

Exercises on Description Logics

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1. Interpretations and models

1. Prove that the DL expressions

$$E1 = \forall r.(C \sqcup D)$$

$$E2 = (\forall r.C) \sqcup (\forall r.D)$$

are not equivalent by finding an interpretation $\mathcal{I} = (C^{\mathcal{I}}, D^{\mathcal{I}}, r^{\mathcal{I}})$ such that $E1^{\mathcal{I}} \neq E2^{\mathcal{I}}$.

2. Find an interpretation \mathcal{I} that satisfies the following axioms:

1. $A \equiv \forall s.L$,
2. $L \sqsubseteq (\exists r.A) \sqcap (\forall r.(\exists s.\top))$,
3. $L(u)$

3. Find an interpretation \mathcal{I} that satisfies the following axioms:

1. $A \sqsubseteq \forall s.L$,
2. $L \sqsubseteq (\exists r.A) \sqcap (\forall r.(\exists s.\top))$,
3. $L(u)$

4. Find an acyclic interpretation that satisfies the above axioms.

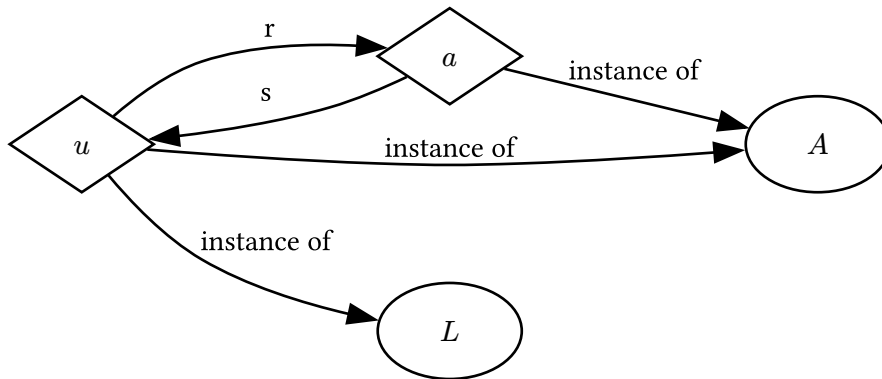
Solutions

1. With the interpretation $\Delta^{\mathcal{I}} = \{a, c, d\}$, $C^{\mathcal{I}} = \{c\}$, $D^{\mathcal{I}} = \{d\}$, $r^{\mathcal{I}} = \{(a, c), (a, d)\}$ we have

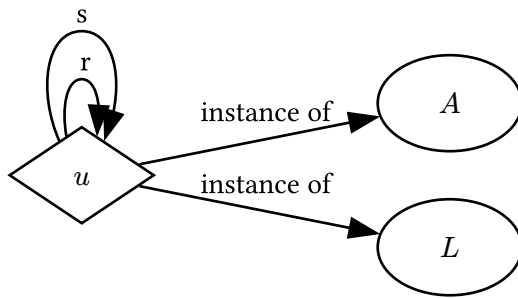
$$A^{\mathcal{I}} = (\forall r.(C \sqcup D))^{\mathcal{I}} = \{a\} \text{ and}$$

$$B^{\mathcal{I}} = ((\forall r.C) \sqcup (\forall r.D))^{\mathcal{I}} = ((\forall r.C))^{\mathcal{I}} \cup (\forall r.D))^{\mathcal{I}} = \emptyset \cup \emptyset = \emptyset$$

2. An interpretation that satisfies the axioms: $\Delta^{\mathcal{I}} = \{u, a\}$, $L^{\mathcal{I}} = \{u\}$, $A^{\mathcal{I}} = \{u, a\}$, $r^{\mathcal{I}} = \{(u, a)\}$, $s^{\mathcal{I}} = \{(u, u)\}$. Note that u must be in $A^{\mathcal{I}}$ because it trivially satisfies $\forall s.L$.

