

## IX. APPENDIX

The appendix presents two supporting examples covered in the two subsections below. The first example demonstrates how LCBs are computed. The second example illustrates the interdependence of LCBs.

### A. Example LCB Calculation

To illustrate our approach for computing LCBs, consider the following example. Assume we have two tasks,  $\tau_1$  and  $\tau_2$  with UCBs and ECBs for each listed in Figure 9. Please note that we have not shown the ECBs, UCBs, or AUCBs in the following figures for  $\delta_i^0$  as it is a dummy basic block with no elements for each of the aforementioned sets.

Task ID	Evicting Cache Blocks $ECB(\delta_i^j)$				
	$\delta_i^1$	$\delta_i^2$	$\delta_i^3$	$\delta_i^4$	$\delta_i^5$
$\tau_1$	{1,2}	{3,4,8}	{4,5,6,8}	{1,2,7,8}	{1,2,7,8}
$\tau_2$	{1,9}	{3,10}	{11,12}	{5,7,13}	{1,3,7,8}

Task ID	Useful Cache Blocks $UCB_{out}(\delta_i^j)$				
	$\delta_i^1$	$\delta_i^2$	$\delta_i^3$	$\delta_i^4$	$\delta_i^5$
$\tau_1$	{1,2}	{1,2,4,8}	{1,2,8}	{1,2,7,8}	{1,2,7,8}
$\tau_2$	{1}	{1,3}	{1,3}	{1,3,7}	{1,3,7,8}

Fig. 9: Taskset ECBs and UCBs.

The computation for LCBs uses the accessed useful cache blocks (AUCBs) since the cache blocks that are re-loaded during execution of the non-preemptive region between preemption points is a function of the memory that is explicitly accessed by the preempted task. The computed AUCBs for each task is shown in Figure 10. In accordance with Equation 7 one can readily see that the AUCBs are simply the intersection of the UCBs and ECBs for each basic block.

Task ID	Accessed Useful Cache Blocks $AUCB_{out}(\delta_i^j)$				
	$\delta_i^1$	$\delta_i^2$	$\delta_i^3$	$\delta_i^4$	$\delta_i^5$
$\tau_1$	{1,2}	{4,8}	{8}	{1,2,7,8}	{1,2,7,8}
$\tau_2$	{1}	{3}	{1,3}	{7}	{1,3,7,8}

Fig. 10: Taskset AUCBs.

In our example, assume preemptions are taken at basic blocks  $\delta_1^2$  and  $\delta_1^4$  for task  $\tau_1$ . For simplicity, we calculate the LCBs associated with these two preemption points. For  $LCB(\delta_1^2, \delta_1^4)$ , we have  $UCB(\delta_1^2) = \{1, 2, 4, 8\}$ . The second expression is the set of memory that is accessed in basic blocks  $\delta_1^3$  and  $\delta_1^4$ , namely  $\{8\} \cup \{1, 2, 7, 8\} = \{1, 2, 7, 8\}$  comprising the set of AUCBs. The third expression is the set of ECBs for task  $\tau_2$  where  $ECB(\tau_2) = \{1, 3, 5, 7, 8, 9, 10, 11, 12, 13\}$ . Thus,  $LCB(\delta_1^2, \delta_1^4)$  is given by the intersection of the three sets:

$$LCB(\delta_1^2, \delta_1^4) = \{1, 2, 4, 8\} \cap \{1, 2, 7, 8\} \cap \{1, 3, 5, 7, 8, 9, 10, 11, 12, 13\} = \{1, 8\}$$

The preemption cost  $\gamma(\delta_1^2, \delta_1^4)$  for a  $BRT = 390\mu s$  is given by:

$$\gamma(\delta_1^2, \delta_1^4) = |\{1, 8\}| \cdot 390 = 780\mu s$$

Using the same method,  $LCB(\delta_1^4, \delta_1^5)$  is given by:

$$LCB(\delta_1^4, \delta_1^5) = \{1, 2, 7, 8\} \cap \{1, 2, 7, 8\} \cap \{1, 3, 5, 7, 8, 9, 10, 11, 12, 13\} = \{1, 7, 8\}$$

The preemption cost  $\gamma(\delta_1^4, \delta_1^5)$  for a  $BRT = 390\mu s$  is given by:

$$\gamma(\delta_1^4, \delta_1^5) = |\{1, 7, 8\}| \cdot 390 = 1170\mu s$$

Using the method illustrated here, the preemption cost matrix entries for each pair of basic blocks are computed in a similar fashion and used as input to our preemption point placement algorithm.

### B. Example of LCB Interdependence

To further exemplify the interdependence of preemption points, consider the example shown below. In order to account for all re-loaded cache blocks (LCBs), preemptions are always included at the first basic block  $\delta_i^0$  and the last basic block  $\delta_i^{N_i}$  as shown in Figure 11. This is commensurate with the preemptions that occur before and after the task executes. Assume we have two tasks where  $\tau_2$  contains four basic blocks which may be preempted by task  $\tau_1$ . For simplification, assume that the ECBs of task  $\tau_1$  evicts all UCBs of task  $\tau_2$ . Let us further assume that we have  $\rho_2 = \{\delta_2^0, \delta_2^1, \delta_2^2, \delta_2^4\}$ . Using our LCB computation approach, only the re-loaded lines as captured in the terms  $LCB(\delta_2^1, \delta_2^2)$  and  $LCB(\delta_2^2, \delta_2^4)$  are included in the  $C_2$  computation.  $LCB(\delta_2^0, \delta_2^1) = 0$  as no LCBs have been cached until after execution of basic block  $\delta_2^1$ .

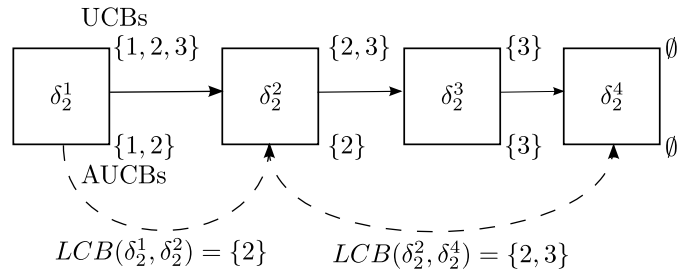


Fig. 11: LCB Interdependence