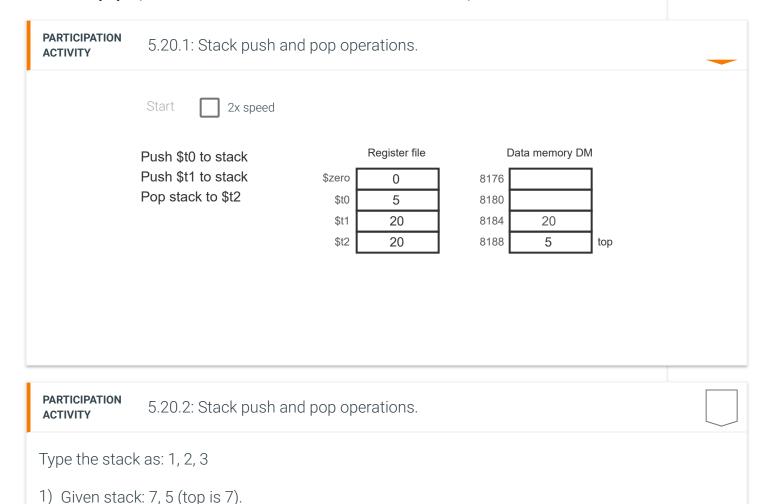
# 5.20 Subroutines and the program stack

### Stack

A **stack** is a data structure in which items are inserted on or removed from the top of the stack. A stack **push** operation inset the top of the stack. A stack **pop** operation removes and returns the item at the top of the stack.



|   | Type the stack after the following push operation: Push 8 to stack                       |                       |  |  |
|---|--|-----------------------|--|--|
|   | Check  | Show answer           |  |  |
| 2)  | Given stack: 3   | 34, 20 (top is 34)    |  |  |
|   | Type the stack after the following two push operations: Push 11 to stack Push 4 to stack |                       |  |  |
|   |  |                       |  |  |
|   | Check  | Show answer           |  |  |
| 3)  | Given stack: 5   | s, 9, 1 (top is 5)    |  |  |
|   | What is \$t1 af operation? Pop stack to \$   | ter the following pop |  |  |
|   | Check  | Show answer           |  |  |
| 4)  | Given stack: 5   | i, 9, 1 (top is 5)    |  |  |
| Type the stack after the following pop operation: Pop stack to \$t1 |  |                       |  |  |

|                                   | Check  | Show answer           |  |  |  |
|-----------------------------------|--|-----------------------|--|--|--|
| 5)                                | 5) Given stack: 2, 9, 5, 8, 1, 3 (top is 2). |                       |  |  |  |
|                                   |  | ter the following pop |  |  |  |
|                                   | operations?                                  |                       |  |  |  |
|                                   | Pop stack to \$t1                            |                       |  |  |  |
|                                   | Pop stack to \$                              |                       |  |  |  |
|                                   |  |                       |  |  |  |
|                                   | Check  | Show answer           |  |  |  |
| 6) Given stack: 41, 8 (top is 41) |  |                       |  |  |  |
|                                   | Type the stack after the following pop       |                       |  |  |  |
|                                   | operations:                                  |                       |  |  |  |
|                                   | Pop stack to \$t0 Push 2 to stack            |                       |  |  |  |
|                                   | Push 15 to stack                             |                       |  |  |  |
|                                   | Pop stack to \$t5                            |                       |  |  |  |
|                                   |  |                       |  |  |  |
|                                   |  |                       |  |  |  |
|                                   | Check  | Show answer           |  |  |  |
|                                   |  |                       |  |  |  |
|                                   |  |                       |  |  |  |

