

REGULAR LANGUAGE

Q1. Consider the following grammar G:

$S \rightarrow aSa \mid$
 aAa
 $A \rightarrow Bb$
 $B \rightarrow c$

The grammar G belongs to which type of Chomsky classification?

- (A) Type 3
- (B) Type 2 but not type 3
- (C) Type 1 but not type 2
- (D) Type 0 but not type 1

Q2. Consider the following grammar G with start symbol S

$S \rightarrow a \ SC \mid T$
 $T \rightarrow bTd \mid R$
 $dC \rightarrow Cd$
 $RC \rightarrow Cr$
 $R \rightarrow \epsilon$

What is the highest type number that can be assigned to the following grammar?

- (A) Type 0
- (B) Type 1
- (C) Type 2
- (D) Type 3

Q3. Consider the following grammar G with start symbol S

$S \rightarrow Aa,$
 $A \rightarrow Ba,$
 $B \rightarrow abc.$

What is the highest type number that can be assigned to the following grammar?

- (A) Type 0
- (B) Type 1
- (C) Type 2
- (D) Type 3

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| <p>Q4.</p> | <p>What is the highest type number that can be assigned to the following grammar?</p> <div style="border: 1px solid black; padding: 10px; margin: 10px 0;"> $S \rightarrow aA \mid bB$ $A \rightarrow aB \mid bC$ $B \rightarrow aB \mid bB$ $C \rightarrow aC \mid bC \mid \lambda$ </div> <p>(A) Type 0 (B) Type 1 (C) Type 2 (D) Type 3</p> |
| <p>Q5.</p> | <p>Consider the following grammar G with starting symbol S:</p> <div style="border: 1px solid black; padding: 10px; margin: 10px 0;"> $S \rightarrow bTbb$ $T \rightarrow bTbb \mid Acccb$ $A \rightarrow aAc \mid \lambda$ </div> <p>What is the highest type number that can be assigned to the following grammar?</p> <p>(A) Type 0 (B) Type 1 (C) Type 2 (D) Type 3</p> |
| <p>Q6.</p> | <p>The following grammar</p> <div style="border: 1px solid black; padding: 10px; margin: 10px 0;"> $G = (N, T, P, S)$ $N = \{S, A, B, C, D, E\}$ $T = \{a, b, c\}$ $P : S \rightarrow aAB$ $AB \rightarrow CD$ $CD \rightarrow CE$ $C \rightarrow aC$ $C \rightarrow b$ $bE \rightarrow bc \text{ is}$ </div> <p>(A) is type 3 (B) is type 2 but not type 3 (C) is type 1 but not type 2 (D) is type 0 but not type 1</p> |

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| <p>Q7.</p> | <p>The following grammar</p> <div style="border: 1px solid black; padding: 10px; margin: 10px 0;"> $G = (N, T, P, S)$ $N = \{S, A, B, C, D, E\}$ $T = \{a, b, c\}$ $P : S \rightarrow ABCD$ $BCD \rightarrow DE$ $D \rightarrow aD$ $D \rightarrow a$ $E \rightarrow bE$ $E \rightarrow c$ </div> <p>(A) is type 3</p> <p>(B) is type 2 but not type 3</p> <p>(C) is type 1 but not type 2</p> <p>(D) is type 0 but not type 1</p> |
| <p>Q8.</p> | <p>A grammar has the following productions:</p> $S \rightarrow aSb \mid a \mid bSa$ <p>Which of the following sentences are in the language that is generated by this grammar?</p> <p>(A) aaaaabb</p> <p>(B) aabbaabb</p> <p>(C) bbbaabbaa</p> <p>(D) All of the answers above are correct</p> |
| <p>Q9.</p> | <p>Consider the grammar below, with start symbol S.</p> <div style="border: 1px solid black; padding: 10px; margin: 10px 0;"> $S \rightarrow AS \mid SB \mid \lambda$ $A \rightarrow Aa \mid a$ $B \rightarrow Bb \mid b$ </div> <p>Which of the following strings can't be generated by this grammar?</p> <p>(A) a</p> <p>(B) abb</p> <p>(C) abba</p> <p>(D) aaabbb</p> |

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| <p>Q10.</p> | <p>Consider the following grammar G:</p> $S \rightarrow AB \mid aB \quad A \rightarrow aab \mid \lambda \quad B \rightarrow bbA$ <p>If the language accepted by G is $L(G)$ then which of the following set of string is the subset of the $L(G)$?</p> <p>(A) {bbaa, abb, bb}</p> <p>(B) {bbaa, bba, abb}</p> <p>(C) {aab, abb}</p> <p>(D) None of these</p> |
| <p>Q11.</p> | <p>Consider the language $L = \{w : \text{for some } u \in \Sigma^*, w = u^R u\}$. Which of the following strings belongs to L?</p> <p>(i) aaabbb (ii) abab (iii) abba (iv) λ</p> <p>(A) (i) and (ii)</p> <p>(B) (iii) only</p> <p>(C) (iv) only</p> <p>(D) (iii) and (iv)</p> |
| <p>Q12.</p> | <p>Consider the grammar G:</p> <div style="border: 1px solid black; padding: 10px; margin: 10px 0;"> $\begin{aligned} S &\rightarrow AB \\ A &\rightarrow 0A1 \mid 2 \\ B &\rightarrow 1B \mid 3A \end{aligned}$ </div> <p>Which of the following strings is in $L(G)$?</p> <p>(A) 0211300021</p> <p>(B) 021300211</p> <p>(C) 00213021</p> <p>(D) 0021113002111</p> |
| <p>Q13.</p> | <p>Identify in the list below a sentence of length 6 that is generated by the grammar</p> $S \rightarrow (S)S \mid \lambda$ <p>(A) () () ()</p> <p>(B)))) (((</p> <p>(C))) (() (</p> <p>(D)) (() ()</p> |

Q14. Consider the following grammar G:

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| $S \rightarrow AB$ $A \rightarrow 0A1 \mid 2$ $B \rightarrow 1B \mid 3A$ |
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Which of the following strings are in $L(G)$?

- i) 021300211
- ii) 002111300211
- iii) 00211100211
- iv) 0021113002111
- (A) i and ii
- (B) iii only
- (C) iii and iv
- (D) i, ii and iv

Q15. A grammar is described as follow:

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| $S \rightarrow aS$ $S \rightarrow b$ $S \rightarrow bA$ $A \rightarrow bB$ $B \rightarrow a$ |
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Which of the following strings cannot be derived from the above grammar?

- (A) abba
- (B) abbb
- (C) bba
- (D) aab

Q16. Which of the following strings cannot be derived from the symbol S using the rules

$S \rightarrow SS \mid aaa \mid aaaaa ?$

- (A) aaaaaa
- (B) aaaaaaa
- (C) aaaaaaaaa
- (D) aaaaaaaaaa

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| Q17. | Consider the grammar given below |
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$$S \rightarrow xB \mid yA$$
$$A \rightarrow x \mid xS \mid yAA$$
$$B \rightarrow y \mid yS \mid xBB$$

Consider the following strings.

i. $xxyyxyxy$

ii. $yyxxyyxx$

iii. $yxxxyxy$

iv.xxyyxy

Which of the above strings are generated by the grammar?

(A) i and ii only

(B) ii , iii and iv only

(C) i, ii and iii only

(D) All the above

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| Q18. | Consider the following context-free grammar over the alphabet $\Sigma = \{a, b, c\}$ with S as the start symbol: |
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$$S \rightarrow abScT \mid abcT$$
$$T \rightarrow bT \mid b$$

Which of the following string is not generated by given grammar?

(A) ababcbbcbbbb

(B) ababcbcb

(C) abababcbcbcb

(D) ababcbcbcb

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| Q19. | Consider the following grammar |
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$$S \rightarrow XY \quad | \quad W$$
$$X \rightarrow aXb \mid \lambda$$
$$Y \rightarrow CY \quad | \quad \lambda$$
$$W \rightarrow aWc \quad | \quad Z$$
$$Z \rightarrow bZ \quad | \quad \lambda$$

What is the language generated by this grammar?

$$(A)\{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i = j = k \}$$
$$(B)\{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i = j \text{ or } i = k \}$$
$$(C) \{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i = j \text{ or } j = k \}$$
$$(D)\{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i \neq j \text{ or } i \neq k \}$$

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| <p>Q20.</p> | <p>Consider the following grammar</p> <div style="border: 1px solid black; padding: 5px; margin: 10px 0;"> $S \rightarrow aSc \mid X$ $X \rightarrow bXc \mid \lambda$ </div> <p>What is the language generated by this grammar?</p> <p>(A) $\{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i + j = k \}$</p> <p>(B) $\{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } i = j + k \}$</p> <p>(C) $\{ a^i b^j c^k \mid i, j, k \geq 0, \text{ and } k = i - j \}$</p> <p>(D) None of the above</p> |
| <p>Q21.</p> | <p>Consider the following language</p> <div style="border: 1px solid black; padding: 5px; margin: 10px 0;"> $S \rightarrow AX \mid YC$ $A \rightarrow aA \mid \lambda$ $C \rightarrow cC \mid \lambda$ $X \rightarrow bXc \mid \lambda$ $Y \rightarrow aYb \mid \lambda$ $L(G) \text{ is}$ </div> <p>(A) $L(a^*b^*c^*)$</p> <p>(B) $\{ a^n b^n c^n \mid n \geq 0 \}$</p> <p>(C) $\{ a^i b^j c^k \mid i = j \text{ or } j = k \}$</p> <p>(D) none of the above</p> |
| <p>Q22.</p> | <p>What is the language of the grammar with the following production rules?</p> $S \rightarrow ASb \mid c \quad A \rightarrow a$ <p>(A) $\{ a^n c b^n \mid n \in \mathbb{N} \}$</p> <p>(B) $\{ xcb \mid x \in \{a\}^* \}$</p> <p>(C) $\{ acy \mid y \in \{b\}^* \}$</p> <p>(D) All of the answers above are incorrect</p> |
| <p>Q23.</p> | <p>Which language generates the grammar G given by the productions?</p> <div style="display: flex; justify-content: space-between;"> $S \rightarrow aSa \mid aBa$ $B \rightarrow bB \mid b$ </div> <p>(A) $L(G) = \{ a^n b^m a^n \mid n > 0, m > 0 \}$.</p> <p>(B) $L(G) = \{ a^n b^m a^n \mid n > 0, m < 0 \}$.</p> <p>(C) $L(G) = \{ b a^n b \mid n > 0, m > 0 \}$.</p> <p>(D) None of these</p> |

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| <p>Q24.</p> | <p>Consider the following grammar G:</p> $S \rightarrow aSbb \mid \lambda$ <p>The language generated by the above grammar is:</p> <p>(A) $L = \{a^n b^n : n \geq 0\}$</p> <p>(B) $L = \{a^n b^{2n} : n \geq 0\}$</p> <p>(C) $L = \{a^m b^n : n \geq 0\}$</p> <p>(D) None</p> |
| <p>Q25.</p> | <p>The language generated by the following grammar is:</p> $S \rightarrow 0S1 \mid C \quad C \rightarrow 1C0 \mid \lambda$ <p>(A) $L1 = \{0^n 1^m 0^m 1^n \mid n, m \geq 0\}$</p> <p>(B) $L1 = \{0^m 1^{2n} 0^n 1^n \mid n, m \geq 0\}$</p> <p>(C) $L1 = \{0^n 1^m 0^{2m} 1^n \mid n, m \geq 0\}$</p> <p>(D) $L = \{0^n 1^n 1^m 0^m \mid n, m \geq 0\}$</p> |
| <p>Q26.</p> | <p>Consider the following grammar G with start symbol S over the alphabet $\Sigma = \{a, b\}$</p> $S \rightarrow aSaa \mid B$ $B \rightarrow bB \mid \lambda$ <p>The language generated by G is</p> <p>$L1 = \{a^n b^n a^{2m} \mid n, m \geq 0\}$</p> <p>$L2 = \{a^n b^{2m} \mid n, m \geq 0\}$</p> <p>$L3 = \{a^n b^m a^{2n} \mid n, m \geq 0\}$</p> <p>$L4 = \{a^n b^m a^{2m} \mid n, m \geq 0\}$</p> <p>(A) L1</p> <p>(B) L2</p> <p>(C) L3</p> <p>(D) L4</p> |
| <p>Q27.</p> | <p>Consider the following language:</p> $L = \{a^k b^{2k} \mid k \geq 2\}$ <p>Which of the following grammar generates L?</p> <p>(A) $S \rightarrow aSbb \mid \lambda$</p> <p>(B) $S \rightarrow aSbb \mid abb$</p> <p>(C) $S \rightarrow aSbb \mid aabbbb$</p> <p>(D) $S \rightarrow aX \mid bY \quad X \rightarrow aX \mid \lambda \quad Y \rightarrow bbY \mid \lambda$</p> |

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| <p>Q28.</p> | <p>Consider the following language:</p> $L = \{ab^n \mid n \geq 0\} \cup \{(ba)^m \mid m \geq 0\}$ <p>(A) $S \rightarrow abS \mid baS \mid \lambda$</p> <p>(B) $S \rightarrow aX \mid bY$ $X \rightarrow bX \mid \lambda$ $Y \rightarrow baY \mid \lambda$</p> <p>(C) $S \rightarrow aX \mid baY$ $X \rightarrow bX \mid \lambda$ $Y \rightarrow baY \mid \lambda$</p> <p>(D) None of the above</p> |
| <p>Q29.</p> | <p>Consider the Grammar G, with productions.</p> $S \rightarrow aA \mid \lambda, \quad A \rightarrow bS$ <p>Which of the following language are generated by G?</p> <p>(A) $L = \{a^n b^n \mid n > 0\}$</p> <p>(B) $L = \{a^n b^m \mid n > 0, b > 0\}$</p> <p>(C) $L = \{(ab)^n \mid n \geq 0\}$</p> <p>(D) $L = \{(ab)^n \mid n > 0\}$</p> |
| <p>Q30.</p> | <p>Which language generates the grammar G given by the productions</p> $S \rightarrow aSdd \mid A$ $A \rightarrow bAc \mid bc$ <p>(A) $L(G) = \{a^n b^m c^m d^{2n} \mid n \geq 0, m > 0\}$</p> <p>(B) $L(G) = \{a^m b^m c^n d^{2n} \mid n \geq 0, m > 0\}$</p> <p>(C) $L(G) = \{a^i b^m c^m d^{2n} \mid i > 0, n \geq 0, m > 0\}$</p> <p>(D) $L(G) = \{a^n b^i c^j d^{2n} \mid n \geq 0, i > 0, j > 0\}$</p> |
| <p>Q31.</p> | <p>Consider the following grammar G with start symbol S over the alphabet $\Sigma = \{a, b\}$</p> <div style="border: 1px solid black; padding: 10px; margin: 10px 0;"> $S \rightarrow aXa \mid bXb \mid a \mid b$ $X \rightarrow aX \mid bX \mid \lambda$ </div> <p>The language generated by G is</p> <p>(A) All strings that start and end with the same symbol.</p> <p>(B) All nonempty strings that start and end with the different symbol.</p> <p>(C) All nonempty strings that start and end with the same symbol.</p> <p>(D) None of the above.</p> |

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| <p>Q32.</p> | <p>The language generated by the following grammar is</p> <div style="border: 1px solid black; padding: 5px; margin: 10px 0;"> $\begin{aligned} S &\rightarrow aB \mid bA \\ A &\rightarrow a \mid aS \mid bAA \\ B &\rightarrow b \mid bS \mid aBB \end{aligned}$ </div> <p>(A) Strings contain equal number of a's and equal number of b's. (B) Strings contain odd number of a's and odd number of b's. (C) Strings contain odd number of a's and even number of b's. (D) Strings contain even number of a's and even number of b's.</p> |
| <p>Q33.</p> | <p>Consider the following grammar</p> $S \rightarrow 0S0 \mid 0S1 \mid 1S0 \mid 1S1 \mid 0$ <p>What is the language generated by this grammar?</p> <p>(A) $\{w \in \{0, 1\}^* \mid \text{the length of } w \text{ is odd} \}$ (B) $\{w \in \{0, 1\}^* \mid \text{the length of } w \text{ is odd and the middle symbol is } 0 \}$ (C) $\{w \in \{0, 1\}^* \mid \text{the length of } w \text{ is odd and the middle symbol is } 1 \}$ (D) $\{w \in \{0, 1\}^* \mid w \text{ contains } 0 \text{ in middle} \}$</p> |
| <p>Q34.</p> | <p>Consider the following grammar G with start symbol S over the alphabet $\Sigma = \{a, b\}$</p> <div style="border: 1px solid black; padding: 5px; margin: 10px 0;"> $\begin{aligned} S &\rightarrow Aa \mid MS \mid SMA \\ A &\rightarrow Aa \mid \lambda \\ M &\rightarrow \lambda \mid MM \mid bMa \mid aMb \end{aligned}$ </div> <p>The language generated by G is</p> <p>(A) All strings with more a's than b's. (B) All strings with one more a's than b's. (C) All strings with more b's than a's. (D) All strings with equal a's and b's.</p> |
| <p>Q35.</p> | <p>How many of the following is/are true? _____</p> <p>(i) $baa \in a^*b^*a^*b^*$ (ii) $b^*a^* \cap a^*b^* = a^* \cup b^*$ (iii) $a^*b^* \cap c^*d^* = \emptyset$ (iv) $abcd \in (a(cd)^*b)^*$</p> |

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| Q36. | Determine whether the strings in the table belong to any of the languages described by the following regular expressions |
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(i) For 1001 : T, F, F, F, F

(ii) For 110 : F, F, F, F, T

(a) i only

(b) ii only

(c) Both i & ii

(d) Neither i nor ii

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| Q37. | Which of the following is not in the set of strings denoted by the regular expression $R = (a^* b c^*)^*$? |
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(A) aabc

(B) bacd

(C) abcbc

(D) babbc

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| Q38. | Which of the following strings are generated by the regular expression |
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$$(ab)^*.\lambda.(a+b+\emptyset).ba?$$

(i) λ (ii) aba

(iii) ababba

(iv) abababa

(A) ii, iii and iv only

(B) ii and iv only

(C) i and ii only

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| Q39. | Which of the following strings is a member of the language described by the regular expression: $(a^*ba^*ba^*ba^*)^*$ |
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(A) bbbb

(B) bbaaabb

(C) bbaaabbabb

(D) bbabbbab

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| Q40. | <p>Someone has asserted that the following two regular expressions describe the same language: $R1 = ((ab^*a) + (ba^*b))^*$ and $R2 = ((ab^*a) + b^*)^*$. Which of the following strings is contained in one of the languages but not in the other?</p> <p>(A) ababab (B) bbbbbb (C) abba (D) bbabba</p> |
| Q41. | <p>The regular expression $b^*ab^*ab^*ab^*$ represents the language</p> <p>(A) $L = \{w : w \in \Sigma^*, n_a(w) = 3\}$ (B) $L = \{w : w \in \Sigma^*, n_a(w) \leq 3\}$ (C) $L = \{w : w \in \Sigma^*, n_a(w) \geq 3\}$ (D) none</p> |
| Q42. | <p>The regular expression $b^* + b^*ab^* + b^*ab^*ab^*$ represents the language</p> <p>(A) $L = \{w : w \in \Sigma^*, n_a(w) = 2\}$ (B) $L = \{w : w \in \Sigma^*, n_a(w) \leq 2\}$ (C) $L = \{w : w \in \Sigma^*, n_a(w) \geq 2\}$ (D) none</p> |
| Q43. | <p>Which of the following pair of regular expression is/are true?</p> <p>I. $(01 + 0)^*0 \Leftrightarrow 0(10 + 0)^*$ II. $0(120)^*12 \Leftrightarrow 01(201)^*2$ III. $\phi^* \Leftrightarrow \lambda^*$ IV. $(0^*1^*)^* \Leftrightarrow (1^*0^*)^*$ (A) I, III and IV only (B) II and III only (C) II, III and IV only (D) All the above</p> |
| Q44. | <p>Regular expression $r = (aa^*b)^* (aa^* + \lambda)$ is equivalent to:</p> <p>(I) $(ab)^* (a + \lambda)$ (II) $a^*(ab)^*$ (III) $(b + ba)^*$ (IV) $(a + ab)^*$ (V) $(aa^*b)^*$ (A) IV and V (B) I, II and III (C) II, IV and V (D) IV only</p> |

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| <p>Q45.</p> | <p>Identify the pairs of regular expressions that are equivalent (in that they describe the same sets of strings):</p> <p>(I) $(ab)^+ (ab)^*ab$</p> <p>(II) $ab^* (ab)^*$</p> <p>(III) $(a \mid b)^+ (a \mid (b)^+)$</p> <p>(IV) $a^{++} a^+$</p> <p>(V) $a^*a (ba^*a)^* (a + ab)^*a$.</p> <p>(A) I, II and III only</p> <p>(B) I and III only</p> <p>(C) I, III and V only</p> <p>(D) All are equivalent</p> |
| <p>Q46.</p> | <p>Match the regular expression with its description</p> <p>(1) $(0 \cup 1)^*01(0 \cup 1)^*$ i. All strings which doesn't contain the substring 101.</p> <p>(2) 1^*0^* ii. Strings containing the substring 01.</p> <p>(3) $(10 \cup 0)^*(1 \cup 10)^*$ iii. Strings of the form $111 \dots 000 \dots$, that is, strings that begins with Zero or more ones followed by zero or more zeroes</p> <p>(4) $0^*(1 \cup 000^*)^*0^*$ iv. All strings where each occurrence of 00 precedes all Occurrences of 11</p> <p>(A) 1-ii,2-iii,3-i,4-iv</p> <p>(B) 1-ii,2-iii,3-iv,4-I</p> <p>(C) 1-iv,2-iii,3-ii,4-i</p> <p>(D) 1-iii,2-ii,3-i,4-iv</p> |
| <p>Q47.</p> | <p>How many of the following is/are true? _____</p> <p>(i) $(ab)^*a = a(ba)^*$</p> <p>(ii) $(a \cup b)^* b (a \cup b)^* = a^* b (a \cup b)^*$</p> <p>(iii) $[(a \cup b)^* b (a \cup b)^* \cup (a \cup b)^* a (a \cup b)^*] = (a \cup b)^*$</p> <p>(iv) $[(a \cup b)^* b (a \cup b)^* \cup (a \cup b)^* a (a \cup b)^*] = (a \cup b)^+$</p> <p>(v) $[(a \cup b)^* b a (a \cup b)^* \cup a^*b^*] = (a \cup b)^*$</p> |

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| Q48. | <p>[MSQ]</p> <p>Which of the following regular expressions are equivalent to the regular expression $R = (bba + aab + ab + b + a)^+ + \lambda$</p> <p>(A) $(a^*b^*)^*$ (B) $(a^*b^*)^+ + \lambda$ (C) $(a + b + aa)^*$ (D) $(a + b)^*$</p> |
| Q49. | <p>[MSQ]</p> <p>Which of the following regular expression is equivalent to given regular expression: $\epsilon + 1^*(011)^*(1^*(011)^*)^*$</p> <p>(A) $(1+011)^*$ (B) $(1^*(011)^*)^*$ (C) $(1+(011)^*)^*$ (D) $(1011)^*$</p> |
| Q50. | <p>If r and s are regular expressions, write $r \leq s$ to mean that the language of strings matching r is contained in the language of strings matching s. Then which of the following is/ are true?</p> <p>(i) $r^* s^* \leq (r s)^*$ (ii) $(r s)^* \leq r^* s^*$ (iii) $(r^*s^*)^* \leq (r s)^*$ (iv) $(r s)^* \leq (r^*s^*)^*$ (v) $(rs r)^*r \leq r(sr r)^*$</p> <p>(A) i, ii, iii only (B) i, iii, iv only (C) ii, iii, iv only (D) All except ii</p> |
| Q51. | <p>Which of the following is true?</p> <p>(i) $(01)^*0 = 0(10)^*$ (ii) $(0+1)^*0(0+1)^*1(0+1) = (0+1)^*01(0+1)^*$ (iii) $(0+1)^*01(0+1)^*+1^*0^* = (0+1)^*$ (iv) $(0(0+1)^*1 + 1(0 + 01)^*0)^* = (01 + 10)^*$</p> <p>(A) i only (B) i and iii only (C) i, ii and iii only (D) All the above</p> |

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| <p>Q52.</p> | <p>Consider the following four regular expressions over the alphabet $\{a, b\}$:</p> <p>$E1 = (ab + a^*b^*b^*)^*$</p> <p>$E2 = ((ab)^* (a^*b^*b^*)^*)^*$</p> <p>$E3 = (a + b)^*$</p> <p>$E4 = a(a + b)^*$</p> <p>Which of the following statements is true?</p> <p>(A) $L(E2) = L(E3)$</p> <p>(B) $L(E3) = L(E4)$</p> <p>(C) $L(E1) = L(E4)$</p> <p>(D) $L(E2) = L(E4)$</p> |
| <p>Q53.</p> | <p>What will be the regular expression for language $L = \{xwx : x, w \in \{0, 1\}^*, x \leq 3\}$?</p> <p>(A) $(0+1)^3 (0 + 1)^*(0+1)^3$</p> <p>(B) $((0+1)+(0+1)^2+(0+1)^3) (0 + 1)^* ((0+1)+(0+1)^2+(0+1)^3)$</p> <p>(C) $(\lambda + (0+1)+(0+1)^2+(0+1)^3) (0 + 1)^*(\lambda + (0+1)+(0+1)^2+(0+1)^3)$</p> <p>(D) None of these</p> |
| <p>Q54.</p> | <p>What will be the regular expression for the following language</p> <p>$L1 = \{pwp : p, w \in \{0, 1\}^*, p = 7, k \in I^+\}$?</p> <p>(A) $(0+1)^7 (0 + 1)^*(0+1)^7$</p> <p>(B) $((0+1)+(0+1)^2+(0+1)^3+\dots+(0+1)^7) (0 + 1)^* ((0+1)+(0+1)^2+(0+1)^3+\dots+(0+1)^7)$</p> <p>(C) $(\lambda + (0+1)+(0+1)^2+(0+1)^3+\dots+(0+1)^7) (0 + 1)^*(\lambda + (0+1)+(0+1)^2+(0+1)^3+\dots+(0+1)^7)$</p> <p>(D) Both a and b</p> |
| <p>Q55.</p> | <p>The regular expression $(a + b)^*a(a + b)^*$ represents the language</p> <p>(A) Contains exactly 1 a</p> <p>(B) Contains at least 1 a</p> <p>(C) Contains at most 1 a</p> <p>(D) none</p> |
| <p>Q56.</p> | <p>The regular expression $(b^*ab^*ab^*)^*$ represents the language</p> <p>(A) $L = \{w : w \in \Sigma^*, n_a(w) \text{ is divisible of } 2 \}$</p> <p>(B) every b is followed by at least one a</p> <p>(C) $L = \{w : w \in \Sigma^* n_a(w) \geq 2 \}$</p> <p>(D) None of these.</p> |

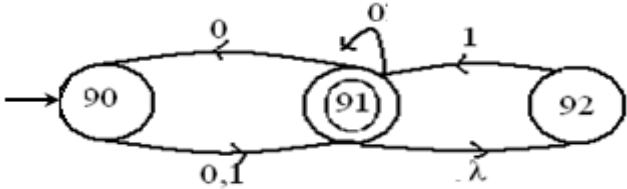
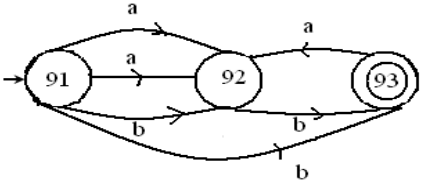
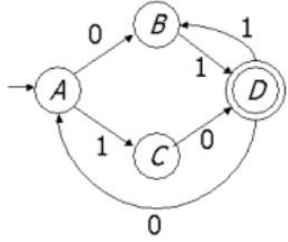
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| Q57. | <p>Which of the following regular expressions generate all the strings with even Numbers of 0?</p> <p>(i) $1^*(00)^*1^*$</p> <p>(ii) $1^*(010)^*1^*$</p> <p>(iii) $(1^*01^*01^*)^*$</p> <p>(iv) $1^*(01^*0)^*1^*$</p> <p>(A) iii and iv only</p> <p>(B) ii, iii and iv only</p> <p>(C) iv only</p> <p>(D) iii only</p> |
| Q58. | <p>Which of the following expression best describes the language $L = \{\alpha \in \{0, 1\}^* \mid \alpha \text{ contains 1's only}\}$</p> <p>(A) 1^*</p> <p>(B) 1^+</p> <p>(C) $0 + 1^*$</p> <p>(D) $1.(1)^+$</p> |
| Q59. | <p>Which of the following expression best describes the language $L = \{\alpha \in \{0, 1\}^* \mid \alpha \text{ contains only 0's or only 1's}\}$</p> <p>(A) $0^* + 1^*$</p> <p>(B) $(00^+)^* + (11^+)^*$</p> <p>(C) $0^+ + 1^+$</p> <p>(D) $0(0 + 1)^*1$</p> |
| Q60. | <p>Let $\Sigma = \{0, 1\}$, and language over Σ, $L = \{\alpha \in \Sigma^* \mid \alpha \text{ contains odd number of 1's}\}$. Which Of the following regular expression best describes the given language?</p> <p>(A) $0^*(10^*10^*)^*10^*$</p> <p>(B) $0^*1(11)^*0^*$</p> <p>(C) $0^*10^*10^*1$</p> <p>(D) $1(11)^*$</p> |
| Q61. | <p>Let $\Sigma = \{0, 1\}$, and language over Σ, $L = \{\alpha \in \Sigma^* \mid \text{any two 0's in } \alpha \text{ are separated by three 1's}\}$. Which of the following regular expression best describes the given language?</p> <p>(A) $(01110)^*$</p> <p>(B) $(01110)^* + 1^*$</p> <p>(C) $1^*(01110)^* + 1^*$</p> <p>(D) $1^*(0111)^*01^* + 1^*$</p> |
| Q62. | <p>Let $\Sigma = \{0, 1\}$, and language over Σ, $L = \{\alpha \in \Sigma^* \mid \alpha \text{ is a binary number divisible by 4}\}$. Which of the following regular expression best describes the given language?</p> <p>(A) $(0+1)^*00$</p> <p>(B) $(0+1)^*0$</p> <p>(C) $(0+1)^*000$</p> <p>(D) $(0+1)^*100$</p> |

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| Q63. | <p>Let $\Sigma = \{0, 1\}$, and language over Σ, $L = \{\alpha \in \Sigma^* \mid \alpha \text{ does not contain } 11\}$. Which of the following regular expression best describes the given language?</p> <p>(A) $0^*(10)^*$ (B) $0^*(10^+)^*(1+\epsilon)$ (C) $(0^*10^*10^*)^*$ (D) $(0+10)^*$</p> |
| Q64. | <p>Which of the following regular expression generate all the strings with odd numbers of 1?</p> <p>(A) $0^*10^*(10^*1)^*$ (B) $0^*10^*(10^*1)^*0^*$ (C) $0^*10^*(101)^*0^*$ (D) $0^*(10^*10^*1)^*0^*$</p> |
| Q65. | <p>Which of the following regular expressions define the same language?</p> <p>(1) $(ab)^*$ (2) $(aa^*b)^* a^*$ (3) a^*b^* (4) $a^* (aba^*)^*$ (A) 1 and 3 only (B) 2 and 4 only (C) 3 and 4 only (D) All define different languages</p> |
| Q66. | <p>Consider the language defined by the regular expression $(a \mid b)^*b^+$. Which of the following regular expression(s) also define that language?</p> <p>(1) $(a^*b^+) \mid (b^*b^+)$ (2) $(ab \mid bb)^* b^*$ (3) $(a \mid b \mid ba)^*b^+$ (A) (1) and (2) (B) (2) and (3) (C) Only (3) (D) Only (2)</p> |
| Q67. | <p>Which of the following defines a language different than the others?</p> <p>(A) The regular expression $(a \mid b)^* a b c$ (B) The regular expression $(a^* b^*)^* a b c$ (C) The regular expression $(a \mid b) (a \mid b)^* c$ (D) All are same</p> |

