

Question 1

a.) $5n^3 + 2n^2 + 3n = O(n^3) \rightarrow 5n^3 + 2n^2 + 3n \leq c \cdot n^3$

$\hookrightarrow 5n^3 + 2n^2 + 3n \leq 5n^3 + 2n^3 + 3n^3 = 10n^3$

$\hookrightarrow 5n^3 + 2n^2 + 3n \leq 10n^3$ for $n \geq 1 \rightarrow c = 10, n_0 = 1$

b.) $\sqrt{7n^2 + 2n - 8} = O(n) \rightarrow c_2 \cdot n \leq \sqrt{7n^2 + 2n - 8} \leq c_1 \cdot n$

$\hookrightarrow 0 \rightarrow \sqrt{7n^2 + 2n - 8} \leq \sqrt{7n^2 + 2n} \leq \sqrt{7n^2 + 2n^2} = \sqrt{9n^2} = 3n$ (Upper)

$\hookrightarrow n \rightarrow \sqrt{7n^2 + 2n - 8} \leq \sqrt{7n^2 + 2n^2 - 8n^2} = \sqrt{n^2} = n$ (Lower)

$\hookrightarrow n \leq \sqrt{7n^2 + 2n - 8} \leq 3n$ for $n \geq 1 \rightarrow c_2 = 1, c_1 = 3, n_0 = 1$

c.) $d(n) = O(f(n)), e(n) = O(g(n)) \rightarrow d(n) \cdot e(n) = O(f(n) \cdot g(n))$

$\hookrightarrow d(n) \leq c_1 \cdot f(n) \quad \hookrightarrow e(n) \leq c_2 \cdot g(n)$

$\hookrightarrow d(n) \cdot e(n) \leq c_1 \cdot f(n) \cdot c_2 \cdot g(n)$

$d(n) \cdot e(n) = O(f(n) \cdot g(n)) \rightarrow$ for $n \geq n_0$ and $c = c_1 \cdot c_2$

Question 2

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1 def example1(lst):
2     n = len(lst) ] → O(1)
3     total = 0 ] → O(1)
4     for j in range(n):
5         for k in range(1+j): ] → O(j)
6             total += lst[k] ] → O(1)
7     return total ] → O(1)

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$$\begin{aligned}
 T(n) &= 2 + (1 + 2 + 3 + 4 + \dots + n) + 1 \\
 &= 3 + \left(\frac{n(n+1)}{2} \right) \\
 &= \frac{n^2}{2} + \frac{n}{2} + 3 \\
 &= \cancel{\frac{1}{2}n^2} + \cancel{\frac{1}{2}n} + \cancel{3} \\
 &= n^2 \\
 &= O(n^2)
 \end{aligned}$$

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1 def example2(lst):
2     n = len(lst) ] → O(1)
3     prefix = 0 ] → O(1)
4     total = 0 ] → O(1)
5     for j in range(n):
6         prefix += lst[j] ] → O(1)
7         total += prefix ] → O(1)
8     return total ] → O(1)

```

$$\begin{aligned}
 T(n) &= 3 + \underbrace{2 + 2 + 2 + 2 + \dots + 2}_{n \text{ times}} + 1 \\
 &= \cancel{\frac{1}{2}n} + \cancel{1} \\
 &= n \\
 &= O(n)
 \end{aligned}$$

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1 def example3(n):
2     i = 1 ] → O(1)
3     sum = 0 ] → O(1)
4     while (i <= n):
5         i *= 2 ] → O(1)
6         sum += i ] → O(1)
7     return sum ] → O(1)

```

$2^0, 2^1, 2^2, 2^3, 2^4, \dots, 2^{k-1}$ where $k = \log_2(n)$

$$\begin{aligned}
 \hookrightarrow T(n) &= 2 + (1 + 2 + 3 + \dots + \log(n)) + 1 \\
 &= 3 + \left(\frac{\log(n) \cdot (\log(n) + 1)}{2} \right) \\
 &= \frac{\log^2(n)}{2} + \frac{\log(n)}{2} + 3 \\
 &= \cancel{\frac{1}{2}\log^2(n)} + \cancel{\frac{1}{2}\log(n)} + \cancel{3} \\
 &= \log(n) \\
 &= O(\log(n))
 \end{aligned}$$