Rscript: a Relational Approach to Program and System Understanding

Paul Klint







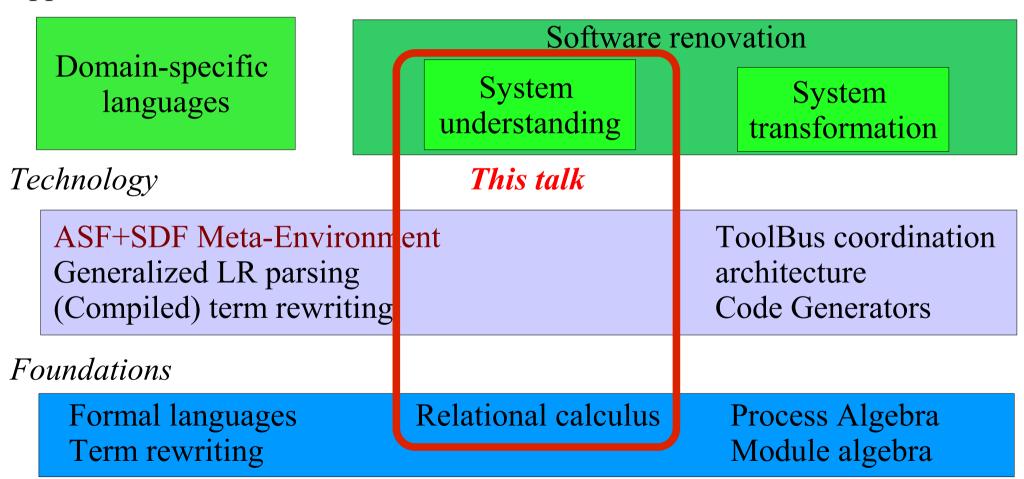
Structure of Presentation

- Background and context
- About program understanding
- Roadmap: Rscript



Background

Application areas





Compilation is a mature area

- Some new developments
 - just-in-time compilation
 - energy-aware code generation
- Many research results are not yet used widely
 - interprocedural pointer analysis
 - slicing
- Why don't we just apply all these techniques to understanding and restructuring?



Compilation is a mature area

- ... of course, we do just that, but ...
- there is a mismatch between
 - standard compilation techniques and
 - the needs for understanding and restructuring



Compilation is ...

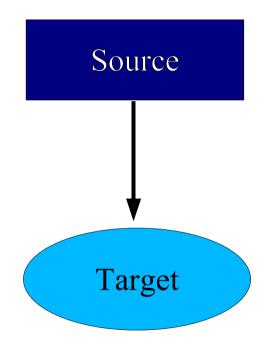
- A well-defined process with well-defined input, output and constraints
- Input: source program in a fixed language with well-defined syntax and semantics
- Output: a fixed target language with well-defined syntax and semantics
- Constraints are known (correctness, performance)
- A batch-like process



Compilation is ...

Single, well defined, source

Single, well defined, target



A batch-like process with clear constraints

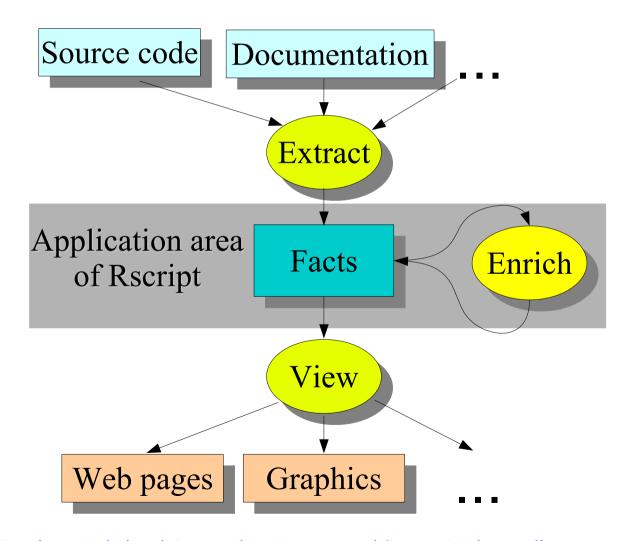


Understanding is ...

- An exploration process with as input
 - system artifacts (source, documentation, tests, ...)
 - implicit knowledge of its designers or maintainers
- There is no clear target language
- An interactive process:
 - Extract elementary facts
 - Abstract to get derived facts needed for analysis
 - View derived facts through visualization or browsing



Extract-Enrich-View Paradigm





Examples of understanding problems

- Which programs call each others?
- Which programs use which databases?
- If we change this database record, which programs are affected?
- Which programs are more complex than others?
- How much code clones exist in the code?



Examples of the results of understanding

- Textual reports indicating properties of system parts (complexity, use of certain utilities, ...)
- Same, but in hyperlinked format
- Graphs (call graphs, use def graphs for databases)
- More sophisticated visualizations



Other aspects of Understanding

- Systems consist of several source languages
- Analysis techniques over multiple language =>
 a language-independent analysis framework is
 needed
- A very close link to the source text is needed



Related approaches

- Generic dataflow frameworks exist but are not used widely
- Relations have been used for querying of software (Rigi, GROK, RPA, ...)
 - All based on untyped, binary, relation algebra
 - Mostly used for architectural, coarse grain, queries



Relation-based analysis

- What happens if we use relations for fine grain software analysis (ex: find uninitialized variables)
- What happens if we use a relational calculus (as opposed to the relational algebra approaches)?
- What happens if we use term rewriting as basic computational mechanism?
 - relations can represent graphs in the rewriting world
- Could yield a unifying framework for analysis and transformation



Roadmap

- Rscript in a nutshell
- Example 1: call graph analysis
- Example 2: component structure
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Rscript in a Nutshell

- Basic types: bool, int, str, loc (text location in specific file with comparison operators)
- Sets, relations and associated operations (domain, range, inverse, projection, ...)
- Comprehensions
- User-defined types
- Fully typed
- Functions and sets of equations over the above



Rscript: examples

- Set: {3, 5, 7}
 - type: set[int]
- Set: {"y", "x","z"}
 - type: set[str]
- Relation: {<"y",3>, <"x",3>, <"z",5>}
 - type: rel[str,int]



Rscript: examples

• rel[str,int]
$$U = \{\langle y'', 3 \rangle, \langle x'', 3 \rangle, \langle z'', 5 \rangle\}$$

