

# Chapter 2

## PRAM: Matrix

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# Matrix addition

$A[n \times n]$

|          |          |     |          |
|----------|----------|-----|----------|
| $a_{11}$ | $a_{12}$ | ... | $a_{1n}$ |
| $a_{21}$ | $a_{22}$ | ... | $a_{2n}$ |
| ...      | ...      | ... | ...      |
| $a_{n1}$ | $a_{n2}$ | ... | $a_{nn}$ |

+

$B[n \times n]$

|          |          |     |          |
|----------|----------|-----|----------|
| $b_{11}$ | $b_{12}$ | ... | $b_{1n}$ |
| $b_{21}$ | $b_{22}$ | ... | $b_{2n}$ |
| ...      | ...      | ... | ...      |
| $b_{n1}$ | $b_{n2}$ | ... | $b_{nn}$ |

=

$C[n \times n]$

|          |          |     |          |
|----------|----------|-----|----------|
| $c_{11}$ | $c_{12}$ | ... | $c_{1n}$ |
| $c_{21}$ | $c_{22}$ | ... | $c_{2n}$ |
| ...      | ...      | ... | ...      |
| $c_{n1}$ | $c_{n2}$ | ... | $c_{nn}$ |

$C[n \times n]$

|                            |                            |     |                            |
|----------------------------|----------------------------|-----|----------------------------|
| $c_{11} = a_{11} + b_{11}$ | $c_{12} = a_{12} + b_{12}$ | ... | $c_{1n} = a_{1n} + b_{1n}$ |
| $c_{21} = a_{21} + b_{21}$ | $c_{22} = a_{22} + b_{22}$ | ... | $c_{2n} = a_{2n} + b_{2n}$ |
| ...                        | ...                        | ... | ...                        |
| $c_{n1} = a_{n1} + b_{n1}$ | $c_{n2} = a_{n2} + b_{n2}$ | ... | $c_{nn} = a_{nn} + b_{nn}$ |

# Matrix addition: PRAM with $n \times n$ processors

$C[n \times n]$

|                            |                            |     |                            |
|----------------------------|----------------------------|-----|----------------------------|
| $c_{11} = a_{11} + b_{11}$ | $c_{12} = a_{12} + b_{12}$ | ... | $c_{1n} = a_{1n} + b_{1n}$ |
| $P_{11}$                   | $P_{12}$                   |     | $P_{1n}$                   |
| $c_{21} = a_{21} + b_{21}$ | $c_{22} = a_{22} + b_{22}$ | ... | $c_{2n} = a_{2n} + b_{2n}$ |
| $P_{21}$                   | $P_{22}$                   |     | $P_{2n}$                   |
| ...                        | ...                        | ... | ...                        |
| ...                        | ...                        | ... | ...                        |
| $c_{n1} = a_{n1} + b_{n1}$ | $c_{n2} = a_{n2} + b_{n2}$ | ... | $c_{nn} = a_{nn} + b_{nn}$ |
| $P_{n1}$                   | $P_{n2}$                   |     | $P_{nn}$                   |

➤ All ( $n \times n$ ) processors run '+' in parallel:  $O(1)$   $\Rightarrow$

▪  $P_{ij}$  processor:

- Read:  $a_{ij}$  &  $b_{ij}$
- Write:  $c_{ij}$

No overlapping data

- CRCW:  $O(1)$
- EREW:  $O(1)$



# Matrix addition: PRAM with $n$ processors

➤ All  $n$  processors run '+' in parallel in  $n$  steps:  $O(n)$  ??? =>

$C[n \times n]$

| Step 1                     | Step 2                     | ... | Step $n$                   |
|----------------------------|----------------------------|-----|----------------------------|
| $c_{11} = a_{11} + b_{11}$ | $c_{12} = a_{12} + b_{12}$ | ... | $c_{1n} = a_{1n} + b_{1n}$ |
| $P_1$                      | $P_1$                      | ... | $P_1$                      |
| $c_{21} = a_{21} + b_{21}$ | $c_{22} = a_{22} + b_{22}$ | ... | $c_{2n} = a_{2n} + b_{2n}$ |
| $P_2$                      | $P_2$                      | ... | $P_2$                      |
| ...                        | ...                        | ... | ...                        |
| ...                        | ...                        | ... | ...                        |
| $c_{n1} = a_{n1} + b_{n1}$ | $c_{n2} = a_{n2} + b_{n2}$ | ... | $c_{nn} = a_{nn} + b_{nn}$ |
| $P_n$                      | $P_n$                      | ... | $P_n$                      |

▪  $P_{ij}$  processor:

○ Read:  $a_{ij}$  &  $b_{ij}$

○ Write:  $c_{ij}$

No overlapping data

○ CRCW:  $O(n)$

○ EREW:  $O(n)$

Your algorithm with  
 $k$  processors ( $k \ll n$ )?



