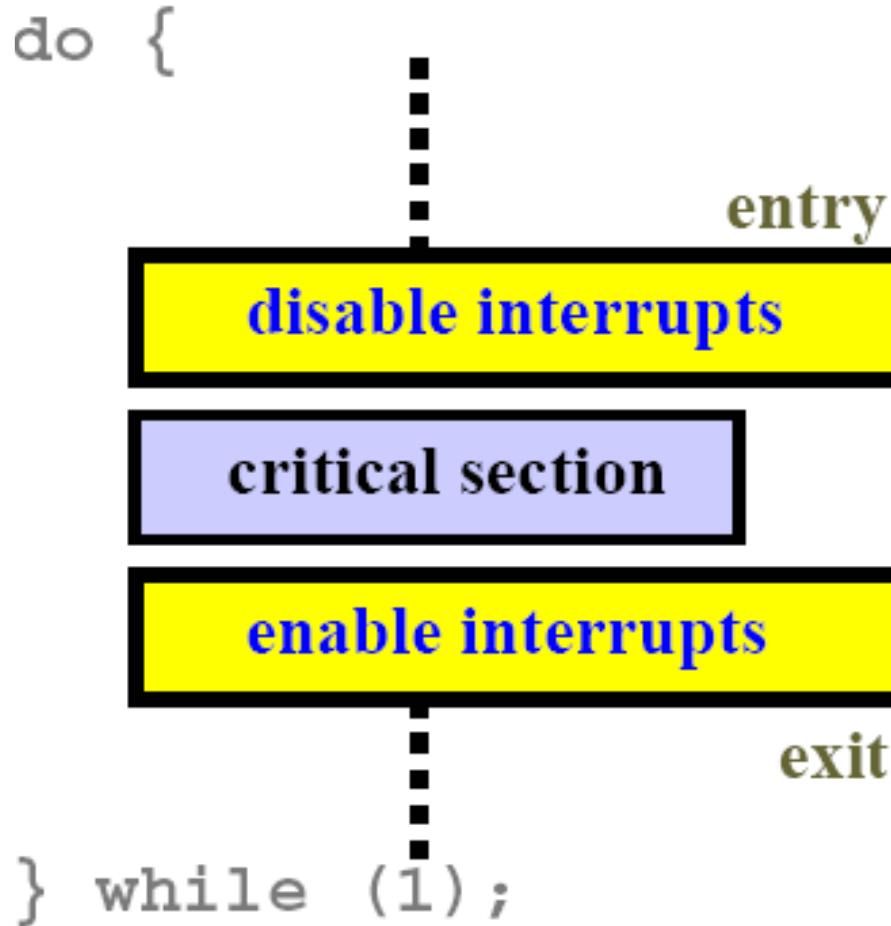
There are two types of hardware synchronization supports:

- ◆ Disabling/Enabling interrupts: This is slow and difficult to implement on multiprocessor systems.
- ◆ Special machine instructions:
 - ✓ Test and set (TAS)
 - ✓ Swap
 - ✓ Atomic fetch-and-add





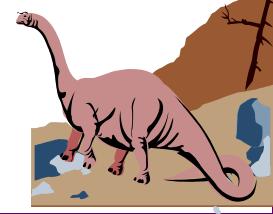
Interrupt Disabling



Because interrupts are disabled, no context switch will occur in a critical section.

Infeasible in a multiprocessor system because all CPUs must be informed.

Some features that depend on interrupts (e.g., clock) may not work properly.

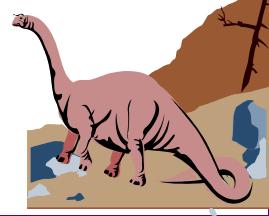




Test-and-Set (TAS)

Test and modify the content of a machine word atomically

```
boolean TestAndSet(boolean &target) {
    boolean rv = target;
    target = true;
    return rv;
}
```





Mutual Exclusion with Test-and-Set

Shared data:

boolean lock = false;

Process P_i

do {

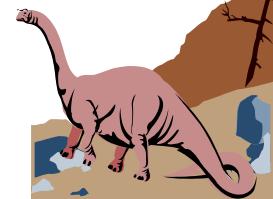
while (TestAndSet(lock)) ;

critical section

lock = false;

remainder section

} while(1);

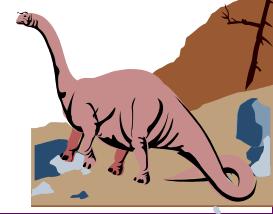




Atomic Swap

Atomically swap two variables.

```
void Swap(boolean &a, boolean &b) {  
    boolean temp = a;  
    a = b;  
    b = temp;  
}
```





Mutual Exclusion with Swap

Shared data (initialized to **false**):

boolean lock = false;

local variable

boolean key;

Process P_i or Interrupt Handler TH_i

do {

key = true;

while (key == true) Swap(lock,key);

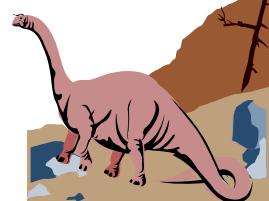
critical section

lock = false;

remainder section

} while(1);

**Cannot satisfy
bounded waiting.
Why?**





Another Atomic CPU

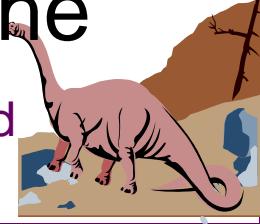
Instruction Fetch-and-add

fetch-and-add instruction performs the operation

```
    << atomic >>
function FetchAndAdd(address location, int inc)
{
    int value := *location
    *location := value + inc
    return value
}
```

can be used to implement concurrency control structures such as mutex locks and semaphores.

An atomic `fetch_add` function appears in the C++11 standard <https://en.wikipedia.org/wiki/Fetch-and-add>





Bounded Waiting Mutual

Exclusion with TestAndSet

Shared data (initialized to false):

boolean lock = false; boolean waiting[n]; //init to false

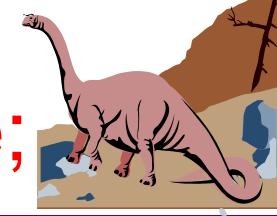
local variable: **boolean key;**

Enter Critical Section (Lock)

```
waiting[i] = true;  
key = true;  
while (waiting[i] && key)  
    key=TestAndSet(lock);  
waiting[i] = false;
```

Leave Critical Section (unlock)

```
j = (i+1)%n  
while ((j!=i) && !waiting[j])  
    j = (j+1)%n;  
if (j == i)  
    lock = false;  
else  
    waiting[j] = false;
```





Question: What is a spinlock? Why spinlocks are not appropriate for single-processor systems yet are often used in multiprocessor system?

用忙等待方式实现的信号量称为自旋锁。自旋锁等待进入临界区需要占有CPU周期。

在单处理器系统中，这将导致已进入临界区的进程得不到机会执行，反而使想进临界区的进程等待更长时间。

在多处理器系统中，当临界区很短时自旋锁是合适的。

由于有多个处理器，忙等待的进程不影响在临界区中的进程在其他处理器上执行。由于临界区很短，在临界区里的进程很快就能离开临界区，其他忙等待的进程就可以进入它的临界区。这种情况下反而避免了由于阻塞和唤醒导致的上下文切换开销。

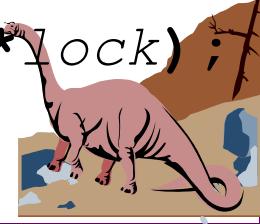




Spin Locks

A spinlock is a lock which causes a thread trying to acquire it to simply **wait in a loop ("spin")** while repeatedly checking if the lock is available. Since the thread remains active but is not performing a useful task, the use of such a lock is a kind of **busy waiting**.

```
#include <pthread.h>
int pthread_spin_lock(pthread_spinlock_t *lock);
int pthread_spin_trylock(pthread_spinlock_t
*lock);
int pthread_spin_unlock(pthread_spinlock_t*lock);
```



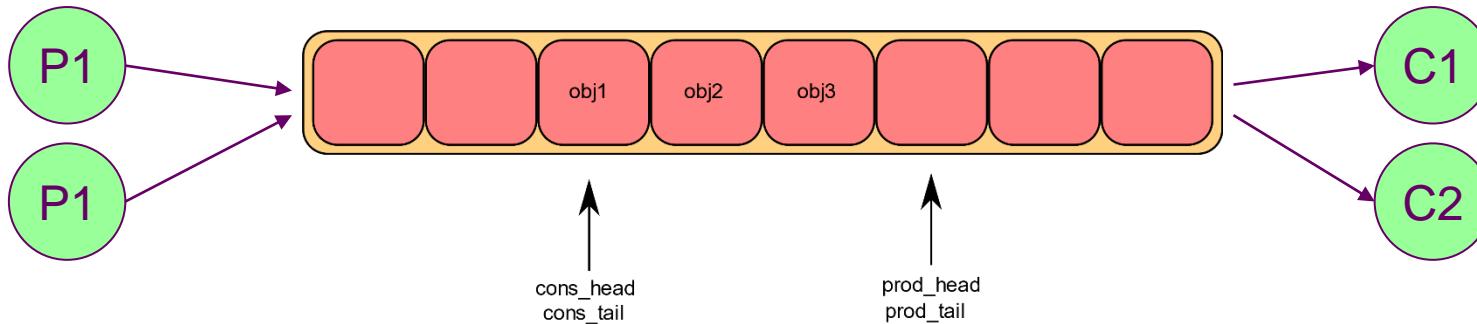


An Introduction to Lock-Free Programming

People often describe **lock-free programming** as programming without kernel-provided mutexes, which are also referred to as locks.

