

PBKDF2: performance matters

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PBKDF2: quick intro



Purpose

Slowly convert a password + salt into a symmetric key of some length

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Origin

RSA labs, 1999. Described in PKCS#5 and then RFC2898

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Simplification

PBKDF2 can produce arbitrary length output.

We're going to ignore this capability from here on in: only considering the first block of output.

PBKDF2: how it was described



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and typically

$$\text{PRF}(\text{pw}, x) = \text{HMAC-H}(\text{pw}, x)$$

$H = \text{SHA-1, SHA-256, SHA-512, or ...}$

Zoom, enhance

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Therefore, we need to compute 4i SHA-256 blocks.

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Actually, we only need compute $2 + 2i$ SHA-256 blocks.

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- ▶ Python (pypi pbkdf2)
- ▶ Ruby (pbkdf2 gem)
- ▶ Go (go.crypto)
- ▶ OpenBSD
- ▶ PolarSSL
- ▶ CyaSSL
- ▶ Java (OpenJDK)
- ▶ Common Lisp (ironclad)
- ▶ Perl (Crypt::PBKDF2)
- ▶ PHP
- ▶ C#

Selected performance measurements

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- ▶ Let's measure PBKDF2-HMAC-SHA1 for large iteration count (2^{22})

Selected performance measurements



Figure : PBKDF2-HMAC-SHA1, one block output, 2^{22} iterations

Selected performance measurements

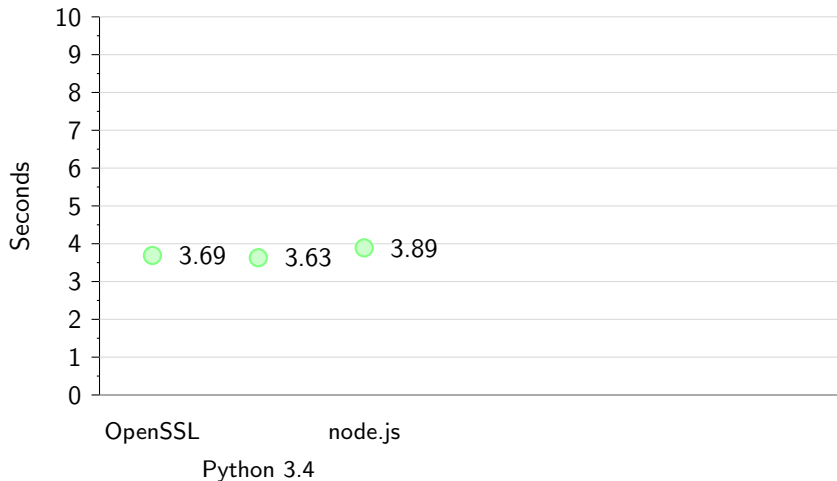


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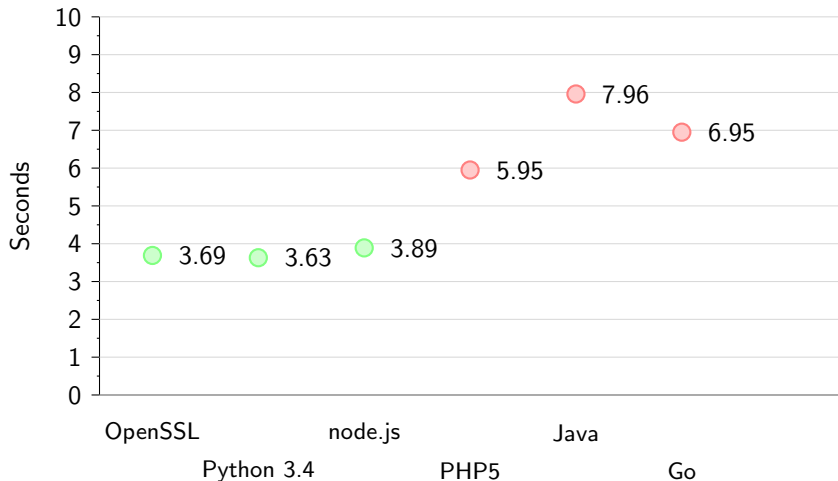


