

Don't be a bean COUNTER

10/2: Discussion 1

- `export` is used to set an environmental variable
 - Can be used to prepend to a `PATH`
- `vi`
 - `~/.profile` is the doc for your default customization in a SEASnet server
 - `i` to insert text, `Esc` to escape from insert
 - `:wq` to quit out of `vi`
- `mv` - rename a file
- `cp` - copy a file
- Moving in Emacs
 - `C-b`: back (`M-b` for one word)
 - `C-f`: forward (`M-f` for one word)
 - `C-p`: previous line
 - `C-n`: next line
 - `C-e`: end of line
 - `C-a`: beginning of line
 - `C-v`: scroll to next screen
 - `M-v`: scroll to previous screen
- Deleting in Emacs
 - `C-k`: Deletes a line, forward
- Formatting in Emacs
 - `C-o`: inserts newline
- Transferring files from server to local
 - `scp` charlesx@lnxsrv06.seas.ucla.edu: /u/cs/ugrad/charlesx/file_name C:\Users\chuck\Downloads
- I/O
 - `d`
- Shell Scripting
 - A shell is a UI that allows access to an OS's services
 - Common Unix Shells: Bash, zsh, sh, csh

10/6: The Command Line

- C++: objects working with operations
 - C++ programs operate in a process
 - These processes are an object in the OS, which contain the program and data necessary for the program
 - Command line commands work the same way: `ls`, `mv`, etc.
 - Processes may run the same program with different data

- The OS also contains files separate from running processes
 - Files contain data with no program or anything that's actually running
- High-level view: processes can be thought of as operations and files can be thought of as objects
 - Processes and files are really both objects
 - Big difference is power draw → shutting off power kills processes while file data remains
 - Processes are not persistent while files are persistent
 - Persistence requires tradeoff of efficiency in modern technology
 - Changing data would require writing to disk, slow read/write, etc.
- Applications must ideally be as persistent as possible, as efficient as possible, and understandable
- Ctrl is a mask of the least significant 5 bits of an ASCII char
- C-h k [key] tells us what key does
 - Press Enter on function name to see source code
- C-x o - switch to other buffer
- C-x 1 - look at just this buffer
- C-x 2 - split buffer in half
- C-x 3 - split buffer in half vertically
- C-x 5 - create a new window
- C-b NAME - switch this frame to buffer NAME
- C-x C-b - put buffer listing into a new buffer and display as other buffer
- C-x d RET - create a buffer containing a list of what's in current directory
 - Every directory has 2 entries
 - . - current directory
 - Pointer to itself
 - .. - parent directory
 - Pointer to parent
 - Typing g refreshes it in the file system
 - Emacs automatically saves backup files when writing
- C-g - stop the current command
- C-x C-f FILE RET - start editing FILE
- Emacs is a modal editor - actions upon typing depends on the modes Emacs is in
- M-q - reformat the current paragraph
- 2 types of files:
 - Regular files are sequences of bytes
 - Directories are mappings from file names to files
- M-x shell RET - opens the shell in Emacs
- ls -l - dumps files in directory with meta-information
 - 1st char is a file type: '-' for reg. file, 'd' for directory
 - Next 3 chars are owner permissions - 'r' for readable, 'w' for writable, 'x' for executable or searchable
 - Next 3 chars are group permissions

- Next 3 chars are public permissions
- Next is a number - number of directory entries that point to this file
 - When you create a file, its owner is the creator and the group is inherited from the group of the parent directory
- Followed by owner
- Followed by group
- Followed by file size in bytes
- Followed by last modification date
- Followed by file name within parent directory
- ls -a - includes output of files that start with .
- Emacs convention for file names
 - FOO~ is a backup file for FOO
 - #FOO# is a saved version of FOO while FOO is benign edited
 - Emacs tries to get around persistence (buffers are not persistent) by saving the contents of the buffer into the # file → autosave
- C-j evaluates commands in the scratch buffer
- ln A B - creates a new name B for the existing file A → A and B are equal afterwards
 - ln -s A B - creates a symbolic link from B to A
 - Symbolic links are neither regular files nor directories → just file name redirections
 - .#FOO is a symbolic link to nothing, just tells Emacs who's editing FOO for overwrite protection
- rm - remove a file
- cat A B C - copy the contents of A, B, C to output
- head -N A - copy the first N lines of A to output
- Why have hard links?
 - Hard links efficient sharing of data
- Why have symbolic links?
 - For meta-information
 - When your file has the wrong name

10/8: REPL

- Emacs and the shell are both instances of the same pattern
 - The Read-Eval-Print loop pattern
 - Program deals with the outside world as follows:
 - 1. reads a command from user input
 - 2. evaluates that command using the syntax and semantics of a particular programming language
 - 3. prints out the result of the command
- To get details about a character: C-u C-x = 'char'
- UTF-8 encoding
 - ASCII characters represent themselves → ASCII is a subset of UTF-8
 - Succeeding parts of a multi-byte encoding begin with bits 1 and 0
 - Therefore, has 6 bits available for encoding

- Everything else is the beginning byte of a multi-byte encoding
 - Begins with as many 1s as the encoding length
- Most commonly used chars have the lowest numerical representation for space efficiency
- Multi-byte encodings don't overlap bounds for security reasons/ambiguity
- Emacs REPL
 - Emacs is modeful - C-x puts you in a different context for commands
 - C-h b lists all keybindings in the *Help* buffer
 - Can use to get command's documentation
 - Can think of Emacs as having 2 REPL
 - At a low level of abstraction - reads a character, executes the function designated by the character, and then displays the resulting screen
 - In the *scratch* buffer, at a higher level of abstraction - Waits for the user to type C-j, reads the contents of the buffer, looking for an expression, evaluates the found expression, then displays the answer in the *scratch* buffer
- Shell REPL
 - Some commands
 - true
 - : (like 'true')
 - false
 - Commands can succeed or fail, you can tell the difference by looking at the command's exit status → the integer that 'main' function returns (0 = success, anything else = failure)
 - You can look at this return value using shell variables
 - Specifically \$? - the exit status of the most recent command
 - || - takes the first command, executes it, if it fails, executes the next command
- DNS - Domain Name System (Paul Mockapetris)
 - Maps www.ucla.edu to IPv4/IPv6 address
- Emacs language is Emacs Lisp (+ 2 3) etc.
- Shell language is the shell language
 - Why the Shell?
 - You have an OS running multiple programs
 - The shell is a thin layer that connects these programs
 - It's thin because it doesn't do anything except be a gateway to the important stuff
 - It's a program that has a REPL and can parse input
 - It has its own programming system, and can be bulked up by the programmer
 - It's just another command, for convenience → takes the capabilities of the OS and connects them together
 - Some common commands:
 - Nondestructive -

- true, false, : - placeholders for control flow
- echo a b c - write 'a b c' to output
- cat a b c - write the contents of the files 'a b c' to output
- exit - exit the shell
- C-c C-c - gets to the top level of the shell
- ps - process status
 - Lists status of all statuses ps thinks you're interested in
 - -e does something
- tr - transliterates inputs (find and replace)
- man X - gives documentation on command X
- grep - looks for patterns in input and copies lines with that input to output
- which - gives the location of a file
- ls - list file information
 - Defaults to base directory '.'
- cd - change directory
 - /a/b/c/ file names are absolute → context independent
 - Begin with '/'
 - a/b/c file names are relative → context dependent
 - Don't begin with '/'
 - Meaning depends on the current working directory

10/13: Scripting

- Client-server application:
 - Several different ways to hook small parts of an app into larger ones
 - Subroutines and main program (function/method calls)
 - Primary/controller and worker nodes (multiple machines)
 - One machine is in charge
 - Other machines accept tasks from the primary, do them, and then come back to ask for more work
 - Client/server - single server that maintains centralized state + clients that talk to users, get requests, ship off to server, get response back, show that response to user
 - Clients are in charge in some sense, server waits passively for request to come
 - Typically the client is a program and the user is a person
 - Good approach for reasonably small applications (my.ucla.edu)
 - Very often, the server has a database as a component, but this is not required
 - States may be kept in RAM
- Scripting
 - Shell, Lisp, and Python
 - Why have multiple languages?

- Different strengths and weaknesses → none of them dominate everywhere, none of them are ridiculously bad everywhere
- Ease of learning (be careful of inertia)
- Performance
- Ease of/flexibility of writing and maintaining code
 - Ideally a scripting language lets you write in one line what would take like 30 lines in C++
 - Flexible notation makes it easier on the programmer
- Reliability
 - Scripting languages are going to be more reliable in terms of avoiding low level errors (bad pointers, subscript errors, etc.)
 - Can still occur, but can recover instead of core dumping
 - Typically, compiled languages have better compilers that examine the code more carefully that can catch errors that the scripting language won't (better static checkers)
 - Dynamic checking - you never know if the program will bug out, runtime behavior varies from run to run
- Scripting languages don't scale as well as traditional languages
 - Designed more for smaller applications
 - Scaling issue arises due to diseconomies of scale
 - Adam Smith - a pin factory is a very efficient way to make lots of pins → not needed for a small village → saves money for a large city
 - Diseconomy of scale → the bigger your system gets, the higher cost per unit
 - Ex) a typical Python program → all languages have this problem, scripting languages have it worse
 - 10 lines - easy to understand, low memory usage
 - 10 mil lines - hard to understand, won't fit in machine
 - Zillions of connections between modules, lots of things that can go wrong
 - Writing another 10 lines will take much longer (10 lines/day is a standard average rate of production)
 - Isolation of code is easier in compiled languages, so they scale a little better
- Performance and flexibility conflict
 - Scripting gives up performance for flexibility
 - Values human time over machine time
- Shell Scripting
 - The shell has several metacharacters that need quotes to stand for themselves
 - `=!#$%&*() \ | " ' ~ < > ? [{ ;` space tab newline

- POSIX - Portable Operating System Interface X
 - Derived from an OS called Unix → Unix derived from research/production system called Multics
 - Unix is a stripped down Multics
 - Multics written by 2 guys in spare time → Turing award winners
 - Very simple compared to other OSes
 - Easy to hook together applications out of code
 - GNU/Linux is a “clone”/“imitation” of Unix
 - Linux is the kernel, GNU is the OS
 - FreeBSD / macOS, OpenBSC, NetBSD are other imitations of Unix
 - How do you write an application that will run on any of these systems?
 - POSIX attempts to answer this question at 2 levels
 - Higher level - the shell → there is a POSIX standard for the shell language and for the applications and utilities that you can run from the shell
 - Lower level - the libraries and system calls → called from C/C++, specifies things like <stdio.h>, <unistd.h> what they contain, and how you use their functions
 - Tension/design decision to make here
 - Should you try to do everything in one language?
 - Should you write a multilingual application, using a language that is well-designed for each individual part of the application?
 - POSIX (GNU/Linux, Unix) is aimed at this option
 - Idea is to use the best available tool for the job, whether it's a program or a programming language
 - “Software Tools” approach
 - “Little Languages” approach - don't build a single language to solve every problem, use small languages for each task
 - C, sh, sed (stream editor), awk (text processing - sed on steroids, syntax is like C), grep (searcher)
 - Perl = the union of all the above little languages, designed by someone who wanted to unify the language
 - “The Hedgehog and the Fox” - Isaiah Berlin

