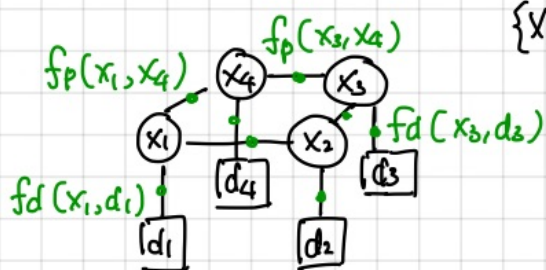


LECTURE 9: MRF OPTIMIZATION

[1] Motivation

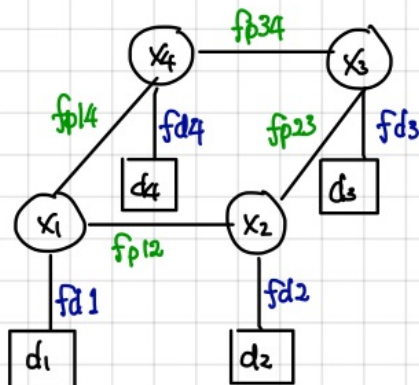
MRF :



$$\{x_1, \dots, x_n\}^* = \underset{\{x\}}{\operatorname{argmin}} \underbrace{\sum_i f_d(x_i, d_i) + \sum_{j \in N_i} f_p(x_i, x_j)}$$

the value of each term indicates the values of the edges in the graph.

Having defined the data term (f_d) & the prior term (f_p) for every node (pixel), we can now have a complete graph:



Q: How can we optimize the graph so that we can predict the labels of x_1, x_2, x_3, x_4 ?

Q: What does it mean by 'optimizing' a graph?

A: Example: $(x_1) - (x_2)$ & $x_i \in \{0, 1\}$

$$\begin{aligned} \{x_1, x_2\}^* &= \underset{\{x_1, x_2\}}{\operatorname{argmin}} f_p(x_1, x_2) \\ &= \underset{\{x_1, x_2\}}{\operatorname{argmin}} \left[f_p(x_1=0, x_2=0), f_p(x_1=1, x_2=0), \right. \\ &\quad \left. f_p(x_1=1, x_2=1), f_p(x_1=0, x_2=1) \right] \end{aligned}$$

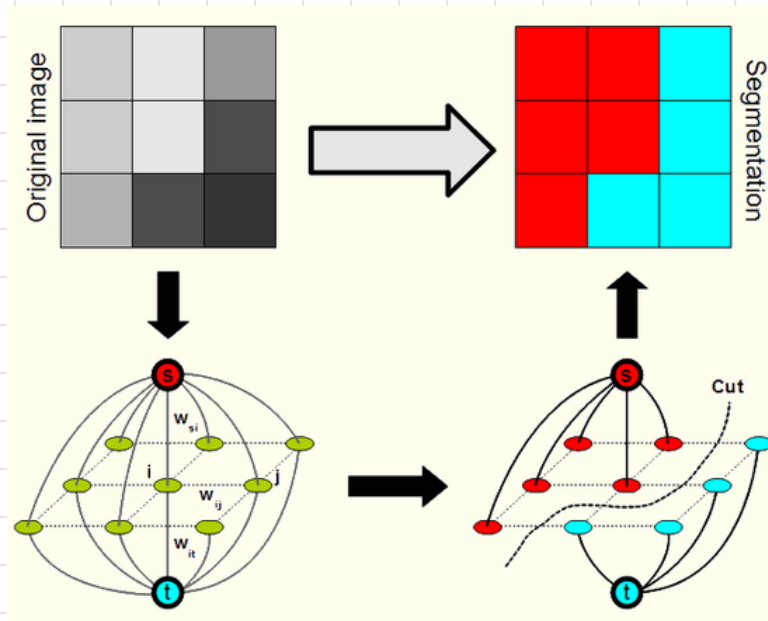
We find the values of x_1 & x_2 that minimize $f_p()$, or to choose one of four possible configurations.

[2] Graphcuts

#2

There are two methods of graph optimization: Graphcuts & Belief Propagation. Here, we focus on graphcuts, particularly on binary graphcuts.

Basic concept: Imagine we have a 3×3 image. Some of the pixels belong to the foreground & some belong to the background. However, due to noise/outlier and color variations, we are not sure & thus want to determine which pixels belongs to which label.

