计算机系统体系结构 Project5

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实验环境

Windows 10 下使用 VMWare Workstation 15 Player 创建和运行虚拟机,虚拟机环境是 Linux 发行版 Ubuntu16.04.6 LTS。

1 线程池

这部分要实现一个线程池,提供给用户几个 API: 1) void pool_init(),2) int pool_submit(void (*somefunction)(void *p), void *p), 3) void pool_shutdown(void)。按照书上的描述,实现各自功能,并使用 semaphore 和 mutex 来实现同步和避免 race condition。

1.1 设计思路

- 1. **client.c.** 需要对原始代码进行一些修改,原本只调用了一次 pool_submit(),这里需要添加多个任务到线程池中来展示其功能。
- 2. threadpool.c. 这里按照用户调用的顺序说明其需要实现的功能。
 - (a) 需要一个队列来保存提交的任务,这里使用数组实现。
 - (b) 首先用户初始化线程池,需要初始化信号量(这里使用了命名的信号量)、互斥锁以及创建特定数量个线程。创建的线程传入 worker 函数,worker 在线程的生命周期中一直循环,每次循环等待表示有任务可执行的信号量,然后执行一次任务。这里取出任务的 dequeue 操作需要加锁,防止多个线程同时更改。
 - (c) 接下来用户多次提交任务,在添加任务到队列的 enqueue 操作中,要对队列加锁,防止多个线程同时对其更改;若提交成功了(队列有位置可以添加)则调用 sem_post 以通知等待的线程有任务可取。然后 worker 会立即取出任务并执行。
 - (d) 最后用户要关闭线程池,对每个线程调用 pthread cancel。

1.2 核心代码解释

1. 声明的全局变量

```
// the work queue
task worktodo[QUEUE_SIZE];

// the worker bee
pthread_t bee[NUMBER_OF_THREADS];

int numTask = 0;
pthread_mutex_t mutex;
sem_t *sem;

// the work queue
task worktodo[QUEUE_SIZE];

// the worker bee
pthread_t bee[NUMBER_OF_THREADS];

// int numTask = 0;
// sem_t *sem;

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//
```

2. 初始化线程池

```
// initialize the thread pool
void pool_init(void)

{
    sem = sem_open("SEM", O_CREAT, 0666, 0);
    pthread_mutex_init(&mutex, NULL);
    for (int i=0; i<NUMBER_OF_THREADS; ++i){
        pthread_create(&bee[i],NULL,worker,NULL);
    }
}
</pre>
```

3. 入队列操作 enqueue

```
// insert a task into the queue
   // returns 0 if successful or 1 otherwise,
3 int enqueue(task t)
   {
4
       pthread_mutex_lock(&mutex);
       if(numTask >= QUEUE SIZE)
           pthread_mutex_unlock(&mutex);
           return 1;
       worktodo[numTask++]=t;
10
       pthread mutex unlock(&mutex);
11
       return 0;
12
13
```

4. 出队列操作 dequeue

```
1  // remove a task from the queue
2  task *dequeue()
3  {
4     pthread_mutex_lock(&mutex);
5     if (numTask<=0){
6       pthread_mutex_unlock(&mutex);
7       return NULL;
8     }
9     task *next_work = &worktodo[--numTask];
10     pthread_mutex_unlock(&mutex);
11     return next_work;
12  }</pre>
```

5. 工作线程

```
// the worker thread in the thread pool
void *worker(void *param)

// execute the task
while(TRUE){
sem_wait(sem);
```

```
task *newTask = dequeue();

if (newTask == NULL) continue;

execute(newTask->function, newTask->data);

pthread_exit(0);

pthread_exit(0);
```

6. 提交任务

```
int pool_submit(void (*somefunction)(void *p), void *p)

task newTask;
newTask.function = somefunction;
newTask.data = p;
int res = enqueue(newTask);
if (!res) sem_post(sem);
return res;
}
```

7. 关闭线程池

```
' // shutdown the thread pool
void pool_shutdown(void)

for(int i=0; i < NUMBER_OF_THREADS; i++){

pthread_cancel(bee[i]);

// pthread_join(bee[i], NULL);

}

</pre>
```

2 生产者-消费者问题

这里要使用两个信号量 empty 和 full 以及一个互斥锁 mutex 来解决生产者消费者问题。

2.1 设计思路

主函数的内容: 首先获得命令行输入的 sleep time, 生产者数量,消费者数量; 然后初始化 buffer; 然后分别创建若干个生产者线程和消费者线程,并让他们运行生产者/消费者的操作; 然后主进程 sleep输入的时间后终止所有线程。

