$\boxed{\Psi \vdash e : \sigma}$ $\quad \Psi ::= \emptyset \mid \Psi, x : \sigma$

$$\text{S-T-Var}\ \frac{x : \sigma \in \Psi}{\Psi \vdash x : \sigma} \qquad \text{S-T-Abs}\ \frac{\Psi, x : \sigma_1 \vdash e : \sigma_2}{\Psi \vdash \lambda x.e : \sigma_1 \to \sigma_2} \qquad \text{S-T-App}\ \frac{\Psi \vdash e_1 : \sigma_2 \to \sigma \qquad \Psi \vdash e_2 : \sigma_2}{\Psi \vdash e_1(e_2) : \sigma}$$

$$\text{S-T-Stringin-I}\ \frac{s \in \mathcal{L}\{r\}}{\Psi \vdash \mathsf{rstr}[s] : \mathsf{stringin}[r]} \qquad \text{S-T-Concat}\ \frac{\Psi \vdash e_1 : \mathsf{stringin}[r_1] \qquad \Psi \vdash e_2 : \mathsf{stringin}[r_2]}{\Psi \vdash \mathsf{rconcat}(e_1; e_2) : \mathsf{stringin}[r_1 \cdot r_2]}$$

$$\text{S-T-Case}\ \frac{\Psi \vdash e_1 : \mathsf{stringin}[r] \qquad \Psi \vdash e_2 : \sigma \qquad \Psi, x : \mathsf{stringin}[\mathsf{lhead}(r)], y : \mathsf{stringin}[\mathsf{ltail}(r)] \vdash e_3 : \sigma}{\Psi \vdash \mathsf{rstrcase}(e_1; e_2; x, y.e_3) : \sigma}$$

$$\text{S-T-Replace}\ \frac{\Psi \vdash e_1 : \mathsf{stringin}[r_1] \qquad \Psi \vdash e_2 : \mathsf{stringin}[r_2] \qquad \mathsf{lreplace}(r; r_1; r_2) = r'}{\Psi \vdash \mathsf{rreplace}[r](e_1; e_2) : \mathsf{stringin}[r']}$$

$$\text{S-T-SafeCoerce}\ \frac{\Psi \vdash e : \mathsf{stringin}[r'] \qquad \mathcal{L}\{r'\} \subseteq \mathcal{L}\{r\}}{\Psi \vdash \mathsf{rcoerce}[r](e) : \mathsf{stringin}[r]}$$

$$\text{S-T-Check}\ \frac{\Psi \vdash e_0 : \mathsf{stringin}[r_0] \qquad \Psi, x : \mathsf{stringin}[r] \vdash e_1 : \sigma \qquad \Psi \vdash e_2 : \sigma}{\Psi \vdash \mathsf{rcheck}[r](e_0; x.e_1; e_2) : \sigma}$$

$\boxed{e \Downarrow v}$

$$\text{S-E-Abs}\ \frac{}{\lambda x.e \Downarrow \lambda x.e} \qquad \text{S-E-App}\ \frac{e_1 \Downarrow \lambda x.e_3 \qquad e_2 \Downarrow v_2 \qquad [v_2/x]e_3 \Downarrow v}{e_1(e_2) \Downarrow v} \qquad \text{S-E-RStr}\ \frac{}{\mathsf{rstr}[s] \Downarrow \mathsf{rstr}[s]}$$

$$\text{S-E-Concat}\ \frac{e_1 \Downarrow \mathsf{rstr}[s_1] \qquad e_2 \Downarrow \mathsf{rstr}[s_2]}{\mathsf{rconcat}(e_1; e_2) \Downarrow \mathsf{rstr}[s_1 s_2]} \qquad \text{S-E-Case-}\epsilon\ \frac{e_1 \Downarrow \mathsf{rstr}[\epsilon] \qquad e_2 \Downarrow v_2}{\mathsf{rstrcase}(e_1; e_2; x, y.e_3) \Downarrow v_2}$$

