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PARALLEL AND DISTRIBUTED ASSOCIATION RULE MINING ALGORITHMS: A RECENT SURVEY

Sudarsan Biswas¹, Neepa Biswas², Kartick Chandra Mondal²

- ¹Department of Information Technology RCC Institute of Information Technology Kolkata, India
- ²Department of Information Technology Jadavpur University Kolkata, India.
- *Corresponding Author Email: <u>biswas.sudarsan@gmail.com, biswas.neepa@gmail.com, kartickjgec@gmail.com</u>

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ARTICLE DETAILS

ABSTRACT

Article History:

Received 27 July 2019 Accepted 29 August 2019 Available online 05 September 2019 Data investigation is an essential key factor now a days due to rapidly growing electronic technology. It generates a large number of transactional data logs from a range of sources devices. Parallel and distributed computing is a useful approach for enhancing the data mining process. The aim of this research is to present a systematic review of parallel association rule mining (PARM) and distributed association rule mining (DARM) approaches. We have observed that the parallelized nature of Apriori, Equivalence class, Hadoop (MapReduce), and Spark proves to be very efficient in PARM and DARM environment. We conclude that this comprehensive review, references cited in this article will convey foremost hypothetical issues and a guideline to the researcher an interesting research direction. The most important hypothetical issue and challenges include the large size of databases, dimensionality of data, indexing schemes of data in the database, data skewness, database location, load balancing strategies, methods of adaptability in incremental databases and orientation of the database.

KEYWORDS

D-sampling, Equivalence Class, MapReduce, Parallel Apriori, Spark.

1. INTRODUCTION

Databases information in this era is intrinsically distributed. This trend of sharing various source of data and generation of immense data volume is certainly showing the way to ridiculous communications costs. Association rule mining (ARM) explores the relationship for finding the meaningful correlation between those large databases. It has been paid more attention to both data mining users and database researchers in the last decade. ARM techniques represent a better choice as they are scalable, highly flexible and can efficiently manage data heterogeneity. It was formally proposed by Agrawal et al. [2] having two counter parts named as PARM and DARM [3]. DARM has become essential for huge and multi scenario database requiring resources, which are heterogeneous and distributed [28]. Recently due to the massive growth of data in organizations, extensive data processing is an essential point of Information Technology. It is not feasible to store and in memory processing of those massive amount of transactional data in organizational databases. ARM is innately disk I/O concentrated task. The I/O costs may be decreased using two approaches by minimizing database scan, or parallelization. This minimization is very essential because the database is generally very large also it is stored in secondary memory. This difficulty is exercised in DARM problem. The aim of this type of DARM problem is solved by parallelizing the disk I/O. In DARM, the database is partitioned between several sites which can perform autonomous parallel computations as well as communication with one site to other sites efficiently [70]. ARM practice may be recognized as distributed and centralized depending on the position of data. For centralized data processing, all data are kept on a single location. In distributed processing, data are placed in several locations. Data are separately accessed from each location [6]. In distributed processing, a collection of sensors, computers, and devices communicate within each other's.

Among the ARM techniques shown in Figure 1, PARM and DARM conveys the most significant role towards rule generation. Prediction and analysis of enormous data volume due to

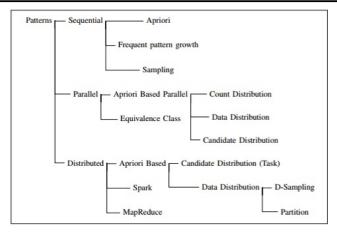


Figure 1: Taxonomy of Pattern Mining Algorithms

rapidly growing of electronic technologies is a challenging task in recent days. Those huge number of transactional data logs generated from several sensors, Internet relay chats, Network logs, Twitter, Facebook, Online banking, ATM transactions etc. in everyday life. Prediction of those data for finding meaningful patterns often became very much helpful in many recent fields included network attack prediction in intrusion detection system [47] [37], locate the causal relationship between drugs and their associated adverse drug reactions [19], detection of suspicious activities in web application [75], advertisement, market basket analysis [2], analysis of business risk, bioinformatics [65], epidemiology, sports, social networking, fluid dynamics, fraud detection [23], [61], crime prevention and prevention [63], telecommunication network, statistical disclosure risk assessment [45], cross marketing, crowd mining and cyber security [63], catalog design, weather prediction [48], recommendation system [21] etc. Also, ARM will examine customers buying behavior that assist the retailers in making advertising strategies, marketing policies, catalog design, store layout orientation, etc.

In a distributed computing environment, there are m sites $\{S_1, S_2, \ldots, S_m\}$ with transactional database TD which is partitioned among m sites into $\{TD_1, TD_2, \ldots, TD_{n_j}\}$ respectively. Assuming the dimension of partitions are TD_i be D_i , where $i = \{1, 2, 3, \ldots, m\}$. In addition, local support count or global support count can be measured $A.sup_i$ of A at site S_i respectively [13]. Let user assigned minimum support constraint is s. An itemset A noted as large if $A.sup \ge s \times D$ globally whereas itemset A is reported as large at site Si locally, if $A.supi \ge s \times D_i$ [28], [69].

1.1 Characteristic of PARM

Mainly architecture of parallel computing can be divided into two branches: distributed memory system (share nothing architectures) and shared memory system [10], [93]. The parallel computing framework includes the following:

- (i) Type of parallelism is required (task or data).
- (ii) Hardware platform of the system.
- (iii) Type of load balancing approaches are used.

1.2 Characteristic of DARM

Discover useful correlation among two nonempty frequent itemsets from distributed large databases.

- (i) A distributed system is loosely coupled that may consist cluster of nodes situated in various places, connected via LAN or internet.
- (ii) Message passing interface design is used for communication between multiple distributed sites that increase the scalability [6].

This overall paper is structured as follows. Section 2 presents the contribution to this paper. Section 3 gives a comprehensive survey of parallel pattern mining algorithms. Section 4 presents a comprehensive survey of distributed pattern mining algorithms. Section 5 presents a review and year wise comparative study among PARM and DARM algorithms. The conclusion of this survey is included in Section 6.