对于生产者线程,将一直循环等待 buffer 有无空位,即 sem_wait(&empty); 若有空,则加锁,然后向 buffer 插入一条 item; 随后解锁,并且对 full 信号量 sem_post(&full)。对于消费者线程,则相反,一直等待 buffer 有无 item,即 sem_wait(&full); 然后取出一条 item,随后解锁、sem_post(&empty);。

另外还要实现 buffer 的 insert 和 remove 操作,这里用了数组实现的循环队列。

2.2 核心代码解释

1. 一些变量的声明。

```
buffer_item buffer[BUFFER_SIZE]; //manipulated as a circular queue
int rear = 0;
int front = 0;
pthread_mutex_t mutex;
sem_t empty;
sem_t full;
```

2. 首先是 buffer 的插入和删除。

```
int insert item(buffer item item) {
        /* insert item into buffer
2
       return 0 if successful, otherwise
       return -1 indicating an error condition */
        if ((rear+1)%BUFFER_SIZE==front){
           return -1;
       }
       rear = (rear+1)\%BUFFER\_SIZE;
       buffer[rear] = item;
       return 0;
10
   }
11
12
    int remove_item(buffer_item *item) {
13
       /* remove an object from buffer
14
       placing it in item
15
       return 0 if successful, otherwise
16
       return -1 indicating an error condition */
17
        if(front = rear)
18
           return -1;
19
20
       front = (front+1)\%BUFFER\_SIZE;
21
       item = \&buffer[front];
22
       return 0;
23
24
```

3. 然后是生产者的工作函数。消费者类似,这里就不展示了。

```
void *producer(void *param) {
       buffer item item;
2
       while (1) {
            /* sleep for a random period of time */
           sleep (rand()\%3);
           /* generate a random number */
           item = rand();
           sem_wait(&empty);
           pthread_mutex_lock(&mutex);
           if (insert_item(item))
10
               printf("insert_error_condition\n");
12
               printf("producer_produced_m%d\n",item);
13
           pthread_mutex_unlock(&mutex);
14
           sem_post(&full);
15
```

```
16 }
17 }
```

4. 最后是主函数。注意 empty 初始值应为 BUFFER_SIZE-1, 因为循环队列需要区别空和满两种情况的 front 和 rear 关系,有一个空存储单元不能使用。

```
int main(int argc, char *argv[]) {
       1. Get command line arguments argv[1],argv[2],argv[3] */
       int sleep_time, producer_threads, consumer_threads;
        if (argc != 4)
       {
            fprintf(stderr, "Input\_form: < sleep\_time > \_ < producer\_threads\_number > \_
                <consumer_threads_number>\n");
           return -1;
       sleep\_time = atoi(argv[1]);
       producer threads = atoi(argv[2]);
10
       consumer\_threads = atoi(argv[3]);
11
12
      2. Initialize buffer */
13
       pthread_mutex_init(&mutex, NULL);
14
       sem_init(&full, 0, 0);
15
       sem_init(&empty, 0, BUFFER_SIZE-1);
16
    /* 3. Create producer thread(s) */
18
       pthread_t producer_ids[producer_threads];
19
       for(int i=0;iproducer_threads;++i){
20
           pthread create(&producer ids[i], NULL, &producer, NULL);
21
22
23
    /* 4. Create consumer thread(s) */
24
       pthread t consumer ids[consumer threads];
25
        for(int i=0;i<consumer_threads;++i){</pre>
26
           pthread create(&consumer ids[i], NULL, &consumer, NULL);
27
29
    /* 5. Sleep */
30
       sleep(sleep_time);
31
    /* 6. Exit */
33
       return 0;
34
35
36
   }
```

3 实验结果

```
xcwang@ubuntu:~/Documents/project5/ThreadPool$ ./example
I add two values 0 and 1 result = 1
I add two values 2 and 3 result = 5
I add two values 1 and 2 result = 3
I add two values 4 and 5 result = 9
I add two values 3 and 4 result = 7
```

Figure 1: ThreadPool

```
xcwang@ubuntu:-/Documents/project5/producer_consumer$ ./pcp 4 3 3 producer produced 719883386 producer produced 596516649 producer produced 1025202362 producer produced 1025202362 producer produced 1025202362 producer produced 1025202362 producer produced 1365180540 producer produced 336989172 producer produced 36989172 producer produced 35085211 producer produced 521595368 producer produced 521595368 producer produced 278722862 producer produced 278722862 producer produced 468703135 producer produced 468703135 producer produced 1315634022 producer produced 1315634022 producer produced 1059961393 producer produced 1059961393 producer produced 1056478042 producer produced 1656478042 producer produced 1656478042 producer produced 1656478042
```

Figure 2: Producer consumer