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# Hidden Markov Models

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Slide 1

#### A Markov System

 $\left(\mathbf{s}_{2}\right)$ 

Has N states, called  $s_1$ ,  $s_2$  ...  $s_N$ There are discrete timesteps, t=0, t=1, ...



 $\left(\mathbf{s}_{3}\right)$ 

N = 3

t=0

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## A Markov System

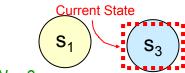
Has N states, called  $s_1$ ,  $s_2$  ..  $s_N$ 



There are discrete timesteps, t=0, t=1, ...

On the t'th timestep the system is in exactly one of the available states. Call it  $q_t$ 

Note:  $q_t \in \{s_1, s_2 ... s_N \}$ 



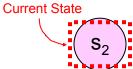
N = 3

t=0

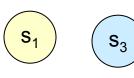
 $q_t = q_0 = s_3$ 

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## A Markov System







N = 3

 $q_t = q_1 = s_2$ 

Has N states, called  $s_1$ ,  $s_2$  ..  $s_N$ 

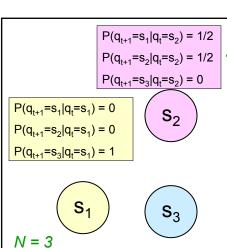
There are discrete timesteps, *t*=0, *t*=1, ...

On the t'th timestep the system is in exactly one of the available states. Call it q,

Note:  $q_t \in \{s_1, s_2 ... s_N\}$ 

Between each timestep, the next state is chosen randomly.

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#### A Markov System

Has N states, called  $s_1$ ,  $s_2$  ..  $s_N$ 

There are discrete timesteps, t=0, t=1, ...

On the t'th timestep the system is in exactly one of the available states. Call it q,

Note: 
$$q_t \in \{s_1, s_2 ... s_N\}$$

Between each timestep, the next  $P(q_{t+1}=s_1|q_t=s_3) = 1/3$  |state is chosen randomly.

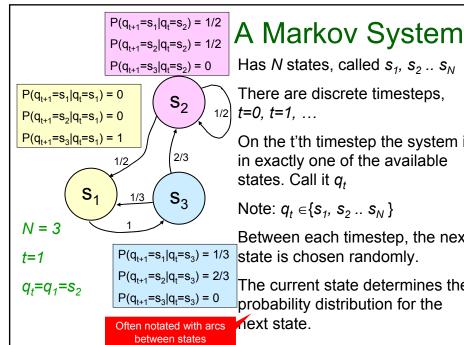
> The current state determines the probability distribution for the next state.

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 $q_t = q_1 = s_2$ 

t=1

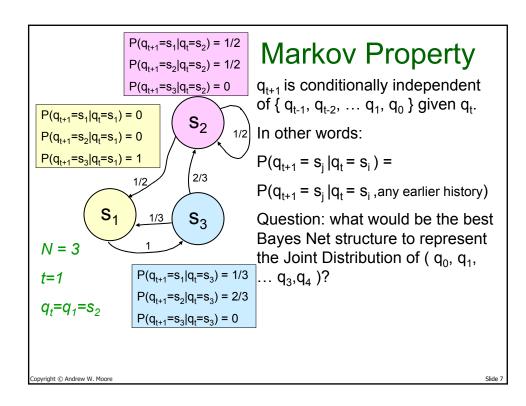


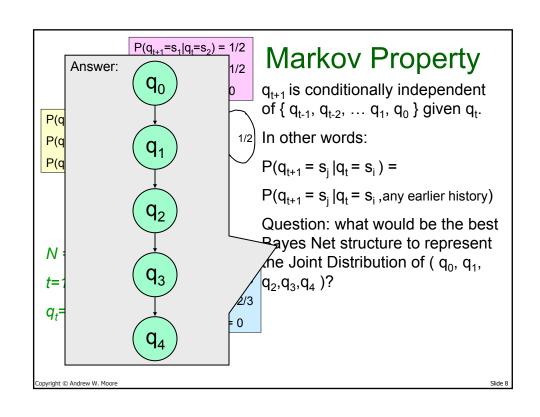
 $P(q_{t+1}=s_2|q_t=s_3) = 2/3$ 

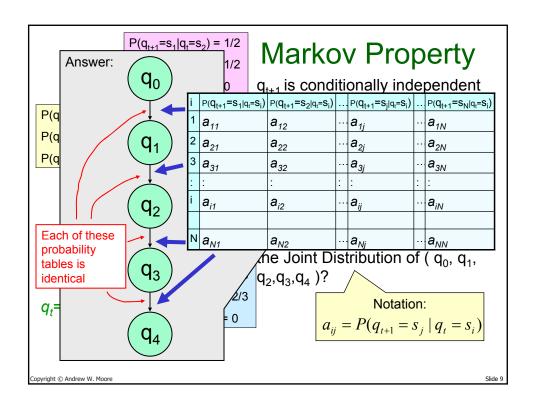
 $P(q_{t+1}=s_3|q_t=s_3)=0$ 

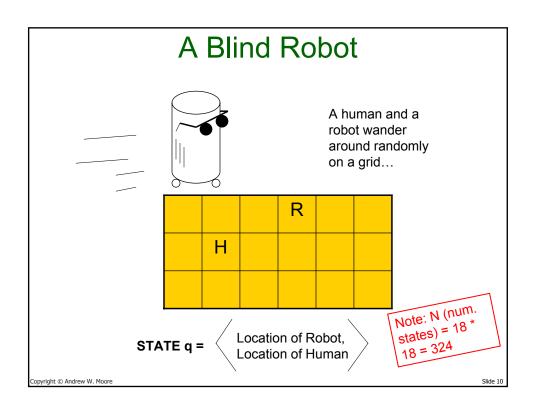
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The current state determines the probability distribution for the

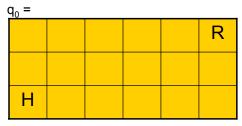








#### Dynamics of System



Each timestep the human moves randomly to an adjacent cell. And Robot also moves randomly to an adjacent cell.

#### **Typical Questions:**

- "What's the expected time until the human is crushed like a bug?"
- "What's the probability that the robot will hit the left wall before it hits the human?"
- "What's the probability Robot crushes human on next time step?"

