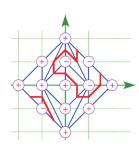
Patchworking on Real Algebraic Curves

A Combinatorial Analogue of Real Algebraic Varieties

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CONTENT

- 1. Logarithm Paper of Polynomials
- 2. Maslov Dequantisation
- 3. Patchworking on Plane Algebraic Curves
- 4. Hilbert's 16th Problem

Our main objective is to present a combinatorial analogue of a special kind of real algebraic curve called patchworking [6], which provides the topology of these polynomials with real varieties.

Here we mainly focus on unary and binary polynomials.

OUTLINE

- Dequantisation of real algebraic geometry on logarithm paper.
- Introduce the method of patchworking, which combines the real algebraic geometry with geometric combinatorics.
- Introduce the first part of Hilbert's 16th problem, and provide a counterexamples to the Ragsdale Conjecture.

Logarithm Paper of Polynomials

VISUALISING ZEROS OF POLYNOMIAL a(x, y) ON $\log \log$ PAPER

In order to visualise an real algebraic curve, we mainly focus on its asymptotic behaviour. The log log paper is utilised for linearisation of a certain real algebraic curve on quadrants. The log log paper is the diffeomorphism given by

an open quadrant
$$\to \mathbb{R}^2$$
, $(x,y) \mapsto (u,v) := (\log |x|, \log |y|)$. (1)

E.g., the log log paper of a real cubic curve is presented in Figure 1.

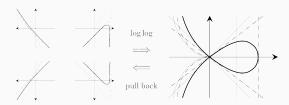


Figure 1: The log log paper of $a(x, y) = 8x^3 - x^2 + 4y^2$.

VISUALISING POLYNOMIAL y = p(x) ON $\log \log \overline{PAPER}$

For simplicity, we settle for polynomials in forms of a(x, y) = y - p(x)with $y = p(x) = \sum_{k=0}^{n} a_k x^k$ as the set of zeros. The behaviour of y = p(x) is determined by each monomials $\{a_k x^k\}_{k=0}^n$, that is, the preimage of asymptomatic lines in the log log paper.

HOW $\log \log PAPER$ CONTROLS y = p(x)

The log log paper manifests the piece-wise asymptomatic behaviour, especially in the vicinity of axes and at infinity, as Figure 2 shows.

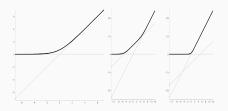


Figure 2: The log log paper of $y = 1 + x^2$ and $y = 1 + e^{\pm 5}x + x^2$ in Q_{++} .

Logarithm Paper of Polynomials

CONTROL THE SET OF ZEROS WITH MAXIMUM OF MONOMIALS

For a unary polynomial $p(x) = \sum_{k=0}^{n} a_k x^k$, we consider the case that $a_k = e^{b_k} > 0$ for each $k = 0, 1, \dots, n$ and the truncation of its plot in $(x, p(x)) \cap Q_{++}$. Let Γ_p denotes the graph of

$$L_p(u) = \ln\left(\sum_{k=0}^n e^{ku+b_k}\right). \tag{2}$$

We define the function of maximal monomial

$$M_p(u) = \max_{k} \{ku + b_k\}_{k=0}^n, \tag{3}$$

which is piecewise linear. It is clear that

$$L_p(u) \le M_p(u) = L_p(u) + R(u),$$
 (4)

for some bounded and $+\infty$ -vanishing function R(u).

Maslov Dequantisation

RESCALING MAP OF y = p(x)

When we consider the rescalings $(u, v) \mapsto (Cu, Cv)$ for large C > 0, it follows that the graph $y = ax^k \Leftrightarrow v = ku + b$ becomes

$$v = ku + b \cdot C \Leftrightarrow y = a^C x^k. \tag{5}$$

Let $h=rac{1}{C}>0$, $p_h(x)=\sum_{k=0}^n a_k^{1/h}x^k$, and Γ_f^h be the graph Γ_p under the change of coordinates $(u, v) \mapsto (hu, hv)$.

It is clear that $\Gamma^h_{M(p_h)}$ does not depend on h. Thus $\lim_{h\to 0} \Gamma^h_{p_h} = \Gamma_{M(p_h)}$ in the C^0 sense. E.g., Figure 3.

