

# Comp 363 - Design and Analysis of Computer Algorithms

---

Spring Semester 2020 - Week 6

Dr Nick Hayward

# Algorithms and Data Structures

---

## recursion and the call stack - part 7

- an example of tail recursion for 3! - factorial(3)

```
def factor(x, tail):  
    print("factor x =",x)  
    if x == 1:  
        print("return from (x == 1) = 1")  
        return tail  
    else:  
        print("x =",x)  
        return factor(x - 1, x * tail)  
  
# set initial tail to 1  
print(factor(3, 1))
```

# Algorithms and Data Structures

---

## *recursion and the call stack - part 8*

- pseudocode outline for pattern of execution for tail recursion
- e.g.  $3!$  - factorial(3)

```
factor(3, 1)
  x = 3, tail = 1
  return factor(3 - 1, 3 * 1) // recurse 1
factor(2, 3)
  x = 2, tail = 3
  return factor(2 - 1, 2 * 3) // recurse 2
factor(1, 6)
  x = 1, tail = 6
  return 6 // pop factor(1, 6) from call stack
return 6 // pop factor(2, 3) from call stack
return 6 // pop factor(3, 1) from call stack

print 6 // stack now clear, execution ends
```

## Video - Algorithms and Data Structures

---

*Recursion, the Call Stack, and Overflow...part II*

!!Con 2019- Tail Call Optimization: The Musical!! by Anjana ...



!!Con 2019- Tail Call Optimization: The Musical!! -  
UP TO 7:58

Source - !!Con 2019- Tail Call Optimization: The  
Musical!! - YouTube

# Algorithms and Data Structures

---

## *recursion and the call stack - part 9*

- stack may be used to represent execution logic
  - *e.g. for a recursive function in JavaScript*
- code example uses a *call stack* to ensure expected execution

```
function findSolution(target) {  
  function find(current, history) {  
    if (current == target) {  
      return history;  
    } else if (current > target) {  
      return null;  
    } else {  
      return find(current + 5, `${history} + 5`) || find(current * 3, `${history} * 3`);  
    }  
  }  
  return find(1, "1");  
}  
  
console.log(findSolution(24));
```

# Algorithms and Data Structures

---

## *recursion and the call stack - part 10*

- initial `findSolution()` function is called
  - *passed parameter of 24 is the value to check*
- function returns an executed `find()` function
  - *initial test values for current and history*
- part of this function's execution
  - *checks initial values until it reaches `else` part of conditional statement*
- returns `find()` function
  - *called recursively*
  - *initially checking against addition of 5*
  - *continues to check possible values with ``+ 5``*
  - *either succeeds or moves onto right side of logical OR, `||`, \*logical OR checks with ``* 3``*
- it will either succeed or fail with these recursive checks
- structure that permits this recursion to execute
  - *structure is the call stack*
- *call stack* provides a defined pattern to execution
  - *pattern allows the code to run as expected*

## Video - Algorithms and Data Structures

---

*Recursion for Fun - part 2*

Recursion - Part 7 of Functional Programming in JavaScript



Recursion and Fun - JavaScript - UP TO 14:46

Source - Recursion and Fun - JavaScript - YouTube

# Algorithms and Data Structures

---

## *stack operations - part 1*

- we may define a stack as a simple list of elements
  - *may be accessed from only one end*
  - *known as the top of the stack*
- i.e. refer to data structure as *last in, first out*
- known limitation of this structure
  - *lack of access to elements not at top of stack*
- a simple difference between structures
  - *i.e. basic list or array and specific stack*
- to access the bottom element
  - *all elements above must first be popped*



# Algorithms and Data Structures

---

## *stack operations - part 2*

- stack operations are simple, e.g.
  - *add elements to the top*
  - *pop elements from the top*
- also means specific restrictions must be in place
  - *ensures only these operations are allowed*
  - *i.e to define usage for the data structure*
  - *if not, it ceases to be a stack*

# Algorithms and Data Structures

---

## *stack operations - part 3*

- complementary operations are commonly available
  - *may vary relative to language implementation*
- e.g. a stack may permit the following
  - *view the top element in the stack*
  - *this is not pop - element is not removed from stack*
  - *operation known as peeking*
  - *clear operation will remove all elements from stack*
  - *Length property returns number of elements in stack*
  - *empty returns whether stack has any values or not*

## Video - Algorithms and Data Structures

---

*Recursion, the Call Stack, and Overflow...part III*

!!Con 2019- Tail Call Optimization: The Musical!! by Anjana ...



