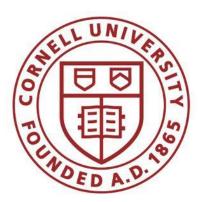
Design Trade-offs for Decentralized Baseband Processing in Massive MU-MIMO Systems

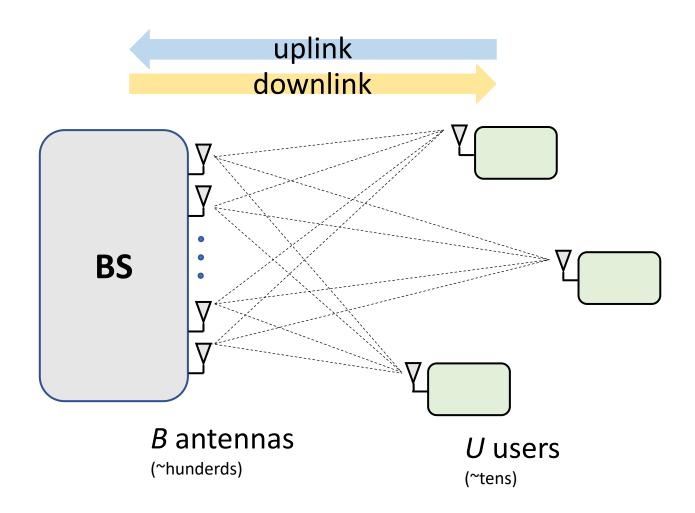
Kaipeng Li, James McNaney, Oscar Castañeda, Chance Tarver, Charles Jeon, Joseph Cavallaro, Christoph Studer

