

Recursion [] In Arrays

1. Display Array

Input:

5
3
1
0
7
5

Output:

3
1
0
7
5

Constraints:

1. $l \leq n \leq 30$.

2. $0 \leq n_1, n_2, \dots$ elements.

Steps:

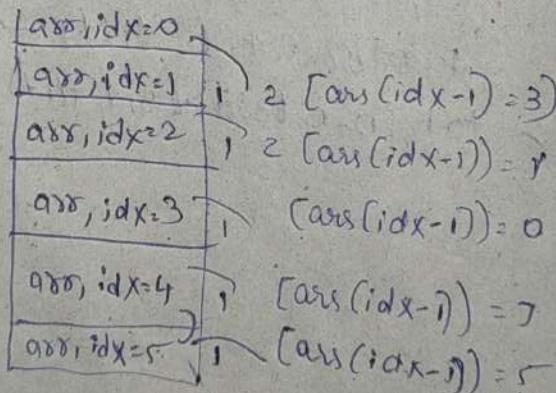
1. Take the inputs as required.

2. We will call the $(idx-1)$ of function. Assuming it has true till it reaches the base condition.

3. Then Print the Values.

Dry Run:

3	1	0	7	5
---	---	---	---	---



Program:

```
displayArr(int arr[], int idx) {  
    if (idx == 0)  
        return;
```

```
    displayArr(arr, idx - 1);
```

```
    pointInt(arr[idx - 1]);
```

}

— 1
— 2

2. Display Array In Reverse.

Input

5
3
1
0
7
5

Output

5
7
0
1
3

Constraints

$1 \leq N \leq 30$
 $0 \leq n_1, n_2 \leq \text{elements} \leq 10$

Steps:

1. Take the inputs as required.
2. We will print $\text{arr}(\text{idx}-1)$ and call the recursive method till the base condition $\text{idx} = 0$.

Program:

```
public static void displayArrReverse(int[] arr, int idx) {
    if (idx == 0)
        return;
    else
        pointIn(arr[idx - 1]); — ①
        displayArrReverse(arr, idx - 1); — ②
}
```

Dry Run:

arr, idx=0	Return
arr, idx=1	1 → print
arr, idx=2	2 → print
arr, idx=3	3 → print
arr, idx=4	4 → print
arr, idx=5	5 → print

3. Max of an Array

constraints

$$1 \leq n \leq 10^4$$

$$O(n) = O(1, 2, \dots, n \text{ elements})$$

$$C = 10^9$$

Input

6
15
30
40
4
11
9

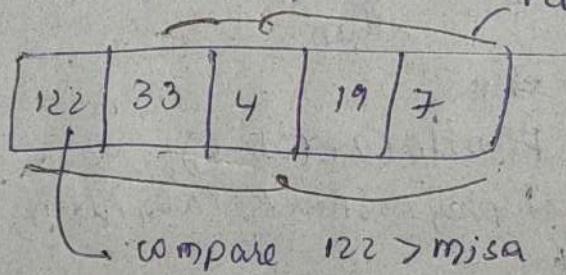
Output

40

Steps:

1. Pass arr, index=0 to the max function.
2. Expectation, $\max(\text{arr}, 0) \rightarrow$ gives the Highest Element.
3. Faith: $\max(\text{arr}, i) \rightarrow$ to end gives the highest element and then compare with idx, which returns the maximum element.

Faith: We know the max element = misa



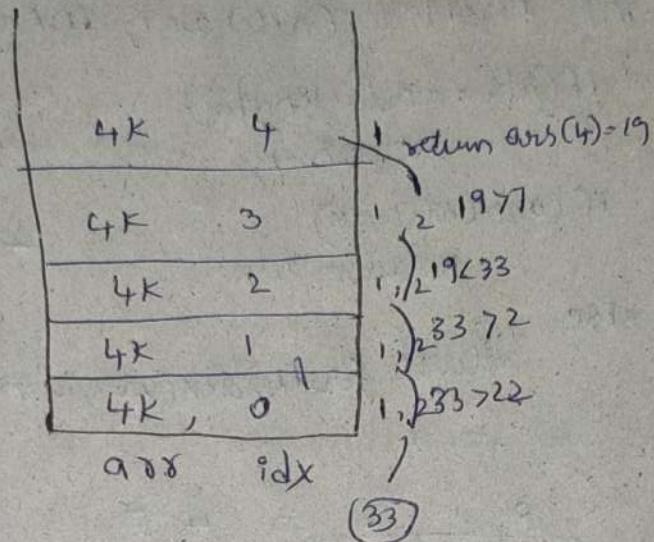
Program:

maxOfArray(a, 0);

```
public static int maxOfArray(int[] arr, int idx) {
    if (idx >= arr.length - 1)
        return arr[idx];
    int misa = maxOfArray(arr, idx + 1); — ①
    if (misa > arr[idx])
        return misa;
    else
        return arr[idx]; } — ②
```

Day Run:

0	1	2	3	4
22	2	33	7	19



4 First Index

Input

6 - n

15
11
40
4
4
9
4-X

elements

Output

3

Constraints

$1 \leq n \leq 10^4$

$0 \leq n_1, n_2, \dots, n_{\text{elements}} \leq 10^3$

$0 \leq x \leq 10^3$

Steps:

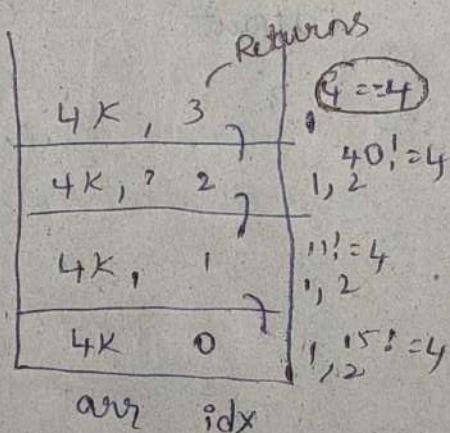
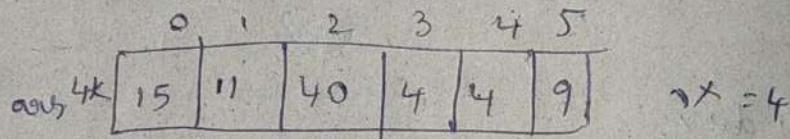
1. Pass arr, idx=0, x
2. If arr(idx)=x, then return idx
3. else Return, next element
4. If it reaches end of Array and x not found, return -1

Program:

```
firstIndex(a, 0, x);
```

```
public static int firstIndex(int[] arr, int idx, int x) {
    if (idx == arr.length)
        return -1;
    if (arr[idx] == x)
        return idx; -①
    else
        return firstIndex(arr, idx + 1, x); -②
}
```

DryRun:



5. Last Index:

Input

5 — n
 15
 11
 40
 4
 4
 9
 4 — x

elements

Output

4

Constraints $1 \leq n \leq 10^4$ $0 \leq n_1, n_2, \dots, n_k \leq n$ $k \leq 10^3$ $0 \leq x \leq 10^3$

Steps:

