

# **Hakim Sabzevari University**

# CTRL+ALT+DEFEAT

# **Team Reference Document**

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https://github.com/ctrl-alt-Defeat-icpc

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## 1. STL

## 1.1. bitscroll

```
__builtin_ctz(x); // first 1 from left (index)
__builtin_popcount(x); // count of 1 in numbers bit
__builtin_ctzll(x); // for long long
```

```
__builtin_popcountll(x); // ...
14
    1.2. 128 bit
     int128 read() {
        int128 x = 0, f = 1;
        char ch = getchar();
        while (ch < '0' || ch > '9') {
            if (ch == '-') f = -1;
            ch = getchar();
        }
        while (ch >= '0' && ch <= '9') {
            x = x * 10 + ch - '0';
            ch = getchar();
        return x * f;
    void print(__int128 x) {
        if (x < 0) {
            putchar('-');
            x = -x;
        if (x > 9) print(x / 10);
22
        putchar(x % 10 + '0');
    bool cmp(\_int128 x, \_int128 y) { return x > y; }
```

# 2. Segment Tree

<u>\_\_int128</u> x = read();

int main() {

print(x);

cout << endl;
return 0;</pre>

## 2.1. easy implementation

const int N = 1e5; // limit for array size

```
int n; // array size
int t[2 * N];
void build() { // build the tree
 for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] +
t[i<<1|1];
}
void modify(int p, int value) { // set value at
position p
 for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] =
t[p] + t[p^1];
}
int query(int 1, int r) { // sum on interval [1, r)
  int res = 0;
  for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1\&1) res += t[1++];
   if (r\&1) res += t[--r];
  }
  return res;
int main() {
  scanf("%d", &n);
 for (int i = 0; i < n; ++i) scanf("%d", t + n + 1
i);
  build();
 modify(0, 2);
  printf("%d\n", query(3, 11));
  return 0;
```

## 2.2. with lazy propagation

```
const int N = 1e5 + 5;
int n;
int seg[2 * N], lazy[2 * N], a[N];
```

```
int segSize;
void build(int u = 1, int ul = 0, int ur = n) {
   if(ur - ul < 2)
        seg[u] = a[ul];
        return;
    int mid = (ul + ur) / 2;
   build(u * 2, ul, mid);
    build(u * 2 + 1, mid, ur);
    seg[u] = seg[u * 2] + seg[u * 2 + 1];
void upd(int u, int ul, int ur, int x){
   lazy[u] += x;
    seg[u] += (ur - ul) * x;
void shift(int u, int ul, int ur){
   int mid = (ul + ur) / 2;
   upd(u * 2, ul, mid, lazy[u]);
   upd(u * 2 + 1, mid, ur, lazy[u]);
   lazy[u] = 0;
void increase(int 1, int r, int x, int u = 1, int ul
= 0, int ur = n){
   if(1 >= ur || ul >= r)return;
   if(1 <= u1 && ur <= r){
        upd(u, ul, ur, x);
        return;
   }
   shift(u, ul, ur);
   int mid = (ul + ur) / 2;
   increase(1, r, x, u * 2, u1, mid);
   increase(l, r, x, u * 2 + 1, mid, ur);
    seg[u] = seg[u * 2] + seg[u * 2 + 1];
int sum(int l, int r, int u = 1, int ul = 0, int ur
= n){
```

```
if(1 \ge ur \mid \mid ul \ge r)return 0;
    if(1 <= ul && ur <= r)return seg[u];</pre>
    shift(u, ul, ur);
    int mid = (ul + ur) / 2;
    return sum(1, r, u * 2, ul, mid) + sum(1, r, u *
2 + 1, mid, ur);
void showSegments() {
    for(int i = 0; i < segSize; i++)</pre>
        cout << seg[i] << ' ';
    cout << endl;</pre>
void Main() {
    cin >> n;
    segSize = 2;
    while(segSize / 2 <= n) segSize *= 2;</pre>
    for(int i = 0; i < n; i++)
        cin >> a[i];
    build();
int main() {
    ios base::sync with stdio(false);
    cin.tie(0); cout.tie(0);
    Main();
    return 0;
```

## 3. Math

## 3.1. choose

```
#define 11 long long
const int N = 2e3 + 5;
const 11 M = 1e9 + 7;
```

```
11 fact[N], inv[N];
int r, n, q;
11 exp(ll b, ll p, ll m) {
    b %= m:
    11 result = 1;
    while(p) {
        if(p % 2)
            result = result * b % m;
        b = b * b % m;
        p /= 2;
    return result;
}
void preProcess() {
    fact[0] = 1;
    for(int i = 1; i < N; i++)</pre>
        fact[i] = fact[i - 1] * i % M;
    inv[N - 1] = exp(fact[N - 1], M - 2, M);
    for(int i = N - 1; i > 0; i--)
        inv[i - 1] = inv[i] * i % M;
}
11 choose(int n, int r) {
    if(r > n) return 0;
    return fact[n] * inv[r] % M * inv[n - r] % M;
}
void Main() {
    cin >> q;
    while(q--) {
        cin \gg n \gg r;
        cout << choose(n, r) << '\n';</pre>
    }
}
int main() {
```

```
ios::sync_with_stdio(false);
    cin.tie(0); cout.tie(0);
    preProcess();
   Main();
   return 0;
3.2. gcd
int gcd (int a, int b) {
    return b ? gcd (b, a % b) : a;
// fast version...
int gcd(int a, int b) {
   if (!a || !b)
        return a | b;
   unsigned shift = builtin ctz(a | b);
    a >>= builtin ctz(a);
    do {
        b >>= builtin ctz(b);
        if (a > b)
            swap(a, b);
        b -= a:
   } while (b);
    return a << shift;</pre>
3.3. compressing
// compressing
sort(temp values, temp values + n);
int numOfUnique = unique(temp values, temp values +
n) - temp values;
for(int i = 0; i < n; i++)</pre>
    h[i] = lower bound(temp values, temp values +
numOfUnique, h[i]) - temp values;
```

```
3.4. lower bound and upper bound int main() {
```

```
int main() {
    vector<int> v = {11, 34, 56, 67, 89};

    // Finding lower bound of 56
    cout << *lower_bound(v.begin(), v.end(), 56)
        << endl;

    // Finding upper bound of 56
    cout << *upper_bound(v.begin(), v.end(), 56);
    return 0;
}
Output:
56
67</pre>
```

