

Data Structures and Algorithms

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Session: Shortest Path Algorithm
(All Pair Shortest Path)

All Pair Shortest Path Problem

1. Given: A directed weighted graph $G(V, E)$, for each edge $(v_1, v_2) \in E$ and an associated weight of an edge $w(v_1, v_2)$
2. Find: A shortest path from v_1 to v_2 for every pair of vertices v_1 and v_2 in V

Floyd-Warshall Algorithm: Setting up Notation

1. Given a directed weighted graph $G(V, E)$, the weight of each edge $(v_1, v_2) \in E$ can be defined using an adjacency-matrix representation W
2. If W is an $n \times n$ matrix,

$$w_{v_i v_j} = \begin{cases} 0 & \text{if } i = j, \\ \text{weight of directed edge } (v_i, v_j) & \text{if } i \neq j \text{ and } (i, j) \in E, \\ \infty & \text{if } i \neq j \text{ OR } (i, j) \notin E. \end{cases}$$

3. Distance matrix $D^0 = W$
4. Predecessor matrix L can be defined as

$$L^0(i, j) = \begin{cases} i & \text{if } (i, j) \in E \\ \text{NULL} & \text{if } i = j \text{ and } (i, j) \notin E. \end{cases}$$

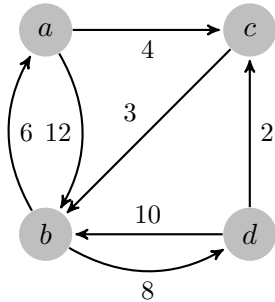
Floyd-Warshall Algorithm

```
Algorithm Floyd-WarshallAPSP( $G, W$ )  
 $n = \text{rows}(W)$   
 $D^0 = W$   
for  $k \in n$  do  
  let  $D^k = d_{ij}^k$  be a new  $n \times n$  matrix  
  for  $i \in n$  do  
    for  $j \in n$  do  
       $d_{ij}^k = \min(d_{ij}^{(k-1)}, d_{ik}^{(k-1)} + d_{kj}^{(k-1)})$   
    end for  
  end for  
end for  
return  $D^n$ 
```

Figure: Floyd-Warshall Algorithm

Example

Graph

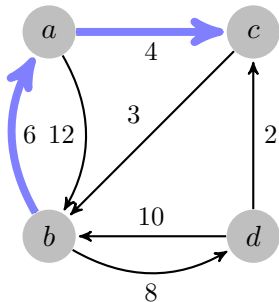


$k=0$

| | a | b | c | d |
|---|----------|----|----------|----------|
| a | 0 | 12 | 4 | ∞ |
| b | 6 | 0 | ∞ | 8 |
| c | ∞ | 3 | 0 | ∞ |
| d | ∞ | 10 | 2 | 0 |

Example

Graph

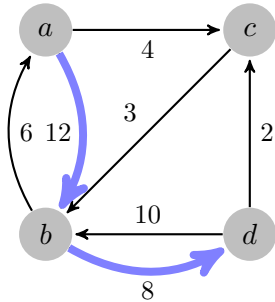


$k = 1 (v_i = a)$

| | a | b | c | d |
|---|----------|----|----|----------|
| a | 0 | 12 | 4 | ∞ |
| b | 6 | 0 | 10 | 8 |
| c | ∞ | 3 | 0 | ∞ |
| d | ∞ | 10 | 2 | 0 |

Example

Graph

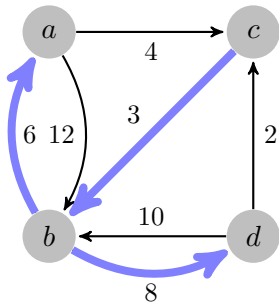


$k = 2$ ($v_i = a$)

| | a | b | c | d |
|---|----------|----|----|-----------|
| a | 0 | 12 | 4 | 20 |
| b | 6 | 0 | 10 | 8 |
| c | ∞ | 3 | 0 | ∞ |
| d | ∞ | 10 | 2 | 0 |

Example

Graph

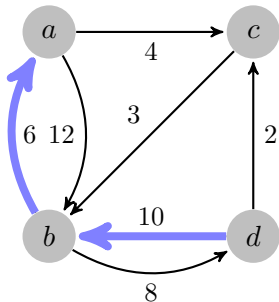


$k = 2$ ($v_i = c$)

| | a | b | c | d |
|---|----------|----|----|-----------|
| a | 0 | 12 | 4 | 20 |
| b | 6 | 0 | 10 | 8 |
| c | 9 | 3 | 0 | 11 |
| d | ∞ | 10 | 2 | 0 |

Example

Graph

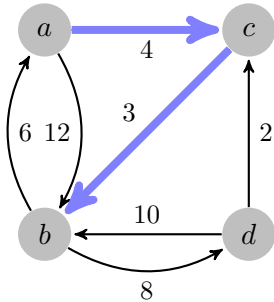


$k = 2$ ($v_i = b$)

| | a | b | c | d |
|---|-----------|----|----|----|
| a | 0 | 12 | 4 | 20 |
| b | 6 | 0 | 10 | 8 |
| c | 9 | 3 | 0 | 11 |
| d | 16 | 10 | 2 | 0 |

Example

Graph

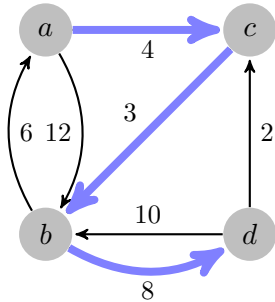


$k = 3$ ($v_i = a$, $v_j = b$)

| | a | b | c | d |
|---|----|----|----|----|
| a | 0 | 7 | 4 | 20 |
| b | 6 | 0 | 10 | 8 |
| c | 9 | 3 | 0 | 11 |
| d | 16 | 10 | 2 | 0 |

Example

Graph



$k = 3$ ($v_i = a$, $v_j = d$)

| | a | b | c | d |
|---|----|----|----|----|
| a | 0 | 7 | 4 | 15 |
| b | 6 | 0 | 10 | 8 |
| c | 9 | 3 | 0 | 11 |
| d | 16 | 10 | 2 | 0 |

Thank you