!!Con 2019- Tail Call Optimization: The Musical!! -  
UP TO END

Source - !!Con 2019- Tail Call Optimization: The  
Musical!! - YouTube

# Algorithms and Data Structures

---

## *example implementations*

- choose various existing data structures to define custom stack
- e.g. might use an array or list
- choice will often depend on support in chosen programming language
- e.g. in JS - common option is an array object
- define constructor for stack object
  - *then extend prototype for custom properties and methods*

# Algorithms and Data Structures

---

## *stack constructor*

- initial constructor is as follows

```
// CONSTRUCTOR = Stack object
function Stack() {
  /* define instance properties for stack
   * - empty array for instantiated stack
   * - options might include max length, restricted data type &c.
   */
  this.store = [];
}
```

- instantiate a basic Stack object
  - *simply defining an empty array*
  - *use array as store for Stack data structure*
- constructor may be updated to include
  - *type checks and restrictions*
  - *initial values for Stack*
  - *required access context*
  - *...*

# Algorithms and Data Structures

---

## *extend the prototype*

- initially extend Prototype for this object
- add required functionality for a basic stack
- e.g. define functions for the following
  - *add data*
  - *delete data*
  - *get size of Stack*

# Algorithms and Data Structures

---

## *prototype - add data*

- *add data* function
  - *add passed data to top of stack*

```
// PROTOTYPE - add method for value pushed to top of stack
Stack.prototype.add = function (value) {
  this.store.push(value);
  console.log(`value added = ${value}`);
}
```

- underlying store object is an array for Stack
  - *use default push() method to add required data*

# Algorithms and Data Structures

---

## *prototype - delete data*

- *delete data* function defined as follows

```
Stack.prototype.delete = function () {  
  const deletedValue = this.store.pop();  
  console.log(`last value deleted = ${deletedValue}`);  
}
```

- same as *add data*
  - use default *pop()* method for store array
  - added to custom *Stack* data structure



# Algorithms and Data Structures

---

## *prototype - size of Stack*

- also define `size()` function for Stack
  - *use built-in Array property for Length*

```
Stack.prototype.size = function () {  
  const size = this.store.length;  
  console.log(`store size = ${size}`);  
}
```

# Algorithms and Data Structures

---

## *prototype - peek Stack*

- useful option is *peeking* at the top of Stack
- e.g.

```
Stack.prototype.peek = function () {  
    const peekValue = this.store[(this.store.length-1)]  
    console.log(`top value = ${peekValue}`);  
}
```

- function will return copy of top value
  - *will not delete item from Stack and underlying store array*

# Algorithms and Data Structures

---

## *prototype - clear stack*

- common operation for a Stack is to clear all entries,
  - *yet preserve the Stack itself*
- i.e. resetting store array for instantiated Stack object
- e.g.

```
Stack.prototype.clear = function () {  
    // resets Stack's array store - clears all items  
    this.store = [];  
}
```

# Algorithms and Data Structures

---

## *prototype - check empty stack - part 1*

- check an instantiated Stack object for entries
  - *i.e. determine if stack is empty or not*

```
Stack.prototype.empty = function () {  
  if (this.store.length === 0) {  
    return true;  
  } else {  
    return false;  
  }  
}
```

# Algorithms and Data Structures

---

## *prototype - check empty stack - part 2*

- conditional logic has been placed in this function
  - *i.e. not passed down chain of logic to requesting application call*
  - *means function is self-contained*
- function returns valid response regardless of execution context
- as we develop Stack's Prototype methods
  - *add further restrictions and controls*
  - *clearly defines how to use this data structure*
- also define what and how may be returned
  - *custom data structure customised to context, usage...*

