



ECE408 / CS483 / CSE408  
Summer 2025

Applied Parallel Programming

Lecture 9: Tiled Convolution

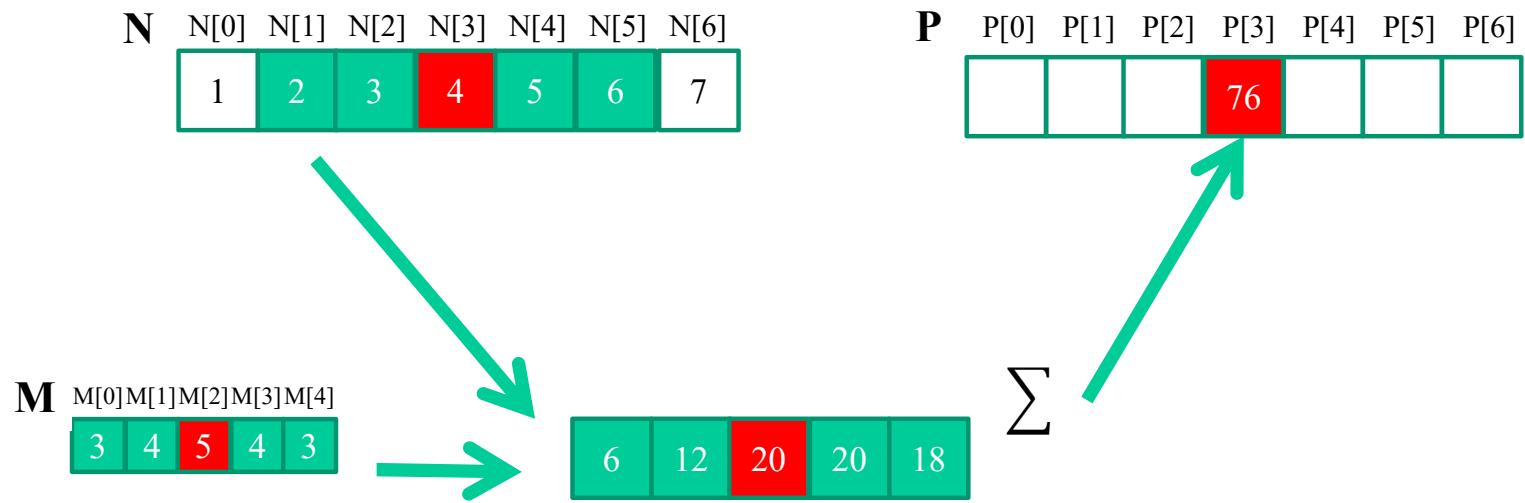


# What Will You Learn Today?

- tiled convolution algorithms
  - some intricate aspects of tiling algorithms
  - output tiles versus input tiles
  - three different styles of input tile loading
- prepare for MP-4: tiled 3D convolution

# Recall the 1D Convolution Operation

- Calculation of  $P[3]$



# N And the 2D Convolution Operation P

1	2	3	4	5	6	7
2	3	4	5	6	7	8
3	4	5	6	7	8	9
4	5	6	7	8	5	6
5	6	7	8	5	6	7
6	7	8	9	0	1	2
7	8	9	0	1	2	3

M

1	2	3	2	1
2	3	4	3	2
3	4	5	4	3
2	3	4	3	2
1	2	3	2	1




$\sum$

1	4	9	8	5
4	9	16	15	12
9	16	25	24	21
8	15	24	21	16
5	12	21	16	5

# Are We Memory-Limited?

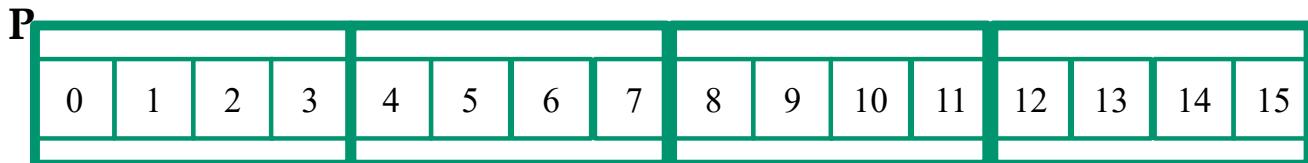
For the 1D case, every output element requires

- $2 * \text{MASK\_WIDTH}$  loads (of M and N each) and
- $2 * \text{MASK\_WIDTH}$  floating-point operations.
- **Memory-limited.**

