

Current Limitations of LLM Guided Program Synthesis

First Author¹, Second Author², Third Author^{2,3}, Fourth Author⁴

¹First Affiliation

²Second Affiliation

³Third Affiliation

⁴Fourth Affiliation

{first, second}@example.com, third@other.example.com, fourth@example.com

Abstract

Large Language Models (LLMs) exhibit strong performance for program synthesis tasks including text-to-SQL, up to 86.6% accuracy on the Spider 1.0 benchmark. The top-performing methods bridge the semantic gap between natural language and SQL by using prompt engineering, fine-tuning, or in-context learning techniques; however, these purely neural techniques remain prone to hallucinations. Existing symbolic approaches to improve accuracy on program synthesis tasks generally rely on post-hoc verification, which fails to constrain unproductive path exploration, or constrained decoding techniques that rely on the underlying distribution of the model. The effectiveness of symbolic guidance under complex queries and the semantic ambiguity inherent in text-to-SQL benchmarks, however, remains unclear. In this work, we study the “ground truth” errors made by LLMs in the text-to-SQL task looking at both fine-tuned models and general-purpose models optimized for code generation. We observe that errors can be grouped into five error categories, two of which are amenable to simple symbolic fixes. Importantly, we find that for many queries, correct predictions do not rank near the top of the LLMs likely predictions, implying that *constrained decoding alone is insufficient* to guide synthesis on hard problems. Our findings provide interpretable insight into where LLMs fail and suggest promising, neuro-symbolic directions for improving LLM guided program synthesis.

1 Introduction

[LLMs are used for program synthesis, but errors can compound if the LLM starts off on the wrong track.] Large Language Models (LLMs) are increasingly used as program generators to synthesize working code from natural language prompts. In contrast to other program synthesis techniques, LLMs are scalable to larger and more complex programs; however, due to their autoregressive nature, they can get stuck down an incorrect path as errors can propagate throughout the decoding process which ultimately produces invalid programs.

[Constrained decoding has shown to improve generation of syntactically valid and typesafe programs by enforcing constraints at decoding time, but constrained decoding techniques with richer semantic constraints, such as those likely required to effectively synthesize SQL queries, have yet to be developed.]

[In this paper we study the use of LLMs to synthesize SQL queries; categorizing the mistakes they make to evaluate whether constrained decoding or other symbolic reasoning techniques could steer LLMs to correct queries.]

[LLMs predict queries token by token and constrained decoding enforces constraints at decoding time, so we focus our study on the incremental errors LLMs make.] Recent LLM augmented program synthesis techniques have focused on enforcing constraints at decoding time to ensure certain properties (Nagy et al. 2026). These techniques are often evaluated in terms of final outputs. However, constrained decoding operates at the lexeme level. To understand its effectiveness, we must analyze the model’s learned distribution for lexemes during decoding, since a constraint can only be enforced if the appropriate lexeme appears within the model’s top-k candidates.

[Existing SQL error taxonomies are based on fully generated queries, so we introduce a new taxonomy focusing on the incremental errors LLMs make.] We conduct our lexeme-level analysis on a subspace of the program synthesis domain, the text-to-SQL task, as the correctness of an SQL query can be accurately assessed. Existing taxonomies analyze errors of fully generated queries (Qu et al. 2024). However, SQL synthesis is particularly dependent on lexeme-level correctness. In addition to overall query structure, SQL queries are dependent on correct schema references and correct operator usage.

[A taxonomy of incremental SQL prediction errors then allows us to observe the potential benefit of symbolic repair and the limitations of constrained decoding for LLM-driven SQL synthesis.] We evaluate the ability of LLMs to synthesize correct SQL queries at the lexeme level and find that we can broadly categorize lexeme level errors into six categories. We find that many errors within these categories can be mitigated with simple symbolic fixes and constrained decoding. We also observe limitations of constrained decoding. It cannot correct errors if the correct lexeme is absent from the model’s top-k candidates.

