

# SFWR ENG 4E03

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Note: material covered in [Stats 3Y03 Summary](#) will not be covered in this summary

## Contents

Statistics .....	1
Variance .....	1
Exponential .....	2
Uniform .....	2
Binomial .....	2
Operations Analysis .....	2
Visitation Trick .....	3
DTMC.....	4
Matrices .....	4
CTMC.....	4
Poisson Process.....	4
Kendall notation.....	4

## Statistics

**Poisson parameter**  $[\lambda]$ : rate

**Service rate**  $[\mu]$ :

**Continuous Random Variable (CRV):**

$$\binom{n}{k} = \frac{n!}{k!(n-k)!}$$

*Think chemistry, i.e. cancelling units*

## Variance

$$\text{var}(X) = E[X^2] - (E[X])^2$$

- Don't change probability, but square X for calculation only

$$\begin{aligned}
 \text{var}(x) &= E[(X - \mu)^2] \\
 &= \sigma^2 \\
 &= \int_{-\infty}^{\infty} (x - \mu)^2 f(x) dx \\
 &= \int x^2 f(x) dx - \mu^2
 \end{aligned}$$

## Exponential

- **Mean**  $E[X]$ :  $1/\lambda$ 
  - a.k.a. Expected value
- **Variance**:
- **Probability Distribution Function (PDF)**  $P(X)$ :  $\lambda e^{-\lambda x}$
- **Cumulative Distribution Function (CDF)**  $f(x)$ :  $\text{CDF} = \int \text{PDF}$ , i.e.  $1 - e^{-\lambda x}$
- Memoryless
- not always for time

## Uniform

- **Mean**:  $(b-a)/2$
- **Variance**:  $(a+b)/2$
- **PDF**:  $1 / (b-a)$ ,  $a \leq x \leq b$
- **CDF**: 1
- **Uniform Distribution**: no memoryless property

## Binomial

- **Mean**  $E[X]$ :  $n \times \text{probability}$
- **Variance**:  $n \times p \times (1 - p)$
- **Probability Distribution Function (PDF)**  $P(X)$ :  $(n \text{ c } x) p^x (1-p)^{n-x}$
- **Cumulative Distribution Function (CDF)**  $f(x)$ :  $\sum_{i=0}^{\lfloor x \rfloor} \binom{n}{i} p_i (1-p)^{n-i}$

## Operations Analysis

**Device**  $[i]$ : units that are in terms of  $i$  are specific to an individual device or node within a system

**Total devices**  $[k]$ :

**Service Time**  $[S]$ : time per specific job

**Visitation**  $[V]$ : given or projected visits/jobs (closed system); cannot be calculated; basically a probability

$E(V)$ : calculated visit/job ratio

$P(\text{visit}) \cdot \text{total visits in previous node}$

**Demand**  $[D]$ : total service time for all jobs

$$D_i = E[S_i] \cdot V_i$$

$$D = \sum_{i=0}^k D_i$$

**Bottleneck**  $[D_{\max}]$ : device with largest demand

**Time in system**  $[T]$ : time the job is in the system

$$E[T] = \frac{N}{X}$$

$$E[T] \geq \max(D, ND_{\max} - E[Z])$$

If  $E[Z] = 0$ ,  $T = R$

**Response Time**  $[R]$ : time the job is *being processed* in the system

If  $E[Z] = 0$ ,  $R = T$

**Users**  $[M]$ :

**Optimal users**  $[M^*]$ :

$$M^* = \frac{D + E[Z]}{D_{\text{bottleneck}}}$$

**Total Jobs**  $[N]$ :  $N=M$  in a closed system

- Little's Law:  $E[N] = \lambda E[T]$ ,  $\lambda = X$
- $E[N] = \lambda E[R]$ ,  $\lambda = X$

**Think time**  $[Z]$ : time it takes the user to put a request in and start, it's kinda like the frequency that users put in requests (seconds / request)

$$E[Z] = E[T] - E[R]$$

**Throughput**  $[X]$ : out-rate, max jobs / hour of full system

$$X \leq \min\left(\frac{1}{D_{\max}}, \frac{N}{D + E[Z]}\right)$$

Note:  $\frac{1}{D_{\max}}$  and  $\frac{N}{D + E[Z]}$  converge at their lowest point, so equate them

$$X_i = E[V_i] X$$

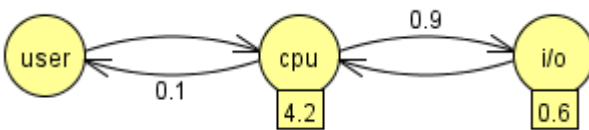
**Utilization**  $[\rho]$ : ratio that the time is busy

$$\rho_i = X_i E[S_i]$$

$$\rho_i = X D_i$$

Visitation Trick

If determining visitation at a node, establish a reference node from one of the incoming nodes, usually the user node, that has a returning percentage



$$V_{\text{user}} = 1 = 0.1 \cdot V_{\text{cpu}}$$

## DTMC

### Discrete Time Markov Chains (DTMC):

**Geometric Series:**  $\sum_{i=0}^{\infty} r^i = \frac{1}{1-r}$ , where  $0 \leq r \leq 1$  (because otherwise it would be unstable)

$$\sum_{i=1}^{\infty} r^i = \frac{r}{1-r}$$

**Geometric Sequence:**  $S_n = \sum_{i=0}^n r^i = \frac{1-r^{n+1}}{1-r}$

$$S_n = \sum_{i=1}^n r^i = \frac{r(1-r^n)}{1-r}$$

Steady state:  $n \rightarrow \infty$

For discrete: use the sum of the  $X$ 's, so  $E[X] = \sum (P(X=i) \cdot X_i)$  and  $E[X^2] = \sum (P(X=i) \cdot X_i^2)$

## Matrices

Rows: equations for nodes going out (add up to 1)

Columns: equations for nodes coming in

## CTMC

### Poisson Process

#### Poisson Process:

- Uses Exponential Distribution
- $\lambda_{\text{total}} = \sum \lambda_i$ 
  - you can also split up  $\lambda$  into multiple  $\lambda$ s
- Not only do you see each second as time independent, each stream of probabilities is independent
- $P(x; \lambda) = e^{-\lambda} \lambda^x / x!$ 
  - $[x]$ : things will happen
  - $[\lambda]$ : rate;  $\lambda = \alpha t$
- $[\alpha]$ : expected number of events during unit interval
- $[t]$ : time interval length
- $P_x(t) = \frac{e^{-\alpha t} \cdot (\alpha t)^x}{x!}$

## Kendall notation

**Job Processing time**  $[\mu]$ : rate of jobs leaving system

M/M/1 Queue

$[M]$ : time between arrivals is Markovian (Memoryless)  $\sim \exp(\lambda)$

$[M]$ : job processing times are Markovian (Memoryless)  $\sim \exp(\mu)$

$[1]$ : single server

Attributes:

- FIFO
- Infinite buffer

#### Variations

- M/M/2 Queue: same, except 2 servers
- M/M/C Queue:  $C$  servers
- M/E<sub>k</sub>/C: Erlang  $k$ , i.e. series of exponential
- H()/M/C: hyperexponential distribution
- PH/M/C: phase type, i.e. any combination of any number of exponentials with any rate
- M/G/C: General distribution
- G/G/1: has not been solved yet
- M/M/1/1: 1 server, 1 job

**Erlang-C Equation:**  $P(\text{job has to wait in queue}) = \sum_{i=0}^{\infty} \pi_i$

$$= \frac{1}{c!} \left( \frac{\lambda}{\mu} \right)^c \left( \frac{1}{1-\rho} \right) \pi_0$$