

Mimblewimble

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Outline



What is Mimblewimble?

- Mimblewimble is a completely new blockchain design that offers several benefits:
 - Privacy by default
 - Massively Prunable
 - Scalable
 - Relies on strong and proven cryptography (ECC)
- Can be implemented as a side chain or as an brand new cryptocurrency
- Output-based like Bitcoin but without script

History

2016

▶ On Tuesday, August 2 at 4:30 UTC, an individual logged on a Bitcoin IRC research channel and posted a link to a paper on a tor hidden service: Mimblewimble by Tom Elvis Jedusor. He then signed off and he never came back.

```
04:35 <majorplayer> hi, i have an idea for improving privacy in bitcoin. my friend who knows technology says this channel would have interest <a href="http://5pdcbqndmprm4wud.onion/mimblewimble.txt">http://5pdcbqndmprm4wud.onion/mimblewimble.txt</a>
MIMBLEWIMBLE
Tom Elvis Jedusor
19 July, 2016

\****/
Introduction
```

- The next week, after discussion on Reddit between Andrew Poelstra, Gregory Sanders and others Bitcoin developers, the idea is considered worth pursuing.
- ▶ On October 10, Andrew Poelstra publish a follow-up paper explaining in details Mimblewimble.

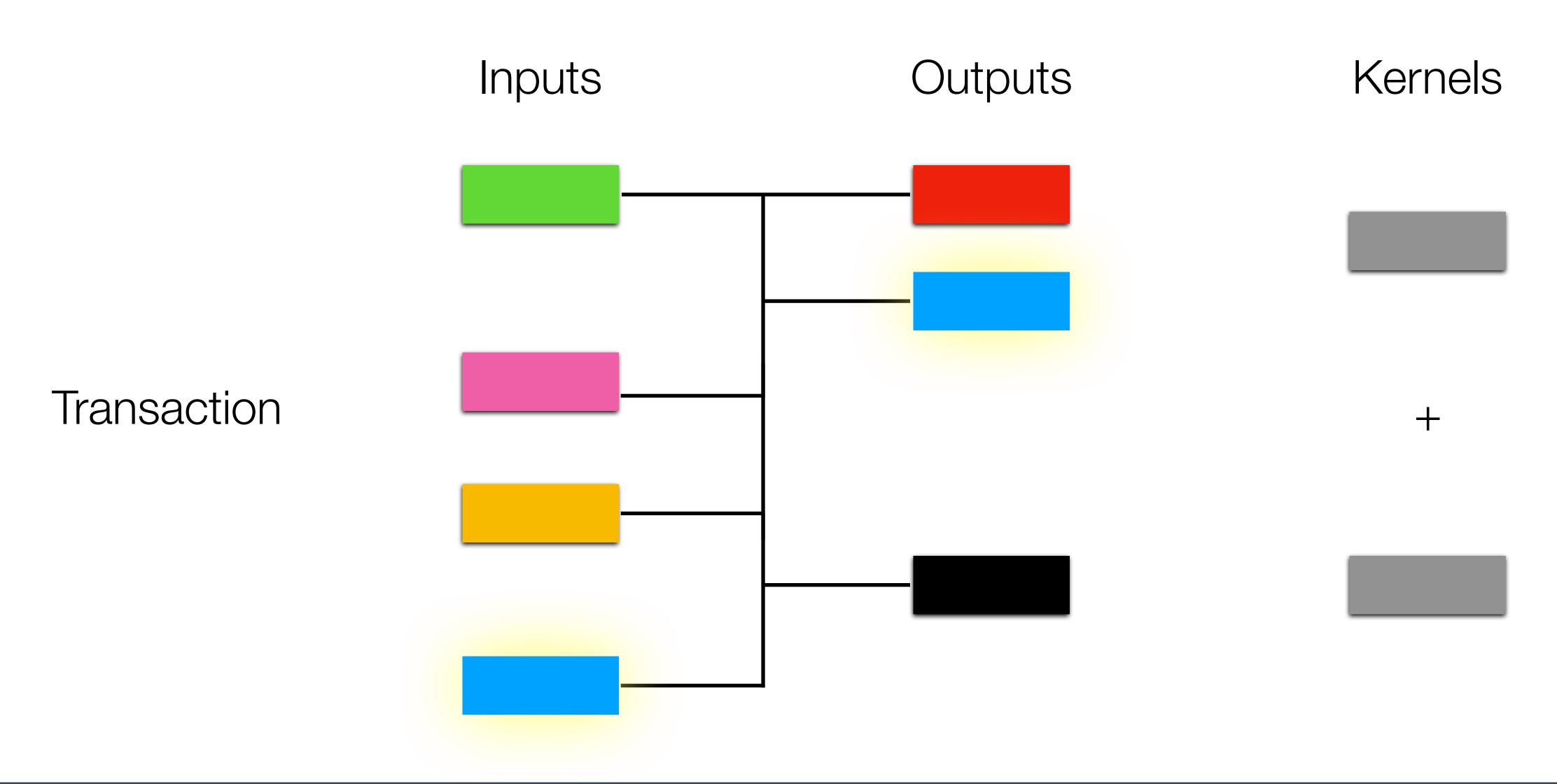
- In Bitcoin, every output has a script (script pub key) attached to it. In order to spend one of them, you must verify the conditions in the script.
- In Mimblewimble outputs only have public keys: no script.
- Hence Mimblewimble transactions are scriptless
 no payment channels, no atomic swap and multisignature at first glance