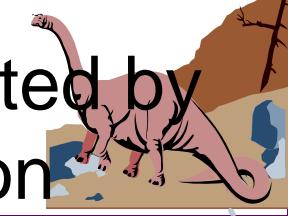
An Example: DPDK Lockless Ring Buffer

- ◆ Communication between multiple applications running on different CPU cores



http://dpdk.org/doc/guides/prog_guide/ring_lib.html

- ◆ buffer enqueue and dequeue are implemented by atomic Compare-And-Swap (CAS) instruction





Chapter 6: Process Synchronization

Background

The Critical-Section Problem

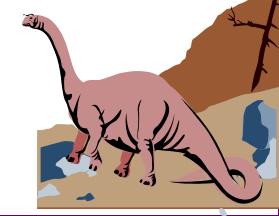
Synchronization Hardware

Semaphore as a generic synchronization tool

Classical Problems of Synchronization

Conditional Variables and Monitors

Synchronization Examples





Dijkstra



Edsger Wybe Dijkstra

- ◆ (Dutch: ['ɛtsxər 'vibə 'dɛikstra])
- ◆ 11 May 1930 - 6 August 2002
(aged 72)

Known for

- ◆ Dijkstra's algorithm (single-source shortest path problem)
- ◆ Structured programming, First implementation of ALGOL 60 ("Goto Statements Considered Harmful")
- ◆ **Semaphores**, Layered approach to operating system design, software-based paged virtual memory in
 - ✓ THE multiprogramming system





Concept of Semaphore

It is a synchronization tool that does not require busy waiting, but needs the support from kernel

semop() - Unix, Linux System Call

http://www.tutorialspoint.com/unix_system_calls/semop.htm

Semaphore S — an integer variable

It can only be accessed via two indivisible
(atomic) operations: **wait** and **signal**

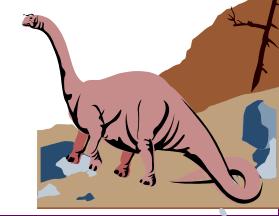
They are functionally equivalent to the following
busy-waiting operations.

wait (S):

while $S \leq 0$ do no-op;
 $S--;$

signal (S):

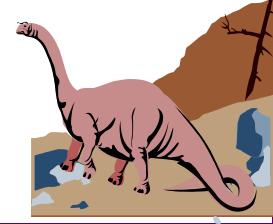
$S++;$





Concept of Semaphore

A useful way to think of a semaphore as used in the real-world systems is as a record of how many units of a particular resource are available,





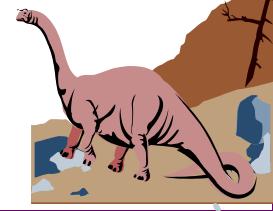
Semaphore Implementation

Define a semaphore as a record

```
typedef struct {  
    int counter;  
    struct process *L;  
    an in-kernel exclusive lock;  
} semaphore;
```

Assume two simple operations:

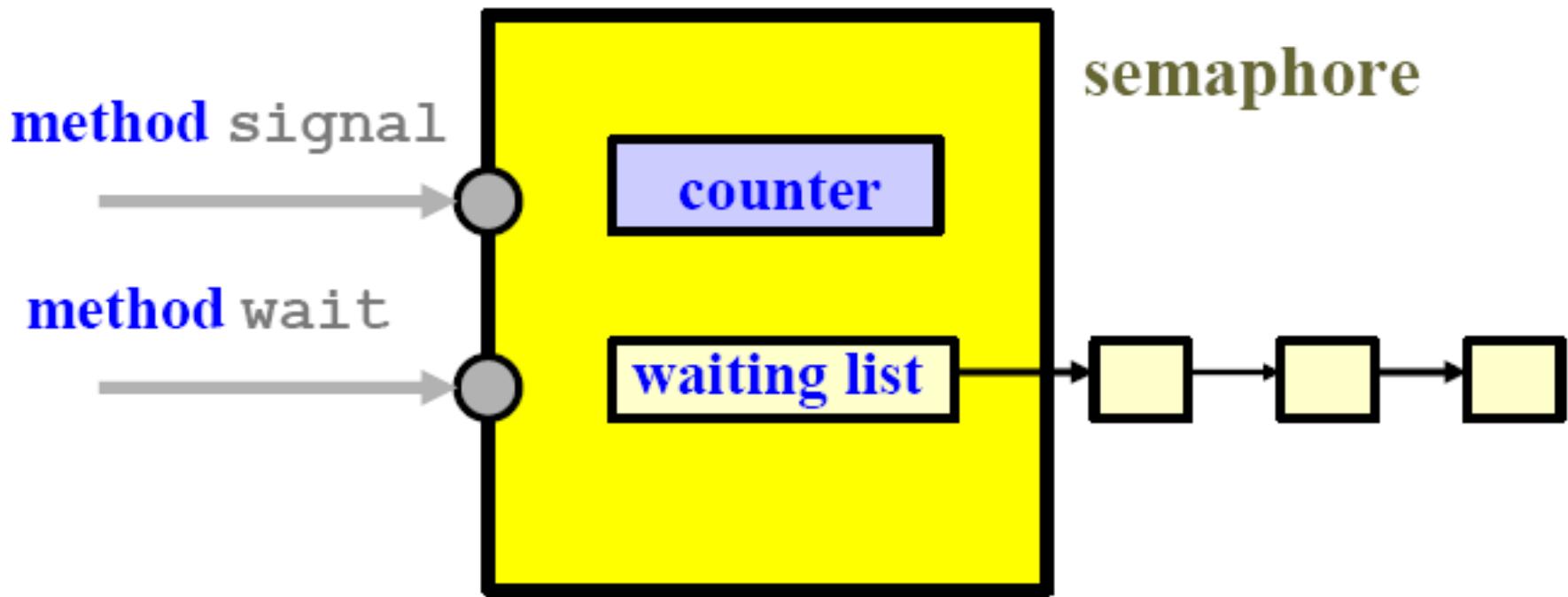
- ◆ **block**: block the process that invokes it.
- ◆ **wakeup(P)**: resumes the execution of a blocked process P .





Semaphore Schematics

Semaphore = counter + kernel mutex + waiting list



A useful way to think of a semaphore as used in the real-world systems is as a record of how many units of a particular resource are available,





Semaphore Implementation (1)

Semaphore operations now defined as follows

- ◆ Uniprocessor solution: disable interrupts. Recall that the only way the dispatcher regains control is through interrupts or through explicit requests.

wait(S):

```
Disable interrupts;  
S.counter--;  
if (S.counter < 0) {  
    add this process to S.L;  
block;  
}  
Enable interrupts;
```

signal(S):

```
Disable interrupts;  
S.counter++;  
if (S.counter <= 0) {  
    remove a process P  
    from S.L;  
wakeup(P);  
}  
Enable interrupts;
```



Semaphore Implementation (2)

What do we do in a multiprocessor platform to implement *wait(S)* and *signal(S)*?

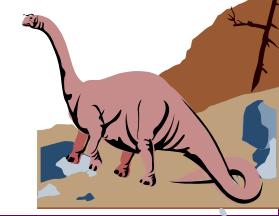
- ◆ Can't turn off interrupts to get low-level mutual exclusion
- ◆ Suppose hardware provides atomic test-and-set instruction

wait(S):

```
while(TAS(S.lock));  
S.counter--;  
if (S.counter < 0) {  
    add this process to S.L;  
    block;  
}  
S.lock = 0;
```

signal(S):

```
while(TAS(S.lock));  
S.counter++;  
if (S.counter <= 0) {  
    remove a process P  
    from S.L;  
    wakeup(P);  
}  
S.lock = 0;
```

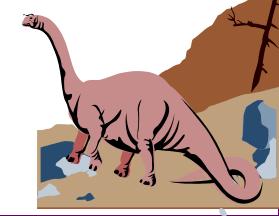




Semaphore Implementation (3)

Use atomic
TAS instrument
to implement
spin lock

```
while (TestAndSet(S.lock)) ; // spin lock
    critical section
    S.lock = false; // spin_unlock
    remainder section
}
wait(S):
    S.counter--;
    if (S.counter < 0) {
        add this process to S.L;
        block;
    }
}
signal(S):
    S.counter++;
    if (S.counter <= 0) {
        remove a process P
        from S.L;
        wakeup(P);
    }
}
```





POSIX Library's Support of Semaphore

All POSIX semaphore functions and types are prototyped or defined in semaphore.h

```
#include <semaphore.h>
```

To define a semaphore object, use

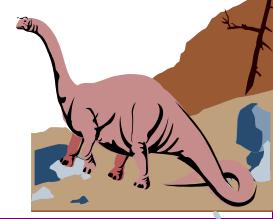
```
sem_t sem_name;      OR      sem_t * sem_pointer;
```

For initialization, use either of the following APIs.

```
int sem_init (sem_t *sem, int pshared, unsigned int initial_value);  
sem_t * sem_open (const char * name, int oflag, unsigned int  
initial_value);
```

To increment/decrement the value of a semaphore,

```
int sem_wait (sem_t * sem_pointer);  
int sem_post (sem_t * sem_pointer);
```





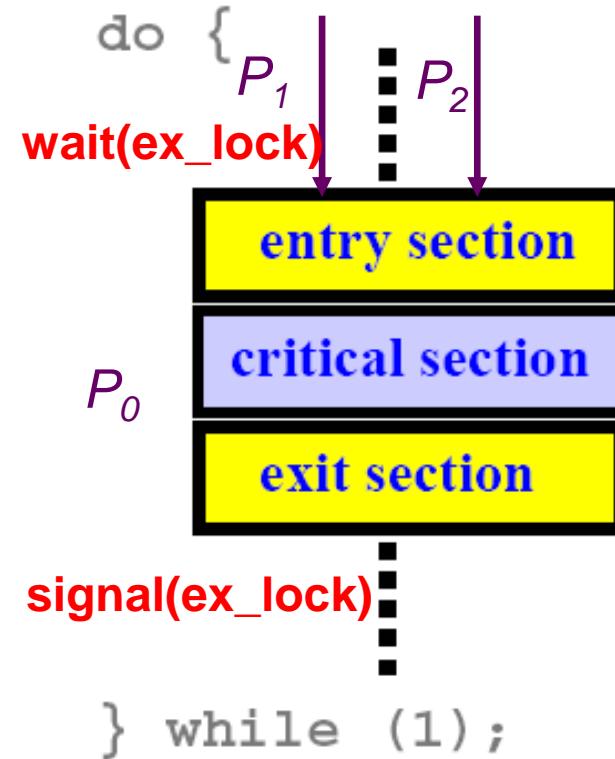
Applications of Binary Semaphore:

