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# Chapter 1. EASY Meta-Programming with Rascal

## *Leveraging the Extract-Analyze-Synthesize Paradigm for Meta-Programming*

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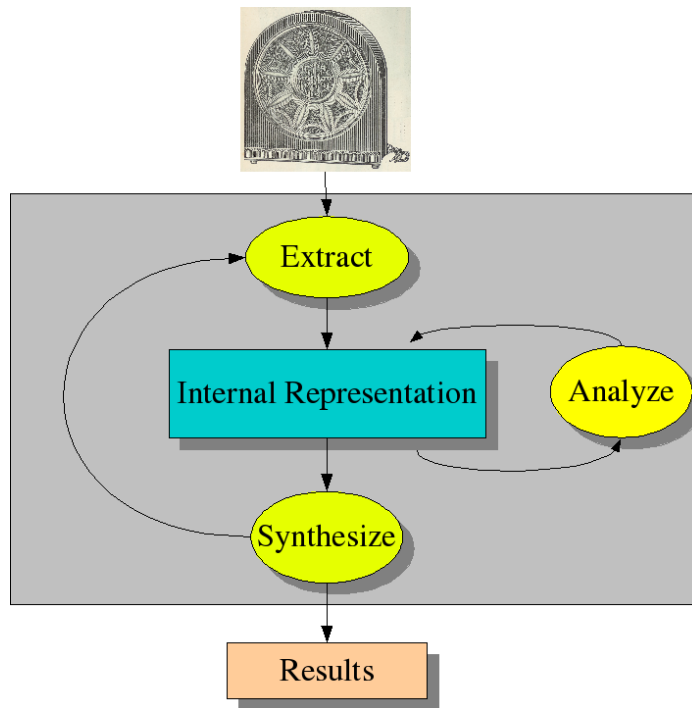
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## Why a new Language?

Rascal is a new language for *meta-programming* and is intended for writing programs that analyze or transform or generate other programs. Why do we need such a new language?

## EASY Programming

Many programming problems follow a fixed pattern. Starting with some *object-of-interest*, first relevant information is extracted from it and stored in an internal representation. This internal representation is then analyzed and used to synthesize results. If the synthesis indicates this, these steps can be repeated over and over again. These steps are shown in Figure 1.1, “EASY: the Extract-Analyze-Synthesize Paradigm” (page 3) (yes some people do like old radios!).

**Figure 1.1. EASY: the Extract-Analyze-Synthesize Paradigm**

This is an abstract view on solving programming problems, but is it uncommon? No, so let's illustrate it with a few examples.

### Example 1.1. Finding security breaches

Alice is system administrator of a large online marketplace and she is looking for security breaches in her system. The objects-of-interest are the system's log files. First relevant entries are extracted. This will include, for instance, messages from the SecureShell demon that reports failed login attempts. From each entry login name and originating IP address are extracted and put in a table (the internal representation in this example). These data are analysed by detecting duplicates and counting frequencies. Finally results are synthesized by listing the most frequently used login names and IP addresses.

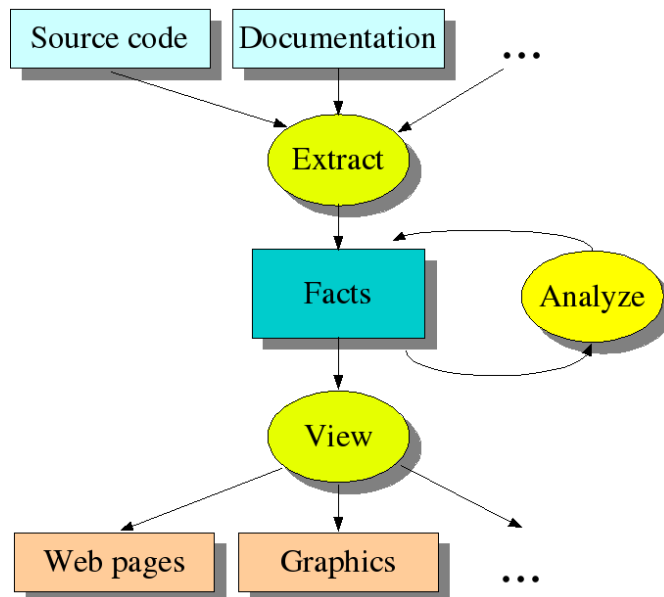
### Example 1.2. A Forensic DSL compiler

Bernd is a senior software engineer working at the Berlin headquarters of a forensic investigation lab of the German government. His daily work is to find common patterns in files stored on digital media that have been confiscated during criminal investigations. Text, audio and video files are stored in zillions of different data formats and each data format requires its own analysis technique. For each new investigation ad hoc combinations of tools are used. This makes the process very labour-intensive and error-prone. Bernd convinces his manager that designing a new domain-specific language (DSL) for forensic investigations may relieve the pressure on their lab. After designing the DSL---let's call it DERRICK---he makes an EASY implementation for it. Given a DERRICK program for a specific case under investigation, he first extracts relevant information from it and analyzes it: which media formats are relevant? Which patterns to look for? How should search results be combined? Given this new information, Java code is synthesized that uses the various existing tools and combines their results.

### Example 1.3. Renovating Financial Software

Charlotte is software engineer at a large financial institution in Paris and she is looking for options to connect an old and dusty software system to a web interface. She will need to analyze the sources of that system to understand how it can be changed to meet the new requirements. The objects-of-interest are in this case the source files, documentation, test scripts and any other available information. They have to be parsed in some way in order to extract relevant information, say the calls between various parts of the system. The call information can be represented as a binary relation between caller and callee (the internal representation in this example). This relation with 1-step calls is analyzed and further extended with 2-step calls, 3-step calls and so on. In this way call chains of arbitrary length become available. With this new information, we can synthesize results by determining the entry points of the software system, i.e. the points where calls from the outside world enter the system. Having completed this first cycle, Charlotte may be interested in which procedures can be called from the entry points and so on and so forth. Results will be typically represented as pictures that display the relationships that were found. In the case of source code analysis, a variation of our workflow scheme is quite common. It is then called the extract-analyze-view paradigm and is shown in Figure 1.2, “The extract-analyze-view paradigm” (page 4).

**Figure 1.2. The extract-analyze-view paradigm**



### Example 1.4. Finding Concurrency Errors

Daniel is concurrency researcher at one of the largest hardware manufacturers worldwide. He is working from an office in the Bay Area. Concurrency is the big issue for his company: it is becoming harder and harder to make CPUs faster, therefore more and of them are bundled on a single chip. Programming these multi-core chips is difficult and many programs that worked fine on a single CPU contain hard to detect concurrency errors due to subtle differences in the order of execution that results from executing the code on more than one CPU. Here is where Daniel enters the picture. He is working on tools for finding concurrency errors. First he extracts facts from the code that are relevant for concurrency problems and have to do with calls, threads, shared variables and locks. Next, he analyzes these facts and synthesizes an abstract model that captures the essentials of the concurrency behaviour of the program. Finally he runs a third-party verification tool with this model as input to do the actual verification.

## Rascal

With these examples in mind, you have a pretty good picture how EASY applies in different use cases. All these cases involve a form of *meta-programming*: software programs are the objects-of-interest

that are being analyzed and transformed. The Rascal language you are about to learn is designed for meta-programming following the EASY paradigm. It can be applied in domains ranging from compiler construction and implementing domain-specific languages to constraint solving and software renovation.

Since representation of data is central to the approach, Rascal provides a rich set of built-in data types. To support extraction and analysis, parsing and advanced pattern matching are provided. High-level control structures make analysis and synthesis of complex datastructures simple.

## Benefits of Rascal

Before you spend your time on studying the Rascal language it may help to first hear our elevator pitch about the main benefits offered by the language:

- **Sophisticated built-in data types** provide standard solutions for many programming problems.
- **Safety** is achieved by finding most errors even before the program is executed and by making common errors like missing initializations or wrong pointers impossible.
- **Pattern matching** is used to analyze even the most complex datastructures.
- **Syntax definitions** make it possible to define new and existing languages and to write tools for them.
- **Visiting** makes it easy to traverse datastructures and to extract information from them or to synthesize results.
- **Functions as values** permit programming styles with high re-use.
- **Generic types** allow writing functions that are applicable for many different types.
- **Local type inference** makes local variable declarations redundant.
- **Familiar syntax** in a *what-you-see is-what-you-get style* is used even for sophisticated concepts and this makes the language easy to learn and easy to use.
- **Eclipse integration** makes Rascal programming a breeze. All familiar tools are at your fingertips.

Interested? Read on!

## Aim and Scope of this Book

**Aim.** The aim of this book is to give an easy to understand but comprehensive overview of the Rascal language and to offer problem solving strategies to handle realistic, large, problems that require meta-programming. Problems may range from security analysis and model extraction to software renovation, domain-specific languages and code generation.

**Audience.** The book is intended for students, practitioners and researchers who want to solve meta-programming problems.

**Background.** Readers should have some background in computer science, software engineering or programming languages. Familiarity with several main stream programming languages and experience with larger software projects will make it easier to appreciate the relevance of the meta-programming domain that Rascal is addressing. Some familiarity with concepts like sets, relations and pattern matching is assumed.

**Scope.** The scope of the book is limited to the Rascal language and its applications but does not address implementation aspects of the language.

## Installing and Running Rascal



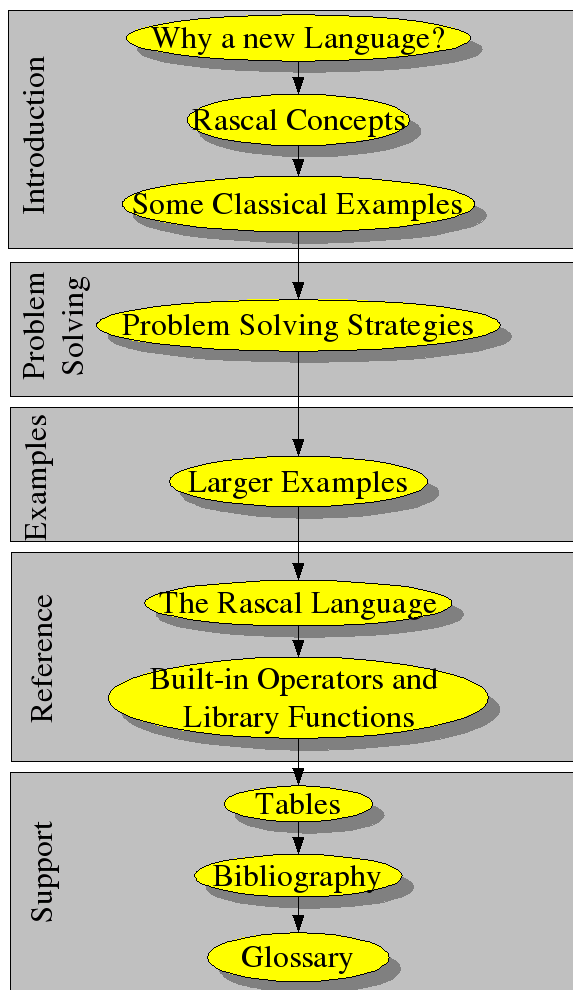
Some topics that are not yet described:

- Where to obtain a copy of Rascal.
- How to install it.
- How to run from the command line.
- Command line options
- How to run from Eclipse.
- Setting the search path.
- How to set up a Rascal project.

For the moment see <http://www.meta-environment.org> for information.

## Reading Guide

**Figure 1.3. Structure of this Book**



The structure of this book is shown in Figure 1.3, “Structure of this Book” (page 6). It consists of five parts:

- **Introduction:** gives a high-level overview of Rascal and consists of the section called “Why a new Language?” (page 2) and the section called “Rascal Concepts”(page 7) It also presents some simple examples in the section called “Some Classical Examples” (page 11).
- **Problem Solving:** describes the major problem solving strategies in Rascal's application domain, see the section called “Problem Solving Strategies” (page 14).
- **Examples:** gives a collection of larger examples, see the section called “Larger Examples” (page 20).
- **Reference:** gives a detailed description of the Rascal language, see the section called “The Rascal Language” (page 38) and all built-in operators and library functions, see the section called “Built-in Operators and Library Functions” (page 68).
- **Support:** gives tables with operators, see Table 1.35, “All Operators” (page 85) and functions, see Table 1.36, “All Functions” (page 87) a bibliography, see the section called “Bibliography” (page 90) and a glossary, see Glossary of Terminology (page 91) that explains many concepts that are used in this book and tries to make the book self-contained.

## Typographic Conventions

Rascal code fragments are always shown as a listing like this:

```
.. here is some Rascal code ...
```

Interactive sessions are shown as a screen like this (*make a more clear typographic distinction between listings and screens*):

```
rascal> Command;  
Type : Value
```

where:

- `rascal>` is the prompt of the Rascal system.



The prompt is not yet fixed. Currently, only `>` is displayed as prompt.

- **Command** is an arbitrary Rascal statement or declaration typed in by the user.
- `Type` is the type of the answer computed by Rascal.
- `Value` is the value of that answer.



For typographic reasons the output is abbreviated or slightly edited in some examples.

## Rascal Concepts

Before explaining the Rascal language in more detail, we detail our elevator pitch a bit and give you a general understanding of the concepts on which the language is based.

### Values

Rascal is a value-oriented language. This means that values are immutable and are always freshly constructed from existing parts and that sharing and aliasing problems are completely avoided. The

language does provide assignment to variables either as the result of an explicit assignment statement or as the result of a successful match.

## Data structures

Rascal provides a rich set of datatypes. From booleans, infinite precision integers and reals to strings and source code locations. From lists, (optionally labelled) tuples and sets to (optionally labelled) maps and relations. From untyped tree structures to fully typed datastructures. Syntax trees that are the result of parsing source files are represented as datatypes. There is a wealth of built-in operators and library functions available on the standard datatypes. A fragment of the abstract syntax for statements in a programming language would look as follows:

```
data STAT = asgStat(Id name, EXP exp)
           | ifStat(EXP exp, list[STAT] thenpart, list[STAT] elsepart)
           | whileStat(EXP exp, list[STAT] body)
           ;
```

## Pattern Matching

Pattern matching is *the* mechanism for case distinction and search in Rascal and a very rich pattern language is provided that includes string matching based on regular expressions, list (associative) and set (associative, commutative, identity) matching, matching of abstract pattern, and matching of concrete syntax patterns. All these forms of matching can be used in a single pattern and can be nested. Patterns may contain variables that are bound when the match is successful. Anonymous (don't care) positions are indicated by an underscore (\_). The following abstract pattern matches the while statement defined above:

```
whileStat(EXP Exp, list[STAT] Stats)
```

Variables in a pattern are either explicitly declared in the pattern itself---as done in the example---or they may be declared in the context in which the pattern occurs. So-called multi-variables in list and set patterns are declared by a \* suffix. The above pattern can then be written as

```
whileStat(EXP Exp, Stats*)
```

or, if you are not interested in the actual value of the statements as

```
whileStat(EXP Exp, _*)
```

Patterns can also be used in an explicit match operator := and can then be part of larger boolean expressions. Since a pattern match may have more than one solution, local backtracking over the alternatives of a match is provided.

## Parsing

All source code analysis projects need to extract information directly from the source code. There are two main approaches to this:

- Use regular expression to extract useful, but somewhat superficial, information. This can be achieved using regular expression patterns.
- Use syntax analysis to extract the complete, nested, structure of the source code in the form of a syntax tree.

In Rascal, we reuse the Syntax Definition Formalism (SDF) and its tooling. SDF modules define grammars and these modules can be imported in a Rascal module. These grammar rules can be applied in writing concrete patterns to match parts of parsed source code. Here is an example of the same pattern we saw above, but now in concrete form:



```
while <EXP Exp> do <list[STAT] Stats> od
```

## Enumerators

Enumerators enumerate the values in a given (finite) domain, be it the elements in a list, the substrings of a string, or all the nodes in a tree. Each value that is enumerated is first matched against a pattern before it can possibly contribute to the result of the enumerator. Examples are:

```
int x <- { 1, 3, 5, 7, 11 }  
int x <- [ 1 .. 10 ]  
asgStat(Id name, _) <- P
```

The first two produce the integer elements of a set of integers, respectively, a range of integers. The third enumerator traverses the complete program *P* (that is assumed to have a *PROGRAM* as value) and only yields statements that match the assignment pattern. Note the use of an anonymous variable at the *EXP* position in the pattern.

## Comprehensions

Rascal generalizes comprehensions in various ways. Comprehensions exist for lists, sets and maps. A comprehension consists of an expression that determines the successive elements to be included in the result and a list of enumerators and tests (boolean expressions). The enumerators produce values and the tests filter them. A standard example is

```
{ x * x | int x <- [1 .. 10], x % 3 == 0 }
```

which returns the set {9, 36, 81}, i.e., the squares of the integers in the range [ 1 .. 10 ] that are divisible by 3. A more intriguing example is

```
{name | asgStat(Id name, _) <- P}
```

which returns a list of all identifiers that occur on the lefthand side of assignment statements in program *P*.

## Control structures

Combinations of generators and boolean expressions also drive the control structures. For instance,

```
for(asgStat(Id name, _) <- P, size(Id) > 10){  
    println(Id);  
}
```

prints all identifiers in assignment statements that consist of more than 10 characters.

## Switching

The switch statement as known from C and Java is generalized: the subject value to switch on may be an arbitrary value and the cases are arbitrary patterns. When a match fails, all its side-effects are undone and when it succeeds the statements associated with that case are executed.

## Visiting

Visiting the elements of a datastructure is one of the most common operations in our domain and we give it first class support by way of visit expressions that resemble the switch statement. A visit expression consists of an expression that may yield an arbitrarily complex subject value and a number of cases. All the elements of the subject are visited and when one of the cases matches the statements associated with that case are executed. These cases may:

- cause some side effect, i.e., assign a value to local or global variables;
- execute an `insert` statement that replaces the current element;
- execute a `fail` statement that causes the match for the current case to fail (and undoing all side-effects due to the successful match itself and the execution of the statements so far).

The value of a visit expression is the original subject value with all replacements made as dictated by matching cases. The traversal order in a visit expressions can be explicitly defined by the programmer.

## Functions

Functions are explicitly declared and are fully typed. Here is an example of a function that counts the number of assignment statements in a program:

```
int countAssignments(PROGRAM P){
    int n = 0;
    visit (P){
        case asgStat(Id name, _):
            n += 1;
    }
    return n;
}
```

## Rewrite Rules

Rewrite rules are the only implicit control mechanism in the language and are used to maintain invariants during computations. For example, in a package for symbolic differentiation it is desirable to keep expressions in simplified form in order to avoid intermediate results like `sum(product(1, x), product(0, y))` that can be simplified to `x`. The following rules achieve this:

```
rule simplify1 product(1, Expression e) => e;
rule simplify2 product(Expression e, 1) => e;
rule simplify3 product(0, Expression e) => 0;
rule simplify4 product(Expression e, 0) => 0;
rule simplify5 sum(0, Expression e)      => e;
rule simplify6 sum(Expression e, 0)      => e;
```

Whenever a new expression is constructed during symbolic differentiation, these rules are applied to that expression and all its subexpressions and when a pattern at the left-hand side of a rule applies the matching subexpression is replaced by the right-hand side of the rule. This is repeated as long as any rule can be applied.

Since rewrite rules are activated automatically, one may always assume that expressions are in simplified form.

## Constraint solving

Many problems can be solved by forms of constraint solving. Rascal provides a `solve` statement that helps writing constraint solvers. A typical example is dataflow analysis where the propagation of values through a program can be described by a set equations. Their solution can be found with the `solve` statement.

## Typechecking

Rascal has a statically checked type system that prevents type errors and uninitialized variables at runtime. There are no runtime type casts as in Java and there are therefore less opportunities for

run-time errors. The language provides *higher-order*, *parametric*, *polymorphism*. A type aliasing mechanism allows documenting specific uses of a type. Builtin operators are heavily overloaded. For instance, the operator `+` is used for addition on integers and reals but also for list concatenation, set union etc.

## Execution

Following the what-you-see-is-what-you-get paradigm, control flow is completely explicit. Boolean expressions determine choices that drive the control structures. Rewrite rules form the only exception to the explicit control flow principle. Only local backtracking is provided (no suprise) in the context of boolean expressions and pattern matching; side effects are undone in case of backtracking.

## Some Classical Examples

The following simple examples will help you to grasp the main features of Rascal quickly. You can also look ahead and consult the section called “The Rascal Language” (page 38) for details of the language or the section called “Built-in Operators and Library Functions” (page 68) for specific operators or functions.

## Hello

The ubiquitous hello world program looks in Rascal as follows:

```
rascal> import IO;
void: null

rascal> println("Hello, this is my first Rascal program");
Hello, this is my first Rascal program
void: null
```

First, the library module `IO` (see the section called “`IO`” (page 73)) is imported since hello world requires printing. Next, we call `println` and proudly observe our first Rascal output!

A slightly more audacious approach is to wrap the print statement in a function and call it:

```
rascal> void hello() {
    println("Hello, this is my first Rascal program");
}
Rascal.Function::=Rascal.Function.Closure(): void hello() { ...

rascal> hello();
Hello, this is my first Rascal program
void: null
```

The summit of hello-engineering can be reached by placing all the above in a separate module:

```
module demo::Hello

public void hello() {
    println("Hello, this is my first Rascal program");
}
```

Note that we added a public modifier to the definition of `hello`, since we want it to be visible outside the `Hello` module. Using this `Hello` module is now simple:

```
rascal> import demo::Hello;
void: null
```

```
rascal> hello();
Hello, this is my first Rascal program
void: null
```



All examples in this book can be found in the demo directory of the Rascal distribution. That is why we prefix all names of examples modules with `demo :`.

## Factorial

Here is another classical example, computing the factorial function:

```
module demo::Factorial

public int fac(int N)
{
    return N <= 0 ? 1 : N * fac(N - 1);
}
```

It uses a conditional expression of the form  $Exp_1 \ ? \ Exp_2 \ : \ Exp_3$  to distinguish cases and here is how to use it:

```
rascal> import demo::Factorial;
void: null

rascal> fac(47);
int: 258623241511168180642964355153611979969197632389120000000000
```

Indeed, Rascal has infinite length integers.

## Colored Trees

Suppose we have binary trees---trees with exactly two children---that have integers as their leaves. Also suppose that our trees can have red and black nodes. Such trees can be defined as follows:

```
module demo::ColoredTrees

data ColoredTree =
    leaf(int N)
  | red(ColoredTree left, ColoredTree right)
  | black(ColoredTree left, ColoredTree right);
```

We can use them as follows:

```
rascal> import demo::ColoredTrees;
void: null
rascal> rb = red(black(leaf(1), red(leaf(2),leaf(3))));
ColoredTree: red(black(leaf(1), red(leaf(2),leaf(3))))
```

We define two operations on ColoredTrees, one to count the red nodes, and one to add all leaves:

```
// continuing module demo::ColoredTrees

public int cntRed(ColoredTree t){
    int c = 0;
    visit(t) {
        case red(_,_): c = c + 1;❶
```

```

    };
    return c;
}

public int addLeaves(ColoredTree t){
    int c = 0;
    visit(t) {
        case leaf(int N): c = c + N;❷
    };
    return C;
}

```

- ❶ Visit all the nodes of the tree and increment the counter `c` for each red node.
- ❷ Visit all nodes of the tree and add the integers in the leaf nodes.

This can be used as follows:

```

rascal> cntRed(rb);
int: 2
rascal> addLeaves(rb);
int: 6

```

A final touch to this example is to introduce green nodes and to replace all red nodes by green ones:

```

// continuing module demo::ColoredTrees

data ColoredTree = green(ColoredTree left, ColoredTree right);❶

public ColoredTree makeGreen(ColoredTree t){
    return visit(t) {
        case red(l, r) => green(l, r);    ❷
    };
}

```

- ❶ Extend the `ColoredTree` datatype with a new green constructor.
- ❷ Visit all nodes in the tree and replace red nodes by green ones. Note that the variables `l` and `r` are introduced here without a declaration.

This is used as follows:

```

rascal> makeGreen(rb);
ColoredTree: green(black(leaf(1), green(leaf(2),leaf(3))))

```

## Word Replacement

Suppose you are in the publishing business and are responsible for the systematic layout of publications. Authors do not systematically capitalize words in titles---"Word replacement" instead of "Word Replacement"--- and you want to correct this. Here is one way to solve this problem:

```

module demo::WordReplacement
import String;

public str capitalize(str word)
{
    if(/^<letter:[a-z]><rest:.*$/ := word) ❶
        return toUpperCase(letter) + rest;❷
    else
        return word;❸
}

```

}

- ❶ The function `capitalize` takes a string as input and capitalizes its first character if that is a letter. This is done using a regular expression match that anchors the match at the beginning (^), expects a single letter and assigns it to the variable `letter` (`letter:[a-z]`) followed by an arbitrary sequence of letters until the end of the string that is assigned to the variable `rest` (`<rest:.*$>`).
- ❷ If the regular expression matches we return a new string with the first letter capitalized.
- ❸ Otherwise we return the word unmodified.