- Hedgehog knows one big idea and tries to put everything in the world under that idea
 - Foxes are always running around chasing new ideas, don't try to integrate them
 - Point is that Leo Tolstoy was a fox who wanted to be a hedgehog
- Pros and cons - you have to learn each little language
- Typically better for large applications, since no single language is appropriate, not a big deal to add a little language to the mix
- Little languages have a similarly little featureset, the spec is simpler
- A little language:
 - grep - defined by POSIX: GNU grep is an extension to POSIX grep
 - Comes from g/re/p - globally look for regular expression and print
 - grep 'BRE'
 - Reads input
 - If a line matches BRE, copies that to output
 - There is a little language for Basic Regular Expressions (BREs)
 - If BRE is a pattern, defined recursively as follows
 - x - x is any ordinary character, a simple pattern that matches only itself
 - . - matches any single character
 - BRE* - Zero or more concatenated instances of BRE
 - BRE1BRE2 - an instance of BRE1 followed by an instance of BRE2

- [abcdef] - matches any single character in the set abcdef
- [^abcdef] - matches any single character that's not in the set abcdef
- [^a-z] = matches any single character that's not in the set a-z
- a[bc]*d - matches any string of 2 or more chars, the first is a, the last is d, the remaining are all either b or c
- [~[:alpha:]/] - matches any single alphabetic character or ~ or /
- ^BRE - matches any instance of BRE that starts a line
- BRE\$ - same but ends a line
- \x - x is a special character, matches x
- grep x - copy to stdout every input line that contains the letter 'x'
- grep 'x*' - copy stdin to stdout
- grep xyz - match only lines containing 'xyz' exactly
- grep 'xy*z' - matches only lines containing x, followed by zero or more ys followed by z
- grep '^a.*z\$' - matches any line that starts with a and ends with z
- Early in grep development, there was a syntax dispute, so now there are 2 syntaxes in POSIX for regex
 - BRE is simpler and less powerful
 - ERE (Extended Regular Expressions) is more complex and powerful
 - ERE1 | ERE2 - either ERE1 or ERE2
 - ERE? - matches 0 or 1 instances of ERE
 - ERE+ - matches 1 or more instances

- (ERE) matches ERE
- A major problem in software construction is configuration
 - If you configure, say, SAP wrong, LAUSD won't be able to pay its employees

10/15: Scripting Continued

- Scripting is just programming - but it's also a way to think about gluing together large programs out of small ones
 - Maybe the script is at the top level
 - PyTorch (ML scripting) Python script + C++ modules
 - Maybe the script is a small part of a large rapp
 - Your web browser
- Example of scripting languages
 - sh - POSIX scripting language (top level is common)
 - Lisp - we'll look at this as a small part of a larger app (Emacs)
 - Python - top level in our example
 - JavaScript - used both ways
 - Server-side code in which JS is in charge
 - Client-side code in which JS is a subroutine
- The shell (sh, POSIX shell - Bash as an example)
 - Not indent-sensitive
 - Bash is a heavyweight shell with a lot of features vs. a smaller, faster subset of bash (sh)
 - It's a full-fledged language - lots of stuff in it
 - Contains various control structures (while loops, for loops, if-else statements, etc.)
 - while cmd1; do cmd2; done
 - for i in \$v; do echo \$i; done
 - if cmd2; then cmd2; else cmd3; fi
 - if grep {thing} /etc/passwd > /dev/null; then
 - echo 'statement 1'
 - else
 - echo 'statement 2'
 - fi
 - Can define functions
 - function f() { body of function; }
 - f a b c
 - Meta-execution → crucial part of software construction, using software to create software / program writes part of itself
 - You can put a shell script into a file and the file becomes a command
 - Use parameters with \$1, \$2, ... , \${10}, ...
 - \$(CMD) means execute CMD, capture its output, and make it a part of your shell program
 - echo "blah: \$(grep {thing})"

- eval command
 - eval “string” → treat the string as code and execute it
 - Compute the string from various methods and use eval to run it as code
 - Above methods are listed in increasing order of power
- Lisp as a Scripting Language
 - Lisp originally not intended as a scripting language
 - History lesson:
 - 2nd oldest language behind Fortran
 - Arose in the 1950s as part of AI research
 - LISP - LISt Processing - idea that the simple data structure of a dynamically allocated list could be used to implement a basic building structure for writing programs to play chess, NLP, etc.
 - Considered to be “classic AI” as opposed to modern ML
 - Lists are built from pairs and they can represent arbitrary data structures (a b c)
 - Some pairs are not part of a list → (a . b)
 - Empty list → ()
 - Can contain other lists → (a (b c) ((e)))
 - Parentheses represent levels of the list
 - Memory was really expensive → these lists were built dynamically → concern for bloating of memory causes resistance
 - First time pointers were used very extensively
 - Many variants: Lisp 1, 1.5 (1950s - 1960s), Scheme (1970s -), Common Lisp (1980s -), Emacs Lisp (1980s -), Racket (1990s -), Clojure (2000s -) → runs atop Java Virtual Machine, Hy (2010s) → runs atop Python AST
 - All somewhat niche, none are dominant in the world today
 - Sign of success and its simplicity
 - Emacs Lisp (Elisp) data structures and functions
 - Numbers
 - Evaluates to itself
 - May have rounding errors for certain levels of precision
 - Symbols (like identifiers)
 - Have values
 - Are objects in their own right
 - Data structures
 - ‘(a (b c) ((e)))
 - The empty list ‘() is called nil
 - nil represents the empty list and false
 - Implemented as the nullptr in early implementations
 - Function Calls
 - (a (b c) ((e))) → same syntax as data structures

- In C, you'd write "a(b(c), e())"
 - Whatever e() returns must be a function
- Same notation as for data structures
 - How do you tell the difference?
 - You quote it
 - Very easy to express code as data → very important strength of Lisp
- Quoting
 - Semicolon for comments
 - Normally you're thinking of your code as something you'd be executing, the Emacs interpreter executes code
 - (cos 3) → call the cosine function
 - Treat parens as data → you quote an expression
 - '(cos 3) → create and return that data structure → 2 item list with cos as one item and 3 as the other
- Standard functions in Elisp
 - (car L) - first item in the list L
 - (cdr L) - list of the remaining items in L (everything except the 1st item)
 - (append L₁ ... L_n) - create a new list, containing the concatenation of the contents of L₁ ... L_n
- Variables
 - We can create a global variable at any time and start using it
 - (setq abc (cos -1)) → prints cos -1 and modifies global variable abc to cos -1
 - Can access by typing abc
 - setq is an assignment statement
 - Avoid variables, focus on function calls
 - Can have local variables using (let ((a 13) (b -9)))
 - (+ (* a a) (* b b)) → $a^2 + b^2$
 - let is a local initialization: { int a = 13; int b = -9; return a*a + b*b ; }
- Syntax for function calls
 - (F A B C): F is the function, A B C are args
 - (setq a b)
 - Functions evaluate the parameters and operate on them
 - a would fail when evaluated → setq is not a function since it runs
 - setq is a "special form" → same syntax as functions, but not a function
 - In place where you'd normally find a function, you write a keyword
 - Differentiate by knowing what's a special form
 - setq is a keyword used for assignment

- Lisp experts tend to avoid `setq` and prefer `let`
 - `let` is a keyword used for initialization
 - `(if A B C)` → if A is true, do B and yield it, otherwise evaluate C
 - Only evaluates the A part and the B part
 - `(defun f (x) (+ x 1))` → sets the expression `x + 1` to the function `f`
- Built-in Functions → evaluate their args and execute
 - `(car L)`
 - `(cdr L)`
 - `(append A B C)`
 - `(cons A B)` → creates a dotted pair `(A . B)`
 - Can use with a list of size `n` to create an `n + 1` sized list
 - The above are standard for any Lisp implementation
- Emacs Built-ins (`C-x C-e` to execute commands outside of `*scratch*` buffer)
 - `(message "abcdef")` → creates a message to display to the user in the message line
 - Emacs equivalent of `printf()`
 - Can execute code from within shell using Emacs with `emacs -batch -eval (message "hi")`
 - `(current-buffer)` → returns current buffer, which is an object
 - Prints a brief summary of the complex buffer object
 - `(other-buffer)` → returns the buffer that isn't the current buffer that the user cares the most about
 - `(switch-to-buffer B)` → turns Emacs' attention to buffer `B`
 - `(point)` → where are we in the buffer?
 - Returns the number character you're cursor is currently on
 - `(buffer-size)` → returns how many characters are in the buffer
 - `(point-min)`, `(point-max)` → minimum and maximum values for point in the current buffer
 - `(point-min)` is always 1, unless Emacs is restricting your control of the buffer
 - `(point-max)` is 1 greater than the buffer size → cursor can be 1 char past the end of the buffer
 - `(goto-char P)` → move the cursor to position `P` in the buffer
 - `C-h f FUNCTION` → tells you about a given function
- In any Emacs buffer, you can evaluate Lisp with `C-x C-e`
 - In the `*scratch*` buffer, evaluating is so common that `C-j` is added as a command
- Errors (divide int by 0)
 - Emacs pops up a debugging window `*Backtrace*`
 - Lists the stack of functions being evaluated when the error occurred
 - To exit debugger → `C-]`
 - `C-h m` - tutorial for the debugger

- Trivia: dividing 1 by 0 doesn't error since floating point arithmetic overflows to infinity

10/16: Discussion 2

- Lisp
 - High level programming language that used parenthesized prefix notation
 - Common dialects today include Racker, Common Lisp, Scheme, and Clojure
 - We're using Emacs Lisp
- Lisp Basics
 - A function call $f(x)$ looks like $(f\ x)$ in Lisp
 - $(+ x_1 x_2 \dots x_n) \rightarrow$ summing numbers $x_1 \rightarrow x_n$
 - $(abs\ x) \rightarrow$ absolute value function
 - $(max\ x_1 x_2 \dots x_n) \rightarrow$ maximum function
- Lisp Lists
 - Three ways to create a list
 - $(cons\ object1\ object2)$
 - Object1 is the head and object2 is the tail
 - $(list\ objects)$
 - Any number of objects
 - $(make-list\ length\ object)$
 - Makes a list of objects with specified length
 - $(car\ list)$
 - Returns head of list
 - $(cdr\ list)$
 - Returns tail of list
- Lisp Functions
 - Defining a function: $(defun\ name\ (args)\ body)$
 - Ex) $(defun\ foo\ (ab))$
 - $(+ a\ b))$
- Python
 - If-else/for/etc.
 - Indent-sensitive

