

Hakim Sabzevari University

CTRL+ALT+DEFEAT

Team Reference Document

Ali Ghanbari, Amirreza Zeraati, Rahmat Ansari

https://github.com/ctrl-alt-Defeat-icpc

List of content

1. STL	2
1.1. bitscroll	2
1.2. 128 bit	2
2. Segment Tree	3
2.1. easy implementation	3
2.2. with lazy propagation	3
3. Math	4
3.1. choose	4
3.2. gcd	4
3.3. compressing	4
3.4. lower bound and upper bound	4
4. Graph	
4.1. BFS	4
4.2. bipartite	
4.3. cycle finding	
4.4. DFS	
4.5. floyd-warshall	
4.6. prim	
4.7. shortest cycle	
4.8. topologycal sort	
4.9. lowest common Ancestor	
4.10. lowest common Ancestor (binary lifting)	
4.11. hungarian algorithm (assignment problem)	
4.12. 2SAT	
4.13. Heavy-light decomposition	
5. Data Structures	
5.1. Array	
5.2. bitset	
5.3. dequeue	
5.4. link list	
5.5. Map	11

	5.7. queue	12
	5.8. set	12
	5.9. stack	12
	5.10. unordered map	12
	5.11. unordered set	13
	5.12. vector	13
	5.13. DS cheatsheet	13
6.	Dynamic Programming	14
	6.1. counting paths matrix	14
	6.2. edit distance	15
	6.3. Egg Dropping	15
	6.4. fibonacci	15
	6.5. knapsack 01	16
	6.6. LCS	16
	6.7. LIS	16
	6.8. LPS	17
	6.9. MCM	17
	6.10. Minimum Coin Change	17
	6.11. optimal BST	18
	6.12. partition problem	18
	6.13. Regular Expression Matching	18
	6.14. ROD cutting	19
	6.15. subset sum	19
	6.16. two player game	20
	6.17. word break	20
	6.18. dp cheatsheet	20
7.	More	22
	7.1. longest polindrom substring (Mancher)	22
	7.2. median of 2 soted array	22
10	. other	23
	10.1. useful geo	23

10.2. number of primes	23
10.3. Factorials	23
10.4. power of 3	23
10.5. C(2n, n)	23
10.6. Most Divisor	23

1. STL

1.1. bitscroll

```
__builtin_ctz(x); // first 1 from left (index)
__builtin_popcount(x); // count of 1 in numbers bit
__builtin_ctzll(x); // for long long
__builtin_popcountll(x); // ...
```

1.2. 128 bit

```
__int128 read() {
    _{\text{int128}} x = 0, f = 1;
    char ch = getchar();
    while (ch < '0' || ch > '9') {
        if (ch == '-') f = -1;
        ch = getchar();
    while (ch >= '0' && ch <= '9') {
        x = x * 10 + ch - '0';
        ch = getchar();
    return x * f;
void print(__int128 x) {
    if (x < 0) {
        putchar('-');
        x = -x;
    if (x > 9) print(x / 10);
    putchar(x % 10 + '0');
```

```
bool cmp(__int128 x, __int128 y) { return x > y; }
int main() {
    __int128 x = read();
    print(x);
    cout << endl;</pre>
    return 0;
```

2. Segment Tree

2.1. easy implementation

```
const int N = 1e5; // limit for array size
int n; // array size
int t[2 * N];
void build() { // build the tree
 for (int i = n - 1; i > 0; --i) t[i] = t[i << 1] +
t[i<<1|1];
}
void modify(int p, int value) { // set value at
position p
 for (t[p += n] = value; p > 1; p >>= 1) t[p>>1] =
t[p] + t[p^1];
}
int query(int 1, int r) { // sum on interval [1, r)
  int res = 0;
  for (1 += n, r += n; 1 < r; 1 >>= 1, r >>= 1) {
   if (1&1) res += t[1++];
   if (r\&1) res += t[--r];
  return res;
int main() {
```

```
scanf("%d", &n);
  for (int i = 0; i < n; ++i) scanf("%d", t + n +</pre>
i);
 build();
 modify(0, 2);
 printf("%d\n", query(3, 11));
 return 0;
```

2.2. with lazy propagation

```
const int N = 1e5 + 5;
int n;
int seg[2 * N], lazy[2 * N], a[N];
int segSize;
void build(int u = 1, int ul = 0, int ur = n) {
    if(ur - ul < 2)
        seg[u] = a[ul];
        return;
    }
    int mid = (ul + ur) / 2;
    build(u * 2, ul, mid);
    build(u * 2 + 1, mid, ur);
    seg[u] = seg[u * 2] + seg[u * 2 + 1];
}
void upd(int u, int ul, int ur, int x){
    lazy[u] += x;
    seg[u] += (ur - ul) * x;
void shift(int u, int ul, int ur){
    int mid = (ul + ur) / 2;
    upd(u * 2, ul, mid, lazy[u]);
    upd(u * 2 + 1, mid, ur, lazy[u]);
    lazy[u] = 0;
```

```
void increase(int 1, int r, int x, int u = 1, int ul
= 0, int ur = n){
    if(1 \ge ur \mid \mid ul \ge r)return;
    if(1 <= u1 && ur <= r){
        upd(u, ul, ur, x);
        return;
    shift(u, ul, ur);
    int mid = (ul + ur) / 2;
    increase(l, r, x, u * 2, ul, mid);
    increase(1, r, x, u * 2 + 1, mid, ur);
    seg[u] = seg[u * 2] + seg[u * 2 + 1];
int sum(int l, int r, int u = 1, int ul = 0, int ur
= n){
    if(1 \ge ur \mid \mid ul \ge r)return 0;
    if(1 <= u1 && ur <= r)return seg[u];</pre>
    shift(u, ul, ur);
    int mid = (ul + ur) / 2;
    return sum(1, r, u * 2, ul, mid) + sum(1, r, u *
2 + 1, mid, ur);
void showSegments() {
    for(int i = 0; i < segSize; i++)</pre>
        cout << seg[i] << ' ';</pre>
    cout << endl;</pre>
void Main() {
    cin \gg n;
    segSize = 2;
    while(segSize / 2 <= n) segSize *= 2;</pre>
    for(int i = 0; i < n; i++)</pre>
        cin \gg a[i];
    build();
```

```
int main() {
    ios_base::sync_with_stdio(false);
    cin.tie(0); cout.tie(0);
    Main();
    return 0;
}
```

3. Math

3.1. choose

```
#define ll long long
const int N = 2e3 + 5;
const 11 M = 1e9 + 7;
11 fact[N], inv[N];
int r, n, q;
11 exp(ll b, ll p, ll m) {
    b \% = m;
    11 \text{ result} = 1;
    while(p) {
        if(p % 2)
            result = result * b % m;
        b = b * b % m;
        p /= 2;
    return result;
}
void preProcess() {
    fact[0] = 1;
    for(int i = 1; i < N; i++)</pre>
        fact[i] = fact[i - 1] * i % M;
    inv[N - 1] = exp(fact[N - 1], M - 2, M);
    for(int i = N - 1; i > 0; i--)
        inv[i - 1] = inv[i] * i % M;
}
```

```
11 choose(int n, int r) {
    if(r > n) return 0;
    return fact[n] * inv[r] % M * inv[n - r] % M;
void Main() {
    cin >> q;
    while(q--) {
        cin \gg n \gg r;
        cout << choose(n, r) << '\n';</pre>
}
int main() {
    ios::sync with stdio(false);
    cin.tie(0); cout.tie(0);
    preProcess();
    Main();
    return 0;
3.2. gcd
int gcd (int a, int b) {
    return b ? gcd (b, a % b) : a;
// fast version...
int gcd(int a, int b) {
    if (!a || !b)
        return a | b;
    unsigned shift = builtin ctz(a | b);
    a >>= builtin ctz(a);
    do {
        b >>= builtin ctz(b);
        if (a > b)
            swap(a, b);
```

```
b -= a:
   } while (b);
   return a << shift;
3.3. compressing
// compressing
sort(temp values, temp values + n);
int numOfUnique = unique(temp values, temp values +
n) - temp values;
for(int i = 0; i < n; i++)
   h[i] = lower bound(temp values, temp values +
numOfUnique, h[i]) - temp values;
3.4. lower bound and upper bound
int main() {
   vector<int> v = \{11, 34, 56, 67, 89\};
      // Finding lower bound of 56
    cout << *lower_bound(v.begin(), v.end(), 56)</pre>
      << endl;
     // Finding upper bound of 56
    cout << *upper_bound(v.begin(), v.end(), 56);</pre>
   return 0:
Output:
56
67
4. Graph
4.1. BFS
#define distance d
const int maxN = 1e5 + 10, oo = 1e9;
vector <int> adj[maxN];
```

