Random Variable

Definition

A random variable X on a sample space Ω is a real-valued probability function on Ω ; that is, $X:\Omega\to\mathcal{R}$. A discrete random variable is a random variable that takes on only a finite or countably infinite number of values.

Examples:

- In rolling a dice, the number that comes up is a random variable.
- 2 Consider a gambling game in which a player flips two coins, if he gets head in both coins we wins \$3, else he losses \$1. The payoff of the game is a random variable.

Independence

Definition

Two random variables X and Y are independent if and only if

$$Pr((X = x) \cap (Y = y)) = Pr(X = x) \cdot Pr(Y = y)$$

for all values x and y. Similarly, random variables $X_1, X_2, ... X_k$ are mutually independent if and only if for any subset $I \subseteq [1, k]$ and any values $x_i, i \in I$,

$$\Pr\left(\bigcap_{i\in I}X_i=x_i\right) = \prod_{i\in I}\Pr(X_i=x_i).$$

Expectation

Definition

The expectation of a discrete random variable X, denoted by E[X], is given by

$$\mathbf{E}[X] = \sum_{i} i \Pr(X = i),$$

where the summation is over all values in the range of X. The expectation is finite if $\sum_{i} |i| \Pr(X = i)$ converges; otherwise, the expectation is unbounded.

The expectation (or mean or average) is a weighted sum over all possible values of the random variable.

Examples:

The expected value of one dice roll is:

$$E[X] = \sum_{i=1}^{6} iPr(X=i) = \sum_{i=1}^{6} \frac{i}{6} = 3\frac{1}{2}.$$

 The expectation of the random variable X representing the sum of two dice is

$$\mathbf{E}[X] = \frac{1}{36} \cdot 2 + \frac{2}{36} \cdot 3 + \frac{3}{36} \cdot 4 + \dots + \frac{1}{36} \cdot 12 = 7.$$

• Let X take on the value 2^i with probability $1/2^i$ for i = 1, 2, ...

$$\mathbf{E}[X] = \sum_{i=1}^{\infty} \frac{1}{2^i} 2^i = \sum_{i=1}^{\infty} 1 = \infty.$$

Median

Definition

The **median** of a random variable X is a value m such

$$Pr(X < m) \le 1/2$$
 and $Pr(X > m) < 1/2$.

Consider a game in which a player chooses a number in [1, ..., 6] and then rolls 3 dice.

The player wins \$1 for each dice the matches the number, he losses \$1 if no dice matches the number.

What is the expected outcome of that game:

$$-1(\frac{5}{6})^3+1\cdot 3(\frac{1}{6})(\frac{5}{6})^2+2\cdot 3(\frac{1}{6})^2(\frac{5}{6})+3(\frac{1}{6})^3=-\frac{17}{216}.$$

Linearity of Expectation

Theorem

For any two random variables X and Y

$$E[X+Y] = E[X] + E[Y].$$

$$E[X + Y] =$$

$$\sum_{i \in range(X)} \sum_{j \in range(Y)} (i + j)Pr((X = i) \cap (Y = j)) =$$

$$\sum_{i} \sum_{j} iPr((X = i) \cap (Y = j)) +$$

$$\sum_{i} \sum_{j} jPr((X = i) \cap (Y = j)) =$$

 $\sum_{i} iPr(X=i) + \sum_{i} jPr(Y=j).$

(Since we sum over all possible choices of i(j).)

Lemma

For any constant c and discrete random variable X,

$$\mathbf{E}[cX] = c\mathbf{E}[X].$$

Proof.

The lemma is obvious for c = 0. For $c \neq 0$,

$$\mathbf{E}[cX] = \sum_{j} j \Pr(cX = j)$$

$$= c \sum_{j} (j/c) \Pr(X = j/c)$$

$$= c \sum_{k} k \Pr(X = k)$$

$$= c \mathbf{E}[X].$$

Examples:

- The expectation of the sum of two dice is 7, even if they are not independent.
- The expectation of the outcome of one dice plus twice the outcome of a second dice is $10\frac{1}{2}$.
- Assume that we flip N coins, what is the expected number of heads?

Using linearity of expectation we get $N \cdot \frac{1}{2}$.

By direct summation we get $\sum_{i=0}^{N} i {N \choose i} 2^{-N}$.