# Matrix multiplication

$$C_{ij} = a_{i1} * b_{1j} + a_{i2} * b_{2j} + \dots + a_{in} * b_{nj}$$

A[nxn]

|          |          |     |     |          |
|----------|----------|-----|-----|----------|
| $a_{11}$ | $a_{12}$ | ... | ... | $a_{1n}$ |
| ...      | ...      | ... | ... | ...      |
| $a_{i1}$ | $a_{i2}$ | ... | ... | $a_{in}$ |
| ...      | ...      | ... | ... | ...      |
| $a_{n1}$ | $a_{n2}$ | ... | ... | $a_{nn}$ |

B[nx n]

|          |     |          |     |          |
|----------|-----|----------|-----|----------|
| $b_{11}$ | ... | $b_{1j}$ | ... | $b_{1n}$ |
| ...      | ... | $b_{2j}$ | ... | ...      |
| ...      | ... | ...      | ... | ...      |
| ...      | ... | ...      | ... | ...      |
| $b_{n1}$ | ... | $b_{nj}$ | ... | $b_{nn}$ |

C[nxn]

|          |     |          |     |          |
|----------|-----|----------|-----|----------|
| $c_{11}$ | ... |          | ... | $c_{1n}$ |
| ...      | ... |          | ... | ...      |
|          |     | $c_{ij}$ |     |          |
| ...      | ... | ...      | ... | ...      |
| $c_{n1}$ | ... | $c_{nj}$ | ... | $c_{nn}$ |

- $C_{ij} = a_{i1} * b_{1j} + a_{i2} * b_{2j} + \dots + a_{in} * b_{nj}$
- Vector:  $A[i^{\text{th}} \text{ row}] \times B[j^{\text{th}} \text{ column}]$
- Number of operations =  $\underbrace{n * \text{opt}(' * ')}_{O(1)} + \underbrace{(n-1) * \text{opt}(' + ')}_{O(\log(n))}$
- $n$  processors  $\sim O(\log(n))$

$C[n \times n]$

|          |     |          |     |          |
|----------|-----|----------|-----|----------|
| $C_{11}$ | ... |          | ... | $C_{1n}$ |
| ...      | ... |          | ... | ...      |
|          |     | $C_{ij}$ |     |          |
| ...      | ... | ...      | ... | ...      |
| $C_{n1}$ | ... | $C_{nj}$ | ... | $C_{nn}$ |



# PRAM

## Matrix multiplication

- $C_{ij}$  using  $n$  processors:  $O(\log(n))$
- $C[n \times n]: (n*n)C_{ij}$ 
  - $n$  processors:  $O(n^2 * \log(n))$
  - $n^3$  processors:  $O(\log(n))$
  - $n^2$  processors???

$C[n \times n]$

|          |     |          |     |          |
|----------|-----|----------|-----|----------|
| $C_{11}$ | ... |          | ... | $C_{1n}$ |
| ...      | ... |          | ... | ...      |
|          |     | $C_{ij}$ |     |          |
| ...      | ... | ...      | ... | ...      |
| $C_{n1}$ | ... | $C_{nj}$ | ... | $C_{nn}$ |

- $C_{ij}$  using  $n$  processors:  $O(\log(n))$
- $C[n \times n]$ :  $(n \times n)C_{ij}$ 
  - $n$  processors:  $O(n^2 * \log(n))$
  - $n^3$  processors:  $O(\log(n))$
  - $n^2$  processors???

=> Difference between **CRCW** & **EREW**???

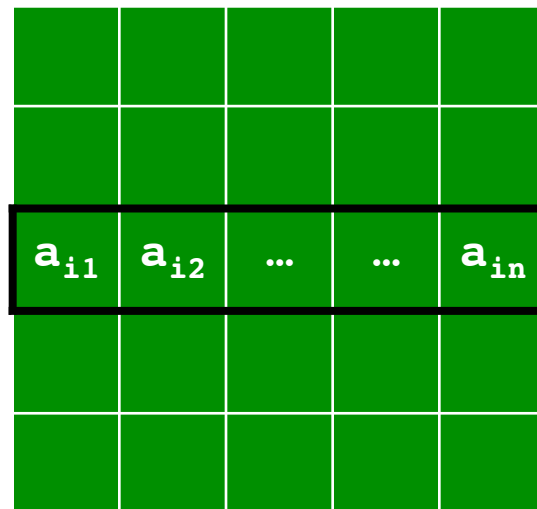
$C[n \times n]$

|          |     |          |     |          |
|----------|-----|----------|-----|----------|
| $C_{11}$ | ... |          | ... | $C_{1n}$ |
| ...      | ... |          | ... | ...      |
|          |     | $C_{ij}$ |     |          |
| ...      | ... | ...      | ... | ...      |
| $C_{n1}$ | ... | $C_{nj}$ | ... | $C_{nn}$ |

