1/6: Introduction

```
• Start early, develop incrementally
• scp Desktop/Archive.zip <a href="mailto:charlesx@cs32.seas.ucla.edu">charlesx@cs32.seas.ucla.edu</a>:
• ssh <a href="mailto:charlesx@cs32.seas.ucla.edu">charlesx@cs32.seas.ucla.edu</a>
• unzip -j -a -d hw5 Archive.zip
• cd hw3
• ls
• g32 -o test *.cpp
#include <iostream>
#include <cstdlib>
#include <cmath>

    using namespace std;

• const double PI = 4 * atan(1.0);
class Circle {
       o public:
              ■ Circle(double x, double y, double r);
              void scale(double factor);
              ■ void draw() const;
              double radius() const;
       o private:
              ■ // Class Invariants: m_r > 0
              double m_x;
              double m_y;
              double m r;
};
double area(const Circle& x);
• Circle::Circle(double x, double y, double r)
• : m_x(x), m_y(y), m_r(r)) { // member initialization list
       \circ if (r <= 0) {
              ■ cout << "Cannot create circle with radius " << r << endl:
              exit(1);
       0 }
• }
bool Circle::scale(double factor) {
       o if (factor <= 0) {</pre>
              ■ return false;
       0 }
```

```
o m_r *= factor;
      return true;
• }
void Circle::draw() const {
      o ... code to draw the circle ...
• }
• double Circle::radius() const {
      o return m_r;
• }
• double area(const Circle& x) {
      o return PI * x.radius() * x.radius();
• }
• int main() {
      Circle blah(8, -3, 3.7);
      Circle c(2, -5, 10);
      c.scale(2);
      c.draw();
      cout << area(c) << endl;</p>
      o cout << c.radius() << endl;</pre>
      double ff;
      cin >> ff;
      o if (!c.scale(ff)) {
             ■ ... error ...
• } // guaranteed to return 0 if no number specified
```

- Interface (what we can do) vs. Implementation (how do we do it)

1/8: Linking

```
myapp.cpp:
       #include <iostream>

    using namespace std;

       void f(double x); // needed to declare
       o int main() {
              ■ double a = 3;
              ■ f(a);
              ■ cout << a;</p>
       } // defines main
       // needs f, cout, <<</li>
stuff.cpp
       #include <cmath>

    using namespace std;

       void f(double x) {
              ...
              \blacksquare sin(x)
       o } // defines f
       o // needs sin
• Libraries - collections of object files with code for library functions

    Defines sin, cos, exponential, cout, cin, >>, <<, etc</li>

    Needs nothing

• Object files combine with libraries to create an executable file

    Object file → executable is called the linker

• Nothing can be defined more than once → linker error
• Every need must be satisfied by some definition
• There must be exactly one main routine

    Convention - 2 files per C++ class

myapp.cpp:
       o #include "circle.h"
       o int main() {
              ■ Circle c(-2, 5, 10);
              c.draw();
       0 }
• circle.h: // class declaration
       o class Circle {
              public:
              ■ Circle(,,);
              ■ void draw() const;
              private:
              double m_x;
```

```
o };
      0
      double area(Circle x);
• circle.cpp: // implementation
      o #include "circle.h"
      o #include <cmath>
      using namespace std;
      Circle::Circle(,,)
      0 :__{{
      0 }
      0
      ○ Circle::void draw() const {
            ■ sin(...);
      0 }
      o double area(Circle x) {
            ...
      0 }
• -o myapp myapp.cpp stuff.cpp (no header files listed)
```

• Don't #include ___.cpp

1/13: Steps of Construction and Resource Management

- Steps of construction:
 - 0 1. -----
 - 2. Construct the data members, using the member initialization list; if a member is not listed:
 - If a data member is of a built in type, it is left uninitialized
 - If a data member is of a class type, the default constructor is called for it; error if there is none
 - If you declare no constructor for a class, the compiler writes a default constructor for you
 - o 3. Execute the body of the constructor

```
• class Circle {
      o public:

    Circle(double x, double y, double r);

      0 ...
      o private:
      double m_x;
      o double m y;
      o double m r;
};

    Circle::Circle(double x, double y, double r) : m_x(x), m_y(y), m_r(r) {

      o if (r <= 0) {
             ■ cerr << "...." << endl;</p>
             exit(1);
      0 }
• }
• class StickFigure {
      o public:

    StickFigure(double bl, double headDiameter, string nm, double hx, double

         hy);
      0 ...
      o private:
      string m_name;
      Circle m_head;
      double m_bodyLength;
};
```