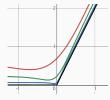


Figure 3: The limit $v = h \ln(e^{2u/h} - 2e^{u/h} + 2)$ as $h \to 0$ in C^0 sense

THE IDEA OF MACLOV DEQUANTISATION

Consider a bunch of homomorphisms¹ with index $h \in [0, +\infty)$

$$D_h: \mathbb{R}[x_1, \dots, x_n] \to \{\log \log \text{ papers}\}, p \mapsto \Gamma_{p_h}^h, \quad x_i \mapsto h \ln x_i.$$
 (6)

Then
$$D_h(a+b) = D_h(a) \oplus_h D_h(b)$$
, $D_h(a \cdot b) = D_h(a) \odot_h D_h(b)$.

From the perspective of Maslov [5], we shall construct a set of semirings $\{(S_h(=\mathbb{R}), \oplus_h, \odot_h)\}_{h\in[0,\infty)}$. Then \odot_h is + in normal sense, and

$$a \oplus_h b = \begin{cases} h \ln(e^{a/h} + e^{b/h}) & h > 0, \\ \max\{a, b\} & h = 0. \end{cases}$$
 (7)

Established by Maslov, the theory idempotent analysis [3] is a systematic illustration of such kind of dequantisation.

¹This is indeed an isomorphism when $h \neq 0$.

More about MacLov Dequantisation

Intuitively, such \odot_h can be extended to $\frac{1}{h}$ -norm for each $h \in [0,1]$.

If one focuses on the limit at h=0 (in sense of C^0 continuity), then we have the following correspondence

$(\mathscr{F},+,\cdot)$	$(\mathscr{F}_0,\oplus_0,\odot_0)$
$p(x) = \sum_{k} a_k x^k$	$M_p(u) = \max_k \{ku + \ln a_k\}$
Integral	essential supremum
$\int_X f$	$\operatorname{esssup}_X\{f(x)\}$
Fourier transform	Legendre transform
$\tilde{f}(\xi) = \int e^{ix\xi} f(x) dx$	$\tilde{f}(\xi) = \operatorname{esssup}\{x \cdot \xi - f(x)\}$
Linear problem	Convex optimisation problem

If fact, the \odot_h can be regarded as the degraded forms of +.

Patchworking on Plane Algebraic Curves

Similarly, the monomial ax^ky^l with a>0 under the $\log\log$ paper is a plane $w=ku+lv+\ln a$. The graph Γ_p lies in the neighbourhood of such surfaces. We can also define $\Gamma_{M(p)}$, Γ_f^h , the continuous deformation $\Gamma_{p_h}^h$ for $h\in(0,1]$, etc..

The set of zeros $V_{Q_{++}}(p)$ is also a continuous deformation of $\Gamma_{p_+}\cap\Gamma_{p_-}$ projected on $S_0^2(=\mathbb{R}^2)$. In order to analysis the topology of p(x,y) with digree m, one just need to consider the supremum of the set of planes $U_+:=\cup_{a^{kl}>0}\{w=ku+lv+\ln a^{kl}\}$ that derives from monomials in $p_+(x,y)$, i.e., $\{a^{kl}x^ky^l\}_{a^{kl}>0}$, as well as the $U_-:=\cup_{a^{kl}<0}\{w=ku+lv+\ln(-a^{kl})\}$ from $p_-(x,y)$. The intersection of $\Gamma_{p_+}:=\sup U_+$ and $\Gamma_{p_-}:=\sup U_-$ is what we desired.

VISUALISE THE PLANE ALGEBRAIC CURVE

We take $a(x,y)=x^3+2y+x+xy-(3x^2+y^2+1)=0$ as an example. The topology of $V_{Q_{++}}(a)$ is shown in Figure 4.

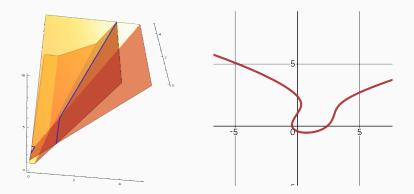


Figure 4: The $\log \log$ visualisation versus the real curve of $V_{Q_{++}}(a)$

CHART OF A POLYNOMIAL

The chart of a polynomial is a set of simplicia with codimension 1 on the components of a Newton polygon² $\Delta(a)$. The chart on a given quadrant Q is homeomorphic (or homotopic for degenerated cases) to $V_O(p)$.

Especially for degenerated cases, the Newton polygon of a non-singular quasi-homogeneous polynomial³ is a line, e.g. Figure 5.

