## **Empreinte Sociométrique**

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#### **Abstract**

We describe an algorithm that takes contact data to represent the social footprint of a person in a secure and universal way. It allows recognizing an identity without sharing or exposing any personal information to any potential eavesdropper. It could evolve with time and may carry the whole contact history of a person, making it possible to verify if he or she matches with even very old data, all without needing any kind of disclosure between parties. Last update: this technique is pending patent.

## I. Introduction

Ike a tyre or a shoe leaving a distinctive mark on the ground, or a finger on a glass, each one of us leaves a trace the more specific the richer our social interactions are. In particular, the history of our contact data is every day more distinct to someone else's. Even before we move from our parents house, we start leaving personal trails (a first cell phone, a pseudo we use for a game, ...) and, of course, our full civil status (names, date of birth, etc.).

Of course, we wouldn't want to share all these information to everyone. So, ensuring maximum security is obviously mandatory when it comes to manipulating personal data.

We herein describe a way to build such a safe footprint that we call *Empreinte Sociométrique*[1] and that is both totally secure, thanks to the use of strong pseudonymization techniques, and particularly effective, that is it ensures that, even though two different *Empreintes Sociométriques* don't look quite alike, they might be representing the same person, but no external stakeholder may ever know.

Embedded in a QR Code or used as is, it could become an assistant to any identification device or to further data manipulation, such as secure deduplication and anonymous enrichment, or even blind comparison of customer databases, all in full compliance with the latest regulations (GDPR, CCPA, ...).

### II. FORMAL DESCRIPTION

#### 1. General Algorithm

An *Empreinte Sociométrique*<sup>1</sup> takes an input data characterized by its source and its type and operates several operations to form the final elements: its corpus and its signature concatenated into one long string.

**Definition 1** (Source Data). A source data  $\varsigma$  is the actual contact data we want to print in the *Empreinte Sociométrique*, eg. "Cyril".

It is defined in the *words* space:  $\omega$ .

**Definition 2** (Data Type). We define the data type  $\tau$  as a code defining which kind of source data we are dealing with, eg. "firstname".

It is defined in a set of data types  $\mathcal{T}^2$ .

**Definition 3** (Input Data). We define an input data  $d_i$  as a tuple of data type and source data:

$$d_i := [d_i^{\tau}, d_i^{\varsigma}] \text{ with } \left\{ egin{array}{l} d_i^{\tau} \in \mathcal{T}, ext{the data type} \\ d_i^{\varsigma} \in \omega, ext{the source data} \end{array} 
ight.$$

**Definition 4** (Recombined Contact). A recombined contact is a final contact data potentially made out of different input data.

For example, you can create a recombined address4 by concatenating a streetName with a streetNumber.

Let  $v: \omega \to \omega$  be a normalization function that takes an input data and returns its normalized counterpart.

 $<sup>^{*} \</sup>mbox{filed}$  under registration number FR1905778 at INPI on May 29, 2019

<sup>&</sup>lt;sup>1</sup>a literate translation being a *Sociometric Imprint* 

 $<sup>^2</sup>$ see Table 1 for available values in  $\mathcal T$ 

And let  $\rho: \omega^n \to \omega$  be a recombination function that takes several input data to build a missing recombined contact.

Finally, let  $\mathfrak{h}()$  be a cryptographic hashing function<sup>3</sup>,  $\mathfrak{c}(msg, key)$  an encryption function and  $\zeta()$  a compression algorithm.

#### Algorithm 1: Overall construction

```
Input: A vector \mathbf{d} := \{d_1, d_2, \dots, d_n\} of input data, a key K
```

**Output:** The *Empreinte Sociométrique* or an error

1 if  $d = \emptyset$  then

3 initialize the set of normalized data  $\mathcal{D} \leftarrow \emptyset$ ;

4 for 
$$i \leftarrow 0$$
 to  $n$  by 1 do  
5 | if  $d_i^{\tau} \notin \mathcal{T}$  then  
6 | continue;

normalize input data: 
$$d_{Norm} \leftarrow \nu(d_i)$$
;

8 if 
$$d_{Norm} \neq \emptyset$$
 then  
9  $\mathcal{D} \leftarrow d_{Norm}$ ;

10 create the set of recombined contacts  $\mathcal{R}$  from the normalized data:  $\mathcal{R} \leftarrow \rho(\mathcal{D})$ ;

11 initialize the vector of ciphered contacts  $\mathcal{C} \leftarrow \emptyset$ ;

12 for 
$$i \leftarrow 0$$
 to  $|\mathcal{R}|$  by 1 do  
13  $|\mathcal{C} \leftarrow \mathfrak{h}(\mathcal{R}_i);$ 

14 initialize the sets of categorized contacts:

 $\mathcal V$  the variants, and  $\overline{\mathcal V}$  the invariants;