Asilomar Conference on Signals, Systems, and Computers November 5, 2019

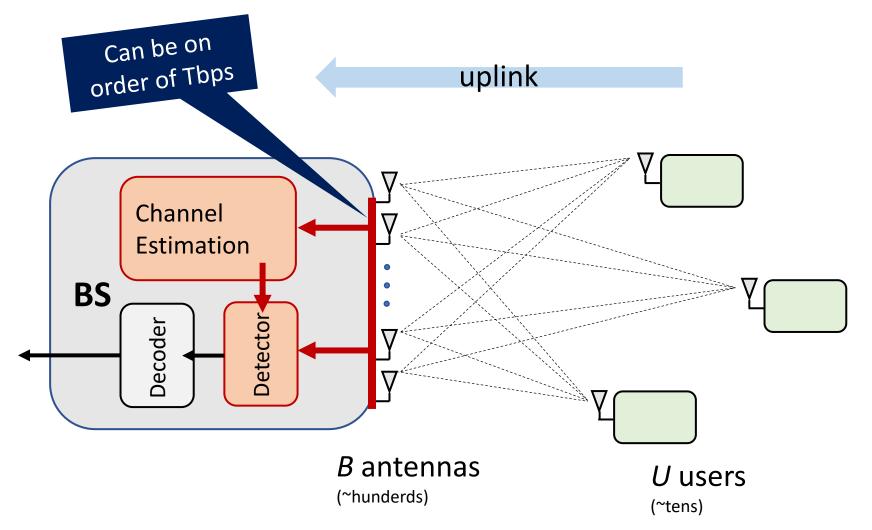




Massive MU-MIMO systems



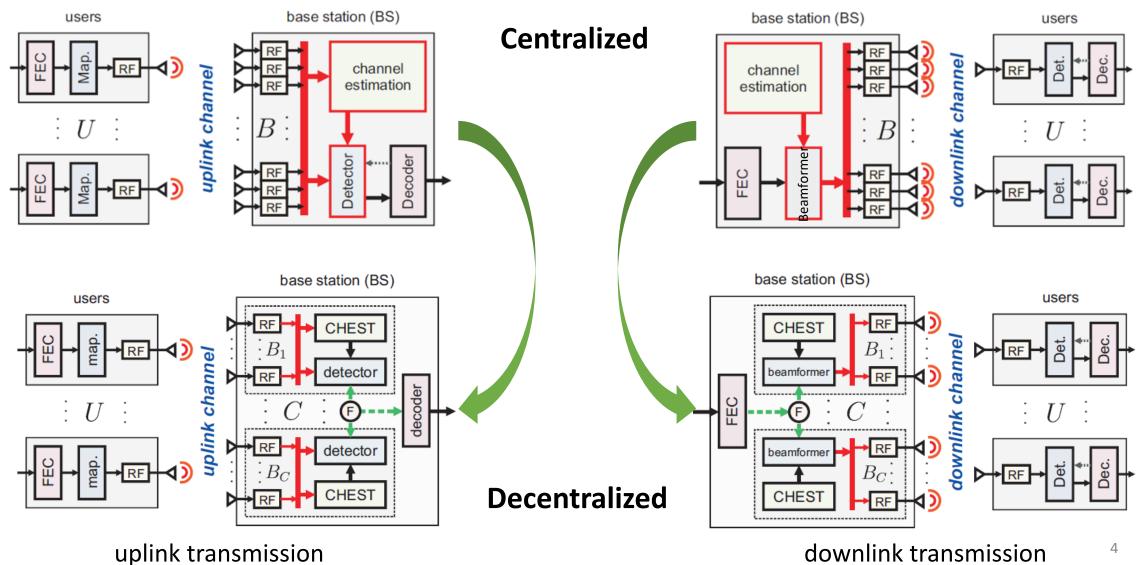
How do we handle this much data?



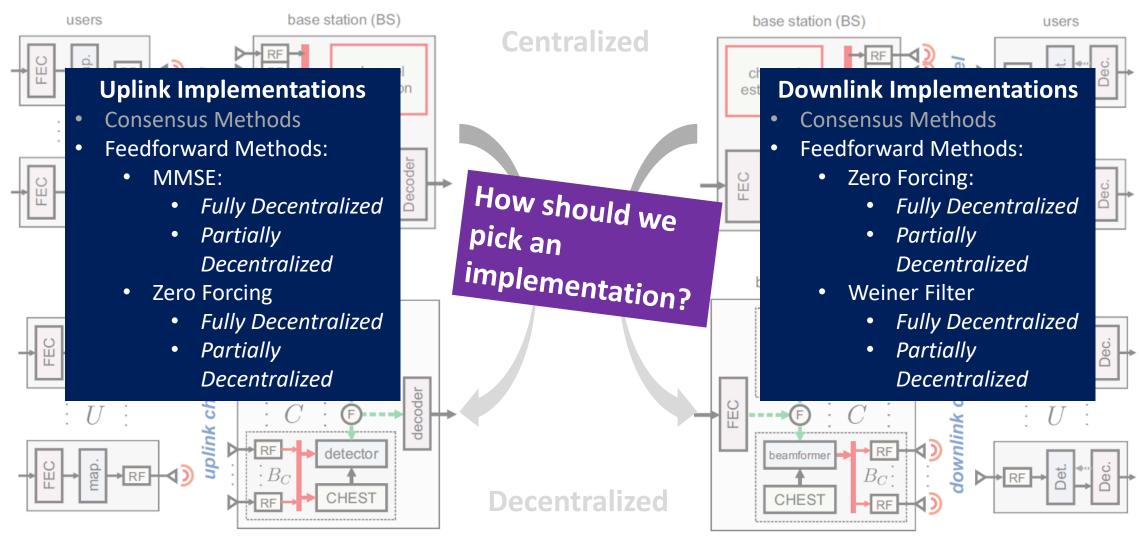
Possible Limitations:

- Chip I/O and interconnection bandwidth
- On-chip memory and storage
- Computing capability of modern computing fabrics