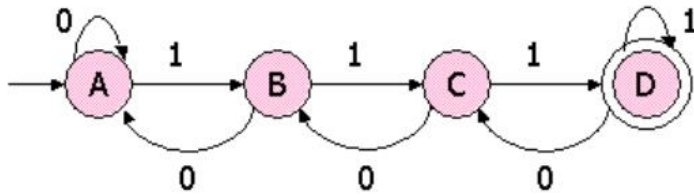
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| <p>Q68.</p> | <p>Which of the following statement is / are correct?</p> <p>I. $(a^*b^*)abc \equiv (a b)^*abc$</p> <p>II. $(a^*b^*)^*abc \equiv (aa+ba+a^*b+(bb)^*a)^*abc$</p> <p>(A) I only</p> <p>(B) II only</p> <p>(C) Both I & II</p> <p>(D) Neither I nor II</p> |
| <p>Q69.</p> | <p>What is the best description of languages denoted by the following regular expression $R = 0(0+1)^*0$?</p> <p>(A) strings of zeros and ones with zeros occurring more frequently than ones</p> <p>(B) strings from alphabet $\{0,1\}$ which begin and end with a zero</p> <p>(C) strings from alphabet $\{0,1\}$ which begin and end with a zero and have an even number of zeroes</p> <p>(D) strings from alphabet $\{0,1\}$ which begin with one or more zeros, followed by zero or more 3ones, followed by a zero.</p> |
| <p>Q70.</p> | <p>What is the best description of the languages denoted by the following regular expressions $R = ((11 + 0)^*)^*$</p> <p>(A) strings from the alphabet $\{0,1\}$ in which there are an even number of 1s</p> <p>(B) strings from the alphabet $\{0,1\}$ in which ones always appear in pairs</p> <p>(C) strings from the alphabet $\{0,1\}$ in which ones occur twice as frequently as zeros</p> <p>(D) strings from the alphabet $\{0,1\}$ in which there are an even number of 1s and an odd number of zeroes.</p> |
| <p>Q71.</p> | <p>Which of the following languages denoted by the regular expressions $R = 0^*10^*10^*10^*$</p> <p>(A) strings from the alphabet $\{0, 1\}$ in which there are an odd number of 1s</p> <p>(B) strings from the alphabet $\{0, 1\}$ in which ones never appear together</p> <p>(C) strings from the alphabet $\{0, 1\}$ in which there are exactly three ones.</p> <p>(D) strings from the alphabet $\{0, 1\}$ in which there are exactly three ones or no ones</p> |

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| Q72. | <p>How many of the regular expression(s) which matches all strings of 0's and 1's that do not contain the substring 011? _____</p> <p>(1) $((011)0^*1^*)^*$ (2) $(0^*1)^+0^* + 1^*0^*$</p> <p>(3) $(01 \mid 010 \mid 100 \mid 110)^*$ (4) $(0+1)^*0^* \mid 1^*0^*$</p> |
| Q73. | How many strings of length less than 4 contains the language described by the regular expression $(a + d)^*b(a + bc)^*$? _____ |
| Q74. | How many strings are there in the language defined by regular expression? _____ $((\emptyset^* \cap a) \cup (\emptyset \cup b^*)) \cap \emptyset^*$ |
| Q75. | How many strings are there in the language defined by regular expression? _____ $((\emptyset^* \cup b) \cap (b^* \cup \emptyset))$ |
| Q76. | How many strings of length less than 5 contains the language described by the regular expression $0^*1(0 + 1)^*01^*$? _____ |
| Q77. | Let L be the language generated by regular expression $((a + b)^*b(a + ab)^*)$. How many strings of length less than four are there in L? _____ |
| Q78. | <p>What is the regular expression for the language generated by the following grammar?</p> <p>$S \rightarrow Aab \quad A \rightarrow Aab \mid B \quad B \rightarrow a$</p> <p>(A) $aab(ba)^*$</p> <p>(B) $aab(ab)^*$</p> <p>(C) $aa(ab)^*$</p> <p>(D) $ab(ab)^*$</p> |
| Q79. | <p>Consider the following grammar G:</p> <p>$S \rightarrow AabB, \quad A \rightarrow b \mid bA \mid \lambda, \quad B \rightarrow aB \mid aa \mid \lambda$</p> <p>Which of the following regular expression is equivalent to the language generated by G?</p> <p>(A) b^*abaa^*</p> <p>(B) a^*abb^*</p> <p>(C) b^*aba^*</p> <p>(D) b^+aba^+</p> |

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| <p>Q80.</p> | <p>Consider the following grammar G:</p> $S \rightarrow AB \mid aAb \quad A \rightarrow b \mid baA \mid \lambda \quad B \rightarrow aB \mid abB \mid \lambda$ <p>Which of the following regular expression is equivalent to the language generated by G?</p> <p>(A) $(ba)^*(b + \lambda)(a + b)^*$</p> <p>(B) $(ba)^*(b + \lambda)(a + ab)^*$</p> <p>(C) $(ba)^*(b + \lambda)(a + ab)^* + a(b + ab)^*b$</p> <p>(D) $a((ba)^*(b + \lambda))b + (ba)^*(b + \lambda)(a + ab)^*$</p> |
| <p>Q81.</p> | <p>Which of the following grammar generates the language generated by given regular expression $R = a(a+b)^*a + b(a+b)^*b$</p> <p>(A) $S \rightarrow aSa \mid bSb \mid aS \mid bS \mid \lambda$</p> <p>(B) $S \rightarrow aAa \mid bAb \quad A \rightarrow aA \mid bA \mid \lambda$</p> <p>(C) $S \rightarrow aAa \mid bAb \quad A \rightarrow aA \mid bA \mid aAa \mid bAb \mid \lambda$</p> <p>(D) Both b & c</p> |
| <p>Q82.</p> | <p>Which of the following grammar generates the language generated by given regular expression $R = aaa^+ + (ba+bb)^*a$</p> <p>(A) $S \rightarrow aaX \mid Ya \quad X \rightarrow aX \mid \lambda \quad Y \rightarrow baYbbY \mid \lambda$</p> <p>(B) $S \rightarrow aX \mid Ya \quad X \rightarrow aX \mid a \quad Y \rightarrow baY \mid bbY \mid \lambda$</p> <p>(C) $S \rightarrow aaX \mid Ya \quad X \rightarrow aX \mid a \quad Y \rightarrow baY \mid bbY \mid \lambda$</p> <p>(D) None of the above</p> |
| <p>Q83.</p> | <p>Which of following Grammar generates the language $L(a^*a(a+ba)^*)$?</p> <p>(A) $S \rightarrow S_1 \mid S_2, \quad S_1 \rightarrow aA, \quad A \rightarrow aA \mid \lambda, \quad S_2 \rightarrow aS_2 \mid baS_2 \mid \lambda$</p> <p>(B) $S \rightarrow S_1S_2 \mid S_1, \quad S_1 \rightarrow aA, \quad A \rightarrow aA \mid \lambda, \quad S_2 \rightarrow aS_2 \mid baS_2 \mid \lambda$</p> <p>(C) $S \rightarrow S_1S_2 \mid a, \quad S_1 \rightarrow aA, \quad A \rightarrow aA \mid a, \quad S_2 \rightarrow aS_2 \mid baS_2 \mid \lambda$</p> <p>(D) Both B and C</p> |

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| <p>Q84.</p> | <p>Which of the following grammars is a right regular grammar with the same language as the regular expression $(a^* \cup b^*) \cap (a^* b^*)$?</p> <p>(A) $S \rightarrow AB$ $A \rightarrow aA \mid \varepsilon$ $B \rightarrow bB \mid \varepsilon$</p> <p>(B) $S \rightarrow A \mid B$ $A \rightarrow Aa \mid \varepsilon$ $B \rightarrow Bb \mid \varepsilon$</p> <p>(C) $S \rightarrow A \mid B$ $A \rightarrow aA \mid a$ $B \rightarrow Bb \mid a$</p> <p>(D) $S \rightarrow A \mid B$ $A \rightarrow aA \mid \varepsilon$ $B \rightarrow bB \mid a$</p> |
| <p>Q85.</p> | <p>Consider the following NFA</p>  <p>Which of the following string is not accepted by the following NFA?</p> <p>(A) 00000000 (B) 01001</p> <p>(C) 111110010 (D) None</p> |
| <p>Q86.</p> | <p>Which of the following string is accepted by given DFA?</p>  <p>(A) aaabbbbbbb (B) aababb</p> <p>(C) abab (D) babaa</p> |
| <p>Q87.</p> | <p>Consider the following DFA:</p>  <p>If the input is 011100101, which edge of the automaton is NOT traversed?</p> <p>(A) CD</p> <p>(B) AC</p> <p>(C) BC</p> <p>(D) BD</p> |

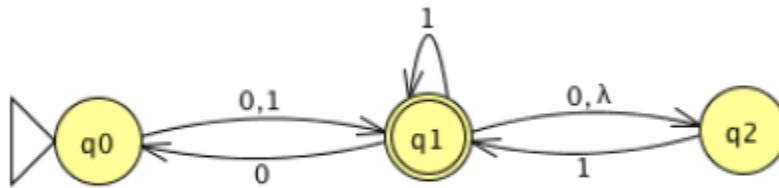
Q88. Consider the following DFA:



If string s is accepted by this DFA, which of the following strings cannot be a suffix of s ?

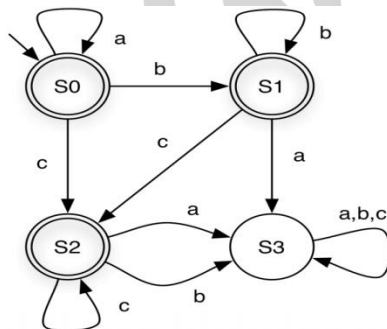
- (A) 111
- (B) 111001
- (C) 11011
- (D) 0010111

Q89. How many of the following strings are accepted by the NFA given below? _____



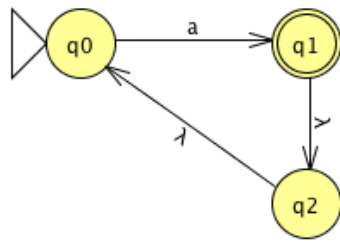
- (i) 00
- (ii) 01001
- (iii) 10010
- (iv) 000
- (v) 0000

Q90. Which of the following strings is accepted by DFA given below?



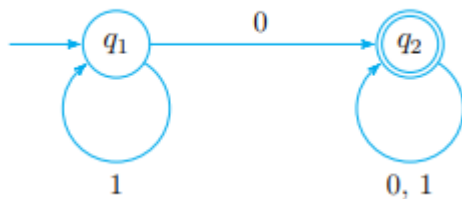
- (A) aabbbcca
- (B) aabbbcccca
- (C) abbbccccc
- (D) bccccabc

Q91. What is the complement of the language accepted by the following NFA? ($\Sigma = \{a\}$)



- (A) \emptyset
- (B) $\{\lambda\}$
- (C) a^*
- (D) a^+

Data for next two questions: Consider the following DFA D: (New Question)



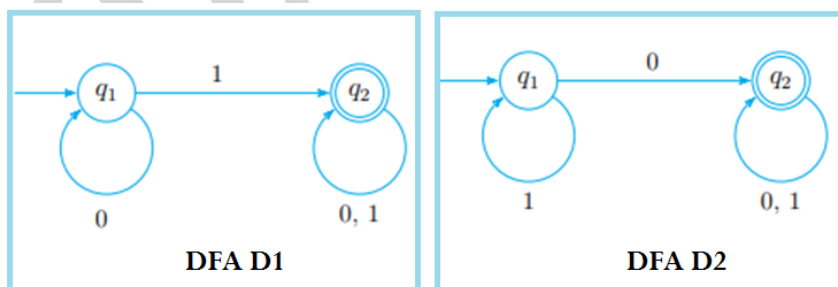
Q92. Let the language accepted by G is $L(D)$ then complement of $L(D)$ is equivalent to

- (A) 1^+
- (B) 1^*
- (C) $1^* + 0^*$
- (D) $1(0+1)^*$

Q93. Let the language accepted by G is $L(D)$ then reverse of $L(D)$ is equivalent to

- (A) 01^*
- (B) $1^*0(0+1)^*$
- (C) $(0+1)^*01^*$
- (D) $(0+1)^* + 01^*$

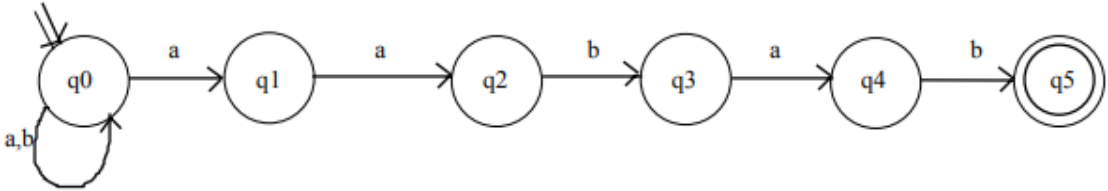
Data for next twenty questions: Consider the following two DFA's D1 and D2:



It is given the language accepted by D1 and D2 are $L(D1)$ and $L(D2)$, respectively.

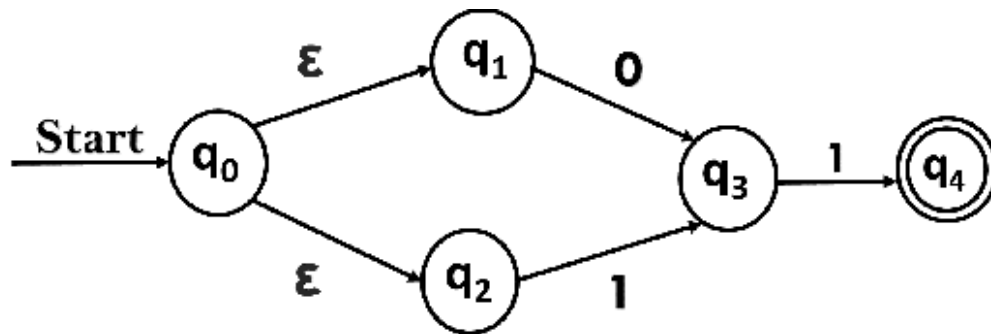
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| Q94. | <p>[MSQ]</p> <p>If $L(D) = L(D1) \cup L(D2)$, then which of the following regular expression is/are equivalent to $L(D)$?</p> <p>(A) $(0 + 1)^*$ (B) $(01 + 10)(0 + 1)^*$</p> <p>(C) $(0^*1 + 1^*0)(0 + 1)^*$ (D) $(0 + 1)^+$</p> |
| Q95. | <p>If $L(D) = L(D1) \cap L(D2)$, then the regular expression for $L(D)$ is equivalent to</p> <p>(A) \emptyset (B) $(01 + 10)(0 + 1)^*$</p> <p>(C) $(00^*1 + 11^*0)(0 + 1)^*$ (D) None of these</p> |
| Q96. | <p>If $L(D) = L(D1) - L(D2)$, then the regular expression for $L(D)$ is equivalent to</p> <p>(A) 1^* (B) 1^+</p> <p>(C) $1(0 + 1)^*$ (D) None of these</p> |
| Q97. | <p>If $L(D) = L(D1).L(D2)$, then the regular expression for $L(D)$ is equivalent to</p> <p>(A) $0^*11^*0(0 + 1)^*$ (B) $0^*1(0+1)^*1^*0(0+1)^*$</p> <p>(C) $(0 + 1)^*$ (D) $(0 + 1)^+$</p> |
| Q98. | <p>If $L(D) = L(D1)^R.L(D2)^R$, then the regular expression for reverse of $L(D)$ is</p> <p>(A) $(0 + 1)^*$ (B) $(0 + 1)^*1^*0(0 + 1)^*10^*$</p> <p>(C) $0^*1(0 + 1)^*1^*0(0 + 1)^*$ (D) $(0 + 1)^*01^*(0 + 1)^*10^*$</p> |
| Q99. | <p>If $L(D) = L(D1) \cup L(D2)$, then the regular expression for complement of $L(D)$ is</p> <p>(A) λ (B) $0^+ + 1^+$</p> <p>(C) $0^* + 1^*$ (D) 0^*1^*</p> |
| Q100. | <p>If $L(D) = L(D1) \cap L(D2)$, then the regular expression for complement of $L(D)$ is</p> <p>(A) λ (B) $0^+ + 1^+$</p> <p>(C) $0^* + 1^*$ (D) $(0+1 + 1+0)(0+1)^*$</p> |
| Q101. | <p>[MSQ]</p> <p>If $L(D) = L(D1)^* \cup L(D2)^*$, then which of the following regular expression is/are equivalent to $L(D)$?</p> <p>(A) $(01 + 10)^*$ (B) $(0 + 1)^*$</p> <p>(C) $0^*1^*(0+1)^*$ (D) none of these</p> |

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| Q102. | Minimum how many states are required to construct the DFA for $L(D1) \cup L(D2)$? ____ |
| Q103. | Minimum how many states are required to construct the DFA for $L(D1) \cap L(D2)$? ____ |
| Q104. | Minimum how many states are required to construct the DFA for $\overline{L(D1)} \cap \overline{L(D2)}$? ____ |
| Q105. | Minimum how many states are required to construct the DFA for $\overline{L(D1)} \cup \overline{L(D2)}$? ____ |
| Q106. | Minimum how many states are required to construct the DFA for $\overline{L(D1)} \cdot \overline{L(D2)}$? ____ |
| Q107. | Minimum how many states are required to construct the DFA for $L(D1) \cdot L(D2)$? ____ |
| Q108. | <p>[MSQ]</p> <p>Which of the following Grammars accept the language $L(D1) \cup L(D2)$?</p> <p>(A) $S \rightarrow 0S \mid 1S \mid 0 \mid 1$</p> <p>(B) $S \rightarrow 0A \mid 1B$, $A \rightarrow 0A \mid \lambda$, $B \rightarrow 1B \mid \lambda$</p> <p>(C) $S \rightarrow 0S1 \mid 1S0 \mid 0S \mid 1S \mid 0 \mid 1$</p> <p>(D) None of the above</p> |
| Q109. | <p>[MSQ]</p> <p>Which of the following Grammars accept the language $L(D1) \cap L(D2)$?</p> <p>(A) $S \rightarrow 0X1A \mid 1Y0A$, $X \rightarrow 0X \mid \lambda$, $Y \rightarrow 1Y \mid \lambda$, $A \rightarrow 0A \mid 1A \mid \lambda$</p> <p>(B) $S \rightarrow XY$, $X \rightarrow 0A1 \mid 1B0$, $A \rightarrow 0A \mid \lambda$, $B \rightarrow 1B \mid \lambda$, $Y \rightarrow 0Y \mid 1Y \mid \lambda$</p> <p>(C) $S \rightarrow 0S \mid 1S \mid 0 \mid 1$</p> <p>(D) None of the above</p> |
| Q110. | <p>[MSQ]</p> <p>Which of the following Grammars accept the language $L(D1)^R \cup L(D2)^R$?</p> <p>(A) $S \rightarrow 0S \mid 1S \mid 0 \mid 1$</p> <p>(B) $S \rightarrow 0A \mid 1B$, $A \rightarrow 0A \mid \lambda$, $B \rightarrow 1B \mid \lambda$</p> <p>(C) $S \rightarrow 0S1 \mid 1S0 \mid 0S \mid 1S \mid 0 \mid 1$</p> <p>(D) None of the above</p> |

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| Q111. | <p>[MSQ]</p> <p>Which of the following Grammars accept the language $L(D1) \cdot L(D2)$?</p> <p>(A) $S \rightarrow AB, A \rightarrow X1Z, \quad X \rightarrow 0X \mid \lambda, B \rightarrow Y0Z, \quad Y \rightarrow 1Y \mid \lambda, \quad Z \rightarrow 0Z \mid 1Z \mid \lambda$</p> <p>(B) $S \rightarrow 0S \mid 1S \mid 0 \mid 1$</p> <p>(C) $S \rightarrow 0S1 \mid 1S0 \mid 0S \mid 1S \mid 0 \mid 1$</p> <p>(D) None of the above</p> |
| Q112. | <p>[MSQ]</p> <p>Which of the following Grammar(s) accept the language $\overline{L(D1)} \cdot \overline{L(D2)}$?</p> <p>(A) $S \rightarrow 0S \mid 1S$</p> <p>(B) $S \rightarrow \lambda$</p> <p>(C) $S \rightarrow 0S \mid 1S \mid 0 \mid 1$</p> <p>(D) $S \rightarrow XY \quad X \rightarrow 0X \mid \lambda \quad Y \rightarrow 1S \mid \lambda$</p> |
| Q113. | <p>Which of the following Grammar accepts the language $\overline{L(D1)} \cap \overline{L(D2)}$?</p> <p>(A) $S \rightarrow 0S \mid 1S$</p> <p>(B) $S \rightarrow \lambda$</p> <p>(C) $S \rightarrow 1S \mid 0S \mid 0 \mid 1$</p> <p>(D) None of the above</p> |
| Q114. | <p>Consider the following NFA M: :</p>  <p>If the language accepted by M is $L(M)$ then the regular expression for reverse of $L(M)$ is</p> <p>(A) $(a + b)^* \{ \lambda + a + aa + aab + aaba + aabbab \}$</p> <p>(B) $(aabab)(a + b)^*$</p> <p>(C) $babaa (a + b)^*$</p> <p>(D) None of the above</p> |

Q115. [MSQ]

Consider the following NFA with ϵ (epsilon):

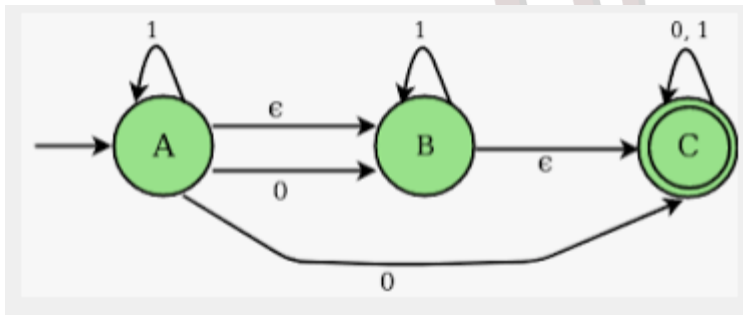


Which of the following statement is/are true?

- (A) ϵ -closure $\{q_0\} = \{q_0, q_1, q_2\}$
- (B) ϵ -closure $\{q_1\} = \{q_1, q_2\}$
- (C) ϵ -closure $\{q_2\} = \{q_0, q_2\}$
- (D) ϵ -closure $\{q_3\} = \{q_3\}$

Q116. [MSQ]

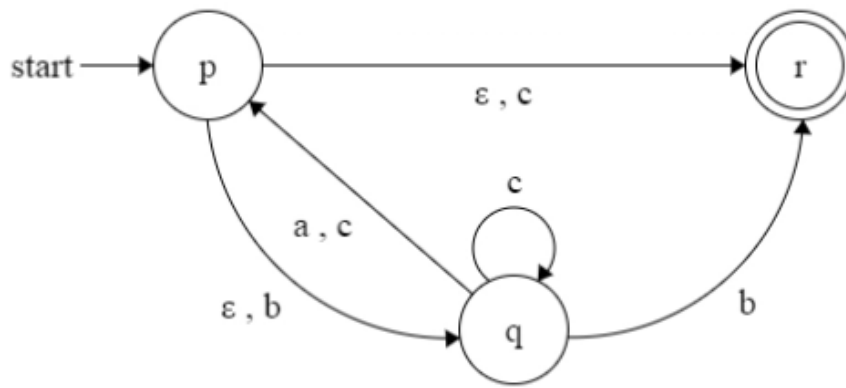
Consider the following NFA with ϵ (epsilon):



Which of the following statement is/are true?

- (A) ϵ -closure $\{A\} = \{A, B, C\}$
- (B) ϵ -closure $\{B\} = \{B, C\}$
- (C) ϵ -closure $\{C\} = \{C\}$
- (D) ϵ -closure $\{A\} = \{A, B\}$

Q117. Consider the following NFA with ϵ (epsilon): :

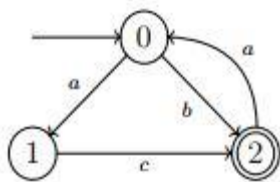


Which of the following option(s) is/are true about ϵ -closure of different states of given NFA?

| | |
|--|---|
| $\epsilon\text{-closure}(p) = \{p, q, r\}$ $\epsilon\text{-closure}(q) = \{p, q, r\}$ $\epsilon\text{-closure}(r) = \{p, q, r\}$ | $\epsilon\text{-closure}(p) = \{p, q, r\}$ $\epsilon\text{-closure}(q) = \{q, r\}$ $\epsilon\text{-closure}(r) = \{r\}$ |
| $\epsilon\text{-closure}(p) = \{p, q, r\}$ $\epsilon\text{-closure}(q) = \{q\}$ $\epsilon\text{-closure}(r) = \{r\}$ | None of these |

Q118. Consider the following NFA with ϵ (epsilon):

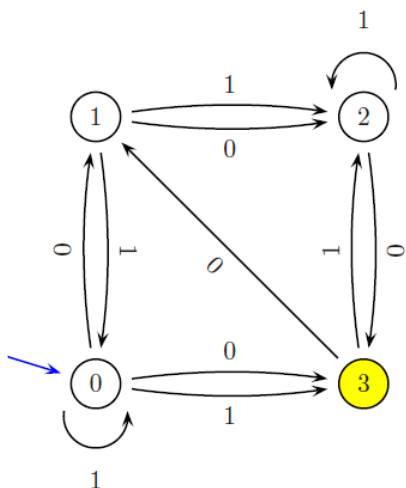
The regular expression of the following non-deterministic finite automaton is



- (A) $ac + b(ab)^*$
- (B) $(ac + b)(aac + ab)^*$
- (C) $(aca + ba)(b + ac)^*$
- (D) Both b and c

Q119. [MSQ]

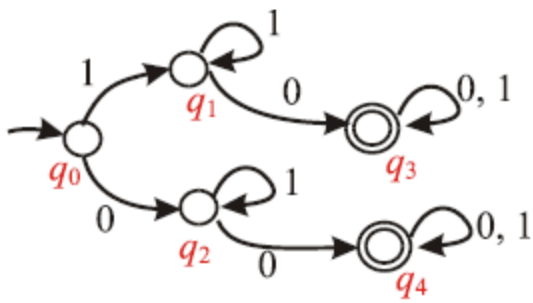
Consider the following FSA: (Where state '3' is final state) :



Which of the following statements is/are true?

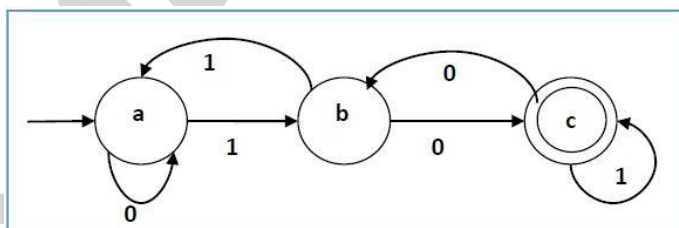
- (a) The FSA accepts 011111101.
- (b) The FSA accepts 11101000.
- (c) The FSA rejects 0000.
- (d) The FSA accepts all bit strings with an odd number of 0.

Q120. What is the language accepted by the following DFA?



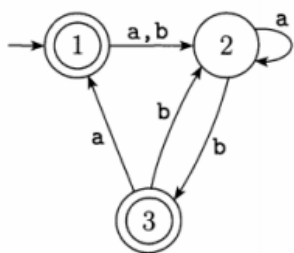
- (A) 0^*1^*
- (B) 1^*0^*
- (C) $(1+0)^*1^*(0+1)^*$
- (D) $(1+0)1^*0(0+1)^*$

Q121. The regular expression for following automata will be:



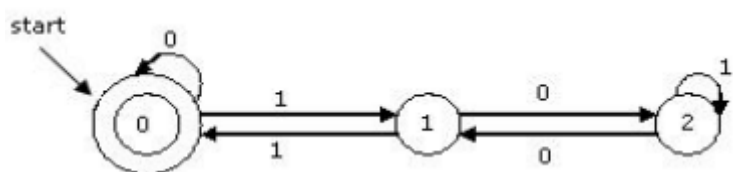
- (A) $(011)^*10(00+1)^*$
- (B) $0^*11^*10(00+1)^*$
- (C) $(0+11)^*10(00)^*1^*$
- (D) $(0+11)^*10(00+1)^*$

Q122. What is the regular expression corresponding to the language accepted by the following finite state automata?



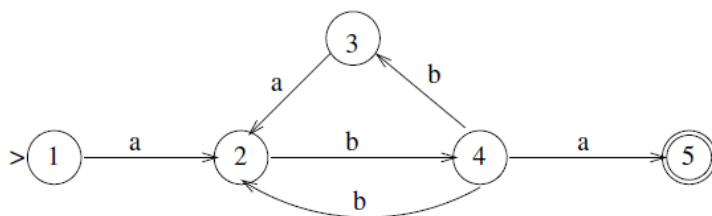
- (A) $((a+b)(a^*(bb)^*)^*b(\lambda+a))^*$
- (B) $(a(a^*(bb)^*)^*b+b(a^*(bb)^*)^*a)$
- (C) $(a+b)(a^*(bb)^*)^*a$
- (D) None

Q123. How many of the following regular expressions that denotes a subset of the language recognized by the given DFA? _____



- (1) $0^*(11)^*0^*$
- (2) $0^*1(10^*1)^*1$
- (3) $0^*1(10^*1)^*10^*$
- (4) $0^*1(10^*1)0(100)^*$
- (5) $(0^*1(10^*1)^*10^* + 0^*)^*$

Q124. What is the language accepted by the following finite state automata?



- (i) $a(bb + bba)^*ba$
- (ii) $ab(bb + bab)^*a$
- (iii) $ab(b+ba)^*a$

(A) i only

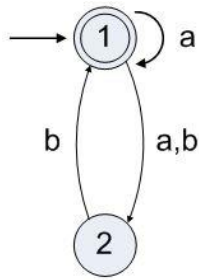
(B) ii and iii only

(C) All of these

(D) i and iii only

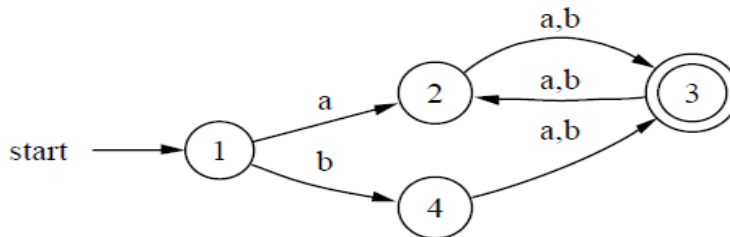
Q125. [MSQ]

What will be the regular expression for following automata?



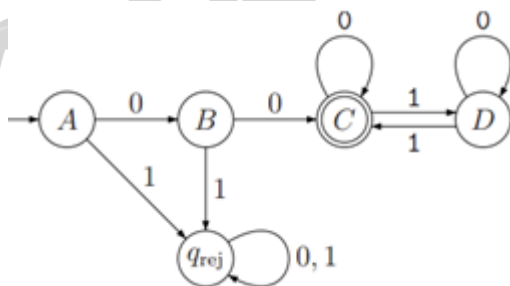
- (A) $a^*(a+b)b$
- (B) $a^*(a+b)ba^*$
- (C) $a^*\{((a+b)b)^*a^* + \lambda\}$
- (D) $(a + (a+b)b)^*$

Q126. What is the language accepted by following automata?



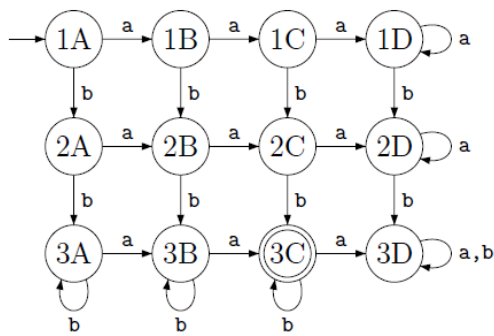
- (A) Any string starts or ends with b
- (B) The strings with an even number of characters
- (C) The strings with an even number of characters and length of at least 2
- (D) The strings with an even number of a's or b's and length of at least 2

Q127. The language accepted by the following DFA is:



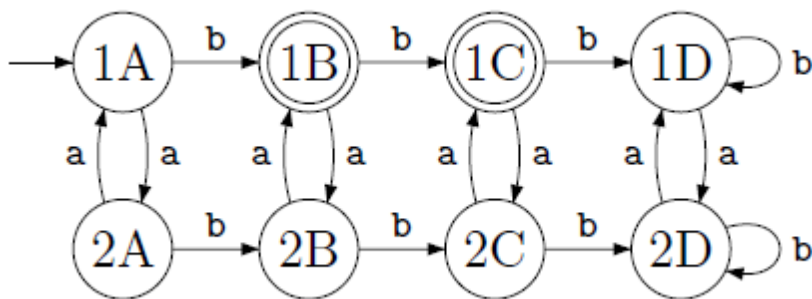
- (A) All strings of the form 0^+w , where w contains an even number of ones.
- (B) All strings of the form 000^+w , where w contains an even number of ones.
- (C) All strings of the form 00^+w , where w contains an odd number of ones.
- (D) All strings of the form 00^+w , where w contains an even number of ones.

