Computer-Aided VLSI System Design

Homework 2: Simple RISC-V CPU

Graduate Institute of Electronics Engineering, National Taiwan University



Goal



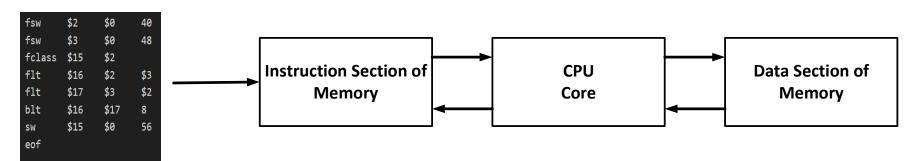
- In this homework, you will learn
 - How to write testbench
 - How to design FSM
 - How to use Memory IP
 - Generate patterns for testing
 - How to implement some RISC-V instruction operations, so designware is not allowed in this assignment

Introduction



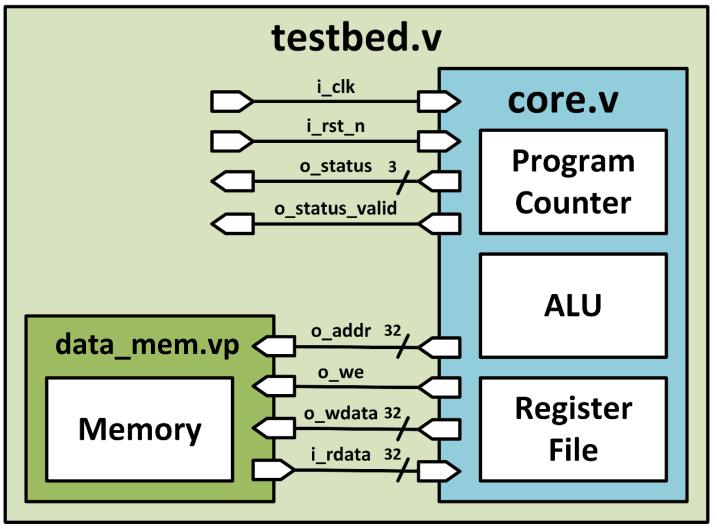
Central Processing Unit (CPU) is the important core in the computer system. In this homework, you are asked to design a simple RISC-V CPU [1], which contains the basic module of program counter, ALU and register files. The instruction set of the simple CPU is similar to RISC-V structure.

Instruction set



Block Diagram





Input/Output



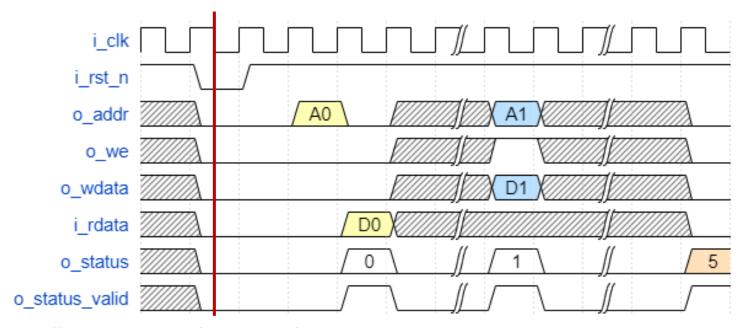
Signal Name	I/O	Width	Simple Description
i_clk	- 1	1	Clock signal in the system.
i_rst_n	I	1	Active low asynchronous reset.
o_we	0	1	Write enable of memory Set low for reading mode, and high for writing mode
o_addr	0	32	Address for memory
o_wdata	0	32	Data input to memory
i_rdata	I	32	Data or instruction output from memory
o_status	0	3	Status of core processing to each instruction
o_status_valid	0	1	Set high if ready to output status



- All outputs should be synchronized at clock rising edge.
- Memory is provided. All values in memory are reset to be zero.
- You should create 32 signed 32-bit registers and 32 singleprecision floating-point registers in register file.
- Less than 1024 instructions are provided for each pattern.
- The whole processing time can't exceed 120000 cycles for each pattern.



