Comp 460 - Algorithms & Complexity

Spring Semester 2020 - Week 9 Dr Nick Hayward

DEV Week Assessment

Course total = 15%

- continue development of application
- built from scratch
- continue design and development of initial project outline and design
- working app (as close as possible...)
- NO blogs, to-do lists, note-taking...
- ,,,
- outline research conducted
- describe data chosen for application
- define algorithms and data structures used in app
- why choose these options?
- how have they been used?
- define current performance &c.?
- define testing of implementation & usage
- show any prototypes, patterns, and designs

DEV Week Demo

DEV week assessment will include the following:

- brief presentation or demonstration of current project work
- ~5 to 10 minutes per group
- analysis of work conducted so far
- o e.g. during semester & DEV week
- presentation and demonstration
- outline current state of app
- explain what works & does not work
- show implemented designs since project outline & mockup
- show latest designs and updates
- due Tuesday 17th March 2020 @ 4.15pm

Video - Algorithms and Data Structures

quicksort - part 3



Quicksort - UP TO 8:42

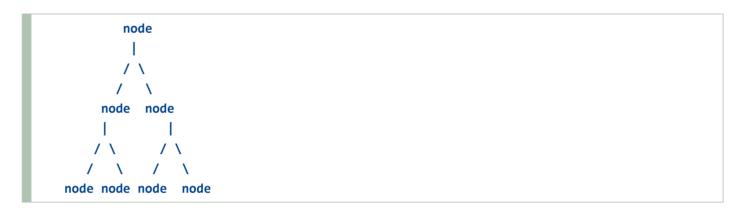
Source - Quicksort - Java - YouTube

binary search tree - intro - part 1

- binary search tree (BST) is a binary tree
- each node has a Comparable key (and an associated value)
- satisfies a defined restriction
- e.g. key in any node is
- larger than the keys in all nodes in that node's left subtree
- smaller than the keys in all nodes in that node's right subtree
- comparison is context specific relative to the current node
- binary search need to work with sorted data sets
- when we add or update the data
- need to re-sort list before using binary search
- if we're working with sorted lists of data
- · e.g. an array of books
- quickly encounter a problem
 - list may need to be updated item in the array is deleted
 - then need to add a value to index where last deleted item stored...

binary search tree - intro - part 2

- may now update list to meet specific criteria
- removing need to repeatedly sort dataset
- binary search tree data structure
 - basic tree data structure



Video

Trees and parsing - part 1

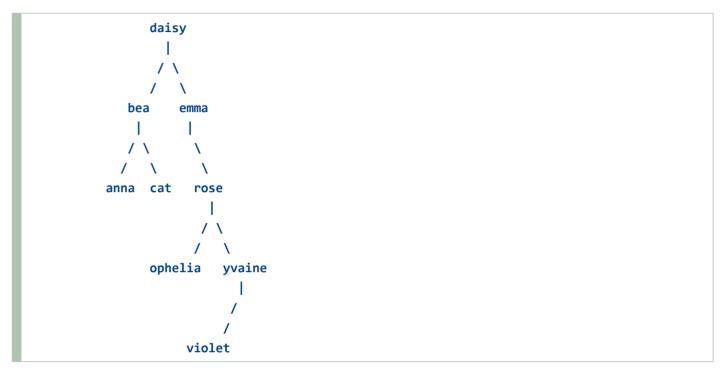


How the browser renders a website - UP TO 7:40

Source - So how does the browser actually render a website - YouTube

binary search tree - intro - part 3

sample binary search tree may be structured as follows



- for every node in this tree
 - nodes to left of current node are smaller
 - nodes to right of current node are larger
- e.g. search for violet begin at root then use following path
- V is after D traverse right side of tree
 - current node = emma
- V is after E continue down right side
- current node = rose
- V is before Y continue down left side
- violet node found...
- searching for a node in a binary search tree takes O(log n) on average
- O(n) for worst cases
- sorted array, by contrast, takes O(log n) in worst case scenarios