Decentralized to resolve bottlenecks



Decentralized to resolve bottlenecks

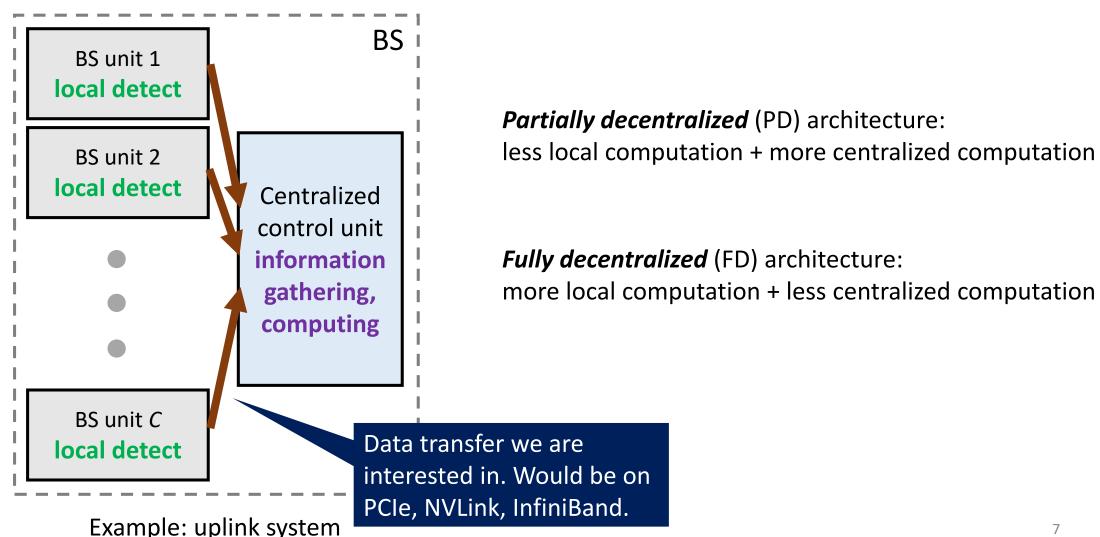


Outline

- Overview of decentralized architectures and algorithms
- Architecture trade-offs
- Algorithm trade-offs
- Precision trade-offs
- Practical design flow
- Conclusion

Decentralized feedforward architecture

Feedforward local information *only once* instead of multiple rounds to centralized unit

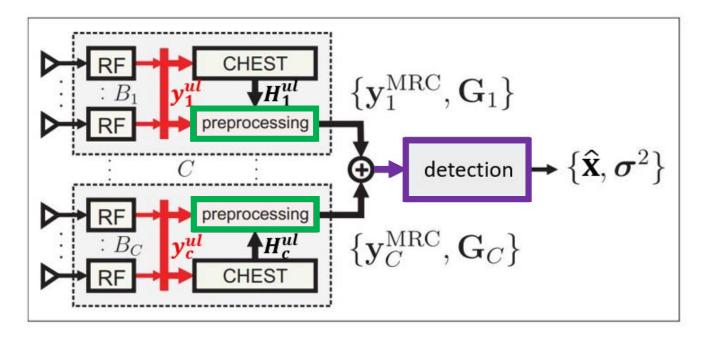


Uplink linear MMSE detection

Centralized linear MMSE (C-LMMSE) detection:

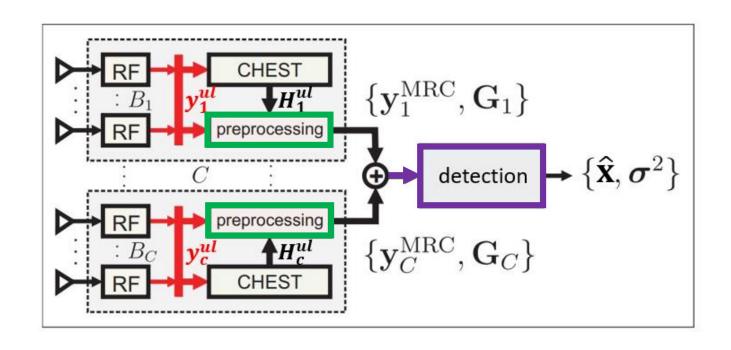
$$\widehat{\boldsymbol{x}} = (\boldsymbol{H}^{H}\boldsymbol{H} + \frac{N_{0}}{E_{x}}\boldsymbol{I})^{-1}\boldsymbol{H}^{H}\boldsymbol{y}$$
$$= (\boldsymbol{G} + \frac{N_{0}}{E_{x}}\boldsymbol{I})^{-1}\boldsymbol{y}^{MRC}$$

Partially decentralization: decentralized matrix preprocessing + centralized detection



 σ^2 : error variance

Uplink linear MMSE detection



PD-LMMSE obtains the same \widehat{x} as C-LMMSE

Complexity: $O(B_cU^2) + O(U^3)$ mults.