- int Usize = #U
 - 3
- rel[int,str] Uinv = inv(U)/

domain:

all elements in lhs of pairs range:

all elements in rhs of pairs carrier:

all elements in lhs or rhs of pairs

set[str] Udom = domain(U)



Comprehensions

- Comprehensions: {Exp | Gen1, Gen2, ... }
 - A generator is an enumerator or a test
 - Enumerators: V: SetExp or <V1,V2>: RelExp
 - Tests: any predicate
 - consider all combinations of values in Gen1, Gen2,...
 - if some Gen is false, reject that combination
 - compute Exp for all legal combinations



Comprehensions

- {X | int X : {1,2,3,4,5}}yields {1,2,3,4,5}
- $\{X \mid int X : \{1,2,3,4,5\}, X > 3\}$
 - yields **{4,5}**
- {<Y, X> | <int X, int Y> : {<1,10>,<2,20>}}
 - yields {<10,1>,<20,2>}



Functions

- rel[int, int] inv(rel[int,int] R) =
 { < Y, X > | < int X, int Y > : R }
 - inv({1,10>, <2,20>} yields {<10,1>,<20,2>}
- rel[&B, &A] inv(rel[&A, &B] R) =
 {<Y, X> | <&A X, &B Y> : R}
 - inv({<1,"a">, <2,"b">}) yields {<"a",1>,<"b",2>}

&A, &B indicate *any* type and are used to define polymorphic functions



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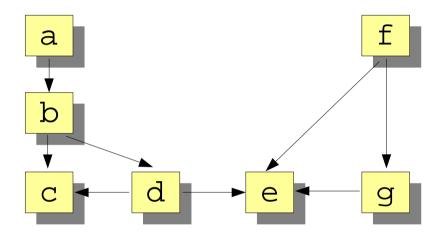


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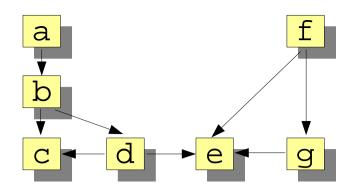


Analyzing the call structure of an application





Some questions



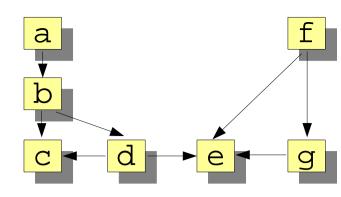
- How many calls are there?
 - int ncalls = # calls
 - 8

Number of elements

- How many procedures are there?
 - int nprocs = # carrier(calls)
 - **-7**

All elements in domain or range of a relations

Some questions



- What are the entry points?
 - set[str] entryPoints = top(calls)
 - {"a", "f"}
- What are the leaves?

The *roots* of a relation (viewed as a graph)

- set[str] bottomCalls = bottom(calls)
- {"c", "e"}

The *leaves* of a relation (viewed as a graph)

Intermezzo: Top

- The roots of a relation viewed as a graph
- top({<1,2>,<1,3>,<2,4>,<3,4>}) yields {1}
- Consists of all elements that occur on the lhs but not on the rhs of a tuple
- set[&T] top(rel[&T, &T] R) = domain(R) \ range(R)

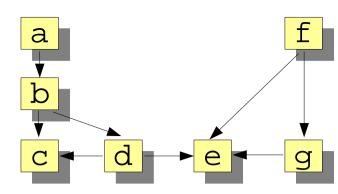


Intermezzo: Bottom

- The leaves of a relation viewed as a graph
- bottom({<1,2>,<1,3>,<2,4>,<3,4>}) yields {4}
- Consists of all elements that occur on the rhs but not on the lhs of a tuple
- set[&T] bottom(rel[&T, &T] R) = range(R) \ domain(R)



Some questions



- What are the indirect calls between procedures?
 - rel[str,str] closureCalls = calls+

- What are the calls from entry point **a**?
 - set[str] calledFromA = closureCalls["a"]
 - {"b", "c", "d", "e"}



domain value "a"

Intermezzo: right image

- Right-image of a relation: all elements that have a given value as left element (resembles array access)
- Notation: relation followed by [Value]
- Ex. Rel = $\{<1,10>,<2,20>,<1,11>,<3,30>,<2,21>\}$
- Rel[1] yields {10,11}
- Rel[{1,2}] yields {10, 11, 20, 21}

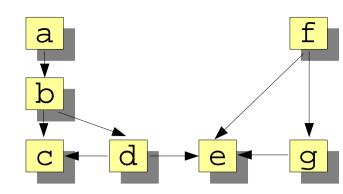


Intermezzo: left image

- Left-image of a relation: all elements that have a given value as right element
- Notation: relation followed by [-,Value]
- Ex. Rel = $\{<1,10>,<2,20>,<1,11>,<3,30>,<2,21>\}$
- Rel[-,10] yields {1}
- Rel[-,{10,20}] yields {1,2}



Some questions

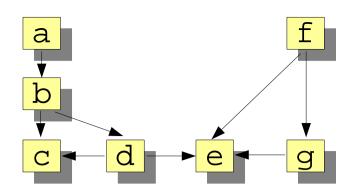


• What are the calls to procedure **e**?

- {"a", "b", "d", "f", "g"}

The domain of image value "e"

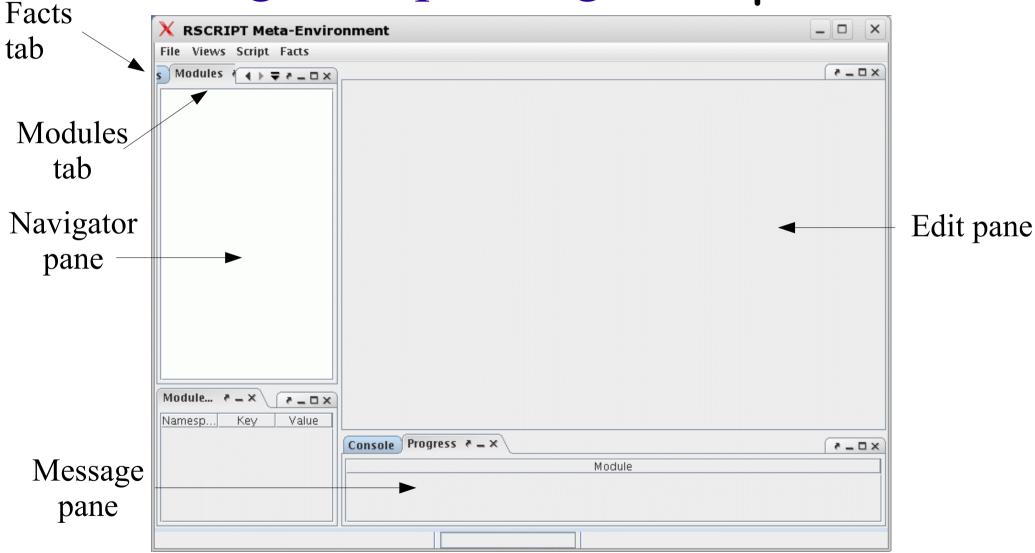
Some questions



- What are the calls from entry point **f**?
 - set[str] calledFromF = closureCalls["f"]
 - {"e", "g"}
- What are the common procedures?
 - set[str] commonProcs =
 calledFromA inter calledFromF
 - {"e"}

