

Cyber Physical System Security

SC4015/CE/CZ4055

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CPSSECURITY



Contents

➔ *Discussions from Last Week*



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Automata and *Examples*

Automata theory

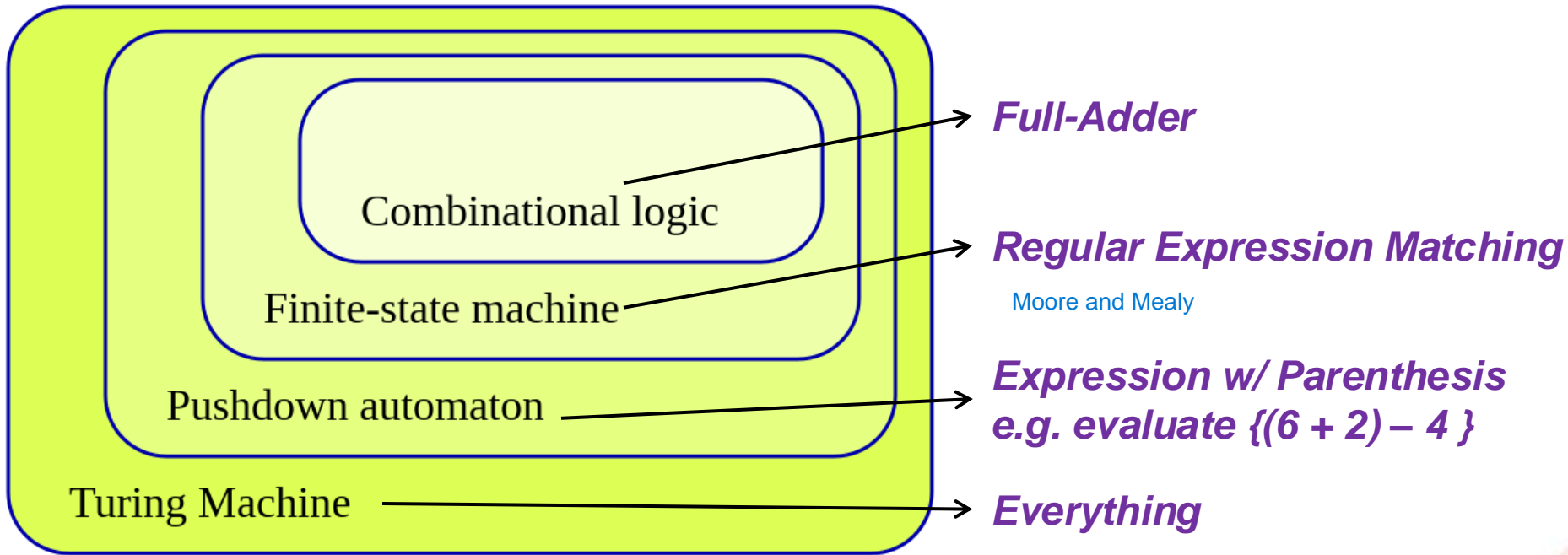


Image source: wikipedia

Trusted Service Manager

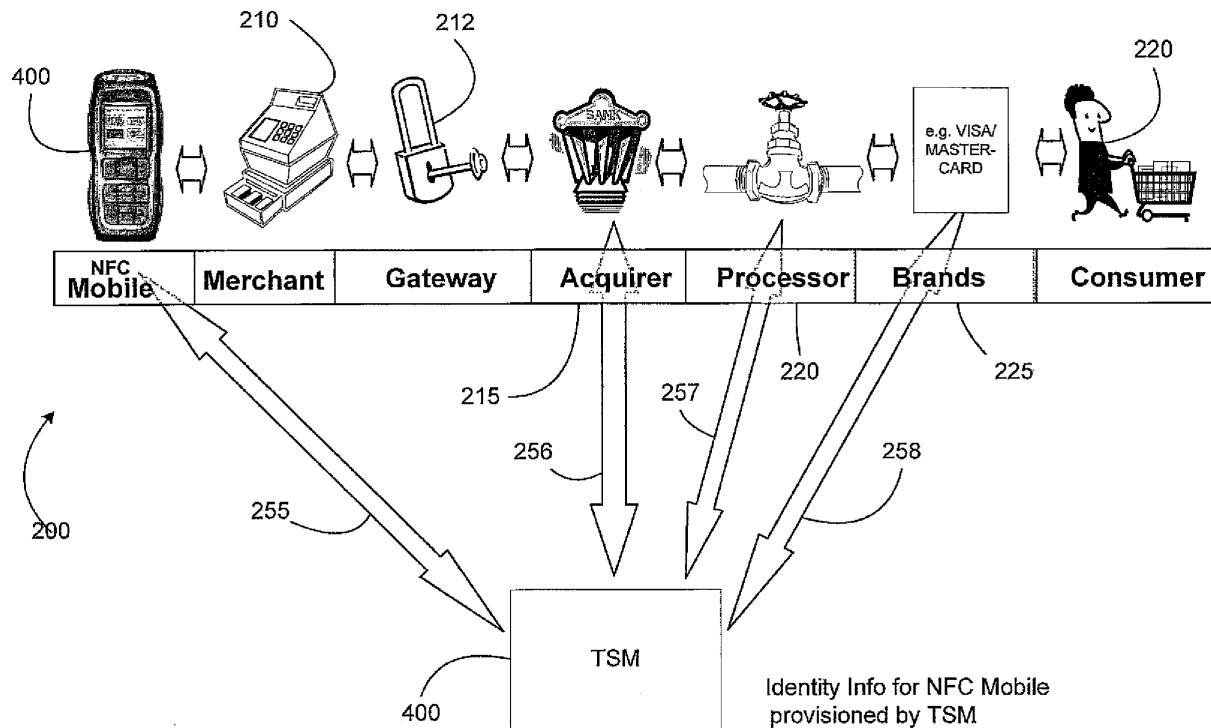


FIG. 2

Trusted Service Manager

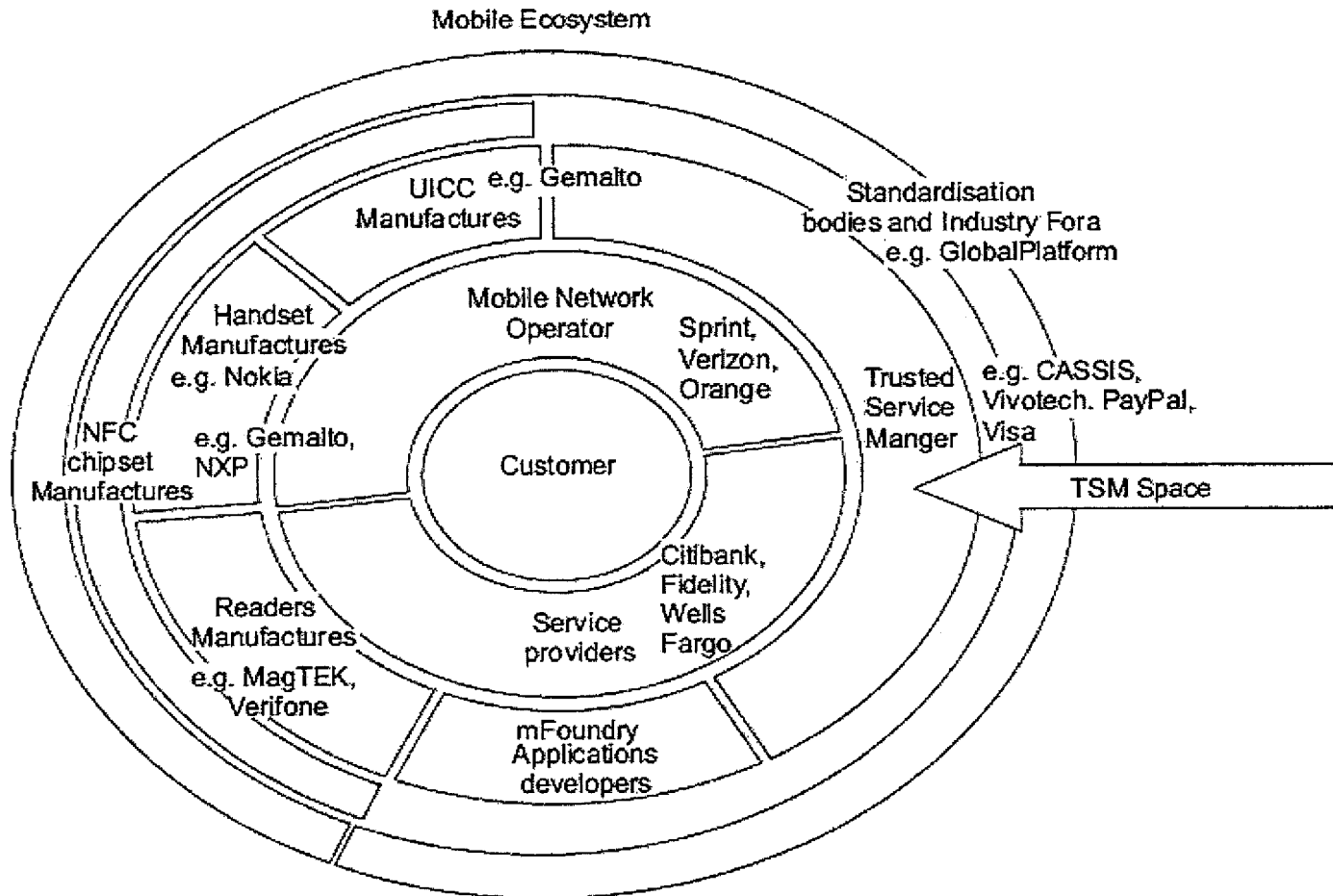


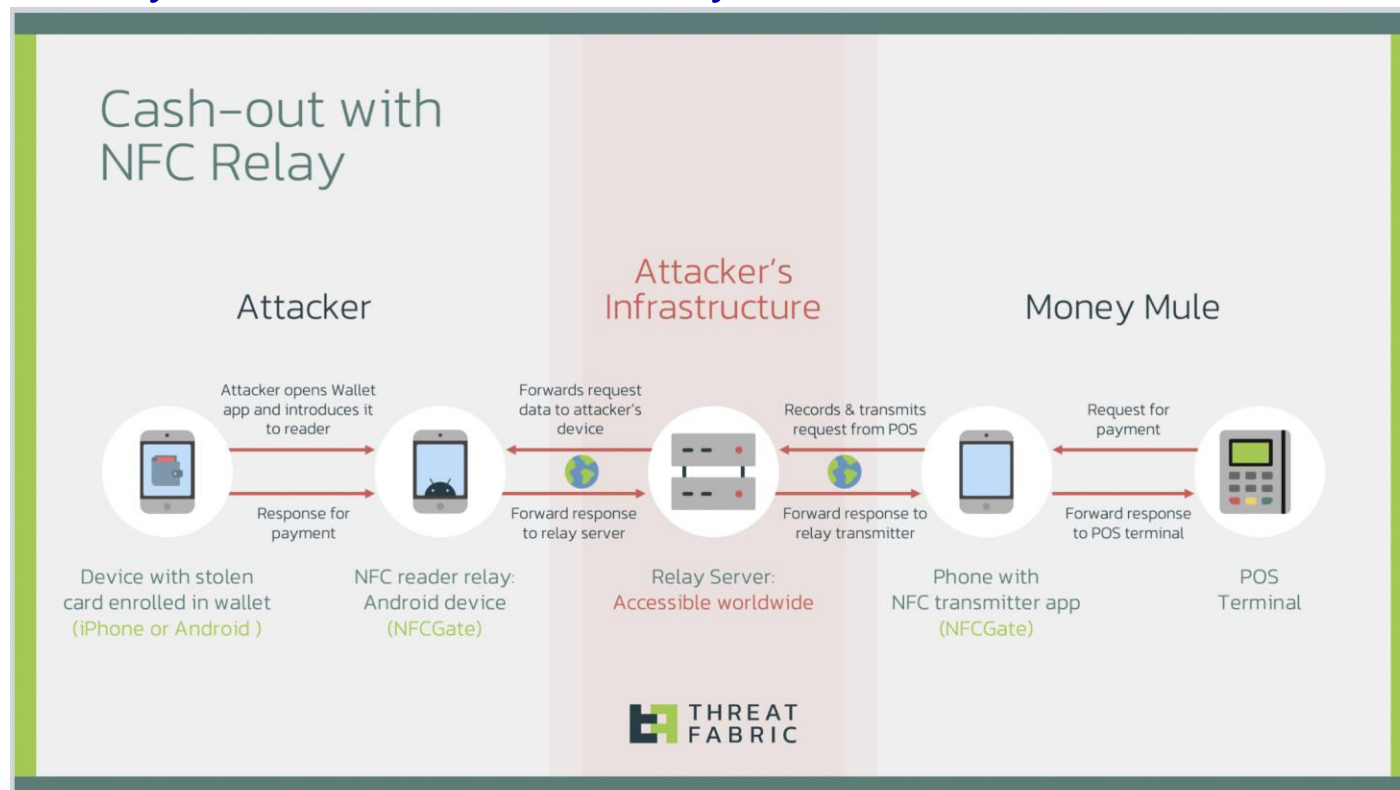
FIG. 1

Trusted Service Manager: Tasks

1. **User registration:** Corresponding to a bank
 2. **Managing Digital Certificates:** User identity
 3. **Authentication and Verification:** Transactions of the mobile wallet with a payment terminal
- The above services are delivered using the underlying platform, in particular
 - **SE:** Storage of keys, passwords, identity
 - **Secure Communication:** NFC
 - **TEE:** Secure Microprocessor for providing security services (e.g., encryption, authentication)