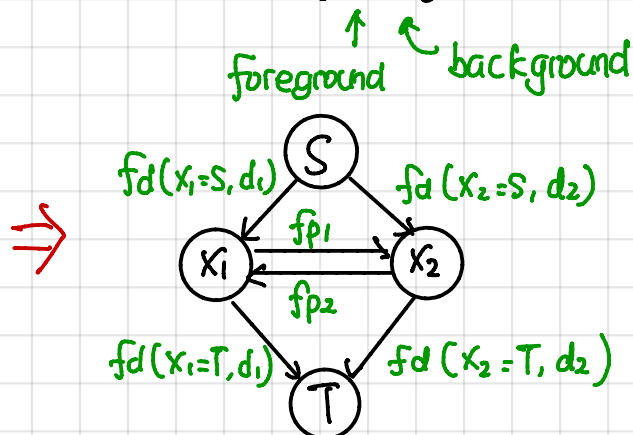
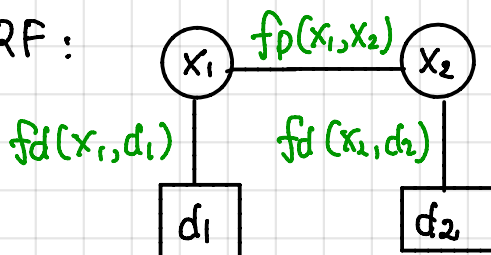


[3] Graphcuts : Data & Prior Terms

Q: How to assign the data & prior terms in a graph?

A: Suppose we have: x_1 & x_2 where $x_i = \{S, T\}$ and d_1, d_2

MRF:

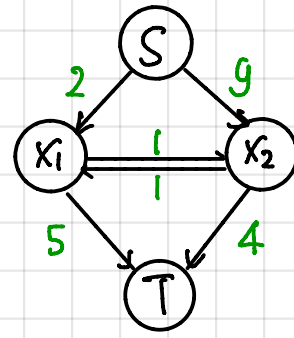
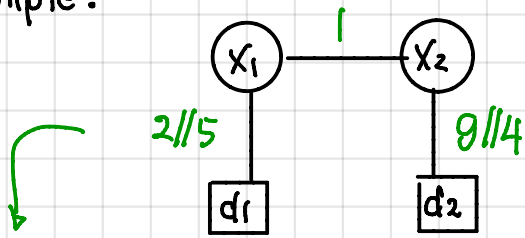


Assume:

$$f_p(x_1=T, x_2=S) = f_p(x_1=S, x_2=T) \rightarrow f_{p1} = f_{p2}$$

$$f_p(x_1=T, x_2=T) = f_p(x_1=S, x_2=S) = 0$$

Example:



$$fd(x_1 = S, d_1) = 2 \quad ; \quad fd(x_2 = S, d_2) = 9$$

$$fd(x_1 = T, d_1) = 5 \quad ; \quad fd(x_2 = T, d_2) = 4$$

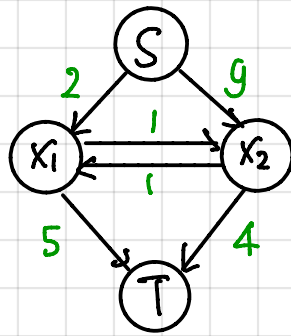
$$fp(x_1 = S, x_2 = T) = fp(x_1 = T, x_2 = S) = 1$$

$$fp(x_1 = S, x_2 = S) = fp(x_1 = T, x_2 = T) = 0$$

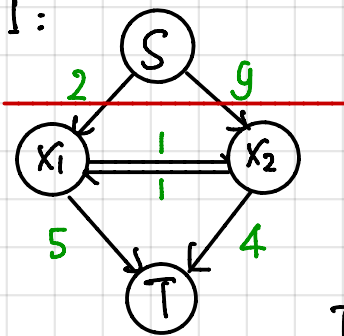
[4] Graphcuts: Max-Flow / Min-Cut

Q: How to optimize the graph, so that we know the labels of x_1 and x_2 ?

A:



cut 1:



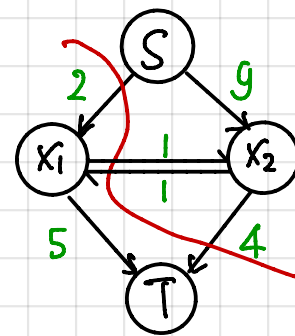
cut cost:
 $9 + 2 = 11$

This cut implies:

This is for minimization. For maximization: $x_1 = T$
 $x_2 = T$

$$\begin{cases} x_1 = S \\ x_2 = S \end{cases}$$

cut 2:



cut cost:
 $2 + 1 + 4 = 7$

Implies:

$$\begin{cases} x_1 = S \\ x_2 = T \end{cases}$$

Q: Which cut is better?

A: Cut 2 is cheaper than cut 1!