1. Solve the Critical Section Problem

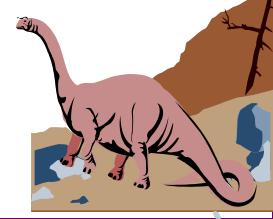
Shared data: **semaphore ex_lock; // initialize =1**

Process P_i :

```
do {  
    wait(ex_lock);  
    critical section  
    signal(ex_lock);  
    remainder section  
} while (1);
```



Give a demonstration





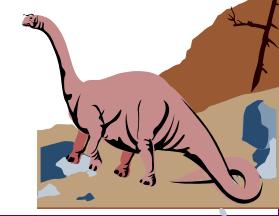
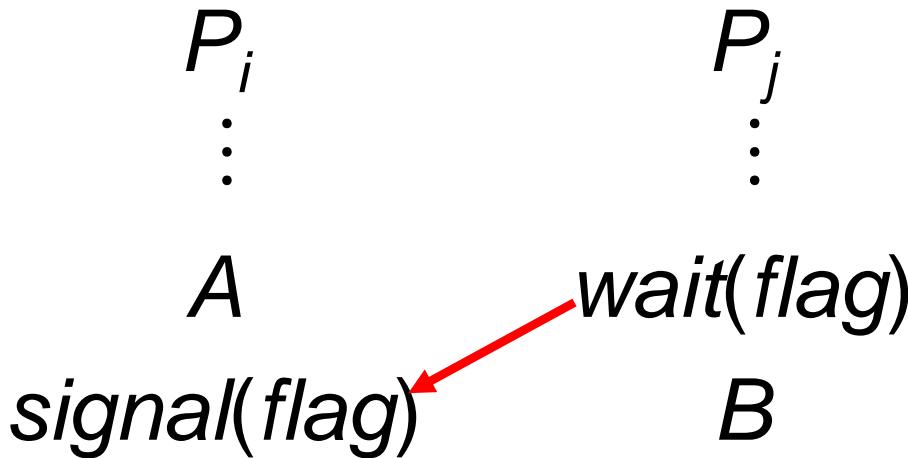
Applications of Binary Semaphore:

2. Act as an Event Notification Tool

Execute B in P_j only after A executed in P_i

Use semaphore *flag*, which is initialized to 0

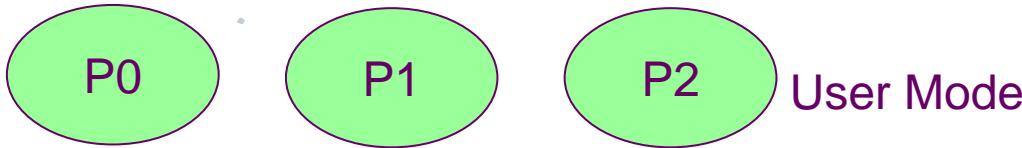
Shared data: **semaphore flag; // initialize =0**





An example of using semaphore to solve critical section problem

Shared data: **semaphore ex_lock; // initialize to 1**



```
Semaphore ex_lock {  
    int counter;  
    struct process * L;  
    an in-kernel exclusive lock;  
}
```

Kernel Mode

```
wait(S):  
    while(TAS(S.lock));  
    S.counter--;  
    if (S.counter < 0) {  
        add this process to S.L;  
        block;  
    }  
    S.lock = 0;
```

```
signal(S):  
    while(TAS(S.lock));  
    S.counter++;  
    if (S.counter <= 0) {  
        remove a process P  
        from S.L;  
        wakeup(P);  
    }  
    S.lock = 0;
```

进程P1和P2同时竞争ex_lock锁的控制权，可能产生竞争条件。因为都试图修改ex_lock内部的S.counter变量。

此时必须依靠内核信号量S的S.lock自旋锁，将wait(S)和signal(S)方法都实现为关键代码段。



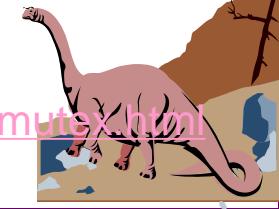


Difference between Binary Semaphore and Mutex

Question: Is there any difference between binary semaphore and mutex, or they are essentially same?

Answer: They're semantically the same, but in practice you will notice weird differences

- ◆ Semaphore is implemented by process/thread blocking and wakeup
- ◆ Mutex may be internally implemented by some kernels as spin locks, which could be more efficient on multi-processor systems but will slow down a single processor machine

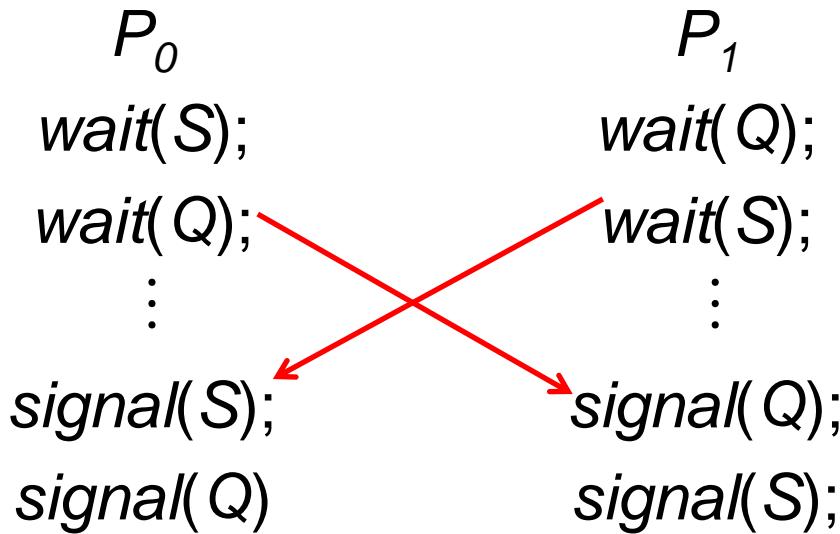




Side Effect of Semaphore: Deadlock and Starvation

Deadlock – two or more processes are waiting indefinitely for an event that can be caused by only one of the waiting processes.

- ◆ Example: Let S and Q be two semaphores initialized to 1



Starvation – indefinite blocking. A process may never be removed from the semaphore queue in which it is waiting.





Chapter 6: Process Synchronization

Background

The Critical-Section Problem

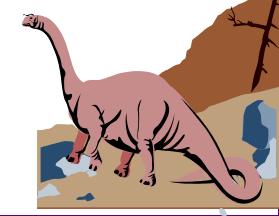
Synchronization Hardware

Semaphores

Classical Problems of Synchronization

Conditional Variables and Monitors

Synchronization Examples



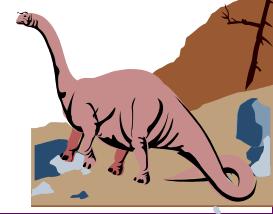


Classical Problems of Synchronization

Bounded-Buffer Problem (or called Producer-Consumer Problem)

Readers and Writers Problem (or called Shared-Lock Problem)

Dining-Philosophers Problem





Solution 1 for Producer-Consumer

Shared variables besides the shared buffer

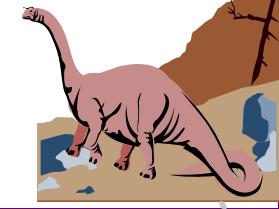
semaphore fillCount, emptyCount;

Initially:

fillCount = 0, emptyCount = n

fillCount: the number of items in the buffer

emptyCount: the number of empty slots in the buffer





Solution 1 for Producer-Consumer

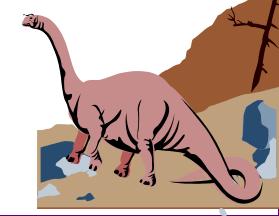
Producer:

```
do {  
    ...  
    produce an item in nextp  
    ...  
    wait(emptyCount);  
    add nextp to buffer  
    signal(fillCount);  
} while (1);
```

Consumer:

```
do {  
    wait(fillCount);  
    remove an item from  
    buffer to nextc  
    signal(emptyCount);  
    ...  
    consume the item in nextc  
    ...  
} while (1);
```

Question: Is this solution correct?



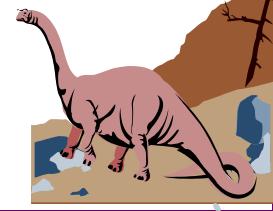


Solution 1 for Producer-Consumer

This solution contains a serious race condition that can result in two or more processes modifying into the same cursor *in* or *out* at the same time.

To understand how this is possible, recall how the procedures “add **nextp** to buffer” and “remove an item from buffer” are implemented, by

in = (*in*+1)%*BUF_SIZE*, and
out = (*out*+1)%*BUF_SIZE*.