Q128. The language accepted by the following DFA is:



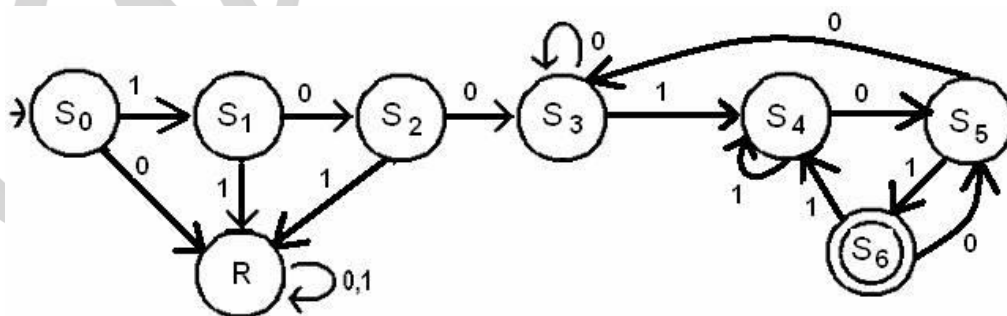
- (A) $L = \{w \mid w \text{ has number of 'a' divisible of 4 and number of 'b' divisible of 3}\}$
- (B) $L = \{w \mid w \text{ has exactly two a's and at least two b's}\}$
- (C) $L = \{w \mid w \text{ has exactly two a's and number of b's divisible by 2}\}$
- (D) $L = \{w \mid w \text{ has at least two a's and number of b's divisible by 2}\}$

Q129. The language accepted by the following DFA is:



- (A) $L = \{w \mid w \text{ has an even number of a's and one or two b's}\}$
- (B) $L = \{w \mid w \text{ has an even number of a's and at least two b's}\}$
- (C) $L = \{w \mid w \text{ has an at least two a's and one or two b's}\}$
- (D) $L = \{w \mid w \text{ has an at least two a's and one or two b's}\}$

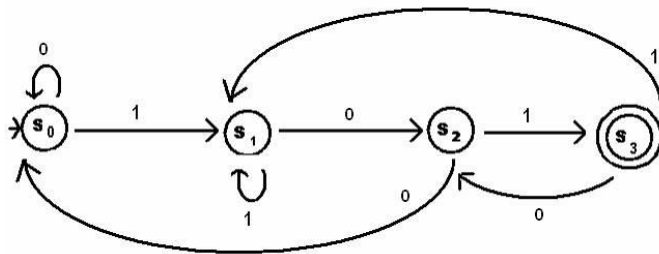
Q130. Consider the following DFA D:



If $L(D)$ is language accepted by DFA D , then which of the following language is exactly equivalent to $L(D)$?

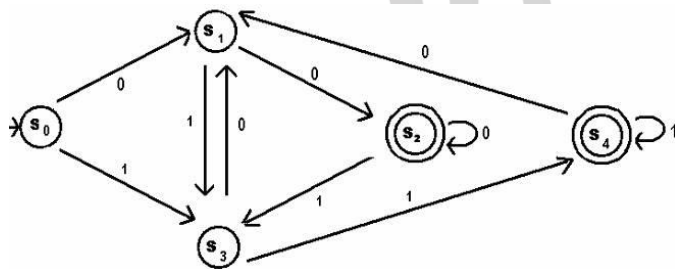
- (A) All the strings containing substring 100101
- (B) All the strings start with 10 and end with 01
- (C) All the strings start with 100 and end with 101
- (D) All the strings of length greater than or equal to six.

Q131. The following DFA accepts



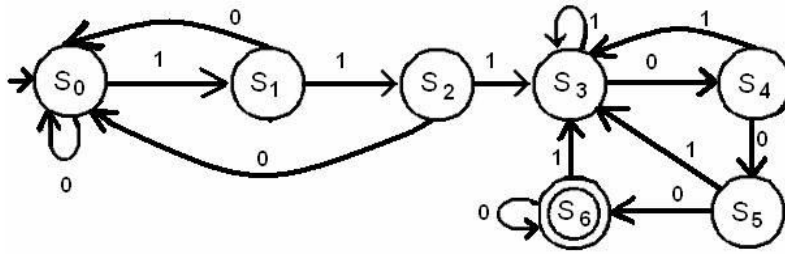
- (A) All the strings that contains 101 as substring
- (B) All the strings that ends with 101
- (C) All the strings does not end with 0
- (D) All the strings that ends with 01

Q132. The following DFA accepts



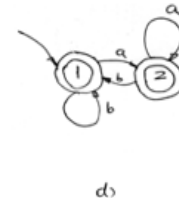
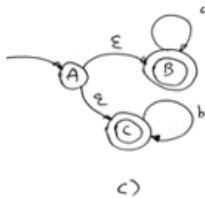
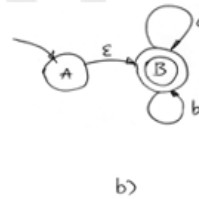
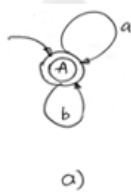
- (A) $(00+11)^+$
- (B) All the strings containing 00 and 11 as substring
- (C) All the strings end with 00 or 11
- (D) All the strings of length greater than or equal to two

Q133. The following DFA accepts

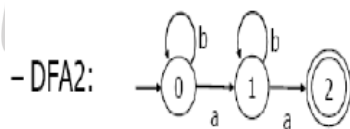
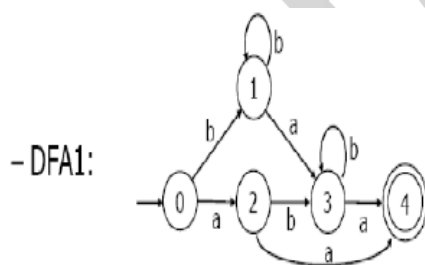


- (A) All the strings containing 111 as the substring.
- (B) All the strings containing 111 and 000 as the substring.
- (C) All the strings start with 111 and ends with 000.
- (D) All the strings containing 111 as a substring and ends with 000.

Q134. Which of the following finite automata accepts different language than other three?



Q135. Consider the following two DFAs

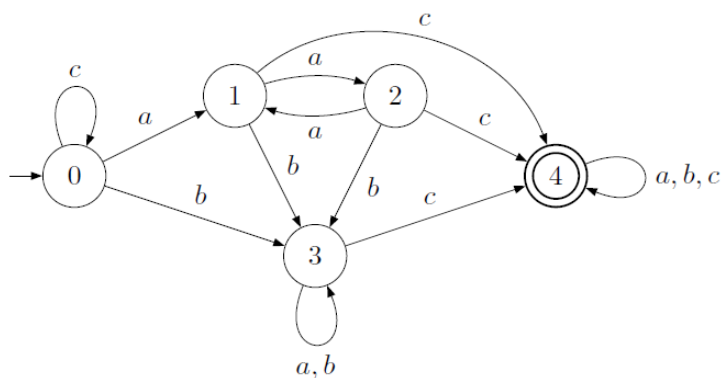


Let the language accepted by DFA1 is L_1 and that of DFA2 is L_2 . Which of the following is true?

- (A) $L_1 \subset L_2$
- (B) $L_2 \subset L_1$
- (C) $L_1 = L_2$
- (D) $L_1 \neq L_2$

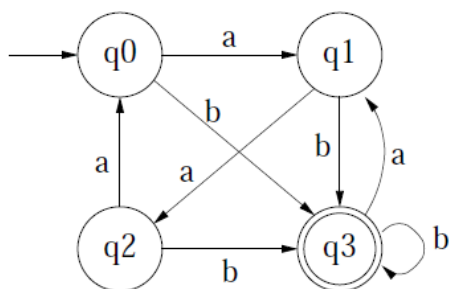
| | |
|--------------|---|
| Q136. | Minimum number of states in a deterministic finite state automaton (FSA) that recognizes all bit strings with a multiple of three 1's. (For example, the following strings are in the language: 111, 111111, 1110, 0111, 10011, but not 1, 11, 1111, 0110. _____ |
| Q137. | Minimum number of states in a deterministic finite state automaton (FSA) that recognizes all strings over a and b that has at least three a's and at least two b's i.e. $L = \{w \mid w \text{ has at least three a's and at least two b's}\}$ _____ |
| Q138. | Consider the following language: $L = \{w \in \{a, b\}^* \mid \text{every a's in } w \text{ is followed immediately by the string } bb\}$. Minimum number of states in a deterministic finite state automaton for L is _____ |
| Q139. | Minimum number of states in a DFA for the language $L = \{ab^na^m \mid n > 3, m > 2\}$ is____ |
| Q140. | The minimum number of state in the DFA for the language $L = \{w \mid (n_a(w) + 2n_b(w)) \bmod 3 < 2\}$ is _____ |
| Q141. | How many number of states are required to construct the minimized DFA for following languages over {a, b} whose languages of accepted strings (exactly) are: (i) {a, aa, aaa, aaaa}. _____ (ii) all strings not in {a, aa, aaa, aaaa}. _____ (iii) all strings whose length is divisible by 2 or 3. _____ (iv) all strings matching the regular expression $(aa \mid b)^*(bb \mid a)^*$.____ (v) all strings not matching the regular expression $(\phi^*)^*$ _____ |

Q142. Consider the following DFA D over the alphabet $\Sigma = \{a, b, c\}$:



The number of state states in minimized DFA will be _____

Q143. Consider the following DFA over $\{a, b\}$



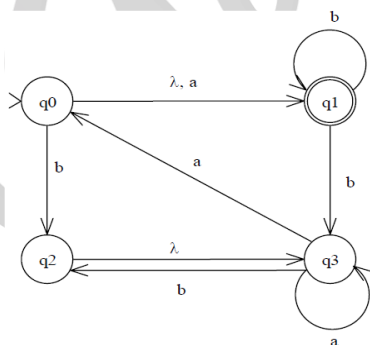
How many states does the minimized DFA have? _____

Q144. Consider the finite automaton $MDL = (\{S, S1, S2, S3, D1, D2\}, \{a, b, c\}, \delta, S, \{D1, D2\})$, where δ is defined as follows:

| | | | |
|--------------------------------|--------------------------------|--------------------------------|--------------------------------|
| $\delta(S, a) \rightarrow S1$ | $\delta(S, b) \rightarrow S2$ | $\delta(S, c) \rightarrow S3$ | $\delta(S1, a) \rightarrow D1$ |
| $\delta(S1, b) \rightarrow S2$ | $\delta(S1, c) \rightarrow S3$ | $\delta(S2, a) \rightarrow S1$ | $\delta(S2, b) \rightarrow D2$ |
| $\delta(S2, c) \rightarrow S3$ | $\delta(S3, a) \rightarrow S1$ | $\delta(S3, b) \rightarrow S2$ | $\delta(S3, c) \rightarrow D2$ |
| $\delta(D1, a) \rightarrow D1$ | $\delta(D1, b) \rightarrow D1$ | $\delta(D1, c) \rightarrow D1$ | $\delta(D2, a) \rightarrow D2$ |
| $\delta(D2, b) \rightarrow D2$ | $\delta(D2, c) \rightarrow D2$ | | |

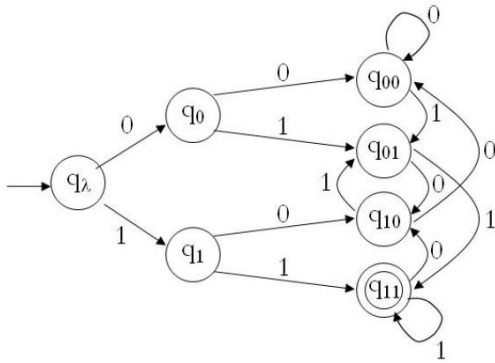
Minimum number of states in DFA of given MDL is _____

Q145. Consider the following NFA



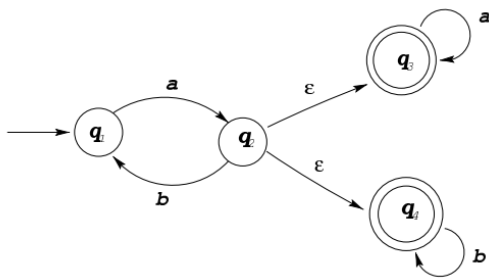
The number of states in the equivalent minimized DFA will be? _____

Q146. Consider the following NFA



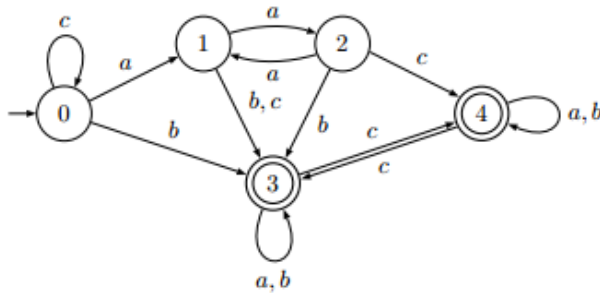
The number of states in the equivalent minimized DFA will be? _____

Q147. Consider the following NFA



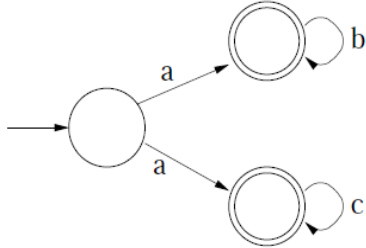
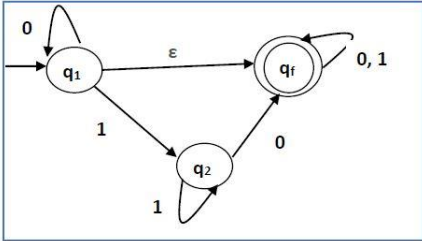
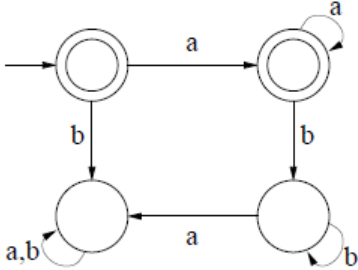
The number of states in the equivalent minimized DFA will be? _____

Q148. Consider the following DFA D over the alphabet $\Sigma_D = \{a, b, c\}$:



The numbers of states in the equivalent minimized DFA is _____

Q149. Find out how many states exist when we draw minimized DFA for the regular language given below $L = \{w x w^r \mid w \in (a+b)^+ \text{ and } x \in (a+b)^*\}$? _____

| | |
|---------------------|---|
| <p>Q150.</p> | <p>What is the minimum number of states of an equivalent DFA corresponding to Following NFA? (Assume $\Sigma = \{a, b, c\}$)</p>  <p>(A) 2 (B) 3 (C) 4 (D) 5</p> |
| <p>Q151.</p> | <p>Minimum numbers of states required to construct the equivalent DFA of the following NDFA? _____</p>  |
| <p>Q152.</p> | <p>Suppose we apply minimization to the following DFA over $\{a, b\}$:</p>  <p>Which of the following correctly describes the resulting DFA M?</p> <p>(A) M has 3 states, 1 of which is accepting. (B) M has 3 states, 2 of which are accepting. (C) M has 2 states, 1 of which is accepting. (D) M has 4 states, 2 of which are accepting</p> |
| <p>Q153.</p> | <p>Suppose that there are two DFA's D1 and D2. DFA D1 has 7 states out of which 3 states are final states. DFA D2 has 6 states out of which 4 states are final states. In the product DFA for the intersection of their languages, maximum how many final states Will be there?</p> <p>(A) 12 (B) 9 (C) 3 (D) 1</p> |

| | |
|--------------|---|
| Q154. | <p>Suppose that there are two DFA's D1 and D2. DFA D1 has 7 states out of which 3 states are final states. DFA D2 has 6 states out of which 4 states are final, In the product DFA for the union of their languages, how many final states will be there?</p> <p>(a) 42 (b) 34 (c) 2. (d) 18</p> |
| Q155. | <p>Suppose we have a DFA $D = (Q, \Sigma, \delta, q_0, F)$ and know that D accepts every string. What can we infer about D?</p> <p>(A) Every state in D is a final state. (B) There is at least 1 state in D that is not final. (C) Every reachable state from q_0 in D is a final state. (D) There is only 1 character in the alphabet.</p> |
| Q156. | <p>Let $A = \{0^n 1^m \mid n \leq m\}$ and $B = \{0^n 1^m \mid m \leq n\}$ then which options are incorrect?</p> <p>(i) $A \cup B$ is regular. (ii) $A \cap B$ is regular. (iii) Both A and B are regular.</p> <p>(A) ii & iii only (B) i & ii only (C) i & iii only (D) All</p> |
| Q157. | <p>Which of the following languages is/are regular?</p> <p>(A) $L_1 = \{wz : w = z , w \in (a + b)^* \text{ and } z \in (b + c)^*\}$. (B) $L_2 = \{w : \text{every } a \text{ in } w \text{ is followed by at least one } b \text{ and at least one } c\}$. (C) $L_3 = \{w : w \text{ does not have the same number of } a\text{'s, } b\text{'s, and } c\text{'s}\}$. (D) $L_4 = \{w : w \text{ contains the same number of patterns } ac \text{ and } abc\}$.</p> |
| Q158. | <p>Which of the following is/are regular?</p> <p>(i) $L_1 = \{a^k b^m c^n \mid (k = m \text{ or } m = n) \text{ and } k + m + n \geq 2\}$ (ii) $L_2 = \{a^k b^m c^n \mid (k = m \text{ or } m = n) \text{ and } k + m + n \leq 2\}$ (iii) $L_3 = \{a^k b^m c^n \mid k + m + n \geq 2\}$</p> <p>(A) i, ii, iii only (B) ii, iii only (C) ii only (D) iii only</p> |

| | |
|--------------|--|
| Q159. | <p>Which of the following languages over $\Sigma = \{a, b\}$ is/are not regular?</p> <p>(i) $\{a^m b^n \mid m, n \in \mathbb{N}\}$; (ii) $\{a^m b^n \mid m \leq n\}$; (iii) $\{a^m b^n \mid m + n \leq 4\}$; (iv) $\{w \in \Sigma^* \mid w \notin L\}$; where L is some given language which is regular. (v) $\{w \in \Sigma^* \mid w \notin L\}$; where L is some given language which is not regular.</p> <p>(A) i and iii only (C) ii and v only</p> <p>(B) i, iii and iv only (D) ii, iv and v only</p> |
| Q160. | <p>For which of the following languages over the alphabet $\{a, b\}$ is /are not regular?</p> <p>(i) $L_1 = \{w \mid w \text{ is not a palindrome}\}$ (ii) $L_2 = \{a^k \mid k \text{ is multiple of } 4\}$ (iii) $L_3 = \{a^k \mid k \bmod 6 = 1 \text{ or } 5\}$ (iv) $L_4 = \{wxw \mid 'x' \text{ can be any non-empty string and } w \leq 3\}$</p> <p>(A) i and iv only (C) iii and iv only</p> <p>(B) ii and iii only (D) i only</p> |
| Q161. | <p>For which of the following languages over the alphabet $\{a, b\}$ is/are regular?</p> <p>(i) $L_1 = \{ww \mid w \in \{a\}^*\}$ (ii) $L_2 = \{ww \mid w \in \{a, b\}^*\}$ (iii) $L_3 = \{w_1 w_2 \mid w_1 \in \{a\}^* \text{ and } w_2 \in \{b\}^*\}$ (iv) $L_4 = \{w \mid w \in \{a, b\}^* \text{ and } w \text{ contains the same number of a's and b's}\}$ (v) $L_5 = \{w \mid w \in \{a, b\}^* \text{ and } w \text{ contains the same number of a's and b's and that number is no more than } 128\}$</p> <p>(A) i, iii and v only (C) ii and iv only</p> <p>(B) iii and v only (D) i, ii and iii only</p> |
| Q162. | <p>Which of the following languages is/are regular?</p> <p>(A) $L_1 = \{wz \mid w = z , w \in (a + b)^* \text{ and } z \in (b + c)^*\}$. (B) $L_2 = \{w \mid \text{every } a \text{ in } w \text{ is followed by at least one } b \text{ and at least one } c\}$. (C) $L_3 = \{w \mid w \text{ does not have the same number of } a\text{'s, } b\text{'s, and } c\text{'s}\}$. (D) $L_4 = \{w \mid w \text{ contains the same number of patterns } ac \text{ and } abc\}$.</p> |

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| Q163. | <p>Which of the following language is/ are not regular?</p> <p>(i) $L1 = \{a^k! : k \geq 1\}$</p> <p>(ii) $L2 = \{a^k : k \text{ is perfect square}\}$</p> <p>(iii) $L3 = \{w \in \{a, b\}^* : w = 4 * n_a(w)\}$</p> <p>(A) L1 and L2 only</p> <p>(B) L2 and L3 only</p> <p>(C) L1 and L3 only</p> <p>(D) All the above</p> |
| Q164. | <p>Which of the following language is/ are regular?</p> <p>i) $L1 = \{w \in \{0, 1\}^* \mid (n_0(w) - n_1(w)) \bmod 3 = 0\}$.</p> <p>ii) $L2 = \{w \in \{0, 1\}^* \mid (n_0(w) - n_1(w) \bmod 3) = 1\}$.</p> <p>(A) I only</p> <p>(B) ii only</p> <p>(C) Both L1 and L2</p> <p>(D) Neither L1 nor L2</p> |
| Q165. | <p>Which of the following languages is/are regular?</p> <p>i) $L1 = \{a^n b^l : n \geq 100, l \leq 100\}$.</p> <p>ii) $L2 = \{u w w^R v : u, v, w \in \{a, b\}^+\}$</p> <p>iii) $L3 = \{u u^R v : u, v \in \{a, b\}^+\}$</p> <p>(A) i and ii only</p> <p>(B) i and iii only</p> <p>(C) i, ii & iii</p> <p>(D) None of three</p> |
| Q166. | <p>Which of the following languages is/are regular?</p> <p>$L1 = \{u w w^R v : u, v, w \in \{a, b\}^+\}$</p> <p>$L2 = \{u w w^R v : u, v, w \in \{a, b\}^+, u \geq v \}$</p> <p>$L3 = \{w w^R v : v, w \in \{a, b\}^+\}$</p> <p>(A) L1 only</p> <p>(B) L2 and L3 only</p> <p>(C) L1 and L2 only</p> <p>(D) L1 and L3 only</p> |
| Q167. | <p>Which of the following language are not regular?</p> <p>$L1 = \{a^n b^l a^k : k \geq n+1\}$</p> <p>$L2 = \{a^n b^l : n+l \geq 0\}$</p> <p>$L3 = \{a^n b^k : n \geq 100 \text{ and } k \leq 100\}$</p> <p>(A) L1 only</p> <p>(B) L2 and L3 only</p> <p>(C) L1 and L2 only</p> <p>(D) L1 and L3 only</p> |

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| <p>Q168.</p> | <p>Which of the following languages is/are regular?</p> <p>$L1 = \{a^k b^k : k \geq 1\} \cup \{a^k b^l : k \geq 1, l \geq 1\}$</p> <p>$L2 = \{a^{k+1} b^{k+1} : k \geq 0\} \cup \{a^{k+1} b^{k+3} : k \geq 0\}$</p> <p>$L3 = \{ww : w \in \{a\}^+\}$</p> <p>(A) All the above</p> <p>(B) L2 and L3 only</p> <p>(C) L1 and L2 only</p> <p>(D) L1 and L3 only</p> |
| <p>Q169.</p> | <p>Which of the following language is/ are regular?</p> <p>1. $L = \{a^n b^m a^n : m, n \geq 0\}$</p> <p>2. $L = \{a^n a^n a^n b^m b^m b^m : m, n \geq 0\}$</p> <p>3. $L = \{a^n b^m : m \neq n \text{ and } m \leq 10, n \geq 0\}$</p> <p>(A) 2 only</p> <p>(B) 2 & 3 only</p> <p>(C) 1 only</p> <p>(D) none of these</p> |
| <p>Q170.</p> | <p>Which of the following language is/ are not regular?</p> <p>1. $L = \{xy : x \in L1 \text{ and } y \notin L1; \text{ where } L1 \text{ is regular}\}$</p> <p>2. $L = \{w1w2 : w1, w2 \in L2; \text{ where } L2 \text{ is regular}\}$</p> <p>3. $L = \{a^{i^3} : i \geq 0 \text{ and } \Sigma = \{a\}\}$</p> <p>(A) 1 & 3 only</p> <p>(B) 2 only</p> <p>(C) 3 only</p> <p>(D) All</p> |
| <p>Q171.</p> | <p>[MSQ]</p> <p>Which of the following languages is/are regular?</p> <p>(A) $L1 : \{wxw^R \mid w, x \in \{a, b\}^*\}$, w^R is the reverse of string w</p> <p>(B) $L2 : \{a^n b^{2m} \mid m \leq 1000 \text{ or } n \geq 1000\}$</p> <p>(C) $L3 : \{a^p b^q c^r \mid p < 10 \text{ and } q > 100 \text{ and } r \geq 0\}$</p> <p>(D) $L4 : \{a^n b^n w \mid n \geq 0 \text{ and } w \in \{a, b\}^*\}$</p> |

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| Q172. | <p>[MSQ]</p> <p>Which of the following languages is/are not regular?</p> <p>(A) $L1 = \{w^R x w : w, x \in \Sigma^*\}$</p> <p>(B) $L2 = \{w w^R x : w, x \in \Sigma^+\}$</p> <p>(C) $L3 = \{w x w^R : w, x \in \Sigma^+\}$</p> <p>(D) $L4 = \{w x w : w, x \in \Sigma^+\}$</p> |
| Q173. | <p>[MSQ]</p> <p>Which of the following languages is/are not regular?</p> <p>(A) $L = \{w : n_a(w) \neq n_b(w)\}$</p> <p>(B) $L = \{a^i b^j c^k : i \geq j+k\}$</p> <p>(C) $L = \{a^i b^j c^k : j \neq 2i+k\}$</p> <p>(D) $L = \{a^i b^j c^k : i = j \text{ or } j \neq k\}$</p> |
| Q174. | <p>If $L1 = \{1^P : P \text{ is a prime number}\}$ and $L2 = \{1^{2^i} : i \geq 0\}$ then which of the following statement is/ are true?</p> <p>I. $\{(L1)^+ \cup \phi^*\}$ is regular.</p> <p>II. $(L2)^*$ is also regular.</p> <p>III. $(L1.(L2 \cup \lambda))^*$ is also regular.</p> <p>(A) I & II only</p> <p>(B) II & III only</p> <p>(C) I & III only</p> <p>(D) All the above</p> |
| Q175. | <p>[MSQ]</p> <p>Let $L1$ and $L2$ are two regular languages over Σ then which of the following is/are regular?</p> <p>(A) $L = \{x \in \Sigma^* \mid \text{either } x \in L1 \text{ or } x \in L2\}$</p> <p>(B) $L = \{x \in \Sigma^* \mid x \in L1 \text{ but } x \notin L2\}$</p> <p>(C) $L = \{x \in \Sigma^* \mid xy \in L1 \text{ or } xy \in L2; \text{ where } y \in \Sigma^*\}$</p> <p>(D) $L = \{x \in \Sigma^* \mid xy \in L1 \text{ and } x = y \}$</p> |

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| Q176. | <p>Which of the following language is regular?</p> <p>(A) $L = \{a^k : k \text{ is not perfect square}\}$</p> <p>(B) $L = \{a^k : k \text{ is perfect cube}\}$</p> <p>(C) $L = \{a^k : k \text{ is either prime or product of two or more prime numbers}\}$</p> <p>(D) $L = \{a^k : k = 2^i \text{ for some } i \geq 0\}$</p> |
| Q177. | <p>[MSQ]</p> <p>Which of the following language/s is/are not regular?</p> <p>(A) $L = \{a^i b^j c^k : i+j+k > 4\}$</p> <p>(B) $L = \{a^i b^j c^k : i < 10, j > 5 \text{ and } k > i\}$</p> <p>(C) $L = \{a^i b^j : i+2j \text{ is a prime number}\}$</p> <p>(D) $L = \{a^i b^j : i - j = 3\}$</p> |
| Q178. | <p>[MSQ]</p> <p>Which of the following language/s is are regular?</p> <p>(A) $L = \{a^i b^j : i = j \text{ or } i < j \text{ or } i > j\}$</p> <p>(B) $L = \{a^n b^n : n \geq 1\} \cup \{a^n b^m : n, m \geq 1\}$</p> <p>(C) $L = \{w_1 w_2 : w_1 = w_2 \text{ and } w_1, w_2 \in \Sigma^*\} \cap \{a^n b^n : n \geq 1\}$</p> <p>(D) $L = \{a^n : n = k^3 \text{ for some } k \geq 0\}$</p> |
| Q179. | <p>Which of the following is/are true?</p> <p>(a) Union of two non-regular languages is always non-regular.</p> <p>(b) Union of a regular language with a disjoint non-regular language is always non-regular.</p> <p>(c) $L((ab^*ba^*) \cap (ba^*ab^*)) = \{\epsilon\}$.</p> <p>(d) $L = \{1^n : n \leq 1000 \text{ and } n \text{ is prime}\}$. A DFA accepting L may have less than 900 states.</p> |
| Q180. | <p>Which of the following languages are not regular?</p> <p>1. $L = \{www^R : w \in \{a, b\}^*\}$</p> <p>2. $L = \{w \in \{a, b\}^* \mid w = w^R\}$</p> <p>3. $L = \{w \in \{a, b\}^* \mid w \text{ has more } a\text{'s than } b\text{'s}\}$</p> <p>4. $L = \{a^{2n} b^{4n} a^n\}$</p> <p>(A) 1 & 2 only</p> <p>(B) 2, 3 & 4 only</p> <p>(C) 1, 3 & 4 only</p> <p>(D) All the above</p> |

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| <p>Q181.</p> | <p>Given that L is regular and</p> $L1 = \{wx: w \in \Sigma^* \text{ and } x \in L\},$ $L2 = \{w: wx \in L \text{ and } w = x \}, \text{ and}$ $L3 = \{w: wx \in L1, \text{ for some } x \in L2\}$ <p>Which of the above language is/are regular?</p> <p>(A) L1, L2 only</p> <p>(B) L2, L3 only</p> <p>(C) L1, L3 only</p> <p>(D) All are regular</p> |
| <p>Q182.</p> | <p>Which of the following statement/s is are true?</p> <p>I. If A is a non-regular language and B is a language such that $B \subseteq A$, then B must be Non-regular.</p> <p>II. If $(L1.L2 \cup L3)$ is regular, L3 is regular and complement of L2 is regular then L1 must be regular.</p> <p>(A) I only</p> <p>(B) II only</p> <p>(C) Both</p> <p>(D) Neither I nor II</p> |
| <p>Q183.</p> | <p>Select the correct statement</p> <ol style="list-style-type: none"> 1. A DFA with n states must accept at least one string of length greater than n. 2. A DFA with n states that accepts an infinite language must accept at least one string x such that $2n < x \leq 3n$. 3. If R is a regular language and L is some language, and $L \cup R$ is a regular language, then L must be a regular language. 4. If F is a finite language and L is some language, and $L - F$ is a regular language, then L must be a regular language. <p>(A) 2 only</p> <p>(B) 1 and 3 only</p> <p>(C) 2 and 4 only</p> <p>(D) None</p> |

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| <p>Q184.</p> | <p>Which of following is/are correct?</p> <p>S1: Let r_1 and r_2 be two regular expressions. Then $L((r_1 + r_2)^*) = L((r_1^* r_2^*)^*)$.</p> <p>S2: Let $L_4 = L_1 \cdot L_2 \cdot L_3$. If L_1 and L_2 are regular and L_3 is not regular, it is possible that L_4 is regular.</p> <p>(A) Only S1</p> <p>(B) Only S2</p> <p>(C) Both S1&S2</p> <p>(D)) Neither S1 nor S2</p> |
| <p>Q185.</p> | <p>Which of following is/are correct?</p> <p>S1: $L_1 \subseteq L \subseteq L_2$ where L_1 and L_2 are regular, then L must be regular.</p> <p>S2: $\{w = xyzzy \mid x, y, z \in \{0, 1\}^+\}$ is regular.</p> <p>(A) Only S1</p> <p>(B) Only S2</p> <p>(C) Both S1&S2</p> <p>(D)) Neither S1 nor S2</p> |
| <p>Q186.</p> | <p>Which of following statement is/are correct?</p> <p>S1: $(\emptyset^* \cdot \emptyset)^* \cdot \emptyset^* = \emptyset$</p> <p>S2: $\{xyx^R \mid x, y \in \{a, b\}^+\}$ is regular.</p> <p>(A) Only S1</p> <p>(B) Only S2</p> <p>(C) Both S1&S2</p> <p>(D) Neither S1 nor S2</p> |
| <p>Q187.</p> | <p>Which of following statement is/are correct?</p> <p>S1: If $L_1.L_2$ is regular then at least one of them (L_1 or L_2) is regular.</p> <p>S2: If $L_1.L_2$ is non-regular then at least one of them (L_1 or L_2) is non-regular.</p> <p>(A) Only S1</p> <p>(B) Only S2</p> <p>(C)Both S1&S2</p> <p>(D)None of them</p> |

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| Q188. | <p>Select the correct statement:</p> <p>(A) If L_1 is regular and $L_2 \subseteq L_1$, then L_2 is regular as well.</p> <p>(B) If L_1 is regular and L_2 is not regular, then $L_1 \cup L_2$ is not regular.</p> <p>(C) If L_1 is regular and $L_1 \cup L_2$ is not regular, then L_2 is not regular.</p> <p>(D) If L_1 is regular and L_2 is not regular, then $L_1 \cap L_2$ is not regular.</p> |
| Q189. | <p>Select the correct statements</p> <p>(1) $L_1 = L_2$ if and only if $L_1^* = L_2^*$</p> <p>(2) For any languages L_1, L_2 and L_3, $L_1 (L_2 \cap L_3) \subseteq (L_1 L_2) \cap (L_1 L_3)$</p> <p>(3) For any languages L_1, L_2 and L_3, $(L_1 L_2) \cap (L_1 L_3) \subseteq L_1 (L_2 \cap L_3)$.</p> <p>(A) 1 and 3</p> <p>(B) 2 and 3</p> <p>(C) 2 only</p> <p>(D) None.</p> |
| Q190. | <p>Consider the following statements:</p> <ol style="list-style-type: none"> 1. An infinite language can have an infinite complement. 2. All infinite languages have infinite complements. 3. The union of infinitely many regular languages is always regular. <p>Which of the above statements is/are true?</p> <p>(A) 1 only</p> <p>(B) 1 and 2 only</p> <p>(C) 1 and 3 only</p> <p>(D) 2 and 3 only</p> |
| Q191. | <p>Consider the following statements</p> <ol style="list-style-type: none"> 1. If L is regular then so is $\{xx : x \in L\}$. 2. If L is regular then so is $\{xy : x, y \in L\}$. 3. Let $A = \{1^{2^p} : p \text{ is prime}\}$. Then A^* is regular. <p>Which of the above statements is true?</p> <p>(A) 1 only</p> <p>(B) 1 and 2 only</p> <p>(C) 1 and 3 only</p> <p>(D) 2 and 3 only</p> |