Contributions:

- We present a hierarchical taxonomy of error categorization for lexeme level LLM mistakes during the decoding process.
- We identify error categories that are amenable to symbolic

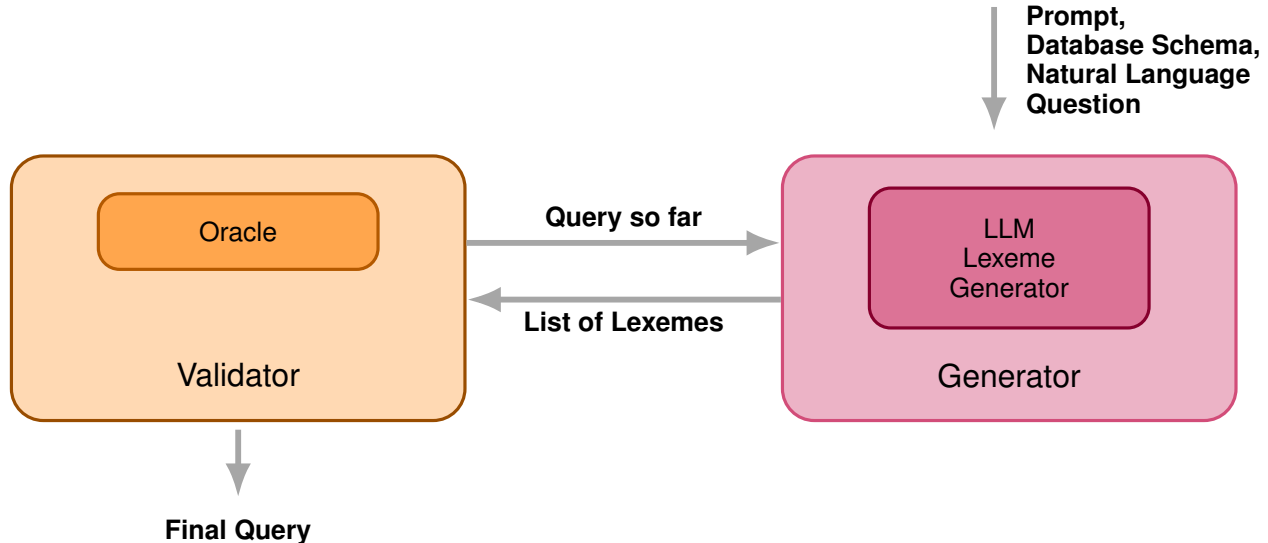


Figure 1: The overall system architecture. The Language Model Generator is prompted to generate a SQL query given a database schema and natural language question. The generator returns a list of lexemes to the Validator which incrementally validates the most probable lexeme against the gold one and corrects the generator if there is a mismatch.

repair and quantify an upper bound on automatically correctable failures.

- We introduce a constrained decoding methodology, *schema constrained decoding* to address schema-based errors.

2 Related Work

[Text-to-SQL is a task in which SQL queries are generated from NL questions, and many high performing methods use techniques such as prompt engineering and in-context learning.] Text-to-SQL is a program synthesis task concerned with SQL generation for natural language questions. Recent research has focused on leveraging Large Language Models (LLMs) for this task and have performed well empirically, reaching test suite accuracies of up to 86.6% (Gao et al. 2024) against the Spider 1.0 dataset (Yu et al. 2018). Many of the methods that top the Spider 1.0 leaderboard involve prompt engineering and in-context learning strategies. These methods often treat LLMs as black box synthesizers but are able to scale to complex queries.

[Traditional program synthesis techniques provide guarantees, but don't scale.] Before the advent of LLMs, symbolic approaches to SQL synthesis explored enumerative search and program sketching to guarantee correct query generation (Wang, Cheung, and Bodik 2017; Takenouchi et al. 2021). These works focus on pruning the search space either through additional information provided by the user, such as the constants that are allowed to appear in a query (Wang, Cheung, and Bodik 2017; Takenouchi et al. 2021), or through pruning unproductive paths (Wang, Cheung, and Bodik 2017; Takenouchi et al. 2021). These traditional methods often include soundness and bounded completeness guarantees, but struggle to scale to larger, more complex queries.