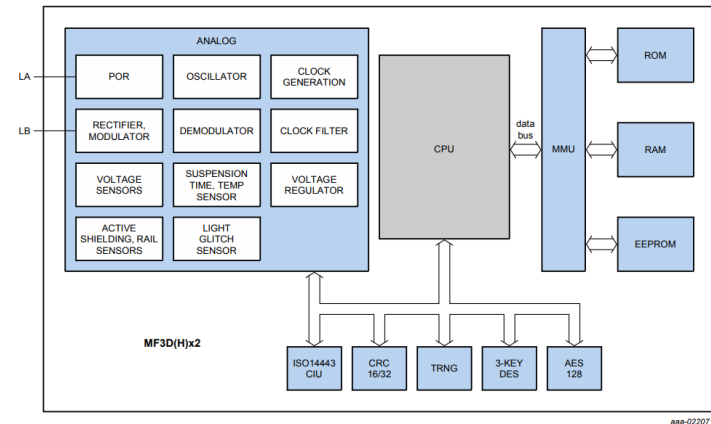
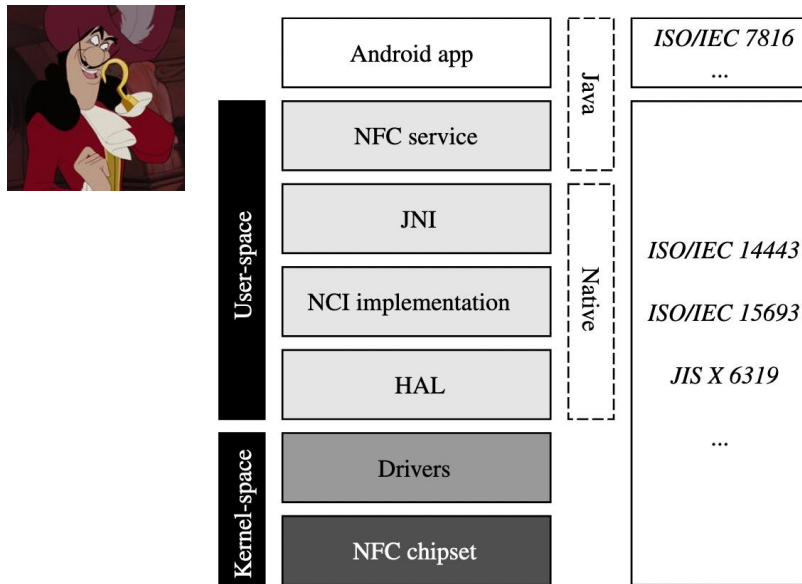
NFC Security (1/2)

- NFCGate: Android toolkit for NFC traffic (and security) analysis
- Used by hackers to steal money



NFC Security (2/2)

- Through *Java symbolic hooks* → from other applications running in the background of the same platform
- Through *native instruction hooks* → alter the underlying platform behavior



Reverse engineer packets, as NFC IC is known


NFC Security: Countermeasures

- Preventing Relay Attack
- Do not let *untrusted apps* running in the background
- Distance Bounding Protocols
 - To ensure that the endpoints of a communication are within a specified distance
- Mandatory Response Timeouts
 - To ensure that the reader ends a communication when there is a significant time-lag (introduced due to the relay)

Contents

➔ Security Triad

- Cryptographic Primitives
- Discussion



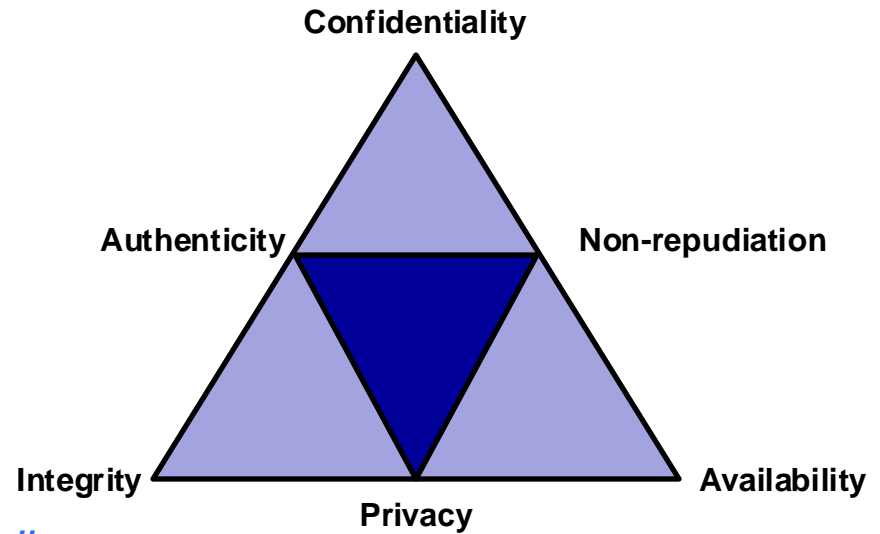
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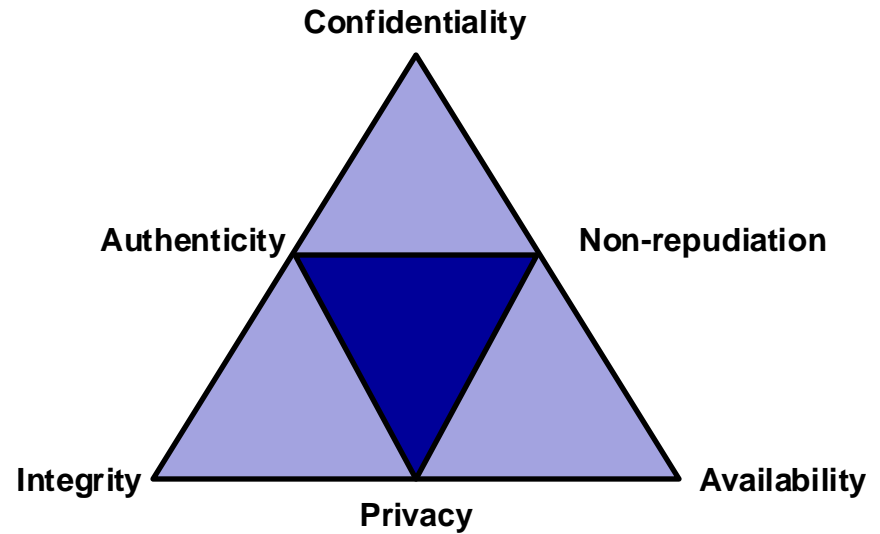