## 4. Graph

## 4.1. BFS

```
#define distance d
const int maxN = 1e5 + 10, oo = 1e9;
vector <int> adj[maxN];
int distance[maxN];
queue<int> q;
void BFS(int n, int r) {
    for (int i=1; i<=n; i++) distance[i] = oo;</pre>
    distance[r] = 0;
    q.push(r);
    while(q.size()) {
        int v = q.front();
        q.pop();
        for (auto u : adj[v])
            if(distance[u] > distance[v] + 1) {
                distance[u] = distance[v] + 1;
                q.push(u);
```

```
int main() {
    ios base::sync_with_stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {</pre>
        int u, v; cin >> u >> v;
        adj[u].push back(v);
        adj[v].push back(u);
    }
    BFS(n, 1);
    for (int i=1; i<=n; i++)</pre>
        cout << i << ':' << distance[i] << '\n';</pre>
}
4.2. bipartite
const int maxN = 1e5 + 10;
vector <int> adj[maxN];
bool mark[maxN];
int color[maxN];
bool bipartite = true;
void DFS(int v, int parent) {
    mark[v] = true;
    if(parent != -1) color[v] = 1 - color[parent];
    else color[v] = 1;
    for (auto u : adj[v]) {
        if(!mark[u])
            DFS(u, v);
        else if(color[u] == color[v])
            bipartite = false;
    }
```

```
int main() {
    ios base::sync with stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {
        int u, v; cin >> u >> v;
        adj[u].push back(v);
        adj[v].push_back(u);
    for (int i=1; i<=n; i++) {
        if(mark[i]) continue;
        DFS(i, -1); //root does not have parent.
    }
    if(bipartite) cout << "Graph Is Bipartite\n";</pre>
    else cout << "Graph Is Not Bipartite\n";</pre>
4.3. cycle finding
const int maxN = 1e5 + 10;
vector <int> adj[maxN];
bool mark[maxN];
bool cycle found = false;
void DFS(int v, int parent) {
    mark[v] = true;
    for (auto u : adj[v]) {
        if(!mark[u]) DFS(u, v); //u's parent is v.
        else if(u != parent) cycle found = true;
    }
int main() {
```

ios base::sync with stdio(0); cin.tie(0);

```
int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {
        int u, v; cin >> u >> v;
        adj[u].push_back(v);
        adj[v].push_back(u);
    }
    for (int i=1; i<=n; i++) {
        if(mark[i]) continue;
        DFS(i, -1); //root does not have parent.
    }
    if(cycle found) cout << "Graph has Cycle\n";</pre>
    else cout << "Graph does not have Cycle\n";</pre>
4.4. DFS
const int maxN = 1e5 + 10;
vector <int> adj[maxN];
bool mark[maxN];
vector <int> component;
void DFS(int v) {
    mark[v] = true;
    component.push back(v);
    for (auto u : adj[v])
        if(!mark[u]) DFS(u);
int main() {
    ios base::sync with stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {
        int u, v; cin >> u >> v;
        adj[u].push back(v);
        adj[v].push back(u);
    }
    for (int i=1; i<=n; i++) {
        if(mark[i]) continue;
```

## 4.5. floyd-warshall

```
// Implementing floyd warshall algorithm
void floydWarshall(int graph[][nV]) {
   int matrix[nV][nV], i, j, k;
   for (i = 0; i < nV; i++)
      for (j = 0; j < nV; j++)
      matrix[i][j] = graph[i][j];
   // Adding vertices individually
   for (k = 0; k < nV; k++) {
      for (i = 0; i < nV; i++) {
        if (matrix[i][k] + matrix[k][j] <
      matrix[i][j])
           matrix[i][j] = matrix[i][k] +
   matrix[k][j];
      }
   }
   // printMatrix(matrix);
}</pre>
```

## **4.6.** prim

```
// Function to find sum of weights of edges of the
Minimum Spanning Tree.
int spanningTree(int V, int E, vector<vector<int>>
&edges) {
    // Create an adjacency list representation of
the graph
    vector<vector<int>> adj[V];
```

```
// Fill the adjacency list with edges and their
weights
   for (int i = 0; i < E; i++) {
        int u = edges[i][0];
        int v = edges[i][1];
        int wt = edges[i][2];
        adj[u].push back({v, wt});
        adj[v].push_back({u, wt});
   // Create a priority queue to store edges with
their weights
   priority queue<pair<int,int>,
vector<pair<int,int>>, greater<pair<int,int>>> pq;
   // Create a visited array to keep track of
visited vertices
   vector<bool> visited(V, false);
   // Variable to store the result (sum of edge
weights)
   int res = 0;
   // Start with vertex 0
   pq.push({0, 0});
   // Perform Prim's algorithm to find the Minimum
Spanning Tree
   while(!pq.empty()){
        auto p = pq.top();
        pq.pop();
        int wt = p.first; // Weight of the edge
        int u = p.second; // Vertex connected to
the edge
        if(visited[u] == true){
            continue; // Skip if the vertex is
already visited
```

```
res += wt; // Add the edge weight to the
result
        visited[u] = true; // Mark the vertex as
visited
        // Explore the adjacent vertices
        for(auto v : adj[u]){
            // v[0] represents the vertex and v[1]
represents the edge weight
            if(visited[v[0]] == false){
                pq.push(\{v[1], v[0]\}); // Add the
adjacent edge to the priority queue
        }
    return res; // Return the sum of edge weights
of the Minimum Spanning Tree
int main() {
    vector<vector<int>> graph = {{0, 1, 5},
                                  \{1, 2, 3\},\
                                   \{0, 2, 1\}\};
    cout << spanningTree(3, 3, graph) << endl;</pre>
    return 0;
4.7. shortest cycle
//this code works for simple graphs.
const int maxN = 1010, oo = 1e9;
vector <int> adj[maxN];
int deleted, distances[maxN];
queue<int> q;
void BFS(int n, int r) {
    for (int i=1; i<=n; i++) distances[i] = oo;</pre>
```

distances[r] = 0;

```
q.push(r);
    while(q.size()) {
        int v = q.front();
        q.pop();
        for (auto u : adj[v]) {
            if(v == r && u == deleted) continue;
//ignore deleted edge.
            if(distances[u] > distances[v] + 1) {
                 distances[u] = distances[v] + 1;
                q.push(u);
            }
        }
    }
}
int main() {
    ios base::sync with stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {</pre>
        int u, v; cin \gg u \gg v;
        adj[u].push back(v); adj[v].push back(u);
    int length = oo;
    for (int i=1; i<=n; i++) {
        for (auto u : adj[i]) {
            deleted = u;
            BFS(n, i);
            length = min(length, distances[u] + 1);
        }
    }
    if(length == oo) cout << "Graph Does Not Have</pre>
Cycle\n";
    else cout << "Minimum Cycle Length is : " <<</pre>
length << '\n';</pre>
```

## 4.8. topologycal sort int n; // number of vertices vector<vector<int>> adj; // adjacency list of graph vector<bool> visited; vector<int> ans; void dfs(int v) { visited[v] = true; for (int u : adj[v]) { if (!visited[u]) dfs(u); ans.push back(v); void topological sort() { visited.assign(n, false); ans.clear(); for (int i = 0; i < n; ++i) { if (!visited[i]) { dfs(i); reverse(ans.begin(), ans.end()); 4.9. lowest common Ancestor struct LCA { vector<int> height, euler, first, segtree; vector<bool> visited; int n; LCA(vector<vector<int>> &adj, int root = 0) { n = adj.size(); height.resize(n);

first.resize(n);
euler.reserve(n \* 2);

```
visited.assign(n, false);
        dfs(adj, root);
        int m = euler.size();
        segtree.resize(m * 4);
        build(1, 0, m - 1);
    }
    void dfs(vector<vector<int>> &adj, int node, int
h = 0) {
        visited[node] = true;
        height[node] = h;
        first[node] = euler.size();
        euler.push back(node);
        for (auto to : adj[node]) {
            if (!visited[to]) {
                 dfs(adj, to, h + 1);
                 euler.push back(node);
            }
    }
    void build(int node, int b, int e) {
        if (b == e) {
            segtree[node] = euler[b];
        } else {
            int mid = (b + e) / 2;
            build(node << 1, b, mid);</pre>
            build(node << 1 | 1, mid + 1, e);
            int 1 = segtree[node << 1], r =</pre>
segtree[node << 1 | 1];</pre>
            segtree[node] = (height[1] < height[r])</pre>
? 1 : r;
    }
    int query(int node, int b, int e, int L, int R)
        if (b > R | | e < L)
```

```
return -1;
        if (b >= L \&\& e <= R)
            return segtree[node];
        int mid = (b + e) \gg 1;
        int left = query(node << 1, b, mid, L, R);</pre>
        int right = query(node << 1 | 1, mid + 1, e,
L, R);
        if (left == -1) return right;
        if (right == -1) return left;
        return height[left] < height[right] ? left :</pre>
right;
    }
    int lca(int u, int v) {
        int left = first[u], right = first[v];
        if (left > right)
            swap(left, right);
        return query(1, 0, euler.size() - 1, left,
right);
};
```