$$\text{S-E-Case-Concat}\ \frac{e_1 \Downarrow \mathsf{rstr}[as] \qquad [\mathsf{rstr}[a], \mathsf{rstr}[s]/x, y]e_3 \Downarrow v_3}{\mathsf{rstrcase}(e_1; e_2; x, y.e_3) \Downarrow v_3}$$

$$\text{S-E-Replace}\ \frac{e_1 \Downarrow \mathsf{rstr}[s_1] \qquad e_2 \Downarrow \mathsf{rstr}[s_2] \qquad \mathsf{subst}(r; s_1; s_2) = s}{\mathsf{rreplace}[r](e_1; e_2) \Downarrow \mathsf{rstr}[s]} \qquad \text{S-E-SafeCoerce}\ \frac{e \Downarrow \mathsf{rstr}[s]}{\mathsf{rcoerce}[r](e) \Downarrow \mathsf{rstr}[s]}$$

$$\text{S-E-Check-Ok}\ \frac{e \Downarrow \mathsf{rstr}[s] \qquad s \in \mathcal{L}\{r\} \qquad [\mathsf{rstr}[s]/x]e_1 \Downarrow v}{\mathsf{rcheck}[r](e; x.e_1; e_2) \Downarrow v}$$

$$\text{S-E-Check-NotOk}\ \frac{e \Downarrow \mathsf{rstr}[s] \qquad s \notin \mathcal{L}\{r\} \qquad e_2 \Downarrow v}{\mathsf{rcheck}[r](e; x.e_1; e_2) \Downarrow v}$$

Figures 3 and 4: Typing Rules and Big step semantics for λ_{RS}

$$\begin{array}{lcll}
\langle \sigma \rangle & ::= & \sigma \to \sigma & \text{source types} \\
& | & \mathsf{stringin}[r] & \\
\\
\langle e \rangle & ::= & x & \text{source terms} \\
& | & \lambda x.e & \\
& | & e(e) & \\
& | & \mathsf{rstr}[s] & s \in \Sigma^* \\
& | & \mathsf{rconcat}(e; e) \mid \mathsf{rstrcase}(e; e; x, y.e) & \\
& | & \mathsf{rreplace}[r](e; e) & \\
& | & \mathsf{rcoerce}[r](e) \mid \mathsf{rcheck}[r](e; x.e; e) & \\
\\
\langle v \rangle & ::= & \lambda x.e \mid \mathsf{rstr}[e] \mid s & \text{source values} \\
\\
\langle \tau \rangle & ::= & \tau \to \tau & \text{target types} \\
& | & \mathsf{string} & \\
& | & \mathsf{regex} & \\
\\
\langle \iota \rangle & ::= & x & \text{target terms} \\
& | & \lambda x.\iota & \\
& | & \iota(\iota) & \\
& | & \mathsf{str}[s] & \\
& | & \mathsf{rx}[r] \mid \mathsf{concat}(\iota; \iota) \mid \mathsf{strcase}(\iota; \iota; x, y.\iota) \mid \mathsf{replace}(\iota; \iota; \iota) \mid \mathsf{check}(\iota; \iota; \iota; \iota) & \\
\\
\langle \dot{v} \rangle & ::= & \lambda x.\iota \mid \mathsf{str}[s] \mid \mathsf{rx}[r] & \text{target values}
\end{array}$$

Figures 2 and 5: Syntax of λ_{RS} and λ_P.

$\boxed{\Theta \vdash \iota : \tau}$ $\quad \Theta ::= \emptyset \mid \Theta, x : \tau$

$$\text{P-T-Var}\ \frac{x : \tau \in \Theta}{\Theta \vdash x : \tau} \qquad \text{P-T-Abs}\ \frac{\Theta, x : \tau_1 \vdash \iota_2 : \tau_2}{\Theta \vdash \lambda x.\iota_2 : \tau_1 \to \tau_2} \qquad \text{P-T-App}\ \frac{\Theta \vdash \iota_1 : \tau_2 \to \tau \qquad \Theta \vdash \iota_2 : \tau_2}{\Theta \vdash \iota_1(\iota_2) : \iota}$$

