

# Analysis and Attacks of decentralized content curation platforms

Andr s Monteoliva Mosteiro and Orfeas Stefanos Thyfronitis Litos

University of Edinburgh  
o.thyfronitis@ed.ac.uk, a.monteoliva@serious.server

**Abstract.** We will attack Steem.

## 1 Introduction

Steem is not incentive-compatible.

## 2 Related Work

Many people have done many similar things.

## 3 Model

### 1 Properties

**Players** There are  $N \in \mathbb{N}^*$  players in the model. A player  $u$  is defined by its Steem Power  $SP$ , its Voting Power  $VP$ , its Likability Distribution  $L$  and its Strategy  $S$ . Let the  $i$ th player be characterized with the following tuple  $u_i = (SP_i, VP_i, L_i, S_i)$  where  $SP \in \mathbb{N}$ ,  $VP \in [0, 100]$ ,  $L \in \mathcal{D}([0, 1]^N)$  and  $S \in \{H, G\} \times [100] \times 2^{[N]}$ . Lets describe each of the properties of a player in more detail:

- The Steem Power funds of player  $U_i$  are defined as  $sp_i \in \mathbb{N}$ . The vector of Steem Power funds for the  $N$  players is defined as  $\mathcal{SP} = (sp_1, \dots, sp_i, \dots, sp_n) \forall i \in N$ .
- The Voting Power...
- A player  $u_i$  will have a "Likability Distribution"  $L_i$  which determines how likely is the  $j$ th to like a post created by  $u_i$ , defined as  $L_i \in \mathcal{D}([0, 1]^N)$ . The Likability Distribution for all the players will be  $\mathcal{L} = (L_1, \dots, L_i, \dots, L_n) \forall i \in N$

- The strategy of player  $u_i$  is defined as  $s_i \in (\{H, G\}, Min, R)$ , where  $H \equiv honest$  and  $G \equiv greedy$ ,  $Min \in [0, 1]$  is the Minimum Voting Power the player can reach and  $r_i \in \mathbb{N}$  is the size of  $u_i$  Voting Ring. If player  $u_i$  is *Honest*, her Voting Ring  $r_i = \emptyset$ . If player  $u_j$  is *Greedy*, her Voting Ring  $r_j = p_1, \dots, p_k, \dots, p_r$  where  $p_k \in [1, N]$ . The vector of the strategies for the  $N$  players is defined as  $\mathcal{S} = (s_1, \dots, s_i, \dots, s_n) \forall i \in N$

The set of players is defined as  $\mathcal{U} = (u_1, \dots, u_i, \dots, u_n) \forall i \in [N]$ .

## Posts

- Each player can create one post. The post created by  $U_i$  is defined as  $p_i$ . The set of all posts will be then  $\mathcal{P} = \bigcup_{i=1}^N p_i$ .
- The likability of a post will be defined as  $l_i \in [0, 1]^N$ , where  $l_i \sim L_i$  is retrieved at random following the "Likability Distribution" of player  $U_i$ .

## 2 Game Execution

---

**Algorithm 1** Each player creates a post according to Likability Distribution

---

```

1: function POSTGENERATION( $N, \mathcal{L}$ )
2:    $\mathcal{P} = \emptyset$  ▷ List of posts
3:   for  $i \in N$  do
4:      $l_i \xleftarrow{R} \mathcal{L}_i$  ▷ Get likability of posts
5:      $p \leftarrow (i, l_i)$ 
6:      $\mathcal{P} \leftarrow \mathcal{P} \parallel p$  ▷ Add post to list of Posts
7:   end for
8:    $\mathcal{P} \leftarrow \text{SHUFFLE}(\mathcal{P})$  ▷ Shuffle the order of Posts
9:   return  $\mathcal{P}$ 
10: end function

```

---

---

**Algorithm 2** Player casts votes according to her strategy and until she reaches her Min Voting Power

---

```

1: function VOTE( $player, \mathcal{P}, s, sp$ )
2:   switch  $s$  do
3:     case Honest
4:       for  $p \in \mathcal{P}$  do ▷ Iterate over all the posts
5:         ▷ If user likes the post and has not reached min VPower
6:         if  $l_p > 0 \wedge p.VPower > s.Min$  then
7:            $voteValue \leftarrow p.VPower \cdot l_p \cdot sp$ 
8:            $p \leftarrow p.votes + voteValue$ 
9:         end if
10:      end for
11:    end case
12:    case Greedy
13:      for  $p \in \mathcal{P}$  do ▷ Iterate over all the posts
14:        ▷ If post belongs to voting ring and not reached min VPower
15:        if  $p \in s.R \wedge p.VPower > s.Min$  then
16:           $voteValue \leftarrow p.VPower \cdot weight \cdot sp$ 
17:           $p \leftarrow p.votes + voteValue$ 
18:        end if
19:      end for
20:    end case
21:  end switch
22:   $\mathcal{P} \leftarrow \text{ORDER}(\mathcal{P}.votes)$  ▷ Order posts by value of votes received
23:  return  $\mathcal{P}$ 
24: end function

```

---



---

**Algorithm 3** Players cast votes over the maxRounds

---

```

1: function CURATION( $N, S, \mathcal{P}, SP, maxR$ )
2:   round = 0
3:   while round <  $maxR$  do ▷ Vote while there are rounds left
4:      $\mathcal{N} \leftarrow \text{SHUFFLE}(\mathcal{N})$  ▷ Randomize order of player participation
5:     for  $i \in \mathcal{N}$  do
6:        $\mathcal{P} \leftarrow \text{VOTE}(i, \mathcal{P}, s, sp)$  ▷ Player i votes for posts
7:     end for
8:     round  $\leftarrow$  round + 1
9:   end while
10:  return  $\mathcal{P}$ 
11: end function

```

---

---

**Algorithm 4** Main protocol of curation of posts in Steem. It represents one week of voting and each player creates only one post. At the end of the week, the list of ordered posts by value of votes received will be returned.

---

```

1: function PROTOCOL( $N, \mathcal{S}, \mathcal{L}, \mathcal{SP}, \max R$ )
2:    $\mathcal{P} \leftarrow \text{POSTGENERATION}(N, \mathcal{L})$ 
3:    $\mathcal{P} \leftarrow \text{CURATION}(N, \mathcal{S}, \mathcal{P}, \mathcal{SP}, \max R)$ 
4:                                      $\triangleright$  Discuss payout part
5: end function

```

---

## 4 Results

Steem won't achieve high quality posts.

## 5 Further Work

Posts at any time

## 6 Conclusion

Keep inventing new decentralized content curation platforms.

## 7 Acknowledgements

We thank Prof. Aggelos Kiayias for constructive conversations, @serious-poster for their invaluable posts analyzing Steem and our mums for the cookies.