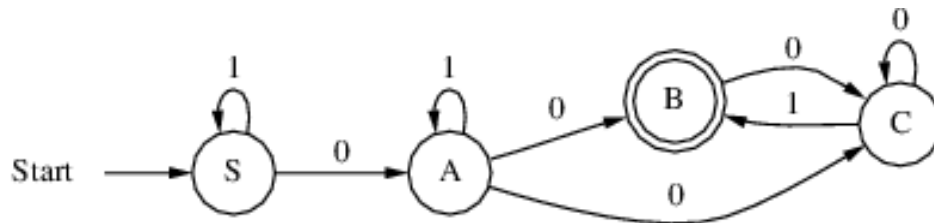
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| <p>Q192.</p> | <p>Select correct statements:</p> <p>S1: If L_1 and L_2 are non-regular languages then $L_1 \cup L_2$ is also non-regular.</p> <p>S2: If L_1 is non-regular language and L_2 is regular language then $L_1.L_2$ may be regular language.</p> <p>S3: If $L_1 \cup L_2$ and L_2 are regular then L_1 may not be regular.</p> <p>(A) S1 and S2 only</p> <p>(B) S2 and S3 only</p> <p>(C) S1 and S3 only</p> <p>(D) All the above.</p> |
| <p>Q193.</p> | <p>Select correct statements:</p> <p>S1: For every regular language L, every subset of L is regular as well.</p> <p>S2: Every non-regular language is infinite.</p> <p>S3: The intersection of any two non-regular languages is non-regular.</p> <p>S4: If each of the languages L_1, L_2, \dots is regular, then $\bigcup_{i=1}^{\infty} L_i$ is regular as well.</p> <p>(A) S1 and S2 only</p> <p>(B) S2 and S3</p> <p>(C) S3 and S4</p> <p>(D) S2 and S4</p> |
| <p>Q194.</p> | <p>Let L_1 and L_2 are two languages over Σ and it is given that $L_1.L_2$ is non-regular languages. Then which of the following statement is not always true?</p> <p>(A) $\overline{L_1}$ is not regular.</p> <p>(B) $L_1 \cup \overline{L_1}$ is regular.</p> <p>(C) $L_1, L_2 \subseteq \Sigma^*$</p> <p>(D) $L_1 \cup L_2$ contain infinitely many strings.</p> |
| <p>Q195.</p> | <p>Let $L = 01^* + 10^*$. Which of the following is regular expression of L^R (reverse of L)?</p> <p>(A) $0^*1 + 1^*0$</p> <p>(B) $1^*0 + 01^*$</p> <p>(C) $1^*0 + 0^*1$</p> <p>(D) $(10)^* + (01)^*$</p> |

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| Q196. | The set of non-regular languages is closed under which of the following operations? (A) Complement (B) Union (C) Intersection (D) Concatenation |
| Q197. | Let $L1 = \{a^n b^n c^m \mid n, m \geq 0\}$ and $L2 = \{a^n b^m c^m \mid n, m \geq 0\}$ then what is $\overline{\overline{L1} \cap \overline{L2}}$? (A) $L1 \cup \overline{L2}$ (B) $\Sigma^* - \{a^n b^n c^n \mid n \geq 0\}$ (C) $\{a^n b^n c^n \mid n \geq 0\}$ (D) $a^* b^* c^*$ |
| Q198. | The right quotient of a language $L1$ with respect to $L2$ is defined as $L1/L2 = \{x: y \in L2, xy \in L1\}$ Let $L1 = L(a^*ba^+)$ and $L2 = \{aba^*\}$ then what will be the $L1/L2$? (A) a^* (B) a^+ (C) a^*b (D) a^+b |
| Q199. | Which of the following strings is NOT in the Kleene closure of the language $\{abb, ba, bba\}$? (A) abbbabbaabb (B) abbbba (C) abbbbababb (D) abbbabbababbaba |
| Q200. | Let $L1 = \{a, ab, c, d, \lambda\}$, $L2 = \{d\}$ and $L3 = L1.L2$. Which string is not in $L3$? (A) a (B) abd (C) cd (D) d |
| Q201. | If $L = \{001, 1101, 101\}$ then the prefix of L is (A) $\{\lambda, 0, 00, 001, 1, 11, 110, 1101, 1, 10, 101\}$ (B) $\{\lambda, 0, 00, 001, 1, 11, 110, 1101, 10, 101\}$ (C) $\{\lambda, 1, 01, 001, 101, 1101\}$ (D) None of the above |

Q202. If $L = \{001, 1101, 101\}$ then the suffix of L is

- (A) $\{\lambda, 0, 00, 001, 1, 11, 110, 1101, 1, 10, 101\}$
- (B) $\{\lambda, 1, 01, 001, 101, 1101, 10, 110\}$
- (C) $\{\lambda, 1, 01, 001, 101, 1101\}$
- (D) None of the above

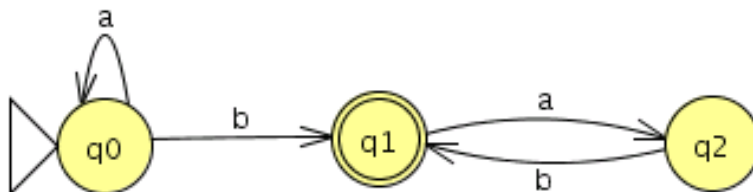
Q203. Consider the following NFA:



Which of the following grammar generates the language accepted by NFA given above?

- (A) $S \rightarrow 1S/0A, A \rightarrow 1A/0B/0C, B \rightarrow 0C, C \rightarrow 1B/0C$
- (B) $S \rightarrow 1S/0A, A \rightarrow 1A/0B, B \rightarrow 0C/\lambda, C \rightarrow 1B/0C$
- (C) $S \rightarrow 1S/0A, A \rightarrow 1A/0B/0C, B \rightarrow 0C/\lambda, C \rightarrow 1B/0C/\lambda$
- (D) $S \rightarrow 1S/0A, A \rightarrow 1A/0B/0C, B \rightarrow 0C/\lambda, C \rightarrow 1B/0C$

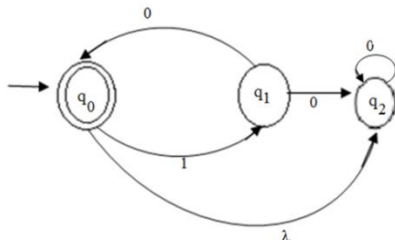
Q204. Consider the following NFA



Which Regular grammar corresponds to the following NFA?

- (A) $S \rightarrow bA \mid aS, A \rightarrow aB \mid \lambda, B \rightarrow bA$
- (B) $S \rightarrow bA \mid aS, A \rightarrow bB \mid \lambda, B \rightarrow bA$
- (C) $S \rightarrow Ba, A \rightarrow aB \mid \lambda, B \rightarrow bA$
- (D) None

Q205. Consider the following NFA.



Which of the following is a left linear grammar for the language accepted by M?

- (A) $S \rightarrow 1S0 \mid \lambda$
- (B) $S \rightarrow 10S \mid \lambda$
- (C) $S \rightarrow 1A \mid \lambda, A \rightarrow 0S$
- (D) $S \rightarrow A0 \mid \lambda, A \rightarrow S1$

Q206. How many of the following statement(s) is/are true? _____

- (i) The language generated by the grammar $G = (\{S\}, \{a\}, S, \{S \rightarrow Saaa \mid aS \mid a\})$ is not regular.
- (ii) The language generated by the grammar $G = (\{S\}, \{a\}, S, \{S \rightarrow aSSa \mid aS \mid a\})$ is regular.
- (iii) The language generated by the grammar $G = (\{S\}, \{a\}, S, \{S \rightarrow bSb \mid aSa \mid a \mid b\})$ is regular.
- (iv) The language generated by the grammar $G = (\{S\}, \{a\}, S, \{S \rightarrow bSa \mid \lambda \mid a\})$ is not regular.

Q207. Consider a generated grammar G with production

$S \rightarrow abA$
 $A \rightarrow baB$
 $B \rightarrow aA \mid bb$

How many states are required to construct minimized DFA for language accepted by G?

- (A) 2
- (B) 3
- (C) 4
- (D) more than 4

Data for next two question: Consider a language $L = \{a^n b^m : n \geq 2, m \geq 3\}$

Q208. Which of the following grammar is left-linear grammar for given L?

- (A) $S \rightarrow S_1 S_2, S_1 \rightarrow aaA, A \rightarrow aA \mid \lambda, S_2 \rightarrow bbbB, B \rightarrow bB \mid \lambda$
- (B) $S \rightarrow S_1 bbb, S_1 \rightarrow S_1 b \mid A, A \rightarrow aaB, B \rightarrow aB \mid \lambda$
- (C) $S \rightarrow aaA, A \rightarrow aA \mid B, B \rightarrow bbbC, C \rightarrow bC \mid \lambda$
- (D) none of these

Q209. Which of the following grammar is right-linear grammar for given L?

- (A) $S \rightarrow S_1 S_2, S_1 \rightarrow aaA, A \rightarrow aA \mid \lambda, S_2 \rightarrow bbbB, B \rightarrow bB \mid \lambda$
- (B) $S \rightarrow S_1 bbb, S_1 \rightarrow S_1 b \mid A, A \rightarrow aaB, B \rightarrow aB \mid \lambda$
- (C) $S \rightarrow aaA, A \rightarrow aA \mid B, B \rightarrow bbbC, C \rightarrow bC \mid \lambda$
- (D) None of these

Q210. Consider following grammar G

$S \rightarrow aA \mid bB \mid \lambda,$

$A \rightarrow bC \mid aS,$

$B \rightarrow aC \mid Bs$

$C \rightarrow aB \mid bA$

Which of the following language is generated by G?

- (A) $L = \{w : n_a(w) + n_b(w) \text{ is even}\}$
- (B) $L = \{w : |n_a(w) - n_b(w)| \text{ is even}\}$
- (C) $L = \{w : n_a(w) \text{ and } n_b(w) \text{ are both even}\}$
- (D) $L = \{w : n_a(w) \text{ and } n_b(w) \text{ are both odd}\}$

Q211. Consider following grammar G

$S \rightarrow aA \mid bB$

$A \rightarrow abA \mid \lambda$

$B \rightarrow ccB \mid \lambda$

Which of the following regular expression is for language generated by G?

- (A) $a(ab)^*b(cc)^*$
- (B) $a(ab)^+b(cc)^+$
- (C) $a(ab)^*+b(cc)^*$
- (D) $a(ab)^++b(cc)^+$

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| Q212. | <p>The language generated by the grammar with productions where $\Sigma = \{a, b\}$ is?</p> <p>$S \rightarrow AaAaAaA,$ $A \rightarrow aA \mid bA \mid \lambda.$</p> <p>(A) All the strings over Σ^* with at least three a's. (B) All the strings over Σ^* with at most three a's (C) All the strings over Σ^* with at least three a's or three b's (D) None of these</p> |
| Q213. | <p>Consider the following grammar G:</p> <p>$S \rightarrow XY$ $X \rightarrow aX \mid bX \mid a$ $Y \rightarrow Ya \mid Yb \mid b$</p> <p>The regular expression for the language generated by G is</p> <p>(A) $(a + b)^+$ (B) $(a + b)^*ab(a + b)^*$ (C) $(a + b)^+ab(a + b)^+$ (D) $(a + b)^+(a + b)^+$</p> |
| <p>For the next two questions: Let $\Sigma_1 = \{0, 1, 2\}$ and $\Sigma_2 = \{a, b, c, d\}$ and define $h(0) = aab$, $h(1) = aabc$, $h(2) = cccd$.</p> | |
| Q214. | <p>Then holomorphic image of $L = \{010, 102, 1011, 0100\}$ is</p> <p>(A) $h(L) = \{aabaabcaab, aabcaabcccd, aabcaabaabaab, aabaabcaabaab\}$ (B) $h(L) = \{aabaabcaab, aabcaabcccd, aabcaabaabaabc, aabaabcaabaab\}$ (C) $h(L) = \{aabaabcaab, aabcaabcccd, aabcaabaabaab, aabaabcaabaabc\}$ (D) $h(L) = \{aabaabcaab, aabcaabcccd, aabcaabaabcaabc, aabaabcaabaab\}$</p> |
| Q215. | <p>Let $h(L) = \{aabaabaabcccdcaab, cccdaabcccdcaabaabc\}$ then L will be</p> <p>(A) $\{00020, 21201\}$ (B) $\{00120, 20201\}$ (C) $\{00120, 21201\}$ (D) $\{00020, 20201\}$</p> |

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| Q216. | <p>Let $A = \{xx \mid x \in \{a, b\}^*\}$. Consider homomorphism $h : \{a, b\}^* \rightarrow \{0, 1\}^*$ with $h(a) = 00$, $h(b) = \epsilon$. What is $h(A)$?</p> <p>(A) $\{0^{4n} \mid n \geq 0\}$</p> <p>(B) $\{0^n \mid n \geq 0\}$</p> <p>(C) $\{0^{2n} \mid n \geq 0\}$</p> <p>(D) none of these</p> |
| <p>For next three questions suppose h is the homomorphism from $\{0,1,2\}$ to $\{a,b\}$ defined by $h(0) = a$; $h(1) = ab$; $h(2) = ba$.</p> | |
| Q217. | <p>What is $h(21120)$</p> <p>(A) baababbaa</p> <p>(B) bababbaa</p> <p>(C) bbaababbaa</p> <p>(D) abaababbaa</p> |
| Q218. | <p>If $L = 01^*2$, then what is $h(L)$?</p> <p>(A) $aab(ab)^*ba$</p> <p>(B) $(ab)^*ba$</p> <p>(C) $a(ab)^*ba$</p> <p>(D) $aa(ab)^*ab$</p> |
| Q219. | <p>If $L = a(ba)^*$, then what is $h^{-1}(L)$?</p> <p>(A) 02^*1^*0</p> <p>(B) 02^*</p> <p>(C) 1^*0</p> <p>(D) $02^* \cup 1^*0$</p> |
| Q220. | <p>The pumping lemma for regular languages implies that</p> <p>(A) Every regular language contains a string that can be pumped.</p> <p>(B) All strings in a regular language can be written as uvw so that uv^iw is also in the language when $i \geq 0$.</p> <p>(C) A regular language is infinite if and only if it contains a string that can be pumped</p> <p>(D) Regular languages are closed under the regular operations</p> |

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| Q221. | <p>Let $L = \{a^n b^m \mid n > m \geq 0\}$ is</p> <p>(A) not regular because $a^{p+1}b^p$ cannot be pumped</p> <p>(B) regular because it is a subset of a^*b^*</p> <p>(C) regular because it is described by a regular expression</p> <p>(D) not regular because $a^p b^p$ cannot be pumped.</p> |
| Q222. | <p>[MSQ]</p> <p>Consider the language $L = \{a^i b^j \mid j < i\}$. Which of the following string can be used in a pumping lemma proof that L is not regular? [Assume n is pumping length]</p> <p>(A) $a^n b^n$</p> <p>(B) $a^{2n} b^n$</p> <p>(C) $a^{2n+1} b^{2n}$</p> <p>(D) $(ab)^n a$</p> |
| Q223. | <p>The pumping lemma for regular languages can be proved by</p> <p>(A) showing that an NFA can be converted to an equivalent DFA.</p> <p>(B) showing that the regular languages are closed under the regular operations.</p> <p>(C) showing that an NFA computation must repeat a state.</p> <p>(D) showing that a DFA computation must repeat a state.</p> |
| Q224. | <p>Consider the language $L = \{w \in \Sigma^* \mid w \text{ is a palindrome}\}$ is not regular.</p> <p>Which of the following are good choices of a string to pick to show that L is not regular with the help of pumping lemma?</p> <p>(A) $a^p b b a^p$</p> <p>(B) a^p</p> <p>(C) $b a^p b$</p> <p>(D) All of the above are good choices of strings</p> |
| Q225. | <p>$\{a^n b^m : n > m > 0\}$ is</p> <p>(A) not regular because $a^{p+1}b^p$ cannot be pumped</p> <p>(B) regular because it is a subset of a^*b^*</p> <p>(C) not regular because $a^p b^p$ cannot be pumped</p> <p>(D) regular because it is described by a regular expression</p> |

CONTEXT FREE LANGUAGE

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| Q1. | <p>Which of the following languages is/are context free languages?</p> <p>(I) $L = \{a^n b^j \mid n \leq j^2\}$</p> <p>(II) $L = \{a^n b^j c^k \mid k = j^* n\}$</p> <p>(III) $L = \{a^n b^j c^k \mid n < j, n \leq k \leq j\}$</p> <p>(A) I & II only</p> <p>(B) II & III only</p> <p>(C) I & III only</p> <p>(D) None of the above</p> |
| Q2. | <p>Which of the following language is/are NCFL?</p> <p>(I) $L = \{a^m b^n c^p d^q \mid n = q \text{ or } m \leq p \text{ and } m + n = p + q\}$</p> <p>(II) $L = \{a^m b^n c^p \mid m = n \text{ or } n = p \text{ or } m = p\}$</p> <p>(III) $L = \{a^m b^n c^p \mid m = n \text{ and } n = p \text{ and } m = p\}$</p> <p>(IV) $L = \{a^m b^n c^p \mid m \neq n \text{ or } n \neq p \text{ or } m \neq p\}$</p> <p>(A) II & IV only</p> <p>(B) I & IV only</p> <p>(C) I & II only</p> <p>(D) None of above</p> |
| Q3. | <p>Which of the following language is/are context free?</p> <p>$L_1 = \{a^n b^j \mid n \leq j^2\}$</p> <p>$L_2 = \{a^n \mid n \text{ is prime}\}$</p> <p>$L_3 = \{a^n w w^R a^n \mid n \geq 0 \text{ and } w \in \{a, b\}^*\}$</p> <p>$L_4 = \{a^n b^j a^i b^n \mid n \leq j\}$</p> <p>$L_5 = \{xy \mid x, y \in \{0, 1\}^* \text{ and } x \neq y\}$</p> <p>(A) L_3 and L_4 only</p> <p>(B) L_1 and L_4 only</p> <p>(C) L_2 and L_5 only</p> <p>(D) L_3 and L_5 only</p> |
| Q4. | <p>Which of the following language is/are context free?</p> <p>$L_1 = \{a^n b^n a^n b^n \mid n \geq 0\}, \Sigma = \{a, b\}.$</p> <p>$L_2 = \{w^R \# z \mid w \text{ is a substring of } z, w, z \in \{a, b\}^*\}, \Sigma = \{a, b, \#\}.$</p> <p>$L_3 = \{w \# z \mid w \text{ is a substring of } z, w, z \in \{a, b\}^*\}, \Sigma = \{a, b, \#\}.$</p> <p>$L_4 = \{x+y=z \mid x+y=z \text{ in unary where } x, y, z \in 1^*\}, \Sigma = \{1, +, =\}.$</p> <p>(A) L_3 and L_4 only</p> <p>(B) L_2 and L_4 only</p> <p>(C) L_4 only</p> <p>(D) None</p> |

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| <p>Q5.</p> | <p>Which of the following language is/are NCFL but not DCFL?</p> <p>$L_1 = \{wyw^R : \text{the length of } y \text{ is even}\}, \Sigma = \{a, b\}.$</p> <p>$L_2 = \{w : w \text{ has the same number of } a\text{'s } b\text{'s and } c\text{'s together}\}, \Sigma = \{a, b, c\}.$</p> <p>$L_3 = \{a^i b^j c^k : i > j \text{ or } j > k, \text{ where } i, j, k \geq 0\}, \Sigma = \{a, b, c\}.$</p> <p>$L_4 = \{xy : x = y \text{ and } x \neq y\}, \Sigma = \{a, b\}.$</p> <p>(A) L_3 only (B) L_1 and L_2 only (C) L_2 and L_4 only (D) L_1, L_3 and L_4 only</p> |
| <p>Q6.</p> | <p>Which of the following language is/are DCFL?</p> <p>$L_1 = \{ww^R \mid w \in a^*b^*\}$</p> <p>$L_2 = \{w \mid w \in \{a, b\}^* \text{ and each prefix of } w \text{ has at least as many } a\text{'s as } b\text{'s}\}$</p> <p>$L_3 = \{w \mid w \in \{a, b\}^* \text{ and } m \times n_a(w) = k \times n_b(w) \text{ for arbitrary constants } m \text{ and } k\}$</p> <p>(A) Only L_2 (B) L_2 and L_3 (C) L_1 and L_3 (D) All of the above</p> |
| <p>Q7.</p> | <p>Which of the following language is/are NCFL?</p> <p>(i) $\{0^n 1^n \mid n > 0\} \cup \{0^n 1^{2n} \mid n \geq 0\}$</p> <p>(ii) $\{1^n 0^n 1^m \mid m, n > 0\} \cup \{1^m 0^n 1^{2n} \mid m, n \geq 0\}$</p> <p>(iii) $\{0^n 1^n \mid n > 0\} \cup \{1^{2n} 0^n \mid n \geq 0\}$</p> <p>(A) i only (B) i and ii only (C) iii only (D) ii only</p> |
| <p>Q8.</p> | <p>Which of the following language is/are context free?</p> <p>(i) $L = \{a^n b^n : n \geq 0 \text{ and } n \text{ is not a multiple of } 5\}$</p> <p>(ii) $\{1^k 0^i 1^j 0^l 1^k \mid i, j, k > 0\}$</p> <p>(iii) $\{w \# x \mid w \text{ is a substring of } x, \text{ where } w, x \text{ are in } \{0, 1\}^*\}.$</p> <p>(iv) $\{0^i 1^j 0^i \mid i, j > 0\}$</p> <p>(v) $\{(0^n 1^n)^m \mid m, n > 0\}.$</p> <p>(vi) Complement of $\{(0^n 1^n)^m \mid m, n > 0\}.$</p> <p>(A) i, ii and vi only (B) i and ii only (C) iii and v only (D) i, ii, iii and v only</p> |

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| <p>Q9.</p> | <p>Which of the following language is/are context free?</p> <p>$L1 = \{w \in \{a, b, c\}^* \mid n_a(w) = n_b(w) \leq n_c(w)\}$</p> <p>$L2 = \{a^n b^m \mid n \text{ and } m \text{ are both prime}\}$</p> <p>$L3 = \{a^n b^j a^k b^l \mid n \leq k, j \leq l\}$</p> <p>$L4 = \{w \in \{a, b\}^* \mid n_a(w) = n_b(w), w \text{ does not contain a substring } aab\}$</p> <p>(A) L2 only (B) L1 and L3 only</p> <p>(C) L4 only (D) None</p> |
| <p>Q10.</p> | <p>The language $L = \{(a+b+c)^* w c w^R, \text{ where } w \in (a+b)^+\}$ is</p> <p>(A) Regular (B) Deterministic CFL</p> <p>(C) Non- Deterministic CFL (D) Not a CFL</p> |
| <p>Q11.</p> | <p>The language $L = \{w^R c w (a + b + c)^*, \text{ where } w \in (a+b)^+\}$ is</p> <p>(A) Regular (B) Deterministic CFL</p> <p>(C) Non- Deterministic CFL (D) Not a CFL</p> |
| <p>Q12.</p> | <p>Which of the following is/are not CFL?</p> <p>(i) $L = \{www \mid w \in \{0, 1\}^*\}$</p> <p>(ii) $L = \{0^n \mid n \text{ is perfect square}\}$</p> <p>(iii) $L = \{0^n \mid n \text{ is either prime number or odd number or both}\}$</p> <p>(iv) $L = \{0^n \mid n \text{ is not perfect cube}\}$</p> <p>(A) i only (B) ii and iv only</p> <p>(C) ii, iii and iv only (D) i and iv only</p> |
| <p>Q13.</p> | <p>[MSQ]</p> <p>Which of the following language is/are context free?</p> <p>(a) $L = \{a^n b^n : n \geq 0 \text{ and } n \text{ is neither prime nor composite}\}$</p> <p>(b) $L = \{a^p \mid p \text{ is prime number or } p = 0\}. \{a^p \mid p \text{ is not a prime number}\}$</p> <p>(c) $L = \{0^i 1^i 0^i 1^i \mid i \text{ is a prime number and } 0 < i < 100\}$</p> <p>(d) $L = \{(0^n 1^m)^k \mid m, n \geq 0\}.$</p> |

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| <p>Q14.</p> | <p>[MSQ]</p> <p>Which of the following language is/are context free?</p> <p>(A) $L1 = \{w \in \{a, b, c\} \mid n_a(w) = n_b(w) = n_c(w)\}$</p> <p>(B) $L2 = \{a^n b^m \mid n/m \text{ is integer ;where } n \text{ \& } m \text{ are prime numbers}\}$</p> <p>(C) $L3 = \{a^n b^j a^k b^l \mid n = k \text{ and } j = 2l\}$</p> <p>(D) $L4 = \{ww \mid w \in \{a, b\}^* \text{ and } n_a(w) < 2n_b(w) < 100\}$</p> |
| <p>Q15.</p> | <p>[MSQ]</p> <p>Which of the following language is/are context free?</p> <p>(A) $L1 = \{a^n b^j \mid n = j^2\}$</p> <p>(B) $L2 = \{a^n \mid n \text{ is multiple of a prime}\}$</p> <p>(C) $L3 = \{www^R w^R \mid w \in \{a\}^*\}$</p> <p>(D) $L4 = \{a^i a^n b^j b^n \mid n = j \text{ and } n, j \geq 0\}$</p> |
| <p>Q16.</p> | <p>[MSQ]</p> <p>Which of the following language is/are DCFL?</p> <p>(A) $L = \{wxcx^R \mid w \in \{a, b, c\}^* \text{ and } x \in \{a, b\}^*\}$</p> <p>(B) $L = \{xcx^R w \mid w \in \{a, b, c\}^* \text{ and } x \in \{a, b\}^*\}$</p> <p>(C) $L = \{xwx^R \mid w \in \{a, b, c\}^* \text{ and } x \in \{a, b\}^*\}$</p> <p>(D) $L = \{xcx^R w \mid w, x \in \{a, b\}^*\}$</p> |
| <p>Q17.</p> | <p>Which of the following language is/are DCFL?</p> <p>1. $L = \{www \mid w \in 0^*\}$</p> <p>2. $L = \{wxw \mid w \in 0^* 1^* \text{ and } x \in \{0, 1\}^*\}$</p> <p>3. $L = \{wxw^R \mid w \in 0^* 1^* \text{ and } x \in \{0, 1\}^*\}$</p> <p>(A) 1 only</p> <p>(B) 1 & 2 only</p> <p>(C) 2 & 3 only</p> <p>(D) All the above</p> |

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| Q18. | <p>Which of the following descriptions best fits the language</p> $L = \{x \in \{a, b, c\}^* \mid \#a(x) = \#b(x) + \#c(x)\}$ <p>(A) L is a regular language</p> <p>(B) L is a context-free language that is not regular</p> <p>(C) L is a context-sensitive language that is not context-free</p> <p>(D) L is not a context-sensitive language</p> |
| Q19. | <p>[MSQ]</p> <p>Let L_1 be a context-free language and $L_2 \subseteq L_1$ then which of the following statement is/are true?</p> <p>(A) L_2 is definitely context-free.</p> <p>(B) L_2 may not be a context-free language.</p> <p>(C) $L_1 \cup L_2$ is definitely context-free.</p> <p>(D) L_1^R is context sensitive.</p> |
| Q20. | <p>Let L language collection of strings over $\{0, 1\}$ containing ten more 0 than 1. These language is</p> <p>(A) Regular</p> <p>(B) Context free but not regular</p> <p>(C) Recursive but not contexts free</p> <p>(D) Recursively enumerable but not recursive</p> |
| Q21. | <p>Let L language collection of strings over $\{0, 1\}$ containing ten more 0 than 1 or containing ten more 1 than 0. These language is</p> <p>(A) Regular</p> <p>(B) Context free but not regular</p> <p>(C) Recursive but not context free</p> <p>(D) Recursively enumerable but not recursive</p> |

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| <p>Q22.</p> | <p>Which of the following languages is/are context free?</p> <ol style="list-style-type: none"> 1. $\{a^i \mid i \text{ is prime}\}$ 2. $\{(a^n b)^n \mid n \geq 1\}$ 3. $\{(a^n b)^m \mid m, n \geq 1\}$ <p>(A) Only 1&2 (B) Only 2&3 (C) Only 1&3 (D) None</p> |
| <p>Q23.</p> | <p>Consider the following language $L = \{a^n b^n a^n b^n \mid n \geq 0\}$. Which of the following statement is true about L?</p> <p>(A) Regular (B) Context-free but not regular, (C) Context-sensitive but not context-free, (D) Recursively enumerable but not context-sensitive.</p> |
| <p>Q24.</p> | <p>Which of the following languages is/ are not context free?</p> <ol style="list-style-type: none"> (i) $L = \{a^i b^j c^k d^l \mid i+j = k+l \text{ and } i, j, k, l \in \mathbb{N}\}$ (ii) $L = \{0^n \# 0^{2n} \# 0^{3n} \mid n \geq 0\}$ (iii) $L = \{uawb \mid u, w \in \{a,b\}^* \text{ and } u = w \}$ (iv) $L = \{w \in \{a, b\}^* : w \text{ has twice as many b's as a's}\}$ (v) $L = \{a^m b^{2n} c^{3n} d^p \mid p > m, \text{ and } m, n \geq 1\}$. <p>(A) iv and v only (B) ii only (C) ii and v only (D) i, iii, iv, v</p> |
| <p>Q25.</p> | <p>Which of the following languages is/are CFLs?</p> <p>$L_1 = \{0^i 1^j 2^k \mid i < j < k\}$.</p> <p>$L_2 = \{w \in (0+1+2)^* \mid w \text{ does not contain the same number of all 3 symbols}\}$.</p> <p>$L_3 = \{uawb \mid u, w \in (a+b)^*, u = w \}$</p> <p>$L_4 = \{b_i \# b_{i+1} \mid b_i \text{ is } i \text{ in binary } i \geq 1\}$. The alphabet here is $\{0,1,\#\}$.</p> <p>(A) L_1 and L_2 only (B) L_2, L_3 and L_4 only (C) L_3 only (D) L_3 and L_4 only</p> |

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| Q26. | <p>Which of the following languages is/are context free?</p> <p>$L1 = \{a^i b^k c^j \mid j = \max(i, k)\}$</p> <p>$L2 = \{x \in \{a, b\}^* : x \text{ is even and the first half of } x \text{ has one more 'a' than does the second half}\}$.</p> <p>$L3 = \{w \in \{a, b\}^* : \text{the first, middle, and last characters of } w \text{ are identical}\}$.</p> <p>$L4 = \{a^i b^j \mid i = kj \text{ for some positive integer } k\}$</p> <p>(A) L2 only (B) L3 only</p> <p>(C) L1 & L2 (D) L2 & L3</p> |
| Q27. | <p>Consider the following language $L = \{w^R \# z : w \text{ is a substring of } z, \text{ and } w, z \in \{a, b\}^*\}$. Which of the Following is true about L?</p> <p>(A) Regular</p> <p>(B) Context-free but not regular,</p> <p>(C) Context-sensitive but not context-free,</p> <p>(D) Recursively enumerable but not context-sensitive.</p> |
| Q28. | <p>Consider the following language $L = \{w \# z : w \text{ is a substring of } z \text{ and } w, z \in \{a, b\}^*\}$. Which of the Following is true about L?</p> <p>(A) Regular</p> <p>(B) Context-free but not regular</p> <p>(C) Context-sensitive but not context-free</p> <p>(d) Recursively enumerable but not context-sensitive.</p> |
| Q29. | <p>The language $L = \{0^i 1^j 2^i 3^j \mid i, j \geq 0\}$ is</p> <p>(A) Regular</p> <p>(B) Deterministic CFL</p> <p>(B) Non-Deterministic CFL</p> <p>(D) Not a CFL</p> |
| Q30. | <p>The language $L = \{0^i 1^j 2^i 3^j \mid i, j \geq 0\}$ is</p> <p>(A) Regular (B) Deterministic CFL</p> <p>(C) Non-Deterministic CFL (D) Not a CFL</p> |