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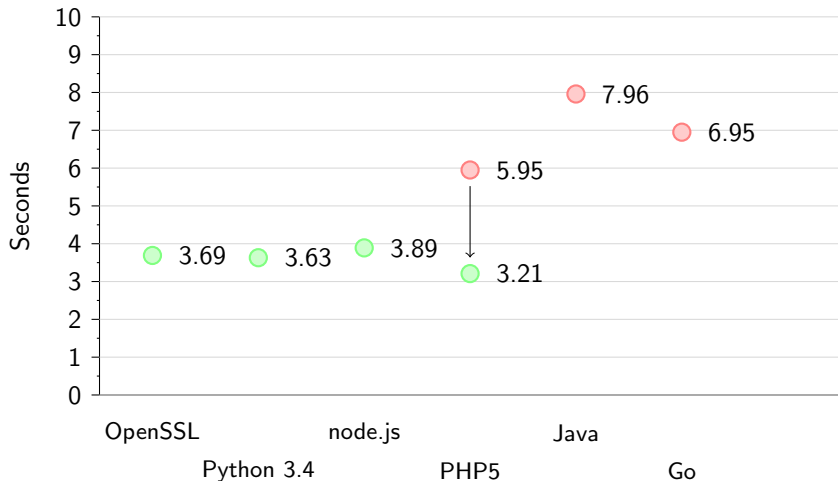


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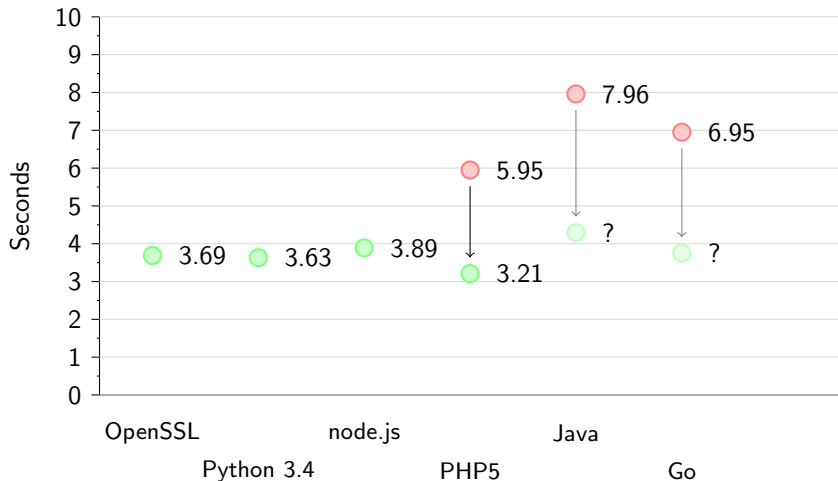


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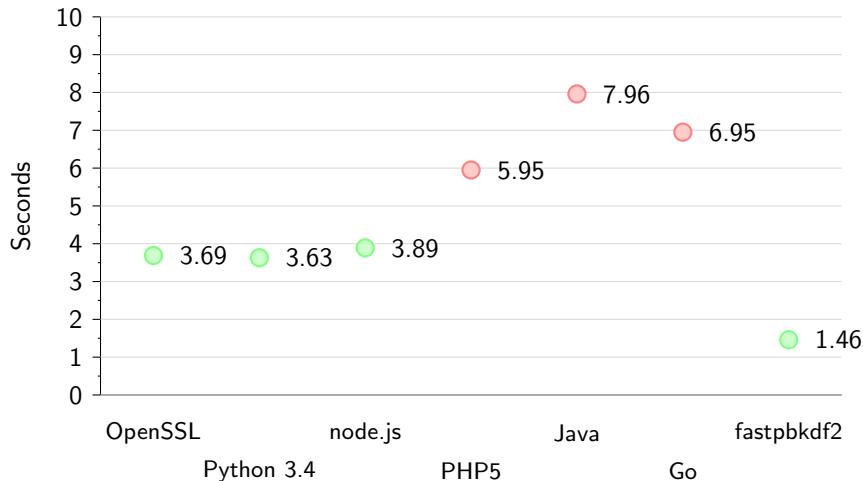


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fastpbkdf2

A faster PBKDF2-HMAC- $\{\text{SHA-1}, \text{SHA-256}, \text{SHA-512}\}$ for defenders.

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- ▶ <https://github.com/ctz/fastpbkdf2/>

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- ▶ PBKDF2 is a poor design, and described in an unhelpful way by its authors.
- ▶ Most implementations waste time and power.
- ▶ If you use PBKDF2, you can probably drop in a faster implementation and either increase security margin, or improve time/power performance.

Thank you!

Questions?

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Mail: jbp@jbp.io

Web: <https://jbp.io/>

Slides: <https://github.com/ctz/talks/>

Code: <https://github.com/ctz/fastpbkdf2/>