Security Triad

- IT security
 - Confidentiality
 - *Data is secured*
 - Integrity
 - *Data is trusted*
 - Availability
 - *Data is accessible*
 - Non-repudiation
 - *Service has a trusted audit-trail*
 - Authenticity
 - *Components have provable identity*
 - Privacy
 - *Service does not see customer activities*

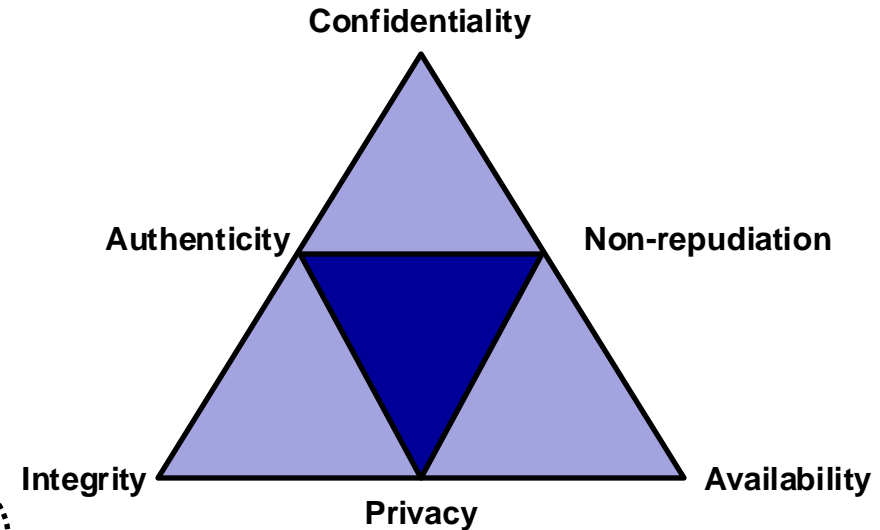
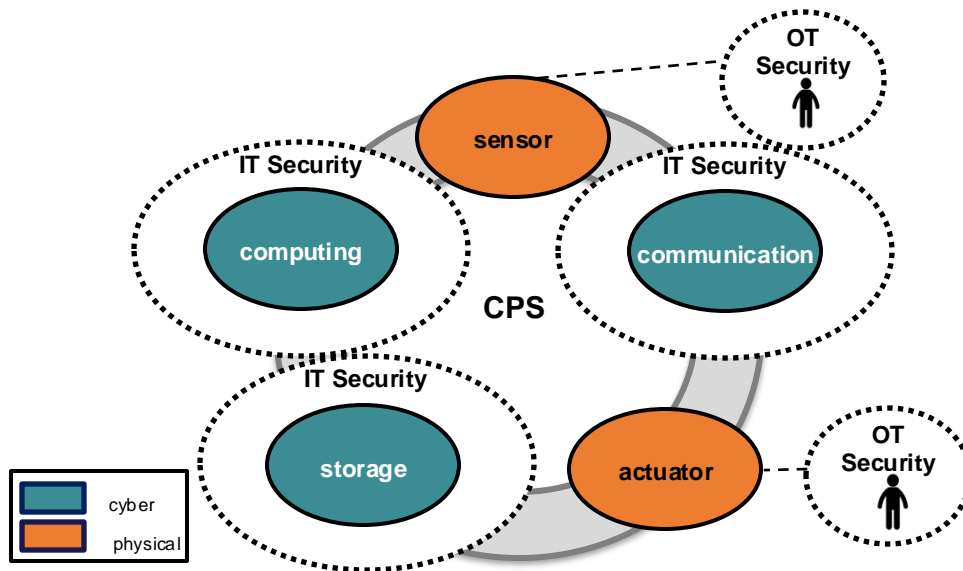


Security Triad: Smart Card

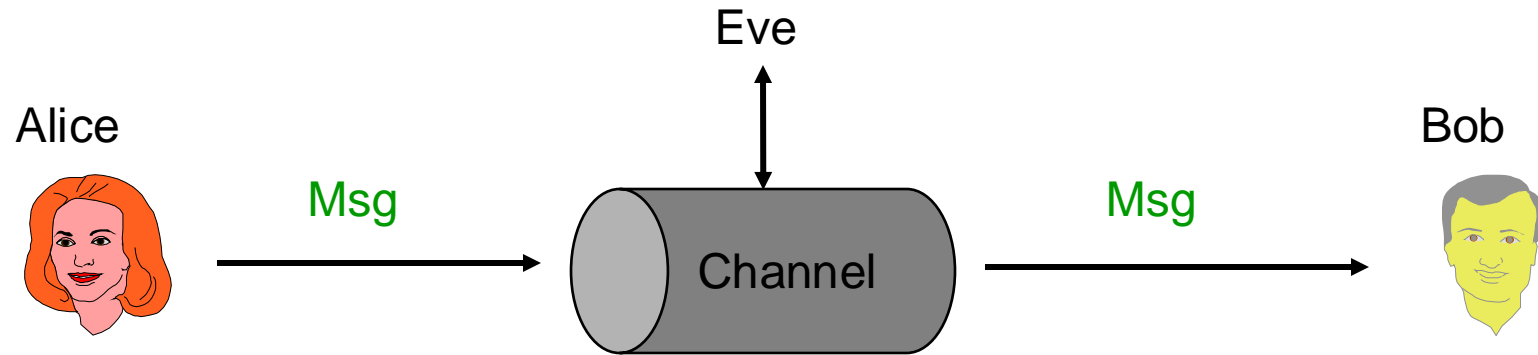


- Smart Card security
 - Confidentiality: *Data is secured*
 - Integrity: *Data is trusted*
 - Availability: *Data is accessible*
 - Non-repudiation: *Service has a trusted audit-trail*
 - Authenticity: *Components have provable identity*
 - Privacy: *Service does not see customer data*

Security Triad for CPS



Network Security Model

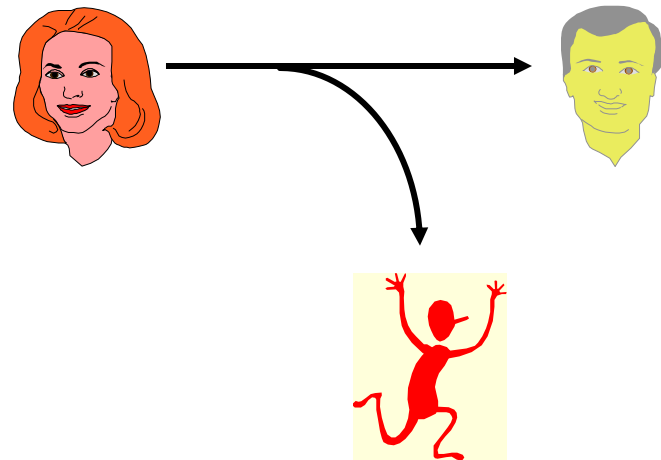


Confidentiality (Secrecy)

