

# SFWR ENG 4E03

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Note: material covered in [Stats 3Y03 Summary](#) will not be covered in this summary

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## Statistics

**Expected Value**  $[\mu]$ : definition of expected (NOT RIGHT!!)

**Poisson parameter**  $[\lambda]$ :

**Exponential distribution**: not always for time

**Probability Distribution Function (PDF)**:

**Cumulative Distribution Function (CDF)**:

**Uniform Distribution**: no memoryless property

**Exponential Distribution**:

- Memoryless
- Either CDF or PDF of original equation  $F = 1 - e^{-\lambda x}$

*Think chemistry, i.e. cancelling units*

## Variance

$$\text{var}(X) = E[X^2] - (E[X])^2$$

- Don't change probability, but square X for calculation only
- For discrete: use the sum of the X's, so  $E[X] = \sum(P(X=i) \cdot X_i)$  and  $E[X^2] = \sum(P(X=i) \cdot X_i^2)$

**Continuous Random Variable (CRV)**:

$$\begin{aligned}
 \text{Var}(x) &= E[(X - \mu)^2] \\
 &= \sigma^2 \\
 &= \int_{-\infty}^{\infty} (x - \mu)^2 f(x) dx \\
 &= \int x^2 f(x) dx - \mu^2
 \end{aligned}$$

## Exponential

- **Mean**  $[E[X]]$ :  $1/\lambda$
- **Variance**  $[E[X]]$ : a.k.a. Expected value
- **Probability Distribution Function (PDF)**  $[P(X)]$ :  $\lambda e^{-\lambda x}$
- **Cumulative Distribution Function (CDF)**  $[f(x)]$ :  $\text{CDF} = \int \text{PDF}$ , i.e.  $1 - e^{-\lambda x}$

## Uniform

- **Mean**:  $(b-a)/2$
- **Variance**:  $(b-a)^2/12$
- **PDF**:  $1/(b-a)$ ,  $a \leq x \leq b$
- **CDF**:  $x - a$

## Operations Analysis

**Device**  $[i]$ : units that are in terms of  $i$  are specific to an individual device or node within a system

**Total devices**  $[k]$ :

**Service Time**  $[S]$ : time per specific job

**Visitation**  $[V]$ : given or projected visits/jobs (closed system); cannot be calculated; basically a probability

$[E(V)]$ : calculated visit/job ratio

$P(\text{visit}) \cdot \text{total visits in previous node}$

**Demand**  $[D]$ : total service time for all jobs

$$D_i = E[S_i] \cdot V_i$$

$$D = \sum_{i=0}^k D_i$$

**Time in system**  $[T]$ : time the job is in the system

$$E[T] = \frac{N}{X}$$

$$E[T] \geq \max(D, ND_{\max} - E[Z])$$

**Response Time**  $[R]$ : time the job is *being processed* in the system

**Users**  $[M]$ :

**Optimal users**  $[M^*]$ :

$$M^* = \frac{D + E[Z]}{D_{\text{bottleneck}}}$$

**Total Jobs [N]:**  $N=M$  in a closed system

$$E[N] = \lambda E[T], \lambda = X$$

If  $E[Z] = 0$ ,  $R = N$

$$E[N] = \lambda E[R], \lambda = X$$

**Think time [Z]:** time it takes the user to put a request in and start, it's kinda like the frequency that users put in requests (seconds / request)

$$E[Z] = E[T] - E[R]$$

**Throughput [X]:** out-rate, jobs / hour of full system

$$X \leq \min\left(\frac{1}{D_{\max}}, \frac{N}{D + E[Z]}\right)$$

Note:  $\frac{1}{D_{\max}}$  and  $\frac{N}{D + E[Z]}$  converge at their lowest point, so equate them

$$X_i = E[V_i] X$$

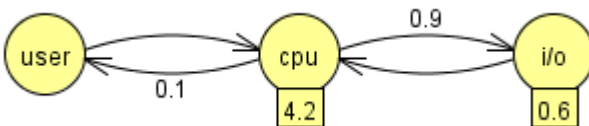
**Utilization [ $\rho$ ]:** ratio that the time is busy

$$\rho_i = X_i E[S_i]$$

$$\rho_i = X D_i$$

Visitation Trick

If determining visitation at a node, establish a reference node from one of the incoming nodes, usually the user node, that has a returning percentage



$$V_{\text{user}} = 1 = 0.1 \cdot V_{\text{CPU}}$$

DTMC

**Discrete Time Markov Chains (DTMC):**

**Geometric Series:**  $\sum_{i=0}^{\infty} r^i = \frac{r}{1-r}$

**Geometric Sequence:**  $S_n = \sum_{i=0}^n r^i = \frac{1-r^{n+1}}{1-r}$

$$S_n = \sum_{i=1}^n r^i = \frac{r(1-r^n)}{1-r}$$