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| Q31. | <p>The language $L = \{0^i1^j \mid i, j \geq 0 \text{ and } i = 3j\}$ is</p> <p>(A) Regular (B) Deterministic CFL</p> <p>(C) Non-Deterministic CFL (D) Not a CFL</p> |
| Q32. | <p>Which of the following language/s is/are not CFL?</p> <p>1. $L = \{a^n b^m c^m a^m b^n c^n \mid n \geq 0, m \geq 0\}$</p> <p>2. $L = \{a^n b^m c^n \mid n = 3m \text{ \& } n \geq 0, m \geq 0\}$</p> <p>3. $L = \{a^m b^n c^m \# w w^R \# w c w^R \mid w \in \{a, b\}^* \text{ and } w^R \text{ is reverse string of } w\}$</p> <p>(A) 1 only (B) 2 & 3 only</p> <p>(C) 1&2 only (D) none</p> |
| Q33. | <p>Let $L_1 = \{w : w \text{ is palindrome \& } w \in \{a, b\}^*\}$ and $L_2 = \{w w^R \# w : w \in \{a, b\}^*\}$, then.</p> <p>Which of the following statement is not true?</p> <p>(A) L_1 is CFL but L_2 is DCFL.</p> <p>(B) L_1 is DCFL but L_1 is not DCFL.</p> <p>(C) Both L_1 and L_2 are CFL.</p> <p>(D) L_1 is CFL but L_2 is not CFL.</p> |
| Q34. | <p>Which of the following languages on $\Sigma = \{a, b\}$ is/are context- free?</p> <p>(i) $L = \{a^i b^j \mid i \leq j^3\}$</p> <p>(ii) $L = \{a^i b^j \mid i + j \text{ is an even number}\}$</p> <p>(iii) $L = \{a^i b^j \mid i * j \text{ is composite number}\}$</p> <p>(iv) $L = \{a^i b^j \mid i + j \text{ is an odd number}\}$</p> <p>(A) Only i, iii (B) ii, iv only</p> <p>(C) Only i, ii and iv (D) all</p> |
| Q35. | <p>Which of the following languages on $\Sigma = \{a, b\}$ is/are not context- free?</p> <p>(i) $L = \{a^i b^j \mid i \text{ and } j \text{ are both prime}\}$</p> <p>(ii) $L = \{a^i b^j \mid i \text{ is prime or } j \text{ is prime}\}$</p> <p>(iii) $L = \{a^i b^j \mid i \text{ is prime but } j \text{ is not prime}\}$</p> <p>(iv) $L = \{a^i b^j \mid \text{neither } i \text{ is prime nor } j \text{ is prime}\}$</p> <p>(A) i, ii only (B) iii, iv only</p> <p>(C) ii, iv only (D) All the above</p> |

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| Q36. | <p>Which of the following language is/are DCFL?</p> <p>(A) $L = \{a^n b^m \mid n \neq m\}$</p> <p>(B) $L = \{w \# w^R \mid w^R \text{ is reverse of } w\}$</p> <p>(C) $L = \{a^n b^m \mid n \leq m \leq 3n\}$</p> <p>(D) both a & b</p> |
| Q37. | <p>Which of the following language is/are not CFL?</p> <p>1. $L = \{a^n b^m c^k \mid k = \text{HCF}(m, n)\}$</p> <p>2. $L = \{a^n b^m c^k \mid k = \text{LCM}(m, n)\}$</p> <p>3. $L = \{a^n b^m c^k \mid k = \max(m, n)\}$</p> <p>4. $L = \{a^n b^m c^k \mid k = \min(m, n)\}$</p> <p>(A) 1 & 2 only</p> <p>(B) 3 & 4 only</p> <p>(C) 2 & 4 only</p> <p>(D) All the above</p> |
| Q38. | <p>Which of the following language/s is/are CFL?</p> <p>1. If $L_1 = \{ww^R \mid w^R \text{ is reverse of } w\}$ then $L_1.L_1.L_1$ is CFL.</p> <p>2. If $L_2 = \{ww^R ww^R ww^R \mid w^R \text{ is reverse of } w\}$ then L_2 is CFL.</p> <p>3. If $L_3 = \{ww^R \mid w^R \text{ is reverse of } w\}$, then complement of L_3 is NCFL.</p> <p>(A) 1 only</p> <p>(B) 1 & 3 only</p> <p>(C) 2 & 3 only</p> <p>(D) All the above</p> |
| Q39. | <p>Which of the following language is/are DCFL?</p> <p>1. $L = \{uv : u \neq v^R \text{ where } u, v \in \Sigma^* \text{ and } v^R \text{ is reverse of } v\}$.</p> <p>2. $L = \{uv : u \text{ contains even number of } 0 \text{ and } v \text{ contains odd number of } 1;$ Where $u, v \in \{0, 1\}^*\}$.</p> <p>3. $L = \{uv : u \in L_1 \& v \notin L_1; \text{ where } L_1 \text{ is DCFL}\}$</p> <p>4. 1. $L = \{u \# v : u \neq v^R \text{ where } u, v \in \Sigma^* \text{ and } v^R \text{ is reverse of } v\}$.</p> <p>(A) 1, 3 only</p> <p>(B) 1, 3, 4 only</p> <p>(C) 2, 3, 4 only</p> <p>(D) All the above</p> |

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|--|--|------------------------------------|-------------------------|--|--|----------------------------------|--------------------|------------------------------------|----------------------|---|---|--------------------------|----------------------|---|---|--|-------------------------|
| Q40. | <p>Which of the following languages are NOT context-free?</p> <p>$L_1 = \{0^{n+m}1^n0^m \mid n, m \geq 0\}$,</p> <p>$L_2 = \{0^{n+m}1^{n+m}0^m \mid n, m \geq 0\}$, and</p> <p>$L_3 = \{0^{n+m}1^n0^{n+m} \mid n, m \geq 0\}$</p> <p>(A) L_1 only</p> <p>(B) L_3 only</p> <p>(C) L_1 and L_2</p> <p>(D) L_2 and L_3</p> | | | | | | | | | | | | | | | | |
| Q41. | <p>A context-free grammar G is ambiguous if and only if</p> <p>(A) Some string $w \in L(G)$ has at least two different derivations.</p> <p>(B) Some string $w \in L(G)$ has at least two different parse trees.</p> <p>(C) Every string $w \in L(G)$ has at least two different parse trees.</p> <p>(D) Every string $w \in L(G)$ has at least two different derivations.</p> | | | | | | | | | | | | | | | | |
| Q42. | <p>Consider the following grammar: (Assume starting symbol is S)</p> <table><tr><td>G_1</td><td>G_2</td></tr><tr><td>$S \rightarrow aA \mid bB \mid cS \mid aS$</td><td>$S \rightarrow aS \mid bSa \mid aSb \mid bS$</td></tr><tr><td>$A \rightarrow bAa \mid \lambda$</td><td>$\mid \lambda$</td></tr><tr><td>$B \rightarrow bB \mid a$</td><td></td></tr></table> <p>Which of the following grammar is/are ambiguous?</p> <p>(a) G_1 only</p> <p>(b) G_2 only</p> <p>(c) Both G_1 & G_2</p> <p>(d) Neither G_1 nor G_2</p> | G_1 | G_2 | $S \rightarrow aA \mid bB \mid cS \mid aS$ | $S \rightarrow aS \mid bSa \mid aSb \mid bS$ | $A \rightarrow bAa \mid \lambda$ | $\mid \lambda$ | $B \rightarrow bB \mid a$ | | | | | | | | | |
| G_1 | G_2 | | | | | | | | | | | | | | | | |
| $S \rightarrow aA \mid bB \mid cS \mid aS$ | $S \rightarrow aS \mid bSa \mid aSb \mid bS$ | | | | | | | | | | | | | | | | |
| $A \rightarrow bAa \mid \lambda$ | $\mid \lambda$ | | | | | | | | | | | | | | | | |
| $B \rightarrow bB \mid a$ | | | | | | | | | | | | | | | | | |
| Q43. | <p>[MSQ]</p> <p>Which of the following grammars have more than one parse tree for the string '1001'?</p> <table><tr><td>(a)</td><td>(b)</td><td>(c)</td><td>(d)</td></tr><tr><td>$S \rightarrow XY$</td><td>$S \rightarrow XY$</td><td>$S \rightarrow 0S \mid 1S \mid SS$</td><td>$S \rightarrow 1S0S$</td></tr><tr><td>$X \rightarrow 0X \mid 1X \mid \lambda$</td><td>$X \rightarrow 0X1 \mid 1X0 \mid \lambda$</td><td>$S \rightarrow 0 \mid 1$</td><td>$S \rightarrow 0S1S$</td></tr><tr><td>$Y \rightarrow 1Y1 \mid 0Y0 \mid \lambda$</td><td>$Y \rightarrow 1Y \mid 0Y \mid \lambda$</td><td></td><td>$S \rightarrow \lambda$</td></tr></table> | (a) | (b) | (c) | (d) | $S \rightarrow XY$ | $S \rightarrow XY$ | $S \rightarrow 0S \mid 1S \mid SS$ | $S \rightarrow 1S0S$ | $X \rightarrow 0X \mid 1X \mid \lambda$ | $X \rightarrow 0X1 \mid 1X0 \mid \lambda$ | $S \rightarrow 0 \mid 1$ | $S \rightarrow 0S1S$ | $Y \rightarrow 1Y1 \mid 0Y0 \mid \lambda$ | $Y \rightarrow 1Y \mid 0Y \mid \lambda$ | | $S \rightarrow \lambda$ |
| (a) | (b) | (c) | (d) | | | | | | | | | | | | | | |
| $S \rightarrow XY$ | $S \rightarrow XY$ | $S \rightarrow 0S \mid 1S \mid SS$ | $S \rightarrow 1S0S$ | | | | | | | | | | | | | | |
| $X \rightarrow 0X \mid 1X \mid \lambda$ | $X \rightarrow 0X1 \mid 1X0 \mid \lambda$ | $S \rightarrow 0 \mid 1$ | $S \rightarrow 0S1S$ | | | | | | | | | | | | | | |
| $Y \rightarrow 1Y1 \mid 0Y0 \mid \lambda$ | $Y \rightarrow 1Y \mid 0Y \mid \lambda$ | | $S \rightarrow \lambda$ | | | | | | | | | | | | | | |

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| Q44. | [MSQ] |
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Consider the following grammar G with starting symbol S:

$$S \rightarrow aSb \mid bSa \mid aSa \mid bSb \mid AB \mid BA$$
$$A \rightarrow aA \mid \lambda$$
$$B \rightarrow bB \mid \lambda$$

For which of the following string(s), G has more than one parse tree?

- (A) ab
(B) aabb
(C) abba
(D) abab

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| Q45. | Consider the following grammar G with starting symbol S: |
|------|--|

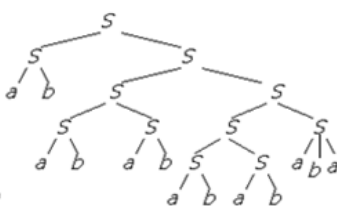
$$S \rightarrow aSb \mid bSa \mid aSa \mid bSb \mid AB \mid BA$$
$$A \rightarrow aA \mid \lambda$$
$$B \rightarrow bB \mid \lambda$$

How many different parse trees are possible for string 'abba'? _____

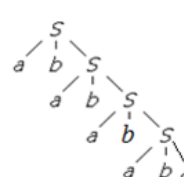
| | |
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| Q46. | Which of the parse trees below yield the same word? |
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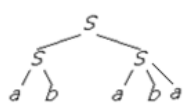
(*i*)



(ii)



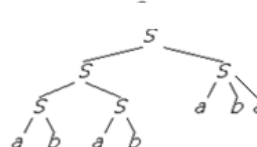
(iii)



(iv)



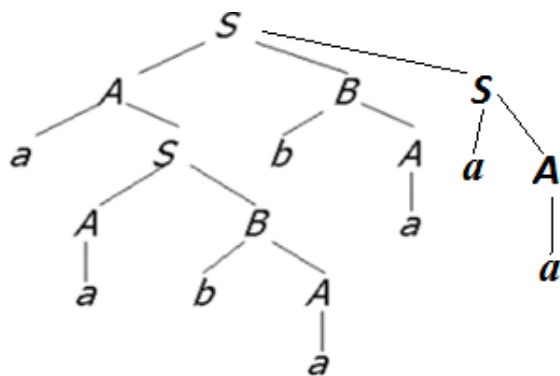
(v)



(vi)

- (A) iii and iv only
(B) i, ii and v only
(C) ii and v only
(D) iii, iv and vi only

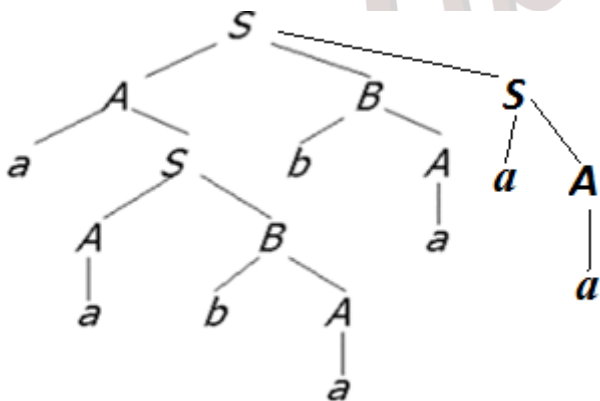
Q47. Here is a parse tree that is derived from some unknown grammar G .



Which of the following productions is not surely for grammar G ?

- (A) $S \rightarrow ABS$
- (B) $A \rightarrow aB$
- (C) $S \rightarrow AB$
- (D) $S \rightarrow aA$

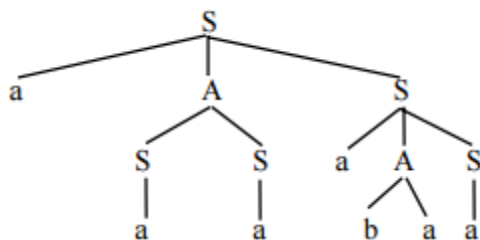
Q48. The parse tree below represents a rightmost derivation according to the grammar $S \rightarrow ABS \mid AB \mid aA$, $A \rightarrow aS \mid a$, $B \rightarrow bA$



Which of the following is a right-sentential form in this derivation?

- (I) $ABaA$
- (II) $bAba$
- (III) $aaBB$
- (IV) $aabAba$
- (A) I & II only
- (B) III & IV only
- (C) I only
- (D) All except III

Q49. Consider the following context free grammar: $G = (\{S, A, a, b\}, \{a, b\}, R, S)$, where $R = \{S \rightarrow aAS \ S \rightarrow a \ A \rightarrow SbA \ A \rightarrow SS \ A \rightarrow ba\}$ and parse tree for the string “aaaabaa” is shown below:



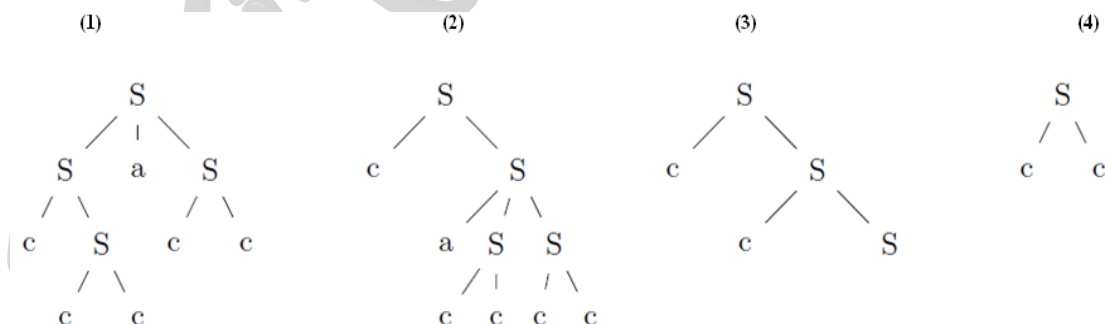
Which of the following are left-sentential forms corresponding to this derivation?

- (i) aAS
- (ii) aSS
- (iii) aaSS
- (iv) aaS
- (v) aaaaAS
- (A) i, iii and v only
- (B) ii and iv only
- (C) ii, iii and v only
- (D) All the above

Q50. Consider the grammar (with start variable S and terminals a and c).

$S \rightarrow SaS \mid cS \mid cc$

Which of following parse trees matches the grammar?



- (A) i ,iii,iv
- (B) iii,iv
- (C) ii
- (D) i and iv

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| <p>Q51.</p> | <p>Consider the following CFG</p> $S \rightarrow aB \mid bA$ $B \rightarrow b \mid bS \mid aBB$ $A \rightarrow a \mid aS \mid bAA$ <p>Now consider the following derivation of above CFG</p> $S \Rightarrow aB \Rightarrow aaBB \Rightarrow aaBb \Rightarrow aabSb \Rightarrow aabbAb \Rightarrow aabbab$ <p>This derivation is</p> <p>(A) a leftmost derivation</p> <p>(B) a rightmost derivation</p> <p>(C) both leftmost and rightmost derivation</p> <p>(D) neither leftmost nor rightmost derivation</p> |
| <p>Q52.</p> | <p>Given the simple CFG with non-terminals $\{S, A, B\}$, terminals $\{a, b\}$, and productions $S \rightarrow AB, A \rightarrow AB \mid a, B \rightarrow BA \mid b$.</p> <p>Which of the following derivation is NOT a derivation of the string abab?</p> <p>(A) $S \Rightarrow AB \Rightarrow Ab \Rightarrow ABb \Rightarrow ABAb \Rightarrow AbAb \Rightarrow Abab \Rightarrow abab$</p> <p>(B) $S \Rightarrow AB \Rightarrow ABB \Rightarrow ABAB \Rightarrow aBaB \Rightarrow abab$</p> <p>(C) $S \Rightarrow AB \Rightarrow ABA \Rightarrow ABAB \Rightarrow ABAb \Rightarrow AbAb \Rightarrow Abab \Rightarrow abab$</p> <p>(D) $S \Rightarrow AB \Rightarrow aB \Rightarrow aBA \Rightarrow abA \Rightarrow abAB \Rightarrow abaB \Rightarrow abab$</p> |
| <p>Q53.</p> | <p>Consider the following grammar G with start symbol S:</p> $S \rightarrow \text{if id then } S \text{ else } S$ $S \rightarrow \text{if id then } S$ $S \rightarrow \text{id}$ <p>Which of following is/are true?</p> <p>(i) G is Type 3 grammar.</p> <p>(ii) G is Type 2 grammar</p> <p>(iii) G is unambiguous grammar.</p> <p>(iv) G is ambiguous grammar.</p> <p>(A) i, ii and iv only</p> <p>(b) ii and iv only</p> <p>(C) ii and iii only</p> <p>(D) i, ii and iii only</p> |

For the next two questions:

Consider the following grammar G

$S \rightarrow 0A \mid 1B \mid \lambda$

$A \rightarrow 0AA \mid 1S \mid 1$

$B \rightarrow 1BB \mid 0S \mid 0$

Let $L(G)$ is the language accepted by G.

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| <p>Q54.</p> | <p>Which of the following statements is/are true about G?</p> <p>(i) G is Type 2.</p> <p>(ii) G is ambiguous.</p> <p>(iii) $0011010110 \notin L(G)$ and $1111100101000 \in L(G)$.</p> <p>(A) i & ii only</p> <p>(B) ii & iii only</p> <p>(C) i only</p> <p>(D) All are true</p> |
| <p>Q55.</p> | <p>The language generated by G is</p> <p>(A) The set of all the strings that don't contain substring 00.</p> <p>(B) The set of all the strings that contains equal numbers of zero's followed by equal number of one's or equal numbers of one's followed by equal number of zero's i.e. $L(G) = \{0^n 1^n \mid n > 0\} \cup \{0^n 1^n \mid n > 0\}$</p> <p>(C) The set of all the strings which have equal number of zero's and one's.</p> <p>(D) None of these</p> |
| <p>Q56.</p> | <p>Consider the following statements about the context free grammar</p> <p>$G = \{S \rightarrow SS, S \rightarrow ab, S \rightarrow ba, S \rightarrow \epsilon\}$</p> <p>Which of the following statements is/are true about G</p> <p>I. G is ambiguous.</p> <p>II. G produces all strings with equal number of a's and b's.</p> <p>III. $L(G)$ can be accepted by a deterministic PDA.</p> <p>(A) I only</p> <p>(B) I and III only</p> <p>(C) II and III only</p> <p>(D) I, II and III</p> |

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| <p>Q57.</p> | <p>Consider the grammar G given by the productions</p> $S \rightarrow aSa \mid aBa \mid b \mid \lambda$ $B \rightarrow bB \mid b$ <p>Then which of the following is/are true about G?</p> <p>(i) G is context sensitive grammar.</p> <p>(ii) G accept the language L where $L = \{a^i b^j a^k \mid i = k \text{ and } i, j, k \geq 0\}$.</p> <p>(iii) G is not unambiguous.</p> <p>(A) ii only</p> <p>(B) ii & iii only</p> <p>(C) i & ii only</p> <p>(D) All are true</p> |
| <p>Q58.</p> | <p>Consider the following grammar G</p> $S \rightarrow TC \mid AR$ $T \rightarrow aTb \mid \lambda$ $C \rightarrow Cc \mid \lambda$ $R \rightarrow bRc \mid C$ $A \rightarrow Aa \mid \lambda$ <p>The language is generated by G is : (Assume $i, j, k \geq 0$)</p> <p>(A) $L = \{a^i b^j c^k \mid k = j \text{ and } i \geq j\}$</p> <p>(B) $L = \{a^i b^j c^k \mid i = j \text{ and } k \geq j\}$</p> <p>(C) $L = \{a^i b^j c^k \mid i = j \text{ or } k \geq 0\}$</p> <p>(D) $L = \{a^i b^j c^k \mid i = j \text{ or } k \geq j\}$</p> |
| <p>Q59.</p> | <p>Consider three language L1, L2, L3 and three grammar G1, G2, G3:</p> $L1 = \{a^i b^j c^k \mid k = i + 2*j\}$ $L2 = \{a^i b^j c^k \mid k \geq 3\},$ $L3 = \{a^i b^j c^k \mid k = i + j\}$ <div style="display: flex; justify-content: space-around; margin-top: 20px;"> <div style="text-align: center;"> <p>G1:</p> $S \rightarrow aSc \mid T$ $T \rightarrow bTc \mid \lambda$ </div> <div style="text-align: center;"> <p>G2:</p> $S \rightarrow aSc \mid T$ $T \rightarrow bTcc \mid \lambda$ </div> <div style="text-align: center;"> <p>G3:</p> $S \rightarrow TC$ $T \rightarrow aTb \mid \lambda$ $C \rightarrow Cc \mid ccc$ </div> </div> <p>Which of the following language, grammar pair (L, G) are equivalent?</p> <p>(a) (L1, G1), (L2, G2), (L3, G3)</p> <p>(b) (L1, G3), (L2, G2), (L3, G1)</p> <p>(c) (L1, G2), (L2, G2), (L3, G3)</p> <p>(d) (L1, G2), (L2, G3), (L3, G1)</p> |

| Q60. | Match the column which CFG corresponds to which language? | | | | | | | | |
|--|--|----------|---------|--|---|--|---|--|---|
| | <table border="1"> <thead> <tr> <th data-bbox="225 152 890 219">Language</th><th data-bbox="890 152 1425 219">Grammar</th></tr> </thead> <tbody> <tr> <td data-bbox="225 219 890 409">1. $\{0^n 1^n \mid n > 0\} \cup \{0^n 1^{2n} \mid n > 0\}$</td><td data-bbox="890 219 1425 409"> i. $S \rightarrow 0A1 \mid 0B11$ $A \rightarrow 0A1 \mid \lambda$ $B \rightarrow 0B11 \mid \lambda$ </td></tr> <tr> <td data-bbox="225 409 890 539">2. Binary strings with twice as many 1s as 0s.</td><td data-bbox="890 409 1425 539"> ii. $S \rightarrow 0A0 \mid 1A1$ $A \rightarrow 0A \mid 1A \mid \lambda$ </td></tr> <tr> <td data-bbox="225 539 890 683">3. $\{w \in \{0,1\}^* \mid w \text{ starts and ends with the same symbol}\}$</td><td data-bbox="890 539 1425 683"> iii. $S \rightarrow 0A11 \mid \lambda$ $A \rightarrow 0A11 \mid \lambda$ </td></tr> </tbody> </table> | Language | Grammar | 1. $\{0^n 1^n \mid n > 0\} \cup \{0^n 1^{2n} \mid n > 0\}$ | i. $S \rightarrow 0A1 \mid 0B11$ $A \rightarrow 0A1 \mid \lambda$ $B \rightarrow 0B11 \mid \lambda$ | 2. Binary strings with twice as many 1s as 0s. | ii. $S \rightarrow 0A0 \mid 1A1$ $A \rightarrow 0A \mid 1A \mid \lambda$ | 3. $\{w \in \{0,1\}^* \mid w \text{ starts and ends with the same symbol}\}$ | iii. $S \rightarrow 0A11 \mid \lambda$ $A \rightarrow 0A11 \mid \lambda$ |
| Language | Grammar | | | | | | | | |
| 1. $\{0^n 1^n \mid n > 0\} \cup \{0^n 1^{2n} \mid n > 0\}$ | i. $S \rightarrow 0A1 \mid 0B11$ $A \rightarrow 0A1 \mid \lambda$ $B \rightarrow 0B11 \mid \lambda$ | | | | | | | | |
| 2. Binary strings with twice as many 1s as 0s. | ii. $S \rightarrow 0A0 \mid 1A1$ $A \rightarrow 0A \mid 1A \mid \lambda$ | | | | | | | | |
| 3. $\{w \in \{0,1\}^* \mid w \text{ starts and ends with the same symbol}\}$ | iii. $S \rightarrow 0A11 \mid \lambda$ $A \rightarrow 0A11 \mid \lambda$ | | | | | | | | |
| | <p>(A) 1-i, 2-iii, 3-ii</p> <p>(B) 1-ii, 2-i, 3-iii</p> <p>(C) 1-ii, 2-iii, 3-I</p> <p>(D) 1-i, 2-ii, 3-iii</p> | | | | | | | | |
| Q61. | <p>Let $V = \{S\}$, $\Sigma = \{a, b\}$. Which of the following grammars is such that the language Generated by it is $\{ww^r : w \in \{a, b\}^*\}$.</p> <p>(A) $G_1 = (V, \Sigma, R_1, S)$, where $R_1 = \{(S \rightarrow aSa), (S \rightarrow bSb), (S \rightarrow \lambda)\}$</p> <p>(B) $G_2 = (V, \Sigma, R_2, S)$, where $R_2 = \{(S \rightarrow aSa), (S \rightarrow bSb), (S \rightarrow a), (S \rightarrow b), (S \rightarrow \lambda)\}$</p> <p>(C) $G_3 = (V, \Sigma, R_3, S)$, where $R_3 = \{(S \rightarrow Sa), (S \rightarrow Sb), (S \rightarrow \lambda)\}$</p> <p>(D) $G_4 = (V, \Sigma, R_3, S)$, where $R_3 = \{(S \rightarrow aSb), (S \rightarrow \lambda)\}$</p> | | | | | | | | |
| Q62. | <p>Which of the following CFG generates all the strings contains more 1's than 0's?</p> <p>(A) $S \rightarrow 0T \mid 11T \quad T \rightarrow 0S \mid 11S \mid \lambda$</p> <p>(B) $S \rightarrow 0S1 \mid 1S0 \mid 1S1 \mid 1$</p> <p>(C) $S \rightarrow TS \mid 1T \mid 1S \quad T \rightarrow TT \mid 0T1 \mid 1T0 \mid \lambda$</p> <p>(D) $S \rightarrow TS \mid 1T \mid 11 \quad T \rightarrow TT \mid 0T1 \mid 1T0 \mid 1$</p> | | | | | | | | |

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| Q63. | Consider the following Pushdown Automaton (PDA) P1: |
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$P1 = (Q = \{q0, q1\}, \Sigma = \{a, b, c\}, \Gamma = \{0, 1, \#\}, \delta, q0, Z0 = \#, F = \{q1\})$ where the transition function δ is given by

$$\delta(q_0, a, \#) = \{(q_0, 0\#), (q_0, 11\#)\}$$
$$\delta(q_0, a, 0) = \{(q_0, 00), (q_0, 110)\}$$
$$\delta(q_0, a, 1) = \{(q_0, 01), (q_0, 111)\}$$
$$\delta(q_0, b, 0) = \{(q_0, \lambda)\}$$
$$\delta(q_0, c, 1) = \{(q_0, \lambda)\}$$
$$\delta(q_0, \lambda, \#) = \{(q_1, \#)\}$$

$\delta(q, w, z) = \emptyset$ everywhere else Acceptance is by final state.

Which of the following words are accepted by the PDA P1?

(i) ab

(ii) aababcc

(iii) ac

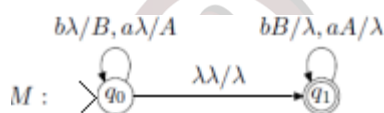
(A)i only

(B) i and ii only

(C)i and iii only

(D)iii only

| | |
|------|--------------------|
| Q64. | Let M be the PDA |
|------|--------------------|



Which of the following strings are accepted by M?

(A)aaaa

(B) baab

(C) aaa

(D) ab

| | |
|------|--|
| Q65. | Consider the pushdown automaton with the following transition rules: |
|------|--|

$$\delta(q, 0, Z_0) = \{(q, XZ_0)\}$$
$$\delta(q, 0, X) = \{(q, XX)\}$$
$$\delta(q, 1, X) = \{(q, X)\}$$
$$\delta(q, \varepsilon, X) = \{(p, \varepsilon)\}$$
$$\delta(p, \varepsilon, X) = \{(p, \varepsilon)\}$$
$$\delta(p, 1, X) = \{(p, XX)\}$$
$$\delta(p, 1, Z_0) = \{(p, \varepsilon)\}$$

The start state is q . For which of the following inputs can the PDA first enter state p with the input empty and the stack containing XXZ_0 [i.e., the ID (p, ϵ, XXZ_0)]?

(A) 001111

(B) 0101010

(C) 001110

(D)111001

Q66. Here are the transitions of a deterministic pushdown automaton. The start state is q_0 , and f is the accepting state.

| State-Symbol | a | b | ϵ |
|--------------|------------------|------------------|------------------|
| q_0-Z_0 | (q_1,AAZ_0) | (q_2,BZ_0) | (f,ϵ) |
| q_1-A | (q_1,AAA) | (q_1,ϵ) | - |
| q_1-Z_0 | - | - | (q_0,Z_0) |
| q_2-B | (q_3,ϵ) | (q_2,BB) | - |
| q_2-Z_0 | - | - | (q_0,Z_0) |
| q_3-B | - | - | (q_2,ϵ) |
| q_3-Z_0 | - | - | (q_1,AZ_0) |

Identify below the one input string that the PDA accepts.

- (A) bbbab
- (B) abbbab
- (C) abbbabb
- (D) bbaabab

Q67. Consider the pushdown automaton with the following transition rules:

$$\delta(q,0,Z_0) = \{(q,XZ_0)\}; \delta(q,0,X) = \{(q,XX)\}; \delta(q,1,X) = \{(q,X)\}$$

$$\delta(q,\epsilon,X) = \{(p,\epsilon)\}$$

$$\delta(p,\epsilon,X) = \{(p,\epsilon)\}$$

$$\delta(p,1,X) = \{(p,XX)\}$$

$$\delta(p,1,Z_0) = \{(p,\epsilon)\}$$

From the $(p, 1101, XXZ_0)$, which of the following states can NOT be reached?

- (A) $(p,101,XXZ_0)$
- (B) $(p,1101,XZ_0)$
- (C) (p,ϵ,XZ_0)
- (D) $(p,101,XXXZ_0)$

PDA to Language

Q68.

Let M be the PDA defined by

$$Q = \{q_0, q_1, q_2\} \quad \Sigma = \{a, b\} \quad \Gamma = \{A\} \quad F = \{q_1, q_2\}$$

$$\delta(q_0, a, \lambda) = \{[q_0, A]\}$$

$$\delta(q_0, \lambda, \lambda) = \{[q_1, \lambda]\}$$

$$\delta(q_0, b, A) = \{[q_2, \lambda]\}$$

$$\delta(q_1, \lambda, A) = \{[q_1, \lambda]\}$$

$$\delta(q_2, b, A) = \{[q_2, \lambda]\}$$

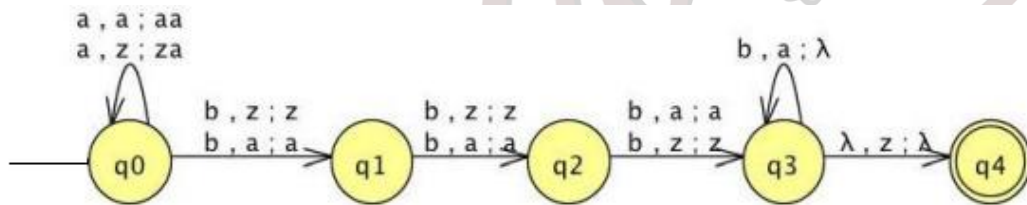
$$\delta(q_2, \lambda, A) = \{[q_2, \lambda]\}$$

The language accepted by M is?

- (A) The PDA M accepts the language $\{a^i b^j \mid 0 \leq i, j\}$
- (B) The PDA M accepts the language $\{a^i b^j \mid 0 \leq i \leq j\}$
- (C) The PDA M accepts the language $\{a^i b^j \mid 0 \leq j \leq i\}$
- (D) The PDA M accepts the language $\{a^i b^j \mid i = j\}$

Q69.