For the 2D case, every output element requires

- $2 * \text{MASK\_WIDTH}^2$  loads and
- $2 * \text{MASK\_WIDTH}^2$  floating-point operations.
- **Also memory-limited.**

# Tiled 1D Convolution Basic Idea



N

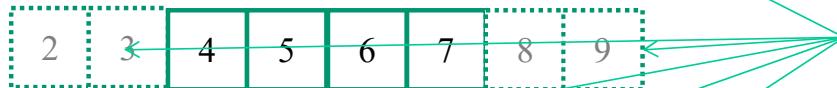


N

Tile 0



Tile 1



Tile 2



Tile 3





# What Shall We Parallelize?

In other words,

**What should one thread do?**

**One answer:**

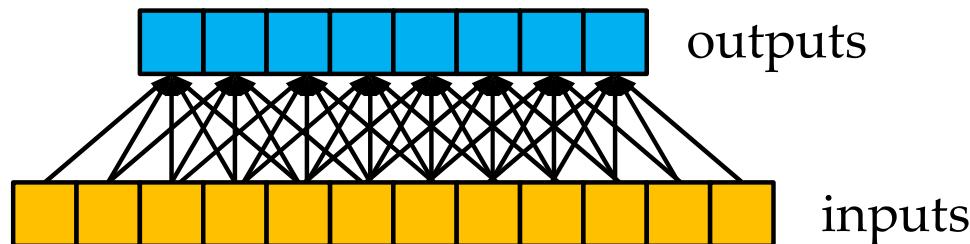
- (same as with vector sum and matrix multiply)
- **compute an output element!**

# Should We Use Shared Memory?

In other words,

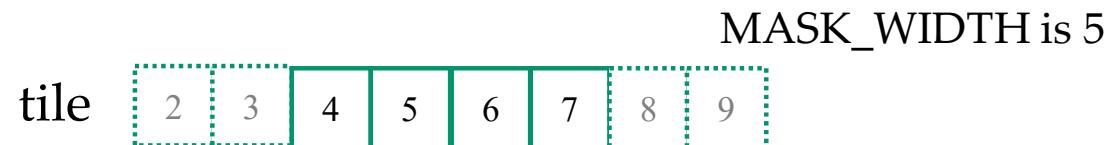
**Can we reuse data read from global memory?**

Let's look at the computation again...



Reuse reduces global memory bandwidth,  
so **let's use shared memory**.

# How Much Reuse is Possible?

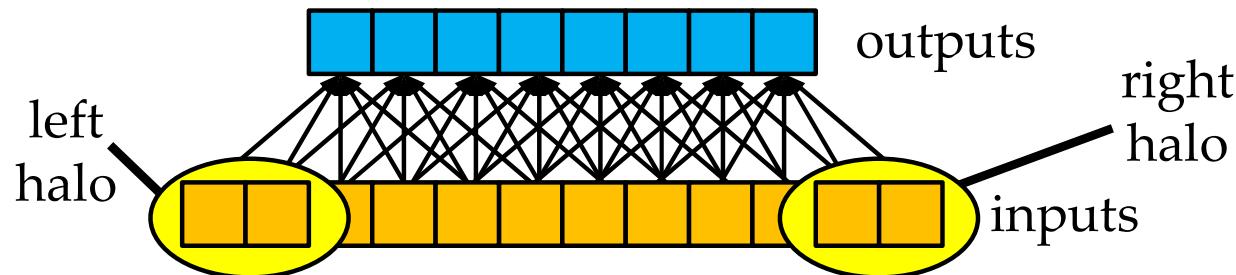


- Element 2 is used by thread 4 (1×)
- Element 3 is used by threads 4, 5 (2×)
- Element 4 is used by threads 4, 5, 6 (3×)
- Element 5 is used by threads 4, 5, 6, 7 (4×)
- Element 6 is used by threads 4, 5, 6, 7 (4×)
- Element 7 is used by threads 5, 6, 7 (3×)
- Element 8 is used by threads 6, 7 (2×)
- Element 9 is used by thread 7 (1×)

# What About the Halos?

In other words,

**Do we also copy halos into shared memory?**



Let's **consider both** possible answers.

# Can Access Halo from Global Mem.

One answer: no,

- threads **read halo values**
- directly **from global memory**.

Advantage:

- optimize reuse of shared memory
- (halo reuse is smaller).