int distance[maxN];

queue<int> q;

```
void BFS(int n, int r) {
    for (int i=1; i<=n; i++) distance[i] = oo;</pre>
    distance[r] = 0;
    q.push(r);
    while(q.size()) {
        int v = q.front();
        q.pop();
        for (auto u : adj[v])
            if(distance[u] > distance[v] + 1) {
                distance[u] = distance[v] + 1;
                q.push(u);
            }
    }
}
int main() {
    ios base::sync with stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {
        int u, v; cin >> u >> v;
        adj[u].push back(v);
        adj[v].push back(u);
    }
    BFS(n, 1);
    for (int i=1; i<=n; i++)</pre>
        cout << i << ':' << distance[i] << '\n';</pre>
}
4.2. bipartite
const int maxN = 1e5 + 10;
vector <int> adj[maxN];
bool mark[maxN];
```

```
int color[maxN];
bool bipartite = true;
```

```
void DFS(int v, int parent) {
    mark[v] = true;
    if(parent != -1) color[v] = 1 - color[parent];
    else color[v] = 1;
    for (auto u : adi[v]) {
        if(!mark[u])
            DFS(u, v);
        else if(color[u] == color[v])
            bipartite = false;
    }
int main() {
    ios base::sync with stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {
        int u, v; cin >> u >> v;
        adj[u].push back(v);
        adj[v].push back(u);
    for (int i=1; i<=n; i++) {
        if(mark[i]) continue;
        DFS(i, -1); //root does not have parent.
    if(bipartite) cout << "Graph Is Bipartite\n";</pre>
    else cout << "Graph Is Not Bipartite\n";</pre>
4.3. cycle finding
const int maxN = 1e5 + 10;
vector <int> adj[maxN];
```

bool mark[maxN];

```
bool cycle found = false;
void DFS(int v, int parent) {
    mark[v] = true;
    for (auto u : adj[v]) {
        if(!mark[u]) DFS(u, v); //u's parent is v.
        else if(u != parent) cycle_found = true;
    }
int main() {
    ios base::sync with stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {
        int u, v; cin >> u >> v;
        adj[u].push back(v);
        adj[v].push_back(u);
    }
    for (int i=1; i<=n; i++) {
        if(mark[i]) continue;
        DFS(i, -1); //root does not have parent.
    if(cycle found) cout << "Graph has Cycle\n";</pre>
    else cout << "Graph does not have Cycle\n";</pre>
4.4. DFS
const int maxN = 1e5 + 10;
vector <int> adj[maxN];
bool mark[maxN];
vector <int> component;
void DFS(int v) {
    mark[v] = true;
    component.push back(v);
```

CTRL+ALT+DEFEAT - Hakim Sabzevari University

```
for (auto u : adj[v])
        if(!mark[u]) DFS(u);
}
int main() {
    ios base::sync with stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {</pre>
        int u, v; cin >> u >> v;
        adj[u].push back(v);
        adj[v].push back(u);
    for (int i=1; i<=n; i++) {
        if(mark[i]) continue;
        component.clear();
        DFS(i);
        for (auto v : component)
            cout << v << ' ';
        cout << '\n';</pre>
    }
```

4.5. floyd-warshall

```
// Implementing floyd warshall algorithm
void floydWarshall(int graph[][nV]) {
   int matrix[nV][nV], i, j, k;
   for (i = 0; i < nV; i++)
      for (j = 0; j < nV; j++)
      matrix[i][j] = graph[i][j];
   // Adding vertices individually
   for (k = 0; k < nV; k++) {
      for (i = 0; i < nV; i++) {
        for (j = 0; j < nV; j++) {
            if (matrix[i][k] + matrix[k][j] <
            matrix[i][j])
            matrix[i][j] = matrix[i][k] +
   matrix[k][j];</pre>
```

```
}
    printMatrix(matrix);
4.6. prim
// Function to find sum of weights of edges of the
Minimum Spanning Tree.
int spanningTree(int V, int E, vector<vector<int>>
&edges) {
    // Create an adjacency list representation of
the graph
    vector<vector<int>> adj[V];
   // Fill the adjacency list with edges and their
weights
    for (int i = 0; i < E; i++) {
        int u = edges[i][0];
        int v = edges[i][1];
        int wt = edges[i][2];
        adj[u].push back({v, wt});
        adj[v].push back({u, wt});
   }
    // Create a priority queue to store edges with
their weights
    priority queue<pair<int,int>,
vector<pair<int,int>>, greater<pair<int,int>>> pq;
    // Create a visited array to keep track of
visited vertices
    vector<bool> visited(V, false);
    // Variable to store the result (sum of edge
weights)
    int res = 0;
    // Start with vertex 0
    pq.push({0, 0});
```

```
// Perform Prim's algorithm to find the Minimum
Spanning Tree
    while(!pq.empty()){
        auto p = pq.top();
        pq.pop();
        int wt = p.first; // Weight of the edge
        int u = p.second; // Vertex connected to
the edge
        if(visited[u] == true){
            continue; // Skip if the vertex is
already visited
        res += wt; // Add the edge weight to the
result
        visited[u] = true; // Mark the vertex as
visited
        // Explore the adjacent vertices
        for(auto v : adj[u]){
            // v[0] represents the vertex and v[1]
represents the edge weight
            if(visited[v[0]] == false){
                pq.push(\{v[1], v[0]\}); // Add the
adjacent edge to the priority queue
    return res; // Return the sum of edge weights
of the Minimum Spanning Tree
int main() {
    vector<vector<int>> graph = {{0, 1, 5},
                                  \{1, 2, 3\},\
                                  \{0, 2, 1\}\};
    cout << spanningTree(3, 3, graph) << endl;</pre>
```

```
return 0;
4.7. shortest cycle
//this code works for simple graphs.
const int maxN = 1010, oo = 1e9;
vector <int> adj[maxN];
int deleted, distances[maxN];
queue<int> q;
void BFS(int n, int r) {
    for (int i=1; i<=n; i++) distances[i] = oo;</pre>
    distances[r] = 0;
    q.push(r);
    while(q.size()) {
        int v = q.front();
        q.pop();
        for (auto u : adj[v]) {
            if(v == r && u == deleted) continue;
//ignore deleted edge.
            if(distances[u] > distances[v] + 1) {
                distances[u] = distances[v] + 1;
                q.push(u);
            }
    }
}
int main() {
    ios_base::sync_with_stdio(0); cin.tie(0);
    int n, m; cin >> n >> m;
    for (int i=0; i<m; i++) {
        int u, v; cin >> u >> v;
        adj[u].push_back(v); adj[v].push_back(u);
    int length = oo;
```

```
for (int i=1; i<=n; i++) {
        for (auto u : adj[i]) {
            deleted = u;
            BFS(n, i);
            length = min(length, distances[u] + 1);
        }
        if(length == oo) cout << "Graph Does Not Have
Cycle\n";
        else cout << "Minimum Cycle Length is : " <<
length << '\n';
}</pre>
```