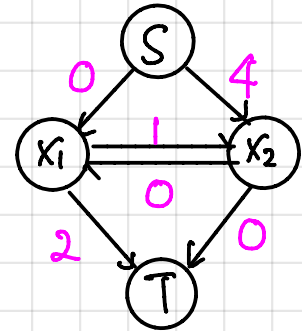
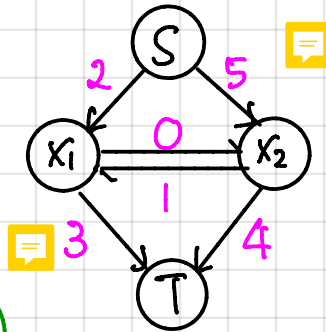
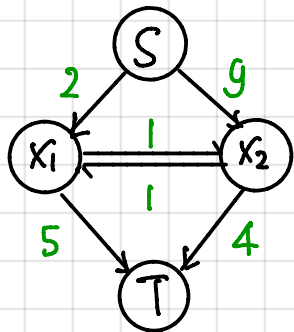
Q: How to find the minimum cut for all possible cuts?

[5] Graphs: Augmenting Path (the Ford - Fulkerson method)

#4

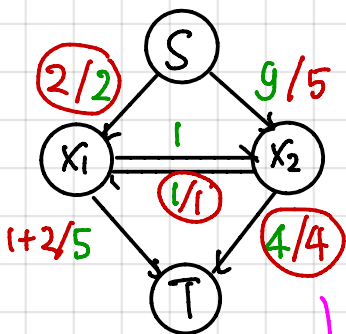
Basic components

- (1) Capacity graph
- (2) Flow graph
- (3) Residual graph

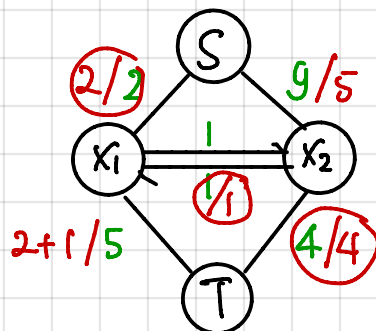


Calculation:

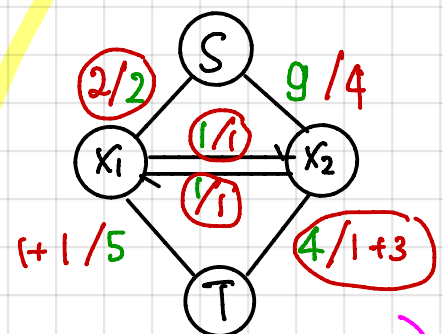
Start from 2:



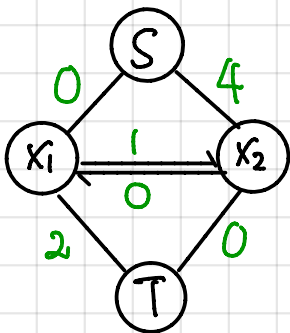
Alternative 1 (Start from 9):



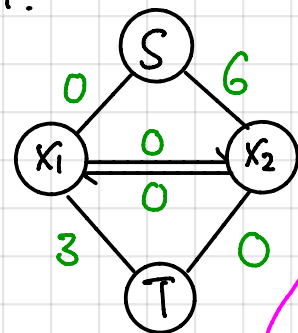
Alternative 2
(from 2 then 1):



Residual:



Residual:



Cut cost = $2 + 1 + 4 = 7$
(Min-cut)

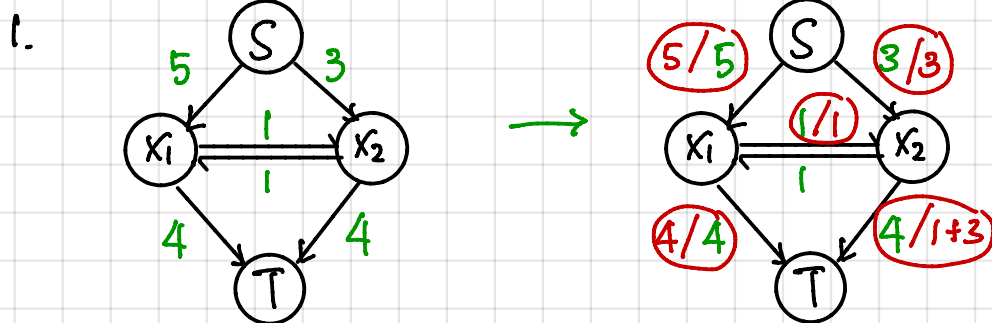
Max-flow = $2 + 5 = 3 + 4 = 7$

Must be the same!

