

# From Timestamping to Blockchain

## Introduction to Computer Security

week 14



暨南大學  
JINAN UNIVERSITY

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  - ▶ copyright
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- ▶ Notary or lawyers testify that a letter or document existed on a certain date
- ▶ What about in the digital world?
- ▶ Document can be copied and modified endlessly
- ▶ It is trivial to change the date stamp on a computer file

# Timestamping requirements

- ▶ The data must be timestamped regardless of the physical medium on which it is stored
- ▶ It must be impossible to change data without the change being apparent
- ▶ It must be impossible to change timestamp of a document with a different date and time from the present one

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- ▶ Problems
  - ▶ Privacy: Alice has to send a copy of the document to Trent
  - ▶ Trent requires a huge database to store the copies, and the bandwidth requirement is high
  - ▶ The database may be attacked by an electromagnetic bomb
  - ▶ Trent may not be that honest, or he may collude with Alice?

# An improved arbitrated solution (credit. B. Schneider's book — Applied Cryptography)

- ▶  $A \rightarrow T : \text{hash}(M)$
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- ▶ Trent may still be dishonest or collude with Alice

# Linking protocol

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- ▶  $A \rightarrow T : H_n$
- ▶  $T \rightarrow A : T_n = [n, \text{Alice}, H_n, t_n, \text{Bob}, H_{n-1}, T_{n-1}, L_n]_{sk(T)}$   
where  $L_n = \text{hash}(\text{Bob}, H_{n-1}, T_{n-1}, L_{n-1})$
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where  $L_n = \text{hash}(\text{Bob}, H_{n-1}, T_{n-1}, L_{n-1})$
- ▶ After Trent certifies someone after Alice, he sends Alice  $T_{n+1}$ .
- ▶ If Someone challenges Alice, she may contact the originators of the previous and following documents.

# Advantages of the linking protocol

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- ▶ The same logic applies, and Trent needs to collude with almost every people linked on the chain (everyone after Alice, and someone before Alice)
- ▶ Note that  $L_n = \text{hash}(\text{Bob}, H_{n-1}, T_{n-1}, L_{n-1})$  will be different if either  $t_{n-1}$  or  $L_{n-1}$  is modified.

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- ▶ Another way is to publish the ongoing chain on a public web site. (So that the entire chain is verifiable by everyone)

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  - ▶ tamper-resistant (cryptographic hash function)
  - ▶ availability (millions of copies on the Internet)
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- ▶ Let's start with the first successful Blockchain application as a crypto-currency

# A centralized crypto-currency protocol

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# A centralized crypto-currency protocol

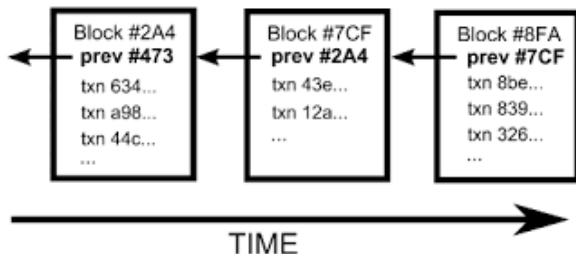
- ▶ If Jinan University issues a crypto-currency JN-coin for students
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  - ▶ Students buy credit from the university
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- ▶ If we crypto-lized the entire procedure
  - ▶ The University has its public key and private key pair
  - ▶ Let  $pk(A)$ ,  $pk(B)$  be Alice and Bob's public key
  - ▶  $sk(A)$ ,  $sk(B)$  be Alice and Bob's private key
  - ▶  $TX_i = [\text{issue to } pk(A) \text{ } 10\text{JN}]_{sk(U)}$
  - ▶  $TX_j = [pk(A) \text{ pays } pk(B) \text{ } 5\text{JN}]_{sk(A)}$  if Alice pays Bob 5JN
  - ▶ Every 10 seconds the university publish a new block including all transactions that have happened during the last 10 seconds

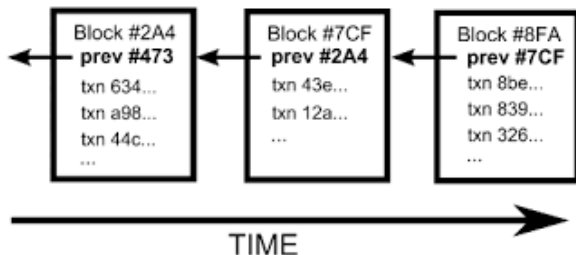


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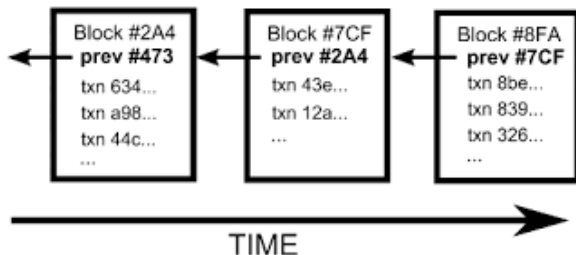
- ▶ It is viable to check which person (represented by his public key) owns how many coins at each moment of time
- ▶ People are anonymized (pseudonymized)