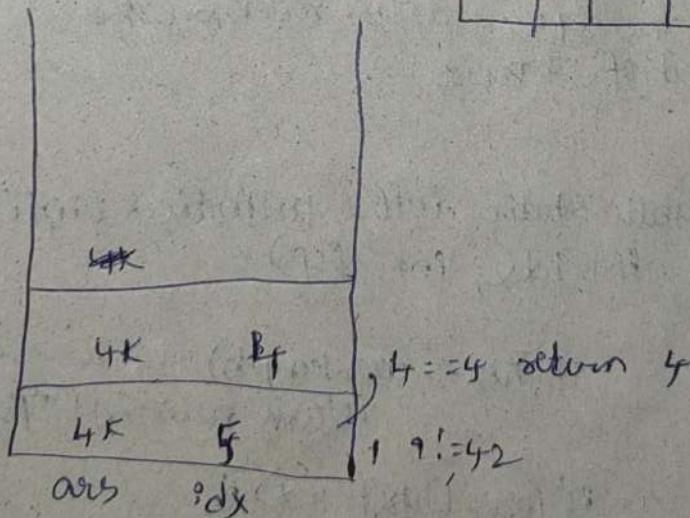
- 1) Pass arr, idx-1, x
- 2) Check if arr(idx) == x else do idx-1 and check
- 3) If it reaches the first element not found return -1;

Program:

```
lastIndex(a, n-1, x);
```

```
public static int lastIndex(int[] arr, int idx, int n) {
    if (idx == 0)
        return -1;
    if (arr[idx] == x) —(1)
        return idx;
    else
        return lastIndex(arr, idx-1, x); —(2)
```

Dry Run:



0	1	2	3	4	5
15	11	40	4	4	9

$x = 40$

6. All Indices of Array

Sample Input

5 - 7
15
11
40
4
4
9
4 → X

Sample Output

3
4

Constraints

1 ≤ n ≤ 10⁴

0 ≤ n₁, n₂, ..., n_k ≤ elements

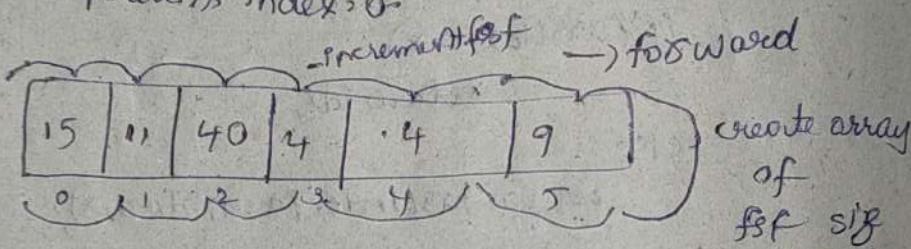
≤ 10³

0 ≤ x ≤ 10³

Steps:

- fsf → indicates the total number of 'x' element got repeated in Array and idx represents index = 0.

2.



- ~~Return~~ Create
3. Create the array after reaching the end of elements.

3 | 4
return thisArray

Program:

```
public static int[] allIndices (int[] arr, int x,
                               int idx, int fsf)
```

{

if (idx == arr.length)

return new int[fsf];

if (arr[idx] == x) {

int[] lq = allIndices (arr, x, idx + 1,

fsf + 1);

lq[fsf] = idx;

return lq;

}

else if

int C3(a: allIndices(ar, x, idx+1, fsf); } -2
return a;

3

5

	0	1	$2^{x=6}$	3	4	5
4K	15	11	40	4	4	9

Dry Run:

4x , 4 , 6 , 2	1, 2 else	$\boxed{3 4}$	$b = 6$, Return arrays of size of fsf
4x , 4 , 5 , 2	1, 2 else	$\boxed{3 4}$	(size of fsf)
4x , 4 , 4 , 1	1, 2 else	$\boxed{3 4}$	
4x , 4 , 3 , 0	1, 2 else	$\boxed{3 4}$	
4x , 4 , 2 , 0	1, 2 else	$\boxed{3 4}$	
4x , 4 , 1 , 0	1, 2 else	$\boxed{3 4}$	
4x , 4 , 0 , 0	1, 2 else	$\boxed{3 4}$	
args x idx fsf	$b = 5$	$\boxed{3 4}$	

Recursion with Array List

1. Get Subsequence

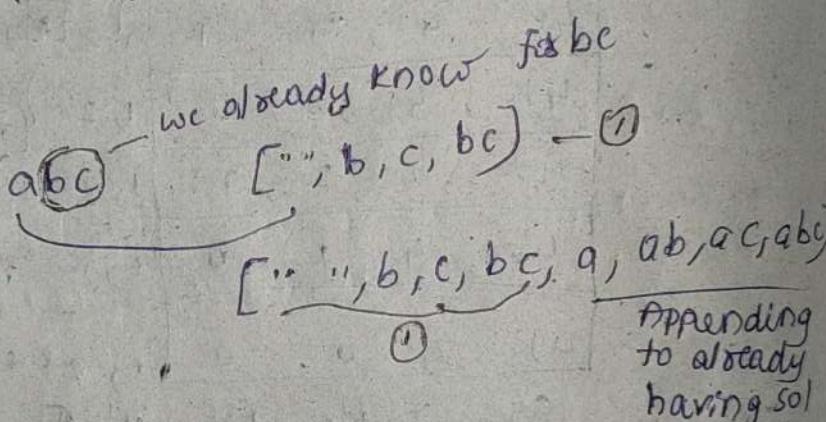
Input
abc

Output
[, c, b, bc, a, ac, ab, abc]

Constraints

$$0 \leq s < \text{length}_s$$

Steps:

1. Consider abc
2. We keep faith that we already have subsequences of bc then we will find abc.
3. 

we already know $f(bc)$

[, b, c, bc] → ①

[, , b, c, bc, a, ab, ac, abc]

①

Appending to already having sol
4. We get starting character and will pass the rest of string to the function
5. For the base condition, return the Array list with empty string.
- 6) for the new Array list, append the result and also concatenate the character with the result.

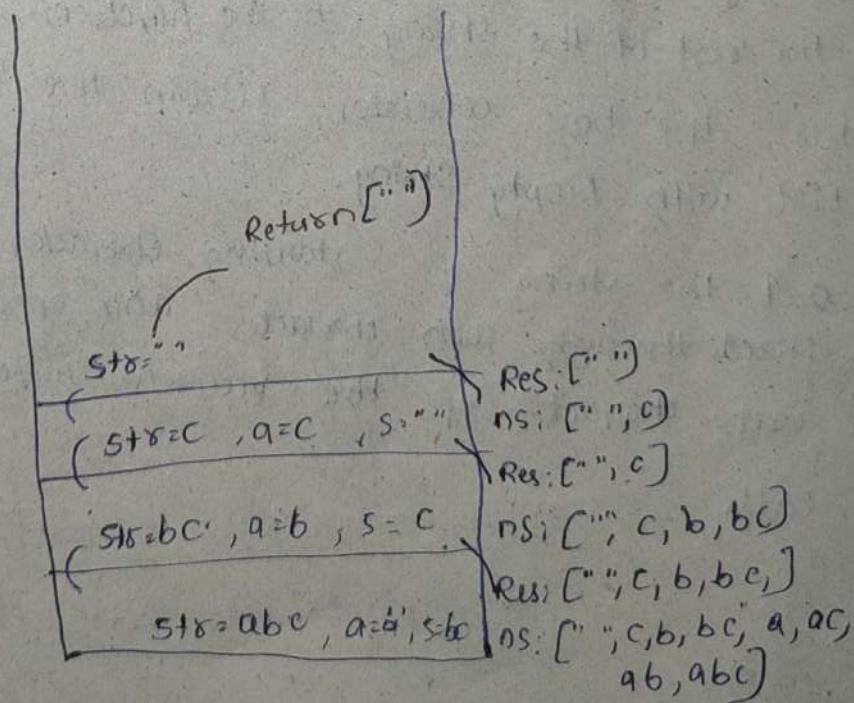
Program:

```

public static ArrayList<String> gss(String str) {
    if(str.length() == 0) {
        ArrayList<String> bs = new ArrayList<String>();
        bs.add("");
        return bs;
    }
    char a = str.charAt(0);
    String s = str.substring(1); ArrayList<String> result = gss(s);
    ArrayList<String> ns = new ArrayList<String>();
    for(String i : result)
        ns.add(i);
    for(String i : result)
        ns.add(a+i);
    return ns;
}