- might initially consider arrays as preferable option
- sorted binary search tree is, on average, faster for insertions and deletions...

binary search tree - issues

- binary search trees do not provide random access
- performance times are averages
- rely on a balanced tree
- specialist trees may provide self-balancing mechanisms
- e.g. might use a *red-black* tree

Video - Algorithms and Data Structures

binary search trees - part 1



Binary Search Trees - UP TO 1:36

Source - Trees - Java - YouTube

binary search tree - basic logic implementation - part 1

- symbol-table API for a binary search tree
- common option for implementing this type of traversal and search
- symbol-table is also known as a map, dictionary, associative array &c.
- may vary by programming language
- general concept is as follows
- abstraction of key/value pairs
 - e.g. insert a value with a specified key
 - o given key, search for corresponding value
- sample usage may include the following

application	search	key	value
dictionary	search for a definition	word	definition
index	search for a given reference, e.g book page	term	e.g. list of page numbers for a book
compiler	search for props of variables	variable name	type and value

- for binary search tree logic may begin with custom function to define nodes
- function includes props required for a node, e.g.
 - key
- value
- left link
- right link
- node count

Video

Trees and rendering - part 2



A common working example - How the browser renders a website - UP TO 17:17

Source - So how does the browser actually render a website - YouTube

binary search tree - order-based methods - part 1

- common reason for working with binary search trees (BST)
- they keep the keys in order
- may use BSTs in many disparate API contexts
 - ensure consistent I/O structure
- e.g. might consider a custom symbol-table API
- some sample methods and usage
- min & max
 - o if left link of root is null
 - · smallest key in BST must be root node
 - if left link is not null
 - · smallest key is in subtree referenced by left link
 - · repeats for each left link in each subtree...
- floor
- if a given key is less than key at root of BST
- floor of key must now be added to left subtree...
- if key is greater than root
- o floor can be in right link but only if there is a smaller or equal existing key
- if not, floor of key is root...
- ceiling
 - same pattern as floor
 - except check relative to right link
- selection
- o e.g. seek key of rank k
- key such that precisely k other keys in BST are smaller
- if number of keys t in left subtree is larger than k
- look (recursively) for key of rank k in left subtree
- o if t is equal to k
- return key at root
- if t is smaller than k
- look (recursively) for key of rank k t 1 in right subtree

binary search tree - order-based methods - part 2

- range search
- to implement keys () method returns all keys in a given range
- begin with basic recursive BST traversal method known as inorder traversal
- e.g. to show this order traversal
- o first print all keys in left side of BST all less than root
- then print root key
- then print all keys in right side of BST
- keys() method
- define code to add each key that is in range to a *Queue*
- skip recursive calls for subtrees that cannot contain keys in range

rank

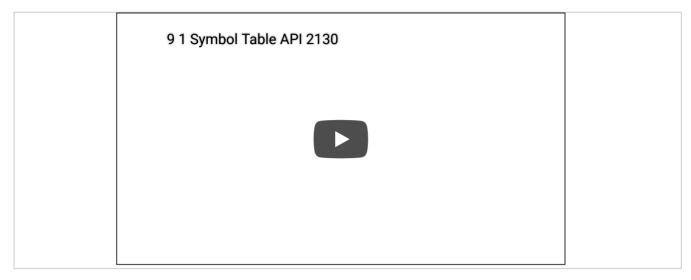
- if given key is equal to key at root
- o return number of keys t in left subtree
- if given key is less than key at root
- o return rank of key in left subtree
- if given key is larger than key at root
- o return t plus one (to count key at root) plus rank of key in right subtree

binary search tree - order-based methods - part 3

- delete min & max
- delete minimum
- go left until finding a node that has a null left link
- · then replace link to that node by its right link
- symmetric method works for delete maximum
- delete
- proceed in a similar manner to delete a node with one or null childrem
- for two or more start by replacing current node with its successor
- successor is node with smallest key in its right subtree
- accomplish task of replacing x by its successor in four easy steps
 - 1. save a link to node to be deleted in t
 - 2. set x to point to its successormin(t.right)
 - 3. set right link of x
 - supposed to point to BST containing all keys larger than x.key
 - ∘ to deleteMin(t.right)
 - the link to BST containing all keys that are larger than x.key after deletion
 - 4. set left link of x to t. Left
 - $\circ\;$ all keys that are less than both deleted key and its successor...