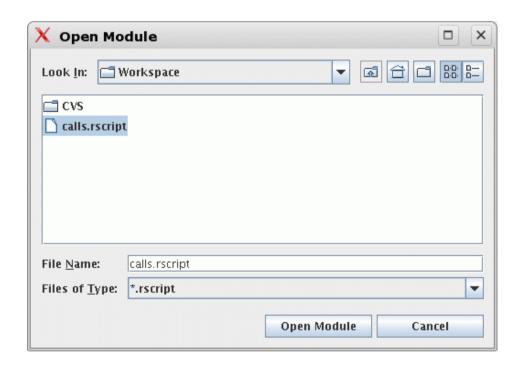
Intersection

Running Rscript using rscript-meta

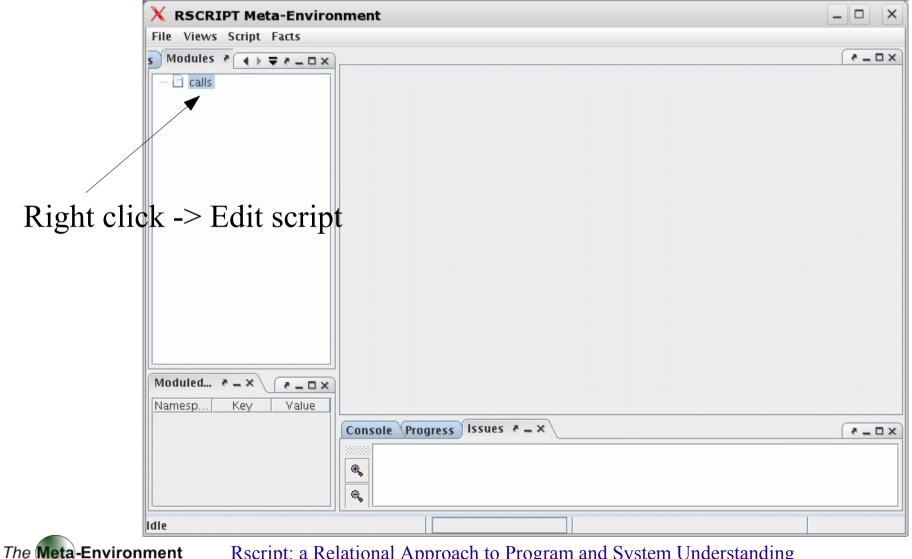




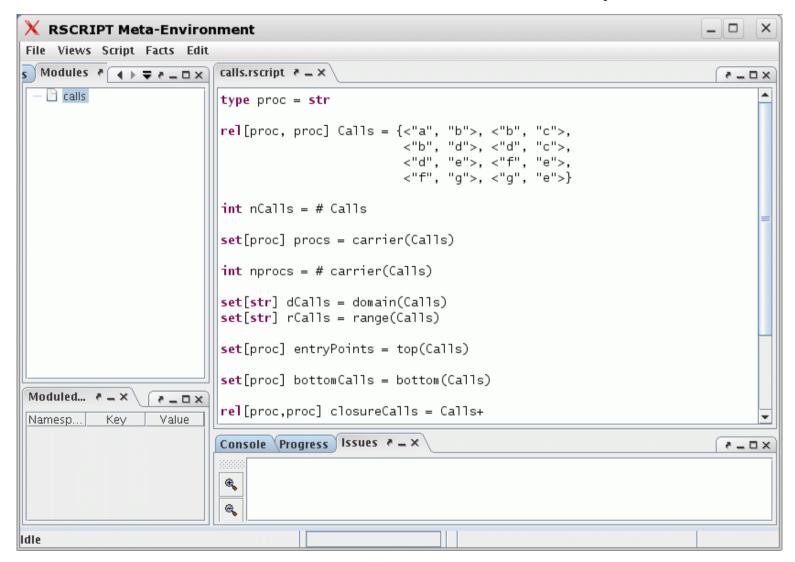
Script -> Open...



