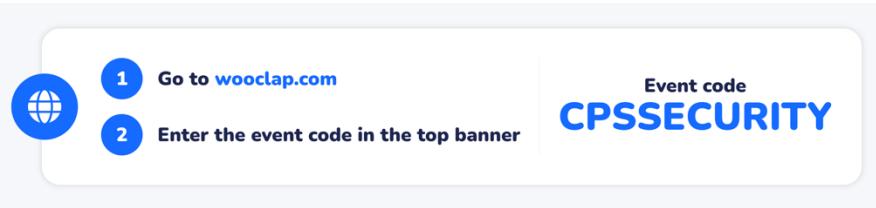


# **Cyber Physical System Security**

## **SC4015/CE/CZ4055**

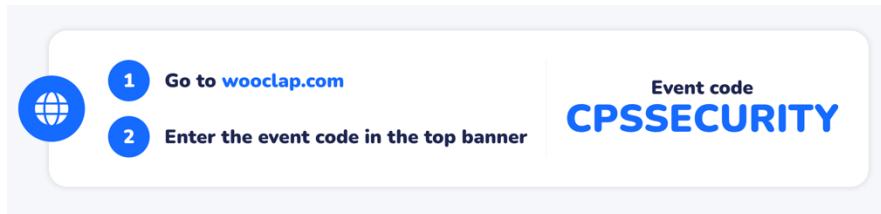
Anupam Chattopadhyay  
CCDS, NTU



# Contents



*Discussions from Last Week*



# Automata and Examples

Automata theory

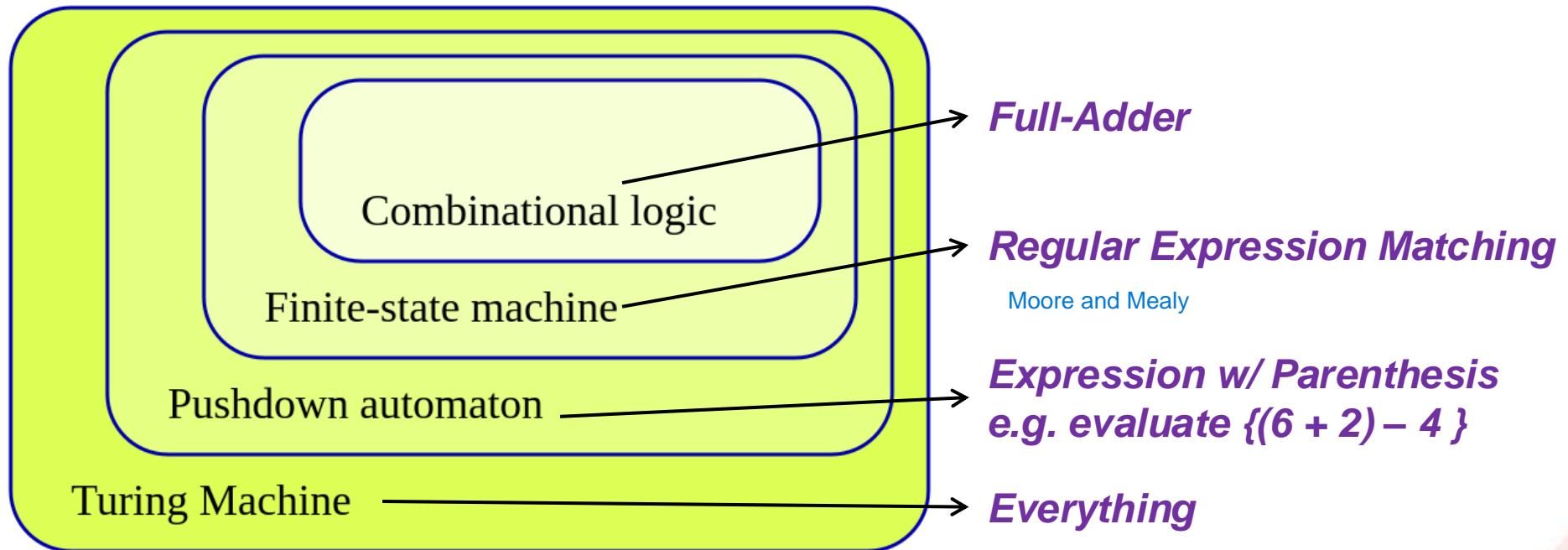


Image source: wikipedia

# Trusted Service Manager

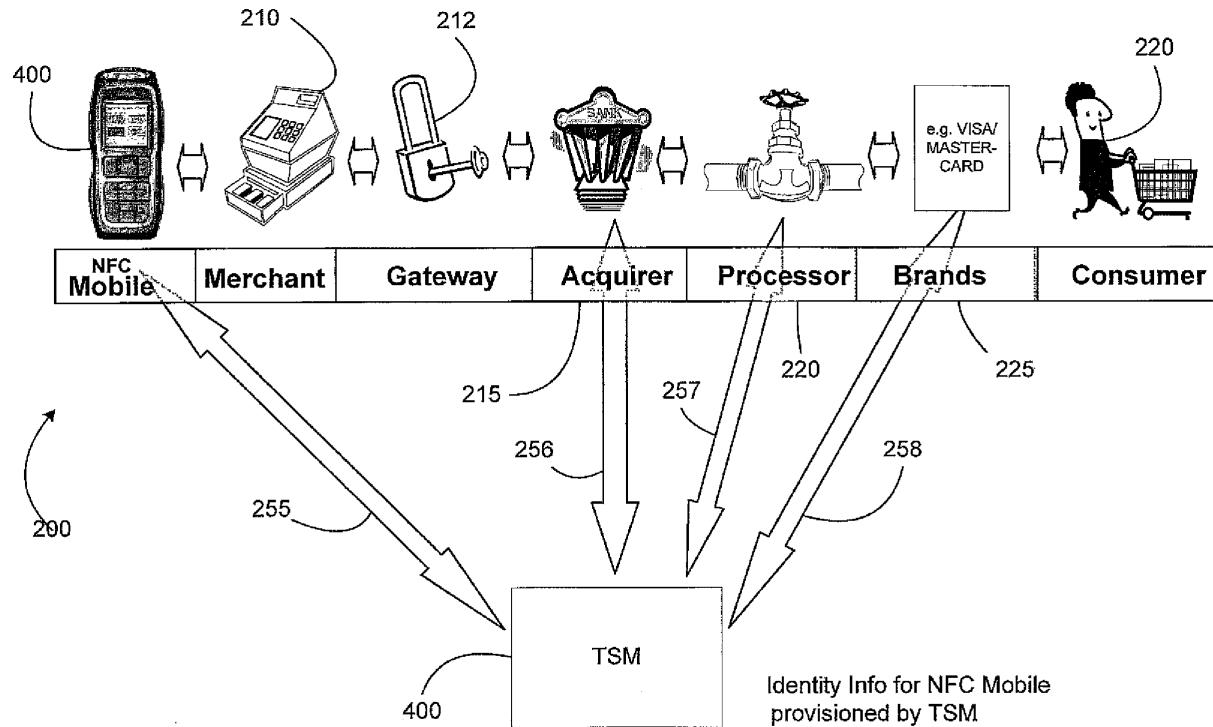


FIG. 2

# Trusted Service Manager

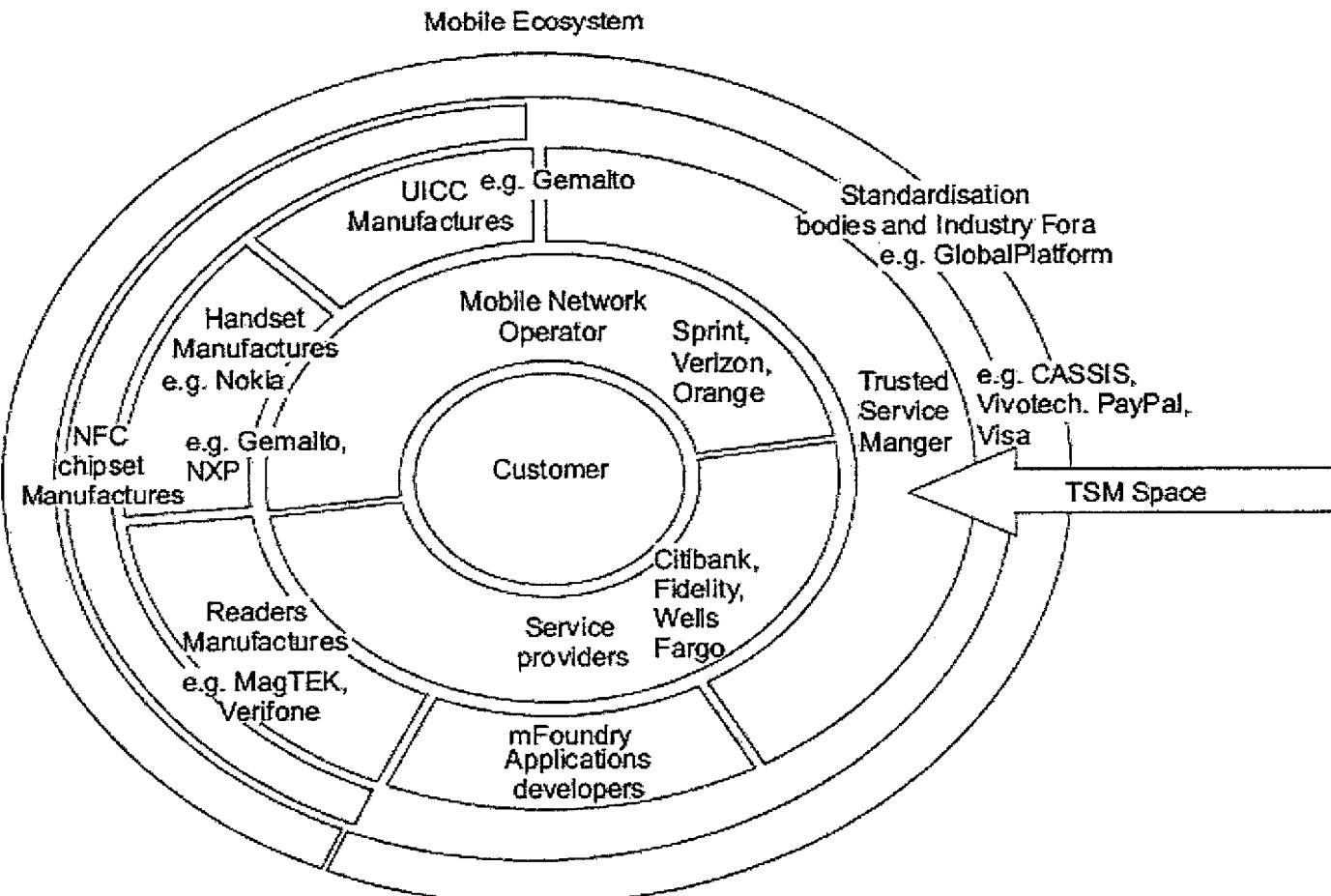


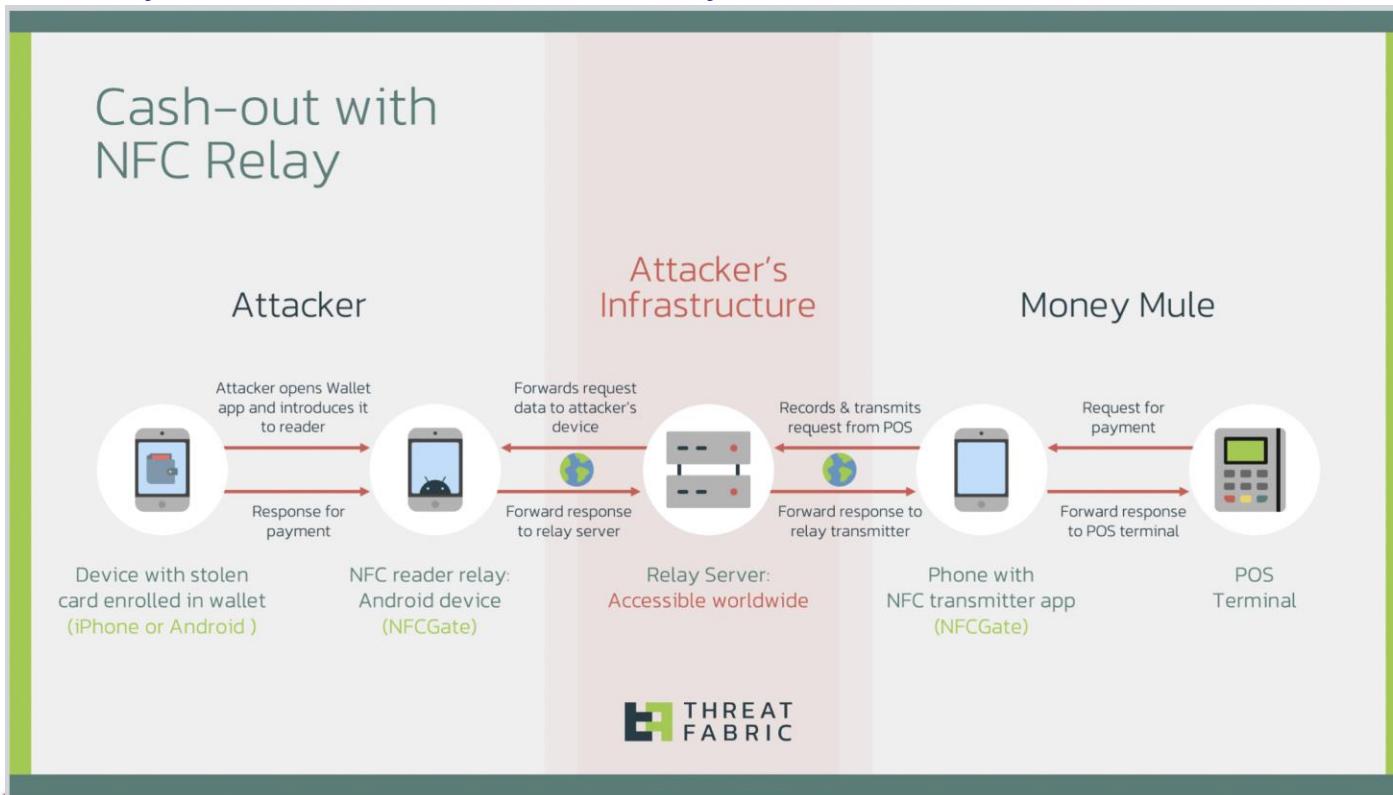
FIG. 1

# Trusted Service Manager: Tasks

1. **User registration:** Corresponding to a bank
  2. **Managing Digital Certificates:** User identity
  3. **Authentication and Verification:** Transactions of the mobile wallet with a payment terminal
- 
- The above services are delivered using the underlying platform, in particular
    - **SE:** Storage of keys, passwords, identity
    - **Secure Communication:** NFC
    - **TEE:** Secure Microprocessor for providing security services (e.g., encryption, authentication)