The language accepted by following PDA is



- (A) $L1 = \{a^n b^{3n} \mid n \geq 0\}$
- (B) $L2 = \{a^n b^n \mid n \geq 0\}$
- (C) $L3 = \{a^n b^{n+3} \mid n \geq 0\}$
- (D) $L4 = \{a^n b^m \mid n \geq 0, m > n + 3\}$

Q70.

Let M be the PDA defined by

$$Q = \{q_0, q_1, q_2\} \quad \Sigma = \{a, b\} \quad \Gamma = \{A\} \quad F = \{q_1, q_2\}$$

$$\delta(q_0, a, \lambda) = \{[q_0, A]\}$$

$$\delta(q_0, \lambda, \lambda) = \{[q_1, \lambda]\}$$

$$\delta(q_0, b, A) = \{[q_2, \lambda]\}$$

$$\delta(q_1, \lambda, A) = \{[q_1, \lambda]\}$$

$$\delta(q_2, b, A) = \{[q_2, \lambda]\}$$

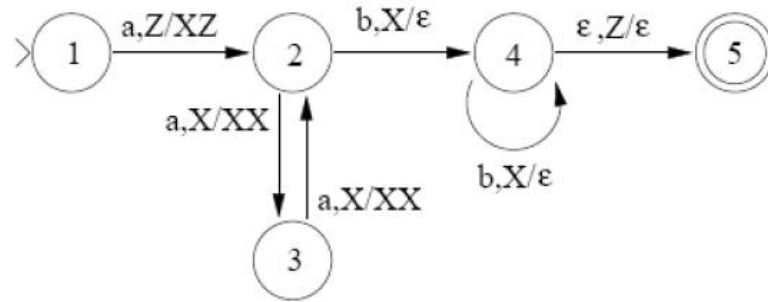
$$\delta(q_2, \lambda, A) = \{[q_2, \lambda]\}$$

The language accepted by M is:

- (A) $\{a^i b^j \mid 0 \leq j \leq i\}$
- (B) $\{a^i b^j \mid 0 \leq i \leq j\}$
- (C) $\{a^n b^n \mid n \geq 0\}$
- (D) $\{(ab)^n \mid n \geq 0\}$

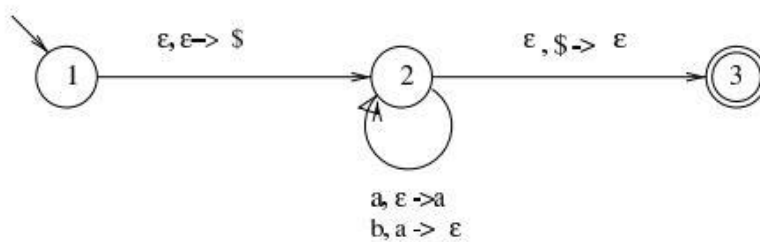
| | |
|--------------------|---|
| <p>Q71.</p> | <p>The language accepted by following PDA is:</p> <div style="text-align: center;"> </div> <p>(A) $\{a^i b^j \mid 0 \leq j \leq i\}$ (B) $\{a^i b^i \mid 0 \leq i \leq j\}$ (C) $\{a^n b^n \mid n \geq 0\}$ (D) $\{(ab)^n \mid n \geq 0\}$</p> |
| <p>Q72.</p> | <p>The language accepted by following PDA</p> <div style="text-align: center;"> </div> <p>(A) $\{a^i b^j c^k \mid i + k = j\}$ (B) $\{a^i b^j c^k \mid i < j = k\}$ (C) $\{a^i b^j c^i \mid i = j\}$ (D) $\{a^i b^j c^k \mid i < j \text{ and } k > 0\}$</p> |
| <p>Q73.</p> | <p>The language accepted by following PDA is:</p> <div style="text-align: center;"> </div> <p>(A) $\{a^i c^j b^i \mid i, j \geq 0\}$ (B) $\{a^i c^j b^k \mid i, j, k \geq 0\}$ (C) $\{a^i c^j b^i \mid i, j > 0\}$ (D) $\{a^i b^i \mid i \geq 0\}$</p> |

Q74. The language accepted by following PDA is



- (A) $L(G) = \{a^{2n}b^{2n}, n \geq 1\}$
- (B) $L(G) = \{a^{2n+1}b^{2n+1}, n \geq 1\}$
- (C) $L(G) = \{a^{2n}b^{2n}, n \geq 0\}$
- (D) $L(G) = \{a^{2n+1}b^{2n+1}, n \geq 0\}$

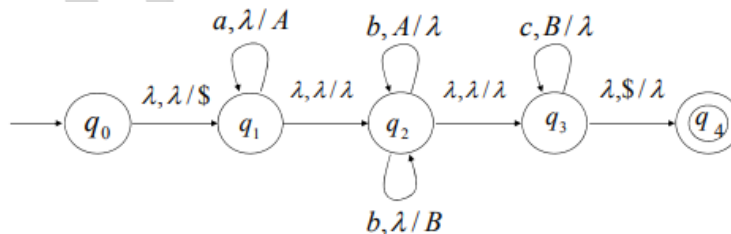
Q75. Consider the following PDA



Which of the following grammar is equivalent to given PDA?

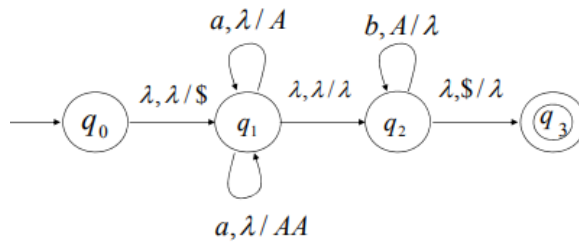
- (A) $S \rightarrow SS \mid aSb \mid ab$
- (B) $S \rightarrow SS \mid ab \mid \lambda$
- (C) $S \rightarrow SS \mid aSb \mid \lambda$
- (D) $S \rightarrow aSb \mid \lambda$

Q76. The language accepted by following automata is



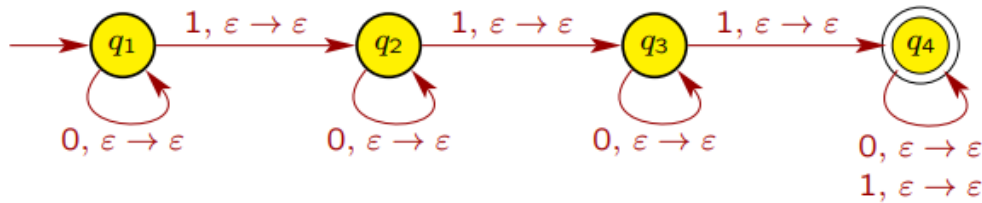
- (A) $L = \{a^i b^j c^k \mid i = j \text{ or } j = k\}$
- (B) $L = \{a^i b^j c^k \mid i = j \text{ and } j = k\}$
- (C) $L = \{a^i b^j c^k \mid i = j + k\}$
- (D) $L = \{a^i b^j c^k \mid j = i + k\}$

Q77. The language accepted by following automata is



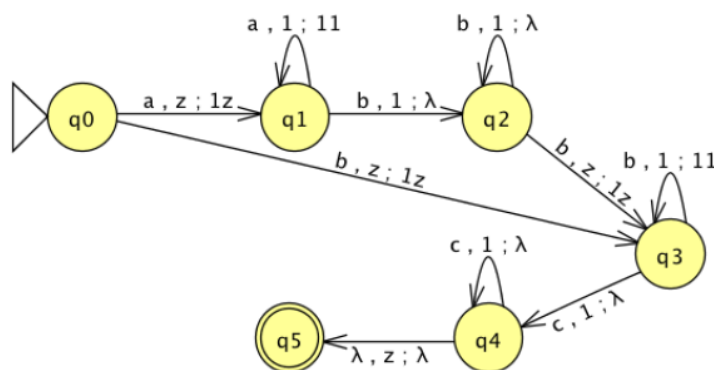
- (A) $L = \{a^i b^j \mid i = j\}$
- (B) $L = \{a^i b^j \mid i = 2*j\}$
- (C) $L = \{a^i b^j \mid i = j \text{ and } i = 2*j\}$
- (D) $L = \{a^i b^j \mid i \leq j \leq 2i\}$

Q78. The language accepted by following automata is



- (A) set of all strings over $\{0, 1\}$ which contain at most three 1's
- (B) set of all strings over $\{0, 1\}$ which contain exactly three 1's
- (C) set of all strings over $\{0, 1\}$ which contain at least three 1's
- (D) set of all strings over $\{0, 1\}$

Q79. Consider the following PDA



The language accepted by following PDA is?

- (A) $L = \{a^n b^n c^m : n \geq 0, m \geq 0\}$
- (B) $L = \{a^n b^m c^m : n \geq 0, m \geq 0\}$
- (C) $L = \{a^n b^{m+n} c^m : n \geq 0, m \geq 0\}$
- (D) $L = \{a^n b^n c^m : n \geq 0, m \geq 0\}$

Q80. The language is accepted by the following PDA is

$M = (q_0, q_1, \dots, q_5, \{a, b\}, \{0, 1, a, z\}, \delta, z, q_0, \{q_5\})$,

$\delta(q_0, b, z) = \{(q_1, 1z)\}$,

$\delta(q_1, b, 1) = \{(q_1, 11)\}$,

$\delta(q_2, a, 1) = \{(q_3, \lambda)\}$,

$\delta(q_3, a, 1) = \{(q_4, \lambda)\}$

$\delta(q_4, a, z) = \{(q_4, z), (q_5, z)\}$

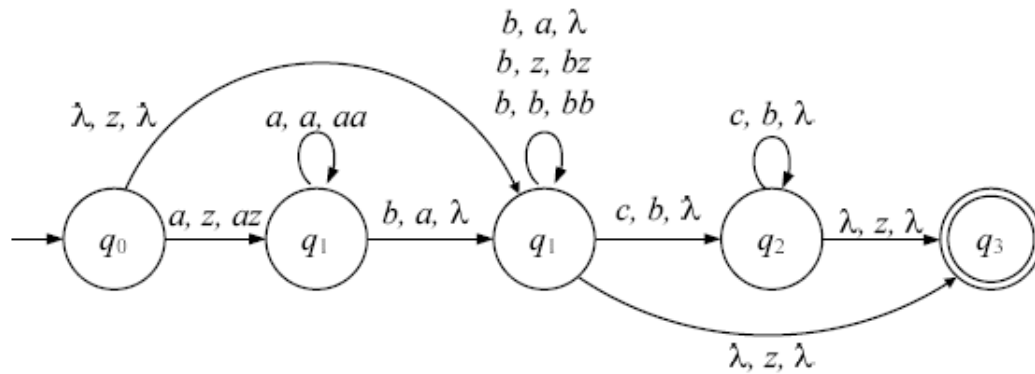
(A) $\{b^i a^j \mid i \geq 1, j \geq 0\}$

(B) $\{b^i a^j \mid i \geq 1, j \geq 1\}$

(C) $\{\lambda\}$

(D) \emptyset

Q81. The language accepted by following PDA is



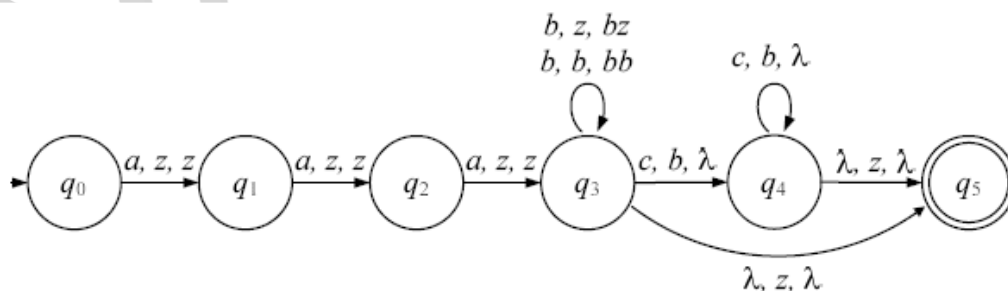
(A) $L = \{a^n b^m c^{n+m} \mid n \geq 0, m \geq 0\}$.

(B) $L = \{a^n b^{n+m} c^m \mid n \geq 0, m \geq 0\}$

(C) $L = \{a^n b^m c^m \mid n \geq 0, m \geq 0\}$

(D) $L = \{a^n b^m c^n \mid n \geq 0, m \geq 1\}$

Q82. The language accepted by following PDA is



(A) $L = \{a^n b^m c^m \mid n \geq 0, m \geq 0\}$

(B) $L = \{a^3 b^m c^m \mid m \geq 0\}$

(C) $L = \{a^{3n} b^m c^m \mid n \geq 0, m \geq 0\}$

(D) None

For next two Question: Push down automata for

$P = \{Q = \{q_0, q_1\}, \Sigma = \{[,]\}, \Gamma = \{A, B\}, q_0, F = \{q_1\}, \delta\}$

$\delta(q_0, [, \epsilon) = \{(q_0, A)\}$

$\delta(q_0,], \epsilon) = \{(q_0, B)\}$

$\delta(q_0, \epsilon, \epsilon) = \{(q_1, \epsilon)\}$

$\delta(q_1, [, A) = \{(q_1, \epsilon)\}$

$\delta(q_1,], B) = \{(q_1, \epsilon)\}$

$\delta(q, a, x) = \emptyset$ otherwise

Q83. Suppose the current instantaneous description of P is $\langle q_1, AAAAAA \rangle$ and the unread portion of the input is $[[]]$. The instantaneous description after one step is

(A) The machine crashes

(B) $\langle q_1, AAAAAA \rangle$

(C) $\langle q_1, AAAA \rangle$

(D) $\langle q_1, AAA \rangle$

Q84. The language recognized by PDA is

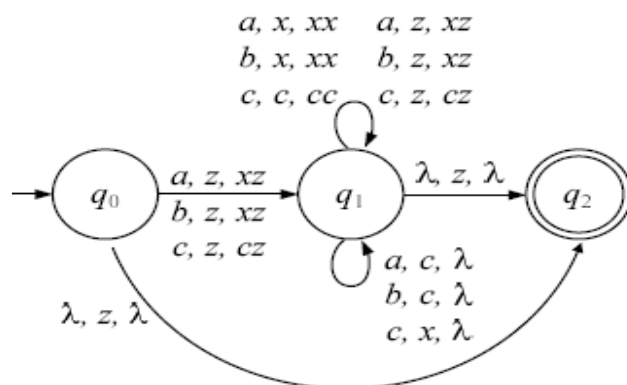
(A) $\{[{}^n n \mid n \geq 0\}$

(B) $\{[{}^i]^j : i \geq j \geq 0\}$

(C) $\{wu \mid w, u \in \{[,]\}^* \text{ and } u \text{ is prefix of } w^R\}$

(D) None of these

Q85. The language accepted by following PDA is



(A) $L = \{n_a(w) + n_b(w) = n_c(w)\}$

(B) $L = \{a^n b^m c^m \mid m, n \geq 0\}$

(C) $L = \{n_a(w) + n_c(w) = n_b(w)\}$

(D) None

Closure properties of CFL, DCFL, (DPDA, NPDA)

Q86. Let L_1 be a context-free language and let $L_2 \subseteq L_1$.

(A) L_2 is definitely context-free.

(B) L_2 may be a context free language.

(C) L_1 and L_2 is definitely not context-free.

(D) L_2 never be a regular

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| Q87. | <p>Let L_1 and L_2 be two DCFL's. Then $L_1 \cap L_2$ is guaranteed to be:</p> <p>(A) Context-Free (B) Deterministic Context-Free</p> <p>(C) Regular (D) can't say</p> |
| Q88. | <p>Let L_1 and L_2 be two DCFL's. Then $L_1 \cup L_2$ is guaranteed to be:</p> <p>(A) Context-Free (B) Deterministic Context-Free</p> <p>(C) Regular (D) can't say</p> |
| Q89. | <p>Let L be a DCFL. Then $\Sigma^* - L$ is guaranteed to be</p> <p>(A) Non-Deterministic Context-Free (B) Deterministic Context-Free</p> <p>(C) Regular (D) can't say</p> |
| Q90. | <p>Which of following is/are correct?</p> <p>S1: If L_1 and $L_1 \cup L_2$ are context free, then L_2 must be context free.</p> <p>S2: If L_1 is context free and L_2 is regular, then $L_1 - L_2$ is context free.</p> <p>(A) Only S1 (B) Only S2</p> <p>(C) Both S1&S2 (D) None of them</p> |
| Q91. | <p>Which of following is/are correct?</p> <p>S1: If L_1 is regular and L_2 is context free, then $L_1 - L_2$ is context free.</p> <p>S2: If L_1 is regular and L_2 is context-free, then $L_1 \cap L_2$ must be a CFL.</p> <p>(A) Only S1</p> <p>(B) Only S2</p> <p>(C) Both S1&S2</p> <p>(D) None of them</p> |
| Q92. | <p>Which of following is/are correct?</p> <p>S1: Every infinite set of strings over a single letter alphabet $\Sigma = \{a\}$ contains an infinite Context free subset.</p> <p>S2: Every infinite context-free set contains an infinite regular subset.</p> <p>(A) Only S1</p> <p>(B) Only S2</p> <p>(C) Both S1&S2</p> <p>(D) Neither S1 nor S2</p> |

| | |
|-------------|--|
| Q93. | <p>Which of following is/are correct?</p> <p>S1: If both L and \bar{L} are context-free, then L must be regular.</p> <p>S2: There is a language L which is context-free but not regular such that \bar{L} is also Context-free.</p> <p>(A) Only S1 (B) Only S2</p> <p>(C) Both S1&S2 (D) Neither S1 nor S2</p> |
| Q94. | <p>Consider the CFL $L1 = \{a^n b^m c^m \mid m, n > 0\}$. Choose a context-free language $L2$ that ensures that $L1 \cap L2$ is not context free?</p> <p>(A) $L2 = \{a^n b^m c^n \mid n, m > 0\}$ (B) $L2 = \{a\}^* \cup \{b\}^* \cup \{c\}$</p> <p>(C) $L2 = \{a, b\}^* L1$ (D) $L2 = \{a^n b^n c^n \mid n > 0\}$</p> |
| Q95. | <p>Which of following is/are correct statement?</p> <p>(i) The language $\{0^n 1^n \cup 0^m 1^m \mid 0 \leq n \leq 1000, m > 0\}$ is regular.</p> <p>(ii) If a language is context-free, then it must be regular</p> <p>(iii) If a language is regular, then it may be infinite.</p> <p>(A)(i) and (ii) (B) (iii) only</p> <p>(C)(i) and (iii) (D)(ii) and (iii)</p> |
| Q96. | <p>The language $L = \{0^i 1^i \cup 1^i 0^i \mid i \geq 0\}$ is</p> <p>(A) Regular (B) Deterministic CFL</p> <p>(C) Non-Deterministic CFL (D) Not a CFL</p> |
| Q97. | <p>Let $L1 = \{b^n a^n \mid n \geq 0\}$ and $L2 = \{(ba)^n \mid n \geq 1\}$. What is $L1 \cap L2$?</p> <p>(A) 0 (B) 1 (C) 2 (D) infinite</p> |
| Q98. | <p>Let $L1 = \{b^n \mid n \geq 1\}$. What is $L1.L1$?</p> <p>(A) $\{b^n b^n \mid n \geq 1\}$ (B) $\{b^n \mid n > 1\}$</p> <p>(C) $\{b^n b^m \mid m, n > 1\}$ (D) $(bb)^+$</p> |
| Q99. | <p>Let $L1 = \{a^n b^n : n \geq 0\}$. Then $L1.L1$ is:</p> <p>(A) $\{a^n b^n : n \geq 0\}$ (B) $\{a^n b^n a^m b^m : m, n \geq 0\}$</p> <p>(C) $\{a^{2n} b^{2n} : n \geq 0\}$ (D) $\{a^n b^n a^n b^n : n \geq 0\}$</p> |

| | | |
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| Q100. | Let $L1 = \{a^n b^n : n \geq 0\}$, and $L2 = \{b^n a^n : n \geq 0\}$. Then $L1.L2$ is: (A) $\{a^n b^{2m} a^n : n, m \geq 0\}$ (C) $\{a^m b^{2m} a^m : m \geq 0\}$ | (B) $\{a^n b^n a^m b^m : m, n \geq 0\}$ (D) $\{a^n b^{n+m} a^m : m, n \geq 0\}$ |
| Q101. | Let $L1 = \{w \mid w \in \Sigma^* \text{ and } w = w^r, \Sigma = \{a, b\}\}$ and $L2 = a^* b^* a^*$. Then $L1 \cap L2 =$ (A) $\{a^n b^n a^m \mid n, m \geq 0\}$ (C) $\{a^m b^n a^n \mid n, m \geq 0\}$ | (B) $\{a^n b^m a^n \mid n, m \geq 0\}$ (D) $\{a^n b^n a^n \mid n, m \geq 0\}$ |
| Q102. | Let $L1 = \{(ab)^n a \mid n \geq 0\}$ and $L2 = \{a(ba)^n \mid n \geq 0\}$. What is $ L1 \cap \overline{L2} $? (A) 0 (C) 2 | (B) 1 (D) infinite |
| Q103. | Let L be a language. Then $\text{symmetric}(L) = \{w : w \in L \text{ and } w^r \in L\}$. Which of the following two statements is/are true? S1: If L is regular, $\text{symmetric}(L)$ is also regular. S2: If L is context-free, $\text{symmetric}(L)$ is also context-free. (A) Only S1 (C) both S1 and S2 | (B) Only S2 (D) Neither S1 nor S2 |
| Q104. | Let $L1 = \{a^n b^n c^p \mid n, p \geq 0\}$ and $L2 = \{a^m b^n c^n \mid m, n \geq 0\}$ then $\overline{(L1 \cup L2)}$ will be (A) \emptyset (B) $a^* b^* c^*$ (C) $a^n b^n c^n$ (D) $\{\Sigma^* - a^* b^* c^*\} \cup \{a^i b^j c^k \mid i \neq j \text{ and } j \neq k\}$ | |
| Q105. | Let $L1 = \{a^n b^n c^p \mid n, p \geq 0\}$ and $L2 = \{a^m b^n c^n \mid m, n \geq 0\}$ then $L1 \cap L2$ is (A) \emptyset (C) $a^n b^n c^n$ | (B) $a^* b^* c^*$ (D) $\{\Sigma^* - a^* b^* c^*\}$ |
| Q106. | Let $L1 = \{w1w2 : w1 = w2 , w1 \in (a, b)^* \text{ and } w2 \in (c, d)^*\}$ and $L2 = \{a, c\}^*$. Then $L1 \cap L2$ will be: (A) \emptyset (C) $a^* c^*$ | (B) $(a + c)^*$ (D) $\{a^n c^n \mid n \geq 0\}$ |

| | |
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| Q107. | <p>Suppose L_1 is a context-free language and L_2 is a non-context-free language, such that $L_1 \cap L_2 = \emptyset$. Then $L_1 \cup L_2$ is</p> <p>(A) Necessarily context-free (B) Necessarily non-context-free</p> <p>(C) May or may not be context-free (D) None of above</p> |
| Q108. | <p>Select the correct statement</p> <p>(i) If L_1 is context-free and L_2 is regular, then $L_1 - L_2$ is context free.</p> <p>(ii) If L_1 is regular and L_2 is context free, then $L_1 - L_2$ is context free.</p> <p>(iii) $L_1 = L_2$ if and only if $L_1^* = L_2^*$</p> <p>(A) Only i (B) ii & iii</p> <p>(C) i & iii (D) None</p> |
| Q109. | <p>Select the correct statement (Assume that L_1, L_2, L_3 are TM acceptable)</p> <p>(I) For any languages L_1, L_2 and L_3, $L_1(L_2 \cap L_3) \subseteq (L_1L_2) \cap (L_1L_3)$</p> <p>(II) For any languages L_1, L_2 and L_3, $(L_1L_2) \cap (L_1L_3) \subseteq L_1(L_2 \cap L_3)$.</p> <p>(A) only I (B) only II</p> <p>(C) Both I&II (D) None of them</p> |
| Q110. | <p>Let A and B be two languages over $\{0,1\}$ such that A is a subset of B. Consider the following statements:</p> <p>(1) If B is finite, then A is finite.</p> <p>(2) If B is context-free, A is context-free.</p> <p>(3) If the complement of B is context-free, then the complement of A is context-free.</p> <p>Which of the above statements is true?</p> <p>(A) 1 only (B) 1 and 2 only</p> <p>(C) 1 and 3 only (D) 2 and 3 only</p> |
| Q111. | <p>How many of the following is/are correct statement?</p> <p>(i) For any CFG G there's a CFG G_0 such that G_0 is not ambiguous and $L(G) = L(G_0)$.</p> <p>(ii) The CFLs are closed under symmetric difference.</p> <p>(iii) If L is context free, the even-length strings of L are context free.</p> <p>(iv) If the prime-length strings of L are context free and the composite-length strings of L are context free then L itself is context free.</p> |

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| <p>Q112.</p> | <p>Which one of the following statements is FALSE?</p> <p>(A) There exist context-free languages such that all the context-free grammars generating them are ambiguous.</p> <p>(B) An unambiguous context free grammar always has a unique parse tree for each string of the language generated by it</p> <p>(C) Both deterministic and non-deterministic pushdown automata always accept the same set of languages</p> <p>(D) A finite set of string from one alphabet is always a context-free language.</p> |
| <p>Q113.</p> | <p>Consider the set of grammars in which every rule that has two symbols on the RHS Must have exactly one variable and one terminal. These grammars can generates</p> <p>(A) It generates all regular languages but no others.</p> <p>(B) It generates some languages that are not context-free.</p> <p>(C) It generates all regular languages and some others, but not all of the CFLs.</p> <p>(D) It can generates all CFLs</p> |
| <p>Q114.</p> | <p>Which of the following statement is/are true?</p> <p>S1: Any regular language can be generated by a context-free grammar</p> <p>S2: Some non-regular languages cannot be generated by any CFG</p> <p>S3: The intersection of a CFL and regular set is a CFL</p> <p>S4: All non-regular languages can be generated by CFGs</p> <p>(A) S1, S2 and S4 only</p> <p>(B) S1, S3 and S4 only</p> <p>(C) S2, S3 and S4 only</p> <p>(D) S1, S2 and S3 only</p> |
| <p>Q115.</p> | <p>Which of the following statements is true?</p> <p>(A) Both the regular languages and the context-free languages are closed under the reverse operation.</p> <p>(B) Neither the regular languages nor the context-free languages is closed under the reverse operation.</p> <p>(C) The regular languages are closed under the reverse operation but the context-free Languages are not.</p> <p>(D)The context-free languages are closed under the reverse operation but the regular languages are not.</p> |

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| Q116. | <p>Which of the following is true?</p> <p>(A) If L is context free, then L^* must be regular.</p> <p>(B) If L is an infinite context-free language, then in any context-free grammar generating L there exists at least one recursive rule.</p> <p>(C) Both a and b are true.</p> <p>(D) None of these.</p> |
| Q117. | <p>Which of the following statement is correct?</p> <p>(A) If there is no pushdown automaton accepting L, then L cannot be regular.</p> <p>(B) If L is accepted by a deterministic PDA, then L' (the complement of L) must be Regular.</p> <p>(C) If, for a given L in $\{a, b\}^*$, there exist x, y, z, such that $y \neq \epsilon$ and $xy^n z \in L$ for all $n \geq 0$, then L must be regular.</p> <p>(D) If L is regular and $L = L_1 \cap L_2$ for some L_1 and L_2, then at least one of L_1 and L_2 must be regular.</p> |
| Q118. | <p>Which of the following is false?</p> <p>(A) If L is an infinite context-free language, then there is some context-free grammar generating L that has no rule of the form $A \rightarrow B$, where A and B are nonterminal symbols.</p> <p>(B) Every context-free grammar can be converted into an equivalent regular grammar.</p> <p>(C) Given a context-free grammar generating L, every string in L has a right-most derivation.</p> <p>(D) All are false.</p> |
| Q119. | <p>Select the correct statements</p> <p>(i) Language $\{0, 1\}$ is context-free.</p> <p>(ii) There is a deterministic Push-down automaton for the language of even length Palindromes.</p> <p>(iii) There are some ambiguous context-free grammars that can be disambiguated (rewritten to become unambiguous).</p> <p>(iv) Every context-free grammar can be disambiguated.</p> <p>(A) i only</p> <p>(B) i and iii only</p> <p>(C) ii and iii only</p> <p>(D) All are true</p> |

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| Q120. | <p>For every context-free grammar, there is a</p> <ul style="list-style-type: none"> (i) Language equivalent left-linear context-free grammar. (ii) Language equivalent deterministic finite automaton. (iii) Language equivalent push down automaton. (iv) Language equivalent Turing machine. <p>(A) i and iii only (B) i, iii and iv only (C) iii and iv only (D) all except i</p> |
| Q121. | <p>Which of the following statement is /are true for an arbitrary context free language?</p> <ul style="list-style-type: none"> (i) Every context-free language A is regular. (ii) Every context-free language is recursive enumerable. (iii) Every context-free language is recursively enumerable but not recursive. (iv) Every context-free language is recursive. <p>(A) ii and iv only (B) i, ii and iv only (C) ii and iii only (D) None</p> |
| Grammar to PDA, | |
| Q122. | <p>Consider the following grammar G:</p> $S \rightarrow aSb \mid aSbb \mid ab$ <p>Suppose you want to construct PDA M (Q, Σ, Γ, δ, z, F) for language generated by given grammar G.</p> <p>M is defined by Q = {q0, q1, qf}, Σ = {a, b} Γ = {S, A, B, z}, F = {qf} and δ is not known.</p> <p>Which of the following is correct transition rule (δ) for M?</p> <p>(a) $\delta(q0, \lambda, S) = \{(q1, Sz)\}$, $\delta(q1, a, S) = \{(q1, SB), (q1, SBB), (q1, B)\}$, $\delta(q1, b, B) = \{(q1, \lambda)\}$, $\delta(q1, \lambda, z) = \{(q1, z)\}$</p> <p>(b) $\delta(q0, \lambda, S) = \{(q1, Sz)\}$, $\delta(q1, a, S) = \{(q1, SB), (q1, SBB), (q1, B)\}$, $\delta(q1, b, B) = \{(q1, \lambda)\}$, $\delta(q1, \lambda, z) = \{(qf, z)\}$</p> <p>(c) $\delta(q0, \lambda, S) = \{(q1, Sz)\}$, $\delta(q1, a, S) = \{(q1, SB), (q1, SBB)\}$, $\delta(q1, b, B) = \{(q1, \lambda)\}$, $\delta(q1, \lambda, z) = \{(qf, z)\}$</p> <p>(d) None of the above</p> |

Q123. Consider the following grammar G: (Assume S is a starting symbol of G)

$$s \rightarrow aA \mid bBC \mid cC$$

$$A \rightarrow a$$

$$B \rightarrow bB \mid b$$

$$C \rightarrow cC \mid d$$

Suppose you want to construct PDA M ($Q, \Sigma, \Gamma, \delta, z, F$) for language generated by given grammar G. M is defined by $Q = \{q_0, q_1, q_f\}$, $\Sigma = \{a, b\}$, $\Gamma = \{S, A, B, C, a, b, c, d, z\}$, $F = \{q_f\}$ and δ is not known. Which of the following is correct transition rule (δ) for M?

(A) $\delta(q_0, \lambda, z) = \{(q_1, Sz)\}$, $\delta(q_1, \lambda, S) = \{(q_1, aA), (q_1, bBC), (q_1, cC)\}$,

$\delta(q_1, \lambda, A) = \{(q_1, a)\}$, $\delta(q_1, \lambda, B) = \{(q_1, bB), (q_1, b)\}$,

$\delta(q_1, \lambda, C) = \{(q_1, cC), (q_1, c)\}$, $\delta(q_1, a, a) = \{(q_1, \lambda)\}$,

$\delta(q_1, c, c) = \{(q_1, \lambda)\}$, $\delta(q_1, d, d) = \{(q_1, \lambda)\}$, $\delta(q_1, b, b) = \{(q_1, \lambda)\}$

(B) $\delta(q_0, \lambda, z) = \{(q_1, Sz)\}$, $\delta(q_1, \lambda, S) = \{(q_1, aA), (q_1, bBC), (q_1, cC)\}$,

$\delta(q_1, \lambda, A) = \{(q_1, a)\}$, $\delta(q_1, \lambda, B) = \{(q_1, bB), (q_1, b)\}$,

$\delta(q_1, \lambda, C) = \{(q_1, cC), (q_1, d)\}$, $\delta(q_1, \lambda, z) = \{(q_f, z)\}$

(C) $\delta(q_0, \lambda, z) = \{(q_1, Sz)\}$, $\delta(q_1, \lambda, S) = \{(q_1, aA), (q_1, bBC), (q_1, cC)\}$,

$\delta(q_1, \lambda, A) = \{(q_1, a)\}$, $\delta(q_1, \lambda, B) = \{(q_1, bB), (q_1, b)\}$,

$\delta(q_1, \lambda, C) = \{(q_1, cC), (q_1, d)\}$, $\delta(q_1, \lambda, z) = \{(q_f, z)\}$,

$\delta(q_1, a, a) = \{(q_1, \lambda)\}$, $\delta(q_1, b, b) = \{(q_1, \lambda)\}$, $\delta(q_1, c, c) = \{(q_1, \lambda)\}$

$\delta(q_1, d, d) = \{(q_1, \lambda)\}$

(D) none of the above

PDA of grammar to string

Q124. Consider the following PDA $M (Q, \Sigma, \Gamma, \delta, z, F)$ for a grammar G :
 M is defined by $Q = \{q_0, q_1, q_f\}$, $\Sigma = \{a, b\}$, $\Gamma = \{S, A, B, z\}$, $F = \{q_f\}$ and the transition rules (δ) are:

$$\begin{aligned} \delta(q_0, \lambda, z) &= (q_1, Sz), & \delta(q_1, \lambda, S) &= \{(q_1, aA), (q_1, bB)\}, \\ \delta(q_1, \lambda, A) &= \{(q_1, aA), (q_1, a)\}, & \delta(q_1, \lambda, B) &= \{(q_1, bB), (q_1, b)\}, \\ \delta(q_1, a, a) &= (q_1, \lambda), & \delta(q_1, b, b) &= (q_1, \lambda), \\ \delta(q_1, \lambda, z) &= (q_f, z) \end{aligned}$$

Which of the following string is not accepted by M ?