2. RARENESS OF THIS SURVEY

In this paper, we are presenting a comprehensive survey of most exercised PARM and DARM algorithms. This can be considered both as an introduction as well as a guideline to advances opportunities for research in ARM. We have given a systematic and rigorous assessment of significant developments in this field. Those are the following:

- (i) Making the classification of pattern mining algorithms using information collected from the different literature shown in Figure 1.
- (ii) An extensive scrutiny has been performed for both PARM and DARM type of research and development. We have discussed different algorithmic conceptualization, analytical performance characteristics, year wise research impact among the academic and research world and more.
- (iii) In PARM the review of count distribution, data distribution, candidate distribution, equivalence class based parallel algorithms are reported in table 2, 3, 4, 5 respectively.
- (iv) In DARM the review of candidate distribution, data distribution (Sampling, Partition) and Spark & MapReduce based distributed algorithms are shown in table 6, 7, 8, 9 respectively.
- (v) Survey on PARM and DARM mining techniques and presents an algorithm wise comprehensive comparison between those algorithms using different key factor like algorithms names, database orientation, name of parent algorithm, the type of data structure or approaches are used, references, year of publication, citations of those literature is shown in table 10 and table 11 respectively.
- (vi) Presents year wise comparative study between PARM and DARM research approaches and identification research growth form last two decades in this area are shown in Figure 2
- (vii) Presents year wise comparative study between PARM and DARM algorithmic count to recognize the research challenges, important issues in this domain shown in Figure 3.

3. SEQUENTIAL PATTERN MINING ALGORITHMS

3.1 Apriori (Candidate Generation)

Finding rules from largely distributed data sources is a demanding task [80] where ARM is used to generate hidden rules and the relationship between data [2]. The model of association rules is represented as knowledge based on an unsupervised learning method. A formal definition of association rule $(A \Rightarrow B)$ where $A \subseteq T$, $B \subseteq T$ and $(A \cap B) = \varphi$. Rules must convey a certain amount of support and confidence constraint. However, initially, the task of rule mining is to discover the itemsets that are generated frequently from transactional database. Where support $(A \Rightarrow B) \ge \min$ minimum support. Secondly, project all rules from frequent itemsets with confidence $(A \Rightarrow B) \ge \min$ minimum confidence [58]. An item or itemset is a collection of patterns that belongs to $I = \{i_1, i_2, i_3, \ldots, i_n\}$. An

itemsets X with k distinct items is known as {k-itemsets} [2]. Support $(A \Rightarrow B) = s\%$, is the frequency of item's occurrence {AB} from a transactional database. Confidence $(A \Rightarrow B) = c\%$, represent strength of the rules. A single minsup and minconf constraint are reported by the user to identify those rules. Agrawal et al. proposed an Apriori [4] is well exercise traditional sequential technique to select ARM from those transactional databases. Apriori principals:

- (i) Itemset Generation: Any subset of frequent items must be frequent. Also, it will significantly minimize the item set search space. For example, if {X, Y, Z} are frequent then {X, Y} and {Y, Z} also be frequent.
- (ii) Itemset Pruning: Any infrequent patterns superset will not be checked for itemset generation.

3.2 Frequent Pattern Growth

Frequent Pattern Growth (FP growth) approach was suggested by Han et al. [30] that generates a list of frequent patterns except generating any candidate similar to Apriori [4]. The weakness of Apriori has overcome by FP Growth using divide and conquer approach. It takes only two database scans for frequent pattern generation by downward closure property. During the first scan, it counts the number of occurrences for each itemsets. The second scan builds the initial FP tree that keeps the frequency details about the original database [18]. FP growth includes two stages for mining the frequent pattern.

- (i) Initially construct compressed data structure known as FP tree. It is an improved version of bidirectional prefix tree structure that allows bottom up scanning. Every branch of the tree represents a frequent itemsets. Prefixes of the corresponding branches are usually represented by every overlapping itemsets.
- (ii) Next, acquire the frequent itemsets directly from the FP Tree. Create a conditional pattern tree extracted from an initial suffix pattern. After that recursively create its conditional FP tree as per the equal order of magnitude.

Advantages of FP growth:

- Used compact data structure.
- No candidate generation and test process.
- · Building of FP tree without pattern matching.
- Counting frequent pattern.
- Two scans over the database reduces repeated costly database scan rather than Apriori.

 $Disadvantages\ of\ FP\ growth:$

- For large pattern, FP tree construction complexity is very high.
- It does not accommodate into main memory with bigger FP tree.
- Generation of FP tree is expensive.

3.3 Sampling

Sampling is a greedy approach to find out frequent item sets from transactional database suggested by Toivonen et al. [78]. It has taken a random sample from transactions in the database efficiently for selecting the ARM. Where sample size is a prime factor to provide the good approximation of frequent sets. Zaki et al. shown in [91] how random sampling of transactions is useful for finding ARM. However, they have considered the MethodA algorithm for identifying samples of the database. MethodA considerably speeds up the sampling process by well determining the number of records to overlook before chosen the new samples.

Advantage of Sampling:

- Single scan is required to generate rules very efficiently from a large database.
- Reducing disk, I/O cost by significantly decreasing the number of transactions to be measured.
- Sampling can enhance the rules generation process by higher than the order of its previous scale.

Disadvantage of Sampling:

• Second pass is required to find missing frequent sets.

4. PARALLEL PATTERN MINING ALGORITHMS

Two major approaches for exploiting parallelism within ARM algorithms are recognized as task parallelism and data parallelism [6].

- 1) Data parallelism: In this model, the training set is divided among each processor either horizontally or vertically segments. All processors create the same model synchronously where each processor works on a different portion of the database with essential message exchange.
- 2) Task parallelism: In this model, total works are distributed among the processors of a parallel machine with each processor performing a unique part of a learning model prior to synchronizing with the other processors to form a global model.