## 4.10. lowest common Ancestor (binary lifting)

```
int n, 1;
vector<vector<int>> adj;
int timer;
vector<int> tin, tout;
vector<vector<int>> up;
void dfs(int v, int p)
    tin[v] = ++timer;
    up[v][0] = p;
    for (int i = 1; i <= 1; ++i)
```

```
up[v][i] = up[up[v][i-1]][i-1];
    for (int u : adj[v]) {
        if (u != p)
            dfs(u, v);
   }
    tout[v] = ++timer;
bool is ancestor(int u, int v)
    return tin[u] <= tin[v] && tout[u] >= tout[v];
int lca(int u, int v)
    if (is ancestor(u, v))
        return u;
    if (is ancestor(v, u))
        return v;
    for (int i = 1; i >= 0; --i) {
        if (!is_ancestor(up[u][i], v))
            u = up[u][i];
   }
   return up[u][0];
void preprocess(int root) {
    tin.resize(n);
    tout.resize(n);
    timer = 0;
   1 = ceil(log2(n));
    up.assign(n, vector<int>(l + 1));
    dfs(root, root);
```

## 4.11. hungarian algorithm (assignment problem)

```
vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
for (int i=1; i<=n; ++i) {
    p[0] = i;
    int j0 = 0;
    vector<int> minv (m+1, INF);
    vector<bool> used (m+1, false);
    do {
        used[j0] = true;
        int i0 = p[j0], delta = INF, j1;
        for (int j=1; j<=m; ++j)</pre>
            if (!used[i]) {
                int cur = A[i0][j]-u[i0]-v[j];
                if (cur < minv[j])</pre>
                     minv[j] = cur, way[j] = j0;
                if (minv[j] < delta)</pre>
                     delta = minv[j], j1 = j;
        for (int j=0; j<=m; ++j)</pre>
            if (used[j])
                u[p[j]] += delta, v[j] -= delta;
            else
                 minv[j] -= delta;
        j0 = j1;
    } while (p[j0] != 0);
    do {
        int j1 = way[j0];
        p[j0] = p[j1];
        j0 = j1;
    } while (j0);
4.12. 2SAT
```

```
struct TwoSatSolver {
   int n_vars;
    int n vertices;
```

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```
vector<vector<int>> adj, adj_t;
    vector<bool> used;
    vector<int> order, comp;
    vector<bool> assignment;
    TwoSatSolver(int n vars) : n vars( n vars),
n vertices(2 * n vars), adj(n vertices),
adj_t(n_vertices), used(n_vertices), order(),
comp(n vertices, -1), assignment(n vars) {
        order.reserve(n vertices);
   }
    void dfs1(int v) {
        used[v] = true;
        for (int u : adj[v]) {
            if (!used[u])
                dfs1(u);
        }
        order.push back(v);
    }
    void dfs2(int v, int cl) {
        comp[v] = cl;
        for (int u : adj t[v]) {
            if (comp[u] == -1)
                dfs2(u, c1);
        }
    }
    bool solve 2SAT() {
        order.clear();
        used.assign(n vertices, false);
        for (int i = 0; i < n vertices; ++i) {</pre>
            if (!used[i])
                dfs1(i);
        }
        comp.assign(n vertices, -1);
```

```
for (int i = 0, j = 0; i < n vertices; ++i)
{
            int v = order[n vertices - i - 1];
            if (comp[v] == -1)
                dfs2(v, j++);
        }
        assignment.assign(n_vars, false);
        for (int i = 0; i < n vertices; i += 2) {
            if (comp[i] == comp[i + 1])
                return false;
            assignment[i / 2] = comp[i] > comp[i +
1];
        }
        return true;
   }
    void add disjunction(int a, bool na, int b, bool
nb) {
        // na and nb signify whether a and b are to
be negated
        a = 2 * a ^ na;
        b = 2 * b ^ nb;
        int neg a = a ^ 1;
        int neg b = b ^ 1;
        adj[neg a].push back(b);
        adj[neg_b].push_back(a);
        adj_t[b].push_back(neg_a);
        adj t[a].push back(neg b);
   }
    static void example usage() {
        TwoSatSolver solver(3); // a, b, c
        solver.add disjunction(0, false, 1,
true); // a v not b
        solver.add disjunction(0, true, 1,
true); // not a v not b
```

```
int dfs(int v, vector<vector<int>> const& adj) {
   int size = 1;
   int max c size = 0;
   for (int c : adj[v]) {
       if (c != parent[v]) {
            parent[c] = v, depth[c] = depth[v] + 1;
            int c size = dfs(c, adj);
            size += c size;
            if (c size > max c size)
                max c size = c size, heavy[v] = c;
   }
   return size;
void decompose(int v, int h, vector<vector<int>>
const& adj) {
   head[v] = h, pos[v] = cur pos++;
   if (heavy[v] != -1)
       decompose(heavy[v], h, adj);
   for (int c : adj[v]) {
        if (c != parent[v] && c != heavy[v])
```

```
decompose(c, c, adj);
    }
}
void init(vector<vector<int>> const& adj) {
    int n = adj.size();
    parent = vector<int>(n);
    depth = vector<int>(n);
    heavy = vector<int>(n, -1);
    head = vector<int>(n);
    pos = vector<int>(n);
    cur pos = 0;
    dfs(0, adj);
    decompose(0, 0, adj);
int query(int a, int b) {
    int res = 0;
    for (; head[a] != head[b]; b = parent[head[b]])
        if (depth[head[a]] > depth[head[b]])
            swap(a, b);
        int cur heavy path max =
segment tree query(pos[head[b]], pos[b]);
        res = max(res, cur heavy path max);
    if (depth[a] > depth[b])
        swap(a, b);
    int last heavy path max =
segment tree query(pos[a], pos[b]);
    res = max(res, last heavy path max);
    return res;
}
```

## 5. Data Structures

## **5.1.** Array

b1.flip();

```
int main() {
    array<int, 5 > arr = \{1, 2, 3, 4, 5\};
    // Accessing elements
    cout << "Element at index 2: " << arr[2] <<</pre>
endl:
    // Size of the array
    cout << "Size of array: " << arr.size() << endl;</pre>
    // Fill array with a value
    arr.fill(10);
    cout << "Array after fill: ";</pre>
    for (int num : arr)
        cout << num << " ";</pre>
    cout << endl;</pre>
    return 0;
5.2. bitset
int main() {
    bitset<8> b1: // All bits initialized to 0
    bitset<8> b2("11001010");
    // Set bit at index 3 to 1
    b1.set(3);
    cout << "b1 after set(3): " << b1 << endl;</pre>
    // Reset all bits of b2
    b2.reset();
    cout << "b2 after reset: " << b2 << endl;</pre>
    // Flip all bits of b1
```