• StickFigure::StickFigure(double bl, double headDiameter, string nm, double hx, double hy): m_name(nm), m_head(hx, hy, headDiameter/2), m_bodyLength(bl) {

// must have m_head in initialization list, since Circle has no default constructor, and will error out if in the body of the constructor, m_name is in this list for performance reasons

```
if (bl <= 0) {</li>
cerr << "....." << endl;</li>
exit(1)
}
```

Common to use member initialization lists to do as much of the initialization as possible

```
void f(String t) {
      String u("Wow");
• }
• void h() {
      String s = "Hello";
      o f(s); // constructing a String using another String, needs a copy constructor
• }
int main() {
      String t;
      o for (...)
             ■ h();
      o ...
 }
class String {
      o public:
      String(const char* value = ""); // makes value "" if there is no argument
          provided (default parameter value), only works for parameters at the end
      String(const String& other); // copy constructor
      String& operator=(const String& rhs);
      ~String();

    private: // Class Invariants: m_text points to a dynamically allocated array of

          m_len+1 characters
      o m len >= 0
      o m_text[m_len] == '\0'
      char* m_text;
      o int m len;
};
```

- String::String(const char* value) { // do not repeat the default parameter value(s) here

 m_len = strlen(value);
 m_text = new char[m_len + 1]; // must dynamically allocate a new array, otherwise you risk the original array being changed, therefore changing string's m_text's value unintentionally
 strcpy(m_text, value);

 String::String(const String& other) {

 m_len = other.m_len; // allowed
 m_text = new char [m_len+1];
 strcpy(m_text, other.m_text)
- String& String::operator=(const String& rhs) { // by convention, this returns the new value of the left hand side, return a reference to the object

```
delete [] m_text;
m_len = rhs.m_len;
m_text = new char[m_len+1];
strcpy(m_text, rhs.m_text);
return *this;
```

- String::~String() { // ensure that all dynamically allocated objects are deleted when their local counterparts move out of scope
 - delete [] m_text; // make sure you match the form of delete to the form of new used (plain vs. array)
- }

• }

1/15: Copy Constructors and Assignment Operators

- If you don't declare a copy constructor for a type, the compiler writes one for you
 - This copy constructor is made by copying all of the members
 - Creates problems with pointers:
 - Changing a value being pointed to inside a function that uses a copy also changes the original object
 - When other functions return, the destructor is called, and the value being pointed at is destructed
- If you write a destructor for a class, you likely have to write a copy constructor and assignment operator too
- Ok for any member functions to access private members of other objects of that class
- The copy constructor cannot be passed by value, as passing by value uses the copy constructor
- Assignment is not the same as the copy constructor

1/17: Discussion 1

- leetcode.com
- Dynamic Memory Allocation
 - Static memory allocation is inflexible
 - Assigns memory when compiling
 - new will automatically allocate the sequential memory space for the requested data type and size and return the starting address of the allocated memory space
 - Variables allocated with new will remain in the memory until deleted with delete
 - Dynamically allocated objects do not get deleted by the default destructor
- Copy Constructors

```
School::School(const School &aSchool) {
            ■ m name = aSchool.m name;
            m_numStudents = aSchool.m_numStudents;
            m_students = new Students[m_numStudents]
            for (int i = 0; i < m_numStudents; i++)</pre>
                  m_students[i] = aSchool.m_students[i];
      0 }
• Assignment Operator
      School& School::operator=(const School &aSchool) {
            ■ if (this != &aSchool) {
                  m_name = aSchool.m_name;
                  m_numStudents = aSchool.m_numStudents;
                  • delete [] m students;
                  m_students = new m_students[m_numStudents]
                  for (int i = 0; i < m_numStudents; i++)</pre>
                        o m_students[i] = aSchool.m_students[i];
            ■ return *this;
      0 }
```