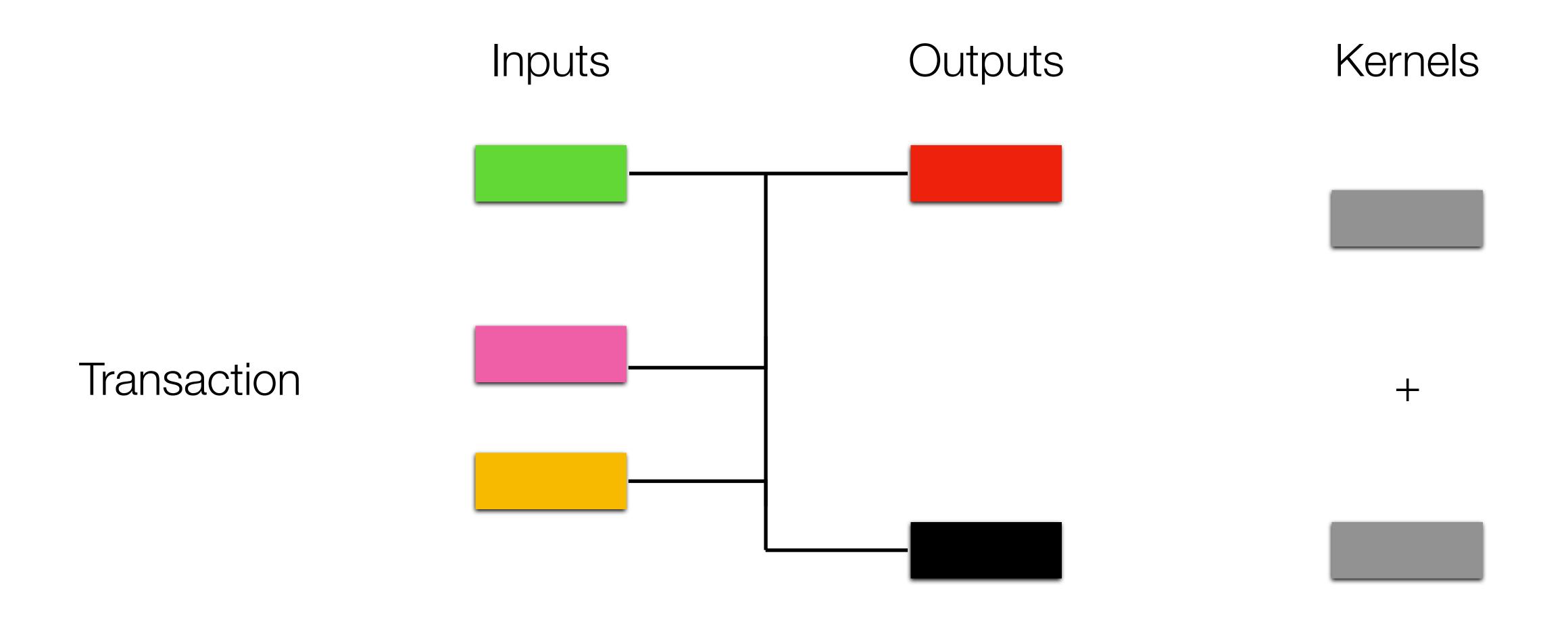
Mimblewimble transactions contains the following parts:

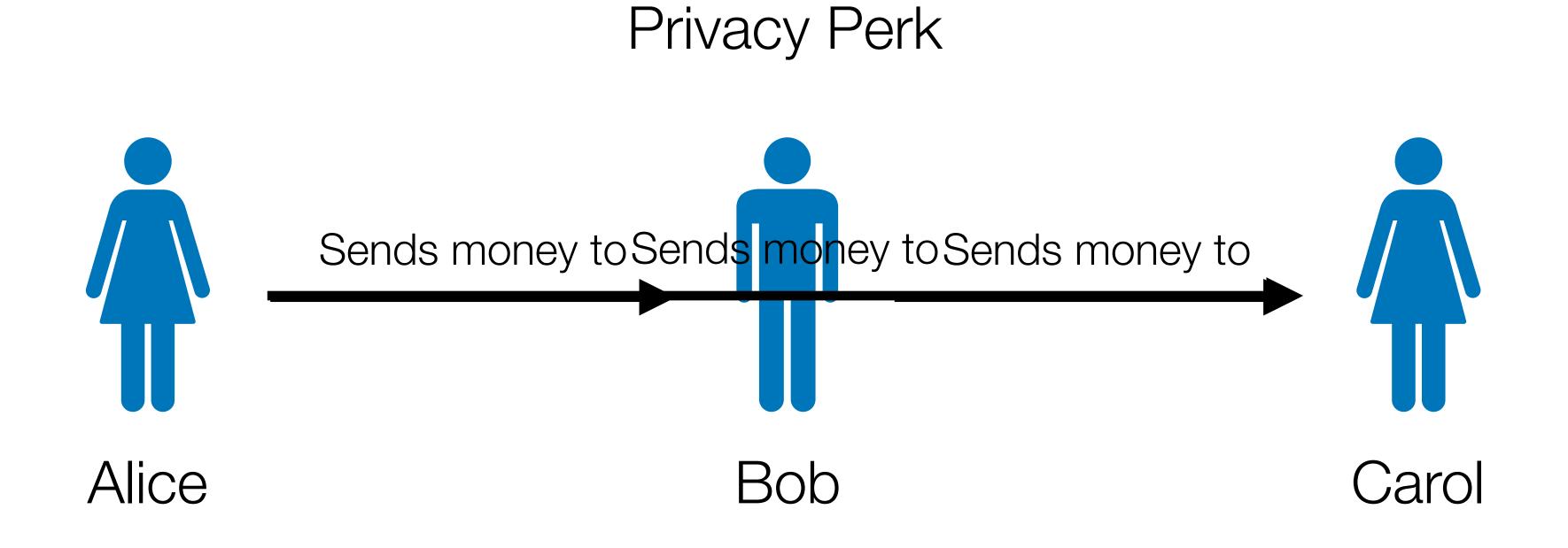
- Inputs (reference to old outputs)
- Outputs: confidential transactions (amounts are blinded) + rangeproofs (cryptographic proof that no money was created)
- ▶ Kernel: difference between outputs and inputs + mining fee + signature