4.8. topologycal sort

```
int n; // number of vertices
vector<vector<int>> adj; // adjacency list of graph
vector<bool> visited;
vector<int> ans;
void dfs(int v) {
    visited[v] = true;
    for (int u : adj[v]) {
        if (!visited[u])
            dfs(u);
    }
    ans.push_back(v);
void topological sort() {
    visited.assign(n, false);
    ans.clear();
    for (int i = 0; i < n; ++i) {
        if (!visited[i]) {
            dfs(i);
        }
    reverse(ans.begin(), ans.end());
```

```
4.9. lowest common Ancestor
struct LCA {
    vector<int> height, euler, first, segtree;
    vector<bool> visited;
   int n;
    LCA(vector<vector<int>> &adj, int root = 0) {
        n = adj.size();
        height.resize(n);
        first.resize(n);
        euler.reserve(n * 2);
        visited.assign(n, false);
        dfs(adj, root);
        int m = euler.size();
        segtree.resize(m * 4);
        build(1, 0, m - 1);
   }
    void dfs(vector<vector<int>> &adj, int node, int
h = 0) {
        visited[node] = true;
        height[node] = h;
        first[node] = euler.size();
        euler.push back(node);
        for (auto to : adj[node]) {
            if (!visited[to]) {
                dfs(adj, to, h + 1);
                euler.push back(node);
    void build(int node, int b, int e) {
        if (b == e) {
```

segtree[node] = euler[b];

```
} else {
             int mid = (b + e) / 2;
             build(node << 1, b, mid);</pre>
             build(node << 1 | 1, mid + 1, e);</pre>
             int l = segtree[node << 1], r =</pre>
segtree[node << 1 | 1];</pre>
             segtree[node] = (height[1] < height[r])</pre>
? 1 : r;
    }
    int query(int node, int b, int e, int L, int R)
{
         if (b > R \mid\mid e < L)
             return -1;
         if (b >= L \&\& e <= R)
             return segtree[node];
         int mid = (b + e) \gg 1;
         int left = query(node << 1, b, mid, L, R);</pre>
         int right = query(node << 1 | 1, mid + 1, e,</pre>
L, R);
         if (left == -1) return right;
         if (right == -1) return left;
         return height[left] < height[right] ? left :</pre>
right;
    }
    int lca(int u, int v) {
         int left = first[u], right = first[v];
        if (left > right)
             swap(left, right);
         return query(1, 0, euler.size() - 1, left,
right);
};
```

4.10. lowest common Ancestor (binary lifting)

```
int n, 1;
vector<vector<int>> adj;
int timer;
vector<int> tin, tout;
vector<vector<int>> up;
void dfs(int v, int p)
    tin[v] = ++timer;
    up[v][0] = p;
    for (int i = 1; i <= 1; ++i)
        up[v][i] = up[up[v][i-1]][i-1];
    for (int u : adj[v]) {
        if (u != p)
            dfs(u, v);
   }
    tout[v] = ++timer;
bool is ancestor(int u, int v)
    return tin[u] <= tin[v] && tout[u] >= tout[v];
int lca(int u, int v)
    if (is_ancestor(u, v))
        return u;
    if (is_ancestor(v, u))
        return v;
    for (int i = 1; i >= 0; --i) {
        if (!is ancestor(up[u][i], v))
```

```
u = up[u][i];
}
return up[u][0];
}

void preprocess(int root) {
   tin.resize(n);
   tout.resize(n);
   timer = 0;
   l = ceil(log2(n));
   up.assign(n, vector<int>(l + 1));
   dfs(root, root);
}
```

4.11. hungarian algorithm (assignment problem)

```
vector<int> u (n+1), v (m+1), p (m+1), way (m+1);
for (int i=1; i<=n; ++i) {</pre>
   p[0] = i;
   int j0 = 0;
   vector<int> minv (m+1, INF);
   vector<bool> used (m+1, false);
   do {
        used[j0] = true;
       int i0 = p[j0], delta = INF, j1;
       for (int j=1; j<=m; ++j)
            if (!used[j]) {
                int cur = A[i0][j]-u[i0]-v[j];
                if (cur < minv[j])</pre>
                    minv[j] = cur, way[j] = j0;
                if (minv[j] < delta)</pre>
                    delta = minv[j], j1 = j;
        for (int j=0; j<=m; ++j)</pre>
            if (used[j])
                u[p[j]] += delta, v[j] -= delta;
            else
```

```
minv[j] -= delta;
        j0 = j1;
    } while (p[j0] != 0);
    do {
        int j1 = way[j0];
        p[j0] = p[j1];
        j0 = j1;
    } while (j0);
}
4.12. 2SAT
struct TwoSatSolver {
    int n_vars;
    int n vertices;
    vector<vector<int>> adj, adj_t;
    vector<bool> used;
    vector<int> order, comp;
    vector<bool> assignment;
    TwoSatSolver(int _n_vars) : n_vars(_n_vars),
n vertices(2 * n vars), adj(n vertices),
adj t(n vertices), used(n vertices), order(),
comp(n vertices, -1), assignment(n vars) {
        order.reserve(n vertices);
    }
    void dfs1(int v) {
        used[v] = true;
        for (int u : adj[v]) {
            if (!used[u])
                dfs1(u);
        order.push back(v);
    }
    void dfs2(int v, int cl) {
        comp[v] = cl;
        for (int u : adj t[v]) {
```

```
if (comp[u] == -1)
                dfs2(u, c1);
        }
    }
    bool solve 2SAT() {
        order.clear();
        used.assign(n_vertices, false);
        for (int i = 0; i < n vertices; ++i) {</pre>
            if (!used[i])
                dfs1(i);
        }
        comp.assign(n vertices, -1);
        for (int i = 0, j = 0; i < n vertices; ++i)
            int v = order[n vertices - i - 1];
            if (comp[v] == -1)
                dfs2(v, j++);
        }
        assignment.assign(n vars, false);
        for (int i = 0; i < n vertices; i += 2) {
            if (comp[i] == comp[i + 1])
                return false;
            assignment[i / 2] = comp[i] > comp[i +
1];
        }
        return true;
    }
    void add disjunction(int a, bool na, int b, bool
nb) {
        // na and nb signify whether a and b are to
be negated
        a = 2 * a ^ na;
        b = 2 * b ^ nb;
        int neg a = a ^ 1;
```

```
int neg_b = b ^ 1;
       adj[neg_a].push_back(b);
       adj[neg b].push back(a);
       adj_t[b].push_back(neg_a);
       adj_t[a].push_back(neg_b);
   }
   static void example_usage() {
       TwoSatSolver solver(3); // a, b, c
       solver.add disjunction(0, false, 1,
true); // a v not b
        solver.add disjunction(0, true, 1,
true); // not a v not b
       solver.add disjunction(1, false, 2, false);
       b v c
//
       solver.add disjunction(0, false, 0, false);
//
       a v
       assert(solver.solve 2SAT() == true);
       auto expected = vector<bool>(True, False,
True);
       assert(solver.assignment == expected);
};
```