[Recent work to aid LLM synthesis focuses on enforcing constraints at decoding time.] Recent work has focused on neuro-symbolic approaches to bridge the scalability of LLM aided program generation and more symbolic approaches. Constrained decoding methods aim to enforce certain properties during the decoding process. Techniques such as PICARD (Scholak, Schucher, and Bahdanau 2021) enforce SQL-specific constraints during token generation. Other methods can enforce semantic constraints (Nagy et al. 2026) or type safety (Mündler et al. 2025) more generally in program synthesis. Constrained decoding approaches improve syntactic and semantic correctness but depend on the correct token appearing within the model’s top-k predictions based on its learned distribution.

[There are error taxonomies for text-to-SQL, but they consider fully generated queries.] Prior work has been done to analyze errors for text-to-SQL, but they are concerned with categorizing fully generated queries (Qu et al. 2024; Shen et al. 2025). These taxonomies do not capture the lexeme-level mistakes that arise during the decoding process.

[Our work looks at the lexeme level efficacy of constrained decoding and symbolic fixes.] Our work addresses the gap of understanding the effectiveness of constrained decoding by analyzing lexeme-level errors made by LLMs for text-to-SQL. By categorizing incremental mistakes and identifying those amenable to symbolic repair, we provide insight into where constrained decoding can succeed and where simple symbolic fixes can aid LLM guided program synthesis.

3 Our Approach

Our overall system is shown in Figure 1.

3.1 Language Model (LM) Generator

Consider \mathcal{L} to be an (autoregressive) Language Model created for a sequence generation task that has a predefined vocabulary Σ , consisting of some k tokens which are used as part of the tokenization. At each timestep $i \in [1, n]$, the model generates a sequence of these tokens, (t_1, t_2, \dots, t_n) based on an autoregressive decoding process where the token with the highest probability is selected. This is usually part of a heuristic or sampling based search procedure on the set of k tokens to predict the next most probable token. Now, consider the LM to be finetuned on a new DSL and the tokens in the vocabulary are not changed. This means that the LM is now capable of producing the probability distribution over the same tokens but conforming to the changed ordering of our DSL *i.e.* if an original ordering of tokens $w \in \Sigma$ exists, then a new ordering of tokens $w^* \in \Sigma$ should also be produced by the LM. A \mathcal{L} generates tokens and concatenates them by taking the token with the highest probability and generates a new token and continues until a whitespace token is generated or a max heap size is reached.

Example 3.1. Consider the original token sequence $w = \text{SELECT}, *, \text{FROM}, \text{books}$, drawn from SQL queries present in the model’s training data. Using our DSL, this sequence is transformed into $w^* = \text{PROJECT}, *, \text{FROM}, \text{books}$. In this transformation, the only substituted token is **PROJECT**, and we therefore expect the model’s vocabulary to assign a non-zero probability p to this token.

For commonly occurring words, we can fully expect the new word to also exist in the vocabulary, however, if the word does not exist we know that some n -gram combination of the word does exist. When for example the token, **REGENERATE** $\notin \Sigma$ then we can form the same word by some combination like $\{\text{RE}, \text{GEN}, \text{E}, \text{RATE}\} \in \Sigma$. In our approach, we use QWEN-7B (Bai et al. 2023), which uses an open-source fast Byte Pair Encoding (BPE) for tokenization and $\text{len}(\Sigma) \in 151642$.

Since we use the language model as an autoregressive next token predictor, we would also like to harness the power of the probabilistic nature in which the tokens are ranked. Usually, since the top-1 token is taken as the next prediction, there is an uncertainty that whether this next token is part of the valid query that satisfies the question that is requested by the user. Hence, we would also like to store, some top- m tokens at each time step and perform a beam search when on these token sets when there is a mismatch in the token requested and the query validated.