10/20: Python Scripting

- Python is often used as a string processing language
 - When you define a class/use an integral data type in languages like C++, you use a string in Python
- `line = 'F,100,15.49'`
 - line declared automatically, type depends on value assigned to it \rightarrow can change later
- `types = [str, int, float]` \rightarrow list of objects, you wouldn't be able to use this structure in C++ \rightarrow these types are objects, Python is a dynamic language
 - A lot of notions that are compile-time in C/C++ are run-time in Python
 - `spl = line.split(',')` \rightarrow splits the string, yields ['F', '100', '15.49'] as a list

- `zspl = zip(types, spl) → yields [(str,'F'), (int,100), (float,15.49)] → list of tuples`
- `a = str('F') → converts argument to a string`
- `b = int('100') → converts argument to a 100`
- `c = float('15.49') → converts argument to a float`
- `fields = [ty(val) for (ty, val) in zspl] → iterates through zspl with the first item in each pair called ty and the second called value, and calls ty(val) on each pair, resulting in ['F', 100, 15.49]`
- Python motivation and history
 - BASIC - developed in 1960s
 - Very popular teaching language: simple, like FORTRAN, scientific
 - Still popular in Microsoft
 - Problem with BASIC - very traditional and low-level language
 - Good at loops, functions, arrays
 - Not so good at other stuff
 - Project at CWI to replace BASIC in the 1980s → unteach bad habits
 - ABC - designed for 1980s computers, had its own development environment
 - Language is always properly indented because the compiler required it and IDE helped you out
 - Build simple stuff like hashing, sorting, etc. into the language so that students can write new programs, not just the same old stuff
 - Flopped
 - Perl - first general purpose scripting language that succeeded, made in US in 1980s
 - Take shell + grep + sed + awk
 - Take little languages and put them into a big language
 - Write code the way you talk
 - Hard to be a disciplined language with multiple ways to do things
 - Python - fixes Perl so there's only one way to do things
 - Indenting is done correctly
 - Think of it as Perl + ABC
 - First came out in a standard implementation that has evolved
 - C-Python → interpreter is written in C, most popular
 - There are other implementations → run atop Java Virtual Machine (from Oracle, etc.)
 - PyPy → Python running on top of Python
 - Indenting - blocks must be indented evenly, and more than their parents
 - String syntax - 'x' and "x" are the same, 'xyz' and "xyz" are the same
 - "xyz\ndef" - string of length 7, containing a newline
 - r"xyz\ndef" - raw string, string of length 8, containing a newline
 - Numbers in Python use same syntax as C, but complex nums are available
 - Python Objects
 - Every value is an object and has an identity, type, and value

- Identity and type cannot be changed once you've created the object, but the value can be changed if the object is mutable
- `id(a)` - returns the identity of `a` as an integer
- `isinstance(o, c)` - returns true if `o` is an instance of `c`
- Functions
 - `def f(x):`
 - `return x + 1`
 - Defines a function and assigns it to `f`
 - You don't need to name functions
 - `h = lambda x: x + 10`
- Classes in Python
 - `class a(b,c):` // `b` and `c` are the parent classes, multiple inheritance
 - `var = 12`
 - `def method(self, x, y):` // `self` is the object this method is being called on behalf of (this)
 - `return x + y + self.m2(var)`
 - Variables are looked up by a depth-first, left-to-right traversal across the parent hierarchy
 - Namespace control
 - A class is an object
 - It has a member `__dict__`, that contains the class' members as a dictionary (data type that maps name to values)
 - By convention, names starting with `__` are private → made private by mangling the names to outside
 - Names that start and end with `__` are reserved for Python internal use
 - Everything is dynamic: classes are objects, you can have lists of classes, you can poke inside them, you can modify them if you know what you're doing
 - Python is willing to give up safety for flexibility
 - Find your bugs by running the program, not by compiling it
 - Like `sh` and `Elisp`
- Why was Python successful
 - Partly came from ABC's built in operations
 - Higher-level than what you'd normally see → performance isn't as valued due to the scripting nature of the language
- Major categories
 - `None` (special value like `nullptr` in C++)
 - Numbers - `int`, `float`, `complex`, `boolean`
 - Sequences - strings, lists, tuples, buffers, ranges
 - Tuples are immutable while lists are mutable
 - Buffers are like strings, but they're mutable
 - Ranges are like ranges of integers

- Mappings - dictionaries (sets of name-value pairs)
- Callables - functions, classes, methods
- Internal

10/22: Python Operations

- Claim: Python standard types and operations on those types are at least partially responsible for its success
- When you talk about writing code in Python: language + library
- Sequences - commonly used types in Python
 - Sequences include lists, strings, etc.
 - Operations on sequences
 - `s[i]` - returns the *i*th element of *s* → valid indexes are `-len(s), ..., 0, 1, 2, ..., len(s) - 1`
 - `s[-1]` means the same thing as `s[len(s) - 1]`, etc.
 - Goes to `s[-len(s)]` as `s[0]`
 - *i* can be any expression yielding an integer
 - *s* can be any expression yielding a sequence (may need to be parenthesized)
 - `s[i:j]` - returns a subsequence `s[i], s[i + 1], ..., s[j - 1]`
 - `s[1:]` - everything in *s* except its first element, equivalent to `s[1:len(s)]`
 - `s[:j]` - equivalent to `s[0:j]`
 - If `i == j`, it's the empty sequence
 - `i <= j` should be the case after accounting for negative values
 - If `i < 0` or `j < 0`, they count backwards from the end
 - `s[0:-1]` - all of *s* except its last element
 - `len(s)` - number of elements in *s*
 - `min(s)` - minimum value in *s* → uses comparison
 - `max(s)` - maximum value in *s*
 - `list(s)` - constructs a *list* with elements equal to those of *s*
 - Not every sequence is a list
 - Lists are the most convenient sequence
 - Strings are not mutable
 - Can be concatenated with '+'
 - New object is constructed by the function
 - Operations on mutable sequences
 - `s[i] = v` - assignment to individual element
 - `s[i:j] = a` - changes a subsequence to the contents of *a*, may change the length of *s*
 - *a* must be iterable
 - `del s[i]` - deletes a sequence member, shrinking the length of *s* by 1
 - `del s[i:j]` - deletes `s[i], ..., s[j - 1]`, so `len(s)` decreases by `j - i`
 - Operations on lists
 - Every list is a mutable sequence, reverse isn't true

- `s.append(v)` - appends an item to a list
 - `len(s)` increases by 1
 - `s[len(s):len(s)] = [v]`, but it's fast
 - Lists have a pointer to a piece of storage
 - Object also contains the length of the list and the size of the list (amount of space allocated for the list)
 - append just has to place value at end (load/store) and increment length
 - When list gets too long, a larger area of memory is allocated and the list object is updated
 - Worst case: `s.append(v)` is $O(n)$
 - `s.append(v)` is $O(1)$ amortized
- `s.extend(a)` - append every element of `a` to `s`
 - Costs $O(\text{len}(a))$ amortized
- `s.insert(i,v)` - insert the value `v` just before `s[i]`
 - `s[i:i] = [v]`, but faster
- `s.pop(i)` - delete `s[i]`, return its previous value
 - `t = s[i]; del s[i]; return t`
- `s.pop()` - `s.pop(len(s) - 1)` - deletes the last element of the list and returns it
- `s.count(v)` - return a count of all members of `s` equal to `v`
- `s.index(v)` - returns the index of the first occurrence of `v` in `s`
- `s.remove(v)` - removes first element of `s` equal to `v`
 - `s.pop(s.index(v))`
- `s.reverse()` - trade `s[0]` with `s[-1]`, `s[1]` with `s[-2]`, etc.
- `s.sort()` - you don't have to implement quicksort
- Operations on strings
 - `s.join(t)` - joins the strings in `t`, using `s` as a separator
 - `s.split(sep)` - splits string into a list of words, using `sep` as a separator
 - `s.split(sep, maxsplit)` - live above, except `maxsplit` bounds the number of words
- Mapping types
 - Dictionaries
 - Indexed by arbitrary immutable keys as opposed to sequences, indexed by integers
 - `d['eggert'] = 27` // dictionaries are mutable even though keys are not
 - `d['eggert']` - yields the value in the dictionary whose key is 'eggert'
 - A dictionary is a partial function from keys to value
 - Typically starts being empty, change over time through assignment
 - From a user's point of view, straightforward generalization of lists
 - A list is like a dictionary where the keys are `0, ..., len(s) - 1`

- They're implemented via hash tables that the Python programmer doesn't see directly
- Curly braces mean dictionary, square brackets mean lists
 - `d = {}` - creates an empty dictionary
 - `e = { 'eggert':27, 'paul':'xyz' }` - creates dictionary of size 2
 - `f = {27:'eggert', 'paul':{}}`
 - `g = {{}:'eggert'}` ← not allowed since `{}` is mutable
- Why not allow keys to be mutable?
 - Since dicts are implemented with hash tables, if keys could mutate, whenever a key was changed, the hash table(s) containing the key would have to be rehashed
 - Hash tables were invented in the 1950s by an IBM programmer
 - Hash function $h(k)$ gives you an integer, mod that int with the hash table size, use the result as an index into an array to find the key
- Operations
 - `d[k]` - look up `k` in `d`, return the corresponding value
 - If the value is absent, a key error exception is raised
 - `d[k] = v` - store `v` as the value corresponding to `k` in `d`
 - `del d[k]` - remove key-value pair from `d`
 - `len(d)` - gives the number of items in the dict
 - `d.clear()` - discard everything from `d`
 - `d.copy` - clone the dict `d`
 - `d.has_key(k)` - true if `k` is a key in `d`
 - Like `d[k]`, but no `KeyError`
 - `d.keys()` - list of keys in dict
 - `d.values()` - list of values in dict
 - `d.items()` - list of key-value pairs in dict
 - `d.update(d1)` - merge `d1` into `d` (`d1` wins if conflicts occur)
 - `d.popitem()` - removes a randomish key-value pair from `d` and returns it
 - `d.get(k[, v])` - returns `d[k]` if it exists, `v` otherwise
 - `v` defaults to `None`

10/23: Discussion 3

- Regular Expressions
 - `'+'` - one or more
 - `'*'` - zero or more
 - `'?'` - zero or one
 - `'(' and ')'` - captures a group
 - `'{i}'` - match exactly this number of instances
 - `'{i, j}'` - match anywhere between `i` and `j` instances, inclusive
- What is React?
 - Open source JS library

- Used for building user interface
 - Backed by Facebook
- DOM (Document Object Model)
- Features of React
 - Declarative → you don't have to give instructions step-by-step
 - Flexible
 - Efficient → uses virtual DOM