Data transfer size: O(U2) samples / cluster

Centralized LMMSE (C-LMMSE):

$$G = H^{H}H \quad y^{MRC} = H^{H}y$$

$$\widehat{x} = (G + \frac{N_0}{E_x}I)^{-1}y^{MRC}$$



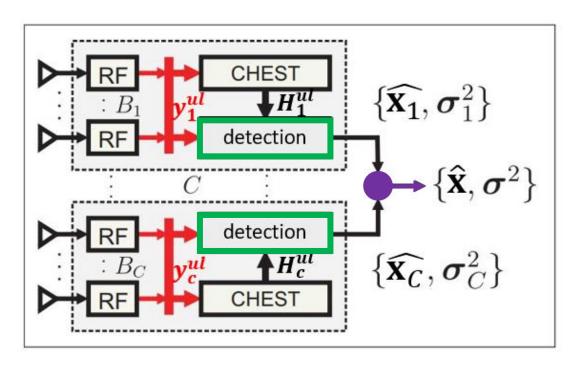
Partially Decentralized:

$$G_c = H_c^H H_c$$
 $y_c^{\text{MRC}} = H_c^H y_c$

$$G = \sum_{c=1}^{C} G_c \qquad y^{\text{MRC}} = \sum_{c=1}^{C} y_c^{\text{MRC}}$$

$$\widehat{\boldsymbol{x}} = (\boldsymbol{G} + \frac{N_0}{E_x} \boldsymbol{I})^{-1} \boldsymbol{y}^{\text{MRC}}$$

Fully decentralized (FD-) LMMSE detection



Decentralized local detection

$$\widehat{\boldsymbol{x}}_{\boldsymbol{c}} = (\boldsymbol{G}_{\boldsymbol{c}} + \frac{N_0}{E_{\boldsymbol{x}}}\boldsymbol{I})^{-1}\boldsymbol{y}_{\boldsymbol{c}}^{\mathrm{MRC}}$$

Fusion of local \widehat{x}_c using weights, λ_c :

$$\widehat{\mathbf{x}} = \sum_{c=1}^{C} \lambda_c \widehat{\mathbf{x}}_c$$

Optimal λ_c is a function of σ_c

Complexity: $O(B_cU^2) + O(U^3)$ mults.

Data transfer size: **O(U)** samples / cluster

Downlink Beamforming

Linear beamforming:

- Power constraint: $\mathbf{E}[\|\mathbf{x}\|^2] \le \rho^2$

Precoding matrix: P

- Linear precoding: x = Ps

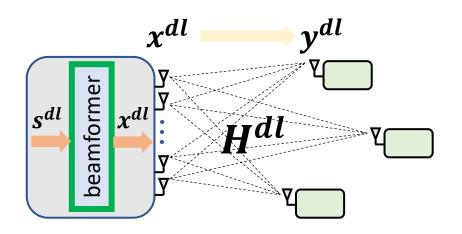
Zero-Forcing beamforming:

- Precoding Matrix: $H^H(HH^H)^{-1} = H^HG^{-1}$

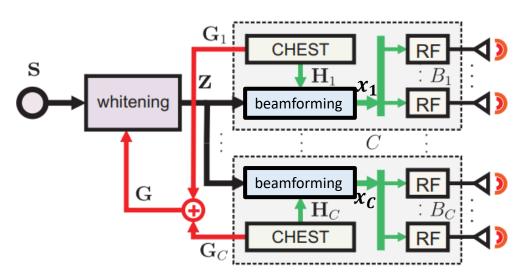
- Power constraint: $\hat{x} = \rho \|\hat{x}\|_2$

Channel reciprocity:

- TDD Transmission: $H^{dl} = (H^{ul})^H$



Decentralized feedforward ZF beamforming



Partially decentralized ZF beamforming:

Set
$$\rho_c^2 = \rho^2/C$$

$$G_c = H_c H_c^H$$

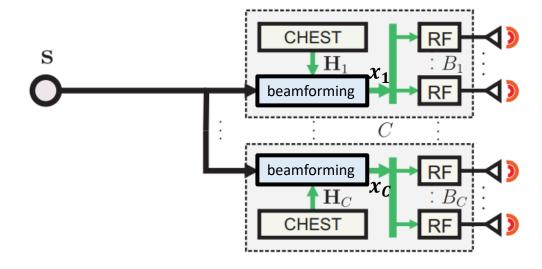
$$G = \sum_{c=1}^C G_c \quad z = G^{-1} s$$

Broadcast z to local clusters

$$\widehat{\boldsymbol{x}}_c = \boldsymbol{H}_c^H \boldsymbol{z}, \quad \widehat{\boldsymbol{x}}_c = \rho_c \|\widehat{\boldsymbol{x}}_c\|_2$$

Complexity: $O(B_cU^2) + O(U^3)$ mults.