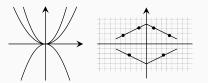


Figure 5: The Newton polygon and charts of a(x, y) = $(x^2 - y)(x^2 + y)(2x^2 - y)y$

The log log paper is a set of parallel lines perpendicular to $\Delta(a)$.

²Newton polygon $\Delta(a)$ of a = a(x, y) is the convex hull of $\{(l, k) : a_{k,l} \neq 0\}$ for $a(x,y) = \sum_{0 \le k+l \le m} a_{k,l} x^k y^l.$ ³which is in forms of $a(x,y) = \sum_{i=0}^k \alpha_i x^{pi+p_0} y^{q(k-i)+q_0}$.

CHART OF A NON-QUASI-HOMOGENEOUS POLYNOMIAL

We first consider the charts of non-quasi-homogeneous trinomial $8x^{3} - x^{2} + 4y^{2}$ shown in **Figure 6**.

For simplicity, let the charts be the mid-segments. Here the signs are deliberately marked to distinguish the connected components.

One can verify that such construction fits all trinomials.

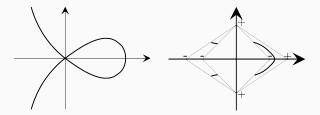


Figure 6: The charts of $8x^3 - x^2 + 4y^2$ on each quadrant.

CHART OF A NON-QUASI-HOMOGENEOUS POLYNOMIAL

Consider a set of real trinomials $\{a_k(x,y)\}_{k=1}^s$, provided that

$$[\operatorname{Int}\Delta(a_i)] \cap [\operatorname{Int}\Delta(a_j)] = 0$$
, whence $i \neq j$. (8)

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Patchworking on Plane Algebraic Curves

We suppose that the coefficients of $x^k y^l$ in $\{a_r : (k, l) \in \Delta(a_r)\}$ are the same, so that $a = \bigcup_{k=1}^{s} a_k$ is well-defined.

Let $\nu:\Delta(a)\to\mathbb{R}$ be such that

- \cdot ν are piecewise linear on $\Delta(a)$, while not linear on the union of any pairs of different triangles.
- $\cdot \nu : \Delta(a) \cap \mathbb{Z}^2 \to \mathbb{Z}.$

PREREQUISITES OF PATCHWORKING THEOREM

- Patchworking for trinomials,
- The bonding lemma for set of disjoint trinomials.

Construct Patchworking Segments in Q_{++}

The initial data for patchworking in $\overline{Q_{++}}$ includes several steps as follows

Patchworking on Plane Algebraic Curves

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- m, the degree of the polynomial,
- Δ , the Newton polygon⁴.
- $\{\sigma_{k,l} \in \{\pm 1\} : (k,l) \in \overline{\Delta}\}.$
- τ , the convex triangulation of Δ with vertices in integer coordinates.
- $\cdot \nu : \Delta \to \mathbb{R}_+$, a piecewise linear function which is not linear on the union of any two triangles.

For instance, the patchworking of a conic with given m=2, triangulation τ , and $\{\sigma_{k,l}\}$ in $\overline{Q_{++}}$ is shown in the first plot of **Figure**

 $[\]overline{\mathbf{7}}_{\text{4}}$ The convex hull of $\{(0,0),(0,m),(m,0)\}$ is commonly utilised.

PROCEDURE FOR CONSTRUCTING PATCHWORKING SEGMENTS

The basic procedure for constructing the patchworking segments (L,Δ) includes:

- 1. Construct $(L, \overline{Q_{++}})$ with initial values.
- 2. Reflect Δ as well as the triangulation τ with x and y axes.
- 3. Extend the sign function onto $\Delta_*{}^5$ in the light of

$$\sigma_{\varepsilon i,\delta j} = \begin{cases} \sigma_{i,j} & \|(i,j) - (\varepsilon i,\delta j)\|_1 = 4k, \\ -\sigma_{i,j} & \|(i,j) - (\varepsilon i,\delta j)\|_1 = 4k + 2. \end{cases}$$
(9)

- 4. Whence a triangle in the triangulation τ has vertices of different signs, draw mid-segment(s) separating pluses from minuses.
- 5. Let L be the union of all segments constructed in the previous step. Then (Δ_*, L) is called the result of combinatorial patchworking (shown in Figure 7).