1/22: Linked Lists

```
• Arrays limited by methods of resizing and inserting items into the middle of the array
struct Node {
      o int value;
      Node* next;
};
Node* head;
• Node* ptr = head
• while (ptr != nullptr) {
      cout << ptr->value << endl;</p>
      o ptr = ptr->next;
• }
for (Node* p = head; p != nullptr; p = p->next) {
      cout << p->value << endl;</li>
• }
• Usually preceded by a pointer (head) to the first element of the list

    Square bracket operators not valid

• When following a pointer, make sure it has previously been given a value and make sure
   that p is not the null pointer
Node* p;
• for (p = head; p != nullptr && p->value != 42; p = p->next) // evaluate left to right
      0;
• if (p != nullptr) {
      cout << p->value << endl;</li>
} else {
      cout << "There's no 42" << endl;</li>
• }
Node* p;
for (p = head; p != nullptr && p->value != 42; p = p->next)
• if (p != nullptr) {
      ○ Node* newGuy = new Node
      o newGuy->value = 37;
      o newGuy->next = p->next;
      o p->next = newGuy; // order matters, set new node's values first
• }
```

1/24: Linked Lists (cont.) and Aliasing

• Check normal case (middle), empty list, and boundary cases (first and last)

```
• struct Node { // doubly linked list
      o int value;
      Node* next;
      Node* prev;
• }
• void transfer(Account& from, Account& to, double amt) {
      o if (&from != &to) {
             ■ if (from.balance >= amt) {
                   from.debit(amt);
                   to.credit(amt);
                   • if (amt >= 10000) {
                          o fileForm(...) {
                   • }
             ■ } else {
                   • cout << .... << endl;
             • }
      0 }
• }
```

1/27: Stacks and Queues

- Stack/LIFO:
 - Create an empty stack
 - Push an item onto the stack
 - Pop an item from the stack, undefined on an empty stack
 - The top of the stack keeps the same position
 - Look at the top item in the stack, undefined on an empty stack
 - o Is the stack empty?
 - Maybe:
 - How many items are in the stack (C++ has)
 - Look at any item in the stack (C++ doesn't have)
- #include <stack>
- using namespace std;
- int main() {
 - o stack<int> s;
 - s.push(10);
 - o s.push(20); // adds 20 to top of stack
 - o int n = s.top() // returns the top item of the stack
 - o **s.pop();** // removes the top item of the stack, no return value
 - o if (!s.empty()) // asks if the stack is empty
 - cout << s.size() << endl;</p>
- }
- Using a stack signifies that you are only using the top element
- Queue/FIFO:
 - Has a front/head and back/tail, 2 active ends
 - o Create an empty queue
 - Enqueue an item (put it in the queue), can fail if there is a limited capacity
 - Dequeue an item (remove it from the queue), can fail if the queue is empty
 - Look at the front item in the gueue
 - Is the queue empty
 - Maybe:
 - How many items are in the queue
 - Look at any item in the queue
- #include <queue>
- using namespace std
- int main() {
 - o queue<int> q;
 - o q.push(10); // enqueues item
 - q.push(20);
 - o int n = q.front(); // n is 10

```
    q.pop();
    if(!q.empty())
    cout << q.size() << endl;</li>
    cout << q.back() << endl;</li>
```

- Queues and stacks have no limited capacity in C++
- String → numerical operations
 - Convert infix to postfix
 - If you hit an operand you append it to the result sequence
 - If you hit an operator and the operator stack is empty, push it onto the stack
 - Else if the top of the stack is an open parenthesis, push the operator
 - Else if the precedence of the current operator is greater than the operator at the top of the stack, push it
 - Else pop the top of the stack, appending it to the result sequence, and check again
 - Open parenthesis: always push
 - Close parenthesis: while top of stack is not an open parenthesis (operator) pop it and append it to the result sequence
 - Pop the open parenthesis
 - At end, pop each operator and append it to the sequence
 - Push operands onto a stack
 - When you hit an operator, you pop the operands and then push the result onto the stack

1/29: Inheritance

- Linked List vs. Array implementation of stacks
 - Harder to expand capacity with the array version
 - The array version takes up less memory
 - List version takes part in a lot of allocations and deallocations
- Linked List vs. Array implementation of queues
 - Use a doubly linked list for list implementation
 - Array:
 - Head and a tail (1 past the last element)
 - # items in queue is tail head
 - When hit end of array, you can shift everything back to the beginning of the array
 - Downside if there are a lot of items constantly, many shifts
 - Can also wrap around to beginning of array circular array
 - Head = tail if the list is empty or full
 - Make a size variable
 - Lists are easier when implemented as a doubly linked list
- Circular doubly linked lists are better for when you need to be able to insert/remove anywhere
 - Every node was guaranteed to have a node before and after it
 - Not very beneficial for queue implementation
- Inheritance:

```
class Circle {

    void move(double xnew, double ynew);

      void draw() const;
     o double m x;
     double m_y;
     double m_r;
};
• class Rectangle {
```