Language Model Validator. With each generated lexeme, we validate the generated lexeme by evaluating it against the expected lexeme of the gold query. If the most probable lexeme does not match the gold lexeme, we consider it to be a mistake. We compare the generated lexeme with the gold lexeme while considering issues that arise with aliasing for column names by only considering the column name. For example *T1.BehaviorMonitoring* and *BehaviorMonitoring* would both be considered correct lexemes compared to the gold lexeme *BehaviorMonitoring*. If the top most probable lexeme does not match the gold lexeme, we guide the

Language Model by correcting it with the gold lexeme and continue generating. We incrementally validate each generated lexeme to guide the Language Model towards more productive generation paths rather than allow it to fully generate a query before evaluating its correctness.

Our system architecture is outlined in Algorithm 1.

Algorithm 1 System Algorithm

```

1: Require: Language Model  $\mathcal{L}$ , beam width  $m$ , gold
   query  $g$  and some string prompt  $\mathcal{P}$ , number of mistakes
    $n$ 
2: procedure LANGUAGEMODELVALIDATOR( $\mathcal{P}, w, g, n$ )
3:   if  $g == w$  returns TRUE then
4:      $\mathcal{P} \leftarrow \mathcal{P} \parallel w$ 
5:     return  $\mathcal{P}, n$ 
6:   else
7:      $\mathcal{P} \leftarrow \mathcal{P} \parallel g$ 
8:      $n \leftarrow n + 1$ 
9:     return  $\mathcal{P}, n$ 
10:  end if
11: end procedure
12: procedure ORCHESTRATOR()
13:    $n \leftarrow 0$ 
14:   for  $i = 1, 2, \dots, \text{len}(\text{gold\_query})$  do
15:      $w_i \leftarrow \mathcal{L}(\mathcal{P})$ 
16:      $g_i \leftarrow \text{GOLDLEXEME}(g, i)$ 
17:      $\mathcal{P}, n \leftarrow \text{LANGUAGEMODELVALIDATOR}(\mathcal{P}, w_i, g_i, n)$ 
18:   end for
19:   return  $n$ 
20: end procedure

```

4 Experimental Setup

4.1 Implementation

Our system is implemented in Python and leverages open-source LLMs hosted on HuggingFace for text-to-SQL generation, including XiYanSQL-QwenCoder-7B-2504, deepseek-coder-6.7b-instruct, and XiYanSQL-QwenCoder-14B-2504. Models are loaded in 4-bit quantized mode when possible to reduce GPU memory usage, and generation is performed on A100 NVIDIA GPUs with 40GB VRAM for XiYanSQL-QwenCoder-7B-2504, deepseek-coder-6.7b-instruct and 80GB for XiYanSQL-QwenCoder-14B-2504 with automatic device mapping and offloading for larger models.

4.2 Model Choice

Finetuned models XiYanSQL-QwenCoder-7B-2504 and XiYanSQL-QwenCoder-14B-2504 were chosen due to their high performance on the Spider dataset (Liu et al. 2025). We also chose a general purpose model deepseek-coder-6.7b-instruct.

4.3 Dataset

From the test split of the Spider 1.0 dataset (Yu et al. 2018), we evaluate the queries deemed hard and extra as we are most interested in errors that arise from complex queries.

| Error Category | deepseek-coder-6.7b-instruct | XiYanSQL-7B | XiYanSQL-14B |
|-------------------------------------|------------------------------|-------------|--------------|
| Structural / Clause Ordering | 1385 | 583 | 472 |
| Limit Errors | 245 | 235 | 243 |
| Other Structural Errors | 1140 | 348 | 229 |
| Schema Grounding | 938 | 1179 | 1255 |
| Value Errors | 26 | 58 | 107 |
| Other Schema Errors | 912 | 1121 | 1148 |
| Operator Errors | 1447 | 312 | 202 |
| Where and Having Errors | 658 | 272 | 179 |
| Other Operator Errors | 789 | 40 | 26 |
| Aggregate Function Errors | 214 | 205 | 190 |
| Alias Errors | 34 | 45 | 864 |
| Other / Unclassified | 419 | 415 | 235 |
| Total Errors | 4437 | 2739 | 3221 |

Table 1: Lexeme-level error distribution across models.