- Protect transmitted data
- Protect against traffic analysis



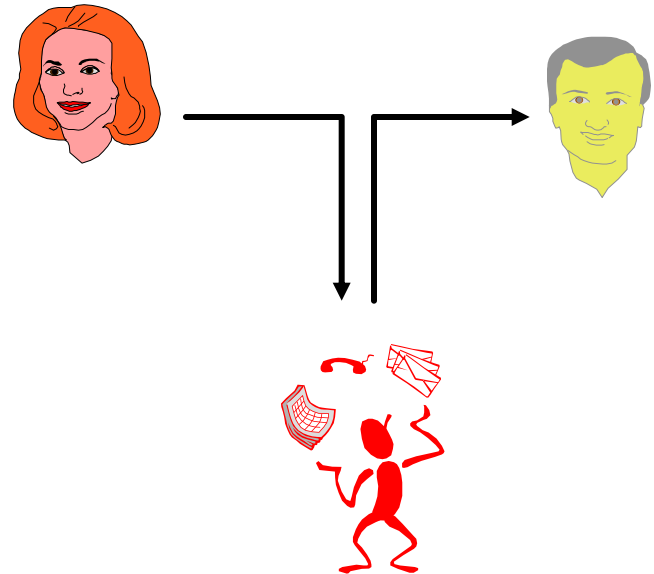
- **INTERCEPTION**
Unauthorised party gains access to data



Integrity

- Message is received as sent
- Modification
- Also interested in replay, re-ordering, deletion, delay

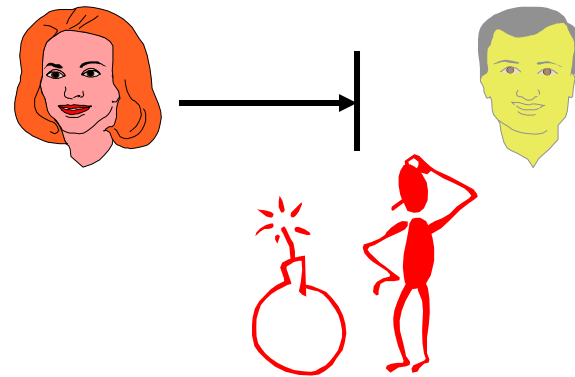
- **MODIFICATION**
Gain access and “tamper” with messages



Availability

- Complete loss of availability
- Reduction/Degradation in availability

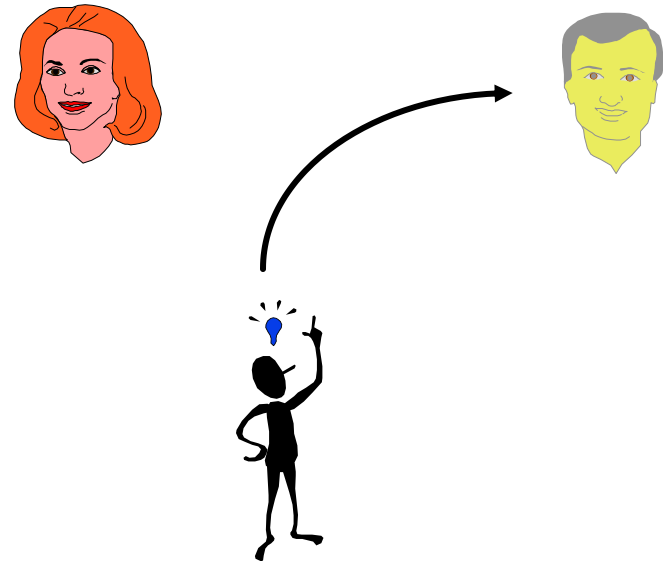
- **INTERRUPTION**
Loss of communication (cut the cable)
- **DENIAL OF SERVICE**
Noisy comms (physical noise, spurious messages)



Authentication

- Assurance that message is from proper source
- Protect from third party masquerade

- **FABRICATION**
Insertion of “counterfeit” messages



Non-repudiation

- Prevents parties from denying they sent or received a message; i.e. concerned with protecting against legitimate protocol participants, not with protection from external source
- Receiver can verify **and prove** who sent a message
- Sender can verify **and prove** who received a message

- **REPUDIATION ATTEMPT**
Party anonymously publishes his or her message/key(s) and falsely claims that they were stolen.

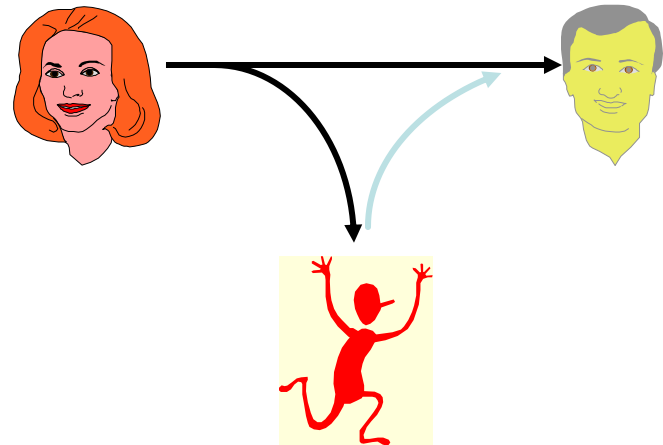


Privacy

- Limit and control access to host system/services
- Limit and control access to networks
- Authenticate each party so that access rights can be assigned
- More fine-grained solutions, e.g. Digital Rights Management

- **REPLAY**

Record a legitimate message e.g. a login, and replay later



Contents

- Security Triad

➔ *Cryptographic Primitives*

- ***Private-Key Cryptography***
 - Hash Function, Message Authentication Code
 - Public-Key Cryptography
 - Digital Signature
- Discussion



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Kerckhoffs's Principle

- An encryption scheme should be secure even if enemy knows everything about it except the key
 - Attacker knows all algorithms
- Do not rely on secrecy of the algorithms (“security by obscurity”)



SYMMETRIC-KEY CRYPTOGRAPHY

Symmetric-key cryptography started thousands of years ago when people needed to exchange secrets (for example, in a war).



7th Century BC, Ancient Greece

Types of Symmetric-Key Cipher

- **Substitution Cipher**

- A substitution cipher replaces one symbol with another.
- Can be mono/poly-alphabetic

The following figures show two examples of a plaintext and its corresponding ciphertext. Which cipher is mono-alphabetic?

Plaintext: HELLO
Ciphertext: KHOOR

Plaintext: HELLO
Ciphertext: ABNZF

Types of Symmetric-Key Cipher *(contd.)*

- **Shift Cipher**
 - The shift cipher is also referred to as the Caesar cipher.