- void move(double xnew, double ynew);
- void draw() const;
- double m_x;
- o double m y;
- o double m dx;
- o double m dy
- };
- Circle ca[100];
- Rectangle ra[100];
- for (int i = 0; i < ...; i++)

```
o ca[i].draw();
• for (int i = 0; i < ...; i++) {
      o ra[i].draw();
void f(Circle& x) {
      o x.move(10, 20);
      o x.draw();
• }
• void f(Rectangle& x) {
      o x.move(10, 20);

    x.draw();

• }
• This code contains a lot of duplicate code, and adding new classes is tedious and error
   prone
• class Shape { // Shape is a superclass/base class of Circle and Rectangle

    void move(double xnew, double ynew);

      void draw() const;
      double m_x;
      double m_y;
};

    class Circle: public Shape { // Circle and Rectangle are a subclass/derived classes of

   Shape
      void draw() const;
      double m_r;
};
• class Rectangle : public Shape {
      void draw() const;
      double m_dx;
      o double m_dy
};
void Shape::move(double xnew, double ynew) {
      o m_x = xnew;
      o m_y = ynew;
• }

    Shape* pic[100];
```

- pic[0] = new Circle; // if you have a pointer/reference of a derived type, it can be automatically converted to a pointer/reference of the base type
- pic[1] = new Rectangle;
- pic[2] = new Circle; // heterogeneous collection

- Circle c;
- **f(c)**; // x has to be another name for a Shape object
- c.move(30, 40); // Circle doesn't have a move function, compiler moves up and searches its base class

1/31: Discussion 2

- Linked Lists
 - Minimum:
 - Key component as unit: Node with a value and pointer to the next Node
 - Head pointer → points to the first term
 - Loop-free (except in circular linked list)
 - Regular Operations:
 - Insertion, Search, Removal
 - Pros and cons:
 - Efficient insertion, flexible memory allocation, simple implementation
 - Arrays have to occupy contiguous memory locations
 - High complexity of search
 - Insertion:
 - Steps: Create a new Node, make its next pointer point to the first item, make the head pointer point to the new Node
 - Search:
 - Steps: Find matching Node and return, if no match then return
 - Removal:
 - Remember to set the previous Node's next pointer to point to the Node following the deleted Node
 - Consider what happens if the selected Node is the head or the last Node in the list
 - Suggestions:
 - Draw pictures and carefully trace through your code
 - Check any operations for the beginning of the list, end of the list, empty lists, and one-element lists
- Stack
 - o FILO: First in, last out
 - A standard stack implementation:
 - push() and pop()
 - Other methods: top(), count()
 - Applications:
 - Stack memory: function call
 - Check expressions: matching brackets
 - Depth-first graph search
 - How do you implement stacks with linked lists?
 - push(): Insert Node before head
 - pop(): Remove head and return the head value
 - top(): Read head Node
 - count(): Maintain a private int data member to keep track
- Queue
 - o FIFO: First in, first out
 - Basic methods:

- enqueue(), dequeue()
- front(), back()
- count()
- Applications:
 - Data streams
 - Process scheduling (DMV service request)
 - Breadth-first graph search
- o How to implement queue with linked lists?
- Suggestions on stacks and queues
 - o Draw pictures and carefully trace the code
 - o Infix to postfix conversion is very important
 - Stacks and queues can be applied to trees/graphs
 - You can use the given Standard Template Library to implement stacks and queues

2/3: Inheritance (cont.)

