

# Fourier Transform Background

Monday, October 28, 2024 9:25 AM

normalization conventions differ, not important as long as you're consistent

## FOURIER TRANSFORM "FT"

All variables continuous

$$\hat{f}(\omega) = \frac{1}{\sqrt{2\pi}} \int_{\mathbb{R}} e^{-i\omega t} f(t) dt$$

" $\hat{f} = F(f)$ "

$$f(t) = \frac{1}{\sqrt{2\pi}} \int_{\mathbb{R}} e^{i\omega t} \hat{f}(\omega) d\omega$$

[\* if not integrable, define weakly (for distributions) or via density (for  $L^2$ )]

## FOURIER SERIES "FS"

Time is periodic  
Frequency is discrete

$$\hat{f}(\omega) = \frac{1}{\sqrt{2\pi}} \int_{-\pi}^{\pi} e^{-i\omega t} f(t) dt$$

$$f(t) = \frac{1}{\sqrt{2\pi}} \sum_{\omega=-\infty}^{\infty} e^{i\omega t} \hat{f}(\omega)$$

All 4 versions share many similar properties,

eg. all linear,

make differentiation multiplication by a given function, ...

Ex: CONVOLUTION, continuous version

$$(f * g)(x) := \int_{\mathbb{R}} f(y) g(x-y) dy$$

Thm  $\widehat{(f * g)} = \hat{f} \cdot \hat{g}$  i.e. pointwise multiplication.

eg.  $f * g = F^{-1}(F(f) \cdot F(g))$   
via DFT in practice

"DTFT"

## DISCRETE-TIME FOURIER TRANS.

Time is discrete  
Frequency periodic

$$\hat{f}(\omega) = \frac{1}{\sqrt{2\pi}} \sum_{t=-\infty}^{\infty} e^{-i\omega t} f(t)$$

$$f(t) = \frac{1}{\sqrt{2\pi}} \int_{-\pi}^{\pi} e^{i\omega t} \hat{f}(\omega) d\omega$$

$t \in \mathbb{Z}$

"DFT"

## DISCRETE FOURIER TRANSFORM

Time periodic and discrete, i.e. finite  
Frequency periodic and discrete, i.e. finite

$$\hat{f}(\omega) = \frac{1}{\sqrt{n}} \sum_{t=1}^n e^{-i\omega t \frac{2\pi}{n}} f(t)$$

$$f(t) = \frac{1}{\sqrt{n}} \sum_{\omega=1}^n e^{i\omega t \frac{2\pi}{n}} \hat{f}(\omega)$$

Sometimes used to approximate (cts.  $\rightarrow$  FT)

DFT defines a  $n \times n$  orthogonal

matrix

eg.  $\begin{bmatrix} f \end{bmatrix} = \frac{1}{\sqrt{n}} \begin{bmatrix} \text{DFT} \end{bmatrix} \cdot \begin{bmatrix} \hat{f} \end{bmatrix}$

So you can compute it in  $O(n^2)$  time.

But... the Fast Fourier Transform

is a (collection of) algorithms to do it in  $O(n \log n)$  time.

"Top 10 Algs" from 20th century

(modern version: 1965 Cooley + Tukey, but ideas back to Gauss 1805)

proof of convolution theorem (in limited setting)

$$\hat{f}(\omega) = \int f(x) e^{-2\pi i \omega x} dx$$

using wikipedia convention for constants

$$\hat{g}(\omega) = \int g(x) e^{-2\pi i \omega x} dx$$

let  $h = f * g$ ,  $h(x) = \int f(y) g(x-y) dy$

$$\hat{h}(\omega) = \int e^{-2\pi i \omega x} \underbrace{\int f(y) g(x-y) dy}_{h} dx$$

interchange limits  
via Fubini's theorem  
(if  $f, g \in L^1$ )

$$= \int f(y) \left( \int g(x-y) e^{-2\pi i \omega x} dx \right) dy$$

change of variables:  $\bar{x} = x - y$ ,  $x = \bar{x} + y$   
 $d\bar{x} = dx$

$$= \int \underbrace{g(\bar{x}) e^{-2\pi i \omega \bar{x}}}_{=\hat{g}(\omega)} d\bar{x} \cdot e^{-2\pi i \omega y}$$

$$= \int f(y) \hat{g}(\omega) e^{-2\pi i \omega y} dy$$

$$= \hat{g}(\omega) \underbrace{\int f(y) e^{-2\pi i \omega y} dy}_{=\hat{f}(\omega)} dy$$

$$= \hat{g}(\omega) \hat{f}(\omega) \quad \square$$

w/ different normalization  
conventions, you might  
get a factor of  $\sqrt{2\pi}$   
somewhere...

Circulant Matrix

$$\begin{bmatrix} c_0 & c_{n-1} & \dots & \dots \\ c_1 & c_0 & c_{n-1} & \dots \\ \vdots & \vdots & \ddots & \vdots \\ c_{n-1} & c_1 & \dots & c_0 \end{bmatrix}$$

Toeplitz Matrix

$$\begin{bmatrix} c_0 & c_{-1} & \dots & \dots \\ c_1 & c_0 & \dots & \dots \\ \vdots & \vdots & \ddots & \vdots \\ c_{n-1} & c_{n-2} & \dots & c_0 \end{bmatrix}$$

... diagonalized by DFT  
(i.e. can invert, find eigenvalues, ...)