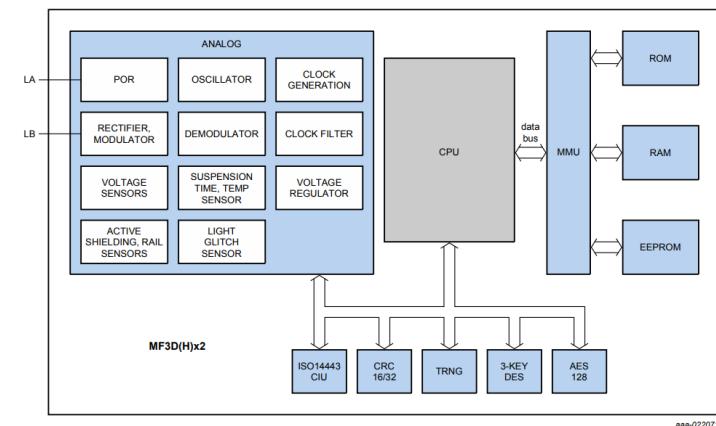
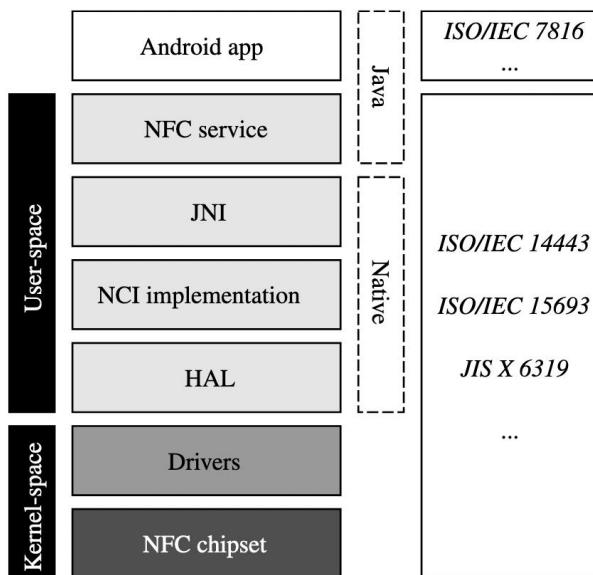
# NFC Security (1/2)

- NFCGate: Android toolkit for NFC traffic (and security) analysis
- Used by hackers to steal money



# NFC Security (2/2)

- Through *Java symbolic hooks* → from other applications running in the background of the same platform
- Through *native instruction hooks* → alter the underlying platform behavior



Reverse engineer packets, as NFC IC is known

# NFC Security: Countermeasures

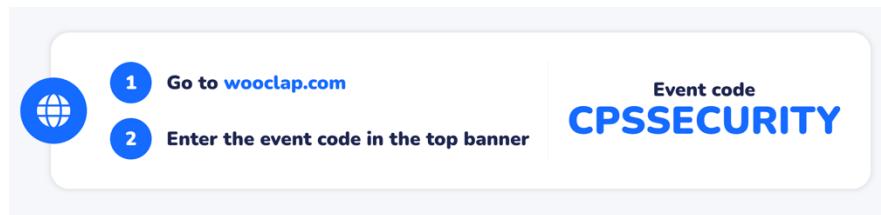
- Preventing Relay Attack
- Do not let *untrusted apps* running in the background
- Distance Bounding Protocols
  - To ensure that the endpoints of a communication are within a specified distance
- Mandatory Response Timeouts
  - To ensure that the reader ends a communication when there is a significant time-lag (introduced due to the relay)

# Contents



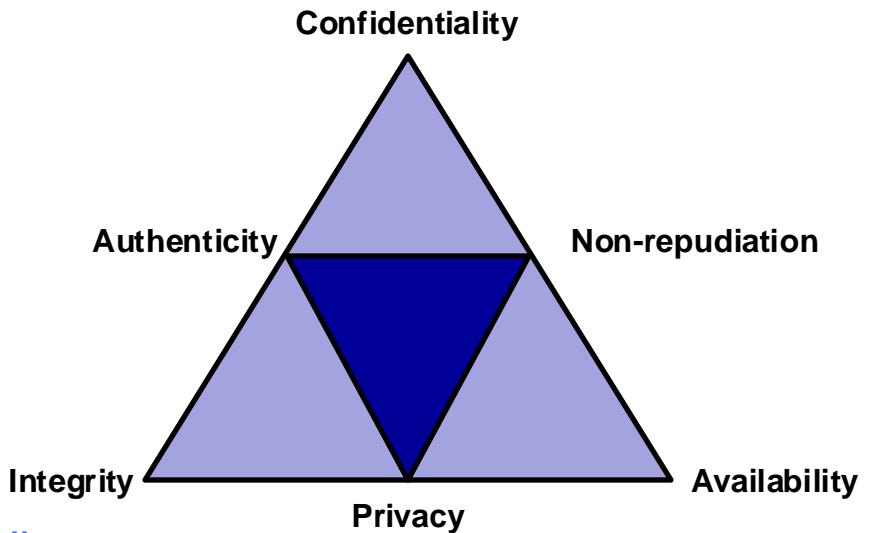
## *Security Triad*

- Cryptographic Primitives
- Discussion

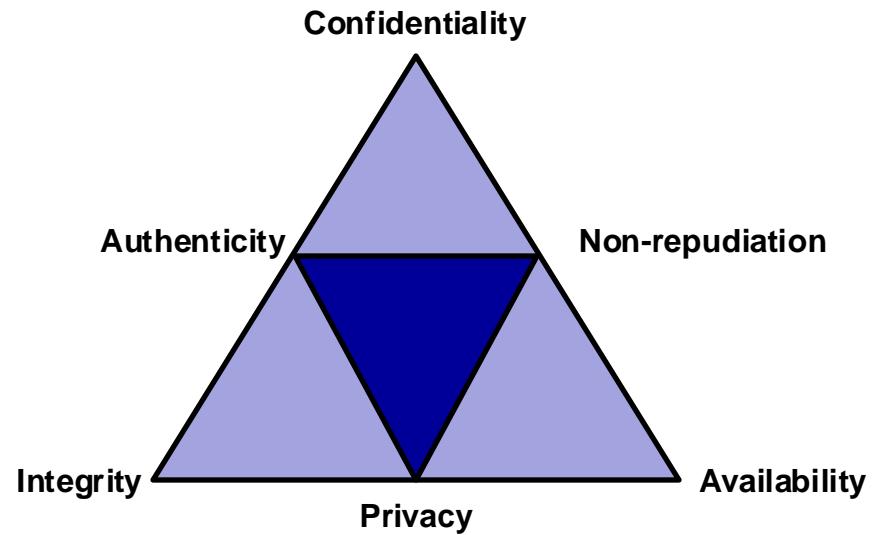


# Security Triad

- IT security
  - Confidentiality
    - *Data is secured*
  - Integrity
    - *Data is trusted*
  - Availability
    - *Data is accessible*
  - Non-repudiation
    - *Service has a trusted audit-trail*
  - Authenticity
    - *Components have provable identity*
  - Privacy
    - *Service does not see customer activities*