Data transfer: O(U²) samples / cluster



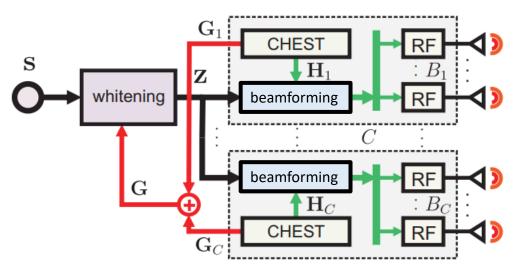
Fully decentralized ZF beamforming:

Broadcast s and set $\rho_c^2 = \rho^2/C$ $\hat{x}_c = H_c^H (H_c H_c^H)^{-1} s$ $\hat{x}_c = \rho_c ||\hat{x}_c||_2$

Complexity: $O(B_cU^2) + O(U^3)$ mults.

Data transfer: **O(U)** samples / cluster

Decentralized feedforward Wiener Filter (WF) beamforming



Partially decentralized WF beamforming

Set
$$\rho_c^2 = \rho^2/C$$

$$G_C = H_C H_C^H$$

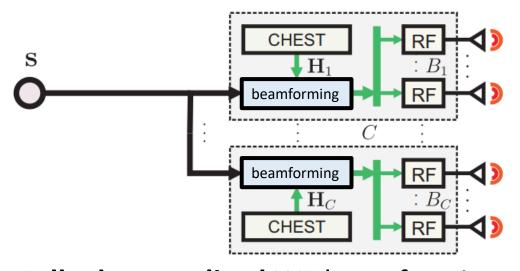
$$G = \sum_{c=1}^{C} G_c \quad z = \frac{1}{\beta} (G + \gamma I_U)^{-1} s$$

Broadcast z to local BS unit

$$\widehat{\boldsymbol{x}}_{c} = \boldsymbol{H}_{c}^{H} \boldsymbol{z}$$

Complexity: $O(B_cU^2) + O(U^3) + O(\beta)$ mult.

Data transfer: O(U²) samples / cluster



Fully decentralized WF beamforming

Broadcast s and set $\rho_c^2 = \rho^2/C$

$$\boldsymbol{P}_{c} = \frac{1}{\beta_{c}} \boldsymbol{H}_{c}^{H} (\boldsymbol{H}_{c} \boldsymbol{H}_{c}^{H} + \gamma \boldsymbol{I}_{U})^{-1} \boldsymbol{s}$$

$$\widehat{\boldsymbol{x}}_{c} = \boldsymbol{P}_{c}\boldsymbol{s}$$

Complexity: $O(B_cU^2) + O(U^3) + O(\beta)$ mult.

Data transfer: O(U) samples / cluster

Architecture Trade-offs: PD vs. FD

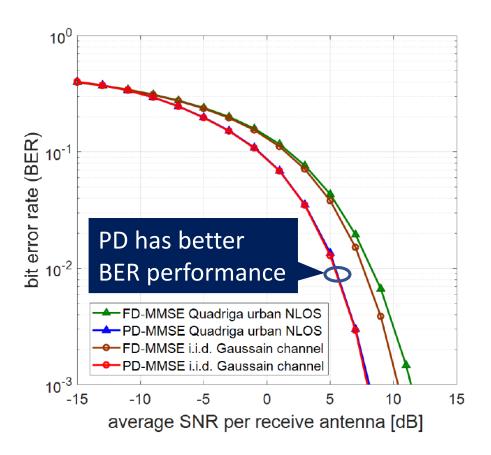
PD-MMSE and FD-MMSE: Data transfer Depends on channel coherency

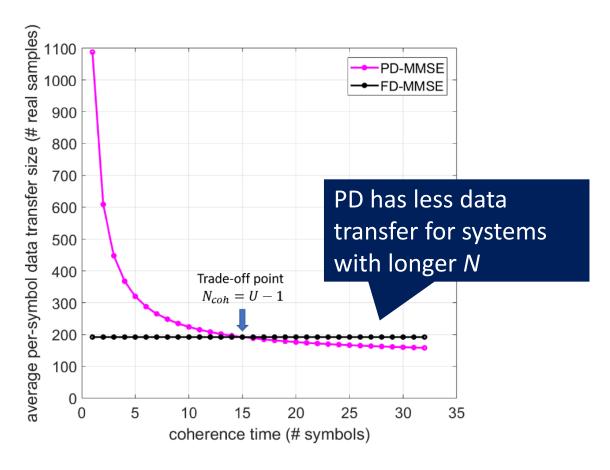
- BER: Centralized MMSE = PD-MMSE, FD-MMSE sacrifices BER
- Computation (timing) complexity: PD-MMSE = FD-MMSE
- N_{coh} : Period in which we update channel state information