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Slide 11

#### **Example Question** "It's currently time t, and human remains uncrushed. What's the probability of crushing occurring at time t + 1?" If robot is blind: We'll do this first We can compute this in advance. If robot is omnipotent: Too Easy. We (I.E. If robot knows state at time t), won't do this can compute directly. If robot has some sensors, but Main Body incomplete state information ... of Lecture Hidden Markov Models are applicable! opyright © Andrew W. Moor

#### What is $P(q_t = s)$ ? slow, stupid answer

# Step 1: Work out how to compute P(Q) for any path $Q = q_1 q_2 q_3 ... q_t$

Given we know the start state  $q_1$  (i.e.  $P(q_1)=1$ )

$$\begin{split} P(q_1 \; q_2 \; ... \; q_t) &= P(q_1 \; q_2 \; ... \; q_{t-1}) \; P(q_t | q_1 \; q_2 \; ... \; q_{t-1}) \\ &= P(q_1 \; q_2 \; ... \; q_{t-1}) \; P(q_t | q_{t-1}) \qquad \textit{WHY?} \\ &= P(q_2 | q_1) P(q_3 | q_2) ... P(q_t | q_{t-1}) \end{split}$$

Step 2: Use this knowledge to get  $P(q_t = s)$ 

$$P(q_t = s) = \sum_{Q \in Paths \text{ of length } t \text{ that end in s}} P(Q)$$

Computation is exponential in t

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Slide 13

#### What is $P(q_t = s)$ ? Clever answer

- For each state  $s_i$ , define  $p_t(i)$  = Prob. state is  $s_i$  at time t =  $P(q_t = s_i)$
- · Easy to do inductive definition

$$\forall i \quad p_0(i) =$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

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# What is $P(q_t = s)$ ? Clever answer

For each state s<sub>i</sub>, define
 p<sub>t</sub>(i) = Prob. state is s<sub>i</sub> at time t

$$= P(q_t = s_i)$$

• Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

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Clido 15

#### What is $P(q_t = s)$ ? Clever answer

• For each state  $s_i$ , define

$$p_t(i)$$
 = Prob. state is  $s_i$  at time  $t$   
=  $P(q_t = s_i)$ 

· Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

$$\sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) =$$

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#### What is $P(q_t = s)$ ? Clever answer

- For each state s<sub>i</sub>, define  $p_t(i)$  = Prob. state is  $s_i$  at time t $= P(q_t = s_i)$
- · Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{aligned} \forall j \quad & p_{t+1}(j) = P(q_{t+1} = s_j) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_j \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_j \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_{t+1} = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_t = s_i \mid q_t = s_i) P(q_t = s_i) = \\ & \sum_{i=1}^{N} P(q_t = s_i \mid q_t = s_i) P(q_t = s_$$

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#### What is $P(q_t = s)$ ? Clever answer

- For each state s<sub>i</sub>, define  $p_t(i)$  = Prob. state is  $s_i$  at time t $= P(q_t = s_i)$
- · Easy to do inductive definition

$$\forall i \quad p_0(i) = \begin{cases} 1 & \text{if } s_i \text{ is the start state} \\ 0 & \text{otherwise} \end{cases}$$

$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) = \sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) = \sum_{i=1}^{N} P(q_{t+1} = s_i) = \sum_{i=1}^{N$$

- · Computation is simple.
  - Just fill in this table in this

t	$p_t(1)$	$p_t(2)$		$p_t(N)$
0	0 —	1	1	0
1				_
:	4			
t <sub>final</sub>				<b>→</b>

 $\sum_{i=1}^{N} P(q_{t+1} = s_j \mid q_t = s_i) P(q_t = s_i) = \sum_{i=1}^{N} a_{ij} p_t(i)$ 

#### What is $P(q_t = s)$ ? Clever answer

- For each state  $s_i$ , define  $p_t(i) = \text{Prob. state is } s_i$  at time  $t = P(q_t = s_i)$
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$$\forall j \quad p_{t+1}(j) = P(q_{t+1} = s_j) =$$

$$\sum_{i=1}^{N} P(q_{t+1} = s_j \land q_t = s_i) =$$

 It was meant to warm you up to this trick, called *Dynamic Programming*, because HMMs do many tricks like this.

Cost of computing P<sub>t</sub>(i) for all

states S<sub>i</sub> is now O(t N<sup>2</sup>)

The stupid way was O(Nt)
This was a simple example

 $\sum_{i=1}^{N} P(q_{t+1} = s_j \mid q_t = s_i) P(q_t = s_i) = \sum_{i=1}^{N} a_{ij} p_t(i)$ 

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Slide 19

#### **Hidden State**

"It's currently time t, and human remains uncrushed. What's the probability of crushing occurring at time t + 1?"

If robot is blind:

We'll do this first

If robot is omnipotent:

(I.E. If robot knows state at time t), won't do this can compute directly.

If robot has some sensors, but

incomplete state information ...

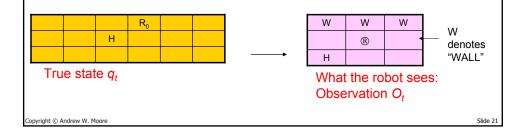
Hidden Markov Models are applicable!

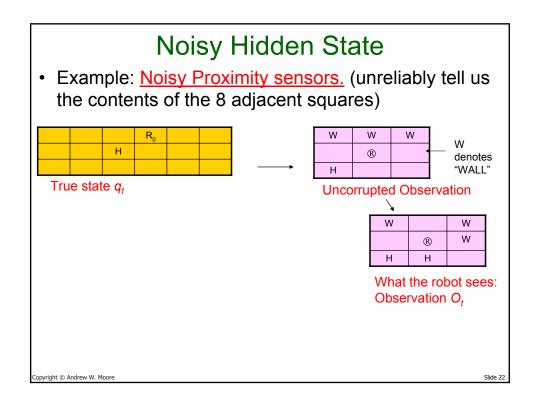
Main Body of Lecture

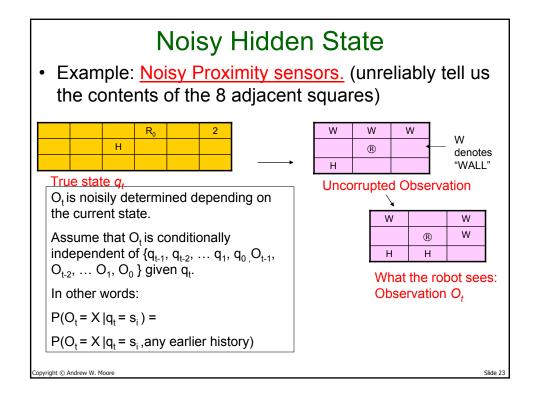
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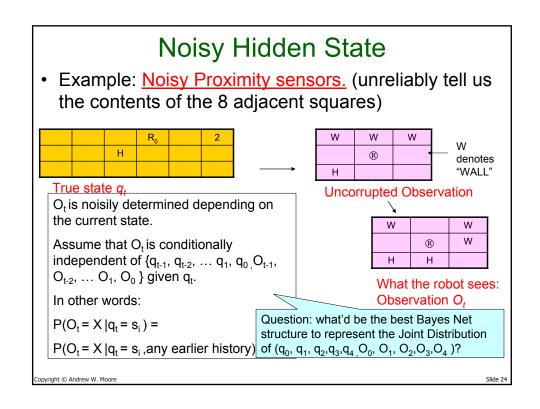
#### **Hidden State**