4.13. Heavy-light decomposition

```
vector<int> parent, depth, heavy, head, pos;
int cur_pos;

int dfs(int v, vector<vector<int>> const& adj) {
    int size = 1;
    int max_c_size = 0;
    for (int c : adj[v]) {
        if (c != parent[v]) {
            parent[c] = v, depth[c] = depth[v] + 1;
            int c_size = dfs(c, adj);
            size += c_size;
            if (c_size > max_c_size)
```

```
max c size = c size, heavy[v] = c;
        }
    }
    return size;
}
void decompose(int v, int h, vector<vector<int>>
const& adj) {
    head[v] = h, pos[v] = cur pos++;
    if (heavy[v] != -1)
        decompose(heavy[v], h, adj);
    for (int c : adj[v]) {
        if (c != parent[v] && c != heavy[v])
            decompose(c, c, adj);
    }
}
void init(vector<vector<int>> const& adj) {
    int n = adj.size();
    parent = vector<int>(n);
    depth = vector<int>(n);
    heavy = vector<int>(n, -1);
    head = vector<int>(n);
    pos = vector<int>(n);
    cur pos = 0;
    dfs(0, adj);
    decompose(0, 0, adj);
}
int query(int a, int b) {
    int res = 0;
    for (; head[a] != head[b]; b = parent[head[b]])
        if (depth[head[a]] > depth[head[b]])
            swap(a, b);
        int cur heavy path max =
segment tree query(pos[head[b]], pos[b]);
```

```
res = max(res, cur_heavy_path_max);
}
if (depth[a] > depth[b])
    swap(a, b);
int last_heavy_path_max =
segment_tree_query(pos[a], pos[b]);
    res = max(res, last_heavy_path_max);
    return res;
}
```

5. Data Structures

5.1. Array

```
int main() {
    array<int, 5> arr = {1, 2, 3, 4, 5};

    // Accessing elements
    cout << "Element at index 2: " << arr[2] <<
endl;

    // Size of the array
    cout << "Size of array: " << arr.size() << endl;

    // Fill array with a value
    arr.fill(10);
    cout << "Array after fill: ";
    for (int num : arr)
        cout << num << " ";
    cout << endl;

    return 0;
}</pre>
```

5.2. bitset

```
int main() {
   bitset<8> b1; // All bits initialized to 0
   bitset<8> b2("11001010");
```

```
// Set bit at index 3 to 1
    b1.set(3);
    cout << "b1 after set(3): " << b1 << endl;</pre>
    // Reset all bits of b2
    b2.reset();
    cout << "b2 after reset: " << b2 << endl;</pre>
    // Flip all bits of b1
    b1.flip();
    cout << "b1 after flip: " << b1 << endl;</pre>
    // Access the bit at index 2
    cout << "b1[2]: " << b1[2] << endl;</pre>
    // Test if bit at index 2 is set to 1
    if (b1.test(2))
    {
        cout << "Bit 2 is set to 1." << endl;</pre>
    }
    // Count the number of 1's in b1
    cout << "Number of 1's in b1: " << b1.count() <<</pre>
endl;
    return 0;
5.3. dequeue
int main() {
    deque<int> deq = \{1, 2, 3, 4, 5\};
    // Adding elements
    deq.push front(0); // Add at front
```

deq.push_back(6); // Add at back

CTRL+ALT+DEFEAT - Hakim Sabzevari University

```
// Removing elements
    deq.pop front(); // Remove from front
    deq.pop back(); // Remove from back
    // Accessing elements
    cout << "First element: " << deq.front() <<</pre>
endl:
    cout << "Last element: " << deq.back() << endl;</pre>
    // Iterating through deque
    cout << "Deque elements: ";</pre>
    for (int x : deq)
        cout << x << " ";
    cout << endl;</pre>
    // Size of deque
    cout << "Deque size: " << deq.size() << endl;</pre>
    return 0;
}
5.4. link list
int main() {
    list<int> lst = \{1, 2, 3, 4, 5\};
    // Adding elements
    lst.push back(6); // Add to the end
    lst.push front(0); // Add to the front
    // Removing elements
    lst.pop back(); // Remove from the end
    lst.pop_front(); // Remove from the front
    // Iterating through list
    cout << "List elements: ";</pre>
    for (int x : lst)
        cout << x << " ";
```

```
cout << endl;</pre>
    // Size of list
    cout << "List size: " << lst.size() << endl;</pre>
    return 0;
5.5. Map
int main() {
    map<string, int> m;
    // Insert key-value pairs
    m["apple"] = 3;
    m["banana"] = 2;
    m["orange"] = 5;
    // Access value by key
    cout << "Value for apple: " << m["apple"] <<</pre>
endl;
    // Iterate through map
    cout << "Map elements: ";</pre>
    for (auto &pair : m)
         cout << pair.first << ": " << pair.second <<</pre>
    cout << endl;</pre>
    // Remove an element
    m.erase("banana");
    cout << "Map after erase: ";</pre>
    for (auto &pair : m)
         cout << pair.first << ": " << pair.second <<</pre>
\mathbf{u} = \mathbf{u}
```

```
cout << endl;</pre>
    // Check if key exists
    if (m.find("orange") != m.end())
        cout << "Orange is in the map." << endl;</pre>
    return 0;
5.6. priority queue
int main() {
    priority_queue<int> pq;
    // Insert elements into the priority queue
    pq.push(10);
    pq.push(30);
    pq.push(20);
    cout << "Top element (max priority): " <<</pre>
pq.top() << endl;
    // Pop the top element
    pq.pop();
    cout << "Top element after pop: " << pq.top() <<</pre>
endl;
    // Size of the priority queue
    cout << "Size of priority queue: " << pq.size()</pre>
<< endl;
    return 0;
```

```
5.7. queue
int main() {
    queue<int> q;
    // Push elements onto the queue
    q.push(10);
    q.push(20);
    q.push(30);
    cout << "Front element: " << q.front() << endl;</pre>
    // Pop an element from the queue
    q.pop();
    cout << "Front element after pop: " << q.front()</pre>
<< endl:
    // Size of the queue
    cout << "Size of queue: " << q.size() << endl;</pre>
    // Check if the queue is empty
    if (q.empty()) {
        cout << "Queue is empty." << endl;</pre>
    }
    else {
        cout << "Queue is not empty." << endl;</pre>
    }
    return 0;
5.8. set
int main() {
    set<int> s;
    // Insert elements
    s.insert(10);
    s.insert(20);
    s.insert(15);
```

```
// Print elements in sorted order
    cout << "Set elements: ";</pre>
    for (int num : s)
        cout << num << " ";
    cout << endl;</pre>
    // Check if an element exists
    if (s.find(15) != s.end())
    {
        cout << "15 is in the set." << endl;</pre>
    // Remove an element
    s.erase(10);
    cout << "Set after erase: ";</pre>
    for (int num : s)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Size of the set
    cout << "Size of set: " << s.size() << endl;</pre>
    return 0;
5.9. stack
int main() {
    stack<int> s;
    // Push elements onto the stack
    s.push(10);
    s.push(20);
    s.push(30);
    cout << "Top element after push: " << s.top() <<</pre>
endl;
```

```
// Pop an element from the stack
    s.pop();
    cout << "Top element after pop: " << s.top() <<</pre>
endl:
    // Check if stack is empty
    if (s.empty()) {
        cout << "Stack is empty." << endl;</pre>
    }
    else {
        cout << "Stack is not empty." << endl;</pre>
    // Size of the stack
    cout << "Size of stack: " << s.size() << endl;</pre>
    return 0;
5.10. unordered map
int main() {
    unordered map<int, string> um;
    // Inserting elements
    um[1] = "apple";
    um[2] = "banana";
    um[3] = "cherry";
    // Accessing elements
    cout << "Key 2 maps to: " << um[2] << endl;</pre>
    // Checking if key exists
    if (um.find(4) == um.end())
        cout << "Key 4 not found!" << endl;</pre>
    // Iterating through unordered map
    cout << "Unordered map elements: ";</pre>
```

```
for (auto &pair : um)
         cout << pair.first << " -> " << pair.second</pre>
<< " | ";
    }
    cout << endl;</pre>
    return 0:
}
5.11. unordered set
int main() {
    unordered set<int> us;
    // Insert elements
    us.insert(10);
    us.insert(20);
    us.insert(15);
    // Print elements
    cout << "Unordered set elements: ";</pre>
    for (int num : us)
        cout << num << " ":
    cout << endl;</pre>
    // Check if an element exists
    if (us.find(15) != us.end())
         cout << "15 is in the unordered set." <<</pre>
endl:
    // Remove an element
    us.erase(10);
    cout << "Unordered set after erase: ";</pre>
    for (int num : us)
         cout << num << " ";
```

```
cout << endl;</pre>
    return 0;
5.12. vector
int main() {
    // Create a vector of integers
    vector<int> v;
    // Add elements using push back
    v.push back(10);
    v.push back(20);
    v.push back(30);
    cout << "Vector after push back: ";</pre>
    for (int num : v)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Accessing elements using at() and indexing
    cout << "Element at index 1: " << v.at(1) <<</pre>
endl:
    cout << "Element at index 0: " << v[0] << endl;</pre>
    // Pop an element from the back
    v.pop back();
    cout << "Vector after pop back: ";</pre>
    for (int num : v)
        cout << num << " ";</pre>
    cout << endl:</pre>
    // Insert an element at a specific position
    v.insert(v.begin() + 1, 25); // Insert 25 at
index 1
    cout << "Vector after insert: ";</pre>
    for (int num : v)
        cout << num << " ";
```

```
cout << endl;</pre>
    // Remove an element from the vector
    v.erase(v.begin() + 1); // Remove the element at
index 1
    cout << "Vector after erase: ";</pre>
    for (int num : v)
        cout << num << " ";</pre>
    cout << endl;</pre>
    // Resize the vector
    v.resize(5, 50); // Resize to size 5, fill new
elements with 50
    cout << "Vector after resize: ";</pre>
    for (int num : v)
        cout << num << " ";
    cout << endl;</pre>
    // Get the size of the vector
    cout << "Size of vector: " << v.size() << endl;</pre>
    // Clear the vector
    v.clear();
    cout << "Vector after clear: " << v.size() << "</pre>
(size is now zero)" << endl;</pre>
    return 0;
```