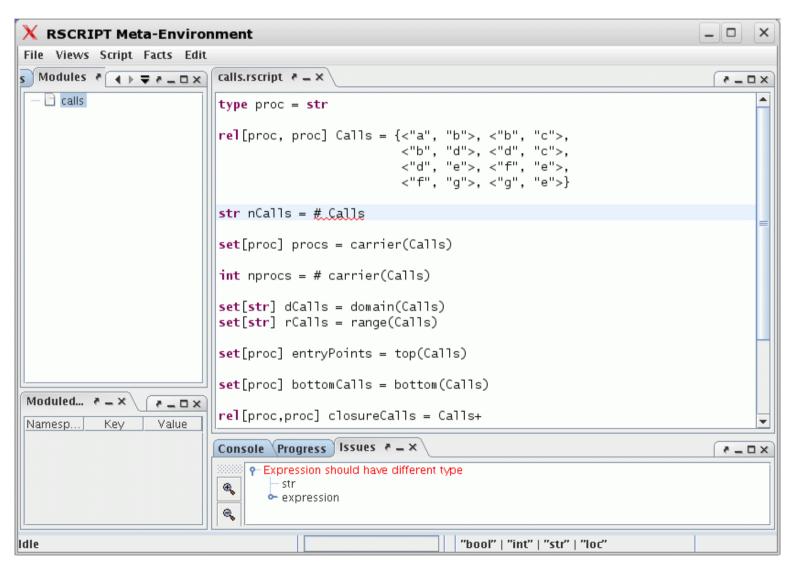
File calls has been opened



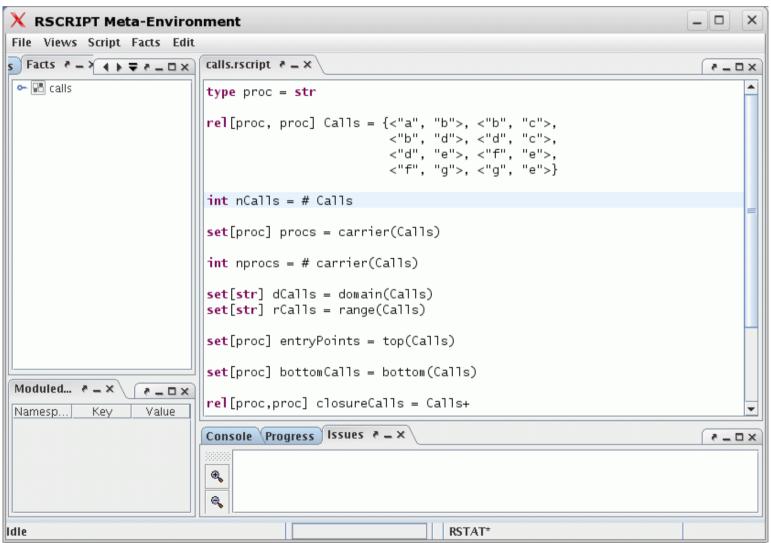
Editing calls.rscript



Making errors ...

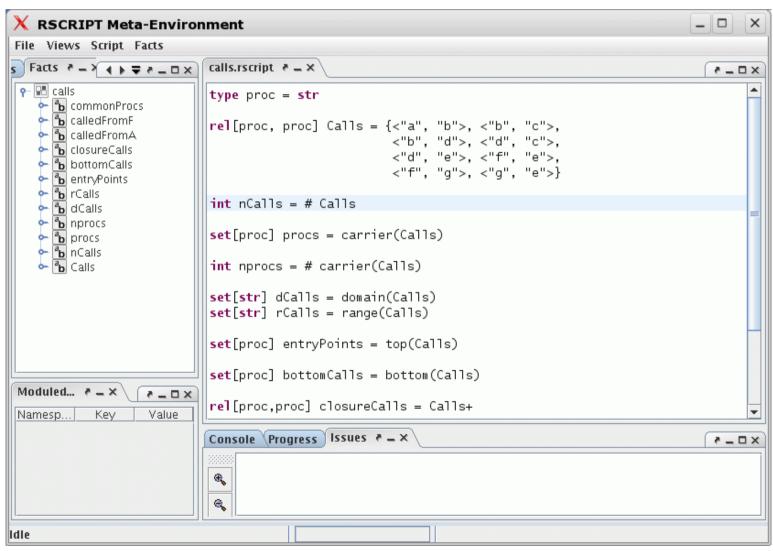


Script -> Run

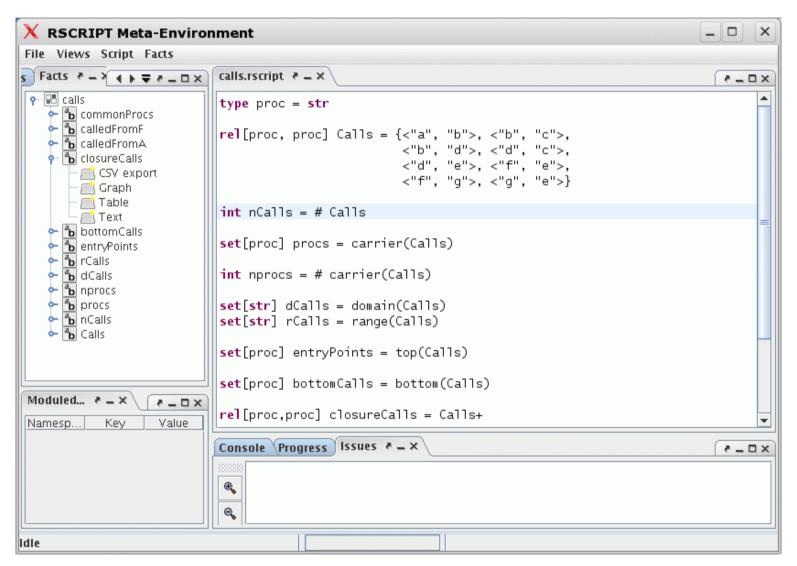




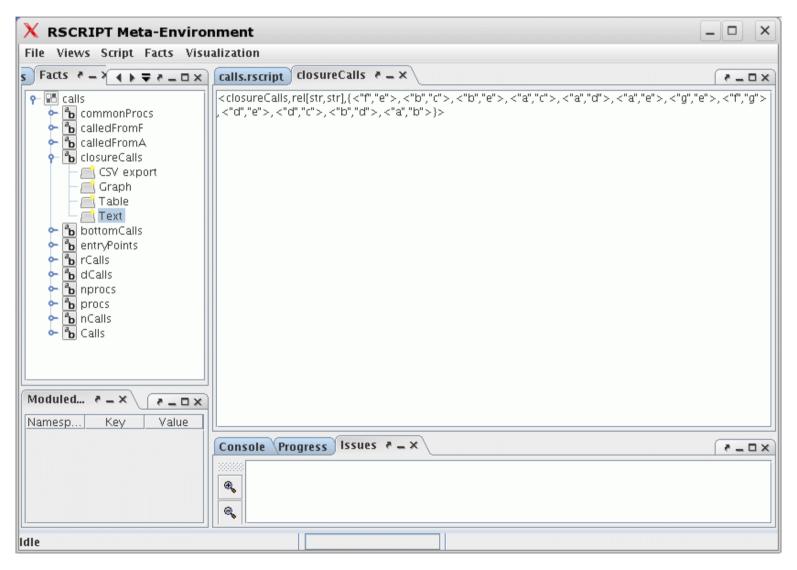
Unfolding the rstore ...



