## **Inverse Power Method**

Let  $A \in \mathbb{R}^{n \times n}$  be a matrix with eigenvalues  $\lambda_i$  for  $j \in \{1, \dots, n\}$  with the property

$$|\lambda_n| \ge \cdots \ge |\lambda_{i+1}| > |\lambda_i| > |\lambda_{i-1}| \ge \cdots \ge |\lambda_1|.$$

Let  $\hat{\lambda} \approx \lambda_i$  be a good guess for the eigenvalue  $\lambda_i$ , which means

$$0 < |\lambda_i - \hat{\lambda}| < |\lambda_i - \hat{\lambda}| \quad \forall \ i \neq j. \tag{1}$$

Or in other words,  $\hat{\lambda}$  is closer to  $\lambda_i$  (in absolute terms) than any other eigenvalue. Let us define

$$B := (A - \hat{\lambda}I).$$

- 1. Show that  $(\lambda_i \hat{\lambda})^{-1}$  is (in magnitude) the largest eigenvalue of the matrix  $B^{-1}$ .
- 2. Let  $x^0 \in \mathbb{R}^n$ . With respect to A, the iteration

$$x^{k+1} := \frac{B^{-1}x^k}{\|B^{-1}x^k\|} \tag{2}$$

for  $k \ge 0$  is called *inverse power iteration*. Under which condition on  $x^0$  does this iteration converge and what is the limit?

Hint: Apply the theorem about the power method from the lecture.

## Solution:

The message here is: Under the assumption of having a "good" guess for  $\lambda_i$ , we find the "exact" eigenvector (and -value) with the help of the *inverse power iteration*. However note that it is computationally more expensive, because you have to solve a linear system in each iteration.

We have  $A \in \mathbb{R}^{n \times n}$  with eigenvalues  $|\lambda_n| \ge \cdots \ge |\lambda_{i+1}| > |\lambda_i| > |\lambda_{i-1}| \ge \cdots \ge |\lambda_1|$  and corresponding eigenvectors  $v_i \ne 0$ . Also we have  $\hat{\lambda} \approx \lambda_i$  so that

$$0 < |\lambda_i - \hat{\lambda}| < |\lambda_i - \hat{\lambda}| \quad \forall \ i \neq j \ (*).$$

Now define  $B := (A - \hat{\lambda}I)$ .

- 1. We show, that  $(\lambda_i \hat{\lambda})^{-1}$  is the largest eigenvalue of  $B^{-1}$ .
  - (i) For all j we have that  $(\lambda_i \hat{\lambda})$  is an eigenvalue of B to the same eigenvector  $v_i$  (previous sheets).
  - (ii) Also B is invertible, otherwise there would exist  $v \neq 0$ , so that Bv = 0, which would imply  $\hat{\lambda}$  eigenvalue of A, which would contradict the assumption:  $|\lambda_j \hat{\lambda}| > 0 \ \forall j \ \Rightarrow \ |\lambda_j \hat{\lambda}| \neq 0 \ \forall j$ . Thus  $\frac{1}{\lambda_j \hat{\lambda}}$  is an eigenvalue of  $B^{-1}$  for all j to the same eigenvector  $v_j$  (see previous sheets). By (\*) we then find that

$$\frac{1}{|\lambda_i - \hat{\lambda}|} > \frac{1}{|\lambda_j - \hat{\lambda}|} \ \forall j \ \ (\sharp).$$

2. Let  $v_1$  be an eigenvector to the eigenvalue  $(\lambda_i - \hat{\lambda})^{-1}$ , and  $x^0$  an orthogonal initial guess, i.e.,  $(x^0, v_1) \neq 0$ , then due to  $(\sharp)$  we find by the theorem about the power method that

$$\frac{B^{-1}x^k}{\|B^{-1}x^k\|} \to v_1.$$

Note:  $v_1$  is eigenvector of

- ... A to the eigenvalue  $\lambda_i$ .
- ... $B = A \hat{\lambda}_i I$  to the eigenvalue  $\lambda_i \hat{\lambda}$ .
- ... $B^{-1}=(A-\hat{\lambda}_iI)^{-1}$  to the eigenvalue  $(\lambda_i-\hat{\lambda})^{-1}.$