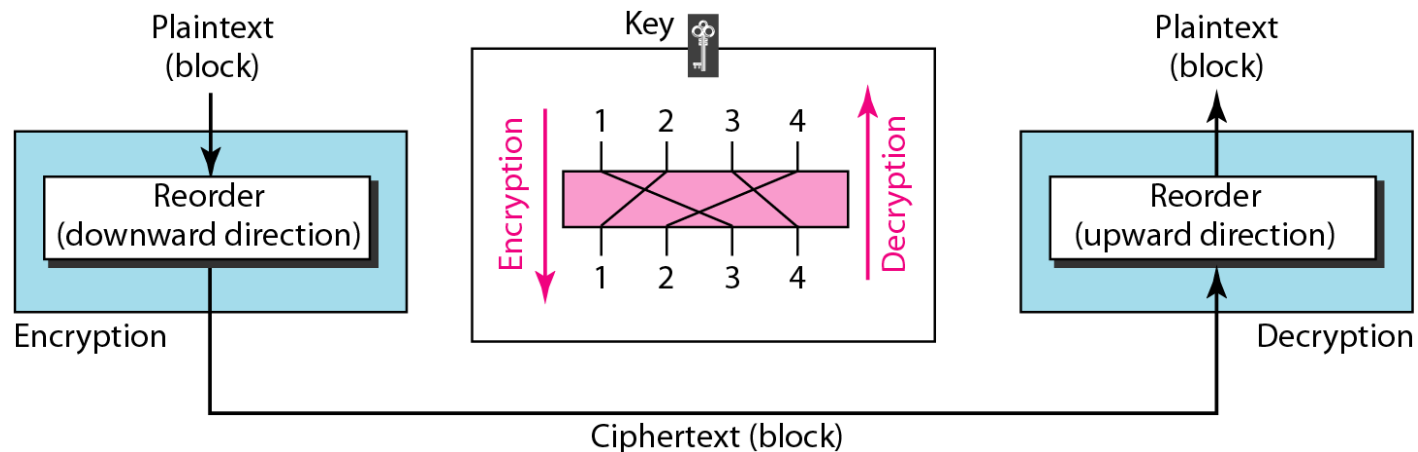
Use the shift cipher with key = 15 to encrypt the message “HELLO.”

Solution

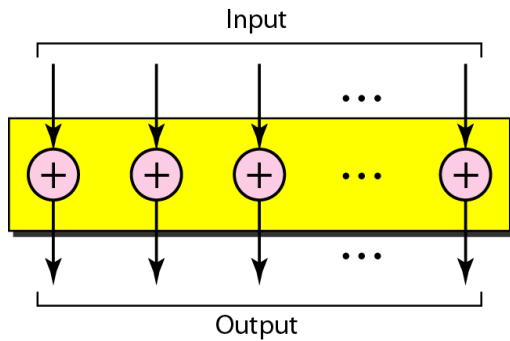
We encrypt one character at a time. Each character is shifted 15 characters down. Letter H is encrypted to W. Letter E is encrypted to T. The first L is encrypted to A. The second L is also encrypted to A. And O is encrypted to D. The cipher text is **WTAAD**.

Types of Symmetric-Key Cipher *(contd.)*

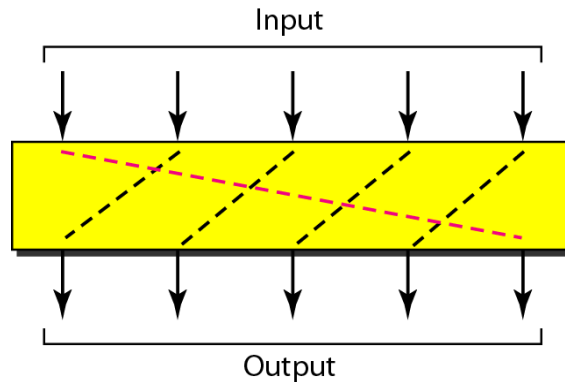
- **Transposition Cipher**
- A transposition cipher reorders (permutes) symbols in a block of symbols.



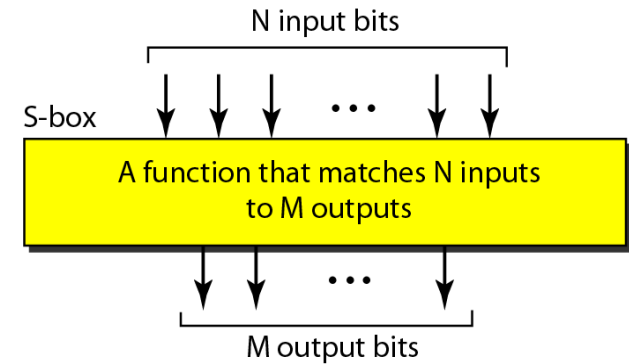
Components of Symmetric-Key Cipher



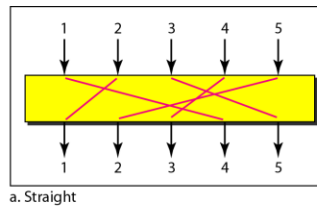
XOR substitution



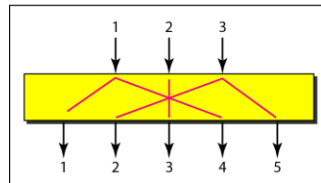
Rotation substitution



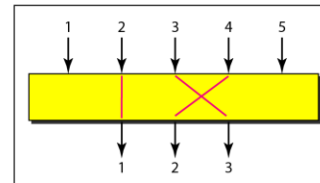
Substitution Box



a. Straight



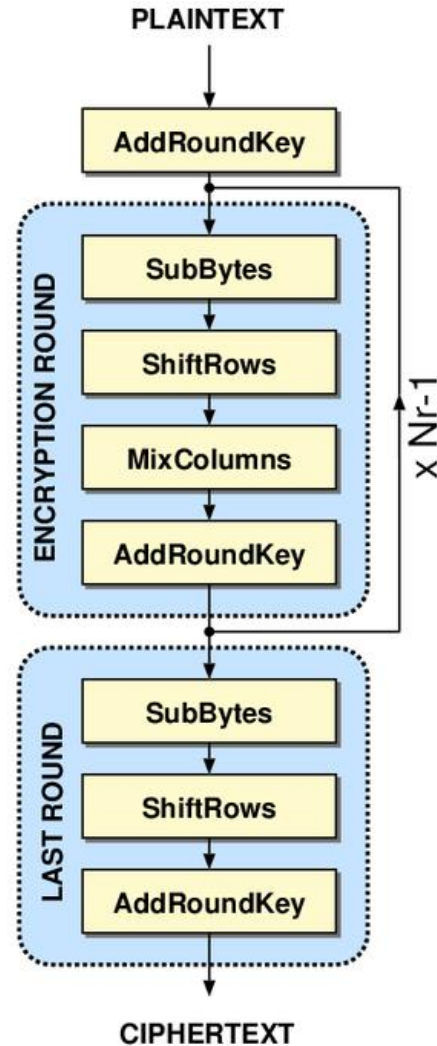
b. Expansion



c. Compression

Permutation Box

Block Cipher: AES



One-Time Pad (Vernam Cipher)



= 10111101...