Unfolding closureCalls

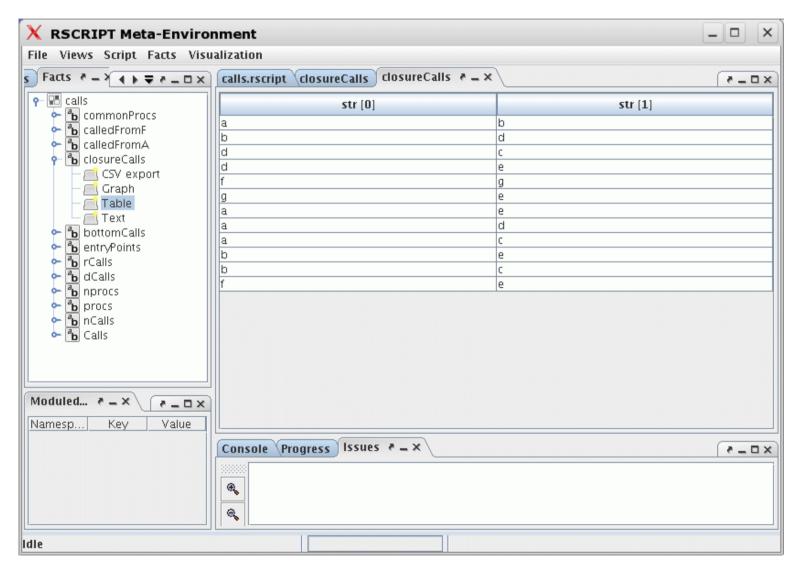


closureCalls as Text



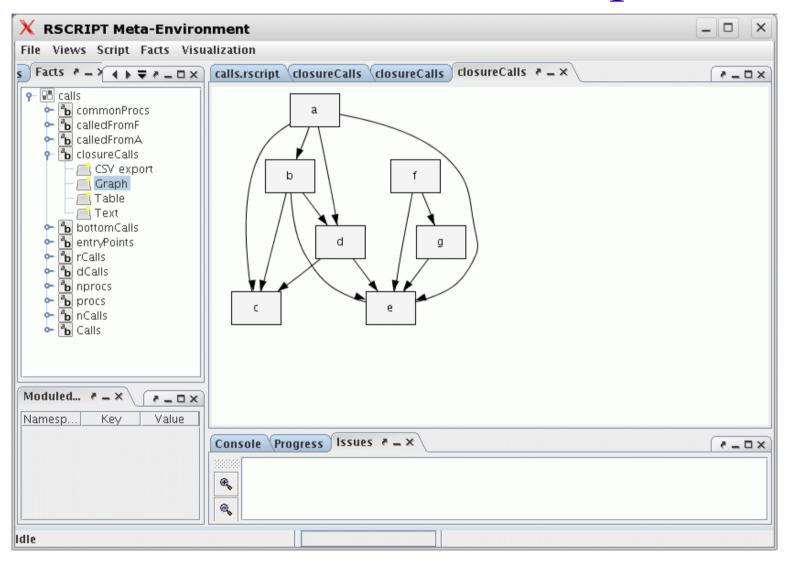


closureCalls as Table





closureCalls as Graph





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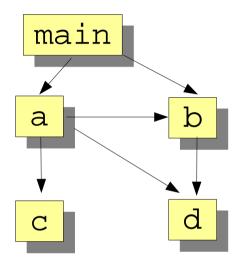


Component Structure of Application

- Suppose, we know:
 - the call relation between procedures (Calls)
 - the component of each procedure (PartOf)
- Question:
 - Can we lift the relation between procedures to a relation between components (Component Calls)?
- This is usefull for checking that real code conforms to architectural constraints



Calls



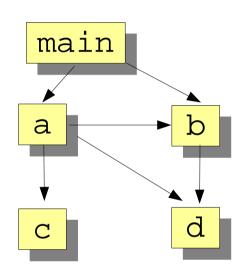
```
type proc = str

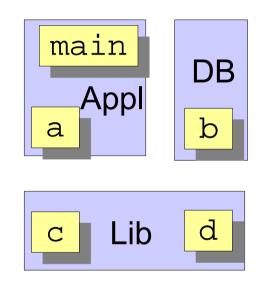
type comp = str

rel[proc,proc] Calls = {<"main", "a">, <"main", "b">, <"a", "b">, <"a", "d">, <"b", "d">}
```



PartOf

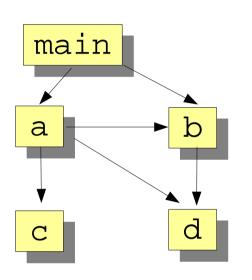


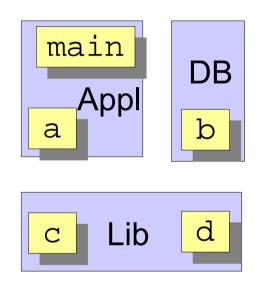


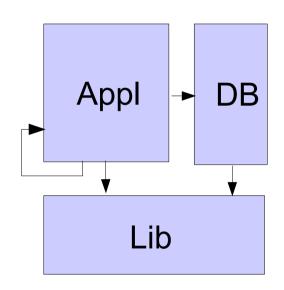
set[comp] Components = {"Appl", "DB", "Lib"}



lift







<comp C1, comp C2> : aPartOf[P1] \times aPartOf[P2] }

rel[comp,comp] ComponentCalls = lift(Calls2, PartOf)

Result: {<"DB", "Lib">, <"Appl", "Lib">, <"Appl", "DB">, <"Appl", "Appl">}



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Cyclic Dependencies

- A class uses (directly or indirectly) itself
- Use = methods calls, inheritance, containment

```
class ContainedClass { }
class SuperClass {}
class SubClass extends SuperClass {
    ContainedClass C;
}

Example of
a contained class
}
```

Motivation: cyclic class dependencies are difficult to understand/maintain



Cyclic Dependencies: Examples

```
class A { B B1; ... }
class B extends A { ... }
```

```
class A { C C1; ... }
class B extends A{ ... }
class C { B B1; ...}
```

Java analysis: classes in cycles

- Assume the following extracted information:
 - rel[str,str] CALL
 - method call from first class to the second
 - rel[str,str] INHERITANCE
 - extends and implements
 - rel[str,str] CONTAINMENT
 - attribute of first class is of the type of the second class
- Question: which classes occur in a cyclic dependency?



Java analysis: cycles in classes

- Define the USE relation between two classes:
 - rel[str,str] USE = CALL union CONTAINMENT union INHERITANCE
 - set[str] ClassesInCycle =
 {C1 | <str C1, str C2> : USE+, C1 == C2}
- In this way we get a set of classes that occur in a cyclic dependency, but ...
- ... which classes are in the cycle?