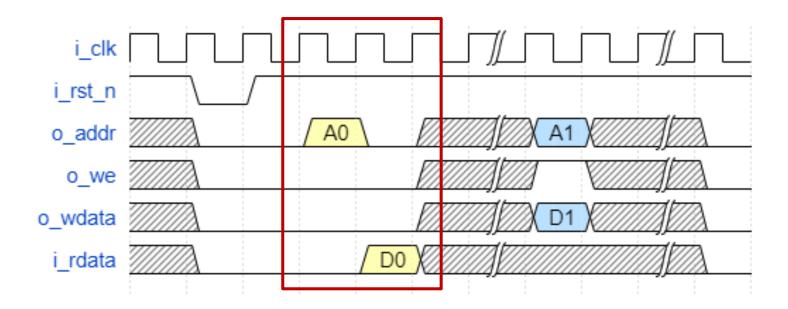
 You should set all your outputs and register file to be zero when i_rst_n is low. Active low asynchronous reset is used.



All output must be zero when reset

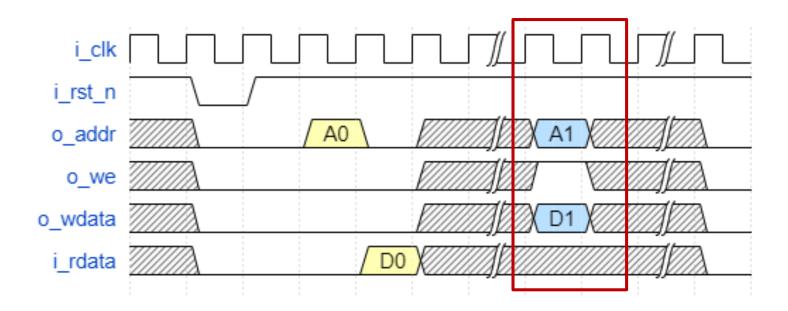


 To load data or instruction from the memory, set o_we to 0 and o_addr to relative address value. i_rdata can be received at the next rising edge of the clock.





 To save data to the memory, set o_we to 1, o_addr to relative address value, and o_wdata to the written data.



Instruction mapping



R-type

[31:25]	[24:20]	[19:15]	[14:12]	[11:7]	[6:0]
funct7	r2/f2	r1/f1	funct3	rd/fd	opcode

I-type

[31:20]	[19:15]	[14:12]	[11:7]	[6:0]
imm[11:0]	r1/f1	funct3	rd/fd	opcode

S-type

[3	1:25]	[24:20]	[19:15]	[14:12]	[11:7]	[6:0]
imn	n[11:5]	r2/f2	r1/f1	funct3	imm[4:0]	opcode

Instruction mapping (cont'd)



B-type

[31]	[30:25]	[24:20]	[19:15]	[14:12]	[11:8]	[7]	[6:0]
imm[12]	imm[10:5]	r2/f2	r1/f1	funct3	imm[4:1]	imm[11]	opcode

EOF

[31:7]			
Not used	opcode		

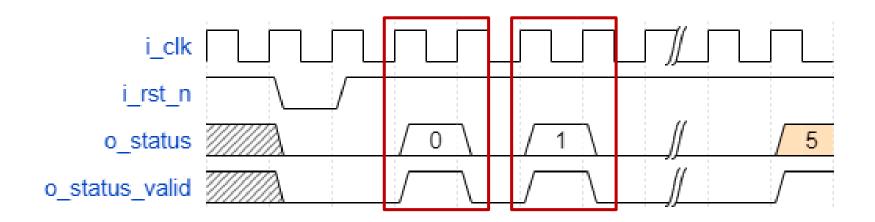
Status

7 statuses of o_status

o_status[2:0]	Definition
3'd0	R_TYPE_SUCCESS
3'd1	I_TYPE_SUCCESS
3'd2	S_TYPE_SUCCESS
3'd3	B_TYPE_SUCCESS
3'd4	U_TYPE_SUCCESS
3'd5	INVALID_TYPE
3'd6	EOF_TYPE