- (A) aaaaaa (B) b
 (C) bbaa (D) bbb

Data for next three questions: Consider the following PDA $M (Q, \Sigma, \Gamma, \delta, z, F)$ for a grammar G : M is defined by $Q = \{q_0, q_1, q_f\}$, $\Sigma = \{a, b\}$, $\Gamma = \{S, A, B, C, z\}$, $F = \{q_f\}$ and the transition rules (δ) are:

$$\begin{aligned} \delta(q_0, \lambda, z) &= \{(q_1, Sz)\}, & \delta(q_1, a, S) &= \{(q_1, A)\}, \\ \delta(q_1, a, A) &= \{(q_1, ABC), (q_1, \lambda)\}, & \delta(q_1, b, B) &= \{(q_1, \lambda)\}, \\ \delta(q_1, b, A) &= \{(q_1, B)\}, & \delta(q_1, c, C) &= \{(q_1, \lambda)\}, & \delta(q_1, \lambda, z) &= (q_f, z) \end{aligned}$$

Q125. How many of the following string is/are accepted by M ? _____

- (i) aaabc (ii) aabbc (iii) aabc (iv) aac (v) aa

Q126. [MSQ]

Which of the following grammar is/are equivalent to G ?

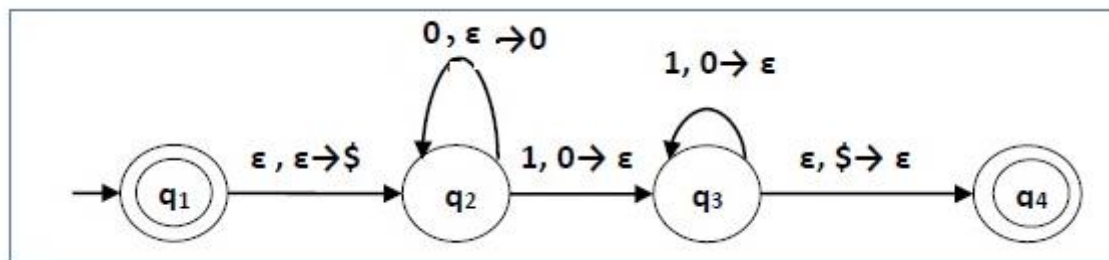
- (a) $S \rightarrow aA$, $A \rightarrow aAbc \mid bb \mid a$
 (b) $S \rightarrow aA$, $A \rightarrow aABc \mid bB \mid a$, $B \rightarrow b$
 (c) $S \rightarrow aA$, $A \rightarrow aABC \mid bB \mid a$, $B \rightarrow b$, $C \rightarrow c$
 (d) $S \rightarrow aS \mid aSbc \mid bb \mid a$

Q127. Which of the following language is equivalent to language generated by G ?

- (A) $L = \{aa^*(a + bb)(bc)^*\}$
 (B) $L = \{aa^k w(bc)^k \mid k \geq 0 \text{ and } w = \{a, bb\}\}$
 (C) $L = \{a^k abb(bc)^k \mid k \geq 0\}$
 (D) None of these

PDA to Grammar

Q128. Consider the following PDA $M (Q, \Sigma, \Gamma, \delta, z)$



Which of the following grammar is/are equivalent to given PDA?

(A) $S \rightarrow 0S1 \mid \lambda$

(B) $S \rightarrow 0S1 \mid 01$

(C) $S \rightarrow 0S1 \mid 0011$

(D) None of the above

(CNF and GNF)

Q129. Which of the following grammar doesn't have useless production?

(A) $S \rightarrow aAa \mid BC, \quad A \rightarrow aa \mid Aa, \quad B \rightarrow cC \mid c$

(B) $S \rightarrow aAb \mid Bc \mid aA, \quad A \rightarrow BC, \quad B \rightarrow bB \mid AC, \quad C \rightarrow cC \mid c$

(C) $S \rightarrow ABb \mid aA, \quad A \rightarrow BC \mid c, \quad B \rightarrow bB \mid AC, \quad C \rightarrow cC \mid CC$

(D) $S \rightarrow aSb \mid \lambda \mid aA, \quad A \rightarrow BC \mid \lambda, \quad B \rightarrow bB \mid AC \mid b, \quad C \rightarrow cC \mid \lambda$

Q130. Consider following grammar G with productions

$S \rightarrow aS \mid BA \mid CD, A \rightarrow a \mid Da, B \rightarrow aa \mid aC, C \rightarrow aCb \mid cC \mid CC$

After removing all useless production from G , which of the following grammar is equivalent to G ?

(A) $S \rightarrow aS \mid a \mid aa$

(B) $S \rightarrow aS \mid aaa$

(C) $S \rightarrow aS \mid aA, A \rightarrow aa$

(D) $b \& c$

Q131. Consider following grammar G with productions

$S \rightarrow Aa \mid C, A \rightarrow bA \mid B, B \rightarrow aB \mid C \mid b, C \rightarrow cC \mid c$

After removing all unit-production from G , which of the following grammar is Equivalent to G ?

(A) $S \rightarrow Aa \mid c, \quad A \rightarrow bA \mid aB \mid cC \mid b \mid c, \quad B \rightarrow aB \mid cC \mid b \mid c, \quad C \rightarrow cC \mid c$

(B) $S \rightarrow Aa \mid cC \mid c, \quad A \rightarrow bA \mid aB \mid cC \mid b \mid c, \quad B \rightarrow aB \mid cC \mid b \mid c$

(C) $S \rightarrow Aa \mid cC \mid c, \quad A \rightarrow bA \mid aB \mid b \mid c, \quad B \rightarrow aB \mid cC \mid b \mid c, \quad C \rightarrow cC \mid c$

(D) $S \rightarrow Aa \mid cC \mid c, \quad A \rightarrow bA \mid aB \mid cC \mid b \mid c, \quad B \rightarrow aB \mid cC \mid b \mid c, \quad C \rightarrow cC \mid c$

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| Q136. | <p>Let G be a grammar in Chomsky Normal Form. Let $w_1, w_2 \in L(G)$ such that $w_1 = w_2$. Which of the following is true?</p> <p>(A) Any derivation of w_1 has exactly the same number of steps as any derivation of w_2.</p> <p>(B) Some derivations of w_1 may be shorter than some derivations of w_2.</p> <p>(C) Different derivations of w_1 have different lengths.</p> <p>(D) Some derivations of w_1 maybe longer than some derivations of w_2.</p> |
| Q137. | <p>[MSQ]</p> <p>Which of the following rule(s) doesn't belong to CNF of any CFG?</p> <p>(A) $S \rightarrow BC$</p> <p>(B) $B \rightarrow Bc$</p> <p>(C) $B \rightarrow a$</p> <p>(D) $S \rightarrow B$</p> |
| <p>For the next three questions, consider the following grammar:</p> <p>$G_1 = (\{S, T\}, \{a, b\}, \{S \rightarrow aT, T \rightarrow aT, T \rightarrow b\}, S)$</p> <p>$G_2 = (\{S, S_1, S_2\}, \{a, b\}, \{S \rightarrow aS_1, S_1 \rightarrow aS_1, S_1 \rightarrow bS_2, S_2 \rightarrow a\}, S)$</p> <p>$G_3 = (\{S, S_1, A, B\}, \{a, b\}, \{S \rightarrow AS_1, S_1 \rightarrow AS_1, S_1 \rightarrow bB, A \rightarrow a, B \rightarrow b\}, S)$</p> <p>$G_4 = (\{S, S_1, S_2, A, B\}, \{a, b\}, \{S \rightarrow AS_1 \mid \lambda, S_1 \rightarrow AS_1 \mid BS_2, A \rightarrow a, B \rightarrow b, S_2 \rightarrow a\}, S)$</p> | |
| Q138. | <p>Which of the above grammars are regular grammars?</p> <p>(A) G_1 and G_2</p> <p>(B) G_2 and G_3</p> <p>(C) G_3 and G_4</p> <p>(D) G_4 and G_1</p> |
| Q139. | <p>Which of the above grammars is in Greibach Normal Form?</p> <p>(A) G_1</p> <p>(B) G_1 and G_2</p> <p>(C) G_3</p> <p>(D) G_4</p> |
| Q140. | <p>Which of the above grammars is in Chomsky Normal Form?</p> <p>(A) G_1</p> <p>(B) G_2</p> <p>(C) G_3</p> <p>(D) G_4</p> |

Q141. Which of the following sets of productions is not in Chomsky normal form?

G1 : $S \rightarrow AB$ $A \rightarrow AB|a$ $B \rightarrow Ba|b$

G2 : $S \rightarrow AB$ $A \rightarrow AB|a$ $B \rightarrow BA|b$

G3 : $S \rightarrow A|B$ $A \rightarrow AB|a$ $B \rightarrow BA|b$

(A) G2

(B) G1, G2, G3

(C) G2 and G3

(D) G1 and G3

Q142. Consider a grammar G

$S \rightarrow abSb | bSaa | aa$

Which of the following grammar are in GNF and equivalent to G?

(A) $S \rightarrow aBSb | bSAA | aA$, $A \rightarrow a$, $B \rightarrow b$

(B) $S \rightarrow aBSb | bSAa | aA$, $A \rightarrow a$, $B \rightarrow b$

(C) $S \rightarrow aBSB | bSAA | A$, $A \rightarrow aa$, $B \rightarrow b$

(D) $S \rightarrow aBSB | bSAA | aA$, $A \rightarrow a$, $B \rightarrow b$

CYK Algorithm

For next four questions: Consider the following grammar G:

$S \rightarrow AB | BC$

$A \rightarrow BA | a$

$B \rightarrow CC | b$

$C \rightarrow AB | a$

We are applying CYK algorithm for the string 'baaba' and given grammar G. The incomplete triangular table is shown below:

| | | | | | |
|---|-----------|-----------|--------|--------|--------|
| 1 | {S, A, C} | | | | |
| 2 | | {S, A, C} | | | |
| 3 | | | | | |
| 4 | {S, A} | {B} | {S, C} | {S, A} | |
| 5 | {B} | {A, C} | | {B} | {A, C} |
| | b | a | a | b | a |
| | 1 | 2 | 3 | 4 | 5 |

For each of following questions select the correct one.

Q143. The cell [2, 1] contains

(A) {B}

(B) {A, B}

(C) \emptyset

(D) {S, A}

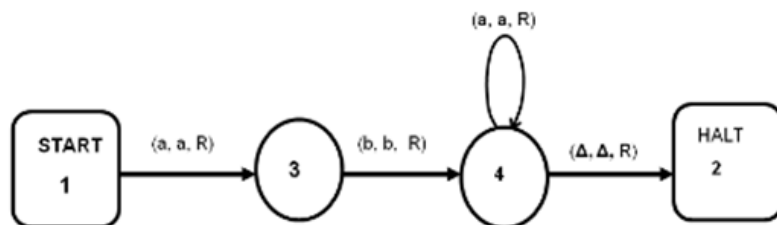
| | |
|------------------------------|---|
| Q144. | <p>The cell [3, 1] contains</p> <p>(A) {B} (B) {A, B}</p> <p>(C) \emptyset (D) {S, A}</p> |
| Q145. | <p>The cell [3, 2] contains</p> <p>(A) {B} (B) {A, B}</p> <p>(C) \emptyset (D) {S, A}</p> |
| Q146. | <p>The cell [5, 3] contains</p> <p>(A) {B} (B) {A, C}</p> <p>(C) \emptyset (D) {S, A}</p> |
| Pumping lemma for CFL | |
| Q147. | <p>Consider the CFG used in the proof of the pumping lemma: we chose a string long enough so that it is accepted by the CFG and repeats a variable. What if such a string does not exist for this CFG?</p> <p>(A) There is no possibility of this happening: there always is such a string.</p> <p>(B) The CFG's language must be finite.</p> <p>(C) The CFG's language may or may not be infinite.</p> <p>(D) Then there is at most 1 variable that produces a single terminal in the CFG.</p> |
| Q148. | <p>Let L be a context free language and $w = uvxyz$ is any string in L such that $w \geq m$; where m is some constant positive integer then which of the following conditions must be satisfied in order to apply pumping lemma?</p> <p>(i) $vxy \leq m$</p> <p>(ii) $vy \neq \epsilon$</p> <p>(iii) $vxy \geq m$</p> <p>(iv) $vy = \epsilon$</p> <p>(A) i and ii only</p> <p>(B) iii and iv only</p> <p>(C) i and iv only</p> <p>(D) ii and iii only</p> |

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| Q149. | <p>Consider the language $L = \{a^i b^j a^i b^j \mid i, j \geq 0\}$. Consider the following "proof" that L satisfies the pumping lemma. Let p be the pumping length. Choose $Z = a^p b a^p b$. Consider a division of Z, where $u = a^{p-1}$, $v = a$, $w = b$, $x = a$ and $y = a^{p-1} b$. Clearly $uv^i wx^i y$ is in L for every i.</p> <p>(A) This is an incorrect proof because all divisions of Z have not been considered.</p> <p>(B) This is an incorrect proof because all possible Z have not been considered.</p> <p>(C) This is a correct proof.</p> <p>(D) None of above</p> |
| Q150. | <p>Consider the language $L = \{a^n b^n c^n \mid n \geq 0\}$. Consider the following "proof" that L does not satisfy the pumping lemma. Let $p \geq 1$ be the pumping length. Choose $z = a^p b^p c^p$. Consider the division of z, where $u = \lambda$, $v = a$, $w = \lambda$, $x = \lambda$, and $y = a^{p-1} b^p c^p$. Clearly $uv^0 wx^0 y$ is not in L.</p> <p>(A) This is an incorrect proof because all divisions of z have not been considered.</p> <p>(B) This is an incorrect proof because all possible z have not been considered.</p> <p>(C) This is a correct proof.</p> <p>(D) None of above</p> |
| Q151. | <p>Select the correct statements</p> <p>(I) Since the pumping lemma applies to all context-free languages, all context-free languages contain an infinite number of strings.</p> <p>(II) The stack alphabet of pushdown automata must be the same as the input alphabet.</p> <p>(III) The stack alphabet of pushdown automata must be different from the input alphabet.</p> <p>(IV) If a language L satisfies the conditions stated in the pumping lemma for CFLs, then L is context free.</p> <p>(A) Only I&II</p> <p>(B) Only I & III</p> <p>(C) Only III&IV</p> <p>(D) None of above</p> |
| Decidability of CFL & CFG | |
| Q152. | <p>Which of the following is/are decidable properties of context free?</p> <p>(I) For context-free grammar G find if string $w \in L(G)$</p> <p>(II) For context-free grammar G find if $L(G) = \emptyset$</p> <p>(III) For context-free grammar G find if $L(G)$ is infinite.</p> <p>(A) Only I</p> <p>(B) Only I&II</p> <p>(C) I, II & III</p> <p>(D) Only II & III</p> |

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| <p>Q153.</p> | <p>Which of the following language is/are Turing decidable?</p> <ol style="list-style-type: none"> 1. $L = \{ \langle G_1, G_2 \rangle \mid G_1 \text{ \& } G_2 \text{ are regular grammar and } L(G_1) \subseteq L(G_2) \}$ 2. $L = \{ \langle G, R \rangle \mid G \text{ is a CFG \& } R \text{ is a regular expression and } L(G) \subseteq L(R) \}$ 3. $L = \{ \langle G, R \rangle \mid G \text{ is a CFG \& } R \text{ is a regular expression and } L(R) \subseteq L(G) \}$ <p>(A) 1 only (B) 1 \& 2 only (C) 2 \& 3 only (D) 1 \& 3 only</p> |
| <p>Q154.</p> | <p>Which of the following language is/are NOT Turing decidable?</p> <ol style="list-style-type: none"> 1. $L = \{ \langle G_1, G_2 \rangle \mid G_1 \text{ \& } G_2 \text{ are CFG and } L(G_1) \cap L(G_2) \text{ is empty} \}$ 2. $L = \{ \langle G_1, G_2 \rangle \mid G_1 \text{ \& } G_2 \text{ are CFG and } L(G_1) \cap L(G_2) \text{ is CFL} \}$ 3. $L = \{ \langle G_1, G_2 \rangle \mid G_1 \text{ \& } G_2 \text{ are CFG and } L(G_1) \cap L(G_2) \text{ is regular} \}$ 4. $L = \{ \langle G_1, G_2 \rangle \mid G_1 \text{ \& } G_2 \text{ are DCFG and } L(G_1) \cap L(G_2) \text{ is empty} \}$ <p>(A) 1 \& 2 only (B) 1, 2 \& 3 only (C) 1, 2 \& 4 only (D) All the above</p> |
| <p>Q155.</p> | <p>Which of the following language is/are not decidable?</p> <ol style="list-style-type: none"> 1. $L = \{ \langle G \rangle \mid G \text{ is a CFG and } \lambda \notin L(G) \}$ 2. $L = \{ \langle G, w \rangle \mid G \text{ is a CFG and } w \notin L(G) \}$ 3. $L = \{ \langle G \rangle \mid G \text{ is a CFG and } G \text{ is unambiguous} \}$ 4. $L = \{ \langle G \rangle \mid G \text{ is a CFG and } L(G) \neq \Sigma^* \}$ <p>(A) 1 \& 2 only (B) 3 \& 4 only (C) 3 only (D) 4 only</p> |
| <p>Q156.</p> | <p>[MSQ]</p> <p>Which of the following languages are undecidable?</p> <p>(A) $L = \{ \langle G, \Sigma \rangle \mid L(G) = \Sigma^*; \text{ where } G \text{ is a CFG and } \Sigma \text{ is finite set of alphabets} \}$</p> <p>(B) $L = \{ \langle G, R \rangle \mid L(G) = R; \text{ where } G \text{ is CFG and } R \text{ is regular language} \}$</p> <p>(C) $L = \{ \langle G_1, G_2 \rangle \mid L(G_1) = L(G_2); \text{ where } G_1 \text{ and } G_2 \text{ are two CFG's} \}$</p> <p>(D) $L = \{ \langle G_1, G_2 \rangle \mid L(G_1) \subseteq L(G_2); \text{ where } G_1 \text{ and } G_2 \text{ are two CFG's} \}$</p> |

TURING MACHINE

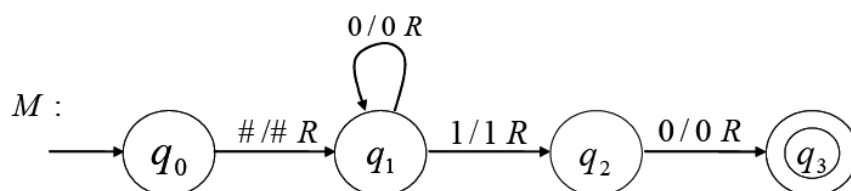
Q1. The language accepted by the following Turing machine is:



- (A) a^*
(C) ba^*

- (B) a^*ba^*
(D) aba^*

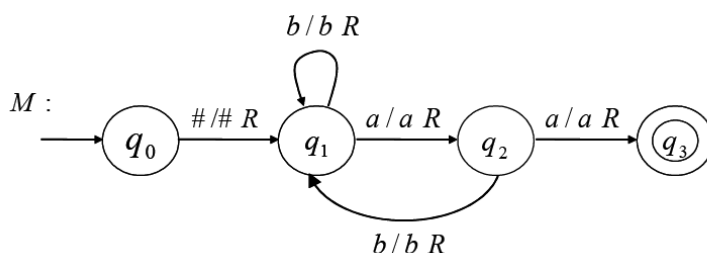
Q2. Given a Turing Machine M , what's $L(M)$?



- (A) 0^*10
(C) $(0+1)^*$

- (B) $0^*10(0+1)^*$
(D) \emptyset

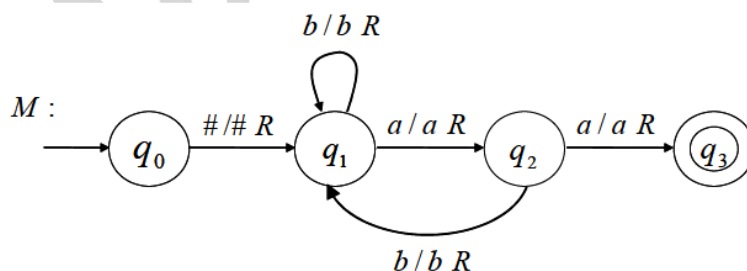
Q3. Given a Turing Machine M , what's $L(M)$?



- (A) $(b+ab)^*aa$
(C) $b^*(ab)^*aa(a+b)^*$

- (B) $b^*aa(a+b)^*$
(D) $(b+ab)^*aa(a+b)^*$

Q4. Given a Turing Machine M , what's $L(M)$?



- (A) $L(M) = (a+b)^*aa(a+b)^*$
(C) $L(M) = b^*ab^*a(a+b)^*$

- (B) $L(M) = (a+b)^*aa$
(D) $L(M) = b^*a(a+b)^*a$

| | |
|-------------------|--|
| <p>Q5.</p> | <p>Consider the following Turing Machine:</p> $M = (\{q_0, q_1, q_2, q_{acc}, q_{rej}\}, \{a, b\}, \{a, b, B\}, \delta, q_0, q_{acc}, q_{rej})$, where $\delta(q_0, a) = (q_1, b, R)$ $\delta(q_1, b) = (q_2, a, L)$ $\delta(q_2, b) = (q_0, b, R)$ $\delta(q_1, B) = (q_{acc}, B, R)$ <p>As always, we assume for cases not mentioned above $\delta(q, a) = (q_{rej}, B, R)$ What can we say about the Turing machine M?</p> <p>(A) M halts on all inputs</p> <p>(B) M never halts on some inputs</p> <p>(C) M does not halt on any input</p> <p>(D) None of these</p> |
| <p>Q6.</p> | <p>What language does the following Turing Machine (TM) accept?</p> $M = (Q = \{q_0, q_1, q_2\}, \Sigma = \{a, b\}, \Gamma = \{a, b, B\}, \delta, q_0, B, F = \{q_2\})$; where the transition function δ is given by $\delta(q_0, a) = \{(q_0, a, R)\}$ $\delta(q_0, b) = \{(q_1, b, R)\}$ $\delta(q_1, B) = \{(q_2, B, R)\}$ $\delta(q, x) = \text{stop everywhere else}$ <p>(A) a^*b (B) aa^*b</p> <p>(C) \emptyset (D) $a^*b(a+b)^*$</p> |
| <p>Q7.</p> | <p>Consider the Turing machine $M = (Q, \Sigma, \Gamma, \delta, q_0, q_{acc}, q_{rej})$; where</p> <ul style="list-style-type: none"> $Q = \{q_0, q_1, q_2, q_3, q_{acc}, q_{rej}\}$ $\Sigma = \{0, 1\}$, and $\Gamma = \{0, 1, A, B, t\}$ δ is given as follows: $\delta(q_0, 0) = (q_1, A, R) \quad \delta(q_0, B) = (q_3, B, R)$ $\delta(q_1, 0) = (q_1, 0, R) \quad \delta(q_1, B) = (q_1, B, R)$ $\delta(q_1, 1) = (q_2, B, L) \quad \delta(q_2, B) = (q_2, B, L)$ $\delta(q_2, 0) = (q_2, 0, L) \quad \delta(q_2, A) = (q_0, A, R)$ $\delta(q_3, B) = (q_3, B, R) \quad \delta(q_3, t) = (q_{acc}, t, R)$ <p>In all other cases, $\delta(q, X) = (q_{rej}, t, R)$. So for example, $\delta(q_0, 1) = (q_{rej}, t, R)$.</p> <p>How many of the following strings are rejected by M? _____</p> <p>i. 10001 ii. 111000 iii. 0010 iv. 10100 v. 101010</p> |

| | | | | | | | | | | | | | | | | | |
|----------------|---|----------------|---------------|---|---|-------|---------------|---|---|-------|---|---------------|---------------|-------|---|---|---|
| Q8. | <p>A Turing machine M with start state q_0 and accepting state q_f has the following transition function:</p> <table><tr><td>$\delta(q, a)$</td><td>0</td><td>1</td><td>B</td></tr><tr><td>q_0</td><td>$(q_1, 0, R)$</td><td>-</td><td>-</td></tr><tr><td>q_1</td><td>-</td><td>$(q_0, 1, R)$</td><td>(q_f, B, R)</td></tr><tr><td>q_f</td><td>-</td><td>-</td><td>-</td></tr></table> <p>Which of the following strings are accepted by TM?</p> <p>(A) 1110 (B) 0110 (C) 01011 (D) 0101</p> | $\delta(q, a)$ | 0 | 1 | B | q_0 | $(q_1, 0, R)$ | - | - | q_1 | - | $(q_0, 1, R)$ | (q_f, B, R) | q_f | - | - | - |
| $\delta(q, a)$ | 0 | 1 | B | | | | | | | | | | | | | | |
| q_0 | $(q_1, 0, R)$ | - | - | | | | | | | | | | | | | | |
| q_1 | - | $(q_0, 1, R)$ | (q_f, B, R) | | | | | | | | | | | | | | |
| q_f | - | - | - | | | | | | | | | | | | | | |
| Q9. | <p>The language accepted by the following Turing machine is</p> <p>$M1 = (\{q_0, q_{\text{accept}}, q_{\text{reject}}\}, \{0, 1\}, \{0, 1, \Delta\}, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$</p> <p>$\delta(q_0, 0) = (q_{\text{accept}}, 0, R)$ $\delta(q_0, 1) = (q_{\text{reject}}, 1, R)$ $\delta(q_0, \Delta) = (q_{\text{reject}}, 1, R).$</p> <p>(A) $0^+ 1^*$ (B) $0^* 1^*$ (C) $(0+1)^*$ (D) $0 (0^*1^*)^*$</p> | | | | | | | | | | | | | | | | |
| Q10. | <p>Consider the Turing machine $M = (Q, \Sigma, \Gamma, \delta, q_0, q_{\text{acc}}, q_{\text{rej}})$ where</p> <ul style="list-style-type: none">$Q = \{q_0, q_1, q_2, q_3, q_{\text{acc}}, q_{\text{rej}}\}$$\Sigma = \{0, 1\}$, and $\Gamma = \{0, 1, A, B, t\}$$\delta$ is given as follows <p>$\delta(q_0, 0) = (q_1, A, R)$ $\delta(q_0, B) = (q_3, B, R)$ $\delta(q_1, 0) = (q_1, 0, R)$ $\delta(q_1, B) = (q_1, B, R)$ $\delta(q_1, 1) = (q_2, B, L)$ $\delta(q_2, B) = (q_2, B, L)$ $\delta(q_2, 0) = (q_2, 0, L)$ $\delta(q_2, A) = (q_0, A, R)$ $\delta(q_3, B) = (q_3, B, R)$ $\delta(q_3, t) = (q_{\text{acc}}, t, R)$</p> <p>In all other cases, $\delta(q, X) = (q_{\text{rej}}, t, R)$. So for example, $\delta(q_0, 1) = (q_{\text{rej}}, t, R)$.</p> <p>Which of the following strings are accepted by M?</p> <p>i) 0011 ii) 0101 iii) 01 iv) 10100 v) 1010</p> <p>(A) i and iii (B) iii and v (C) i, ii and iv (D) ii and iii</p> | | | | | | | | | | | | | | | | |

Q11. The following Turing machine

$M_2 = (\{q_0, q_1, q_{\text{accept}}, q_{\text{reject}}\}, \{0, 1\}, \{0, 1, \#, \Delta\}, \delta, q_0, q_{\text{accept}}, q_{\text{reject}})$

$\delta(q_0, 0) = (q_{\text{accept}}, 0, R)$

$\delta(q_0, 1) = (q_1, 1, R)$

$\delta(q_0, \#) = (q_{\text{reject}}, 1, R)$

$\delta(q_0, \Delta) = (q_{\text{reject}}, 1, R).$

$\delta(q_1, 0) = (q_1, \#, R)$

$\delta(q_1, 1) = (q_1, \#, R)$

$\delta(q_1, \#) = (q_1, \#, R)$

$\delta(q_1, \Delta) = (q_1, \#, R)$

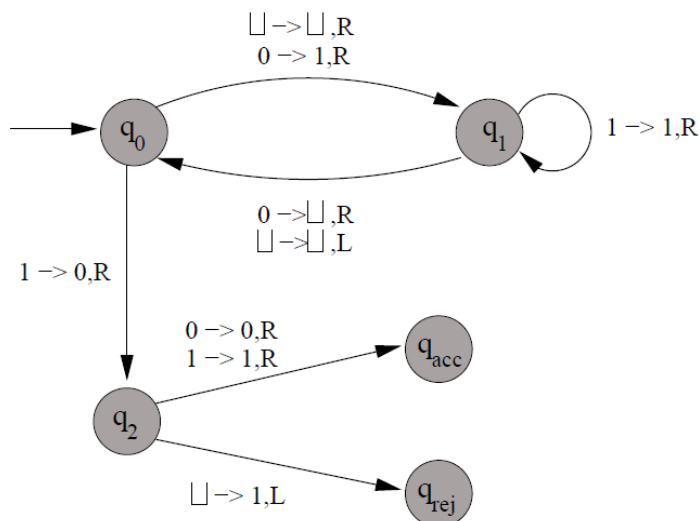
(A) halts on all the string starts with 0

(B) does not halts on λ

(C) halts on all the string starts with 1

(D) does not halts on any string.