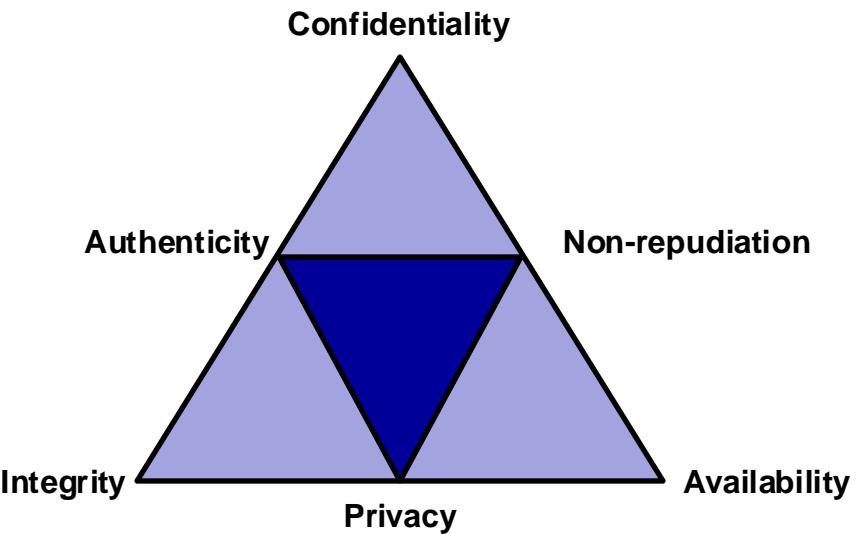
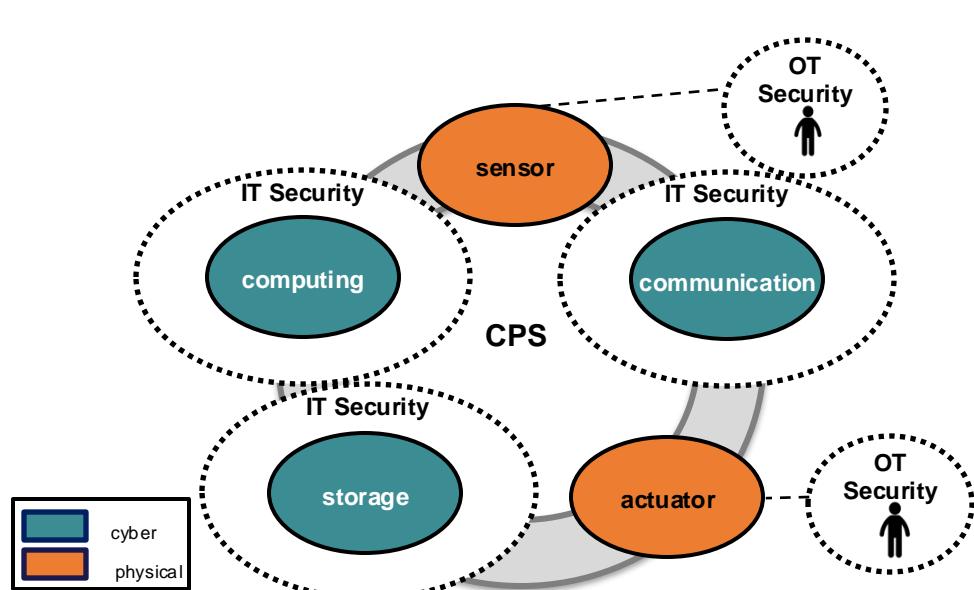


# Security Triad: Smart Card

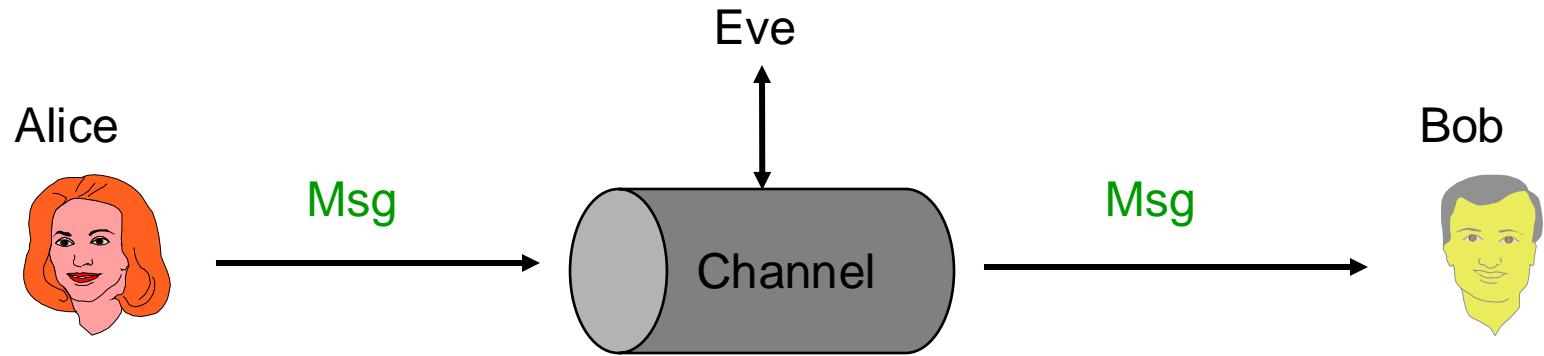


- Smart Card security
  - Confidentiality: *Data is secured*
  - Integrity: *Data is trusted*
  - Availability: *Data is accessible*
  - Non-repudiation: *Service has a trusted audit-trail*
  - Authenticity: *Components have provable identity*
  - Privacy: *Service does not see customer data*

# Security Triad for CPS

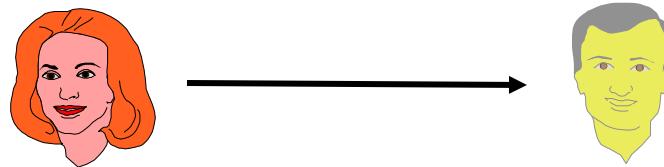


# Network Security Model

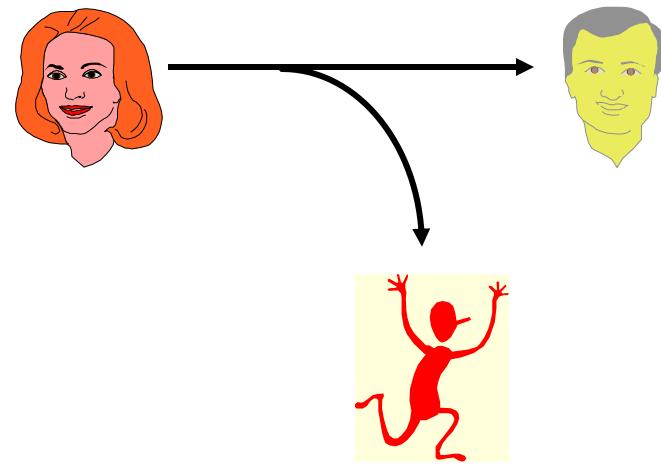


# Confidentiality (Secrecy)

- Protect transmitted data
- Protect against traffic analysis



- **INTERCEPTION**  
Unauthorised party gains access to data

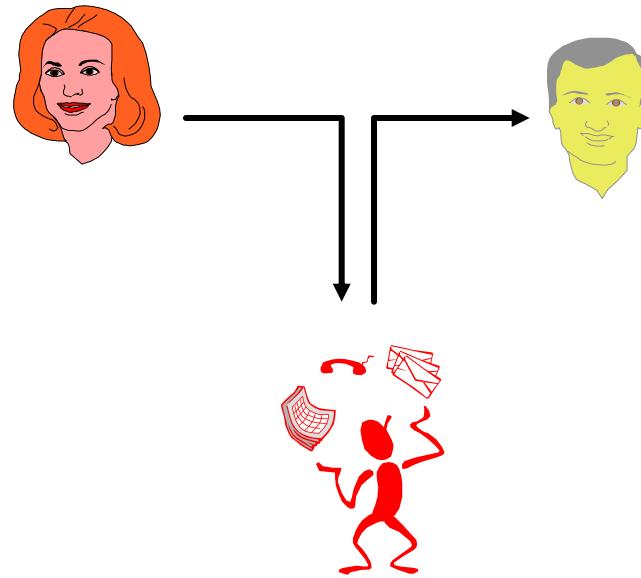


# Integrity

- Message is received as sent
- Modification
- Also interested in replay, re-ordering, deletion, delay