$$\text{P-T-String}\ \frac{}{\Theta \vdash \mathsf{str}[s] : \mathsf{string}} \qquad \text{P-T-Regex}\ \frac{}{\Theta \vdash \mathsf{rx}[r] : \mathsf{regex}} \qquad \text{P-T-Concat}\ \frac{\Theta \vdash \iota_1 : \mathsf{string} \qquad \Theta \vdash \iota_2 : \mathsf{string}}{\Theta \vdash \mathsf{concat}(\iota_1; \iota_2) : \mathsf{string}}$$

$$\text{P-T-Case}\ \frac{\Theta \vdash \iota_1 : \mathsf{string} \qquad \Theta \vdash \iota_2 : \tau \qquad \Theta, x : \mathsf{string}, y : \mathsf{string} \vdash \iota_3 : \tau}{\Theta \vdash \mathsf{strcase}(\iota_1; \iota_2; x, y.\iota_3) : \tau}$$

$$\text{P-T-Replace}\ \frac{\Theta \vdash \iota_1 : \mathsf{regex} \qquad \Theta \vdash \iota_2 : \mathsf{string} \qquad \Theta \vdash \iota_3 : \mathsf{string}}{\Theta \vdash \mathsf{replace}(\iota_1; \iota_2; \iota_3) : \mathsf{string}}$$

$$\text{P-T-Check}\ \frac{\Theta \vdash \iota_r : \mathsf{regex} \qquad \Theta \vdash \iota_1 : \mathsf{string} \qquad \Theta \vdash \iota_2 : \sigma \qquad \Theta \vdash \iota_3 : \sigma}{\Theta \vdash \mathsf{check}(\iota_r; \iota_1; \iota_2; \iota_3) : \sigma}$$

$\boxed{\iota \Downarrow \dot{v}}$

$$\text{P-E-Abs}\ \frac{}{\lambda x.\iota \Downarrow \lambda x.\iota} \qquad \text{P-E-App}\ \frac{\iota_1 \Downarrow \lambda x.\iota_3 \qquad \iota_2 \Downarrow \dot{v}_2 \qquad [\dot{v}_2/x]\iota_3 \Downarrow \dot{v}_3}{\iota_1(\iota_2) \Downarrow \dot{v}_3} \qquad \text{P-E-Str}\ \frac{}{\mathsf{str}[s] \Downarrow \mathsf{str}[s]}$$

$$\text{P-E-Rx}\ \frac{}{\mathsf{rx}[r] \Downarrow \mathsf{rx}[r]} \qquad \text{P-E-Concat}\ \frac{\iota_1 \Downarrow \mathsf{str}[s_1] \qquad \iota_2 \Downarrow \mathsf{str}[s_2]}{\mathsf{concat}(\iota_1; \iota_2) \Downarrow \mathsf{str}[s_1 s_2]} \qquad \text{P-E-Case-}\epsilon\ \frac{\iota_1 \Downarrow \mathsf{str}[\epsilon] \qquad \iota_2 \Downarrow \dot{v_2}}{\mathsf{strcase}(\iota_1; \iota_2; x, y.\iota_3) \Downarrow \dot{v_2}}$$

$$\text{P-E-Case-Concat}\ \frac{\iota_1 \Downarrow \mathsf{str}[as] \qquad [\mathsf{str}[a], \mathsf{str}[s]/x, y]\iota_3 \Downarrow \dot{v}}{\mathsf{strcase}(\iota_1; \iota_2; x, y.\iota_3) \Downarrow \dot{v}}$$

$$\text{P-E-Replace}\ \frac{\iota_1 \Downarrow \mathsf{rx}[r] \qquad \iota_2 \Downarrow \mathsf{str}[s_2] \qquad \iota_3 \Downarrow \mathsf{str}[s_3] \qquad \mathsf{subst}(r; s_2; s_3) = s}{\mathsf{replace}(\iota_1; \iota_2; \iota_3) \Downarrow \mathsf{str}[s]}$$