$$m_{PD} = \frac{C \times (U^2 - U + 2N_{coh}U)}{N_{coh}}. \quad m_{FD} = \frac{C \times 3N_{coh}U}{N_{coh}} = 3CU.$$

PD vs. FD trade-off on BER and data transfer

C=4, U=16, $B_c=32$, B=128, 16QAMSimple i.i.d. Gaussian channel and Quadriga NLOS urban campus channel



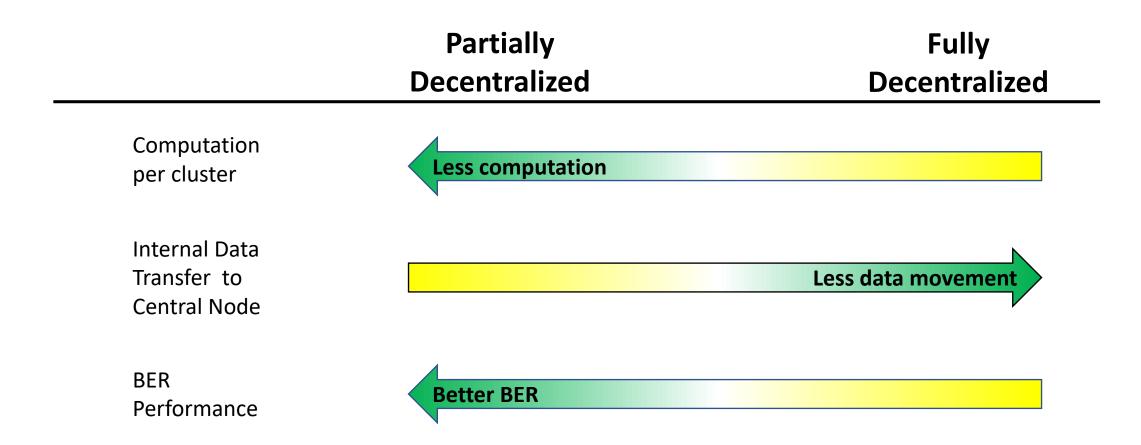


(a) BER: PD-MMSE vs. FD-MMSE

(b) Data transfer size vs. N_{coh}

Decentralized feedforward architecture

Feedforward local information *only once* instead of multiple rounds to centralized unit

