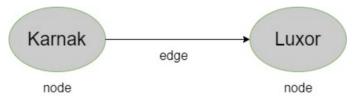
Comp 363 - Design and Analysis of Computer Algorithms

Spring Semester 2020 - Week 13 - Part 1

Dr Nick Hayward

graphs - intro recap

- graph data structure in computer science
- a way to model a given set of connections
- commonly use a graph to model patterns and connections for a given problem
- e.g. connections may infer relationships within data
- graph includes nodes and edges
- help us define such connections
- e.g. we have two nodes with a single edge

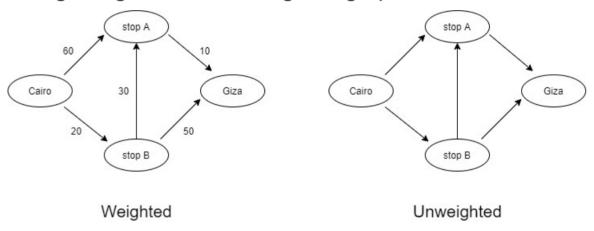


Graph Nodes and Edge

- each node may be connected to many other nodes in the graph
 - commonly referenced as neighbour nodes

graphs - Dijkstra's algorithm - consideration of terminology - part 1

- key concept for working with Dijkstra's algorithm is the association of values,
- e.g. numbers for each edge in the graph
- valyes are the weights assigned to the edge in the graph
- when we assign weights to an edge
- creating a weighted graph
- if we do not assign weights to edges
- defining an unweighted graph
- e.g. weighted and unweighted graphs



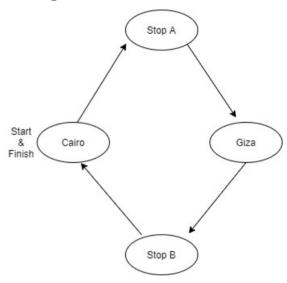
Graphs - Weighted and Unweighted

graphs - Dijkstra's algorithm - consideration of terminology - part 2

- choice of weighted versus unweighted
- · also affecta our choice of algorithm
- and the context of its usage
- e.g. if we want to calculate the shortest path in an unweighted graph
- we may use the breadth-first search algorithm
- to perform a similar calculation for a weighted graph
 - we may use Dijkstra's algorithm...

graphs - Dijkstra's algorithm - graph cycles - part 1

- may also encounter a graph with cycles
- i.e. we can cycle from Stop A back around to Stop A
- e.g.



Graph - Cycle

we may start and finish at the same node in the graph

graphs - Dijkstra's algorithm - graph cycles - part 2

- if we consider a graph with a cycle segment
- need to calculate shortest path between two defined nodes
- for most calculations commonly choose a path that avoids the cycle
- cycle will usually add greater weight to the calculation
- if we then follow cycle more than once
- simply adding extra weight to calculation for each completed cycle
- if we consider an undirected graph
 - now working with a cycle
- connected nodes in an undirected graph point to each other
 - effectively a cycle
- each edge will add another cycle to an undirected graph
- Dijkstra's algorithm only works with directed acyclic graphs (DAGs)

graphs - Dijkstra's algorithm - shortest path - part 1

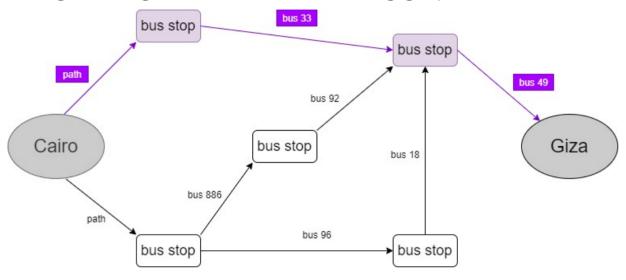
- common requirement for working with graphs
- a consideration of weighted and unweighted edges
- with weighted graphs
- interested in options for assigning more or less weight
- i.e. to edges within the graph
- Dijkstra's algorithm helps us work with queries for paths in our graphs
- e.g. we might need to answer the question

which path is the shortest to a node?

- e.g. which is the shortest path to node A?
- this path may not be fastest
- · but it will be the shortest in the graph
- shortest because it will include least number of edges between nodes

graphs - Dijkstra's algorithm - shortest path - part 2

e.g. we might consider the following graph for the shortest route

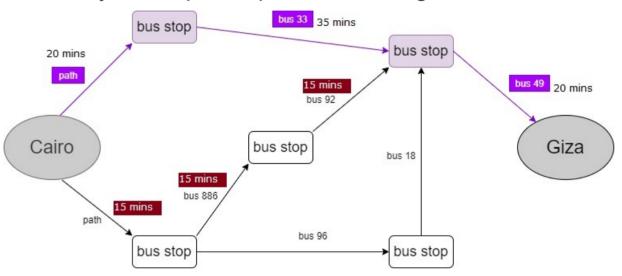


Graph Routes - shortest

- clearly see shortest path with least number of segments, three
- may use breadth-first search to find path with fewest segments

graphs - Dijkstra's algorithm - shortest path - part 3

- if we then added travel times to edges for competing routes
- i.e. costs for each edge
- we may find a quicker path for traveling