```

Dry Run:



2. Get KPC

Input

78

Output

[tv, tw, tx, uv, uw,
ux)

Constraints

0 <= Str.length <= 10
str contains
numbers only

Steps:

1. Store all the number keys in string Array
2. Consider 789
3. we keep faith that we already know the key values of 89, from that 789 can be computed.
4. Already know [vy, vz, wy, wz, xy, xz]
 7 (8 9) → {tvy, tvz, twy, twz, txy, txz,
 uvy, urz, uwz, uxz, uxy, uxz}
 we can calculate 789
5. We get the starting character and will pass the rest of the string to the function.
6. For, the Base condition, return the Array List with Empty string.
7. Get the string of starting character and iterate through each character and append each character of the recursive function call.

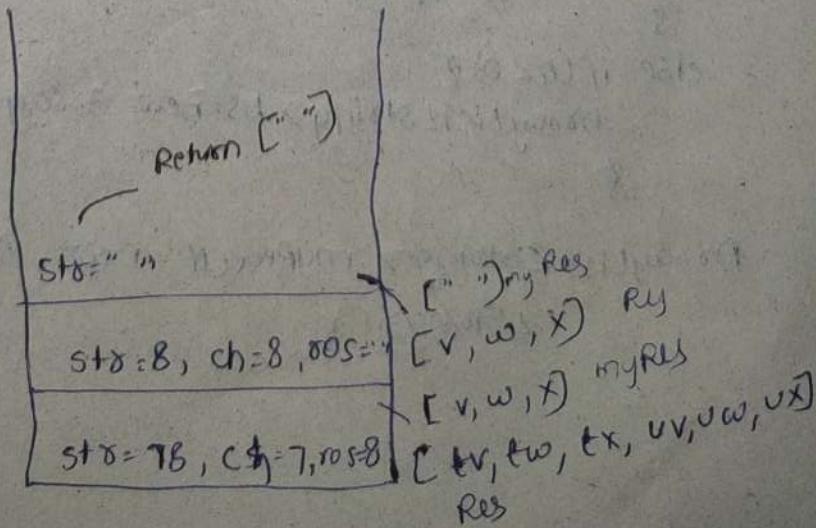
Program:

```

static String[] codes = {"", "abc", "def", "ghi", "jkl",
"mno", "pqrs", "tu", "vwx", "yz"};
public static ArrayList<String> getKPC(String str) {
    if (str.length() == 0) {
        ArrayList<String> base = new ArrayList<String>();
        base.add("");
        return base;
    }
    char ch = str.charAt(0);
    String ros = str.substring(1);
    ArrayList<String> myres = getKPC(ros);
    ArrayList<String> res = new ArrayList<String>();
    String temp = codes[Character.getNumericValue(ch)];
    for (int i = 0; i < temp.length(); i++) {
        char tt = temp.charAt(i);
        for (String j : myres) {
            res.add(tt + j);
        }
    }
    return res;
}

```

Dry Run:



3. Get Stair Paths

Input:

3

Output

(1,1,1,2,2,1,3)

Constraints

0 <= n <= 10

Exit:

Steps:

1. As per the statement, either 1 step / 2 steps / 3 steps at a time possible.
2. We make three recursive calls for (n-1), (n-2), (n-3) and we will get the respect and append 1 step, 2 step, 3 step for respective recursive calls.
3. If n == 0, no more steps possible, so a blank list is returned.
4. If n < 0, return Empty List.

Program:

```

public static ArrayList<String> getStairPaths(int n){
    if(n == 0) {
        ArrayList<String> bs = new ArrayList<String>();
        bs.add("");
        return bs;
    }
    else if(n < 0) {
        ArrayList<String> bs = new ArrayList<String>();
    }

    ArrayList<String> myResult = new ArrayList<String>();
    
```

```
ArrayList<String> sp1 = getStairPaths(n-1); -①
```

```
for (String i: sp1)
```

```
myResult.add("1" + i);
```

```
ArrayList<String> sp2 = getStairPaths(n-2); -②
```

```
for (String i: sp2)
```

```
myResult.add("2" + i);
```

```
ArrayList<String> sp3 = getStairPaths(n-3); -③
```

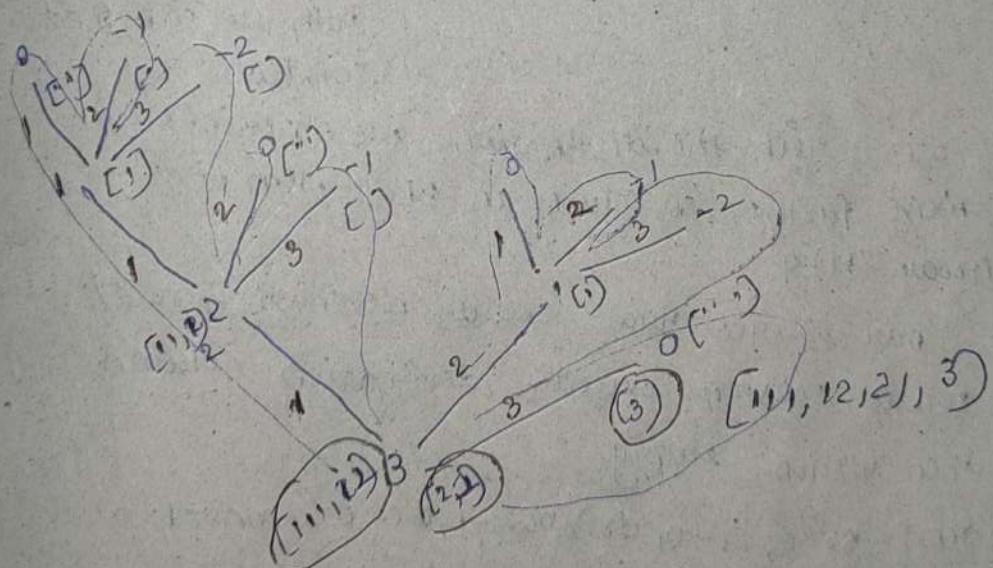
```
for (String i: sp3)
```

```
myResult.add("3" + i);
```

```
return myResult;
```

```
}
```

Day Run:



4. Get Maze Paths

Input

3 → rows
3 → cols

Output

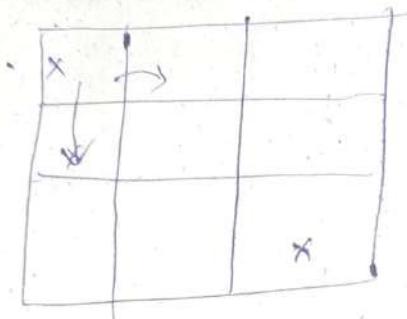
[hhvh, hvvh, vhvh,
vhvh, rvhh]