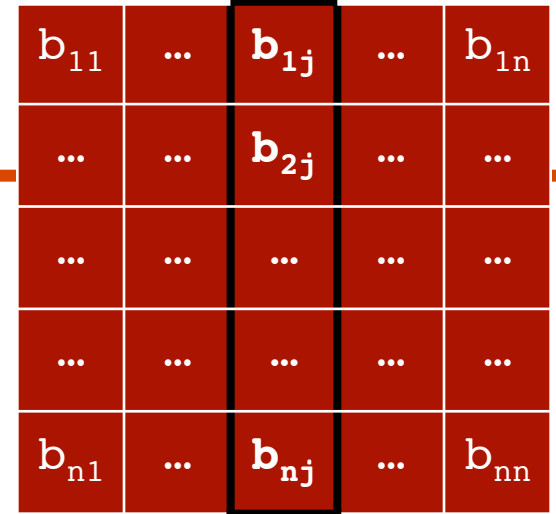


$$\begin{aligned}
 C_{i1} &= a_{i1} * b_{11} + a_{i2} * b_{21} + \dots + a_{in} * b_{n1} \\
 C_{i2} &= a_{i1} * b_{12} + a_{i2} * b_{22} + \dots + a_{in} * b_{n2} \\
 &\dots \\
 C_{in} &= a_{i1} * b_{1n} + a_{i2} * b_{2n} + \dots + a_{in} * b_{nn}
 \end{aligned}$$

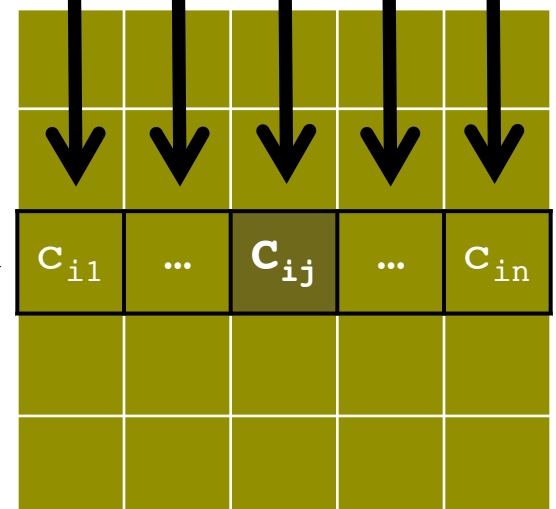
$A[n \times n]$



$B[n \times n]$



$C[n \times n]$



# Concurrent Read: in row

$$\begin{array}{l}
 C_{i1} = a_{i1} * b_{11} + a_{i2} * b_{21} + \dots + a_{in} * b_{n1} \\
 C_{i2} = a_{i1} * b_{12} + a_{i2} * b_{22} + \dots + a_{in} * b_{n2} \\
 \dots \\
 C_{in} = a_{i1} * b_{1n} + a_{i2} * b_{2n} + \dots + a_{in} * b_{nn}
 \end{array}$$

$P_1 \dots P_n$   
 In parallel

CR: working  
 ER: problem

CW: working  
 EW: working



# Column

$$\begin{aligned} C_{1j} &= a_{11} * b_{1j} + a_{12} * b_{2j} + \dots + a_{1n} * b_{nj} \\ C_{2j} &= a_{21} * b_{1j} + a_{22} * b_{2j} + \dots + a_{2n} * b_{nj} \\ &\dots \\ C_{nj} &= a_{n1} * b_{1j} + a_{n2} * b_{2j} + \dots + a_{nn} * b_{nj} \end{aligned}$$

A[nxn]

|          |          |     |     |          |
|----------|----------|-----|-----|----------|
| $a_{11}$ | $a_{12}$ | ... | ... | $a_{1n}$ |
| ...      | ...      | ... | ... | ...      |
| $a_{i1}$ | $a_{i2}$ | ... | ... | $a_{in}$ |
| ...      | ...      | ... | ... | ...      |
| $a_{n1}$ | $a_{n2}$ | ... | ... | $a_{nn}$ |

B[nxn]

|          |     |          |     |          |
|----------|-----|----------|-----|----------|
| $b_{11}$ | ... | $b_{1j}$ | ... | $b_{1n}$ |
| ...      | ... | $b_{2j}$ | ... | ...      |
| ...      | ... | ...      | ... | ...      |
| ...      | ... | ...      | ... | ...      |
| $b_{n1}$ | ... | $b_{nj}$ | ... | $b_{nn}$ |



C[nxn]

|  |  |          |  |  |
|--|--|----------|--|--|
|  |  | $C_{1j}$ |  |  |
|  |  | ...      |  |  |
|  |  | $C_{ij}$ |  |  |
|  |  | ...      |  |  |
|  |  | $C_{nj}$ |  |  |

# Column

$$\begin{array}{l}
 C_{1j} = a_{11} * b_{1j} + a_{12} * b_{2j} + \dots + a_{1n} * b_{nj} \\
 C_{2j} = a_{21} * b_{1j} + a_{22} * b_{2j} + \dots + a_{2n} * b_{nj} \\
 \dots \\
 C_{nj} = a_{n1} * b_{1j} + a_{n2} * b_{2j} + \dots + a_{nn} * b_{nj}
 \end{array}$$