5.13. DS cheatsheet

1. Vector

- Description: Dynamic array that allows fast random access.
- Methods:
 - push back(x): Add element x to the end.
- pop back(): Remove the last element.
- at(i): Access the element at index i (bounds checked).
- operator[]: Access the element at index i (no bounds check).
- size(): Return the number of elements.

CTRL+ALT+DEFEAT – Hakim Sabzevari University

- empty(): Check if the vector is empty.
- resize(n, x): Resize the vector to size n and fill new elements with x.
- clear(): Remove all elements.

2. Stack

- Description: Last-In-First-Out (LIFO) structure, used for backtracking problems.
- Methods:
- push(x): Add element x to the top.
- pop(): Remove the top element.
- top(): Get the top element.
- size(): Return the number of elements.
- empty(): Check if the stack is empty.

3. Oueue

- Description: First-In-First-Out (FIFO) structure, ideal for problems involving processing in order.
- Methods:
- push(x): Add element x to the back.
- pop(): Remove the front element.
- front(): Get the front element.
- back(): Get the back element.
- size(): Return the number of elements.
- empty(): Check if the queue is empty.

4. Priority Queue (Max-Heap by default)

- Description: A heap-based structure that always gives the maximum element.
- Methods:
 - push(x): Add element x to the gueue.
- pop(): Remove the largest element.
- top(): Get the largest element.
- size(): Return the number of elements.
- empty(): Check if the queue is empty.

5. Set

- Description: Collection of unique elements in sorted order.
- Methods:
- insert(x): Add element x.
- erase(x): Remove element x.
- find(x): Check if element x exists.
- size(): Return the number of elements.
- empty(): Check if the set is empty.
- clear(): Remove all elements.

6. Map

- Description: Stores key-value pairs in sorted order based on keys.
- Methods:
- insert({key, value}): Add key-value pair.
- erase(key): Remove element by key.
- find(key): Check if a key exists.
- operator[]: Access the value associated with a key.
- size(): Return the number of elements.
- empty(): Check if the map is empty.
- clear(): Remove all elements.

7. Unordered Set

- Description: Collection of unique elements with no specific order.
- Methods:
- insert(x): Add element x.
- erase(x): Remove element x.
- find(x): Check if element x exists.
- size(): Return the number of elements.
- empty(): Check if the unordered set is empty.

8. Unordered Map

- Description: Stores key-value pairs with no specific order.
- Methods:
- insert({key, value}): Add key-value pair.
- erase(key): Remove element by key.
- find(key): Check if a key exists.
- operator[]: Access the value associated with a key.
- size(): Return the number of elements.
- empty(): Check if the unordered map is empty.

9. Bitset

- Description: A space-efficient container for a fixed-size sequence of bits (0 or 1).
- Methods:
- set(i): Set bit at index i to 1.
- reset(i): Set bit at index i to 0.
- flip(i): Toggle the bit at index i.
- test(i): Check if the bit at index i is 1.
- count(): Count the number of bits set to 1.
- size(): Return the number of bits.
- operator[]: Access the bit at index i.
- to string(): Convert bitset to string.

10. Array

- Description: Fixed-size array used for fast access, but size cannot be changed after initialization.
- Methods:
- fill(x): Fill all elements with the value x.
- size(): Return the number of elements.
- operator[]: Access element at index i.
- at(i): Access element at index i with bounds checking.
- front(): Get the first element.
- back(): Get the last element.

11. Deque

- Description: Double-ended queue that allows fast insertion and removal at both ends.
- Methods:
- push front(x): Add element x to the front.
- push back(x): Add element x to the back.
- pop front(): Remove the front element.
- pop back(): Remove the back element.
- front(): Get the front element.
- back(): Get the back element.
- size(): Return the number of elements.
- empty(): Check if the deque is empty.

12. Linked List (Using STL List)

- Description: Doubly linked list that allows fast insertion and deletion at both ends.
- Methods:
- push back(x): Add element x to the back.
- push front(x): Add element x to the front.
- pop back(): Remove the last element.
- pop front(): Remove the first element.
- size(): Return the number of elements.
- empty(): Check if the list is empty.
- front(): Get the first element.
- back(): Get the last element.
- clear(): Remove all elements.

6. Dynamic Programming

6.1. counting paths matrix

```
\ensuremath{//} Function to count the number of unique paths in a \ensuremath{\mathsf{matrix}}
```

```
int countPaths(int n, int m) {
```

```
vector<vector<int>> dp(n, vector<int>(m, 0));
    // Starting point: only one way to be at the
start
    dp[0][0] = 1;
    // Fill the DP table for first row and first
column
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < m; j++) {
            if (i > 0)
                dp[i][j] += dp[i - 1][j]; // From
top
            if(j>0)
                dp[i][j] += dp[i][j - 1]; // From
left
    }
    return dp[n - 1][m - 1]; // Return the number of
paths to bottom-right corner
int main() {
    int n = 3, m = 3;
                                      // Example
matrix dimensions
    cout << countPaths(n, m) << endl; // Output the</pre>
result
    return 0;
}
6.2. edit distance
// Function to compute the Edit Distance
int editDistance(string str1, string str2, int m,
int n) {
```

```
vector<vector<int>> dp(m + 1, vector<int>(n +
1));
```

```
for (int i = 0; i <= m; i++) {
        for (int j = 0; j <= n; j++) {
            if (i == 0)
                dp[i][j] = j;
            else if (i == 0)
                dp[i][j] = i;
            else if (str1[i - 1] == str2[j - 1])
                dp[i][j] = dp[i - 1][j - 1];
            else
                dp[i][j] = 1 + min({dp[i - 1][j -
1], dp[i][j - 1], dp[i - 1][j]});
    return dp[m][n];
int main() {
    string str1 = "sitting", str2 = "kitten";
    int m = str1.length(), n = str2.length();
    cout << "Edit Distance: " << editDistance(str1,</pre>
str2, m, n) << endl;
    return 0;
6.3. Egg Dropping
// Function to find the minimum number of attempts
needed
int eggDrop(int eggs, int floors)
    vector<vector<int>> dp(eggs + 1,
vector<int>(floors + 1, 0));
```