## Video - Algorithms and Data Structures

---

*Prototype in JavaScript - part 1*

Prototypes in JavaScript - FunFunFunction #16



Prototype in JavaScript - UP TO 1:00

Source - Prototypes in JavaScript - YouTube

## Video - Algorithms and Data Structures

---

*Prototype in JavaScript - part 2*

Prototypes in JavaScript - FunFunFunction #16



Prototype in JavaScript - UP TO 6:41

Source - Prototypes in JavaScript - YouTube

# Algorithms and Data Structures

---

## *control access to the stack - part 1*

- Stack object and methods
  - *now working as expected*
- Stack is still open to mis-use
  - *due to array object in the Stack*
- restrict and control access to this Stack data structure
  - *e.g. using a Proxy*



# Algorithms and Data Structures

---

## *control access to the stack - part 2*

- to use a Proxy with our Stack constructor
  - *define a custom construct trap*
- may also use Reflect API to define defaults for handlers
- construct trap intercepts calls
  - *i.e. to defined new operator for a given constructor*

# Algorithms and Data Structures

---

## *control access to the stack - part 3*

- define initial Proxy wrapper for passed constructor

```
/*
 * PROXY
 */
function proxyConstruct(constructor) {

  const handler = {
    construct(constructor, args) {
      console.log('proxy constructor...');
      // const stack = Reflect.construct(constructor, args);
      return new constructor(...args);
    }
  };

  return new Proxy(constructor, handler);
}
```

# Algorithms and Data Structures

---

## *control access to the stack - part 4*

- then pass basic Stack constructor to the proxy

```
// proxy wrapper for Stack constructor
const proxiedStack = new proxyConstruct(Stack);
// instantiate proxied Stack & check store...
console.log(new proxiedStack().store);
```

# Algorithms and Data Structures

---

## *control access to the stack - part 5*

- instantiation of a proxied Stack object
  - *allows us to wrap constructor for Stack*
- may still use prototype methods for instantiated Stack object
- benefit of using a proxy for the constructor
  - *control of initial object instantiation*
- i.e. if object cannot be instantiated
  - *access to Prototype methods becomes irrelevant*

## Video - Algorithms and Data Structures

---

*Proxy in JavaScript*



Using a JavaScript Proxy - UP TO 3:28

Source - Proxy in JavaScript - YouTube

# Algorithms and Data Structures

---

## *Fibonacci*

- fun way to test recursion and stacks (i.e. call stack)
  - *problem of searching Fibonacci series of numbers*
- Fibonacci series is simply an ordered sequence of numbers
  - *each number is the sum of the preceding two...*
- e.g.

```
[0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89]
```

- might also see the series beginning with 1 instead of 0
- function should return n-th entry in sequence.
  - *e.g. 5th index entry will return 5*
- Fibonacci may be solved using various techniques and algorithms
  - *e.g. iteration and recursion...*
- a good test of runtime speed and complexity

# Algorithms and Data Structures

---

## *Fibonacci - iteration example*

- initially test an iterative solution
  - *check and return values in the Fibonacci series*
- e.g.

```
function fib(n) {  
  // pre-populate array - allow calculation with two initial values  
  const result = [0, 1];  
  // i starts at index 2...  
  for (let i = 2; i <= n; i++) {  
    // get the previous two results in array  
    const a = result[i-1];  
    const b = result[i-2];  
    // calculate next value in series & push to result array  
    result.push(a + b);  
  }  
  // get result at specified index posn in series...  
  return result[n-1]; // -1 due to array index starting at 0...  
}  
// log to console...  
console.log('index posn 8 in fibonacci series = ', fib(8));
```

- $O(n)$  - linear time for iteration
  - *assuming constraints of memory for 64bit system*
- beyond memory bounds and complexity becomes quadratic
  - $O(n^2)$  or  $O(n^2)$