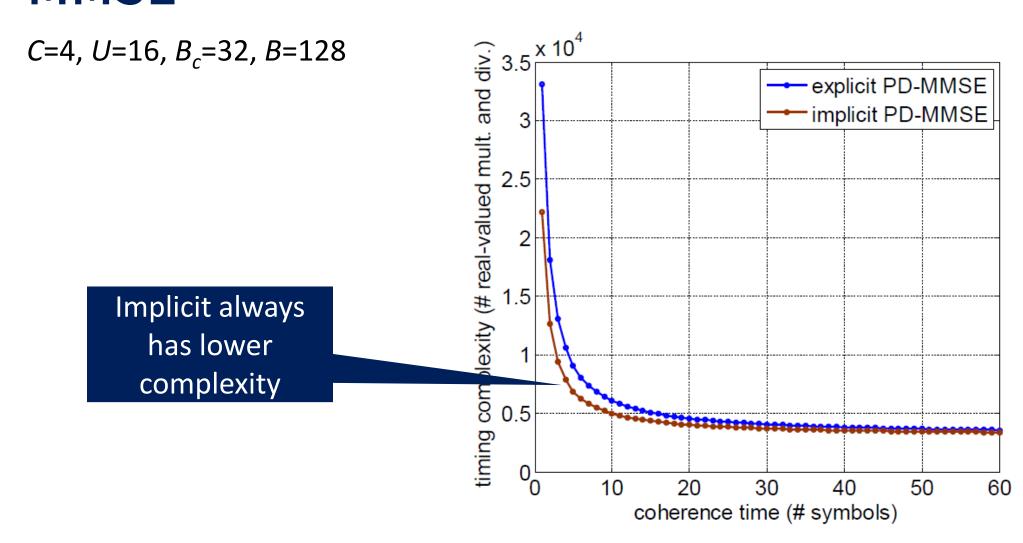


Algorithm Trade-offs: Explicit vs. Implicit method

- Example: PD-MMSE with explicit matrix inversion vs. implicit matrix inversion
- Implicit matrix inversion $A^{-1} = (G + \frac{N_0}{E_x}I)^{-1}$ for PD-MMSE
 - $A = LL^H$ (Cholesky decomposition, L is lower triangular matrix)
 - Get z by solving $Lz = y^{MRC}$ using forward substitution
 - Get \widehat{x} by solving $L^H \widehat{x} = z$ using backward substitution
- ullet Per-symbol complexity of explicit and implicit methods depend on N_{coh}

$$n_{ex} = (2B_c U^2 + \frac{10}{3}U^3 - \frac{1}{3}U)/N_{coh} + 4B_c U + 4U^2$$
$$n_{im} = (2B_c U^2 + \frac{2}{3}U^3 + \frac{1}{3}U)/N_{coh} + 4B_c U + 4U^2$$

Complexity of explicit vs. implicit PD-MMSE

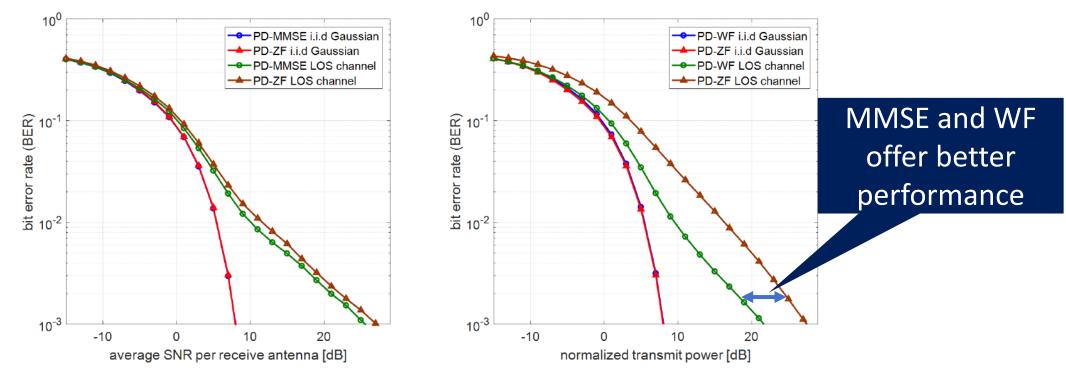


Reusing Uplink (UL) Results for Downlink (DL)

- Channel reciprocity in TDD system: $H^{UL} = (H^{DL})^H$
- Gram matrix: $G^{DL} = H^{DL}(H^{DL})^H = (H^{UL})^H H^{UL} = G^{UL}$
- Store and reuse computed uplink results for downlink to reduce complexity
- UL MMSE detection + DL WF beamforming can only reuse $m{G}^{UL}$
- UL ZF detection + DL ZF beamforming can even reuse $(G^{UL})^{-1}$

UL and DL integration trade-offs on BER and complexity

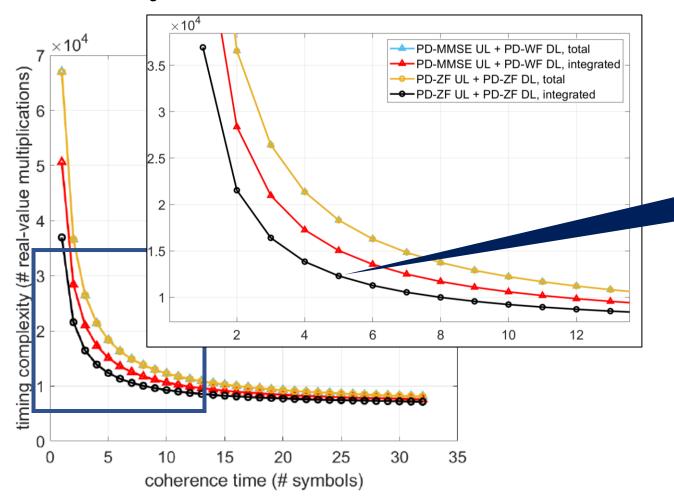
Example: UL *PD-MMSE* + *DL PD-WF* integration vs. *UL ZF* + *DL ZF* integration C=4, U=16, $B_c=32$, B=128, 16QAM, LOS channel