Kernels Inputs Outputs Transaction 1 Transaction 2

Kernels Outputs Inputs Transaction +



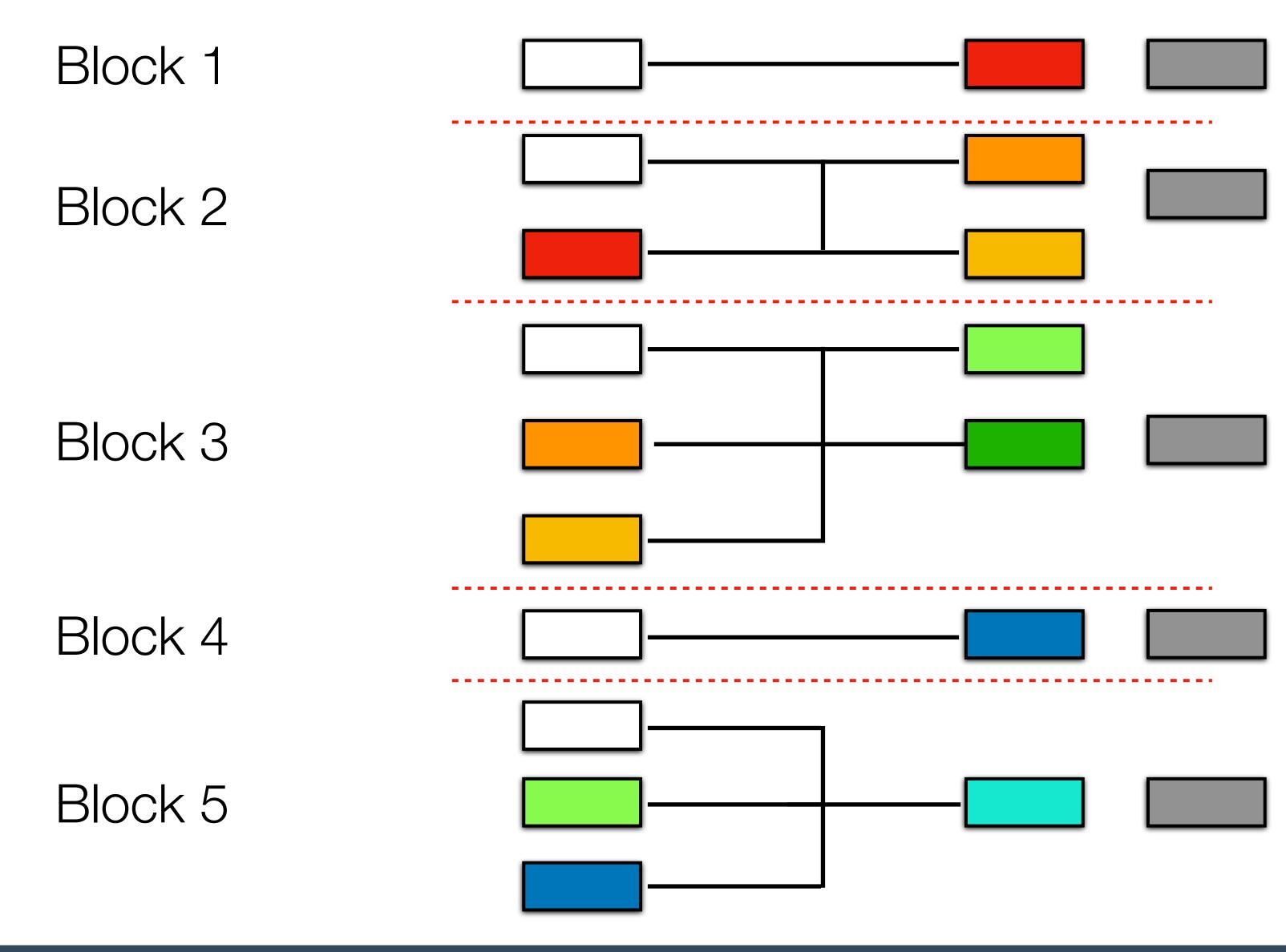


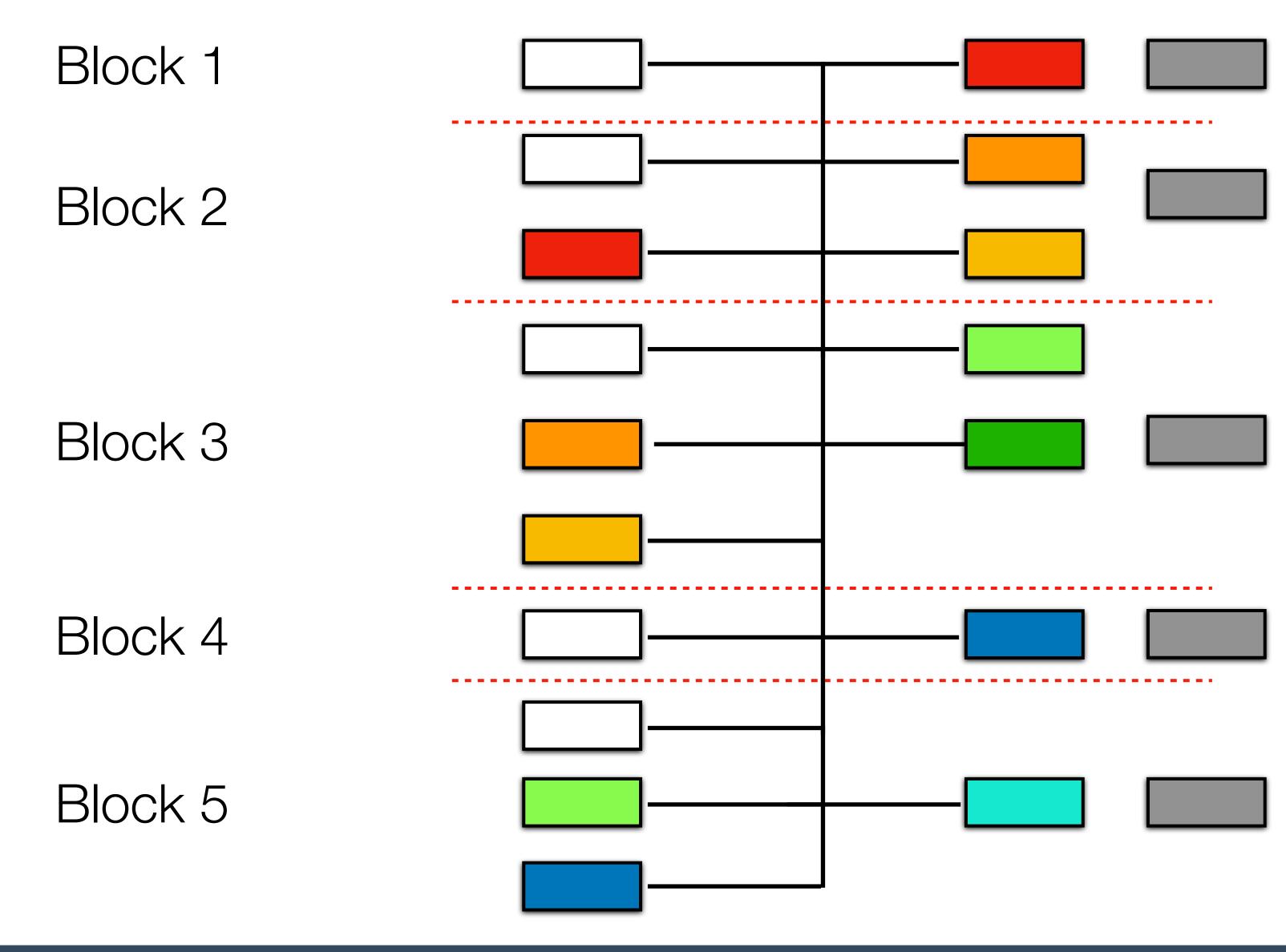


Mimblewimble transactions are a variant of a non-interactive Coinjoin transaction (every participant doesn't need to be online in order to group the transactions)

