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Pseudocode

```
Function groupScheduler(mySchedule, myDailyAct, otherSchedules, otherDailyActs, duration) {  
    // Combine all busy times from everyone's schedules  
    combinedBusyTimes = []  
  
    for each time slot in mySchedule:  
        add time slot to combinedBusyTimes  
  
    for each schedule in otherSchedules:  
        for each time slot in schedule:  
            add time slot to combinedBusyTimes  
  
    // Sort combinedBusyTimes  
    sort combinedBusyTimes by start time  
  
    // Merge overlapping busy times  
    mergedBusyTimes = []  
    add combinedBusyTimes[0] to mergedBusyTimes  
  
    for i from 1 to size of combinedBusyTimes:  
        // Check if mergedBusyTimes' end time is earlier than combinedBusyTimes' Start time  
        if combinedBusyTimes[i] overlaps with the last time in mergedBusyTimes:  
            // merge by replacing the mergedBusyTime's end time with the latest max end time.  
            merge them (update end time) using the max of the two end times.  
        else:  
            add combinedBusyTimes[i] to mergedBusyTimes  
  
    // Find the earliest and latest working hours for everyone  
    earliestStart = later start between myDailyAct and otherDailyActs[0]'s start time  
    latestEnd = earlier end between myDailyAct and otherDailyActs[0]'s end time  
  
    // Find free times between merged busy times  
    availableTimes = []  
    previousEnd = earliestStart  
  
    for each busy time in mergedBusyTimes:  
        if there is free time between previousEnd and this busy time:  
            if free time is long enough (greater than or equal to duration):  
                add that free time to availableTimes  
            update previousEnd to the end of this busy time  
  
    // Check if there is free time after the last busy period  
    if there is enough free time between the end of the last busy time and latestEnd:  
        add that free time to availableTimes  
  
    return availableTimes }
```

Proving the Efficiency of Pseudocode Using Step Count

Step 1: Combining all busy times

```
combinedBusyTimes = [] // 1

for each time slot in mySchedule: // n + 1
    add time slot to combinedBusyTimes // n

for each schedule in otherSchedules: // s + 1
    for each time slot in schedule: // t + 1
        add time slot to combinedBusyTimes // t
```

The nested loop iterates over differently sized vectors, therefore it's not going to have n^2 operations.

$$f(n) = 1 + ((n + 1) + n) + ((s + 1) + s(2t + 1))$$

$$f(n) = 1 + (2n + 1) + (2s + 2st + 1)$$

$$f(n) = (n + s + st) \Rightarrow \text{Worst Case Time Complexity is } O(n + s + st)$$

Step 2: Sorting CombinedBusyTimes

Calling `std::sort` for the *combinedBusyTimes* vector has a **worst case time complexity of $O(n \log n)$** .

Step 3: Merge overlapping busy times

```
mergedBusyTimes = [] // 1
add combinedBusyTimes[0] to mergedBusyTimes // 1

for i from 1 to size of combinedBusyTimes: // (n - 1) + 1
    if combinedBusyTimes[i] overlaps with the last time in mergedBusyTimes: // 1
        merge them (update end time) using the max of the two end times. // 1
    else:
        add combinedBusyTimes[i] to mergedBusyTimes // n
```

$$f(n) = (((n-1) + 1) + 1 + 1 + n) + 1 + 1 = 2n + 1 + 1 + 1$$

$$f(n) = 2n + 3 \Rightarrow \text{Worst Case Time Complexity is } O(n)$$

Step 4: Find the earliest and latest working hours for everyone

```
earliestStart = later start between myDailyAct and otherDailyActs[0]'s start time // 1
latestEnd = earlier end between myDailyAct and otherDailyActs[0]'s end time // 1
```

Calling `std::max` and `std::min` would have the **worst time complexity of $O(1)$** because we are just comparing two values.

$$f(n) = 1 + 1$$

$$f(n) = 2 \Rightarrow \text{Worst Case Time Complexity is } O(1)$$

Step 5: Find free times between merged busy times

```
availableTimes = [] // 1
previousEnd = earliestStart // 1

for each busy time in mergedBusyTimes: // n + 1
    if there is free time between previousEnd and this busy time: // 1
        if free time is long enough (greater than or equal to duration): // 1
            add that free time to availableTimes // n
        update previousEnd to the end of this busy time // n

// Check if there is free time after the last busy period
if there is enough free time between the end of the last busy time and latestEnd: // 1
    add that free time to availableTimes // n

return availableTimes // 1
```

$$f(n) = 1 + 1 + ((n + 1) + 1 + 1 + n + n) + 1 + n + 1$$

$$f(n) = 2 + (3n + 3) + 1 + n + 1$$

$$f(n) = 4n + 7 \quad \Rightarrow \quad \text{Worst Case Time Complexity is } O(n)$$

Overall Time Complexity of the groupScheduler Function

$$f(n) = O(n + s + st) + O(n \log n) + O(n) + O(1) + O(n)$$

Thus, the **overall worst case time complexity is $O(n \log n)$** for the groupScheduler function.