- **MODIFICATION**

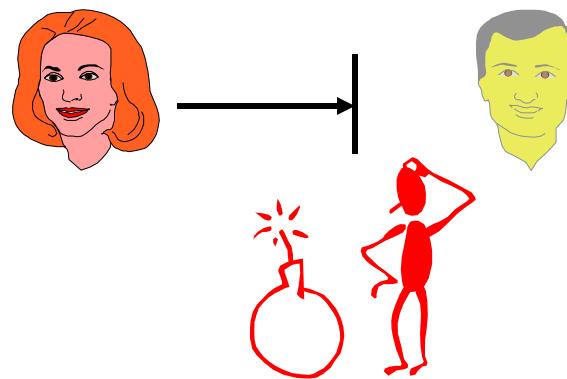
Gain access and “ tampers” with messages



# Availability

- Complete loss of availability
- Reduction/Degradation in availability

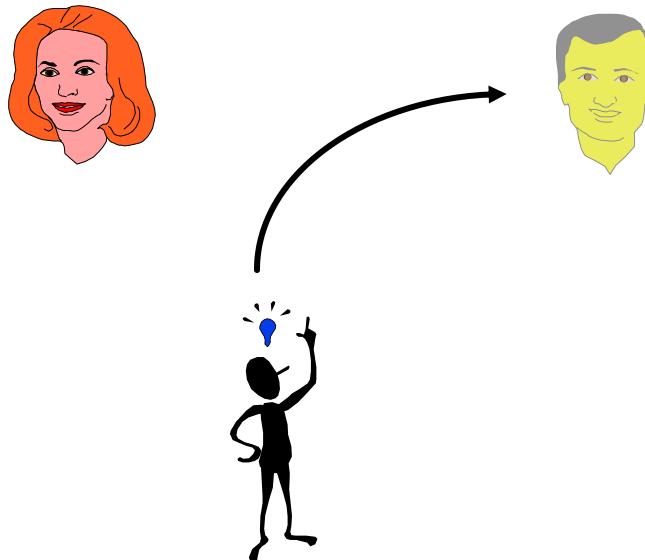
- **INTERRUPTION**  
Loss of communication (cut the cable)
- **DENIAL OF SERVICE**  
Noisy comms (physical noise, spurious messages)



# Authentication

- Assurance that message is from proper source
- Protect from third party masquerade

- **FABRICATION**  
Insertion of “counterfeit” messages



# Non-repudiation

- Prevents parties from denying they sent or received a message; i.e. concerned with protecting against legitimate protocol participants, not with protection from external source
- Receiver can verify **and prove** who sent a message
- Sender can verify **and prove** who received a message

- **REPUDIATION ATTEMPT**

Party anonymously publishes his or her message/key(s) and falsely claims that they were stolen.

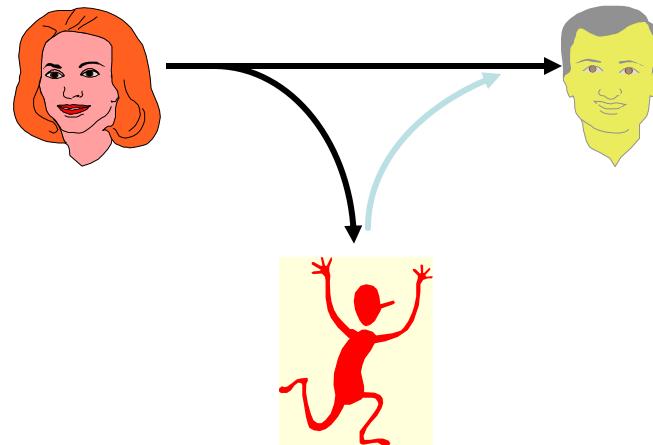


# Privacy

- Limit and control access to host system/services
- Limit and control access to networks
- Authenticate each party so that access rights can be assigned
- More fine-grained solutions, e.g. Digital Rights Management

- **REPLAY**

Record a legitimate message e.g. a login, and replay later



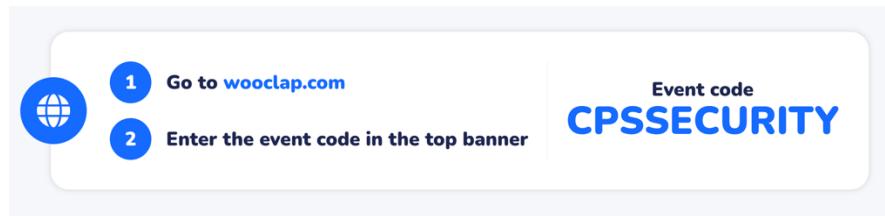
# Contents

- Security Triad



## *Cryptographic Primitives*

- **Private-Key Cryptography**
  - Hash Function, Message Authentication Code
  - Public-Key Cryptography
  - Digital Signature
- Discussion



# Kerckhoffs's Principle

- An encryption scheme should be secure even if enemy knows everything about it except the key
  - Attacker knows all algorithms
- Do not rely on secrecy of the algorithms (“security by obscurity”)



# SYMMETRIC-KEY CRYPTOGRAPHY

Symmetric-key cryptography started thousands of years ago when people needed to exchange secrets (for example, in a war).



7<sup>th</sup> Century BC, Ancient Greece

# Types of Symmetric-Key Cipher

- **Substitution Cipher**
  - A substitution cipher replaces one symbol with another.
  - Can be mono/poly-alphabetic

The following figures show two examples of a plaintext and its corresponding ciphertext. Which cipher is mono-alphabetic?

**Plaintext:** HELLO

**Ciphertext:** KHOOR

**Plaintext:** HELLO

**Ciphertext:** ABNZF

# Types of Symmetric-Key Cipher (contd.)

- **Shift Cipher**
  - The shift cipher is also referred to as the Caesar cipher.

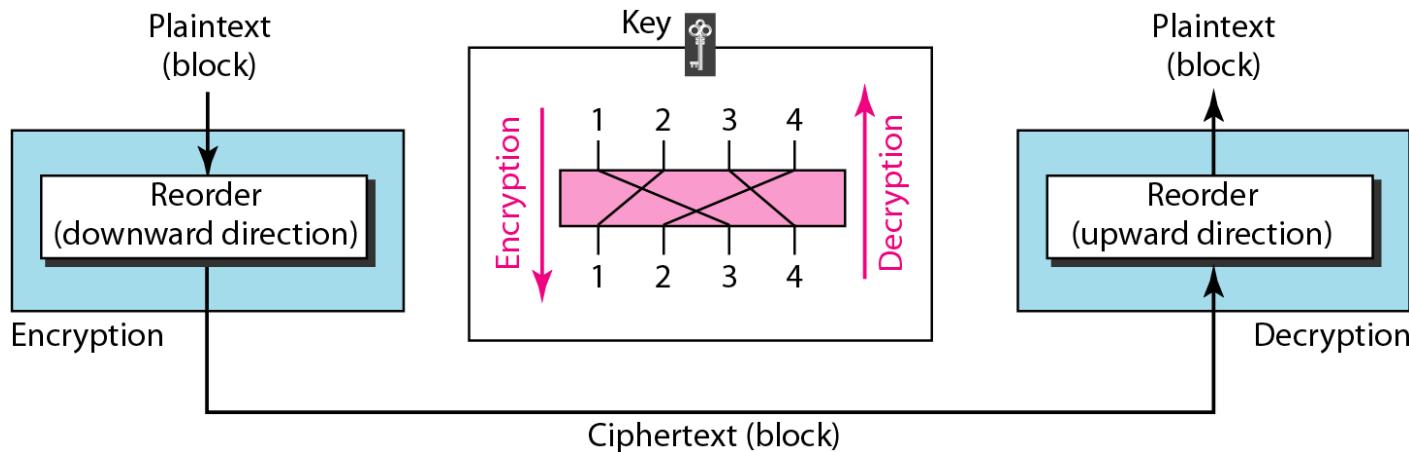
Use the shift cipher with key = 15 to encrypt the message “HELLO.”