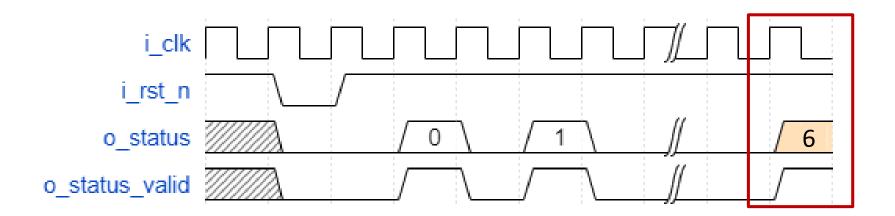


- Your o_status_valid should be turned to high for only one cycle for every o_status.
- The testbench will get your output at negative clock edge to check the o_status if your o_status_valid is high.



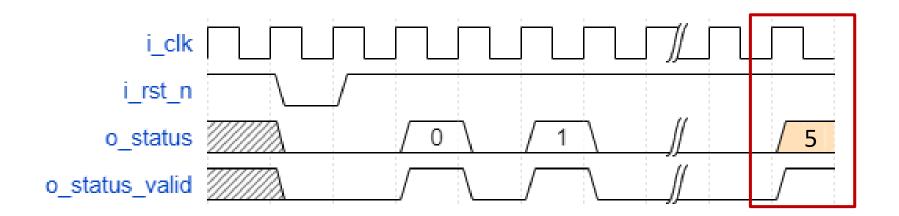


 When you set o_status_valid to high and o_status to 6, stop processing. The testbench will check your memory value with golden data.





• If invalid operation happened (see p.24~26), stop processing and raise o_status_valid to high and set o_status to 5. The testbench will check your memory value with golden data.

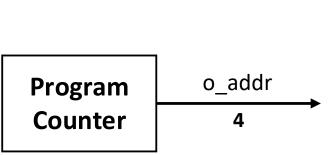


Program Counter



Program counter is used to control the address of memory for instruction.

\$pc = \$pc + 4 for every instruction (except beq, blt, jalr)



Memory					
Addr.	Instruction				
0	addi \$1 \$0 20				
4	addi \$2 \$0 12				
••	:				

Instruction



Operation	Assemble	Туре	Meaning	Note
Subtract	sub	R	\$rd = \$r1 - \$r2	Signed Operation
Add immediate	addi	I	\$rd = \$r1 + im	Signed Operation
Load word	lw	I	\$rd = Mem[\$r1 + im]	Signed Operation
Store word	SW	S	Mem[\$r1 + im] = \$r2	Signed Operation
Branch on equal	beq	В	if(\$r1==\$r2), \$pc = \$pc + im; else, \$pc = \$pc + 4	PC-relative Signed Operation
Branch less than	blt	В	if($r1 < r2$), $pc = pc + im$; else, $pc = pc + 4$	PC-relative Signed Operation
Jump and link register	jalr	l	rd = pc + 4; $pc = (r1 + im) & (\sim 0x1)$	PC-relative Signed Operation
Add upper immediate to PC	auipc	U	\$rd = \$pc + (im << 12)	PC-relative Signed Operation
Set on less than	slt	R	if(\$r1<\$r2), \$rd = 1; else, \$rd = 0	Signed Operation
Shift right logical	srl	R	\$rd = \$r1 >> \$r2	Unsigned Operation

Instruction (cont'd)