Solution 2 for Producer-Consumer

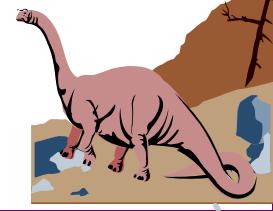
Shared data

semaphore fillCount, emptyCount, mutex;

Initially:

fillCount = 0, emptyCount = n, mutex = 1

mutex: guarantee the **mutual exclusive access** of the shared buffer





Solution 2 for Producer-Consumer

Producer:

```
do {
```

```
    ...
```

produce an item in **nextp**

```
    ...
```

```
    wait(mutex);
```

```
    wait(emptyCount);
```

add **nextp** to buffer

```
    signal(fillCount);
```

```
    signal(mutex);
```

```
} while (1);
```

Consumer:

```
do {
```

```
    wait(mutex);
```

```
    wait(fillCount);
```

remove an item from
buffer to **nextc**

```
    signal(emptyCount);
```

```
    signal(mutex);
```

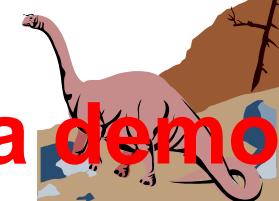
```
    ...
```

consume the item in **nextc**

```
    ...
```

```
} while (1);
```

Question: Is this solution correct? Give a demo





Solution 3 for Producer-Consumer

Producer:

```
do {
```

```
    ...
```

produce an item in **nextp**

```
    ...
```

```
    wait(emptyCount);
```

```
    wait(mutex);
```

add **nextp** to buffer

```
    signal(mutex);
```

```
    signal(fillCount);
```

```
} while (1);
```

Consumer:

```
do {
```

```
    wait(fillCount);
```

```
    wait(mutex);
```

remove an item from
buffer to **nextc**

```
    signal(mutex);
```

```
    signal(emptyCount);
```

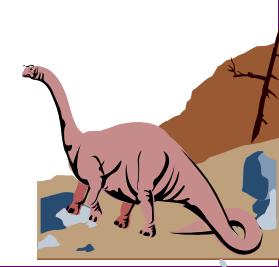
```
    ...
```

consume the item in **nextc**

```
    ...
```

```
} while (1);
```

Give a code demonstration





A Short Summary

结论1：需要用**mutex**确保对关键共享资源的互斥访问，比如 **shared bounded buffer**

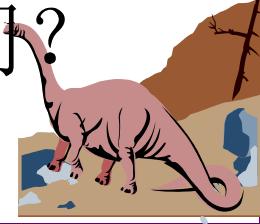
结论2：信号量**wait**的顺序很重要

◆ 例子：如果**wait(mutex)**错误放在了**wait(fillCount)**或者**wait(emptyCount)**之前，会导致死锁

问题1：信号量**signal**的顺序重要吗？

◆ 例子：**signal(mutex)**和**signal(fillCount)**可以交换吗？

问题2：能否把**produce an item**和**consume an item**放到**wait(mutex)**和**signal(mutex)**之间？





When the buffer size is only one, can we remove the mutex variable?

Initially: **semaphore** **fillCount = 0**, **emptyCount = 1**

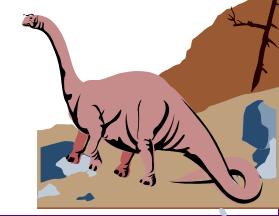
Producer:

```
do {  
    ...  
    produce an item in nextp  
    ...  
    wait(emptyCount);  
    add nextp to buffer  
    signal(fillCount);  
} while (1);
```

Consumer:

```
do {  
    wait(fillCount);  
    remove an item from  
    buffer to nextc  
    signal(emptyCount);  
    ...  
    consume the item in nextc  
    ...  
} while (1);
```

Please give out your reasons.





信号量：生产者消费者习题

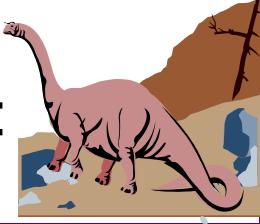
考虑三个吸烟者进程和一个经销商进程的系统

- ◆ 每个吸烟者连续不断地做烟卷并抽他做好的烟卷，做一支烟卷需要烟草、纸和火柴三种原料。
- ◆ 这三个吸烟者分别掌握有烟草、纸和火柴。
- ◆ 经销商源源不断地提供上述三种原料，但他只随机的将其中的两种原料组合（A:烟草+纸，B:纸+火柴，C:烟草+火柴）放在桌上，具有另一种原料的吸烟者就可以做烟卷并抽烟，且在做完后给经销商发信号，然后经销商再拿出两种原料放在桌上，如此反复

基于信号量设计一个同步算法描述他们的活动

✓ Semaphore empty=1, fullA=0, fullB=0, fullC=0;

✓ Broker: SmokerA: SmokerB: SmokerC:



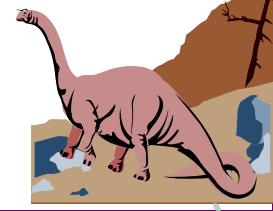


信号量：生产者消费者习题

可以考虑：设置三个信号量fullA、fullB和fullC，分别代表三种原料组合，初值均为0，即

- ◆ fullA表示烟草和纸的组合，
- ◆ fullB表示纸和火柴的组合，
- ◆ fullC表示烟草和火柴的组合。

桌面上一次只能放一种组合，可以看作是只能放一个产品的共享缓冲区，设置信号量empty初值为1，控制经销商往桌子上放原料





信号量：生产者消费者习题

算法

Semaphore fullA=fullB=fullC=0, empty=1;

process smokerA() {

do {

wait(fullA);

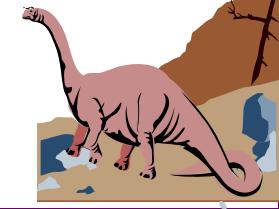
take tobacco and paper from the table;

signal(empty); // signal an empty table event

make cigarette;

smoke cigarette;

} while (1);

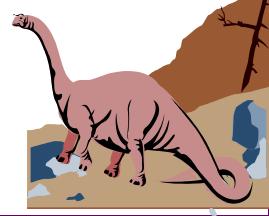




信号量：生产者消费者习题

```
process smokerB() {  
    do {  
        wait(fullB);  
        take paper and match  
        from the table;  
        signal(empty);  
        make cigarette;  
        smoke cigarette;  
    } while (1);  
}
```

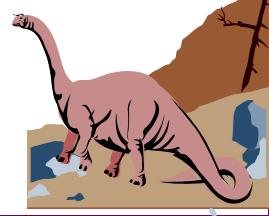
```
process smokerC() {  
    do {  
        wait(fullC);  
        take tobacco and match  
        from the table;  
        signal(empty);  
        make cigarette;  
        smoke cigarette;  
    } while (1);  
}
```





信号量：生产者消费者习题

```
process provider( ) {  
    integer i;  
    do {  
        i = random() % 3; // produce a combination  
        wait(empty); // wait for an empty table event  
        switch(i) {  
            case 0: put T&P on table; signal(fullA); break;  
            case 1: put P&M on table; signal(fullB); break;  
            case 2: put T&M on table; signal(fullC); break;  
        }  
    } while(1);
```

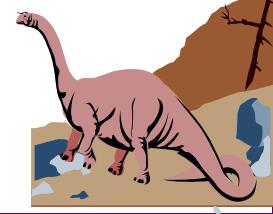




Reader-Writer Locks

Imagine a number of concurrent operations, including **reads** and **writes**.

- ◆ Writes change the state of the data
- ◆ Reads do not.
 - ✓ Many reads can proceed concurrently, as long as we can guarantee that no write is on-going.





Readers-Writers Problem (or Shared-Lock Problem)

Shared data

semaphore mutex, wrt;

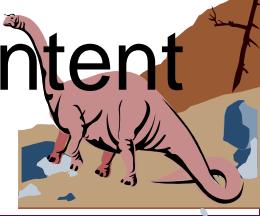
int readcount;

Initially **mutex = 1, wrt = 1, readcount = 0**

readcount: the number of readers browsing the shared content

mutex: guarantee the mutual exclusive access to the readcount variable

wrt: the right of modifying the shared content





Readers-Writers Problem (solution 1)

Writer Process

```
wait(wrt);
```

...

writing is performed

...

```
signal(wrt);
```

Question: Is this solution correct?

Reader Process

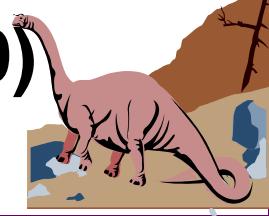
```
wait(mutex);  
readcount++;  
signal(mutex);  
if (readcount == 1)  
    wait(wrt);
```

...

reading is performed

...

```
wait(mutex);  
readcount--;  
signal(mutex);  
if (readcount == 0)  
    signal(wrt);
```





Readers-Writers Problem (solution 2)

Writer Process

```
wait(wrt);
```

...

writing is performed

...

```
signal(wrt);
```

Reader Process

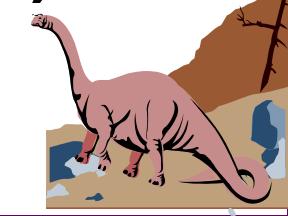
```
wait(mutex);  
readcount++;  
if (readcount == 1)  
    wait(wrt);  
signal(mutex);
```

...

reading is performed

...

```
wait(mutex);  
readcount--;  
if (readcount == 0)  
    signal(wrt);  
signal(mutex);
```



```
typedef struct _rwlock_t {  
    sem_t * writelock;  
    sem_t * lock;  
    int readers;  
} rwlock_t;  
  
void rwlock_acquire_readlock(rwlock_t * rw) {  
    sem_wait(rw->lock);  
    rw->readers++;  
    if (rw->readers == 1)  
        sem_wait(rw->writelock);  
    sem_post(rw->lock);  
}  
  
void rwlock_release_readlock(rwlock_t * rw) {  
    sem_wait(rw->lock);  
    rw->readers--;  
    if (rw->readers == 0)  
        sem_post(rw->writelock);  
    sem_post(rw->lock);  
}  
  
void rwlock_acquire_writelock(rwlock_t *rw) {  
    sem_wait(rw->writelock);  
}  
  
void rwlock_release_writelock(rwlock_t *rw) {  
    sem_post(rw->writelock);  
}
```

Give a demo



Exercise

由于读者优先，存在写者饥饿问题

R3: read 21	R2: read 21	R3: read 21	R1: read 21
R1: read 21	R4: read 21	R1: read 21	R2: read 21
R2: read 21	R5: read 21	R2: read 21	R4: read 21
R4: read 21	R3: read 21	R4: read 21	R2: done 21
R5: read 21	R1: read 21	R5: read 21	R5: read 21
R3: read 21	R2: read 21	R3: read 21	R1: read 21
R1: read 21	R4: read 21	R1: read 21	R4: read 21
R2: read 21	R5: read 21	R2: read 21	R1: done 21
R4: read 21	R3: read 21	R4: read 21	R5: read 21
R5: read 21	R1: read 21	R5: read 21	R4: done 21
R3: read 21	R2: read 21	R3: read 21	R5: done 21
R1: read 21	R4: read 21	R1: read 21	W2: write 22
R2: read 21	R5: read 21	R2: read 21	W1: write 23
R4: read 21	R3: read 21	R4: read 21	W2: write 24
R5: read 21	R1: read 21	R5: read 21	W1: write 25
R3: read 21	R2: read 21	R3: read 21	W2: write 26
R1: read 21	R4: read 21	R1: read 21	W1: write 27

用信号量解决无饥饿的读者——写者问题。



More Info about Reader-Writer Locks

The first readers-writers problem

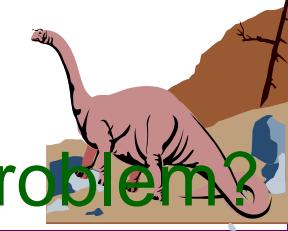
- ◆ requires that, no reader be kept waiting unless a writer has already obtained permission to use the shared object.