Solution

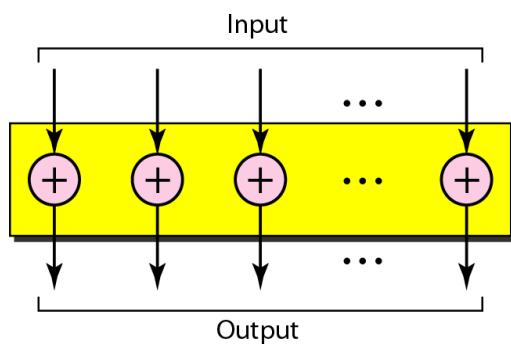
We encrypt one character at a time. Each character is shifted 15 characters down. Letter H is encrypted to W. Letter E is encrypted to T. The first L is encrypted to A. The second L is also encrypted to A. And O is encrypted to D. The cipher text is WTAAD.

# Types of Symmetric-Key Cipher (contd.)

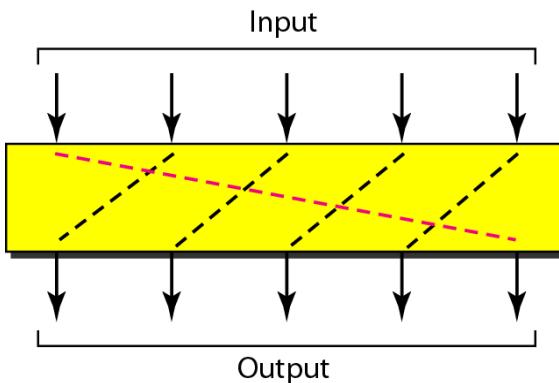
- **Transposition Cipher**
- A transposition cipher reorders (permutes) symbols in a block of symbols.