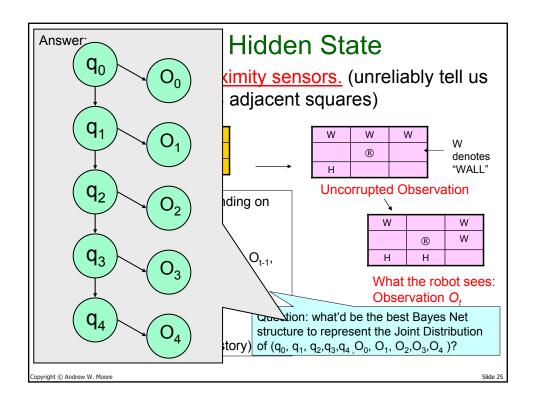
- The previous example tried to estimate  $P(q_t = s_i)$  unconditionally (using no observed evidence).
- Suppose we can observe something that's affected by the true state.
- Example: <u>Proximity sensors.</u> (tell us the contents of the 8 adjacent squares)

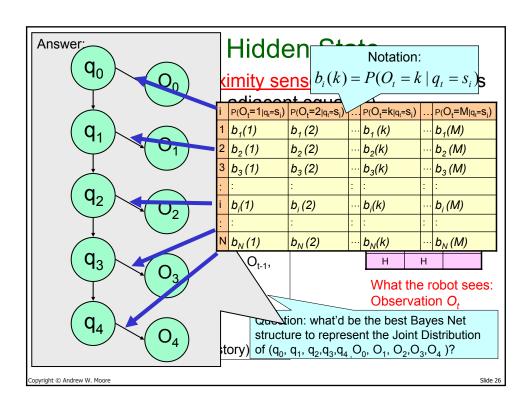












#### Hidden Markov Models

Our robot with noisy sensors is a good example of an HMM

Question 1: State Estimation
 What is P(q<sub>T</sub>=S<sub>i</sub> | O<sub>1</sub>O<sub>2</sub>...O<sub>T</sub>)

It will turn out that a new cute D.P. trick will get this for us.

· Question 2: Most Probable Path

Given  $O_1O_2...O_T$ , what is the most probable path that I took? And what is that probability?

Yet another famous D.P. trick, the VITERBI algorithm, gets this.

Question 3: Learning HMMs:

Given  $O_1O_2...O_T$ , what is the maximum likelihood HMM that could have produced this string of observations?

Very very useful. Uses the E.M. Algorithm

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#### Are H.M.M.s Useful?

#### You bet !!

- Robot planning + sensing when there's uncertainty (e.g. Reid Simmons / Sebastian Thrun / Sven Koenig)
- Speech Recognition/Understanding
   Phones → Words, Signal → phones
- Human Genome Project
   Complicated stuff your lecturer knows nothing about.
- Consumer decision modeling
- Economics & Finance.

Plus at least 5 other things I haven't thought of.

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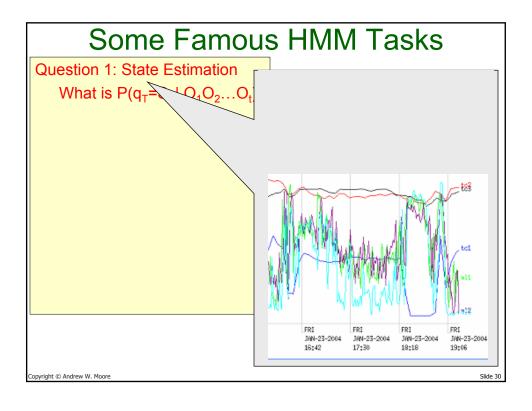
lide 28

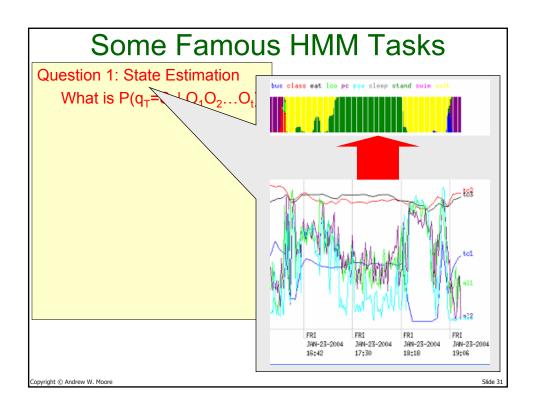
# Some Famous HMM Tasks

Question 1: State Estimation What is  $P(q_T=S_i \mid O_1O_2...O_t)$ 

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lide 29

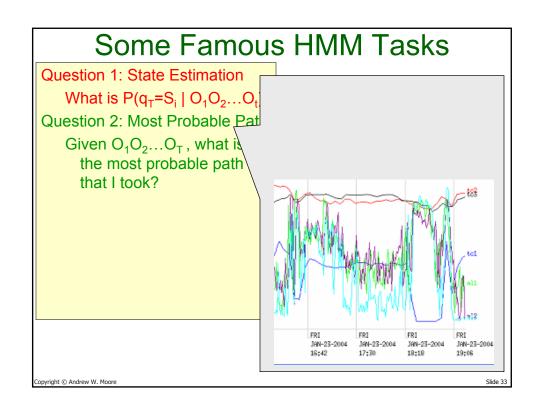


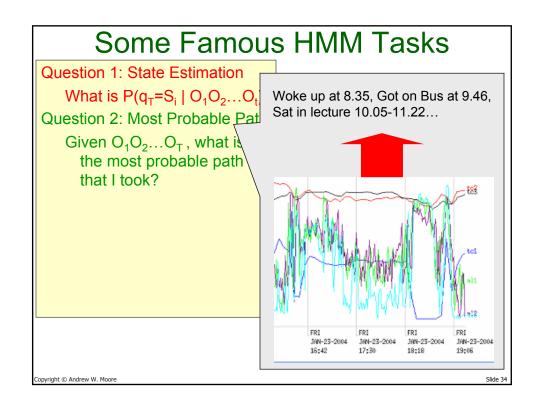


# Question 1: State Estimation What is $P(q_T=S_i \mid O_1O_2...O_t)$ Question 2: Most Probable Path Given $O_1O_2...O_T$ , what is the most probable path that I took?

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Some Famous HMM Tasks





#### Some Famous HMM Tasks

Question 1: State Estimation

What is  $P(q_T=S_i \mid O_1O_2...O_t)$ 

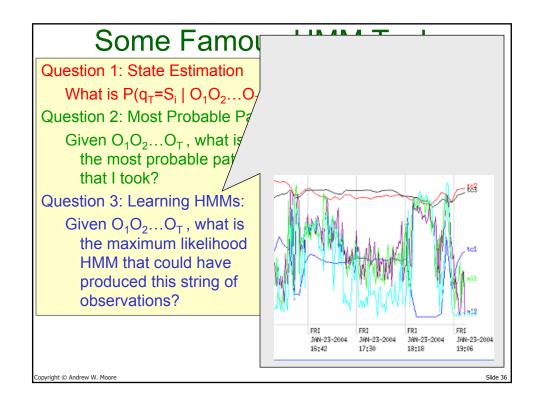
Question 2: Most Probable Path

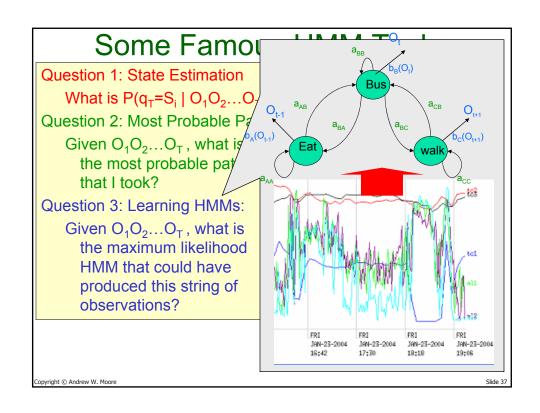
Given O<sub>1</sub>O<sub>2</sub>...O<sub>T</sub>, what is the most probable path that I took?