Constraints

$0 \leq n \leq 10$
 $0 \leq m \leq 10$

Steps:

1. High level thinking



, we assume we take one horizontal step and keep faith from remaining path we have solution

Similarly, if we take one vertical step and keep from remaining path we can find solution

2. If we reach the destination, we no need to move further, so return an blank item of Array string

3. If our source row exceeds destination row/ source column exceeds destination column then return empty

Program: $sr=0, sc=0, dr=rows-1, dc=columns-1$

```
public static ArrayList<String> getMazePaths (int sr)
    int sc, int dr, int dc) {
    if (sr > dr || sc > dc) {
        return new ArrayList<String>();
    }
}
```

```

if (sx == dx && sy == dy) {
    ArrayList<String> base = new ArrayList<String>();
    base.add("");
    return base;
}

```

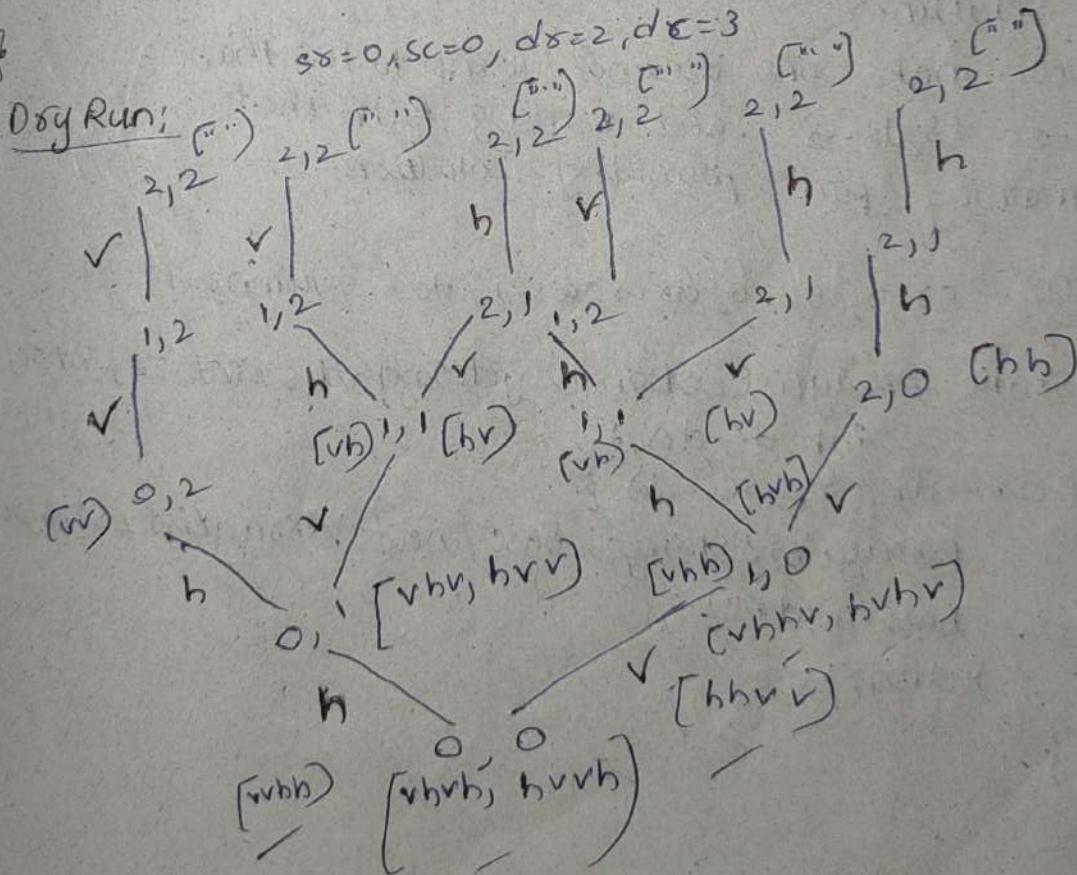
```

ArrayList<String> myResult = new ArrayList<String>();
ArrayList<String> hr = getMazePaths(sx, sy, sc, dc);
for (String i : hr)
    myResult.add("h" + i);

ArrayList<String> vc = getMazePaths(sx, sy, sc, dc);
for (String i : vc)
    myResult.add("v" + i);

return myResult;
}

```



5. Get Maze Path with Jumps

Input:

2 - rows

2 - columns

Output

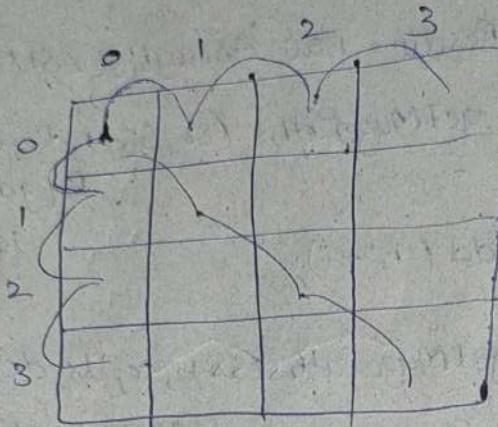
[h1v1, v1h1, d1]

Constraints

$$0 \leq n \leq 10$$

$$0 \leq m \leq 10$$

Steps:



with High level thinking

We can make one horizontal jump to less than column no of jumps and keeping faith with remaining (we have solution) to end point.

- 1) We can make one vertical jump to less than number of rows and keeping faith that remaining path providing solution.
- 2) We can make one diagonal jump to less than $i_c - sc$ & $dc - sr$ and keeping faith that remaining path providing solution.
- 3) We can make one horizontal jump to less than $dc - sc$ and keeping faith that remaining path providing solution.

Program: $sr = 0, sc = 0, dr = \text{rows} - 1, dc = \text{columns} - 1$

```
public static ArrayList<String> getMazePaths (int sr, int sc)
{
    int dr, dc;
    if (sr == dr && sc == dc) {
        ArrayList<String> base = new ArrayList<String>;
        base.add ("");
        return base;
    }
}
```

1) for horizontal jumps

```
for (int i=1; i <= dc - sc; i++) {
```

```
    ArrayList<String> hr = getMazePaths(sr, sc+i, dr, dc);
```

```
    for (String j: hr)
```

```
        myResult.add("h"+i+j);
```

}

2) for vertical jumps

```
for (int i=1; i <= dr - sr; i++) {
```

```
    ArrayList<String> vr = getMazePaths(sr+i, sc, dr, dc);
```

```
    for (String j: vr)
```

```
        myResult.add("v"+i+j);
```

}

3) for diagonal jumps

```
for (int i=1; i <= dc - sc && i <= dr - sr; i++) {
```

```
    ArrayList<String> dg = getMazePaths(sr+i, sc+i, dr, dc);
```

```
    for (String j: dg)
```

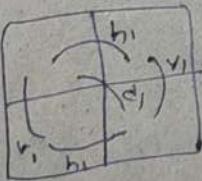
```
        myResult.add("d"+i+j);
```

}

```
return myResult;
```

}

Dry Run:



h1 v1 , v1 h1 , d1

Recursion On the Way-Up

1. Point Subsequence

<u>Input:</u>	<u>Output:</u>	<u>constraints</u>
yvTA	yvTA	$0 \leq \text{str.length} \leq 7$
	yvT	
	yvA	
	yr	
	yTA	
	YT	
	yA	
	y	
	vTA	
	VT	
	VA	
	V	
	TA	
	T	
	A	