10/27: Python Modularization

- Modularization and Packaging - meta-tools for writing software
 - Techniques for managing your code
 - Management is a big deal - can take a big chunk of your development costs
- Functions in Python
 - Functions are objects
 - Lots of languages are like this (not C/C++)
 - E.g., if you define a function $f(x, y)$, it creates a function object and assigns it to f
 - $g = f$ is valid
 - Functions have the same lifetime as a list, etc - lasts until it gets garbage collected
 - Python functions can have a varying number of arguments:
 - `def printf(format, *args):` // represents a placeholder for a tuple
format will be bound to the first arg
args will be bound to a tuple of the remaining args (`args[0]`, `args[1]`, ...)
 - Python functions can have named arguments:
 - `def arctan(x, y):`
computer the arctangent of y with respect to x
 - Can call with `arctan(y = 1.5, x = 2.7)`
 - Not assignment statements, just allows reordering of parameters
 - Can combine these things:
 - `def foo(x, y, **rest):` // represents a placeholder for dict
...
 - x is bound to 1st arg
 - y is bound to 2nd arg
 - rest is bound to a dictionary of the remaining args
 - `foo(3, 9, alpha = 0.1, beta = 9.3)`
 - rest is bound to `{'alpha': 0.1, 'beta': 9.3}`
 - Helps Python code be more extensible
 - Can later extend foo with a new keyword argument:
 - `def foo(x, y, z, **rest):`
 - Callers that do this: `foo(27, 19, z = 12)` will work with both versions of the code

- Functions can also have attributes:
 - `foo.secure = 1` # where foo is a function
- Classes and Typing
 - Recall that Python does dynamic type checking, not static
 - Dynamic - during runtime not compile time
 - `a = ...`
 - `b = ...`
 - `return a + b`
 - So, how does it work?
 - `a.__add__(b)` → invoking add method on a with parameter b
 - Leading and trailing `__` means the Python interpreter reserves these names
 - `class c:`

```
def __add__(self, other):
    return self.name + "+" + other.name
```
 - `x=c()`
 - `return x + y`
 - So, what does it mean to have a “type error” in Python?
 - C++ - compile error, Python - runtime error
 - duck typing - if you want to add numbers, run `x+y`, and if it works, this means x and y both waddled and quacked like ducks, so they must be ducks
 - Type checking is done by runtime behavior checking: if you don't get an error, it must be ok
 - This gives you a lot of flexibility: your code can work in a lot of environments
 - It also encourages error-prone code, as errors can easily slip through
 - Many other builtin method names
 - `def __init__(self a, b, c):`
Used by constructors `c(1, 2, 3)`
 - `__del__(self)` when your object is deleted
 - `__repr__(self)` create a string representation of the object (full version)
 - `__str__(self)` same thing, except shorter, might be abbreviated
 - `__hash__(self)` used when your object is a key in a dictionary
 - `__nonzero__(self)` used for 'if o: ...'; this calls `o.nonzero__()`
 - `__cmp__(self, other)` returns -1, 0, or 1 depending on `<`, `=`, `>`
 - Why does 'self' not exist in C++?

- C++ methods work by passing a pointer to the object to the method as a hidden argument, and you can see that argument in the method using a keyword
- Software Construction Management
 - Make it easy to plug things together or tear things apart later
 - Python modules (lowest level)
 - Typical modules are a single file with Python code to be executed at the right time
 - ocean.py file contains:
 - `abc = 27`
 - `def f(x):`
 - `return x + 2`
 - `class c:`
 - `def __init__(self):`
 - `self.val = 0`
 - `def bar(self, y):`
 - `return self.val + y`
 - ...
 - The right time occurs when you execute a statement to access the module
 - `if x < 0:`
 - `import ocean`
 - Up to the caller to determine the “right time”
 - Certain things happen when “import FOO” is executed
 - The Python interpreter creates a new namespace
 - Read the file ocean.py and execute its code in the context of that new namespace
 - Add a name FOO to the current namespace
 - FOO is bound to the newly created namespace
 - How to run a module from the top level → when Python starts up
 - `$ python3 modulename a b c ...`
 - imports module named ‘modulename’ with `__name__ == ‘__main__’`
 - Lots of modules are not intended to be top-level programs, they’re intended to be used as parts of other programs
 - Still it’s helpful to use this convention as a way of testing a module that isn’t top-level
 - foo.py:
 - definitions of some sort
 - `if __name__ == ‘__main__’:`
 - test cases for foo

- If foo isn't intended to be a standalone program, you turn it into one that runs test cases for foo
 - Test-first software development → first write the test cases for a module, then write the module's code
 - Why is this a good idea?
 - You have to write the test cases anyways
 - Test cases are easier to write than code
 - Let's you debug your module design faster
 - API has to be good enough to be tested
- Searching for Modules
 - Where to look for modules when you do import?
 - A Python installation consists not just of /usr/bin/python executable, but also of a bunch of files somewhere in the filesystem → where should it look?
 - Complicated answer because it's a big configuration problem
 - One part of the answer is PYTHONPATH
 - Environment variable in POSIX systems
 - Environment variables are global variables, set in the shell and their names and values are exported to subsidiary programs
 - Names are arbitrary shell identifiers, values are arbitrary strings
 - 'env' command lists your current environment
 - There is a module hierarchy as well as a class hierarchy
 - class c(a):
 - This means c is a subclass of a
 - class d(c):
 - D's grandparent is a
 - There's a tree of classes, in which the parent node is the parent class
 - In modules:
 - import ocean.island
 - import ocean.island.hawaii
 - Acts by reading 'ocean/island/hawaii.py' from some directory in your PYTHONPATH
 - We also have a hierarchy in modules, because the directory hierarchy is a tree and modules live in that tree
 - Keep these two hierarchies distinct in your mind:
 - Class hierarchy is about behavior
 - Child objects act sort of like parent objects (they should be compatible)

- Module hierarchy is about maintenance
 - It's typical for a single dev org to be in charge of a particular directory of the module hierarchy, and the module will be upgraded as a unit
 - You could have a.b.c be a subclass of d.e.f:
 - a/b.py contains:
 - class c(d.e.f):
- Python packages (a higher level than modules)
 - A package is implemented by having a directory
 - Contains module files (m1.py, m2.py, etc.) along with one extra metafile (`__init__.py`) that tells Python "this is a package" and is read whenever you import the package
 - It could be empty
 - More commonly, it at least defines `__all__` = [list of modules to be imported when the user wants to grab them all]
 - User does this by saying `from packagename import *`
 - `*` in imports is considered bad style by some → package may define names that you wanted to use
 - Maybe better to use `from packagename import a, b, c`
 - You can use relative names
 - `from . import x` → import from same package
 - `from .. import x` → import from parent package
 - `from ..z import x` → import from uncle package z
 - How do you put packages in the right place in the filesystem?
 - Can be done by hand → tedious and error prone
 - Standard for Python package installation
 - Continually evolving with time, more rapidly than the rest of Python
 - Assume Python 3.9
 - `$ pip install somepackage` → arranges for the package's files to be put in the proper spot so that it will be found
 - `$ pip uninstall somepackage` → undoes install
 - `$ pip list` → lists your current packages
 - `$ pip show --files somepackage` → lists files installed as part of somepackage
 - 3 logical places to get packages from
 - Get it from the Python installation
 - Traditional default, very simple
 - Shared by everybody who runs that Python
 - Get it from a standard place under your home directory
 - Newer approach, but it still has a problem: sometimes you'll need incompatible packages just for your own stuff
 - Get it from a standard place in your application's virtual environment

- Can install modules that disagree with each other without clashing

10/29: React and JS

- Basic idea for Python
 - We want a language core (reasonably small, general-purpose) + extensions via:
 - Python source code that you put into a library package
 - Code in some other language (C, C++, Fortran, etc.) that may be lower level → may have access to lower-level facilities or may be more efficient
 - Some combination of both
 - Sometimes not good enough
 - You can change Python (with effort)
 - PEP (python Enhancement Proposal)
 - You implement a change to Python (you have the source code) and propose it to the community → if accepted, it may appear in a future version of Python
- In short, Python is evolving
 - Every successful software technology is evolving
 - You need to adapt by knowing the evolution techniques
 - In the Python world, PEP is an extreme because it lets you change the language
 - More commonly, you extend Python instead of changing it
 - It's a continuum in practice
 - Some packages get used so much that they migrate into the Python core → dateutils
- Sources for Python packages + 3 ways to install them into the environment
 - Put it into /usr/lib/python/whatever → everybody can use package
 - Use PYTHONPATH to continue specify an alternate location, such as your home directory (per-user)
 - Virtual environments - lets you create a separate environment for each application (per application, all your procedures are on the same page)
- __pycache__ directory can contain these files to cache the result of compilation
- We're building an application by running a bunch of pip commands
 - No code, we're just figuring out what other packages we need
 - Mistakes will be made → need to be able to uninstall stuff, upgrade, tc.
 - You can save your state with pip freeze >reqs.txt
 - You can restore it later by doing pip install -r reqs.txt
 - File remembers all the stuff you needed to do to configure your application → like a spec
- How do you create packages?
 - You have to write some code
 - This code will have some requirements, which you'll have to tell users
 - Usually your code will have some legal requirements, such as a software license
 - LICENSE - big deal
 - README.md - "elevator pitch"

- md is short for Markdown
 - popular formatting language
 - yourpackage/code.py - source for a module (several of these)
 - yourpackage/__init__.py - code to run when your package is pulled in
 - setup.py - Python code to be run when your package is installed
 - Lots of stuff here
 - Let's focus on dependencies
 - tests/ - test cases (written in Python)
- Dependency Management
 - When one part of your software assumes another part
 - It is important in many phases of software construction
 - Build-time dependencies in traditional Linux/Unix apps
 - 'make' does this
 - Example 'make' rule in a file Makefile:
 - # foo.c contains '#include "stat.h"'
 - # stat.h contains '#include "sticks.h"'
 - foo.o: foo.c stat.h sticks.h
 - gcc -c foo.c # simple way to build foo.o
 - foo.o is the target → the file that you want to exist and keep up to date
 - stat.h and sticks.h are the dependencies → the things that the target depends on
 - You may need to build some of them because you have indirect dependencies
 - 'gcc -c foo.c' is the command - executing this command will fix any problem with foo being out of date with respect to its dependencies
 - Why not just use a shell script?
 - Shell scripts are not flexible enough - they don't capture the notion of dependencies very well
 - 'buildit' always starts from scratch
 - 'make' operates incrementally; it does the minimal set of commands needed to satisfy the dependencies
 - It can restart from a partially failed computation without doing all the work all over again
 - How does 'make' record whether a file is up to date?
 - Looks at a file's timestamps
 - Looks at checksums
 - 'make' remembers checksums the last time it is used and uses the checksums instead of timestamps
 - Where do you store the checksums?
 - Installation-time dependencies in Python/JavaScript/Unix/Linux/etc.
 - You have a package P that depends on package Q already being installed