Question 3: Learning HMMs:

Given O<sub>1</sub>O<sub>2</sub>...O<sub>T</sub>, what is the maximum likelihood HMM that could have produced this string of observations?

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### Basic Operations in HMMs

For an observation sequence  $O = O_1 ... O_T$ , the three basic HMM operations are:

Problem	Algorithm	Complexity
		+
Evaluation:	Forward-Backward	$O(TN^2)$
Calculating $P(q_t=S_i \mid O_1O_2O_t)$		
Inference:	Viterbi Decoding	O(TN <sup>2</sup> )
Computing $Q^* = argmax_Q P(Q O)$		7
Learning:	Baum-Welch (EM)	O(TN <sup>2</sup> )
Computing $\lambda^* = \arg\max_{\lambda} P(O \lambda)$		

T = # timesteps, N = # states

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## **HMM Notation** (from Rabiner's Survey) Hidden Markov Models and Selected Applications in Speech

The states are labeled  $S_1$   $S_2$  ..  $S_N$  Recognition," Proc. of the IEEE, Vol.77, No.2, pp.257-286, 1989.

For a particular trial....

Let T be the number of observations

> is also the number of states passed through

 $O = O_1 O_2 ... O_T$  is the sequence of observations

 $Q = q_1 q_2 ... q_T$  is the notation for a path of states

 $\lambda = \langle N, M, \{\pi_{i,}\}, \{a_{ij}\}, \{b_i(j)\} \rangle$ is the specification of an **HMM** 

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#### **HMM Formal Definition**

An HMM, λ, is a 5-tuple consisting of

- N the number of states
- M the number of possible observations
- $\{\pi_1, \, \pi_2, \, ... \, \pi_N\}$  The starting state probabilities

 $P(q_0 = S_i) = \pi_i$ 

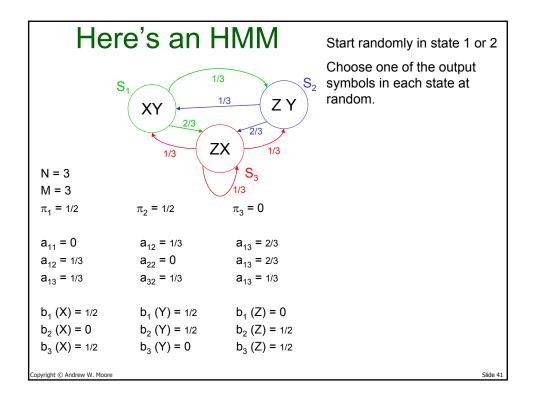
This is new. In our previous example, start state was deterministic

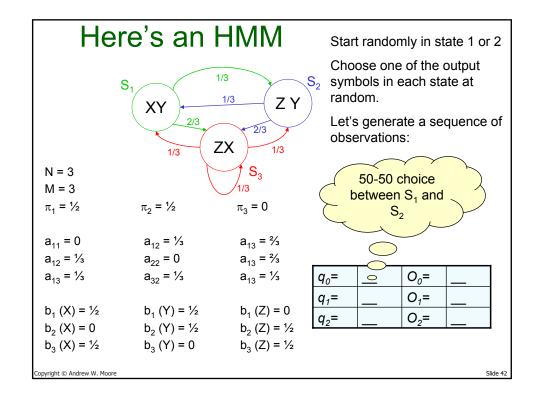
- The state transition probabilities  $P(q_{t+1}=S_i | q_t=S_i)=a_{ii}$

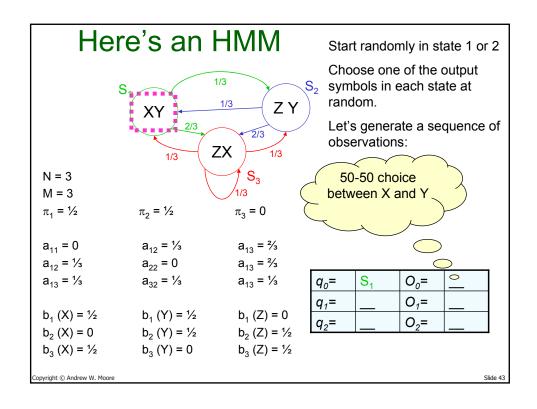
b<sub>1</sub>(2) ...  $b_1(1)$  $b_1(M)$  $b_2(M)$  $b_2(1)$  $b_N(2)$  ...  $b_N(M)$  $b_{N}(1)$ 

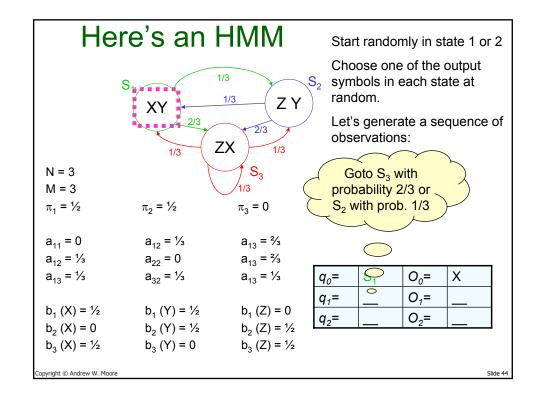
The observation probabilities  $P(O_t=k \mid q_t=S_i)=b_i(k)$ 

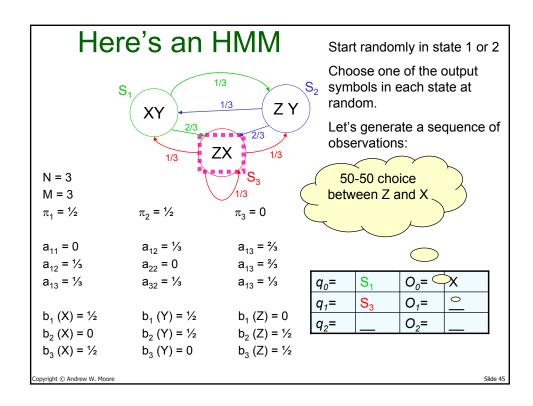
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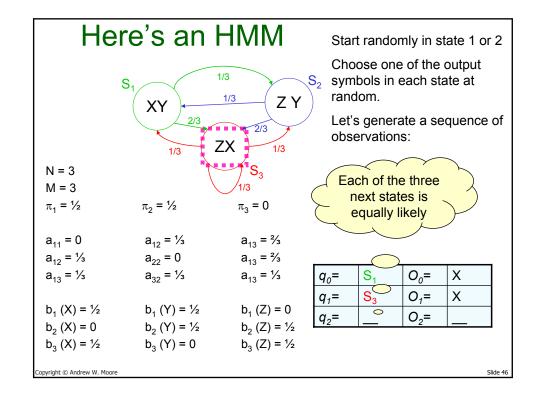