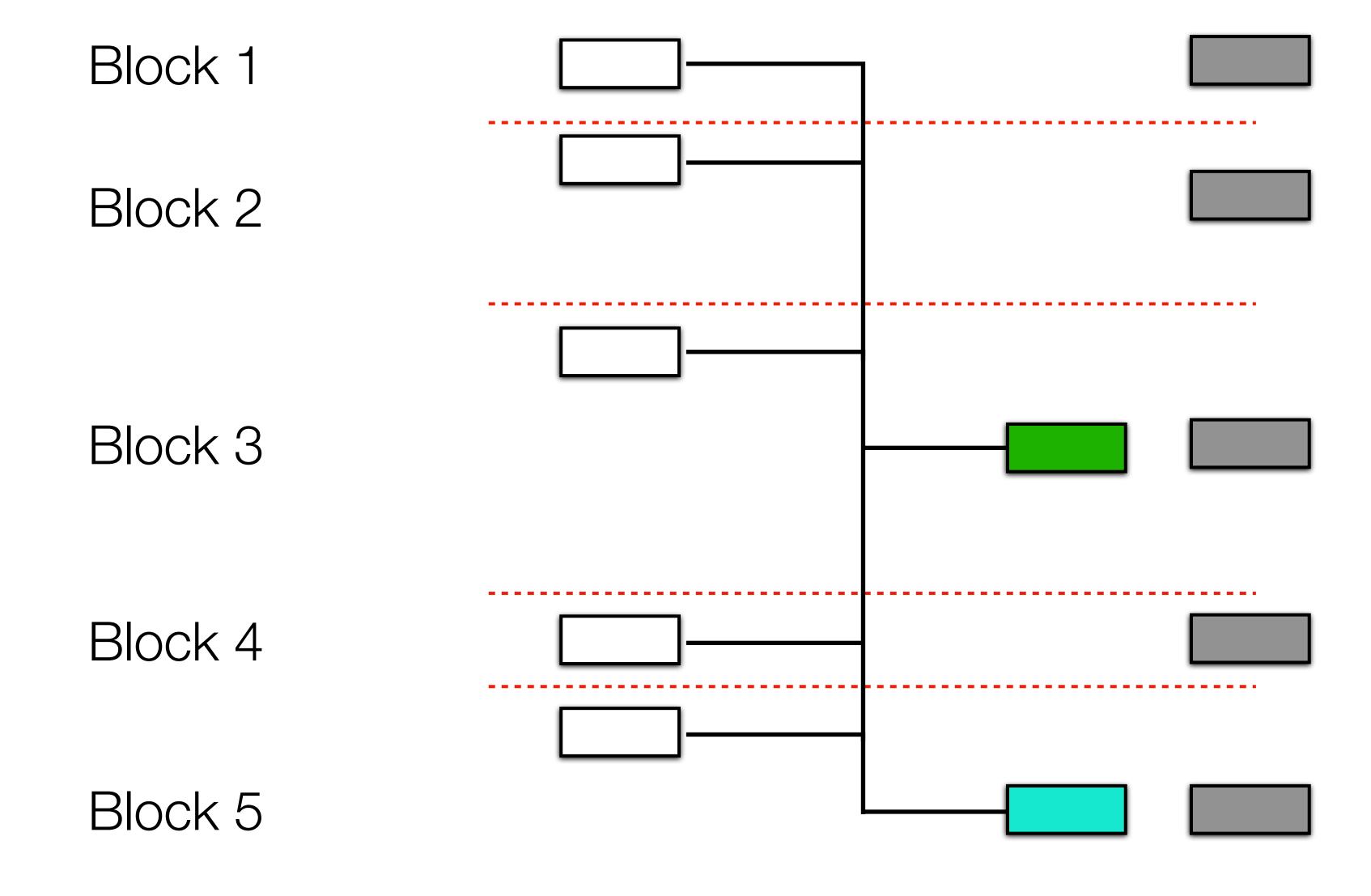
There is no addresses

▶ There is no amounts





Block 1 Block 2 Block 3 Block 4 Block 5



Applied to Bitcoin

- ▶ 150 millions transactions with approximately 400 millions outputs, 65 millions of which are unspent
- Currently 180 Gb; if you add CT 450 Gb.
- ▶ Mimblewimble with Bitcoin:
 - 18 Gb of transactions kernels, headers...
 - 2 Gb of UTXO
 - 45 Gb of UTXO rangeproofs.

Trust Model

- A transaction is valid if:
 - Signed by the owner of the inputs
 - The sum of outputs minus inputs always equals zero proving that the transaction did no create new funds
- Once in a block, transactions cannot be reversed without doing enough work
- Current state reflects zero net theft and inflation
- No need to know the exact sequence of transaction in order to verify that the blockchain is correct

Scriptless Script

A first glance, scripts are not possible on the MW blockchain. However, using some magic (with Adaptor Signature derived from Schnorr signature) its possible to do:

Mimblewimble

- ▶ Multi-signature transactions.
- Atomic swaps.
- Time-locked transactions and outputs.
- Lightning Network

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Recap

- Mimblewimble is a new blockchain design proposed by an anonymous
- No amounts and no addresses
- ▶ Transactions are aggregated in block to form one unique transactions: removing intermediaries.
- ▶ Resulting in a massively prunable blockchain

```
core::{Transaction, Input, Output, OutputFeatures, SwitchCommitHash, COINBASE_OUTPUT, DEFAULT_OUTPUT};
core::hash::Hash;
keychain;
keychain::{Keychain, BlindSum, BlindingFactor, Identifier};
util::LOGGER;
Context information available to transaction combinators.
struct Context<'a> {
   keychain: &'a Keychain,
Function type returned by the transaction combinators. Transforms a
(Transaction, BlindSum) pair into another, provided some context.
type Append = for<'a> Fn(&'a mut Context, (Transaction, BlindSum)) -> (Transaction, BlindSum);
Adds an input with the provided value and blinding key to the transaction
being built.
uild_input(
   value: u64,
   features: OutputFeatures,
   out_block: Option<Hash>,
   key_id: Identifier,
Box<Append> {
   Box::new(move |build, (tx, sum)| -> (Transaction, BlindSum) {
           let commit = build.keychain.commit(value, &key_id).unwrap();
           let input = Input::new(
                   features,
                   commit,
                   out_block,
           );
           (tx.with_input(input), sum.sub_key_id(key_id.clone()))
   })
Adds an input with the provided value and blinding key to the transaction