The second readers-writers problem

- ◆ requires that, once a writer is ready, that writer perform its write as soon as possible.

Discussion:

- ◆ Which problem is solved by the previous codes?
- ◆ The first readers-writers problem.
- ◆ How to solve the second readers-writers problem?





The (No-starve) Readers-Writers Problem

```
semaphore lock= 1;  
semaphore writelock=1;  
int read_count = 0;  
semaphore wflock =1;
```

```
void write() {  
    do {  
        wait(wflock);  
        wait(writelock);  
        /* writing */  
        signal(writelock);  
        signal(wflock);  
    }  
    while (1);  
}
```

```
void read() {  
    do {  
        wait(wflock);  
        signal(wflock);  
        wait(lock);  
        read_count ++;  
        if (read_count == 1)  
            wait(writelock);  
        signal(lock);  
        /* reading */  
        wait(lock);  
        read_count --;  
        if (read_count == 0)  
            signal(writelock);  
        signal(lock);  
    }  
    while (1);  
}
```



```
typedef struct  
_rwlock_t {  
    sem_t * writelock;  
    sem_t * lock;  
    int readers;  
    sem_t * wflock;  
} rwlock_t;  
  
void  
rwlock_init(rwlock_t *  
rw) {  
    rw->readers = 0;  
    rw->lock =  
sem_open(..., 1);  
    rw->writelock =  
sem_open(..., 1);  
    rw->wflock =  
sem_open(..., 1);  
}
```

```
void rwlock_acquire_readlock(rwlock_t * rw) {  
    sem_wait(rw->wflock);  
    sem_post(rw->wflock);  
    sem_wait(rw->lock);  
    rw->readers++;  
    if (rw->readers == 1)  
        sem_wait(rw->writelock);  
    sem_post(rw->lock);  
}  
  
void rwlock_release_readlock(rwlock_t * rw) {  
    sem_wait(rw->lock);  
    rw->readers--;  
    if (rw->readers == 0)  
        sem_post(rw->writelock);  
    sem_post(rw->lock);  
}  
  
void rwlock_acquire_writelock(rwlock_t *rw) {  
    sem_wait(rw->wflock);  
    sem_wait(rw->writelock);  
}  
  
void rwlock_release_writelock(rwlock_t *rw) {  
    sem_post(rw->writelock);  
    sem_post(rw->wflock);  
}
```

Give a demo



The (Writer-priority) Readers-Writers Problem

```
void write() {  
    do {  
        wait(writecount_lock);  
        write_count++;  
        if (write_count == 1)  
            wait(readlock);  
        signal(writecount_lock);  
        wait(writelock);  
  
        /* writing */  
  
        signal(writelock);  
        wait(writecount_lock);  
        write_count--;  
        if (write_count == 0)  
            signal(readlock);  
        signal(writecount_lock);  
    }  
    while (1);  
}
```

Give a demo

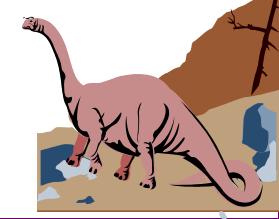
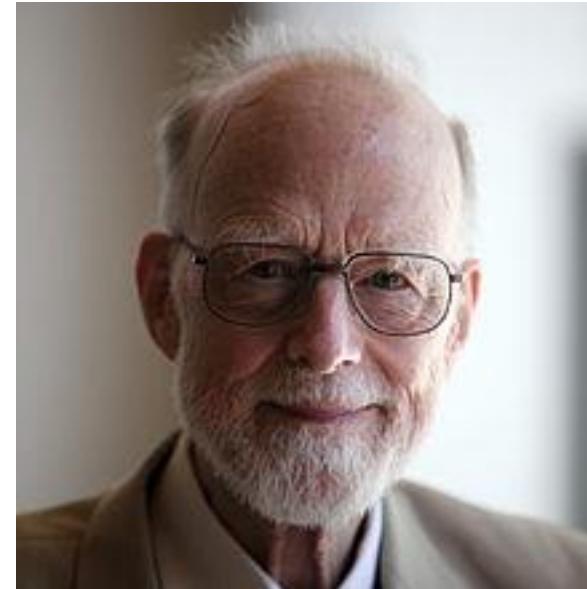
```
int write_count = read_count = 0;  
semaphore readcount_lock= 1;  
semaphore writecount_lock= 1;  
semaphore readlock=1; //0表示不能读  
semaphore writelock=1; //0表示不能写
```

```
void read() {  
    do {  
        wait(readlock);  
        wait(readcount_lock);  
        read_count++;  
        if (read_count == 1)  
            wait(writelock);  
        signal(readcount_lock);  
        signal(readlock);  
  
        /* reading */  
  
        wait(readcount_lock);  
        read_count--;  
        if (read_count == 0)  
            signal(writelock);  
        signal(readcount_lock);  
    }  
    while (1);  
}
```



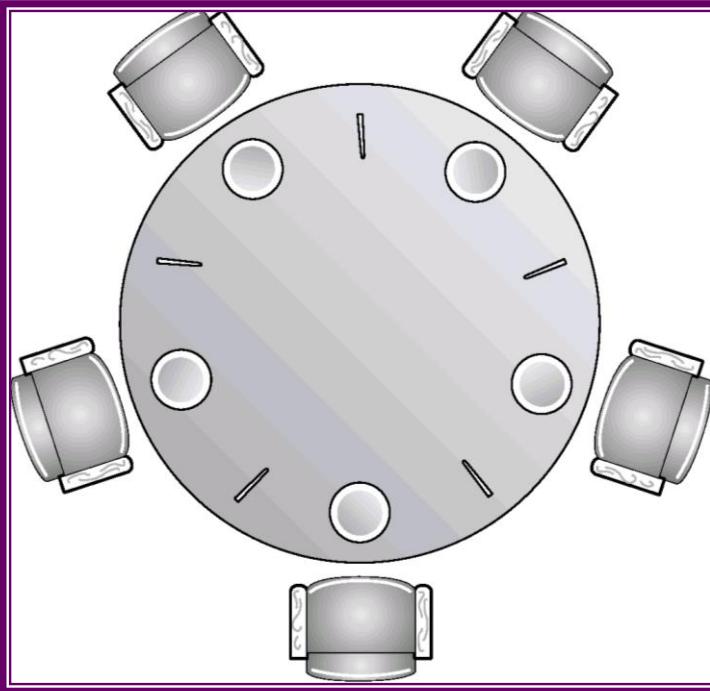
The Dining Philosophers

- Originally formulated in 1965 by Edsger Dijkstra
- Tony Hoare gave the problem its present formulation





Dining-Philosophers Problem



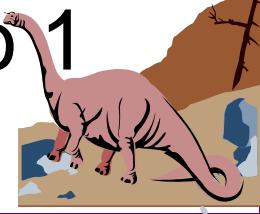
Here is the basic loop of each philosopher:

```
while (1) {  
    think();  
    getforks();  
    eat();  
    putforks();  
}
```

Shared data

semaphore chopstick[5];

Initial values of all semaphores are set to 1





Dining-Philosophers Problem

Philosopher i :

do {

wait(chopstick[i]);

wait(chopstick[(i+1) % 5]);

 ...

 eat

 ...

signal(chopstick[i]);

signal(chopstick[(i+1) % 5]);

 ...

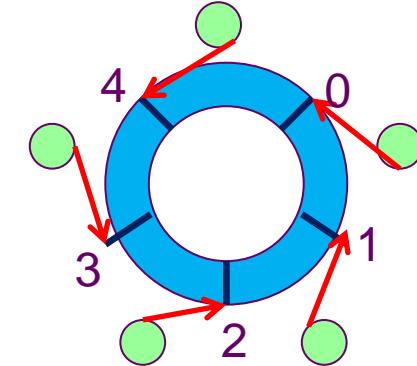
 think

 ...

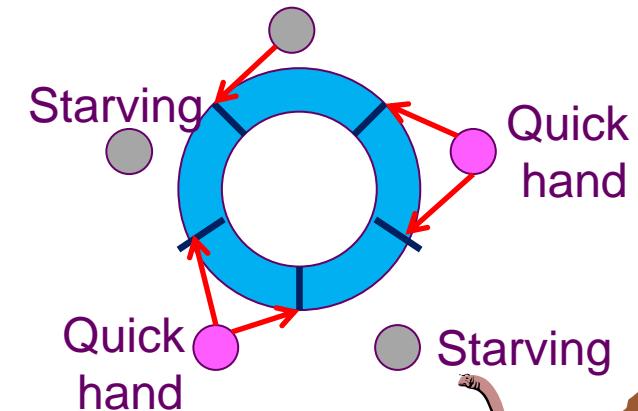
} while (1);

Challenges

◆ Deadlock



◆ Starvation



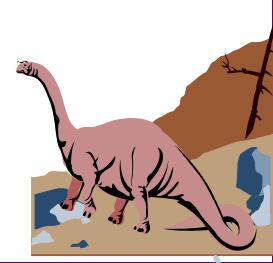
◆ Lack of Fairness





Semaphore 学习的四重境界

1. 理解基础概念
2. 熟练掌握经典问题（PC, RW, DP）。
3. 熟悉经典问题的变种，能够将应用题恰当的归约到某个经典问题的变种。
4. 能够将经典问题灵活组合应用，随心所欲，信手拈来。





Chapter 6: Process Synchronization

Background

The Critical-Section Problem

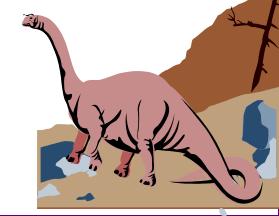
Synchronization Hardware

Semaphores

Classical Problems of Synchronization

Condition Variables and Monitors

Synchronization Examples





Condition Variable

Semaphore and condition variables are very similar and are used mostly for the same purposes.