for (int i = 1; $i \leftarrow eggs$; i++)

for (int j = 0; j <= floors; j++)</pre>

dp[i][0] = 0;

```
dp[1][j] = j;
    for (int i = 2; i \leftarrow eggs; i++)
        for (int j = 2; j <= floors; j++)</pre>
            dp[i][j] = INT MAX;
            for (int x = 1; x <= j; x++)
                 dp[i][j] = min(dp[i][j], 1 +
\max(dp[i - 1][x - 1], dp[i][j - x]));
        }
    return dp[eggs][floors];
int main()
    int eggs = 2, floors = 10;
    cout << "Minimum attempts: " << eggDrop(eggs,</pre>
floors) << endl;
    return 0;
6.4. fibonacci
// Fibonacci sequence using dynamic programming
(Memoization)
// Function to compute Fibonacci number
int fib(int n, vector<int> &dp)
    // Base cases
    if (n <= 1)
        return n;
```

// If the value is already computed, return it

6.5. knapsack 01

```
return dp[n][W];
int main()
                                            11
    int W = 100;
Capacity of the knapsack
    vector<int> val = {80, 24, 23, 22, 21}; //
Values of the items
    vector<int> wt = {80, 25, 25, 25, 25}; //
Weights of the items
    int n = val.size();
                                            //
Number of items
    cout << "Maximum value in knapsack: " <<</pre>
knapsack(W, wt, val, n) << endl;</pre>
    return 0;
6.6. LCS
// Function to compute the length of the Longest
Common Subsequence
int lcs(string X, string Y, int m, int n)
    vector<vector<int>> dp(m + 1, vector<int>(n + 1,
0)); // DP table
    for (int i = 1; i <= m; i++)
        for (int j = 1; j <= n; j++)
            if (X[i-1] == Y[j-1])
                dp[i][j] = 1 + dp[i - 1][j - 1];
```

dp[i][j] = max(dp[i - 1][j], dp[i][j]

- 1]);

```
return dp[m][n];
int main()
    string X = "AGGTAB", Y = "GXTXAYB";
    int m = X.length(), n = Y.length();
    cout << "Length of Longest Common Subsequence: "</pre>
<< lcs(X, Y, m, n) << endl;
    return 0:
6.7. LIS
// Function to find the length of the Longest
Increasing Subsequence
int lis(vector<int> &arr, int n)
    vector<int> dp(n, 1); // DP array, initialized
to 1
    for (int i = 1; i < n; i++)
        for (int j = 0; j < i; j++)
            if (arr[i] > arr[j])
                dp[i] = max(dp[i], dp[j] + 1);
    return *max_element(dp.begin(), dp.end());
int main()
```

CTRL+ALT+DEFEAT – Hakim Sabzevari University

```
vector\langle int \rangle arr = {10, 22, 9, 33, 21, 50, 41,
60};
    int n = arr.size();
    cout << "Length of Longest Increasing</pre>
Subsequence: " << lis(arr, n) << endl;
    return 0;
6.8. LPS
// Function to compute the length of the Longest
Palindromic Subsequence
int lps(string s)
    int n = s.length();
    vector<vector<int>> dp(n, vector<int>(n, 0));
    for (int i = 0; i < n; i++)
        dp[i][i] = 1; // Single character is a
palindrome
    for (int len = 2; len <= n; len++)</pre>
        for (int i = 0; i < n - len + 1; i++)
            int j = i + len - 1;
            if (s[i] == s[i])
                 dp[i][j] = 2 + dp[i + 1][j - 1];
            else
                 dp[i][j] = max(dp[i + 1][j], dp[i][j]
- 1]);
    return dp[0][n - 1];
int main()
```

```
string s = "bbabcbcab";
    cout << "Length of Longest Palindromic</pre>
Subsequence: " << lps(s) << endl;</pre>
    return 0;
6.9. MCM
// Function to compute the minimum number of scalar
multiplications
int matrixChainMultiplication(vector<int> &dims, int
n) {
    vector<vector<int>> dp(n, vector<int>(n, 0)); //
DP table
    for (int len = 2; len \langle n; len++)
    {
        for (int i = 1; i < n - len + 1; i++)
            int j = i + len - 1;
            dp[i][j] = INT MAX;
            for (int k = i; k < j; k++)
                int q = dp[i][k] + dp[k + 1][j] +
dims[i - 1] * dims[k] * dims[j];
                dp[i][j] = min(dp[i][j], q);
    return dp[1][n - 1];
}
int main()
    vector<int> dims = {10, 20, 30, 40, 30};
    int n = dims.size();
```

```
cout << "Minimum number of scalar</pre>
multiplications: " <<</pre>
matrixChainMultiplication(dims, n) << endl;</pre>
    return 0;
6.10. Minimum Coin Change
// Function to find the minimum number of coins
required to make a total
int minCoins(const vector<int> &coins, int total)
    int n = coins.size();
    vector<int> dp(total + 1, INT MAX); //
Initialize DP array with infinity
    dp[0] = 0;
                                         // Base
case: 0 coins needed to make total 0
    // Fill DP array
    for (int i = 0; i < n; i++)
        for (int j = coins[i]; j <= total; j++)</pre>
            if (dp[j - coins[i]] != INT MAX)
                dp[j] = min(dp[j], dp[j - coins[i]]
+ 1); // Minimize coin count
    return dp[total] == INT MAX ? -1 : dp[total]; //
Return -1 if not possible
int main()
```

```
vector<int> coins = {1, 2, 5};
                                             //
Example denominations
    int total = 11;
                                             //
Target total
    cout << minCoins(coins, total) << endl; //</pre>
Output the result
    return 0;
6.11. optimal BST
// Function to calculate the minimum search cost for
an optimal BST
int optimalBST(const vector<int> &freq)
    int n = freq.size();
    vector<vector<int>> dp(n, vector<int>(n, 0));
    // Fill the DP table for subarrays of increasing
length
    for (int len = 1; len <= n; len++)</pre>
    { // len is the range length
        for (int i = 0; i \le n - len; i++)
            int j = i + len - 1;
            dp[i][j] = INT_MAX;
            int sum = 0;
            // Calculate sum of frequencies from i
to j
            for (int k = i; k <= j; k++)
                sum += freq[k];
            // Try each k as the root and calculate
the minimum cost
            for (int k = i; k <= j; k++)
```

```
int cost = (k == i ? 0 : dp[i][k -
1]) + (k == j ? 0 : dp[k + 1][j]);
                dp[i][j] = min(dp[i][j], cost +
sum);
           }
        }
   return dp[0][n - 1]; // The minimum cost for the
entire range
int main()
   vector<int> freq = {34, 8, 50, 13}; // Example
frequencies of keys
    cout << optimalBST(freq) << endl; // Output</pre>
the minimum cost
   return 0;
6.12. partition problem
// Function to determine if a given set can be
partitioned into two subsets
bool canPartition(vector<int> &nums)
   int sum = 0;
   for (int num : nums)
        sum += num;
   if (sum % 2 != 0)
        return false;
   int target = sum / 2;
   vector<bool> dp(target + 1, false);
   dp[0] = true;
```