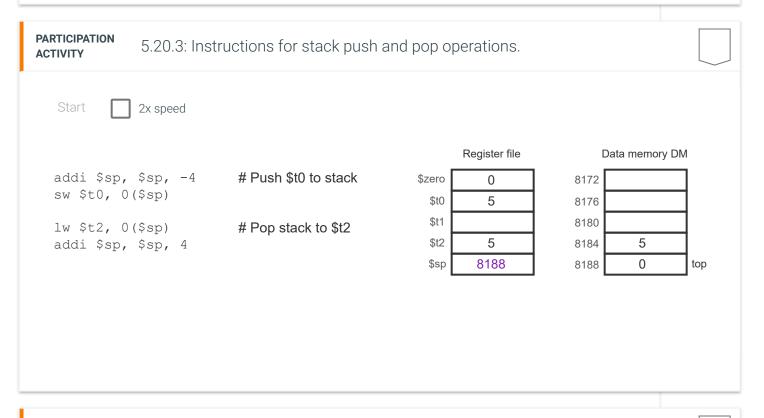
# Program stack and stack pointer

The **program stack** is a stack used by a program to store data for subroutines. The **stack pointer** (\$sp) register is used to h of the top of the program stack. In MIPSzy, the \$sp register is automatically initialized to the last data memory location, wh 8188. The MIPszy program stack is limited in size to 1KB, or 256 words. The stack grows toward decreasing memory address.

value to the stack first decrements \$sp by 4 and then copies the value held in a register to data memory at address \$sp. Pc from the stack first copies the top of the stack to a register and then increments \$sp by 4.

## Stack overflow

A **stack overflow** occurs when the number of values pushed to the stack exceeds the size allocated for the stack. Ex: Pushing 1005 values to the MIPSzy results in a stack overflow, as the stack size is limited to 1000 entries. A processor may have special circuitry to detect a stack overflow, allowing the system to execute special operation to handle the overflow, such as terminating the program.



PARTICIPATION ACTIVITY

5.20.4: MIPS program stack.

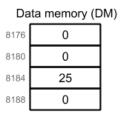
1) Complete the assembly to push the

value held in \$t3 to the stack. sw \$t3, 0(\$sp) Check **Show answer** 2) Complete the assembly to pop a value from the top of the stack to \$t4. addi \$sp, \$sp, 4 Check **Show answer** 3) Given the registers and stack content below, what is DM[8184] after the following instructions? addi \$sp, \$sp, -4 sw \$t0, 0(\$sp) Register file Data memory (DM) \$zero 0 8176 0 20 0 \$t0 8180 \$t1 15 0 8184 \$t2 43 8188 0 71 \$t3 8188 Check **Show answer** 

4) Given registers and stack content below,

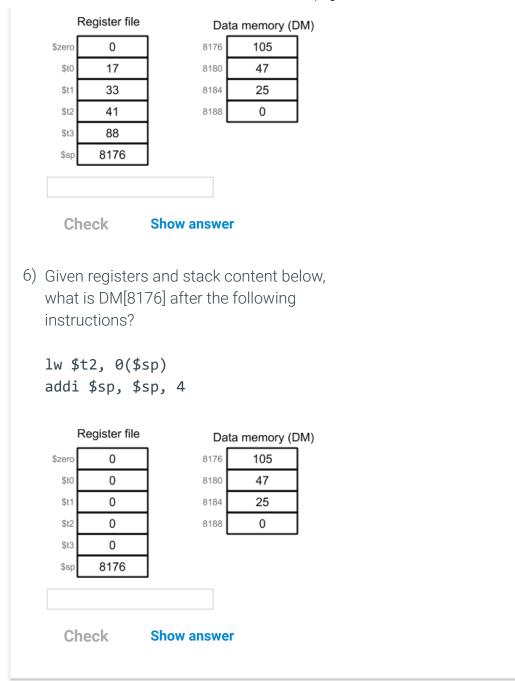
what is \$sp after the following instructions?