5 Taxonomy

In this section, we outline our error taxonomy.

5.1 Schema-based Errors

Schema-based errors are defined as errors in which the model fails to correctly reference the database schema. These errors typically involve either selecting an incorrect table or column name, hallucinating schema attributes, or choosing an inappropriate value from the database for a condition.

Incorrect Naming These errors are defined as ones where the model generates a lexeme that does not correspond to the correct schema element. For example, the model might predict `FIXME` instead of the correct column `FIXME`.

Symbolic Fix: Easy To address the generation of hallucinated schema values, we propose a symbolic fix *schema-based constrained decoding* in which the first valid schema element that appears in the top-k is chosen. If no valid element is initially predicted, one can force the model to select from the set of valid schema elements, often yielding the correct token. Based on empirical results presented in `FIXME`, the highest valid schema element is `FIXME` percentage of the time the correct one. Such a technique would fix `FIXME` number of errors.

Incorrect Values These errors typically occur in `WHERE` conditions, where the model selects an incorrect literal value, even though the query structure and schema reference are correct. For instance, given the question:

“What are the addresses, towns, and county information for all customers who live in the United States?”

the gold query is:

```
SELECT address_line_1, town_city, county
FROM Customers
WHERE Country = 'USA';
```

A model might predict ‘United’ as the value, which reflects a plausible misunderstanding of the question’s phrasing. Correctly mapping natural language values to the database representation often requires explicitly grounding the question in the schema or the database contents.

Symbolic Fix: Moderate One approach is to align natural language mentions with schema values, using either string similarity or enumeration over the database domain. This allows the model to select or rank candidates that correspond to actual entries in the database, mitigating errors in value selection.

Wrong Naming these are errors like wrong column or table name (the gold is present in the schema but the predicted candidate is not)

Symbolic fix: Easy- if valid schema element in top-k then you use a schema based constrained decoding to select (usually the gold is the top most). If there is no valid schema element, have the LLM pick from or rank valid schema elements

Wrong Values For where conditions - the wrong value is selected for example - “What are the addresses, towns, and county information for all customers who live in the United States?” but the gold query is

```
SELECT $address_line_1$ , $town_city$ ,
county FROM Customers WHERE Country = 'USA'
```

The LLM predicts ‘United’ which makes sense given the phrasing of the question but involves parsing out all the values of the schema to “infer” that “United States” maps to “USA”

Symbolic Fix: Medium - We may be able to infer the correct condition but we need a way to map that question to the actual values in the database

5.2 Structural based errors

These are elements that are a violation of the expected SQL structure (ie clause based errors).

LIMIT errors The LLM doesn't predict a LIMIT at the end of a full query.

Symbolic Fix: Easy - we basically just look at the output and if it matches except the number of rows we just add a LIMIT

Everything else

5.3 Where Conditions

Wrong operator wrong operator (sometimes the LLM will correctly predict bigger/ greater but not >

Symbolic fix: easy for wrong operator - if LLM predicts bigger then just replace it with >

upon further inspection there aren't many of these that can be easily fixable (only about 2 were bigger/greater instead of >) but I'll leave this category in for now

5.4 Aggregate Errors

Wrong aggregate functions like COUNT, MAX, etc

Symbolic fix: hard

5.5 Other

catch all

includes syntactically/ semantically "similar" ie they look right....but aren't

Symbolic fix: hard

6 Conclusion

LLMs are being increasingly used for program synthesis but solely relying on their probabilistic token ranking can lead to incorrect programs. In this work, we investigate these limitations and propose a lexeme level validation methodology to guide the LLM towards semantically valid queries. We find that many errors can be corrected by using a constraint decoding and beam-search guided reasoning technique to rerank the generated lexemes. For errors that cannot be corrected using these methods, purely symbolic techniques can be used to aid program synthesis.

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