Cut cost = $2 + 1 + 1 + 4 = 8$
(Min-cut)

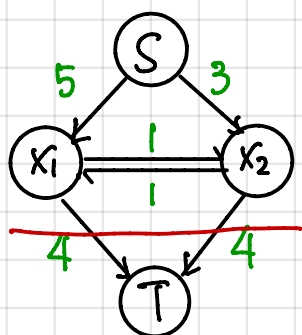
Max-flow = $2 + 4 = 1 + 1 + 1 + 3 = 6$

Other examples:



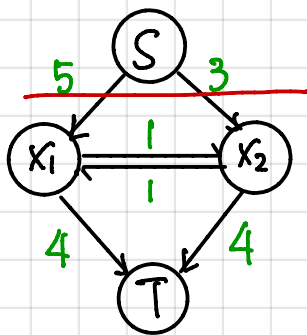
Min cut = 8
Max flow = 8

For this case, there are a few different ways of cutting:



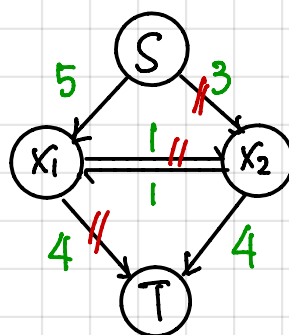
cut cost = 8

$x_1 = x_2 = T$



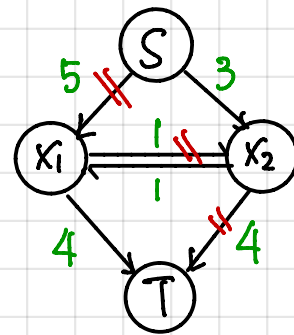
cut cost = 8

$x_1 = x_2 = S$



cut cost = 8

$x_1 = T; x_2 = S$



cut cost = 10

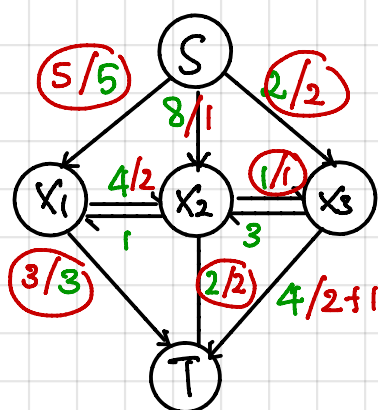
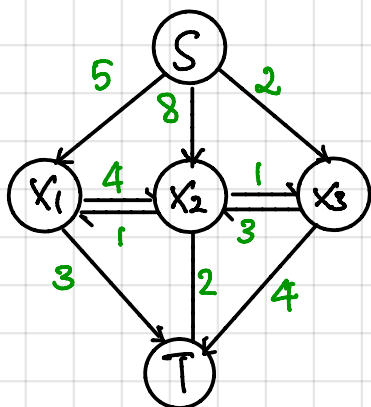
$x_1 = S, x_2 = T$

(not preferred because the cost is high)

all these labelings are correct.

This implies there are ambiguities in the labeling.

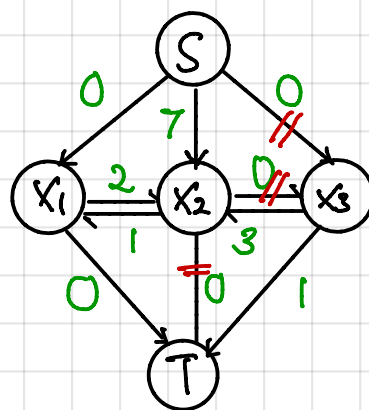
2.



$x_2 = T, x_3 = S \rightarrow \text{total cost} = 2 + 1 + 2 = 5$

$x_1 = \begin{cases} S \rightarrow \text{additional cost} = 5 \\ T \rightarrow \text{additional cost} = 3 \end{cases}$

We know the Max-flow = 8, thus $x_1 = T$



Aside from graphcuts, BP is another popular method for graph optimization. It's more general than graphcuts (can accept many labels) but much slower.

Formulation: $\xrightarrow{\text{Scalar Data term}}$ to prevent a very small / large value due to the product of m_{ji}

Belief: $b_i(x_i) = k \phi_i(x_i) \prod_{j \in N_i} m_{ji}(x_i)$

Message passing: $m_{ij}(x_j) = \sum_{x_i} \phi_i(x_i) \psi_{ij}(x_i, x_j) \prod_{k \in N_i/j} m_{ki}(x_i)$

\uparrow
Prior term

\downarrow
Neighbors of i except j .