The next challenge is how to capitalize all the words in a string. Here are two solutions:

```
// continuing module demo::WordReplacement

public str capAll1(str S)
{
    result = "";
    while (/^<before:\W*><word:\w+><after:.*$>/ := S) { ❶
        result = result + before + capitalize(word);
        S = after;
    }
    return result;
}

public str capAll2(str S)
{
    return visit(S){❷
        case /<word:\w+>/i ❸ => capitalize(word)❹
    };
}
```

- ❶ In the first solution `capAll1` we just loop over all the words in the string and capitalize each word. The variable `result` is used to collect the successive capitalized words.
- ❷ In the second solution we use a visit expression to visit all the substrings of `S`. Each matching case advances the substring by the length of the pattern it matches and replaces that pattern by another string. If no case matches the next substring is tried.
- ❸ The single case matches a word (note that `\w` matches a word character).
- ❹ When the case matches a word, it is replaced by a capitalized version.

We can apply this all as follows:

```
rascal> import demo::WordReplacement;
void: null

rascal> capitalize("rascal");
str: "rascal"

rascal> capAll1("rascal is great");
str: "Rascal Is Great"
```

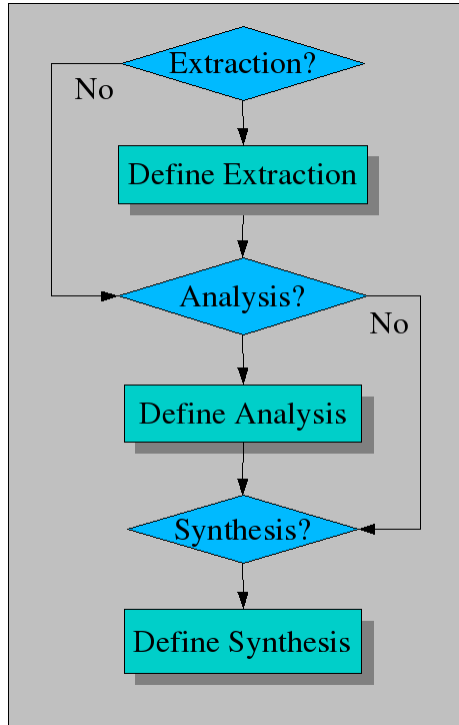
## Problem Solving Strategies

Before we study more complicated examples, it is useful to discuss some general problem solving strategies that are relevant in Rascal's application domain. The overall workflow is shown in Figure 1.4, “General 3-Phased Problem Solving Workflow” (page 15) and consists of three optional phases:

- Is extraction needed to solve the problem, then define the extraction phase, see the section called “Defining Extraction” (page 16).

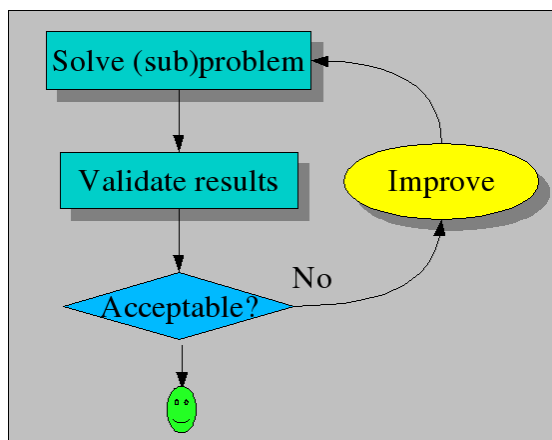
- Is analysis needed, then define the analysis phase, see the section called “Defining Analysis” (page 18).
- Is synthesis needed, then define the synthesis phase, see the section called “Defining Synthesis” (page 19).

**Figure 1.4. General 3-Phased Problem Solving Workflow**



Each phase is subject to a validation and improvement workflow as shown in Figure 1.5, “Validation and Improvement Workflow” (page 15). Each individual phase as well as the combination of phases may introduce errors and has thus to be carefully validated. In combination with the detailed strategies for each phase, this forms a complete approach for problem solving and validation using Rascal.

**Figure 1.5. Validation and Improvement Workflow**

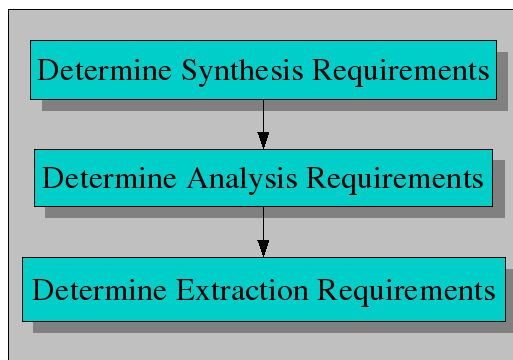


A major question in every problem solving situation is how to determine the requirements for each phase of the solution. For instance, how do we know what to extract from the source code if we do not know what the desired end results of the project are? The standard solution is to use a workflow for

requirements gathering that is the inverse of the phases needed to solve the complete problem. This is shown in Figure 1.6, “Requirements Workflow” (page 16) and amounts to the phases:

- Determine the requirements of the synthesis phase. This amounts to making an inventory of the desired results of the whole project and may include generated source code, abstract models, or visualizations.
- Determine the requirements of the analysis phase. Once these results of the synthesis phase are known, it is possible to list the analysis results that are needed to synthesize desired results. Possible results of the analysis phase include type information, structural information of the original source.
- Determine the requirements of the extraction phase. As a last step, one can make an inventory of the facts that have to be extracted to form the starting point for the analysis phase. Typical facts include method calls, inheritance relations, control flow graphs, usage patterns of specific library functions or language constructs.

**Figure 1.6. Requirements Workflow**



When these requirements have been established, it becomes much easier to actually carry out the project using the three phases of Figure 1.4, “General 3-Phased Problem Solving Workflow” (page 15).

## Defining Extraction

How can we extract facts from the *System under Investigation* (SUI) that we are interested in? The extraction workflow is shown in Figure 1.7, “Extraction Workflow” (page 18) and consists of the following steps:

- First and foremost we have to determine which facts we need. This sounds trivial, but it is not. The problem is that we have to anticipate which facts will be needed in the next---not yet defined---analysis phase. A common approach is to use look-ahead and to sketch the queries that are likely to be used in the analysis phase and to determine which facts are needed for them. Start with extracting these facts and refine the extraction phase when the analysis phase is completely defined.
- If relevant facts are already available (and they are reliable!) then we are done. This may happen when you are working on a system that has already been analyzed by others.
- Otherwise you need the source code of the SUI. This requires:
  - Checking that all sources are available (and can be compiled by the host system on which they are usually compiled and executed). Due to missing or unreliable configuration management on the original system this may be a labour-intensive step that requires many iterations.
  - Determining in which languages the sources are written. In larger systems it is common that three or more different languages are being used.



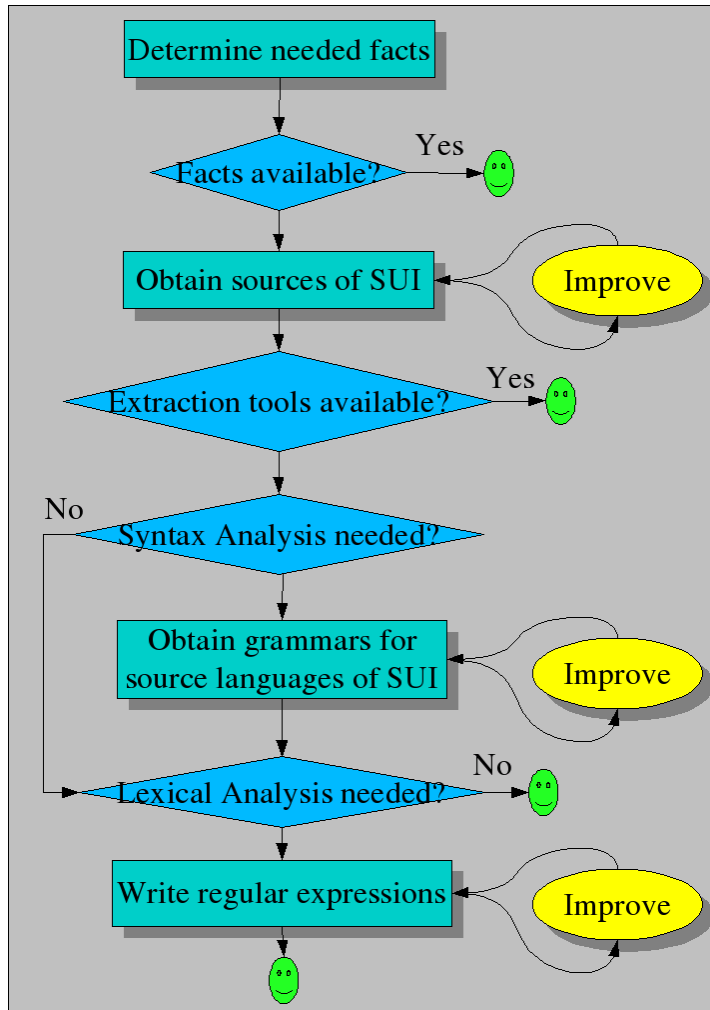
- If there are reliable third-party extraction tools available for this language mix, then we only have to apply them and we are done. Here again, validation is needed that the extracted facts are as expected.
- The extraction may require syntax analysis. This is the case when more structural properties of the source code are needed such as the flow-of-control, nesting of declarations, and the like. There two approaches here:
  - Use a third-party parser, convert the source code to parse trees and do the further processing of these parse trees in Rascal. The advantage is that the parser can be re-used, the disadvantage is that data conversion is needed to adapt the generated parse tree to Rascal. Validate that the parser indeed accepts the language the SUI is written in, since you will not be the first who is biting by the language dialect monster when it turns out that the SUI uses a local variant that slightly deviates from a mainstream language.
  - Use an existing SDF definition of the source language or write your own definition. In both cases you can profit from Rascal's seamless integration with SDF. Be aware, however, that writing a grammar for a non-trivial language is a major undertaking and may require weeks to month of work. Whatever approach you choose, validate that the result.
- The extraction phase may only require lexical analysis. This happens when more superficial, textual, facts have to be extracted like procedure calls, counts of certain statement statement and the like. Use Rascal's full regular expression facilities to do the lexical analysis.

It may happen that the facts extracted from the source code are *wrong*. Typical error classes are:

- Extracted facts are *wrong*: the extracted facts incorrectly state that procedure P calls procedure Q but this is contradicted by a source code inspection.
- Extracted facts are *incomplete*: the inheritance between certain classes in Java code is missing.

The strategy to validate extracted facts differ per case but here are three strategies:

- Postprocess the extracted facts (using Rascal, of course) to obtain trivial facts about the source code such as total lines of source code and number of procedures, classes, interfaces and the like. Next validate these trivial facts with tools like `wc` (word and line count), `grep` (regular expression matching) and others.
- Do a manual fact extraction on a small subset of the code and compare this with the automatically extracted facts.
- Use another tool on the same source and compare results whenever possible. A typical example is a comparison of a call relation extracted with different tools.

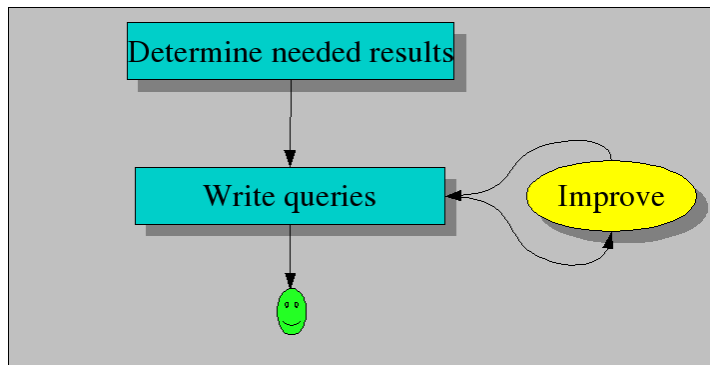
**Figure 1.7. Extraction Workflow**

## Defining Analysis

The analysis workflow is shown in Figure 1.8, “Analysis Workflow” (page 19) and consists of two steps:

- Determine the results that are needed for the synthesis phase.
- Write the Rascal code to perform the analysis. This may amount to:
  - Reordering extracted facts to make them more suitable for the synthesis phase.
  - Enriching extracted facts. Examples are computing transitive closures of extracted facts (e.g., A may call B in one or more calls), or performing data reduction by abstracting away details (i.e., reducing a program to a finite automaton).
  - Combining enriched, extracted, facts to create new facts.

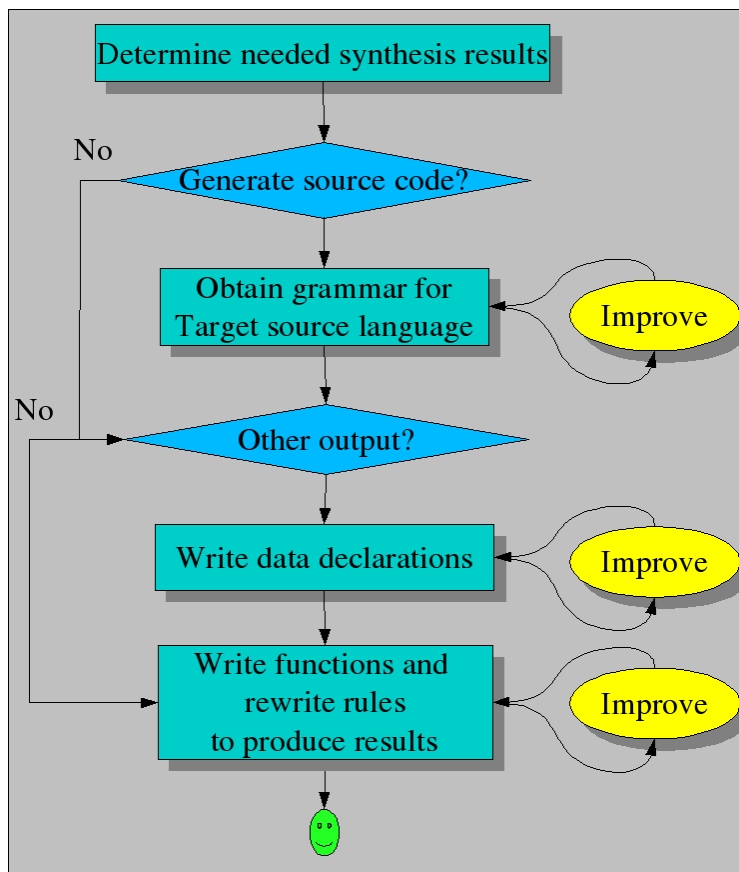
As before, validate, validate and validate the results of analysis. Essentially the same approach can be used as for validating the facts. Manual checking of answers on random samples of the SUI may be mandatory. It also happens frequently that answers inspire new queries that lead to new answers, and so on.

**Figure 1.8. Analysis Workflow**

## Defining Synthesis

Results are synthesized as shown in Figure 1.9, “Synthesis Workflow” (page 20) This consists of the following steps:

- Determine the results of the synthesis phase. Wide range of results is possible including:
  - Generated source code.
  - Generated abstract representations, like finite automata or other formal models that capture properties of the SUI.
  - Generated data for visualizations that will be used by visualization tools.
- If source is to be generated, we need a grammar of the target source language. Note that this approach guarantees the generation of syntactically correct source code as opposed to code generation using print statements
- If other output is needed (e.g., an automaton or other formal structure) write data declarations to represent that output.
- Finally, write functions and rewrite rules that generate the desired results.

**Figure 1.9. Synthesis Workflow**

## Larger Examples

Now we will have a closer look at some larger applications of Rascal. We start with a call graph analysis in the section called “Call Graph Analysis” (page 20) and then continue with the analysis of the component structure of an application in the section called “Analyzing the Component Structure of an Application” (page 24) and of Java systems in the section called “Analyzing the Structure of Java Systems” (page 25). Next we move on to the detection of initialized variables in the section called “Finding Uninitialized and Unused Variables in a Program” (page 26). As an example of computing code metrics, we describe the calculation of McCabe's cyclomatic complexity in the section called “McCabe Cyclomatic Complexity”. Several examples of dataflow analysis follow in the section called “Dataflow Analysis” (page 28). A description of program slicing concludes the chapter (the section called “Program Slicing” (page 35)).

## Call Graph Analysis

Suppose a mystery box ends up on your desk. When you open it, it contains a huge software system with several questions attached to it:

- How many procedure calls occur in this system?
- How many procedures does it contain?
- What are the entry points for this system, i.e., procedures that call others but are not called themselves?
- What are the leaves of this application, i.e., procedures that are called but do not make any calls themselves?
- Which procedures call each other indirectly?

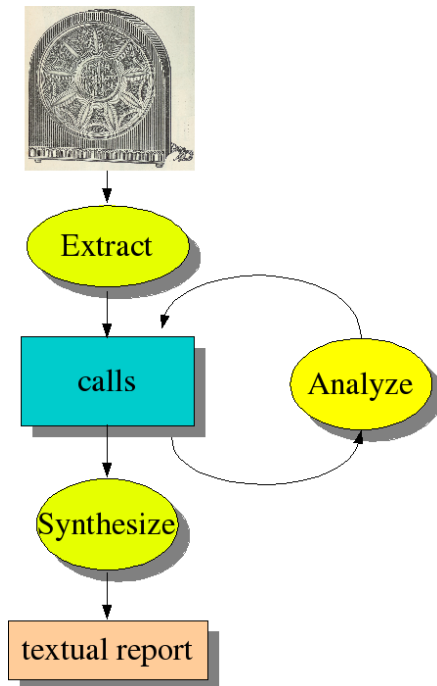
- Which procedures are called directly or indirectly from each entry point?
- Which procedures are called from all entry points?

There are now two possibilities. Either you have this superb programming environment or tool suite that can immediately answer all these questions for you or you can use Rascal.

## Preparations

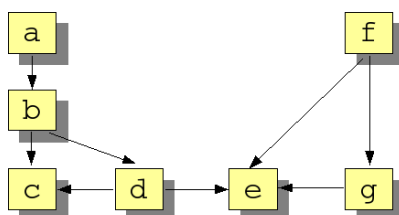
To illustrate this process consider the workflow in Figure 1.10, “Workflow for analyzing mystery box” (page 21) First we have to extract the calls from the source code. Rascal is very good at this, but to simplify this example we assume that this call graph has already been extracted. Also keep in mind that a real call graph of a real application will contain thousands and thousands of calls. Drawing it in the way we do later on in Figure 1.11, “Graphical representation of the `calls` relation”(page 21) makes no sense since we get a uniformly black picture due to all the call dependencies. After the extraction phase, we try to understand the extracted facts by writing queries to explore their properties. For instance, we may want to know *how many calls* there are, or *how many procedures*. We may also want to enrich these facts, for instance, by computing who calls who in more than one step. Finally, we produce a simple textual report giving answers to the questions we are interested in.

**Figure 1.10. Workflow for analyzing mystery box**



Now consider the call graph shown in Figure 1.11, “Graphical representation of the `calls` relation” (page 21) This section is intended to give you a first impression what can be done with Rascal. Please return to this example when you have digested the detailed description of Rascal in the section called “The Rascal Language” (page 38) and the section called “Built-in Operators and Library Functions” (page 68).

**Figure 1.11. Graphical representation of the `calls` relation**



Rascal supports basic data types like integers and strings which are sufficient to formulate and answer the questions at hand. However, we can gain readability by introducing separately named types for the items we are describing. First, we introduce therefore a new type `proc` (an alias for strings) to denote procedures:

```
rascal> alias proc = str;
void: null
```

Suppose that the following facts have been extracted from the source code and are represented by the relation `Calls`:

```
rascal> rel[proc, proc] Calls =
  { <"a", "b">, <"b", "c">, <"b", "d">, <"d", "c">, <"d", "e">,
    <"f", "e">, <"f", "g">, <"g", "e">
  };
rel[proc,proc]: { <"a", "b">, <"b", "c">, <"b", "d">, <"d", "c">,
                  <"d", "e">, <"f", "e">, <"f", "g">, <"g", "e"> }
```

This concludes the preparatory steps and now we move on to answer the questions.

## Questions

### How many procedure calls occur in this system?

To determine the numbers of calls, we simply determine the number of tuples in the `Calls` relation, as follows. First, we need the Relation library (described in the section called “Relation” (page 77) so we import it:

```
rascal> import Relation;
void: null
```

next we describe a new variable and calculate the number of tuples:

```
rascal> nCalls = size(Calls);
int: 8
```

The library function `size` determines the number of elements in a set or relation and is explained in the section called “Relation” (page 77). In this example, `nCalls` will get the value 8.

### How many procedures are contained in it?

We get the number of procedures by determining which names occur in the tuples in the relation `Calls` and then determining the number of names:

```
rascal> procs = carrier(Calls);
set[proc]: { "a", "b", "c", "d", "e", "f", "g" }

rascal> nprocs = size(procs);
int: 7
```

The built-in function `carrier` determines all the values that occur in the tuples of a relation. In this case, `procs` will get the value `{ "a", "b", "c", "d", "e", "f", "g" }` and `nprocs` will thus get value 7. A more concise way of expressing this would be to combine both steps:

```
rascal> nprocs = size(carrier(Calls));
int: 7
```

### What are the entry points for this system?

The next step in the analysis is to determine which *entry points* this application has, i.e., procedures which call others but are not called themselves. Entry points are useful since they define the external

interface of a system and may also be used as guidance to split a system in parts. The `top` of a relation contains those left-hand sides of tuples in a relation that do not occur in any right-hand side. When a relation is viewed as a graph, its `top` corresponds to the root nodes of that graph. Similarly, the `bottom` of a relation corresponds to the leaf nodes of the graph. See the section called “Graph” (page 71) for more details. Using this knowledge, the entry points can be computed by determining the `top` of the `Calls` relation:

```
rascal> import Graph;
void: null

rascal> entryPoints = top(Calls);
set[proc]: {"a", "f"}
```

In this case, `entryPoints` is equal to `{"a", "f"}`. In other words, procedures “a” and “f” are the entry points of this application.

### What are the leaves of this application?

In a similar spirit, we can determine the *leaves* of this application, i.e., procedures that are being called but do not make any calls themselves:

```
rascal> bottomCalls = bottom(Calls);
set[proc]: {"c", "e"}
```

In this case, `bottomCalls` is equal to `{"c", "e"}`.

### Which procedures call each other indirectly?

We can also determine the *indirect calls* between procedures, by taking the transitive closure of the `Calls` relation:

```
rascal> closureCalls = Calls+;
rel[proc, proc]: {<"a", "b">, <"b", "c">, <"b", "d">, <"d", "c">,
                  <"d", "e">, <"f", "e">, <"f", "g">, <"g", "e">,
                  <"a", "c">, <"a", "d">, <"b", "e">, <"a", "e">}
```

### Which procedures are called directly or indirectly from each entry point?

We now know the entry points for this application (“a” and “f”) and the indirect call relations. Combining this information, we can determine which procedures are called from each entry point. This is done by indexing `closureCalls` with appropriate procedure name. The index operator yields all right-hand sides of tuples that have a given value as left-hand side. This gives the following:

```
rascal> calledFromA = closureCalls["a"];
set[proc]: {"b", "c", "d", "e"}
```

and

```
rascal> calledFromF = closureCalls["f"];
set[proc]: {"e", "g"}
```

### Which procedures are called from all entry points?

Finally, we can determine which procedures are called from both entry points by taking the intersection of the two sets `calledFromA` and `calledFromF`:

```
rascal> commonProcs = calledFromA & calledFromF;
set[proc]: {"e"}
```

In other words, the procedures called from both entry points are mostly disjoint except for the common procedure "e".

## Wrap-up

These findings can be verified by inspecting a graph view of the calls relation as shown in Figure 1.11, “Graphical representation of the `calls` relation” (page 21) Such a visual inspection does *not* scale very well to large graphs and this makes the above form of analysis particularly suited for studying large systems.



Add a screen dump here of the Eclipse browser.

## Analyzing the Component Structure of an Application

A frequently occurring problem is that we know the call relation of a system but that we want to understand it at the component level rather than at the procedure level. If it is known to which component each procedure belongs, it is possible to *lift* the call relation to the component level as proposed in [Kri99]. Actual lifting, amounts to translating each call between procedures by a call between components. This described in the following module:

```
module demo::Lift

alias proc = str;
alias comp = str;

public rel[comp,comp] lift(rel[proc,proc] aCalls,
                           rel[proc,comp] aPartOf){
    return { <C1, C2> | <proc P1, proc P2> <- aCalls,
                       <comp C1, comp C2> <- aPartOf[P1] *
                                               aPartOf[P2]
    };
}
```

Let's now apply this. First import the above module, and define a call relation and a partof relation:

```
rascal> import demo::Lift;
void: null

rascal> Calls = {<"main", "a">, <"main", "b">, <"a", "b">,
                <"a", "c">, <"a", "d">, <"b", "d">
                };
rel[str,str] : {<"main", "a">, <"main", "b">, <"a", "b">,
                <"a", "c">, <"a", "d">, <"b", "d">
                }

rascal> Components = {"Appl", "DB", "Lib"};
set[str] : {"Appl", "DB", "Lib"}

rascal> PartOf = {<"main", "Appl">, <"a", "Appl">,
                  <"b", "DB">, <"c", "Lib">, <"d", "Lib">};
rel[str,str] : {<"main", "Appl">, <"a", "Appl">,
                  <"b", "DB">, <"c", "Lib">, <"d", "Lib">}
```

The lifted call relation between components is now obtained by:

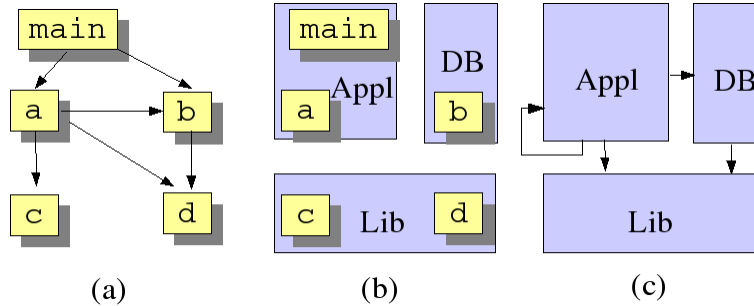
```
rascal> ComponentCalls = lift(Calls, PartOf);
rel[str,str] : {<"DB", "Lib">, <"Appl", "Lib">, <"Appl", "DB">,
```



```
<"Appl", "Appl">}
```

The relevant relations for this example are shown in Figure 1.12, “(a) Calls; (b) PartOf; (c) ComponentCalls.” (page 25).