# Algorithms and Data Structures

---

## *Fibonacci - recursion example - part 1*

- also consider a solution using *recursion*
- e.g.

```
function fib(n) {  
  // base case  
  if (n < 2) {  
    console.log(n);  
    return n;  
  }  
  // dynamic calculation of number in sequence  
  return fib(n-1) + fib(n-2);  
}  
  
console.log('index posn 5 in fibonacci series = ', fib(5));
```



# Algorithms and Data Structures

---

## *Fibonacci - recursion example - part 2*

- add some logging for this recursion
  - *e.g.*

```
function fib(n, r) {  
  console.log(`n = ${n} and r = ${r}`);  
  // base case  
  if (n < 2) {  
    console.log(n);  
    return n;  
  }  
  // dynamic calculation of number `n` in sequence and recursive call `r`...  
  return fib(n-1, 1) + fib(n-2, 2);  
}  
  
console.log('index posn 5 in fibonacci series = ', fib(5, 0));
```

# Algorithms and Data Structures

---

## *Fibonacci - recursion example - part 3*

- sample output to help track recursive calls and addition
  - *e.g.*

```
n = 5 and r = 0
n = 4 and r = 1
n = 3 and r = 1
n = 2 and r = 1
n = 1 and r = 1
return base = 1
n = 0 and r = 2
return base = 0
n = 1 and r = 2
return base = 1
n = 2 and r = 2
n = 1 and r = 1
return base = 1
n = 0 and r = 2
return base = 0
n = 3 and r = 2
n = 2 and r = 1
n = 1 and r = 1
return base = 1
n = 0 and r = 2
return base = 0
n = 1 and r = 2
return base = 1
index posn 5 in fibonacci series = 5
```

# Algorithms and Data Structures

---

## *Fibonacci - recursion example - part 4*

- recursive pattern may be defined as follows

```
fib(5)
  n = 5
  return fib(5-1) + fib(5-2) // recurse
fib(5-1)
  n = 4
  return fib(4-1) + fib(4-2) // recurse
fib(4-1)
  n = 3
  return fib(3-1) + fib(3-2) // recurse
fib(3-1)
  n = 2
  return fib(2-1) + fib (2-2) // recurse
fib(2-1)
  n = 1
  return 1 // base returned - recurse
fib(2-2)
  n = 0
  return 0 // base returned - recurse
fib(3-2)
  n = 1
  return 1 // base returned - recurse
fib(4-2)
  n = 2
  return fib(2-1) + fib(2-2) // recurse
fib(2-1)
  n = 1
  return 1 // base returned
fib(2-2)
  n = 0
  return 0 // base returned
fib(5-2)
  n = 3
  return fib(3-1) + fib(3-2) // recurse
fib(3-1)
  n = 2
  return fib(2-1) + fib(2-2) // recurse
fib(2-1)
  n = 1
  return 1 // base returned
fib(2-2)
  n = 0
  return 0 // base returned
```

```
fib(3-2)
  n = 1
  return 1 // base returned
```

```
return 5 // sum return values for base
```

- follow pattern of recursion and base case returns
  - *shows return values needed to calculate index position 5 in Fibonacci series*
  - *i.e.*

```
// Fibonacci series to index 5
[0,1,1,2,3,5]
```

## Video - Algorithms and Data Structures

---

### *Recursion and Fibonacci*

Algorithms: Recursion



Recursion - UP TO 4:30

Source - Recursion & Fibonacci - YouTube

# Algorithms and Data Structures

---

## *Fibonacci - recursion example - part 5*

- why does this JavaScript recursive solution actually work as expected?
- as function is called recursively
  - *only returns value for base case*
  - *i.e. either 0 or 1*
- as it continues down from value of passed n-th position in series,
  - *it is storing each return*
  - *then returns total for that position in series*
- for current JavaScript example
  - *may consider execution of functions to better understand pattern*
- e.g. function where another function is called
  - *paused whilst inner execution is completed*
  - *i.e. outer will be paused as inner is executed...*