Operation	Assemble	Туре	Meaning	Note
Floating-point substract	fsub	R	\$fd = \$f1 - \$f2	Floating-point Operation
Floating-point multiply	fmul	R	\$fd = \$f1 * \$f2	Floating-point Operation
Floating-point to signed integer conversion	fcvt.w.s	R	\$rd = s32f32(\$f1)	Floating-point Operation
Load floating-point	flw	I	\$fd = Mem[\$r1 + im]	Signed Operation
Store floating-point	fsw	S	Mem[\$r1 + im] = \$f2	Signed Operation
Floating-point classify	fclass	R	\$rd = fclass(\$f1)	Classify floating-point format
End of File	eof	EOF	Stop processing	Last instruction in the pattern

Note: The notation of **im** in instruction is **2's complement.**

Note: The \$r notes that the data is read/written to **integer register file**; the \$f notes that the data is read/written to **floating-point register file**.

Note: Set the result of fsub and fmul to +0 if the arithmetic result is 0

Floating Point



- For instructions fsub, fmul, fcvt.w.s, fclass, you will have to implement operations with floating point format
- IEEE-754 single precision format [2]
 - 1 signed bit
 - 8 exponent bit
 - 23 mantissa bit

[31]	[30:23]	[22:0]
sign	exponent	mantissa

31

IEEE-754 Single Precision Format



[31]	[30:23]	[22:0]
sign	exponent	mantissa

31

Single-Format Bit Pattern	Value	
0 < e < 255	$(-1)^s \times 2^{e-127} \times 1.m$ (normal numbers)	
$e = 0$; $m \ne 0$ (at least one bit in f is nonzero)	$(-1)^s \times 2^{-126} \times 0.m$ (subnormal numbers)	
e = 0; m = 0 (all bits in f are zero)	$(-1)^s \times 0.0$ (signed zero)	
s = 0; e = 255; m = 0 (all bits in f are zero)	+INF (positive infinity)	
s = 1; e = 255; m = 0 (all bits in f are zero)	-INF (negative infinity)	
e = 255; m \neq 0 (at least one bit in f is nonzero)	NaN (Not-a-Number)	

Round to Nearest Even



• For instructions **fsub**, **fmul**, **fcvt.w.s**, you will have to round the mantissa or decimal with **round to nearest even** [3]

Rounding

1.BBGRXXX

Guard bit: LSB of result -

Sticky bit: OR of remaining bits

Round bit: 1st bit removed

Round up conditions

- Round = 1, Sticky = 1 → > 0.5
- Guard = 1, Round = 1, Sticky = 0 → Round to even

Value	Fraction	GRS	Incr?	Rounded
128	1.0000000	000	N	1.000
15	1.1010000	100	N	1.101
17	1.0001000	010	N	1.000
19	1.0011000	110	Y	1.010
138	1.0001010	011	Y	1.001
63	1.1111100	111	Y	10.000

Floating Point Classification



For instruction **fclass**, you will have to classify the floatingpoint number stored in registers and a bitmask with the corresponding bit set should be written to the destination

register

rd bit	Meaning		
0	Negative infinite		
1	Negative normal number		
2	Negative subnormal number		
3	Negative zero		
4	Positive zero		
5	Positive subnormal number		
6	Positive normal number		
7	Positive infinite		
8	Signaling NaN		
9	Quiet NaN		

Memory IP



- Size: 2048 × 32 bit
- i_addr[12:2] for address mapping in memory
- Instructions are stored in address 0 address 4095 of memory
- Data should be read from and written to address 4096 address
 8191 of memory

