Data Structures and Algorithms

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Session: Shortest Path Algorithm (Dijkstra's Algorithm)



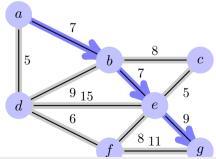
Finding Shortest Paths in Graphs

- 1. Shortest Path in Weighted Graphs
- 2. Shortest Path Properties in Weighted Graphs
- 3. Dijkstra's algorithm
- 4. Edge Relaxaton



Shortest Path in Weighted Graphs

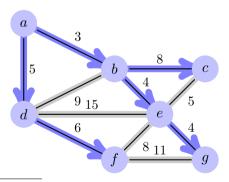
- 1. Given two vertices a and g in the weighted graph below, find a path of minimum total weight between them.
- 2. Length of path = sum of weight of constituent edges
- 3. Applications: Routing of vehicles, flights, internet packets





Shortest Path Properties in Weighted Graphs

- 1. **Optimal Substructure:** A subpath of a shortest path is also a shortest path
- 2. Tree of shortest paths exists 1 from a start vertex s to every other vertex



¹Illustrated below is the tree of shortest paths from a



Dijkstra's Algorithm and Shortest Path Properties

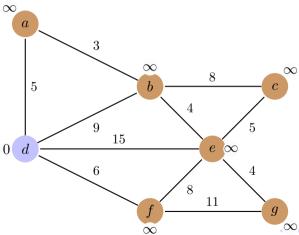
- 1. **Optimal Substructure:** A subpath of a shortest path is also a shortest path
- 2. Tree of shortest paths exist from a start vertex s to every other vertex
- 3. Dijkstra's Algorithm combines the two:
 - Store for each v a label d(v) = the distance of <math>v from s
 - Successively computes d(v), starting from the neighbors of s, assuming
 - graph is connected
 - edges are undirected
 - edge weights are nonnegative

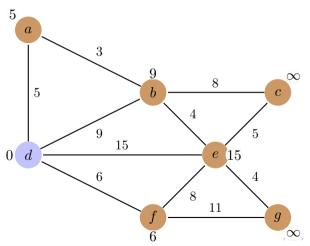


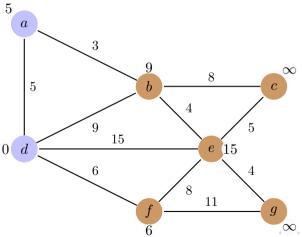
Dijkstra's Shortest Path Algorithm

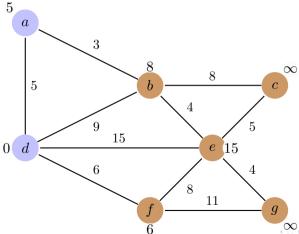
```
Algorithm ComputeDijkstraSPs(G, s)
Output: Array SP[.] with length of shortest path from s to v stored in SP[v].
P \leftarrow \text{new Min-heap}
for v \in G.aetVertices() do
   if v = s then
       SP[v] = 0
   else
       SP[v] = \infty
   end if
    P.insert(v, SP[v])
end for
while P.isNotEmpty() do
    u \leftarrow P.aetMin()
    P.removeMin() //Remove u
   for e \in G.incidentEdges(u) do
       w \leftarrow G.otherVertex(e, u)
       r \leftarrow SP[\mathbf{u}] + weight(e) / | Relax e |
       if r < SP[w] then
           SP[w] = r
           P.undate(w,r)/Update the value for key w
       end if
   end for
end while
```

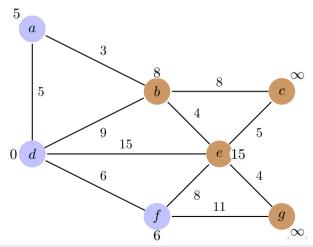
Shortest Path (from d)

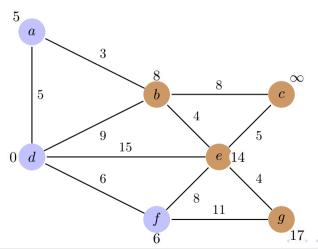


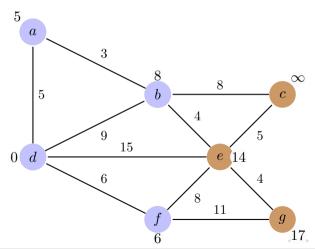


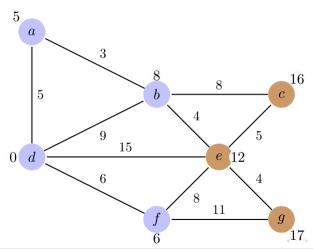


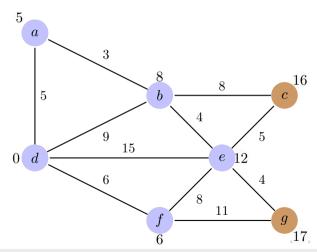


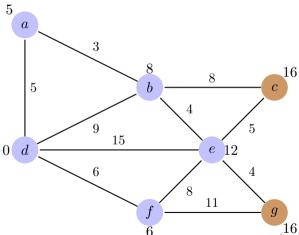


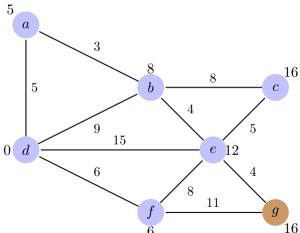


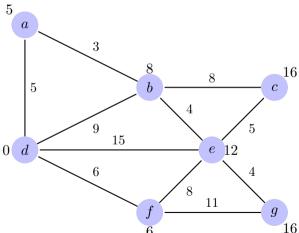












Analysis of Dijkstra's Algorithm

```
Algorithm ComputeDijkstraSPs(G, s)
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P \leftarrow \text{new Min-heap}
for v \in G.aetVertices() do
   if v = s then
       SP[v] = 0
    معام
       SP[v] = \infty
   end if
   P.insert(v, SP[v])
end for \implies c_1 \times |V| times
while P.isNotEmpty() do
   u \leftarrow P. qetMin() \implies c_2 \times \log |V| \text{ times}
    P.removeMin() //Remove u
   for e \in G.incidentEdges(u) do
       w \leftarrow G.otherVertex(e, u)
       r \leftarrow SP[u] + weight(e) / |Relax| e
       if r < SP[w] then
           SP[w] = r
           P.update(w,r)/Update the value for key w \implies c_3 \times \log |V| times
       end if
   end for \implies c_A \times deg(v) times
end while \implies c_5 \times |V| times
```

$$T(n) = c_3 c_4 \log |V| \sum_{v \in V} deg(v) + c_2 c_5 |V| (\log |V|) + c_1 |V| = O(|E| + |V|) \log |V|$$



Thank you