4.1 Apriori Based Parallel

Various algorithms have been formally developed in the parallel frameworks [1], [3], [13], [87], [90]. Agrawal et al. presents a different version of Apriori like count distribution (CD), data distribution (DD) and candidate distribution (Can.D) of sequential Apriori [4]. Those parallel approaches are used for label wise search for candidate generation that has O(m.c) communication cost for each phase. Where m and c is the

number of data sites and size of candidates itemsets respectively. Also, it will take multiple scans of the database for itemset generation. For example, if the database has N number of transactions then it needs (N+1) scan. table 1 describes the taxonomy of parallel Apriori algorithms. Assessment of CD, DD and Can.D approach are discussed based on their parallelism type, number of database scan, database operational strategy, data structure mentioned as DS and their system architecture.

Table 1: Review of Parallel Apriori Algorithms

| Parallelism | Assessment of Parallel Apriori | | | | | | |
|-------------|--------------------------------|------------------------------------|------------|-----------|------------------|--|--|
| | Algo | Algo Scan Database DS Architecture | | | | | |
| Data | CD | N+1 | Partition | Hash tree | ^a DSN | | |
| Task | DD | N+1 | RRP | Hash tree | a DSN | | |
| Task | Can. D | N+1 | Replicated | Hash tree | a DSN | | |

^aDistributed Shared nothing, RRP Round robin partition, .

4.1.1 Count Distribution (CD)

CD is a sequential Apriori [4] based candidates generation approach by redundant computations in parallel, else processors keep idle to avoid communication cost [3]. It has used DSN architecture for mining the rules from distributed databases. The main objective is to be parallelizing the computation either using a horizontal or vertical partition of the database. In horizontal partition contain the whole transaction in one partition whereas vertical partition contains the same transaction in a different partition. It makes database partitioned into equal sized blocks among all the processors (or different sites) [12]. Each site is identified the biased support of candidates itemsets from local partitions. All processors having its own private memory and disks. All sites are communicated using message passing interface (MPI) having n(n-1) message exchange at each pass. Each processor finds out global itemset frequency by interchanging local itemset frequency from other processors [14]. CD can generate disjoint candidate sets independently. Except for global minimization, each processor exercises traditional Apriori on locally loaded transactions. table 2 describes count distribution based PARM type of algorithms. Here we are considering all parallel algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the parallel algorithms are perceptible in this table.

Table 2: Algorithmic Review of PARM

| Parallel Algorithms | Algorithms for CD | | | |
|---------------------|-------------------|--------------------------------|------|--|
| | Parent Algo | Performance Criteria | Year | |
| DHP [57] | CD | Hash table size | 1995 | |
| PDM [58] | DHP | Count exchange | 1995 | |
| NPA, SPA [74] | CD | Memory size | 1996 | |
| CCPD [89] | CD | Adaptive hash table size | 1996 | |
| HPA, HPA-LED [74] | CD | Hash join | 1996 | |
| DIC [9] | CD | Reordering items | 1997 | |
| APM [12] | DIC | Virtual Partition Pruning | 1998 | |
| FPM [14] | CD | Entropy based partition | 1998 | |
| APM-DIC [12] | APM | Impulsive candidate generation | 1998 | |
| NPGM [73] | NPA | Candidate itemset partition | 1998 | |
| APM-AIC [12] | APM | k clustering | 1998 | |
| DAA [46] | FPM | Principal component analysis | 2001 | |
| DCP [60] | CD | Heuristics pruning | 2001 | |
| DCI [52] | DCP | K way itemset intersection | 2001 | |
| PHP [55] | DHP | Hash table size | 2001 | |
| kDCI [43] | DCP | FSC | 2003 | |
| OPT-DIC [56] | DIC | Scan reduction | 2009 | |

Advantages of count distribution:

- Data parallelism algorithm that has superior performance among DD and Can.D [28].
- For minimizing communication cost, exchange only count do not exchange data tuples between the processors.
- FPM has shown better performance than CD always [14].

Disadvantages of count distribution:

- Total memory of the system does not utilize efficiently.
- Candidates replication are required.
- $\bullet \mbox{Synchronization}$ are essential at each and every pass with high communication load.

4.1.2 Data Distribution (DD)

DD approach are exercised in DSN architecture where each processor has its own private memory and disks [3]. Each processor communicates through message passing interface [29]. Databases are partitioned between all the processors in same sized blocks. After every passes n(n-1) message exchange and synchronization required accordingly [89]. During iteration, each processor scans the complete database to identify global support like local and isolated partitions. It is considered as good utilization of the overall system memory and better load balancing strategies with minimizing idle time for processors [83]. Because the number of processors is large, even disagreement is a major difficulty and processors can be idle at the time of communication. table 3 describes data distribution based PARM type of algorithms. Here we are considering all parallel algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the parallel algorithms are perceptible in this table.

Table 3: Algorithmic Review of PARM

| Algorithms | Algorithms for DD | | | |
|----------------------|-------------------|-------------------------|------|--|
| | Parent Algo | Performance Criteria | Year | |
| PCCD [89] | DD | Hash tree balancing | 1996 | |
| HD, IDD [29] | DD | Candidate hash tree | 2000 | |
| PA [82] | DD | Size of trie tree | 2006 | |
| WDPA [83] | PA | TID counts | 2008 | |
| PARMA-P [94] | FP tree | Hash distribution | 2009 | |
| dRAP independent [8] | Partition | ILP classifier | 2012 | |

PAParallel Apriori.

Advantages of data distribution:

- \bullet Task parallelism algorithm distributes candidate item sets within the processors.
- Competent utilization of system memory.

Disadvantages of data distribution:

- Produced poor performance with respect to CD.
- Huge communication cost because of broadcasting of local database portion to other processors and perform some redundant work.
- Distribution of transactional task for each processor is not possible.

4.1.3 Candidate Distribution (Can.D)

Can.D approach is used in distributed shared nothing architecture [3]. All Processors are communicated using message passing interface, at initial passes n(n-1) no of message exchange required. During the first iteration, it partitions the data among each processor where every processor may produce disjoint candidates that is not dependent on other processors. Every processor gets an equal amount of task based on heuristics approach [74]. Synchronization is not required after each pass. Without exchanging local data, only global values are exchanged. Accordingly, data is acknowledged asynchronously between sites. Processors may not wait for the complete pruning information to arrive from all the processors. table 4 describes candidate distribution based PARM type of algorithms. Here we are considering all parallel algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the parallel algorithms are perceptible in the table.