**Figure 1.12. (a) Calls; (b) PartOf; (c) ComponentCalls.**



## Analyzing the Structure of Java Systems

Now we consider the analysis of Java systems (inspired by [BNL03]). Suppose that the type `class` is defined as follows

```
alias class = str;
```

and that the following relations are available about a Java application:

- `rel[class, class] CALL`: If  $\langle C_1, C_2 \rangle$  is an element of `CALL`, then some method of  $C_2$  is called from  $C_1$ .
- `rel[class, class] INHERITANCE`: If  $\langle C_1, C_2 \rangle$  is an element of `INHERITANCE`, then class  $C_1$  either extends class  $C_2$  or  $C_1$  implements interface  $C_2$ .
- `rel[class, class] CONTAINMENT`: If  $\langle C_1, C_2 \rangle$  is an element of `CONTAINMENT`, then one of the attributes of class  $C_1$  is of type  $C_2$ .

To make this more explicit, consider the class `LocatorHandle` from the `JHotDraw` application (version 5.2) as shown here:

```
package CH.ifa.draw.standard;

import java.awt.Point;
import CH.ifa.draw.framework.*;
/**
 * A LocatorHandle implements a Handle by delegating the
 * location requests to a Locator object.
 */
public class LocatorHandle extends AbstractHandle {
    private Locator    fLocator;
    /**
     * Initializes the LocatorHandle with the given Locator.
     */
    public LocatorHandle(Figure owner, Locator l) {
        super(owner);
        fLocator = l;
    }
    /**
     * Locates the handle on the figure by forwarding the request
     * to its figure.
     */
}
```

```

public Point locate() {
    return fLocator.locate(owner());
}

```

It leads to the addition to the above relations of the following tuples:

- To CALL the pairs `<"LocatorHandle", "AbstractHandle">` and `<"LocatorHandle", "Locator">` will be added.
- To INHERITANCE the pair `<"LocatorHandle", "AbstractHandle">` will be added.
- To CONTAINMENT the pair `<"LocatorHandle", "Locator">` will be added.

Cyclic structures in object-oriented systems makes understanding hard. Therefore it is interesting to spot classes that occur as part of a cyclic dependency. Here we determine cyclic uses of classes that include calls, inheritance and containment. This is achieved as follows:

```

rel[class, class] USE = CALL + CONTAINMENT + INHERITANCE;
set[str] ClassesInCycle =
    {C1 | <class C1, class C2> <- USE+, C1 == C2};

```

First, we define the USE relation as the union of the three available relations CALL, CONTAINMENT and INHERITANCE. Next, we consider all pairs `<C1, C2>` in the transitive closure of the USE relation such that `C1` and `C2` are equal. Those are precisely the cases of a class with a cyclic dependency on itself. Probably, we do not only want to know which classes occur in a cyclic dependency, but we also want to know which classes are involved in such a cycle. In other words, we want to associate with each class a set of classes that are responsible for the cyclic dependency. This can be done as follows.

```

rel[class, class] USE = CALL + CONTAINMENT + INHERITANCE;
set[class] CLASSES = carrier(USE);
rel[class, class] USETRANS = USE+;
rel[class, set[class]] ClassCycles =
    {<C, USETRANS[C]> | class C <- CLASSES, <C, C> in USETRANS };

```

First, we introduce two new shorthands: CLASSES and USETRANS. Next, we consider all classes `C` with a cyclic dependency and add the pair `<C, USETRANS[C]>` to the relation `ClassCycles`. Note that `USETRANS[C]` is the right image of the relation `USETRANS` for element `C`, i.e., all classes that can be called transitively from class `C`.

## Finding Uninitialized and Unused Variables in a Program

Consider the following program in the toy language Pico: (This is an extended version of the example presented earlier in [Kli03].)

```

[ 1] begin declare x : natural, y : natural,
[ 2]           z : natural, p : natural;
[ 3]  x := 3;
[ 4]  p := 4;
[ 5]  if q then
[ 6]      z := y + x
[ 7]  else
[ 8]      x := 4
[ 9]  fi;
[10]  y := z
[11] end

```

Inspection of this program learns that some of the variables are being used before they have been initialized. The variables in question are `q` (line 5), `y` (line 6), and `z` (line 10). It is also clear that

variable `p` is initialized (line 4), but is never used. How can we automate these kinds of analysis? Recall from the section called “EASY Programming” (page 2) that we follow the Extract-Analyze-SYNthesize paradigm to approach such a problem. The first step is to determine which elementary facts we need about the program. For this and many other kinds of program analysis, we need at least the following:

- The *control flow graph* of the program. We represent it by a graph `PRED` (for predecessor) which relates each statement with each predecessors.
- The *definitions* of each variable, i.e., the program statements where a value is assigned to the variable. It is represented by the relation `DEFS`.
- The *uses* of each variable, i.e., the program statements where the value of the variable is used. It is represented by the relation `USES`.

In this example, we will use line numbers to identify the statements in the program. Assuming that there is a tool to extract the above information from a program text, we get the following for the above example:

```
module demo::Uninit
import Graph;

alias expr = int;
alias varname = str;

public expr ROOT = 1;

public graph[expr] PRED = { <1,3>, <3,4>, <4,5>, <5,6>, <5,8>,
                           <6,10>, <8,10> };

public rel[varname,expr] DEFS = { <"x", 3>, <"p", 4>, <"z", 6>,
                                  <"x", 8>, <"y", 10> };

public rel[varname, expr] USES = { <"q", 5>, <"y", 6>, <"x", 6>,
                                  <"z", 10> };
```

This concludes the extraction phase. Next, we have to enrich these basic facts to obtain the initialized variables in the program. So, when is a variable `V` in some statement `S` initialized? If we execute the program (starting in `ROOT`), there may be several possible execution path that can reach statement `S`. All is well if *all* these execution path contain a definition of `V`. However, if one or more of these path do *not* contain a definition of `V`, then `V` may be uninitialized in statement `S`. This can be formalized as follows:

```
// module demo::Unit continued
public rel[varname,expr] UNINIT =
  { <V,E> | <varname V, expr E> <- USES,
          E in reachX({ROOT}, DEFS[V], PRED)
  };
```

We analyze this definition in detail:

- `<varname V, expr E> : USES` enumerates all tuples in the `USES` relation. In other words, we consider the use of each variable in turn.
- `E in reachX({ROOT}, DEFS[V], PRED)` is a test that determines whether statement `S` is reachable from the `ROOT` without encountering a definition of variable `V`.
  - `{ROOT}` represents the initial set of nodes from which all path should start.
  - `DEFS[V]` yields the set of all statements in which a definition of variable `V` occurs. These nodes form the exclusion set for `reachX`: no path will be extended beyond an element in this set.

- PRED is the relation for which the reachability has to be determined.
- The result of `reachX({ROOT}, DEFS[V], PRED)` is a set that contains all nodes that are reachable from the ROOT (as well as all intermediate nodes on each path).
- Finally, `E in reachX({ROOT}, DEFS[V], PRED)` tests whether expression E can be reached from the ROOT.
- The net effect is that UNINIT will only contain pairs that satisfy the test just described.

When we execute the resulting Rascal code (i.e., the declarations of ROOT, PRED, DEFS, USES and UNINIT), we get as value for UNINIT:

```
rascal> import demo::Uninit;
void: null

rascal> UNINIT;
rel[varname,expr]: {<"q", 5>, <"y", 6>, <"z", 10>}
```

and this is in concordance with the informal analysis given at the beginning of this example.

As a bonus, we can also determine the *unused* variables in a program, i.e., variables that are defined but are used nowhere. This is done as follows:

```
// module demo::Unit continued

public set[varname] UNUSED = domain(DEFS) - domain(USES)
```

Taking the domain of the relations DEFS and USES yields the variables that are defined, respectively, used in the program. The difference of these two sets yields the unused variables, in this case { "p" }.

## McCabe Cyclomatic Complexity

The *cyclomatic complexity* of a program is defined as  $e - n + 2$ , where  $e$  and  $n$  are the number of edges and nodes in the control flow graph, respectively. It was proposed by McCabe [McC76] as a measure of program complexity. Experiments have shown that programs with a higher cyclomatic complexity are more difficult to understand and test and have more errors. It is generally accepted that a program, module or procedure with a cyclomatic complexity larger than 15 is *too complex*. Essentially, cyclomatic complexity measures the number of decision points in a program and can be computed by counting all if statement, case branches in switch statements and the number of conditional loops. Given a control flow in the form of a predecessor graph `graph[&T] PRED` between elements of arbitrary type &T, the cyclomatic complexity can be computed in Rascal as follows:

```
module demo::McCabe
import Graph;

public int cyclomaticComplexity(graph[&T] PRED){
    return size(PRED) - size(carrier(PRED)) + 2;
}
```

The number of edges  $e$  is equal to the number of tuples in PRED. The number of nodes  $n$  is equal to the number of elements in the carrier of PRED, i.e., all elements that occur in a tuple in PRED.

## Dataflow Analysis

*Dataflow analysis* is a program analysis technique that forms the basis for many compiler optimizations. It is described in any text book on compiler construction, e.g. [ASU86]. The goal of dataflow analysis is to determine the effect of statements on their surroundings. Typical examples are:

- Dominators (the section called “Dominators” (page 29): which nodes in the flow dominate the execution of other nodes?
- Reaching definitions (the section called “Reaching Definitions” (page 31): which definitions of variables are still valid at each statement?
- Live variables (the section called “Live Variables” (page 34): of which variables will the values be used by successors of a statement?
- Available expressions: an expression is available if it is computed along each path from the start of the program to the current statement.

## Dominators

A node  $d$  of a flow graph *dominates* a node  $n$ , if every path from the initial node of the flow graph to  $n$  goes through  $d$  [ASU86] (Section 10.4). Dominators play a role in the analysis of conditional statements and loops. The function `dominators` that computes the dominators for a given flow graph `PRED` and an entry node `ROOT` is defined as follows:

```
module demo::Dominators
import Set;
import Relation;
import Graph;

public rel[&T, set[&T]] dominators(rel[&T,&T] PRED, &T ROOT)
{
    set[&T] VERTICES = carrier(PRED);
    return { <V, (VERTICES - {V, ROOT}) - reachX(PRED, {ROOT}, {V})>
            | &T V <- VERTICES
            };
}
```

First, the auxiliary set `VERTICES` (all the statements) is computed. The relation `DOMINATES` consists of all pairs  $\langle S, \{S_1, \dots, S_n\} \rangle$  such that

- $S_i$  is not an initial node or equal to  $S$ .
- $S_i$  cannot be reached from the initial node without going through  $S$ .

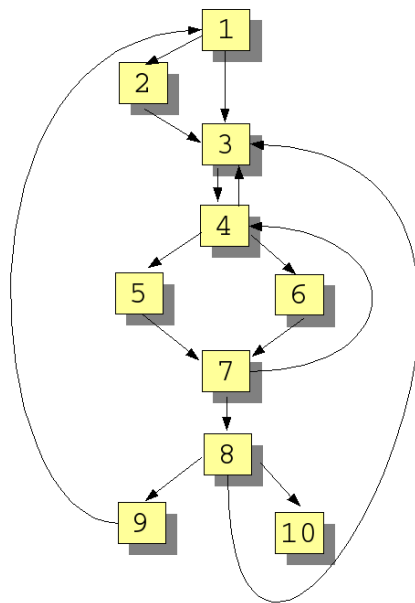
First import the above module and consider the sample flow graph `PRED`:

```
rascal> import demo::Dominators;
void: null

rascal> rel[int,int] PRED = {
<1,2>, <1,3>,
<2,3>,
<3,4>,
<4,3>,<4,5>, <4,6>,
<5,7>,
<6,7>,
<7,4>,<7,8>,
<8,9>,<8,10>,<8,3>,
<9,1>,
<10,7>
};

rel[int,int]: { <1,2>, <1,3>, ...
```

It is illustrated in Figure 1.13, “Flow graph” (page 30)

**Figure 1.13. Flow graph**

(a)

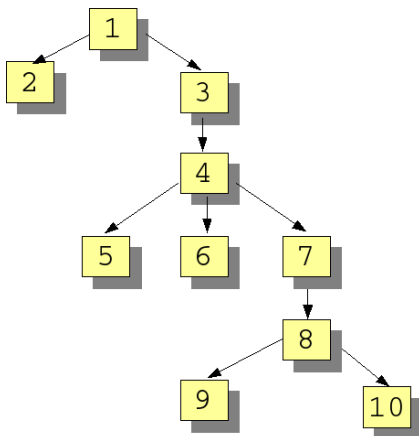
The result of applying dominators to it is as follows:

```

rascal> dominators(PRED);
rel[int,int]: {<1, {2, 3, 4, 5, 6, 7, 8, 9, 10}>,
<2, {}>,
<3, {4, 5, 6, 7, 8, 9, 10}>,
<4, {5, 6, 7, 8, 9, 10}>,
<5, {}>,
<6, {}>,
<7, {8, 9, 10}>,
<8, {9, 10}>,
<9, {}>,
<10, {}>}}

```

The resulting *dominator tree* is shown in Figure 1.14, “Dominator tree” (page 30) The dominator tree has the initial node as root and each node  $d$  in the tree only dominates its descendants in the tree.

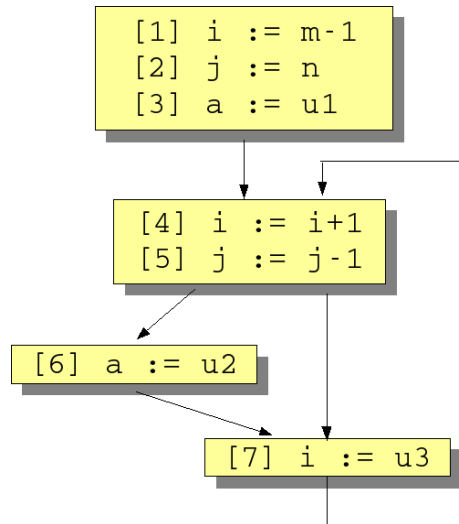
**Figure 1.14. Dominator tree**

(b)

## Reaching Definitions

We illustrate the calculation of reaching definitions using the example in Figure 1.15, “Flow graph for various dataflow problems” (page 31) which was inspired by [ASU86] (Example 10.15).

**Figure 1.15. Flow graph for various dataflow problems**



We assume the following basicdefinitions to represent information about the program:

```

module demo::ReachingDefs

import Relation;
import IO;

public alias stat = int;
public alias var = str;
public alias def = tuple[stat, var];
public alias use = tuple[stat, var];

public rel[stat, def] definition(rel[stat, var] DEFS){
  return {<S, <S, V> | <stat S, var V> <- DEFS};
}

public rel[stat, def] use(rel[stat, var] USES){
  return {<S, <S, V> | <stat S, var V> <- USES};
}

```

Let's use the following values to represent our example:

```

rascal> rel[stat, stat] PRED = { <1, 2>, <2, 3>, <3, 4>, <4, 5>, <5, 6>,
                                <5, 7>, <6, 7>, <7, 4>
                                };
rel[stat, stat]: { <1, 2>, <2, 3>, ...

rascal> rel[stat, var] DEFS = { <1, "i">, <2, "j">, <3, "a">,
                                <4, "i">, <5, "j">, <6, "a">,
                                <7, "i">
                                };
rel[stat, var]: { <1, "i">, <2, "j">, ...

rascal> rel[stat, var] USES = { <1, "m">, <2, "n">, <3, "u1">,

```

```

                                <4, "i">, <5, "j">, <6, "u2">,
                                <7, "u3">
                                };
rel[stat,var]: { <1, "m">, <2, "n">, ...

```

For convenience, we have introduced above a notion `def` that describes that a certain statement defines some variable and we revamp the basic relations into a more convenient format using this new type and the auxiliary functions `definition` and `use`:

```

rascal> definition(DEFS);
rel[stat,def]: { <1, <1, "i">>, <2, <2, "j">>, <3, <3, "a">>,
                  <4, <4, "i">>, <5, <5, "j">>, <6, <6, "a">>,
                  <7, <7, "i">>
                }

rascal> use(USES);
rel[stat,def]: { <1, <1, "m">>, <2, <2, "n">>, <3, <3, "u1">>,
                  <4, <4, "i">>, <5, <5, "j">>, <6, <6, "u2">>,
                  <7, <7, "u3">>
                }

```

Now we are ready to define an important new relation `KILL`. `KILL` defines which variable definitions are undone (killed) at each statement and is defined by the following function `kill`:

```

// continuing module demo::ReachingDefs

public rel[stat,def] kill(rel[stat,var] DEFS) {
    return {<S1, <S2, V>> | <stat S1, var V> <- DEFS,
                           <stat S2, V> <- DEFS,
                           S1 != S2};
}

```

In this definition, all variable definitions are compared with each other, and for each variable definition all *other* definitions of the same variable are placed in its kill set. In the example, `KILL` gets the value

```

rascal> kill(DEFS);
rel[stat,def]:
{ <1, <4, "i">>, <1, <7, "i">>, <2, <5, "j">>, <3, <6, "a">>,
  <4, <1, "i">>, <4, <7, "i">>, <5, <2, "j">>, <6, <3, "a">>,
  <7, <1, "i">>, <7, <4, "i">>
}

```

and, for instance, the definition of variable `i` in statement 1 kills the definitions of `i` in statements 4 and 7.

After these preparations, we are ready to formulate the reaching definitions problem in terms of two relations `IN` and `OUT`. `IN` captures all the variable definitions that are valid at the entry of each statement and `OUT` captures the definitions that are still valid after execution of each statement. Intuitively, for each statement `S`, `IN[S]` is equal to the union of the `OUT` of all the predecessors of `S`. `OUT[S]`, on the other hand, is equal to the definitions generated by `S` to which we add `IN[S]` minus the definitions that are killed in `S`. Mathematically, the following set of equations captures this idea for each statement:

$$IN[S] = \text{UNION}_{P \text{ in predecessor of } S} OUT[P]$$

$$OUT[S] = DEF[S] + (IN[S] - KILL[S])$$

This idea can be expressed in Rascal quite literally:

```

public rel[stat, def] reachingDefinitions(rel[stat,var] DEFS,
                                           rel[stat,stat] PRED){

```



```

set[stat] STATEMENT = carrier(PRED);
rel[stat,def] DEF = definition(DEFS);
rel[stat,def] KILL = kill(DEFS);

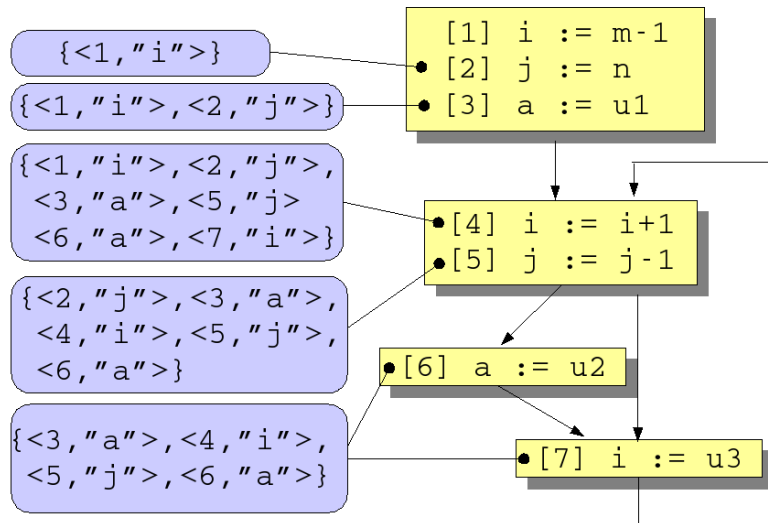
// The set of mutually recursive dataflow equations that has
// to be solved:

with
  rel[stat,def] IN = {};
  rel[stat,def] OUT = DEF;
solve {
  IN = {<S, D> | int S <- STATEMENT,
                stat P <- predecessor(PRED,S),
                def D <- OUT[P]};
  OUT = {<S, D> | int S <- STATEMENT,
                 def D <- DEF[S] + (IN[S] - KILL[S])};
};
return IN;
}

```

First, the relations IN and OUT are declared and initialized. Next, two equations are given that resemble the mathematical equations given above. Note the use of the library function predecessor to obtain the predecessor of a statement for a given control flow graph.

**Figure 1.16. Reaching definitions for flow graph in Figure 1.15, “Flow graph for various dataflow problems” (page 31)**



For our running example (Figure 1.16, “Reaching definitions for flow graph in Figure 1.15, “Flow graph for various dataflow problems”” (page 33)) the results are as follows (see Figure 1.16, “Reaching definitions for flow graph in Figure 1.15, “Flow graph for various dataflow problems”” (page 33)). Relation IN has as value:

```

{
  <2, <1, "i">, <3, <2, "j">, <3, <1, "i">, <4, <3, "a">,
  <4, <2, "j">, <4, <1, "i">, <4, <7, "i">, <4, <5, "j">,
  <4, <6, "a">, <5, <4, "i">, <5, <3, "a">, <5, <2, "j">,
  <5, <5, "j">, <5, <6, "a">, <6, <5, "j">, <6, <4, "i">,
  <6, <3, "a">, <6, <6, "a">, <7, <5, "j">, <7, <4, "i">,
  <7, <3, "a">, <7, <6, "a">
}

```

If we consider statement 3, then the definitions of variables *i* and *j* from the preceding two statements are still valid. A more interesting case are the definitions that can reach statement 4:

- The definitions of variables *a*, *j* and *i* from, respectively, statements 3, 2 and 1.
- The definition of variable *i* from statement 7 (via the backward control flow path from 7 to 4).
- The definition of variable *j* from statement 5 (via the path 5, 7, 4).
- The definition of variable *a* from statement 6 (via the path 6, 7, 4).

Relation OUT has as value:

```
{ <1, <1, "i">>, <2, <2, "j">>, <2, <1, "i">>, <3, <3, "a">>,
  <3, <2, "j">>, <3, <1, "i">>, <4, <4, "i">>, <4, <3, "a">>,
  <4, <2, "j">>, <4, <5, "j">>, <4, <6, "a">>, <5, <5, "j">>,
  <5, <4, "i">>, <5, <3, "a">>, <5, <6, "a">>, <6, <6, "a">>,
  <6, <5, "j">>, <6, <4, "i">>, <7, <7, "i">>, <7, <5, "j">>,
  <7, <3, "a">>, <7, <6, "a">>
}
```

Observe, again for statement 4, that all definitions of variable *i* are missing in OUT[4] since they are killed by the definition of *i* in statement 4 itself. Definitions for *a* and *j* are, however, contained in OUT[4]. The result of reaching definitions computation is illustrated in Figure 1.16, “Reaching definitions for flow graph in Figure 1.15, “Flow graph for various dataflow problems” (page 33)

We will use the function `reachingDefinitions` later on in the section called “Program Slicing” (page 35) when defining program slicing.

## Live Variables

The live variables of a statement are those variables whose value will be used by the current statement or some successor of it. The mathematical formulation of this problem is as follows:

$$IN[S] = USE[S] + (OUT[S] - DEF[S])$$

$$OUT[S] = \text{UNION}_{S' \text{ in successor of } S} IN[S']$$

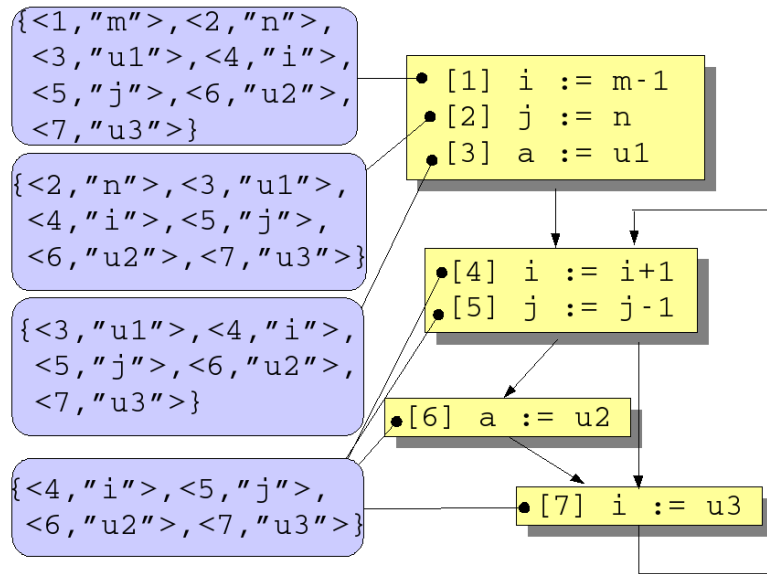
The first equation says that a variable is live coming into a statement if either it is used before redefinition in that statement or it is live coming out of the statement and is not redefined in it. The second equation says that a variable is live coming out of a statement if and only if it is live coming into one of its successors.