```
cout << "b1 after flip: " << b1 << endl;</pre>
    // Access the bit at index 2
    cout << "b1[2]: " << b1[2] << endl;</pre>
    // Test if bit at index 2 is set to 1
    if (b1.test(2))
    {
        cout << "Bit 2 is set to 1." << endl;</pre>
    }
    // Count the number of 1's in b1
    cout << "Number of 1's in b1: " << b1.count() <<</pre>
endl;
    return 0;
5.3. dequeue
int main() {
    deque<int> deq = \{1, 2, 3, 4, 5\};
    // Adding elements
    deq.push front(0); // Add at front
    deq.push_back(6); // Add at back
    // Removing elements
    deg.pop front(); // Remove from front
    deq.pop_back(); // Remove from back
    // Accessing elements
    cout << "First element: " << deq.front() <<</pre>
endl:
    cout << "Last element: " << deq.back() << endl;</pre>
    // Iterating through deque
    cout << "Deque elements: ";</pre>
```

```
for (int x : deq)
        cout << x << " ";
    cout << endl;</pre>
    // Size of deque
    cout << "Deque size: " << deq.size() << endl;</pre>
    return 0;
}
5.4. link list
int main() {
    list<int> lst = \{1, 2, 3, 4, 5\};
    // Adding elements
    lst.push back(6); // Add to the end
    lst.push front(0); // Add to the front
    // Removing elements
    lst.pop back(); // Remove from the end
    lst.pop front(); // Remove from the front
    // Iterating through list
    cout << "List elements: ";</pre>
    for (int x : lst)
        cout << x << " ";
    cout << endl;</pre>
    // Size of list
    cout << "List size: " << lst.size() << endl;</pre>
    return 0;
}
5.5. Map
int main() {
    map<string, int> m;
```

```
// Insert key-value pairs
    m["apple"] = 3;
    m["banana"] = 2;
    m["orange"] = 5;
    // Access value by key
    cout << "Value for apple: " << m["apple"] <<</pre>
endl;
    // Iterate through map
    cout << "Map elements: ";</pre>
    for (auto &pair : m)
    {
        cout << pair.first << ": " << pair.second <<</pre>
" ";
    cout << endl;</pre>
    // Remove an element
    m.erase("banana");
    cout << "Map after erase: ";</pre>
    for (auto &pair : m)
    {
        cout << pair.first << ": " << pair.second <<</pre>
    cout << endl;</pre>
    // Check if key exists
    if (m.find("orange") != m.end())
        cout << "Orange is in the map." << endl;</pre>
    return 0;
```

```
5.6. priority queue
int main() {
    priority_queue<int> pq;
    // Insert elements into the priority queue
    pq.push(10);
    pq.push(30);
    pq.push(20);
    cout << "Top element (max priority): " <<</pre>
pq.top() << endl;
    // Pop the top element
    pq.pop();
    cout << "Top element after pop: " << pq.top() <<</pre>
endl;
    // Size of the priority queue
    cout << "Size of priority queue: " << pq.size()</pre>
<< endl;
    return 0;
5.7. queue
int main() {
    queue<int> q;
    // Push elements onto the queue
    q.push(10);
    q.push(20);
    q.push(30);
    cout << "Front element: " << q.front() << endl;</pre>
    // Pop an element from the queue
    q.pop();
```

```
cout << "Front element after pop: " << q.front()</pre>
<< endl;
    // Size of the queue
    cout << "Size of queue: " << q.size() << endl;</pre>
    // Check if the queue is empty
    if (q.empty()) {
        cout << "Queue is empty." << endl;</pre>
    }
    else {
         cout << "Queue is not empty." << endl;</pre>
    }
    return 0;
5.8. set
int main() {
    set<int> s;
    // Insert elements
    s.insert(10);
    s.insert(20);
    s.insert(15);
    // Print elements in sorted order
    cout << "Set elements: ";</pre>
    for (int num : s)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Check if an element exists
    if (s.find(15) != s.end())
         cout << "15 is in the set." << endl:</pre>
    }
```

```
// Remove an element
    s.erase(10);
    cout << "Set after erase: ";</pre>
    for (int num : s)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Size of the set
    cout << "Size of set: " << s.size() << endl;</pre>
    return 0;
5.9. stack
int main() {
    stack<int> s;
    // Push elements onto the stack
    s.push(10);
    s.push(20);
    s.push(30);
    cout << "Top element after push: " << s.top() <<</pre>
endl;
    // Pop an element from the stack
    5.pop();
    cout << "Top element after pop: " << s.top() <<</pre>
endl;
    // Check if stack is empty
    if (s.empty()) {
        cout << "Stack is empty." << endl;</pre>
    }
    else {
        cout << "Stack is not empty." << endl;</pre>
```

```
// Size of the stack
    cout << "Size of stack: " << s.size() << endl;</pre>
    return 0;
5.10. unordered map
int main() {
    unordered map<int, string> um;
    // Inserting elements
    um[1] = "apple";
    um[2] = "banana";
    um[3] = "cherry";
    // Accessing elements
    cout << "Key 2 maps to: " << um[2] << endl;</pre>
    // Checking if key exists
    if (um.find(4) == um.end())
        cout << "Key 4 not found!" << endl;</pre>
    // Iterating through unordered map
    cout << "Unordered map elements: ";</pre>
    for (auto &pair : um)
        cout << pair.first << " -> " << pair.second</pre>
<< " | ";
    cout << endl;</pre>
    return 0;
5.11. unordered set
int main() {
```

```
unordered set<int> us;
    // Insert elements
    us.insert(10);
    us.insert(20);
    us.insert(15);
    // Print elements
    cout << "Unordered set elements: ";</pre>
    for (int num : us)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Check if an element exists
    if (us.find(15) != us.end())
        cout << "15 is in the unordered set." <<</pre>
endl;
    }
    // Remove an element
    us.erase(10);
    cout << "Unordered set after erase: ";</pre>
    for (int num : us)
        cout << num << " ";</pre>
    cout << endl;</pre>
    return 0;
}
5.12. vector
int main() {
    // Create a vector of integers
    vector<int> v;
    // Add elements using push_back
    v.push back(10);
```

```
v.push back(20);
    v.push back(30);
    cout << "Vector after push back: ";</pre>
    for (int num : v)
        cout << num << " ";
    cout << endl;</pre>
    // Accessing elements using at() and indexing
    cout << "Element at index 1: " << v.at(1) <<</pre>
endl:
    cout << "Element at index 0: " << v[0] << endl;</pre>
    // Pop an element from the back
    v.pop back();
    cout << "Vector after pop back: ";</pre>
    for (int num : v)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Insert an element at a specific position
    v.insert(v.begin() + 1, 25); // Insert 25 at
index 1
    cout << "Vector after insert: ";</pre>
    for (int num : v)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Remove an element from the vector
    v.erase(v.begin() + 1); // Remove the element at
index 1
    cout << "Vector after erase: ";</pre>
    for (int num : v)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Resize the vector
    v.resize(5, 50); // Resize to size 5, fill new
elements with 50
```

```
cout << "Vector after resize: ";
for (int num : v)
      cout << num << " ";
cout << endl;