- With pip, you say this in setup.py
- We have declarations of dependencies:
 - P: "I depend on package Q, version 3 or later"
 - P: "I depend on package R, version 2 or later, but version 4 or earlier"
 - Avoid this, prevents upgrades
 - ...
 - Q: "I depend on package S"
 - ...
- DAG of dependencies → nodes are packages and edges are when a package depends on another
- pip must resolve these dependencies
 - Builds the graph
 - Finds nodes that are already installed
 - Installs nodes that aren't installed in order
 - Makes sure it doesn't install a package unless it has installed all of its prerequisites
- At the high level, pip just looks at the dependency graph
- It's just looking at declarations
- It decides what to do
 - You can specify code to be executed when a package is installed
 - setup.py can contain arbitrary code, it's bad style
 - You want pip to have full control
- Assumption that later versions won't remove functionality from earlier versions
 - Not always true
 - Semantic versioning - version numbers indicate how compatible a package is compared to a previous version
 - P.Q.R. → incrementing P is a big deal - new version is incompatible with the old
 - Incrementing Q - new features, but they're all extensions to the old behavior
 - Incrementing R - no visible change in the API
- React
 - Client-server applications
 - Basic model:
 - Application is split into cooperating pieces
 - Each piece runs independently on its own "computer" (might be virtual)
 - One distinguished piece is called the "server"
 - Central part of the application

- Application's state (contents of variables, files, that tell you the state of the system) is centrally controlled by the server
- The other pieces are called "clients"
 - Peripheral to the application
 - Often talk to human users - have a GUI, touchscreen, etc.
 - Typically they do this:
 - Wait for user request
 - Format it
 - Send formatted request to server
 - Get response back
 - Display it to user
- There are other ways to do distributed applications
 - Peer-to-peer applications
 - Every client talks to every other client
 - No central server
 - More complicated management than client/server
 - With client/server, the server can manage things
 - BitTorrent is an example
 - Primary/secondary approach
 - One piece is primary (in charge of the whole computation; decides what to do next)
 - Others are secondary:
 - Wait for instructions from the primary
 - Do the task that the primary tells you to do
 - Ship the answers back to the primary
 - "Reverse" of client/server
 - Clients are telling the server what to do next
- Client/server performance
 - Throughput
 - How many actions per second can your application do?
 - You can support more users if your throughput is higher
 - Latency
 - What's the delay between a user request, and the response back to the user?
 - Typical user requirement: latency < 1 ms

11/3: Client-Server Apps

- Client server computing
 - Last time - performance throughput and latency
- Correctness issues in client-server computing (arise due to common ways of attacking performance issues)

- Out-of-order execution → client and server (operating on different machines) may not be executing things in the order which you want
 - Client code and server code are not always synchronized
 - We think: C1 C2 S1 S2
 - Actually: C1 S1 C2 S2
 - Code is attempted to run as fast as possible, communication between client and server doesn't always happen
 - Possibly no serial order that makes sense → client and server may be altering shared state
 - Serialization - assume client and server code are running in parallel, explain what happens by a serial order of all their actions, resulting from interleaving the respective actions
 - Coming up with an explanation where all observable behavior by either the client or the server is explainable by the order
 - Might not be what really happened
 - Good enough implementation if works
- Out-of-date caches
 - Clients often cache server state for performance
 - What happens when the server state changes?
 - Cache validation - the client tries to keep its cache synchronized with the server cheaply (in terms of incremental updates) in a timely way
- The Internet
 - Before the Internet the most widely used form of communication was the traditional phone system → used circuit switching
 - Preallocation of circuit → N bits/sec and S ms latency
 - Lot of wasted capacity, inefficient use of the hardware resource
 - Basic idea of the Internet is to use packet switching
 - Instead of preallocating the circuit, you break up the message into packets of fixed-size (relatively small, few kB), then ship it to the network
 - Best-effort transmission for each packet
 - Small packets have more header overhead
 - Large packets have less flexibility, need more memory in routers
 - Best size depends on hardware, latency, and throughput desires
 - More efficient use of the limited capacity, avoid having all the dead-time from previous implementation
 - First proposed by Paul Baran (RAND Corp)
 - Controversial because:
 - Unreliable
 - Billing
 - Packets have 2 parts: headers and payloads of a particular format
 - Headers - meta-information about the packet, overhead
 - Payloads - explains the data within the packet, data intended for the recipient

- Packets are exchanged via protocols (specify packet formats, the order packets can be sent in, why the packets are being sent, etc.)
- Cons:
 - Unreliable
 - Packets can be lost to network congestion
 - Packets can be received out of order due to routing
 - Packets can be duplicated due to misconfiguration or other low-level hardware things
- Internet protocol suite addresses the above issues
 - Basic idea: use layers, don't try to solve it all at once
 - Lowest layer: the link layer - a point to point protocol, hardware-oriented
 - Layer 1: the Internet layer - shipping of packets with the above problems
 - Layer 2: transport layer - channels (or streams of data) implemented by sending packets
 - Transport layer only sees the channels
 - Layer 3: application layer - specific to apps (web browsers, video, etc.)
 - Core protocol is called the Internet Protocol (IP)
 - For level 1
 - Comes in various versions
 - IPv4 specifies packet formats, and it's connectionless, each packet is independent;y generated, sent, received
 - Head contains length, protocol number, source address (32-bit number expressed like 192.168.1.9), destination address, cheap checksum (let's the recipient detect data transmission errors), TTL (time to live) hop count (number of routers the packet has traversed)
 - Apps that really want packets instead of streams typically use UDP
 - User Datagram Protocol
 - Very thin layer over IP
 - "Datagram" is close to packet → one piece of data sent independently of other pieces of data
 - Transmission Control Protocol (TCP)
 - Streams of data that are:
 - Reliable
 - Ordered
 - Error-checked
 - Via:
 - Divide the stream into sequenced packets
 - Flow control
 - Retransmission and reassembly
 - Many application protocols built atop TCP/IP

- Real-time Transport Protocol
 - Runs atop UDP, intended for real-time applications (live video, etc.)
 - Video apps care about timing, if a packet gets lost, we want to keep going
 - TCP would cause jitter
 - HyperText Transfer Protocol (HTTP) runs atop TCP
 - We want a reliable copy of someone else's webpage
 - The World Wide Web
 - Invented by Tim Berners-Lee at CERN
 - A physicist's use of TCP to solve: how to present results on screen?
 - Browser (client) + web server
 - Client talks to server via an app protocol built atop TCP
 - 2 basic components:
 - HTTP protocol
 - Originally very simple
 - Client sends a "GET" request
 - Server responds with the resulting web page
 - Easy to implement, explain
 - Some problems with the protocol:
 - Security: no encryption, minimal checksumming
 - Uncompressed - bloated
 - No server push (client always had the initiative)
 - No pipelining (a single webpage with several components)
 - You want the ability to send several requests before getting the responses back
 - No multiplexing (several windows talking to same server)
 - Several connections from client to server, doing different things
 - Addressed by HTTP/2 by complicating the protocol
 - HTTP/3 (in progress)
 - Fundamental change is that HTTP/3 uses UDP instead of TCP
 - Can do more multiplexing
 - Avoids head-of-the-line blocking delays - first packet in response is delayed
 - HTML data format
 - Borrowed from book publishers
 - SGML - Standard Generalized Markup Language
 - Declarative language for text formatting
 - Document Type Declaration (DTD) specified what markup elements were allowed
- HTML notions and terminology
 - HTML element

- Node in the abstract tree that structures the HTML text document
- Highest level is HTML element (single node root of tree)
- Lowest level will be simple elements
- Each element is surrounded by tags
 - <tag> [contents] </tag>
 - Closing tag can be omitted if obvious
 - Void elements do not have a close by definition, never have child elements
- HTML is a way of sending tree-structured text over the Internet
- Documentation for HTML can be found on MDN
 - Mozilla Development Network
- HTML started off with specific DTDs
 - A DTD is a spec for the trees you can put into a document or a grammar for the HTML language that can be used
 - These DTDs helped specify the original web
 - But, they were too limiting
 - For example, they didn't do videos
 - How do we let DTD's evolve rapidly enough to new apps?
 - Eventually gave up on DTDs as a standardization mechanism
 - Document Object Model (DOM)
 - Each HTML document is a tree
 - Easy and standard way to walk through the tree and to modify the tree
 - Comes with APIs for traversing, updating the tree
 - Callable from any language
 - JS is the most common language

11/10: More Client-Server Model

- HTML
- DOM (Document Object Model)
 - Standard way (APIs) to manipulate what you see in a browser
 - The browser renders the DOM of the webpage
 - Static pages don't do anything
- Cascading Style Sheets (CSS)
 - Goal: separate concerns
 - Content - core part of your tree]
 - Presentation - how to present it to the user
 - Without CSS, you get a default and boring presentation
 - Styles are inherited by subtrees of DOM
 - You don't need to specify a style for everything
 - "Cascading" because a priority scheme is used
 - Get styles from ancestors, browsers, user, authors
 - Idea is to specify style without writing code to implement it
 - Allows more freedom to implement styles

- JavaScript
 - Another scripting language
 - Like Python, even more dynamic
 - Can be hooked into HTML
 - Want our browsers to be programmable
 - Can generate HTML (or DOM) to be rendered by browsers later
 - Not just a subroutine, can be thought of as main program of your webpage'
 - Relatively small and simple language
 - Has to fit into browsers (parse, compile, and execute)
 - Even today, IoT client are underpowered
 - So, you need libraries to build real apps
 - Even with libraries, it can be a pain to generate DOM/HTML
 - You have a bunch of calls to create a tree
 - You wanted to see the tree directly as an HTML-like syntax
- JSX - extension to JavaScript that allows easy generation of DOM/HTML via a syntax that looks like what you're generating
 - `const header = <h1 lang="en">blah</h1>;` → angle brackets signal beginning of JSX
 - Can be used anywhere in JS that a function call could be used
 - JSX produces a React element that implements the JSX
 - Can think of JSX as a preprocessor over JS
 - You have to know both JSX and JS levels to debug a system that you're building
 - Use JS inside of JSX
 - `const language = "en";
const class = 'CS 97';
const n = 3;
const header = <h1 lang={language}>{class} assignment {n + 2}</h1?>;`
- Efficiency issues can be understood by knowing how the browser takes/executes your code - browser rendering pipeline
 - HTML → DOM → ... → pixels on screen
 - Browser starts rendering before it knows how to render everything
 - Some optimizations in this process:
 - Can skip subtrees that don't look like they appear on the screen
 - Can skip JavaScript code inside an element that's low priority
 - Decide the overall geometry of the page and then start rendering some components
 - Routine nowadays, sometimes wrong
 - Resizes are necessary, may need to re-render the page in worst cases
- Notation issue
 - Battle between HTML and JSON
 - JavaScript Object Notation

- Notation that represents JS objects → data, not code
 - Same thing as saying it's a text format for communicating tree-structured objects
- Node.js
 - JS runtime for asynchronous events
 - Runtime - set of cooperative classes and methods for supporting a particular style of programming
 - Callbacks - user-defined functions that are called at particular points during execution
 - Cedes control from a called function back to the caller
 - Event handlers - callbacks executed when particular events occur (e.g. the arrival of a request from browser, press of a button, etc.)
 - Event loop - basic programming construct
 - ```
int main (void)
{
 for (;;) {
 Event e = waitForNextEvent();
 handle(e); // Key point: event handler must finish
 "quickly", can't wait for anything else
 }
}
```
      - Event handlers don't do I/O and don't block or wait for anything
        - Can request, but can't complete
      - No locks (needed for multithreading)
        - Because event handler runs by itself
        - Avoid many race condition problems
      - So, multiple CPUs running in the same address space doesn't work - scaling via multithreading is impossible
        - Can scale with multiple computers (multiple web servers)
        - Can scale with multiple processes on the same computer
- Node.js and React are built atop these ideas
  - Uses JSON, JSX, etc.
- You can use this idea to build lots of cooperating servers or to implement code that runs in browsers
- Node.js also provides packages
  - Lots of choices - active software ecosystem
  - npm manages them
  - The packages your project is using is recorder by npm in a file
    - package.json
      - metadata about your package: name, description, author, dependencies, etc.
  - npm init
- Version control
  - Git basics

- Git state for your project (past and future)
  - Git has an object database that contains the history of your project
  - Index file contains the plans for the future of the project
    - Cache for the source code directories and files
    - Creates commits
    - Handles merges
- Basic commands
  - git init - for projects starting from scratch
  - git clone - the more common case - cloning an existing project
  - git log - outputs log of changes made to the source
    - Reports all changes to the project
    - One entry per commit
    - Commits have auto-generated IDs
    - Abbreviations for IDs commonly used
      - HEAD - most recent commit in the current repository
      - HEAD^ - the just-previous commit to HEAD
  - git diff - compares 2 commits and tells you the differences
- Where Git came from, what problems made Git?