This can be expressed in Rascal as follows:

```
public rel[stat,def] liveVariables(rel[stat,var] DEFS,
                                   rel[stat, var] USES,
                                   rel[stat,stat] PRED){
  set[stat] STATEMENT = carrier(PRED);
  rel[stat,def] DEF = definition(DEFS);
  rel[stat,def] USE = use(USES);
  with
    rel[stat,def] LIN = {};
    rel[stat,def] LOUT = DEF;
  solve {
    LIN = { <S, D> | stat S <- STATEMENT,
                    def D <- USE[S] + (LOUT[S] - (DEF[S])) };
    LOUT = { <S, D> | stat S <- STATEMENT,
                  stat Succ <- successor(PRED,S),
                  def D <- LIN[Succ] };
  }
  return LIN;
}
```

The results of live variable analysis for our running example are illustrated in Figure 1.17, “Live variables for flow graph in Figure 1.15, “Flow graph for various dataflow problems”” (page 35).

**Figure 1.17. Live variables for flow graph in Figure 1.15, “Flow graph for various dataflow problems” (page 31)**



## Program Slicing

Program slicing is a technique proposed by Weiser [Wei84] for automatically decomposing programs in parts by analyzing their data flow and control flow. Typically, a given statement in a program is selected as the *slicing criterion* and the original program is reduced to an independent subprogram, called a *slice*, that is guaranteed to represent faithfully the behavior of the original program at the slicing criterion. An example will illustrate this:

[ 1] read(n)	[ 1] read(n)	[ 1] read(n)
[ 2] i := 1	[ 2] i := 1	[ 2] i := 1
[ 3] sum := 0	[ 3] sum := 0	
[ 4] product := 1		[ 4] product := 1
[ 5] while i<= n do	[ 5] while i<= n do	[ 5] while i<= n do
begin	begin	begin
[ 6] sum := sum + i	[ 6] sum := sum + i	
[ 7] product :=		[ 7] product :=
product * i		product * i
[ 8] i := i + 1	[ 8] i := i + 1	[ 8] i := i + 1
end	end	end
[ 9] write(sum)	[ 9] write(sum)	
[10] write(product)		[10] write(product)
<b>(a)</b> Sample program	<b>(b)</b> Slice for statement [9]	<b>(c)</b> Slice for statement [10]

The initial program is given as (a). The slice with statement [9] as slicing criterion is shown in (b): statements [4] and [7] are irrelevant for computing statement [9] and do not occur in the slice. Similarly, (c) shows the slice with statement [10] as slicing criterion. This particular form of slicing is called *backward slicing*. Slicing can be used for debugging and program understanding, optimization and more. An overview of slicing techniques and applications can be found in [Tip95]. Here we will explore a relational formulation of slicing adapted from a proposal in [JR94]J. The basic ingredients of the approach are as follows:

- We assume the relations PRED, DEFS and USES as before.
- We assume an additional set CONTROL-STATEMENT that defines which statements are control statements.

- To tie together dataflow and control flow, three auxiliary variables are introduced:
  - The variable TEST represents the outcome of a specific test of a conditional statement. The conditional statement defines TEST and all statements that are control dependent on this conditional statement will use TEST.
  - The variable EXEC represents the potential execution dependence of a statement on some conditional statement. The dependent statement defines EXEC and an explicit (control) dependence is made between EXEC and the corresponding TEST.
  - The variable CONST represents an arbitrary constant.

The calculation of a (backward) slice now proceeds in six steps:

- Compute the relation `rel[use,def]` use-def that relates all uses to their corresponding definitions. The function `reaching-definitions` as shown earlier in the section called “Reaching Definitions” (page 31) does most of the work.
- Compute the relation `rel[def,use]` def-use-per-stat that relates the *internal* definitions and uses of a statement.
- Compute the relation `rel[def,use]` control-dependence that links all EXECs to the corresponding TESTs.
- Compute the relation `rel[use,def]` use-control-def combines use/def dependencies with control dependencies.
- After these preparations, compute the relation `rel[use,use]` USE-USE that contains dependencies of uses on uses.
- The backward slice for a given slicing criterion (a use) is now simply the projection of USE-USE for the slicing criterion.

This informal description of backward slicing can now be expressed in Rascal:

```
module demo::Slicing

import Set;
import Relation;
import demo::ReachingDefs;
import demo::Dominators;
import UnitTest;

set[use] BackwardSlice(set[stat] CONTROLSTATEMENT,
                      rel[stat,stat] PRED,
                      rel[stat,var] USES,
                      rel[stat,var] DEFS,
                      use Criterion) {

    rel[stat, def] REACH = reachingDefinitions(DEFS, PRED);

    // Compute the relation between each use and corresponding
    // definitions: use_def

    rel[use,def] use_def =
    {<<S1,V>, <S2,V>> | <stat S1, var V> <- USES,
                      <stat S2, V> <- REACH[S1]};

    // Internal dependencies per statement

    rel[def,use] def_use_per_stat =
```

```

    {<<S,V1>, <S,V2>> | <stat S, var V1> <- DEFS,
                        <S, var V2> <- USES}
    +
    {<<S,V>, <S,"EXEC">> | <stat S, var V> <- DEFS}
    +
    {<<S,"TEST">,<S,V>> | stat S <- CONTROLSTATEMENT,
                        <S, var V> <- domainR(USES, {S})});

// Control dependence: control-dependence

rel[stat, set[stat]] CONTROLDOMINATOR =
domainR(dominators(PRED, 1), CONTROLSTATEMENT);

rel[def,use] control_dependence =
{ <<S2, "EXEC">,<S1,"TEST">>
  | <stat S1, stat S2> <- CONTROLDOMINATOR};

// Control and data dependence: use-control-def

rel[use,def] use_control_def = use_def + control_dependence;
rel[use,use] USE_USE = (use_control_def o def_use_per_stat)*;

return USE_USE[Criterion];
}

```

Let's apply this to the example from the start of this section and assume the following:

```

rascal> import demo::Slicing;
void: null

rascal> rel[stat,stat] PRED = { <1,2>, <2,3>, <3,4>, <4,5>, <5,6>,
                              <5,9>, <6,7>, <7,8>,<8,5>, <8,9>,
                              <9,10>
                              };
rel[stat,stat]: {<1,2>, ...

rascal> rel[stat,var] DEFS = { <1, "n">, <2, "i">, <3, "sum">,
                              <4,"product">, <6, "sum">,
                              <7, "product">, <8, "i">
                              };
rel[stat,var]: {<1, "n">, ...

rascal> rel[stat,var] USES = { <5, "i">, <5, "n">, <6, "sum">,
                              <6,"i">, <7, "product">, <7, "i">,
                              <8, "i">, <9, "sum">,
                              <10, "product">
                              };
rel[stat,var]: { <5, "i"> ...

rascal> set[int] CONTROL-STATEMENT = { 5 };
set[int]: {5}

rascal> BackwardsSlice(CONTROL-STATEMENT,
                      PRED, USES, DEFS, <9, "sum">);
set[use]: { <1, "EXEC">, <2, "EXEC">, <3, "EXEC">, <5, "i">,
            <5, "n">, <6, "sum">, <6, "i">, <6, "EXEC">,
            <8, "i">, <8, "EXEC">, <9, "sum"> }

```

Take the domain of this result and we get exactly the statements in (b) of the example.

# The Rascal Language

A Rascal program consists of one or more modules. Each module may import other modules and declare data types, variables, functions or rewrite rules. We now describe the basic ingredients of Rascal in more detail:

- Types and values, see the section called “Types and Values” (page 38).
- Declarations, see the section called “Declarations” (page 43).
- Expressions, see the section called “Expressions” (page 48).
- Statements, see the section called “Statements” (page 62).

## Types and Values

### Elementary Types and Values

**Void.** Void stands for *nothing* and is represented by the type `void`. It is a type without any values.

**Value.** Value stands for all possible Rascal values and is represented by the type `value`. This type is a container for all other types and does not have any values itself.

**Boolean.** The Booleans are represented by the type `bool` which has two values: `true` and `false`.

**Integer.** The integer values are represented by the type `int` and are written as usual, e.g., 0, 1, or 123. They can be arbitrarily large.

**Real.** The real values are represented by the type `real` and are written as usual, e.g., 1.5, or 3.14e-123. They can have arbitrary size and precision.

**String.** The string values are represented by the type `str` and consist of character sequences surrounded by double quotes. e.g., "a" or "a\nlong\tstring".

String literals permit *interpolation* of variable values: when `<X>` occurs inside a string literal, the value of the variable `X` is converted to a string that replaces `<X>`. As a consequence, the character `<` has to be escaped as `\<` in string literals.

**Location.** Location values are represented by the type `loc` and serve as text coordinates in a specific source file. It is very handy to associate a source code location which extracted facts.

Source locations have the following syntax:

```
loc(Url?offset=O&length=L&begin=BL,BC&end=EL,EC)
```

where:

- `Url` is an arbitrary URL.
- `O` and `L` are integer expressions giving the offset of this location to the begin of file, respectively, its length.
- `BL` and `BC` are integers expressions giving the begin line and begin column.
- `EL` and `EC` are integers expressions giving the end line and end column.

Locations should *always* be generated automatically but for the curious here is an example:

```
loc(file:///home/paulk/pico.trm?offset=0&length=1&begin=2,3&end=4,5)
```

The elements of a location value can be accessed and modified using the standard mechanism of field selection and field assignment. The corresponding field names are:

- url
- offset
- length
- beginLine, beginColumn
- endLine, endColumn.

## List, Set, Map, Tuple, and Relation

**List.** A list is an ordered sequence of values and has the following properties:

- All elements have the same static type.
- The order of the elements matters.
- The list may contain the same element more than once.

Lists are represented by the type `list[T]`, where  $T$  is an arbitrary type. Examples are `list[int]`, `list[tuple[int,int]]` and `list[list[str]]`. Lists are denoted by a list of elements, separated by comma's and enclosed in bracket as in  $[E_1, E_2, \dots, E_n]$ , where the  $E_n$  ( $1 \leq i \leq n$ ) are expressions that yield the desired element type. For example,

- `[1, 2, 3]` is of type `list[int]`,
- `{<1,10>, <2,20>, <3,30>}` is of type `set[tuple[int,int]]`,
- `[1, "b", 3]` is of type `list[value]`,
- `[<"a",10>, <"b",20>, <"c",30>]` is of type `list[tuple[str,int]]`, and
- `[[ "a", "b"], [ "c", "d", "e"]]` is of type `list[list[str]]`.



`[1, 2, 3]` and `[3, 2, 1]` are different lists.



`[1, 2, 3]` and `[1, 2, 3, 1]` are also different lists.

When variables of type list occur inside a list, their elements are automatically spliced into the surrounding list. This can be prevented by surrounding them with extra `[` and `]` brackets.

**Range.** For lists of integers, a special shorthand exists to describe ranges of integers:

- `[F .. L]` ranges from first element  $F$  upto (and including) last element  $L$  with increments of 1.
- `[F, S, .. E]`, ranges from first element  $F$ , second element  $S$  upto (and including) last element  $E$  with increments of  $S - F$ .

**Set.** A set is an unordered sequence of values and has the following properties:

- All elements have the same static type.
- The order of the elements does not matter.
- A set contains an element only once. In other words, duplicate elements are eliminated and no matter how many times an element is added to a set, it will occur in it only once.

Sets are represented by the type `set[T]`, where  $T$  is an arbitrary type. Examples are `set[int]`, `set[tuple[int,int]]` and `set[set[str]]`. Sets are denoted by a list of elements, separated by comma's and enclosed in braces as in  $\{E_1, E_2, \dots, E_n\}$ , where the  $E_n$  ( $1 \leq i \leq n$ ) are expressions that yield the desired element type. For example,

- $\{1, 2, 3\}$  is of type `set[int]`,
- $\{<1,10>, <2,20>, <3,30>\}$  is of type `set[tuple[int,int]]`,
- $\{1, "b", 3\}$  is of type `set[value]`,
- $\{<"a",10>, <"b",20>, <"c",30>\}$  is of type `set[tuple[str,int]]`, and
- $\{\{"a", "b"\}, \{"c", "d", "e"\}\}$  is of type `set[set[str]]`.



$\{1, 2, 3\}$  and  $\{3, 2, 1\}$  are identical sets.



$\{1, 2, 3\}$  and  $\{1, 2, 3, 1\}$  are also identical sets.

In a similar fashion as with lists, sets variables are automatically spliced into a surrounding set. This can be prevented by surrounding them with extra `{}` and `}` brackets.

**Map.** A map is a set of key : value pairs and has the following properties:

- Key and value may have different static types.
- A key can only occur once.

Maps are represented by the type `map[T1, T2]`, where  $T_1$  and  $T_2$  are arbitrary types. Examples are `map[int,int]`, and `map[str,int]`. Maps are denoted by a list of pairs, separated by comma's and enclosed in parentheses as in  $(K_1:V_1, \dots, K_n:V_n)$ , where the  $K_n$  ( $1 \leq i \leq n$ ) are expressions that yield the keys of the map and  $V_n$  ( $1 \leq i \leq n$ ) are expressions that yield the values for each key. Maps resemble functions rather than relations in the sense that only a single value can be associated with each key. For example,

- $(\text{"pear"} : 1, \text{"apple"} : 3, \text{"banana"} : 0)$  is of type `map[str,int]`.

**Tuple.** A tuple is a sequence of elements with the following properties:

- Each element in a tuple (may) have a different type.
- Each element of a tuple may have a label that can be used to select that element of the tuple.

Tuples are represented by the type `tuple[T1L1, T2L2, ..., TnLn]`, where  $T_1, T_2, \dots, T_n$  are arbitrary types and  $L_1, L_2, \dots, L_n$  are optional labels. An example of a tuple type is `tuple[str name, int freq]`. Examples are:

- $<1, 2>$  is of type `tuple[int, int]`,
- $<1, 2, 3>$  is of type `tuple[int, int, int]`,
- $<"a", 3>$  is of type `tuple[str name, int freq]`.

**Relation.** A relation is a set of elements with the following property:

- All elements have the same static tuple type.

Relations are thus nothing more than sets of tuples, but since they are used so often we provide a shorthand notation for them. Relations are represented by the type `rel[T1L1, T2L2, ..., TnLn]`.



$L_n$ ], where  $T_1, T_2, \dots, T_n$  are arbitrary types and  $L_1, L_2, \dots, L_n$  are optional labels. It is a shorthand for  $\text{set}[\text{tuple}[T_1 L_1, T_2 L_2, \dots, T_n L_n]]$ . Examples are  $\text{rel}[\text{int}, \text{str}]$  and  $\text{rel}[\text{int}, \text{set}[\text{str}]]$ . An  $n$ -ary relations with  $m$  tuples is denoted by  $\{\langle E_{11}, E_{12}, \dots, E_{1n} \rangle, \langle E_{21}, E_{22}, \dots, E_{2n} \rangle, \dots, \langle E_{m1}, E_{m2}, \dots, E_{mn} \rangle\}$ , where the  $E_{ij}$  are expressions that yield the desired element type. For example,  $\{\langle 1, "a" \rangle, \langle 2, "b" \rangle, \langle 3, "c" \rangle\}$  is of type  $\text{rel}[\text{int}, \text{str}]$ . Examples are:

- $\{\langle 1, 10 \rangle, \langle 2, 20 \rangle, \langle 3, 30 \rangle\}$  is of type  $\text{rel}[\text{int}, \text{int}]$  (yes indeed, you saw this same example before and then we gave  $\text{set}[\text{tuple}[\text{int}, \text{int}]]$  as its type; remember that these types are interchangeable.),
- $\{\langle "a", 10 \rangle, \langle "b", 20 \rangle, \langle "c", 30 \rangle\}$  is of type  $\text{rel}[\text{str}, \text{int}]$ , and
- $\{\langle "a", 1, "b" \rangle, \langle "c", 2, "d" \rangle\}$  is of type  $\text{rel}[\text{str}, \text{int}, \text{str}]$ .

## Alias Type

Everything can be expressed using the elementary types and values that are provided by Rascal. However, for the purpose of documentation and readability it is sometimes better to use a descriptive name as type indication, rather than an elementary type. The alias declaration

```
alias Name = Type;
```

states that *Name* can be used everywhere instead of the already defined type *Type*. Both types are thus structurally equivalent. For instance,

```
alias ModuleId = str;
alias Frequency = int;
```

introduces two new type names *ModuleId* and *Frequency*, both an alias for the type *str*. The use of type aliases is a good way to document your intentions. Another example is an alias definition for a graph containing integer nodes:

```
alias Graph = rel[int, int];
```

## Data Type

In ordinary programming languages record types or classes exist to introduce a new type name for a collection of related, named, values and to provide access to the elements of such a collection through their name. In Rascal, data declarations provide this facility. The type declaration

```
data Name = Pat1 | Pat1 | ...
```

introduces a new datatype *Name* and *Pat<sub>1</sub>, Pat<sub>2</sub>, ...* are prefix patterns describing the variants of the datatype. For instance,

```
data Bool = T | F | conj(Bool L, Bool R) | disj(Bool L, Bool R);
```

defines the datatype *Bool* that contains various constants (T and F) and constructor functions *conj* and *disj*.

## Type Parameters and Parameterized types

In addition to the types that we have already discussed, a type may also be a *type parameter* of the form

```
&Name
```

A type parameter may occur at every syntactic position where a type is required and turns an ordinary type into a parameterized type. Parameterized types are used to define polymorphic functions and data types, i.e., functions and data types that are applicable for more than one type. Type parameters are

bound to an actual type when the function or data type is applied and further uses of the type parameter are consistently replaced by the actual type.

The following syntactic positions are *binding occurrences* for type parameters:

- Type parameters in the type declaration of a function are bound to the types of the actual parameters in the call of that function. Type parameters that occur in the body of the function are replaced by the corresponding actual type.
- The left-hand side of an alias. The type parameters are bound when the alias is used and occurrences of type parameters in the right hand side are replaced by corresponding actual types.
- The alternatives of a data type. Binding and replacement is identical to that of function declarations.

All other occurrences of type parameters are *using occurrences*. The following rules apply:

- When the same type parameter is used at different binding occurrences it should be bound to the same actual type.
- For every using occurrence of a type parameter there should be a binding occurrence of a type parameter with the same name.

We refer to the section called “Function Declaration” (page 45) for a full description of function declaration, but here

The following function `swap` returns a tuple in which its arguments are swapped and can be applied to arbitrary values in a type safe manner:

```
rascal> tuple[&B, &A] swap(&A a, &B b) { return <b, a>; }
void: null

rascal> swap(1,2);
tuple[int,int]: <2,1>

rascal> swap("abc", 3);
tuple[int,str]: <3, "abc">
```

Observe that the type parameters that occur in the return type should occur in the formal parameter types of the function.

An alias definition may be parameterized. So we can generalize graphs as follows:

```
rascal> alias Graph[&Node] = rel[&Node, &Node];
void: null

rascal> Graph[int] GI = {<1,2>, <3,4>, <4,1>};
Graph[int] : {<1,2>, <3,4>, <4,1>}

rascal> Graph[str] GS = {<"a", "b">, <"c","d">, <"d", "a">};
Graph[str] : {<"a", "b">, <"c","d">, <"d", "a">}
```

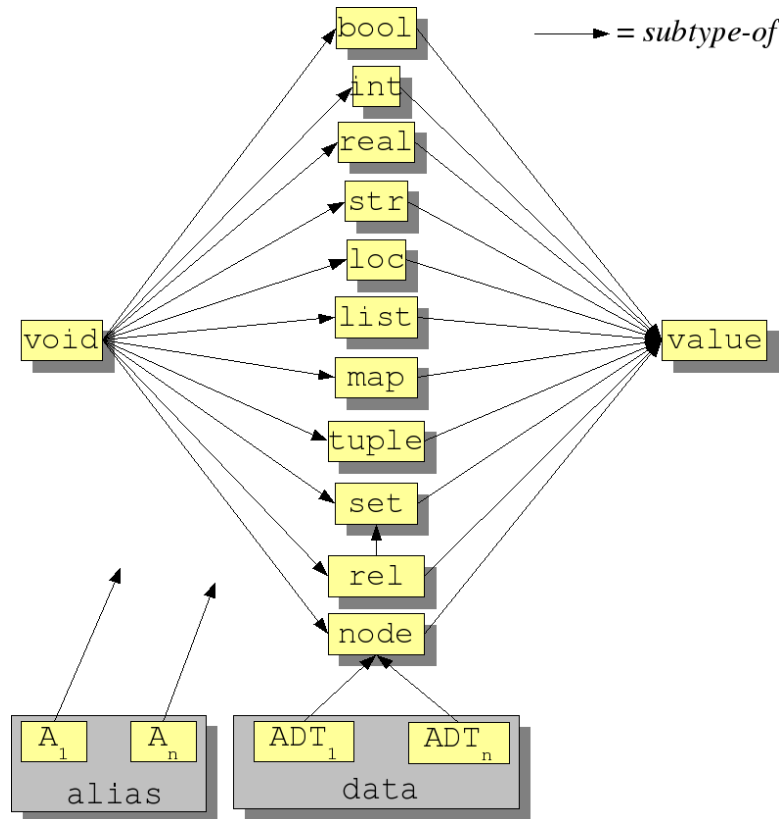
Of course, the type variables that are used in the type in the left part of the alias declaration should occur as parameters in the right part of the definition and vice versa.

## Typing

Rascal is based on static typing, this means that as many errors and inconsistencies as possible are spotted before the program is executed. The types introduced earlier are ordered in a so-called *type lattice* shown in the section called “Typing” (page 42). The arrows describe a *subtype-of* relation between types. The type `void` is the *smallest* type and is included in all other types and the type `value` is the *largest* type that includes all other types. We also see that `rel` is a subtype of `set` and that

each ADT is a subtype of `node`. Finally, each alias is structurally equivalent to one or more specific other types.

**Figure 1.18. The Rascal Type Lattice**



## Declarations

### Rascal Program

A Rascal program consists of a number of modules that may import each other.

### Module

A module declaration has the following form:

```
module Name
Imports;
Declaration1;
...
Declarationn;
```

and consists of a *Name* and zero or more imports of other modules (the section called “Import” (page 44)) and declarations for

- Data type, see the section called “Data Type” (page 41).
- Alias, see the section called “Alias Type” (page 41).
- Variable, see the section called “Variable Declaration” (page 46).
- Function, see the section called “Function Declaration” (page 45).
- Rewrite rule, see the section called “Rewrite Rule Declaration” (page 48).

- Node annotation, see the section called “Node Annotation Declaration” (page 47).
- Declaration tag, see the section called “Declaration Tag” (page 46).

The module name *Name* will be used when the current module is imported in another module. A module is usually a qualified name of the form:

```
Name1 :: Name2 :: ... :: Namen
```

which corresponds to a path relative to the root of the current workspace.



Explain this better.

The constituents of a module are also shown in Figure 1.19, “Constituents of a Module”(page 44)  
The more sophisticated features are shown in a separate color.

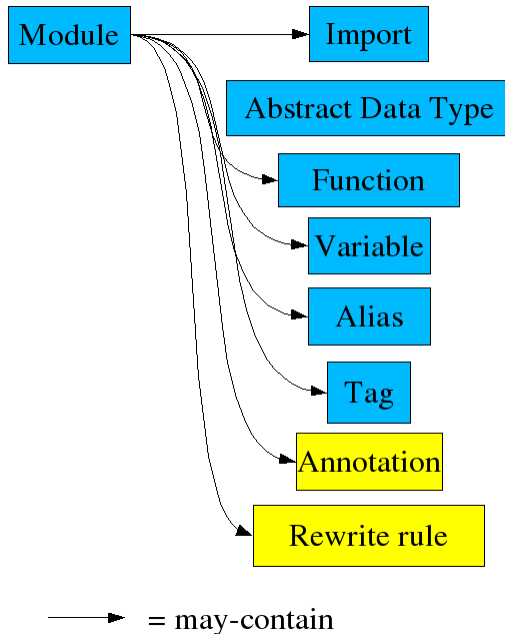
The entities that are *visible* inside a module are

- The private or public entities declared in the module itself.
- The public entities declared in any imported module.

The only entities that are visible outside the module, are the public entities declared in the module itself. The visible entities of an imported may be explicitly qualified with their module name:

```
Module :: Name
```

**Figure 1.19. Constituents of a Module**



## Import

An import has the form:

```
import QualifiedName;
```

and has as effect that all public entities declared in module *QualifiedName* are made available to the importing module. Circular imports are allowed.

## Function Declaration

A function declaration has the form

```
Type Name(Type1 Var1, ..., Typen Varn) Statement
```

Here *Type* is the result type of the function and this should be equal to the type of the result of *Statement* (using the return statement, see the section called “Return Statement” (page 65). Each *Type*<sub>*i*</sub> *Var*<sub>*i*</sub> represents a typed formal parameter of the function. The formal parameters may be used in *Statement* and get their value when function *Name* is invoked. Example:

```
rascal> rel[int, int] invert(rel[int,int] R){
    return {<Y, X> | <int X, int Y> <- R }
}
void: null

rascal> invert({<1,10>, <2,20>});
rel[int,int] : {<10,1>, <20,2>}
```

**Variable argument lists.** A function may have a variable list of arguments, this has the form:

```
Type Name(ordinary parameters, Type Var...) Statement
```

The last parameter of a function may be followed by ... and this has as effect that all remaining actual parameters that occur in a call to this function are collected as list value of the last formal parameter. Inside the function body, the type of this parameter will therefore be `list[Type]`.