for (int num : nums)

```
for (int j = target; j >= num; j--)
            dp[j] = dp[j] \mid | dp[j - num];
    }
    return dp[target];
int main()
    vector\langle int \rangle nums = \{1, 5, 11, 5\};
    cout << "Can partition: " << (canPartition(nums)</pre>
? "Yes" : "No") << endl;
    return 0;
6.13. Regular Expression Matching
bool isMatch(const string &s, const string &p) {
    int m = s.size(), n = p.size();
    // DP table dp[i][j] will be true if s[0...i-1]
matches p[0...j-1]
    vector<vector<bool>> dp(m + 1, vector<bool>(n +
1, false));
    // Base case: empty string matches empty pattern
    dp[0][0] = true;
    // Handle patterns like "a*" or ".*" where "*"
can match 0 occurrence
    for (int j = 1; j <= n; j++)
        if (p[j - 1] == '*')
            dp[0][j] = dp[0][j - 2];
```

```
}
   // Fill the dp table
    for (int i = 1; i <= m; i++)
   {
        for (int j = 1; j <= n; j++)
            if (p[j-1] == s[i-1] || p[j-1] ==
'.')
            {
                dp[i][j] = dp[i - 1][j - 1]; //
Character matches
           else if (p[j - 1] == '*')
                // '*' matches zero occurrence or
one/more occurrences of the preceding character
                dp[i][j] = dp[i][j - 2] | | (dp[i -
1|[j] \&\& (s[i-1] == p[j-2] || p[j-2] == '.'));
           }
    }
    return dp[m][n]; // Final answer whether the
whole string matches the pattern
int main()
    string s = "aab";
    string p = "c*a*b";
   if (isMatch(s, p))
        cout << "The string matches the pattern." <<</pre>
endl;
    else
```

```
cout << "The string does not match the</pre>
pattern." << endl;</pre>
   }
    return 0;
6.14. ROD cutting
// Function to compute the maximum profit from
cutting a rod
int rodCutting(vector<int> &prices, int n)
    vector<int> dp(n + 1, 0); // DP array
   for (int i = 1; i <= n; i++)
        for (int j = 1; j <= i; j++)
            dp[i] = max(dp[i], prices[j - 1] + dp[i]
- j]);
        }
    return dp[n];
int main()
    vector<int> prices = {1, 5, 8, 9, 10, 17, 17,
20);
    int n = prices.size();
    cout << "Maximum profit from rod cutting: " <<</pre>
rodCutting(prices, n) << endl;</pre>
    return 0;
```

6.15. subset sum

```
// Function to determine if there's a subset with
sum equal to the target
bool subsetSum(vector<int> &nums, int sum)
    int n = nums.size();
    vector<vector<bool>> dp(n + 1, vector<bool>(sum
+ 1, false));
    for (int i = 0; i <= n; i++)
        dp[i][0] = true;
    for (int i = 1; i <= n; i++)
        for (int j = 1; j \leftarrow sum; j++)
            if (nums[i - 1] <= j)</pre>
                dp[i][j] = dp[i - 1][j] || dp[i -
1][j - nums[i - 1]];
            else
                dp[i][j] = dp[i - 1][j];
    return dp[n][sum];
int main()
    vector<int> nums = {3, 34, 4, 12, 5, 2};
    int sum = 9;
    cout << "Subset sum possible: " <<</pre>
(subsetSum(nums, sum) ? "Yes" : "No") << endl;
    return 0;
```

6.16. two player game

```
// Function to find the maximum sum player A can get
int maxCoins(const vector<int> &coins)
    int n = coins.size();
    vector<vector<int>> dp(n, vector<int>(n, 0));
    // Base case: when there's only one coin, player
A takes it
    for (int i = 0; i < n; i++)
        dp[i][i] = coins[i];
    // Fill DP table for subarrays of length 2 to n
    for (int len = 2; len <= n; len++)</pre>
        for (int i = 0; i < n - len + 1; i++)
            int j = i + len - 1;
            dp[i][j] = max(coins[i] + min(dp[i +
2|[j], dp[i + 1][j - 1]),
                           coins[j] + min(dp[i +
1][j - 1], dp[i][j - 2]));
    }
    return dp[0][n - 1]; // Maximum sum player A can
get
int main()
    vector<int> coins = {8, 15, 3, 7}; // Example
coins array
    cout << maxCoins(coins) << endl; // Output the</pre>
result
    return 0;
```

```
}
```

6.17. word break

```
// Function to check if a word can be segmented into
words from a dictionary
bool wordBreak(string s, unordered set<string>
&wordDict)
    int n = s.length();
    vector<bool> dp(n + 1, false);
    dp[0] = true;
    for (int i = 1; i <= n; i++)
        for (int j = 0; j < i; j++)
            if (dp[j] && wordDict.find(s.substr(j, i
- j)) != wordDict.end())
                dp[i] = true;
                break;
    return dp[n];
int main()
    string s = "leetcode";
    unordered set<string> wordDict = {"leet",
"code"};
    cout << "Can break: " << (wordBreak(s, wordDict)</pre>
? "Yes" : "No") << endl;
    return 0;
```

6.18. dp cheatsheet

1. Fibonacci Sequence (Memoization)

- Explanation: This algorithm calculates the nth Fibonacci number using memoization to store previously computed results, avoiding redundant calculations.
- When to Use: When you need to compute Fibonacci numbers for large values of `n` efficiently (i.e., for recursive problems that involve overlapping subproblems).

2. 0/1 Knapsack Problem

- Explanation: Given a set of items with weights and values, the goal is to determine the maximum value that can be obtained by putting items in a knapsack without exceeding the weight capacity.
- When to Use: In optimization problems where you need to maximize profit or value while respecting constraints (e.g., weight, space).

3. Minimum Coin Change

- Explanation: This algorithm calculates the minimum number of coins needed to make a given total using a set of coin denominations. It uses dynamic programming to store the results for all possible totals.
- When to Use: When you need to find the fewest coins needed to form a specific amount (e.g., for making change, budget optimization).

4. Longest Common Subsequence (LCS)

- Explanation: This algorithm finds the longest subsequence common to two sequences (strings, arrays, etc.). It uses a dynamic programming table to store intermediate results.
- When to Use: When comparing two sequences (e.g., DNA sequences, text comparison, diff tools) and need the longest subsequence they share.

5. Edit Distance (Levenshtein Distance)

CTRL+ALT+DEFEAT – Hakim Sabzevari University

- Explanation: This algorithm calculates the minimum number of operations (insertions, deletions, substitutions) needed to convert one string into another.
- When to Use: When you need to compare two strings and find the minimum edit operations required (e.g., spell checkers, natural language processing).

6. Longest Increasing Subsequence (LIS)

- Explanation: This algorithm finds the length of the longest increasing subsequence in a sequence of numbers. The subsequence need not be contiguous.
- When to Use: When you need to find the longest increasing subsequence in a sequence of numbers (e.g., stock price prediction, finding trends).

7. Matrix Chain Multiplication

- Explanation: This algorithm calculates the most efficient way to multiply a chain of matrices by minimizing the number of scalar multiplications.
- When to Use: When multiplying multiple matrices and you need to minimize the cost of multiplication (e.g., in computer graphics, optimization problems).

8. Rod Cutting Problem

- Explanation: This algorithm finds the maximum profit you can obtain by cutting a rod of length `n` into smaller pieces and selling them, based on the prices for each length.
- When to Use: When you need to solve problems related to cutting materials into pieces to maximize profit (e.g., resource allocation, profit optimization).

9. Subset Sum Problem

- Explanation: This algorithm checks if there is a subset of a given set of numbers that adds up to a target sum. It uses dynamic programming to keep track of achievable sums.
- When to Use: When you need to check whether a subset exists with a given sum (e.g., partitioning problems, subset analysis).

10. Egg Dropping Problem

- Explanation: This algorithm determines the minimum number of attempts required to find the highest floor from which an egg can be dropped without breaking, given 'k' eggs and 'n' floors.
- When to Use: In optimization problems where you need to minimize the number of trials in a worst-case scenario (e.g., testing, fault tolerance, hardware).