Steps:

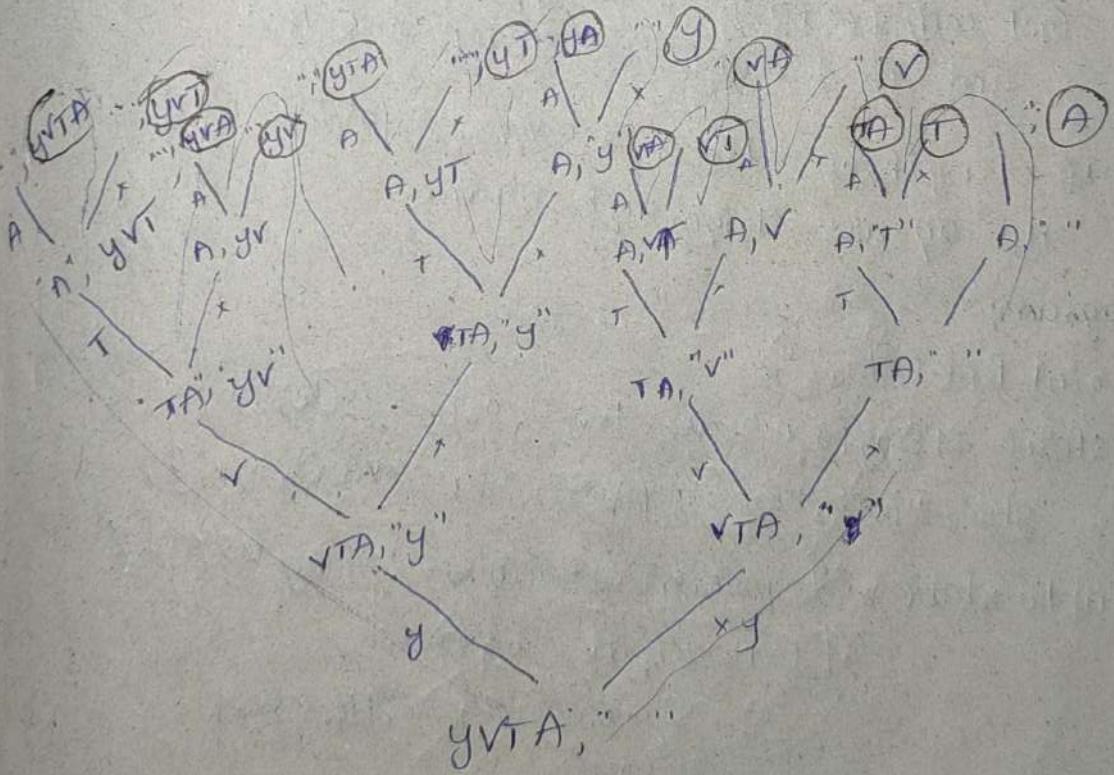
1. We will have an empty string, we will get the first character and rest of the string goes to recursive calls.
2. Two cases: The first character will be considered or not.
3. If the length of string reaches zero point answer and return.

Program:

```
PrintSS(sts, " ");
```

```
public static void PrintSS(string sts, string ans) {
    if(sts.length() == 0) {
        println(ans);
        return;
    }
    char ch = sts.charAt(0);
    String ros = sts.substring(1);
    PrintSS(ros, ans + ch);           // character considered
    PrintSS(ros, ans + "");          // character not considered
}
```

Dry Run:



2 Point KPC

Input

78

Output

tv
tw
tx
uv
uw
ux

constraints

$0 \leq s.length() \leq 10$
str contains numbers
only

Steps:

1. we will have an empty string, we will get the first character and rest of the string goes to recursive calls.
2. Two Cases The first character will be considered and appended to answer string get the respective string from the codes Array.
3. Get each character and append to answer and go through the recursive calls.
4. If length of string reaches 0, print the answer string and return.

Program:

```
printKPC(str, "");
```

static String[] words = {"", "abc", "def",
"ghi", "jkl", "mno", "pqrs", "tu", "vwx", "yz"};

```
public static void printKPC(String str, String ansf) {  
    if (str.length() == 0) {  
        System.out.println(ansf);  
        return;  
    }
```

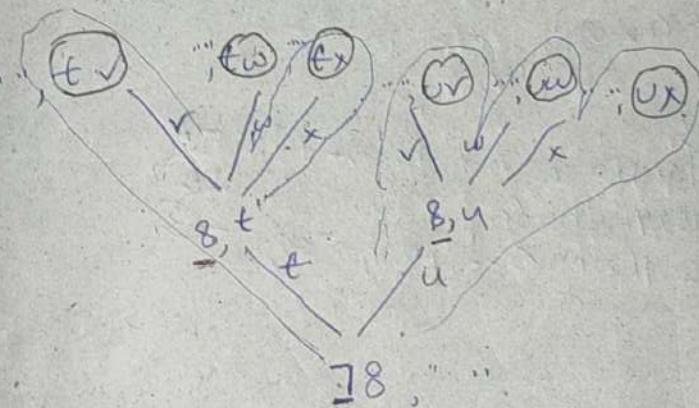
}

```

char ch = str.charAt(0),
String temp = str.substring(1);
String curri = words[ch - '0'];
for (int i = 0; i < curri.length(); i++) {
    char a = curri.charAt(i);
    PointKPC(temp, ansf + a);
}

```

Dry Run:



3. Print Stair Paths

Sample Input

3

Sample Output

111
12
21
3

Constraints

$0 \leq n \leq 10$

Steps:

- As per the statement, we can make either 1/2/3 steps. So we will have 3 recursive calls.
- If we take one step, append 1 to ans in Recursive call, similarly for 2 and 3 steps.

- 3. If we reach the destination, no need to move point ans, return
 - 4. If we make higher number of steps than the actual left over steps, return; [we cannot make solution in that path].

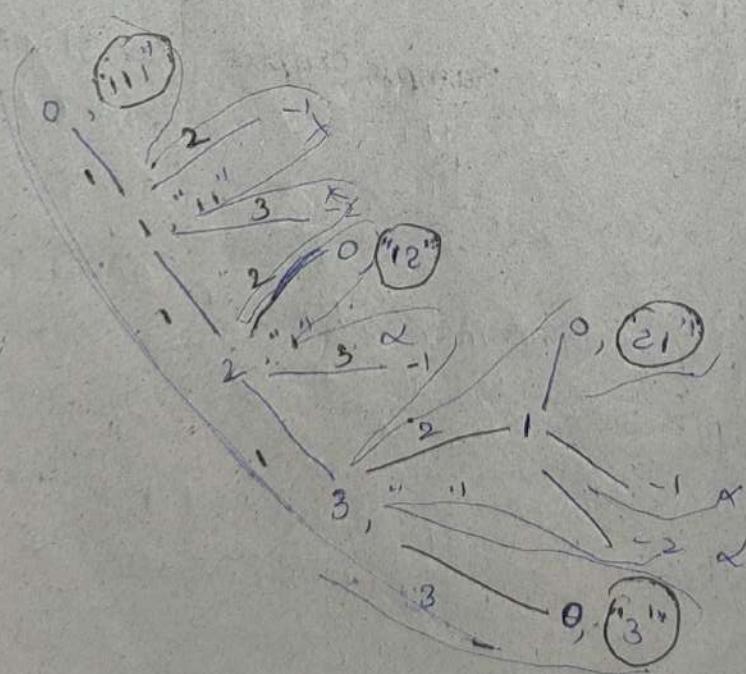
Program:

```

    pointStairPaths(n, "");
public static void pointStairPaths(int n, String path) {
    if (n == 0) {
        System.out.println(path);
        return;
    }
    if (n < 0) {
        return;
    }
    pointStairPaths(n-1, path + "1");
    pointStairPaths(n-2, path + "2");
    pointStairPaths(n-3, path + "3");
}

