

1. Write a function that, given a number  $n$ , returns another number where the  $k^{\text{th}}$  bit from the right is set to 0.

Examples:

`killKthBit(37, 3) = 33` because  $37_{10} = 100\mathbf{1}01_2 \sim> 100\mathbf{0}01_2 = 33_{10}$

`killKthBit(37, 4) = 37` because the 4<sup>th</sup> bit is already 0.

```
int killKthBit(int n, int k) {  
  
    return n & ~(1 << (k - 1));  
  
}
```

2. `mov` vs `leaq` - describe the difference between the following:

```
movq (%rdx), %rax  
leaq (%rdx), %rax
```

`movq` takes the **contents** of what's stored in register `%rdx` and moves it to `%rax`. `leaq` computes the load effective **address** and stores it in `%rax`. `leaq` analogous to returning a pointer, whereas `movq` is analogous to returning a dereferenced pointer.

3. What would be the corresponding instruction to move 64 bits of data from the memory location stored in register `%rax` to register `%rcx`?

```
movq (%rax), %rcx
```

(important part is that you know the suffix of the MOV instruction!)

4.

```

int cool1(int a, int b) {
    if ( b < a )
        return b;
    else
        return a;
}

int cool2(int a, int b) {
    if ( a < b )
        return a;
    else
        return b;
}

int cool3(int a, int b) {
    unsigned ub = (unsigned) b;
    if ( ub < a )
        return a;
    else
        return ub;
}

```

Which of the functions would compile into this assembly code:

```

    movl %esi, %eax
    cmpl %eax, %edi
    jge .L4
    movl %edi, %eax
.L4:  ret

```

cool2

- Arguments passed to a function is stored in the %edi, %esi, etc registers
  - %edi is a and %esi is b
- When comparing, we compare as *cmp Two One*
  - Thus the instruction jge is checking if %edi is greater than or equal to %eax
  - This is essentially checking if  $a \geq b$ , which is the else condition
- We can observe that when we do jump, %eax is not updated
  - We return b in the else case
- If we don't jump, we update %eax to %edi
  - We return a in the if case
- Thus cool2
- This question was inspired by a previous midterm

## 5. Operand Form Practice (see page 181 in textbook)

Assume the following values are stored in the indicated registers/memory addresses.

<u>Address</u>	<u>Value</u>	<u>Register</u>	<u>Value</u>
0x104	0x34	%rax	0x104
0x108	0xCC	%rcx	0x5
0x10C	0x19	%rdx	0x3
0x110	0x42	%rbx	0x4

Fill in the table for the indicated operands:

<u>Operand</u>	<u>Value</u>	<u>Operand</u>	<u>Value</u>
\$0x110	0x110 (immediate value)	3(%rax, %rcx)	0x19 (value in %rax is 0x104, value in %rcx is 0x5, $3 + 0x104 + 0x5 = 0x10C$ , value in 0x10C is 0x19)
%rax	0x104 (value stored in %rax)	256(, %rbx, 2)	0xCC (value in %rbx is 0x4, 256 in hex is 0x100, $0x100 + (0x4 * 2) = 0x108$ , value in memory address 0x108 is 0xCC)
0x110	0x42 (value stored in memory address 0x110)	(%rax, %rbx, 2)	0x19 (value in %rax is 0x104, value in %rbx is 0x4, $0x104 + (0x4 * 2) = 0x10C$ , value in memory address 0x10C is 0x19)
(%rax)	0x34 (%rax holds 0x104,		

memory address 0x104  
holds 0x34)

8(%rax)                   0x19  
(%rax holds 0x104, 8  
+ 0x104 = 0x10C,  
value in memory  
address 0x10C is  
0x19)

(%rax, %rbx)            0xCC  
(value in %rax is  
0x104, value in %rbx  
is 0x4, 0x104 + 0x4  
= 0x108, value in  
memory address 0x108  
is 0xCC)

- \$ denotes immediates
- Note: any numbers starting with "0x" are hexadecimal numbers!!
- All of the operands can be evaluated using the specific formulas on page 181 in the textbook
- More generally, whenever you see an address of the form  $D(r_b, r_i, s)$ , where  $D$  is a number,  $r_b$  and  $r_i$  are registers, and  $s$  is either 1, 2, 4, or 8, you can use the following formula:

$$D + R[r_b] + R[r_i] * s$$

If  $D$  is missing, assume  $D == 0$

If  $r_b$  is missing, assume  $r_b == 0$

If  $r_i$  is missing, assume  $r_i == 0$

If  $s$  is missing, assume  $s == 1$

- For more practice, try practice problem 3.1 on page 182 of the textbook