Video

symbol table API



Symbol table API - UP TO 5:40

Source - Symbol Table API - YouTube

binary search tree - usage - intro

- binary search tree (BST) has a non-linear insertion algorithm
- BST is similar in nature to a doubly-linked list
- a linked data structure
- includes set of sequentially linked records, commonly known as nodes
- each node defines three fields,
- link field previous node in sequence of nodes
- link field next node in sequence of nodes
- one data field
- link fields may be represented using a common example of a convoy, a train &c.
- e.g. linked ships sailing in a convoy...

binary search tree - usage - links and usage - part 1

- binary search tree defines its pointers as left and right to help indicate any duplication of logic &c.
- algorithm may be detected
- providing support for left and right traversal
- binary search tree node's pointers are typically called *left* and *right*
 - indicate subtrees of values relating to current value
- simple JavaScript implementation of such a node is as follows

```
const node = {
    value: 123,
    left: null,
    right: null
}
```

 BST is a unique tree due to its inherent ordering of nodes based on value

binary search tree - usage - links and usage - part 2

- any child nodes in a left subtree are always less than parent node's value
- converse holds for a right subtree
- values in subtree will always be greater
- e.g.

```
7
/\
5 9
/\
2 11
/\\
1 3 12
\
14
/
13
```

- traverse the tree and check key of current node
- if search key is less than current node's value
- follow left link
- otherwise, follow right link
- position of node values is based on a few factors
 - value of node
 - value of root
 - order of insertion
- e.g.
- root is set to 7
- 5 is less than root insert as left link
- 9 is greater than root insert as right link
- 2 is less than root follow left link
 - 2 is less than 5 left link is null insert as left link

• ...

repeat for additional inserts...

Video - Algorithms and Data Structures

binary search trees - part 2



Binary Search Trees - Insert - UP TO 3:00

Source - Trees - Java - YouTube

binary search tree - usage - search pattern

- clearly defined pattern searching a BST becomes easier to reduce to a repeatable pattern
- traverse left for a value less than current node, and right for a greater value
- simple pattern is helped by exclusion of duplicates
- in last example currently have a depth of 5 means furthest we
 need to travel is 5 nodes from root to find a value
- BSTs also have a natural sorted order
- due to the insertion algorithm
- makes BSTs particularly useful for quick searching
- eliminate options at each node

- a few options for implementing BSTs in JavaScript
- e.g. extend an object's prototype to include various custom functions to manage a BST
- a non-map implementation might include the following
- define initial data and props for constructor
- extend Prototype add custom methods
 - add
 - has
 - o delete
 - o size
- expose universal interface

binary search tree - usage - BST with JavaScript - part 2

initial coded example

```
* Constructor BST
*/
function BinarySearchTree() {
    // instantiated object - private prop - root default...
    this._root = null;
}
* Prototype
* -extend with custom functions
* - methods
*/
BinarySearchTree.prototype = {
   // extend - custom functions
    add: function(value) {
    },
    has: function(value) {
    },
    delete: function(value) {
    },
    size: function(
    }
};
```

- easiest to begin such custom methods with an outline of a has()
 method
- may define a general structure for querying a BST
- method defines single parameter for value
- returns true if value is found and false if null found
- logic follows a basic binary search algorithm
- to determine presence of a value

- coded example for has() method
- e.g. consider following initial outline for querying keys

```
has: function(value){
   const found = false,
        current = this._root;
   // check node is available for search...
   while(!found && current){
        // check node - if value less than current node's, go LEFT
        if (value < current.value){</pre>
            // update 'current' prop
            current = current.left;
        // check node - if value greater than current node's, go RIGHT
        } else if (value > current.value){
            // update 'current' prop
            current = current.right;
        //check node - values are equal, found node...
        } else {
            // update boolean...
            found = true;
    // return search status...
    return found;
},
```

- search starts at root
- initially checking there is a root key and key has not been found
- current is initially set to root to begin BST traversal
- sets while loop running
- allows following checks to be executed
- check search value against current node value
 - if less than set current to left link
- check search value again...
- o if greater than set current to right link
- otherwise value has been found
- found updated to true
- while loop broken
- contains now complete...