Java analysis: cyclic classes

```
rel[str,str] USE = CALL union CONTAINMENT union INHERITANCE
```

The Meta-Environment

Each cyclic class is associated with a set of classes that form a cycle

Applications of this approach

- Search for "similar" classes
- Search for design patterns (as characterized by specific relations between the classes in the pattern)

•



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Toy program

```
begin declare x: natural, y: natural,
               z: natural;
                                         y is undefined
   x := 3;
        3 then
          z := y + x
   else
          x := 4
                                          z may be undefined
end
```

Toy program

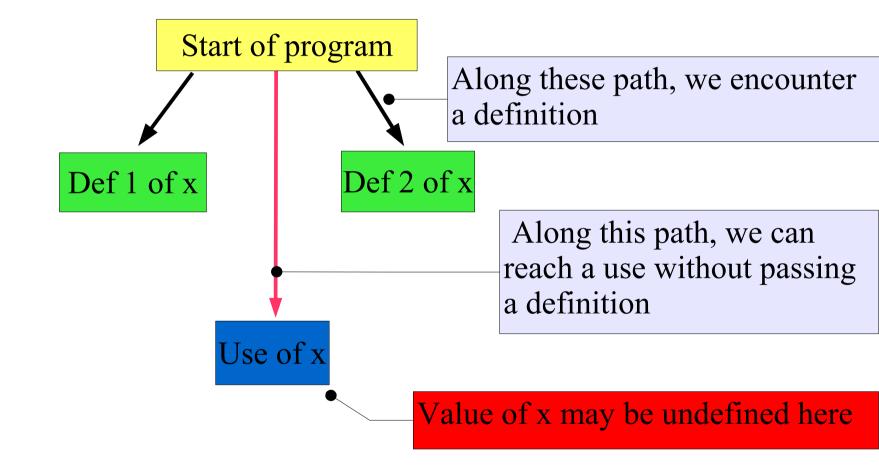
rel[int,str] DEFS = {<1,"x">, <3,"z">, <4,"x">, <5,"y">}

```
begin declare x: natural, y: natural,
                z: natural:
[1] x := 3;
                        rel[int,str] USES = \{<3,"y">, <3,"x">, <5,"z">\}
    if[2] 3 then
      [3] z := y + x
    else
                        rel[int, int] PRED = {<0,1>, <1,2>, <2,3>,<2,4>,
     [4] x := 4
                                       <3,5>,<4,5>}
[5] y := z
```



end

Finding uninitialized variables

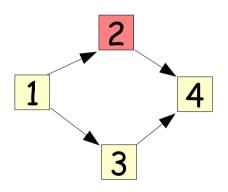




Intermezzo: reachX

- Reachability with exclusion of certain elements
- set[&T] reachX(

set[&T] Start,
set[&T] Excl,
rel[&T,&T] Rel)



reachX({1}, {2}, {<1,2>,<1,3>,<2,4>,<3,4>})yields {<3,4>}



The undefined query

```
rel[int,str] DEFS = ...
                                                       Start from the root
  rel[int,str] USES = ...
                                                                Exclude all
  rel[int,int] PRED = ...
                                                                definitions of V
  rel[int,str] UNINIT =
  { <N,V> | <int N, str V>:USES, N in reachX({0}, DEFS[-,V],PRED)}
There is a path from the root
                                             Use the PRED relation
to N: V is not initialized
                                             Reach exclude
```



Applying the undefined query

```
begin declare x: natural, y: natural,
               z: natural:
[1] x := 3;
                                                 y is undefined
    if[2] 3 then
                                                z may be undefined
      [3] z := y + x
    else
      [4] x := 4
                                            Result:
                                                 {<5,"z">, <3,"y">}
end
```

Some Questions

- There are several additional questions:
 - In the example so far we have worked with statement numbers but how do we make a connection with the source text? (Discussed now)
 - How do we extract relations like PRED and USE from the source text? (Discussed later)



Use locations to connect with the source text

```
rel[int,str] DEFS = ...
rel[int,str] USES = ...
rel[int,int] PRED = ...
```

Use *location* instead of number

Variable occurrence in a statement

```
rel[loc,str] DEFS
rel[loc,str] USES
rel[loc,loc] PRED
rel[str, loc] OCCURS
```



Example Rstore

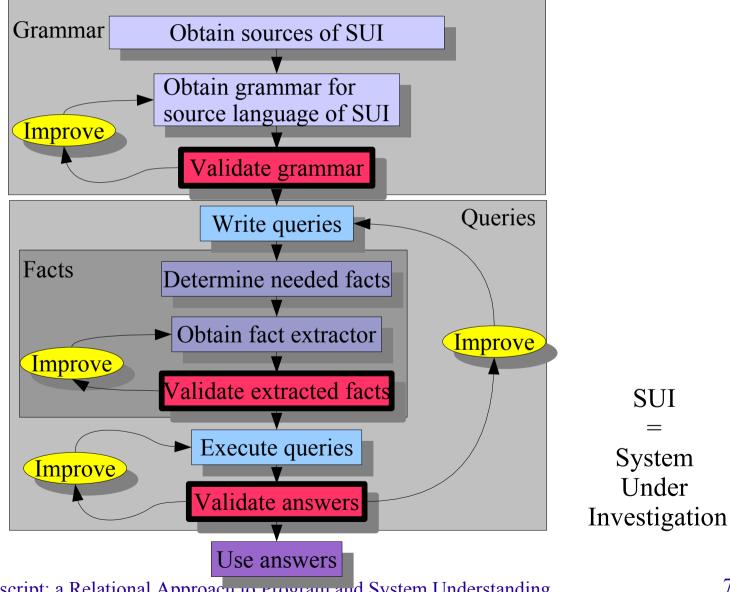
```
rstore(
  <PRED, rel[loc,loc],
          {<area-in-file("/home/paulk/.../example.pico", area(4, 2,4, 8,84, 6)),
            area-in-file("/home/paulk/.../example.pico", area(5, 2,5, 8,94, 6))>,
           <area-in-file("/home/paulk/.../example.pico", area(5, 2,5, 8, 94, 6)),</pre>
            area-in-file("/home/paulk/.../example.pico", area(6, 2,10, 4, 104, 56))>,
           ... }>,
<DEFS, {</pre>
   <OCCURS, rel[str,loc].
           {<"y", area-in-file("/home/paulk/.../example.pico", area(11, 2,11, 3,164, 1))>,
            <"z", area-in-file("/home/paulk/.../example.pico", area(11, 7,11, 8,169, 1))>,
```

Extracting Facts

- Goal: extract facts from source code and use as input for queries
- How should fact extraction be organized?
- How to write a fact extractor?



