

Solution: (a) Let V and E denote the set of vertices and edges in graph [1]. Define a variable x_e for each edge $e \in E$. Then, the maximum-cardinality bipartite matching can be formulated as:

$$\begin{aligned} &\text{Maximize } \sum_{e \in E} x_e \\ &\text{Subject to } \sum_{e \sim v} x_e \leq 1 \text{ for all } v \in V \\ &x_e \in \{0, 1\} \text{ for all } e \in E \end{aligned}$$

where $e \sim v$ denotes the set of edges connected to vertex v . Note that the Linear Program can be written as

$$\begin{aligned} &\text{Maximize } cx \\ &\text{Subject to } Ax \leq b \\ &x \geq 0 \end{aligned}$$

with c a $1 \times |E|$, where all elements are equal to 1 and $|E|$ denotes the total number of edges. Moreover, b is a $|V| \times 1$ vector with all elements equal to 1, where $|V|$ denotes the total number of vertices. The matrix A has a 1 on entry (i, j) if vertex i is connected to edge j and 0 otherwise.

(b) The dual problem is given by

$$\begin{aligned} &\text{Minimize } \sum_{v \in V} y_v \\ &\text{Subject to } \sum_{v \sim e} y_v \geq 1 \text{ for all } e \in E \\ &y_v \in \{0, 1\} \text{ for all } v \in V \end{aligned}$$

where $v \sim e$ denotes the vertices v connected to edge e . The dual variable y_i is equal to 1 if vertex i is matched and 0 otherwise. The objective function minimizes the sum of matched vertices subject to the given constraints. ■

1 References

[1] <http://www.imsc.res.in/meena/matching/lecture5.pdf>