```

3
Doy Run:



Print Maze Paths

Input

2 → rows
2 → columns

Output

hv

vh

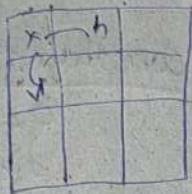
Constraints

$$0 \leq n \leq 10$$

$$0 \leq m \leq 10$$

Steps:

1. we take an empty string psf (path so far)



→ we assume we take one horizontal step and assign to empty string and goes through iterations

→ we assume we take one vertical step and assign to psf string and goes through iterations

2. If we reach the destination, print psf

3. If our source row exceeds destination row
source column exceeds destination column
then return empty.

Program: $sr=0, sc=0, dr=\text{rows}-1, dc=\text{columns}-1$

pointMazePaths(0, 0, n-1, m-1, "");

public static void pointMazePaths(int sr, int sc, int dr)

int dc, String psf) {

if ($sr == dr \text{ and } sc == dc$) {

System.out.println(psf);

return;

}

if ($sr > dr \text{ or } sc > dc$) {

return;

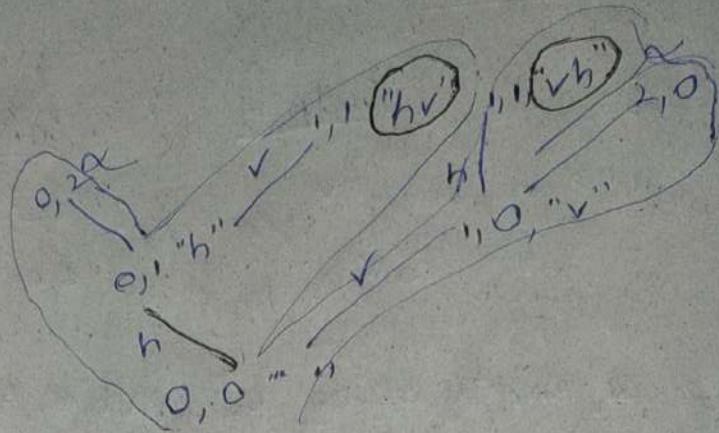
}

pointMazePaths(sr, sc+1, dr, dc, psf + "h");

pointMazePaths(sr+1, sc, dr, dc, psf + "v");

}

D&Q Run;



5. Point Maze Paths with Jumps

Constraints

Input

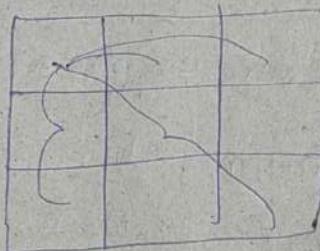
3 - rows

3 - columns

Output

h1h1v1v1	or	$0 \leq n \leq 5$
h1h1v2	v1h1h1v1 / d1h1v1	$0 \leq m \leq 5$
h1v1h1v1	v1h1v1h1 / d1h1h1	
h1v1v1h1	v1h1d1	
h1v1v1h1	v1h2v1	d1d2
b1v1d1	v1d1, h1h1	d2
h1v2d1	v1v1 b2	
h1d1v1	v1v1, b2	
h2v1v1	v2, d1h1	
h2v2	v2 h1h1	
	v2 b2	

Steps:



we can make one horizontal jump to do-5C horizontal jumps and assign 'h' and number to psf.

→ we can make one vertical jump to do-5R and assign rand number to psf.

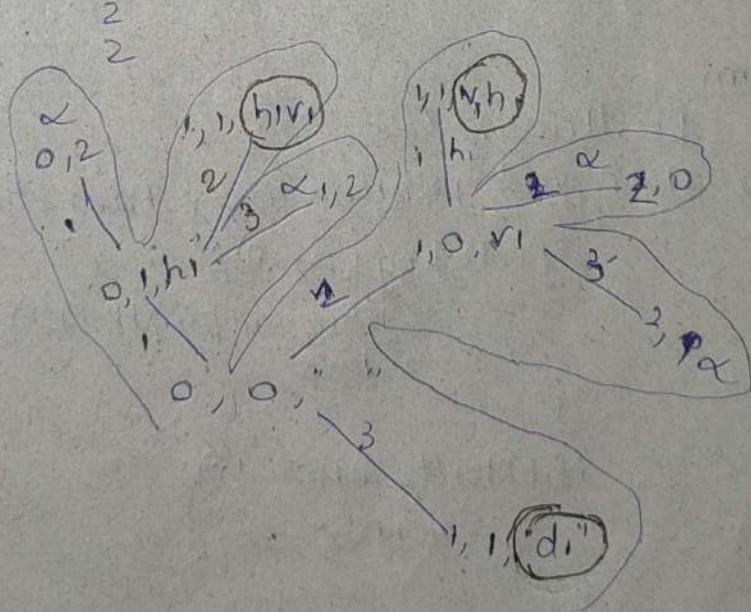
→ Similarly for diagonal(s) $i = dx - sc$; $j = dy - sr$ and assign diagonal and number to psf.

→ If you reach destination point, print psf, return

Program:

```
printMazePaths(0, 0, n-1, m-1, "");  
public static void printMazePaths(int sx, int sc, int dx,  
int dc, String psf) {  
    if (sx == dx && sc == dc) {  
        System.out.println(psf);  
        return;  
    }  
    // horizontal jump  
    for (int i = 1; i <= dc - sc; i++)  
        printMazePaths(sx, sc + i, dx, dc, psf + "h" + i);  
    // vertical jump  
    for (int i = 1; i <= dx - sx; i++)  
        printMazePaths(sx + i, sc, dx, dc, psf + "v" + i);  
    // diagonal jumps  
    for (int i = 1; i <= dc - sc && i <= dx - sx; i++)  
        printMazePaths(sx + i, sc + i, dx, dc, psf + "d" + i);  
}
```

Dry Run:



6. Print Encodings

Input

655196

Output

feeaif

feesf

Constraints

$$0 \leq \text{str.length} \leq 10$$

Steps:

1. Take an Empty string
2. Get starting character from string, and pass rest of string, append [Get integer (From starting character) - 1 + 'a'] to asf (Recursively)
3. If length of string ≥ 2 , then check the first two characters are in the range If they are, then pass rest of string, append [Get integer (From starting 2 characters - 1 + 'a')] to asf Recursively
4. If $\text{str.length} = 0$, then return asf
5. If starting character is '0' then return.