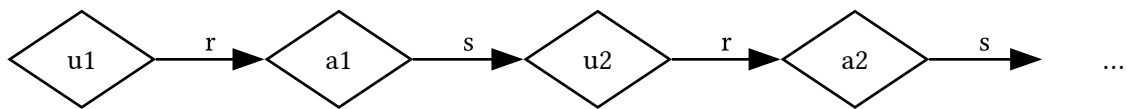


3. A minimal interpretation that satisfies the axioms: $\Delta^{\mathcal{I}} = \{u\}$, $L^{\mathcal{I}} = \{u\}$, $A^{\mathcal{I}} = \{u\}$, $r^{\mathcal{I}} = \{(u, u)\}$, $s^{\mathcal{I}} = \{(u, u)\}$



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4. An acyclic interpretation: $\Delta^{\mathcal{I}} = \{u_1, u_2, \dots, a_1, a_2, \dots\}$, $L^{\mathcal{I}} = \{u_1, u_2, \dots\}$, $A^{\mathcal{I}} = \{a_1, a_2, \dots\}$, $r^{\mathcal{I}} = \{(u_1, a_1), (u_2, a_2), \dots\}$, $s^{\mathcal{I}} = \{(a_1, u_2), (a_2, u_3), \dots\}$. This interpretation is infinite.



2. DL Logical consequences

Source: Rudolph, S. (2011). Foundations of Description Logics. Longman Publishing. <http://www.aifb.kit.edu/images/1/19/DL-Intro.pdf>

Consider the knowledge base \mathcal{KB} with

RBox

- $\text{owns} \sqsubseteq \text{caresFor}$

TBox

- $\text{Healthy} \sqsubseteq \neg \text{Dead}$
- $\text{Cat} \sqsubseteq \text{Dead} \sqcup \text{Alive}$
- $\text{HappyCatOwner} \sqsubseteq \exists \text{ owns} . \text{Cat} \sqcap \forall \text{ caresFor} . \text{Healthy}$

ABox

- $\text{HappyCatOwner}(\text{schrödinger})$
- $\text{Cat}(\text{hbar})$,
- $\text{owns}(\text{schrödinger}, \text{psi}), \text{owns}(\text{schrödinger}, \text{hbar})$

Decide whether the following propositions about the knowledge base are true and give evidence:

1. \mathcal{KB} is satisfiable,
2. $\mathcal{KB} \models \text{Alive}(\text{schrödinger})$,
3. $\mathcal{KB} \models \text{Dead} \sqcap \text{Alive} \sqsubseteq \perp$,
4. $\mathcal{KB} \models \text{Alive}(\text{psi})$,
5. $\mathcal{KB} \models \text{Alive}(\text{hbar})$.

Solutions

1. \mathcal{KB} is satisfiable, with the model \mathcal{M} defined as

- $\Delta^{\mathcal{M}} = \{s, h, p\}$

- $\text{schrödinger}^{\mathcal{M}} = s, \text{psi}^{\mathcal{M}} = p, \text{hbar}^{\mathcal{M}} = h,$
 - $\text{Cat}^{\mathcal{M}} = \{h, p\},$
 - $\text{HappyCatOwner}^{\mathcal{M}} = \{s\}$
 - $\text{Healthy}^{\mathcal{M}} = \{h, p\},$
 - $\text{Alive}^{\mathcal{M}} = \{h\},$
 - $\text{Dead}^{\mathcal{M}} = \{\},$
 - $\text{owns}^{\mathcal{M}} = \{(s, p), (s, h)\},$
 - $\text{caresFor}^{\mathcal{M}} = \{(s, p), (s, h)\},$
2. $\mathcal{KB} \models \text{Alive}(\text{schrödinger})$ is false. For instance in the model \mathcal{M} $s \notin \text{Alive}(\text{schrödinger})^{\mathcal{M}}$
 3. $\mathcal{KB} \models \text{Dead} \sqcap \text{Alive} \sqsubseteq \perp$ is false. If we add s to $\text{Alive}^{\mathcal{M}}$ and $\text{Dead}^{\mathcal{M}}$ we obtain a model of \mathcal{KB} in which the interpretation of $\text{Dead} \sqcap \text{Alive}$ is not empty.
 4. $\mathcal{KB} \models \text{Alive}(\text{psi})$ is false, in the model \mathcal{M} p is not in $\text{Alive}^{\mathcal{M}}$
 5. $\mathcal{KB} \models \text{Alive}(\text{hbar})$ is true. In every model \mathcal{J} of KB h must be in $\text{Cat}^{\mathcal{J}}$ and in $\text{Healthy}^{\mathcal{J}}$ because of $\text{HappyCatOwner} \sqsubseteq \exists \text{ owns}.\text{Cat} \sqcap \forall \text{ caresFor}.\text{Healthy}$ so it must be in $\neg \text{Dead}^{\mathcal{J}}$ because of $\text{Healthy} \sqsubseteq \neg \text{Dead}$ and hence in $\text{Alive}^{\mathcal{J}}$ because $\text{Cat} \sqsubseteq \text{Dead} \sqcup \text{Alive}$