Disadvantages:

- **Branch divergence!** (shared vs. global reads)
- Halo **too narrow to fill** a memory **burst**

A really bad idea on early GPUs, but later GPUs offer larger last-level caches (shared by all SMs), making performance competitive.

```

__global__
void convolution_1D_tiled_cache_kernel(float *N, float *P, int Width)
{
    __shared__ float tile[TILE_WIDTH];

    int This_tile_start_point = blockIdx.x * blockDim.x;
    int i = This_tile_start_point + threadIdx.x;

    tile[threadIdx.x] = N[i]; // boundary checking is missing here

    __syncthreads();

    int radius = MASK_WIDTH / 2;
    int N_start_point = i - radius;

    float Pvalue = 0;
    for (int j = 0; j < MASK_WIDTH; j++) {
        int N_index = N_start_point + j;
        if ((N_index >= 0) && (N_index < Width)) {
            int tile_index = N_index - This_tile_start_point;
            if ((tile_index >= 0) && (tile_index < blockDim.x))
                Pvalue += tile[tile_index] * Mc[j];
            else
                Pvalue += N[N_index] * Mc[j];
        }
    }
    P[i] = Pvalue;
}

```

1D convolution with  
halos read from  
global memory

# Can Load Halo to Shared Mem.

Better answer: yes,

**load halos to shared memory.**

Advantages:

- **Coalesce global memory accesses.**
- **No branch divergence during computation.**

Let's write  
the code!

Disadvantages:

- Some threads must do >1 load, so  
**some branch divergence** in reading data.
- Slightly more shared memory needed.

# Allocate and Initialize Variables

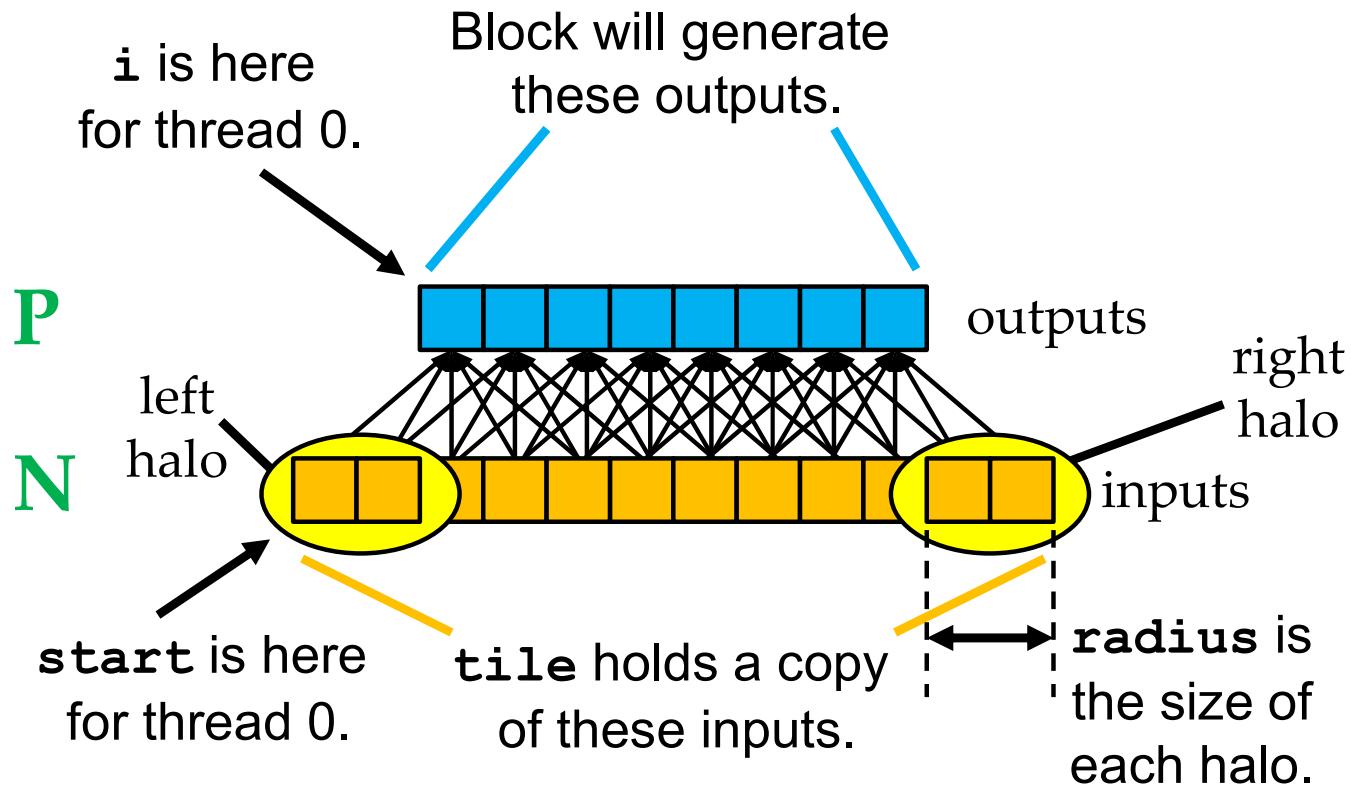
```
__global__ void convolution_1D_tiled_kernel
    (float *N, float *P, int Width)
{
    // shared tile with space for both halos
    __shared__ float tile[TILE_SIZE + MASK_WIDTH - 1];

    int radius = MASK_WIDTH / 2; // a useful constant

    // this thread's index into output P
    int i      = blockIdx.x * blockDim.x + threadIdx.x;

    // this thread's starting element of N
    int start  = i - radius;
```