$P_1 \dots P_n$   
In parallel

**CR: working**  
**ER: problem**

**CW: working**  
**EW: working**

# PRAM

## EREW with $n^2$ processors

|          |          |          |          |          |     |          |
|----------|----------|----------|----------|----------|-----|----------|
| $a_{11}$ | $a_{12}$ | ...      | $a_{1n}$ |          |     |          |
| $a_{2n}$ | $a_{21}$ | $a_{22}$ | ...      | $a_{2n}$ |     |          |
| ...      | ...      | ...      | ...      | ...      | ... |          |
| $a_{n2}$ | ...      | $a_{nn}$ | $a_{n1}$ | $a_{n2}$ | ... | $a_{nn}$ |



|          |          |     |          |
|----------|----------|-----|----------|
| $P_{11}$ | $P_{12}$ | ... | $P_{1n}$ |
| $P_{21}$ | $P_{22}$ | ... | $P_{2n}$ |
| ...      | ...      | ... | ...      |
| $P_{n1}$ | $P_{n2}$ | ... | $P_{nn}$ |

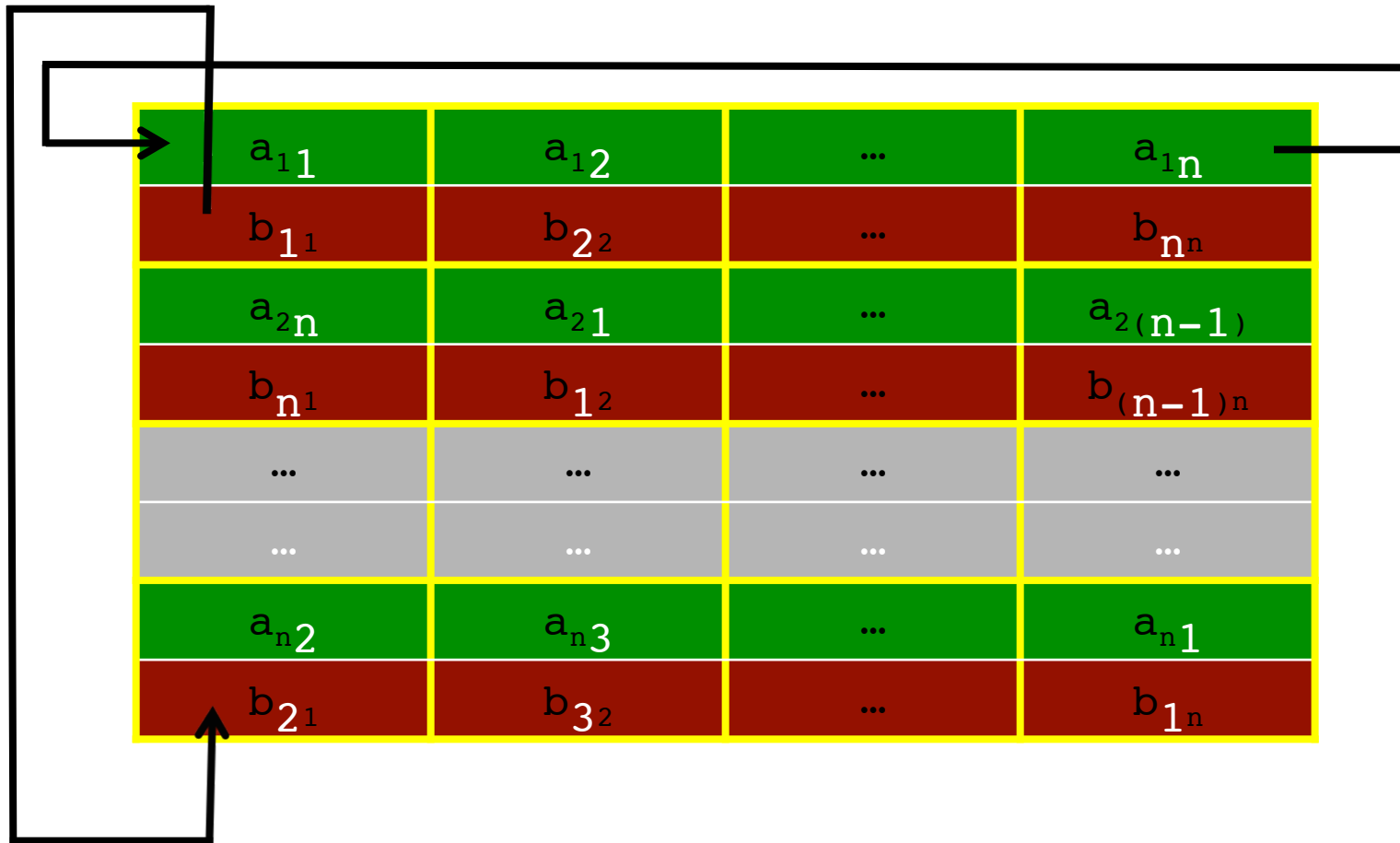
Why not use  $n^3 > n^2$  processors with EREW?

