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COMP 3270 Programming Assignment – Fall 2022

## Algorithm-1

| Step | Cost of each execution | Total # of times executed |
|------|------------------------|---------------------------|
| 1    | 1                      | 1                         |
| 2    | 1                      | n+1                       |
| 3    | 1                      | $(n^2/2)+(3n/2)$          |
| 4    | 1                      | $(n^2/2)+(n/2)$           |
| 5    | 1                      | $(n^3/2)+(5n/2)$          |
| 6    | 6                      | $(n^3/2)+(3n/2)$          |
| 7    | 4                      | $(n^2/2)+(n/2)$           |
| 8    | 2                      | 1                         |

Multiply col.1 with col.2, add across rows and simplify

$$T_1(n) = 1 + n + 1 + (n^2/2) + (3n/2) + (n^2/2) + (n/2) + (n/2) + (5n/2) + 6(n^3/2) + 6(3n/2) + 4(n^2/2) + 4(n/2) + 2 = (7n^3/2) + 3n^2 + (33/2)n + 3 = O(n^3)$$

## Algorithm-2

| Step | Cost of each execution | Total # of times executed |  |
|------|------------------------|---------------------------|--|
| 1    | 1                      | 1                         |  |
| 2    | 1                      | n+1                       |  |
| 3    | 1                      | n                         |  |
| 4    | 1                      | $(n^2/2)+(3n/2)$          |  |
| 5    | 6                      | $(n^2/2)+(n/2)$           |  |
| 6    | 4                      | $(n^2/2)+(n/2)$           |  |
| 7    | 2                      | 1                         |  |

Multiply col.1 with col.2, add across rows and simplify

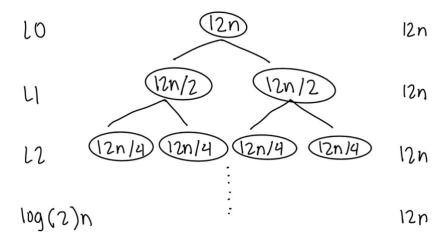
T<sub>2</sub>(n) = 1 + n + 1 + n + (n<sup>2</sup>/2)+(3n/2) + 3n<sup>2</sup> + 3n + 2n<sup>2</sup> + 2n + 2 = n<sup>2</sup>/2 + 3n/2 + 7n + 4 = n<sup>2</sup>/2 + 17n/2 + 4 = 
$$O(n^2)$$

## Algorithm-3

| Step     | Cost of each execution                           | Total # of times executed in any single recursive |
|----------|--|---|
|          |  | call  |
| 1        | 3  | 1   |
| 2        | 3  | 1   |
| Steps e  | xecuted when the input is a base case: 1-2       | 2   |
| First re | currence relation: $T(n=1 \text{ or } n=0) = 12$ |   |
| 3        | 5  | 1   |
| 4        | 2  | 1   |
| 5        | 1  | (n/2)+1   |
| 6        | 6  | n/2   |
| 7        | 4  | n/2   |
| 8        | 2  | 1   |
| 9        | 1  | (n/2)+1   |
| 10       | 6  | n/2   |
| 11       | 4  | n/2   |
| 12       | 4  | 1   |
| 13       | 5  | (cost excluding the recursive call) 1             |

| 14   | 6 | (cost excluding the recursive call) 1 |
|--|---|---------------------------------------|
| 15   | 5 | 1                                     |
| Steps executed when input is NOT a base case: 1, 3-15                                      |   |                                       |
| Second recurrence relation: $T(n>1) = 2T(n/2) + 12n + 37$                                  |   |                                       |
| Simplified second recurrence relation (ignore the constant term): $T(n>1) = 2T(n/2) + 12n$ |   |                                       |

Solve the two recurrence relations using any method (recommended method is the Recursion Tree). Show your work below:



$$T_3(n) = 12n ((log_2n) + 1) = 12nlog_2n + 12n = O(nlog_n)$$

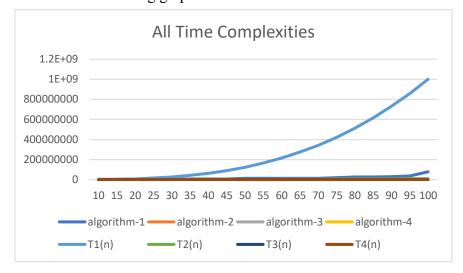
Algorithm-4

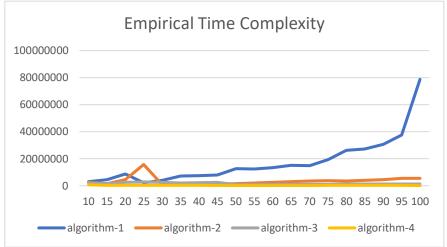
| Step | Cost of each execution | Total # of times executed |  |
|------|------------------------|---------------------------|--|
| 1    | 1                      | 1                         |  |
| 2    | 1                      | 1                         |  |
| 3    | 1                      | n+1                       |  |
| 4    | 8                      | n                         |  |
| 5    | 4                      | n                         |  |
| 6    | 2                      | 1                         |  |

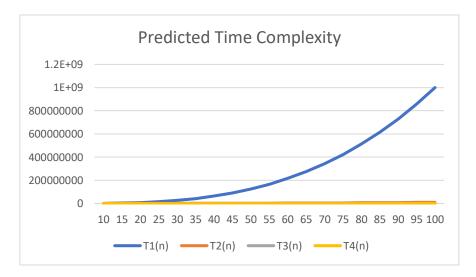
Multiply col.1 with col.2, add across rows and simplify

$$T_4(n) = 1 + 1 + (n+1) + 8n + 4n + 2 = (n+1)+12n+4 = 13n+5 = O(n)$$

After taking the data from UpadhyayYoung\_phw\_output.txt and displaying the data on a spreadsheet, we obtain the following graphs:







## **Explanation of Graphs:**

The above graphs show the rate of growth for each algorithm that is being measured. The empirical and predicted time complexities represent the actual time and predicted time, respectively. There were certain variations in the algorithmic time complexity and the scaled complexity we calculated as a reference value, especially in algorithm 2. Overall, we believe that the algorithm matches the predicted values and had the same behavior as the input sizes increased from 10 to 100.