- Every derived object has a base object embedded in it
 - Base object is constructed first
- Calling a member function of a derived object in a function that takes a base object as the parameter will not compile
 - These kind of functions can only access functions that the base object can definitely access

void Shape::draw() const {

 ... draw a cloud centered at (m_x, m_y) // not the right implementation for derived objects

• }

- void Circle::draw() const {
 - o draw a circle of radius m r centered at (m x, m y)

• }

- void f(Shape& x) {x.move(10, 20);
 - o x.draw();
- }
- Static binding vs. dynamic binding
 - Connecting the name and the function
 - Static binding function decided in compile time uses Shape's draw()
 - C++ default
 - Dynamic binding function decided in runtime uses derived type's draw()
- Use virtual to make function dynamically bound
- class Shape {
 - void move(double xnew, double ynew);
 - virtual void draw() const;
 - o double m x;
 - double m_y;
- }
- **Shape::move(...)** // calling the base class' function
- class Rectangle : public Shape {
 - void draw() const;
 - virtual double diag() const;
 - double m_dx;

```
double m_dy;
• }
• double Rectangle::diag() const {
       o return sqrt(m_x*m_x + m_y*m_y);
• }
class Shape {

    void move(double xnew, double ynew);

       virtual void draw() const = 0; // function introduced as virtual, but not
          implemented, pure virtual function
       o double m x;
       double m_y;
• }

    You cannot create objects that have pure virtual functions

    You can create pointers to these objects

       o If an object has at least one pure virtual function, it is an abstract class cannot be
          constructed)

    Abstract classes still have objects, they are just never created by themselves
```

• class Shape {

```
class Polygon: public Shape {

...
~Polygon() // necessary for Polygon, not Shape
Node* head;
...

};
```

- Shape* sp;
- sp = new Polygon
- **delete sp;** // calls the Shape destructor, not Polygon's, must declare Shape to have a virtual destructor (must also be implemented)
- Destruction: 1) Execute the body of the destructor, 2) Destroy the data members, 3) Destroy the base part
- If a class is designed to be a base class, declare and implement a virtual destructor for that class
- class Shape {

```
public:
Shape(double x, double y)
...
private:
double m_x;
double m_y;
};
Shape::Shape(double x, double y) : m_x(x), m_y(y) {}
class Circle : public Shape {
public:
Circle (double x, double y, double r);
...
private:
double m_r;
};
```

- Circle::Circle (double x, double y, double r): m_x(x), m_y(y), m_r(r) {} // doesn't compile, m_x and m_y are private members of Shape, Circle cannot access them
- Construction: 1) Construct the base part, 2) Construct the data members, 3) Execute the body of the constructor
- Circle::Circle (double x, double y, double r) : Shape(x, y), m_r(r) {} // fixed

2/5: Recursion

- Base cases: can be solved without making a recursive call
- Recursive cases: make one or more recursive calls for strictly smaller (closer to a base case) instances of the problem
- Verify there's at least one base case and that each recursive call works towards that base case

```
    int arr[5];
    ...
    sort(arr, 0, 5);
    void sort(int a[], int b, int e) { // sort a[b] to a[e - 1]
    if (e - b > 1) {
    int mid = (b + e) / 2;
    sort(a, b, mid); // sort left half
    sort(a, mid, e); // sort right half
    merge(a, b, mid, e); // merges sequences b → mid and mid → e that produces a sorted final array
    }
    }
```

- If tracing code, start tracing small, the problem should appear early on
- Divide and conquer solve the problem by breaking it up (see above)
- The first and the rest divide collection into first and the rest, then solve and combine
 - Sometimes more convenient to to last and the rest

```
bool has(const int a[], int n, int target) {
if (n <= 0) return false;</li>
if (a[0] == target) return true;
return has(a + 1, n - 1, target)
```

 When you have a function returning a boolean using recursion above (returning the recursive call) make sure the return is not left off

```
    bool solve(start, goal) { // base cases: at the goal or surrounded by walls
    if (start == goal) return true;
    // mark this position as visited
    // for each direction {
    // if moving one step in that direction is possible and hasn't been visited {
    // if(solve(pos reached by moving, goal) return true;
    }
    return false;
```

2/10: Template Functions and Class Templates

```
• int minimum(int a, int b) {
       o if (a < b) return a;</p>
      o else return b;
• }
• double minimum(double a, double b) { // effectively the same code as the int minimum
   at the C++ language level

    if (a < b) return a;</li>

       o else return b;
• }
• int main() {
       o int k;
       o cin >> k;
       cout << minimum(k, 10) / 2;</li>
       o double x;
       o ...
      o double y = 3 * minimum(x * x, x + 10);
       o ...
       o int z = minimum(0, 3 * k - k + 4);
```