- ◆ Semaphore can be easily understood as an in-kernel counter for the units of a type of resource.
- ◆ Condition is an advanced event notification tech.

However, there are minor differences that could make one preferable.

- ◆ For example, to implement barrier synchronization, you would not be able to use a semaphore. But a condition variable is ideal.





Condition Variable

The condition variable mechanism allows threads to suspend execution and relinquish the processor until some condition is true.

Semaphore = counter + mutex + waiting list

Conditional Variable = waiting list

A problem of semaphore: We cannot read the in-kernel counter hiding inside a semaphore

A condition variable must be used inside a mutex to avoid a race condition created by one thread preparing to wait and another thread which may signal the condition before the first thread actually waits on it resulting in a deadlock.



大家工作中条件变量会用的更多

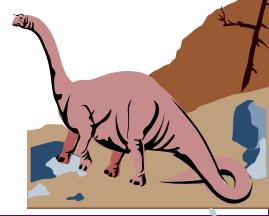
Java在企业开发市场占比80%

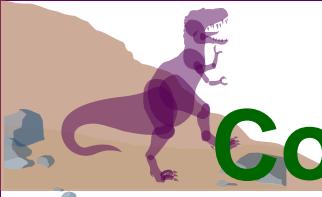
Java最常用的同步机制

- ◆ (1) `synchronized`关键字实现的条件变量。每一个Java对象就有一把看不见的锁，称为内部锁或者Monitor锁，内部。
- ◆ (2) `Lock`接口及其实现类，如
`ReentrantLock`.`ReadLock`和
`ReentrantReadWriteLock`.`WriteLock`。

不可不说的Java“锁”事

<https://tech.meituan.com/2018/11/15/java-lock.html>





Condition Variable vs. Semaphore

Semaphore	Condition Variable
Can be used anywhere	Must be used inside the protection of a mutex
<code>wait()</code> does not always block its caller	<code>wait()</code> always blocks its caller
<code>signal()</code> either releases a process, or increases the semaphore counter	<code>signal()</code> either releases a process, or the signal is lost as if it never occurs
If <code>signal()</code> releases a process, the caller and the released both continue	If <code>signal()</code> releases a process, either the caller or the released continues, but not both



Condition Variable in Pthread Library

Creating/Destroying:

- ◆ `pthread_cond_t cond = THREAD_COND_INITIALIZER;`
- ◆ `pthread_cond_init`
- ◆ `pthread_cond_destroy`

Waiting on condition:

- ◆ `pthread_cond_wait`(`pthread_cond_t *cond, pthread_mutex_t *mutex`) - unlocks the mutex and waits for the condition variable `cond` to be signaled.

Waking thread based on condition:

- ◆ `pthread_cond_signal`(`pthread_cond_t *cond`) - restarts one of the threads that are waiting on the condition variable `cond`.
- ◆ `pthread_cond_broadcast`(`pthread_cond_t *cond`) - wake up all threads blocked by the specified condition variable.

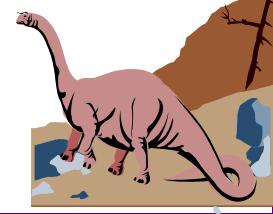




Barrier Problem

Suppose we wanted to perform a multi-threaded calculation that has two stages, but we don't want to advance to the second stage until the first stage is completed.

We could use a synchronization method called a **barrier**. When a thread reaches a barrier, it will wait at the barrier until all the threads reach the barrier, and then they'll all proceed together.



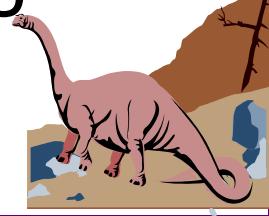


Barrier Problem

Pthreads has a **pthread_barrier_wait()** function that implements this. You'll need to declare a **pthread_barrier_t** variable and initialize it with **pthread_barrier_init()**.

- ◆ **pthread_barrier_init()** takes the number of threads that will be participating in the barrier as an argument.

Now let's implement our own barrier and use it to keep all the threads in sync in a large calculation.

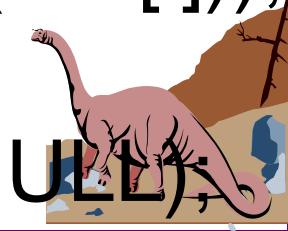




Barrier Implementation by Condition Variable

```
#define N (16)
double data[256][8192] ;
pthread_mutex_t m;
pthread_cond_t cv;
int main() {
    int tids[N], i;
    pthread_mutex_init(&m, NULL);
    pthread_cond_init(&cv, NULL);
    for(i = 0; i < N; i++) { tids[i] = i;
        pthread_create(&ids[i], NULL, calc, &(tids[i]));
    }
    for(i = 0; i < N; i++) pthread_join(ids[i], NULL);
}
```

<https://github.com/angrave/SystemProgramming/wiki/Synchronization%2C-Part-6%3A-Implementing-a-barrier>





Barrier Implementation by Condition Variable

```
double data[256][8192]
```

```
void *calc(void *ptr) {
```

1. Threads do first calculation (use and change values in data)
2. Barrier! Wait for all threads to finish first calculation before continuing
3. Threads do second calculation (use and change values in data)

```
}
```

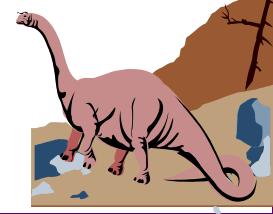
<https://github.com/angrave/SystemProgramming/wiki/Synchronization%2C-Part-6%3A-Implementing-a-barrier> 6.115



If using condition variable, the state of counter can be access. But when using semaphore, the state of inner count cannot be accessed.

```
#int remain = N;  
void *calc(void *ptr) {  
    // The thread does first calculation  
    pthread_mutex_lock(&m);  
    remain--;  
    if (remain == 0) pthread_cond_broadcast(&cv);  
    else {  
        while(remain != 0) pthread_cond_wait(&cv,&m);  
    }  
    pthread_mutex_unlock(&m);  
    // The thread does second calculation
```

Give a demo

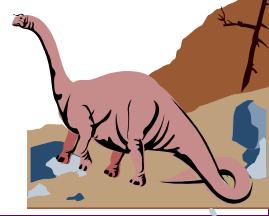




Object-Oriented Monitors

High-level synchronization construct that allows the safe sharing of an abstract data type among concurrent processes.

```
monitor monitor-name
{   shared variable declarations
    procedure body P1 (...) {
        ...
    }
    procedure body P2 (...) {
        ...
    }
    procedure body Pn (...) {
        ...
    }
    { initialization code }
}
```





Monitors: Mutual Exclusion

No more than one process can be executing *within* a monitor. Thus, *mutual exclusion* is guaranteed within a monitor.

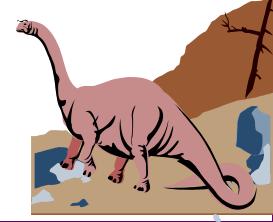
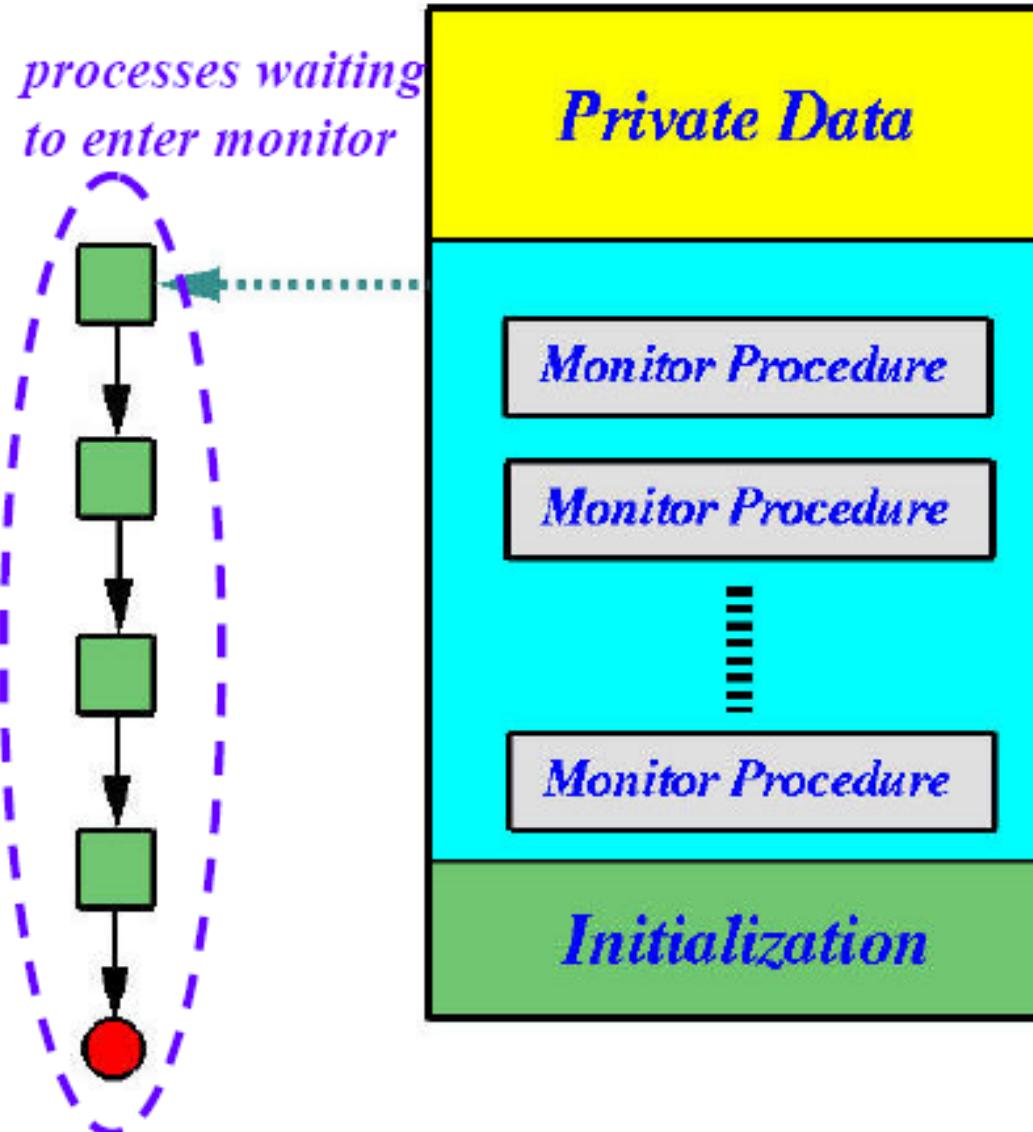
When a process calls a monitor procedure and enters the monitor successfully, it is the *only* process executing in the monitor.