Video - Algorithms and Data Structures

binary search trees - part 3



Binary Search Trees - has/add - UP TO 7:24

Source - Trees - Java - YouTube

- may also use same underlying pattern for node insertion
- defining add() method
- need to modify search for a place to insert a node
- instead of returning an existing value

```
add: function(value){
   // define a new node - placeholder object & props...
    const node = {
            value: value,
            left: null,
            right: null
        },
        // variable for current node - use during BST traversal...
        current;
   // CHECK - no items yet in the BST
    if (this. root === null){
        // ROOT - BST empty - set root to current node -
        this._root = node;
       // update current prop - set to root node
        current = this._root;
       // TRAVERSE - begin traversal of BST from current node - start at root
        while(true){
            // check node - if value less than current node's, go LEFT
            if (value < current.value){</pre>
                // check node - if no node in left link
                if (current.left === null){
                    // update current prop - set `left` to new node
                    current.left = node;
                    // EXIT - node inserted as `left` link
                    break;
                } else {
                    // node set to existing left link...
                    current = current.left;
```

```
// check node - if value greater than current node's, go RIGHT
            } else if (value > current.value){
                // chevk node = if no node in right link
                if (current.right === null){
                   // update current prop - set `right` to new node
                    current.right = node;
                    // EXIT - node inserted as `right` link
                    break;
                } else {
                    // node set to existing right link...
                    current = current.right;
                }
           // if new value = current one - ignore
            } else {
                break;
            }
       }
   }
},
```

- to determine just size of BST
- traverse tree in any order
- may also need to flatten tree to an array, map &c.
- e.g. abstract a traverse() function to ensure ordered traversal

```
traverse: function(process){
   // inner scope helper function - pass node...call recursively
    function inOrder(node){
        // check node exists
        if (node) {
           // check node - if left link exists
            if (node.left !== null){
                // call recursively - pass current node from subtree - checks extent of
         subtree...
                inOrder(node.left);
            // call the passed process method on this node
            process.call(this, node);
            // traverse the right subtree
            if (node.right !== null){
                inOrder(node.right);
            }
        }
    }
    // define start node - pass root
    inOrder(this._root);
},
```

Video - Algorithms and Data Structures

binary search trees - part 4



Binary Search Trees - Traverse - UP TO 3:55

Source - Trees - Java - YouTube

- traverse() function defines single parameter, process
 - passed argument should be a function
 - may be executed on every node in tree
- function inOrder
- used to recursively traverse tree
- n.b. recursion only works when left and right links exist
- rule is designed to reduce reference to null nodes to a bare minimum
- traverse() function begins traversal from root
- passed process() function handles each node...

Video - Algorithms and Data Structures

binary search trees - part 5



Binary Search Trees - traversal function - UP TO 9:27

Source - Trees - Java - YouTube

- now use abstracted function, traverse()
- e.g. use with a custom size() function

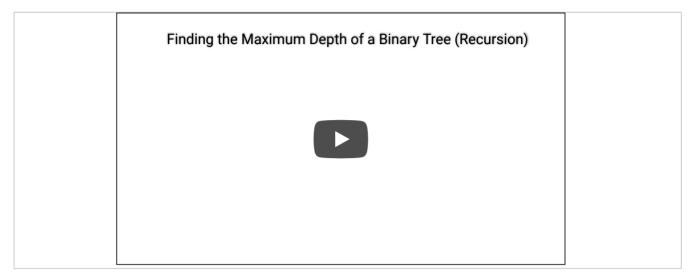
```
size: function(){
   const length = 0;

   this.traverse(function(node){
       length++;
   });

   return length;
},
```