# Algorithms and Data Structures

---

## *Fibonacci - recursion example - part 6*

- recursive solution will produce an *exponential time* for the complexity
- i.e. as n-th value increases
  - *so will time required to find a value in the series...*
- commonly define complexity for a recursive solution as exponential
  - $O(2^n)$
- improvements may be made to this recursive algorithm
  - *e.g. using memoisation*
- due to repetitive calls to same values for fib()
  - *e.g. multiple calls to fib(3)*

## Video - Algorithms and Data Structures

---

*memoisation - part 1*

Algorithms: Memoization and Dynamic Programming



What is Memoisation - UP TO 2:51

Source - Memoisation - YouTube



# Algorithms and Data Structures

---

## *Fibonacci - memoisation - part 1*

- store arguments of a given function call along with computed result
- e.g. when `fib(4)` is first called
  - *computed value will be stored in memory*
  - *a temporary cache in effect*
- then call stored return
  - *e.g. each and every subsequent call to `fib(4)`*

# Algorithms and Data Structures

---

## *Fibonacci - memoisation - part 2*

- now improve performance of recursive algorithm
  - *e.g. for Fibonacci series*
  - *add support for memoisation*
- abstract functionality to a separate, re-usable *memoisation* function
- then use this function to add memoisation to an algorithm, application &c.
- main part of function
  - *records used and repeated functions &c.*
  - *then call again as needed*
- i.e. a *cache* for the function
- derive speed improvement for passed function

## Video - Algorithms and Data Structures

---

*memoisation - part 2*

Algorithms: Memoization and Dynamic Programming



Memoisation and Complexity - UP TO 4:12

Source - Memoisation - YouTube

# Algorithms and Data Structures

---

## *Fibonacci - memoisation - part 3*

e.g. define initial *memoisation* function

```
// pass original function - e.g. slow recursive fibonacci function
function memoise(fn) {
  // temporary store
  const cache = {};
  // return anonymous function - use spread operator to allow variant no. args
  return function(...args) {
    // check passed args in cache - if true, return cached args...
    if (cache[args]) {
      return cache[args];
    }
    // no cached args - call passed fn with args
    const result = fn.apply(this, args);
    // add result for args to the cache
    cache[args] = result;
    // return the result...
    return result;
  };
}
```

# Algorithms and Data Structures

---

## *Fibonacci - memoisation - part 4*

- use memoisation with Fibonacci function

```
function fib(n) {  
  // base case  
  if (n < 2) {  
    console.log(n);  
    return n;  
  }  
  // dynamic calculation of number in sequence  
  return fib(n-1) + fib(n-2);  
}  
  
// reassign memoised fib fn to fib - recursion then calls memoised fib fn...  
fib = memoise(fib);  
console.log('index posn 100 in fibonacci series = ', fib(100));
```

- now able to check higher index values in Fibonacci series
  - *without previous memory issues...*
- e.g. 100th position in the Fibonacci series is,
  - 354224848179262000000
- position 1000 = 4.346655768693743e+208

## Video - Algorithms and Data Structures

---

*Recursion and Fibonacci - memoisation*

Algorithms: Recursion



Recursion - UP TO END

Source - Recursion & Fibonacci - YouTube

# Algorithms and Data Structures

---

## *divide and conquer - intro*

- algorithms and development - often trying to solve a problem in a given context
- many techniques we may consider to solve a problem
  - *might start with a common option to help us get started...*
- *Divide and conquer* is a general technique
  - *e.g. used to solve various problems in application development and data usage*
- *Divide and conquer* is a well known *recursive* technique for solving various problems
  - *e.g. an option for analysing and solving such problems*
- consider use of *divide and conquer* from different perspectives
  - *use various examples to help outline its general usage...*

## Video - Algorithms and Data Structures

---

*Recursion & Divide and Conquer - part 1*

Divide & Conquer (Think Like a Programmer)