= 00110010...

Key is a random bit sequence
as long as the plaintext

Encrypt by bitwise XOR of
plaintext and key:
 $\text{ciphertext} = \text{plaintext} \oplus \text{key}$



10001111...



10111101...

00110010...

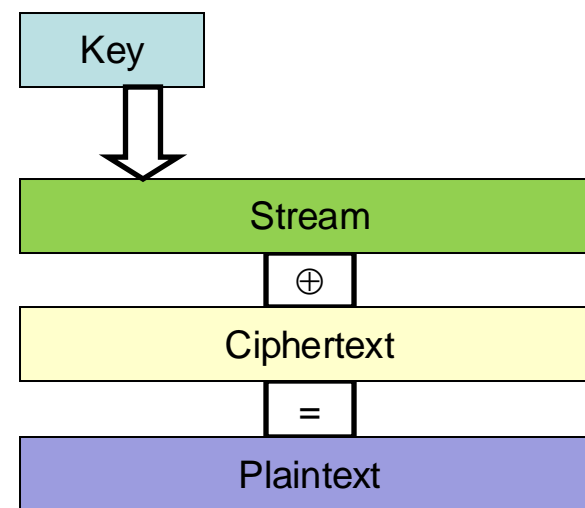
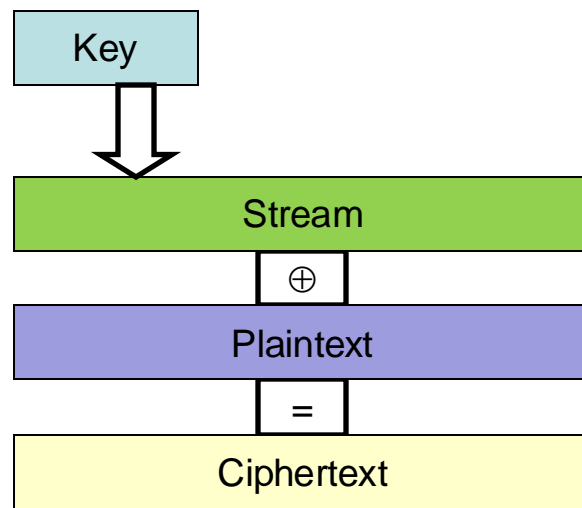


Decrypt by bitwise XOR of
ciphertext and key:
 $\text{ciphertext} \oplus \text{key} =$
 $(\text{plaintext} \oplus \text{key}) \oplus \text{key} =$
 $\text{plaintext} \oplus (\text{key} \oplus \text{key}) =$
 plaintext

Cipher achieves perfect secrecy if and only if there are as many possible keys
as possible plaintexts, and every key is equally likely (Claude Shannon, 1949)

Stream Ciphers

- Start with a secret key (“seed”)
- Generate a keying stream
 - Uses a Pseudo-Random Number Generator (PRNG)
- Combine the stream with the plaintext to produce the ciphertext (typically by XOR)



Pseudo Random Number Generator

- Commonly constructed using Linear Feedback Shift Registers (LFSRs)
- An adversary cannot distinguish a PRNG output from a random sequence
- Example Randomness Tests
 - Monobit tests (frequency of 0s and 1s)
 - Run tests (frequency of runs of different lengths)
- *Pi: A naturally occurring random number*

String 1: 01
String 2: 1100100001100001110111101110110011111010010000100101011110010110

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How to Authenticate Messages?

- content is authentic – bits are as sent
- sequence of messages is proper

Note: Separate problems

- confidentiality that the message is encrypted
- authenticity that the sender is genuine

Authenticating messages - Types

- **Authenticator** - a value that authenticates a message content
- **Message Authentication Code, MAC** (cryptographic checksum)
 - public function, with a given secret key produces fixed length value
- **Hash or Message Digest**
 - public function, which maps message (any length) to fixed-length hash value – no key

Message Authentication Code (MAC)

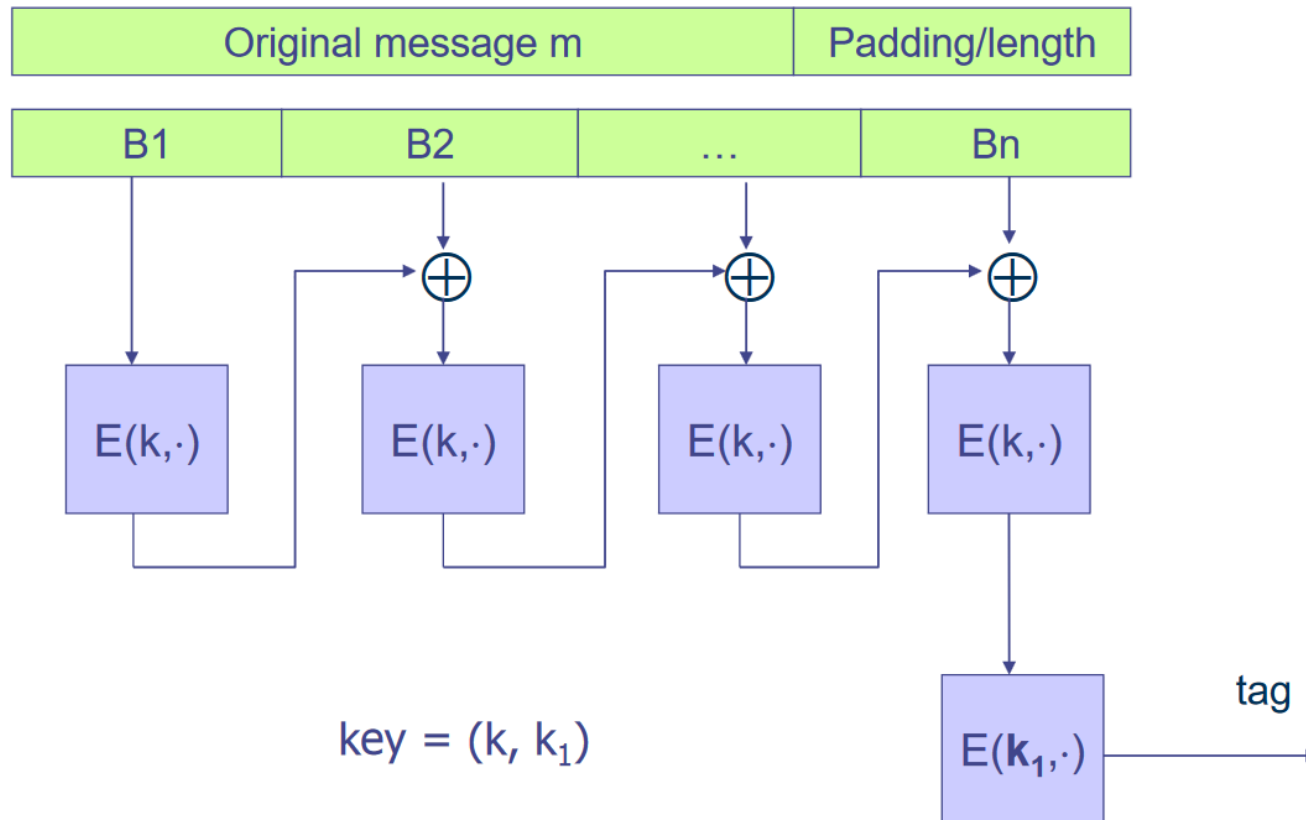
- Objective: assure authentic content
 - Message is not encrypted
- MAC is small fixed-size data block, appended to message M
- Generator and Validator share secret key