Thus we prove

$$\sum_{i=0}^{N} i \binom{N}{i} 2^{-N} = \frac{N}{2}.$$

Assume that N people checked coats in a restaurants. The coats are mixed and each person gets a random coat.

How many people got their own coats?

It's hard to compute $E[X] = \sum_{k=0}^{N} kPr(X = k)$. Instead we define N 0-1 random variables X_i , where $X_i = 1$ iff i got his coat.

$$E[X_i] = 1 \cdot Pr(X_i = 1) + 0 \cdot Pr(X_i = 0) =$$

$$Pr(X_i = 1) = \frac{1}{N}.$$

$$E[X] = \sum_{i=1}^{N} E[X_i] = 1.$$

Bernoulli Random Variable

A Bernoulli or an indicator random variable:

$$Y = \begin{cases} 1 & \text{if the experiment succeeds,} \\ 0 & \text{otherwise.} \end{cases}$$

$$\mathbf{E}[Y] = p \cdot 1 + (1 - p) \cdot 0 = p = \Pr(Y = 1).$$

Binomial Random Variable

Definition

A binomial random variable X with parameters n and p, denoted by B(n,p), is defined by the following probability distribution on $j=0,1,2,\ldots,n$:

$$\Pr(X=j) = \binom{n}{j} p^{j} (1-p)^{n-j}.$$

Expectation of a Binomial Random Variable
$$\mathbf{E}[X] = \sum_{j=0}^{n} j \binom{n}{j} p^{j} (1-p)^{n-j}$$

$$= \sum_{j=0}^{n} j \frac{n!}{j!(n-j)!} p^{j} (1-p)^{n-j}$$

$$= \sum_{j=0}^{n} \frac{n!}{(j-1)!(n-j)!} p^{j} (1-p)^{n-j}$$

$$\sum_{j=0}^{n} \frac{j!(n-j)!}{(j-1)!(n-j)!} p^{j} (1-p)^{n-j}
= np \sum_{j=1}^{n} \frac{(n-1)!}{(j-1)!((n-1)-(j-1))!} p^{j-1} (1-p)^{(n-1)-(j-1)}
= np \sum_{j=0}^{n-1} \frac{(n-1)!}{k!((n-1)-k)!} p^{k} (1-p)^{(n-1)-k}$$

$$= np \sum_{k=0}^{\infty} \frac{1}{k!((n-1)-k)!} p^{n} (1-p)^{(n-1)}$$

$$= np \sum_{k=0}^{n-1} {n-1 \choose k} p^{k} (1-p)^{(n-1)-k} = np.$$

Using linearity of expectations

$$\mathbf{E}[X] = \mathbf{E}\left[\sum_{i=1}^{n} X_i\right] = \sum_{i=1}^{n} \mathbf{E}[X_i] = np.$$

Randomized Algorithms:

- Analysis is true for any input.
- The sample space is the space of random choices made by the algorithm.
- Repeated runs are independent.

Probabilistic Analysis;

- The sample space is the space of all possible inputs.
- If the algorithm is deterministic repeated runs give the same output.

Algorithm classification

A **Monte Carlo Algorithm** is a randomized algorithm that may produce an incorrect solution.

For decision problems: A **one-side error** Monte Carlo algorithm errs only one one possible output, otherwise it is a **two-side error** algorithm.

A **Las Vegas** algorithm is a randomized algorithm that **always** produces the correct output.

In both types of algorithms the run-time is a random variable.

Compund events:

- A program that has one call to a process S.
- Each call to process S recursively spawns new copies of the process S, where the number of new copies is a binomial random variable with parameters n and p.
- These random variables are independent for each call to \mathcal{S} .
- What is the expected number of copies of the process S generated by the program?

Conditional Expectation

Definition

$$\mathbf{E}[Y \mid Z = z] = \sum_{y} y \Pr(Y = y \mid Z = z),$$

where the summation is over all y in the range of Y.

