

# Analysis and Attacks of decentralized content curation platforms

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**Abstract.** We will attack Steem.

## 1 Introduction

Steem is not incentive-compatible.

## 2 Related Work

Many people have done many similar things.

## 3 Model

### 1 Notation

- We denote the set of all probability distributions on set  $A$  as  $\mathcal{D}(A)$ .
- We denote the powerset of a set  $A$  with  $2^A$ .
- $a||b$  denotes the concatenation of  $a$  and  $b$ .

### 2 Properties of Post Voting Systems

A post voting system has the objective to arrange the posts according to the preferences of the participants. The ideal order is defined based on the likeability matrix for the posts.

**Definition 1 (Post).** Let  $N \in \mathbb{N}^*$ ,  $L_i \in \mathcal{D}([0, 1]^N)$ . A post is defined as  $p = (i, l(v))$ , with  $i \in [N]$ ,  $l \sim L_i$  (and  $v \in \mathbb{R}_+$ ).

- **Author.** The first element of a post is the index of its creator,  $i$ .
- **Likeability.** The likeability of a post is defined as  $l \in [0, 1]^N$  (where  $l$  is drawn from  $L_i$  the Likeability Distribution of its creator  $u_i$ ).

- (**Votes.** A post has an associated “vote” value, which is a real non-negative number. It is initialized at 0 and increases whenever a player votes for the post, as explained later in detail.)

Let  $P \in \mathbb{N}^*$  the number of posts. Then  $\forall j \in [P]$ , let  $\text{creator}_j \in [N]$ ,  $l_j \sim L_{\text{creator}_j}$  and  $p_j = (\text{creator}_j, l_j, 0)$ . The set of all posts is  $\mathcal{P} = \bigcup_{j=1}^P \{p_j\}$ .<sup>1,2</sup>

**Definition 2 (Post score).** Let post  $p = (m, l)$ . We define the score of  $p$  as  $\text{sc}(p) = \sum_{i=1}^N l_i$ .

The score of a post is a single number that represents its overall worth to the community. By using simple summation, we assume that the opinions of all players have the same weight. In an ordered list of posts where higher posts are more visible, the “common interest” would require that a post with higher score appear before another post with a lower score.

**Definition 3 (Ideal Post Order).** Let  $\mathcal{P}$  a set of posts. We define  $\text{IDEALORDER}(\mathcal{P})$  as a list of the posts in  $\mathcal{P}$  such that

$$\forall i < j \in |\mathcal{P}|, \text{sc}(\text{IDEALORDER}(\mathcal{P})[i]) \geq \text{sc}(\text{IDEALORDER}(\mathcal{P})[j]) \quad .$$

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**Definition 4 (Player).** Let  $N \in \mathbb{N}^*$ . A player  $u$  is defined by her unique  $\text{id} \in [N]$ , her *Likeability Distribution*  $L \in \mathcal{D}([0, 1]^N)$  and her *Strategy*  $S \in \{H, G\} \times \mathbb{N}^* \times 2^{[N]}$ . Let the  $i$ th player be represented by the tuple  $u_i = (i, L_i, S_i)$ . The tuple of players is defined as  $\mathcal{U} = (u_1, \dots, u_i, \dots, u_N)$ .

We will now explain each field in  $u_i$  in detail:

- **Likeability Distribution.** The *Likeability Distribution*  $L_i \in \mathcal{D}([0, 1]^N)$  of  $u_i$  is a distribution on how likeable is the content produced by  $u_i$  to the rest of the players. The *Likeability Distribution* for the whole system is  $\mathcal{L} = (L_1, \dots, L_i, \dots, L_n)$ .
- **Strategy.** The strategy of  $u_i$  is defined as  $S_i \in \{H, G\} \times \mathbb{N}^* \times 2^{[N]}$ , where the first element is the player’s core strategy, the second is her attention span and the third is her voting ring.