Advantage of candidate distribution:

· All processors have balanced workload.

Table 4: Algorithmic Review of PARM

| Algorithms | | Algorithms for Can.D | |
|--|-------------|-----------------------------|------|
| | Parent Algo | Performance Criteria | Year |
| HPA [74] | Can.D | Hash join | 1996 |
| HPGM, H-HPGM [73] | НРА | Candidate itemset partition | 1998 |
| ^{нн} -TGD, ^{нн} -PGD, ^{нн} -FGD [73], | НН | Candidate itemset partition | 1998 |

ннН-НРGM.

Disadvantages of candidate distribution:

- · Redistribution of databases includes additional cost.
- Repeated scan of local database partitions.
- · Can.D performs worse than CD.

4.2 Equivalence Class (EC)

Equivalence class transformation (ECLAT) is a parallel ARM algorithm proposed by Zaki et al. [87]. It provides clusters oriented frequent itemsets from vertical transactional databases by equivalence class partitioning. All interconnection between processors accepts a user application that write the memory of remote nodes. That user application makes it convenient to rapid user level messages communication with minimum synchronization costs. The performance of this approach is better than CD algorithm [3]. table 5 describes equivalence class based PARM type of algorithms. Here we are considering all parallel algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the parallel algorithms are perceptible in this table.

Table 5: Algorithmic Review of PARM

| Algorithms for EC | | | | |
|-------------------|---|--|--|--|
| Parent Algo | Performance Criteria | Year | | |
| EC | Tid intersection | 1997 | | |
| Eclat | Bottom up search | 1997 | | |
| Eclat | Hybrid search | 1997 | | |
| Max Eclat | Hybrid search | 1997 | | |
| Itemset lattice | Bottom up search | 1997 | | |
| Clique | Bottom up | 1997 | | |
| Itemset lattice | Top down search | 1997 | | |
| Clique | Hybrid search | 1997 | | |
| Max Clique | Hybrid search | 1997 | | |
| Clique | Hash tree | 1997 | | |
| Snake | Run length encoding | 2000 | | |
| Eclat | Diffset | 2003 | | |
| Eclat | DFS | 2003 | | |
| | EC Eclat Eclat Max Eclat Itemset lattice Clique Itemset lattice Clique Max Clique Clique Snake Eclat | Parent Algo Performance Criteria EC Tid intersection Eclat Bottom up search Eclat Hybrid search Max Eclat Hybrid search Itemset lattice Bottom up search Clique Bottom up Itemset lattice Top down search Clique Hybrid search Max Clique Hybrid search Clique Hybrid search Max Clique Hybrid search Clique Hash tree Snake Run length encoding Eclat Diffset | | |

| hEclat [95] | Eclat | Hash table | 2010 |
|-------------------|-------|----------------------|------|
| Eclat-opt [24] | Eclat | Tow layer hash table | 2013 |
| Bi-Eclat [84] | Eclat | BFS | 2014 |
| P-Eclat [26] | Eclat | Partial BFS | 2016 |
| Eclat-Growth [44] | Eclat | BSRI | 2016 |
| HashEclat [92] | Eclat | Min hash | 2019 |

EcEquivalence class.

These approaches shown in table 5 not only reduce I/O costs by producing only a smaller number of databases scan [90]. It has also minimized computation costs by using competent search strategies.

5. DISTRIBUTED PATTERN MINING ALGORITHMS

Generation of meaningful correlation from distributed large databases. Majority of DARM algorithms are associated with traditional sequential algorithms [4], [30], [68], [78]. Usually, two approaches used to dispense data for parallel processing can be recognized [10], [28].

5.1 Apriori Based DARM Algorithms

Generally, Taxonomy of parallel Apriori based DARM algorithms [16] includes the following:

a) Candidate Distribution: Task distribution. b) Data Distribution: Sampling, Partition.

5.1.1 Candidate Distribution Based DARM

DARM has been coupled with Candidate Distribution (Can.D) and Data Distribution (DD). Whereas Can.D has mainly chosen the condition when data is evenly partitioned into different data sites. Each data site calculates support counts for the corresponding candidate itemsets and accumulates at a central site for selecting the large itemsets for succeeding pass. The most well known algorithms used in DARM are CD [3], FDM [11], ODAM [6], FPM [14]. table 6 describes candidate distribution based DARM type of algorithms. Here we are considering all Apriori based task distributed algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the distributed algorithms are perceptible in this table.

Table 6: Algorithmic Review of DARM

| Algorithms | Apriori Based Task | | | | |
|----------------|--------------------|----------------------------------|------|--|--|
| | Parent Algo | Performance Criteria | Year | | |
| DMA [13] | CD | Local pruning | 1996 | | |
| FDM [11] | CD | Local & Global pruning | 1996 | | |
| FDM-LB [11] | FDM | Local & Lower bound pruning | 1996 | | |
| FDM-LUB [11] | FDM | Local & Upper bound pruning | 1996 | | |
| FDM-LPP [11] | FDM | Polling site pruning | 1996 | | |
| DMCA [80] | FDM | Local pruning & itemset | 2000 | | |
| DDM [69] | FDM | Priority queue | 2001 | | |
| DDDM [69] | DDM | Local count reduction | 2001 | | |
| MDDM [69] | DDM | Priority queue | 2001 | | |
| PDDM [69] | DDM | R-optimal negotiation | 2001 | | |
| ODAM [6] | CD | Candidate itemset reduction | 2004 | | |
| PPDM [34] | FDM | RSA encryption | 2004 | | |
| D-HOTM [40] | CD | RRE | 2005 | | |
| DiHO [39] | CD | Level wise search | 2005 | | |
| DFDM [27] | FDM | Logistic operations | 2006 | | |
| BFDM [27] | DFDM | Logistic operations | 2006 | | |
| DDRM [77] | CD | Lattice based | 2007 | | |
| EDMA [31] | DMA | CMatrix | 2008 | | |
| AprTidRec [79] | CD | TidRec | 2009 | | |
| DDN [22] | CD | Nonderivable & derivable itemset | 2009 | | |
| LMatrix [66] | FDM | Compressed binary matrix | 2010 | | |
| EDFIM [1] | CD | Local & global pruning | 2013 | | |
| PPFDM [64] | FDM | Secure multi party protocol | 2014 | | |
| FDM-KC [76] | FDM | RSA encryption | 2014 | | |
| FDM-UK [76] | FDM | HMAC encryption | 2014 | | |
| PEMA [51] | CD | Size of partition | 2015 | | |