Example

We role two dice. X_1 be the number that shows on the first die, X_2 be the number on the second die, and X be the sum of the numbers on the two dice.

$$\mathbf{E}[X \mid X_1 = 2] = \sum_{X} x \Pr(X = x \mid X_1 = 2) = \sum_{X=2}^{8} x \cdot \frac{1}{6} = \frac{11}{2}.$$

As another example, consider $E[X_1 \mid X = 5]$.

$$\mathbf{E}[X_1 \mid X = 5] = \sum_{x=1}^{4} x \Pr(X_1 = x \mid X = 5)$$

$$= \sum_{x=1}^{4} x \frac{\Pr(X_1 = x \cap X = 5)}{\Pr(X = 5)}$$

$$= \sum_{x=1}^{4} x \frac{1/36}{4/36}$$

$$= 5/2.$$

Lemma

For any random variables X and Y,

$$\mathsf{E}[Y] = \sum \mathsf{D}_{\mathsf{r}}(V = v)$$

 $\mathbf{E}[X] = \sum_{y} \Pr(Y = y) E[X \mid Y = y],$ where the sum is over all values in the range of **Y**.

Proof.

$$\sum_{y} \Pr(Y = y)E[X \mid Y = y]$$

$$= \sum_{y} \Pr(Y = y) \sum_{x} x \Pr(X = x \mid Y = y)$$

$$= \sum_{x} \sum_{y} x \Pr(X = x \mid Y = y) \Pr(Y = y)$$

$$= \sum_{x} \sum_{y} x \Pr(X = x \cap Y = y)$$

$$= \sum_{x} x \Pr(X = x) = E[X].$$



Conditional Expectation as a Random variable

Definition

The expression $\mathbf{E}[Y \mid Z]$ is a random variable f(Z) that takes on the value $\mathbf{E}[Y \mid Z = z]$ when Z = z.

Consider the outcome of rolling two dice $X_1, X_2, X = X_1 + X_2$.

$$\mathbf{E}[X \mid X_1] = \sum_{x} x \Pr(X = x \mid X_1) = \sum_{x = X_1 + 1}^{X_1 + 6} x \cdot \frac{1}{6} = X_1 + \frac{7}{2}.$$

If $\mathbf{E}[Y \mid Z]$ is a random variable, it has an expectation.

Theorem

$$\mathbf{E}[Y] = \mathbf{E}[\mathbf{E}[Y \mid Z]].$$

$$\mathbf{E}[X \mid X_1] = X_1 + \frac{7}{2}.$$

Thus

$$\mathbf{E}[\mathbf{E}[X \mid X_1]] = \mathbf{E}\left[X_1 + \frac{7}{2}\right] = \frac{7}{2} + \frac{7}{2} = 7.$$

Proof.

 $\mathbf{E}[Y \mid Z] = f(Z)$, where f(Z) takes on the value $\mathbf{E}[Y \mid Z = z]$ when Z = z.

$$\mathbf{E}[\mathbf{E}[Y \mid Z]] = \sum_{z} \mathbf{E}[Y \mid Z = z] \Pr(Z = z)$$

$$= \sum_{z} \left(\sum_{y} y \Pr(Y = y \mid Z = z) \right) \Pr(Z = z)$$
$$= \sum_{z} \sum_{y} y \Pr(Y = y \mid Z = z) \Pr(Z = z)$$

$$= \sum_{z} \sum_{y} y \Pr(Y = y \cap Z = z)$$

$$= \sum_{z}^{z} y \sum_{y} \Pr(Y = y \cap Z = z)$$

$$= \sum y \Pr(Y = y) = \mathbf{E}[Y].$$

Back to the Spawning Process

- The initial process S is in generation 0.
- A process S is in generation i if it was spawned by another process S in generation i-1.
- Let Y_i denote the number of S processes in generation i.
- $Y_0 = 1$, and Y_1 has a binomial distribution.

$$\mathbf{E}[Y_1] = np.$$

- Z_k^{i-1} = number of copies spawned by the kth process spawned in the (i-1)-st generation.
- Z_{ν}^{i-1} is a binomial random variable with parameters n and p.

•

$$\mathbf{E}[Y_{i} \mid Y_{i-1} = y_{i-1}] = \mathbf{E}\left[\sum_{k=1}^{y_{i-1}} Z_{k}\right]$$

$$= \sum_{k=1}^{y_{i-1}} \mathbf{E}[Z_{k}]$$

$$= y_{i-1}np.$$

$$E[Y_i] = E[E[Y_i \mid Y_{i-1}]] = E[Y_{i-1}np] = npE[Y_{i-1}].$$

• By induction on i, and using $Y_0 = 1$,

$$E[Y_i] = (np)^i.$$

$$E[\sum_{i>0} Y_i] = \sum_{i>0} E[Y_i] = \sum_{i>0} (np)^i.$$

• If $np \ge 1$, the expectation is unbounded, and if np < 1, the expectation is 1/(1-np).

