

CS 3530: Assignment 7b

Fall 2023

Problem 7.6 (10 points)

Problem

Show that P is closed under complement.

Hint: As we discussed in class, construct the language (e.g. $PCAT = \{\langle P_1, P_2, w \rangle \mid P_1, P_2 \in P \text{ and } w \text{ is a string, } w = x \cdot y, x \in P_1 \text{ and } y \in P_2\}$), then provide a deterministic machine that decides the language in polynomial time.

Solution to P closed under complement.

assume a language $P_1 \in P$ in which P is the class of languages which are decidable in polynomial time
the TM M which decides $\overline{P_1}$ is as follows

$M =$ "On input w :

1. check if $w \in P_1$
2. if $w \in P_1$, reject
3. if $w \notin P_1$, accept"

Since w must be finite this is decidable in polynomial time

Problem 7.7 (10 points)

Problem

Show that NP is closed under union and concatenation.

Hint: As we discussed in class, construct the language (e.g. $NPCAT = \{\langle NP_1, NP_2, w \rangle \mid NP_1, NP_2 \in NP \text{ and } w \text{ is a string, } w = x \cdot y, x \in NP_1 \text{ and } y \in NP_2\}$), then describe a certificate and a verifier that deterministically verifies the certificate in polynomial time.

Solution to NP closed under union.

$NPUN = \{\langle NP_1, NP_2, w \rangle \mid NP_1, NP_2 \in NP \text{ and } w \text{ is a string, where either } w \in NP_1 \text{ or } w \in NP_2\}$

let T_{NP_1} and T_{NP_2} be the machines that decide NP_1 and NP_2

we will construct a machine T_{NP_u} that decides the union of NP_1 and NP_2

$T_{NP_u} =$ "On input w :

1. run T_{NP_1} on w , if T_{NP_1} accepts, accept
2. run T_{NP_2} on w , if T_{NP_2} accepts, accept
3. else reject"

Since T_{NP_1} and T_{NP_2} are both polynomial the union of the two will also be polynomial

Solution to NP closed under concatenation.

$NPCAT = \{ \langle NP_1, NP_2, w \rangle \mid NP_1, NP_2 \in NP \text{ and } w \text{ is a string, } w = x \cdot y, x \in NP_1 \text{ and } y \in NP_2 \}$

let T_{NP_1} and T_{NP_2} be the machines that decide NP1 and NP2

we will construct a machine T_{NP_c} that decides the concatenation of NP1 and NP2

$T_{NP_c} =$ "On input w :

1. split w into w_1 and w_2 so that $w_1 w_2 = w$
2. run T_{NP_1} on w_1 , if T_{NP_1} rejects, reject
3. run T_{NP_2} on w_2 , if T_{NP_2} rejects, reject
4. else, accept