Graph Routes - shortest

- e.g. we can see time difference between purple and red paths
- from Cairo to Giza
- red path is faster, in spite of one more segment
- to find the fastest path
- may use a different option to breadth-first search
- we may use *Dijkstra's algorithm* to help us query graph for fastest path

Video - Algorithms and Data Structures

graphs - Dijkstra's algorithm - part 1

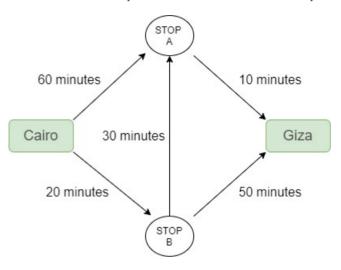


Graphs - Dijkstra's algorithm - intro - UP TO 1:01

Source - Dijkstra's algorithm - YouTube

graphs - Dijkstra's algorithm - working example 1 - part 1

- if we consider the following basic weighted graph
- may initially see how Dijkstra's algorithm works
- i.e. to help us find the fastest path



Dijkstra - example graph

- each segment in this graph has a corresponding cost
- time in minutes
- may use cost with Dijkstra's algorithm
- calculate shortest possible time from Cairo to Giza...

- may initially use the following steps with Dijkstra's algorithm
- i.e. to calculate fastest path from a defined start to finish in the graph
- for the current start node
- 1. identify the cheapest node
- i.e. next node we can reach in least amount of time
- simply check neighbour nodes for current node
- identify path with shortest time
- e.g. from our starting point of Cairo there are two options
 - either Stop A with a time of 60 minutes
- or Stop B with a time of 20 minutes
- we don't know values of other nodes at this point of search
- as we don't yet know how long it will take to get to finish
 - Giza defined with an overall time of infinity...

- we know closest node from start, Cairo
- Stop B with a time of 20 minutes

node	time	
Stop A	60 minutes	
Stop B	20 minutes	
Giza	infinity	

- 2. update any costs for neighbours of this node
- need to calculate all times from Stop B to available neighbour nodes
- in current example
- includes Stop A and our finish node of Giza
- may now update our times from start, Cairo
- update to each node currently known in the graph

node	time
Stop A	50 minutes
Stop B	20 minutes
Giza	70 minutes

- first improvement is a faster time from Cairo to Stop A
- even though we have to go through node Stop B
- with current known neighbour nodes
- may also follow a path from start node, Cairo,
- follow to finish in Giza
- path takes 70 minutes
- currently have a shorter path from Cairo to Stop A
- plus a shorter path to finish from start...

- 3. repeat this pattern for each node in the graph
- may now repeat pattern to check for other neighbour nodes
- potentially faster routes from start to finish...
- repeat first step again
- need to find next node with shortest travel time
- checked all of the neighbour nodes for Stop B
- we can now check next fastest neighbour of start node Cairo
- in current example
- this will be node Stop A...

- don't need to update time from start node
 - i.e. from Cairo to node Stop A
 - already identified a faster route...
- we may check times for quickest route to finish
- now have an extra path to check
- i.e. from Stop A to finish in Giza
- gives us a shorter time from start to finish
- due to its time of 10 minuutes...

graphs - Dijkstra's algorithm - working example 1 - part 8

now update our times as follows

node	time	
Stop A	50 minutes	
Stop B	20 minutes	
Giza	60 minutes	

- i.e. define fastest routes for following paths
 - Cairo to Stop A = 50 minutes
 - Cairo to Stop B = 20 minutes
 - Cairo to Giza = 60 minutes
- able to identify a quicker path from start to Stop A
- and a quicker path from start to finish...

- calculate the time for the final path
- may currently define final path
- calculated fastest time of 60 minutes...
- if we compare this calculation with a search using breadth-first
- may see that it would not have found that path
- breadth-first would have found shorter path
- but a slower path in this example graph...

- in current example, we may see an initial benefit
- benefit relative to this context
- e.g. for a search with Dijkstra's algorithm compared with a breadth-first search
- may use breadth-first search to find shortest path
- e.g. between defined nodes in graph
- may use Dijkstra's algorithm to assign weights to graph
- use to find path with smallest total of calculated weights...

Video - Algorithms and Data Structures

graphs - Dijkstra's algorithm - part 2

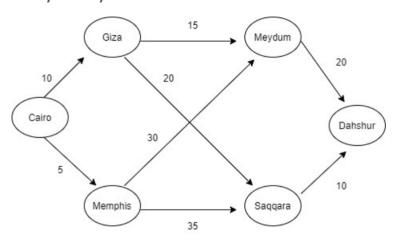


Graphs - Dijkstra's algorithm - example outline - UP TO 3:57

Source - Dijkstra's algorithm - YouTube

graphs - Dijkstra's algorithm - working example 2 - part 1

- consider another working example for a graph with weighted edges
- e.g. a graph with values for cost to travel from one node to another
- perhaps from Cairo to Giza or Giza to Saggara...



