

CS-341 Lecture 8

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Dr. Vickery

Administrivia

- Exam March 9
 - Will cover material “to date”
 - Lectures, readings, and homeworks
 - Short answer
 - Bring a calculator
- Slides available in PDF format

Floating-Point

- Normalized Fractions
 - 1.xxx ...
 - This makes the leading 1 redundant, so it is not stored in the encoded value.
 - Doubles the precision!
- Special values: $\pm\infty$, NaN

Bandwidth

- The *rate* (speed) at which information can be transmitted through a *channel*.
 - *Bits per second*.
- HDTV
 - 30 frames per second
 - 800 x 600 pixels per frame
 - 256 red, green, and blue intensities per pixel
 - 43.2 MB per second
 - Broadcast bandwidth
- Computer monitor
 - 60 frames per second, 1024 x 768 pixels, 2²⁴ colors
 - 1.132 Gbps
 - Framebuffer bandwidth

Error Detection and Correction

- Section 2.2.4 of the textbook.
- Used when transmitting binary data over a “noisy channel.”
- Basic principle is to add redundant information (reducing effective bandwidth).
 - Transmitter adds the extra bits.
 - Receiver uses the extra bits to check the integrity of the data received.
- Detection: Receiver can recognize invalid codewords when they are received.
 - Signal an error or request transmitter to send data again.
- Correction: Enough redundant information is sent so the receiver can determine which bit(s) are in error.
 - Correct by inverting the error bits.

Parity: Single Bit Detection

- Each transmitted codeword consists of a packet of data bits and one additional bit whose value is set so that the code word will have an even number of 1's.
 - Parity bit is the exclusive OR of the data bits.
 - Receiver detects an error if the exclusive OR of the codeword bits is 1.
- Odd parity: Each code word has an odd number of 1's.
 - Parity bit is the exclusive NOR of the data bits.
 - Receiver detects an error if the exclusive OR of the codeword bits is 0.

Exclusive OR and Exclusive NOR

a	b	XOR	XNOR
0	0	0	1
0	1	1	0
1	0	1	0
1	1	0	1

a	b	c	XOR	XNOR
0	0	0	0	1
0	0	1	1	0
0	1	0	1	0
0	1	1	0	1
1	0	0	1	0
1	0	1	0	1
1	1	0	0	1
1	1	1	1	0

$$a \text{ xor } b \text{ xor } c = (a \text{ xor } b) \text{ xor } c$$

Hamming Codes

- Codewords have multiple parity bits.
- $P_0 P_1 P_2 D_3 P_4 D_5 D_6 D_7 P_8 D_9 D_{10} D_{11} D_{12} D_{13} D_{14} D_{15} \dots$
 - Parity bits are in positions that are numbered with powers of two.
 - Data bits fill in other places.
 - Each parity bit produces even parity for all codeword bits that have the parity bit's power of two in their subscripts.
 - P_0 is even parity across the entire codeword.
 - Omitted in the textbook, but needed for double error detection.
- Checking: Pattern of parity errors tells position in which error occurred. Correct it by inverting the value.
 - But if overall parity is correct when there are other errors, a double error was detected and cannot be corrected.