# A centralized crypto-currency protocol



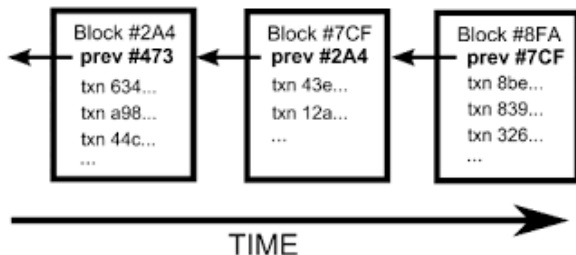
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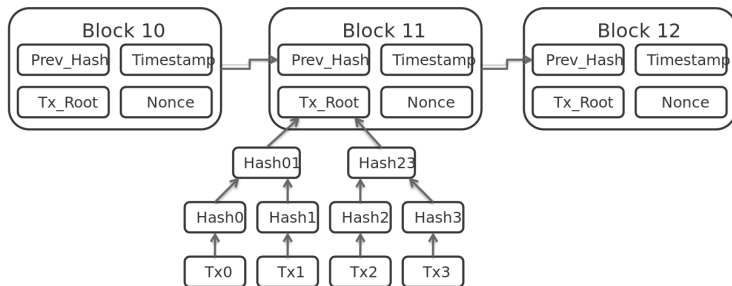
- ▶ Trusting the university is essential
- ▶ What if the university JN-coin server is down
- ▶ What if the university system admin is corrupted

# A centralized crypto-currency protocol



- ▶ Trusting the university is essential
- ▶ What if the university JN-coin server is down
- ▶ What if the university system admin is corrupted
- ▶ Actually we have the same problem with Alipay, wechat-pay (and every centralized system)

# Block organization



- ▶ Prev\_Hash of Block  $n$  is  $hash(Block_{n-1})$
- ▶ Timestamp records the time when the block is generated
- ▶ Merkle tree inside each block bundle transactions to Tx\_Root

# Bitcoin as a centralized crypto-currency

- ▶ Bitcoin: A Peer-to-Peer Electronic Cash System (Satoshi Nakamoto, 2009)
- ▶ A distributed ledger made of Blocks
- ▶ Mining and Proof of Work (PoW)
- ▶ Distributed consensus (to avoid double spending)

- ▶ Hashcash — a denial of service counter-measure (A. Back, 2002)
- ▶ To introduce a cost to spammers by setting up rules for a client's mailbox
  - ▶ Given an email  $M$ , it needs to be of the form  $[M, I, nonce]$
  - ▶  $I$  is the receiver's Id (email address)
  - ▶ needs to satisfy that  $hash([M, I, nonce]) = 0^k \{0, 1\}^{128-k}$  (let's say, md5 hash)
- ▶ For each recipient, the spammer needs to invest a tiny amount of computation time
- ▶ Based on the principle that the hash function is secure (pre-image resistant, and well distributed)

# Proof of Work (PoW) Mining in Bitcoin

- ▶ Miners get a bonus by checking consistency of the existing chain and adding a new block to the chain (1M to 2M in size)
- ▶ The reward is 50 BTC in block #1 and halves every 210,000 blocks (against inflation)