// Get the size of the vector
cout << "Size of vector: " << v.size() << endl;

// Clear the vector
v.clear();
cout << "Vector after clear: " << v.size() << "
(size is now zero)" << endl;

return 0;
}</pre>
```

## 5.13. DS cheatsheet

#### ## 1. Vector

- Description: Dynamic array that allows fast random access.
- Methods:
- push back(x): Add element x to the end.
- pop back(): Remove the last element.
- at(i): Access the element at index i (bounds checked).
- operator[]: Access the element at index i (no bounds check).
- size(): Return the number of elements.
- empty(): Check if the vector is empty.
- resize(n, x): Resize the vector to size n and fill new elements with x.
- clear(): Remove all elements.

#### ## 2. Stack

- Description: Last-In-First-Out (LIFO) structure, used for backtracking problems.
- Methods:
- push(x): Add element x to the top.
- pop(): Remove the top element.
- top(): Get the top element.
- size(): Return the number of elements.
- empty(): Check if the stack is empty.

#### ## 3. Queue

#### CTRL+ALT+DEFEAT – Hakim Sabzevari University

- Description: First-In-First-Out (FIFO) structure, ideal for problems involving processing in order.
- Methods:
- push(x): Add element x to the back.
- pop(): Remove the front element.
- front(): Get the front element.
- back(): Get the back element.
- size(): Return the number of elements.
- empty(): Check if the queue is empty.

#### ## 4. Priority Queue (Max-Heap by default)

- Description: A heap-based structure that always gives the maximum element.
- Methods:
- push(x): Add element x to the queue.
- pop(): Remove the largest element.
- top(): Get the largest element.
- size(): Return the number of elements.
- empty(): Check if the queue is empty.

#### ## 5. Set

- Description: Collection of unique elements in sorted order.
- Methods:
  - insert(x): Add element x.
  - erase(x): Remove element x.
  - find(x): Check if element x exists.
  - size(): Return the number of elements.
  - empty(): Check if the set is empty.
  - clear(): Remove all elements.

#### ## 6. Map

- Description: Stores key-value pairs in sorted order based on keys.
- Methods:
- insert({key, value}): Add key-value pair.
- erase(key): Remove element by key.
- find(key): Check if a key exists.
- operator[]: Access the value associated with a key.
- size(): Return the number of elements.
- empty(): Check if the map is empty.
- clear(): Remove all elements.

#### ## 7. Unordered Set

- Description: Collection of unique elements with no specific order.
- Methods:

- insert(x): Add element x.
- erase(x): Remove element x.
- find(x): Check if element x exists.
- size(): Return the number of elements.
- empty(): Check if the unordered set is empty.

#### ## 8. Unordered Map

- Description: Stores key-value pairs with no specific order.
- Methods
- insert({key, value}): Add key-value pair.
- erase(key): Remove element by key.
- find(key): Check if a key exists.
- operator[]: Access the value associated with a key.
- size(): Return the number of elements.
- empty(): Check if the unordered map is empty.

#### ## 9. Bitset

- Description: A space-efficient container for a fixed-size sequence of bits (0 or 1).
- Methods:
- set(i): Set bit at index i to 1.
- reset(i): Set bit at index i to 0.
- flip(i): Toggle the bit at index i.
- test(i): Check if the bit at index i is 1.
- count(): Count the number of bits set to 1.
- size(): Return the number of bits.
- operator[]: Access the bit at index i.
- to\_string(): Convert bitset to string.

#### ## 10. Array

- Description: Fixed-size array used for fast access, but size cannot be changed after initialization.
- Methods:
- fill(x): Fill all elements with the value x.
- size(): Return the number of elements.
- operator[]: Access element at index i.
- at(i): Access element at index i with bounds checking.
- front(): Get the first element.
- back(): Get the last element.

#### ## 11. Deque

- Description: Double-ended queue that allows fast insertion and removal at both ends.
- Methods:
- push\_front(x): Add element x to the front.
- push\_back(x): Add element x to the back.

- pop front(): Remove the front element.
- pop back(): Remove the back element.
- front(): Get the front element.
- back(): Get the back element.
- size(): Return the number of elements.
- empty(): Check if the deque is empty.

#### ## 12. Linked List (Using STL List)

- Description: Doubly linked list that allows fast insertion and deletion at both ends.
- Methods:
- push\_back(x): Add element x to the back.
- push front(x): Add element x to the front.
- pop back(): Remove the last element.
- pop\_front(): Remove the first element.
- size(): Return the number of elements.
- empty(): Check if the list is empty.
- front(): Get the first element.
- back(): Get the last element.
- clear(): Remove all elements.

# 6. Dynamic Programming

## 6.1. counting paths matrix

#### 6.2. edit distance

```
// Function to compute the Edit Distance
int editDistance(string str1, string str2, int m,
int n) {
   vector<vector<int>> dp(m + 1, vector<int>(n +
1));
    for (int i = 0; i <= m; i++) {
        for (int j = 0; j <= n; j++) {
            if (i == 0)
                dp[i][j] = j;
            else if (j == 0)
                dp[i][j] = i;
            else if (str1[i - 1] == str2[j - 1])
                dp[i][j] = dp[i - 1][j - 1];
            else
                dp[i][j] = 1 + min({dp[i - 1][j -
1], dp[i][j - 1], dp[i - 1][j]});
    }
```

```
return dp[m][n];
}
int main() {
    string str1 = "sitting", str2 = "kitten";
    int m = str1.length(), n = str2.length();

    cout << "Edit Distance: " << editDistance(str1, str2, m, n) << endl;
    return 0;
}</pre>
```

## 6.3. Egg Dropping

```
// Function to find the minimum number of attempts
needed
int eggDrop(int eggs, int floors)
    vector<vector<int>> dp(eggs + 1,
vector<int>(floors + 1, 0));
    for (int i = 1; i \leftarrow eggs; i++)
        dp[i][0] = 0;
    for (int j = 0; j <= floors; j++)</pre>
        dp[1][j] = j;
    for (int i = 2; i <= eggs; i++)
        for (int j = 2; j \leftarrow floors; i++)
             dp[i][j] = INT_MAX;
            for (int x = 1; x <= j; x++)
                 dp[i][j] = min(dp[i][j], 1 +
\max(dp[i - 1][x - 1], dp[i][j - x]));
```

```
return dp[eggs][floors];
}
int main()
{
   int eggs = 2, floors = 10;
   cout << "Minimum attempts: " << eggDrop(eggs, floors) << endl;
   return 0;
}</pre>
```

## 6.4. fibonacci

```
// Fibonacci sequence using dynamic programming
(Memoization)
// Function to compute Fibonacci number
int fib(int n, vector<int> &dp)
   // Base cases
   if (n <= 1)
       return n;
   // If the value is already computed, return it
   if (dp[n] != -1)
        return dp[n];
   // Store the computed value in dp array
   dp[n] = fib(n - 1, dp) + fib(n - 2, dp);
   return dp[n];
int main()
   int n = 10;
   vector<int> dp(n + 1, -1); // Initialize dp
array with -1
```

```
cout << "Fibonacci of " << n << " is " << fib(n,
dp) << endl;
    return 0;
}</pre>
```