Workflow Fact Extraction





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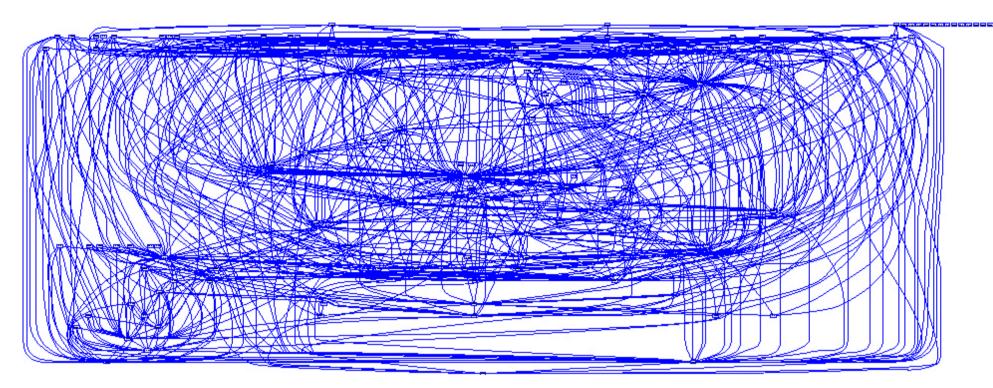
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Issues in Program Visualization

• Small graphs are nice, large graph are a disaster



(Courtesy: Arie van Deursen)