Cheung et al. presented Fast Distributed Mining (FDM) that is CD based candidate's selection that minimize message exchange cost to O(cp.m) where cp=union of locally large itemsets [11]. It has observed that, globally large itemsets is also happened to be locally large at single or multiple sites. FDM doesn't extend well with skewed organization of data. The databases are divided between all processors in equal sized blocks. Generally, ARM algorithms only come across for rules that are globally large. In other words, FDM finds out rules those are locally as well as globally large. FDM has adapted CD to minimize the communication overhead [69]. In FDM the first phase of CD is divided into two rounds of communication. First round finds out locally large in its partition at each site. Second round globally count and sum up of those itemset that is common in all sites. Moreover, interesting associations among locally and globally large itemset are found as smaller set of candidate sets at individual pass. As a result, it cut down the number of messages exchanges. Subsequent to candidates itemset generation global pruning, local pruning is used to prune itemset from each site at each pass. FDM also suggest another optimization count polling.

Advantage of FDM:

- Better Performance than CD.
- Small candidate sets rather than CD [31].
- Three optimizations are used count, local, global polling.
- Less number of candidates are used for counting than CD.

Disadvantage of FDM:

- •CD has shown poor performance in terms of communication for large number of partitions.
- Each iteration polling mechanism need 2 passes of messages.
- This 2 pass scheme for calculating global supports and broadcasting frequent itemsets may reduce performance in parallel environment.

Ashrafi et al. presents ODAM [6], to identify optimized association rules from the distributed database on distributed shared nothing architecture. During first scan counting support of {1-itemsets} from all sites is similar to Apriori [4]. After first scan, it removes all infrequent itemsets and keeps into primary memory. Then broadcasts support count of other sites for finding the frequent {1- itemsets} globally. But it accumulates higher transactions within main memory. This approach has minimized average transactions length and also reduces itemset size apparently. During the new transaction inclusion, it checks the main memory for its presence. If presence increases the transaction counter or include that transaction into memory will increase the counter by one. It uses message exchange optimization and reduction techniques are client-server based. Due to distributed nature, it achieves better performance by reducing the number of message exchanges. Communication and synchronization are required to each and every pass for indirect support or direct counts exchange.

Advantage of ODAM:

- Exchanges fewer messages compared to CD and FDM.
- \bullet Minimize communication cost by 50% to 80% and 20% to 45% respect to CD and FDM respectively
- Instead of sending each support count to individuals' site, FDM sends directly to the polling site.
- Competent method for finding rules from different distributed sites.

Disadvantage of ODAM:

 \bullet Exchanged numerous messages during mining process due to efficient message optimization technique.

Cheung et al. presented FPM [14] to identify association rules on a distributed memory architecture that has taken count distribution approach. Database are partitioned between all the processors in same size blocks. It has included two different types of candidate pruning techniques i.e., global and distributed pruning adopted from FDM [11]. In each iteration, it performs single round of message interchange. It is very competent when data imbalance is large. Global pruning is highly efficient compared to distributed pruning.

Advantage of FPM:

- Better than CD in context to performance.
- Not uses count polling mechanism and subsequentl broadcasts local supports to all processors.
- Distributed pruning is efficient to handling high degree skewness of data.
- Global pruning is effective when mild skewness.
- Selected less number of candidates compared to other.
- High workload balance.

Disadvantage of FPM:

- \bullet Much responsive to workload balance rather than data skewness.
- FPM proves effective by partitioning initially database using balanced k-means clustering.
- It is a variation of CD.

However, Schuster et al. proposed DDM that has reduced communication cost to O(prob.c.n) [69]. Where, prob is the probability of itemset contains support larger then known threshold. Like FDM, candidate's generation in DDM is level wise approach and count from its database locally [11]. After that, the nodes execute a distributed decision protocol for exploring frequent or infrequent. FDM diverge from DDM in such a way that, a locally large itemset is not recognized as globally large itemset till it is demonstrated by messages exchange. This type of behavior has two hypotheses for candidate sets generation. In public hypothesis, the individual node reports global support of the itemset is same as the average local support available for it, or otherwise zero. In private hypothesis, individual node keeps track of its local support and it is available for those nodes that has not preserved their own local support for a candidate. GyHorodi presents a comparative study of DARM algorithms between CD and FDM in [28]. Otey et al. suggested an incremental technique ZIGZAG, a distributed asynchronous approach that provides less communication overhead to mining rules from dynamically distributed datasets [54].

5.1.2 Data Distribution Based DARM

DARM algorithms competence is mostly reliant on DD. Those are concentrated on maximizing parallelism by exchange of the data partitions. DD broadcast candidate itemsets such that individual site calculates a disjoint subset of item sets. Therefore, it can be feasible only for machines with high speed communications between sites. table 7 describes data distribution based DARM type of algorithms. Here we are considering all sampling based data distributed algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the distributed algorithms are perceptible in this table.