Invalid operation

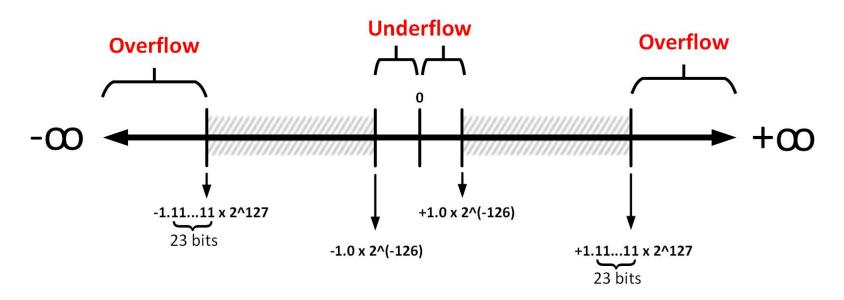


- Invalid operation may happened.
 - Situation1: Overflow happened at integer arithmetic instructions (sub, addi)
 - Situation2: Infinite, NaN happened at floating-point arithmetic instructions (fsub, fmul, fcvt.w.s)
 - For fcvt.w.s instruction, additional consideration is required when the floating-point value being converted to a 32-bit signed integer exceeds the representable range of a 32-bit signed integer
 - Do not consider when loading/storing infinite or NaN numbers from memory
 - Do not consider when executing fclass on infinite or NaN numbers

Invalid operation



- Invalid operation may happened.
 - <u>Situation3</u>: Overflow and underflow result happened at floating-point arithmetic instructions (**fsub**, **fmu1**)
 - Consider the overflow and underflow before rounding arithmetic result
 - Underflow does not include zero



Invalid operation



- Invalid operation may happened.
 - Situation4: If output address are mapped to unknown address in memory.
 - Consider the case when trying to load/store the address of memory for instruction
 - Consider the case when program counter is fetching instruction from the address of memory for data
 - Do not consider the case if instruction address is beyond eof, but the address mapping is in the size of memory for instruction

rtl.f



Filelist

core.v





```
nodule core #( // DO NOT MODIFY INTERFACE!!!
   parameter DATA WIDTH = 32,
   parameter ADDR WIDTH = 32
   input i clk,
   input i rst n,
  output [2:0] o_status,
   output
               o_status_valid,
   // Memory IOs
   output [ADDR_WIDTH-1:0] o_addr,
   output [DATA_WIDTH-1:0] o_wdata,
   output
                           o_we,
   input [DATA_WIDTH-1:0] i_rdata
```

define.v

- Do not modify
- Define all the 17 instruction patterns

```
DO NOT MODIFY THIS FILE
// status definition
define R_TYPE 0
define I_TYPE 1
define S TYPE 2
define B TYPE 3
define U TYPE 4
define INVALID_TYPE 5
define EOF_TYPE 6
// opcode definition
define OP_SUB
                 7'b0110011
define OP ADDI 7'b0010011
define OP LW
                7'b0000011
define OP SW
                 7'b0100011
define OP BEQ
                7'b1100011
define OP_BLT
                7'b1100011
define OP JALR
               7'b1100111
define OP_AUIPC 7'b0010111
define OP_SLT
                 7'b0110011
define OP SRL
                 7'b0110011
define OP FSUB
               7'b1010011
define OP FMUL
                7'b1010011
define OP_FCVTWS 7'b1010011
define OP_FLW
                 7'b0000111
define OP_FSW
                 7'b0100111
define OP_FCLASS 7'b1010011
define OP_EOF
                 7'b1110011
```

```
funct7 definition
define FUNCT7 SUB
                     7'b0100000
define FUNCT7_SLT
                     7'b0000000
define FUNCT7 SRL
                     7'b0000000
define FUNCT7 FSUB
                     7'b0000100
define FUNCT7 FMUL
                     7'b0001000
define FUNCT7 FCVTWS 7'b1100000
define FUNCT7 FCLASS 7'b1110000
/ funct3 definition
define FUNCT3_SUB
                     3'b000
define FUNCT3 ADDI
                     3'b000
define FUNCT3_LW
                     3'b010
define FUNCT3 SW
                     3'b010
define FUNCT3 BEQ
                     3'b000
define FUNCT3 BLT
                     3'b100
define FUNCT3 JALR 3'b000
define FUNCT3_SLT
                     3'b010
define FUNCT3 SRL
                     3'b101
define FUNCT3_FSUB
                     3'b000
define FUNCT3_FMUL
                     3'b000
define FUNCT3_FCVTWS 3'b000
define FUNCT3_FLW
                     3'b010
define FUNCT3 FSW
                     3'b010
define FUNCT3_FCLASS 3'b000
```