|              |     |              |          |
|--------------|-----|--------------|----------|
| $b_{n1}$     |     |              |          |
| $b_{(n-1)1}$ | ... |              |          |
| ...          | ... | $b_{n(n-1)}$ |          |
| $b_{11}$     | ... | ...          | $b_{nn}$ |
| $b_{n1}$     | ... | $b_{2(n-1)}$ | ...      |
| ...          | ... | $b_{1(n-1)}$ | $b_{2n}$ |
| $b_{21}$     | ... | $b_{n(n-1)}$ | $b_{1n}$ |



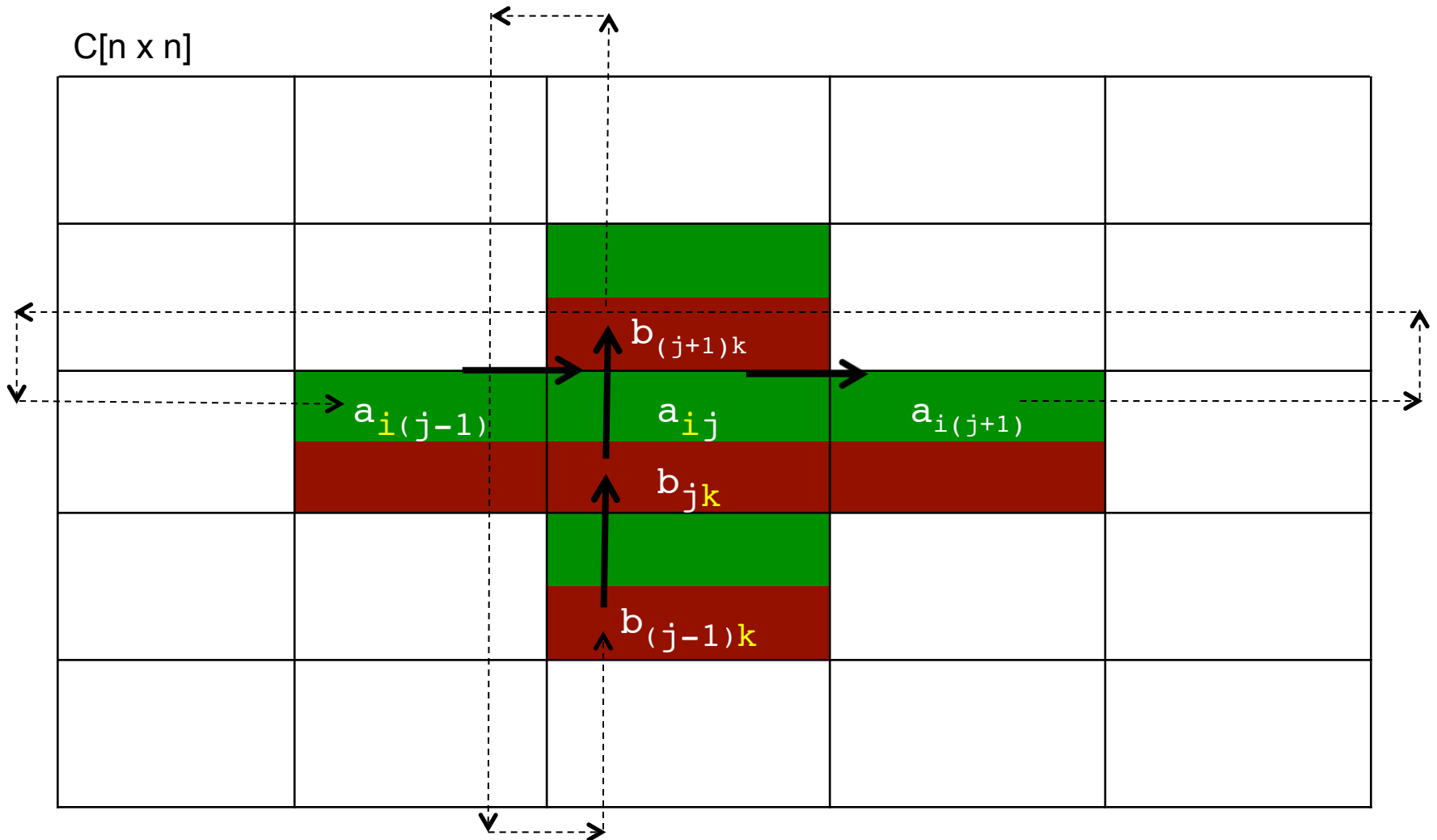


# PRAM EREW with $n^2$ processors





# PRAM EREW with $n^2$ processors



# PRAM

## EREW with $n^2$ processors

|          |          |          |          |          |     |          |
|----------|----------|----------|----------|----------|-----|----------|
| $a_{11}$ | $a_{12}$ | ...      | $a_{1n}$ |          |     |          |
| $a_{2n}$ | $a_{21}$ | $a_{22}$ | ...      | $a_{2n}$ |     |          |
| ...      | ...      | ...      | ...      | ...      | ... |          |
| $a_{n2}$ | ...      | $a_{nn}$ | $a_{n1}$ | $a_{n2}$ | ... | $a_{nn}$ |