```
6.5. knapsack 01
// Function to solve the 0/1 Knapsack problem
int knapsack(int W, vector<int> &wt, vector<int>
&val, int n)
    vector<vector<int>> dp(n + 1, vector<int>(W + 1,
0)); // DP table
    for (int i = 1; i <= n; i++)
        for (int w = 0; w \le W; w++)
            if (wt[i - 1] <= w)
                dp[i][w] = max(dp[i - 1][w], val[i -
1] + dp[i - 1][w - wt[i - 1]]);
            else
                dp[i][w] = dp[i - 1][w];
    return dp[n][W];
int main()
    int W = 100:
                                            //
Capacity of the knapsack
    vector<int> val = {80, 24, 23, 22, 21}; //
Values of the items
    vector<int> wt = {80, 25, 25, 25, 25}; //
Weights of the items
    int n = val.size();
                                            //
Number of items
```

```
cout << "Maximum value in knapsack: " <<</pre>
knapsack(W, wt, val, n) << endl;</pre>
    return 0;
6.6. LCS
// Function to compute the length of the Longest
Common Subsequence
int lcs(string X, string Y, int m, int n)
    vector<vector<int>> dp(m + 1, vector<int>(n + 1,
0)); // DP table
    for (int i = 1; i <= m; i++)
        for (int j = 1; j <= n; j++)
            if (X[i - 1] == Y[j - 1])
                dp[i][j] = 1 + dp[i - 1][j - 1];
            else
                dp[i][j] = max(dp[i - 1][j], dp[i][j]
- 1]);
    return dp[m][n];
int main()
    string X = "AGGTAB", Y = "GXTXAYB";
    int m = X.length(), n = Y.length();
    cout << "Length of Longest Common Subsequence:</pre>
<< lcs(X, Y, m, n) << endl;
    return 0;
```

## 6.7. LIS

```
// Function to find the length of the Longest
Increasing Subsequence
int lis(vector<int> &arr, int n)
    vector<int> dp(n, 1); // DP array, initialized
to 1
    for (int i = 1; i < n; i++)
        for (int j = 0; j < i; j++)
            if (arr[i] > arr[j])
                 dp[i] = max(dp[i], dp[j] + 1);
    return *max_element(dp.begin(), dp.end());
int main()
    vector\langle int \rangle arr = {10, 22, 9, 33, 21, 50, 41,
60);
    int n = arr.size();
    cout << "Length of Longest Increasing</pre>
Subsequence: " << lis(arr, n) << endl;</pre>
    return 0;
6.8. LPS
// Function to compute the length of the Longest
```

```
// Function to compute the length of the Longes
Palindromic Subsequence
int lps(string s)
{
  int n = s.length();
```

```
vector<vector<int>> dp(n, vector<int>(n, 0));
    for (int i = 0; i < n; i++)</pre>
        dp[i][i] = 1; // Single character is a
palindrome
    for (int len = 2; len <= n; len++)</pre>
    {
        for (int i = 0; i < n - len + 1; i++)
            int j = i + len - 1;
            if (s[i] == s[j])
                 dp[i][j] = 2 + dp[i + 1][j - 1];
            else
                 dp[i][j] = max(dp[i + 1][j], dp[i][j]
- 1]);
    return dp[0][n - 1];
}
int main()
    string s = "bbabcbcab";
    cout << "Length of Longest Palindromic</pre>
Subsequence: " << lps(s) << endl;</pre>
    return 0:
}
6.9. MCM
```

```
// Function to compute the minimum number of scalar
multiplications
int matrixChainMultiplication(vector<int> &dims, int
n) {
    vector<vector<int>> dp(n, vector<int>(n, 0)); //
DP table
```

```
for (int len = 2; len < n; len++)</pre>
        for (int i = 1; i < n - len + 1; i++)
            int j = i + len - 1;
            dp[i][j] = INT MAX;
            for (int k = i; k < j; k++)
                int q = dp[i][k] + dp[k + 1][j] +
dims[i - 1] * dims[k] * dims[j];
                dp[i][j] = min(dp[i][j], q);
        }
    return dp[1][n - 1];
int main()
    vector<int> dims = {10, 20, 30, 40, 30};
    int n = dims.size();
    cout << "Minimum number of scalar</pre>
multiplications: " <<</pre>
matrixChainMultiplication(dims, n) << endl;</pre>
    return 0;
```

## 6.10. Minimum Coin Change

```
// Function to find the minimum number of coins
required to make a total
int minCoins(const vector<int> &coins, int total)
{
   int n = coins.size();
   vector<int> dp(total + 1, INT_MAX); //
Initialize DP array with infinity
```

```
dp[0] = 0;
                                         // Base
case: 0 coins needed to make total 0
    // Fill DP array
    for (int i = 0; i < n; i++)
        for (int j = coins[i]; j <= total; j++)</pre>
            if (dp[j - coins[i]] != INT MAX)
                 dp[j] = min(dp[j], dp[j - coins[i]]
+ 1); // Minimize coin count
    }
    return dp[total] == INT MAX ? -1 : dp[total]; //
Return -1 if not possible
int main()
    vector<int> coins = {1, 2, 5};
                                              //
Example denominations
    int total = 11;
                                              //
Target total
    cout << minCoins(coins, total) << endl; //</pre>
Output the result
    return 0;
```

## 6.11. optimal BST

```
// Function to calculate the minimum search cost for
an optimal BST
int optimalBST(const vector<int> &freq)
{
   int n = freq.size();
```

```
vector<vector<int>> dp(n, vector<int>(n, 0));
    // Fill the DP table for subarrays of increasing
length
    for (int len = 1; len <= n; len++)</pre>
    { // len is the range length
        for (int i = 0; i <= n - len; i++)
            int j = i + len - 1;
            dp[i][j] = INT MAX;
            int sum = 0;
            // Calculate sum of frequencies from i
to j
            for (int k = i; k \leftarrow j; k++)
                sum += freq[k];
            // Try each k as the root and calculate
the minimum cost
            for (int k = i; k \le j; k++)
                int cost = (k == i ? 0 : dp[i][k -
1]) + (k == j ? 0 : dp[k + 1][j]);
                dp[i][j] = min(dp[i][j], cost +
sum);
            }
    return dp[0][n - 1]; // The minimum cost for the
entire range
int main()
    vector<int> freq = {34, 8, 50, 13}; // Example
frequencies of keys
    cout << optimalBST(freq) << endl; // Output</pre>
the minimum cost
```

```
return 0:
6.12. partition problem
// Function to determine if a given set can be
partitioned into two subsets
bool canPartition(vector<int> &nums)
    int sum = 0;
    for (int num : nums)
        sum += num;
    if (sum % 2 != 0)
        return false;
    int target = sum / 2;
    vector<bool> dp(target + 1, false);
    dp[0] = true;
    for (int num : nums)
        for (int j = target; j >= num; j--)
            dp[j] = dp[j] \mid \mid dp[j - num];
    return dp[target];
int main()
    vector<int> nums = {1, 5, 11, 5};
    cout << "Can partition: " << (canPartition(nums)</pre>
? "Yes" : "No") << endl;
    return 0;
```

## 6.13. Regular Expression Matching