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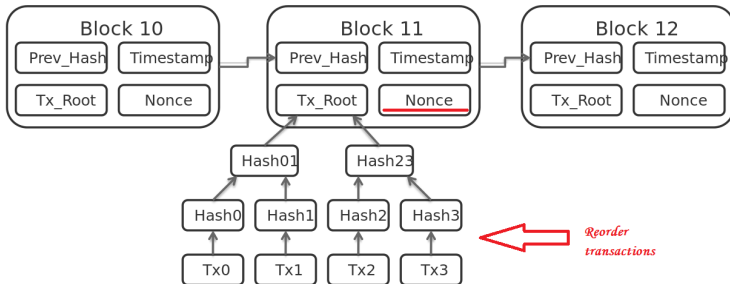
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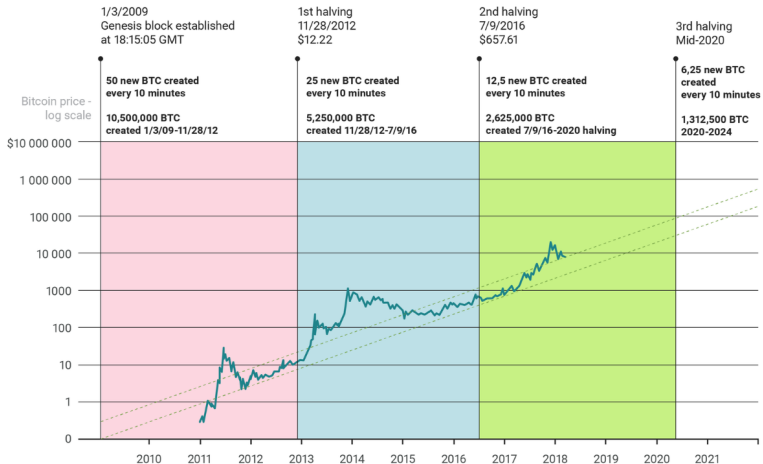
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# Proof of Work (PoW) Mining in Bitcoin

## Bitcoin price history with reward halving days marked



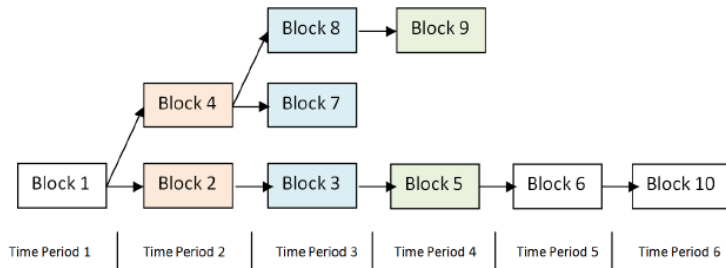
Note: The dotted lines show an average price appreciation of 200% per year (or 3x year/year).

# Proof of Work (PoW) Mining— Miner Competition

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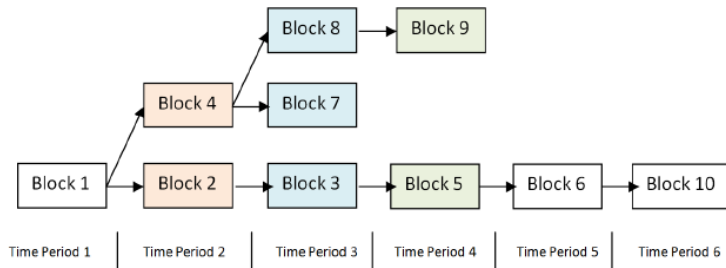
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- ▶ Security guaranteed by the assumption that the majority of miners are honest

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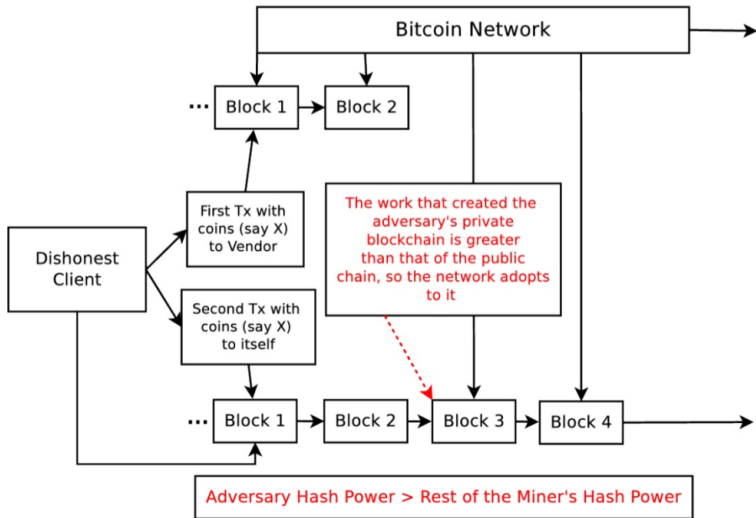
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- ▶ Then Alice tries to generate a new block that retains her coin, and attaches it to the block immediately before the block that contains  $TX_A$ .
- ▶ Alice has a good chance to succeed if her hashing power is more than the rest of the Bitcoin network (why?)