|          |          |     |          |
|----------|----------|-----|----------|
| $P_{11}$ | $P_{12}$ | ... | $P_{1n}$ |
| $P_{21}$ | $P_{22}$ | ... | $P_{2n}$ |
| ...      | ...      | ... | ...      |
| $P_{n1}$ | $P_{n2}$ | ... | $P_{nn}$ |

- Each  $P_{ij}$ 
  - Run "\*" in  $n$  steps sequentially
  - Run "+" in  $n$  steps
- $O(n)$



|              |     |              |          |
|--------------|-----|--------------|----------|
| $b_{n1}$     |     |              |          |
| $b_{(n-1)1}$ | ... |              |          |
| ...          | ... | $b_{n(n-1)}$ |          |
| $b_{11}$     | ... | ...          | $b_{nn}$ |
| $b_{n1}$     | ... | $b_{2(n-1)}$ | ...      |
| ...          | ... | $b_{1(n-1)}$ | $b_{2n}$ |
| $b_{21}$     | ... | $b_{n(n-1)}$ | $b_{1n}$ |

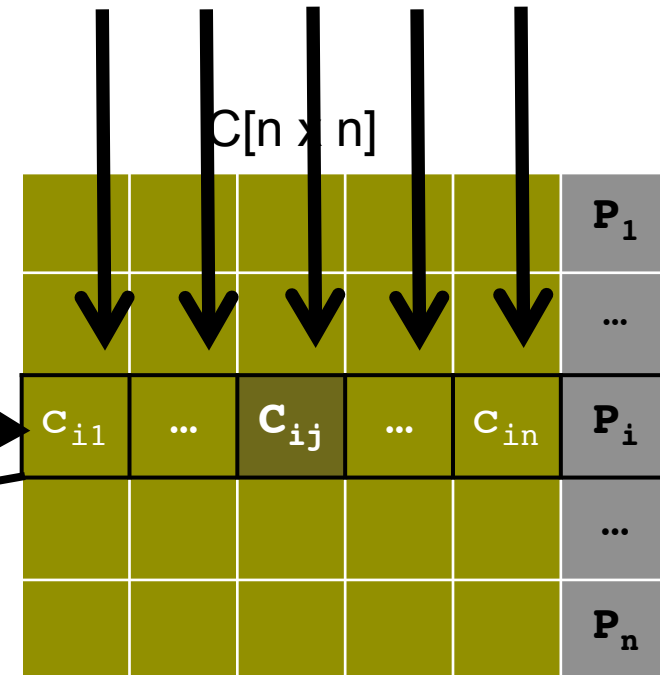
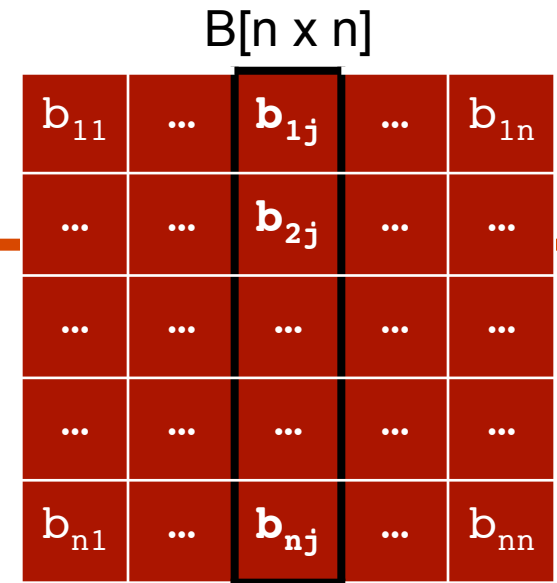
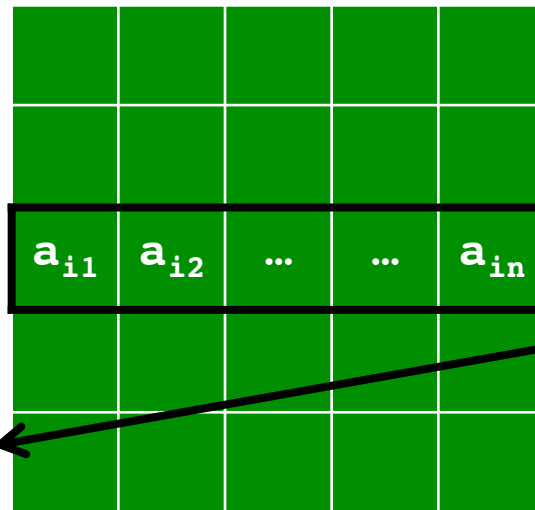
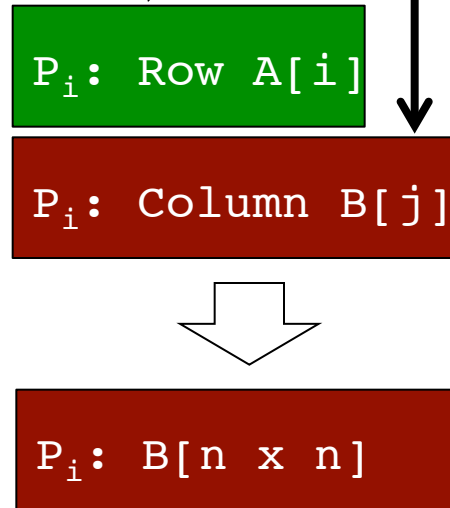