Table 7: Algorithmic Review Of DARM

| Algorithms | Algorithm for Sam | Algorithm for Sampling Based DARM | | | | |
|-----------------|------------------------------|-----------------------------------|------|--|--|--|
| | Parent Algo Performance Crit | | | | | |
| Sampling [78] | Negative Border | Sample size | 1996 | | | |
| D-Sampling [70] | Sampling | Size of tries | 2003 | | | |
| SEE [15] | Sampling | Self similarity curve | 2005 | | | |
| Par-Fp [18] | Sampling | Selective sampling | 2005 | | | |

5.1.2.1 D-Sampling

D-Sampling was proposed by Schuster et al. [71], a parallel approach of sampling algorithm [78]. It performs loading a sample into memory that used clusters of shared nothing architecture. All samples are represented by a trie structure known as lexicographic tree and examined entire subroutines repeatedly. Individual node of trie represents more structural information like parents, descendants, etc., that are connected with these nodes. The first level of trie created from samples and intersection is used to create new trie node from the TID list [15]. After creation of candidates set all individual partition is scanned exactly once to count the frequency of individual candidate in parallel. This approach uses M max algorithm to enhance a number of frequent itemsets by a given factor rather than reducing support constraint by a random value. However, creating a candidate's sets takes distributed samples using modified DDM [69].

Advantage of D-Sampling:

- \bullet Single database scan is required to itemset generation.
- Superior speed up and performance over the sampling [78].
- \bullet Combined memory is utilized to linearly enlarge the sample size.

Disadvantage of D-Sampling:

• Increasing communication overhead for large number of patterns.

5.1.2.2 Partition

Partition algorithm was suggested by Savasere et al. [68] that minimizes the database size. It can be used to extract rules in a large database by dividing the database among a set of sites. It implements the task in a distributed way. It has selected all frequent itemsets using two database scans by the level wise approach [8]. It has partitioned the entire database until small enough non overlapping partition that may be handled in main memory. The several processors are working on a different subset of database partitions [50]. The databases are either horizontal or vertical segmented. In the vertical segment, databases are divided based on column number whereas horizontal segment databases are divided based on row number. For {k-itemsets} generation like {1- itemsets}, {2-itemsets}, {3-itemsets} a different partition will be formed and comparing the minimum support

constraint to find out the frequent itemsets from different partitions. This concept is very essential to enhance the execution speed to minimize disk I/O cost [17]. table 8 describes data distribution based DARM type of algorithms. Here we are considering all lattice partition based data distributed algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the distributed algorithms are perceptible in this table.

Table 8: Algorithmic Review Of DARM

| Algorithms | Algorithm for Lattice Partition | | | |
|----------------------------|---------------------------------|-----------------------------|------|--|
| | Parent Algo | Performance Criteria | Year | |
| SPTID SPINC, SPEAR [50] | Partition | Prefix trees | 1995 | |
| PPAR, PEAR [50] | SPEAR | Size of messages | 1995 | |
| AS-CPA [41] | Partition | Prior knowledge of sampling | 1998 | |
| Apriori T [17] | Apriori | No. of message exchange | 2006 | |

5.2 Spark & MapReduce

MapReduce is a programming model for parallel processing of large scale data on Hadoop platforms. It has provided distributed data processing capabilities presented by Google [20]. It consists of two phases, mapper and reducer. First phase does the filtering job by splitting the input based on requisite output keys and values. Second phase takes values from the first phase, perform grouping based on key values and finally aggregate all output keys as well as values [49]. MapReduce can operate on GFS, NDFS or another distributed file system also. Apache Hadoop provides an ecosystem for open source distributed project development suitable for process and storage of large dataset [53]. Where MapReduce supervise the data processing task and HDFS (Hadoop distributed file system) does the data storage job. It can store structured, semi structured and unstructured types of data. Basically, hadoop framework is suitable for handling big data problems. Spark is another Scala based Apache framework targeting real time data analytics [85]. Sing spark in memory computation can be performed for rapid data processing over MapReduce. For large scale data processing, its performance is much faster than hadoop. Spark provides high level library support for Python, R, SQL, Java etc. which support flawless integration within any complex workflow [33]. Moreover, it promotes various service integration like GraphX, SQL, Dataframe, MLlib, Streaming etc.

A research effort has been enhanced of Apriori based and other sequential ARM algorithms by changing them into distributed versions using the Spark and MapReduce. Currently, several DARM algorithms are designed based on Hadoop or Spark framework shown in table 9. Here we are considering all spark and MapReduce based algorithms and their corresponding parent algorithms, performance criteria including their year of publication. Relationship hierarchy of all the distributed algorithms are perceptible in this table.

Table 9: Algorithmic Review of DARM

| Algorithms | Algorithms for Spark & MapReduce | | | | |
|-----------------|----------------------------------|-----------------------------------|------|--|--|
| | Parent Algo | Performance Criteria | | | |
| PFP [38] | MapReduce | Group dependent transactions | 2008 | | |
| PARMA [67] | MapReduce | Sampling&Aggregation | 2012 | | |
| АН [53] | MapReduce | Size of the node clusters | 2013 | | |
| BigFIM [49] | MapReduce | Block partitioning& Prefix tree 2 | | | |
| Dist-Eclat [49] | MapReduce | Round Robin& Prefix tree size | 2013 | | |
| MR-Apriori [42] | MapReduce | Candidate itemset partitions | 2014 | | |
| RuleMR [35] | MapReduce | Entropy minimization | 2014 | | |
| DFIMA [93] | Spark | Boolean vectors size | 2015 | | |
| PaMPa-HD [5] | MapReduce | Enumeration tree partitions | 2015 | | |
| DWS [36] | Spark | LRU partitions | 2015 | | |
| MRH-mine [25] | MapReduce | Queue size | 2015 | | |
| R-Apriori [63] | Spark | Bloom filter | 2015 | | |

| FiDoop [81] | MapReduce | Metric of workload balance | 2016 |
|-------------|-----------|----------------------------|------|
| YAFIM [33] | Spark | Buffer size | 2016 |
| AM [62] | Spark | Bloom filter | 2018 |

DWSDistributed Weka Spark, AHApache Hadoop, AMAdaptive Miner

6. COMPARATIVE ANALYSIS OF PATTERN MINING ALGORITHMS

The main objective of this section is to represent a comparative study between parallel and distributed oriented approaches. An extensive scrutiny has been performed for both type of research and development. This section will discuss different algorithmic conceptualization, analytical performance characteristics, year wise research impact among the academic and research world and more. We are presenting the characteristic of all PARM type approaches discussed before into a summarized format in table 10. Here we are considering some essential parameters like name of the algorithms, database orientation, name of parent algorithm, the type of data structure or approaches are used, references, year of publication, citations of those literature respectively.