- template<typename T> // declare T to be used in template, T is now a placeholder for a type name
- T minimum(T a, T b) { // both functions fit into this template
 - if (a < b) return a;
 - o else return b;
- } // this function can replace both the int and double declarations above
- "Instantiate the template"
 - You have to match some template
 - The instantiated template compiles
 - The instantiated template has to do what you want
- You cannot mix types in the example above (passing an int and a double doesn't work)
- template<typename T1, typename T2>
- bool isEqualTo(T1 a, T2 b) { // this function can take 2 different types as parameters
 - return a == b;
- }

• }

- template<typename T1, typename T2>
- ??? minimum (T1 a, T2 b) { // no way to give a correct return type

```
if (a < b) return a;</li>else return b;
```

- Conversions:
 - Boolean compared to int/double → boolean converted to int/double
 - Int compared to double → int converted to double
- Chicken c1, c2;
- minimum(c1, c2); // how do you compare 2 chickens? fits template but probably won't compile
- char ca1[100];
- char ca2[100];
- **char* ca3 = minimum(ca1, ca2);** // minimum would compare the addresses of the pointers, not the values they point to, will compile, wouldn't do what you want
- Given a function vs. a template for a function, the language will choose the function
- Passing by const reference is common
- template<typename T>
- T sum(const T a[], int n) {
 - T total = 0; // won't compile with string call, no conversion between string and int, solution: builtInType() gives the "default" value, change to T total = T();

• }

• int main() {

```
double da[100];
...
cout << sum(da, 100);</li>
...
string sa[10] = {"This ", "is ", "a ", "test."};
string s = sum(sa, 4)
```

2/12: STL

<vector>

- push_back(n) // adds an index to the vector that contains n
- pop_back() // deletes the last index of the vector
- .size(), .front(), .back() // same as queues
- o [] // access elements, undefined for subscripts out of range
- o .at() // access elements, throws an exception for subscripts out of range
- vector<double> vd(10) // 10 doubles, each is 0.0
- vector<string> vs(10, "Hello") // 10 strings, each is "Hello"
- o vector<int> vx(a, a+5) // vx.size() is 5, v[0] is 10, vx[1] is 20, ..., vx[4] is 50
- vector<int>::iterator q = vi.insert(p, 40); // inserts 40 before p, returns iterator to new element
 - vector iterators support using +=/-= i
- p = vi.erase(q); // erases the value at q, returns iterator to new element in that spot

- o list<int> li;
- li.push_back(20);
- li.push_back(30);
- li.push_front(10);
- o cout << li.size(); // writes 3</p>
- cout << li.front(); // writes 10
- cout <back(); // writes 30
- li.push_front(40);
- o li.pop_front();
- o for (list<int>::iterator p = li.begin(); p != li.end(); p++)
 - **cout** << *p << endl; // writes 10 20 30
- o list<double> ld(10); // ld.size() is 10, each element is 0.0
- vector<string> ls(10, "Hello"); // ls.size() is 10, each element is "Hello"
- vector<string> vs(ls.begin(), ls.end()); // vs.size() is 10, each element is "Hello"
- o list<int>::iterator p = li.end();
- o p--;
- o p--;
- o p += 2; // won't compile
- list<int>::iterator q = li.insert(p, 40); // inserts 40 before p, returns iterator to new element
- list<int>::iterator q = li.erase(p); // erases element at p and using p is now undefined, returns iterator to element after