# PRAM with $n$ processors

$$\begin{aligned} C_{i1} &= a_{i1} * b_{11} + a_{i2} * b_{21} + \dots + a_{in} * b_{n1} \\ &\dots \\ C_{ij} &= a_{i1} * b_{1j} + a_{i2} * b_{2j} + \dots + a_{in} * b_{nj} \\ &\dots \\ C_{in} &= a_{i1} * b_{1n} + a_{i2} * b_{2n} + \dots + a_{in} * b_{nn} \end{aligned}$$

$A[n \times n]$





# PRAM

## CRCW with $n$ processors

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Each  $P_i$ :

- $C_{ij} = a_{i1} * b_{1j} + a_{i2} * b_{2j} + \dots + a_{in} * b_{nj} \Rightarrow O(n)$
- Row  $C[i]$ :  $O(n^2)$



# PRAM

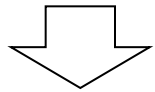
## EREW with $n$ processors

$$\begin{aligned} C_{i1} &= a_{i1} * b_{11} + a_{i2} * b_{21} + \dots + a_{in} * b_{n1} \\ &\vdots \\ C_{ij} &= a_{i1} * b_{1j} + a_{i2} * b_{2j} + \dots + a_{in} * b_{nj} \\ &\vdots \\ C_{in} &= a_{i1} * b_{1n} + a_{i2} * b_{2n} + \dots + a_{in} * b_{nn} \end{aligned}$$

$A[n \times n]$

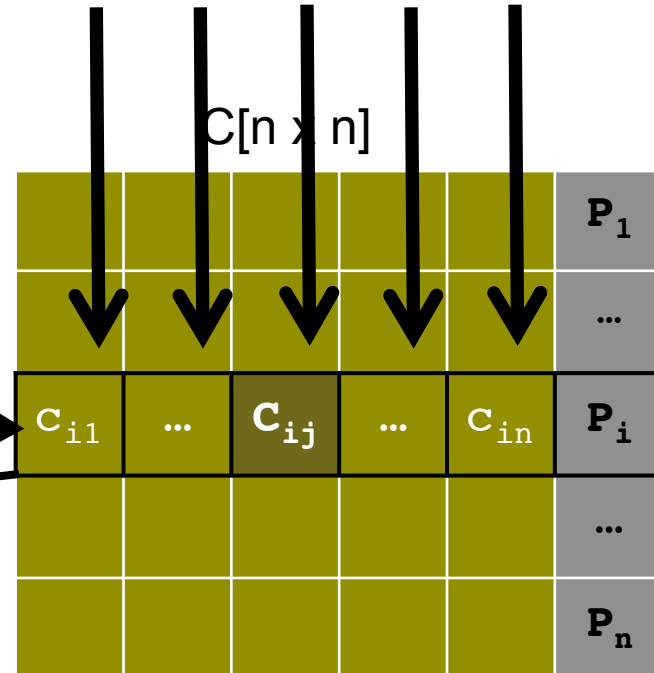
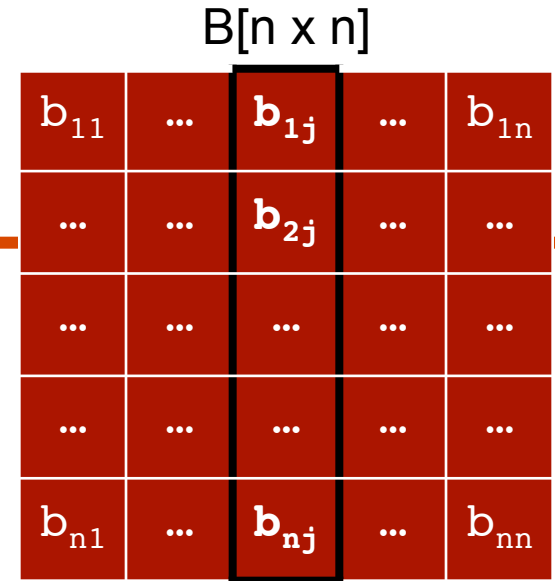
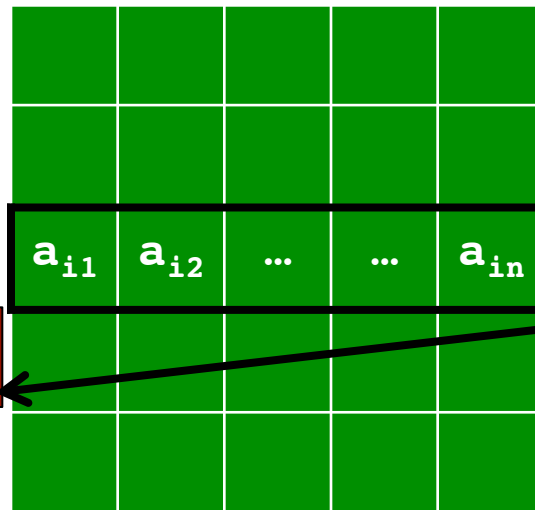
$P_i$ : Row  $A[i]$