- size() function calls above abstracted traverse() function
- passing a custom function for counting nodes

binary search trees - maximum depth

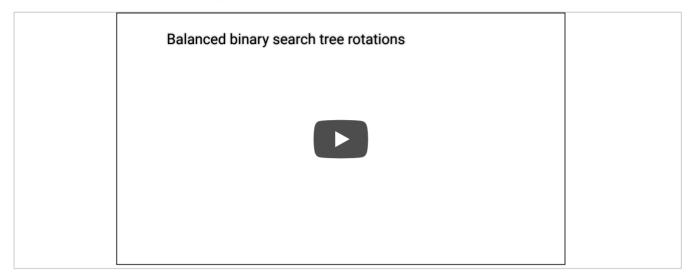


Maximum depth of binary tree with recursion - UP TO 2:34

Source - Trees - Max height using recursion - YouTube

- removing or deleting nodes from a BST can become complex
- due to necessary balancing of tree
- for each node removed
- need to check if it's root
- removal of root node is handled in a similar manner to other nodes
- except it will also need to be replaced
- simple matter of tree integrity
- may be handled as a special case in the logic

binary search trees - integrity and balance - part 1



Trees - Balancing - UP TO 1:22

Source - Trees - Balancing - YouTube

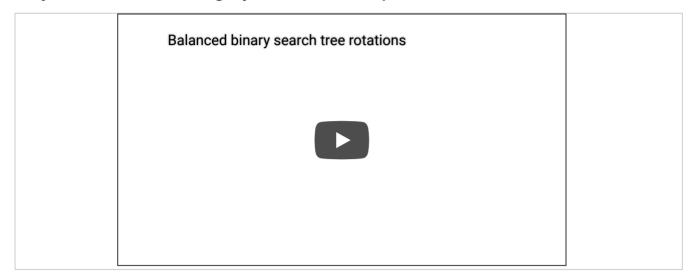
binary search tree - usage - BST with JavaScript - part 11 delete() - part 1

first part of a node's removal is checking it exists in defined BST

```
delete: function(value){
    let found = false,
        parent = null,
        current = this._root,
        childCount,
        replacement,
        replacementParent;
   // check node - if not found & node still exists
    while(!found && current){
        // check value - if less than current - traverse left
        if (value < current.value){</pre>
            parent = current;
            current = current.left;
        // check value - if greater than current - travers right
        else if (value > current.value){
            parent = current;
            current = current.right;
        // value found...
        else {
            found = true;
    }
   // continue - value found...
    if (found){
       // continue
    }
},
```

- may check required node using a standard binary search
- traverse left if value is less than current node
- traverse right if value is greater
- as part of traversal
- monitor parent node as passed node will need to be removed from its parent
- when requested node is found
 - current defines node to remove
 - initial part of defining a remove or delete option for binary search trees

binary search trees - integrity and balance - part 2



Trees - Balancing and Rotation - UP TO 4:21

Source - Trees - Balancing & Rotation-YouTube

binary search tree - usage - BST with JavaScript - part 13 delete() - part 2

- node in BST is foound
- consider options for removing node
- i.e. three applicable conditions we need to consider
 - a leaf node
 - a node with one child
 - a node with two children
- first two cases are easy to implement
- leaf node may simply be deleted from tree
- a child may replace parent leaf node upon deletion
- third condition is more involved in its modification of the tree
- need to determine if selected node has children, how many, and if it's root...
- if leaf node selected for deletion is root
- updated decision tree is relatively simple

binary search tree - usage - BST with JavaScript - part 14

for example,

```
// root node - special case
if (current === this._root){
   // check no. of child nodes - execute matching case
    switch(childCount){
        // no children - erase root
        case 0:
            this._root = null;
            break;
       // one child - child is now root
        case 1:
            this._root = (current.right === null ?
                            current.left : current.right);
            break;
       // two children -
        case 2:
            //T0D0
       // no default - one of above cases always matched...
    }
```