Program:

```

printEncodings(str, "0");
public static void printEncodings (String str, String asf)
{
    if (str.length() == 0)
    {
        System.out.println(asf);
        return;
    }
    if (str.charAt(0) != '0')
    {
        return;
    }
}

```

String $s = str.substring(0, i)$;

String $sos = str.substring(i)$;

PrintEncodings(sos, asf + (char)(Integer.parseInt(s) - 1 + 'a'));

if(str.length() >= 2) {

String $s1 = str.substring(0, 2)$;

String $sos1 = str.substring(2)$;

if(Integer.parseInt(s1) <= 26) {

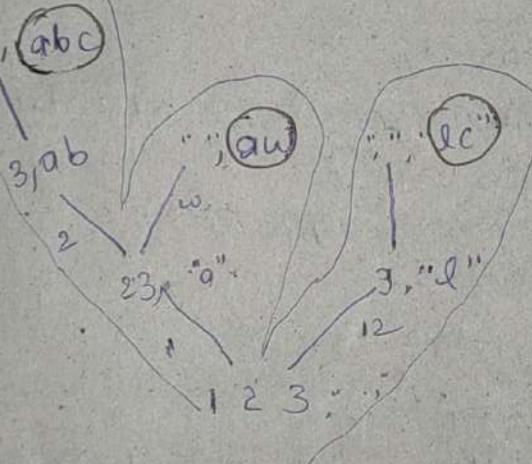
PrintEncodings(sos1, asf + (char)(Integer.parseInt(s1) - 1 + 'a'));

}

}

Dry Run:

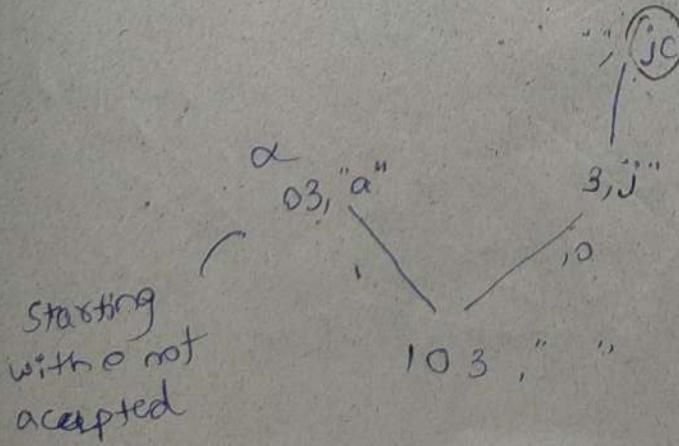
①



123

②

103



7. Print Permutations

<u>Sample Input</u>	<u>Sample Output</u>	<u>constraints</u>
abc	abc acb bac bca cab cba	$0 < \text{str.length} \leq 10$

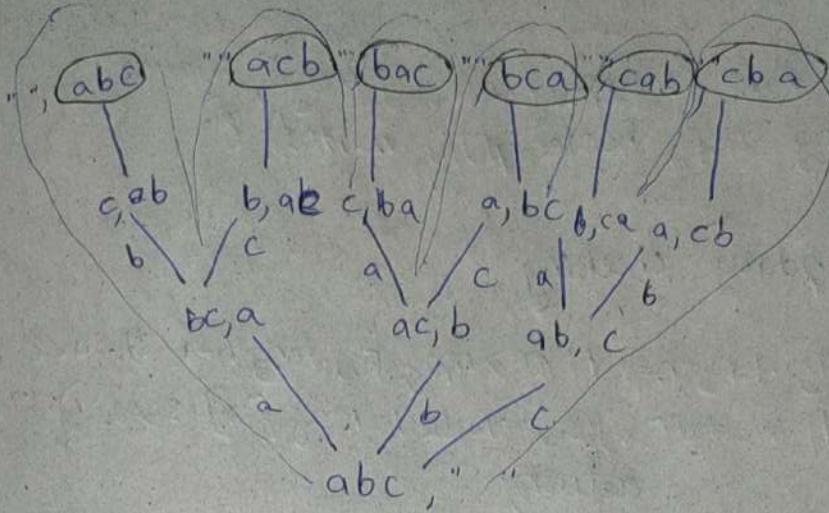
Steps:

1. Take an empty string;
2. Iterate through the string, get first character, assign the $\text{ros} \rightarrow$ (by deleting the first character), append $\text{asf} + \text{ch}$ \rightarrow pass Recursively.
3. If length of string reaches 0, print the asf

Program:

```
PointPermutations(str, "");  
public static void pointPermutations(String str, String asf){  
    if(str.length() == 0) {  
        System.out.println(asf);  
        return;  
    }  
    for(int i=0; i<str.length(); i++) {  
        char ch = str.charAt(i);  
        String ros = str.substring(0, i) + str.substring(i+1, str.length());  
        pointPermutations(ros, asf + ch);  
    }  
}
```

try Run!



Recursion Backtracking

Flood Fill

Sample Input

```
3 3  
0 0 0  
1 0 1  
0 0 0
```

Sample Output

8dd8

Constraints

$1 \leq n \leq 10$
 $1 \leq m \leq 10$

e_1, e_2, \dots, e_m elements
belongs to set $\{0, 1\}$

Steps:

- As per the statement we need to use operations in order of t, l, d, r $\begin{cases} \text{top} \\ \text{left} \\ \text{down} \\ \text{right} \end{cases}$
- We keep track an visited Array, if element is visited, path found make true else after performing all recursive calls no path found then make visited = false
- Check the boundary conditions and if element is , \rightarrow return, already visited \rightarrow return
- If you reach the destination, print ans \rightarrow return

Program:

```
public static void floodfill(int[][] maze, int sr,  
    int sc, String asf, boolean[][] visited) {
```

// boundary checking

```
if (sr > maze.length - 1 || sc > maze[0].length - 1 || sr < 0  
    || sc < 0 || visited[sr][sc] == true || maze[sr][sc] == -1)  
    return;
```

```
if (sr == maze.length - 1 || sc == maze[0].length - 1) {  
    System.out.println(asf);  
    return;
```

}

```
visited[sr][sc] = true;
```

```
floodfill(maze, sr - 1, sc, asf + "t", visited); -1
```

```
floodfill(maze, sr, sc - 1, asf + "l", visited); -2
```

```
floodfill(maze, sr + 1, sc, asf + "d", visited); -3
```

```
floodfill(maze, sr, sc + 1, asf + "r", visited); -4
```

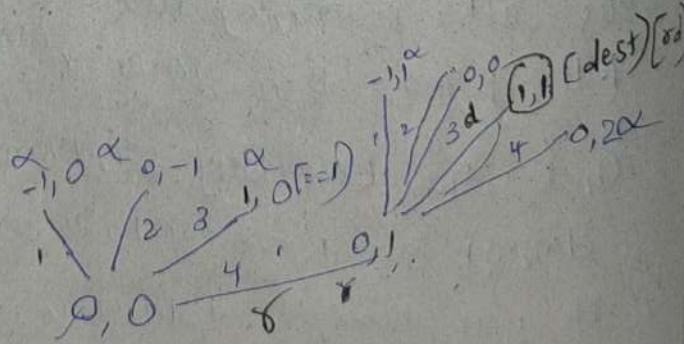
```
visited[sr][sc] = false;
```

}

Dry Run:

$$\begin{bmatrix} \alpha & \alpha \\ \alpha & \alpha \end{bmatrix}$$

sol: (rd)



2. Target Sum Subsets

Input

5	-	n
10		
20		
30		
40		
50		
60	-	sum

elements

Output

10, 20, 30,
10, 50,
20, 40,

Constraints

$$0 \leq n \leq 30$$

$$0 \leq n_1, n_2, \dots, n_{\text{element}} \leq 20$$

~~$$0 \leq \text{target} \leq 50$$~~

Steps:

1. Consider the element of the array and add to asf, and targetsum / Don't add the element to asf and targetsum
2. If targetsum exceeds target return,
3. If all elements are covered and $\text{targetsum} = \text{target}$, point asf \rightarrow Return.