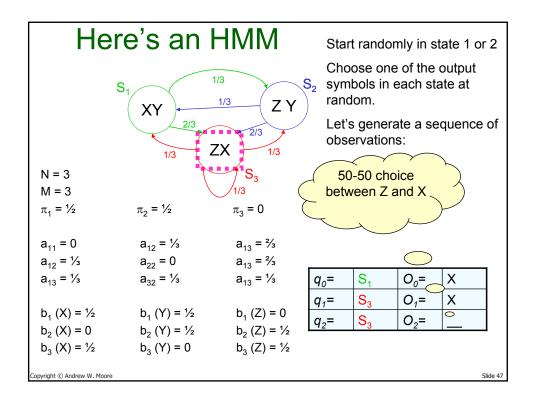


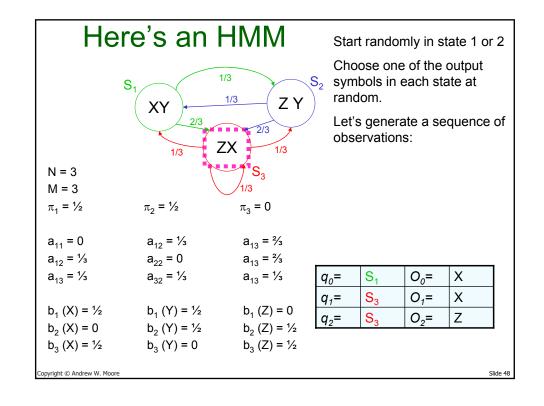




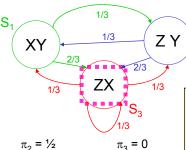












$$\pi_1 = \frac{1}{2}$$
  $\pi_2 = \frac{1}{2}$   $\pi_3 = 0$ 
 $a_{11} = 0$   $a_{12} = \frac{1}{3}$   $a_{13} = \frac{2}{3}$ 
 $a_{12} = \frac{1}{3}$   $a_{22} = 0$   $a_{13} = \frac{2}{3}$ 
 $a_{13} = \frac{1}{3}$   $a_{32} = \frac{1}{3}$   $a_{13} = \frac{1}{3}$ 

Start randomly in state 1 or 2

Choose one of the output symbols in each state at random.

Let's generate a sequence of observations:

This is what the observer has to work with...

	7		
$q_o =$	?	O <sub>0</sub> =	Χ
q <sub>1</sub> =	?	O <sub>1</sub> =	Χ
q <sub>2</sub> =	?	O <sub>2</sub> =	Z

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N = 3M = 3

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#### Prob. of a series of observations

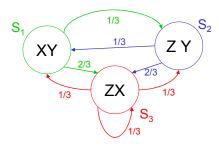
What is 
$$P(\mathbf{O}) = P(O_1 O_2 O_3) = P(O_1 = X ^ O_2 = X ^ O_3 = Z)$$
?

Slow, stupid way:

$$P(\mathbf{O}) = \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \wedge \mathbf{Q})$$
$$= \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \mid \mathbf{Q}) P(\mathbf{Q})$$

How do we compute P(Q) for an arbitrary path Q?

How do we compute P(O|Q) for an arbitrary path Q?



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#### Prob. of a series of observations

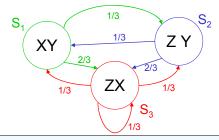
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How do we compute P(Q) for an arbitrary path Q?

How do we compute P(O|Q)for an arbitrary path Q?



 $P(Q) = P(q_1, q_2, q_3)$ 

 $=P(q_1) P(q_2,q_3|q_1)$  (chain rule)

 $=P(q_1) P(q_2|q_1) P(q_3|q_2,q_1)$  (chain)

 $=P(q_1) P(q_2|q_1) P(q_3|q_2) (why?)$ 

Example in the case  $Q = S_1 S_3 S_3$ :

=1/2 \* 2/3 \* 1/3 = 1/9

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#### Prob. of a series of observations

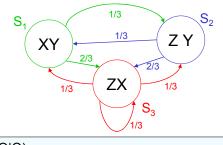
What is 
$$P(\mathbf{O}) = P(O_1 O_2 O_3) = P(O_1 = X ^ O_2 = X ^ O_3 = Z)$$
?

Slow, stupid way:

$$P(\mathbf{O}) = \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \wedge \mathbf{Q})$$
$$= \sum_{\mathbf{Q} \in \text{Paths of length 3}} P(\mathbf{O} \mid \mathbf{Q}) P(\mathbf{Q})$$

How do we compute P(Q) for  $= P(O_1 O_2 O_3 | q_1 q_2 q_3)$ an arbitrary path Q?

How do we compute P(O|Q)for an arbitrary path Q?



P(O|Q)

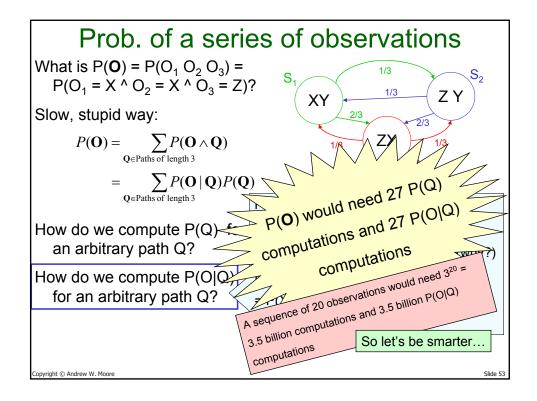
 $= P(O_1 | q_1) P(O_2 | q_2) P(O_3 | q_3) (why?)$ 

Example in the case  $Q = S_1 S_3 S_3$ :

 $= P(X|S_1) P(X|S_3) P(Z|S_3) =$ 

=1/2 \* 1/2 \* 1/2 = 1/8

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# The Prob. of a given series of observations, non-exponential-cost-style

Given observations  $O_1 O_2 \dots O_T$ 

Define

$$\alpha_t(i) = P(O_1 \ O_2 \ \dots \ O_t \ \land q_t = S_i \mid \lambda) \qquad \text{ where } 1 \le t \le T$$

 $\alpha_t(i)$  = Probability that, in a random trial,

- We'd have seen the first t observations
- We'd have ended up in S<sub>i</sub> as the t'th state visited.

In our example, what is  $\alpha_2(3)$ ?

