

CS 3530: Assignment 7a

Fall 2023

Problem 7.6 (10 points)

Problem

Show that P is closed under union and concatenation.

Hint: As we discussed in class, construct the language (e.g. $PCAT = \{\langle P_1, P_2, w \rangle \mid P_1, P_2 \in P \text{ and } w \text{ is a string, } w = x \cdot y, x \in P_1 \text{ and } y \in P_2\}$), then provide a deterministic machine that decides the language in polynomial time.

Solution to P closed under union.

$PUN = \{\langle P_1, P_2, w \rangle \mid P_1, P_2 \in P \text{ and } w \text{ is a string, where either } w \in P_1, \text{ or } w \in P_2\}$

M = "On input w :

1. Check if $w \in P_1$, if accept, then accept w
2. If not check if $w \in P_2$, if accept then accept w
3. If both reject, reject input w "

Since each check requires polynomial time the overall time is polynomial

Solution to P closed under concatenation.

$PCAT = \{\langle P_1, P_2, w \rangle \mid P_1, P_2 \in P \text{ and } w \text{ is a string, } w = x \cdot y, x \in P_1 \text{ and } y \in P_2\}$

M = "On input w on length n :

1. w can be split into two strings in n different ways
2. for each split w_1 , and w_2
 - a. check if $w_1 \in P_1$
 - b. check if $w_2 \in P_2$
3. If any split succeeds we accept, else we reject if all splits have been tried

Problem 7.5 (10 points)

Is the following formula satisfiable? (Give your reasoning.)

$$(x \vee y) \wedge (x \vee \bar{y}) \wedge (\bar{x} \vee y) \wedge (\bar{x} \vee \bar{y})$$

Solution

1. $x = T, y = F$

$$(T \vee F) \wedge (T \vee \overline{F}) \wedge (\overline{T} \vee F) \wedge (\overline{T} \vee \overline{F})$$

$$T \wedge (T \vee \overline{F}) \wedge (\overline{T} \vee F) \wedge (\overline{T} \vee \overline{T})$$

$$T \wedge T \wedge F \wedge T$$

$$T \wedge F$$

$$F$$

$$\textit{!Satisfiable}$$

$$2. \text{ x} = \text{F}, \text{ y} = \text{T}$$

$$(F \vee T) \wedge (F \vee \overline{T}) \wedge (\overline{F} \vee T) \wedge (\overline{F} \vee \overline{T})$$

$$T \wedge (F \vee \overline{T}) \wedge (\overline{F} \vee T) \wedge (\overline{F} \vee \overline{T})$$

$$T \wedge F \wedge T \wedge T$$

$$F \wedge T$$

$$F$$

$$\textit{!Satisfiable}$$