15 for 
$$i \leftarrow 0$$
 to  $|\mathcal{C}|$  by 1 do  
16 | if  $\mathcal{C}_i$  is invariant then  
17 |  $\overline{\mathcal{V}} \leftarrow \mathcal{C}_i$ ;  
18 | else  
19 |  $\mathcal{V} \leftarrow \mathcal{C}_i$ ;

20 build the corpus c using  $\mathcal V$  and  $\overline{\mathcal V}$ ;

21 encrypt it and compress it to build the *Empreinte Sociométrique*:

$$ES \leftarrow \zeta \circ \mathfrak{c}(c,K);$$

22 return ES;

Algorithm 1 describes the general steps to take that leads from a set of input data to its

Empreinte Sociométrique.

In a nutshell, the whole process could be summed up as follows:

- · Data handling;
- Normalization of contact data;
- Recombination to maximize endpoints;
- Ciphering of each endpoint;
- Organization of the corpus;
- Encryption;
- Compression;
- Signature and final packaging.

To ensure maximum security, the whole operation should take place on the data owner side.

#### 2. Step-by-Step Process

#### 2.1 Normalization

Lorem ipsum ...

#### 2.2 Recombination

Lorem ipsum ...

#### 2.3 Ciphering

Lorem ipsum ...

#### 2.4 Corpus

Lorem ipsum ...

#### 2.5 Signature

Lorem ipsum ...

#### 2.6 Output

Lorem ipsum ...

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<sup>&</sup>lt;sup>3</sup>set as a system parameter

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## REFERENCES

[1] Cyril Dever. Système de traitement d'une base de données personnelle par un opérateur extérieur, patent pending FR1905778, 2019.

# Appendix

 Table 1: Input data types

Code	Description	Examples
gender	Title or gender	M, Mr., 1,
firstname	First name or given name	John
middle	Middle initials or other names	J., John
birthname	Birth name	Kennedy
lastname	Last name or married name	Kennedy
suffix	Suffix	Jr.
birthdate	Date of birth	
birthplace	Place of birth	
addresses	List of postal address	(see Table 2)
aliases	List of aliases	(see Table 3)
emails	List of e-mail addresses	(see Table 4)
ids	List of official IDs	(see Table 5)
mobiles	List of mobile phones	(see Table 6)
phones	List of telephone numbers	(see Table 7)
socials	List of social media	(see Table 8)
updated Unix timestamp of collect		1544529071

 Table 2: Address data type

Field Definition		Possible values (or examples)
type Address type		"birth" \rangle "home" \rangle "work"
address2 Additional name		c/o Mme Dupont
address3	Additional address	Apt. 123
address4	Street number and name	1600 Pennsylvania Ave NW
address5	PO Box or locality	BP 987
address6 City and ZIP code		Washington, DC 20500
address7	International destination	U.S.A.
city	City	Washington
country	Country	United States of America
fullAddress	Full address	1600 Pennsylvania Ave NW, Washington, DC 20500
streetName	Street name	Pennsylvania Ave NW
streetNumber	Street number	1600
zip	ZIP code	DC 20500
updated	Time of collect	1544529071

 Table 3: Alias data type

Field	Definition	Possible values (or examples)
type	Alias type	
value	Alias value	John John
updated	Time of collect	1544529071

**Table 4:** *E-mail data type* 

Field	Definition	Possible values (or examples)
type	E-mail type	"business"∨"personal"
value E-mail address		john@john.com
isHash If is already hashed		true V false
engine Hashing algorithm		"blake2"∨"md5"
updated Time of collect		1544529071

 Table 5: ID data type

Field Definition		Possible values (or examples)	
type	ID type		
value	ID value	1234567890abcdef	
isHash	If is already hashed	true $\lor$ false	
engine	Hashing algorithm	"blake2" ∨ "md5"	
updated	Time of collect	1544529071	

 Table 6: Mobile phone data type

Field Definition		Possible values (or examples)
type	Mobile phone type	"business"∨"personal"
value	Mobile phone number	+33 (0) 623 456 789
isHash If is already hashed		true V false
engine Hashing algorithm		"blake2"∨"md5"
format Format		"+dd (d) ddd ddd"
updated Time of collect		1544529071

 Table 7: Telephone data type

Field	Definition	Possible values (or examples)
type Telephone type		"business"∨"personal"
value	Phone number	+33 (0) 123 456 789
isHash If is already hashed		true V false
engine Hashing algorithm		"blake2"∨"md5"
format Format		"+dd (d) ddd ddd"
updated Time of collect		1544529071

 Table 8: Social media data type

Field	Definition	Possible values (or examples)
type	Alias type	"facebook" \rangle "linkedin" \rangle "twitter" \rangle "youtube"
value	Alias value	@john
updated	Time of collect	1544529071