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### $\alpha_t(i)$ : easy to define recursively

 $\alpha_t(i) = P(O_1 \ O_2 \ \dots \ O_T \ \land \ q_t = S_i \mid \lambda) \ \text{($\alpha_t(i)$ can be defined stupidly by considering all paths length "t". How?)}$ 

$$\alpha_{1}(i) = P(O_{1} \wedge q_{1} = S_{i})$$

$$= P(q_{1} = S_{i})P(O_{1}|q_{1} = S_{i})$$

$$= what?$$

$$\alpha_{t+1}(j) = P(O_{1}O_{2}...O_{t}O_{t+1} \wedge q_{t+1} = S_{j})$$

$$=$$

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## $\alpha_t(i)$ : easy to define recursively

 $\alpha_{t}(i) = P(O_{1} \ O_{2} \ \dots \ O_{T} \ \land \ q_{t} = S_{i} \mid \lambda) \ (\alpha_{t}(i) \ \text{can be defined stupidly by considering all paths length "t". How?})$ 

$$\alpha_{1}(i) = P(O_{1} \wedge q_{1} = S_{i})$$

$$= P(q_{1} = S_{i})P(O_{1}|q_{1} = S_{i})$$

$$= what?$$

$$\alpha_{t+1}(j) = P(O_{1}O_{2}...O_{t}O_{t+1} \wedge q_{t+1} = S_{j})$$

$$= \sum_{i=1}^{N} P(O_{1}O_{2}...O_{t} \wedge q_{t} = S_{i} \wedge O_{t+1} \wedge q_{t+1} = S_{j})$$

$$= \sum_{i=1}^{N} P(O_{t+1}, q_{t+1} = S_{j}|O_{1}O_{2}...O_{t} \wedge q_{t} = S_{i})P(O_{1}O_{2}...O_{t} \wedge q_{t} = S_{i})$$

$$= \sum_{i} P(O_{t+1}, q_{t+1} = S_{j}|q_{t} = S_{i})\alpha_{t}(i)$$

$$= \sum_{i} P(q_{t+1} = S_{j}|q_{t} = S_{i})P(O_{t+1}|q_{t+1} = S_{j})\alpha_{t}(i)$$

$$= \sum_{i} a_{ij}b_{j}(O_{t+1})\alpha_{t}(i)$$

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$$\alpha_{t}(i) = P(O_{1}O_{2}..O_{t} \land q_{t} = S_{i}|\lambda)$$

$$\alpha_{1}(i) = b_{i}(O_{1})\pi_{i}$$

$$\alpha_{t+1}(j) = \sum_{i} a_{ij}b_{j}(O_{t+1})\alpha_{t}(i)$$

$$XY$$

$$2/3$$

$$ZX$$

$$1/3$$

$$ZX$$

$$1/3$$

$$S_{3}$$

WE SAW  $O_1 O_2 O_3 = X X Z$ 

$$\alpha_{1}(1) = \frac{1}{4}$$
 $\alpha_{1}(2) = 0$ 
 $\alpha_{1}(3) = 0$ 
 $\alpha_{2}(1) = 0$ 
 $\alpha_{2}(2) = 0$ 
 $\alpha_{2}(3) = \frac{1}{12}$ 
 $\alpha_{3}(1) = 0$ 
 $\alpha_{3}(2) = \frac{1}{72}$ 
 $\alpha_{3}(3) = \frac{1}{72}$ 

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#### **Easy Question**

We can cheaply compute

$$\alpha_t(i)=P(O_1O_2...O_t \land q_t=S_i)$$

(How) can we cheaply compute

$$P(O_1O_2...O_t)$$
 ?

(How) can we cheaply compute

$$P(q_t=S_i|O_1O_2...O_t)$$

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#### **Easy Question**

We can cheaply compute

$$\alpha_t(i)=P(O_1O_2...O_t \land q_t=S_i)$$

(How) can we cheaply compute

$$P(O_1O_2...O_t)$$
 ?  $\sum_{i=1}^{N} \alpha_i(i)$ 

$$P(q_t = S_i | O_1 O_2 ... O_t)$$

(How) can we cheaply compute 
$$P(\mathbf{q_t} = \mathbf{S_i} | \mathbf{O_1} \mathbf{O_2} ... \mathbf{O_t}) \qquad \frac{\alpha_{t}(i)}{\sum\limits_{j=1}^{N} \alpha_{t}(j)}$$

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#### Most probable path given observations

What's most probable path given  $O_1O_2...O_T$ , i.e.

What is 
$$\underset{Q}{\operatorname{argmax}} P(Q|O_1O_2...O_T)$$
?

Slow, stupid answer:

$$\underset{Q}{\operatorname{argmax}} P(Q|O_1O_2...O_T)$$

$$= \underset{Q}{\operatorname{argmax}} \frac{P(O_1 O_2 ... O_T | Q) P(Q)}{P(O_1 O_2 ... O_T)}$$

$$= \underset{Q}{\operatorname{argmax}} P(O_1 O_2 ... O_T | Q) P(Q)$$

#### **Efficient MPP computation**

We're going to compute the following variables:

$$\delta_t(i) = \max_{\substack{q_1 q_2 ... q_{t-1} \\ q_1 q_2 ... q_{t-1}}} P(q_1 \ q_2 \, ... \ q_{t-1} \wedge q_t = S_i \wedge O_1 \, ... \, O_t)$$

= The Probability of the path of Length t-1 with the maximum chance of doing all these things:

...OCCURING

and

...ENDING UP IN STATE S<sub>i</sub>

and

...PRODUCING OUTPUT O<sub>1</sub>...O<sub>t</sub>

DEFINE:  $mpp_t(i) = that path$ 

So:  $\delta_t(i) = \text{Prob}(\text{mpp}_t(i))$ 

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#### The Viterbi Algorithm

$$\delta_t(i) = q_1 q_2 \dots q_{t-1} P(q_1 q_2 \dots q_{t-1} \land q_t = S_i \land O_1 O_2 \dots O_t)$$

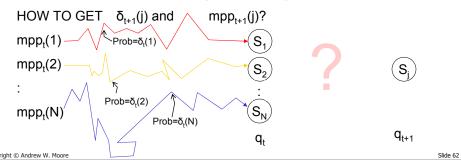
argmax

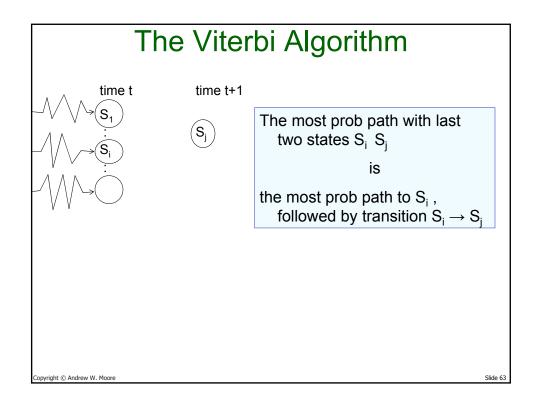
$$mpp_{t}(i) = q_{1}q_{2}...q_{t-1} P(q_{1}q_{2}...q_{t-1} \wedge q_{t} = S_{i} \wedge O_{1}O_{2}..O_{t})$$