$$\text{MAC}_K[M] = F(M, \text{key})$$

Why use MAC, given encryption?

- Message may not need to be secret, but must be authentic
- Broadcast - only have one site responsible for monitoring authenticity. Broadcast plaintext plus MAC
- Overload - send plaintext messages (i.e. receiver need not decrypt). Authenticate selectively

MAC Construction using Encryption



Cipher Block Chaining Message Authentication Code (CBC-MAC)

MAC usage

- message authentication (no confidentiality)
 - $A \rightarrow B$: $[M, \text{MAC}_K[M]]$
- authentication and confidentiality
 - $A \rightarrow B$: $E_{K_1}[M, \text{MAC}_K[M]]$
- authentication and confidentiality
 - $A \rightarrow B$: $[E_{K_1}[M], \text{MAC}_K[E_{K_1}[M]]]$

MAC Usage (contd.)

Encryption key K_E MAC key = K_I

Option 1: MAC-then-Encrypt (SSL)



Option 2: Encrypt-then-MAC (IPsec)



Option 3: Encrypt-and-MAC (SSH)



Hash function (Message Digest)

- no key
- like MAC, small amount of data; hash of message gives fixed-size value
- define hash function so that change of any one bit of message will result in different hash value
- hash function is not secret
- one-way; receiver re-computes hash function

Simple Message Digest: Parity

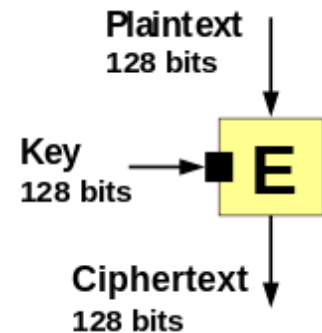
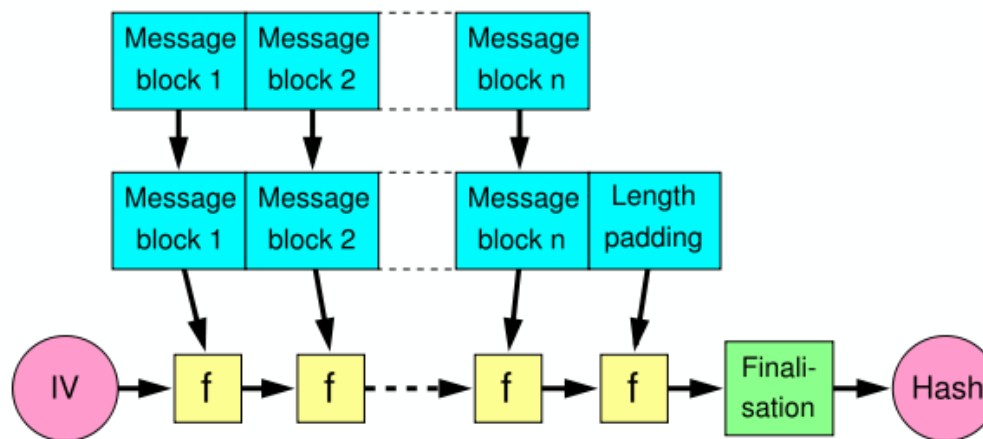
- parity, or longitudinal redundancy check
- process one bit at a time
- use XOR

$$C = b_1 \oplus b_2 \oplus \cdots \oplus b_{n-1} \oplus b_n$$

- input – arbitrary length message
- size of processed block – one bit
- output – one bit

Merkle–Damgård Construction for Hash Functions

- Message is divided into fixed-size blocks and padded
- Use a one-way function f , which takes a chaining variable (of size of hash output) and a message block, and outputs the next chaining variable
- Final chaining variable is the hash value



f built out of block cipher

$$M = m_1 m_2 \dots m_n; C_0 = IV, C_{i+1} = f(C_i, m_i); H(M) = C_n$$

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- Security Triad
- Cryptographic Primitives

➔ *Discussion*

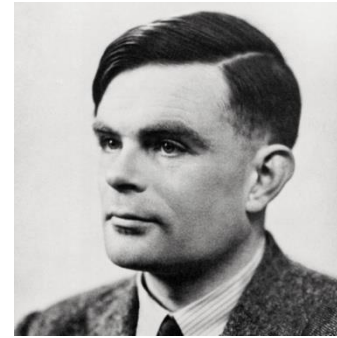


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What did we learn?

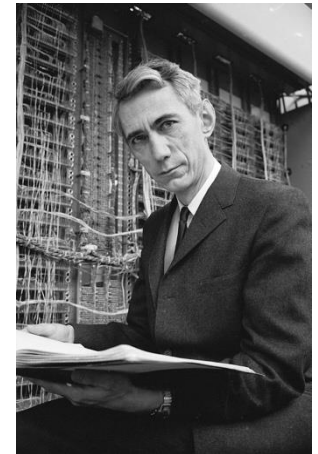
- What is Security?



- Confidentiality, Integrity, Availability
- *Authenticity, Privacy, Non-repudiation*

- What are the building blocks of security?

- *Upcoming Lectures*
- Private-Key Cryptography, Hash, MAC
- Public-Key Cryptography, Digital Signature



Further Reading

- Security Engineering
 - by Ross Anderson, available online - <http://www.cl.cam.ac.uk/~rja14/book.html>
- Handbook of Cryptography
 - by Alfred J. Menezes, Paul C. van Oorschot and Scott A. Vanstone, available online - <http://cacr.uwaterloo.ca/hac/>
- Applied Cryptography
 - by Dan Boneh and Victor Shoup, available online - <http://toc.cryptobook.us/>
- Leisure Reading - Simon Singh: *The Code Book*, Fourth Estate 1999

The End