## **11/12: Version Control**

- Basic Git continued
  - index - says what changes you're planning to make in the next commit/set of commits
  - You're not just writing a program, you're developing a set of changes to a program
  - You want the best changes you can come up with
    - To make the program better
    - To "sell" your changes to fellow developers
      - Make your changes convincing
      - These changes are reliable → fix a new feature, fix a bug, etc.
      - These changes do what you say they do
- A few more Git command examples
  - git ls-files → outputs all the names of files currently available
  - git grep → calls grep directly off the repository
    - More efficient than normal grep for large programs
  - git config → tells you about the configuration of the repository
  - git show → tells us more information about a commit + meta-information
    - Tool for looking at lower-level information about a commit
  - git pull → updates the current repository by copying all upstream commits
  - git blame → tells you who wrote a specific line of source code, helps find out why lines are the way they are
  - git reset → go back to the previous version, discarding all changes since then
- Git internals

- Each revision is an object in the Git repository
  - Each stored as a set of changes from the previous revision
  - Contains pointers to previous revision
  - HEAD points to latest version in the repository
  - May not be linear, another branch making independent changes
  - Main default branch is called the master branch
- `git branch -m master master-bad` → rename master branch to master-bad
- `git branch --track master origin/master` → creates another branch called master to what's upstream
- Merging involves at least 4 commits
  - Common ancestor
  - Branch A
  - Branch B
  - Merged descendant
- Low-level programming
  - Commonly done in C or C++
  - Not used much to do quick/user-facing applications
    - Used to access hardware features, do system calls, bypass efficiency gotchas in Java, Python, JS
    - Used to be software components or low-level tools that end users don't see directly
  - Examples
    - C-Python interpreter, Linux kernel, Emacs interpreter - C
    - Chromium, Firefox, JavaScript V8 - C++
  - C++ builds an abstraction layer atop of C
    - Classes - inheritance, encapsulation, polymorphism
      - Not in C
      - Data objects can be abstract - we don't know how they're implemented and we don't care → more of a pain in C, not done often
    - Namespace control → better modularity
    - Overloading is easy → badly supported in C, rarely used
    - Exception handling → C has low quality exception handling
    - Heap memory build in → C, it's just functions (`malloc()`, `realloc()`, `free()`, `<stdlib.h>`)
    - `cin` and `cout` → abstract ways to do I/O → C just has functions (`printf()`, etc.) in `<stdio.h>` and `<unistd.h>`
  - Architecture of a C/C++ environment
    - Compiling, linking, executing split into pieces
      - `$ gcc -E dumb.c` → just run preprocessor (expands macros)
      - `$ gcc -S dumb.i` → compiles from macro-free C to assembly language

## **11/17: Low-Level Programming**

- Low level programming
  - We want to have the ability to write some low level code
  - Some of the low-level stuff is fairly advanced, as far as software construction tools go
  - Some tools are even nicer than Python, JS, etc. to some extent
    - Of course, C/C++ are far less convenient
- What's between C and the machine level?
  - Can see by looking at what gcc generates
  - Can ask to show the assembly language code
    - `gcc -S -O2 prog-name`
      - `-S` → generate assembly language into a .s file instead of generating an executable
      - `-O2` → optimization option
        - Code is not optimal → code is just better
        - Harder to understand, longer to compile, sometimes worsens performances
    - We see that `sqrt(x)` is usually implemented by a single instruction, although a function exists as a fallback → efficiency
- Path between source code and what's actually running in the machine
  - `foo.c` (source code file - `foo.cc` if C++)
  - `foo.s` (assembly language - textual representation of machine language)
    - `gcc -S foo.c`
  - `foo.o` (object code - binary representation of the machine language)
    - `gcc -c foo.c`
- How does an executable work?
  - A copy of most of the program is put into main memory
  - This contains instructions and data
  - The OS jumps into that copy
  - The instructions in the program are now in charge of CPU
  - The first thing these instructions do is dynamically link in whatever libraries you asked for
    - A copy of the dynamic linker is available in the program → this like in the libraries you need using system calls
      - Ordinary function calls work by executing single instructions in your program
      - When you call a function, it may execute many instructions before returning, but these are all part of your program
        - The function can't do anything that the caller can't do
      - When your program runs, it's walled off from other programs for security and walled off from your computer for security
      - So there has to be an escape hatch, where your program can do something dangerous under the supervision of the OS
      - This is called a system call → it looks like a function call, but it's not: it's a single weird instruction that causes the program to

temporarily suspend while the OS does the real work in a safe way

- strace is a standard utility that tells you what system calls a program does
  - Logs all the system calls
- Other tools (besides compilers) for doing low-level development
  - Operations maintenance tools (used for SEASnet ops staff)
    - ps, top, etc.
  - Developer tools (used by developers)
    - Used to find out what's wrong with your program
    - time → runs a program for you and gives time values
    - strace/ltrace (library calls)
    - valgrind → looks for probable mistakes in the execution of your program
    - GDB → debugger
    - GCC → compiler
  - DevOps says developers and op staff should be interchangeable
- Let's look at compilers
  - 2 major free compilers: GCC (GNU/Linux) and Clang (macOS)
  - We'll look at GCC
    - GCC internals manual describes how it's implemented
  - How is GCC built?
    - Portability - we want GCC to work on lots of platforms (ARM, SPARC, x86-64, etc.)
      - GCC should be able to run on SPARC architecture, for example
      - GCC should be able to generate machine code to run on ARM architecture, for example
      - GCC developers distinguish among:
        - Target - machine that GC will produce code for
        - Build - machine that you're compiling GCC on
        - Host - machine that GCC will run on
          - Could be all different architectures, generally at most 2
    - All GCC targets have flat address spaces - all pointers are the same size and have the same interpretation (indexes into a large array of bytes, which represent our program/data)
      - Describes most machines nowadays
    - How can GCC generate code for all these architectures?
      - It'd be too much work to rewrite/maintain GCC from scratch for every target
      - GCC is split into a machine-independent part (executed regardless of target) and a machine-dependent part (specific to the target, hopefully small)
      - Does this by having GCC developers write a machine description file that describes the target machine

- In a high level way for most things, but with hooks (calls to some machine-specific C++ code) for machine quirks
  - This machine description file is a separate programming language
  - If you write a new file, you can generate code for a new machine
- This is backend stuff, there's a similar thing for frontend stuff
  - Which language you're compiling
  - You can specify which language you're compiling, it'll give you a parser that'll feed in to the core part of GCC, the core will use the .md file to generate the assembly language
- What is GCC useful for once it's build
  - Generate machine code of course'

### **11/19: More Low-Level Programming**

- What GCC is good for, besides compiling your program
  - GCC is good for security improvement
    - Stack overflow is a classic way to break into a system
    - gcc -fstack-protector → tells GCC to generate extra code for each function in which stack overflow is a real problem
      - Compiles with canary protections
      - Slows down your code a bit
      - On by default in some systems
  - Performance improvement
    - -O optimize
    - -O2 optimize some more → make the compiler a lot slower, to make the program a bit faster
    - -Os optimize for space, not time → small code
      - Can inhibit function inlining (common optimization)
        - Caller →  $y = p(x)$
        - Callee →  $p(x)$  is a simple function
        - Inlining → pretends the caller just did the simple function without the function call
    - Optimization should not change the meaning of a valid program
      - But in practice, many programs are not valid → larger program = more likely to be bugged
    - Program might be busted and you want to debug it
      - Debugging optimized code is tricky → compile without optimization?
    - Modifying program to improve performance
      - Best is to use a better algorithm

- Help the compiler by writing the program in a better way
  - `__builtin_unreachable()` → GCC builtin
    - Compiler can assume that this function cannot possibly be called
- `__attribute__` as a way of improving performance
  - `__attribute__((aligned(x)))`
    - Tells gcc to align based on address
    - Tells GCC to perhaps waste memory
      - Proper alignment takes advantage of how memory is accessed/cached
  - `__attribute__((cold))`
    - Function is rarely used
    - GCC can put f into “cold” sections of the code
      - Stays in RAM/Disk → more effective use of caching
  - `__attribute__((hot))`
    - Function is used a lot
    - Cold and hot require work by the programmer → nukes performance if wrong
      - Better way → profiling
        - Run program, measure parts that are cold and hot
        - Recompile with this
- `gcc -flto`
  - Enables link time optimization
  - When GCC compiles a .c file into a .o file, it puts more stuff into the .o file
    - Enough so that you can inspect the .o file and reconstruct the source code easily
  - When you eventually link a bunch of .o files using this flag. GCC reconstructs the original source info and then optimizes your whole program
  - If program is large, this takes a long time, can still be worthwhile
  - This flag is less explored, can be buggier
- Static checking → debugging program before runtime
  - `_Static_assert`
    - Directive to the compiler, tells the compiler the expression must be true, writes a message and refuses the compiler otherwise
    - `assert` is done at runtime
    - More efficient (zero runtime costs), more reliable (guaranteed if your program compiles)
    - Downside: argument must be evaluated at compile-time to a constant
  - `gcc -Wall` → generate some sane warnings
    - Uninitialized variables



- Includes a wide variety of options, which you can disable or enable individually
- -Wcomment → warns about nested comments
- -Wparentheses → warns about precedence rules in C
- -Waddress → warns about addressing errors with pointers
- -Wstrict-aliasing → warns about aliasing
- -Wmaybe-uninitialized → is there a path through the function that might use a local variable without initializing it?
- gcc -Wextra → more controversial warnings (less likely to be useful to everybody)
  - -Wtype-limits → warns about sign conventions