testbed_temp.v



- Things to add in your testbench
 - Clock
 - Reset
 - Waveform file
 - Function test
 - **—** ...

```
module testbed;

wire clk, rst_n;
wire dmem_we;
wire [ 31 : 0 ] dmem_addr;
wire [ 31 : 0 ] dmem_wdata;
wire [ 31 : 0 ] dmem_rdata;
wire [ 1 : 0 ] mips_status;
wire mips_status_valid;
```

```
core u core (
    .i_clk(),
   .i_rst_n(),
    .o_status(),
    .o_status_valid(),
   .o_we(),
    .o_addr(),
    .o_wdata(),
    .i_rdata()
data_mem u_data_mem (
    .i_clk(),
    .i_rst_n(),
    .i_we(),
    .i_addr(),
    .i_wdata(),
    .o_rdata()
```

Protected Files

- The following files are protected
 - data_mem.vp

```
module data mem (
                      i clk,
    input
    input
                      i_rst_n,
    input
                      i we,
    input [ 31 : 0 ] i addr,
    input [ 31 : 0 ] i wdata,
    output [ 31 : 0 ] o rdata
 protected
&6JU@A,>B[ZKNH#f\\dWJ5ZgKY/4LTZcTK[9H@IT99E YU\L\&A8-)gLM#\H80&9
CAINT2\;]80c#b5-A;1-?4M?C77#/U@ 1&DWDI#/gT[Vd?L&5U#I6::&,-e822f.
dPcB[;AOLA8FQd+Td+L2#YEY+#D1JX1Q#6TF0N^_2@aJc(RIWe8:AN=DV.0XBTP-
B,<E/\4X\GAJbWfYF)g07^)83,802)?K+>I,9M(UXOSg2?g4RW:^,Y^?JH28>J=8
2FK>6\HU(3?LIBQQK9(:WZ+e/KCQgI/<T8FPN0KCIcU/1.=L;VQcB03PPV+G_:1\
8N,g9>],5^](9f(g?^R[DW>/[/OTa>S).K4-C=85)5S>FC6LaO\2g9Q+,Ad7fBF?
b6XA=:M7[ 3COF+ 59;H6E-Dfc7#U+&/A/A]WDWU>QUW.124=b>LE5EE04f6J:W)
44Za5?:](CHHVagBN[2/dBWMJ?2NgZ6,WN^P[W@YaI+,0]=Yb_W+?5AK/\a>SBF-
Z6M; KM/.e05RCFK+ M?\IJII8)@,@J1N^DOE033(<Rg3df<=W#b|EB3dc0g[TOb
09CRJ3G3+DbS=;VI?_&/1f-VHY/5:WE,U<3g;#d]0eRaUU4-BDZ9P-@U\Q_4&W[B
IEB(fLJM45&JGf.&MX@=N#QdV1@;gc#d0ZR/Kc@6+PfE17d.+S0f6L(+(QON-KUM
4FHe<QSVE;JNgd1U(Z0D1B57Z]RZWU^L>;>ZITDL1T?-)\E=KEF<8]5I019@fZA-
f4;NUL/a9(7</dS#+;: 9aX4P&UC^8:=1g-,b&F4I5=P e[6+99gHL+a]W/R8C()
P40gM; E>@Y1V1d9/fIP7PN1:#ffG-FUS=@?bU9SE(>=^dL,; D00X0RU00ZKaX, \
@,GLKWM,gX:DcdF2W@8M92XHHdcN>Q?M03I,C9HLE(@3=G/bb[J;TB=gLTSBB>f2
0S0; V?<6, FW=NI@^H#<aM]@)29VETb]B1Cg7gN(9CC@-2TR/;NDFdF=gM$
 endprotected
```