When a process calls a monitor procedure and the monitor has a process running, the caller will be blocked *outside of the monitor*.





Schematic View of a Monitor





Monitors: Event Notification

To allow a process to wait within the monitor, a **condition** variable must be declared, as

condition x, y;

Condition variable can only be used with the operations **wait** and **signal**.

- ◆ The operation

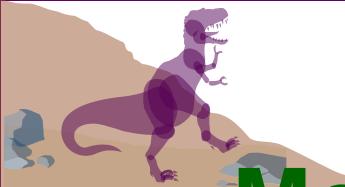
x.wait();

means that the process invoking this operation is blocked until another process invokes

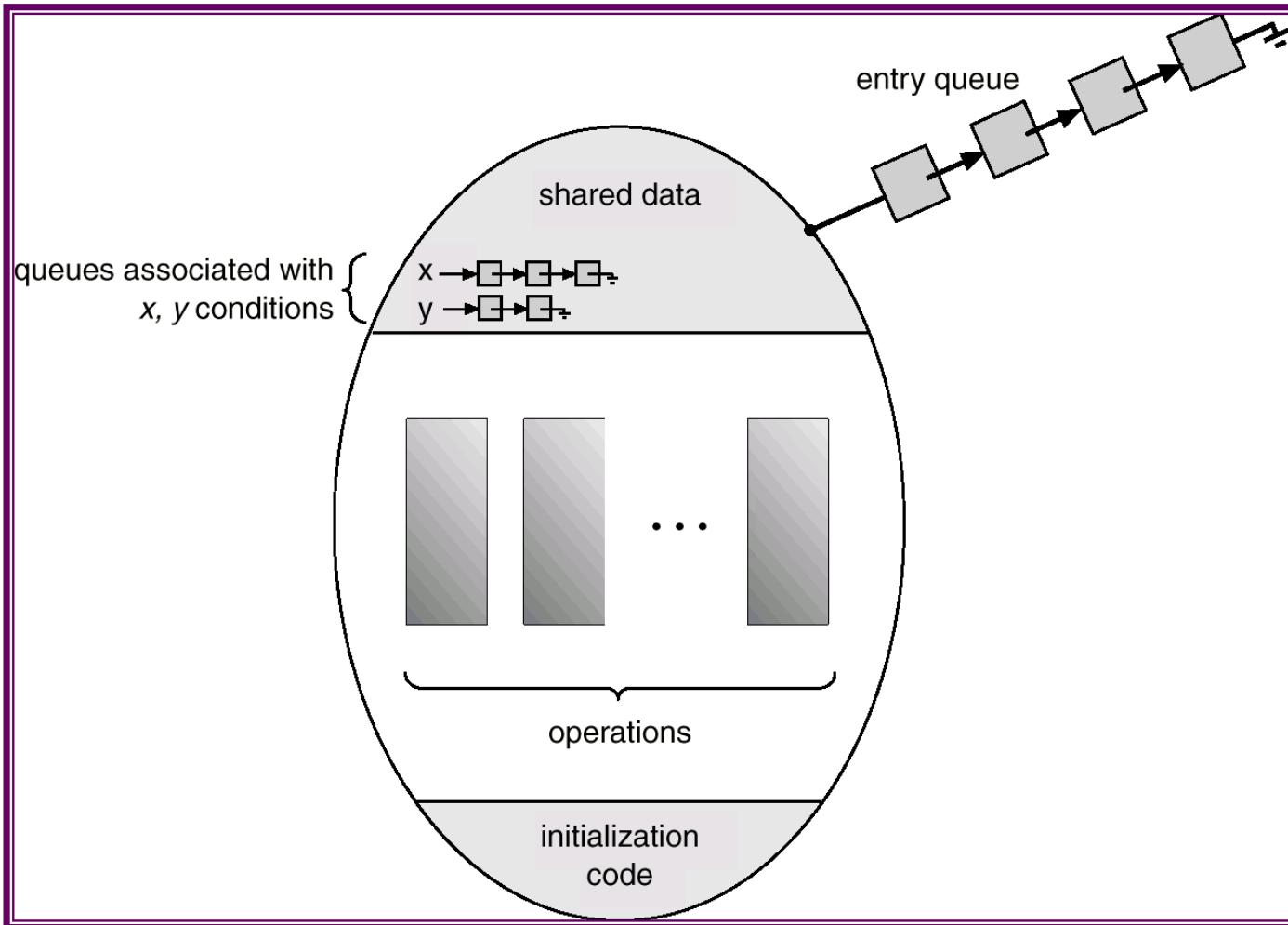
x.signal();

- ◆ The **x.signal** operation wakeup exactly one blocked process. If no process is waiting for the condition, then the **signal** operation has no effect.





Schematic View of a Monitor With Condition Variables





A Subtle Issue of Condition Variable

Consider the released process (from the signaled condition) and the process that signals. There are **two processes executing in the monitor, and mutual exclusion is violated!**

There are two common and popular approaches to address this problem:

- ◆ The released process takes over the monitor and the signaling process waits somewhere.
- ◆ The released process waits somewhere and the signaling process continues to use the monitor.





Java's Monitor Supports

Synchronized methods for mutual exclusion

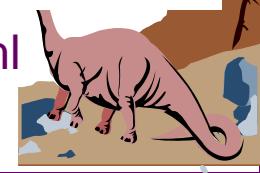
```
class classname {  
    synchronized return_type methodname() {.....}  
}
```

Coordination support for event notification

Method	Description
void Object.wait();	Enter a monitor's wait set until notified by another thread
void Object.wait(long timeout);	Enter a monitor's wait set until notified by another thread or timeout milliseconds elapses
void Object.notify();	Wake up one thread waiting in the monitor's wait set. (If no threads are waiting, do nothing.)
void Object.notifyAll();	Wake up all threads waiting in the monitor's wait set. (If no threads are waiting, do nothing.)

<http://www.ibm.com/developerworks/cn/java/j-lo-synchronized/index.html>

Operating System Concepts 6.123 <http://www.artima.com/insidejvm/ed2/threadsynchP.html> Southeast University





Producer-Consumer Example

```
procedure producer() {  
    do {  
        item = producItem();  
        PCbuffer.add(item);  
    } while (true);  
}  
  
procedure consumer() {  
    do {  
        item = PCbuffer.remove();  
        consumeItem(item);  
    } while (true);  
}
```

```
monitor PCbuffer {  
    int itemCount; // <= BUFSIZE  
    condition full, empty;  
    putItemIntoBuffer(item) {...}  
    Item removeItemFromBuffer()  
    {...}  
    procedure void add(item) {  
        ... // how to implement?  
    }  
    procedure item remove() {  
        ... // how to implement?  
    }  
}
```





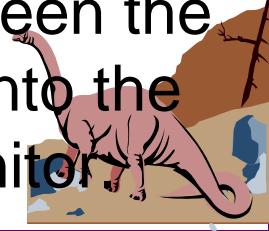
Producer-Consumer Example

```
procedure void add(item) {  
    if (itemCount == BUFSIZE)  
        full.wait();  
  
    putItemIntoBuffer(item);  
    itemCount = itemCount + 1;  
    if (itemCount == 1)  
        empty.signal();  
  
    return;  
}
```

```
procedure item remove() {  
    if (itemCount == 0)  
        empty.wait();  
  
    item = removeItemFromBuffer();  
    itemCount = itemCount - 1;  
    if (itemCount == BUFSIZE - 1)  
        full.signal();  
  
    return item;  
}
```

Note that **if** statement has been used in the above code, both when testing if the buffer is full or empty.

With multiple consumers, there is a race condition between the consumer who gets notified that an item has been put into the buffer and another consumer who is waiting on the monitor.





Producer-Consumer Example

```
procedure void add(item) {  
    while (itemCount == BUFSIZE)  
        full.wait();  
  
    putItemIntoBuffer(item);  
    itemCount = itemCount + 1;  
    if (itemCount == 1)  
        empty.signal();  
  
    return;  
}
```

```
procedure item remove() {  
    while (itemCount == 0)  
        empty.wait();  
  
    item = removeItemFromBuffer();  
    itemCount = itemCount - 1;  
    if (itemCount == BUFSIZE - 1)  
        full.signal();  
  
    return item;  
}
```

Note that **while** statement has been used in the above code, both when testing if the buffer is full or empty.

With multiple consumers, there is a race condition between the consumer who gets notified that an item has been put into the buffer and another consumer who is waiting on the monitor.





Producer-Consumer Example

```
procedure void add(item) {  
    while (itemCount == BUFSIZE)  
        full.wait();  
    putItemIntoBuffer(item);  
    itemCount = itemCount + 1;  
    if (itemCount == 1)  
        empty.signal();  
    return;  
}
```

```
procedure item remove() {  
    while (itemCount == 0)  
        empty.wait();  
    item = removeItemFromBuffer();  
    itemCount = itemCount - 1;  
    if (itemCount == BUFSIZE - 1)  
        full.signal();  
    return item;  
}
```

With multiple producers, there is also a [race condition](#) between the producer who gets notified that the buffer is no longer full and another producer is already waiting on the monitor.

If the **while** was instead an **if**, too many items might be put into the buffer or a remove might be attempted on an empty buffer.



