

# 计算机系统体系结构 Project8

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## 实验环境

Windows 10 下使用 VMWare Workstation 15 Player 创建和运行虚拟机，虚拟机环境是 Linux 发行版 Ubuntu16.04.6 LTS。

## 1 虚拟内存管理

编写一个程序，将大小为  $2^{16} = 65536$  字节的虚拟地址空间映射物理地址空间。读入的文件包含一些 32 位的地址，但我们只用低 16 位。要使用到 TLB 和页表，这里 TLB 有 16 个 entry，页表大小为 256 entry\*256 byte，内存大小相同。

### 1.1 设计思路

1. **初始化页表和 TLB**：将页表和 TLB 的每个条目赋值-1，表示没有内容。
2. **从输入数字中提取页号和偏移**：页号：将读入的数字右移页表大小的位数，这里是 8 位，即 offset 的位数；再与 255 按位与，就获得了 8 位的页号。偏移：读入的数字直接用 255 来按位与，就获得了偏移。
3. **查询步骤**：获得页号与偏移后，首先根据页号在 TLB 中查找，若命中，则直接获得页框号，从而有物理地址在内存中查找；否则，根据页号访问页表相应页，若未发生页缺失，则同样获得页框号，得到物理地址；若发生了页缺失，就要从 backing store 文件中取得相应页并存入到内存的某个页框中；最后根据 TLB 命中情况以及页缺失情况更新页表及 TLB。
4. **当物理地址空间小于虚拟地址空间**：此时查询的虚拟地址不能一一对应到物理地址，所以采用 FIFO 的页替换算法。当要更新页表而物理内存已满时，新页将把内存中最早进入的页框覆盖掉，同时页表和 TLB 中映射到被替换掉的页框的条目将被清除。这样的 FIFO 物理内存很好实现，维护一个地址，指向内存的下一处存储位置，每次内存被存入后，该地址  $addr := (addr + FRAME\_SIZE) \% MEM\_SIZE$ ，即可循环使用内存空间。

### 1.2 核心代码解释

1. 一些常量和全局变量

```
1 #define PAGE_SIZE 256
2 #define PAGE_ENTRIES 256
3 #define FRAME_SIZE 256
4 #define FRAME_ENTRIES 256
5 #define MEM_SIZE (FRAME_SIZE * FRAME_ENTRIES)
6 #define VIRTUAL_SIZE (PAGE_SIZE * PAGE_ENTRIES)
7 #define TLB_ENTRIES 16
8 #define OFFSET_BITS 8
9 #define OFFSET_MASK 255
```

```

10
11 int page_table[PAGE_ENTRIES];
12 int tlb[TLB_ENTRIES][2];
13 char memory[MEM_SIZE];
14 int mem_ptr = 0;
15 int faultNum = 0;
16 int tlbHit = 0;
17 int tlb_ptr = 0;

```

## 2. 初始化

```

1  for (int i=0; i<PAGE_ENTRIES; i++){
2      page_table[i] = -1;
3  }
4
5  for (int i=0; i < TLB_ENTRIES; i++){
6      tlb[i][0] = -1;
7      tlb[i][1] = -1;
8  }
9
10 int addrNum = 0;
11 int physical;
12 int value;
13 char buf[10];
14
15 const char *store_file = "BACKING_STORE.bin";
16 const char *input_file = argv[1];
17 const char *out_file = "output.txt";
18
19 int store_id = open(store_file, O_RDONLY);
20 FILE *input_fp = fopen(input_file, "r");
21 FILE *output_fp = fopen(out_file, "a");
22 char *store_ptr = mmap(0, VIRTUAL_SIZE, PROT_READ, MAP_SHARED,
    store_id, 0);

```

## 3. 查询步骤

```

1  while(fgets(buf, 10, input_fp)!=NULL){
2      addrNum++;
3      int logical_addr = atoi(buf);
4      int offset = get_offset(logical_addr);
5      int page_number = get_page_number(logical_addr);
6
7      int frame_number = search_tlb(page_number);
8      if(frame_number!=-1){ // tlb hit
9          physical = frame_number + offset;
10         value = memory[physical];
11     }
12     else{
13         frame_number = search_page_table(page_number);
14         if(frame_number!=-1){