**Exceptions.** The exceptions that can be thrown by a function can be (optionally) declared as follows:

```
Type Name(Type1 Var1, ..., Typen Varn)
    throws Exception1, Exception2, ...
```

See the section called “Try Catch Statement” (page 66) and the section called “Throw Statement” (page 66) for details.

**Parameterized types in function declaration.** The types that occur in function declarations may also contain *type variables* that are written as & followed by an identifier. In this way functions can be defined for arbitrary types. In the following example, we declare an inversion function that is applicable to any binary relation. :

```
rascal> rel[&T2, &T1] invert2(rel[&T1,&T2] R){
    return {<Y, X> | <&T1 X, &T2 Y> <- R };
}
void: null

rascal> invert2({<1,10>, <2,20>});
rel[int,int] : {<10,1>, <20,2>}

rascal> invert2({<"mon", 1>, <"tue", 2>});
rel[int,str] : {<1, "mon">, <2, "tue">}
```

Here we declare a function that can be used to swap the elements of pairs of arbitrary types:

```
rascal> <&T2, &T1> swap(&T1 A, &T2 B) { return <B, A>; }
void: null

rascal> swap(<1, 2>);
tuple[int,int] : <2,1>

rascal> swap(<"wed", 3>);
```

```
tuple[int,str] : <3, "wed">
```

## Variable Declaration

A variable declaration has the form

```
Type Var = Exp
```

where *Type* is a type, *Var* is a variable name, and *Exp* is an expression that should have type *Type*. The effect is that the value of expression *Exp* is assigned to *Var* and can be used later on as *Var*'s value. The following rules apply:

- Double declarations in the same scope are not allowed.
- The type of *Exp* should be compatible with *Type*.

As a convenience, also declarations without an initialization expression are permitted inside functions (but not at the module level) and have the form

```
Type Var
```

and only introduce the variable *Var*. When a variable is declared, it has as scope the nearest enclosing block (see the section called “Block Statement” (page 68)).

Rascal provides *local type inference*, which allows the implicit declaration for variables that are used locally in functions. The following rules apply:

- An implicitly declared variable is declared at the level of the function body.
- An implicitly declared variable gets as type the type of the first value that is assignment to it.
- If a variable is implicitly declared in different execution path of a function, all these implicit declarations should result in the same type.
- All uses of an implicitly declared variable must be compatible with its implicit type.

Examples.

```
rascal> int max = 100;
int: 100

rascal> min = 0;
int : 0

rascal> day = {<"mon", 1>, <"tue", 2>, <"wed",3>,
               <"thu", 4>, <"fri", 5>, <"sat",6>, <"sun",7>}
rel[str,int]: {<"mon", 1>, <"tue", 2>, <"wed",3>,
               <"thu", 4>, <"fri", 5>, <"sat",6>, <"sun",7>}
```

## Declaration Tag



Tags are not yet implemented.



The syntax of tags has to be aligned with the syntax of annotations. This is done in the examples below but not yet in the syntax.

Tags are intended to add metadata to a Rascal program and allow to influence the execution of the Rascal program, for instance, by adding memoization hints or database mappings for relations.

All declarations in a Rascal program may contain (in fixed positions depending on the declaration type) one or more declaration tags (`tag`). A tag is defined by declaring its name, the declaration type to which it can be attached, and the name and type of the annotation. The declaration type `all`, makes the declaration tag applicable for all possible declaration types. All declaration tags have the generic format `@Name{ ... }`, with arbitrary text between the brackets that is further constrained by the declared type. Here is an example of a license tag:

```
tag str license on module;
```

This will allow to write things like:

```
module Booleans
@license{This module is distributed under the GPL}
...
```

Other examples of declaration tags are:

```
tag str todo on all          %% todo note for all types
tag void deprecated on function %% marks a deprecated function
tag int memo on function      %% bounded memoization of
                              %% function calls
tag str doc on all           %% documentation string
tag str primitive on function %% a primitive, built-in, function
```

Here is an example of a documentation string as used in the Rascal standard library:

```
public &T max(set[&T] R)
  @doc{Maximum of a set: max}
{
  &T result = arb(R);
  for(&T E : R){
    result = max(result, E);
  }
  return result;
}
```

## Node Annotation Declaration

An annotation may be associated with any node value. Annotations are intended to attach application data to values, like adding position information or control flow information to source code or adding visualization information to a relation. An annotation has a name and the type of its value is explicitly declared. Any value of any named type can be annotated and the type of these annotations can be declared precisely.

For instance, we can add to certain syntactic constructs of programs (e.g., `EXPRESSION`) an annotation with name `posinfo` that contains location information:

```
anno loc EXPRESSION @ posinfo;
```

or location information could be added for all syntax trees:

```
anno loc node @ posinfo;
```

We can add to the graph datatype introduced earlier, the annotation with name `LayoutStrategy` that defines which graph layout algorithm to apply to a particular graph, e.g.,

```
data LayoutStrategy = "dot" | "tree" | "force" |
                     "hierarchy" | "fisheye";
anno LayoutStrategy Graph @ strategy;
```

The following constructs are provided for handling annotations:

- `Val @ Anno`: is an expression that retrieves the value of annotation *Anno* of value *Val* (may be undefined!).
- `Val1[@Anno = Val2]`: is an expression that set the value of annotation *Anno* of the value *Val1* to *Val2* and return XXX as result.
- `Var @ Anno = Val`: is an assignment statement that sets the value of annotation *Anno* of the value of variable *Var* to *Val*.

## Rewrite Rule Declaration

Functions are the workhorses of Rascal. They can have any value as parameter or result and are explicitly called by the user. Also, functions are declared inside modules and their visibility can be controlled.

Rewrite rules, on the other hand, operate only on nodes and defined datatypes, they are implicitly applied when a new value (we refer to this as the *subject value*) is constructed. The scope of rewrite rules is the whole Rascal program. Rewrite rules are applied to the subject value in a bottom-up fashion. As a result, the subject value may be changed. This process is repeated as long as there are rules that can be applied to the current subject value. Technically, this is called *innermost rewriting*. When done, the result of rewriting the original subject value is used instead of that original value.

Rules have the general form:

```
rule Name PatternWithAction
```

where *Name* is the name of the rule and *PatternWithAction* is the body of the rule consisting of a pattern and an associated action (see the section called “PatternWithAction” (page 58) for a detailed description).

Here is an example for a user-defined type Booleans:

```
rascal>
data Bool = btrue;
data Bool = bfalse;
data Bool = band(Bool left, Bool right);
data Bool = bor(Bool left, Bool right);

rule a1 band(btrue, Bool B)    => B;
rule a2 band(bfalse, Bool B)   => bfalse;
void: null
rascal> band(band(btrue,btrue),band(btrue, bfalse));
Bool: bfalse
```

During execution of rules the following applies:

- Rules are applied non-deterministically, and in any order of matching.
- The right-hand side of rules can contain fail statements, which cause backtracking over the alternative matches or alternative rules for a certain constructor.
- When the right-hand side is a statement, an `insert` statement determines the value of the actual replacement.

## Expressions

The Expression is the basic unit of evaluation and may consist of (see Figure 1.20, “Expression Forms” (page 49)):

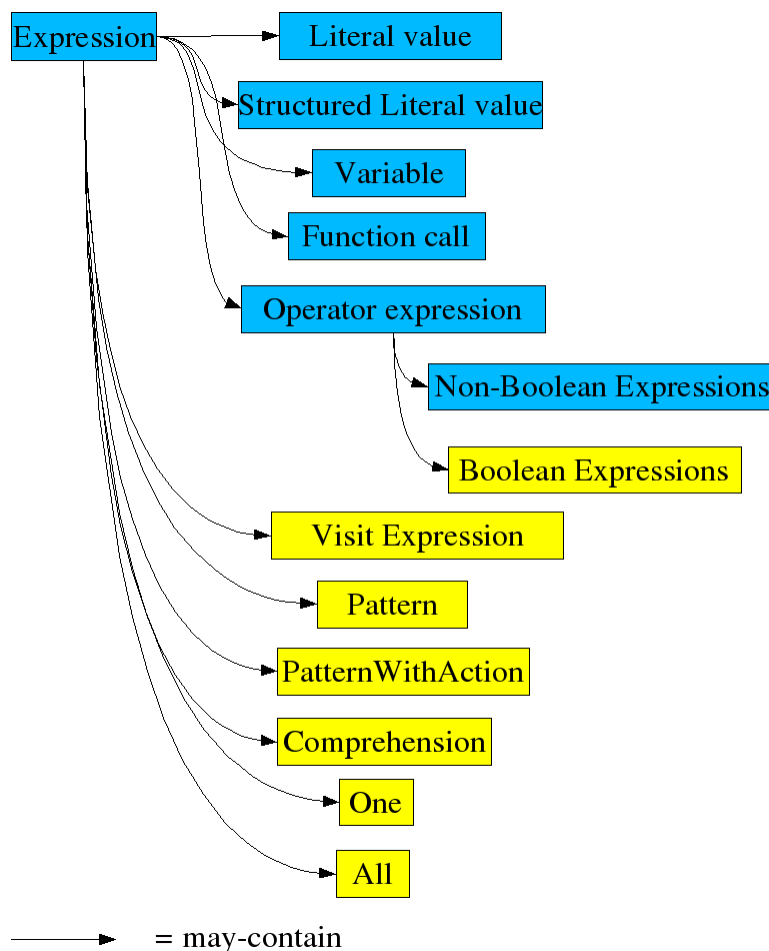
- An *elementary literal value*, e.g. constants of type `bool`, `int`, `real`, `str` or `loc`. Elementary literals evaluate to themselves.



- A *structured literal value* for list, set, map, tuple, or rel. The elements are first evaluated before the structured value is built.
- A *variable*, evaluates to its current value.
- A *function call*. First the arguments are evaluated and the corresponding function is called. The value returned by the function is used as value of the function call.
- A *constructor*. First the arguments are evaluated and then a data value is constructed for the corresponding type. This data value is used as value of the constructor.
- An *operator expression*. The operator is applied to the arguments; the evaluation order of the arguments depends on the operator. The result returned by the operator is used as value of the operator expression.
- A visit expression, see the section called “Visit expression” (page 60).
- A pattern, see the section called “Patterns” (page 54).
- A pattern with associated action, see the section called “PatternWithAction” (page 58).
- A comprehension, see the section called “Comprehension Expression” (page 58).
- A one expression, see the section called “One Expression” (page 62).
- An all expression, see the section called “All Expression” (page 62).

We explain the various operator expressions in more detail. Non-Boolean and Boolean operator expressions behave somewhat differently and therefore we describe separately.

**Figure 1.20. Expression Forms**



## Non-Boolean Operator Expressions

The non-Boolean operators are summarized in Table 1.1, “Non-Boolean Operators” (page 50). All operators are highly overloaded and we refer to the section called “Built-in Operators and Library Functions” (page 68) for a description of their meaning for each specific type. The following rules apply:

- All *name* arguments stand for themselves and are not evaluated.
- For all operators, except, IfThenElse, the argument expressions are first evaluated before the operator is applied.
- The operators in Table 1.1, “Non-Boolean Operators” (page 50) are listed from high precedence to low precedence. In other words, operators listed higher in the table bind stronger.

**Table 1.1. Non-Boolean Operators**

Operator	Name	Description
$Exp . Name$	Field selection	$Exp$ should evaluate to a tuple or datatype with field <i>Name</i> ; return the value of that field
$exp_1 [ Name = Exp_2 ]$	Field assignment	$Exp_1$ should evaluate to a tuple or datatype with a field <i>Name</i> ; assign value $Exp_2$ to that field
$Exp < field, \dots >$	Field projection	$Exp$ should evaluate to a tuple or relation, and field should be a field name or an integer constant. A new tuple or relation is returned that only contains the listed fields.
$Exp_1 [ Exp_2 ]$	Subscription	The value of $Exp_2$ is used as index in $Exp_1$ 's value. On list, tuple return the element with given index; for map return the value associated with $Exp_2$ 's value. On relation ...
$- Exp$	Negation	Negative of $Exp$ 's integer or real value
$Exp +$	Transitive closure	Transitive closure on relation
$Exp *$	Reflexive transitive closure	Reflexive transitive closure on relation
$Exp @ Name$	Annotation selection	Value of annotation <i>Name</i> of $Exp$ 's value
$Exp_1 [ @ Name = Exp_2 ]$	Annotation replacement	Assign value of $Exp_2$ to annotation <i>Name</i> of $Exp_1$ 's value
$Exp_1 \circ Exp_2$	Composition	$Exp_1$ and $Exp_2$ should evaluate to a relation; return their composition
$Exp_1 / Exp_2$	Division	Divide two integers and reals
$Exp_1 \% Exp_2$	Modulo	Modulo on integer
$Exp_1 * Exp_2$	Multiplication / Product	Multiply integers or real; product of list, set, relation
$Exp_1 \& Exp_2$	Intersection	Intersection of list, set, map or relation
$Exp_1 + Exp_2$	Addition / concatenation / union	Add integer and real; concatenate string, list or tuple; union on set, map, or relation
$Exp_1 - Exp_2$	Subtraction / difference	Subtract integer or real; difference of list, set, map, or relation
$Exp_1 \text{ join } Exp_2$	Join	Join on relation

Most of these operators are described in detail in the section called “Built-in Operators and Library Functions” (page 68) for specific types. Here we describe the remaining generic operators.

**Field Selection and Field Assignment.** Field selection and field assignment apply to all values that have named components like tuples and relations with named elements, data types, and locations. Field selection returns the value of the named component and Field assignment returns a new value in which the named component has been replaced by a new value. We illustrate this for tuples:

```
rascal> tuple[int key, str val] T = <1, "abc">;
tuple[int, str] : <1, "abc">

rascal> T.val;
str : "abc"

rascal> T[val = "def"];
tuple[int, str] : <1, "def">

rascal> T;
tuple[int, str] : <1, "abc">
```

Observe that field assignment creates a *new* value with an updated field. The old value remains unchanged as can be seen from the unchanged value of T in the above example. In the section called “Assignment Statement” (page 63) we explain how to change a field of variable value.

**Field projection.** Field projection applies to tuples and relations and may contain element names or integer constants that refer to elements in the order in which they occur in the original value (counting from 0). A field projection returns a new value that consists of the selected elements. Suppose we have a relation with traffic information that records the name of the day, the day number, and the length of the traffic jams at that day.

```
rascal> rel[str day, int daynum, int length] Traffic =
    {<"mon", 1, 100>, <"tue", 2, 150>, <"wed", 3, 125>,
     <"thur", 4, 110>, <"fri", 5, 90>};
rel[str, int, int]: {<"thur", 4, 110>, <"tue", 2, 150>, <"wed", 3, 125>,
    <"fri", 5, 90>, <"mon", 1, 100>}

rascal> Traffic<length, daynum>;
rel[int, int]: {<110, 4>, <150, 2>, <90, 5>, <100, 1>, <125, 3>}

rascal> Traffic<2, day>;
rel[int, str]: {<125, "wed">, <110, "thur">, <100, "mon">, <90, "fri">,
    <150, "tue">}
```

Field projection thus selects parts from a larger value that has a fixed number of parts. The selection is based on position and not on value and can be used to completely reorder or remove the parts of a larger value.

**Subscription.** Subscription selects values with a given computed index from a larger value that has a variable number of elements. For lists and tuples a single integer expression is allowed as index and the returned value is the element with that index (counting from 0). For maps, the index type should correspond to the key type of the map and the value associated with the index is returned. For relations, more than one index expression is allowed and as value a new, reduced, relation is returned with all elements that contained the index values at the corresponding tuples positions (but these values are removed, hence a reduced relation). Some examples illustrate this:

```
rascal> L = [10, 20, 30];
list[int] : [10, 20, 30]

rascal> L[1];
int : 20
```

```

rascal> T = <"mon", 1>;
tuple[str,int] : <"mon", 1>;

rascal> T[0];
str : "mon"

rascal> colors = ("hearts":"red", "clover":"black",
                  "trumps":"black", "clubs":"red");
map[str,str] : ("hearts":"red", "clover":"black",
                "trumps":"black", "clubs":"red")

rascal> colors["trumps"];
str: "black"

rascal> rel[str country, int year, int amount] GDP =
    {<"US", 2008, 14264600>, <"EU", 2008, 18394115>,
     <"Japan", 2008, 4923761>, <"US", 2007, 13811200>,
     <"EU", 2007, 13811200>, <"Japan", 2007, 4376705>};
rel[str,int,int] :
    {<"US", 2008, 14264600>, <"EU", 2008, 18394115>,
     <"Japan", 2008, 4923761>, <"US", 2007, 13811200>,
     <"EU", 2007, 13811200>, <"Japan", 2007, 4376705>}

rascal> GDP["Japan"];
rel[int,int] : {<2008, 4923761>, <2007, 4376705>}

rascal> GDP["Japan", 2008];
set[int] : {4923761}

```

**Other Non-Boolean Operator Expressions.** The other non-Boolean operator expressions are explained in more detail for each datatype, see the section called “Built-in Operators and Library Functions” (page 68).

## Boolean Operator Expressions

The Boolean operators are summarized in Table 1.2, “Boolean Operators” (page 53) All operators are highly overloaded and we refer to the section called “Built-in Operators and Library Functions” (page 68) for a description of their meaning for each specific type. Most operators are self-explanatory except the Match and NoMatch operators that are also the main reason to treat Boolean operator expressions separately. Although we describe patterns in full detail in the section called “Patterns” (page 54), a preview is useful here. A pattern can

- match (or not match) any arbitrary value (that we will call the *subject value*);
- during the match variables may be bound to subvalues of the subject value.

**Match Operator.** The match operator

```
Pat := Exp
```

is evaluated as follows:

- *Exp* is evaluated, the result is a subject value;
- the subject value is matched against the pattern *pat*;
- if the match succeeds, any variables in the pattern are bound to subvalues of the subject value and the match expression yields `true`;
- if the match fails, no variables are bound and the match expression yields `false`.

This looks and *is* nice and dandy, so why all this fuss about Boolean operators? The catch is that-- as we will see in detail in the section called “Patterns” (page 54)--a match need not be unique. This means that there may be more than one way of matching the subject value resulting in different variable bindings. A quick example. Consider the following match of a list

```
[1, list[int] L, 2, list[int] M] := [1,2,3,2,4]
```

There are two solutions for this match:

- $L = []$  and  $M = [2, 3, 2, 4]$ ; and
- $L = [2, 3]$  and  $M = [4]$ .

Depending on the context, only the first solution of a match expression is used, respectively all solutions are used. If a match expression occurs in a larger Boolean expression, a subsequent subexpression may yield `false` and -- depending on the actual operator -- evaluation *backtracks* to a previously evaluated match operator to try a next solution. Let's illustrate this by extending the above example:

```
[1, list[int] L, 2, list[int] M] := [1,2,3,2,4] && size(L) > 0
```

where we are looking for a solution in which  $L$  has a non-empty list as value. Evaluation proceeds as follows:

- The left argument of the `&&` operator is evaluated: the match expression is evaluated resulting in the bindings  $L = []$  and  $M = [2, 3, 2, 4]$ ;
- The right argument of the `&&` operator is evaluated: `size(L) > 0` yields `false`;
- Backtrack to the left argument of the `&&` operator to check for more solutions: indeed there are more solutions resulting in the bindings  $L = [2, 3]$  and  $M = [4]$ ;
- Proceed to the right operator of `&&`: this time `size(L) > 0` yields `true`;
- The result of evaluating the complete expression is `true`.

This behaviour is applicable in the context of the following Rascal constructs:

- Comprehensions, see the section called “Comprehension Expression” (page 58).
- Tests in `for`, `one`, `all` statements, see the section called “Statements” (page 62).



Make the above more specific.

**Table 1.2. Boolean Operators**

Operator	Name	Description
<code>! Exp</code>	Negation	Negate <i>Exp</i> 's boolean value
<code>Exp ?</code>	IsDefined	true if <i>Exp</i> has a well-defined value
<code>Exp<sub>1</sub> in Exp<sub>2</sub></code>	ElementOf	Element of
<code>Exp<sub>1</sub> not in Exp<sub>2</sub></code>	NotElementOf	Not element of
<code>Exp<sub>1</sub> &lt;= Exp<sub>2</sub></code>	LessThanOrEqual	Less than or equal on bool, int, real or string; sublist on list; subset on set, map or relation
<code>Exp<sub>1</sub> &lt; Exp<sub>2</sub></code>	LessThan	Less than on bool, int, real or string; strict sublist on list; strict subset on set, map or relation

Operator	Name	Description
$Exp_1 \geq Exp_2$	GreaterThanOrEqual	Greater than or equal on bool, int, real or string; superlist on list; superset on set, map or relation
$Exp_1 > Exp_2$	GreaterThan	Greater than on bool, int, real or string; strict superlist on list; strict superset on set, map or relation
$Pat := Exp$	Match	Value of $Exp$ matches with pattern $Pat$
$Pat !:= Exp$	NoMatch	Value of $Exp$ does not match with pattern $Pat$
$Exp_1 == Exp_2$	Equal	Equality
$Exp_1 != Exp_2$	NotEqual	Inequality
$Exp_1 ? Exp_2$	IfDefinedElse	the value of $Exp_1$ if it is well-defined, otherwise the value of $Exp_2$
$Exp_1 ? Exp_2 : Exp_3$	IfThenElse	Conditional expression
$Exp_1 ==> Exp_2$	Implication	true, unless the value of $Exp_1$ is true and that of $Exp_2$ is false
$Exp_1 <==> Exp_2$	Equivalence	true if $Exp_1$ and $Exp_2$ have the same value
$Exp_1 \&\& Exp_2$	And	true if the value of both $Exp_1$ and $Exp_2$ is true
$Exp_1, Exp_2, \dots, Exp_n$	MultiCondition	Equivalent to: $Exp_1 \&\& Exp_2 \&\& \dots \&\& Exp_n$
$Exp_1    Exp_2$	Or	true if the value of either $Exp_1$ or $Exp_2$ is true
$Pat <- Exp$	Enumerator	true for every element in $Exp$ 's value that matches $pat$

## Patterns

Patterns come in three flavours:

- *Regular expression patterns* to do string matching with regular expressions, see the section called “Regular Expression Patterns” (page 54).
- *Abstract patterns* to matching on arbitrary values, see the section called “Abstract Patterns” (page 55).
- *Concrete syntax patterns* to match syntax trees that are the result of parsing, see the section called “Concrete Syntax Patterns” (page 57).

## Regular Expression Patterns

Regular expression patterns are ordinary regular expressions that are used to match a string value and to decompose it in parts and also to compose new strings. Regular expression patterns bind variables of type `str` when the match succeeds, otherwise they do not bind anything. Their syntax and semantics parallels abstract and concrete syntax patterns as much as possible. This means that they can occur in cases of `visit` and `switch` statements, on the left-hand side of the match operator (`:=` or `!:=`) and as declarator in generators.

We use a regular expression language that slightly extends the Java Regex language with the following exceptions:

- Regular expression are delimited by `/` and `/` optionally followed by a modifier (see below).

- We allow named groups, syntax `<Name : Regexp>`, which introduce a variable of type `str` named `Name`. Currently, these names have to be unique in the pattern.
- Java regular expressions might have optional groups, which may introduce null bindings. Since uninitialized variables are not allowed in Rascal, we limit the kinds of expressions one can write here by not allowing nesting of named groups.
- Named groups have to be outermost, such that they can only bind in one way.
- Unlike Perl, Java uses the notation `(?Option)` inside the regular expression to set options like multi-line matching `(?m)`, case-insensitive matching `(?i)` etc. We let these options follow the regular expression.
- We allow name use in a regular expression: `<Name>` which inserts the string value of `Name` in the pattern.

Here are some examples of regular patterns.

```
/\brascal\b/i
```

does a case-insensitive match (`i`) of the word `rascal` between word boundaries (`\b`). And

```
/^.*?(word:\w+)<rest:.*$/m
```

does a multi-line match (`m`), matches the first consecutive word characters (`\w`) and assigns them to the variable `word`. The remainder of the string is assigned to the variable `rest`.

## Abstract Patterns



Change this name. E.g., "General pattern", "Value Pattern"???

An abstract pattern is recursively defined and may contain the following elements:

- *Literal* of one of the basic types `bool`, `int`, `real`, `str`, or `loc`. A literal pattern matches with a value that is identical to the literal.
- A *variable declaration pattern*

```
Type Var
```

A variable declaration introduces a new variable that matches any value of the given type. That value is assigned to `Var` when the whole match succeeds.

- A *multi-variable pattern*

```
Var*
```

A multi-variable is an abbreviation for a variable declaration pattern. It can occur in a list pattern or set pattern and can match zero or more list or set elements.

- A *variable pattern*

```
Var
```

A variable pattern can act in two roles:

- If `Var` has already a defined value then it matches with that value.
- If `Var` has not been defined before (or it has been declared but not initialized) then it matches any value. That value is assigned to `Var`.