being built.
fn input(
   value: u64,
   out_block: Hash,
   key_id: Identifier,
Roy-Annends (
```

Grin

What is Grin?

- ► On October 20, 2016, "Ignotus Peverell" advised Andrew Poelstra that he was working on an implementation of the Mimblewimble blockchain: Grin. https://github.com/mimblewimble/grin
- From Gringotts, the wizard bank
- ▶ Shortly after, he was joined by "Antioch Peverell", "Garrick Olivander" and others HP characters.
- As of today 579 commits, 42 contributors and 778 stars on Github

What is Grin?

First implementation of Mimblewimble as a brand new cryptocurrency

Minimal implementation of the Mimblewimble protocol

In Rust (a programming language focused on safety, speed, and concurrency)



Use Proof-of-Work: cuckoo cycle



Currently under development, slides might be obsolete soon



What is Grin?

Development Funding

- ▶ Grin is a free open source software (FOSS)
- No Grin foundation and not a company, voluntary based development
- No ICO
- ▶ 100% community driven model
- Currently one developer full time funded Michael Cordner a.k.a Yeastplume

Cuckoo Cycle

- Created by John Tromp at the end of 2013
- First graph-theoretic proof-of-work system
- Memory-bound algorithm
- Designed to be ASIC resistant
- Only CPU and GPU (hopefully)

Cuckoo Cycle: a memory bound graph-theoretic proof-of-work

John Tromp

July 24, 2015

Abstract

We introduce the first graph-theoretic proof-of-work system, based on finding small cycles or other structures in large random graphs. Such problems are trivially verifiable and arbitrarily scalable, presumably requiring memory linear in graph size to solve efficiently. Our cycle finding algorithm uses one bit per edge, and up to one bit per node. Runtime is linear in graph size and dominated by random access latency, ideal properties for a memory bound proof-of-work. We exhibit two alternative algorithms that allow for a memory-time trade-off (TMTO)—decreased memory usage, by a factor k, coupled with increased runtime, by a factor $\Omega(k)$. The constant implied in $\Omega(k)$ gives a notion of memory-hardness, which is shown to be dependent on cycle length, guiding the latter's choice. Our algorithms are shown to parallelize reasonably well.

1 Introduction

A proof-of-work (PoW) system allows a verifier to check with negligible effort that a prover has expended a large amount of computational effort. Originally introduced as a spam fighting measure, where the effort is the price paid by an email sender for demanding the recipient's attention, they now form one of the cornerstones of crypto currencies.