- for a root leaf node
 - deletion is simple to handle and implement

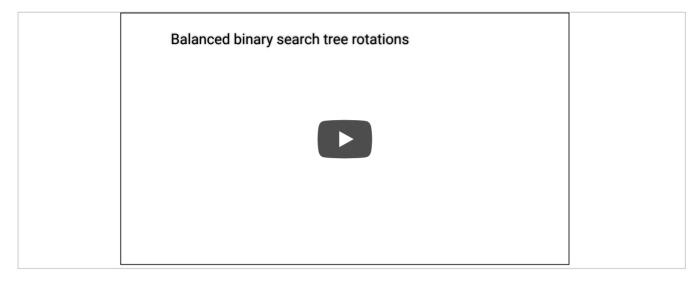
- for a child leaf node
 - · need to check and update tree
- e.g.
- value lower than parent
- o left pointer must be reset
- o reset to null (no children) or node's left child pointer
- value higher than parent
- right pointer must be reset
- o reset to null (no children) or node's right child pointer

- for a child node
 - initially check for 0 or 1 children of current selected node

```
switch (childCount){
   // no children - delete from tree
    case 0:// check delete value relative to parent
        if (current.value < parent.value){</pre>
            // value < parent - null parent's left pointer</pre>
            parent.left = null:
        } else {
            // else - null parent's right pointer
            parent.right = null;
        }
        break;
   // one child - replace deleted parent node
    case 1: // check value relative to parent
        if (current.value < parent.value){</pre>
            // value < parent - reset left pointer
            parent.left = (current.left === null ?
                             current.right : current.left);
        } else {
            // value > parent - reset right pointer
            parent.right = (current.left === null ?
                             current.right : current.left);
        break;
```

- need to update pointer on parent based on value of node to delete
- if deleted node's value was less than parent
 - reset left pointer either to null or to left pointer of deleted node
- if deleted node's value was greater than parent
 - need to reset right pointer

binary search trees - integrity and balance - part 3



Trees - Balancing and Rotation - UP TO 6:08

Source - Trees - Balancing & Rotation-YouTube

binary search tree - usage - BST with JavaScript - part 17 delete() - part 3

- most complex deletion is a node with two children
- e.g. consider following tree issue with deletion of node 2

```
8
/\
6 9
/\
2 11
/\\
1 4 12
/\\
3 5 14
/\
13
```

issue is how we now update tree based on child nodes of deleted node

- two common options to consider for such trees relative to deleted node
- in-order predecessor left child
- in-order successor left-most child of sright subtree
- may end up with either 1 or 3 replacing 2 in tree
- either option is acceptable
- i.e. may be used to update tree
- e.g. consider the following
- in-order predecessor = value before deleted value
 - o examine left subtree of deleted node and select right-most descendant
- in-order successor = value immediately after deleted value
- o examine right subtree of deleted node and select left-most descendant
- n.b. each option also requires traversal of tree to find required node

binary search trees - in-order traversal



Binary Search Trees - In-Order Traversal - UP TO 2:48

Source - Trees - In-Order Traversal - YouTube

binary search tree - usage - BST with JavaScript - part 19

for a root node with two children

```
/* root - two children
* - in-order predecessor
    - check left subtree
        - select right most descendant
              8
            / \
           6 9
          / \ \
          2
            7 11
                    12
          3 5
                     14
                     /
                    13
case 2: // e.g. delete root node - 8
       // check left subtree - get left of root (6)
   replacement = this._root.left;
   // check right-most child node - if not null
   while (replacement.right !== null){ // (7)
        replacementParent = replacement; // (6)
        replacement = replacement.right; // (7)
   }
   // check replacement parent
    if (replacementParent !== null){ // (6)
       // check for left node of replace - if exists, move to right of parent
        replacementParent.right = replacement.left; // (null)
       // new root - update with child nodes from existing root node
        replacement.right = this._root.right;
        replacement.left = this._root.left;
   } else {
       // new root - assign existing root's child nodes
        replacement.right = this._root.right;
```

```
}
// new root - UPDATE root value after deletion of root...
this._root = replacement;
```