# Double spending attack in Bitcoin

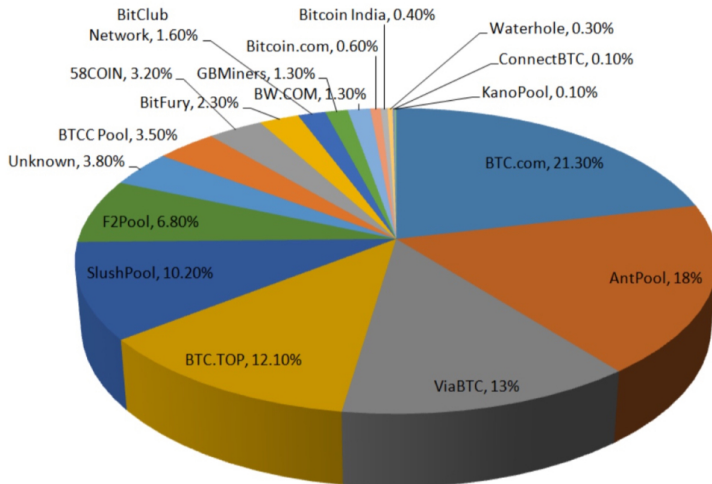


# Other attacks in Bitcoin

- ▶ **Finney attack:** a form of double spending attack in which a dishonest client pre-mines a block (or a sequence of blocks) containing a transaction paying the coins to an address under their control.
- ▶ **Goldfinger attack:** if a miner (mining pool) can get to 51% of the overall hashrate, then all bets are off. This is known as a 51% attack (it enables you to double spend at will for example). Now the motivation is not to profit directly through Bitcoin, but instead to bring down the currency or network.
- ▶ **Block discarding / selfish mining:** In block discarding, a dishonest miner (or colluding set of miners) working in a pool withholds a block once found. They keep working on the private chain, and publish their mined blocks immediately before honest forks get back to the same length.
- ▶ **Bribery attacks, Wallet theft, DDoS to miners**

# Mining pools in the Bitcoin network

- ▶ The market share of hashrate for mining pools as of December 2017.



# Other cryptocurrency systems than Bitcoin





















- ▶ Litecoin, Namecoin . . .
- ▶ Ethereum (Crypto-currency plus smart contract)
- ▶ IOTA (Based on a Directed Acyclic Graph)

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- ▶ Ethereum (Crypto-currency plus smart contract)
- ▶ IOTA (Based on a Directed Acyclic Graph)
- ▶ Some interesting places to check (price, exchange rates, market cap etc.)
  - ▶ <https://coinmarketcap.com/>
  - ▶ <https://www.coinbase.com/>



## Top 100 Cryptocurrencies by Market Capitalization

Cryptocurrencies ▾		Exchanges ▾	Watchlist		USD ▾		Next 100 →	View All
#	Name	Market Cap	Price	Volume (24h)	Circulating Supply	Change (24h)	Price Graph (7d)	
1	 Bitcoin	\$60,762,890,529	\$3,488.73	\$4,900,331,896	17,416,912 BTC	-3.26%		...
2	 XRP	\$12,639,509,054	\$0.308831	\$401,981,861	40,926,963,305 XRP *	-0.63%		...
3	 Ethereum	\$9,483,071,106	\$91.42	\$1,655,144,473	103,725,395 ETH	-2.79%		...
4	 Stellar	\$2,262,128,286	\$0.118032	\$98,315,870	19,165,371,791 XLM *	-3.93%		...
5	 Tether	\$1,881,827,942	\$1.01	\$3,145,400,537	1,856,421,736 USDT *	-0.32%		...
6	 Bitcoin Cash	\$1,855,246,091	\$105.99	\$68,797,089	17,503,913 BCH	-1.69%		...
7	 EOS	\$1,741,449,091	\$1.92	\$857,349,601	906,245,118 EOS *	-3.26%		...
8	 Bitcoin SV	\$1,621,506,908	\$92.64	\$57,653,631	17,503,611 BSV	-7.12%		...
9	 Litecoin	\$1,477,317,823	\$24.82	\$417,556,570	59,519,217 LTC	-2.79%		...
10	 TRON	\$885,486,196	\$0.013367	\$53,903,239	66,246,286,486 TRX *	-1.04%		...

# Summary of the day

- ▶ Timestamping services (proof of time and integrity)
- ▶ Arbitration and Linking
- ▶ Centralized control vs de-centralized control
- ▶ Block as a techniques used for cryptocurrencies
  - ▶ Block + Hash  $\Rightarrow$  Integrity and tamper-resistance
  - ▶ PoW Mining
  - ▶ (Probabilistic) distributed consensus (double spending attack)
  - ▶ Altcoins