## **11/24: Debugging**

- Using GCC, you're only ever looking at a single function
  - Don't look at a single function, look at the entire compilation unit (the whole source code file → includes .h files, etc.)
  - gcc -fanalyzer
    - Like -Wmaybe-uninitialized → look for use of uninitialized variables
      - Also finds this when crossing function boundaries
    - Interprocedural flow analysis
    - Can find more errors, can be a lot more expensive → can greatly slow down compilation
- Changing the source code to make it easier to check
  - `__attribute__((pure))`
    - Side-effect free function (no I/O, no visible storage modified), the return value depends only on the argument variables and on the contents of storage
    - Why do this?
      - Optimization → compiler can cache values into registers without worrying that the pure function has modified those values
      - Clarity / communicating the API to the users of the function
  - `__attribute__((const))`
    - Stricter version of pure
    - Writes down more of the function's API formally so that the compiler can check it and take advantage of it
  - Downsides:
    - `__attribute__` is less portable → only works with GCC-compatible
    - More work to change your code, decorate it with `__attribute__`s
  - Point is to evolve languages gently to make them more reliable/efficient/useful
  - Set of GCC attributes evolves
  - Don't go overboard, too much work to use all of them and sometimes causes GCC to issue unnecessary warnings
    - Static checking either misses errors that are actually in the program or cry false alarms

- Compilers cannot in general predict what a program will do
      - Halting problem is undecidable
- Dynamic checking (runtime checking)
  - gcc -fsanitize=address
    - Tells GCC to catch out-of-bounds address access
    - Generates extra code to subscript check when it can
      - Slows down program
      - Makes program crash reliably when it does the wrong thing
  - gcc -fsanitize=undefined
    - Catches undefined behavior other than address errors
      - Mutually incompatible with -fsanitize=address)
      - Ex) integer overflow
    - Ensures reliable crashes
  - gcc -fsanitize=leak
    - Catches memory leaks
  - gcc -fsanitize=thread
    - Catches common errors in multi-threaded code
      - Multiple instruction pointers for multiple threads of execution sharing the same variable
      - + performance due to parallelism
      - - logs of bugs due to race conditions
    - Operates by slowing down the application and looks for other ways to interleave threads
  - Not perfect
    - Only works for that particular run - bug may be input-dependent
    - gcc, clang, valgrind can miss some errors for efficiency reasons
    - Can be severe performance penalties
      - Not a big deal with Python, JS, etc.
- Common theme - GCC generates different code that's slower but has more checks in it
- API - Application Programming Interface
  - For a function, contract between the caller and callee
  - Caller must implement the API correctly
  - Callee must rely on the API\
- Portability checking
  - Your program works fine on Ubuntu but fails on Fedora
  - You can compile and run on one platform, but it doesn't mean it'll work on a different platform
  - How to deal with this?
    - Know what platforms you might run on/try them out
    - Cross-compile for other platforms
    - Run on other platforms
  - Not just a low-level problems
    - JS code must run on Chromium, Firefox, Edge, Safari, etc.
      - Also run on phone screens, laptop screens, etc.

- Debugging
  - Don't do it if you can avoid it - it's an inefficient way to find and fix bugs
    - In badly-run projects, it can consume more than 50% of your developers' time
    - Debugging efficiently can give you an advantage on your competition
    - Need to be proactive
      - Prevent bugs from happening in the first place
      - Make bugs easier to detect when they do happen
      - Technology: static checking, dynamic checking, test cases
      - Especially true for low-level code
      - Use a better platform, port to worse platform later
        - Or pick a better language
    - Defensive programming
      - Traces and logs (print statements, output put into log files)
      - Checkpoint restart
        - Every now and then, save entire app state into file in a format so that you can later restore that state
      - Assertions

## **12/1: GDB**

- Defensive Programming
  - Trying to prevent disasters
  - Shouldn't limit self to single technique
  - Traces and logs
    - Put in print statements
    - Organize and analyze the output
    - Web server logs, for example
  - Checkpoint/restart
    - Program periodically saves its entire state into a file
    - Program has an option to start with a saved state
    - This lets you restart the program if it failed
      - Hardware failure - restart with fixed hardware
      - Software crash - it may crash the same way, run it a bit differently
  - Assertions
    - Have to be relatively efficient, can't have side-effects
  - Exception handling
    - Try ... catch ...
  - Barricades
    - Divide your data into 2 major regions
      - Safe data - generated internally, can be trusted
      - Tainted data - data from the outside world
        - Can't be sure what the data is
    - Set of software routines that convert the tainted data to safe data
      - Filter out/edit issues

- Must be systematic
- Interpreters
  - When you don't entirely trust your program
  - Can check each statement as it's executed and catch errors, preventing bad behavior
  - Adding checks slows the interpreter down, interpreters are slow anyways
  - Defense against bugs in your programs
    - Ex: Throw a bug on a runtime error and let your program recover/do something more useful than crashing
- Virtual machines
  - Enlist hardware to build interpreters that are faster while still having safety
  - Insulate the real machine from your application → barricade
- Defensive techniques do not always work
- Debugging strategy - assumption is: your program doesn't work
  - Don't guess at random what the bug is
    - Horribly inefficient, especially for larger programs
    - Your program's features combine
    - When debugging, you're exploring a combinatorially-explosive space
  - A more systematic approach for larger systems with bugs:
    - Stabilize the failure - make it reproducible
      - Many failures are randomish
    - Locate the failure's source
      - From symptoms to cause
      - Requires real understanding of how the program works
      - Debuggers can help here
      - Terminology:
        - Error - mistake made by the developer
        - Fault - latent problem in the program
          - In principle, it's static
        - Failure - observed bad behavior by program during a run
          - Prevent failures by fixing faults and errors
- Debuggers
  - Your program runs under a debugger's control
    - Either a virtual environment like GDB or interpreter
  - The debugger can stop your program's execution and examine your program's state (contents of variables, registers, ip, sp, etc.)
  - The debugger can change your program's state
  - Since the state includes the ip (instruction pointer), it can run program's code that otherwise wouldn't run
  - If your program is optimized, it can be harder to debug because machine code is in a different order than the source code
- GDB
  - r ARGS <INPUT >OUTPUT - starts a program

- start ARGS <INPUT >OUTPUT - does the same as r, but sets a breakpoint on main
    - attach PID - lets GDB take control of an already running program
      - detach - lets the program run free again
  - set env PATH "PATH" - setting up environment variables
  - set cwd "DIRECTORY" - sets current working directory
  - set disable-randomization on/off - modifies Address Space Layout Randomization (ASLR)
    - Causes low-level primitives to create objects at random-ish locations
    - Makes it harder to exploit bugs in your program
    - Commonly enabled in modern era, makes program behavior irreproducible
    - Disabled by default
  - bt - generates a backtrace
  - b LOCATION - set a breakpoint at LOCATION, execution stops if it reaches there
    - LOCATION can be a function name, filename, etc.
    - info break - gives information on existing breakpoints
  - c - continue execution
  - step - single step to the next line, descending into any subroutines and stopping there
  - next - single step to the next line, let the subroutines execute
  - fin - continue until the current function finishes, then stop
  - u LOCATION - continue until we exit the current function or we reach LOCATION
  - stepi - like step, except executes just one machine code instruction
  - rc - reverse continuation: like c, except it goes backwards in your execution history
  - watch EXPR - stop the program if EXPR changes
    - Can slow down execution greatly, but if you keep EXPR simple, it'll run at full speed
  - checkpoint - saves the state and outputs ID for the saved state
  - restart CHECKPOINT\_ID
  - p EXPR - prints the value of an expression
    - p VAR = VALUE - assigns VALUE to VAR, prints result
  - p/x EXPR - print the value of EXPR in hex
- Debugging targets
    - You can use GDB on one machine X to debug a program running on a different machine Y
      - X and Y don't need to have the same architecture
    - Common for IoT or embedded systems, where GDB is too big and unwieldy to run on a small device
  - Back to version control
    - DevOps - combination of developers and operations staff
      - Traditionally different people
      - Nowadays, same person can have both roles

- Break down the boundaries between these 2 organizations
- Version control becomes more important in this world
  - You are developing both programs and program configurations at the same time
  - Use version control to make sure this stuff stays in sync
- Version control needs
  - Backups in case source code gets messed up
  - History for code/configuration historians
    - Developers or ops staff wondering why the code is the way it is
    - Looking here is often the most efficient way to understand a program

### **12/3: DevOps and Version Control**

- Version control needs
  - Backups - you've messed things up
  - History - for knowing why the code the way it is
    - Relationship between bug reports and feature requests and changes you make to the source code
      - Bug report → ... → weird source code
    - What's the relationship[ between 2 projects that get integrated/split apart
    - Review of the source code by the original developers, but they forgot why they did things
      - The why is important, the what is easy
  - The future
    - In software development, there are many possible evolution paths
      - Sometimes future features don't interact, so the development group can work on them independently
      - Too often, they do interact
        - You need plans to minimize problems due to these interactions
        - A good version control tool will support planning for the future, as well as looking into the past
- Backups and disaster recovery
  - You must be prepared for bad things to happen
  - Data/systems will be lost, minimize the damage when that happens
  - Periodically make a copy of everything, restore that from copy on disaster
  - In some sense, inverse of caches
    - Caches help performance by making throwaway copies in faster memory
    - Backups hurt performance by making permanent copies in stable storage, gaining reliability
  - If you're designing a backup system:
    - You must have a failure model
      - What things can go wrong
      - How likely they are to go wrong

- To prioritize when and what to back up
  - Examples:
    - Your flash drive fails in your laptop
    - A disk drive fails in your servers
    - You delete or trash files by mistake
    - An outside/inside attacker trashes files on purpose
- A failure model with a few more details:
  - Flash drives AFR (annualized failure rate)
    - Say it's 1% that the whole drive fails
  - If you have a backup policy: "I back up every file once per day", what's the possibility you'll lose data sometime in the next year?
  - Make sure your key assumptions are valid
  - Your failure model must be end-to-end, which means you must worry about recovery
- What to backup?
  - File data - just the contents
  - File metadata - information about each file
    - Last-modified time, last-accessed time, ownership, permissions
    - Directory organization
  - Filesystem metadata
    - Keeps track of the OS infrastructure lying underneath all the files on your system
  - Alternate possibility: just worry about the underlying hardware - just copy all the hardware blocks of data on the hardware
- When to backup
  - Do you back up every change to the system, or just some changes?
    - You can just back up every now and then, omitting intermediate states
    - You can back up just a part of your system, omitting less-important parts of the machine state
  - If you generate a lot of backups, when do you reclaim storage that's no longer needed?
- How to do backups cost-effectively
  - Do them less often
  - Do them to a cheaper device with poorer performance
    - Flash → disk
    - Disk → Magtape
  - Remote backups → let someone else do the work
    - Some trust issues here
  - Incremental backups
    - Instead of making a copy of all your files, just copy the changes you've made since the previous backup
      - Need a way of computing the delta from the previously backed-up system