Program:

```

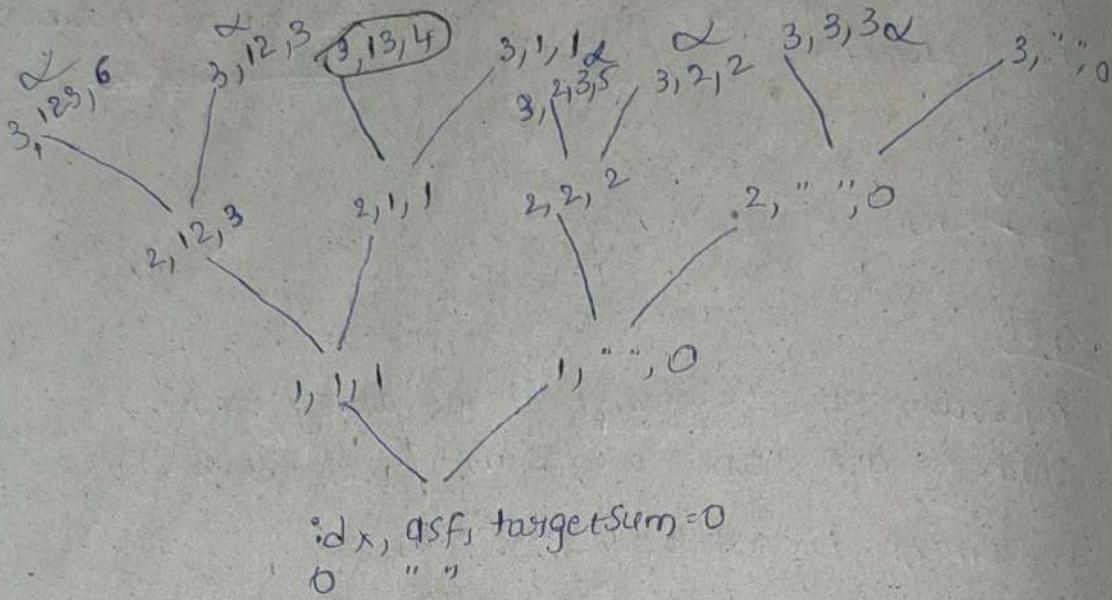
public static void printTargetSumSubset (int arr[],
int idx, String asf, int targetsum, int target) {
    if (targetsum > target)
        return;
    if (idx == arr.length) {
        if (targetsum == target)
            System.out.println(asf + ":");
        return;
    }
    printTargetSumSubset (arr, idx + 1, asf + arr[idx] + ", " + targetsum + arr[idx], target);
    printTargetSumSubset (arr, idx + 1, asf, targetsum, target);
}

```

Day Run:

1	2	3
---	---	---

target=4
targetSum=0



3. N Queens

Input

4

Output

0-1, 1-3, 2-0, 3-2,
0-2, 1-0, 2-3, 3-1,

Constraints

$1 \leq n \leq 10$

Steps:

1. Go through the number of columns
2. Check placing the queen is safe or not with that col and row
3. Place the Queen
4. Make Recursive call with $\text{row} + 1$
5. If placed queen, doesn't make solution, we need to do Backtrack, unplace the queen.

for Queen Safe: check in the row wise, and column wise, diagonal wise^(both), if element found return false.

Program:

```
boolean[][] chess = new boolean[n][n];  
printNQueens(chess, "", 0);
```

1) check Queen is safe to place

```
public static boolean isQueenSafe(boolean[][] chess, int row, int col) {
```

 1) check in col

```
        for (int i = row - 1; i >= 0; i--) {  
            if (chess[i][col] == true)  
                return false;
```

}

 2) check in diag

```
        for (int i = row, j = col - 1; i >= 0 && j >= 0; i--, j--) {  
            if (chess[i][j] == true)  
                return false;
```

}

 3) check in diag2

```
        for (int i = row + 1, j = col + 1; i >= 0 && j < chess.length - 1; i++, j++) {  
            if (chess[i][j] == true)  
                return false;
```

}

 return true;

}

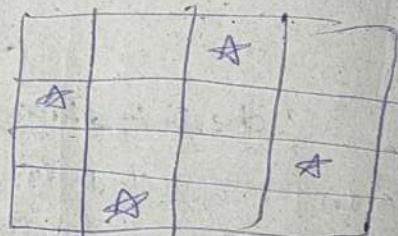
```

public static void printNQueens(boolean[][] chess,
String asf, int row) {
    if (row == chess.length) {
        System.out.println(asf + ".");
        return;
    }
    for (int col = 0; col < chess.length; col++) {
        if (isQueenSafe(chess, row, col) == true) {
            // place Queen
            chess[row][col] = true;
            // make Recursive call
            printNQueens(chess, asf + row + "-" + col, "row+1);
            // unplace the Queen
            chess[row][col] = false;
        }
    }
}

```

Dry Run; — Visualize by Recursion Tree

For $n=4 \rightarrow 2$ possible solutions



4. Knights Tour

Sample Input

5
2
0

Sample Output:

25	2	13	8	23
12	7	24	3	14
1	18	15	22	9
6	11	20	17	4
19	16	5	10	21

Constraints

n=5

0 <= row < n

0 <= col < n

(continues with
all patterns)

Steps:

1. Take initially stepNo=1, and assign it to starting position.
2. Go through all 8 directions, mentioned in statement.
3. If next element not covered, backtrack, and make visited element = 0.
4. If stepNo equal to ArrayLength \rightarrow then assign left over position to stepNo, print the board, and reassign to 0 \Rightarrow to get other possible solutions \Rightarrow return.

Program:

```
PrintKnightsTour(chess, r, c, i);
```

```
Public static void printKnightsTour (int[] chess, int r, int c,
int stepNo) {
    if ((r < 0) || (c < 0) || (r >= chess.length) || (c >= chess.length) || chess[r][c] != 0)
        return;
    if (stepNo == chess.length * chess.length) {
        chess[r][c] = stepNo;
        displayBoard(chess);
        chess[r][c] = 0;
        return;
    }
}
```

$\text{chess}[\gamma][c] = \text{stepnot})$

$\text{PrintKnightsTour}(\text{chess}, \gamma-2, c+1, \text{stepnot}),$

$\text{PrintKnightsTour}(\text{chess}, \gamma-1, c+2, \text{stepnot}),$

$\text{PrintKnightsTour}(\text{chess}, \gamma+1, c+2, \text{stepnot}),$

$\text{PrintKnightsTour}(\text{chess}, \gamma+2, c+1, \text{stepnot}),$

$\text{PrintKnightsTour}(\text{chess}, \gamma+1, c-2, \text{stepnot}),$

$\text{PrintKnightsTour}(\text{chess}, \gamma-1, c-2, \text{stepnot}),$

$\text{PrintKnightsTour}(\text{chess}, \gamma-2, c-1, \text{stepnot}),$

$\text{chess}[\gamma][c] = 0$

3

Dry Run:

Eg for 1 pattern : 2,0 starting position

	0	1	2	3	4	
0	25	2(3)	13(3)	8	27(6)	$\gamma-3, c-1$
1	12(2)	7(4)	287(3)	3(4)	14(5)	$\gamma-1, c-2$
2	1(1)	18	15(1)	25(1)	9(5)	
3	6	"(8)	20(6)	19(6)	4(6)	$\gamma+1, c-2$
4	9	16(1)	5(1)	10(7)	21	$\gamma+2, c-1$