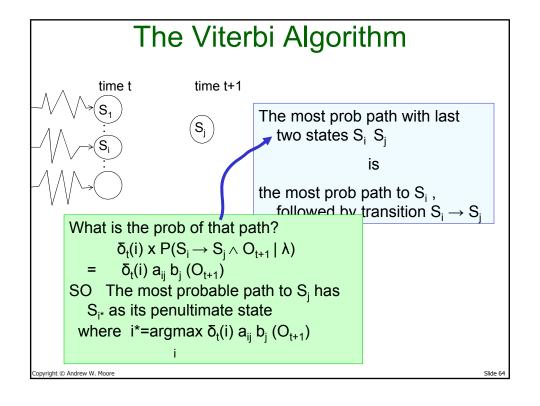
$$\delta_1(i) = \text{ one choice } P(q_1 = S_i \wedge O_1)$$
$$= P(q_1 = S_i)P(O_1|q_1 = S_i)$$

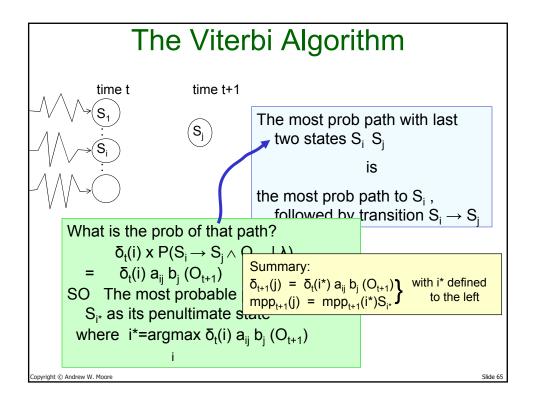
$$=\pi_i b_i(O_1)$$

Now, suppose we have all the  $\delta_t(i)$ 's and  $mpp_t(i)$ 's for all i.









#### What's Viterbi used for?

Classic Example

Speech recognition:

Signal  $\rightarrow$  words

HMM → observable is signal

→ Hidden state is part of word formation

What is the most probable word given this signal?

#### **UTTERLY GROSS SIMPLIFICATION**

In practice: many levels of inference; not one big jump.

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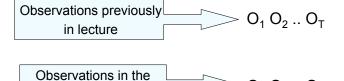
#### HMMs are used and useful

But how do you design an HMM?

Occasionally, (e.g. in our robot example) it is reasonable to deduce the HMM from first principles.

But usually, especially in Speech or Genetics, it is better to infer it from large amounts of data.  $O_1 O_2 ... O_T$  with a big "T".

>>> 0<sub>1</sub> 0<sub>2</sub> .. 0



next bit

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#### Inferring an HMM

Remember, we've been doing things like

$$\mathsf{P}(\mathsf{O}_1\,\mathsf{O}_2 \ldots \mathsf{O}_T \mid \lambda \,)$$

That " $\lambda$ " is the notation for our HMM parameters.

Now We have some observations and we want to estimate  $\lambda$  from them.

AS USUAL: We could use

- (i) MAX LIKELIHOOD  $\lambda = \underset{\lambda}{\operatorname{argmax}} P(O_1 ... O_T | \lambda)$
- (ii) BAYES

Work out P( $\lambda \mid O_1 ... O_T$ )

and then take E[ $\lambda$ ] or max P(  $\lambda$  | O<sub>1</sub> .. O<sub>T</sub> )

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#### Max likelihood HMM estimation

#### Define

$$\begin{split} & \gamma_t(i) = P(q_t = S_i \mid O_1O_2...O_T , \lambda ) \\ & \epsilon_t(i,j) = P(q_t = S_i \land q_{t+1} = S_j \mid O_1O_2...O_T , \lambda ) \end{split}$$

 $y_t(i)$  and  $\varepsilon_t(i,j)$  can be computed efficiently  $\forall i,j,t$ (Details in Rabiner paper)

$$\sum_{t=1}^{T-1} \gamma_t(i) = \sum_{t=1}^{T-1} \gamma_t(i)$$
 Expected number of transitions out of state i during the path

$$\sum_{t=1}^{T-1} \gamma_t(i) = \sum_{t=1}^{T-1} \varepsilon_t(i) = \sum_{t=1}^{T-1} \varepsilon_t(i,j) = \sum_{t=1}^{T-1} \varepsilon_t(i,j)$$

$$\begin{split} & \gamma_t(i) = \mathrm{P}\big(q_t = S_i \big| O_1 O_2 ... O_T, \lambda \big) \\ & \varepsilon_t\big(i,j\big) = \mathrm{P}\big(q_t = S_i \wedge q_{t+1} = S_j \big| O_1 O_2 ... O_T, \lambda \big) \\ & \sum_{t=1}^{T-1} \gamma_t\big(i\big) = \mathrm{expected\ number\ of\ transitions\ out\ of\ state\ i\ during\ path} \\ & \sum_{t=1}^{T-1} \varepsilon_t\big(i,j\big) = \mathrm{expected\ number\ of\ transitions\ out\ of\ i\ and\ into\ j\ during\ path} \end{split}$$

#### **HMM** estimation

Notice 
$$\frac{\sum_{t=1}^{T-1} \varepsilon_t(i, j)}{\sum_{t=1}^{T-1} \gamma_t(i)} = \frac{\left(\begin{array}{c} \text{expected frequency} \\ i \to j \end{array}\right)}{\left(\begin{array}{c} \text{expected frequency} \\ \text{i} \end{array}\right)}$$

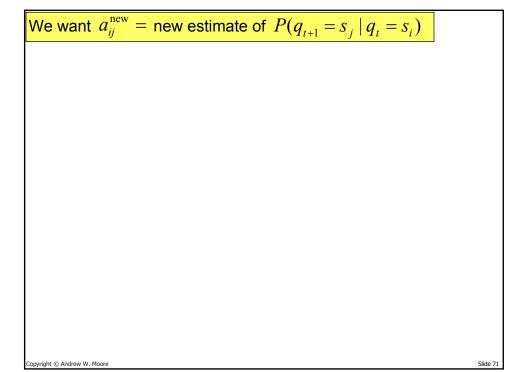
= Estimate of Prob(Next state  $S_i$ |This state  $S_i$ )

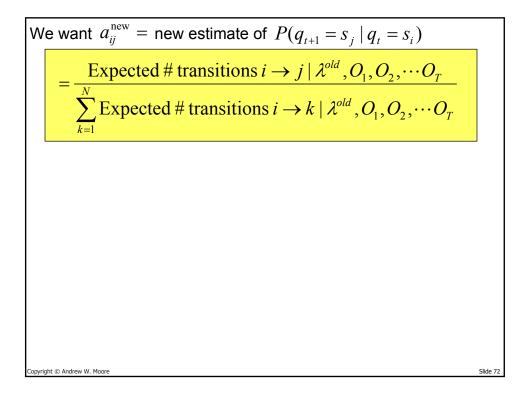
We can re - estimate

$$\mathbf{a}_{ij} \leftarrow \frac{\sum \varepsilon_t(i,j)}{\sum \gamma_t(i)}$$

We can also re - estimate

$$b_j(O_k) \leftarrow \cdots$$
 (See Rabiner)