Explain scope.

- A *list pattern*

```
[ Pat1, Pat2, ..., Patn ]
```

A list pattern matches a list value, provided that  $Pat_1, Pat_2, \dots, Pat_n$  match the elements of that list in order. Two special cases exist when one of the patterns  $Pat_i$  is

- a variable declaration pattern with a list type that is identical to the type of the list that is being matched.
- a variable pattern, where the variable has been declared, but not initialized, outside the pattern with a list type that is identical to the type of the list that is being matched.

In both cases list matching is applied and the variable can match an arbitrary number of elements of the subject list.

- A *set pattern*

```
{ Pat1, Pat2, ..., Patn }
```

A set pattern matches a set value, provided that  $Pat_1, Pat_2, \dots, Pat_n$  match the elements of that set in any order. Completely analogous to list patterns, there are two special cases when one of the patterns  $Pat_i$  is

- a variable declaration pattern with a set type that is identical to the type of the set that is being matched.
- a variable pattern, where the variable has been declared, but not initialized, outside the pattern with a set type that is identical to the type of the set that is being matched.

In both cases set matching is applied and the variable can match an arbitrary number (in arbitrary order!) of elements of the subject set.

- A *tuple pattern*

```
< Pat1, Pat2, ..., Patn >
```

A tuple pattern matches a tuple value, provided that  $Pat_1, Pat_2, \dots, Pat_n$  match the elements of that tuple in order.

- A *node pattern*

```
Name ( Pat1, Pat2, ..., Patn )
```

A node pattern matches a node value or a datatype value, provided that *Name* matches with the constructor symbol of that value and  $Pat_1, Pat_2, \dots, Pat_n$  match the children of that value in order.

- A *descendant pattern*

```
/ Pat
```

performs a deep match of the pattern *Pat*. In other words, it matches when any element of the subject at any depth matches *Pat* and is used to match, for instance, tree nodes at an arbitrary distance from the root.

- A *labelled pattern*



```
Var : Pat
```

A labelled pattern matches the same values as *Pat*, but has as side-effect that the matched value is assigned to *Var*.

- A *typed, labelled, pattern*

```
Type Var : Pat
```

A typed, labelled, pattern matches when the subject value has type *Type* and *Pat* matches. The matched value is assigned to *Var*.

- A *type constrained pattern*

```
[Type] Pat
```

matches provided that the subject has type *Type* and *Pat* matches.



Map patterns are currently not supported.

## Concrete Syntax Patterns



Concrete patterns are not yet implemented.

A *concrete pattern* is a (possibly quoted) concrete syntax fragment that may contain variables. We want to cover the whole spectrum from maximally quoted patterns that can unambiguously describe **any** syntax fragment to minimally quoted patterns as we are used to in ASF+SDF. A concrete pattern may have the following forms:

- A *typed variable pattern*

```
<Type Var>
```

- A *variable pattern*

```
<Var>
```

- A *quoted pattern*

```
[ | Token1 Token2 ... Tokenn | ]
```

Inside a quoted pattern arbitrary lexical tokens may occur, but the characters `<`, `>` and `|` have to be escaped as `\<`, `\>`, `\|`. Quoted patterns may contain variable declaration patterns and variable patterns.

- A *typed quoted pattern*

```
Symbol [ | Token1 Token2 ... Tokenn | ]
```

is a quoted pattern that is preceded by an SDF symbol to define its desired syntactic type.

- An *unquoted pattern*

```
Token1 Token2 ... Tokenn
```

is a quoted pattern without the surrounding quotes.

- Inside syntax patterns, layout is ignored.

Examples:

- Quoted syntax pattern with two pattern variable declarations:

```
[ | while <EXP Exp> do <{STATEMENT ";"}* Stats> od | ]
```

- Quoted syntax pattern with two pattern variable uses:

```
[ | while <Exp> do <Stats> od | ]
```

- Identical to the previous example, but with a declaration of the desired syntactic type:

```
STATEMENT [ | while <Exp> do <Stats> od | ]
```

- Unquoted syntax pattern with two pattern variable declarations:

```
while <EXP Exp> do <{STATEMENT ";"}* Stats> od
```

- Unquoted syntax pattern with two pattern variable uses:

```
while <Exp> do <Stats> od
```

Obviously, with less quoting and type information, the probability of ambiguities increases. Our assumption is that a type checker can resolve most of them.

## PatternWithAction

Patterns can be used in various contexts, but a common context is a *PatternWithAction*, which in its turn, may be used in a visit expression (see the section called “Visit expression” (page 60), a switch statement (see the section called “Switch Statement” (page 63), or a rewrite rule (see the section called “Rewrite Rule Declaration” (page 48)).

A *PatternWithAction* can have one of the following forms:

- *Pat => Exp*

When the subject matches *Pat*, the expression *Exp* is evaluated. The use of the resulting value depends on the context and is described in the relevant section.

- *Pat : Statement*

This is the most general case. When the subject matches *Pat*, the *Statement* is executed. The execution of *Statement* should, depending on the context, lead to one of the following:

- Execution of a `return` statement that returns a value from the enclosing function.
- Execution of a `fail` statement: all side effects of *Statement* are undone and the *PatternWithAction* as a whole fails.
- None of the above: execution continues with the statement following the switch.

## Comprehension Expression

We will use the familiar notation for *list comprehension*

```
[Exp1, ..., Expn | Gen1, ..., Genm]
```

to denote the construction of a list consisting of the successive values of the *contributing expressions* *Exp<sub>1</sub>, ..., Exp<sub>n</sub>*. The values and the resulting list are determined by *Exp<sub>1</sub>, ..., Exp<sub>n</sub>* and the *generators* *Gen<sub>1</sub>, ..., Gen<sub>m</sub>*. *Exp<sub>1</sub>, ..., Exp<sub>n</sub>* are computed for all possible combinations of values produced by the generators. Each generator may introduce new variables that can be used in

subsequent generators as well as in the expressions. A generator can use the variables introduced by preceding generators. Generators may enumerate all the values in a set or relation, or they may perform an arbitrary test.

In addition to list comprehensions, Rascal also supports *set comprehension*

```
{Exp1, ..., Expn | Gen1, ..., Genm}
```

that also serve as relation comprehension in the case that *Exp* is of a tuple type.

Finally, *map comprehensions* are written as:

```
(Exp1 : Exp2 | Gen1, ..., Genm)
```

Since the entries in a map require both a key and a value for each entry, two expressions are needed in this case.

## Enumerator

An enumerator generates all the values in a given list, set, map, tuple, relation or datatype. It has the following form:

```
Pat <- Exp
```

where *Pat* is a pattern and *Exp* is an expression. An enumerator is evaluated as follows:

- Expression *Exp* is evaluated and may have an arbitrary value *V*.
- The elements of *V* are enumerated one by one.
- Each element value is matched against the pattern *Pat*. There are two cases:
  - The match succeeds, any variables in *Pat* are bound, and the next generator in the comprehension is evaluated. The variables that are introduced by an enumerator are only available to generators that appear later (i.e., to the right) in the comprehension. When this enumerator is the last generator in the comprehension its contributing expression is evaluated.
  - The match fails, no variables are bound. If *V* has more elements, a next element is tried. Otherwise, a previous generator (i.e., to the left) is tried. If this enumerator is the first generator in the comprehension, the evaluation of the comprehension is complete.

These are examples of enumerators:

- `int N <- {1, 2, 3, 4, 5},`
- `str K <- KEYWORDS,` where `KEYWORDS` should evaluate to a value of `set[str]`.
- `<str K, int N> <- {"a",10>, <"b",20>, <"c",30>}`.
- `<str K, int N> <- FREQUENCIES,` where `FREQUENCIES` should evaluate to a value of type `rel[str,int]`.
- `<str K, 10> <- FREQUENCIES,` will only generate pairs with 10 as second element.



Type information will be used to check the plausibility of an enumerator and guard you against mistakes. An impossible enumerator like `int N <- {"apples", "oranges"}` will be flagged as an error since the pattern can never match.



An enumerator may be preceded by a *strategy indication*:

- top-down
- bottom-up (this is the default)

These take only effect for enumerators that produce the elements of a data type and determine the order in which the elements are enumerated.

## Test

A test is a boolean-valued expression. If the evaluation yields `true` this indicates that the current combination of generated values up to this test is still as desired and execution continues with subsequent generators. If the evaluation yields `false` this indicates that the current combination of values is undesired, and that another combination should be tried by going back to the previous generator.

Examples:

- `N >= 3` tests whether `N` has a value greater than or equal 3.
- `S == "coffee"` tests whether `S` is equal to the string `"coffee"`.

In both examples, the variable (`N`, respectively, `S`) should have been introduced by a generator that occurs earlier in the comprehension.

## Examples of Comprehensions

Here are some examples of comprehensions:

```
rascal> {x | int x : {1, 2, 3, 4, 5}, x >= 3};
set[int] : {3,4,5}

rascal> {<x, y> | int x : {1, 2, 3}, int y : {2, 3, 4}, x >= y};
rel[int,int] : {<2, 2>, <3, 2>, <3, 3>}

rascal> {<y, x> | <int x, int y> : {<1,10>, <2,20>}};
rel[int,int] : {<10,1>, <20,2>}

rascal> {x, x * x | x : {1, 2, 3, 4, 5}, x >= 3};
set[int] : {3,4,5,9,16,25}
```

## Visit expression

Visiting the nodes in a tree is a very common task in the EASY domain. In many cases (but certainly not all) the tree is a syntax tree of some source code file and the nodes correspond to expressions or statements. Computing metrics or refactoring are examples of tasks that require a tree visit. In object-oriented programming, the *visitor pattern* is in common use for this. There are three frequently occurring scenarios:

- Accumulator: traverse the tree and collect information.
- Transformer: traverse the tree and transform it into another tree.
- Accumulating Transformer: traverse the tree, collect information and also transform the tree.

The `visit` expression in Rascal can accommodate all these (and more) use cases and has the form:

```
Strategy visit ( Exp ) {
case PatternWithAction1;
case PatternWithAction2;
...
default: ...
}
```

Given a subject term (the current value of `Exp`) and a list of cases (consisting of `PatternWithActions`, see the section called “`PatternWithAction`” (page 58) it traverses the term. Depending on

the precise actions it may perform replacement (mimicking a transformer), update local variables (mimicking an accumulator) or a combination of these two (accumulating transformer). If **any** of the actions contains an `insert` statement, the value of the visit expression is a new value that is obtained by successive insertions in the subject term by executing one or more cases. Otherwise, the original value of the subject term is returned.

The visit expression is optionally preceded by one of the following strategy indications that determine the traversal order of the subject:

- `top-down`: visit the subject from root to leaves.
- `top-down-break`: visit the subject from root to leaves, but stop at the current path when a case matches.
- `bottom-up`: visit the subject from leaves to root (this is the default).
- `bottom-up-break`: visit the subject from leaves to root, but stop at the current path when a case matches.
- `innermost`: repeat a bottom-up traversal as long as the traversal changes the resulting value (compute a fixed-point).
- `outermost`: repeat a top-down traversal as long as the traversal changes the resulting value (compute a fixed-point).

The execution of the cases has the following effect:

- `PatternWithActions` of the form  $Pat \Rightarrow Exp$  insert their result in the subject.
- For `PatternWithActions` of the form  $Pat : Statement$ , executing *Statement* should lead to one of the following:
  - Execution of an `insert` statement of the form

```
insert Expr
```

The value of *Exp* replaces the subtree of the subject that is currently being visited. Note that a copy of the subject is created at the start of the visit statement and all insertions are made in this copy. As a consequence, insertions cannot influence matches later on.



An `insert` statement may only occur inside a visit expression.

- Execution of a `fail` statement: all side effects of *Statement* are undone, no insertion is made, and the next case is tried.
- Execution of a `return` statement that returns a value from the enclosing function.

The precise behaviour of the visit expression depends on the type of the subject:

- For type `node` or `ADT`, all nodes of the tree are visited (in the order determined by the strategy). Concrete patterns and abstract patterns directly match tree nodes. Regular expression patterns match only values of type string.
- For structured types (`list`, `set`, `map`, `tuple`, `rel`), the elements of the structured type are visited and matched against the cases. When inserts are made, a new structured value is created.



Have strategies any effect for non-tree subjects?

## One Expression

```
one ( Exp1 , Exp2 , ... , Expn )
```

The one expression yields true when one combination of values of *Exp*<sub>i</sub> is true.



The status of one is under discussion.

## All Expression

```
all ( Exp1 , Exp2 , ... , Expn )
```

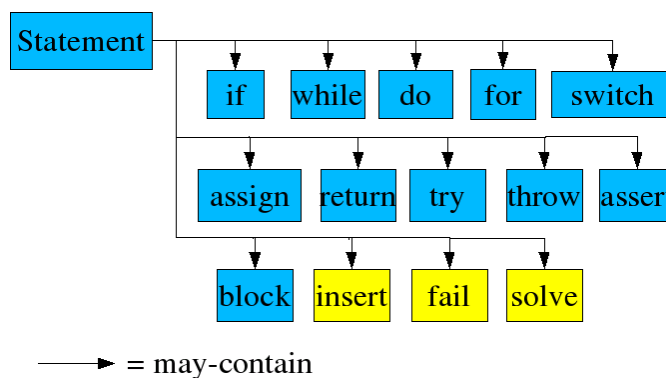
The all expression yields true when all combinations of values of *Exp*<sub>i</sub> are true.



The status of all is under discussion.

## Statements

**Figure 1.21. Various Statement Forms**



The various statements forms are summarized in Figure 1.21, “Various Statement Forms” (page 62). The more advanced statements are show in a different color.

## If Statement

The if-statement is completely standard and comes in an if-then and an if-then-else variant:

```
if ( Bool ) Statement
```

```
if ( Bool ) Statement1 else Statement2
```

In both cases the test *Bool* is evaluated and its outcome determines the statement to be executed. Recall from the section called “Boolean Operator Expressions”(page 52)that boolean expression maybe multi-valued. In this case only the first true value (if any) is used.

## While Statement

The while-statement is completely standard:

```
while ( Bool ) Statement
```

The test *Bool* is evaluated repeatedly and *Statement* is executed when the test is true. Execution ends the first time that the test yields false. The test *Bool* is executed anew in each repetition and only the first true value (if any) is used.

## Do Statement

The do-while-statement is completely standard:

```
do Statement while ( Bool )
```

Statement is executed repeatedly, as long as the test *Bool* yields true. The test *Bool* is executed anew in each repetition and only the first true value (if any) is used.

## For Statement

The for-statement is more exciting:

```
for ( Exp1 , Exp2 , ... , Expn ) Statement
```

It generates executes *Statement* for all possible combinations of values of the expressions *Exp*<sub>i</sub>. Note that is one of the expressions is a boolean expression, we do try all its possible values.

## Switch Statement

A switch statement is similar to a switch statement in C or Java and has the form:

```
switch ( Exp ) {  
case PatternWithAction1;  
case PatternWithAction2;  
...  
default: ...  
}
```

The value of the expression *Exp* is the subject term that will be matched by the successive *PatternWithActions* (see the section called “*PatternWithAction*” (page 58)) in the switch statement. The switch statement provides **only** matching at the top level of the subject term and does not traverse it. The type of the pattern in each case must be identical to the type of the subject term (or be a subtype of it). If no case matches, the switch acts as a dummy statement. There is no fall through from one case to the next.

## Assignment Statement

The purpose of an assignment is to assign a new value to a simple variable or to an element of a more complex data structure. The most general form of an assignment statement is

```
Assignable AssignmentOp Exp
```

where *AssignmentOp* may be =, +=, -=, \*=, /=, or ?=. Here = is the ordinary assignment operators and the other forms can be derived from it according to Table 1.3, “Assignment Operators” (page 63).

**Table 1.3. Assignment Operators**

Assignment Operator	Equivalent to
<i>Assignable</i> += <i>Exp</i>	<i>Assignable</i> = <i>Assignable</i> + <i>Exp</i>
<i>Assignable</i> -= <i>Exp</i>	<i>Assignable</i> = <i>Assignable</i> - <i>Exp</i>
<i>Assignable</i> *= <i>Exp</i>	<i>Assignable</i> = <i>Assignable</i> * <i>Exp</i>

Assignment Operator	Equivalent to
<i>Assignable</i> /= <i>Exp</i>	<i>Assignable</i> = <i>Assignable</i> / <i>Exp</i>
<i>Assignable</i> ?= <i>Exp</i>	<i>Assignable</i> = <i>Assignable</i> ? <i>Exp</i>

An assignable is either a single variable, (the *base variable*), optionally followed by subscriptions or field selections. The assignment statement always results in assigning a completely new value to the base variable. We distinguish the following forms of assignment:

- *Var* = *Exp*

The expression *Exp* is evaluated and its value is assigned to the base variable *Var*.

- *Assignable* [ *Exp*<sub>1</sub> ] = *Exp*<sub>2</sub>

First the value *V* of *Assignable* is determined. Next the value of *Exp*<sub>1</sub> is used as index in *V* and the value of *Exp*<sub>2</sub> replaces the original value at that index position. The result is a new value *V'* that is assigned to the *Assignable*.

- *Assignable* . *Name* = *Exp*

The value *V* of *Assignable* is determined and should be of a type that has a field *Name*. The value of that field is replaced in *V* by the value of *Exp* resulting in a new value *V'* that is assigned to *Assignable*.

- < *Assignable*<sub>1</sub>, *Assignable*<sub>2</sub>, ..., *Assignable*<sub>n</sub> > = *Exp*

First the value *Exp* is determined and should be a tuple of the form <*V*<sub>1</sub>, *V*<sub>2</sub>, ..., *V*<sub>n</sub>>. Next the assignments *Assignable*<sub>i</sub> = *V*<sub>i</sub> are performed for 1 ≤ i ≤ n.

- *Assignable* ? *Exp*<sub>1</sub> = *Exp*<sub>2</sub>

First the value of *Exp*<sub>2</sub> is determined and if that is defined it is assigned to *Assignable*. Otherwise, the value of *Exp*<sub>1</sub> is assigned to *Assignable*.

- *Assignable* @ *Name* = *Exp*

The value *V* of *Assignable* is determined and should be of a type that has an annotation *Name*. The value of that annotation is replaced in *V* by the value of *Exp* resulting in a new value *V'* that is assigned to *Assignable*.

- *Name* ( *Assignable*<sub>1</sub>, *Assignable*<sub>2</sub>, ... ) = *Exp*

First the value *Exp* is determined and should be a data value of the form *Name* (*V*<sub>1</sub>, *V*<sub>2</sub>, ..., *V*<sub>n</sub>). Next the assignments *Assignable*<sub>i</sub> = *V*<sub>i</sub> are performed for 1 ≤ i ≤ n.



Constructor assignable is not yet implemented.

Here are some examples:

```
rascal> N = 3;
int : 3

rascal> N;
int : 3

rascal> L = [10,20,30];
list[int] : [10,20,30]
```



```
rascal> L[1] = 200;
list[int] : [10,200,30]

rascal> M = ("abc": 1, "def" : 2);
map[str,int] : ("abc": 1, "def" : 2)

rascal> M["def"] = 3;
map[str,int] : ("abc": 1, "def" : 3)

rascal> T = <1, "abc", true>;
tuple[int,str,bool] : <1, "abc", true>

rascal> T[1] = "def";
tuple[int,str,bool] : <1, "def", true>

rascal> data FREQ = wf(str word, int freq);
void: null

rascal> W = wf("rascal", 1000);
FREQ : wf("rascal", 1000)

rascal> W.freq = 100000;
FREQ : wf("rascal",100000)

rascal> <A, B, C> = <"abc", 2.5, [1,2,3]>;
tuple[str,real,list[int]] : <"abc", 2.5, [1,2,3]>

rascal> A;
str : "abc"

rascal> B;
real : 2.5

rascal> C;
list[int] : [1,2,3]

rascal> ** good V ? E1 = E2 example here **

rascal> anno str FREQ@color;
void: null

rascal> W @ color = "red";
FREQ: wf("rascal",100000)[@color="red"]

rascal> wf(S, I) = W;
void: null

rascal> S;
str : "rascal"

rascal> I;
int : 100000
```

## Return Statement

A return statement has either the form

```
return;
```

or

```
return Exp;
```

both end the execution of the current function. The first form applies to functions with `void` as return type. The second form applies to non-void functions and returns the value of *Exp* as result of the function. The following rules apply:

- The static type of *Exp* should be compatible with the declared return type of the function in which the return statement occurs.
- In each function with a return type that is not `void`, every possible execution path through the body of the function end in a return statement.
- In each function with a return type that is `void`, a return statement is implicitly assumed at the end of each execution path through the function body.

## Try Catch Statement

A try catch statement has the form

```
try
  Statement1;
catch PatternWithAction1;
catch PatternWithAction2;
...
catch: Statement2;
finally: Statement3;
```

and has as purpose to catch any exceptions that are raised during the execution of *Statement*<sub>1</sub>. These exceptions may be caused by:

- The execution of an explicit throw statement, see the section called “Throw Statement” (page 66).
- The Rascal system that discover an abnormal condition, e.g., an out of bounds error when accessing a list element.

Note that all elements of the try catch statement are optional but that at least one has to be present. Their meaning is as follows:

- If a pattern of some *PatternWithAction*<sub>*i*</sub> matches, the corresponding action is executed.
- Otherwise, *Statement*<sub>2</sub> is executed (when present).
- Before leaving the try catch statement *Statement*<sub>3</sub> is always executed (when present).

## Throw Statement

A throw statement has the form

```
throw Exp;
```

and causes the immediate abortion of the execution of the current function with *Exp*'s value as exception value. The exception can be caught by a try catch statement (the section called “Try Catch Statement” (page 66) in the current function or in one of its callers. If the exception is not caught, the execution of the Rascal program is terminated. The following rules apply:

- The static type of *Exp* should be `RuntimeException`, see the section called “Exception” (page 71).
- The Rascal program may contain data declarations that extend the type `RuntimeException`.

## Assert Statement

An assert statement may occur everywhere where a declaration is allowed. It has two forms:

```
assert Exp1
```

and

```
assert Exp1 : Exp2
```

where *Exp*<sub>1</sub> is a boolean-value expression and *Exp*<sub>2</sub> is a string-valued expression that serves as a identifying message for this assertion. During execution, a list of true and false assertions is maintained. When the script is executed as a *test suite* a summary of this information is shown to the user. When the script is executed in the standard fashion, the assert statement has no affect.



Update the above description.

Example:

```
rascal> assert {1, 2, 3, 1} == {3, 2, 1, 1} : "Equality on Sets";  
bool: true
```

## Insert Statement

An insert statement has the form

```
insert Exp;
```

and replaces the value of the *current subject* (see below) by the value of *Exp*. An insert statement may only occur in the action part of a `PatternWithAction` (see the section called “PatternWithAction” (page 58)), more precisely in

- A case in a visit expression, see the section called “Visit expression” (page 60) The current subject is the value matched by the pattern of this case.
- An action of a rewrite rule, see the section called “Rewrite Rule Declaration” (page 48) The current subject is the value matched by the pattern of the rewrite rule.

The following rule applies:

- The static type of *Exp* and of the current subject should be comparable.

## Fail Statement

A fail statement has the form

```
fail;
```

and may only occur in the action of `PatternWithAction` (see the section called “PatternWithAction” (page 58)). The fail statement forces the failure of that action. Any bindings caused by the pattern or side-effects caused by the action are undone.

## Block Statement

A block consists of a sequence of statements separated by semi-colons:

```
{ Statement1; ... Statementn }
```

Since a block is itself a statement, it may be used in all places where a statement is required. A block also introduces a new scope and variables that are declared in the block are local to that block.

## Solve Statement

Rascal provides a solve statement for performing arbitrary fixed-point computations. This means, repeating a certain computation as long as it causes changes. This can, for instance, be used for the solution of sets of simultaneous linear equations but has much wider applicability. The format is:

```
with {
  Type1 Var1 = Exp1;
  Type2 Var2 = Exp2;
  ...
} solve
  Statement
```

The solve statement consists of an initialization section and a statement. Optionally, an expression directly following the solve keyword, gives an upperbound on the number of iterations.

In the initial section, the variables *Var*<sub>*i*</sub> are declared and initialized. In the solve section, the statement can use and modify these variables. The statement is executed, assigning new values to the variables *Var*<sub>*i*</sub>, and this is repeated as long as the value of any of the variables was changed compared to the previous repetition.

Let's consider transitive closure as an example (transitive closure is already available as built-in operator, we use it here just as a simple illustration). Transitive closure of a relation is usually defined as:

$$R^+ = R + (R \circ R) + (R \circ R \circ R) + \dots$$

This can be expressed as follows:

```
rascal> rel[int,int] R = {<1,2>, <2,3>, <3,4>};
rel[int,int] : {<1,2>, <2,3>, <3,4>}
rascal>
  with
    rel[int,int] T = R;
  solve
    T = T + (T o R1);
rel[int,int] : {<1,2>, <1,3>, <1,4>, <2,3>, <2,4>, <3,4>}
```

## Built-in Operators and Library Functions

The built-in operators and library functions can be subdivided in the following categories:

- Benchmark: measuring functions, see the section called “Benchmark” (page 70).
- Boolean: operators and functions on Boolean values, see the section called “Boolean” (page 70).
- Exception: data definition of all soft exceptions that can be caught by Rascal programs, see the section called “Exception” (page 71).