 $^{^5}$ We denote $\Delta_* := \cup_{arepsilon, \delta \in \{\pm 1\}} \Delta_{arepsilon, \delta}$

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The corresponding one-parameter family of polynomials on affine plane $\mathbb{A}^2_{\mathbb{R}}$ is shown as

$$b_t(x,y) = \sum_{(k,l)\in\mathcal{V}(\tau)} \sigma_{k,l} \cdot t^{\nu(k,l)} \cdot x^k y^l. \tag{10}$$

Q: Does $b_t(x,y)$ has the same topology as the segments?

This is true when $t \to 0^+$.

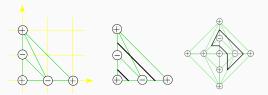


Figure 7: Patchwork a conic in $\mathbb{A}^2_{\mathbb{R}}$

Main Theorem: Patchworking Theorem

Patchworking theorem There exists $t_0 > 0$ such that for any $t \in (0, t_0]$ the equation $b_t(x,y) = 0$ defines in the plane c_t a curve such that the pair $(\mathbb{A}^2_{\mathbb{R}}, c_t)$ is homeomorphic to the pair (Δ_*, L) .

Patchworking on Plane Algebraic Curves

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The equation $B_t(x_0, x_1, x_2)$ defines in the real projective plane $\mathbb{P}^2_{\mathbb{R}}$ a curve C_t such that $(\mathbb{P}^2_{\mathbb{P}}, C_t)$ is homeomorphic to $(\overline{\Delta}_*, L)$.

REMARK:

- $\bullet \ B_t(x_0,x_1,x_2) = \sum_{k,l} \sigma_{k,l} \cdot t^{\nu(k,l)} \cdot x_0^{m-k-l} x_1^k x_2^l \text{ on } \mathbb{P}^2_{\mathbb{R}} \text{ is the}$ $(k,l)\in\mathcal{V}(\tau)$ homeogeneous form of $b_t(x,y) = \sum_{k,l} \sigma_{k,l} \cdot t^{\nu(k,l)} \cdot x^k y^l$ on $\mathbb{A}^2_{\mathbb{R}}$. $(k,l)\in\mathcal{V}(\tau)$
- c_t (or C_t) is the set of zeros $V_{\mathbb{A}^2_n} b_t$ (or $V_{\mathbb{P}^2_n} B_t$).

Hilbert's 16th Problem

INTRODUCTION TO HILBERT'S 16TH PROBLEM

Harnack (1876) prove that the number of components of a real projective curve of degree m is at most $\binom{m-1}{2} + 1$. We call a curve M-curve whenever the equality holds.

The first part of Hilbert's 16th problem is to describe which real schemes of ovals can be realised by a real algebraic curve of degree m. The full classification M-curve of degree 6 is shown in Figure 8.



Figure 8: Topology of M-curve of degree 6.

HARNACK'S CONSTRUCTION

Harnack's construction is to

- 1. Perturbs the union of a line and a circle.
- 2. Combined the result curve with the previous line for a new perturbation.
- 3. And so forth...

A Harnack curve of degree 2k has

- $p=3\binom{k}{2}+1$ outer ovals,
- $n = \binom{k-1}{2}$ inner ovals.

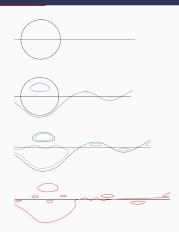


Figure 9: The idea of Harnack's construction

In the light of Harnack's construction, the sign distribution $\sigma_{k,l}$ in $\overline{Q_{++}}$ takes minus only on even grids, i.e., Figure 10.

We utilise the symbol in [4] to denote such scheme by $\langle 9 \coprod 1\langle 1 \rangle \rangle$.

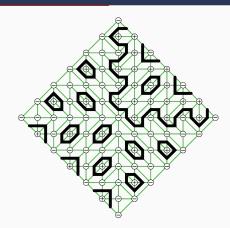


Figure 10: Harnack's sign distribution for m=6.

HILBERT'S CONSTRUCTION

Hilbert's construction begin with the union of two ellipses...

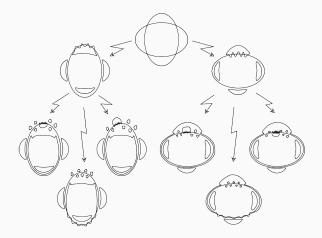


Figure 11: The idea of Hilbert's construction.

GUDKOV'S EUREKA

The preexisting schemes are coloured black in **Figure 12**. In order to make it more symmetric, Gudkov (1971) found the final answer.

Figure 12: Schemes of curves with degree 6.