We want 
$$a_{ij}^{\text{new}} = \text{new estimate of } P(q_{t+1} = s_j \mid q_t = s_i)$$

$$= \frac{\text{Expected \# transitions } i \to j \mid \lambda^{old}, O_1, O_2, \cdots O_T}{\sum_{k=1}^{N} \text{Expected \# transitions } i \to k \mid \lambda^{old}, O_1, O_2, \cdots O_T}$$

$$= \frac{\sum_{t=1}^{T} P(q_{t+1} = s_j, q_t = s_i \mid \lambda^{\text{old}}, O_1, O_2, \dots O_T)}{\sum_{k=1}^{N} \sum_{t=1}^{T} P(q_{t+1} = s_k, q_t = s_i \mid \lambda^{\text{old}}, O_1, O_2, \dots O_T)}$$

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We want 
$$a_{ij}^{\text{new}} = \text{new estimate of } P(q_{t+1} = s_j \mid q_t = s_i)$$

$$= \frac{\text{Expected \# transitions } i \to j \mid \lambda^{old}, O_1, O_2, \cdots O_T}{\sum_{k=1}^{N} \text{Expected \# transitions } i \to k \mid \lambda^{old}, O_1, O_2, \cdots O_T}$$

$$= \frac{\sum_{t=1}^{T} P(q_{t+1} = s_j, q_t = s_i \mid \lambda^{\text{old}}, O_1, O_2, \cdots O_T)}{\sum_{k=1}^{N} \sum_{t=1}^{T} P(q_{t+1} = s_k, q_t = s_i \mid \lambda^{\text{old}}, O_1, O_2, \cdots O_T)}$$

$$= \frac{S_{ij}}{\sum_{k=1}^{N} S_{ik}} \text{ where } S_{ij} = \sum_{t=1}^{T} P(q_{t+1} = s_j, q_t = s_i, O_1, \dots O_T \mid \lambda^{\text{old}})$$

$$= \text{What?}$$

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We want 
$$a_{ij}^{\text{new}} = \text{new estimate of } P(q_{t+1} = s_j \mid q_t = s_i)$$

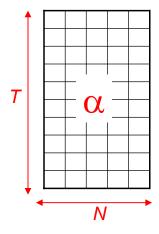
$$= \frac{\text{Expected \# transitions } i \to j \mid \lambda^{old}, O_1, O_2, \cdots O_T}{\sum_{k=1}^{N} \text{Expected \# transitions } i \to k \mid \lambda^{old}, O_1, O_2, \cdots O_T}$$

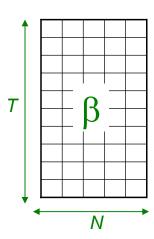
$$= \frac{\sum_{k=1}^{T} P(q_{t+1} = s_j, q_t = s_i \mid \lambda^{\text{old}}, O_1, O_2, \cdots O_T)}{\sum_{k=1}^{N} \sum_{t=1}^{T} P(q_{t+1} = s_k, q_t = s_i \mid \lambda^{\text{old}}, O_1, O_2, \cdots O_T)}$$

$$=\frac{S_{ij}}{\sum\limits_{k=1}^{N}S_{ik}} \text{ where } S_{ij} = \sum\limits_{t=1}^{T}P(q_{t+1}=s_j,q_t=s_i,O_1,\cdots O_T\mid \lambda^{\text{old}})$$
 
$$= a_{ij}\sum\limits_{t=1}^{T}\alpha_t(i)\beta_{t+1}(j)b_j(O_{t+1})$$

We want 
$$a_{ij}^{\text{new}} = S_{ij} \bigg/ \sum_{k=1}^N S_{ik}$$
 where  $S_{ij} = a_{ij} \sum_{t=1}^T \alpha_t(i) \beta_{t+1}(j) b_j(O_{t+1})$ 

We want 
$$a_{ij}^{\text{new}} = S_{ij} \bigg/ \sum_{k=1}^N S_{ik}$$
 where  $S_{ij} = a_{ij} \sum_{t=1}^T \alpha_t(i) \beta_{t+1}(j) b_j(O_{t+1})$ 





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#### **EM for HMMs**

If we knew λ we could estimate EXPECTATIONS of quantities such as

Expected number of times in state i

Expected number of transitions  $i \rightarrow j$ 

If we knew the quantities such as

Expected number of times in state i

Expected number of transitions  $i \rightarrow j$ 

We could compute the MAX LIKELIHOOD estimate of

$$\lambda = \langle \{a_{ij}\}, \{b_i(j)\}, \pi_i \rangle$$

Roll on the EM Algorithm...

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#### EM 4 HMMs

- Get your observations O<sub>1</sub> ...O<sub>T</sub>
- 2. Guess your first  $\lambda$  estimate  $\lambda(0)$ , k=0
- 3. k = k+1
- 4. Given  $O_1 ... O_T$ ,  $\lambda(k)$  compute  $\gamma_t(i) , \ \epsilon_t(i,j) \quad \forall \ 1 \le t \le T, \quad \forall \ 1 \le i \le N, \quad \forall \ 1 \le j \le N$
- 5. Compute expected freq. of state i, and expected freq. i→j
- 6. Compute new estimates of  $a_{ij}$ ,  $b_j(k)$ ,  $\pi_i$  accordingly. Call them  $\lambda(k+1)$
- 7. Goto 3, unless converged.
- Also known (for the HMM case) as the BAUM-WELCH algorithm.

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#### **Bad News**

There are lots of local minima

#### **Good News**

 The local minima are usually adequate models of the data.

#### **Notice**

- EM does not estimate the number of states. That must be given.
- Often, HMMs are forced to have some links with zero probability. This is done by setting a<sub>ij</sub>=0 in initial estimate λ(0)
- Easy extension of everything seen today: HMMs with real valued outputs

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Dad Name

Trade-off between too few states (inadequately modeling the structure in the data) and too many (fitting the noise).

There are lots of Thus #states is a regularization parameter.

Blah blah blah... bias variance tradeoff...blah blah...cross-validation...blah blah....AIC, BIC....blah blah (same ol' same ol')

The local minim data.

**votice** 

- EM does not estimate the number of states. That must be given.
- Often, HMMs are forced to have some links with zero probability. This is done by setting a<sub>ij</sub>=0 in initial estimate λ(0)
- Easy extension of everything seen today: HMMs with real valued outputs

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DON'T PANIC: starts on p. 257.

#### What You Should Know

- What is an HMM?
- Computing (and defining)  $\alpha_t(i)$
- The Viterbi algorithm

Outline of the EM algorithm

 To be very happy with the kind of maths and analysis needed for HMMs

Fairly thorough reading of Rabiner\* up to page 266\*
 [Up to but not including "IV. Types of HMMs"].

\*L. R. Rabiner, "A Tutorial on Hidden Markov Models and Selected Applications in Speech Recognition," Proc. of the IEEE, Vol.77, No.2, pp.257--286, 1989.

http://ieeexplore.ieee.org/iel5/5/698/00018626.pdf?arnumber=18626

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