The major key issues are associated with the parallel rule mining framework includes:

- (i) I/O minimization, load balancing among processors.
- (ii) Reducing communication, synchronization cost between nodes.
- (iii) Efficient choice database layout, search strategies, correct decomposition of data.
- (iv) Minimizing duplication of task or overlapping.

Most of the literature has suggested I/O costs that can be minimized by different approaches are given below.

- · Minimizing database scan.
- Parallelization the mining process to speed up the rule generation.
- Partitioning the database between individual processor and perform a distributed computing.

We are presenting the characteristic of all DARM type approaches discussed before into a summarized format in table 11. Here we are considering some essential

Table 10: Comparison Of PARM Algorithms

| Algorithms | Important features of PARM Algorithms | | | | |
|----------------|---------------------------------------|-----------|----------------------------|------|--------|
| | DB | Parent | Approach/Features | Year | Impact |
| | | Algo | | | |
| DHP [57] | Н | CD | bit vector | 1995 | 380 |
| PDM [58] | Н | DHP | Hash table | 1995 | 2371 |
| HPA [74] | Н | Can D | Message broadcast | 1996 | 185 |
| CCPD [89] | Н | CD | Hash Tree Balancing | 1996 | 188 |
| NPA [74] | Н | CD | Candidate distribution | 1996 | 185 |
| PCCD [89] | Н | DD | Synchronization | 1996 | 188 |
| DIC [9] | Н | CD | Trie with Counter | 1997 | 2705 |
| Eclat [87] | V | Partition | Equivalence class | 1997 | 82 |
| Par Eclat [90] | V | Eclat | Itemset clustering | 1997 | 344 |
| Max Eclat [88] | V | Eclat | Bottom up search | 1997 | 1513 |
| APM [12] | Н | DIC | without synchronization | 1998 | 53 |
| FPM [14] | Н | CD | Global&Distributed | 1998 | 92 |
| HPGM [73] | Н | НРА | Load skew minimization | 1998 | 90 |
| NPGM [73] | Н | NPA | Minimize load skew | 1998 | 90 |
| HD,IDD [29] | Н | DD | Hash tree | 2000 | 252 |
| VIPER [72] | V | Snake | Bit vectors | 2000 | 357 |
| DAA [46] | Н | FPM | PCA | 2001 | 15 |
| DCP [60] | Н | CD | Candidate itemset sorting | 2001 | 135 |
| DCI [52] | Н | DCP | Itemset intersecting | 2001 | 07 |
| PHP [55] | Н | DHP | Perfect hashing | 2001 | 43 |
| kDCI [43] | Н | DCP | Prefix sharing | 2003 | 34 |
| U-Eclat [86] | V | Eclat | Automatic pruning | 2003 | 698 |

| PA [82] | Н | DD | Trie structure | 2006 | 145 |
|----------------------|---|---------|------------------------------|------|-----|
| WDPA [83] | Н | PA | Tid with metadata | 2008 | 28 |
| OPT-DIC [56] | Н | DIC | Itemsets counting | 2009 | 02 |
| PARMA-P [94] | Н | FP tree | Hash assignment strategy | 2009 | 04 |
| hEclat [95] | V | Eclat | Boolean matrix | 2010 | 11 |
| P-Mine [7] | Н | FP Tree | HY Tree | 2013 | 16 |
| Eclat-opt [24] | V | Eclat | Suffix based | 2013 | 05 |
| Bi-Eclat [84] | V | Eclat | Sorted support | 2014 | 15 |
| P-Eclat [26] | V | Eclat | Partial BFS | 2016 | 01 |
| Eclat-Growth [44] | V | Eclat | Increased search strategy | 2016 | 06 |
| HashEclat [92] | V | Eclat | Min Hash | 2019 | - |

^HHorizontal, ^VVertical, ^RRandom, ^{Hy}Hybrid, ^{ME}MessageExchange, ^{PA}Parallel Apriori

parameters like name of the algorithms, database orientation, name of parent algorithm, the type of data structure or approaches are used, references, year of publication, citations of those literature respectively.

Table 11: Comparison Of Distributed ARM Algorithms

| Algorithms | Important features of DARM | | | | | | | |
|--------------------------------|----------------------------|-------------------------|--|------|--------|--|--|--|
| | DB | Parent Algo | Approach/Features | Year | Impact | | | |
| DMA [13] | Н | Apriori | ME optimization | 1996 | 541 | | | |
| FDM [11] | Н | CD | PVM,ME | 1996 | 670 | | | |
| FDM-LP | Н | FDM | Local pruning | | | | | |
| FDM-LUB FDM-LPP | H H | FDM FDM | Local& Upper bound pruning Local& Polling site | | | | | |
| TD1/ 54 43 | | TDM OD | pruning | 4000 | 0.0 | | | |
| FPM [14] | Н | FDM,CD | Simple messaging schemes | 1998 | 93 | | | |
| DMCA [80] | Н | Apriori | DNF,Guided set, ME | 2000 | 09 | | | |
| DDM [69] | Н | FDM | Message Exchange | 2001 | 155 | | | |
| PDDM | Н | DDM | | | | | | |
| DDDM | Н | DDM | | | | | | |
| MDDM | Н | DDM | | | | | | |
| H-mine [59] | Н | Fp Growth | H Struct | 2001 | 535 | | | |
| ZIGZAG [54] | Н | Backtrack Tree | MFI search | 2003 | 46 | | | |
| Apriori-T [16] | V | Apriori | T tree, ME | 2003 | 42 | | | |
| ODAM [6] | Н | DD | ME optimization | 2004 | 137 | | | |
| PPDM [34] | Н | FDM | Collision probability | 2004 | 1171 | | | |
| D-HOTM [40] | Ну | Apriori | RRE, Hybrid fragmentation | 2005 | 29 | | | |
| DiHO [39] | V | Apriori | Multi relational ARM, Trie | 2005 | 03 | | | |
| D-sampling [71] | R | Sampling | Lexicographic tree | 2005 | 109 | | | |
| BFDM [27] | Н | FDM | Binary code mapping | 2006 | 02 | | | |
| DDRM [77] | Ну | Apriori | WMPI, Lattice Partition | 2007 | 03 | | | |
| EDMA [31] | Н | Apriori | Compressed matrix | 2008 | 13 | | | |
| PFP [38] | Н | FP Growth | MapReduce | 2008 | 453 | | | |
| DDN [22] | Н | Apriori | MPI, DI, NDI | 2009 | 05 | | | |
| AprTidRec [79] | Н | Apriori | TidRec record structure | 2009 | 03 | | | |
| LMatrix [66] | Н | FP tree | FP array technique | 2010 | 10 | | | |
| PARMA [67] | Н | FP Growth | MapReduce | 2012 | 117 | | | |
| EDFIM [1] | Н | Apriori | Node&Global pruning, FPM | 2013 | 01 | | | |
| BigFIM [49] Dist-Eclat [49] | н, V н | Apriori, Eclat Eclat | BFS, MapReduce Diffsets, MapReduce | 2013 | 168 | | | |
| АН [53] | Н | Apriori | MapReduce | 2013 | 20 | | | |