The Geometric Distribution

Definition

A geometric random variable X with parameter p is given by the following probability distribution on n = 1, 2, ...

$$\Pr(X = n) = (1 - p)^{n-1}p.$$

memoryless property

Lemma

For a geometric random variable with parameter p and n > 0,

$$Pr(X = n + k \mid X > k) = Pr(X = n).$$

Proof.

$$\Pr(X = n + k \mid X > k) = \frac{\Pr((X = n + k) \cap (X > k))}{\Pr(X > k)}$$

$$= \frac{\Pr(X = n + k)}{\Pr(X > k)} = \frac{(1 - p)^{n + k - 1} p}{\sum_{i=k}^{\infty} (1 - p)^{i} p}$$

$$= \frac{(1 - p)^{n + k - 1} p}{(1 - p)^{k}} = (1 - p)^{n - 1} p$$

$$= \Pr(X = n).$$

Lemma

Let X be a discrete random variable that takes on only non-negative integer values. Then

$$\mathbf{E}[X] = \sum_{i=1}^{\infty} \Pr(X \ge i).$$

Proof.

$$\sum_{i=1}^{\infty} \Pr(X \ge i) = \sum_{i=1}^{\infty} \sum_{j=i}^{\infty} \Pr(X = j)$$

$$= \sum_{j=1}^{\infty} \sum_{i=1}^{j} \Pr(X = j)$$

$$= \sum_{i=1}^{\infty} j \Pr(X = j) = \mathbf{E}[X].$$

For a geometric random variable X with parameter p,

$$\Pr(X \ge i) = \sum_{i=0}^{\infty} (1-p)^{n-1} p = (1-p)^{i-1}.$$

$$\mathbf{E}[X] = \sum_{i=1}^{\infty} \Pr(X \ge i)$$

$$= \sum_{i=1}^{\infty} (1 - p)^{i-1}$$

$$= \frac{1}{1 - (1 - p)}$$

$$= \frac{1}{p}$$

Alternative Proof

$$Y = 1$$
 if $X = 1$, else $Y = 0$.

$$\begin{aligned} \mathbf{E}[X] &= & \Pr(Y=0)\mathbf{E}[X \mid Y=0] + \Pr(Y=1)\mathbf{E}[X \mid Y=1] \\ &= & (1-\rho)\mathbf{E}[X \mid Y=0] + \rho\mathbf{E}[X \mid Y=1]. \end{aligned}$$

When X > 1, let Z = X - 1.

$$E[X] = (1-p)E[Z+1] + p \cdot 1 = (1-p)E[Z] + 1,$$

By the memoryless property Z is also a geometric random variable with parameter p. Hence $\mathbf{E}[Z] = \mathbf{E}[X]$.

$$\mathbf{E}[X] = (1 - p)\mathbf{E}[Z] + 1 = (1 - p)\mathbf{E}[X] + 1,$$
 which yields $\mathbf{E}[X] = 1/p$.

Example: Coupon Collector's Problem

Suppose that each box of cereal contains a random coupon from a set of n different coupons.

How many boxes of cereal do you need to buy before you obtain at least one of every type of coupon?

Let X be the number of boxes bought until at least one of every type of coupon is obtained.

Let X_i be the number of boxes bought while you had exactly i-1 different coupon.

$$X = \sum_{i=1}^{n} X_i$$

 X_i is a geometric random variable with parameter

$$p_i = 1 - \frac{i-1}{n}$$
.

$$\mathsf{E}[X_i] = \frac{1}{p_i} = \frac{n}{n-i+1}.$$

$$\mathbf{E}[X] = E\left[\sum_{i=1}^{n} X_i\right]$$

$$\mathbf{E}[X] = E\left[\sum_{i=1}^{n} X_{i}\right]$$
$$= \sum_{i=1}^{n} \mathbf{E}[X_{i}]$$

$$= \sum_{i=1}^{n} \mathbf{E}[X_i]$$

$$= \sum_{i=1}^{n} \frac{n}{n-i+1}$$

 $= n \sum_{i=1}^{n} \frac{1}{i} = n \ln n + \Theta(n).$