- You can use timestamps for this (file metadata), backup all files newer than the previous backup
  - You can do something fancier than timestamps by having a delta that records just the changes
  - Idea is that the delta is smaller than the originals
    - You do incur the cost of computing the delta
  - The edit scripts can do the following say:
    - Insert some lines at location N
    - Delete some lines at location M
    - Replace = insert + delete
    - These 2 commands generally suffice
- Deduplication is an optimization of a backup policy
  - Suppose the data are organized as a sequence of blocks
    - cp bigfile bigcopy
      - Supposed to copy all the data from bigfile to bigcopy
      - The file system notices whenever you write a block of data that happens to equal a block of data it already contains
      - It cheats in that case by using a pointer to the already-existing block → hashing
      - cp at the low-level simply creates a file whose metadata are pointers to the other file
  - Problem with this approach → if a block is corrupted, several files may be corrupted
- Compression
  - gzip → backup can be smaller than original
- Encryption
  - For security in case an attacker gets hold of the backups
- Multiplexing (single backup device for many systems)
- Staging
  - Primary copy backed up to secondary copy to tertiary copy, etc.
- How do you know your backups are working
  - Test them by doing test restores
  - If you restore the data, how do you know the restored copy is correct?
    - Checksum your data (store them somewhere else)
- Backups for software developers
  - We're not just talking about code
  - It's for data, documentation, configuration, etc. (you're writing down everything anyways)
  - File systems with versioning
    - Version as part of the file name, application decide when there's a new version
      - open() or write() doesn't change version, close() does



- More typically, you'll need a new API to control this
  - Not every application wants to do this
  - Every version is a separate file
  - There's a limit to a number of outstanding versions a file can have
  - You can preen a filesystem by removing older versions of stuff
- Snapshot approach - filesystem periodically decides to make a snapshot of itself (of a single point in time for all the files)
  - SEASnet → WAFL
  - Backups are often written to different media, while snapshots are often done on the same media, for efficiency
  - Snapshots work better for procedural/human/software failures, not hardware failures
  - Snapshots are commonly implemented via copy-on-write
    - You pretend you made a copy, but you simply create a pointer to the original data, and you don't actually copy until someone modifies either the original or the copy
- Version control systems
  - Originally intended mostly for software development
  - Basic idea is to be smarter than versioning file systems
    - More efficient
    - More useful
      - Navigating through a pile of old versions of software
  - Useful and/or necessary features for version control systems
    - Keep histories indefinitely (at least for source code)
    - Record metadata as well as data
      - Not just file system metadata (last-modified time), but also the why for a change
    - Suppose for example you rename a file
      - You need to be able to link the 2 names together
    - Metadata about the history, not just about the source code
      - Ex: in Git, branch names belong to the history of the source code
    - Atomic commits
      - Change several files or objects all at once so that the history never records just some of the changes, you either have the old version or the new one
      - Essential for collaboration, you don't want to be messing with other developers' code
        - You want the source code to always be in a consistent state, not some random mix of 2 versions
    - Pre-commit hooks
      - Run just before commit
      - Can be used to sanity-check a commit
    - Post-commit hooks
      - Run just after a commit

- Can be used to update your other files
  - Signed commits
    - Cryptographic signature that is hard to forge that only you can create with the possession of a secret key
  - Format conversion
  - Navigation and visualization of complex histories of source code
- History of version control systems
  - SCCS (Source Code Control System)
    - Prototype written in Alish language Snobol
    - Rewritten in C for Unix, mutated as it went
    - Each file F in your source code has a corresponding history file s.F which contained all of F's history
    - Bulk of the history was at the end of the file, containing all of the lines that ever appeared in F
    - Start of the s.F file contained the metadata for F's history
      - This metadata included locations of line numbers in the bulk history
    - To construct an old version of F
      - Read s.F's copy of the metadata
      - Deduce where all the old version's lines are
      - Left-to-right scan through the bulk data, grabbing just the lines you want in order - 1 sequential pass through the data
    - At the end of s.F, there's a checksum
    - Cost to extract F is  $O(|s.F|)$
  - Several successors
    - RCS (Like SCCS, but free software)
      - Its bulk data put the most recent version of F first
      - Cost to extract latest F was  $O(|F|)$
    - CVS
      - RCS front end
      - Key differences:
        - Single commits that crossed file boundaries
        - Client/server
    - Linux kernel started out using CVS
      - It didn't scale to lots of developers
      - Switched to Subversion
      - Switched again to BitKeeper (proprietary from SCCS, etc.)
        - Not open source, so controversial
    - Git arose when BitKeeper started costing money
      - Design a new system from scratch
      - Do the data structure first

## **12/8: Git Internals**

- Git implementation/internals

- Versioning filesystems are built into the OS
  - Hard to change how they work - you must change the OS
- Git is all user-mode code
  - It's a program that you could've written given the time
  - You can evolve it without evolving the OS
- Linux was originally criticized for having too much in the kernel
  - To some extent, this criticism was right
- Built atop the Linux filesystem hierarchy (directories, files) atop compression (save space, can save time in some cases)
  - Brief aside about compression technologies (vast field)
    - 2 basic ideas used by Git:
      - Huffman coding - represent more-common symbols by shorter bit strings, rarer symbols by longer bit strings
        - Compress a string of symbols into a string of bits
        - Ex) Normal English has a lot of e's
        - ASCII uses up 7 bits for each char - let's represent E via a shorter bit string, etc.
        - Shortest way to represent English using this encoding? → look at probability of each English character occurring, build tree structure
        - Downside: assumes we have good estimates for probabilities of letters
          - Cross-language probabilities will fall apart
          - Static probabilities vs. dynamic probabilities
        - We can do better if we want to devote more resources to the problem
      - Dictionary coder - requires more memory in the sender and recipient
        - Create a dictionary of commonly-used words, then transmit and receive indexes into the dictionary
        - If sender and receiver have the same dictionary and words average more than 2 bytes each, this will be a win
        - Harder problem than Huffman coding - what's the best dictionary to use?
        - Static dictionaries vs. dynamic dictionaries
        - Dynamic - the sender builds the dictionary as content is processed, dictionary gets sent to recipient as it's updated
      - These 2 approaches can be combined
        - We can Huffman code the output of the dictionary coder
        - Approach is taken by gzip, zlib
        - Ordinary text in any language will compress well

- Not so much for audio/video
- Git internals
  - Git history is stored into read-only “objects” files, which can be shared among repositories on the same machine
    - Speeds up git clone, which needn’t copy these files, it can just link to them
  - .git subdirectory
    - Branches directory - obsolescent
    - config - configuration information for this particular repository
      - Overrides your Git configuration in ~/.gitconfig
    - HEAD - tells you where you are
    - hooks - scripts that Git executes at crucial points (during a commit, etc.)
    - index - your proposed next commit (“staging area”)
      - “Future” of your current branch
      - Developers are not writing programs, they’re writing changes to programs
        - You want to make the best set of patches you can to improve software/explain the improvement
        - You want to use the staging area to not waste time
    - info/exclude - like .gitignore
    - logs - contains reflogs, records of where branch tips used to be/are. controlled by git configuration
    - objects - where the actual history is kept
    - packed-refs - optimized version of refs
    - refs - keeps track of where your branches are, as well as tags
- Git objects
  - Git repository is a user-mode filesystem built atop Linux/POSIX filesystem
  - It solves many of the same problems ordinary filesystems do
  - It also addresses problems specific to software development
    - Metainformation about changes to the software
    - Commits that atomically change lots of different files simultaneously
  - Simplest Git object - blob of bytes
    - Any sequence of bytes - Git doesn’t care
    - SHA-1 checksum - fingerprint for data, reliance on the hash being unique
  - Trees of objects
    - A Git tree is like a POSIX directory (blobs are like POSIX files)
    - A tree contains a list of entries, each have a name (of the subtree/blob), type (blob or tree), mode (octal number that represents Linux permissions of the corresponding file), and SHA-1 hash of the entry

## **12/10: Git Externals**

- blobs
  - Take the SHA-1 checksum of the blob
  - Compress the blob string using zlib (Huffman + dictionary coding)

- Write that compressed string into a file
- Branches
  - What branches are used for:
    - Mainline and maintenance releases
      - Merging maintenance fixes to master
        - You won't merge every fix - some of the maintenance fixes aren't appropriate
        - One convention (Emacs development) - "Do not merge to master" in a maintenance commit (skip this merge)
        - "Cherry-picking" or "backporting" from mainline to maintenance branches
    - Alternate visions of the future
      - Do we do A or B? Do both, see which is better
      - You can still do merging, cherry-picking, etc. here, but neither branch is "main"
    - Feature branches (more common, cheaper)
      - Let's implement feature A without worrying about feature B + vice versa
      - Teams work somewhat independently, whoever finishes first gets an easy merge into the master
        - Whoever finishes second may get a harder merge → must consider all interactions between A and B
        - Overhead → work of the second merge, some changes may collide
    - Forking because of disputes among developers
  - 2 ways to keep master consistent - merge and rebase
    - Merge makes connection to old versions much clearer
- How do branches work?
  - A branch is a lightweight moveable pointer to a commit
    - It can be changed to point towards a different commit
    - Normally moved by installing a new commit whose parent is pointed to by the branch
    - So, a branch keeps track of the latest commit in a logical sequence
- Merging
  - git merge X - merges (typically a branch) X into the current branch C
  - C and X will have a common ancestor A
  - First, we deduce A, then look at contents of A, C, X
  - Use this to infer the contents of the new commit that contains the changes in both C and X
    - Can do this by computing the difference between A and C and between A and X
    - Find the consensus between all 3 differences, marking merge conflicts
      - This is just a heuristic, just checking for textual overlaps
      - Conflicts fixed manually

- This is just for a 2-way merge, can be generalized for an N-way merge
- Rebasing
  - git rebase X - prepare all the changes from common ancestor and X and apply them to the current branch
  - git rebase -i X - interactive rebase (lets you treat each change differently if you want)
  - Should never rebase a branch that is shared with other developers
  - Can split/merge patches to help understanding
- Git is technology for preparing patches
  - git diff should generate nice output because many developers review patches by reading them
- Remote repositories
  - Git is a distributed version control system
  - No single repository is in charge
  - Repositories for the same project can disagree or can become out of sync
    - This happens all the time
    - Use different branches to minimize this problem
      - Private branches in your own repository → can rebase whenever
      - When you're ready, install into master and push upstream
  - git remote - lists your remote repositories (can be on the same machine)
  - git remote -v - lists a lot more detail
  - git remote show origin - even more
  - git fetch REMOTE - fetch changes from the remote repository into yours, do not merge
  - git pull - git fetch; git merge
    - Merge can make lots of changes, may be dangerous
  - git push - push changes from current repository into the remote
    - Publishing your changes upstream