Recursion and Divide and Conquer - UP TO 4:08

Source - Divide and Conquer - YouTube



# Algorithms and Data Structures

---

## *divide and conquer - part 1*

- start with a common example problem
  - *helps define basic structure and usage of divide and conquer*
- e.g. consider a parcel (plot or lot) of land
  - *need to sub-divide it evenly into square plots*
  - *need these plots of land to be as large as possible*
  - *fit all of the available space in original parcel of land*
- land has been measured to the following size
  - *1680 feet by 640 feet*
  - *approximately same as 6.74 Jumbo Jet planes in length*
  - *or 560 yd ( a decent length par 5 in golf)*
- n.b. to solve this problem effectively
  - *can't simply divide this land in half - not two even squares*
  - *nor 20x20 squares, which are too small...*
- need to ensure we can always find maximum size for a square
  - *then divide the specified parcel of land...*

# Algorithms and Data Structures

---

## *divide and conquer - part 2*

- how do we calculate largest square
  - *i.e. largest used for a defined parcel of land*
- we may use *divide and conquer* to help solve this problem
- divide and conquer is a recursive technique
- divide and conquer algorithms are recursive algorithms...
- begin by defining two initial steps for the algorithm
  - *define base case - should be as simple as possible*
  - *divide and decrease underlying problem until it is base case*

# Algorithms and Data Structures

---

## *divide and conquer - part 3*

- consider the base case:
  - *begin by considering and defining base case for this algorithm*
  - *e.g.*

*What is the largest possible square we may use to divide the land?*

- easiest base case for this type of problem might be as follows
  - *i.e. if one side was a multiple of the other side...*
- e.g. simple box of 50x50, which may be divided as two boxes of 25x25
  - *largest box we may use is 25x25*
  - *this meets requirement for defined base case as well...*

# Algorithms and Data Structures

---

## *divide and conquer - part 4*

- consider the recursive case:
- once we've defined a base case for the problem
  - *need to consider an appropriate recursive case to achieve base case*
- *divide and conquer* proves useful
  - *effectively reducing the problem to meet the base case*
- divide and conquer states - for each recursive call you need to reduce the problem
- for our land - 1680 feet by 640 feet
  - *begin by marking largest boxes we may use to divide this size*
- e.g. two boxes of 640x640 and one remaining box of 640x400

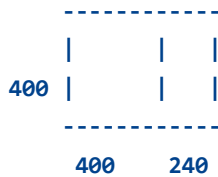


- still have land measuring 640x400 to divide
- division of this land may now follow same underlying pattern as original land size
  - *i.e. find largest box to fill this remaining land of 640x400*
- when we find this size
  - *define largest box for overall land of 1680x640*
- problem has now been reduced from a land size of 1680x640 to 640x400

# Algorithms and Data Structures

## *divide and conquer - part 5*

- now apply same algorithm to this problem - a land size of  $640 \times 400$
- largest box we may define is  $400 \times 400$
- still some land remaining after this division,  $400 \times 240$



- continue to apply this algorithm & reduce the problem as follows



- and then



- finally arrive at the base case...



- now have two evenly sized boxes of  $80 \times 80$  with no land left over
  - *i.e. 80 is a factor of 160*
- we may sub-divide original land of  $1680 \times 640$  into even plots of  $80 \times 80$

## Video - Algorithms and Data Structures

---

*divide and conquer - Euclid's algorithm*



Euclid and the Greatest Common Divisor - UP TO 8:27

Source - Algorithms - YouTube

# Algorithms and Data Structures

---

## *divide and conquer - part 6*

- we may summarise this use of divide and conquer as follows
  - *define a simple case for the base case*
  - *define how to reduce the problem to reach the base case*
- *divide and conquer*
  - *not itself an algorithm or reductive solution*
  - *can't apply as is to solve a given problem...*
- *divide and conquer*
  - *give us a clear way of thinking about a problem to reach a solution*