- Graph: graphs are a special kind of binary relation, see the section called “Graph” (page 71).
- Integer: operators and functions on integers, see the section called “Integer” (page 72).
- IO: simple print functions, see the section called “IO” (page 73).
- Labelled Graph: labelled graphs with addition edge information, see the section called “Labelled Graph” (page 73).
- List: operators and functions on lists, see the section called “List” (page 73).
- Location: operators and functions on source locations, see the section called “Location” (page 75).
- Map: operators and functions on maps, see the section called “Map” (page 75).
- Node: operators and functions on nodes, see the section called “Map” (page 75).
- Real: operators and functions on reals, see the section called “Real” (page 77).
- Relation: operators and functions on relations, see the section called “Relation” (page 77).
- Resource: functions to retrieve resources from an Eclipse workspace, see the section called “Resource (Eclipse only)” (page 79).
- RSF: function for reading files in Rigi Standard Format, see the section called “RSF” (page 79).
- Set: operators and functions on sets, see the section called “Set” (page 80).
- String: operators and functions on strings, see the section called “String” (page 82).
- Tuple: operators and functions on tuples, see the section called “Tuple” (page 83).
- UnitTest: functions for unit testing, see the section called “UnitTest” (page 83).
- ValueIO: functions for reading and writing Rascal values, both in textual and in binary form, see the section called “ValueIO” (page 84).
- View: functions for graphical display of values in Eclipse, see the section called “View (Eclipse only)” (page 84).
- Void: the type void, see

All operators are directly available for each program, but library functions have to be imported in each module that uses them.

We use some notational conventions to describe the argument of operators, as shown in Table 1.4, “Notational conventions” (page 69) When an operator has more than one argument of the same type, they are distinguished by subscripts.

**Table 1.4. Notational conventions**

Argument	Describes expression of type
<i>Bool</i>	bool
<i>Int</i>	int
<i>Real</i>	real
<i>Str</i>	str
<i>Loc</i>	loc
<i>Node</i>	node
<i>List</i>	Any list type

Argument	Describes expression of type
<i>Set</i>	Any set type
<i>Map</i>	Any map type
<i>Tuple</i>	Any tuple type
<i>Rel</i>	Any rel type
<i>Value</i>	value
<i>Elm</i>	Compatible with element type of list, set, map, relation

## Benchmark

**Table 1.5. Benchmark Functions**

Function	Description
<code>real currentTimeMillis()</code>	Current time in milliseconds since January 1, 1970 GMT.
<code>p.m. benchmark</code>	Measure and report the execution time of <code>name:void-closure</code> pairs.

## Boolean

**Table 1.6. Boolean Operators**

Operator	Description
$Bool_1 == Bool_2$	Yields true if both arguments are identical
$Bool_1 != Bool_2$	Yields true if both arguments are not identical
$Bool_1 <= Bool_2$	Yields true if both arguments are identical or $Bool_1$ is false and $Bool_2$ is true
$Bool_1 < Bool_2$	Yields true if $Bool_1$ is false and $Bool_2$ is true
$Bool_1 >= Bool_2$	Yields true if both arguments are identical or $Bool_1$ is true and $Bool_2$ is false
$Bool_1 > Bool_2$	Yields true if $Bool_1$ is true and $Bool_2$ is false
$Bool_1 \&\& Bool_2$	Yields true if both arguments have the value true and false otherwise
$Bool_1    Bool_2$	Yields true if either argument has the value true and false otherwise
$Bool_1 ==> Bool_2$	Yields false if $Bool_1$ has the value true and $Bool_2$ has value false, and true otherwise
$! Bool$	Yields true if Bool is false and true otherwise
$Bool_1 ? Bool_2 : Bool_3$	If $Bool_1$ is true then $Bool_2$ else $Bool_3$

**Table 1.7. Boolean Functions**

Function	Description
<code>bool arbBool()</code>	Arbitrary boolean value
<code>bool fromInt(int i)</code>	Convert an integer to a bool
<code>bool fromString(str s)</code>	Convert the strings "true" or "false" to a bool
<code>int toInt(bool b)</code>	Convert a boolean value to integer
<code>real toReal(bool b)</code>	Convert a boolean value to a real value

Function	Description
<code>str toString(bool b)</code>	Convert a boolean value to a string

## Exception

The following "soft" exceptions are defined:

```
data RuntimeException =
  EmptyList
  | EmptyMap
  | EmptySet
  | IndexOutOfBounds(int index)
  | AssertionFailed
  | AssertionFailed(str label)
  | NoSuchElement(value v)
  | IllegalArgument(value v)
  | IllegalArgument
  | IO(str message)
  | FileNotFound(str filename)
  | LocationNotFound(loc location)
  | PermissionDenied
  | PermissionDenied(str message)
  | ModuleNotFound(str name)
  | NoSuchKey(value key)
  | NoSuchAnnotation(str label)
  | Java(str message)
;
```

## Graph

The graph datatype is a special form of binary relation defined as follows:

```
alias graph[&T] = rel[&T from, &T to];
```

**Table 1.8. Graph Functions**

Function	Description
<code>set[&amp;T] bottom(graph[&amp;T] G)</code>	Bottom nodes of a graph
<code>set[&amp;T] predecessors(grap[&amp;T], &amp;T From)</code>	Direct predecessors of node From
<code>set[&amp;T] predecessors(grap[&amp;T], set[&amp;T] FromSet)</code>	Direct predecessors of all nodes in FromSet
<code>set[&amp;T] reach(graph[&amp;T] G, set[&amp;T] Start)</code>	Reachability from set of start nodes.
<code>set[&amp;T] reachR(graph[&amp;T] G, set[&amp;T] Start, set[&amp;T] Restr)</code>	Reachability from set of start nodes with restriction to certain nodes.
<code>set[&amp;T] reachX(graph[&amp;T] G, set[&amp;T] Start, set[&amp;T] Excl)</code>	Reachability from set of start nodes with exclusion of certain nodes
<code>list[&amp;T] shortestPathPair(graph[&amp;T] G, &amp;T From, &amp;T To)</code>	Shortest path between pair of nodes
<code>set[&amp;T] successors(grap[&amp;T], &amp;T From)</code>	Direct successors of node From
<code>set[&amp;T] successors(grap[&amp;T], set[&amp;T] FromSet)</code>	Direct successors of all nodes in FromSet
<code>set[&amp;T] top(graph[&amp;T] G)</code>	Top nodes of a graph

The following examples illustrate these functions:

```
rascal> top(<1,2>, <1,3>, <2,4>, <3,4>);
set[int] : {1}

rascal> bottom(<1,2>, <1,3>, <2,4>, <3,4>);
set[int] : {4}

rascal> reachR({1}, {1, 2, 3}, {<1,2>, <1,3>, <2,4>, <3,4>});
set[int] : {2, 3}

rascal> reachX({1}, {2}, {<1,2>, <1,3>, <2,4>, <3,4>});
set[int,int] : {3, 4}
```

## Integer

Rascal integers are unbounded in size.

**Table 1.9. Integer Operators**

Operator	Description
$Int_1 == Int_2$	Yields true if both arguments are numerically equal and false otherwise
$Int_1 != Int_2$	Yields true if both arguments are numerically unequal and false otherwise
$Int_1 <= Int_2$	Yields true if $Int_1$ is numerically less than or equal to $Int_2$ and false otherwise
$Int_1 < Int_2$	Yields true if $Int_1$ is numerically less than $Int_2$ and false otherwise
$Int_1 >= Int_2$	Yields true if $Int_1$ is numerically greater than or equal than $Int_2$ and false otherwise
$Int_1 > Int_2$	Yields true if $Int_1$ is numerically greater than $Int_2$ and false otherwise
$Int_1 + Int_2$	Sum of $Int_1$ and $Int_2$
$Int_1 - Int_2$	Difference of $Int_1$ and $Int_2$
$Int_1 * Int_2$	$Int_1$ multiplied by $Int_2$
$Int_1 / Int_2$	$Int_1$ divided by $Int_2$
$Int_1 \% Int_2$	Remainder of dividing $Int_1$ by $Int_2$
$-Int$	Negate sign of $Int$
$Bool ? Int_1 : Int_2$	If $Bool$ is true then $Int_1$ else $Int_2$

**Table 1.10. Integer Functions**

Function	Description
<code>int abs(int N)</code>	Absolute value of integer N
<code>int arbInt()</code>	Arbitrary integer value
<code>int arbInt(int limit)</code>	Arbitrary integer value in the interval [0, limit)
<code>int max(int n, int m)</code>	Largest of two integers
<code>int min(int n, int m)</code>	Smallest of two integers
<code>real toReal(int n)</code>	Convert an integer value to a real value
<code>str toString(int n)</code>	Convert an integer value to a string



## IO

**Table 1.11. IO Functions**

Function	Description
<code>void java println(value V...)</code>	Print a list of values on the output stream
<code>list[str] readFile(str filename)</code> <code>throws NoSuchFileError(str msg),</code> <code>IOException(str msg)</code>	Read a named file as list of strings

## Labelled Graph

The labelled graph datatype is a special form of binary relation with labelled edges and is defined as follows:

```
alias lgraph[&T,&L] = rel[&T from, &L label, &T to];
```

**Table 1.12. Labelled Graph Functions**

Function	Description
<code>set[&amp;T] bottom(lgraph[&amp;T] G)</code>	Bottom nodes of a labelled graph
<code>set[&amp;T] predecessors(grap[&amp;T], &amp;T From)</code>	Direct predecessors of node From
<code>set[&amp;T] predecessors(grap[&amp;T], set[&amp;T] FromSet)</code>	Direct predecessors of all nodes in FromSet
<code>set[&amp;T] reach(lgraph[&amp;T] G, set[&amp;T] Start)</code>	Reachability from set of start nodes.
<code>set[&amp;T] reachR(lgraph[&amp;T] G, set[&amp;T] Start, set[&amp;T] Restr)</code>	Reachability from set of start nodes with restriction to certain nodes.
<code>set[&amp;T] reachX(lgraph[&amp;T] G, set[&amp;T] Start, set[&amp;T] Excl)</code>	Reachability from set of start nodes with exclusion of certain nodes
<code>list[&amp;T] shortestPathPair(lgraph[&amp;T] G, &amp;T From, &amp;T To)</code>	Shortest path between pair of nodes
<code>set[&amp;T] successors(grap[&amp;T], &amp;T From)</code>	Direct successors of node From
<code>set[&amp;T] successors(grap[&amp;T], set[&amp;T] FromSet)</code>	Direct successors of all nodes in FromSet
<code>set[&amp;T] top(lgraph[&amp;T] G)</code>	Top nodes of a labelled graph



shortestPath not yet implemented for lgraph.

## List

**Table 1.13. List Operators**

Operator	Description
<code>List<sub>1</sub> == List<sub>2</sub></code>	Yields true if both arguments have the same elements in the same order
<code>List<sub>1</sub> != List<sub>2</sub></code>	Yields true if both arguments have different elements
<code>List<sub>1</sub> &lt;= List<sub>2</sub></code>	Yields true if both lists are equal or <i>List<sub>1</sub></i> is a sublist of <i>List<sub>2</sub></i>
<code>List<sub>1</sub> &lt; List<sub>2</sub></code>	Yields true if <i>List<sub>1</sub></i> is a sublist of <i>List<sub>2</sub></i>
<code>List<sub>1</sub> &gt;= List<sub>2</sub></code>	Yields true if both lists are equal or <i>List<sub>2</sub></i> is a sublist of <i>List<sub>1</sub></i>

Operator	Description
$List_1 > List_2$	Yields true if $List_2$ is a sublist of $List_1$
$List_1 + List_2$	Concatenation of $List_1$ and $List_2$
$List_1 - List_2$	List consisting of all elements in $List_1$ that do not occur in $List_2$
$List_1 * List_2$	$List_1$ multiplied by $List_2$
$Elm \text{ in } List$	Yields true if $Elm$ occurs as element in $List$
$Elm \text{ not in } List$	Yields true if $Elm$ does not occur as element in $List$
$Bool ? List_1 : List_2$	If $bool$ is true then $List_1$ else $List_2$
$List [ int ]$	Element at position $int$ in $List$

**Table 1.14. List Functions**

Function	Description
<code>&amp;T average(list[&amp;T] lst, &amp;T zero)</code>	Average of elements of a list
<code>list[&amp;T] delete(list[&amp;T] lst, int n)</code>	Delete nth element from list
<code>set[int] domain(list[&amp;T] lst)</code>	Set of all legal index values for a list
<code>&amp;T java head(list[&amp;T] lst) throws EmptyListError</code>	Get the first element of a list
<code>list[&amp;T] head(list[&amp;T] lst, int n) throws IndexOutOfBoundsException</code>	Get the first n elements of a list
<code>&amp;T getOneFrom(list[&amp;T] lst)</code>	Get an arbitrary element from a list
<code>list[&amp;T] insertAt(list[&amp;T] lst, int n, &amp;T elm) throws IndexOutOfBoundsException</code>	Add an element at a specific position in a list
<code>bool isEmpty(list[&amp;T] lst)</code>	Is list empty?
<code>list[&amp;T] mapper(list[&amp;T] lst, &amp;T (&amp;T, &amp;T) fn)</code>	Apply a function to each element of a list
<code>&amp;T max(list[&amp;T] lst)</code>	Largest element of a list
<code>&amp;T min(list[&amp;T] lst)</code>	Smallest element of a list
<code>&amp;T multiply(list[&amp;T] lst, &amp;T unity)</code>	Multiply the elements of a list
<code>set[list[&amp;T]] permutations(list[&amp;T] lst)</code>	All permutations of a list
<code>&amp;T reducer(list[&amp;T] lst, &amp;T (&amp;T, &amp;T) fn, &amp;T unit)</code>	Apply function F to successive elements of a list
<code>list[&amp;T] reverse(list[&amp;T] lst)</code>	Elements of a list in reverse order
<code>int size(list[&amp;T] lst)</code>	Number of elements in a list
<code>list[&amp;T] slice(list[&amp;T] lst, int start, int len)</code>	Sublist from start of length len
<code>list[&amp;T] sort(list[&amp;T] lst)</code>	Sort the elements of a list
<code>&amp;T sum(list[&amp;T] lst, &amp;T zero)</code>	Add elements of a List
<code>list[&amp;T] tail(list[&amp;T] lst)</code>	All but the first element of a list
<code>list[&amp;T] tail(list[&amp;T] lst, int len) throws IndexOutOfBoundsException</code>	Last n elements of a list
<code>tuple[&amp;T, list[&amp;T]] takeOneFrom(list[&amp;T] lst)</code>	Remove an arbitrary element from a list, returns the element and the modified list

Function	Description
<code>map[&amp;A,&amp;B] toMap(list[tuple[&amp;A, &amp;B]] lst)</code>	Convert a list of tuples to a map
<code>set[&amp;T] toSet(list[&amp;T] lst)</code>	Convert a list to a set
<code>str toString(list[&amp;T] lst)</code>	Convert a list to a string

## Location

**Table 1.15. Operations on Locations**

Operator	Description
$Loc_1 == Loc_2$	Yield <code>true</code> if both arguments are identical and <code>false</code> otherwise
$Loc_1 != Loc_2$	Yield <code>true</code> if both arguments are not identical and <code>false</code> otherwise
$Loc_1 <= Loc_2$	Yield <code>true</code> if $Loc_1$ is textually contained in or equal to $Loc_2$ and <code>false</code> otherwise
$Loc_1 < Loc_2$	Yield <code>true</code> if $Loc_1$ is strictly textually contained in $Loc_2$ and <code>false</code> otherwise
$Loc_1 >= Loc_2$	Yield <code>true</code> if $Loc_1$ is textually encloses or is equal to $Loc_2$ and <code>false</code> otherwise
$Loc_1 > Loc_2$	Yield <code>true</code> if $Loc_1$ is textually encloses $Loc_2$ and <code>false</code> otherwise
$Loc . Field$	Retrieve one of the fields of location value

The field names for locations are:

- `url`
- `offset`
- `length`
- `beginLine, beginColumn`
- `endLine, endColumn`.

## Map

**Table 1.16. Map Operators**

Operator	Description
$Map_1 == Map_2$	Yield <code>true</code> if both arguments consist of the same pairs
$Map_1 != Map_2$	Yield <code>true</code> if both arguments have different pairs
$Map_1 <= Map_2$	Yield <code>true</code> if all pairs in $Map_1$ occur in $Map_2$ or $Map_1$ and $Map_2$ are equal
$Map_1 < Map_2$	Yield <code>true</code> if all pairs in $Map_1$ occur in $Map_2$ but $Map_1$ and $Map_2$ are not equal
$Map_1 >= Map_2$	Yield <code>true</code> if all pairs in $Map_2$ occur in $Map_1$ or $Map_1$ and $Map_2$ are equal
$Map_1 > Map_2$	Yield <code>true</code> if all pairs in $Map_2$ occur in $Map_1$ but $Map_1$ and $Map_2$ are not equal
$Map_1 + Map_2$	Union of $Map_1$ and $Map_2$

Operator	Description
$Map_1 - Map_2$	Difference of $Map_1$ and $Map_2$
$Key \text{ in } Map$	Yield <code>true</code> if $Key$ occurs in a key:value pair in $Map$
$Key_1 \text{ not in } Map_2$	Yield <code>true</code> if $Key$ does not occur in a key:value pair in $map$
$Bool ? Map_1 : Map_2$	If $Bool$ is true then $Map_1$ else $Map_2$
$Map [ Key ]$	The value associated with $Key$ in $Map$ if that exists, undefined otherwise

**Table 1.17. Map Functions**

Function	Description
<code>set[&amp;K] domain(map[&amp;K, &amp;V] M)</code>	The domain (keys) of a map
<code>&amp;K getOneFrom(map[&amp;K, &amp;V] M)</code>	Arbitrary key of a map
<code>map[&amp;V, &amp;K] invert(map[&amp;K, &amp;V] M)</code>	Map with key and value inverted
<code>bool isEmpty(map[&amp;K, &amp;V] M)</code>	Is map empty?
<code>map[&amp;K, &amp;V] mapper(map[&amp;K, &amp;V] M, &amp;K (&amp;K) F, &amp;V (&amp;V) G)</code>	Apply two functions to each key/value pair in a map.
<code>set[&amp;V] range(map[&amp;K, &amp;V] M)</code>	The range (values) of a map
<code>int size(map[&amp;K, &amp;V] M)</code>	Number of elements in a map.
<code>list[tuple[&amp;K, &amp;V]] toList(map[&amp;K, &amp;V] M)</code>	Convert a map to a list
<code>rel[&amp;K, &amp;V] toRel(map[&amp;K, &amp;V] M)</code>	Convert a map to a relation
<code>str toString(map[&amp;K, &amp;V] M)</code>	Convert a map to a string.

## Node

**Table 1.18. Node Operators**

Operator	Description
$Node_1 == Node_2$	<code>true</code> if both arguments are identical
$Node_1 != Node_2$	<code>true</code> if both arguments are not identical
$Node_1 <= Node_2$	
$Node_1 < Node_2$	
$Node_1 >= Node_2$	
$Node_1 > Node_2$	
$Bool ? Node_1 : Node_2$	if $Bool$ is true then $Node_1$ else $Node_2$
$Node [ Int ]$	child of $Node$ at position $Int$

**Table 1.19. Node Functions**

Function	Description
<code>int arity(node T)</code>	Number of children of a node
<code>list[value] getChildren(node T)</code>	The children of a node
<code>str getName(node T)</code>	The function name of a node
<code>node makeNode(str N, value V...)</code>	Create a node given its function name and arguments

## Real

**Table 1.20. Real Operators**

Operator	Description
$Real_1 == Real_2$	true if both arguments are numerically equal and false otherwise
$Real_1 != Real_2$	true if both arguments are numerically unequal and false otherwise
$Real_1 <= Real_2$	true if $Real_1$ is numerically less than or equal to $Real_2$ and false otherwise
$Real_1 < Real_2$	true if $Real_1$ is numerically less than $Real_2$ and false otherwise
$Real_1 >= Real_2$	true if $Real_1$ is numerically greater than or equal than $Real_2$ and false otherwise
$Real_1 > Real_2$	true if $Real_1$ is numerically greater than $Real_2$ and false otherwise
$Real_1 + Real_2$	Sum of $Real_1$ and $Real_2$
$Real_1 - Real_2$	Difference of $Real_1$ and $Real_2$
$Real_1 * Real_2$	$Real_1$ multiplied by $Real_2$
$Real_1 / Real_2$	$Real_1$ divided by $Real_2$
$-Real$	Negate sign of $Real$
$Real_1 \% Real_2$	Remainder of dividing $Real_1$ by $Real_2$
$Bool ? Real_1 : Real_2$	If $Bool$ is true then $Real_1$ else $Real_2$

**Table 1.21. Real Functions**

Function	Description
<code>real arbReal()</code>	Arbitrary real value in the interval [0.0,1.0).
<code>real max(real n, real m)</code>	Largest of two reals
<code>int toInteger(real d)</code>	Convert a real to integer.
<code>str toString(real d)</code>	Convert a real to a string.

## Relation

Relation are sets of tuples, therefore all set operators (see, Table 1.26, “Set Operators” (page 80)) apply to relations as well

**Table 1.22. Operations on Relations**

Operator	Description
$Rel_1 \circ Rel_2$	Relation resulting from the composition of the two arguments
$Set_1 \times Set_2$	Relation resulting from the Cartesian product of the two arguments
$Rel +$	Relation resulting from the transitive closure of $Rel$
$Rel *$	Relation resulting from the reflexive transitive closure of $Rel$
$Rel [ elem ]$	Relation resulting from subscription of $Rel$ with
$Rel [ Set ]$	Relation resulting from subscription of $Rel$ with all elements of $Set$ .
$Rel < Index_1, Index_2, \dots >$	Relation resulting from restricting $Rel$ to the columns described by $Index_1, Index_2, \dots$

Examples:

```

rascal> {<1,10>, <2,20>, <3,15>} o {<10,100>, <20,200>};
rel[int,int] : {<1,100>, <2,200>}

rascal> {1, 2, 3} x {9};
rel[int,int] : {<1, 9>, <2, 9>, <3, 9>}

rascal> Rel = {<1,10>, <2,20>, <1,11>, <3,30>, <2,21>};
rel[int,int] : {<1,10>, <2,20>, <1,11>, <3,30>, <2,21>}

rascal> Rel[1];
set[int] : {10, 11}

rascal> Rel[{1}];
set[int] : {10, 11}

rascal> Rel[{1, 2}];
set[int] : {10, 11, 20, 21}

```

**Table 1.23. Relation Functions**

Function	Description
set[&T] carrier (rel[&T,&T] R)	All elements in any tuple in a relation
rel[&T,&T] carrierR (rel[&T,&T] R, set[&T] S)	Relation restricted to tuples with elements in a set S
rel[&T,&T] carrierX (rel[&T,&T] R, set[&T] S)	Relation excluded tuples with some element in S
rel[&T0, &T1] complement (rel[&T0, &T1] R)	Complement of relation
set[&T0] domain (rel[&T0,&T1] R)	First element of each tuple in binary relation
rel[&T0,&T1] domainR (rel[&T0,&T1] R, set[&T0] S)	Restriction of a relation to tuples with first element in S
rel[&T0,&T1] domainX (rel[&T0,&T1] R, set[&T0] S)	Relation excluded tuples with first element in S
ident?	<i>status?</i>
rel[&T1, &T0] invert (rel[&T0, &T1] R)	Inverse the tuples in a relation
rel[&T2, &T1, &T0] invert (rel[&T0, &T1, &T2] R)	All but the first element of each tuples in binary relation
rel[&T0,&T1] ranger (rel[&T0,&T1] R, set[&T2] S)	Restriction of a binary relation to tuples with second element in set S

Examples:

```

rascal> id({1,2,3});
rel[int,int] : {<1,1>, <2,2>, <3,3>}

rascal> id({"mon", "tue", "wed"});
rel[str,str] : {<"mon","mon">, <"tue","tue">, <"wed","wed">}

rascal> invert({<1,10>, <2,20>});
rel[int,int] : {<10,1>, <20,2>}

rascal> compl({<1,10>});

```

```

rel[int,int] : {<1, 1>, <10, 1>, <10, 10>}

rascal> domain({<1,10>, <2,20>});
set[int] : {1, 2}

rascal> domain({<"mon", 1>, <"tue", 2>});
set[str] : {"mon", "tue"}

rascal> range({<1,10>, <2,20>});
set[int] : {10, 20}

rascal> range({<"mon", 1>, <"tue", 2>});
set[int] : {1, 2}

rascal> carrier({<1,10>, <2,20>});
set[int] : {1, 10, 2, 20}

rascal> domainR({<1,10>, <2,20>, <3,30>}, {3, 1});
rel[int,int] : {<1,10>, <3,30>}

rascal> rangeR({<1,10>, <2,20>, <3,30>}, {30, 10});
rel[int,int] : {<1,10>, <3,30>}

rascal> carrierR({<1,10>, <2,20>, <3,30>}, {10, 1, 20});
rel[int,int] : {<1,10>}

rascal> domainX({<1,10>, <2,20>, <3,30>}, {3, 1});
rel[int,int] : {<2, 20>}

rascal> rangeX({<1,10>, <2,20>, <3,30>}, {30, 10});
rel[int,int] : {<2, 20>}

rascal> carrierX({<1,10>, <2,20>, <3,30>}, {10, 1, 20});
rel[int,int] : {<3,30>}

```

## RSF

**Table 1.24. RSF Functions**

Function	Description
map[str, rel[str,str]] readRSF(str nameRSFFile)	Read a file in Rigi Standard Format (RSF).