<set>

- o set<string> s;
- s.insert("Fred");
- s.insert("Ethel");
- s.insert("Fred");

```
o cout << s.size(); // 2</pre>
       s.insert("Desi");
       s.erase("Ethel");
       o for(set<string>::iterator ptr = s.begin(); p!= s.end(); p++) {
              cout << *p << endl; // prints using < (Desi is first)</p>
       0 }
• template<typename T>
• T* find(T* b, T* e, const T& target) {
       o for (; b != e; b++) {
              if (*b == target) break;
      0 }
      o return b;
• }
• int main() {
       o int a[5] = {10, 40, 50, 20, 30};

    int k;

      o cin >> k;
       o int* p = find(a, a + 5, k);
       \circ if (p == a + 5)
       o ... not found ...
       o else
       o ... found, p points to the first element with that value ...
       o string sa[4] = {"Lucy", "Ricky", "Fred", "Ethel"};
       o string* sp = find(sa, sa + 4, "Fred") // won't compile, "Fred" is a pointer to a
          char[]
• }
• template<typename Iter, typename T>
• Iter find(Iter b, Iter e, const T& target) {
       o for (; b != e; b++) {
              ■ if (*b == target) {
                     break;
              • }
      0 }
       return b;
• }
int main() {
       o int a[5] = {10, 40, 50, 20, 30};
      o int k;
       o cin >> k;
```

```
int* p = find(a, a+5, k);
       \circ if (p == a+5)
       o ... not found ...
       o else
       o ... found, p points to the first element with that value ...
       o list<string> ls;
       o list<string>::iterator q = find(ls.begin(), ls.end(), "Fred");
       o if (q == Is.end()) \
       o ... not found ...
       o vector<int> vi;
       vector<int>::iterator r = find(vi.begin(), vi.begin() + 5, 42);
      o if (r == vi.begin() + 5)
      o ... not found ...
• }
• template<typename Iter, typename Func>
• Iter find_if(Iter b, Iter e, Func f) {
       o for (; b != e; b++) {
              ■ if (f(*b)) break;
      0 }
       o return b;
• }
bool isNegative(int k) {
       ○ return k < 0;</li>
• }
bool isEmpty(string s) {
      return s.empty();
• }
• int main() {
       o vector<int> vi;
       0 ...
       vector<int>::iterator p = find_if(vi.begin(), vi.end(), isNegative); // function
          name by itself acts as a pointer to the function
       o if (p == vi.end())
       o ... not found ...
      o list<string> ls;
       o list<string>::iterator q = find_if(ls.begin(), ls.end(), isEmpty)
• }
```

```
2/24: Big-O
    • for (int i = 0; i < N; i++) { // O(N)
            o c[i] = a[i] * b[i]; // O(1)
    • }

    Work from the inside out

    • for (int i = 0; i < N; i++) { // O(N^2)
            o a[i] *= 2; // O(1)
            o for (int j = 0; j < N; j++) { // O(N)
                     \mathbf{d}[\mathbf{i}][\mathbf{j}] = \mathbf{a}[\mathbf{i}] * \mathbf{c}[\mathbf{j}]; // O(1)
            0 }
    • }
    • for (int i = 0; i < N; i++0 \{ // O(N) \}
            o a[i] *= 2; // O(1)
            o for (int j = 0; j < 100; j++) { // O(1) \rightarrow still a constant, not reliant on N
                     \bullet d[i][j] = a[i] * c[j]; // O(1)
            0 }
    • }
    • for (int i = 0; i < N; i++) { // O(N^2)
            o a[i] *= 2; // O(1)
            o for (int j = 0; j < i; j++) { // O(i) \rightarrow proportional to 0 + 1 + 2 + ... + (N-1) =
                (N-1)N/2
                     ■ d[i][j] = a[i] * c[j]; // O(1)
    • }
    • for (int i = 0; i < N; i++) \{ // O(N^2) \}
            o if (std::find(a, a+N, 10*i != a+N) { // takes time proportional to O(N)
                     ■ count++;
            0 }
    • }
    • for (int i = 0; i < N; i++) { // O(N^2 log N)
            o a[i] *= 2;
            o for (int j = 0; j < N; j++) { // O(N log N)
                     ■ d[i][j] = f(a, N); // suppose f(a, N) is O(log N)
            0 }
    • }
```

• for (int i = 0; i < N; i++) { // O(RC log C)

```
o a[i] *= 2;
       o for (int j = 0; j < N; j++) { // O(C log c)
              ■ d[i][j] = f(C); // suppose f(C) is O(log C)
       0 }
• }
• Selection Sort: O(N^2)
```

- Bubble Sort:
 - Worst Case: O(N^2) Average Case: O(N^2)
- Insertion Sort: $O(N^2) \rightarrow best of the N^2 sorts$
- Shell Sort: O(N^1.47) • Merge Sort: O(N log N)