# Variable Meanings for a Block



# Load the Input Data

```
if (0 <= start && Width > start) { // all threads
    tile[threadIdx.x] = N[start];
} else {
    tile[threadIdx.x] = 0.0f;
}
if (MASK_WIDTH - 1 > threadIdx.x) { // some threads
    start += TILE_SIZE;
    if (Width > start) {
        tile[threadIdx.x + TILE_SIZE] = N[start];
    } else {
        tile[threadIdx.x + TILE_SIZE] = 0.0f;
    }
}
__syncthreads(); // OUTSIDE of if's
```

# And Compute an Output Element

```
if (i < Width) {    // only threads computing outputs

    float Pvalue = 0;          // running sum

    // compute output element
    for (int j = 0; MASK_WIDTH > j; j++) {
        Pvalue += tile[threadIdx.x + j] * Mc[j];
    }

    // write to P
    P[i] = Pvalue;
}
```

```

__global__
void convolution_1D_tiled_cache_kernel(float *N, float *P, int Width) {
    __shared__ float tile[TILE_SIZE + MASK_WIDTH - 1];

    int radius = MASK_WIDTH / 2;
    int i = blockIdx.x * blockDim.x + threadIdx.x;
    int start = i - radius;

    if (0 <= start && Width > start) {           // all threads
        tile[threadIdx.x] = N[start];
    } else {
        tile[threadIdx.x] = 0.0f;
    }
    if (MASK_WIDTH - 1 > threadIdx.x) {           // some threads
        start += TILE_SIZE;
        if (Width > start) {
            tile[threadIdx.x + TILE_SIZE] = N[start];
        } else
            tile[threadIdx.x + TILE_SIZE] = 0.0f;
    }
    __syncthreads();
    if (i < Width) {
        float Pvalue = 0.0f;
        for (int j = 0; MASK_WIDTH > j; j++) {
            Pvalue += tile[threadIdx.x + j] * Mc[j];
        }
        P[i] = Pvalue;
    }
}

```

1D convolution with  
halos read into  
shared memory  
(output parallelism)

# Review: What Shall We Parallelize?

In other words,

**What should one thread do?**

**One answer:**

- (same as with vector sum and matrix multiply)
- **compute an output element!**

**Is that our only choice?**

# Parallelize Loading of a Tile

Alternately,

- **each thread loads** one input element, and
- **some threads compute** an output.

(compared with previous approach)

Advantage:

- **No branch divergence for load** (high latency).
- **Avoid narrow global access** ( $2 \times$  halo width).

Disadvantage:

- Branch **divergence for compute** (low latency).

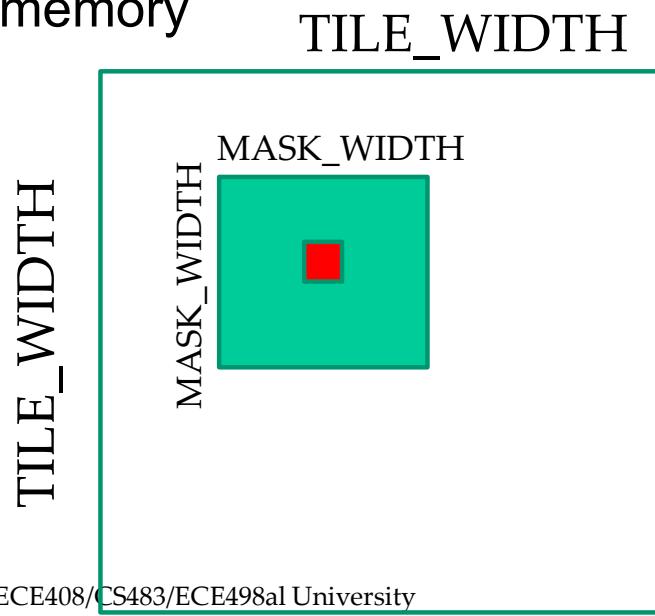
## 2D Example of Loading Parallelization

Let's do an example for 2D convolution.