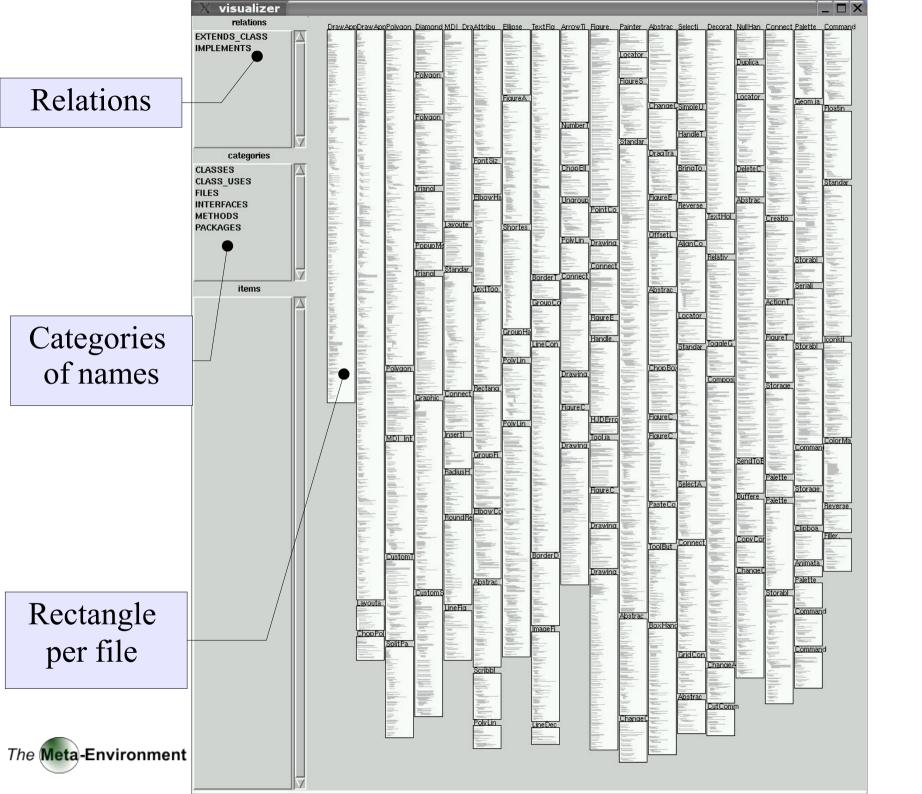


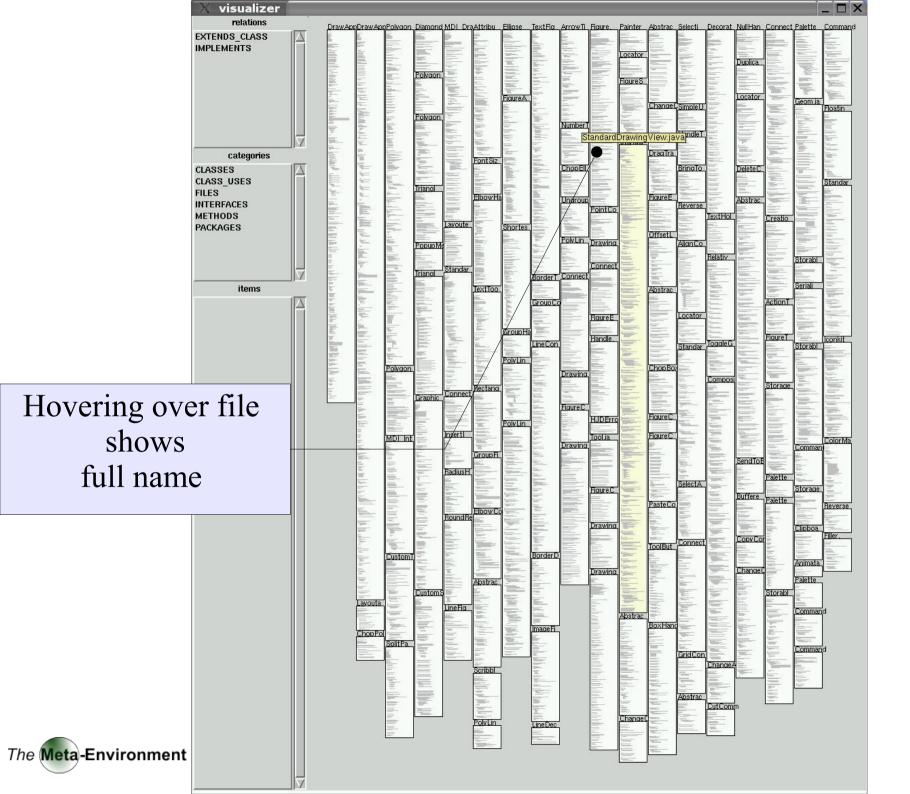
Issues in Program Visualization

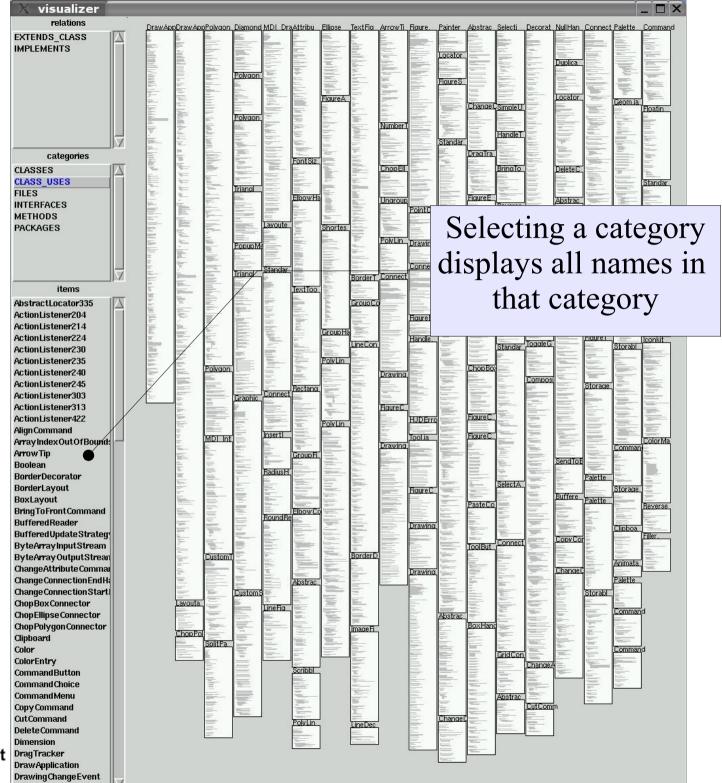
- Howto display information related to source text?
- Approach (Steven Eick): use a pixel-based image of the source text
- Over 100.000 LOC on one screen!
- Experiment: visualize an Rstore for JHotDRaw (15.000 LOC) Extraction by Hayco de Jong a

Extraction by Hayco de Jong and Taeke Kooiker (using ASF+SDF)

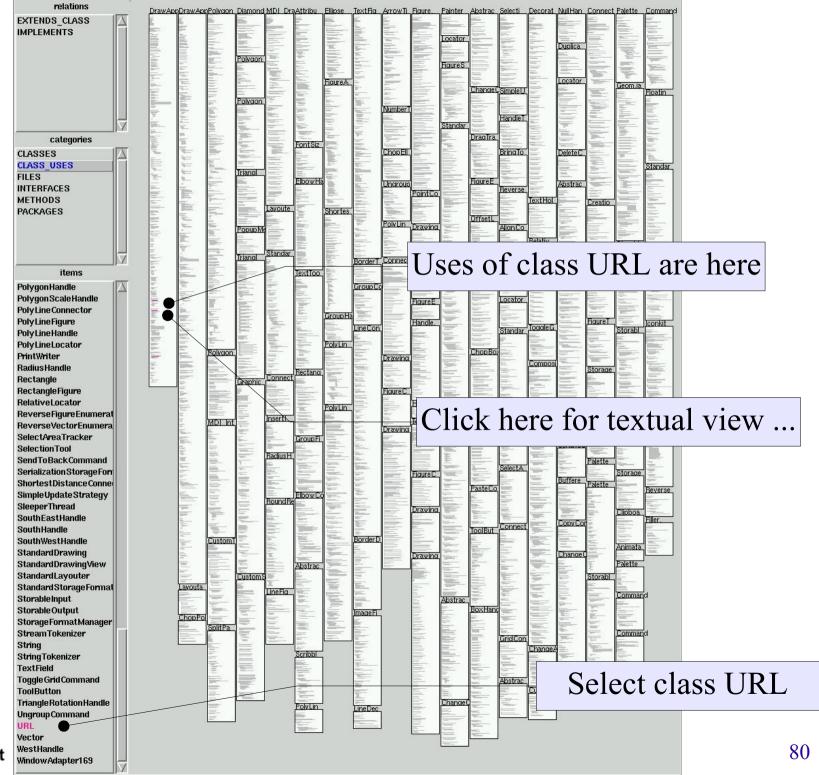














```
home/paulk/software/source/sources/CH/ifa/
   private void readFromStorableInput(String filename) {
      try {
         URL url = new URL(getCodeBase(), filename);
         InputStream stream = url.openStream();
         StorableInput input = new StorableInput(stream);
         fDrawing.release();
        fDrawing = (Drawing)input.readStorable();
fView.setDrawing(fDrawing);
     } catch (IOException e) {
        initDrawing();
showStatus("Error:"+e);
   private void readFromObjectInput(String filename) {
         URL url = new URL(getCodeBase(), filename);
         InputStream stream = url.openStream();
         ObjectInput input = new ObjectInputStream(stream);
         fDrawing.release();
        fDrawing = (Drawing)input.readObject();
         fView.setDrawing(fDrawing);
      } catch (IOException e) {
         initDrawing();
         showStatus("Error: " + e);
     } catch (ClassNotFoundException e) {
         initDrawing();
         showStatus("Class not found: " + e);
   private String guessType(String file) {
      if (file.endsWith(".draw"))
         return "storable";
      if (file.endsWith(".ser"))
         return "serialized";
```

Wrap up: Rscript

- A simple, language-independent, relational calculus
- Fully typed
- Equation solver (=> dataflow equations)
- Areas allow close link with source text
- Implementation: ASF+SDF
- IDE: rscript-meta
 - an instance of The Meta-Environment



Wrap up: Rscript

- Calls analysis
- Lifting of procedure calls to component relations
- Unitialized/unused variables
- McCabe & friends
- Clones in C code

- Dataflow analysis
 - reaching definitions
 - live variables
- Program slicing
- Java &ToolBus analysis
- Feature Descriptions/ package dependencies



Wrap up: visualization

- A lot of work to do but promising start
- Alternative pixel representations?
- Treemaps for directory structure of files?
- Colormaps for displaying metrics?
- Implementation: Tcl/Tk but may change to Swing
- Some simple visualizations are included in rscript-meta



Further reading

- P. Klint, How understanding and restructuring differ from compiling: a rewriting approach, IWPC03
- P. Klint, A tutorial introduction to Rscript on www.meta-environment.org
- www.cwi.nl/~paulk/publications/all.html

