### Data Structures and Algorithms

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Session: String Matching Algorithms



# String Matching Problem<sup>1</sup>

Finding all valid shifts with which a given pattern P occurs in a given text T



<sup>&</sup>lt;sup>1</sup>Chapter 32, CLRS, Third Edition





$$s=1$$
 S T E P S T Y L E S T Y

$$s=2$$
  $\begin{bmatrix} \mathsf{S} & \mathsf{T} & \mathsf{E} & \mathsf{P} & \mathsf{S} & \mathsf{T} & \mathsf{Y} & \mathsf{L} & \mathsf{E} \end{bmatrix}$ 



$$s=1$$
 S T E P S T Y L E S T Y

$$s=2$$
 S T E P S T Y L E

$$s=3$$
 S T E P S T Y L E



$$s=1$$
 S T E P S T Y L E

$$s=2$$
 S T F P S T Y L E

$$s=3$$
 S T E P S T Y L E



$$s=0$$
 S T E P S T Y L E

$$s=3$$
 S T E P S T Y L E



$$s = 0$$
 S T E P S T Y L E

$$s = 4$$
 S T E P S T Y L E

$$s = 5$$
 S T E P S T Y L E

$$s = 6$$
 S T E P S T Y L E



### Naive String Matching Algorithm

```
Algorithm NaiveStringMatchingAlgorithm(T,P) Input Text T of size n and Pattern P of size m Define s as the shift index to T for s\in (0...n-m) do j=0 \qquad \text{while } j< m \ \& \ T[s+j]=P[j] \ \text{do} \\ j=j+1 \qquad \text{end while} \\ \text{if } j=m \ \text{then} \\ \text{print 'Valid at shift } s' \\ \text{end if} \\ \text{end for}
```

Figure: Naive String Algorithm



## Analysis of Naive String Matching Algorithm

```
Algorithm NaiveStringMatchingAlgorithm(T,P)
Input Text T of size n and Pattern P of size m
Define s as the shift index to T
for s \in (0...n-m) do
j=0
while j < m \ \& \ T[s+j] = P[j] do
j=j+1
end while \implies c_1 \times m times
if j=m then
print 'Valid at shift s'
end if
end for \implies c_2 \times (n-m+1) times
```

Figure: Naive String Matching Algorithm

```
T(n) = c_1c_2(n-m+1)m = O((n-m+1)m)
**Note: Preprocessing time is 0
```



## Thank you