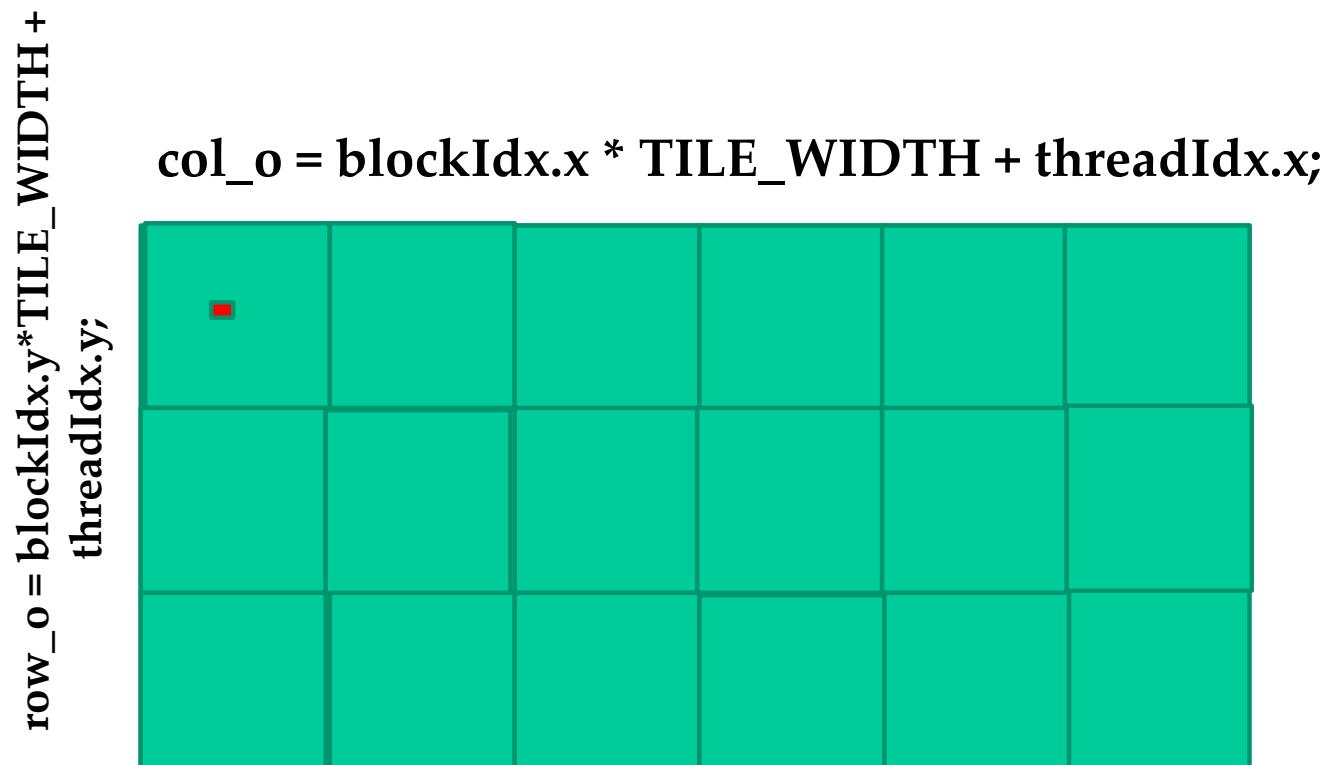
- Thread block matches input tile size.
- Each thread loads one element of input tile.
- Some threads do not participate in calculating output,