Hilbert's 16th Problem

Ragsdale (1906) observed that

- for any Harnack's M-curve of even degree m=2k, we have $p = 3\binom{k}{2} + 1$ and $n = \binom{k-1}{2}$.
- for any Hilbert's M-curve of even degree m=2k, we have $\binom{k-1}{2} + 1 .$

It motivates the following conjecture by Ragsdale, that is,

RAGSDALE'S CONJECTURE

For any curve of even degree m=2k, one has $p \leq 3\binom{k}{2}+1$, and $n \leq 3\binom{k}{2}$.

THE COUNTEREXAMPLE OF RAGSDALE'S CONJUNCTURE

It is proved by Oleg [2] that

 there exists a nonsingular real algebraic plane curve of degree 2k with

$$p = \frac{3k(k-1)}{2} + 1 + \left[\frac{(k-3)^2 + 4}{8}\right]. \tag{11}$$

• there exists a nonsingular real algebraic plane curve of degree 2k with

$$n = \frac{3k(k-1)}{2} + \left[\frac{(k-3)^2 + 4}{8} \right]. \tag{12}$$

Tessellation of $\langle 1 \langle 2 \rangle \rangle$ Hexagon

Consider the patchworking of the hexagon in Figure 13 and the tessellation in $\Delta(x^m + y^m + 1)$ in Figure 14.

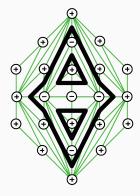


Figure 13: $\langle 1\langle 2 \rangle \rangle$ hexagon.

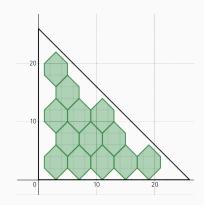


Figure 14: The hexagon tessellation.

For the partition from tessellation of $\langle 1\langle 2 \rangle \rangle$ hexagons, we apply the Harnack's distribution to the uncovered vertices in $\overline{Q_{++}}$. We denote the number of hexagons in $\Delta(x^{2k} + y^{2k} + 1)$ by

$$\alpha := \left\lfloor \frac{(k-3)^2 + 4}{8} \right\rfloor. \tag{13}$$

Then the real scheme outcomes

$$\langle (3\binom{k}{2} - \alpha) \coprod 1 \langle (\binom{k-1}{2} - 4\alpha) \coprod \alpha \langle 2 \rangle \rangle \rangle. \tag{14}$$

It yields that the number of even ovals is

$$p = 3\binom{k}{2} + 1 + \alpha = \frac{3k(k-1)}{2} + 1 + \left\lfloor \frac{(k-3)^2 + 4}{8} \right\rfloor. \tag{15}$$

As shown in Figure 15, we adjust the the signs and triangulation in pleural region of the previous counterexample.

- ▶ hexa-tessellation
- ▶ pleural region
- ▶ sign distribution
- ▶ triangulation

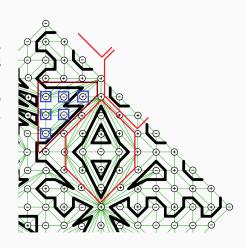


Figure 15: Adjustment in Q_{++} .

Counterexample of Ragsdale's conjecture on n

The patchworking in the previous page yields the real scheme

$$\langle 1\langle 3\binom{k}{2} - \alpha \rangle \coprod \alpha \langle 2 \rangle \coprod (\binom{k-1}{2} - 4\alpha - 1) \rangle.$$
 (16)

It yields that the number of even ovals is

$$p = 3\binom{k}{2} + \alpha = \frac{3k(k-1)}{2} + \left\lfloor \frac{(k-3)^2 + 4}{8} \right\rfloor. \tag{17}$$

IMPROVED UPPER BOUNDS FOR COUNTEREXAMPLES

The improved upper bound [1] of p is presented below

$$p = 3\binom{k}{2} + 1 + \left\lfloor \frac{k^2 - 7k + 16}{6} \right\rfloor.$$

A COUNTEREXAMPLE FOR n OF DEGREE 10

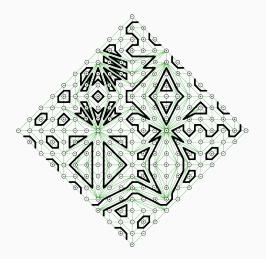


Figure 16: Scheme $\langle 1\langle 29\rangle \coprod 1\langle 2\rangle \coprod 1\rangle$

Hilbert's 16th Problem

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