#### 



**Check** Show answer

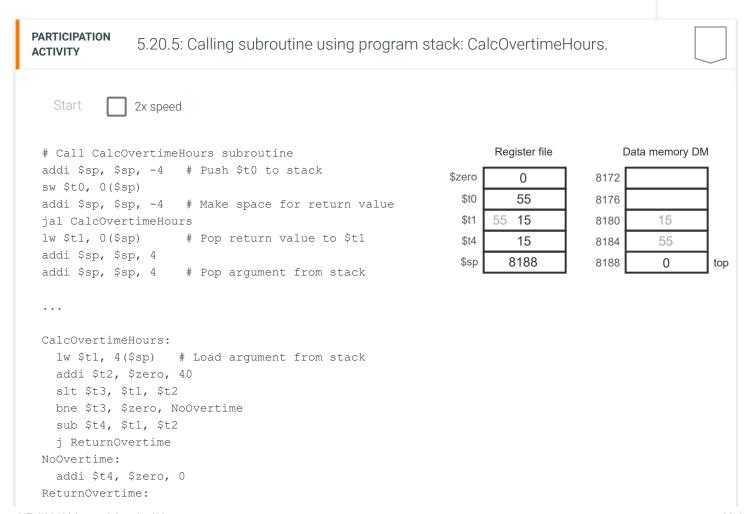
5) Given registers and stack content below, what is \$t3 after the following instructions?



# Using the program stack for subroutines

Because registers are limited, a subroutine call can use the program stack for arguments and return values rather than dire registers. The values stored in the program stack for a subroutine is called a **stack frame**. A subroutine call using the program performs the following steps.

- 1. Push subroutine arguments to program stack
- 2. Reserve space on program stack for the return value.
- 3. Jump to the subroutine.
- 4. Subroutine performs task storing return value to the reserved program stack location.
- 5. Subroutine returns.
- 6. Pop return value from program stack.
- 7. Pop arguments from program stack.



sw \$t4, 0(\$sp) # Copy return value to stack
ir \$ra

In the subroutine, an offset is used in lw or sw instructions to load arguments or store the return value. Ex: For a subroutine argument and 1 return value,  $\theta(\$sp)$  is the address for the return value, and 4(\$sp) is the address for the argument.

PARTICIPATION ACTIVITY

5.20.6: Calling subroutine using program stack.

Assume the program stack is used for all subroutine arguments and return values.

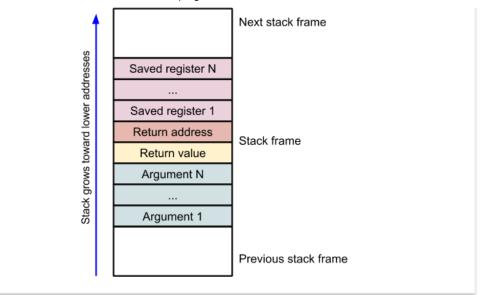
- 1) What is the stack frame size for a subroutine with one argument and one return value?
  - 0 1
  - **O** 2
- 2) What is the stack frame size for a subroutine with one argument and no return values?
  - 0
  - **O** 2
- 3) For a subroutine with 2 arguments and a return value, which instruction loads the first argument to \$11.
  - O lw \$t1, 0(\$sp)
  - O lw \$t1, 4(\$sp)

| Iw \$t1, 8(\$sp)   |  |
|--|--|
| 4) For a subroutine with 1 argument and a return value, which instruction stores \$t4 to the stack entry allocated for the return value. |  |
| O sw \$t4, 8(\$sp)   |  |
| O sw \$t4, 4(\$sp)   |  |
| O sw \$t4, 0(\$sp)   |  |

# Saving return address and registers to the program stack

If a subroutine calls another subroutine, the value held in \$ra must be saved before the second subroutine is called, becaus second subroutine writes a new value to \$ra. So, a subroutine that calls another subroutine will also push \$ra to the program the value held in registers used by a subroutine may still be needed by code that called the subroutine. To avoid overwriting subroutine can save the values held in any registers used by the subroutine to the program stack, and restore them before the subroutine. The following shows the organization for a complete MIPSzy stack frame.

Figure 5.20.1: MIPSzy stack frame.



PARTICIPATION ACTIVITY

5.20.7: Saving and restoring register using the program stack.