Command



- 01_run
 - Usage: ./01_run p0

```
vcs -full64 -R -f rtl.f +v2k -sverilog -debug_access+all +define+$1 | tee sim.log
```

99_clean_up

```
rm -rf *.history *.key *.log
rm -rf novas.rc novas.fsdb novas.conf
rm -rf INCA_libs nWaveLog BSSLib.lib++
```

PATTERN (1/2)



Files in PATTERN are for your references

inst_assembly.dat

R-type	\$rd	\$r1 \$	r2
I-type	\$rd	\$r1 i	m
S-type	\$r2	\$r1 i	m
B-type	\$r1	\$r2 i	m
addi	\$0	\$0	1024
addi	\$1	\$0	32
addi	\$3	\$30	127
addi	\$1	\$31	-64
sub	\$4	\$2	\$0
add	\$5	\$2	\$ 3
beq	\$5	\$0	12
addi	\$6	\$8	2
s11	\$0	\$0	\$ 6
SW	\$4	\$0	0
lw	\$4	\$0	64
slt	\$2	\$8	\$10
add	\$4	\$3	\$3

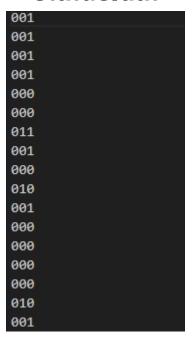
inst.dat

PATTERN (2/2)



- Data in data.dat includes golden data in memory
- You can compare the data in memory with the golden data directly

status.dat



data.dat

Grading Policy



TA will run your code with following command (p0 is example)

vcs -f rtl.f -full64 -sverilog -R -debug_access+all +define+p0 -v2k

- Pass the patterns to get full score
 - Provided pattern: 70% (4 patterns in total)
 - 15% for each test
 - 10% for spyglass check (lint_rtl and lint_rtl_enhanced)
 - Hidden pattern: 30%
 - 20 patterns in total

Grading Policy



- **Deadline:** 2025/10/14 13:59:59 (UTC+8)
- Late submissions are not accepted
 - Any submission after the deadline will receive 0 points
- File corrections after the deadline should be avoided
 - Corrections for the folder name, file name, file hierarchy cause 5-point deduction
- The TA will grade your submissions by using scripts
- No plagiarism
 - Plagiarism in any form, including copying from online sources, is strictly prohibited

Submission



 Create a folder named studentID_hw2, and put all below files into the folder

```
r13943119_hw2/

O1_RTL
core.v
rtl.f
(other design files)
```

- Compress the folder studentID_hw2 in a tar file named studentID_hw2_vk.tar (k is the number of version, k =1,2,...)
 - Use lower case for student ID. (Ex. r13943119_hw2_v1.tar)
- Submit to NTU Cool

Hint

- Design your FSM with following states
 - 1. Idle
 - 2. Instruction Fetching
 - 3. Instruction decoding
 - 4. ALU computing/Load data
 - 5. Data write-back
 - 6. Next PC generation
 - 7. Process end

Discussion



- NTU Cool Discussion Forum
 - For any questions not related to assignment answers or privacy concerns, please use the NTU Cool discussion forum.
 - TAs will prioritize answering questions on the NTU Cool discussion forum
- **Email:** r13943119@ntu.edu.tw
 - Title should start with [CVSD 2025 Fall HW2]
 - Email with wrong title will be moved to trash automatically

Reference



- [1] RISC-V User Manual
 - https://riscv.org/wp-content/uploads/2017/05/riscv-specv2.2.pdf
- [2] IEEE 754 Single Precision Format
 - https://zh.wikipedia.org/zh-tw/IEEE 754
- [3] Round to Nearest Even
 - https://www.cs.cmu.edu/afs/cs/academic/class/15213-s16/www/lectures/04-float.pdf
- [4] F Standard Extension for Single-Precision Floating-Point
 - https://five-embeddev.com/riscv-user-isa-manual/Privv1.12/f.html