# Parallelizing Tile Loading

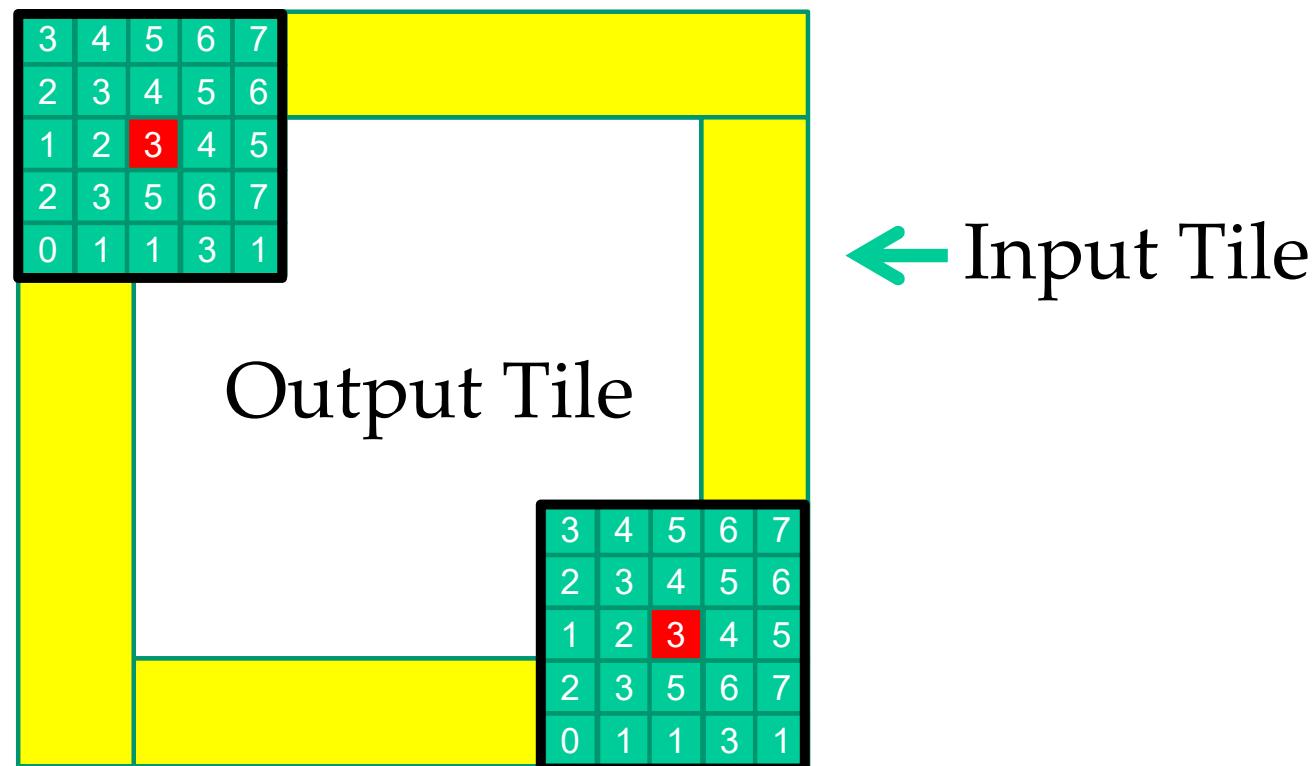
- Load a tile of  $N$  into shared memory
  - All threads participate in loading
  - A subset of threads then use each  $N$  element in shared memory



# Output Tiles Still Cover the Output!



Input tiles need to be larger than output tiles.