11. Partition Problem

- Explanation: This algorithm checks whether a given set can be partitioned into two subsets such that their sums are equal. It uses dynamic programming to check for subset sums.
- When to Use: When dividing a set of numbers into two equal subsets (e.g., load balancing, resource allocation).

12. Longest Palindromic Subsequence (LPS)

- Explanation: This algorithm finds the longest subsequence within a string that is a palindrome. It uses dynamic programming to build a table based on matching characters.
- When to Use: When you need to find the longest palindromic subsequence in a string (e.g., text processing, bioinformatics).

13. Word Break Problem

- Explanation: This algorithm checks if a string can be segmented into a space-separated sequence of words from a dictionary. It uses dynamic programming to store results for substrings.

- When to Use: When you need to determine whether a string can be split into valid words (e.g., for tokenizing sentences, text segmentation).

14. Regular Expression Matching

- Explanation: This algorithm checks if a string matches a pattern with . (any character) and * (zero or more of the previous character). It uses dynamic programming to track matching results for substrings.
- When to Use: When matching strings to patterns with wildcards like . and * (e.g., text matching, search engines, file pattern matching).

15. Optimal Binary Search Tree

- Explanation: This algorithm finds the minimum cost to construct a binary search tree based on the frequencies of elements. It uses dynamic programming to calculate the optimal cost for different subarrays.
- When to Use: When you need to minimize the cost of searching elements with different access frequencies (e.g., database indexing, search optimization).

16. 2 Player Game

- Explanation: This algorithm calculates the maximum sum player A can collect in a turn-based game where players alternate picking coins from either end of the array. It uses dynamic programming to determine the optimal strategy for player A.
- When to Use: When you need to calculate the optimal strategy in a turn-based game with alternating choices (e.g., maximizing outcomes in competitive games).

17. Counting Paths in Matrix

- Explanation: This algorithm calculates the number of unique paths from the top-left corner to the bottom-right corner of an n x m matrix. It uses dynamic programming to count paths by combining the results from adjacent cells.

CTRL+ALT+DEFEAT – Hakim Sabzevari University

- When to Use: When you need to find the number of distinct paths in a grid, where you can only move right or down (e.g., grid-based traversal problems).

7. More

7.1. longest polindrom substring (Mancher)

```
// O(n)
class Solution {
public:
    string longestPalindrome(std::string s) {
        if (s.length() <= 1) return s;</pre>
        // Preprocess the string with '#' characters
to handle even-length palindromes
        string modified s = "#";
        for (char c : s) {
            modified s += c;
            modified s += '#';
        }
        int n = modified s.size();
        vector<int> dp(n, 0); // dp array to store
the radius of the palindrome centered at each
character
        int center = 0, right = 0; // Initialize
center and right boundary
        int max len = 1; // Maximum length of
palindrome found
        string max str = s.substr(0, 1); //
Initialize the max palindrome substring with the
first character
        for (int i = 0; i < n; i++) {
```

```
// If i is within the current right
boundary, use previously calculated values to
minimize comparisons
            if (i < right) {</pre>
                dp[i] = min(right - i, dp[2 * center]
- i]);
            }
            // Expand around center i
            while (i - dp[i] - 1 >= 0 \&\& i + dp[i] +
1 < n \&\& modified s[i - dp[i] - 1] == modified s[i +
dp[i] + 1]) {
                dp[i]++;
            }
            // Update center and right boundary if
we've expanded beyond the current right
            if (i + dp[i] > right) {
                center = i;
                right = i + dp[i];
            }
            // Update max len and max str if a
longer palindrome is found
            if (dp[i] > max len) {
                max len = dp[i];
                max str = modified s.substr(i -
dp[i], 2 * dp[i] + 1);
                max str.erase(remove(max str.begin()
, max str.end(), '#'), max str.end());
        }
        return max str;
    }
};
```

7.2. median of 2 soted array

```
class Solution {
public:
    double findMedianSortedArrays(vector<int>
&nums1, vector<int> &nums2) {
        int n = nums1.size();
        int m = nums2.size();
        if(n > m) return
findMedianSortedArrays(nums2, nums1);
        int size = n + m;
        int left = (size + 1) / 2;
        int low = 0, high = n;
        int 11, 12, r1, r2, mid1, mid2;
        while(low <= high) {</pre>
            mid1 = (low + high) >> 1;
            mid2 = left - mid1;
            r1 = mid1 < n ? nums1[mid1] : 1e9;
            r2 = mid2 < m ? nums2[mid2] : 1e9;
            11 = mid1 - 1 >= 0 ? nums1[mid1 - 1] : -
1e9;
            12 = mid2 - 1 >= 0 ? nums2[mid2 - 1] : -
1e9;
            if(l1 <= r2 && l2 <= r1)
                if(size % 2) return max(11, 12);
                else return
(static cast<double>(max(l1, l2) + min(r1, r2))) /
2.0:
            else if(11 > r2) high = mid1 - 1;
            else low = mid1 + 1;
        return 0;
};
```

10. other

10.1. useful geo

Area of triangle with sides a, b, c: sqrt(S *(S-a)*(S-b)*(S-c)) where S = (a+b+c)/2

Area of equilateral triangle: s^2 * sqrt(3) / 4 where is side lenght

Pyramid and cones volume: 1/3 area(base) * height

if p1=(x1, x2), p2=(x2, y2), p3=(x3, y3) are points on circle, the center is

 $x = -((x2^2 - x1^2 + y2^2 - y1^2)^*(y3 - y2) - (x2^2 - x3^2 + y2^2 - y3^2)^*(y1 - y2)) / (2^*(x1 - x2)^*(y3 - y2) - 2^*(x3 - x2)^*(y1 - y2))$

 $y = -((y2^2 - y1^2 + x2^2 - x1^2)*(x3 - x2) - (y2^2 - y3^3 + x2^2 - x3^2)*(x1 - x2)) / (2*(y1 - y2)*(x3 - x2) - 2*(y3 - y2)*(x1 - x2))$

10.2. number of primes

30: 10 60: 17 100: 25 1000: 168 10000: 1229 100000: 9592 1000000: 78498

10.3. Factorials

10000000: 664579

1: 1 2: 2

3: 6

4: 24

.. - .

5: 120

6: 720

7: 5040

8: 40320

9: 362880

10: 3628800

11: 39916800

12: 479001600

13:6227020800

14:87178291200

15: 1307674368000

10.4. power of 3

1:3

2:9

3: 27

4:81

5: 243

6: 729

7: 2187

8: 6561

9: 19683 10: 59049

11: 177147

12: 531441

13: 1594323

14: 4782969

15: 14348907

16: 43046721

17: 129140163

17. 123140103

18: 387420489

19: 1162261467 20: 3486784401

10.5. C(2n, n)

1: 2

2:6

3:20

4: 70

5: 252

6: 924

7: 3432

8: 12870

9: 48620

10: 184756

11: 705432

12: 2704156

13: 10400600

14: 40116600

15: 155117520

10.6. Most Divisor

<= 1e2: 60 with 12 divisors

<= 1e3: 840 with 32 divisors

<= 1e4: 7560 with 64 divisors

<= 1e5: 83160 with 128 divisors

<= 1e6: 720720 with 240 divisors

<= 1e7: 8648640 with 448 divisors

<= 1e8: 73513440 with 768 divisors

1- 100. 73313440 With 700 divisors

<= 1e9: 735134400 with 1344 divisors

<= 1e10: 6983776800 with 2304 divisors <= 1e11: 97772875200 with 4032 divisors

<= 1e12: 963761198400 with 6720 divisors

<= 1e13: 9316358251200 with 10752 divisors

<= 1e14: 97821761637600 with 17280 divisors

<= 1e15: 866421317361600 with 26880 divisors

<= 1e16: 8086598962041600 with 41472 divisors

<= 1e17: 74801040398884800 with 64512 divisors

<= 1e18: 897612484786617600 with 103680 divisors