<sup>1</sup> TODO: consider removing things in parentheses

<sup>2</sup> TODO: consider dropping  $L_i$  entirely

<sup>3</sup> TODO: discuss list notation

- **Honest/Greedy.**  $H$  corresponds to the honest and  $G$  to the greedy strategy. An honest player  $u_i$  would vote a post  $p$  according to its likeability towards herself,  $l_{p,i}$ , that is to say she votes posts following her preferences. (For honest players, the value of the vote is computed as  $v_{H,i} = a \cdot \mathcal{VP} \cdot l \cdot \mathcal{SP} + b$ , where  $l$  is drawn from the Likeability distribution)<sup>4,5</sup>. In Steem terms,  $l$  can be understood as the weight of a vote.  
A greedy player only votes for posts produced by users of its Voting Ring. (The value of vote for a player if  $u_i$  is greedy is defined as  $v_{G,i} = a \cdot \mathcal{VP}_i \cdot \mathcal{SP}_i$ ,<sup>6</sup> as in our model all greedy votes are executed with full weight.)
- **Attention Span.** This is a positive integer that represents the number of posts a player can consider voting simultaneously. For the benefit of simplicity, we will assume that this number is constant throughout all players.<sup>7</sup>
- **Voting Ring.** If player  $u_i$  is honest, her Voting Ring is  $R_i = \emptyset$ . If  $u_i$  is greedy, her Voting Ring is  $R_i \in 2^{[N]}$ . A voting ring is defined as  $R_i = \{g_1, \dots, g_j, \dots, g_n\}$  where  $g_j \in \mathcal{U}$  is the  $j$ th member of the voting ring and  $n$  is the size of the ring. Two greedy players will either have the same or disjoint voting rings ( $\forall i \neq j \in [N], R_i = R_j \vee (R_i \cap R_j = \emptyset)$ ).

The tuple of all strategies is defined as  $\mathcal{S} = (S_1, \dots, S_i, \dots, S_N)$ .

The set of players is defined as  $\mathcal{U} = (u_1, \dots, u_i, \dots, u_N)$ .

**Definition 5 (Post-Voting System).** An algorithm  $F$  that accepts as inputs:

- list of posts (in some order)
- set of players

All posts must have a likeability vector of length equal to the number of players.

The system returns an ordered list of the input posts, possibly re-ordered.

*Discussion* If the system's description ends here, it is trivial to make a system which always returns the ideal order. Also, the contribution of the

<sup>4</sup> TODO: move to posts section

<sup>5</sup> TODO: explain  $a, b$

<sup>6</sup> TODO: same

<sup>7</sup> TODO: discuss

players is completely unspecified, thus we cannot define what it means to have honest behaviour.

I suggest that we turn to a simulation-based approach. The Environment gives players their likeabilities and puts posts in a global "Feed" functionality. Players can interact with this functionality in very specific ways (e.g. vote up to  $m$  times for each post, pay some money/reduce some value (like voting power) with each vote) and the global functionality re-orders the posts according to the votes from time to time (when?). Then we can talk about honest behaviour (vote according to a specific protocol, e.g. view first  $k$  posts and choose which to vote) and deviation (being malicious) and maybe prove incentive-compatibility. Also we can easily envision a "Vote" functionality that replaces the players and makes the necessary votes so that the final list of posts in the "Feed" functionality matches the ideal.

**Definition 6 (Honest Player v1).** *A player  $u$  is considered honest if her utility is maximized when the result of the post-voting system equals her subjective ordering of the posts:*

$$\forall i < j \in |\mathcal{P}|, l_{u, \text{u-ORDER}(\mathcal{P})[i]} \geq l_{u, \text{u-ORDER}(\mathcal{P})[j]} \quad .$$

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**Definition 7 (Honest Player v2).** *A player  $u$  is considered honest if she always votes the best post of the top  $k$  (her attention span) with a weight equal to how much she likes the post.<sup>9</sup>*

**Definition 8 (Convergence under honesty).** *We say that a post-voting system  $F$   $t$ -converges under honesty if, for every valid input  $(\mathcal{P}, \mathcal{U})$  where all players are honest, it is*

$$F(\mathcal{P}, \mathcal{U})_{1..t} = \text{IDEALORDER}(\mathcal{P})_{1..t} \quad .$$

**Theorem 1.** *The Steem system  $t$ -converges under honesty, assuming the following conditions on  $N, P, \mathcal{SP}, v, \text{att\_span}, a, b, c \dots$ <sup>10</sup>*

The above result is tight. If the conditions are violated the above theorem is not true.