# Setting Block Width

```
dim3 dimBlock(TILE_WIDTH+4,TILE_WIDTH+4, 1);
```

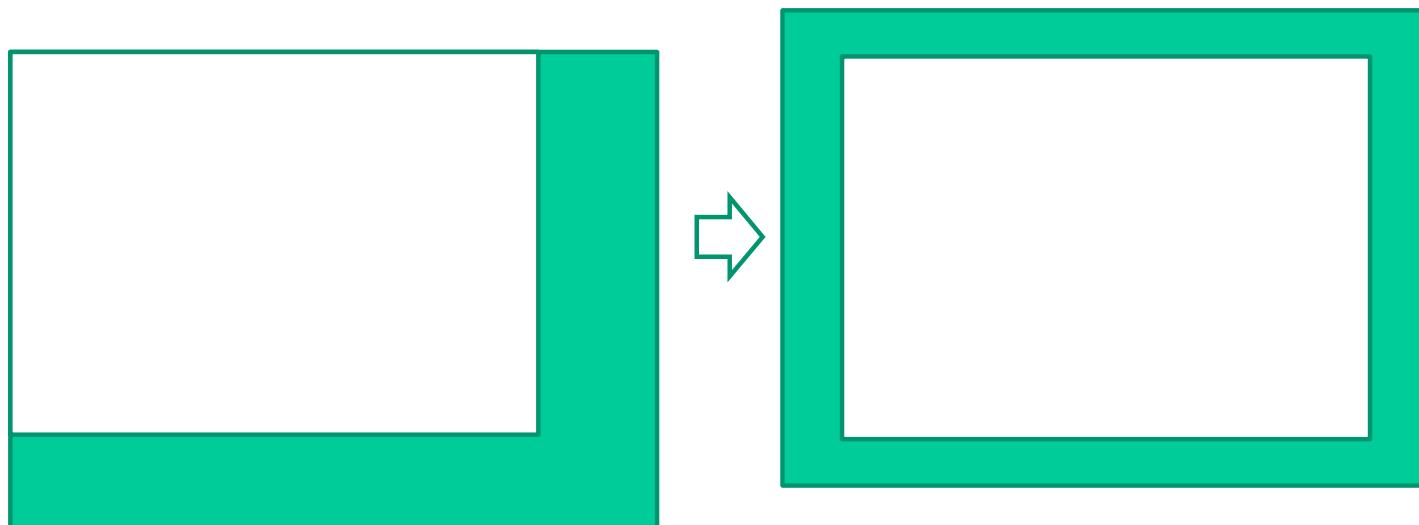
In general, block width should be

$$\text{TILE\_WIDTH} + (\text{MASK\_WIDTH} - 1)$$

```
Dim3 dimGrid(ceil(Width/(1.0*TILE_WIDTH)),  
             ceil(Width/(1.0*TILE_WIDTH)), 1)
```

There need to be enough thread blocks  
to generate all P elements.

# Shifting from output coordinates to input coordinates

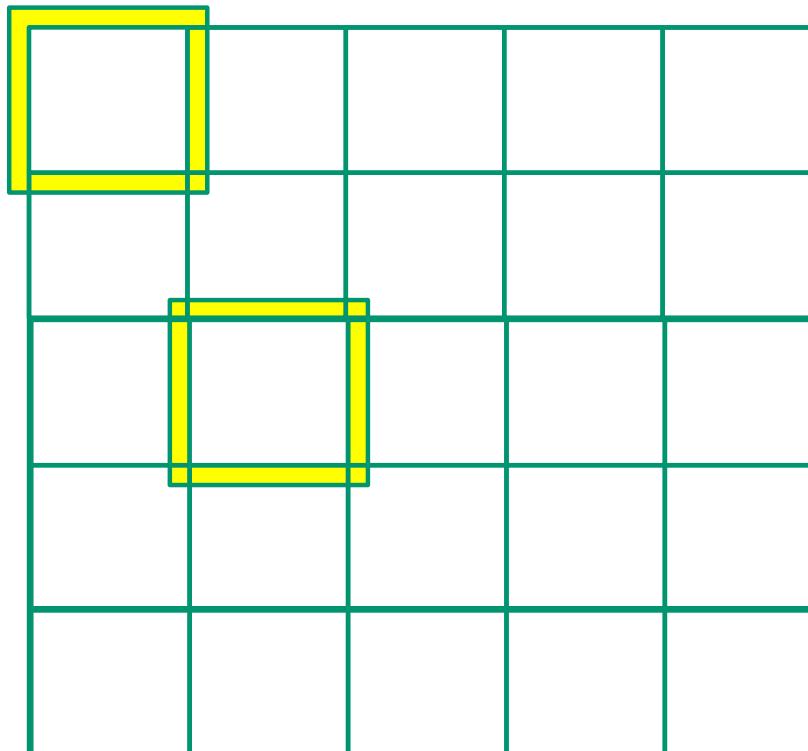


# Shifting from output coordinates to input coordinates

```
int tx = threadIdx.x;
int ty = threadIdx.y;
int row_o =
    blockIdx.y * TILE_WIDTH + ty;
int col_o =
    blockIdx.x * TILE_WIDTH + tx;

int row_i = row_o-2; // MASK_WIDTH / 2
int col_i = col_o-2; // (radius in
                     // prev. code)
```

Threads that loads halos outside N should return 0.0



# Taking Care of Boundaries

```
float Pvalue = 0.0f;  
if((row_i >= 0) && (row_i < Width) &&  
    (col_i >= 0) && (col_i < Width)) {  
    tile[ty][tx] =  
        N[row_i*Width + col_i];  
} else {  
    tile[ty][tx] = 0.0f;  
}  
__sync_threads () ; // wait for tile
```