Dining Philosophers without Deadlock

```
monitor dining_philosopher_sync_table
{
    enum {thinking, hungry, eating} state[5];
    condition self[5];
    procedure void pickup(int i) ; // pick up chopsticks
    procedure void putdown(int i) ; // put down chopsticks
    private void test(int i) ; // test if Pi is eligible for eating
    void init() {
        for (int i = 0; i < 5; i++)
            state[i] = thinking;
    }
}
```



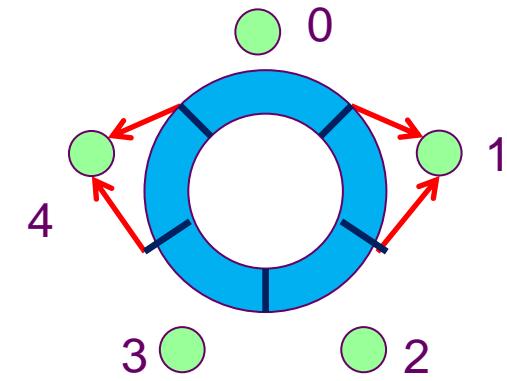
Dining Philosophers without Deadlock

```
void pickup(int i) {  
    state[i] = hungry;  
    test(i);  
    while(state[i] != eating)  
        self[i].wait();  
}  
}
```

```
void putdown(int i) {  
    state[i] = thinking;  
    test((i+4) % 5); // left  
    test((i+1) % 5); // right  
}
```

The code has NO deadlock!!! Why?

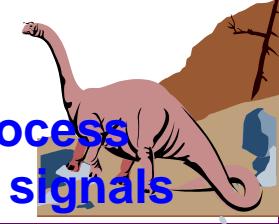
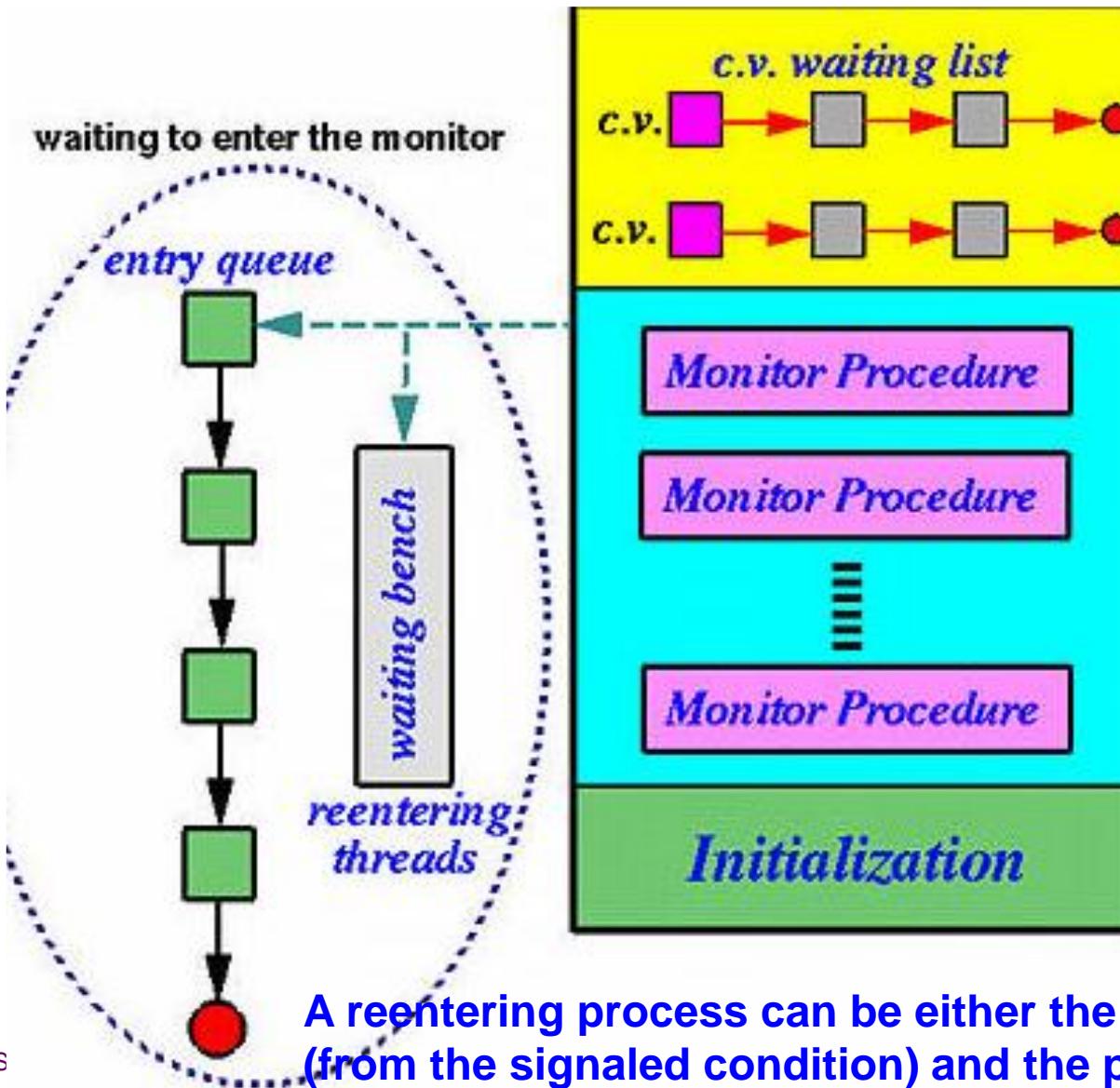
```
void test(int i) {  
    if ( (state[(i + 4) % 5] != eating) &&  
        (state[i] == hungry) &&  
        (state[(i + 1) % 5] != eating)) {  
        state[i] = eating;  
        self[i].signal();  
    }  
}
```



When P_1 and P_4 finish eating at the same time, will P_2 and P_3 compete for their common chopstick after their wakeup?



Another Subtle Issue of Monitor: Queue of Reentering Threads/Proc





For Better Understanding, Let's Implement Monitor by Semaphores

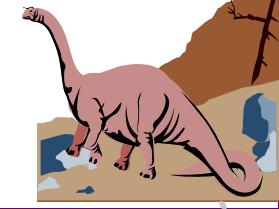
Variables

```
semaphore mutex; // (initially = 1)
semaphore next;   // (initially = 0)
int next-count = 0;
```

Each external procedure F will be replaced by

```
wait(mutex);
    ...
    ... // body of  $F$ ;
if (next-count > 0)
    signal(next);
else
    signal(mutex);
```

Mutual exclusion within a monitor is ensured.





Monitor Implementation Using Semaphores

For each condition variable x , we have:

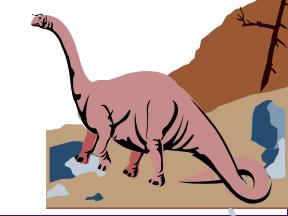
```
semaphore x-sem; // (initially = 0)  
int x-count = 0;
```

The operation **$x.wait$** can be implemented as:

```
x-count++;  
if (next-count > 0)  
    signal(next);  
else  
    signal(mutex);  
wait(x-sem);  
x-count--;
```

The operation **$x.signal$** can be implemented as:

```
if (x-count > 0) {  
    next-count++;  
    signal(x-sem);  
    wait(next);  
    next-count--;  
}
```

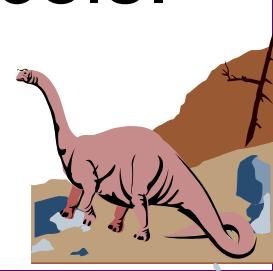




Monitor Implementation (Cont.)

Check two conditions to establish correctness of system:

- ◆ User processes must always make their calls on the monitor in a correct sequence.
- ◆ Must ensure that an uncooperative process does not ignore the mutual-exclusion gateway provided by the monitor, and try to access the shared resource directly, without using the access protocols.

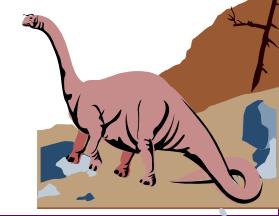




Condition Enhanced with a Priority Number

*Conditional-wait construct: **x.wait(c);***

- ◆ **c** – integer expression evaluated when the **wait** operation is executed.
- ◆ value of **c** (*a priority number*) stored with the name of the process that is suspended.
- ◆ when **x.signal** is executed, process with smallest associated priority number is resumed next.





Chapter 6: Process Synchronization

Background

The Critical-Section Problem

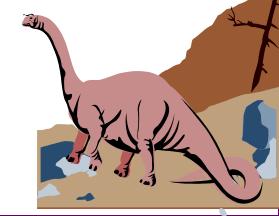
Synchronization Hardware

Semaphores

Classical Problems of Synchronization

Condition Variables and Monitors

Synchronization Examples





Solaris 2 Synchronization

Implements a variety of locks to support multitasking, multithreading (including real-time threads), and multiprocessing.

Uses *adaptive mutexes* for efficiency when protecting data from short code segments.

Uses *condition variables*, *semaphore*, and *readers-writers locks* when longer sections of code need access to data.

Uses *turnstile*s to order the list of threads waiting to acquire either an adaptive mutex or reader-writer lock.

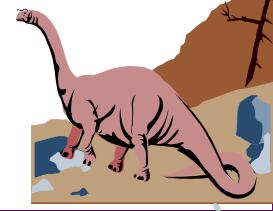




Adaptive Mutex

Most operating systems
(including [Solaris](#), [Mac OS X](#) and [FreeBSD](#))
use a hybrid approach called
"adaptive [mutex](#)". The idea is to use a
spinlock when trying to access a resource
locked by a currently-running thread, but to
sleep if the [thread](#) is not currently running.
(The latter is *always* the case on single-
processor systems.)

<https://en.wikipedia.org/wiki/Spinlock#Alternatives>





Windows XP Synchronization

Uses interrupt masks to protect access to global resources on uniprocessor systems.

Uses *spinlocks* on multiprocessor systems.

Also provides *dispatcher objects* which may act as mutexes and semaphores.

Dispatcher objects may also provide *events*. An event acts much like a condition variable.