```
bool isMatch(const string &s, const string &p) {
    int m = s.size(), n = p.size();
   // DP table dp[i][j] will be true if s[0...i-1]
matches p[0...j-1]
    vector<vector<bool>> dp(m + 1, vector<bool>(n +
1, false));
   // Base case: empty string matches empty pattern
    dp[0][0] = true;
    // Handle patterns like "a*" or ".*" where "*"
can match 0 occurrence
    for (int j = 1; j <= n; j++)
        if (p[j - 1] == '*')
            dp[0][j] = dp[0][j - 2];
    // Fill the dp table
    for (int i = 1; i <= m; i++)
        for (int j = 1; j <= n; j++)
            if (p[j-1] == s[i-1] || p[j-1] ==
'.')
                dp[i][j] = dp[i - 1][j - 1]; //
Character matches
            else if (p[j - 1] == '*')
                // '*' matches zero occurrence or
one/more occurrences of the preceding character
```

```
dp[i][j] = dp[i][j - 2] || (dp[i -
1][j] && (s[i - 1] == p[j - 2] || p[j - 2] == '.'));
    }
    return dp[m][n]; // Final answer whether the
whole string matches the pattern
int main()
    string s = "aab";
    string p = "c*a*b";
    if (isMatch(s, p))
        cout << "The string matches the pattern." <<</pre>
endl;
    }
    else
        cout << "The string does not match the</pre>
pattern." << endl;</pre>
    }
    return 0:
6.14. ROD cutting
// Function to compute the maximum profit from
cutting a rod
int rodCutting(vector<int> &prices, int n)
    vector<int> dp(n + 1, 0); // DP array
```

```
for (int i = 1; i <= n; i++)
```

```
for (int j = 1; j <= i; j++)
            dp[i] = max(dp[i], prices[j - 1] + dp[i]
- j]);
        }
    return dp[n];
int main()
    vector<int> prices = {1, 5, 8, 9, 10, 17, 17,
20}:
    int n = prices.size();
    cout << "Maximum profit from rod cutting: " <<</pre>
rodCutting(prices, n) << endl;</pre>
    return 0;
```

## 6.15. subset sum

```
// Function to determine if there's a subset with
sum equal to the target
bool subsetSum(vector<int> &nums, int sum)
    int n = nums.size();
    vector<vector<bool>> dp(n + 1, vector<bool>(sum
+ 1, false));
    for (int i = 0; i <= n; i++)
        dp[i][0] = true;
    for (int i = 1; i <= n; i++)
        for (int j = 1; j \leftarrow sum; j++)
```

```
if (nums[i - 1] <= j)</pre>
                dp[i][j] = dp[i - 1][j] || dp[i -
1][j - nums[i - 1]];
            else
                dp[i][j] = dp[i - 1][j];
        }
    return dp[n][sum];
int main()
    vector<int> nums = {3, 34, 4, 12, 5, 2};
    int sum = 9;
    cout << "Subset sum possible: " <<</pre>
(subsetSum(nums, sum) ? "Yes" : "No") << endl;
    return 0;
6.16. two player game
```

```
// Function to find the maximum sum player A can get
int maxCoins(const vector<int> &coins)
    int n = coins.size();
    vector<vector<int>> dp(n, vector<int>(n, 0));
    // Base case: when there's only one coin, player
A takes it
    for (int i = 0; i < n; i++)
        dp[i][i] = coins[i];
    // Fill DP table for subarrays of length 2 to n
    for (int len = 2; len <= n; len++)</pre>
```

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```
for (int i = 0; i < n - len + 1; i++)
            int j = i + len - 1;
            dp[i][j] = max(coins[i] + min(dp[i +
2][j], dp[i + 1][j - 1]),
                           coins[i] + min(dp[i +
1][j - 1], dp[i][j - 2]));
    }
    return dp[0][n - 1]; // Maximum sum player A can
get
int main()
    vector<int> coins = {8, 15, 3, 7}; // Example
coins array
    cout << maxCoins(coins) << endl; // Output the</pre>
result
    return 0:
```

## 6.17. word break

## 6.18. dp cheatsheet

## 1. Fibonacci Sequence (Memoization)

- Explanation: This algorithm calculates the nth Fibonacci number using memoization to store previously computed results, avoiding redundant calculations.
- When to Use: When you need to compute Fibonacci numbers for large values of `n` efficiently (i.e., for recursive problems that involve overlapping subproblems).

#### ## 2. 0/1 Knapsack Problem

- Explanation: Given a set of items with weights and values, the goal is to determine the maximum value that can be obtained by putting items in a knapsack without exceeding the weight capacity.
- When to Use: In optimization problems where you need to maximize profit or value while respecting constraints (e.g., weight, space).

#### ## 3. Minimum Coin Change

- Explanation: This algorithm calculates the minimum number of coins needed to make a given total using a set of coin denominations. It uses dynamic programming to store the results for all possible totals.
- When to Use: When you need to find the fewest coins needed to form a specific amount (e.g., for making change, budget optimization).

#### ## 4. Longest Common Subsequence (LCS)

- Explanation: This algorithm finds the longest subsequence common to two sequences (strings, arrays, etc.). It uses a dynamic programming table to store intermediate results.
- When to Use: When comparing two sequences (e.g., DNA sequences, text comparison, diff tools) and need the longest subsequence they share.

#### ## 5. Edit Distance (Levenshtein Distance)

- Explanation: This algorithm calculates the minimum number of operations (insertions, deletions, substitutions) needed to convert one string into another.
- When to Use: When you need to compare two strings and find the minimum edit operations required (e.g., spell checkers, natural language processing).

#### ## 6. Longest Increasing Subsequence (LIS)

- Explanation: This algorithm finds the length of the longest increasing subsequence in a sequence of numbers. The subsequence need not be contiguous.
- When to Use: When you need to find the longest increasing subsequence in a sequence of numbers (e.g., stock price prediction, finding trends).

#### ## 7. Matrix Chain Multiplication

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- Explanation: This algorithm calculates the most efficient way to multiply a chain of matrices by minimizing the number of scalar multiplications.
- When to Use: When multiplying multiple matrices and you need to minimize the cost of multiplication (e.g., in computer graphics, optimization problems).

#### ## 8. Rod Cutting Problem

- Explanation: This algorithm finds the maximum profit you can obtain by cutting a rod of length `n` into smaller pieces and selling them, based on the prices for each length.
- When to Use: When you need to solve problems related to cutting materials into pieces to maximize profit (e.g., resource allocation, profit optimization).

#### ## 9. Subset Sum Problem

- Explanation: This algorithm checks if there is a subset of a given set of numbers that adds up to a target sum. It uses dynamic programming to keep track of achievable sums.
- When to Use: When you need to check whether a subset exists with a given sum (e.g., partitioning problems, subset analysis).

#### ## 10. Egg Dropping Problem

- Explanation: This algorithm determines the minimum number of attempts required to find the highest floor from which an egg can be dropped without breaking, given 'k' eggs and 'n' floors.
- When to Use: In optimization problems where you need to minimize the number of trials in a worst-case scenario (e.g., testing, fault tolerance, hardware).

#### ## 11. Partition Problem

- Explanation: This algorithm checks whether a given set can be partitioned into two subsets such that their sums are equal. It uses dynamic programming to check for subset sums.
- When to Use: When dividing a set of numbers into two equal subsets (e.g., load balancing, resource allocation).

#### ## 12. Longest Palindromic Subsequence (LPS)

- Explanation: This algorithm finds the longest subsequence within a string that is a palindrome. It uses dynamic programming to build a table based on matching characters.
- When to Use: When you need to find the longest palindromic subsequence in a string (e.g., text processing, bioinformatics).