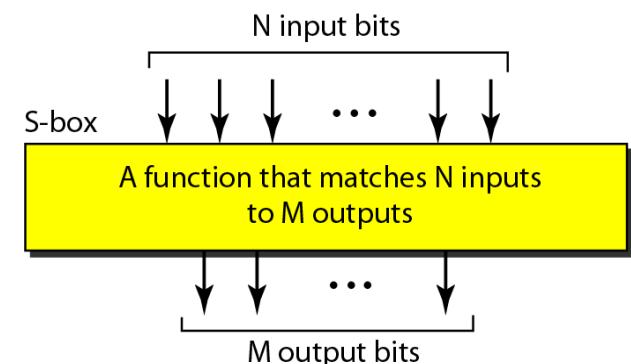
# Components of Symmetric-Key Cipher



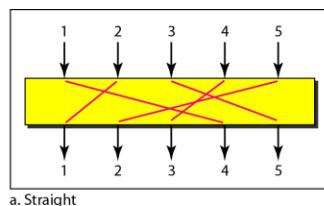
XOR substitution



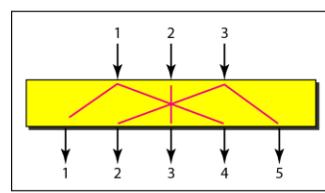
Rotation substitution



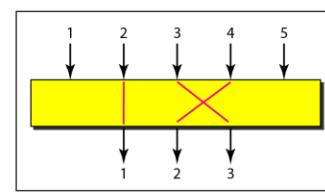
Substitution Box



a. Straight



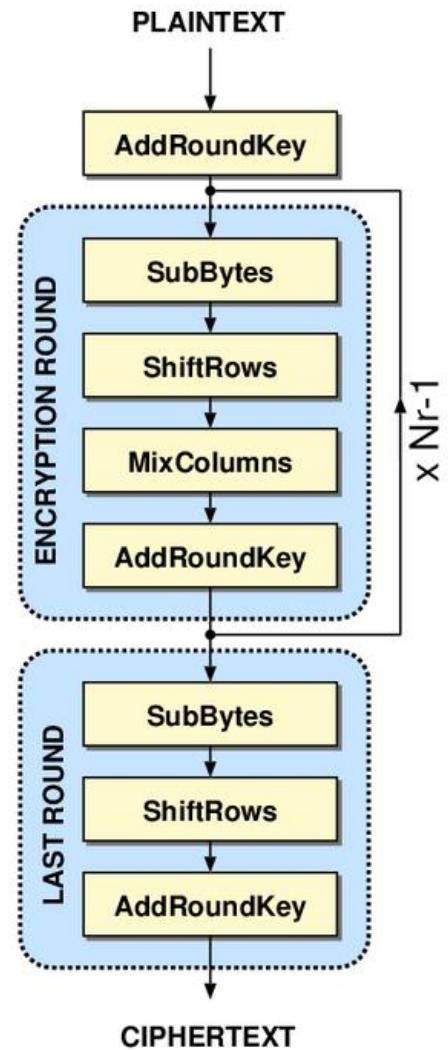
b. Expansion



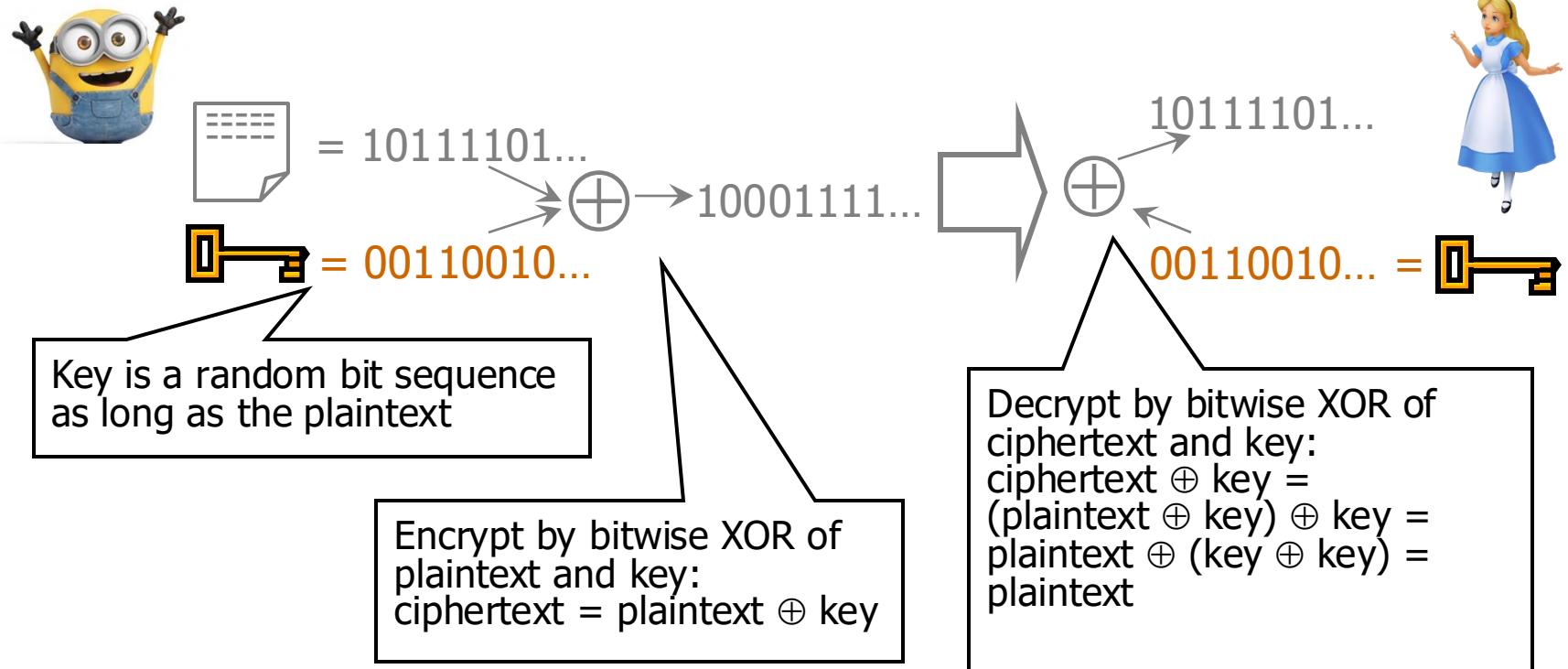
c. Compression

Permutation Box

# Block Cipher: AES



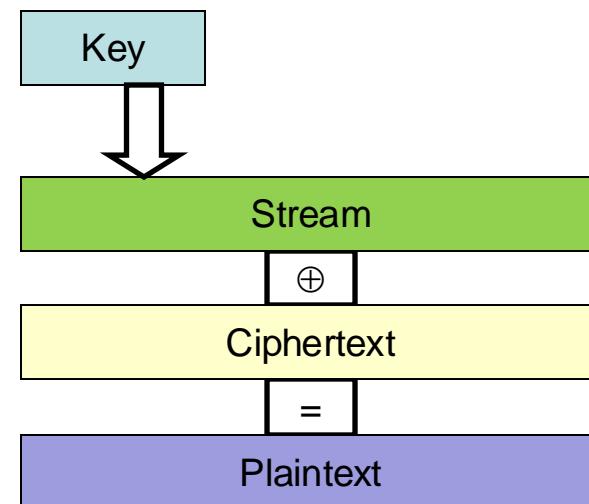
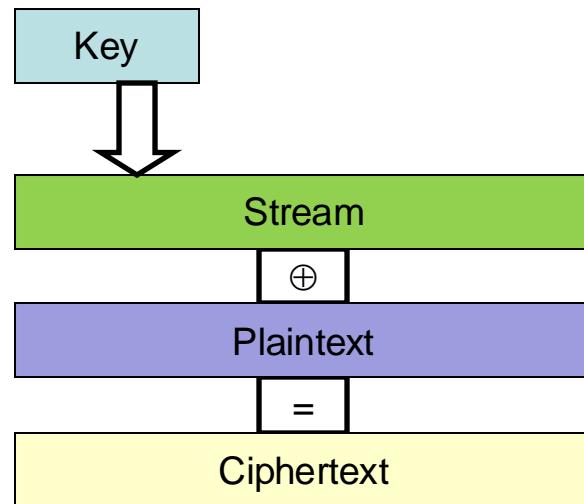
# One-Time Pad (Vernam Cipher)



Cipher achieves perfect secrecy if and only if there are as many possible keys as possible plaintexts, and every key is equally likely (Claude Shannon, 1949)

# Stream Ciphers

- Start with a secret key (“seed”)
- Generate a keying stream
  - Uses a Pseudo-Random Number Generator (PRNG)
- Combine the stream with the plaintext to produce the ciphertext (typically by XOR)



# Pseudo Random Number Generator

- Commonly constructed using Linear Feedback Shift Registers (LFSRs)
- An adversary cannot distinguish a PRNG output from a random sequence
- Example Randomness Tests
  - Monobit tests (frequency of 0s and 1s)
  - Run tests (frequency of runs of different lengths)
- *Pi: A naturally occurring random number*

String 1: 01  
String 2: 11001000011000011101110111011001111101001000010010101110010110

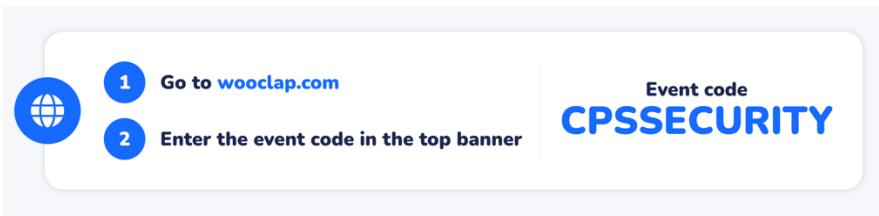
# Contents

- Security Triad



## *Cryptographic Primitives*

- Private-Key Cryptography
  - **Hash Function, Message Authentication Code**
  - Public-Key Cryptography
  - Digital Signature
- 
- Discussion



# How to Authenticate Messages?

- content is authentic – bits are as sent
- sequence of messages is proper

Note: Separate problems

- confidentiality that the message is encrypted
- authenticity that the sender is genuine

# Authenticating messages - Types

- **Authenticator** - a value that authenticates a message content
- **Message Authentication Code, MAC** (cryptographic checksum)
  - public function, with a given secret key produces fixed length value
- **Hash or Message Digest**
  - public function, which maps message (any length) to fixed-length hash value – no key

# Message Authentication Code (MAC)

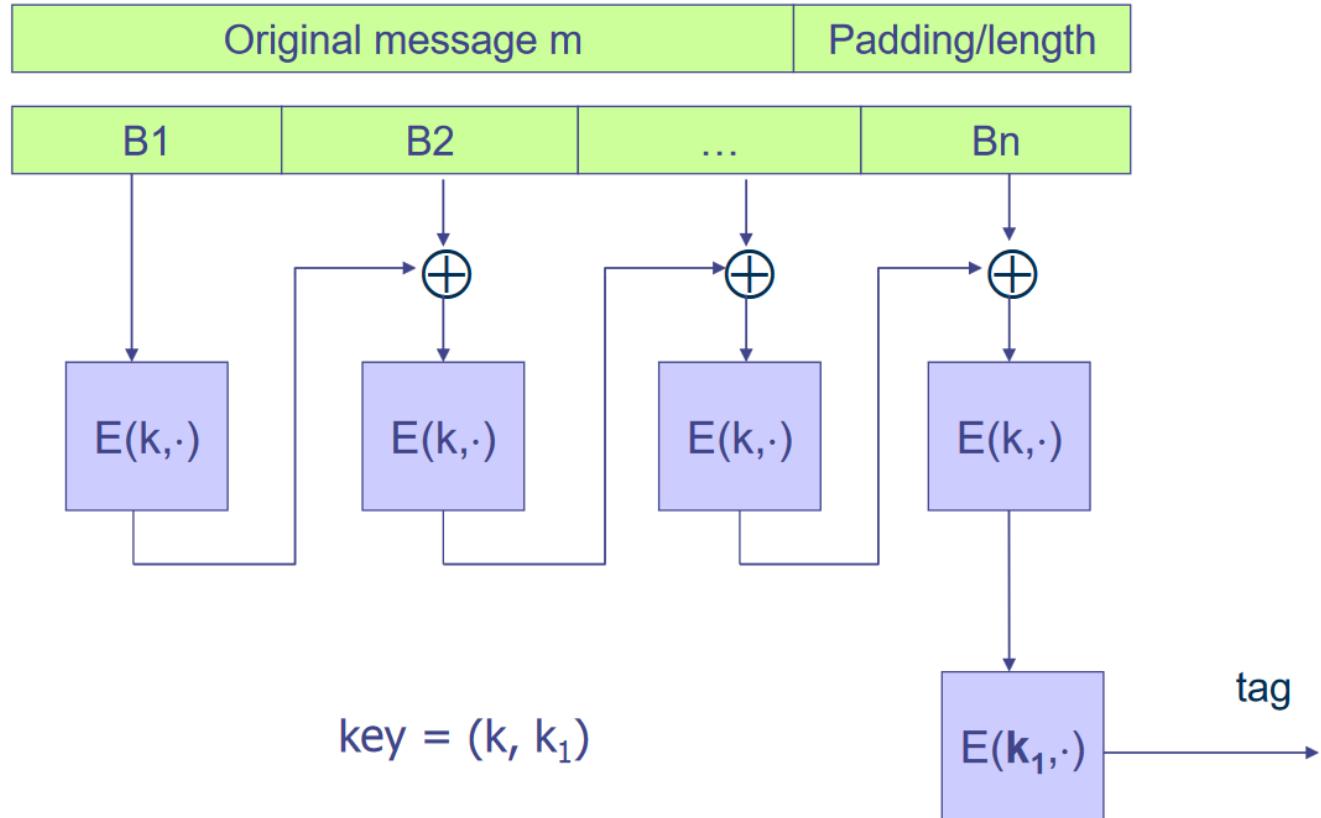
- Objective: assure authentic content
  - Message is not encrypted
- MAC is small fixed-size data block, appended to message M
- Generator and Validator share secret key