3. Inference

Use the tableau DL algorithm to show that $(\forall r.(B \sqcap C)) \sqcap (\exists r.(\neg C))$ is not satisfiable

Solution

Let's start with the ABox

$$\{(\forall r.(B \sqcap C)) \sqcap (\exists r.(\neg C))(a)\}$$

Applying the \sqcap rule adds

$$\begin{aligned} &(\forall r.(B \sqcap C))(a), \\ &(\exists r.(\neg C))(a) \end{aligned}$$

The \exists rule adds

$$r(a, b), (\neg C)b$$

The \forall rule adds

$$(B \sqcap C)(b)$$

The \sqcap rule adds

$$\begin{aligned} &B(b), \\ &C(b) \end{aligned}$$

There is a clash between $\neg C(a)$ and $C(a) \Rightarrow$ the starting concept was inconsistent.

4. OWL-RL inference

Use the OWL-RL inference rules to find the ABox logical consequences of the following knowledge bases

\mathcal{KB}_1 :

1. $\exists \text{ owns} . \text{Cat} \sqsubseteq \text{CatOwner}$
2. $\exists \text{ owns} . \text{Mouse} \sqsubseteq \text{MouseOwner}$

3. $\text{CatOwner} \sqsubseteq \forall \text{ owns. Cat}$
4. $\text{owns}(\text{felix}, \text{max})$,
5. $\text{owns}(\text{felix}, \text{illi})$,
6. $\text{owns}(\text{marta}, \text{chese})$,
7. $\text{owns}(\text{marta}, \text{bob})$,
8. $\text{Cat}(\text{max})$,
9. $\text{Cat}(\text{bob})$
10. $\text{Mouse}(\text{chese})$

Solution

| Rule | Applied on (IF) | Infers (THEN) |
|-------------|-----------------|---------------------------------------|
| some | 1, 4, 8 | 11. $\text{CatOwner}(\text{felix})$ |
| some | 1, 7, 9 | 12. $\text{CatOwner}(\text{marta})$ |
| some | 2, 6, 10 | 13. $\text{MouseOwner}(\text{marta})$ |
| only | 3, 11, 5 | 14. $\text{Cat}(\text{illi})$ |
| only | 3, 12, 10 | 14. $\text{Cat}(\text{chese})$ |

So, chese is both a Cat and a Mouse 🤖

Exercise 6. Closing the world

Consider the knowledge base \mathcal{KB}

TBox

1. $\text{BlueRun} \sqsubseteq \text{SkiRun}$
2. $\text{BlueRun} \sqsubseteq \text{SkiRun}\}$,
3. $\text{RedOnlyPlace} \equiv \text{Place} \sqcap \exists \text{ isEndOf . SkiRun} \sqcap (\forall \text{ isEndOf . RedRun})$

ABox

4. $\text{isEndOf}(\text{p1}, \text{sr1})$,
5. $\text{RedRun}(\text{sr1})$,
6. $\text{Place}(\text{p1})$

What must be added to the ABox and/or TBox to ensure that $\mathcal{KB} \models \text{RedOnlyPlace}(\text{p1})$? Test your answer with a DL reasoner (in Protégé).

Solution

For example:

7. $\text{IsEndOfOnlyOneRun} \equiv \leq_1 \text{ isEndOf . } \top$
8. $\text{IsEndOfOnlyOneRun}(\text{p1})$