Example: $(x_1) \text{---} (x_2) \quad ; \quad x_i = \{F, B\}$

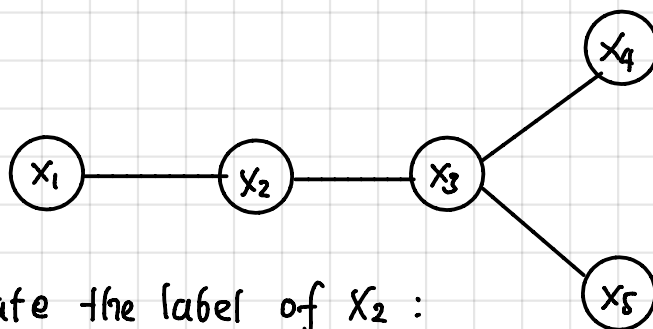
$$b(x_i = F) = k f_d(x_i = F, d) m_{21}(x_i = F)$$

$$\begin{aligned} m_{21}(x_i = F) &= \sum_{x_2 \in \{F, B\}} f_d(x_2, d_2) f_p(x_2, x_i = F) \\ &= f_d(x_2 = F, d_2) f_p(x_2 = F, x_i = F) + f_d(x_2 = B, d_2) f_p(x_2 = B, x_i = F) \end{aligned}$$

Aside from $b(x_i = F)$, we also compute $b(x_i = B)$, and compare $b(x_i = F)$ and $b(x_i = B)$.

If $b(x_i = F) > b(x_i = B)$, then we label $x_i = B$
 \downarrow
 For minimization

Another example:



Suppose we want to calculate the label of x_2 :

$$b_2(x_2) = k \phi_2(x_2) m_{12}(x_2) m_{32}(x_2)$$

$$m_{12}(x_2) = \sum_{x_1} \phi_1(x_1) \psi_{12}(x_1, x_2)$$

$$m_{32}(x_2) = \sum_{x_3} \phi_3(x_3) \psi_{32}(x_3, x_2) m_{43}(x_3) m_{53}(x_3)$$

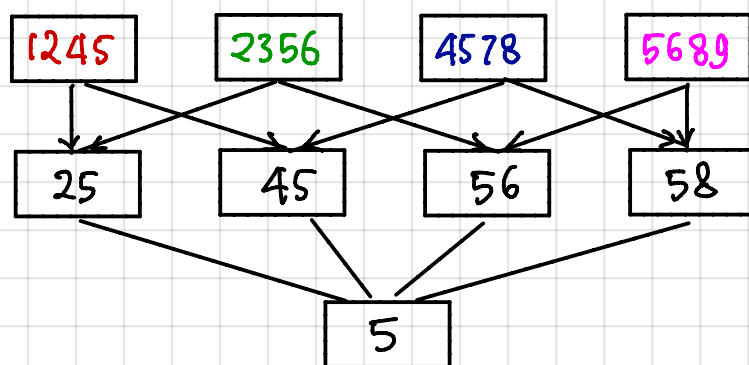
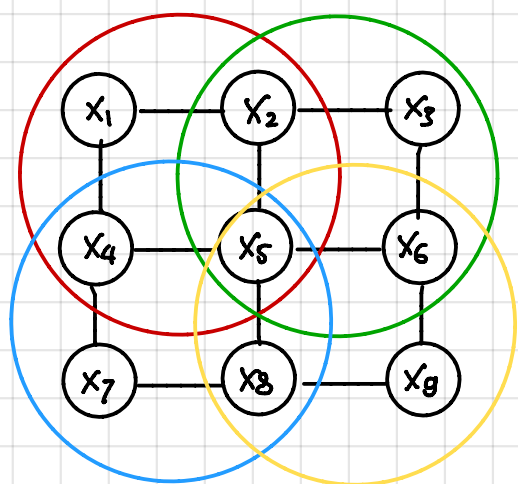
$$m_{43}(x_3) = \sum_{x_4} \phi_4(x_4) \psi_{43}(x_4, x_3)$$

$$m_{53}(x_3) = \sum_{x_5} \phi_5(x_5) \psi_{53}(x_5, x_3)$$

Note:

» The standard BP works only for a non-loopy graph.

For loopy graphs, one of the solutions is the Generalized BP.



$$b_5(x_5) = k \phi_5(x_5) m_{2 \rightarrow 5} m_{4 \rightarrow 5} m_{6 \rightarrow 5} m_{8 \rightarrow 5}$$

$$m_{4 \rightarrow 5}(x_5) = k \sum_{x_4} \phi_4(x_4) \psi_{45}(x_4, x_5) m_{12 \rightarrow 45} m_{78 \rightarrow 45}$$