Data for next two questions: Consider the following TM



Q12. The outcome of the computation of M on input '00' is

(A) ACCEPT

(B) REJECT

(C) LOOP

(D) None

Q13. The outcome of the computation of M on input '01' is

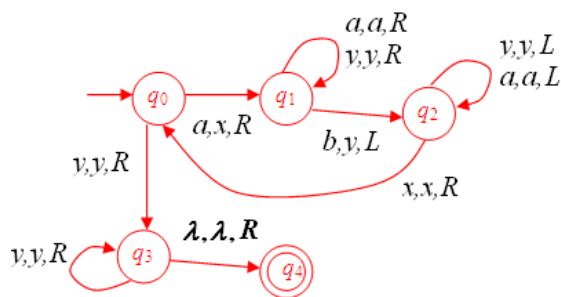
(A) ACCEPT

(B) REJECT

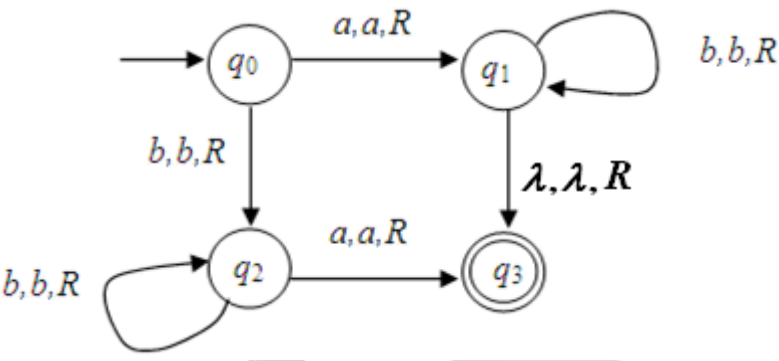
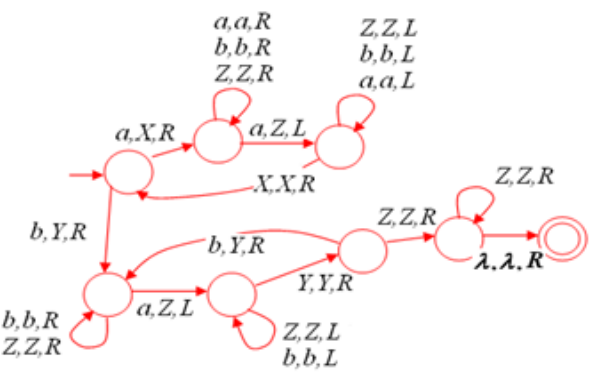
(C) LOOP

(D) None

For Next Three Questions: Consider the Turing machine:

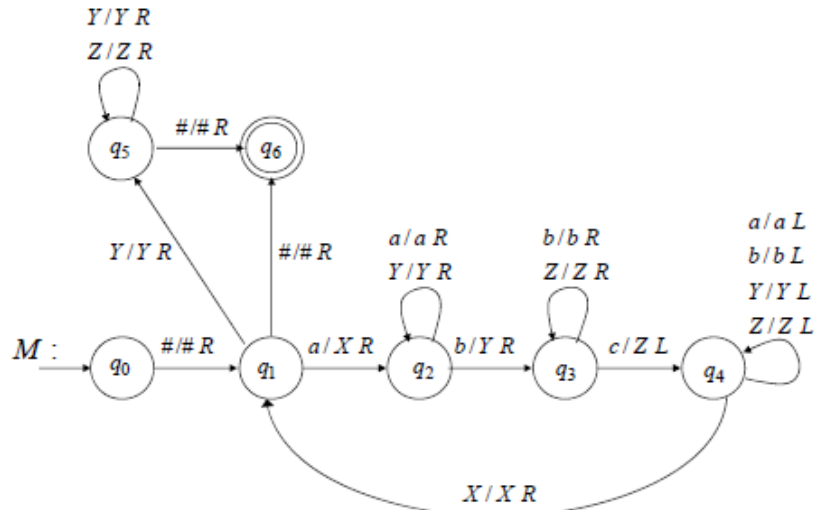


- Q14.** Determine what the Turing machine does when presented with the “aaabbbb”.
- (A) TM halt at final state
 (B) TM does not halt and loop forever
 (C) TM reject the input
 (D) Cannot say about this input
- Q15.** Determine what the Turing machine does when presented with the “aaaaabbbbb”.
- (A) TM halt at final state
 (B) TM does not halt and loop forever
 (C) TM reject the input
 (D) Cannot say about this input
- Q16.** Is there any input for which the Turing machine goes into an infinite loop?
- (A) No.
 (B) Yes, there is an input “abb” for which the Turing machine goes into an infinite loop.
 (C) Yes, there is an input “aaaaabbbbbbbbb” for which the Turing machine goes into an infinite loop.
 (D) None of the above.
- Q17.** What is the smallest number of states that a TM could have?
- (A) 0 (B) 1
 (C) 2 (D) 3
- Q18.** What is the smallest number of tape symbols that a TM with nonempty input alphabet could have?
- (A) 0 (B) 1
 (C) 2 (D) 3

| | |
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| Q19. | <p>Suppose that a Turing Machine is run on an input. Which of the following is false?</p> <p>(A) Infinitely many configurations of the machine can occur.</p> <p>(B) If a configuration of the machine arises twice in a computation on this input, the machine enters an infinite loop.</p> <p>(C) Every configuration contains exactly one state of the machine.</p> <p>(D) A configuration of the machine might have infinite length because the tape is infinite.</p> |
| Q20. | <p>Consider a Turing machine which can move only left, at most 5 times in the course of computation on any input. Consider all TMs that have this property. What languages do they recognize?</p> <p>(A) Exactly the regular languages.</p> <p>(B) Exactly the context-free languages.</p> <p>(C) Exactly the decidable languages.</p> <p>(D) Exactly the recognizable languages</p> |
| Q21. | <p>What language is accepted by the Turing machine whose transition graph is shown below?</p>  <p>(A) The language of the machine is therefore ab^*b</p> <p>(B) The language of the machine is therefore $bb^*a(a+b)^*$</p> <p>(C) The language of the machine is therefore $ab^* + bb^*a(a+b)^*$</p> <p>(D) The language of the machine is therefore $ab^*b + bb^*a(a+b)^*$</p> |
| Q22. | <p>[MSQ]</p> <p>The language accepted by the following Turing machine is</p>  <p>(A) $L = \{a^n b^m a^{n+m} : n \geq 0, m \geq 1\}$</p> <p>(B) $L = \{a^m b^n a^{m+n} : m \geq 0, n \geq 1\}$</p> <p>(C) $L = \{a^n b^m a^{n+m} : n \geq 1, m \geq 1\}$</p> <p>(D) $L = \{a^n b^m a^{n+m} : n \geq 0, m \geq 0\}$</p> |

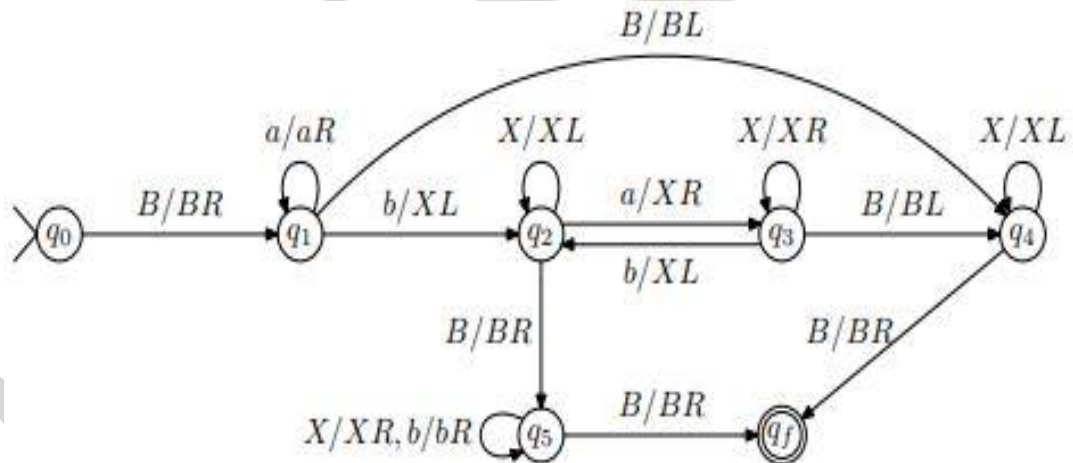
Q23. [MSQ]

Which language does the following TM accept?



- (A) $\{a^k b^k c^k \mid k \geq 0\}$
- (B) $\{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i = k\}$
- (C) $\{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i = k \text{ or } j = k\}$
- (D) $\{a^i b^j c^k \mid i, j, k \geq 0 \text{ \& } i = k \text{ and } j = k\}$

Q24. Construct following Turing machine with input alphabet $\{a, b\}$,



The language accepted by given Turing machine is

- (A) $L = \{a^i b^j \mid i \geq 0, j = i\}$
- (B) $L = \{a^i b^j \mid i \geq 0, j < i\}$
- (C) $L = \{a^i b^j \mid i \geq 0, j \leq i\}$
- (D) $L = \{a^i b^j \mid i \geq 0, j \geq i\}$

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| Q25. | <p>The following language</p> $L1 = \{ \langle M \rangle \mid M \text{ is a TM and there exists an input on which } M \text{ halts in less than } \langle M \rangle \text{ steps} \}$ <p>is:</p> <p>(A) recursive</p> <p>(B) recursively enumerable but not recursive</p> <p>(C) not recursively enumerable</p> <p>(D) Regular</p> |
| Q26. | <p>The language $L2 = \{ \langle M \rangle \mid M \text{ is a TM and } L(M) \leq 3 \}$ is:</p> <p>(A) recursive</p> <p>(B) recursively enumerable but not recursive</p> <p>(C) not recursively enumerable</p> <p>(D) none of the above</p> |
| Q27. | <p>The following language</p> $L2 = \{ \langle M \rangle \mid M \text{ is a TM and } L(M) \geq 3 \}$ <p>is:</p> <p>(A) Recursive</p> <p>(B) recursively enumerable but not recursive</p> <p>(C) not recursively enumerable</p> <p>(D) none of the above</p> |
| Q28. | <p>The following language</p> $L4 = \{ \langle M \rangle \mid M \text{ is a TM that accepts all even numbers} \}$ <p>is:</p> <p>(A) recursive</p> <p>(B) recursively enumerable but not recursive</p> <p>(C) not recursively enumerable</p> <p>(D) none of these</p> |
| Q29. | <p>The following language</p> $L1 = \{ \langle M \rangle \mid M \text{ is a TM and } L(M) \text{ is finite} \}$ <p>is:</p> <p>(A) recursive</p> <p>(B) recursively enumerable but not recursive</p> <p>(C) not recursively enumerable</p> <p>(D) context-Sensitive but not CFL</p> |

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| Q30. | <p>The following language</p> $L1 = \{ \langle M \rangle \mid \text{there exist } x, y \in \Sigma^* \text{ such that either } x \in L(M) \text{ or } y \notin L(M) \}.$ <p>(A) recursive</p> <p>(B) recursively enumerable but not recursive</p> <p>(C) not recursively enumerable</p> <p>(D) context-Sensitive but not CFL</p> |
| Q31. | <p>The following language</p> $L1 = \{ \langle M, x, k \rangle \mid M \text{ is a TM, and } M \text{ does not halt on } x \text{ within } k \text{ steps} \}$ <p>(A) Recursive</p> <p>(B) Recursively enumerable but not recursive</p> <p>(C) Not recursively enumerable</p> <p>(D) Not recursive</p> |
| Q32. | <p>The following language</p> $L1 = \{ \langle M \rangle \mid M \text{ is a TM, and } M \text{ accepts (at least) two strings of different lengths} \}$ <p>(A) Recursive</p> <p>(B) Recursively enumerable but not recursive</p> <p>(C) Not recursively enumerable</p> <p>(D) Not recursive</p> |
| Q33. | <p>Let the language $A_{TM} = \{ \langle M, w \rangle \mid M \text{ is a TM that accepts string } w \}$ then the complement of the language is</p> <p>(A) Decidable</p> <p>(B) Turing Recognizable</p> <p>(C) Not Turing Recognizable</p> <p>(D) None</p> |

Data for next ten questions: Consider the following languages (Q34 to 43 - 2020)

$L1 = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ accepts at most 225 distinct inputs} \}$

$L2 = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ accepts at least 225 distinct inputs} \}$

$L3 = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ accepts exactly 225 distinct inputs} \}$

$L4 = \{ \langle M, w \rangle \mid M \text{ is a TM and } M \text{ doesn't halt on 'w'} \}$.

$L5 = \{ \langle M, w \rangle \mid M \text{ is a TM and } M \text{ halts on 'w'} \}$

$L6 = \{ \langle M, w \rangle \mid M \text{ is a TM and } M \text{ accepts 'w'} \}$

$L7 = \{ \langle M, w \rangle \mid M \text{ is a TM and } M \text{ rejects 'w'} \}$

$L8 = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ has exactly five states} \}$

$L9 = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ has exactly three tapes} \}$

$L10 = \{ \langle M \rangle \mid M \text{ is a TM and } L(M) = \{001, 100\} \}$

| | | | | | |
|-------------|-------------|---------------|-----------|------------|----------------|
| Q34. | Then L1 is | (A) Recursive | (B) RE | (C) Not RE | (D) finite |
| Q35. | Then L2 is | (A) Recursive | (B) RE | (C) Not RE | (D) finite |
| Q36. | Then L3 is | (A) Recursive | (B) RE | (C) Not RE | (D) finite |
| Q37. | Then L4 is | (A) Recursive | (B) RE | (C) Not RE | (D) finite |
| Q38. | Then L5 is | (A) Recursive | (B) RE | (C) Not RE | (D) finite |
| Q39. | Then L6 is | (A) Recursive | (B) RE | (C) Not RE | (D) finite |
| Q40. | Then L7 is | (A) Recursive | (B) RE | (C) Not RE | (D) finite |
| Q41. | Then L8 is | (A) Recursive | (B) RE | (C) Not RE | (D) Co-RE |
| Q42. | Then L9 is | (A) Recursive | (B) RE | (C) Not RE | (D) Co-RE |
| Q43. | Then L10 is | (A) RE | (B) Co-RE | (C) Not RE | (D) Both b & c |

| | |
|-------------------|---|
| Q44. [MSQ] | <p>The language $A_{DFA} = \{ \langle D, w \rangle \mid D \text{ is DFA that accepts } w \}$ is</p> <p>(A) Decidable (B) Turing Recognizable</p> <p>(C) Not Turing Recognizable (D) None</p> |
| Q45. | <p>The language $A_{CFG} = \{ \langle G \rangle \mid G \text{ is CFG and } L(G) = \Sigma^* \}$ is</p> <p>(A) Decidable (B) Turing Recognizable</p> <p>(C) Not Turing Recognizable (D) None</p> |
| Q46. [MSQ] | <p>The language $E_{CFG} = \{ \langle G \rangle \mid G \text{ is a CFG and } L(G) = \emptyset \}$ is</p> <p>(A) Decidable (B) Turing Recognizable</p> <p>(C) Not Turing Recognizable (D) None</p> |
| Q47. | <p>$EQ_{TM} = \{ \langle M_1, M_2 \rangle \mid M_1, M_2 \text{ are TMs with } L(M_1) = L(M_2) \}$?</p> <p>(A) Decidable (B) Turing Recognizable</p> <p>(C) Not Turing Recognizable (D) None</p> |
| Q48. | <p>$EQ_{DFA} = \{ \langle M_1, M_2 \rangle \mid M_1, M_2 \text{ are DFAs with } L(M_1) = L(M_2) \}$?</p> <p>(A) Decidable (B) Turing Recognizable</p> <p>(C) Not Turing Recognizable (D) None</p> |
| Q49. | <p>Consider the following language</p> <p>$L = \{ \langle M \rangle : M \text{ is TM such that } L(M) = \emptyset \}$.</p> <p>(A) Turing recognizable (B) Turing decidable</p> <p>(C) Not Turing recognizable (D) CO-RE</p> |
| Q50. | <p>Which of the following statements is true for every language $L \subseteq \{0, 1\}^*$?</p> <ol style="list-style-type: none"> 1. L is non-empty. 2. L is decidable or L is infinite (or both). 3. L is accepted by some DFA with 42 states if and only if L is accepted by some NFA with 42 states. 4. L is decidable if and only if its complement L is undecidable. <p>(A) 1 and 4 Only (B) 2 and 3 Only</p> <p>(C) 2 Only (D) 1, 2, 3 Only</p> |

Data for next two questions: Let M be a standard Turing machine (with a single one-track tape and a single head) that decides the regular language 0^*1^* .

Q51. Which of the following must be true?

1. Given an empty initial tape, M eventually halts.
2. M accepts the string 1111.
3. M rejects the string 0110.
4. M moves its head to the right at least once, given input 1100.

- (A) 2 & 3 only
(B) 1 & 4 only
(C) 2, 3 & 4 only
(D) All are true

Q52. Which of the following is /are not always true?

1. M moves its head to the right at least once, given input 0101.
2. M never accepts before reading a blank.
3. For some input string, M moves its head to the left at least once.
4. For some input string, M changes at least one symbol on the tape.

- (A) 2, 3 & 4
(B) 3 & 4
(C) 1 & 2
(D) All

Q53. Which of the following problems about Turing machines is/are undecidable?

1. To determine, given a Turing machine M , a state q , and a string w , whether M ever reaches state q when started with input w from its initial state.
2. To determine, given a Turing machine M and a string w , whether M ever moves its head to the left when started with input w .
3. To determine, given Turing machine M and a string w , whether w is accepted by M within some given time constraint T .
4. To determine, given a Turing machine M , and a string w , whether M ever halt for input w .

- (A) 1 & 2 only
(B) 2 and 4 only
(C) 1, 2 & 4 only
(D) all are undecidable

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| Q54. | Which of the following problems about context-free grammars are undecidable? 1. Given a context-free grammar G , is $\varepsilon \in L(G)$? 2. Given a context-free grammar G , is $\{\varepsilon\} = L(G)$? 3. Given two context-free grammars G_1 and G_2 , is $L(G_1) \subseteq L(G_2)$. (A) 1 only (B) 2 only (C) 3 only (D) 1, 2 and 3 |
| Data for next two questions: Let G_1 and G_2 are CFGs and R denotes given regular expression. | |
| Q55. | Which of the following problems are undecidable? 1. $L(G_1) \cap L(G_2) \neq \emptyset$ 2. $L(G_1) \neq L(G_2)$ 3. $L(G_1) \neq L(R)$ (A) 1 only (B) 2 only (C) 1, 3 only (D) 1, 2 and 3 |
| Q56. | Which of the following problems is/are decidable? 1. $L(G_2) - L(G_1) \neq \emptyset$ 2. $L(R) - L(G_1) \neq \emptyset$ 3. $L(R) = \text{finite}$ 4. $L(G_1 \cup G_2) = \Sigma^*$ (A) 1 & 2 only (B) 3 & 4 only (C) 1, 3 only (D) 3 only |
| Q57. | How many of the following languages are Un-decidable? _____ 1. $L = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ accepts at least one string over } \{0, 1\} \}$. 2. $L = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ accept all strings over } \{0, 1\} \}$. 3. $L = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ accepts "Hello"} \}$. 4. $L = \{ \langle M_1, M_2 \rangle \mid M_1 \text{ and } M_2 \text{ are TM's and } M_1 \text{ accept more strings than } M_2 \}$. 5. $L = \{ \langle M \rangle \mid M \text{ is a TM and } M \text{ take more than 1000 steps to process input } w \}$. 6. $L = \{ \langle M_1, M_2 \rangle \mid M_1 \text{ and } M_2 \text{ are TM's and } M_1 \text{ take more steps than } M_2 \text{ to process input } w \}$ |
| Q58. | Suppose that a Turing Machine is run on an input. Which of the following is false? (A) Infinitely many configurations of the machine can occur. (B) If a configuration of the machine arises twice in a computation on this input, the machine enters an infinite loop. (C) Every configuration contains exactly one state of the machine. (D) A configuration of the machine might have infinite length because the tape is infinite |

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| <p>Q59.</p> | <p>Which of the following statement is/are true?</p> <ul style="list-style-type: none"> i. There is a language that is decidable and its complement is not decidable. ii. There is a language that is recognizable and its complement is not recognizable. iii. The problem of determining whether a TM accepts at least 7 strings is undecidable. iv. The problem of determining whether a TM has at least 7 states is undecidable. <p>(A) i and iv only (B) ii and iii only (C) i, ii and iii only (D) iii and iv only</p> |
| <p>Q60.</p> | <p>Consider a Turing machine which can move only left, at most 5 times in the course of computation on any input. Consider all TMs that have this property. What languages do they recognize?</p> <p>(A) Exactly the regular languages. (B) Exactly the context-free languages. (C) Exactly the decidable languages. (D) Exactly the recognizable languages</p> |
| <p>Q61.</p> | <p>Select all the statements that are TRUE. The class of recursively enumerable languages :</p> <ul style="list-style-type: none"> 1. Not closed under union 2. Closed under intersection 3. Closed under complement 4. Not closed under Kleen's star 5. Not closed under infinite union <p>(A) 1, 2 and 4 only (B) 3 and 5 only (C) 2 and 5 only (D) 1 and 4 only</p> |
| <p>Q62.</p> | <p>Which of the following statement/s is/are true?</p> <ul style="list-style-type: none"> 1. If A is regular and B is regular, then $A \cup B^c$ is always Recursive. 2. If A is context-free and B is regular, then $A \cup B^c$ is always Recursively enumerable 3. If A is regular and B is context-free, then $A \cap B^c$ is regular. 4. If A is recursive and B is recursive, then $A \cap B^c$ is always Recursively enumerable but not recursive. <p>(A) 1 only (B) 1 & 2 only (C) 1, 2 & 4 only (D) All the above</p> |

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| <p>Q67.</p> | <p>Consider the following sets of languages over the alphabet $\{0, 1\}$:</p> <ul style="list-style-type: none"> • L_{DTM} is the set of all languages $L \subseteq \{0, 1\}^*$ such that L is accepted by at least one deterministic Turing machine. • L_{NTM} is the set of all languages $L \subseteq \{0, 1\}^*$ such that L is accepted by at least one non-deterministic Turing machine. <p>Which of the following statements is/are true?</p> <p>(A) $L_{DTM} \subseteq L_{NTM}$</p> <p>(B) $L_{DTM} \supseteq L_{NTM}$</p> <p>(C) $L_{DTM} = L_{NTM}$</p> <p>(D) Undecidable problem</p> |
| <p>Q68.</p> | <p>Which of the following statements are correct?</p> <p>i) For every recursive language L, there is a Turing machine M with \bar{L} as its language.</p> <p>ii) For every non-deterministic push-down automaton, there is an equivalent deterministic push-down automaton.</p> <p>iii) Non-deterministic Turing machines are strictly more powerful than Deterministic Turing machines.</p> <p>(A) i only</p> <p>(B) ii only</p> <p>(C) iii only</p> <p>(D) i and iii only</p> |
| <p>Q69.</p> | <p>Which of the following statements are correct?</p> <p>i) DFA reads its input fully before accepting a string.</p> <p>ii) A multi-tape TM is equivalent to a single tape TM.</p> <p>iii) A DTM may accept a string without reading its input.</p> <p>iv) The number of configurations of an LBA is fixed by its number of states Q.</p> <p>(A) ii and iii only</p> <p>(B) ii only</p> <p>(C) i, ii and iii</p> <p>(D) All the above</p> |

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| <p>Q74.</p> | <p>$L1 - R = \emptyset$ is decidable for</p> <div> <div>i) regular languages</div> <div>ii) context-free languages</div> <div>iii) recursive languages</div> <div>iv) recursively enumerable languages</div> <div>(A) Only ii</div> <div>(B) Only i & ii</div> <div>(C) Only iii & iv</div> <div>(D) Only ii & iii</div> </div> |
| <p>Q75.</p> | <p>$\exists y \in L1, y < 5$ is decidable for</p> <div> <div>i) regular languages</div> <div>ii) context-free languages</div> <div>iii) recursive languages</div> <div>iv) recursively enumerable languages</div> <div>(A) Only iii & iv</div> <div>(B) Only ii & iii</div> <div>(C) Only i, ii & iii</div> <div>(D) Only i & iv</div> </div> |
| <p>Q76.</p> | <p>Which of following is/are correct?</p> <p>S1: Deterministic two-stack automata are as powerful as Turing machines.</p> <p>S2: $\{ \langle M \rangle \mid M \text{ is a Turing machine and } M \text{ does not accept } 101 \}$ is recursively enumerable</p> <div> <div>(A) Only S1</div> <div>(B) Only S2</div> <div>(C) Both S1&S2</div> <div>(D) None of them</div> </div> |
| <p>Q77.</p> | <p>Which of the following problem is undecidable?</p> <div> <div>(A) membership problem for CFL</div> <div>(B) membership problem for regular sets</div> <div>(C) membership problem for CSL</div> <div>(D) membership problem for type 0 languages</div> </div> |
| <p>Q78.</p> | <p>Which of the following is/are true about uncountable set?</p> <div> <div>(i) If A is a subset of B and A is uncountable, then so is B.</div> <div>(ii) If A is infinite (even countably infinite) then the power set of A is uncountable</div> <div>(iii) If A is uncountable and B is any set, then the union $A \cup B$ is also uncountable.</div> <div>(iv) If A is uncountable and B is any set, then the Cartesian product $A \times B$ is also uncountable.</div> <div>(A) i, ii only</div> <div>(B) iii, iv only</div> <div>(C) ii, iii, iv only</div> <div>(D) All the above</div> </div> |

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| Q79. | <p>Which of the following statement is/are True?</p> <ol style="list-style-type: none"> 1. The language $(0 \cup 1)^*$ is countable. 2. The set of all possible language over $\{a, b\}$ is countable. 3. The set of recursively enumerable languages is countable. 4. The Cartesian product of a finite number of countable set is countable. <p>(A) 1 & 4 only (B) 2 & 3 only (C) 1, 3 & 4 only (D) All the above</p> |
| Q80. | <p>Which of the following statement is/are NOT true?</p> <ol style="list-style-type: none"> 1. Set of all prime numbers is countable. 2. Set of rational numbers is uncountable. 3. Set of Binary Sequences is uncountable. 4. Set of real Numbers is uncountable. 5. Set of all integers is countable. <p>(A) 1 & 2 only (B) 3, 4 & 5 only (C) 2 only (D) 1, 2 & 3 only</p> |
| Q81. | <p>Which of the following statement is/are true?</p> <ol style="list-style-type: none"> 1. The set of all pairs, (i, j) with i, j positive integers, is countable. 2. The set of all triplets, (i, j, k) with i, j, k positive integers, is countable. <p>(A) 1 only (B) 2 only (C) Both 1 & 2 (D) Neither 1 nor 2</p> |
| Q82. | <p>If S_1 and S_2 are countable set, then how many of the following statement is/are true?</p> <ol style="list-style-type: none"> 1. $S_1 \cup S_2$ is countable. 2. $S_1 \cap S_2$ is countable. 3. $S_1 \oplus S_2$ is countable. 4. $S_1 - S_2$ is countable. |
| Q83. | <p>Which of the following statement is/are true?</p> <ol style="list-style-type: none"> 1. Subset of infinite countable set is countable. 2. Superset of countable set may be uncountable. 3. There exists a countably infinite set which is subset of an uncountable set. 4. Power set of countably infinite set is always countable. <p>(A) 1 & 2 only (B) 3 & 4 only (C) 1, 2 & 3 only (D) All the above</p> |

Closure Properties

| operation | REG | DCFL | CFL | CSL | RC | RE |
|--------------------------------------|-----|------|-----|-----|----|----|
| union | Y | N | Y | Y | Y | Y |
| intersection | Y | N | N | Y | Y | Y |
| set difference | Y | N | N | Y | Y | N |
| complementation | Y | Y | N | Y | Y | N |
| intersection with a regular language | Y | Y | Y | Y | Y | Y |
| concatenation | Y | N | Y | Y | Y | Y |
| Kleene star | Y | N | Y | Y | Y | Y |
| Kleene plus | Y | N | Y | Y | Y | Y |
| reversal | Y | Y | Y | Y | Y | Y |

Closure properties of language families

| Operation | Regular | DCFL | CFL | CSL | Recursive | RE |
|--------------------------------------|---------|------|-----|-----|-----------|-----|
| Union | yes | no | yes | yes | yes | yes |
| Intersection | yes | no | no | yes | yes | yes |
| Complement | yes | yes | no | yes | yes | no |
| Concatenation | yes | no | yes | yes | yes | yes |
| Kleene star | yes | no | yes | yes | yes | yes |
| Homomorphism | yes | no | yes | no | no | yes |
| ϵ -free Homomorphism | yes | no | yes | yes | yes | yes |
| Substitution (ϵ -free) | yes | no | yes | yes | no | yes |
| Inverse Homomorphism | yes | yes | yes | yes | yes | yes |
| Reverse | yes | no | yes | yes | yes | yes |
| Intersection with a regular language | yes | yes | yes | yes | yes | yes |

List of theorems:

1. The class of regular languages is closed under union.
2. The class of regular languages is closed under concatenation.
3. Every NFA has an equivalent DFA.
4. The class of regular languages is closed under Kleene - star.
5. (Kleene's theorem) Language A is regular if A has a regular expression.
6. If A is finite language, then A is regular.
7. The class of regular languages is closed under intersection.
8. The class of regular languages is closed under complementation.
9. If A is regular language, then \exists number p where, if $s \in A$ with $|s| \geq p$, then \exists strings x, y, z such that
 $s = xyz$ and (1) $xy^iz \in A$ for each $i \geq 0$, (2) $|y| > 0$, and (3) $|xy| \leq p$
10. Every CFL can be described by a CFG $G = (V, \Sigma, R, S)$ in Chomsky normal form, i.e., each rule in G has one of two form s: $A \rightarrow BC$ or $A \rightarrow x$, where $(A \in V)$, $(B, C \in V - \{S\})$, $x \in \Sigma$, and we also allow the rule $S \rightarrow \epsilon$.
11. If A is a regular language, then A is also a CFL.
12. A language is context free iff some PDA recognizes it.
13. (Pumping lemma for CFLs) For every CFL L, \exists pumping length p such that \forall strings $s \in L$ with $|s| \geq p$, we can write $s = uvxyz$ with (1) $uv^ixy^iz \in L \forall i \geq 0$, (2) $|v| \geq 1$, (3) $|vxy| \leq p$.
14. The class of CFLs is closed under union.
15. The class of CFLs is closed under concatenation.
16. The class of CFLs is closed under Kleene- star.
17. For every multi- tape TM M, there is a single- tape TM M' such that $L(M) = L(M')$
18. Every NTM has an equivalent deterministic TM.
19. Language L is Turing-recognizable if an NTM recognizes it.
20. A language is enumerable if some enumerator enumerates it

Church-Turing Thesis The information of algorithm is the same as Turing machine algorithm

1. $A_{DFA} = \{ \langle B, w \rangle \mid B \text{ is a DFA that accepts string } w \}$ is Turing-decidable.
2. $A_{NFA} = \{ \langle B, w \rangle \mid B \text{ is an NFA that accepts string } w \}$ is Turing-decidable.
3. $A_{REX} = \{ \langle R, w \rangle \mid R \text{ is a regular expression that generates string } w \}$ is Turing-decidable.
4. $E_{DFA} = \{ \langle B \rangle \mid B \text{ is a DFA with } L(B) = \emptyset \}$ is Turing-decidable.

5. $EQ_{DFA} = \{ \langle A, B \rangle \mid A \text{ and } B \text{ are DFAS with } L(A) = L(B) \}$ is Turing-decidable
6. $A_{CFG} = \{ \langle G, w \rangle \mid G \text{ is a CFG that generates string } w \}$ is Turing-decidable.
7. $E = \{ G \mid G \text{ is a CFG with } L(G) = \emptyset \}$ is Turing-decidable.
8. Every CFL is Turing-decidable
9. $A_{TM} = \{ \langle M, w \rangle \mid M \text{ is a TM that accepts string } w \}$ is undecidable.
10. The set R of all real numbers is uncountable
11. Some languages are not Turing-recognizable.
12. A language is decidable if it is both Turing-recognizable and co-Turing-recognizable.
13. $A_{TM}(\bar{})$ is not Turing-recognizable.
14. $HALT_{TM} = \{ \langle M, w \rangle \mid M \text{ is a TM that halts on } w \}$ is undecidable.
15. $ETM_{TM} = \{ \langle M \rangle \mid M \text{ is a TM with } L(M) = \emptyset \}$ is undecidable.
16. $REG_{TM} = \{ \langle M \rangle \mid M \text{ is a TM and } L(M) \text{ is regular} \}$ is undecidable.
17. $EQ_{TM} = \{ \langle M_1, M_2 \rangle \mid M_1, M_2 \text{ are TMs with } L(M_1) = L(M_2) \}$ is undecidable.
18. (Rice's Thm.) Let P be any subset of the class of Turing-recognizable languages such that $P \neq \emptyset$ and $P(\bar{}) \neq \emptyset$. Then $L_P = \{ \langle M \rangle \mid L(M) \in P \}$ is undecidable.
19. If $A \leq_m B$ and B is Turing-decidable, then A is Turing-decidable.
20. If $A \leq_m B$ and A is undecidable, then B is undecidable.
21. If $A \leq_m B$ and B is Turing-recognizable, then A is Turing-recognizable.
22. If $A \leq_m B$ and A is not Turing-recognizable, then B is not Turing-recognizable.
23. $E_{TM} = \{ \langle M \rangle \mid M \text{ is a TM with } L(M) = \emptyset \}$ is not Turing-recognizable.
24. $EQ_{TM} = \{ \langle M_1, M_2 \rangle \mid M_1, M_2 \text{ are TMs with } L(M_1) = L(M_2) \}$ is neither Turing-recognizable nor co-Turing-recognizable.
25. Let $t(n)$ be a function with $t(n) \geq n$. Then any $t(n)$ -time multi-tape TM has an equivalent $O(t^2(n))$ -time single-tape TM.
26. Let $t(n)$ be a function with $t(n) \geq n$. Then any $t(n)$ -time NTM has an equivalent $2^{O(t(n))}$ -time deterministic 1-tape TM.
27. $PATH \in P$.
28. $REL\ PRIME \in P$.
29. Every CFL is in P .
30. A language is in NP iff it is decided by some nondeterministic polynomial-time TM.
31. $NP = \bigcup_{k \geq 0} NTIME(n^k)$

32. CLIQUE \in NP.
33. SUBSET-SUM \in NP.
34. If $A \leq_P B$ and $B \in P$, then $A \in P$.
35. 3SAT is polynomial-time reducible to CLIQUE
36. If there is an NP-Complete problem B and $B \in P$, then $P = NP$.
37. If B is NP- Complete and $B \leq_P C$ for $C \in NP$, then C is NP- Complete.
38. SAT is NP- Complete.
39. SAT is NP- Complete
40. CLIQUE is NP- Complete.
41. ILP is NP-Complete.

Decidability/Undecidability

- Every question we've asked about the language of finite automata is decidable
Is $w \in L(M)$?, Is $L(M) = \emptyset$, Is $L(M)$ finite, Is $L(M) = \Sigma^*$
- Some questions about the languages of CFGs are decidable
Is $L(G) = \emptyset$, some aren't like Is $L(G) = \Sigma^*$
- Everything we've tried for Turing machines is undecidable.
Is $L(G) = \emptyset$, $L(G) = \Sigma^*$ • Is the language semi decided by M regular? Context-free?

For Context free Grammar it is undecidable

- whether a CFG accepts the language of all strings
 - whether two CFGs describe the same language
 - whether CFG is ambiguous
 - if a context-sensitive grammar describes a context-free language
 - if a CFG, has an equivalent PDA that is deterministic
 - if Given an ambiguous CFG, whether or not there is a different CFG that generates the same language but is not ambiguous
 - whether the complement of a given CFL is context free?
 - whether the intersection of two CFL's is context free?
 - given two CFG's, how can we tell if they have a word in common?
- Decidable Problem for CFG
- DPDA Equality is Decidable
 - Membership is decidable