$P_i$ : Column  $B[j]$



Each  $P_i$ :  $B[n \times n]$

ER???



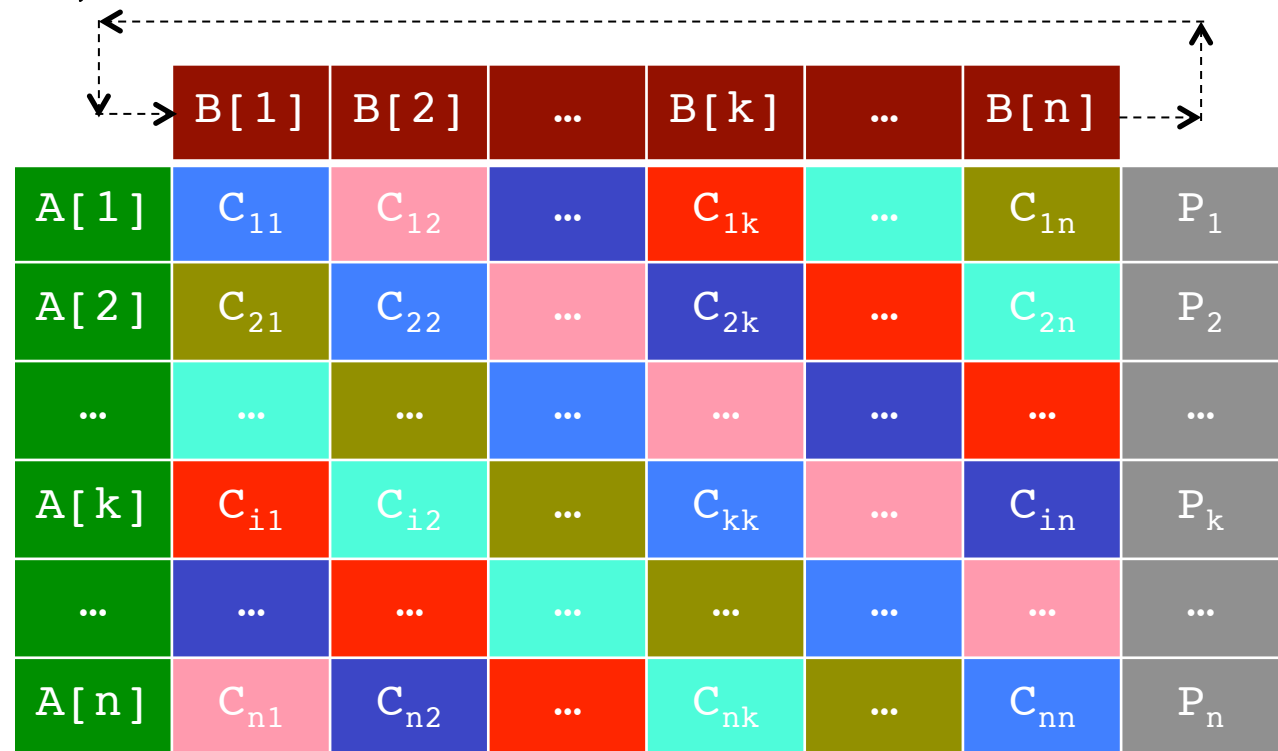


# PRAM

## EREW with $n$ processors

Each  $P_i$ :

- $C_{ij} = a_{i1} * b_{1j} + a_{i2} * b_{2j} + \dots + a_{in} * b_{nj} \Rightarrow O(n)$
- Row  $C[i]$ :  $O(n^2)$





# PRAM

## EREW with $k$ processors ( $k \ll n$ )

Each  $P_i$ :

- 1 step:  $O(k \cdot n) \Rightarrow$
- 1 phase ( $n/k$  steps):  $O(k \cdot n \cdot n/k) \sim O(n^2)$
- $n/k$  phases:  $O(n^2 \cdot n/k) \sim O(n^3/k)$