| PPFDM [64] H PPDM | | PPDM | Local Pruning, Broadcasting sup | 2014 | 03 |
|-------------------|----|-----------------|------------------------------------|------|----|
| FDM-UNIFI-KC | Н | FDM | Private binary vectors, | 2014 | 93 |
| MR-Apriori [42] | Н | Apriori | MapReduce | 2014 | 31 |
| RuleMR [35] | Н | ID3, CN2 | MapReduce | 2014 | 05 |
| PEMA [51] | Ну | Apriori | MAARM, ME | 2015 | 06 |
| MRH-mine [25] | Н | H mine | MapReduce | 2015 | 01 |
| DFIMA [93] | Н | Apriori | Matrix pruning, Spark | 2015 | 33 |
| DWS [36] | Н | CloudWatch | Spark | 2015 | 31 |
| R-Apriori [63] | Н | Apriori | Spark | 2015 | 35 |
| PaMPa-HD [5] | Н | Carpenter | MapReduce | 2015 | 05 |
| FiDoop [81] | Н | FIU tree | MapReduce | 2016 | 48 |
| YAFIM [33] | Н | Faster IAPI | Spark | 2016 | 08 |
| AM [62] | Н | Bloom filter | Spark | 2018 | 03 |
| MAD-ARM [32] | Н | IDMA, AeMSAR | Mobile agent | 2018 | 06 |

^HHorizontal, ^VVertical, RRandom, ^{Hy}Hybrid, ^{ME}Message Exchange, ^{DWS}Distributed Weka Spark.

The major key issues are associated with the distributed rule mining framework includes:

- (i) These approaches are essential to decrease communication costs.
- (ii) May produce enormous communication overhead due to huge information forwarding by the data sites. $\,$
- (iii) The majority of distributed algorithms may not enclose a competent optimization technique.
- (iv) It generates low cost global association rules.

The graph in Figure 2 represents a comparative study between distributed and parallel oriented research approach. Research citation is counted for each five year slap for both cases. A year wise decreasing pattern is observed in both research field. After year 2000 research in distributed gains more impact than parallel approach. Year wise algorithmic count in both parallel and distributed field is represented in Figure 3. Here also it is observed that after 2000 distributed field gain more attention compared to parallel approach. Here for both graph representation, all the observations are thoroughly performed on Google Scholar up to year 2019.

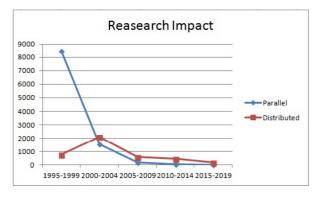


Figure 2: Year wise Research Impact of PARM and DARM

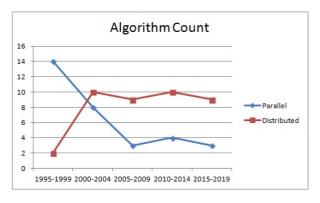


Figure 3: Year wise Algorithms Count of PARM and DARM

Security concerns are associated with privacy, sensitivity, correctness, integrity and false matches. Now a days, secured and privacy related PARM and DARM has become the most essential task practically. Where diverse business organizations attempt to distribute their transaction database with one another. In security domain, usual behavior is mostly frequent but abnormal or suspicious activity is less frequent. For example, consider in a shopping mall event activity analysis, the set of customer's activities is normally represented by frequent patterns, but suspicious activity may be considered as rare patterns. From this review work, various real life application areas are identified and reported. Some important application areas where the results may be useful are distributed network attack prediction, fraud detection, fluid dynamics, risk analysis, bioinformatics, weather analysis and many other fields.

7. CONCLUSION

The significant problems in data mining are finding of association rules from the large databases. This article mainly gives the idea about several algorithms correlated with PARM, DARM and classify them into interrelated techniques.

We have presented a comparative study between distributed and parallel oriented research approach that makes a pathway for future research directions in ARM. Research citation is counted for each five year slap for both cases. A year wise decreasing pattern is observed in both research field. After year 2000 research in DARM gains more impact than PARM approach. Year wise algorithmic count in both PARM and DARM field is represented. Here also it is observed that after 2000 distributed field gain more attention compared to PARM approach. PARM and DARM rule mining techniques required the consistent framework to mining the constructive knowledge from spread database sites. DARM has needed external communications during the entire mining process. The majority of distributed algorithms try to decrease communication cost that is an important issue. Also, we have presented the comparative study between the list of PARM algorithms that are very useful to load balancing strategy, MPI policy among the various sites would be the major demanding area of research in this domain.

The recent trend of research on this filed is paying more attention to advance the algorithm efficiency, high flexibility, scalability, and enhancement of the rule mining process. A research effort has been enhanced of Apriori based and other sequential ARM algorithms by changing them into distributed versions using the Spark and MapReduce. In future, we would like to explore the approaches that partition the set of frequent patterns into several groups and applying a dissimilar interestingness measure on each group. Future target also includes identifying rare pattern approaches to mine infrequent pattern under the parallel and distributed environment.

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