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## Assignment 8

## Exercise 1.

Describe in plain English (a short paragraph with at most 5-6 lines should be enough) an algorithm for the following task:

Input: A directed graph G = (V, E), and a node  $u \in V$ .

Goal: Output 1 if there is a path from every  $v \in G$  to u (so if u is reachable from any other node), and output 0 otherwise.

Your algorithm should have runtime O(n+m). (Hint: Use an idea that we have seen for similar connectivity problems in directed graphs.)

## Exercise 2.

We have seen that Dijkstra's algorithm can be implemented in two ways: Variant (a) uses an array to store the dist[] values of the unknown nodes, and Variant (b) uses a MIN-HEAP to store these values.

- (a) Suppose in your application  $m \leq 3n$ . Which variant gives a faster runtime? Justify your answer.
- (b) Suppose in your application  $m \ge n^2/3$ . Which variant gives a faster runtime? Justify your answer.

Our algorithm starts off by initializing two integer arrays dist and npath.

```
int[] dist = new int[g.n];
int[] npath = new int[g.n];

for (int i = 0; i < g.n; i++) {
    dist[i] = Integer.MAX_VALUE;
    npath[i] = 0;
}</pre>
```

I've set the undiscovered nodes to have a distance of infinity, and in this case, the maximum integer value that Java supports.

The algorithm in its entirety.

```
static void getShortestPath(Adj_List_Graph g, int s) {
  int[] dist = new int[g.n];
  int[] npath = new int[g.n];
  for (int i = 0; i < g.n; i++) {</pre>
      dist[i] = Integer.MAX_VALUE;
      npath[i] = 0;
  }
  Queue < Integer > queue = new LinkedList < Integer > ();
  dist[s] = 0;
 npath[s] = 1;
 queue.add(s);
  while (!queue.isEmpty()) {
      int p = queue.poll();
      List<Integer> neighbors = g.adj.get(p);
      for (int v : neighbors) {
          if (dist[v] == Integer.MAX_VALUE) {
              dist[v] = dist[p] + 1;
              npath[v] = npath[p];
              queue.add(v);
          }
          else if (dist[v] == dist[p] + 1)
              npath[v] += npath[p];
      }
 }
 for (int index = 1; index < g.n; index++)</pre>
      System.out.printf("dist[%d] = %d \t npath[%d] = %d%n",
              index, dist[index], index, npath[index]);
}
```

Program output can be found on the next page.

```
G1 results:
dist[1] = 0
               npath[1] = 1
dist[2] = 1
               npath[2] = 1
dist[3] = 1
               npath[3] = 1
dist[4] = 1
               npath[4] = 1
dist[5] = 2
               npath[5] = 2
dist[6] = 2
               npath[6] = 1
dist[7] = 3
               npath[7] = 3
G2 results:
dist[1] = 0
               npath[1] = 1
dist[2] = 1
               npath[2] = 1
dist[3] = 1
               npath[3] = 1
dist[4] = 1
               npath[4] = 1
dist[5] = 1
               npath[5] = 1
dist[6] = 1
               npath[6] = 1
dist[7] = 2
               npath[7] = 5
dist[8] = 3
               npath[8] = 5
dist[9] = 3
               npath[9] = 5
```

To visually reflect one of our solutions, namely dist[7] and npath[7] of G2, we'll highlight the paths.