As proof-of-work for new blocks of transactions, Bitcoin [1] adopted Adam Back's hashcash [2]. Hashcash entails finding a nonce value such that application of a cryptographic hash function to this nonce and the rest of the block header, results in a number below a target threshold¹. The threshold is dynamically adjusted by the protocol so as to maintain an average block interval of 10 minutes.

Bitcoin's choice of the simple and purely compute-bound SHA256 hash function allowed for an easy migration of hash computation from desktop processors (CPUs) to graphics-card processors (GPUs), to field-programmable gate arrays (FPGAs), and finally to custom designed chips (ASICs), with huge improvements in energy-efficiency at every step.

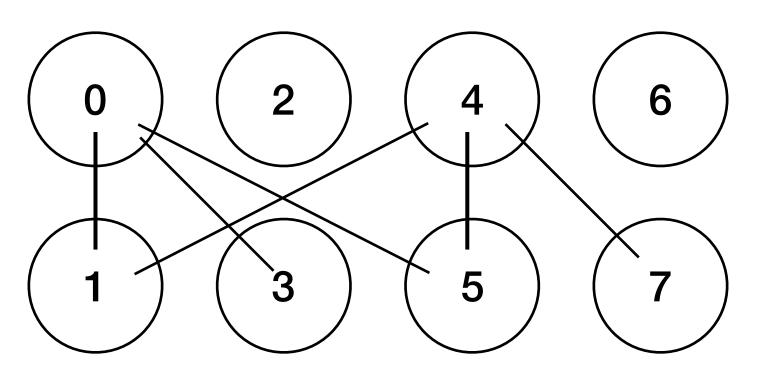
Since Bitcoin, many other crypto-currencies have adopted hashcash, with various choices of underlying hash function. the most well-known being *scrypt* as introduced by Tenebrix [3] (since faded into obscurity) and copied by Litecoin [4]. Scrypt, designed as a sequential memory-hard key derivation function, was specifically chosen to resist the migration away from CPUs and be "GPU-hostile". However, to adapt to the efficient verifiability requirement of proof-of-work, its memory footprint was severely limited, and migration slowed down only slightly.

Primecoin [5] introduced the notion of a number-theoretic proof-of-work, thereby offering the first alternative to hashcash among crypto-currencies. Primecoin identifies long chains of nearly doubled prime numbers, constrained by a certain relation to the block header. Verification of these chains, while very slow compared to bitcoin's, is much faster than attempting to find one. This asymmetry between proof attempt and verification is typical in non-hashcash proofs of work. Recently, two other

1

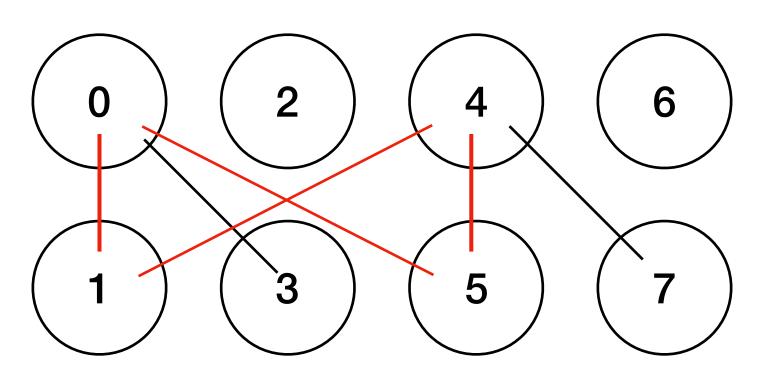
¹or, less accurately, results in many leading zeroes

▶ Goal: find "cycle" in a graph Find a series of connected nodes starting and ending on the same node e.g.: find a cycle of length 4 in the following graph

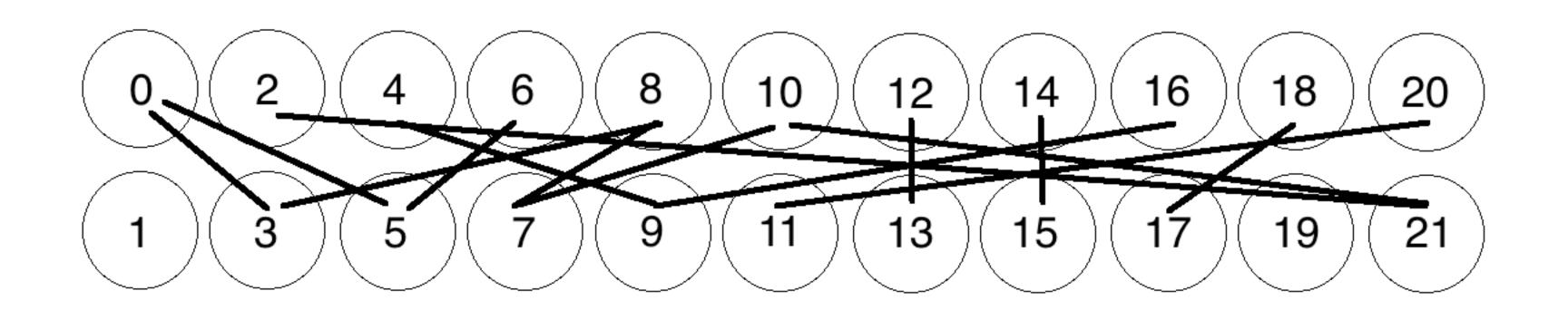


▶ Goal: find "cycle" in a graph Find a series of connected nodes starting and ending on the same node e.g.: find a cycle of length 4 in the following graph