Complete the CalcQuadFunc subroutine to save and restore any registers used by the subroutine using the program stack by determine the missing instructions for the highlighted lines.

```
# Computes x^2 + 2x
CalcQuadFunc:
  # Save $t0, $t1, and $t2 to stack
  addi $sp, $sp, -4
  sw $t0, 0($sp)
  addi $sp, $sp, -4
  sw $t1, 0($sp)
   (a)
   (b)
  lw $t0, 16($sp)
                     # Load argument from stack
  addi $t1, $zero, 2
  mul $t2, $t0, $t1 # Calculate 2*x
  mul $t1, $t0, $t0 # Calculate x*x
  add $t2, $t2, $t1 # Calculate x*x + 2x
  sw $t2, 12($sp)
                     # Copy return value to stack
  # Restore $t0, $t1, and $t2 from stack
  (c)
  addi $sp, $sp, 4
  lw $t1, 0($sp)
  addi $sp, $sp, 4
   (d)
  addi $sp, $sp, 4
  jr $ra
 sw $t2, 0($sp) lw $t2, 0($sp) addi $sp, $sp, -4 lw $t0, 0($sp)
                                           (a)
                                           (b)
                                           (c)
                                           (d)
                                                                               Reset
```

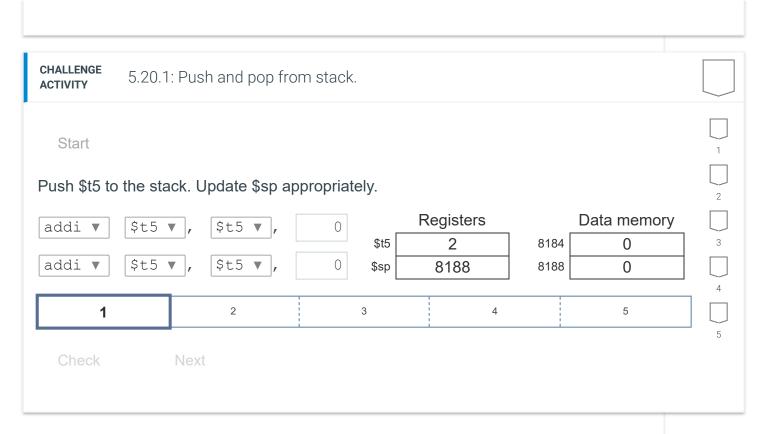
5.20.8: Saving and restoring registers using the program stack.

#### **ACTIVITY**

The CalcOvertimeHours subroutine uses registers \$11, \$12, \$13, and \$14. Modify the subroutine to save and restore those registers.

- 1. At the beginning of the subroutine, push the values held in \$11, \$12, \$13, and \$14 to the stack in that order.
- 2. Because more values have been pushed on the stack, the offsets for the location of the subroutine's argument and return value in the stack have changed. Modify the lw \$t1, 4(\$sp) and sw \$t4, 0(\$sp) instructions to use the correct offsets.
- 3. Just before the jr \$ra instruction, restore the saved registers' values by popping 4 values from the stack into \$t4, \$t3, \$t2, and \$t1, in that order. Note the order is reversed from when the values were pushed on the stack.

#### Registers Assembly 0 \$zero 5000 Line 1 # Load hours worked from DM[5000] 0 \$t0 5004 Line 2 addi \$t6, \$zero, 5000 Line 3 1w \$t0, 0(\$t6)\$t1 0 Line 4 5008 Line 5 # Call CalcOvertimeHours subroutine \$t2 0 Line 6 # Push \$t0 as argument to stack 8180 Line 7 addi \$sp, \$sp, -4 \$t3 0 Line 8 sw \$t0, 0(\$sp)8184 Line 9 # Reserve space for return value 0 \$t4 Line 10 addi \$sp, \$sp, -4 8188 Line 11 jal CalcOvertimeHours Line 12 # Pop return value to \$t1 0 \$t5 Line 13 lw \$t1, 0(\$sp) Line 14 addi \$sp, \$sp, 4 0 \$t6 Line 15 # Pop argument from stack Line 16 addi \$sp, \$sp, 4 8188 \$sp Line 17 line 18 # Store overtime hours to DM[5004] 0 \$ra RUN **ENTER SIMULATION** STEP



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