- this example always looka for *in-order predecessor*
 - check left subtree
- select right most descendant

binary search tree - usage - BST with JavaScript - part 20

if previous tree was modified as follows

```
8
/\
6 9
/\\
2 7 11
/\\
1 4 12
/\\
3 5 14
/\
13
```

- if we deleted root node 8
- update the root to the node 7
- follows the pattern
- check left subtree = 6
- select right most descendant = 7

- if we deleted 6 node using *in-order predecessor*
- check left subtree, node 2
- then traverse following right pointers
- now replace deleted node 6 with node 5
- n.b. if we used *in-order successor*
- end up replacing deleted node in this tree with node 7

binary search trees - review of deletion



Trees - Review of Deletion - UP TO 7:56

Source - Trees - Deletion - YouTube

- for a child node with two children
 - add following case to existing switch statement

```
/* child node - two children
* - in-order predecessor
  - check left subtree
    - select right most descendant
*/
case 2:
   // two children - reset pointers for new traversal
    replacement = current.left;
    replacementParent = current;
   //find the right-most node
    while(replacement.right !== null){
        replacementParent = replacement;
        replacement = replacement.right;
    }
    replacementParent.right = replacement.left;
   // assign - children to replacement
    replacement.right = current.right;
    replacement.left = current.left;
   // add replacement to correct node in tree
    if (current.value < parent.value){</pre>
        // current < parent - add replacement to parent's left pointer
        parent.left = replacement;
    } else {
       // current > parent - add replacement to parent's right pointer
        parent.right = replacement;
    }
```

- either of these options will work to update tree
- e.g. initial tree may be represented as follows

```
initial BST = {
  "_root": {
    "value": 8,
    "left": {
      "value": 6,
     "left": {
        "value": 2,
        "left": {
          "value": 1,
          "left": null,
          "right": null
        },
        "right": {
          "value": 4,
          "left": {
            "value": 3,
            "left": null,
            "right": null
          },
          "right": {
            "value": 5,
            "left": null,
            "right": null
          }
        }
      },
      "right": {
        "value": 7,
        "left": null,
        "right": null
      }
    },
    "right": {
      "value": 9,
      "left": null,
      "right": {
        "value": 11,
        "left": null,
```

```
"right": {
    "value": 12,
    "left": null,
    "right": {
        "value": 14,
        "left": {
             "value": 13,
             "left": null,
             "right": null
             },
             "right": null
        }
    }
}
```

- now delete node 6
 - using in-order predecessor
- updated tree may be rendered as follows

```
after node deletion = {
  "_root": {
    "value": 8,
    "left": {
      "value": 5,
      "left": {
        "value": 2,
        "left": {
          "value": 1,
          "left": null,
          "right": null
        "right": {
          "value": 4,
          "left": {
            "value": 3,
            "left": null,
            "right": null
          },
          "right": null
        }
      },
      "right": {
        "value": 7,
        "left": null,
        "right": null
      }
    },
    "right": {
      "value": 9,
      "left": null,
      "right": {
        "value": 11,
        "left": null,
        "right": {
          "value": 12,
```

```
"left": null,
    "right": {
        "value": 14,
        "left": {
            "value": 13,
            "left": null,
            "right": null
        },
        "right": null
        }
    }
}
```

- n.b. by just using one option exclusively, we may end up with an unbalanced tree
- may consider modifying logic to ensure monitor of tree to maintain a selfbalancing search tree...

Resources

- Quicksort Java YouTube
- So how does the browser actually render a website YouTube
- Symbol Table API YouTube
- Trees Balancing YouTube
- Trees Deletion YouTube
- Trees In-Order Traversal YouTube
- Trees Java YouTube
- Trees Max height using recursion YouTube