ECE/CS 250 Midterm Exam #1, Spring 2017

Name:
Duke students are bound by an academic integrity standard:
1. I will not lie, cheat, or steal in my academic endeavors, nor will I accept the actions
of those who do.
2. I will conduct myself responsibly and honorably in all my activities as a Duke stu-
dent.
Please sign your name below to acknowledge that you follow this standard:

PLEASE CAREFULLY READ THE QUESTIONS.

MAKE SURE YOU ARE ANSWERING THE QUESTIONS AS WRITTEN.

NOTE: This question says "6-bit 2s complement" but should've said "7-bit 2s complement"

1) [2 points] (a) Represent the base-10 number 1	
tation. Put your final answer in the box. (No n	eed to snow work.)
	0010001
(b) [3 points] Represent the base-10 number -55 tion. Show your work and then put your final a	<u>-</u>
tion. Show your work and then put your imar a	iswer in the box.
	1001001
	: 2 1 4 2 5
(c) [5 points] Add these two numbers together us NOT invert the negative number and then add to	_
must just directly add the two numbers. Show	
	γ your work and put the final resul
in the box.	your work and put the final resul
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2) (a) [2 points] Write 187_{10} in hexadecimal (your final result in the box.	base 16). Show your work and put
	0xBB
(b) [8 points] The IEEE 754 floating point st point numbers have one sign bit, an 8-bit expor- significand (with an implicit "1"). What floatin following 32 bits? Show your work and put you	nent (with a bias of 127), and a 23-bit g point number is represented by the
0 01000001 011000000000000000000000	
msb = 0> positive	
exponent = 65 unsigned> exponent is 65 - 127 = -62	
mantissa is 1.011000 0	
This ended up being kind of ugly because of the exponent. I	If you wrote 1.011 x 2^-62, you got credit.
[

3) [5 points	[] (a) You	bought	software	that ru	ins on	your	laptop	that	has	an	Intel
processor.	Why doe	s that so	ftware no	t run or	a sma	artpho	one with	n an A	ARM	[pr	oces-
sor?											

They have different ISAs. A computer can only execute its ISA. (E.g., all AMD processors execute the x86 ISA.)

- (b) [5 points] Which of the following issues are part of the instruction set architecture? Circle the issues that are part of the ISA.
 - The number of registers YES
 - The number of bits in each register YES
 - A detailed specification of the hardware that performs addition NO
 - The memory addressing modes YES
 - The number of memory addresses YES
- (c) [5 points] Are the following statements true or false? Circle one for each.

All ISAs are pretty much like MIPS.	TRUE	FALSE	F
A MIPS jump (j) instruction has a register operand.	TRUE	FALSE	F
In C, all large variables should be malloc'ed.	TRUE	FALSE	F
I can copy a string by copying the pointer to its 1st char.	TRUE	FALSE	F
A carry-out in 2s complement addition always indicates an overflow.	TRUE	FALSE	F

4) (a) [5 points] Write <u>one line</u> of C code to dynamically allocate space for 16 pointers to characters on the heap.

```
char** dan_ptr = (char**)malloc(16*sizeof(char*));
```

(b) [10 points] At the end of the following snippet code, what are the values of A, B, C, D, and E? Put your results in the table to the right. If any line that writes to A, B, C, D, or E causes a seg fault, write "seg fault for that value" and continue with the program (even though a seg fault would normally end a program). Assume a 32-bit machine with 32-bit ints.

int numbers[100]; // assume numbers[0] is at address 1000

```
for (i=0; i<50; i++){
    numbers[i] = i;
}
for (i=50; i<100; i++){
    numbers[i] = (int) & numbers[i];
}
int A = *(numbers+2);
int B = numbers[52];
int* x = (numbers + 10); // assume x is at address 2000
int** y = &x; // assume y is at address 2004
int* C = *y; // assume C is at address 2008
int D = *C; // assume D is at address 2012
int* E = numbers + 100; // assume E is at address 2016</pre>
```

A	2
В	1208
С	1040
D	10
Е	1400

5) [20] Convert the following C code for the function f() into MIPS code. Use appropriate MIPS conventions for procedure calls, including the passing of arguments and return values, as well as the saving/restoring of registers. Assume that there are 2 argument registers (\$a0-\$a1), 2 return value registers (\$v0-\$v1), 2 general-purpose callee saved registers (\$s0-\$s1), and 2 general-purpose caller-saved registers (\$t0-\$t1). Assume \$ra is callee-saved. The C code is obviously somewhat silly and unoptimized, but YOU MAY NOT OPTIMIZE IT -- you must translate it as is.

```
1: int f (int num){
2:
    // set up stack frame
3:
     int x = 0; // x must be in $t0
4:
     int y = 1; // y must be in $s0
     if (num == 0) {
5:
6:
       y = y + x;
7: }
8: x = bar(x,y);
9: y = x + y;
10: return (y + 2);
11: // clean up stack frame
12: }
 int bar (int arg) {
    // don't worry about bar(),
    // but bar could call function
 }
```

C lines	Lines of Assembly
1-2	subi \$sp, \$sp, 12 sw \$ra, 0(\$sp) sw \$s0, 4(\$sp)
3-4	li \$t0, 0 li \$s0, 1
5-7	bneqz \$a0, next add \$s0, \$s0, \$t0 next:
8	sw \$t0, 8(\$sp) move \$a0, \$t0 move \$a1, \$s0 jal bar lw \$t0, 8(\$sp)
9	add \$s0, \$v0, \$s0
10	addi \$v0, \$s0, 2
11-12	lw \$s0, 4(\$sp) lw \$ra, 0(\$sp) addi \$sp, \$sp, 12 jr \$ra