#### ## 13. Word Break Problem

- Explanation: This algorithm checks if a string can be segmented into a space-separated sequence of words from a dictionary. It uses dynamic programming to store results for substrings.
- When to Use: When you need to determine whether a string can be split into valid words (e.g., for tokenizing sentences, text segmentation).

#### ## 14. Regular Expression Matching

- Explanation: This algorithm checks if a string matches a pattern with . (any character) and \\* (zero or more of the previous character). It uses dynamic programming to track matching results for substrings.
- When to Use: When matching strings to patterns with wildcards like . and  $\$ \* (e.g., text matching, search engines, file pattern matching).

#### ## 15. Optimal Binary Search Tree

- Explanation: This algorithm finds the minimum cost to construct a binary search tree based on the frequencies of elements. It uses dynamic programming to calculate the optimal cost for different subarrays.

- When to Use: When you need to minimize the cost of searching elements with different access frequencies (e.g., database indexing, search optimization).

#### ## 16. 2 Player Game

- Explanation: This algorithm calculates the maximum sum player A can collect in a turn-based game where players alternate picking coins from either end of the array. It uses dynamic programming to determine the optimal strategy for player A.
- When to Use: When you need to calculate the optimal strategy in a turn-based game with alternating choices (e.g., maximizing outcomes in competitive games).

#### ## 17. Counting Paths in Matrix

- Explanation: This algorithm calculates the number of unique paths from the top-left corner to the bottom-right corner of an n x m matrix. It uses dynamic programming to count paths by combining the results from adjacent cells.
- When to Use: When you need to find the number of distinct paths in a grid, where you can only move right or down (e.g., grid-based traversal problems).

## 7. More

# 7.1. longest polindrom substring (Mancher)

```
// O(n)

class Solution {
public:
    string longestPalindrome(std::string s) {
        if (s.length() <= 1) return s;

        // Preprocess the string with '#' characters
to handle even-length palindromes
        string modified s = "#";</pre>
```

```
for (char c : s) {
            modified s += c;
            modified s += '#';
        int n = modified_s.size();
        vector<int> dp(n, 0); // dp array to store
the radius of the palindrome centered at each
character
        int center = 0, right = 0; // Initialize
center and right boundary
        int max len = 1; // Maximum length of
palindrome found
        string max str = s.substr(0, 1); //
Initialize the max palindrome substring with the
first character
        for (int i = 0; i < n; i++) {
            // If i is within the current right
boundary, use previously calculated values to
minimize comparisons
            if (i < right) {</pre>
                dp[i] = min(right - i, dp[2 * center]
- i]);
            }
            // Expand around center i
            while (i - dp[i] - 1 >= 0 \&\& i + dp[i] +
1 < n \&\& modified s[i - dp[i] - 1] == modified s[i +
dp[i] + 1]) {
                dp[i]++;
            }
            // Update center and right boundary if
we've expanded beyond the current right
            if (i + dp[i] > right) {
                center = i;
                right = i + dp[i];
```

```
// Update max_len and max_str if a
longer palindrome is found
    if (dp[i] > max_len) {
        max_len = dp[i];
        max_str = modified_s.substr(i -
dp[i], 2 * dp[i] + 1);
        max_str.erase(remove(max_str.begin()),
        max_str.end(), '#'), max_str.end());
    }
}
return max_str;
}
```

## 7.2. median of 2 soted array

```
class Solution {
public:
    double findMedianSortedArrays(vector<int>
&nums1, vector<int> &nums2) {
        int n = nums1.size();
        int m = nums2.size();
        if(n > m) return
findMedianSortedArrays(nums2, nums1);
        int size = n + m;
        int left = (size + 1) / 2;
        int low = 0, high = n;
        int 11, 12, r1, r2, mid1, mid2;
        while(low <= high) {</pre>
            mid1 = (low + high) >> 1;
            mid2 = left - mid1:
            r1 = mid1 < n ? nums1[mid1] : 1e9;
            r2 = mid2 < m ? nums2[mid2] : 1e9;
            11 = mid1 - 1 >= 0 ? nums1[mid1 - 1] : -
1e9:
```

## 10. other

## 10.1. useful geo

Area of triangle with sides a, b, c: sqrt(S \*(S-a)\*(S-b)\*(S-c)) where S = (a+b+c)/2

Area of equilateral triangle:  $s^2 * sqrt(3) / 4$  where is side lenght

Pyramid and cones volume: 1/3 area(base) \* height

if p1=(x1, x2), p2=(x2, y2), p3=(x3, y3) are points on circle, the center is  $x = -((x2^2 - x1^2 + y2^2 - y1^2)^*(y3 - y2) - (x2^2 - x3^2 + y2^2 - y3^2)^*(y1 - y2)) / (2^*(x1 - x2)^*(y3 - y2) - 2^*(x3 - x2)^*(y1 - y2))$   $y = -((y2^2 - y1^2 + x2^2 - x1^2)^*(x3 - x2) - (y2^2 - y3^3 + x2^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y2^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y2^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y2^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y1^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y1^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y1^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y1^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y1^2 - y1^2 - y1^2 + x2^2 - y1^2)^*(x3 - x2) - (y1^2 - y1^2 - y1^2)^*(x1^2 - y1^2)$ 

 $x3^2$ (x1 - x2)) / (2\*(y1 - y2)\*(x3 - x2) - 2\*(y3 - y2)\*(x1 - x2))

## 10.2. number of primes

30: 10 60: 17 100: 25 1000: 168 10000: 1229 100000: 9592 1000000: 78498

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10000000: 664579

## 10.3. Factorials

1: 1 2: 2

2 6

3: 6 4: 24

5: 120

6: 720

7: 5040

8: 40320

9: 362880

10: 3628800

11: 39916800

12: 479001600 13: 6227020800

14: 87178291200

15: 1307674368000

## 10.4. power of 3

1: 3

2: 9 3: 27

4: 81

5: 243

6: 729

7: 2187

8: 6561

9: 19683

10: 59049

11: 177147

12: 531441

13: 1594323 14: 4782969

15: 14348907

16: 43046721

17: 129140163

18: 387420489

19: 1162261467

20: 3486784401

## 10.5. C(2n, n)

1: 2

2: 6

3: 20

4: 70

5: 252

6: 924

7: 3432

8: 12870

9: 48620

10: 184756

11: 705432

12: 2704156

13: 10400600

14: 40116600

15: 155117520

## 10.6. Most Divisor

<= 1e2: 60 with 12 divisors

<= 1e3: 840 with 32 divisors

<= 1e4: 7560 with 64 divisors

<= 1e5: 83160 with 128 divisors

<= 1e6: 720720 with 240 divisors

<= 1e7: 8648640 with 448 divisors

<= 1e8: 73513440 with 768 divisors

<= 1e9: 735134400 with 1344 divisors

<= 1e10: 6983776800 with 2304 divisors

<= 1e11: 97772875200 with 4032 divisors

<= 1e12: 963761198400 with 6720 divisors

<= 1e13: 9316358251200 with 10752 divisors

<= 1e14: 97821761637600 with 17280 divisors

<= 1e15: 866421317361600 with 26880 divisors

<= 1e16: 8086598962041600 with 41472 divisors

<= 1e17: 74801040398884800 with 64512 divisors

<= 1e18: 897612484786617600 with 103680 divisors