$$\text{P-E-Check-OK}\ \frac{\iota_r \Downarrow \mathsf{rx}[r] \qquad \iota \Downarrow \mathsf{str}[s] \qquad s \in \mathcal{L}\{r\} \qquad \iota_1 \Downarrow \dot{v_1}}{\mathsf{check}(\iota_r; \iota; \iota_1; \iota_2) \Downarrow \dot{v_1}}$$

$$\text{P-E-Check-NotOK}\ \frac{\iota_r \Downarrow \mathsf{rx}[r] \qquad \iota \Downarrow \mathsf{str}[s] \qquad s \notin \mathcal{L}\{r\} \qquad \iota_2 \Downarrow \dot{v_2}}{\mathsf{check}(\iota_r; \iota; \iota_1; \iota_2) \Downarrow \dot{v_2}}$$

Figures 6 and 7: Typing rules and big step semantics for λ_P

$\boxed{[\![\sigma]\!] = \tau}$

$$\text{Tr-T-String}\ \frac{}{[\![\mathsf{stringin}[r]]\!] = \mathsf{string}} \qquad \text{Tr-T-Arrow}\ \frac{[\![\sigma_1]\!] = \tau_1 \qquad [\![\sigma_2]\!] = \tau_2}{[\![\sigma_1 \to \sigma_2]\!] = \tau_1 \to \tau_2}$$

$\boxed{[\![\Psi]\!] = \Theta}$

$$\text{Tr-T-Context-Emp}\ \frac{}{[\![\emptyset]\!] = \emptyset} \qquad \text{Tr-T-Context-Ext}\ \frac{[\![\Psi]\!] = \Theta \qquad [\![\sigma]\!] = \tau}{[\![\Psi, x : \sigma]\!] = \Theta, x : \tau}$$

$\boxed{[\![e]\!] = \iota}$

$$\text{Tr-Var}\ \frac{}{[\![x]\!] = x} \qquad \text{Tr-Abs}\ \frac{[\![e]\!] = \iota}{[\![\lambda x.e]\!] = \lambda x.\iota} \qquad \text{Tr-App}\ \frac{[\![e_1]\!] = \iota_1 \qquad [\![e_2]\!] = \iota_2}{[\![e_1(e_2)]\!] = \iota_1(\iota_2)}$$

$$\text{Tr-Case}\ \frac{[\![e_1]\!] = \iota_1 \qquad [\![e_2]\!] = \iota_2 \qquad [\![e_3]\!] = \iota_3}{[\![\mathsf{rstrcase}(e_1; e_2; x, y.e_3)]\!] = \mathsf{strcase}(\iota_1; \iota_2; x, y.\iota_3)} \qquad \text{Tr-String}\ \frac{}{[\![\mathsf{rstr}[s]]\!] = \mathsf{str}[s]}$$

$$\text{Tr-Concat}\ \frac{[\![e_1]\!] = \iota_1 \qquad [\![e_2]\!] = \iota_2}{[\![\mathsf{rconcat}(e_1; e_2)]\!] = \mathsf{concat}(\iota_1; \iota_2)} \qquad \text{Tr-Subst}\ \frac{[\![e_1]\!] = \iota_1 \qquad [\![e_2]\!] = \iota_2}{[\![\mathsf{rreplace}[r](e_1; e_2)]\!] = \mathsf{replace}(\mathsf{rx}[r]; \iota_1; \iota_2)}$$

$$\text{Tr-SafeCoerce}\ \frac{[\![e]\!] = \iota}{[\![\mathsf{rcoerce}[r'](e)]\!] = \iota} \qquad \text{Tr-Check}\ \frac{[\![e]\!] = \iota \qquad [\![e_1]\!] = \iota_1 \qquad [\![e_2]\!] = \iota_2}{[\![\mathsf{rcheck}[r](e; x.e_1; e_2)]\!] = \mathsf{check}(\mathsf{rx}[r]; \iota; (\lambda x.\iota_1)(\iota); \iota_2)}$$

Figure 8: Translation from source terms (e) to target terms (ι).