# Algorithms and Data Structures

---

## *divide and conquer - part 7*

- e.g. if we consider following problem
  - *we may clearly see how useful this approach can be to defining an algorithm*
- e.g. for a defined data structure
  - *[6, 9, 13, 5, 11, 16]*
- need to add all of the values and return the total
- might simply use a loop to sum these values

```
def sum(data):  
    total = 0  
    for x in data:  
        total += x  
    return total  
  
print(sum([6, 9, 13, 5, 11, 16]))
```



# Algorithms and Data Structures

---

## *divide and conquer - part 8*

- might define a solution using recursion for the same array of values
- e.g. define following steps to create a recursive algorithm to solve this problem
- **Step 1 - define base case**
  - *i.e. what's simplest array we may sum*
  - *e.g. an array of size 1 or 0 may be passed to the `sum()` function*
  - *this is easy to sum...base case*
- **Step 2 - recursive calls**
  - *need to reduce problem with each recursive call*
  - *i.e. move closer to defined base case*

## Video - Algorithms and Data Structures

---

*Recursion & Divide and Conquer - part 2*

Divide & Conquer (Think Like a Programmer)



Recursion and Divide and Conquer - UP TO 7:53

Source - Divide and Conquer - YouTube

# Algorithms and Data Structures

---

## *divide and conquer - part 9*

- begin by considering how to sum values in a passed array
- e.g.

```
sum([6, 9, 13, 5, 11, 16])
```

- actually the same as

```
6 + sum([9, 13, 5, 11, 16])
```

- both examples return same summed value
- n.b. second example has started to reduce size of passed array
  - *now reduced size of problem*

# Algorithms and Data Structures

---

## *divide and conquer - part 10*

- define this algorithm as follows
  - *get the passed data*
    - e.g. array of numbers
  - *if data is empty*
    - return zero
  - *else total equals*
    - first number + rest of data
- check expected output as follows

```
sum([6, 9, 13, 5, 11, 16]) - sum = `60`  
  6 + sum([9, 13, 5, 11, 16]) 6 + 54 = return `60`  
    9 + sum([13, 5, 11, 16]) 9 + 45 = return `54`  
      13 + sum([5, 11, 16]) 13 + 32 = return `45`  
        5 + sum([11, 16]) - 5 + 27 = return `32`  
          11 + sum([16]) - 11 + 16 = return `27`  
            sum([16]) - base case & first return from execution - return `16`  
  
print 60
```

# Algorithms and Data Structures

---

## *divide and conquer - part 11*

- now implement `sum()` function using divide and conquer with recursion

```
def sum(data):  
    if len(data) == 1:  
        return data[0]  
    else:  
        return data[0] + sum(data[1:])  
  
print(sum([6, 9, 13, 5, 11, 16]))
```

# Algorithms and Data Structures

---

## *divide and conquer - part 12*

### ■ JavaScript example 1

```
function sum(data) {  
  if (data.length === 1) {  
    return data[0];  
  } else {  
    // slice - return array from index 1 to end...  
    return data[0] + sum(data.slice(1));  
  }  
}  
  
console.log(sum([6, 9, 13, 5, 11, 16]))
```

### ■ JavaScript example 2

```
function sum(data) {  
  if (data.length === 1) {  
    return data[0];  
  } else {  
    // destructure data - get head and return rest  
    const [head, ...rest] = data;  
    return head + sum(rest);  
  }  
}  
  
console.log(`sum of values = ${sum([6, 9, 13, 5, 11, 16])}`);
```

# Resources

---

## *JavaScript*

- MDN - Array
- MDN - Prototype
- MDN - Proxy
- Prototypes in JavaScript - YouTube
- Proxy in JavaScript - YouTube

## *Python*

- Stacks - YouTube

## *Various*

- !!Con 2019- Tail Call Optimization: The Musical!! - YouTube
- Algorithms - YouTube
- Divide and Conquer - YouTube
- Event Driven Architecture - YouTube
- Memoisation - YouTube
- Memory Manager - YouTube
- Recursion & Fibonacci - YouTube
- Recursion and Fun - JavaScript - YouTube
- Recursion and the Call Stack - Java - YouTube