

# PIEO Trees for Fun and Profit

---

We assume familiarity with [MLF+23], adopt its notational conventions, and borrow many of its definitions!

## 1 Structure & Semantics

Our PIEO trees are near identical to PIFO trees, with only two modifications:

- (1) All nodes hold PIEOs instead of PIFOs
- (2) Internal PIEOs hold just enough extra data to compute predicates of packets that live in subtrees.

To make “just enough extra data” precise, let  $f$  be a predicate on  $\mathbf{Pkt}$ . We'll presuppose the existence of a set  $\mathbf{Data}_f$  and associated surjective map  $\text{data}_f : \mathbf{Pkt} \rightarrow \mathbf{Data}_f$ . This map and set are such that, for  $\text{pkt}, \text{pkt}' \in \mathbf{Pkt}$ ,

$$\text{data}_f(\text{pkt}) = \text{data}_f(\text{pkt}') \implies f(\text{pkt}) = f(\text{pkt}')$$

In other words, all preimages of  $d \in \mathbf{Data}_f$  evaluate to the same truth value under  $f$ . For this reason,

**Abuse of Notation 1.1.** We'll often write  $f(d)$  when we mean  $f(\text{pkt})$  for some  $\text{pkt} \in f^{-1}(d)$ .

Broadly, this set  $\mathbf{Data}_f$  is designed to encode just the parts of  $\mathbf{Pkt}$  necessary for computing  $f(\text{pkt})$ .

**Example 1.2.** For any  $f$ , it is legal to set  $\mathbf{Data}_f = \mathbf{Pkt}$  and  $\text{data}_f$  to be the identity map. We'd expect this to work since each  $\text{pkt} \in \mathbf{Pkt}$  certainly holds all the information necessary to compute  $f(\text{pkt})$ .

We can now formally define PIEO trees!

**Definition 1.3** (PIEO tree). Let  $t \in \mathbf{Topo}$  be a topology and  $f$  a predicate over  $\mathbf{Pkt}$ . The set of *PIEO trees* over  $t$  and  $f$ , denoted  $\mathbf{PIEOTree}(t, f)$ , is defined inductively by

$$\frac{p \in \mathbf{PIEO}(\mathbf{Pkt}, f)}{\text{Leaf}(p) \in \mathbf{PIEOTree}(*)} \quad \frac{n \in \mathbb{N} \quad ts \in \mathbf{Topo}^n \quad \begin{array}{l} p \in \mathbf{PIEO}(\{1, \dots, n\} \times \mathbf{Data}_f(\mathbf{Pkt}), f) \\ \forall i \in [1, n]. qs[i] \in \mathbf{PIEOTree}(ts[i]) \end{array}}{\text{Internal}(qs, p) \in \mathbf{PIEO}(\text{Node}(ts))}$$

## References

- [MLF<sup>+</sup>23] Anshuman Mohan, Yunhe Liu, Nate Foster, Tobias Kappé, and Dexter Kozen. Formal abstractions for packet scheduling, 2023.