## Resource (Eclipse only)

```

data Resource = root(set[Resource] projects)
                | project(str name, set[Resource] contents)
                | folder(str name, set[Resource] contents)
                | file(str name, str extension);

```

**Table 1.25. Resource Functions**

Function	Description
set[loc] java files(str project)	The files contained in a project
loc java location(str project)	Source location of given project

Function	Description
set[str] java projects()	The projects in the Eclipse workspace
set[str] java references(str project)	The project references of a given project
Resource java root()	The root of the Eclipse workspace

## Set

**Table 1.26. Set Operators**

Operator	Description
$Set_1 == Set_2$	true if both arguments are equal sets and false otherwise
$Set_1 != Set_2$	true if both arguments are unequal sets and false otherwise
$Set_1 \leq Set_2$	true if $Set_1$ is a subset of $Set_2$ and false otherwise
$Set_1 < Set_2$	true if $Set_1$ is a strict subset of $Set_2$ and false otherwise
$Set_1 \geq Set_2$	true if $Set_1$ is a superset of $Set_2$ and false otherwise
$Set_1 > Set_2$	true if $Set_1$ is a strict superset of $Set_2$ and false otherwise
$Set_1 + Set_2$	Set resulting from the union of the two arguments
$Set_1 - Set_2$	Set resulting from the difference of the two arguments
$Set_1 * Set_2$	Set resulting from the product of the two arguments
$Set_1 \& Set_2$	Set resulting from the intersection of the two arguments
$Elm \text{ in } Set$	true if $Elm$ occurs as element in $Set$ and false otherwise
$Elm \text{ not in } Set$	false if $Elm$ occurs as element in $Set$ and false otherwise
$Set_1 \text{ join } Set_2$	
$Bool ? Set_1 : Set_2$	

Examples:

```

rascal> {1, 2, 3} + {4, 5, 6};
set[int] : {1, 2, 3, 4, 5, 6}

rascal> {1, 2, 3} + {1, 2, 3};
set[int] : {1, 2, 3}

rascal> {1, 2, 3, 4} - {1, 2, 3};
set[int] : {4}

rascal> {1, 2, 3} - {4, 5, 6};
set[int] : {1, 2, 3}

rascal> {1, 2, 3} & {4, 5, 6};
set[int] : { }

rascal> {1, 2, 3} & {1, 2, 3};
set[int] : {1, 2, 3}

rascal> 3 in {1, 2, 3};
bool : true

rascal> 4 in {1, 2, 3};

```



```

bool : false

rascal> 3 notin {1, 2, 3};
bool : false

rascal> 4 notin {1, 2, 3};
bool : true

rascal> <2,20> in {<1,10>, <2,20>, <3,30>};
bool : true

rascal> <4,40> notin {<1,10>, <2,20>, <3,30>};
bool : true

```

**Table 1.27. Set Functions**

Function	Description
&T average(set[&T] st, &T zero)	Average of the elements of a set
&T getOneFrom(set[&T] st)	Pick a random element from a set
bool isEmpty(set[&T] st)	Is set empty?
set[&T] mapper(set[&T] st, &T (&T,&T) fn)	Apply a function to each element of a set
&T max(set[&T] st)	Largest element of a set
&T min(set[&T] st)	Smallest element of a set
&T multiply(set[&T] st, &T unity)	Multiply the elements of a set
set[set[&T]] power(set[&T] st)	All subsets of a set
set[set[&T]] power1(set[&T] st)	All subsets (excluding empty set) of a set
&T reducer(set[&T] st, &T (&T,&T) fn, &T unit)	Apply function F to successive elements of a set
int size(set[&T] st)	Number of elements in a set
&T sum(set[&T] st, &T zero)	Add the elements of a set
tuple[&T, set[&T]] takeOneFrom(set[&T] st)	Remove an arbitrary element from a set, returns the element and the modified set
list[&T] toList(set[&T] st)	Convert a set to a list
map[&A,&B] toMap(rel[&A, &B] st)	Convert a set of tuples to a map
str toString(set[&T] st)	Convert a set to a string

Examples:

```

rascal> power({1, 2, 3, 4});
set[set[int]] : { {}, {1}, {2}, {3}, {4}, {1,2}, {1,3}, {1,4},
                  {2,3}, {2,4}, {3,4}, {1,2,3}, {1,2,4},
                  {1,3,4}, {2,3,4}, {1,2,3,4}
                }

rascal> power1({1, 2, 3, 4});
set[set[int]] : { {1}, {2}, {3}, {4}, {1,2}, {1,3}, {1,4}, {2,3},
                  {2,4}, {3,4}, {1,2,3}, {1,2,4}, {1,3,4},
                  {2,3,4}, {1,2,3,4}
                }

```

```

rascal> size({1,2,3});
int : 3

rascal> size(<1,10>, <2,20>, <3,30>);
int : 3

```

## String

**Table 1.28. Operations on Strings**

Operator	Description
$Str_1 == Str_2$	Yields true if both arguments are equal and false otherwise
$Str_1 != Str_2$	Yields true if both arguments are unequal and false otherwise
$Str_1 <= Str_2$	Yields true if $Str_1$ is lexicographically less than or equal to $Str_2$ and false otherwise
$Str_1 < Str_2$	Yields true if $Str_1$ is lexicographically less than $Str_2$ and false otherwise
$Str_1 >= Str_2$	Yields true if $Str_1$ is lexicographically greater than or equal to $Str_2$ and false otherwise
$Str_1 > Str_2$	Yields true if $Str_1$ is lexicographically greater than $Str_2$ and false otherwise
$Str_1 + Str_2$	Concatenates $Str_1$ and $Str_2$
$Bool ? Str_1 : Str_2$	If $Bool$ is true then $Str_1$ else $Str_2$

**Table 1.29. String Functions**

Function	Description
int charAt(str s, int i) throws out_of_range(str msg)	Character at position i in string s.
bool endsWith(str s, str suffix)	Yields true if string s ends with given string suffix.
str center(str s, int n)	Center s in string of length n using spaces
str center(str s, int n, str pad)	Center s in string of length n using a pad character
bool isEmpty(str s)	Is string empty?
str left(str s, int n)	Left align s in string of length n using spaces
str left(str s, int n, str pad)	Left align s in string of length n using pad character
str right(str s, int n)	Right align s in string of length n using spaces
str reverse(str s)	String with all characters in reverse order.
int size(str s)	Length of string s.
bool startsWith(str s, str prefix)	Yields true if string s starts with the string prefix.
str toLowerCase(str s)	Convert all characters in string s to lowercase.
str toUpperCase(str s)	Convert all characters in string s to uppercase.

## Tuple

**Table 1.30. Tuple Operators**

Operator	Description
$Tuple_1 == Tuple_2$	true if both arguments are identical
$Tuple_1 != Tuple_2$	true if both arguments are not identical
$Tuple_1 <= Tuple_2$	true if both arguments are identical or if the leftmost element in $Tuple_1$ that differs from the corresponding in $Tuple_2$ is smaller than that element in $Tuple_2$
$Tuple_1 < Tuple_2$	true if both arguments are not identical and the leftmost element in $Tuple_1$ that differs from the corresponding in $Tuple_2$ is smaller than that element in $Tuple_2$
$Tuple_1 >= Tuple_2$	true if both arguments are identical or if the leftmost element in $Tuple_1$ that differs from the corresponding in $Tuple_2$ is greater than that element in $Tuple_2$
$Tuple_1 > Tuple_2$	true if both arguments are not identical and the leftmost element in $Tuple_1$ that differs from the corresponding in $Tuple_2$ is greater than that element in $Tuple_2$
$Tuple_1 + Tuple_2$	Concatenates $Tuple_1$ and $Tuple_2$
$Bool ? Tuple_1 : Tuple_2$	If $Bool$ is true then $Tuple_1$ else $Tuple_2$
$Tuple . Name$	Select field $Name$ from $Tuple$
$Tuple [ Int ]$	Select field at position $int$ from $Tuple$

## UnitTest

We provided a very rudimentary library for unit testing that will certainly evolve over time:

**Table 1.31. UnitTest Functions**

Function	Description
<code>void assertTrue(bool outcome)</code>	Check that outcome is true
<code>void assertEquals(value V1, value V2)</code>	Check that two values are equal
<code>bool report()</code>	Print unit test summary
<code>bool report(str msg)</code>	Print unit test summary, including msg

## Value

**Table 1.32. Value Operators**

Operator	Description
$Value_1 == Value_2$	true if both arguments are identical
$Value_1 != Value_2$	true if both arguments are not identical
$Value_1 <= Value_2$	
$Value_1 < Value_2$	
$Value_1 >= Value_2$	
$Value_1 > Value_2$	

Operator	Description
<i>Bool</i> ? <i>Value</i> <sub>1</sub> : <i>Value</i> <sub>2</sub>	if <i>Bool</i> is true then <i>Value</i> <sub>1</sub> else <i>Value</i> <sub>2</sub>

## ValueIO

**Table 1.33. ValueIO Functions**

Function	Description
value readValueFromBinaryFile(str namePBFFile)	Read a value from a binary file in PBF format
value readValueFromTextFile(str namePBFFile)	Read a value from a text file
void writeValueToBinaryFile(str namePBFFile, value val)	Write a value to a binary file in PBF format
void writeValueToTextFile(str namePBFFile, value val)	Write a value to a binary file in PBF format

## View (Eclipse only)

**Table 1.34. View Functions**

Function	Description
void show(value v)	Show value v in a graphical viewer
void browse(value v)	Show value v a graphical browser

## Void

There are no operators or functions defined on the type `void`.

# Table of Built-in Operators

Table 1.35. All Operators

Operator	Description	See Library Module/Section
$Exp_1[Name = Exp_2]$	Field assignment	Non-Boolean Operator Expressions (page 50)
$Exp . Name$	Field selection	Non-Boolean Operator Expressions (page 50), Location (page 75), Tuple (page 83)
$Exp < field, \dots >$	Field projection	Non-Boolean Operator Expressions (page 50), Relation (page 77)
$Exp_1 [ Exp_2 ]$	Index	Non-Boolean Operator Expressions (page 50), List (page 73), Map (page 75), Node (page 76), Relation (page 77), Tuple (page 83)
$Exp ?$	Isdefined: true is Exp has a well-defined value	Boolean Operator Expressions (page 52)
$! Exp$	Negation	Boolean Operator Expressions (page 52), Boolean (page 70)
$- Exp$	Negation	Integer (page 72), Real (page 77)
$Exp +$	Transitive closure	Relation (page 77)
$Exp *$	Reflexive transitive	Relation (page 77)
$Exp @ Name$	Attribute value	
$Exp_1 [ @ Name = Exp_2 ]$	Assign attribute value	
$Exp_1 \circ Exp_2$	Composition	Relation (page 77)
$Exp_1 \times Exp_2$	Cartesian product	Relation (page 77)
$Exp_1 / Exp_2$	Division	Integer (page 72), Real (page 77)
$Exp_1 \% Exp_2$	Modulo	Integer (page 72), Real (page 77)
$Exp_1 * Exp_2$	Multiplication/product	Integer (page 72), List (page 73), Real (page 77), Set (page 80)
$Exp_1 \& Exp_2$	Intersection	Set (page 80)
$Exp_1 + Exp_2$	Addition/ concatenation/union	Integer (page 72), List (page 73), Map (page 75), Real (page 77), Set (page 80), String (page 82), Tuple (page 83)
$Exp_1 - Exp_2$	Subtraction/difference	Integer (page 72), List (page 73), Map (page 75), Real (page 77), Set (page 80)
$Exp_1 \text{ join } Exp_2$	Join	Set (page 80)
$Exp_1 \text{ in } Exp_2$	Element of	Boolean Operator Expressions (page 52), List (page 73), Map (page 75), Set (page 80)
$Exp_1 \text{ not in } Exp_2$	Not element of	Boolean Operator Expressions (page 52), List (page 73), Map (page 75), Set (page 80)
$Exp_1 \leq Exp_2$	Less than/sublist / subset	Boolean Operator Expressions (page 52), Boolean (page 70), Integer (page 72), List (page 73), Location (page 75), Map (page 75), Node (page 76), Real (page 77), Set (page 80), String (page 82), Tuple (page 83), Value (page 83)

Operator	Description	See Library Module/Section
$Exp_1 < Exp_2$	Less than/strict sublist/ strict subset	Boolean Operator Expressions (page 52), Boolean (page 70), Integer (page 72), List (page 73), Location (page 75), Map (page 75), Node (page 76), Real (page 77), Set (page 80), String (page 82), Tuple (page 83), Value (page 83)
$Exp_1 >= Exp_2$	Greater than/superlist/ superset	Boolean Operator Expressions (page 52), Boolean (page 70), Integer (page 72), List (page 73), Location (page 75), Map (page 75), Node (page 76), Real (page 77), Set (page 80), String (page 82), Tuple (page 83), Value (page 83)
$Exp_1 > Exp_2$	Greater than/strict superlist/ strict superset	Boolean Operator Expressions (page 52), Boolean (page 70), Integer (page 72), List (page 73), Location (page 75), Map (page 75), Node (page 76), Real (page 77), Set (page 80), String (page 82), Tuple (page 83), Value (page 83)
$Pat := Exp$	Match	Boolean Operator Expressions (page 52)
$Pat !:= Exp$	No Match	Boolean Operator Expressions (page 52)
$Exp_1 == Exp_2$	Equality	Boolean Operator Expressions (page 52), Boolean (page 70), Integer (page 72), List (page 73), Location (page 75), Map (page 75), Node (page 76), Real (page 77), Set (page 80), String (page 82), Tuple (page 83), Value (page 83)
$Exp_1 != Exp_2$	Inequality	Boolean Operator Expressions (page 52), Boolean (page 70), Integer (page 72), List (page 73), Location (page 75), Map (page 75), Node (page 76), Real (page 77), Set (page 80), String (page 82), Tuple (page 83), Value (page 83)
$Exp_1 ? Exp_2$	Ifdefined Otherwise	Boolean Operator Expressions (page 52)
$Exp_1 ? Exp_2 : Exp_3$	Conditional Expression	Boolean Operator Expressions (page 52), Boolean (page 70), Integer (page 72), List (page 73), Map (page 75), Node (page 76), Real (page 77), Set (page 80), String (page 82), Tuple (page 83), Value (page 83)
$Exp_1 ==> Exp_2$	Implication	Boolean Operator Expressions (page 52), Boolean (page 70)
$Exp_1 <==> Exp_2$	Equivalence	Boolean Operator Expressions (page 52), Boolean (page 70)
$Exp_1 \&\& Exp_2$	Boolean and	Boolean Operator Expressions (page 52), Boolean (page 70)
$Exp_1    Exp_2$	Boolean or	Boolean Operator Expressions (page 52), Boolean (page 70)

# Table of Built-in Functions

**Table 1.36. All Functions**

Function	See Library Module
abs	Integer (page 72)
arbBool	Boolean (page 70)
arbInt	Integer (page 72)
arbReal	Real (page 77)
arity	Node (page 76)
assertEqual	UnitTest (page 83)
assertTrue	UnitTest (page 83)
average	List (page 73), Set (page 80)
bottom	Graph (page 71), Labelled Graph (page 73)
browse	View (Eclipse only) (page 84)
carrier	Relation (page 77)
carrierR	Relation (page 77)
carrierX	Relation (page 77)
center	String (page 82)
charAt ?remove	String (page 82)
complement	Relation (page 77)
currentTimeMillis	Benchmark (page 70)
delete	List (page 73)
domain	List (page 73), Map (page 75), Relation (page 77)
domainR	Relation (page 77)
domainX	Relation (page 77)
endsWith	String (page 82)
files	Resource (Eclipse only) (page 79)
fromInt	Boolean (page 70)
fromString	Boolean (page 70)
getChildren	Node (page 76)
getName	Node (page 76)
getOneFrom	List (page 73), Map (page 75), Set (page 80)
head	List (page 73)
insertAt	List (page 73)
invert	Map (page 75), Relation (page 77)
isEmpty	List (page 73), Map (page 75), Set (page 80), String (page 82)
left	String (page 82)
location ?name	Resource (Eclipse only) (page 79)
makeNode	Node (page 76)
mapper?	List (page 73), Map (page 75), Set (page 80)

Function	See Library Module
max	Integer (page 72), List (page 73), Real (page 77), Set (page 80)
min	Integer (page 72), List (page 73), Real (page 77), Set (page 80)
multiply	List (page 73), Set (page 80)
permutations	List (page 73)
power	Set (page 80)
power1	Set (page 80)
predecessors	Graph (page 71), Labelled Graph (page 73)
println	IO (page 73)
projects	Resource (Eclipse only) (page 79)
range	Map (page 75)
rangeR	Relation (page 77)
reach	Graph (page 71), Labelled Graph (page 73)
reachR	Graph (page 71), Labelled Graph (page 73)
reachX	Graph (page 71), Labelled Graph (page 73)
readFile	IO (page 73)
readRSF	RSF (page 79)
readValueFromBinaryFile	ValueIO (page 84)
readValueFromTextFile	ValueIO (page 84)
reducer	List (page 73), Set (page 80)
references	Resource (Eclipse only) (page 79)
report	UnitTest (page 83)
reverse	List (page 73), String (page 82)
right	String (page 82)
root	Resource (Eclipse only) (page 79)
shortestPathPair	Graph (page 71), Labelled Graph (page 73)
show	View (Eclipse only) (page 84)
size	List (page 73), Map (page 75), Set (page 80), String (page 82)
slice	List (page 73)
sort	List (page 73)
startsWith	String (page 82)
successors	Graph (page 71), Labelled Graph (page 73)
sum	List (page 73)
tail	List (page 73)
takeOneFrom	List (page 73), Set (page 80)
toInt	Boolean (page 70)
toInteger!	Real (page 77)
toList	Map (page 75), Set (page 80)
toLowerCase	String (page 82)
toMap	List (page 73), Set (page 80)



Function	See Library Module
top	Graph (page 71), Labelled Graph (page 73)
toReal	Boolean (page 70), Integer (page 72)
toRel	Map (page 75)
toSet	List (page 73)
toString	Boolean (page 70), Integer (page 72), List (page 73), Map (page 75), Real (page 77), Set (page 80)
toUpperCase	String (page 82)
writeValueToBinaryFile	ValueIO (page 84)
writeValueToTextFile	ValueIO (page 84)

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# Glossary of Terminology

Annotation	An annotation is a (name, value) pair that can be attached to a datatype value. The value can be retrieved and changed via the name. One datatype value can have more than one annotation.
ASF	Algebraic Specification Formalism. This a notation for describing rewrite rules and is mostly used for defining software analysis, fact extraction, and software transformation. It is the predecessor of Rascal.
ASF+SDF	The combination of the formalisms ASF and SDF. ASF+SDF can describe both the syntax of a language and the operations on that language (checking, execution, analysis, transformation). It is the predecessor of Rascal.
AsFix	<p>ASF+SDF Fixed format. A dataformat used to represent parse trees. AsFix is a specialized view on ATerms. Important features are:</p> <ul style="list-style-type: none"><li>• The AsFix format is a full parse tree that contains all the original layout and comments from the original source code program that was parsed.</li><li>• The AsFix format is self-descriptive: each subtree contains information about the exact grammar production that has been used to parse the text that has resulted in that parse tree.</li><li>• The AsFix format does not contain source code coordinates per se, but a separate tool can easily compute these coordinates and add them to the parse tree in the form of annotations.</li></ul> <p>AsFix is the predecessor of the UPTR format used by Rascal</p>
ATerm	<p>Annotated terms. A dataformat used for the internal representation of all data. Distinguishing features are:</p> <ul style="list-style-type: none"><li>• ATerms are language-independent and can be processed by programs in any language.</li><li>• ATerms can be annotated with auxiliary information that does not affect the tree structure.</li><li>• ATerms preserve <i>maximal subterm sharing</i>. This means that common parts of the data are not duplicated but shared. This leads to considerable size-reduction of the data.</li></ul> <p>ATerms are the predecessors of the shared PDB implementation used in Rascal.</p>
Backtracking	Backtracking is a general algorithmic method for finding all (or some) solutions to some computational problem, in the case of Rascal for pattern matching. It incrementally builds candidate matches, and abandons each partial match ("backtracks") as soon as it determines that it cannot possibly be completed to a valid match.

Bag	A bag can contain arbitrary values, without any particular order, and possibly repeated elements. All elements of a bag have the same type. <i>Rascal does not (yet?) support bags.</i>
Comprehension	An algorithmic schema for computing lists, sets or relations.
Constraint solving	The process of finding solutions to a problem that is only specified by the boundary conditions (constraints) that a solution should satisfy.
Control flow	The order in which the statements in a program can be executed. Usually represented as a control flow graph.
Debugger	A debugger allows the step-by-step execution of a program.
Dataflow	The order in which values may be assigned to variables during program execution. Usually represented as a data flow graph.
Eclipse	Interactive Development Environment (IDE) for Java and other languages. The Rascal IDE is built on top of Eclipse and IMP.
Enumerator	A Rascal construct that enumerates all the elements of a given value.
Function (higher order)	A function is a computational entity with typed formal parameters and a typed result. When applied to actual parameters, the function is applied and the result replaces the function call. A higher-order function is a function that has one or more functions as parameter.
Graph	In mathematics a graph is a set of objects (nodes) that may be pairwise connected with links (edges). In Rascal, a graph is a binary relation in which each tuple represents an edge in the corresponding (mathematical) graph.
Interactive Development Environment (IDE)	A programming environment for the interactive design, development, debugging and testing of programs. Eclipse is an example of an IDE.
IDE Meta Platform (IMP)	IMP is an extension of Eclipse and is intended to make it easier to add support for new languages to Eclipse. The Rascal implementation is part of IMP.
Location	Data format to describe locations in source code.
List	A list can contain arbitrary values, with a particular order, and possibly repeated elements. All elements of a list have the same type.
Map	A map is a binary relation consisting of 2-tuples (pairs) of the form (key, value) that associate each key with a unique value.
Meta-Environment	The ASF+SDF Meta-Environment is the IDE for writing ASF+SDF specifications.
Meta-programming	The activity of writing programs that manipulate other programs as data.
Parser	A program that performs syntactic analysis on a given input text and builds a parse tree.
Parse table	A parse table is produced by a parse table generator and is an efficient representation of grammar that can be used by a parser.

Parse table generator	A parse table generator takes a grammar as input and converts it to a parse table to be used by a parser.
Parse tree	Tree-structured representation of a text that has been analyzed by a parser.
Pattern	A structural description of values that is used for matching.
Pattern matching	Matching is an algorithmic method for determining the structural similarity of a pattern and a subject value. The pattern may contain variables that are bound to corresponding parts of the subject value when the match succeeds.
Program Data Base (PDB)	Subsystem of IMP that is responsible for implementing Rascal's datatypes.
Polymorphism	Functions and datatypes that can handle values of different types in a uniform manner.
Prettyprinter	A prettyprinter converts parse trees to text. The prettyprinter uses default rules to insert layout in a parse tree so that its corresponding text is presented in a uniform way. Optionally, user-defined formatting rules can replace this default behaviour.
Relation	A relation is a set of tuples. All tuples of a relation have the same type.
Rscript	A small scripting language for defining relational expressions. Used for the analysis of facts extracted from software. Predecessor of Rascal.
Relational Algebra	An algebra on relations that provides operators on relations that yield new relations. Relational algebra uses no variables and provides an operational view on relations.
Relational calculus	Relational calculus provides a more declarative and descriptive way to formulate operations on relations and depends on the use of variables.
Read-Eval-Print-Loop (REPL)	Command-line loop that reads user input, evaluates it and prints the result.
Rewrite rule	A rewrite rule consists of a pattern and a replacement. A rewrite rule is applied to all subparts of a datastructure and when the pattern matches the matching subpart is replaced. Pattern and replacement may contain variables and variables in the replacement are first replaced by their value before the replacement is applied.
SDF	Syntax Definition Formalism. A notation for describing the grammar of programming and application languages. SDF definitions can be imported in Rascal.
Set	A set can contain arbitrary values, without any particular order, and no repeated elements. All elements of a set have the same type.
Side-effect	In imperative programming languages like C and Java, assignments cause changes in a global state that are visible for other statements and procedures.

Syntax tree	Tree-structured representation of a text that has been analyzed by a parser.
Summary	An error or message summary. A dataformat for the internal representation of errors and messages. Summaries are produced by checker and compilers and are used by the GUI.
Term rewriting	Term rewriting is the act of applying a set of rewrite rules to a subject value. The rules are applied as long as any pattern in some rule matches. The result is the subject value with replacements made as prescribed by the matching rules.
Traversal	The process of visiting all the sub-elements of a larger datastructure. The <code>visit</code> expression and enumerators provide traversal functionality in Rascal.
Tuple	A tuple is an ordered list of values of possibly different type. An element of a tuple may have a label by which that element can be selected or modified.
Type checking	The activity to check that all entities in a program are declared and are used in a way that is consistent with their declared type.
Type inference	The activity to infer the type of program entities from their use.
UPTR	Universal Parse Tree Representation is the parse tree format used in Rascal. It is the successor of AsFix.