# Not All Threads Calculate Output

```
if(ty < TILE_WIDTH && tx <TILE_WIDTH) {  
    for(i = 0; i < 5; i++) {  
        for(j = 0; j < 5; j++) {  
            Pvalue += Mc[i][j] *  
                tile[i+ty][j+tx];  
        }  
    }  
    // if continues on next page
```

# Not All Threads Write Output

```
    if(row_o < Width && col_o < Width)
        P[row_o * Width + col_o] = Pvalue;
    }
} // end of if selecting output
// tile threads
```

# Which Strategy Will You Choose?

- We recommend the input parallelization strategy for MP4.
- Alternatively, you may instead choose to parallelize based on either of the two output parallelization strategies
  - reading halo values from global memory, or
  - Keeping the halos in shared memory and issuing more than one load from some threads.

**YOU MUST NOTE YOUR STRATEGY  
AT THE START OF YOUR `lab4.cu` FILE!**



**QUESTIONS?**

**READ CHAPTER 7!**

# Problem Solving

Let's think a little more deeply about the tradeoffs with the parallelization strategies. Assuming that...

- we use input parallelism
- with **16×16** thread blocks (**256** threads/block)
- And a **5×5** mask (**MASK\_WIDTH** of **5**),

**Q: how many threads compute output elements?**

**A: TILE\_WIDTH is  $16 - 5 + 1 = 12$ , so  $12 \times 12 = 144$  threads.**

# Problem Solving

Assuming that...

- we use input parallelism
- with  $16 \times 16$  thread blocks (256 threads/block)
- And a  $5 \times 5$  mask (MASK\_WIDTH of 5),
- TILE\_WIDTH is 12, so 144 threads compute output elements.

**Q: How many of the 8 warps are active during computation?**

Here was our condition for computation:

```
if (ty < TILE_WIDTH && tx < TILE_WIDTH) {
```

**A: Not as bad as it might seem: warps 6 and 7 have  $ty \geq 12$  (TILE\_WIDTH is 12), so they do nothing during computation.**

# Problem Solving

Assuming input parallelism,  $16 \times 16$  thread blocks, a  $5 \times 5$  mask, `TILE_WIDTH` of 12, and 144 threads computing output elements, warps 6 and 7 do nothing during computation.

Here was our condition for computation:

```
if (ty < TILE_WIDTH && tx < TILE_WIDTH) {
```

**Q: What happens with warps 0 through 5?**

**Warps 0 through 5 have 24 active threads,**

- so **wasting some computation resources**,
- but since we **operate out of shared memory**,
- having **less parallelism** during computation is **probably ok**.

(We could use a more complex mapping to reduce to 5 warps.)

# Problem Solving

**For larger masks, input parallelism may be less attractive:**

- with the same **16×16** thread blocks (**256** threads/block) and
- a **9×9** mask (**MASK\_WIDTH** of **9**),
- **TILE\_WIDTH** is  **$16 - 9 + 1 = 8$** ,
- so output computation requires  **$8 \times 8 = 64$**  threads, or **2** warps,
- but **4** warps are active (with **16** active threads each).

Even if we remap to use **2** warps,

- the other **6** warps still count against warp resources,
- so each SM uses only **1/4 of possible warps for computation.**

# Problem Solving

Using output parallelism with a large mask

- improves computation utilization
- at the expense of loading more data.

For the same **16×16** thread blocks (**256** threads/block)

- and **9×9** mask (**MASK\_WIDTH** of **9**),
- output parallelism gives an **input tile of  $16 + 9 - 1 = 24$** , which
- requires **up to three loads from global memory per thread**.
- But **all threads** are **active during computation** of outputs.

# Problem Solving

We might also **compute multiple elements per thread**.

A **24×24** input tile (floats) requires only **2.25 kB**.

Using the same **16×16** thread blocks,

- we might instead **compute 4 elements per thread**,
- using a **40×40 input tile** (floats),
- which **requires 6.25 kB**,
- allowing **10 thread blocks on an SM**  
with 64kB of shared memory
- (limit on threads will be more restrictive).

# Problem Solving

## Which approach is best?

As you can see, the **answer** often

- **depends on** both the **parameters of the problem and**
- the **resources available** on the GPU
- (SMs, shared memory, caches, threads and thread blocks per SM, and so forth).