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**Algorithm 1** Each player creates one post

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```
1: function GENERATEPOSTS( $\mathcal{U}$ )
2:    $\mathcal{P} = \emptyset$  ▷ List of posts
3:   for  $u_i \in \mathcal{U}$  do
4:      $l \xleftarrow{r} L_i$  ▷ Get likeability of posts
5:      $\mathcal{P} \leftarrow \mathcal{P} \parallel (i, l, 0)$  ▷ Add post to list of Posts
6:   end for
7:    $\mathcal{P} \leftarrow \text{SHUFFLE}(\mathcal{P})$  ▷ Shuffle posts to a random order
8:   return  $\mathcal{P}$ 
9: end function
```

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**Algorithm 2** Player votes for best of  $k$  posts

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```
1: function VOTE( $u_i, \mathcal{P}$ )
2:   switch  $S_i$  do
3:     case honest
4:        $p_j \leftarrow \underset{p \in \mathcal{P}_{1..k}}{\text{argmax}} \{l_{i,p}\}$ 
5:       Parse  $p_j$  as  $(m, l_p, v)$ 
6:        $v' \leftarrow v + a \cdot \text{VP}_i \cdot l_{i,p} \cdot \text{SP}_i + b$ 
7:        $\text{VP}_i \leftarrow \text{VP}_i - (a \cdot \text{VP}_i \cdot l_{i,p} + b)$ 
8:        $\mathcal{P} \leftarrow p_1 \parallel p_2 \parallel \dots \parallel p_{j-1} \parallel (m, l_p, v') \parallel p_{j+1} \parallel \dots \parallel p_N$ 
9:     end case
10:    case Greedy
11:      ▷ If post belongs to voting ring and not reached min VPower
12:      if  $p \in s.R \wedge p.VPower > s.Min$  then
13:         $voteValue \leftarrow p.VPower \cdot weight \cdot sp$ 
14:         $p \leftarrow p.votes + voteValue$ 
15:      end if
16:    end case
17:  end switch
18:  return  $\mathcal{P}$ 
19: end function
```

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### 3 Steem Execution

TODOS<sup>11,12</sup>

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**Algorithm 3** Players cast votes for  $r$  rounds

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```

1: function CURATE( $\mathcal{U}, \mathcal{P}, r$ )
2:   for  $j = 1$  to  $r$  do                                      $\triangleright r$  voting rounds
3:     for  $u_i \in \mathcal{U}$  do
4:       if ISVOTEROUND( $j, S_i, r, N$ ) then
5:          $\mathcal{P} \leftarrow \text{VOTE}(u_i, \mathcal{P})$                         $\triangleright$  Player  $i$  votes zero or one posts
6:       end if
7:     end for
8:      $\mathcal{P} \leftarrow \text{ORDER}(\mathcal{P})$                                 $\triangleright$  Order posts by vote count after each round of votes
9:   end for
10:  return  $\mathcal{P}$ 
11: end function

```

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**Algorithm 4** Calculates whether voting in this round is optimal

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```

1: function ISVOTEROUND( $j, S, r, N$ )
2:
3: end function

```

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**Algorithm 5** Steem( $\mathcal{U}, \mathcal{P}, r$ )

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```

1:  $\mathcal{P} \leftarrow \text{CURATE}(\mathcal{U}, \mathcal{P}, r)$ 
2: return  $\mathcal{P}$ 

```

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### 4 Results

Steem won't achieve high quality posts.

### 5 Further Work

Posts at any time

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<sup>8</sup> TODO: discuss list notation

<sup>9</sup> TODO: discuss definition of honesty

<sup>10</sup> TODO: continue

<sup>11</sup> TODO: Improve argmax notation

<sup>12</sup> TODO: Write Greedy

## **6 Conclusion**

Keep inventing new decentralized content curation platforms.

## **7 Acknowledgements**

We thank @seriousposter for their invaluable posts analyzing Steem and our mums for the cookies.

## **References**