0 - 5 - 4 - 1 -0

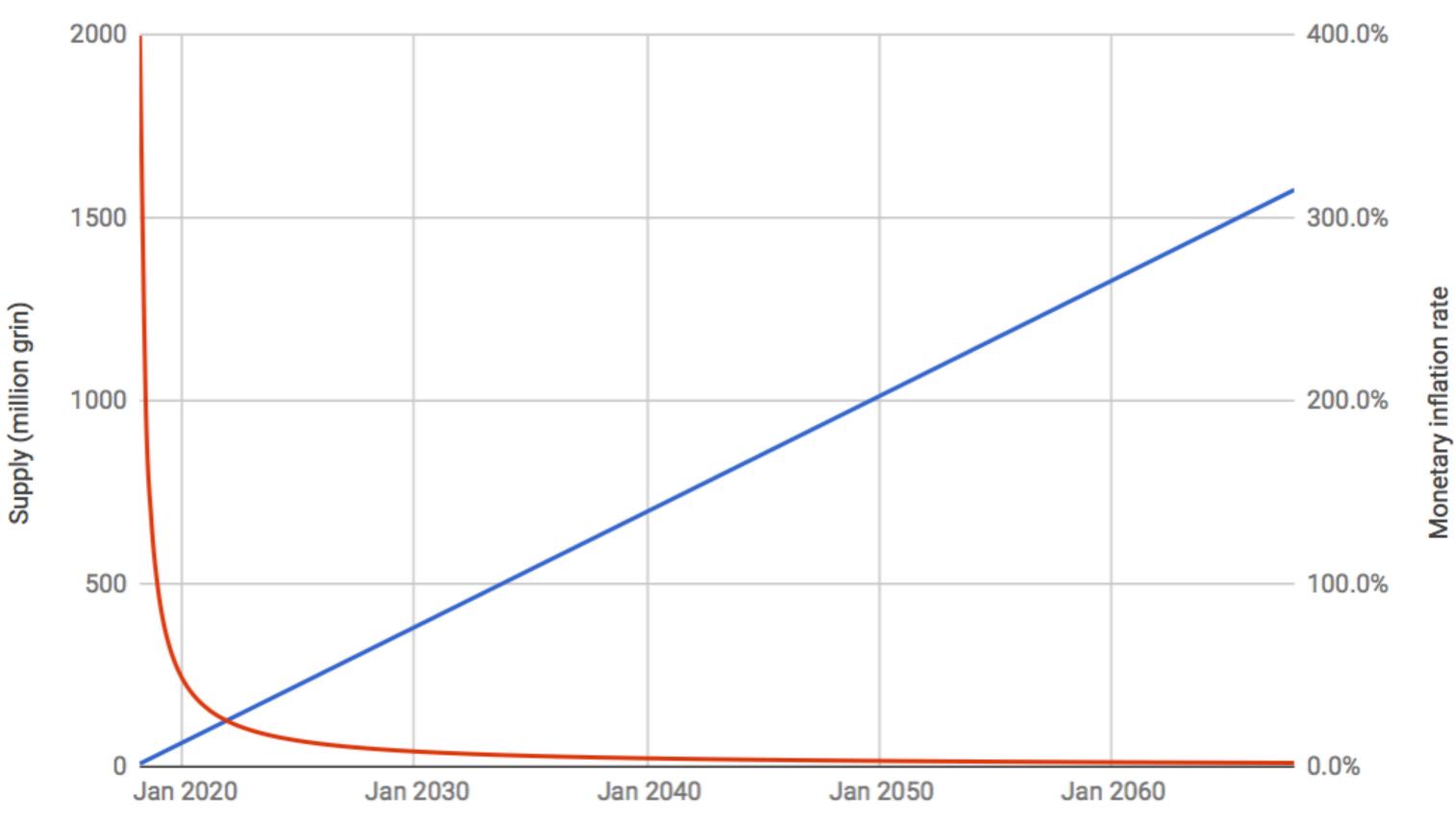


Easy?
Try to find a cycle of length 8 here

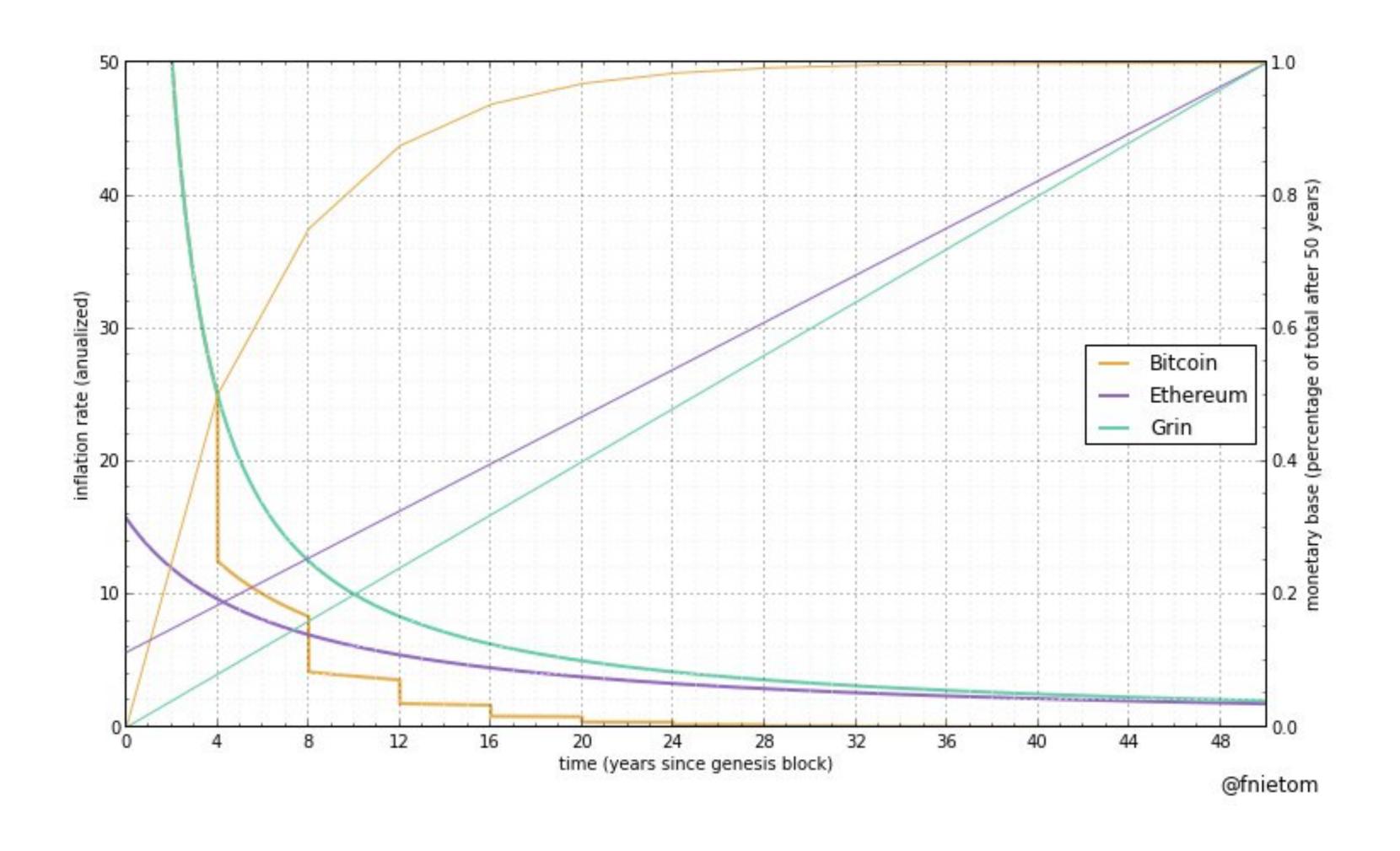


- ▶ 60 grins per block
- Fast block time: 1 block every minute
- ▶ 1 grin per second
- Difficulty adjusted every 23 blocks with difficulty calculation based on both Digishield (Digibyte) and GravityWave (Zcash)





Grin vs Bitcoin vs Ethereum Inflation



Grin vs Bitcoin vs Monero

| | GRIN | | |
|-----------------------------|---------------------------|---------------------------|---------------------------------|