Graph Weighted

- in this graph
- define weights for associated costs of travel along each edge
- i.e. travel from *Memphis* to *Meydum* for 30
- or, perhaps, from Giza to Meydum for only 15...

graphs - Dijkstra's algorithm - working example 2 - part 2

- if we consider this graph
- may need to calculate cheapest route from Cairo
- e.g. start point to an end point of Dahshur...
- may use Dijkstra's algorithm to perform this calculation
- follow defined four steps for this algorithm
- initial costings may be defined as follows

parent	node	cost
Cairo	Giza	10
Cairo	Memphis	5
N/A	Meydum	infinity
N/A	Saqqara	infinity
N/A	Dahshur	infinity

and set an initial parent for each node...

- then update this table as we execute the algorithm
- 1. start by finding cheapest node
- in this graph, from a starting node of Cairo cheapest edge is 5 to Memphis
- we can't make this initial path any cheaper
- cheapest node will be Cairo to Memphis
- 2. then, calculate cost to neighbours of this cheapest node, i.e. from Memphis
- we now have costs for Meydum, 30, and Saqqara, 35
- we can update our table of costs from our starting point to each neighbour
- Cairo to Meydum and Cairo to Saggara
- Cairo -> Memphis -> Meydum
- Cairo -> Memphis -> Saggara

parent	node	cost
Cairo	Giza	10
Cairo	Memphis	5
Memphis	Meydum	35
Memphis	Saqqara	40
N/A	Dahshur	infinity

- now have costs for Meydum and Saqqara
- may define their costs as we travel through Memphis node
- as we can see in the table
- their parent node may also be updated to Memphis...

- may now repeat these two steps for next cheapest node from Cairo
- i.e. Giza at a cost of 10
- update its values in the table as well

parent	node	cost
Cairo	Giza	10
Cairo	Memphis	5
Giza	Meydum	25
Giza	Saqqara	30
N/A	Dahshur	infinity

- the costs for travel from starting point, Cairo, to Meydum and Saqqara updated
 - they are cheaper now
 - so we update the values in the table
- i.e. it's now cheaper to travel to Meydum and Saqqara via Giza...

- we may check cost to travel to end point, Dahshur
- check cheapest node from Giza
- currently Meydum at 15
- may update its neighbours
 - gives us an initial cost for Dahshur of 20
- if we update table at this point
 - we get the following travel cost from Cairo to Dahshur

parent	node	cost
Cairo	Giza	10
Cairo	Memphis	5
Giza	Meydum	25
Giza	Saqqara	30
Meydum	Dahshur	45

- we finally have an initial travel cost from start to finish of 45
 - i.e. from Cairo -> Giza -> Meydum -> Dahshur

- we may also check next cheapest node from Giza, Saqqara
- then, we may travel from Saqqara to Dahshur
- for a total of 40 from Cairo
- i.e. may now update cost of travel from start to finish
- update to a lower overall cost of 40

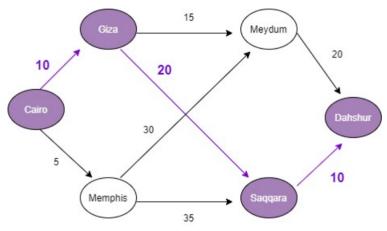
parent	node	cost
Cairo	Giza	10
Cairo	Memphis	5
Giza	Meydum	25
Giza	Saqqara	30
Saqqara	Dahshur	40

parent	node	cost
Cairo	Giza	10
Cairo	Memphis	5
Giza	Meydum	25
Giza	Saqqara	30
Saqqara	Dahshur	40

- we may see that shortest path costs 40
- using this overall cost
 - now define path for travel from start to finish in this graph
- to help with this path definition
- may check parent node set in last table
- i.e. ended up with Saggara as parent for end point Dahshur

graphs - Dijkstra's algorithm - working example 2 - part 9

- we know that we need to travel from Saqqara to Dahshur
- may then follow path to parent of Saqqara, set to Giza
- i.e. to travel to Saqqara we need to begin at Giza
- we follow Giza back to its parent
- starting point at Cairo
- now have a complete route for traveling from start point to end point in least cost, 40



Graph Weighted - Final Costs

Video - Algorithms and Data Structures

graphs - Dijkstra's algorithm - part 3



Graphs - Dijkstra's algorithm - example usage - UP TO 8:48

Source - Dijkstra's algorithm - YouTube

Resources

various

- Dijkstra's algorithm
- Graph abstract data type

videos

Dijkstra's algorithm