(a) BER: PD-MMSE detection vs. PD-ZF detection (b) BER: PD-WF precoding vs. PD-ZF precoding

UL and DL integration trade-offs on BER and complexity

Example: UL *PD-MMSE* + *DL PD-WF* integration vs. *UL ZF* + *DL ZF* integration C=4, U=16, $B_c=32$, B=128, 16QAM, LOS channel

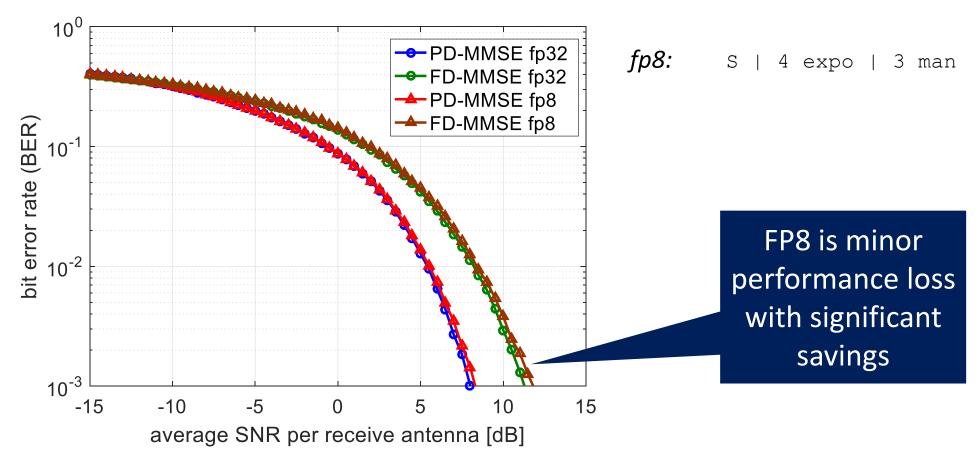


By integrating, ZF only requires 65% of the multiplies

Precision Trade-offs: 32-bit vs. 8bit floating point

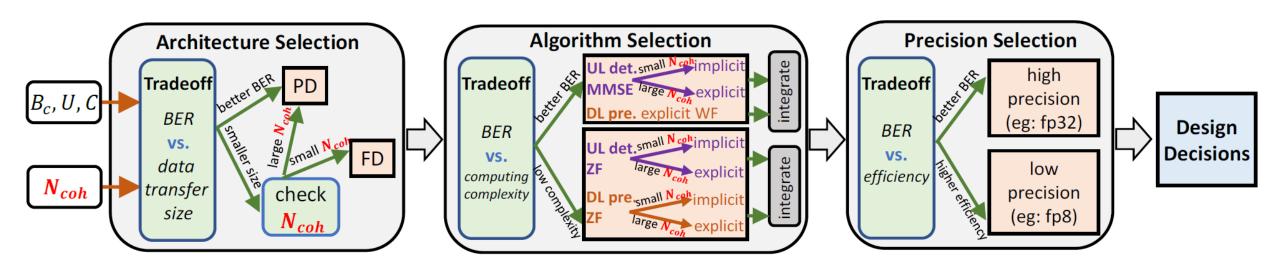
Example: PD-MMSE and FD-MMSE

C=4, U=16, B_c=32, B=128, 16QAM, Quadriga NLOS urban campus channel



8-bit floating point reduces 4x data transfer size compared to 32-bit at only little loss of BER

Summary of Tradeoffs



Conclusion

- Decentralized baseband processing resolves complexity and interconnection bandwidth bottlenecks for massive MU-MIMO
- Practical massive MU-MIMO should leverage design trade-offs at different aspects:
 - Architecture trade-offs of PD and FD on BER vs. data transfer size
 - Unless you expect very low coherence time, choose partially decentralized.
 - Algorithms trade-offs of explicit and implicit methods on BER vs. complexity
 - Use implicit matrix inversions whenever possible. Reuse results from uplink to downlink.
 - Precision trade-offs of various data precision options on BER vs. efficiency
 - Use fp16 or even fp8 unless BER is serious concern.