```

```

15         physical = frame_number+offset;
16         update_tlb(page_number, frame_number);
17         value = memory[physical];
18     }
19     else{ // page fault
20         int page_address = page_number * PAGE_SIZE;
21         memcpy(memory + mem_ptr, store_ptr + page_address, PAGE_SIZE);
22         frame_number = mem_ptr;
23         for(int i=0;i<PAGE_ENTRIES;i++){
24             if(page_table[i]==frame_number){
25                 page_table[i] = -1;
26             }
27         }
28         for (int i=0; i < TLB_ENTRIES; i++){
29             if (tlb[i][1] == frame_number){
30                 tlb[i][0] = -1;
31                 tlb[i][1] = -1;
32             }
33         }
34         physical = frame_number + offset;
35         value = memory[physical];
36         page_table[page_number] = frame_number;
37         update_tlb(page_number, frame_number);
38         mem_ptr = (mem_ptr + FRAME_SIZE) % MEM_SIZE;
39     }
40 }
41
42 fprintf(output_fp, "Virtual_address:_%d_", logical_addr);
43 fprintf(output_fp, "Physical_address:_%d_", physical);
44 fprintf(output_fp, "Value:_%d\n", value);
45
46 }

```

#### 4. 获得页号

```

1 int get_page_number(int virtual) {
2     return (virtual >> OFFSET_BITS) & (PAGE_ENTRIES-1));
3 }

```

#### 5. 获得偏移

```

1 int get_offset(int virtual) {
2     return virtual & OFFSET_MASK;
3 }

```

#### 6. 查询 TLB

```

1 int search_tlb(int page_number) {
2     for (int i = 0; i < TLB_ENTRIES; i++) {
3         if (tlb[i][0] == page_number) {
4             tlbHit++;

```

```

5         return tlb[i][1];
6     }
7 }
8 return -1;
9 }

```

## 7. 查询页表

```

1 int search_page_table(int page_number) {
2     if (page_table[page_number] == -1) {
3         faultNum++;
4         return -1;
5     }
6     return page_table[page_number];
7 }

```

## 8. 更新 TLB

```

1 void update_tlb(int page_number, int frame_number) {
2     tlb[tlb_ptr][0] = page_number;
3     tlb[tlb_ptr][1] = frame_number;
4     tlb_ptr = (tlb_ptr + 1) % TLB_ENTRIES;
5 }

```

## 9. 实现 FIFO 页替换

```

1 memcpy(memory + mem_ptr, store_ptr + page_address, PAGE_SIZE);
2 mem_ptr = (mem_ptr + FRAME_SIZE) % MEM_SIZE;

```

## 1.3 实验结果

当物理地址空间和虚拟地址空间同样大时，由于最初内存里没有存任何页，所以会有页缺失；当所有被查询的页缺失一次后，以后将不再缺失；由于查询是随机的，所以页缺失数量略小于总的页数量；也是因为随机访问，故 TLB 命中数不高。

```

Virtual address: 48065 Physical address: 25793 Value: 0
Virtual address: 6957 Physical address: 26413 Value: 0
Virtual address: 2301 Physical address: 35325 Value: 0
Virtual address: 7736 Physical address: 57912 Value: 0
Virtual address: 31260 Physical address: 23324 Value: 0
Virtual address: 17071 Physical address: 175 Value: -85
Virtual address: 8940 Physical address: 46572 Value: 0
Virtual address: 9929 Physical address: 44745 Value: 0
Virtual address: 45563 Physical address: 46075 Value: 126
Virtual address: 12107 Physical address: 2635 Value: -46
Number of Translated Addresses = 1000
Page Faults = 244
Page Fault Rate = 0.244
TLB Hits = 54
TLB Hit Rate = 0.054

```

Figure 1: 物理地址空间为 256\*256

当物理地址空间只有逻辑地址空间一半时，由于简单的 FIFO 页替换算法导致很多的页缺失；每个虚拟地址对应的物理地址也不再与之前的相同，但值是相同的。

```
Virtual address: 48065 Physical address: 18113 Value: 0
Virtual address: 6957 Physical address: 27693 Value: 0
Virtual address: 2301 Physical address: 21245 Value: 0
Virtual address: 7736 Physical address: 13112 Value: 0
Virtual address: 31260 Physical address: 5148 Value: 0
Virtual address: 17071 Physical address: 5551 Value: -85
Virtual address: 8940 Physical address: 5868 Value: 0
Virtual address: 9929 Physical address: 6089 Value: 0
Virtual address: 45563 Physical address: 6395 Value: 126
Virtual address: 12107 Physical address: 6475 Value: -46
Number of Translated Addresses = 1000
Page Faults = 538
Page Fault Rate = 0.538
TLB Hits = 54
TLB Hit Rate = 0.054
```

---

Figure 2: 物理地址空间为 256\*128