$$\text{MAC}_K[M] = F(M, \text{key})$$

# Why use MAC, given encryption?

- Message may not need to be secret, but must be authentic
- Broadcast - only have one site responsible for monitoring authenticity. Broadcast plaintext plus MAC
- Overload - send plaintext messages (i.e. receiver need not decrypt). Authenticate selectively

# MAC Construction using Encryption



Cipher Block Chaining Message Authentication Code (CBC-MAC)

# MAC usage

- message authentication (no confidentiality)
  - $A \rightarrow B: [M, MAC_K[M]]$
- authentication and confidentiality
  - $A \rightarrow B: E_{K1}[M, MAC_K[M]]$
- authentication and confidentiality
  - $A \rightarrow B: [E_{K1}[M], MAC_K[E_{K1}[M]]]$

# MAC Usage (contd.)

Encryption key  $K_E$     MAC key =  $K_I$

## Option 1: MAC-then-Encrypt (SSL)



## Option 2: Encrypt-then-MAC (IPsec)



## Option 3: Encrypt-and-MAC (SSH)



# Hash function (Message Digest)

- no key
- like MAC, small amount of data; hash of message gives fixed-size value
- define hash function so that change of any one bit of message will result in different hash value
- hash function is not secret
- one-way; receiver re-computes hash function

# Simple Message Digest: Parity

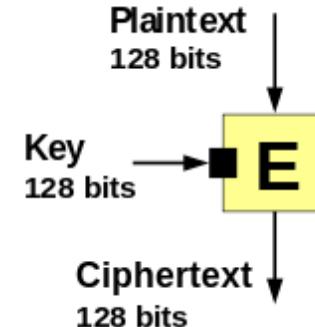
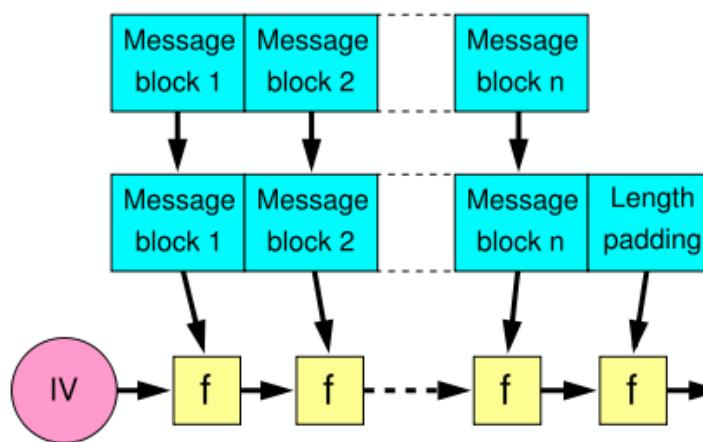
- parity, or longitudinal redundancy check
- process one bit at a time
- use XOR

$$C = b_1 \oplus b_2 \oplus \cdots \oplus b_{n-1} \oplus b_n$$

- input – arbitrary length message
- size of processed block – one bit
- output – one bit

# Merkle–Damgård Construction for Hash Functions

- Message is divided into fixed-size blocks and padded
- Use a one-way function  $f$ , which takes a chaining variable (of size of hash output) and a message block, and outputs the next chaining variable
- Final chaining variable is the hash value



$f$  built out of block cipher

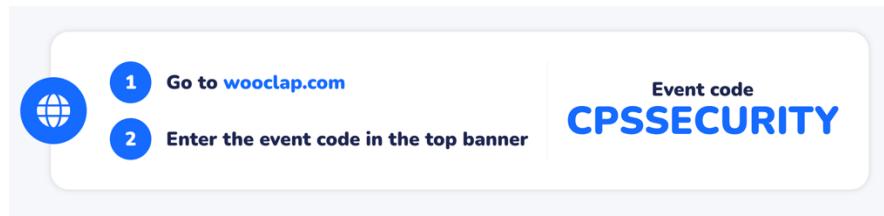
$$M = m_1 m_2 \dots m_n; C_0 = IV, C_{i+1} = f(C_i, m_i); H(M) = C_n$$

# Contents

- Security Triad
- Cryptographic Primitives

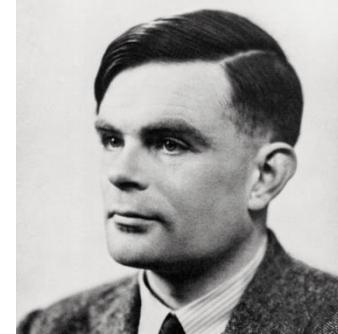
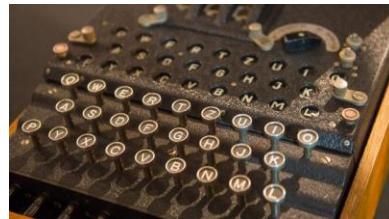


*Discussion*



# What did we learn?

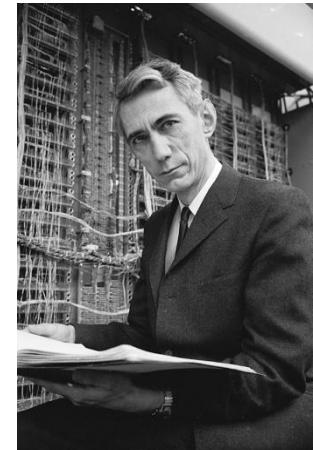
- **What is Security?**



- Confidentiality, Integrity, Availability
  - *Authenticity, Privacy, Non-repudiation*

- **What are the building blocks of security?**

- *Upcoming Lectures*
  - Private-Key Cryptography, Hash, MAC
  - Public-Key Cryptography, Digital Signature



# Further Reading

- Security Engineering
  - by Ross Anderson, available online -  
<http://www.cl.cam.ac.uk/~rja14/book.html>
- Handbook of Cryptography
  - by Alfred J. Menezes, Paul C. van Oorschot and Scott A. Vanstone, available online - <http://cacr.uwaterloo.ca/hac/>
- Applied Cryptography
  - by Dan Boneh and Victor Shoup, available online -  
<http://toccryptobook.us/>
- Leisure Reading - Simon Singh: *The Code Book*, Fourth Estate 1999

# The End