| Transactions Format | Confidential Transactions | Confidential Transactions | Publicly Visible Transactions |
| Addresses | No addresses | Stealth addresses | Public list of addresses |
| Transactions Aggregation | Yes | No | No (but possible with Coinjoin) |
| Blockchain Pruning | Yes | No | No |

What's next?

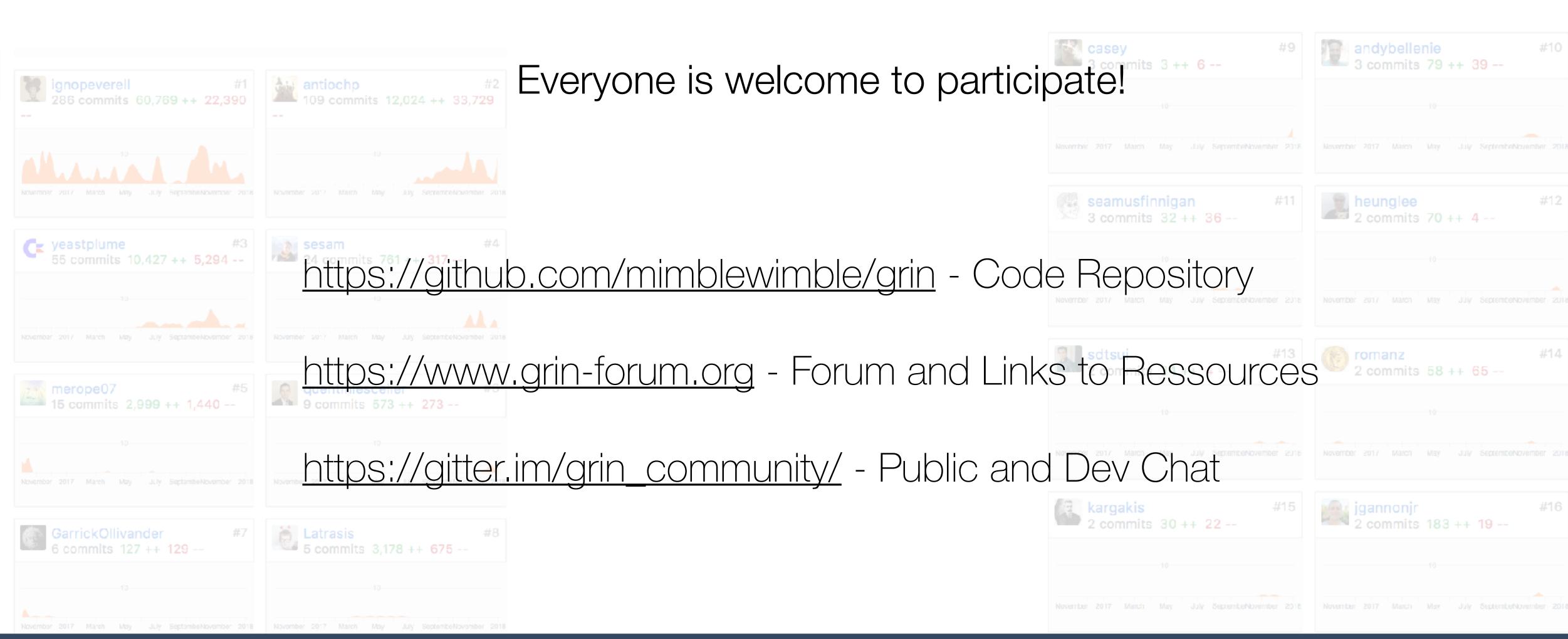
Lots of development is still needed

Currently on testnet1

Soon™ testnet2

Later this year (hopefully) mainnet

Get Involved



Thank you

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