3/2: Hash Tables

- Use a formula to set up a data structure that can contain a lot of data and is inexpensive to search/insert
- Always expect collisions
 - If load factor (# of items/# of buckets) > 1, there is guaranteed to be a collision
- Using a prime number of buckets allows there to be no common factor → more equal distribution
- Hash function
 - unsigned int hash(string s) { // numbers greater than the max size for unsigned ints are just represented by their last 32? digits
 - unsigned int h = 216613261U; // U makes the compiler treat it as an unsigned integer constant

- Hash table with N items: lookup/insert/erase for fixed # of buckets: O(N) with low constant of proportionality
- Hash table with N items + rehashing: lookup/insert/erase: O(1) on average
- unordered_set
- unordered_multiset // allows duplicates
- unordered map
- unordered_multimap // allows duplicate keys
- #include <unordered map>
- using namespace std;
- unordered_map<string, double> ious;
- string name;
- double amt;
- while (cin >> name >> amt) {
 - o ious[name] += amt;
- }
- if (ious.find("ricky") == ious.end()) {
 - cout << "Ricky doesn't owe me anything" << endl;
- }
- #include <functional>
- #include <string>

- string s;
- cin >> s;
- unsigned int h = std::hash<std::string>()(s);
- unsigned int h2 = std::hash<double>()(3.75);
- unsigned int h3 = std::hash<string*>()(&s);

3/4: Binary Trees

- Made up of nodes connected by edges, stemming from a root node, and paths that connect one node to another
 - o There is a unique path from the root node to any other node in the tree
 - Nodes follow a parent/child relationship
 - A leaf node has no children, interior nodes do
 - The depth of a node is determined by how many edges away from the root node a node is
 - Any program with a tree that goes down multiple paths is most likely best implemented with recursion

```
struct Node {
      string data;
      vector<Node*> children;
      o Node* parent; // not always necessary, returning to parent may or may not be
          relevant
 };
Node* root;
• int count(const Node* p) { // base cases: empty tree and leaf nodes
      o if (p == nullptr) {
             ■ return 0;
      0 }
      o int total = 1; // count self
      o for (int k = 0; k != p->children.size(); k++) {
             total += count (p->children[k]);
      0 }
      return total;
cout << count(root) << endl;</li>
void printTree(const Node* p, int depth) {
      o if (p == nullptr) {
             return;
      o cout << string(2 * depth, ' ') << p->data << endl;</pre>
      o for (int i = 0; i < p->children.size(); i++) {
             printTree(p->children[i], depth + 1);
      0 }
• }
```

- Preorder traversal: process the node before you process the children
- Postorder traversal: process the node after you process the children

```
• A binary tree is either:
       Empty

    Or is a node with 2 binary subtrees (a left subtree and a right subtree)

struct Node {
       string data;
       Node* left;
       Node* right;
};
• A binary search tree (BST) is either:
       Empty
       o A node with a left BST and a right BST such that the value at every node on the
          left subtree is < the value at this node, and every node in the right subtree is >
          the value at this node

    Assuming no duplicates

set<string> visitors;
string name;
• visitors.insert(name) // inserts if name is not already in the set
if (visitors.find("Ricky") == visitors.end();
       cout << "not found" << endl;</li>
for (set<string>::iterator p = visitors.begin(); p!= visitors.end(); p++) {
       cout << *p << endl;</li>
• }
set<string>::iterator p = visitors.find("Fred");
• if (p != visitors.end()) {
       cout << "Bye, Fred" << endl;</li>
       visitors.erase(p) // harmless to pass erase the end iterator
• }
visitors.remove("Lucy");
• In-order traversal:
       void printTree(const Node* p) {
              ■ if (p!= nullptr) {
                     printTree(p->left);
                     cout << p->data << endl;</li>
                     printTree(p->right);
```

• }
• }

- Insertion/deletion/search O(log N) in a relatively balanced BST
- Deletion:
 - Leaf node delete the node
 - o Node with one child delete the node, move the child up
 - Node with two children delete the node and take either the greatest value in the left subtree or the least value in the right subtree to replace it
- AVL Tree:
 - A BST where for every node, the height of the left subtree and the height of the right subtree differ by at most 1
 - Guarantees log N behavior
- 2-3 Tree:
 - Every node has 2 or 3 children
 - Nodes with 3 children contain 2 values, the 3rd child is between these 2 values
 - All leaf nodes have the same depth
- Red-black tree:
 - Used by STL for set, map, multiset, multimap
 - o log N data structures